

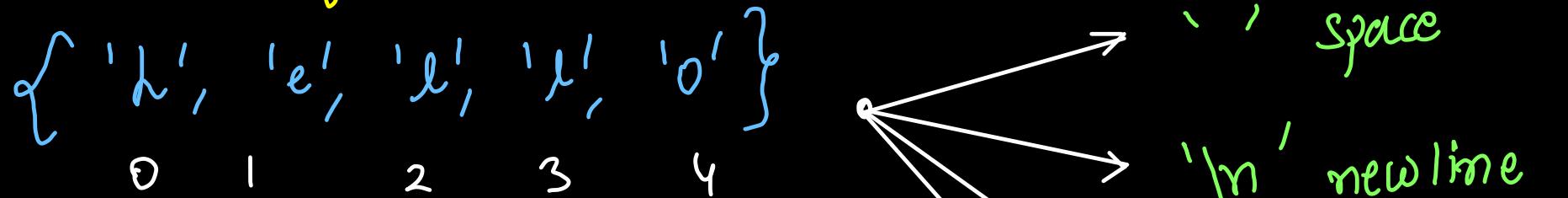
## Strings, StringBuilder & StringBuffer

- ✓ String Literal vs Object Memory Mapping
- ✓ Strings vs Character Array
- ✓ Implicit vs Explicit Null Initialization
- ✓ Inbuilt Methods in String Class  
(charAt, length, Constructors, etc)
- ✓ Concatenation Operation in String (Time)
- ✓ String Pool or String Interning
- ✓ String Immutability  
(Reference Variable vs Object Modification)
- ✓ String vs StringBuilder
- ✓ Operator == vs Method equals()
- ✓ StringBuilder vs StringBuffer
- ✓ Inbuilt Methods in StringBuilder  
(setCharAt, append)
- ✓ Wrapper Classes - Immutability

JSIC & Code  
native intern();

1D array  $\Rightarrow$  character array  
↳ 2 bytes (unicode)

Collection of characters  $\leadsto$  random index  $\rightarrow$  contiguous allocm



Disadvantages?

① size fixed  $\nearrow$  size  $\uparrow \times$   $\searrow$  size  $\downarrow \times$

② functions or methods (missing)  $\nearrow$  append  $\nearrow$  reverse  $\nearrow$  hashCode  $\nearrow$  equals  $\nearrow$  compare  $\nearrow$  concat(+)

# String Literal vs String object.

#  `String` 

`str = "hello world";`

heap area

reference variable (stack)

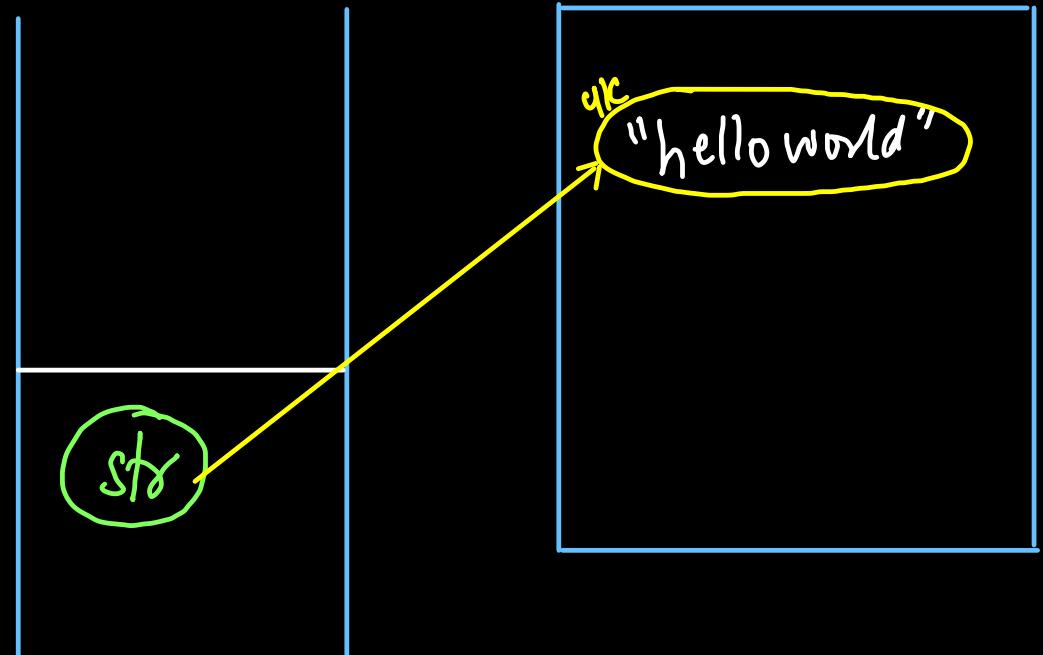
`System.out.println(str);`

↳ hello world

String literal / constant

stack

heap



# String str = new String ("hello world");

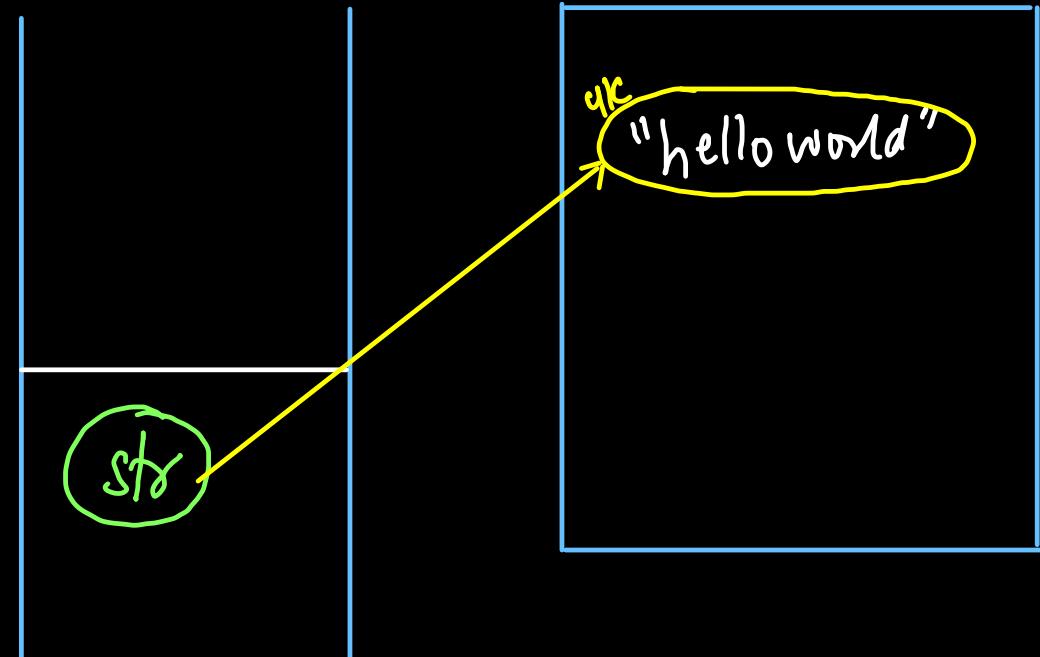
reference variable

string object

System.out.println(str);  
↳ hello world

stack

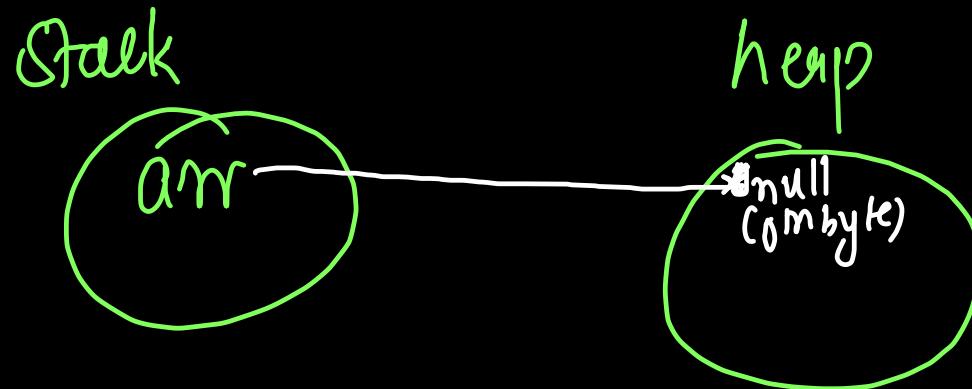
heap





```
int[] arr; // null (implicitly initialized)  
System.out.println(arr); // compilation error
```

```
int[] arr = null; // null (explicit initialization)  
System.out.println(arr); // null will be printed  
System.out.println(arr.length); // Null Pointer Exception
```



```
String str1;  
System.out.println(str1); // compilation error
```

Finished in N/A

Line 28: error: variable str1 might not have been initialized [in Main.java]

```
    System.out.println(str1);  
           ^
```

```
String str2 = null;  
System.out.println(str2);  
System.out.println(str2.length());
```

Finished in N/A

null

java.lang.NullPointerException: Cannot invoke "String.length()" because "<local1>" is null  
at line 32, Main.main

```
0 1 2 3 4  
h e l l o  
String str = "hello";
```

```
System.out.println(str); // hello -> horizontally
```

```
// hello -> vertically
```

```
for(int idx = 0; idx < str.length(); idx++){  
    System.out.println(str.charAt(idx));  
}
```

⇒ hello

⇒ h  
e  
l  
l  
o

String funcs →

length()	<del>length</del>
charAt()	<del>charAt</del>

```

String str = "";  $\rightarrow O(1)$ 
System.out.println(str);  $\rightarrow O(1) \rightarrow O(1)$ 
System.out.println(str.length());

```

$n = 26$

```

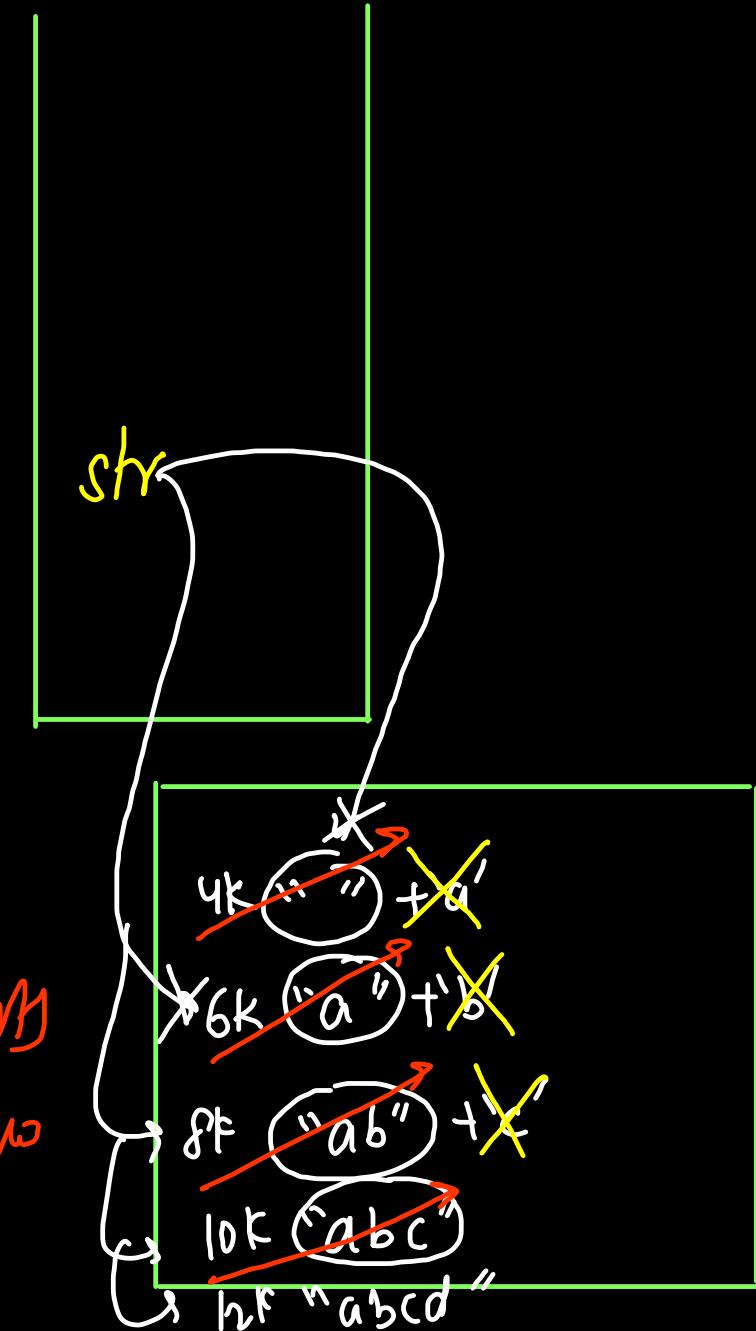
for(char ch = 'a'; ch <= 'z'; ch++){
    str = str + ch; // concatenation
    System.out.println(str);
}

```

~~$O(n^2)$~~   
 ~~$O(n)$~~   
 ~~$O(n)$~~   
 ~~$O(n)$~~   
 ~~$O(n)$~~   
ignoring point }

