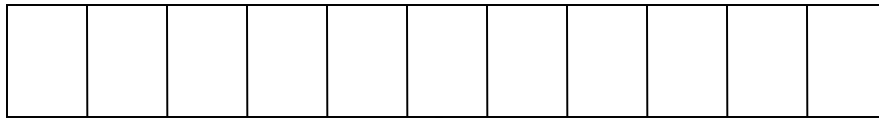


# Pushdown Automata

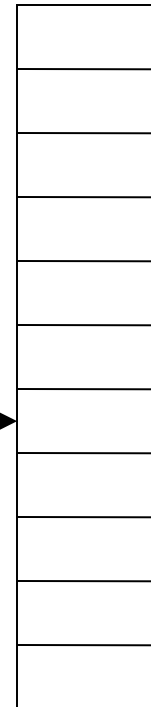
## PDA's

# Pushdown Automaton -- PDA

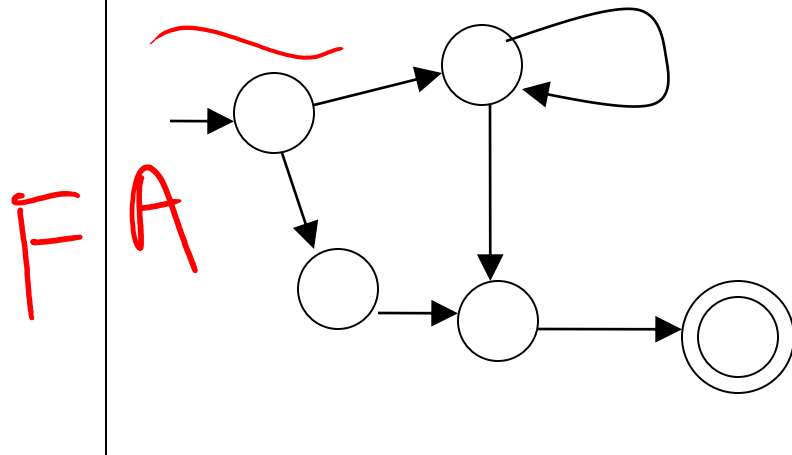
Input String



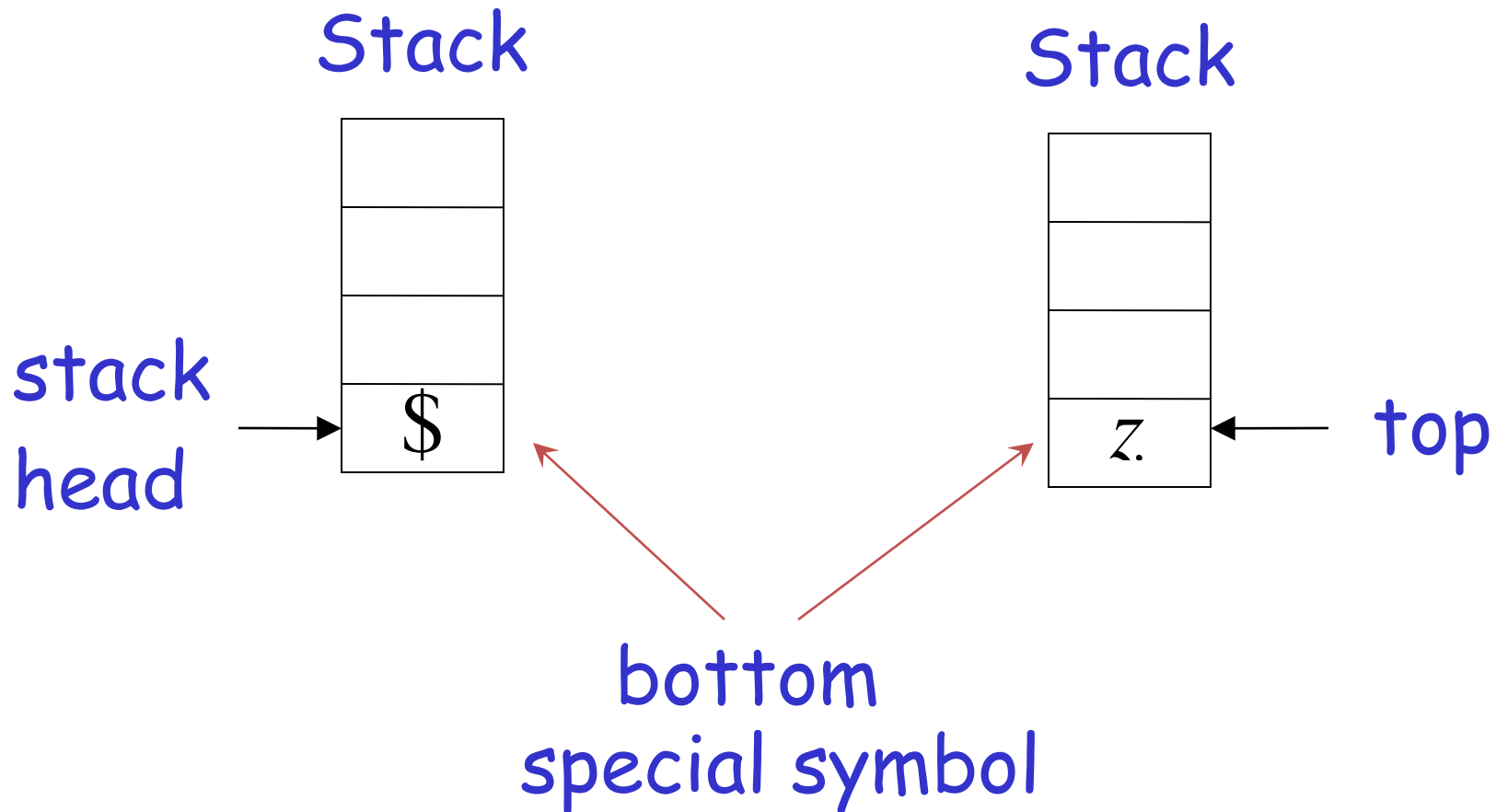
Stack



States

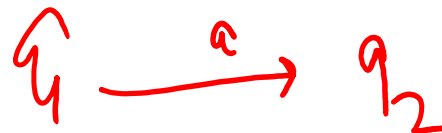
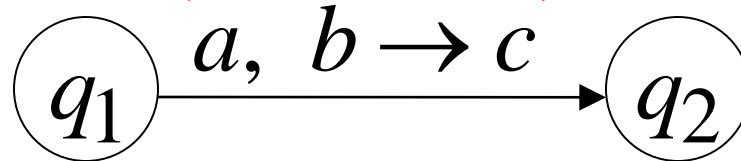


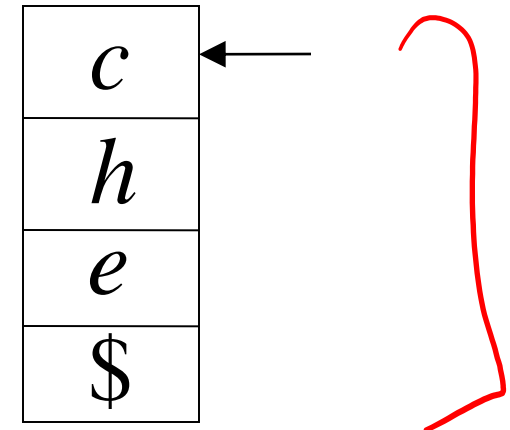
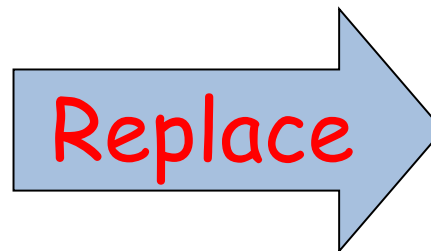
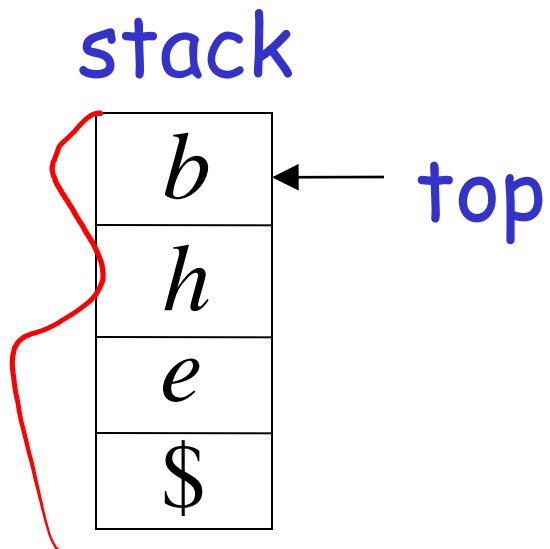
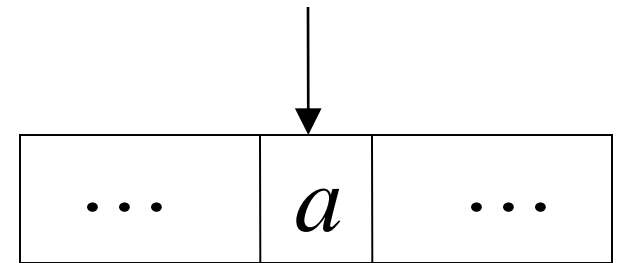
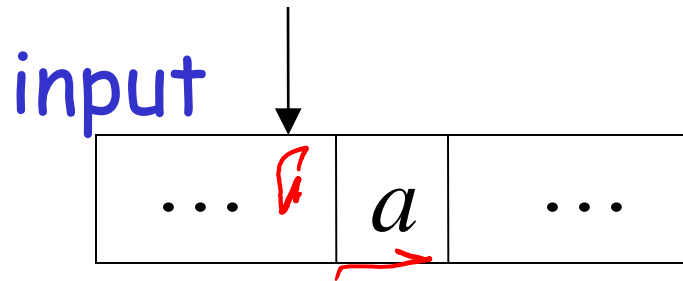
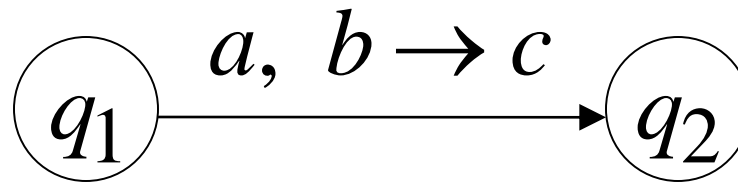
# Initial Stack Symbol

