

Minimization of FA

Ex.		0	1
	$\rightarrow q_0$	q_1	q_5
	q_1	q_6	q_2
	q_2	q_0	q_2
	q_3	q_2	q_6
	q_4	q_7	q_5
	q_5	q_2	q_6
	q_6	q_6	q_4
	q_7	q_6	q_2

Def. Two states q_1 & q_2 are equivalent if both $\delta(q_1, x)$ & $\delta(q_2, x)$ are final states or both of them are non-final states for all $x \in \Sigma^*$.

\Rightarrow Two states are k -equivalent ($k \geq 0$) if both $\delta(q_1, x)$ & $\delta(q_2, x)$ are final or non-final states for all strings x of length k or less.

\Rightarrow 0 equivalent: Any two final states are equivalent and any two non-final are equivalent.

Properties:

1. The rel. defined is Equivalence rel. \rightarrow prove it.

1. The rel. defined is Equivalence rel. \rightarrow prove it!
2. These induces partition of Q , represent as Π_k .
Elements of Π_k are k -equivalent classes.
3. If q_1 & q_2 are $(k+1)$ equivalent, they are k -equivalent.
4. If q_1 & q_2 are k -equivalent $\forall k \geq 0$ then they are equivalent.
5. $\Pi_k = \Pi_{k+1}$ then no new partition is created.

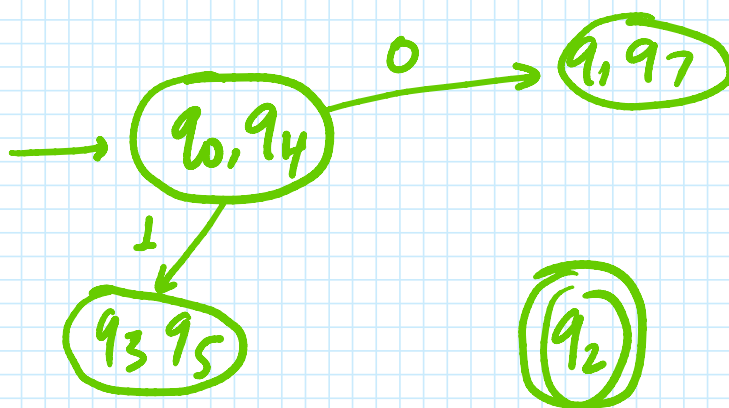
$$\Pi_0 = \{ \{q_2\}, \{q_0, q_1, q_3, q_4, q_5, q_6, q_7\} \}$$

$$\Pi_1 = \{ \{q_2\}, \{q_0, q_4, q_6\}, \{q_1, q_7\}, \{q_3, q_5\} \}$$

$$\Pi_2 = \{ \{q_2\}, \{q_0, q_4\}, \{q_6\}, \{q_1, q_7\}, \{q_3, q_5\} \}$$

$$\Pi_3 = \Pi_2$$

stop



$\downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow$
1 2 3 4 5

$$\frac{2}{\underline{\quad}} \rightarrow \begin{array}{c|c|c} & a & b \\ \hline q_0 & q_1 & q_2 \end{array}$$

2

	a	b
$\rightarrow q_0$	q_1	q_0
q_1	q_0	q_2
q_2	q_3	q_1
$\textcircled{q_3}$	q_3	q_0
q_4	q_3	q_5
q_5	q_6	q_4
q_6	q_5	q_6
q_7	q_6	q_3

$$\pi_0 = \{ \{q_3\} \}$$

$$\pi_1 = \{ \{q_3\} \{q_0, q_1, q_5, q_6\} \{q_2, q_4\} \{q_7\} \}$$

$$\pi_2 = \{ \{q_3\}, \{q_0, q_6\} \{q_1, q_5\} \{q_2, q_4\} \{q_7\} \}$$

$$\underline{\pi_3 = \pi_2}$$

Quiz - 1

1) Design DFA for language

$$L = \{ w = \{a, b, c\}^+ \mid w \text{ starts \& end with same symbols} \}$$