$a \rightarrow ab \rightarrow abc \rightarrow abcd \rightarrow abcde$

time  $\rightarrow$  copying bim part  
& forming new string



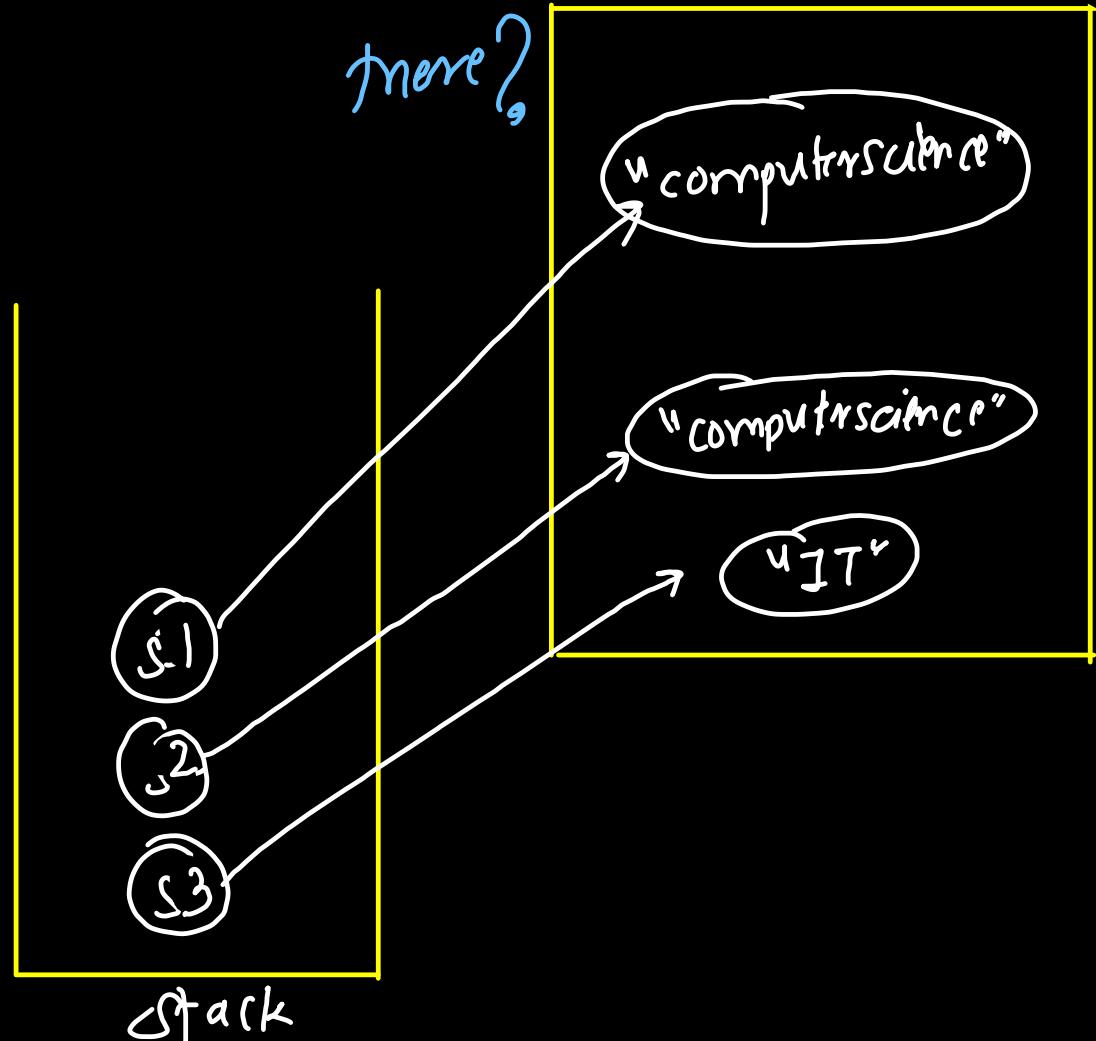
String s1dept = "Computer science"

String s2dept = "computer science"

String s3dept = "IT"

need?  
memory inefficient  
→ String interning / pool

What if string  
interning is not  
there?

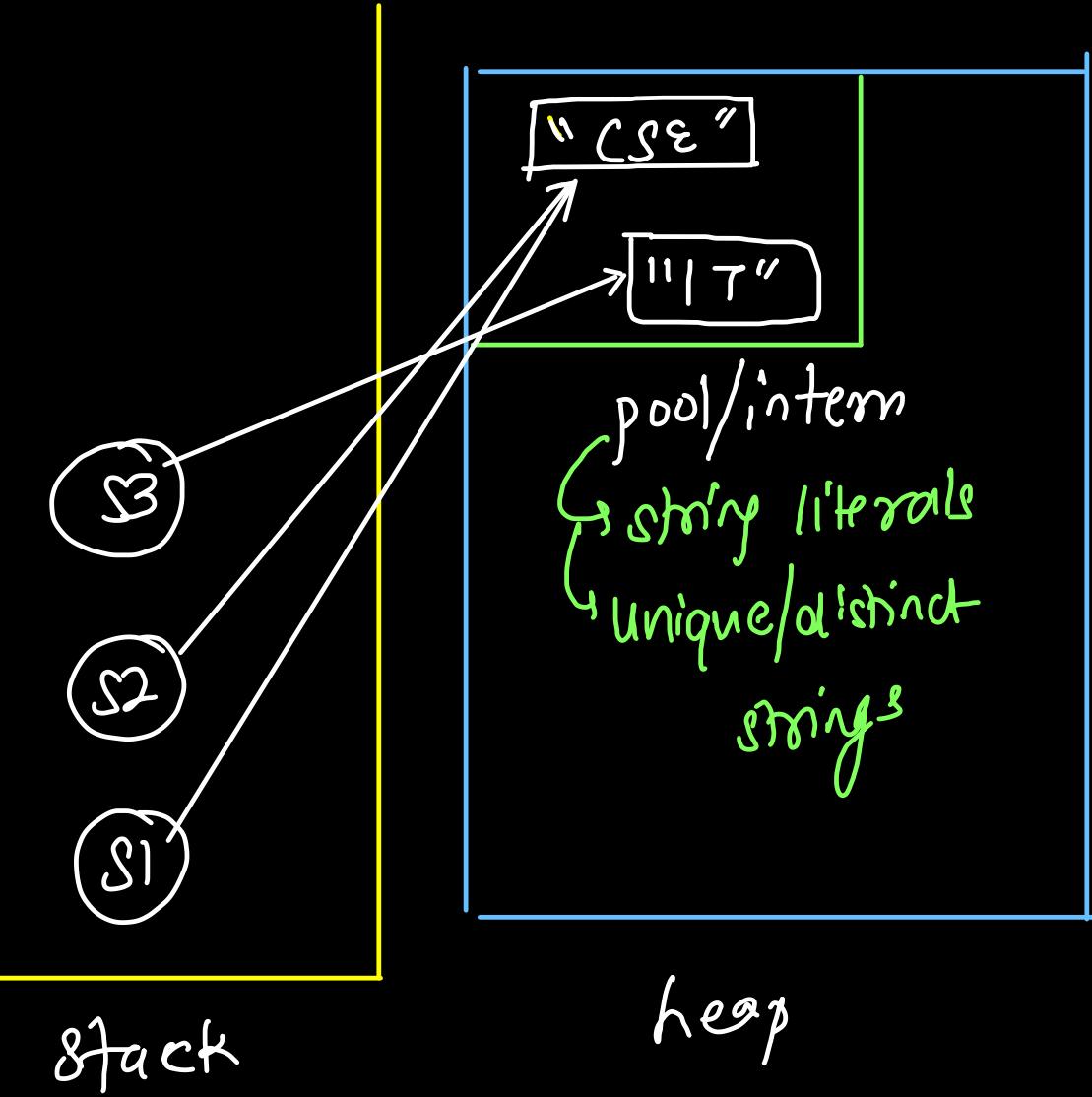


String s1 = "CSE";

String s2 = "CSE";

String s3 = "IT"

#advantage → reusability  
(string literals)  
↓  
memory efficient



# Implication → String Immutability

"Strings are **immutable** in Java"

→ cannot update/change/modify

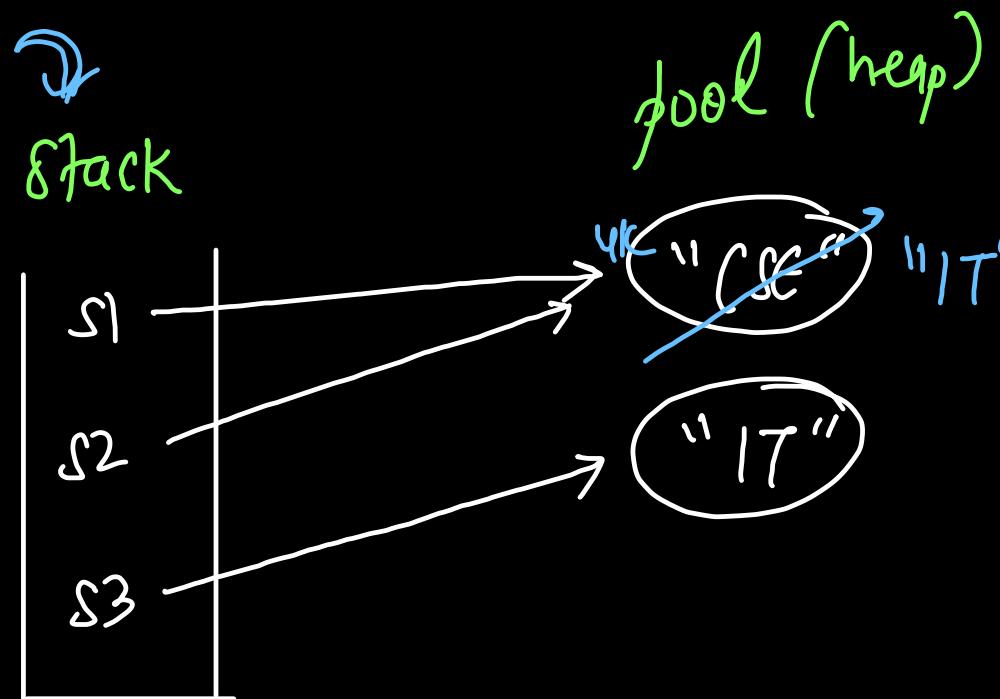
String immutability is not implemented?

String s1 = "CSE";

String s2 = "CSE";

String s3 = "IT";

s2 changed  
but not  
needed  
(inconsistent)



String s1 = "CSE";

String s2 = "CSE";

String s3 = "IT";

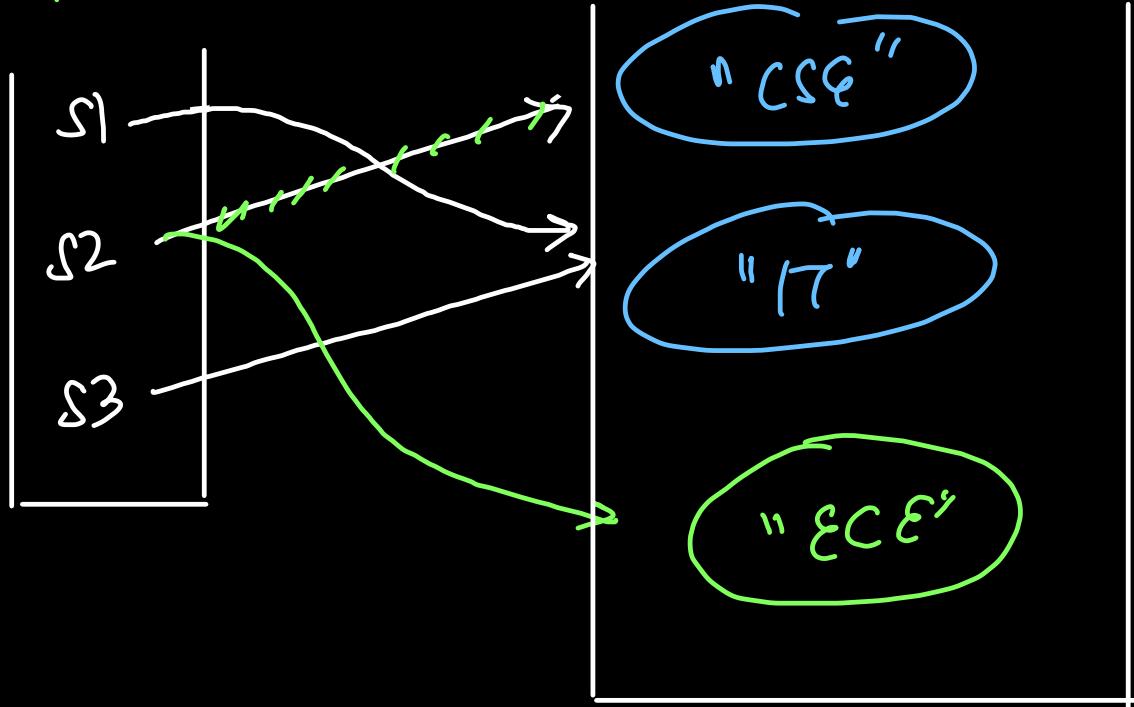
s1 = "IT"; ↵

s2 = "ECE"; ↵

} reference  
variable  
can be updated

stack

String literal/constant  
heap(pool)



memory wastage → inconsistency  
(interning X)

" Reference variables are mutable

but

; instance/object of heap are immutable"

variable ✓ reference ↘  
`str = "hello"` → object X  
`str[0] = 'd'`  
    { characters cannot be set }

) stop  
imm mutability  
new strings are created on every modification (time costly)