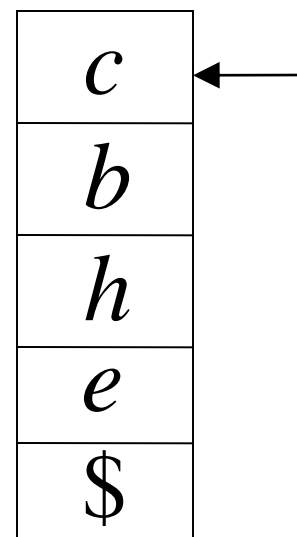
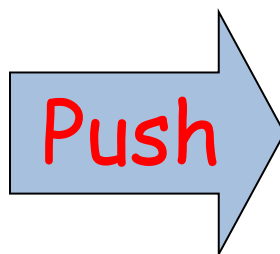
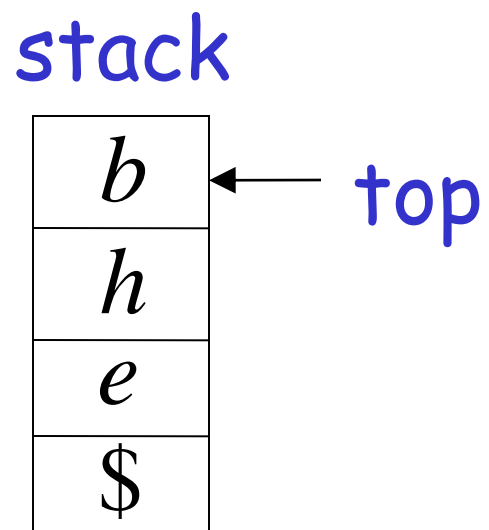
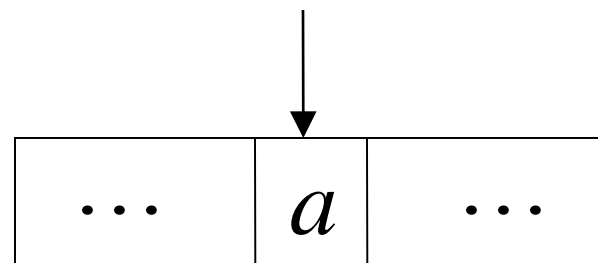
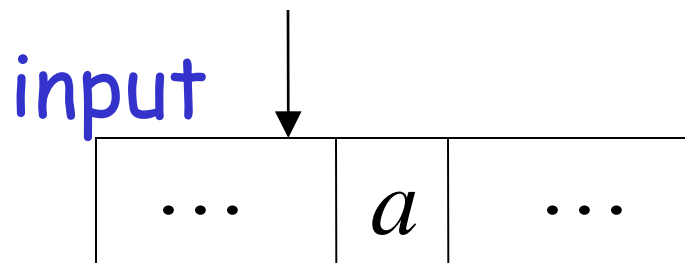
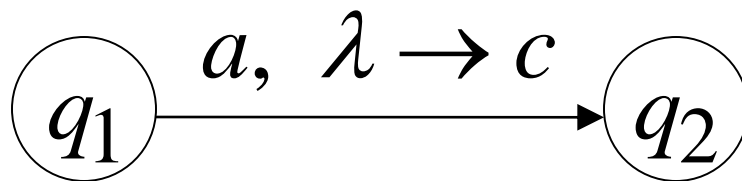


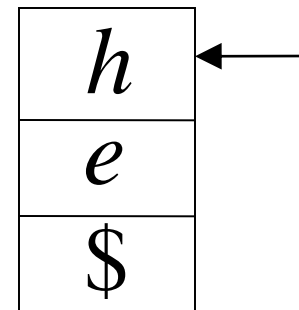
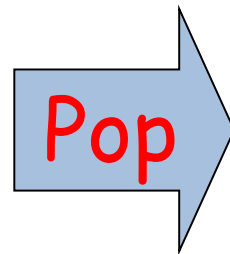
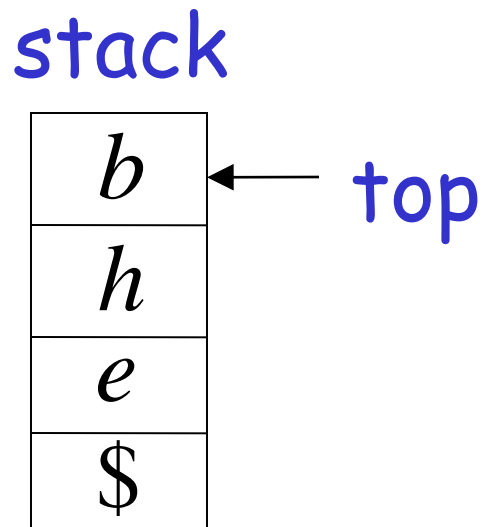
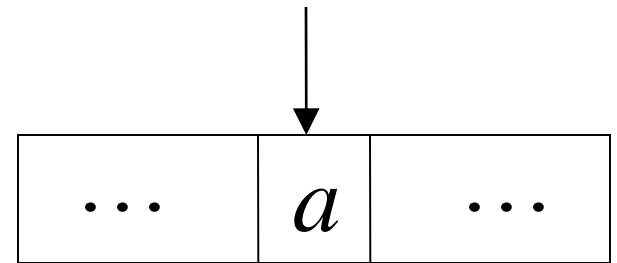
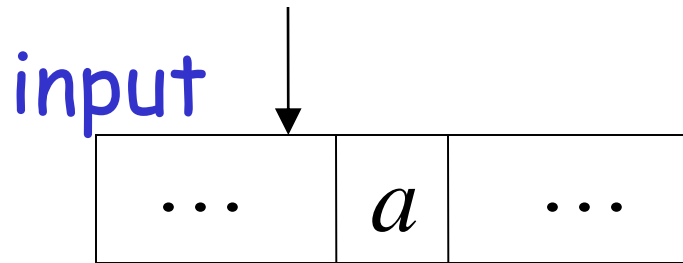
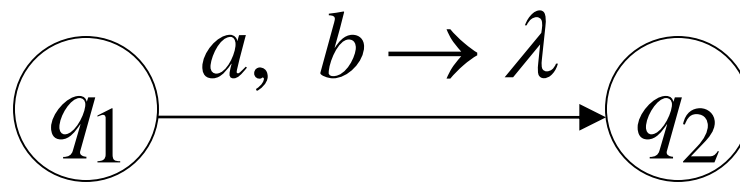
# The States