# StringBuilder & StringBuffer

read only  
characters

↓  
String

read & update  
characters

↓  
StringBuilder

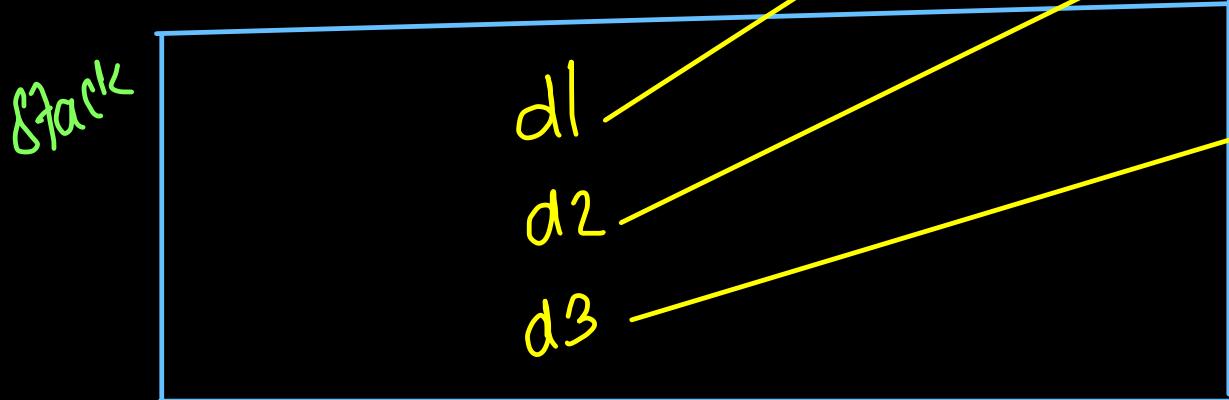
## String Builder

~~memory waste~~

StringBuilder d1 = new StringBuilder("CSE");

StringBuilder d2 = new StringBuilder("CSE");

StringBuilder d3 = new StringBuilder("IT");



# No interning # No immutability

```
StringBuilder d1 = new StringBuilder("CSE");
System.out.println(d1.hashCode());

StringBuilder d2 = new StringBuilder("CSE");
System.out.println(d2.hashCode());

StringBuilder d3 = new StringBuilder("IT");
System.out.println(d3.hashCode());

d1.append(" AI");
System.out.println(d1.hashCode());

System.out.println(d1 + " " + d2);
```

1828972342  
1452126962  
931919113  
1828972342  
CSE AI CSE

```
String s1 = "CSE";
System.out.println(s1.hashCode());

String s2 = "CSE";
System.out.println(s2.hashCode());

String s3 = "IT";
System.out.println(s3.hashCode());

s1 = "COE";
System.out.println(s1.hashCode() + " " + s2.hashCode());
```

67029  
67029  
2347  
66905 67029

```
|  
d1.setCharAt(1, 'O'); // update or set: O(1)  
  
d1.deleteCharAt(0); // remove first char: O(n)  
d1.deleteCharAt(d1.length() - 1); // remove last char: O(1)
```

```
StringBuilder d1 = new StringBuilder("CSE"); // n1 length  
// System.out.println(d1.hashCode());
```

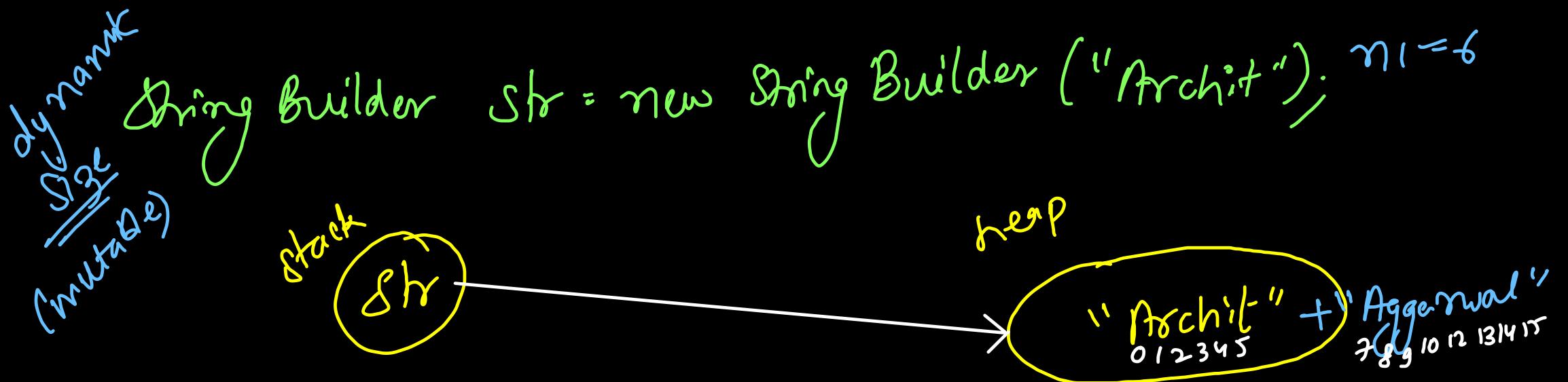
```
d1.append(" AI"); // concatenate n2 length
```

✗ O( $n_1 + n_2$ )

✗ O( $n_1$ )

✗ O( $n_2$ )

✗ O(1)

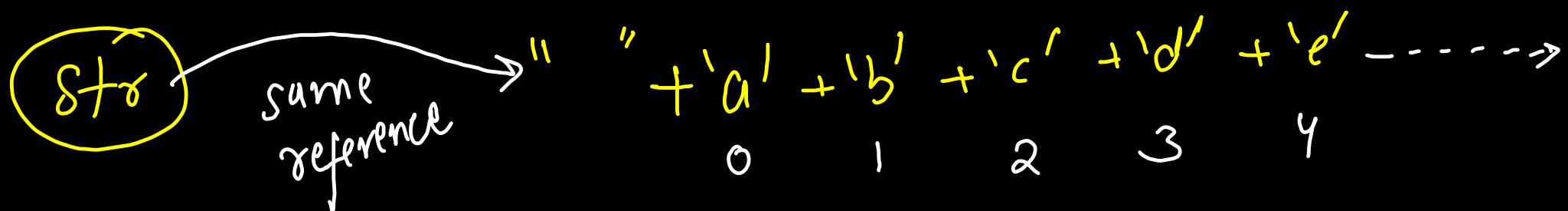


$str.append("Aggarwal"); \Rightarrow O(n^2)$  time  
 $n_2=8$

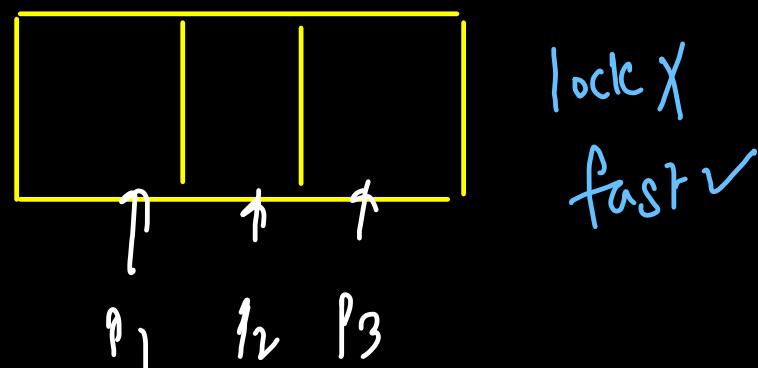
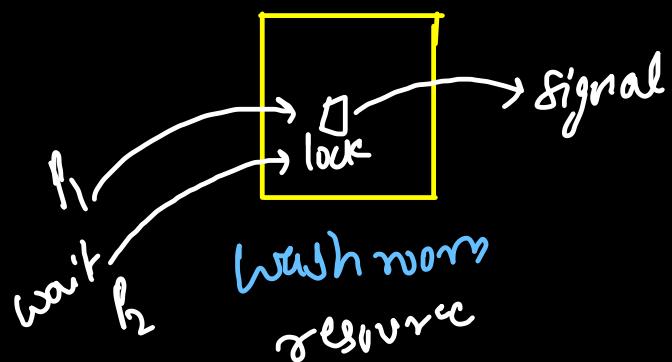
$str.append('Z');$   $\Rightarrow O(1)$  constant

*StringBuilder shr = new StringBuilder();*

```
// N = 26      O(n) time (where n=26)
for(char ch = 'a'; ch <= 'z'; ch++){
    d1.append(ch);
}
|           ↳ O(1) constant
```



No.	<b>StringBuffer</b> <i>(multithreading)</i>	<b>StringBuilder</b>
1)	StringBuffer is <u>synchronized</u> i.e. thread safe. It means two threads can't call the methods of StringBuffer simultaneously.	StringBuilder is <u>non-synchronized</u> i.e. not thread safe. It means two threads can call the methods of StringBuilder simultaneously.
2)	StringBuffer is <u>less efficient</u> than StringBuilder.	StringBuilder is <u>more efficient</u> than StringBuffer.
3)	StringBuffer was introduced in Java 1.0	StringBuilder was introduced in Java 1.5



```

String s1 = "CSE";
String s2 = "CSE";
String s3 = "IT";
} } interning ✓ immutability ✓

```

```

// == operator: string literal 4k == 4k ✓
System.out.println(s1 == s2); // true
System.out.println(s1 == s3); // false
4k ≠ 6k X

```

```

String s4 = new String("CSE");
String s5 = new String("CSE");
String s6 = new String("IT");
} } immutability ✓
} interning X

```

```

System.out.println(s4 == s5); // false
System.out.println(s1 == s4); // false
System.out.println(s4 == s6); // false
8k ≠ 10k
8k ≠ 8k
8k ≠ 12k

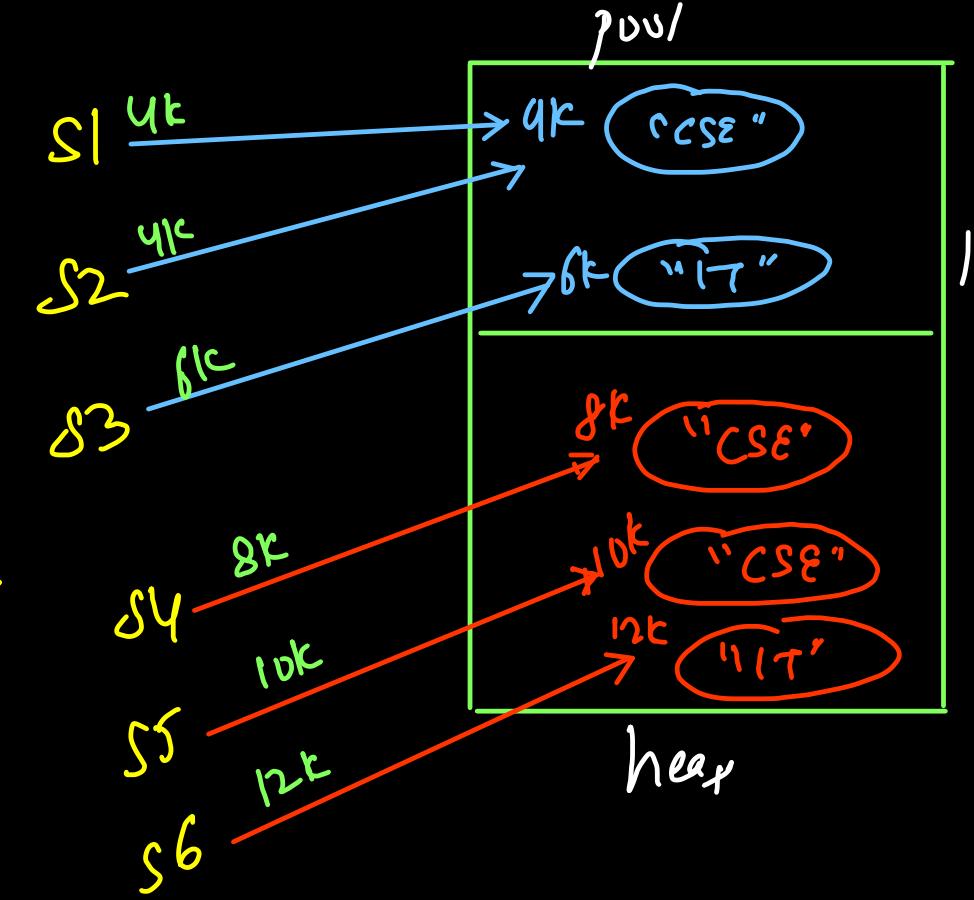
```

$\Rightarrow$

$= =$  operator:  
reference

same  $\rightarrow$  true  
diff  $\rightarrow$  false

need new objects on every new string  $\rightarrow$  backward compatibility



```

String s1 = "CSE";
String s2 = "CSE";
String s3 = "IT";

// == operator: string literal
System.out.println(s1 == s2); // true
System.out.println(s1 == s3); // false

// equals function: data
System.out.println(s1.equals(s2)); // true
System.out.println(s1.equals(s3)); // false

String s4 = new String("CSE");
String s5 = new String("CSE");
String s6 = new String("IT");

System.out.println(s4 == s5); // false
System.out.println(s1 == s4); // false
System.out.println(s4 == s6); // false

System.out.println(s4.equals(s5)); // true
System.out.println(s1.equals(s4)); // true
System.out.println(s4.equals(s6)); // false

```

Finished in 79 ms

true
false
true
false
false
false
false
true
true
false

~~O(1)~~ == operator  
compares on reference  
~~O(n)~~ worst case equals()  
compared first on reference, then on object (data)

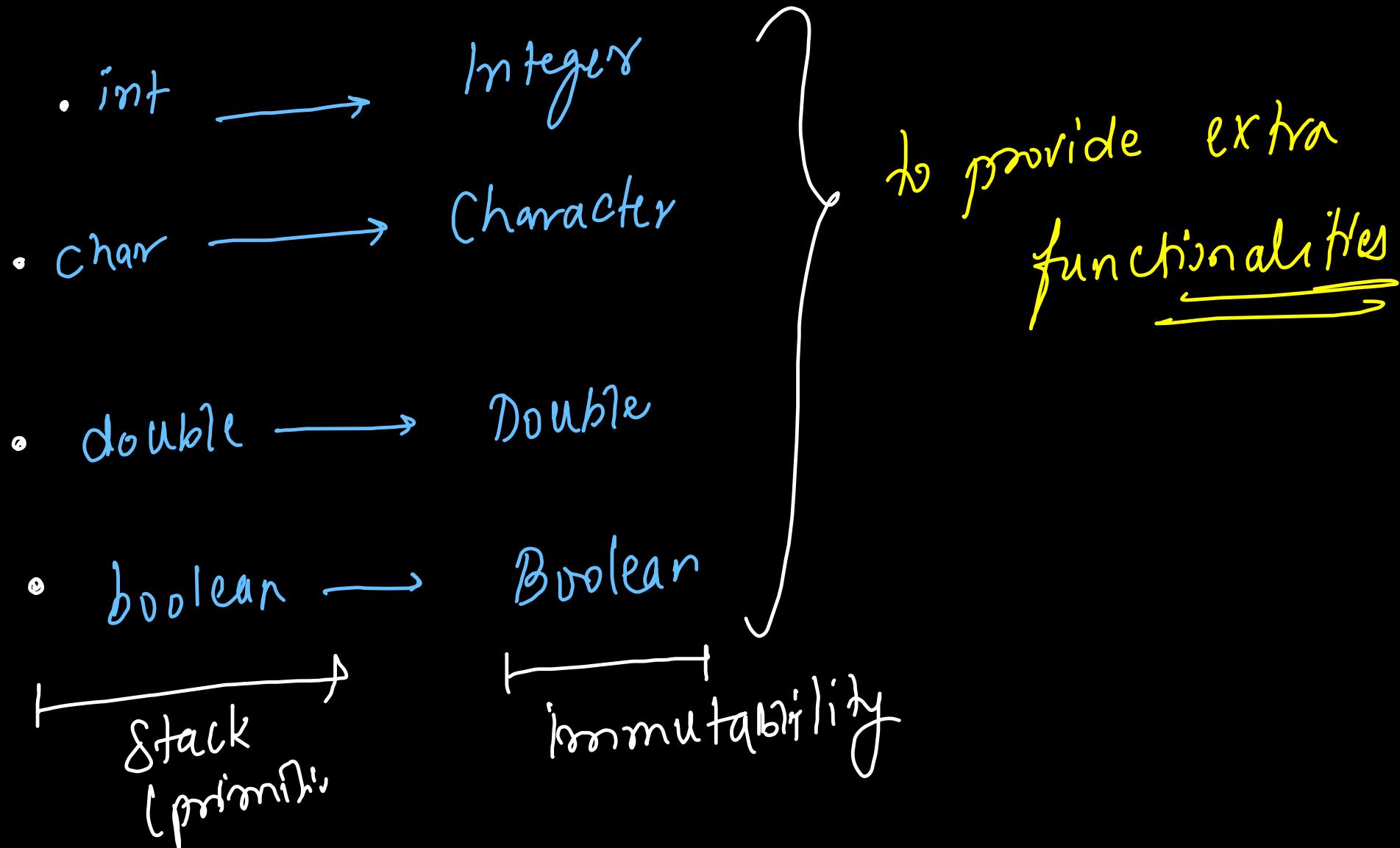
```
public boolean equals( String s1, String s2) {  
    if (s1 == s2) return true; // O(1)  
    if (s1.length() != s2.length()) return false; // O(1)  
    for (int i=0; i<s1.length(); i++) {  
        if (s1.charAt(i) != s2.charAt(i))  
            return false;  
    }  
    return true;  
}
```

*best case reference*

*O(n)*

*data*

Wrapper class (reference  $\Rightarrow$  stack, object  $\Rightarrow$  heap)



```

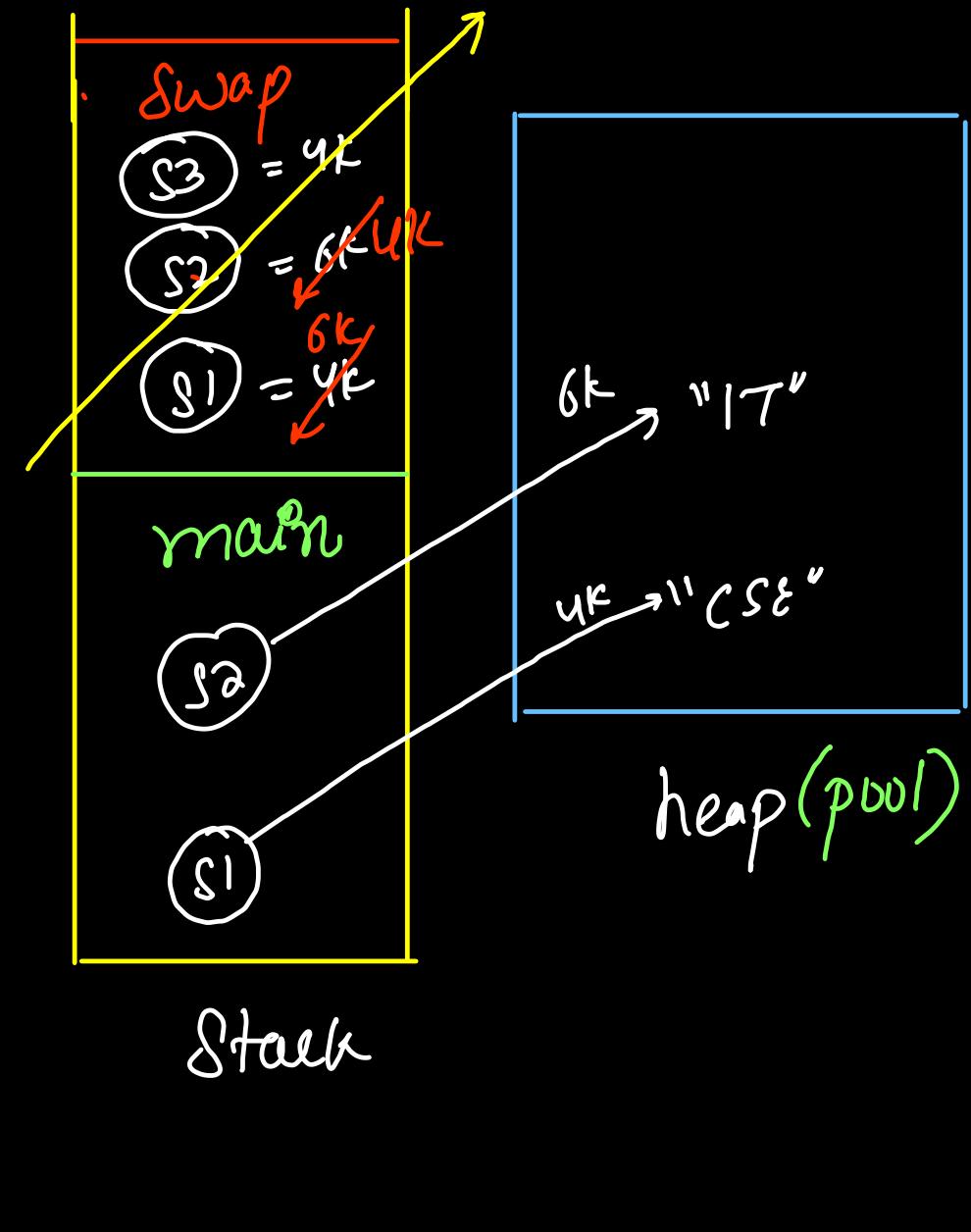
public class Main {
    public static void main(String[] args) {
        String s1 = "CSE";
        String s2 = "IT";
        System.out.println(s1 + " " + s2);
        swap(s1, s2);           no swap → immutable
        System.out.println(s1 + " " + s2);  String
    }
}

```

```

public static void swap(String s1, String s2){
    System.out.println(s1 + " " + s2);
    String s3 = s1;      CSE      IT
    s1 = s2;
    s2 = s3;
    System.out.println(s1 + " " + s2);
}

```



```

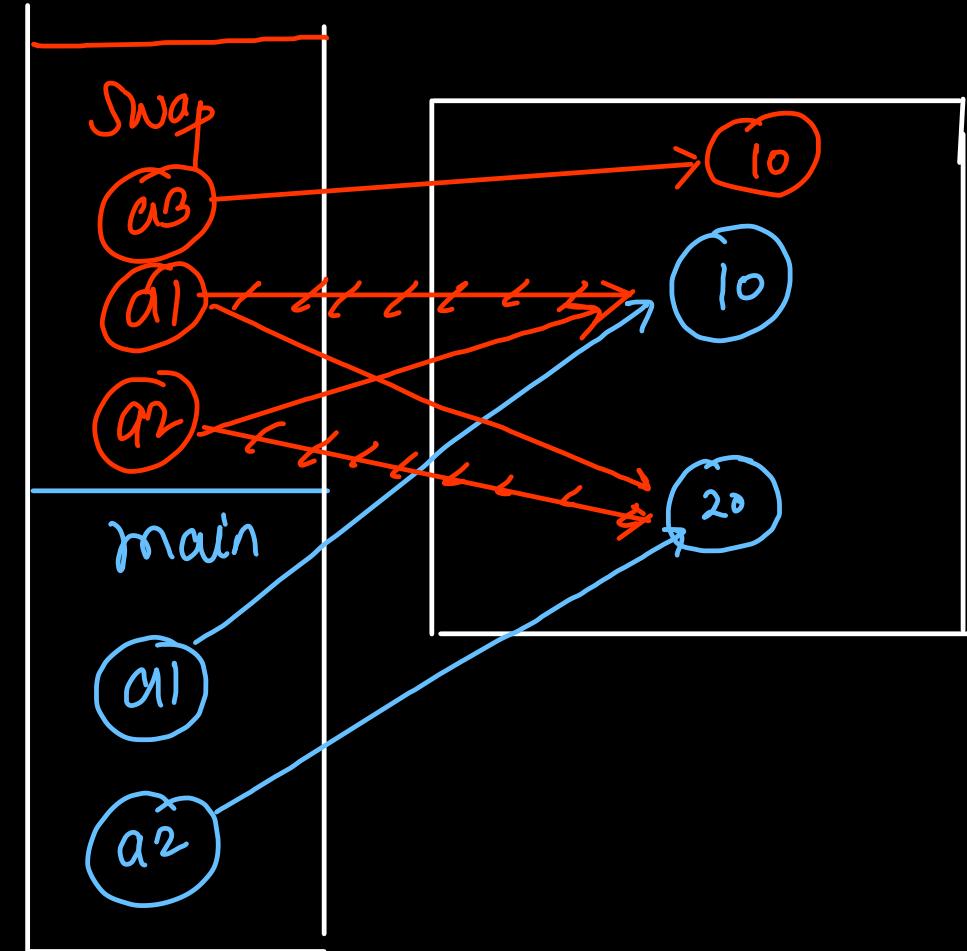
public class Main {
    public static void main(String[] args) {
        Integer a1 = 10;
        Integer a2 = 20;
        System.out.println(a1 + " " + a2);
        swap(a1, a2);
        System.out.println(a1 + " " + a2);
    }
}

public static void swap(Integer a1, Integer a2){
    System.out.println(a1 + " " + a2);
    Integer a3 = new Integer(a1);
    a1 = a2;
    a2 = a3;
    System.out.println(a1 + " " + a2);
}

```

10                  20

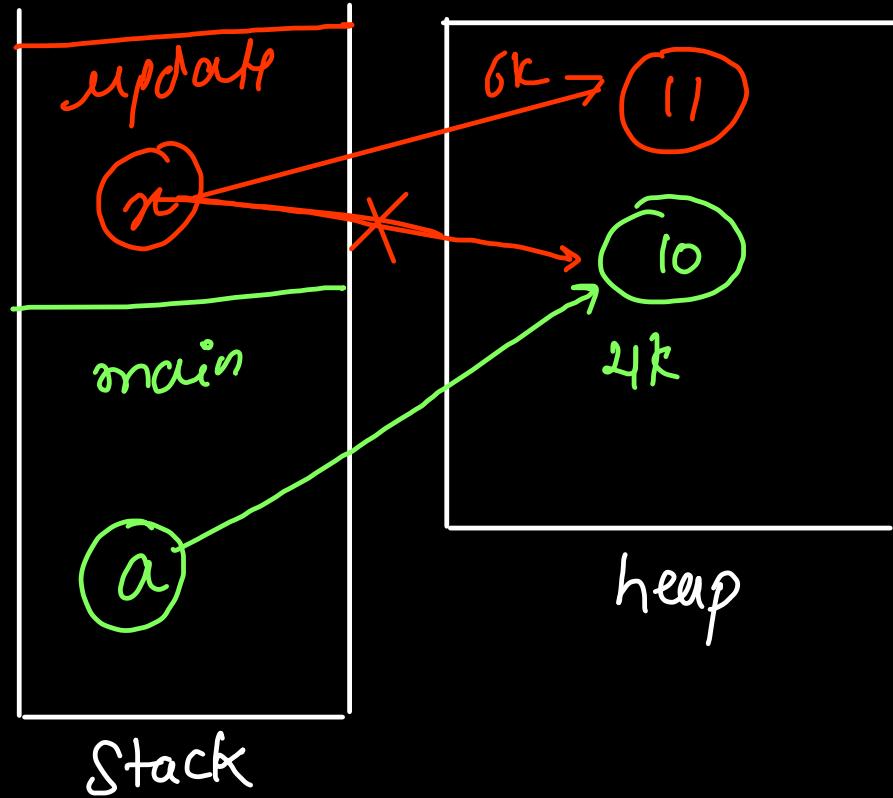
20                  10



"Wrapper classes are immutable in Java"

```
public class Main {  
    public static void main(String[] args) {  
  
        Integer a = 10;  
        System.out.println(a);  
        update(a);  
        System.out.println(a);  
    }  
}
```

```
public static void update(Integer x){  
    System.out.println(x);  
    x++;  
    System.out.println(x);  
}
```