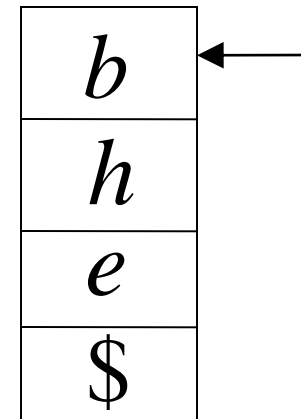
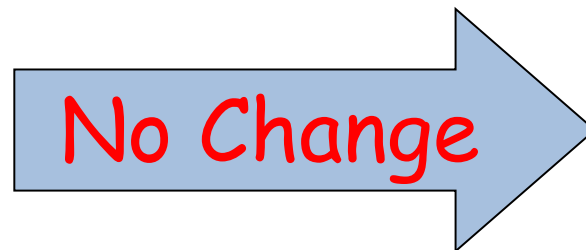
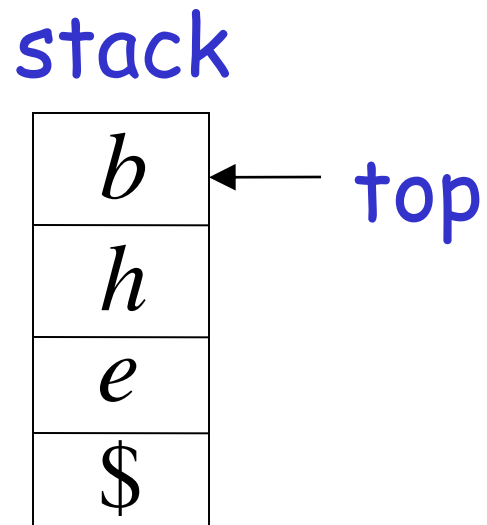
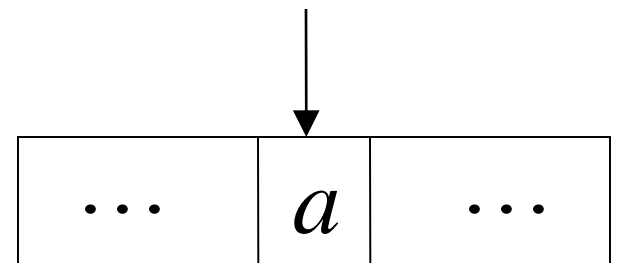
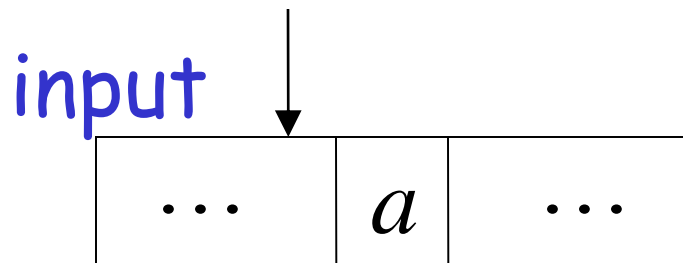
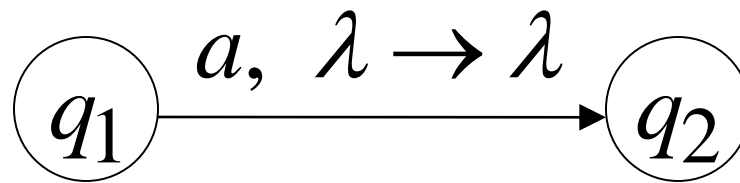
Input symbol      Pop symbol      Push symbol