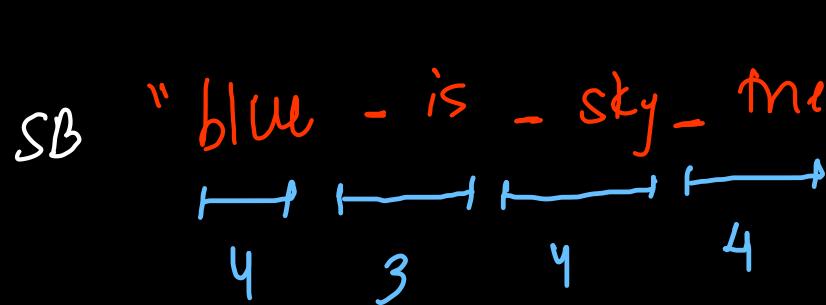
## Lc(51) Reverse Words

eg1) "the\slay\is\blue" → "blue.is.sky.the"

eg2) "    the ... slay .. is ... blue   "  
      ↑  ↑  
      leading  trailing  
      space  spaces

approach ① Split the string based on spaces.

array of strings  $\rightarrow \{ "true", "sky", "is", "blue" \}$   
str.split("\s+");  
0 1 2 3

SB "blue - is - sky - me"  $\Rightarrow$  return as string  


time =  $O(n)$

$n = str.length()$

space =  $O(n)$

extra space  
(array)

```
public String reverseWords(String str) {  
    String[] words = str.trim().split("\\s+");  
    // trim: leading & trailing spaces: discard  
    // split on multiple spaces: single split  
  
    StringBuilder res = new StringBuilder();  
    for(int idx = words.length - 1; idx >= 0; idx--){  
        if(idx < words.length - 1){  
            res.append(" ");  
        }  
        res.append(words[idx]);  
    }  
  
    return res.toString();  
}
```

no space before first word

```
public String reverseWords(String str) {  
    String[] words = str.split("\\s+");  
    StringBuilder res = new StringBuilder();  
    for(int idx = words.length - 1; idx >= 0; idx--){  
        res.append(" " + words[idx]);  
    }  
    return res.toString().trim();  
}
```

StringBuilder  
to  
String  
convert

leading &  
trailing spaces  
are removed

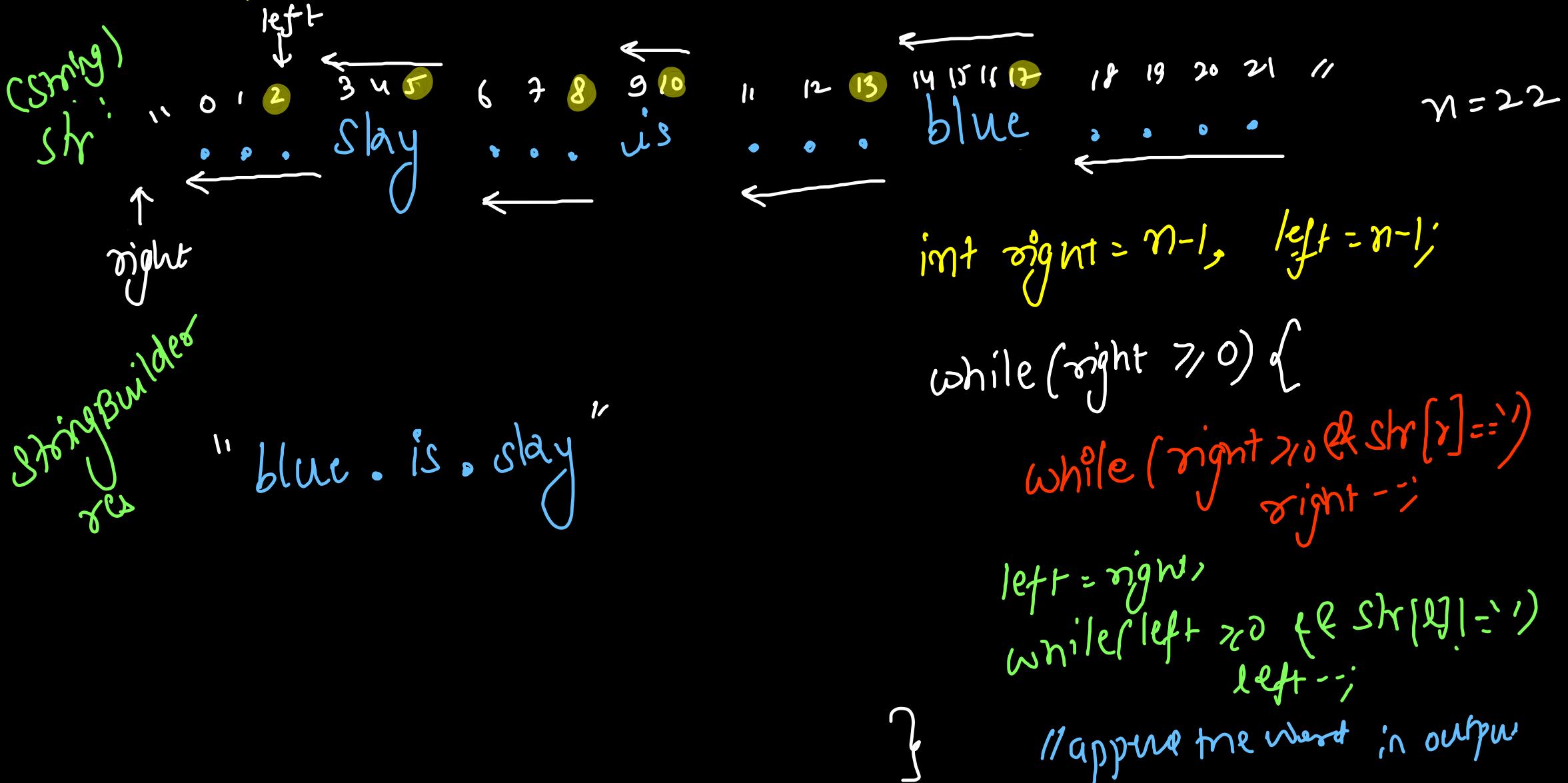
Input split  
on multiple  
spaces.

LC 151

and way of

1st approach,

## and approach : Two pointer Technique



```

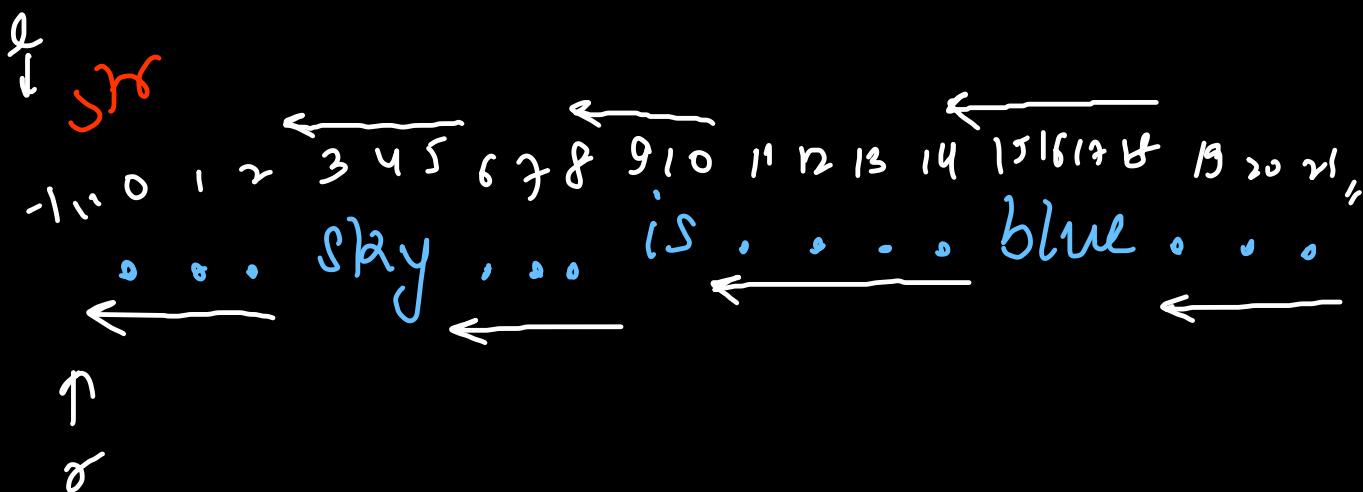
public String reverseWords(String str) {
    int right = str.length() - 1;
    StringBuilder res = new StringBuilder();
    // iterate on all words
    while( right >= 0 ){
        // right will stop at ending index of a word
        while(right >= 0 && str.charAt(right) == ' '){
            right--;
        }
        int left = right;
        // left will stop at starting index of a word - 1
        while(left >= 0 && str.charAt(left) != ' '){
            left--;
        }
        left is word right is empty word
        if(res.length() > 0 && left < right) {
            res.append(" ");
        }
        for(int idx = left + 1; idx <= right; idx++){
            res.append(str.charAt(idx));
        }
        right = left;
    }
    return res.toString();
}

```

*input*

*output*

*Independent loops*



*blue . is . sky*

*2 pointer*

*time  $\Rightarrow \Theta(n)$*

*Space  $\Rightarrow \Theta(1)$*

*no extra space*

## L(C SS7) Reverse words - II

- no leading spaces
- no trailing spaces
- no multiple spaces

Input "fine . sky , is , blue"  
↓  
Output "eht . yks . si . eulb"

Input  
"the . sky , is . blue"

① approach 1) Split

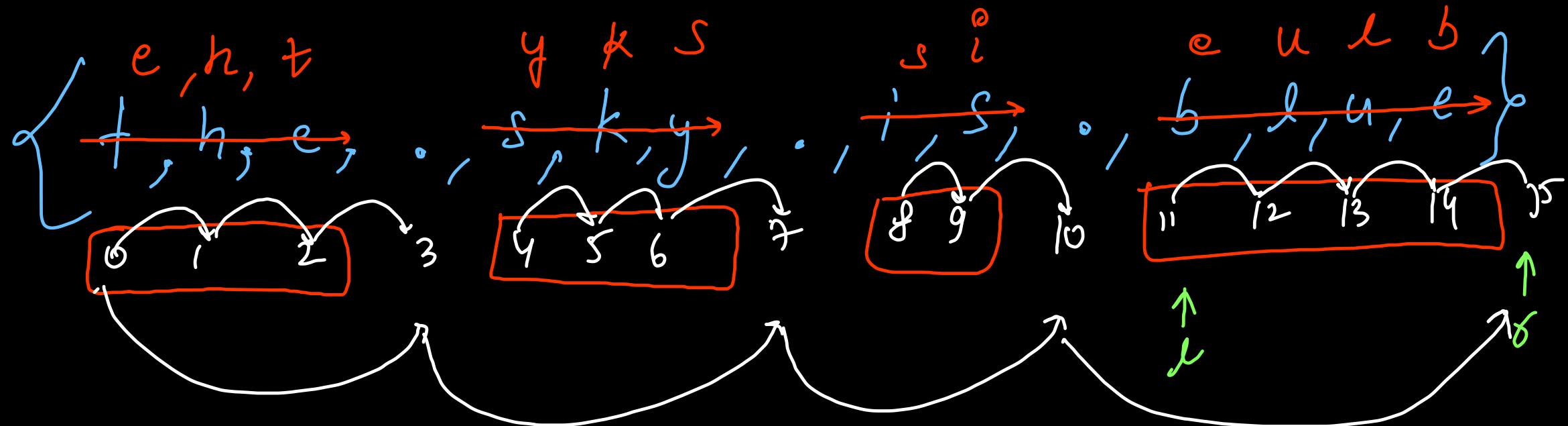
↓  
Output  
"eht . yks . si . eulb"

{ "the", "sky", "is", "blue" }  
[ ] [ ] [ ] [ ]  
"eht" . "yks" . "si" . "eulb"

Time  $\Rightarrow O(n)$  ↪ append <sup>split</sup>

Space  $\Rightarrow O(n)$  (array)

## ② approach 2: character array + two pointer



```
public void reverse(char[] res, int left, int right){  
    while(left <= right){  
        char temp = res[left];  
        res[left] = res[right];  
        res[right] = temp;  
        left++; right--;  
    }  
}  
  
public String reverseWords(String str) {  
    char[] res = str.toCharArray();  
    ← output  
    int left = 0; → all words  
    while(left < str.length()){  
        int right = left; → all characters of word  
        while(right < str.length() && res[right] != ' '){  
            right++;  
        } → reverse word  
        reverse(res, left, right - 1);  
        left = right + 1;  
    }  
  
    return String.valueOf(res);  
}
```