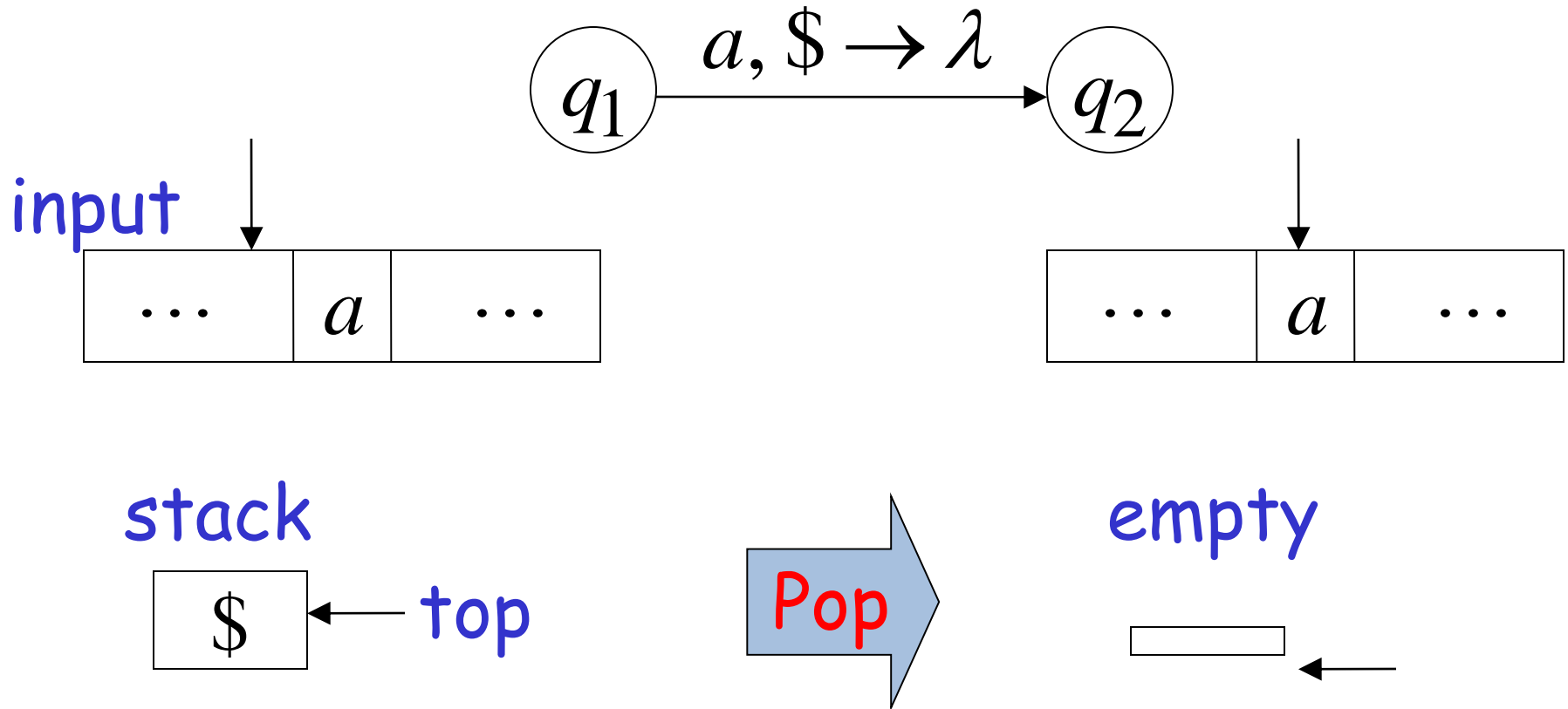




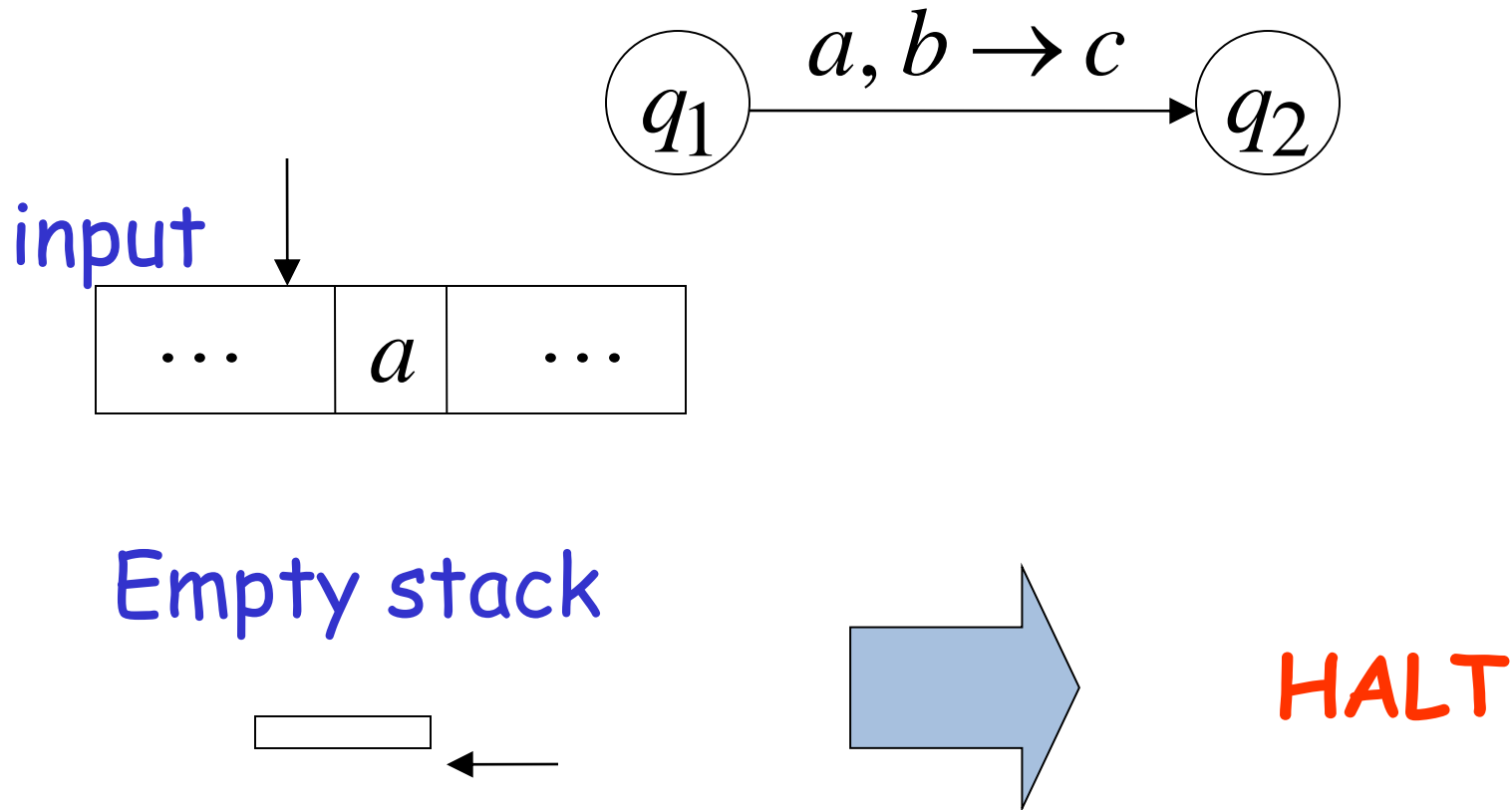




# A Possible Transition

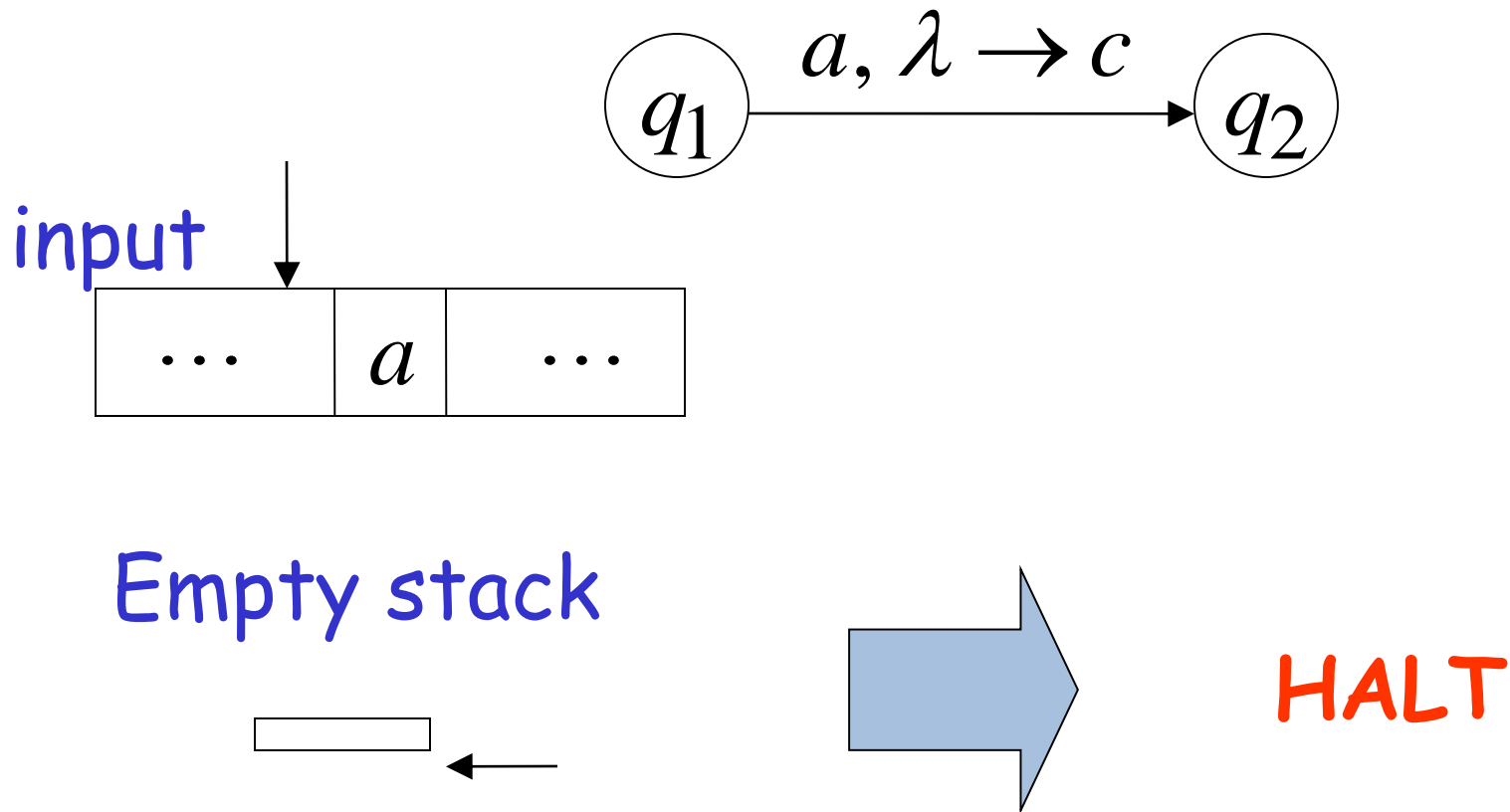


# A Bad Transition



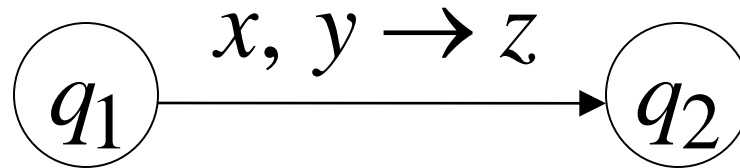
The automaton **Halts** in state  $q_1$   
and **Rejects** the input string

# A Bad Transition

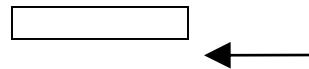


The automaton **Halts** in state  $q_1$   
and **Rejects** the input string

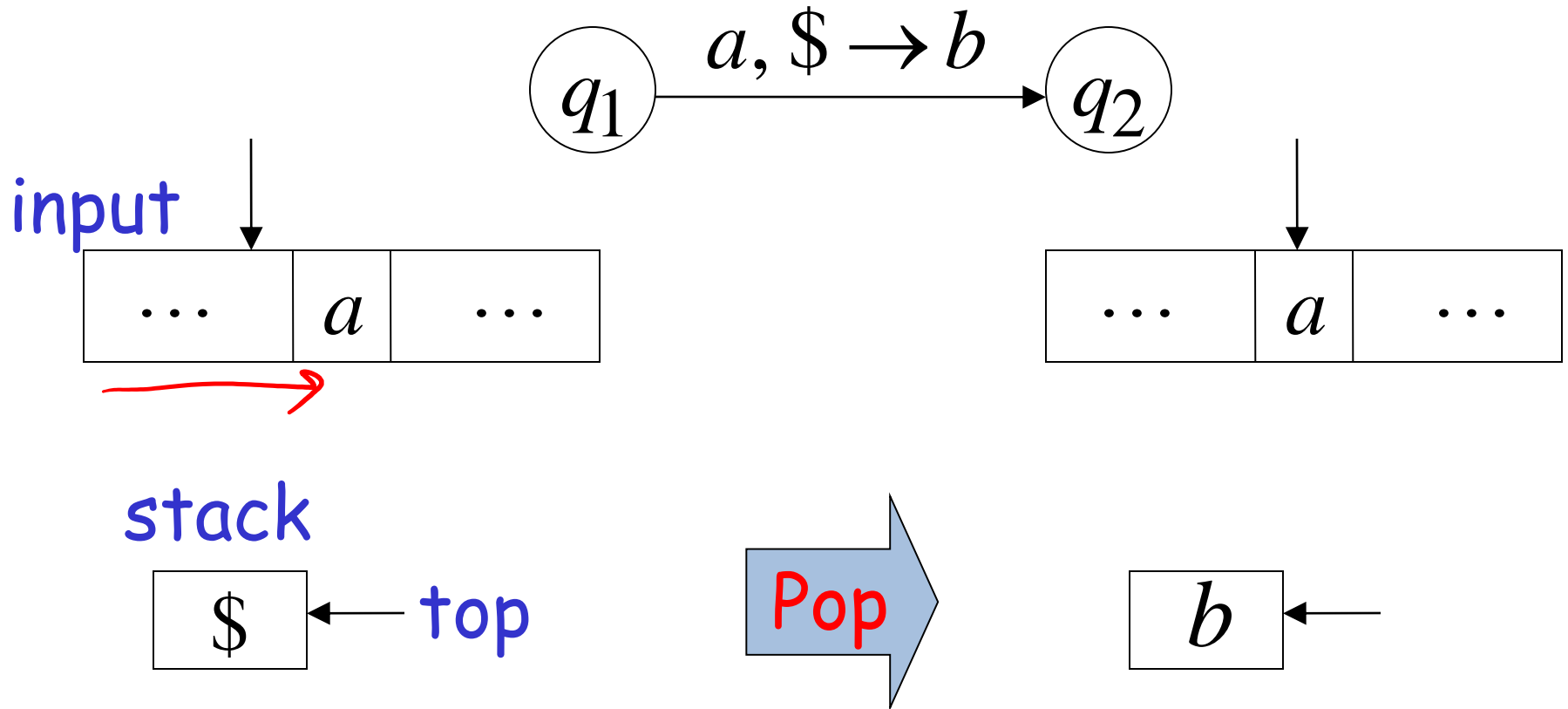
No transition is allowed to be followed  
When the stack is empty



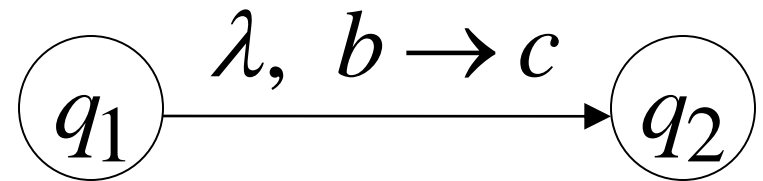
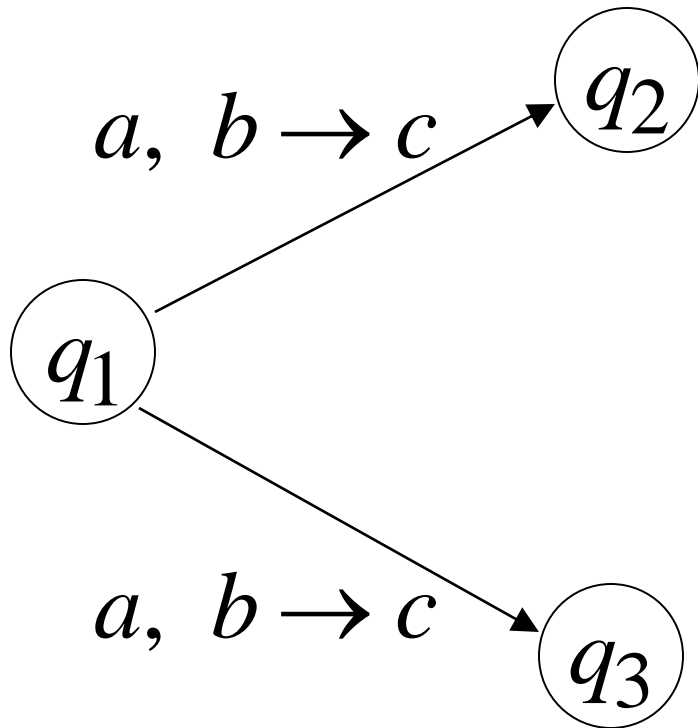
Empty stack



# A Good Transition



# Non-Determinism

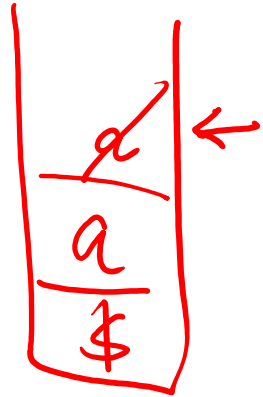
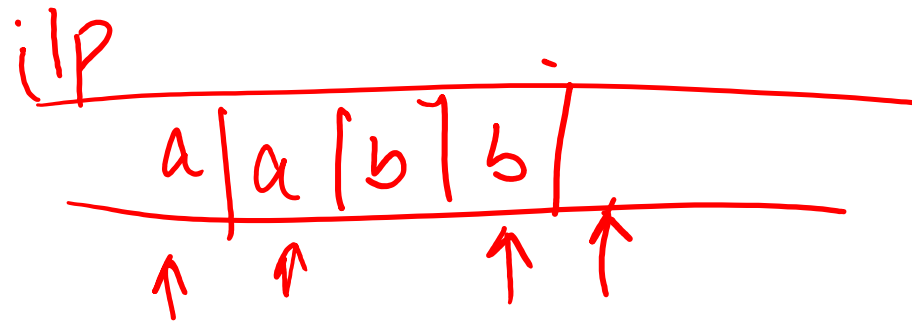


$\lambda$  – transition

These are allowed transitions in a  
Non-deterministic PDA (NPDA)

# NPDA: Non-Deterministic PDA

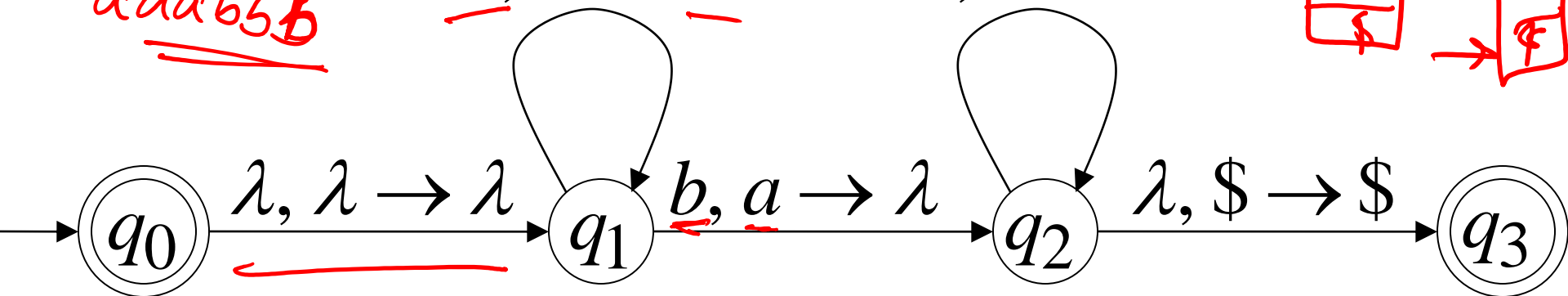
Example:



Handwritten string  $aaabbb$  with red underlines.

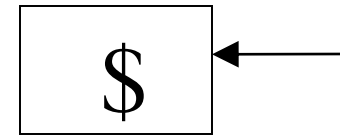
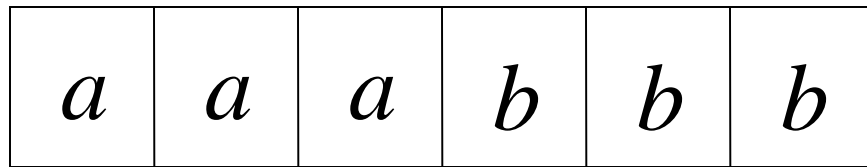
Transition rule:  $a, \lambda \rightarrow a$

Transition rule:  $b, a \rightarrow \lambda$



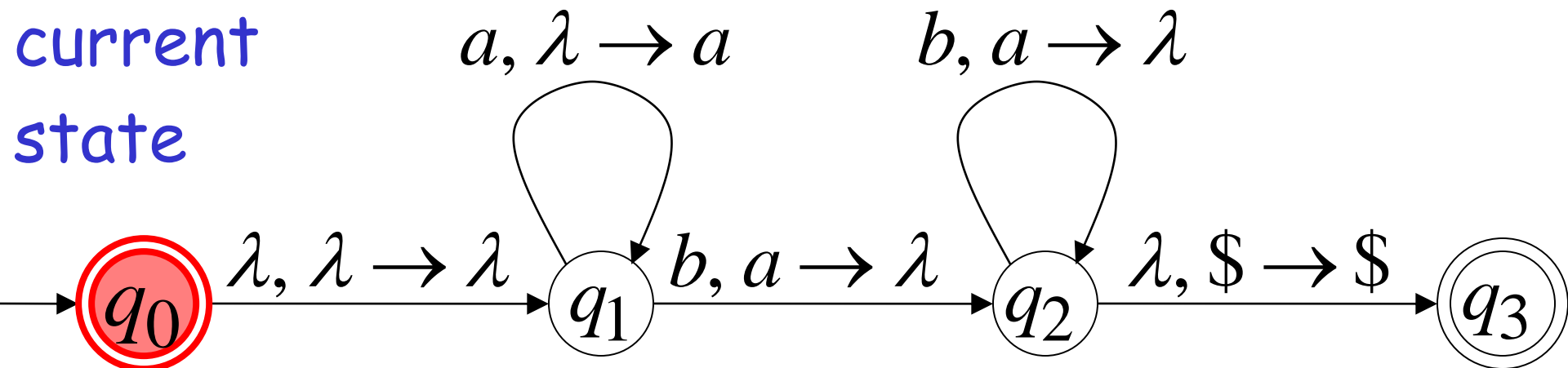
# Execution Example: Time 0

Input



Stack

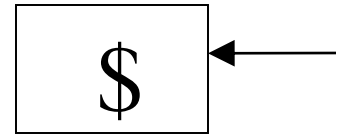
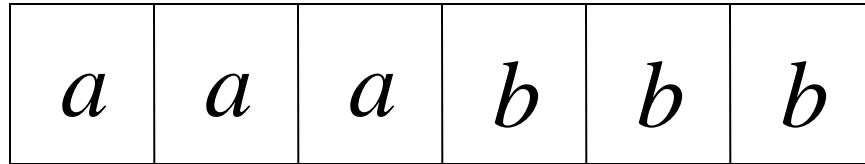
current  
state