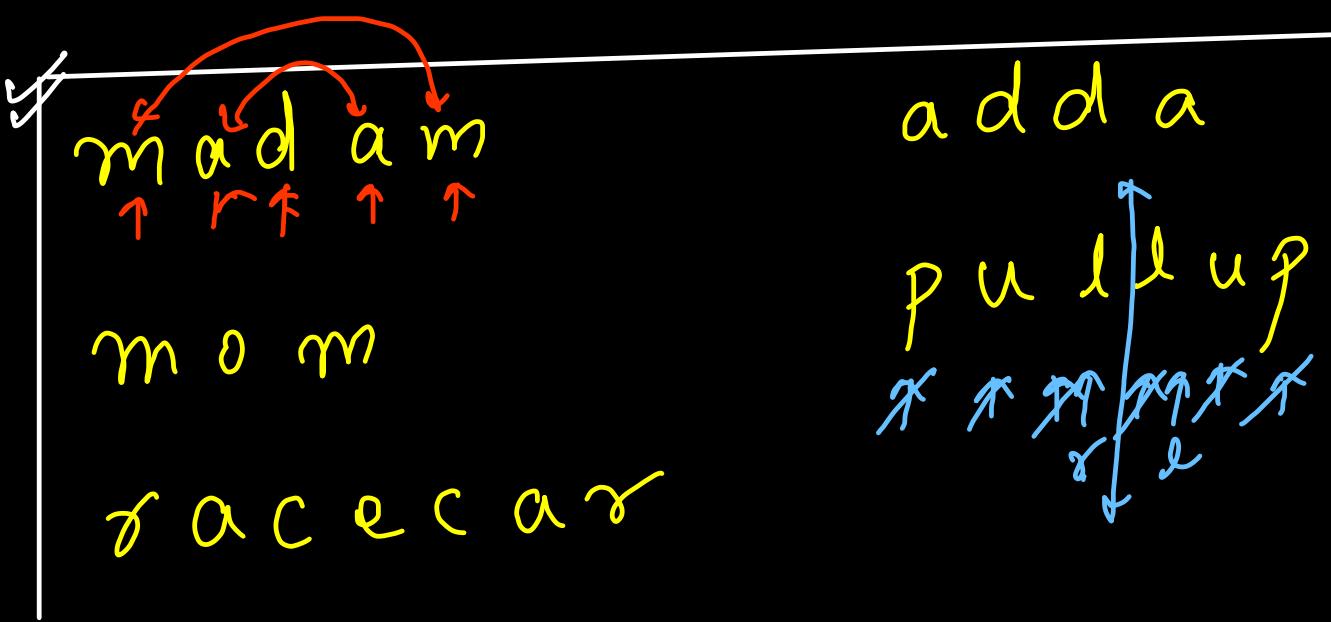
Time  $\Rightarrow O(n)$   
linear

Space  $\Rightarrow O(1)$   
no extra space

# Valid Palindrome

[a c l a s]

palindrome : left to right == right to left



abcabc X  
abcdef X  
abcdba X

" abc<sup>d</sup> "      X X X      0 1 2 3      # \$ @ ! # ... 3 2 1 0 d<sup>c</sup> b a "     

X X Y X      Y Y X      X X Y X      X X X X      X X X X      Y Y Y Y      Y Y Y Y      X X X X

X X X X      Y Y Y Y      X X X X      X X X X      X X X X      Y Y Y Y      Y Y Y Y      X X X X

- convert uppercase to lowercase  $\Rightarrow \text{toLowerCase}()$
- remove all non-alphanumeric  $\Rightarrow$  two pointer (discard)  
`isAlphanumeric()`  
`(`0' - 'g' || 'a' - 'z')`

```

public boolean isAlphaNumeric(char ch){
    if(ch >= 'a' && ch <= 'z') return true; // lowercase
    if(ch >= 'A' && ch <= 'Z') return true; // uppercase
    if(ch >= '0' && ch <= '9') return true; // digit
    return false;
}

public boolean isPalindrome(String str) {
    str = str.toLowerCase();
    int left = 0, right = str.length() - 1;

    while(left <= right){
        if(isAlphaNumeric(str.charAt(left)) == false){
            left++; continue;
        }

        if(isAlphaNumeric(str.charAt(right)) == false){
            right--; continue;
        }

        if(str.charAt(left) != str.charAt(right)){
            return false;
        }
        left++; right--;
    }

    return true;
}

```

*discard non-alpha numeric*

Time =  $O(n)$   
 (Linear)

Space =  $O(1)$   
 (no extra space)

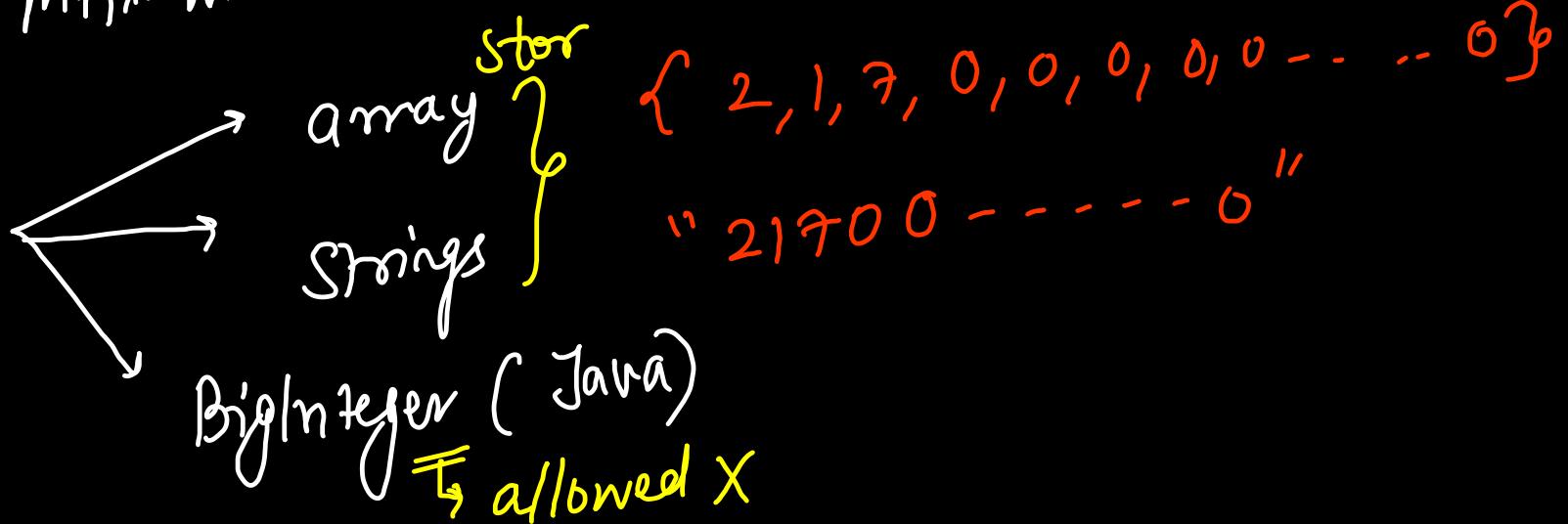
# Big Integers

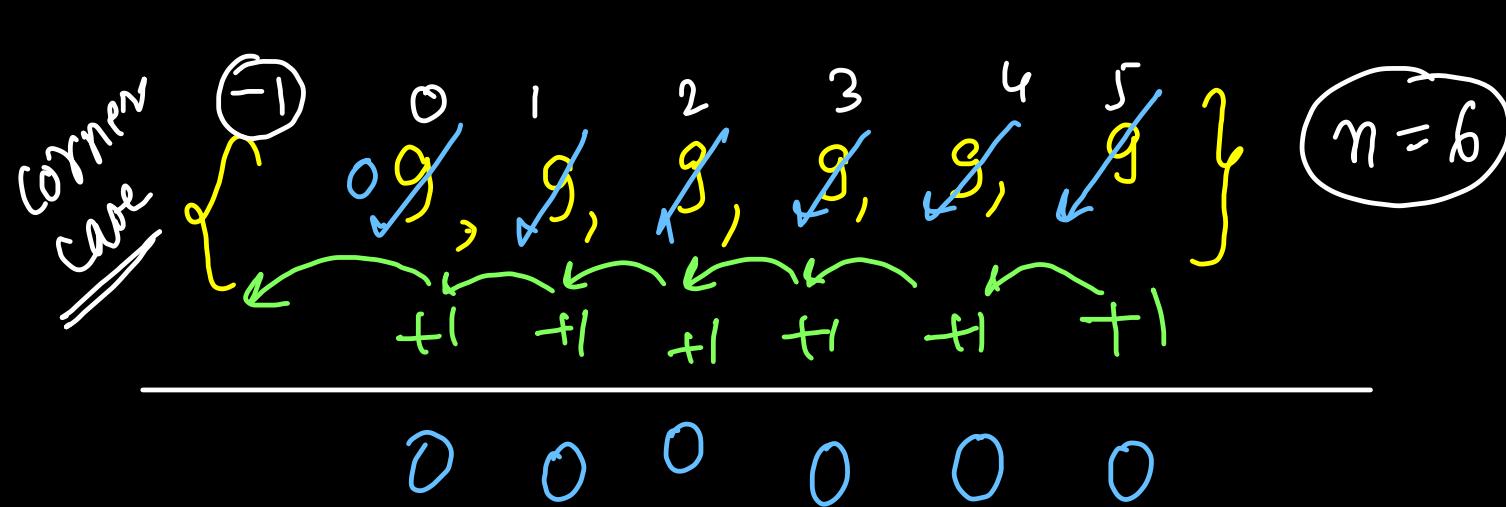
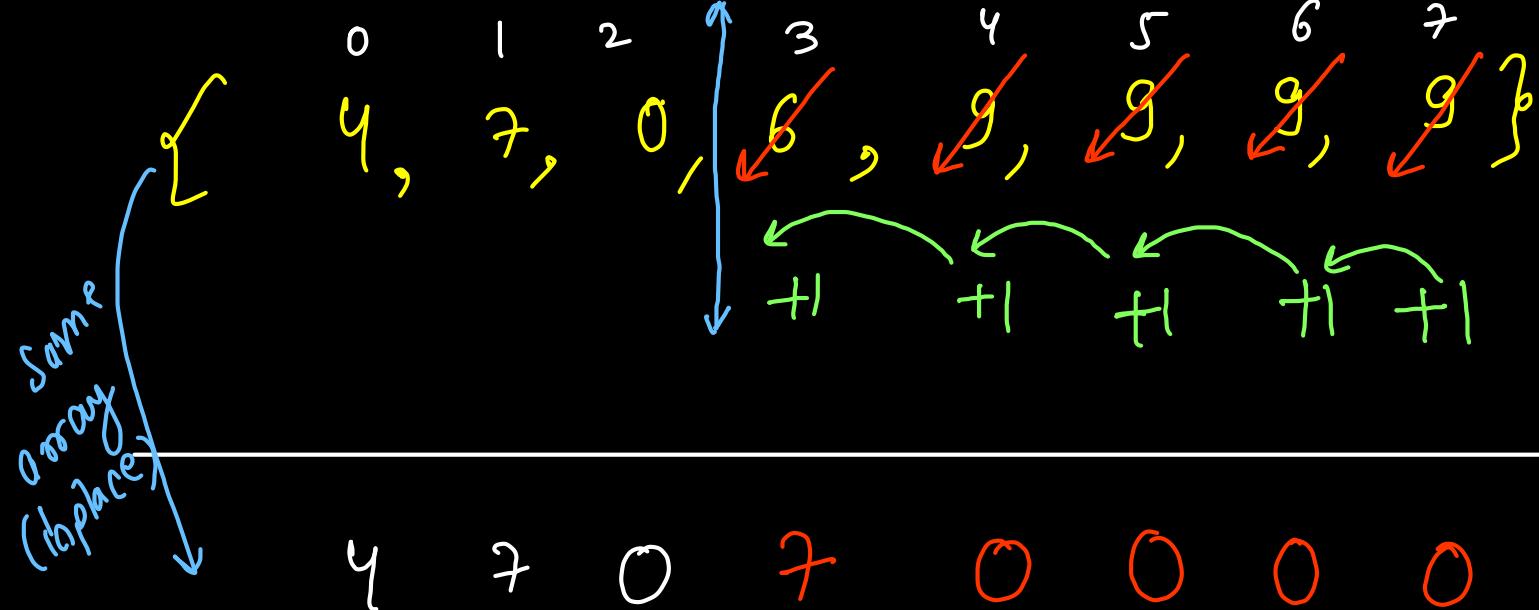
`int` →  $2^{31} - 1$  →  $2 * 10^9$   
`Integer.MAX-VALUE`

$$\text{long} \rightarrow 2^{63} - 1 \rightarrow 10^{15}$$

long. MAX-VALUE

# Bigger Numbers





Resultant size  $\Rightarrow (n+1) = 7$   
 (extra array)