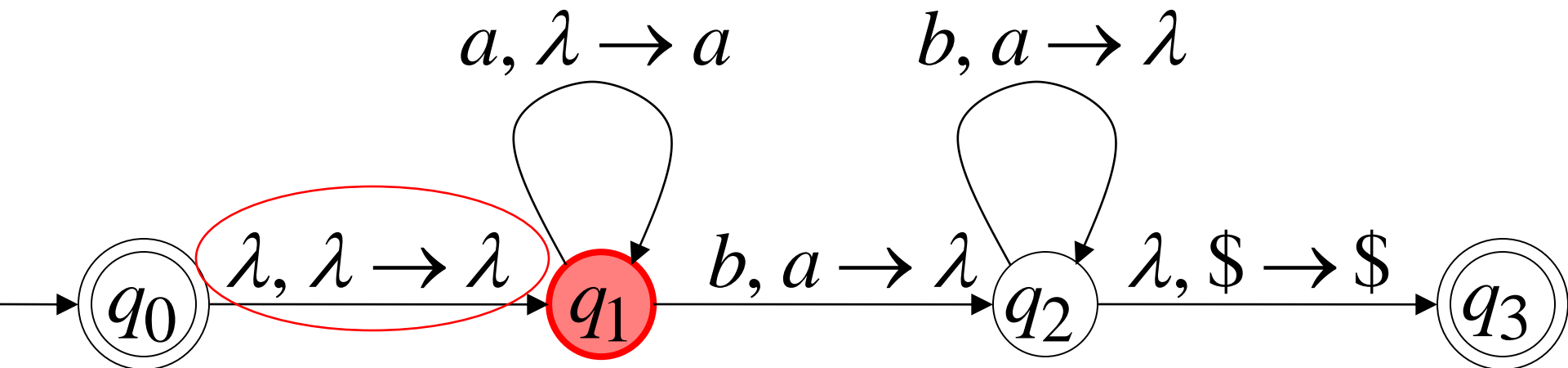


Time 1

Input

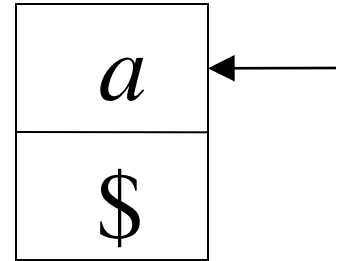
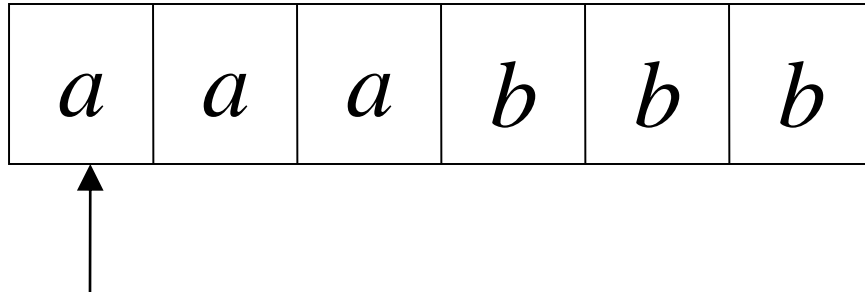


Stack

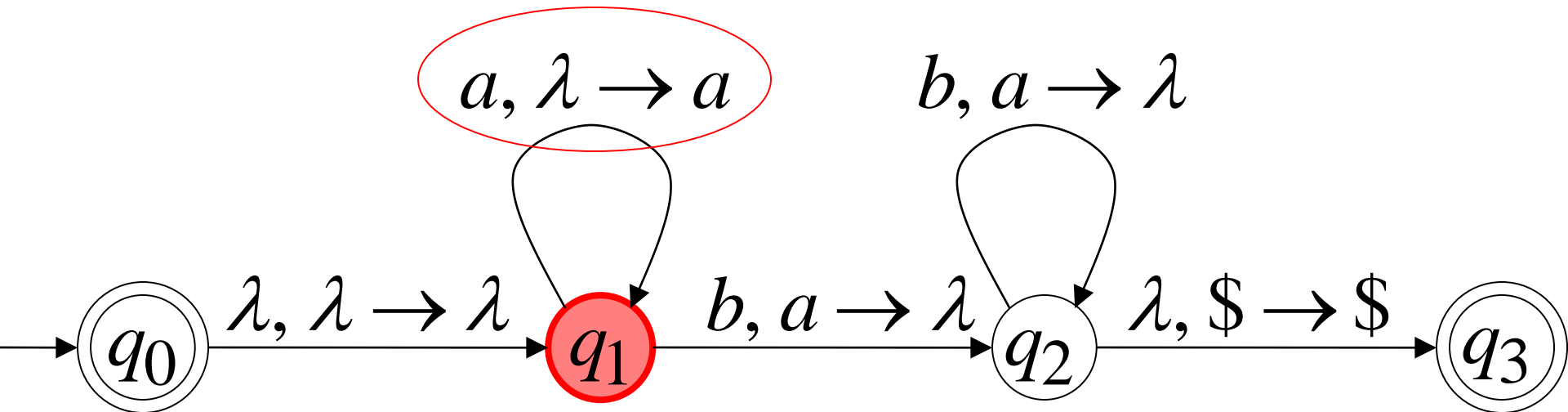


Time 2

Input

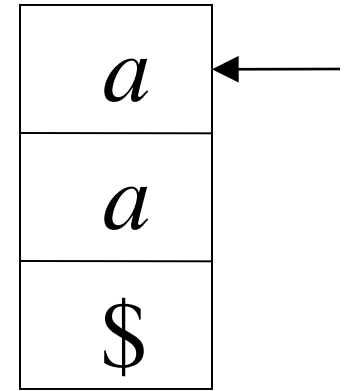
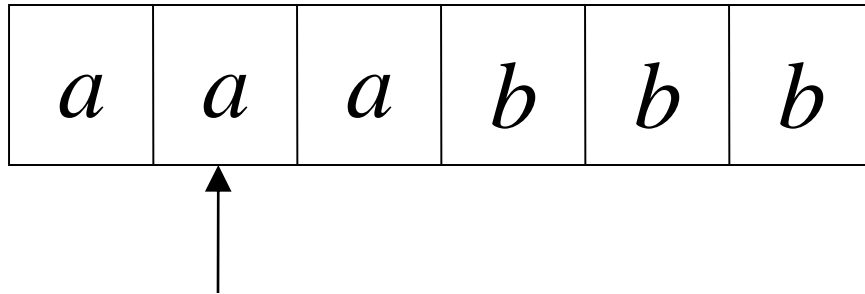


Stack

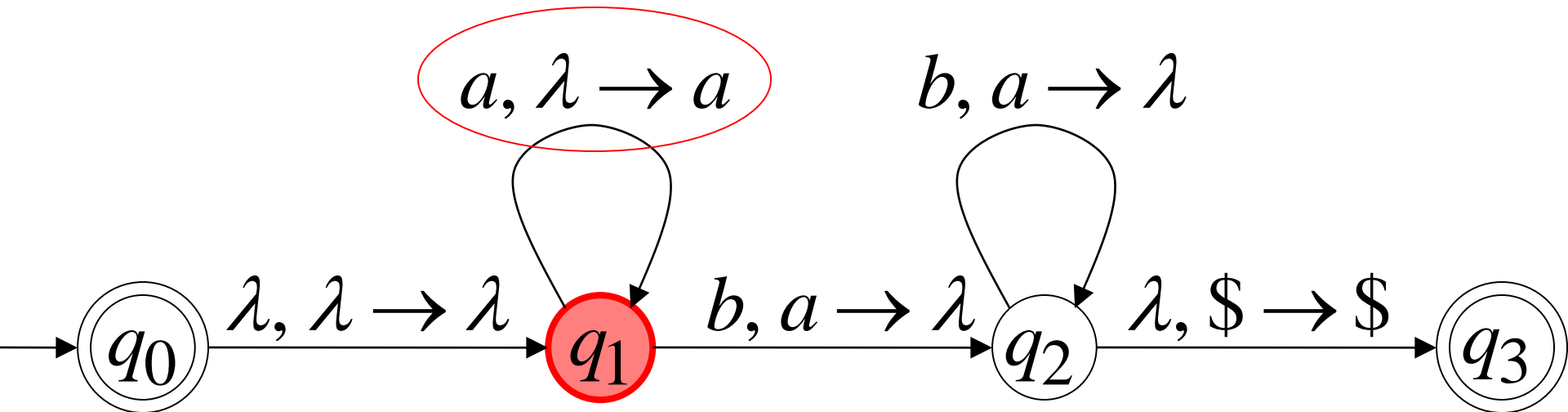


Time 3

Input

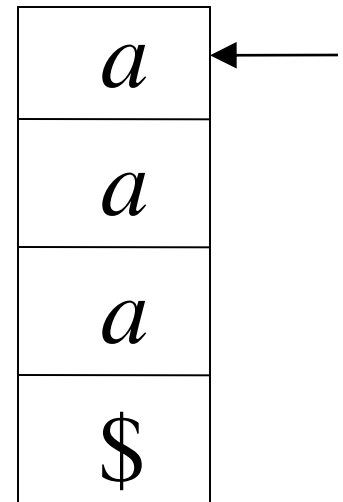
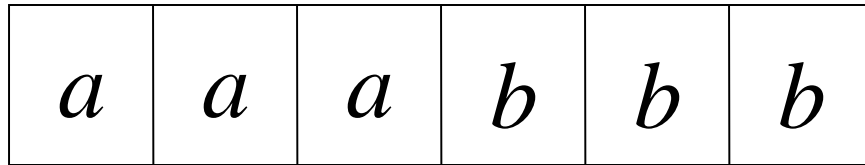