{ 1, 0, 9, 0, 0, 0 }

{ 0, 1, 2, 3, 4, 5 }

```

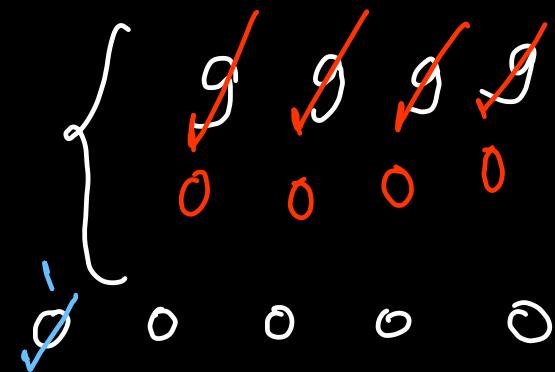
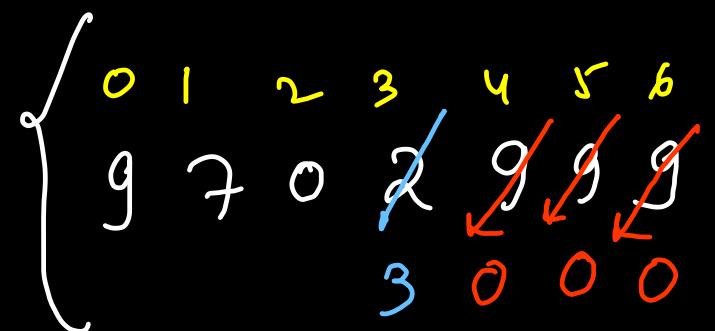
public int[] plusOne(int[] digits) {
    for(int idx = digits.length - 1; idx >= 0; idx--){
        if(digits[idx] == 9){ //carry
            digits[idx] = 0;
        } else {
            digits[idx]++;
            return digits; //usual case  $\Rightarrow O(1)$  no extra
                           //space
                           //inplace
        }
    }

    // corner case: Input has all digits equal to 9
    int[] res = new int[digits.length + 1];
    res[0] = 1; // corner case:  $O(n)$  extra
                // space
    return res;
}

```

Time =  $O(n)$  linear

Space =  $O(n)$  worst  
case  
(corner)



# LC 415 Add 2 Strings

carry = 0 0 0 1 0 1 0 0 0

~~\* Char → int~~

$$q1 = n_1 - 1, q2 = n_2 - 1$$

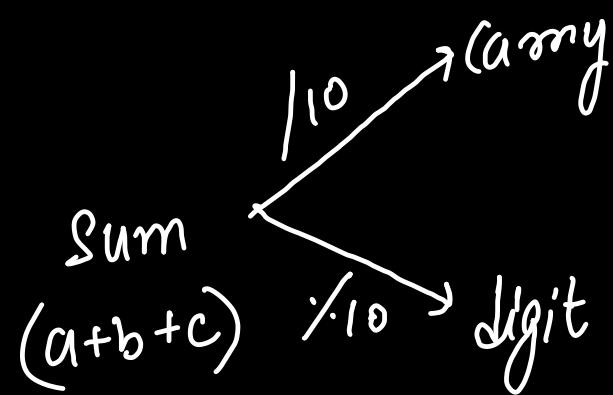
String(a) " 0 1 2 3 0 9 6 5 "

0 4 2 3 0 9 6 5

String(b) " 7 1 6 9 5 4 0 2 "

0 1 2 3 4 5 6 7

String Builder " 7 6 9 2 6 3 6 7 "



$$\begin{array}{r}
 \text{carry} 0 \\
 \begin{array}{r}
 \cancel{-1} \cancel{9} \cancel{3} \cancel{4} \cancel{2} \\
 + \cancel{-1} \cancel{7} \cancel{8} \cancel{6} \cancel{9} \\
 \hline
 \end{array} \\
 \hline
 1 \ 7 \ 2 \ 1 \ 1
 \end{array}$$

$$\begin{array}{r}
 \text{carry} 0 \ 0 \ | \ | \ | \ | \ | \ | \ 0 \\
 \begin{array}{c}
 \text{---} \\
 -1 \ " \ 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ "
 \end{array} \\
 \begin{array}{c}
 \text{---} \\
 \cancel{-1} \ 9 \ 2 \ 4 \ 6 \ 0 \ 7 \ 3 \ "
 \end{array} \\
 \begin{array}{c}
 \text{---} \\
 0 \ 0 \ " \ 7 \ 8 \ 9 \ 9 \ 8 \ "
 \end{array} \\
 \begin{array}{c}
 \text{---} \\
 -2 \ -1 \ 0 \ 1 \ 2 \ 3 \ 4
 \end{array} \\
 \hline
 \end{array} \\
 \hline
 9 \ 3 \ 2 \ 5 \ 0 \ 7 \ 1 \leftarrow$$

```

public String addStrings(String num1, String num2) {
    StringBuilder res = new StringBuilder();
    int p1 = num1.length() - 1, p2 = num2.length() - 1;
    int carry = 0;

    while(p1 >= 0 || p2 >= 0 || carry > 0){
        int d1 = (p1 >= 0) ? num1.charAt(p1) - '0' : 0;
        int d2 = (p2 >= 0) ? num2.charAt(p2) - '0' : 0;
        int sum = d1 + d2 + carry;

        carry = sum / 10;
        res.append(sum % 10);
        p1--; p2--;
    }

    return res.reverse().toString();
}

```

Linear time  
 $O(\max(n_1, n_2))$

Space  $\Rightarrow O(1)$   
Constant extra  
(considering SB as output)

# Subtract Two Numbers

borrow -1 0 -1 0 0  
9 2 1 7 5

- 5 8 0 9 4

---

3 4 0 8 1

---

$$\text{int diff} = (c + d1 - d2)$$

if (diff < 0) {

    diff += 10;

    c = -1;

}

else {

    c = 0;

}

~~LC43~~ Multiply Two Big Integers

$$\begin{array}{r} 2310 \\ 0473 \\ \times 5 \\ \hline 2365 + 42570 + \overbrace{283800}^{\text{circled}} \\ \hline \end{array} \quad \begin{array}{r} 462 \\ 0473 \\ \times 90 \\ \hline 42570 \\ + \end{array} \quad \begin{array}{r} 41 \\ 473 \\ \times 600 \\ \hline 283800 \end{array} = (473 \times 5) * 1 + (473 \times 9) * 10 + (473 \times 6) * 100$$

```

public String addStrings(String num1, String num2) {
    StringBuilder res = new StringBuilder();
    int p1 = num1.length() - 1, p2 = num2.length() - 1;
    int carry = 0;

    while(p1 >= 0 || p2 >= 0 || carry > 0){
        int d1 = (p1 >= 0) ? num1.charAt(p1) - '0' : 0;
        int d2 = (p2 >= 0) ? num2.charAt(p2) - '0' : 0;
        int sum = d1 + d2 + carry;

        carry = sum / 10;
        res.append(sum % 10);
        p1--; p2--;
    }

    return res.reverse().toString();
}

```

$$\begin{array}{r}
 473 \\
 \times 5 \\
 \hline
 2365
 \end{array}$$
  

$$\begin{array}{r}
 473 \\
 \times 9 \\
 \hline
 4257
 \end{array}$$
  

$$\begin{array}{r}
 473 \\
 \times 6 \\
 \hline
 2838
 \end{array}$$

```

public String multiply(String num1, int d2){
    StringBuilder res = new StringBuilder();

    int p1 = num1.length() - 1;
    int carry = 0;

    while(p1 >= 0 || carry > 0){
        int d1 = (p1 >= 0) ? num1.charAt(p1) - '0' : 0;
        int prod = d1 * d2 + carry;

        res.append(prod % 10);
        carry = prod / 10;
        p1--;
    }

    return res.reverse().toString();
}

```

Time =  $O(n^2)$   
 Space =  $O(n^2)$   
 extra space

```

public String multiply(String num1, String num2) {
    String res = "0";
    int count = 0;

    for(int idx = num2.length() - 1; idx >= 0; idx--){
        int d2 = num2.charAt(idx) - '0';
        String temp = multiply(num1, d2);

        for(int c = 0; c < count; c++){
            temp += '0';
        }
    }

    res = addStrings(res, temp);
    count++;

    if(res.charAt(0) == '0') return "0";
    return res;
}

Corner case
"0000"
↓
"0"

```

res "0" + "2365"

$$\begin{aligned}
 &= "2365" \\
 &\quad + "42570" \\
 &\quad + "1283800"
 \end{aligned}$$

# Frequency, Substring & Subsequence

rearrangement → permutations

Leetcode 242) Valid anagram

"anagram"  $\Leftrightarrow$  "nagaram"  
true

"car"  $\Leftrightarrow$  "rat"  
false

①  $sl.length()$   
 $= s2.length()$

② Characters  
(frequency)  
same

brute force  $\Rightarrow$  Generate  
all permutations

Time  $\Rightarrow \mathcal{O}(n!)$   $\Rightarrow$  exponential  
(TLE)

"nagaram"  $\Rightarrow$  "nagaram"

0-127 ASCII

26 distinct letters ('a'-'z')  $\Rightarrow$  26 size

$\text{freq}['a'] \Rightarrow \text{freq}[g\gamma]$

$\text{f}(n)$	'a'	'b'	'c'	'd'	'e'	'f'	'g'	'h'	'i'	'j'	'k'	'l'	'm'	'n'	'o'	'p'	'q'	'r'	's'	't'	'u'	'v'	'w'	'x'	'y'	'z'
	3			-																						
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	

II compare (26 iterations)

$\ell_2 \rightarrow O(n^2)$

A horizontal timeline from 0 to 25 with vertical grid lines every 1 unit. Red numbers 3, 6, 12, 13, and 17 are marked on the grid.

```

public int[] getFrequency(String str){
    int[] freq = new int[26]; → O(n)
    for(char ch : str.toCharArray()){ n = str.length()
        freq[ch - 'a']++;
    }
    return freq;
}

public boolean isAnagram(String s1, String s2) {
    if(s1.length() != s2.length()){
        return false;
    }

    int[] f1 = getFrequency(s1); O(n1)
    int[] f2 = getFrequency(s2); O(n2)

    for(int idx = 0; idx < 26; idx++){
        if(f1[idx] != f2[idx])
            return false;
    }
}

return true;
}

```

$$\text{Time} = O(n_1 + n_2)$$

$$\text{Space} = O(26 + 26)$$

extra space

$\approx O(1)$  inplace

$$\left. \begin{array}{l} \\ \\ \end{array} \right\} O(26) = O(1)$$

LC 1832)

Pangram

Input: sentence = "thequickbrownfoxjumpsoverthelazydog"

→ true

Every character is lowercase English alphabets

- ① frequency array build:  $O(n)$
- ②  $\text{freq}[\text{'a' - 'z'}] \geq 1 \quad \forall \text{'a' - 'z'}$

## LC 205) Isomorphic Strings

"egg"     $\xrightarrow{\text{true}}$     "add"  
 $e \rightarrow a$   
 $g \rightarrow d$   
1:1 ✓

"foo"     $\xrightarrow{\text{false}}$     "bar"  
 $f \rightarrow b$   
 $o \rightarrow a$   
 $o \rightarrow r$   
1: many ✗

"paper"     $\xrightarrow{\text{true}}$     "title"  
 $p \rightarrow t$   
 $a \rightarrow i$   
 $e \rightarrow l$   
 $r \rightarrow e$   
1:1 ✓

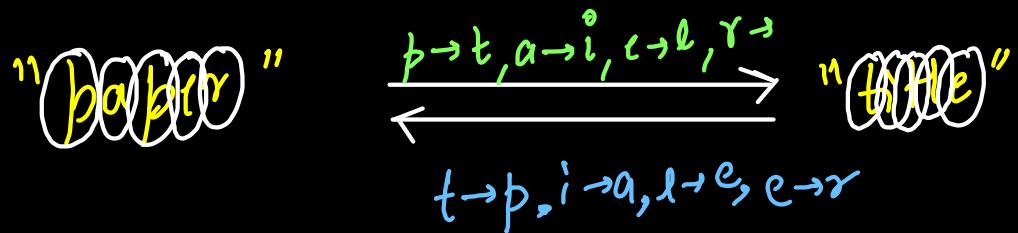
"bar"     $\xrightarrow{\text{false}}$     "foo"  
 $b \rightarrow f$   
 $a \rightarrow o$   
 $r \rightarrow o$   
many:1 ✗

# no two characters should have same mapping

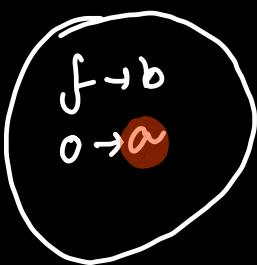
# Character may map to itself.

array ( map)  $\Rightarrow$  character array

'a'	'b'	'c'	'd'	'e'	'f'	'g'	'h'	'i'	'j'	'k'	'l'	'm'	'n'	'o'	'p'	'q'	'r'	's'	't'	'u'	'v'	'w'	'x'	'y'	'z'
0	10	10	10	10	10	10	10	10						10	10										
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	



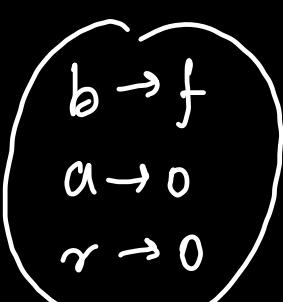
"foo"  
 $\downarrow\downarrow\downarrow$



"bar"  
 $\downarrow\downarrow$

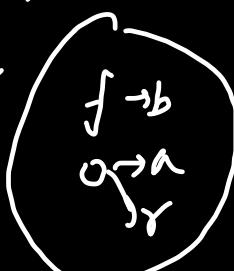
'o'  $\neq$  'r'  $\Rightarrow$  false  
1: many

"bar"



many:  
false

"foo"



```

public boolean mapOneToOne(String s1, String s2){
    char[] map = new char[256]; //ascii code

    for(int idx = 0; idx < s1.length(); idx++){
        char ch1 = s1.charAt(idx);
        char ch2 = s2.charAt(idx);
        if(map[ch1] != '\u0000' && map[ch1] != ch2)
            return false;
        map[ch1] = ch2;
    }

    return true;
}

public boolean isIsomorphic(String s1, String s2) {
    return mapOneToOne(s1, s2) && mapOneToOne(s2, s1);
}