Stack

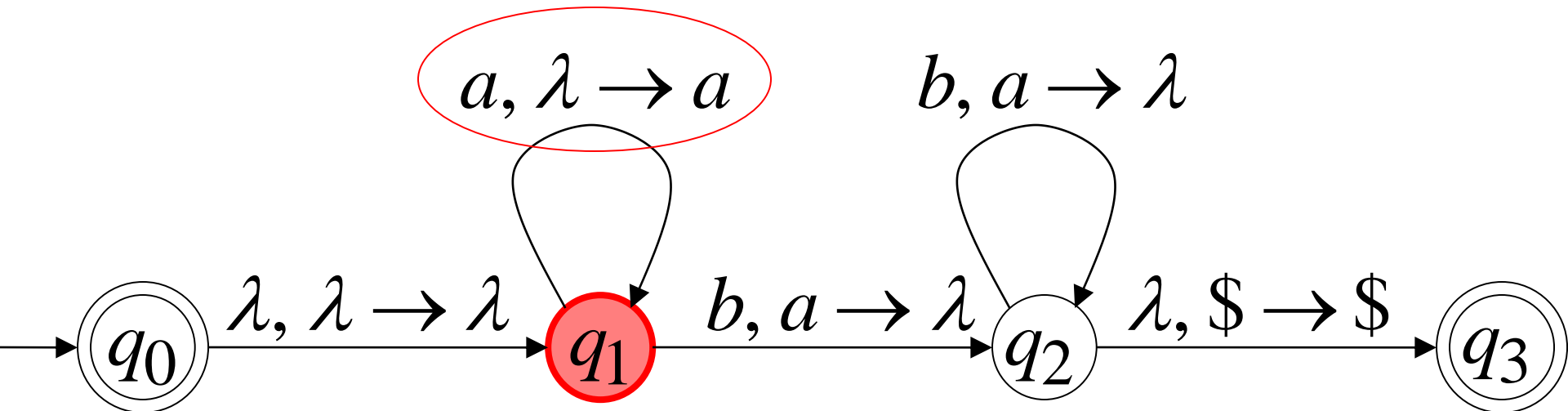


Time 4

Input

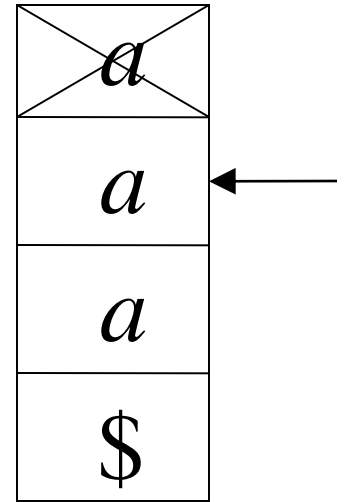
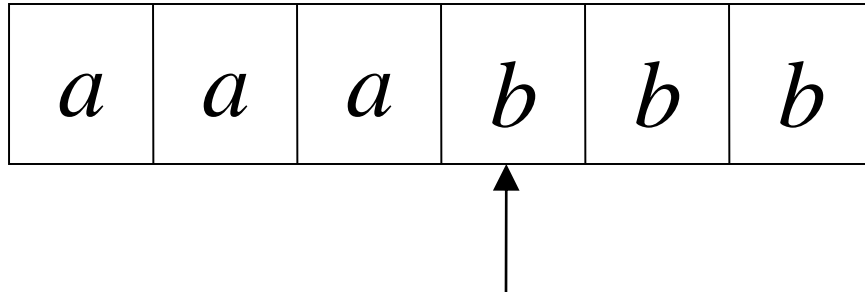


Stack

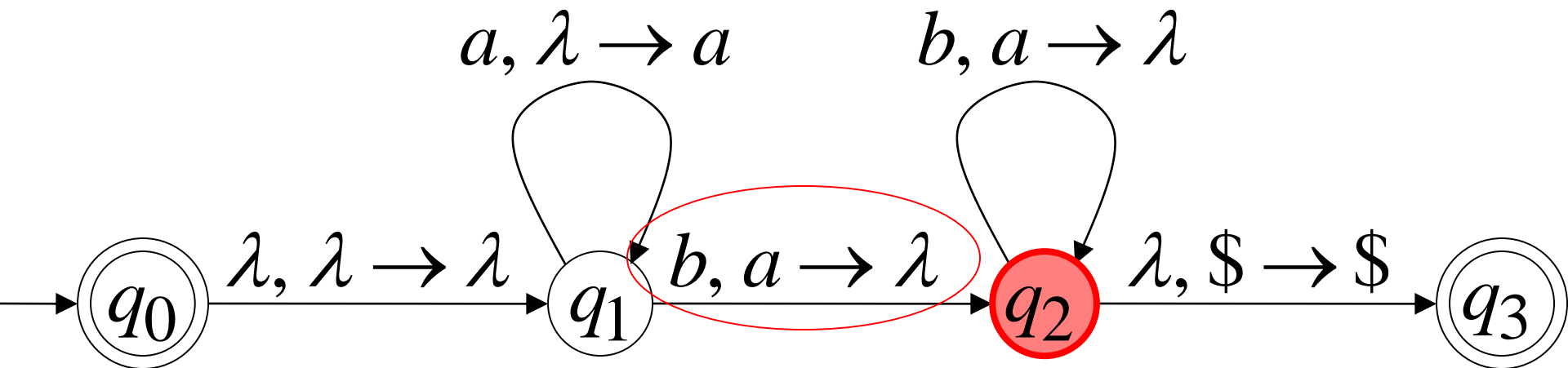


Time 5

Input

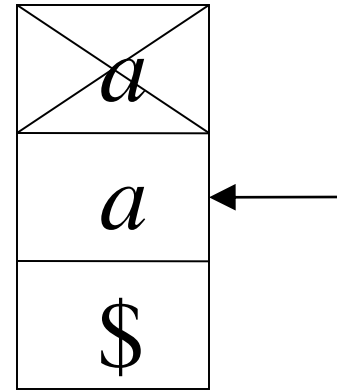
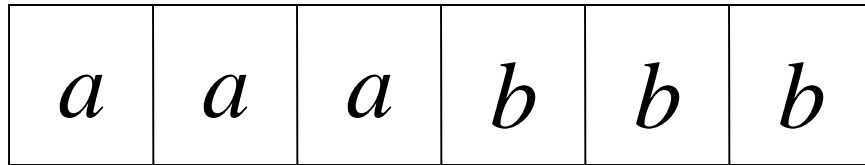


Stack

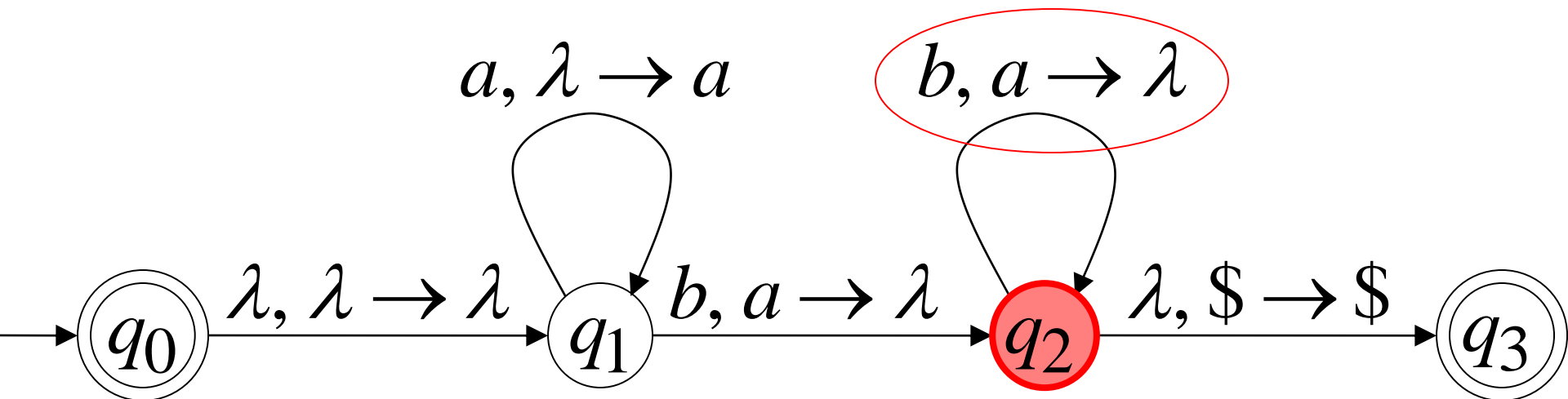


Time 6

Input

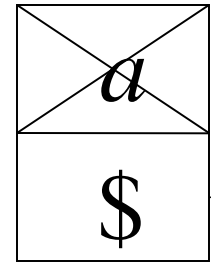
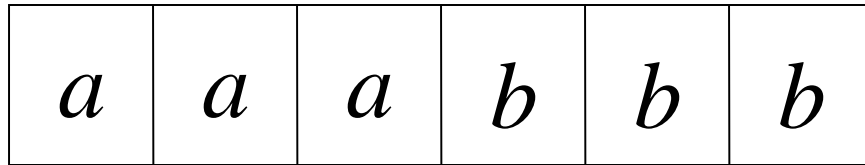