```

ch1  
 ↓ ↓ ↓ ↓  
 paper

b → p  
 a → a  
 e → e  
 l → t

title

ch2  
 ↓ ↓ ↓ ↓  
 title

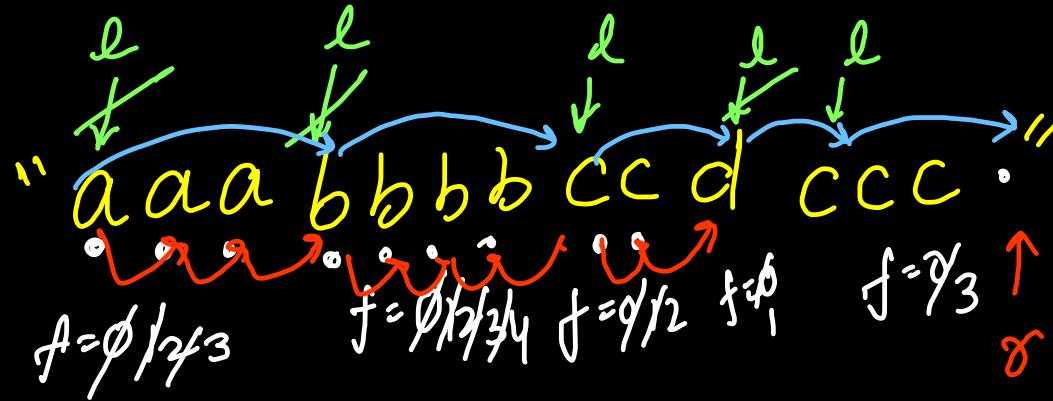
paper

t → p  
 i → a  
 l → e  
 e → r

# String Compression (Run Length Encoding)

input "aaa bbbb cc d ccc aaaaa eff"  
↓ info loss X compress ✓

output "a3 b4 c2 d c3 a5 e f2"



"a3b4c2d c3"



Time  $\Rightarrow O(n)$  time  
linear

Space  $\Rightarrow O(1)$  no extra  
(SB as output)

consecutive  
freq  
cumulative

```
String encode(String str)
{
    StringBuilder res = new StringBuilder();

    int left = 0, right = 0;
    while(left < str.length()){left  $\Rightarrow O(n)$ 
        char ch = str.charAt(left);
        int freq = 0;
        loops  $\Rightarrow O(n)$ 
        while(right < str.length() && str.charAt(right) == ch){ $\Rightarrow O(n)$ 
            freq++; right++;
        }

        res.append(ch);
        res.append(freq);
        left = right;
    }

    return res.toString();
}
```

# brute force → optimized  
nested loops  $\xrightarrow{\text{reason}}$  a pointer

# Time & Space complexity  $\Rightarrow$  dry run  $\xrightarrow{\text{proper reasoning} \checkmark}$

# team work / communication  $\Rightarrow$  interest / focus  
(continuous)

# good companies  $\Rightarrow$  introduction  $\begin{cases} \xrightarrow{\text{short}} \\ \xrightarrow{\text{crisp}} \end{cases} \} \rightarrow 3-5 mins$

Subarray / Substring

part of an array  
# continuous  
# same order

empty subarray  
⋮  
whole array

" chit "  
0 1 2 3  
 $n = 4$

total substrings

$$= \frac{n \times (n+1)}{2}$$

c h i t  
ch hi it  
chi hit  
chit

$$1 + 2 + 3 + \dots + n$$

$$= \frac{n \times (n+1)}{2}$$

```

static void printSubstrings(String str) {
    for(int start = 0; start < str.length(); start++){
        for(int end = start; end < str.length(); end++){
            for(int idx = start; idx <= end; idx++){
                System.out.print(str.charAt(idx));
            }
            System.out.println();
        }
    }
}

```

$O(N^3)$   
Cubic

Str = Code

0 1 2 3

start = 0  
end = 0     $[0,0] \Rightarrow "c"$   
            1     $[0,1] \Rightarrow "co"$   
            2     $[0,2] \Rightarrow "cod"$   
            3     $[0,3] \Rightarrow "code"$

start = 1  
end = 1     $[1,1] \Rightarrow "o"$   
            2     $[1,2] \Rightarrow "od"$   
            3     $[1,3] \Rightarrow "ode"$

start = 2    start = 3  
end = 2     $[2,3] \Rightarrow "d"$     end = 3  
            3     $[3,3] \Rightarrow "de"$      $[3,3] \Rightarrow "e"$

`str.substring(left) :→ [left, n)`

time:  $O(\text{substring length})$

`str.substring(left, right)`

$\rightarrow [left, right)$

included

excluded

$\text{substring length} \Rightarrow right - left$

"code"  
"0123"  
  
 $\text{substring}(0) \Rightarrow "code"$   
 $\text{substring}(1) \Rightarrow "ode"$   
 $\text{substring}(2) \Rightarrow "de"$   
 $\text{substring}(3) \Rightarrow "e"$

$\text{substring}(0,4) \Rightarrow "code" \ 4-0 = 4$

$\text{substring}(1,3) \Rightarrow "od" \ 3-1 = 2$

$\text{substring}(2,2) \Rightarrow " " \ 2-2 = 0$

$\text{substring}(1,4) \Rightarrow "ode" \ 4-1 = 3$

```
static void printSubstrings(String str) {  
    for(int start = 0; start < str.length(); start++){ → N  
        for(int end = start; end < str.length(); end++){ → N  
            for(int idx = start; idx <= end; idx++){ → N  
                System.out.print(str.charAt(idx));  
            }  
            System.out.println();  
        }  
    }  
}
```

$O(N^3)$   
cubic  
(polynomial)

```
static void printSubstrings(String str) {  
    for(int start = 0; start < str.length(); start++){ → N  
        for(int end = start; end < str.length(); end++){ → N  
            System.out.println(str.substring(start, end + 1)); → N  
        }  
    }  
}
```

$O(n^3)$  cubic

Subsequence  
 part of a string  
 need not be contiguous  
 order should be maintained

$$N_{\max} = 15-18$$

$n=4$   
 "abcd"

$2 \times 2 \times 2 \times \dots N \text{ times}$   
 Subsequences  
 (exponential)

$$2^4 = 16 \text{ subsequences}$$

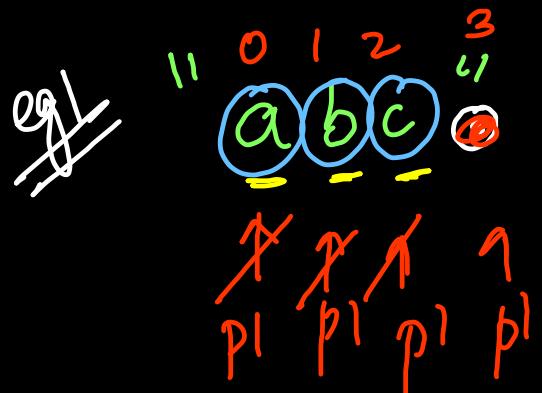
↳ " "	↳ "d"
↳ "a"	↳ "ad"
↳ "b"	↳ "bd"
↳ "ab"	↳ "abd"
↳ "c"	↳ "cd"
↳ "ac"	↳ "acd"
↳ "bc"	↳ "bcd"
↳ "abc"	↳ "abcd"

# Subsequence : "remaining string  
after removal of 0 or more characters  
from the original string"

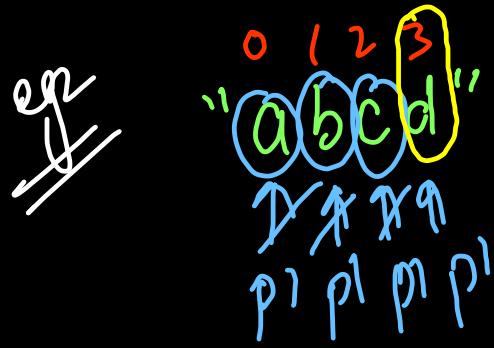
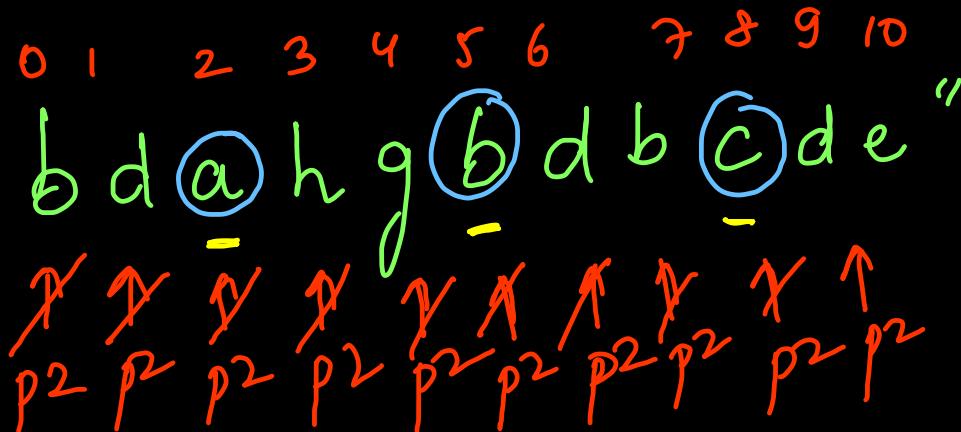
# All substrings are subsequences ✓  
or  
=

All subsequences are substrings ✗

LC 392) Is Subsequence



→ true  
 $p1 = s1.length()$



→ false  
 $p2 = s2.length()$



```

public boolean isSubsequence(String s1, String s2) {
    int p1 = 0, p2 = 0;

    while(p1 < s1.length() && p2 < s2.length()){
        if(s1.charAt(p1) == s2.charAt(p2)){
            p1++; p2++;
        } else {
            p2++;
        }
    }

    return (p1 == s1.length());
}

```

Time  $\Rightarrow O(n_1 + n_2)$

linear

Space  $\Rightarrow O(1)$

constant  
(Inplace)

Two pointer Technique

# Find longest word in Dictionary (GFG)

## Example 1:

**Input:** d = {"ale", "apple", "monkey", "plea"}

S = "abpcplea"

**Output:** "apple"

**Explanation:** After deleting "b", "c"

"a" S became "apple" which is present  
in d.

~~eg~~ S = "abcdefg"

dict = { "xyz", "ade",  
"afg" } ✓

tie breaker  
lexicographical

~~eg~~ S = "abcdefg"  
d, dict = { "axy", "nya", "cdx" }

answer :: "

" a b c d d c x x Y Z "  
0 1 2 3 4 5 6 7 8 9

First Unique Character  
answer = 'a'

First Repeating character  
answer = 'c'

" a b c d c b a d "

First Unique Character  
answer = -1

First Repeating character  
answer = 'a'

" a b c d e f g "

first unique = 'a'      first repeating = -1

# Before: Nested loops  $\rightarrow O(n^2)$  time  
 $O(1)$  space

# frequency array  $\rightarrow$  character?  
unique  $c=1$   
repeating  $>1$

"  
a b c d d c x x Y Z"  
0 1 2 3 4 5 6 7 8 9  
  
 $a \rightarrow 1$ ,  $b \rightarrow 1$ ,  $c \rightarrow 2$ ,  $d \rightarrow 2$ ,  $x \rightarrow 2$   
 $y \rightarrow 1$ ,  $Z \rightarrow 1$

If only english lowercase  $\Rightarrow 26$    
ascii characters  $\Rightarrow 256$    
Space =  $O(1)$  inplace  
(constant)

```

public int firstUniqChar(String str) {
    int[] freq = new int[26];

    // Fill the Frequency Array
    for(char ch: str.toCharArray()){
        freq[ch - 'a']++;
    }

    // Find the First Unique or Non Repeating or Distinct Character
    for(int idx = 0; idx < str.length(); idx++){
        char ch = str.charAt(idx);
        if(freq[ch - 'a'] == 1) return idx;
    }

    return -1;
}

```

*AC 387*  
2 Loops

Time =  $O(n)$  linear  
Space =  $O(1)$   
Constant / in place

*CFS*  
One loop  
done in

```

char firstRep(String str)
{
    int[] freq = new int[26];

    // Fill the Frequency Array
    for(char ch: str.toCharArray()){
        freq[ch - 'a']++;
    }

    // Find the First Unique or Non Repeating or Distinct Character
    for(char ch : str.toCharArray()){
        if(freq[ch - 'a'] > 1) return ch;
    }

    return '#';
}

```

Diracn: Right to Left

" a b c d d c x x Y Z "  
0 1 2 3 4 5 6 7 8 9  
↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑

first repeating character

in single 'ikram'

answer = ~~'d'~~ ~~'c'~~  
~~'d'~~ ~~'c'~~

x:1

d:1/2

y:1

b:1

x:1/2

a:1

c:1/2

```
char firstRep(String str)
{
    int[] freq = new int[26];
    char repeating = '#';

    for(int idx = str.length() - 1; idx >= 0; idx--){
        char ch = str.charAt(idx);
        freq[ch - 'a']++;
        if(freq[ch - 'a'] > 1) repeating = ch;
    }

    return repeating;
}
```

"cabcd dfg"

c:1/2

a:1

b:1

d:1/2

answer = ~~c/a/b/d~~  
backtrack  
X

first unique character  
cannot be solved  
in single loop