Stack

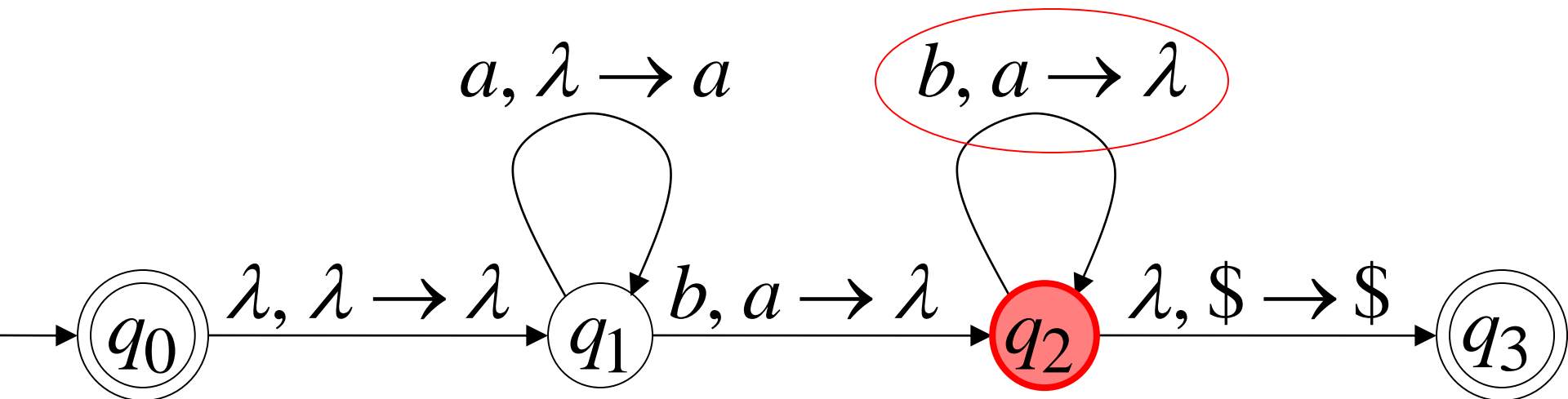


Time 7

Input

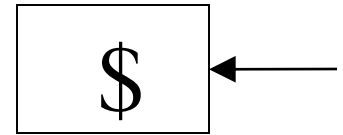
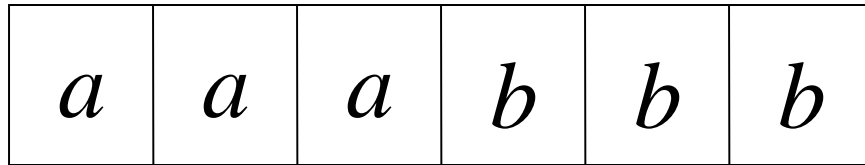


Stack

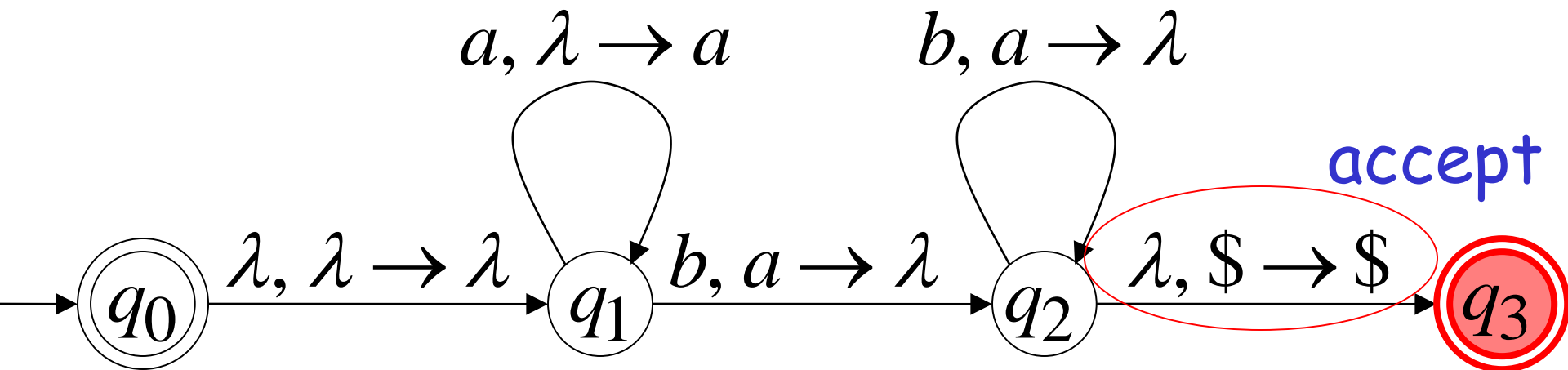


Time 8

Input



Stack





A string is accepted if there is  
a computation such that:

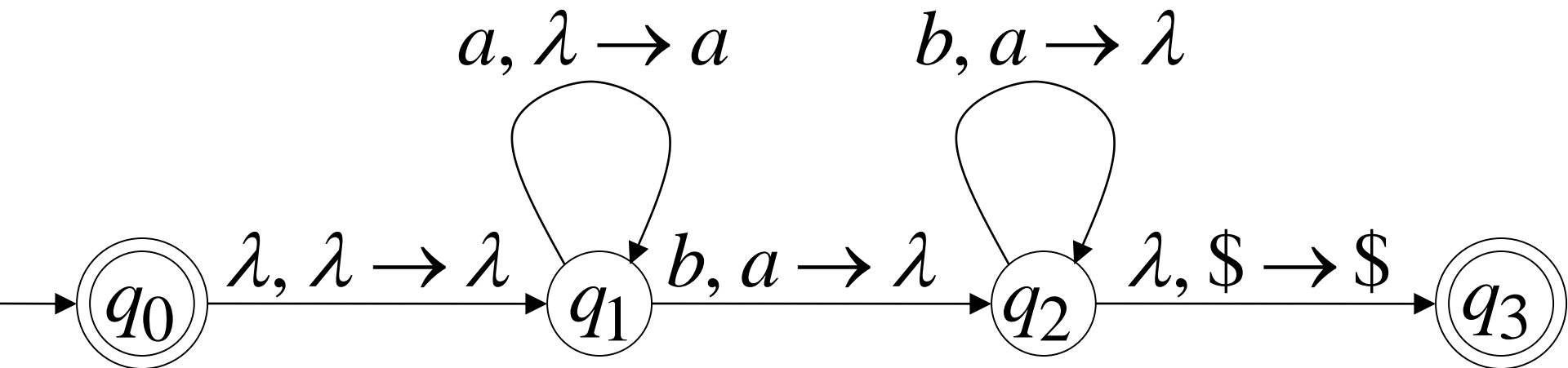
All the input is consumed

**AND**

The last state is a final state

At the end of the computation,  
we do not care about the stack contents

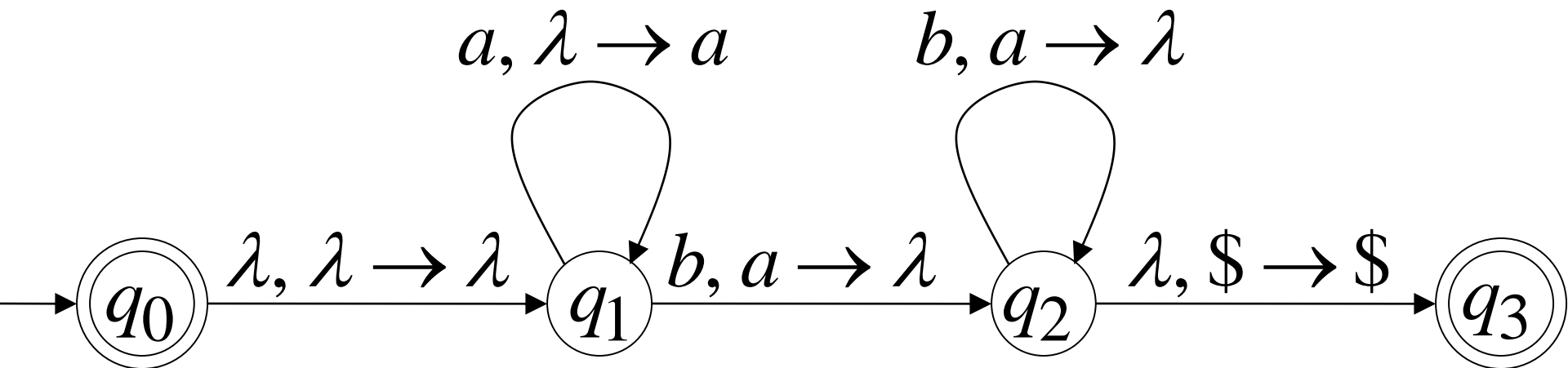
The input string *aaabbb*  
is accepted by the NPDA:



In general,

$$L = \{a^n b^n : n \geq 0\}$$

is the language accepted by the NPDA:



# Another NPDA example

NPDA  $M$

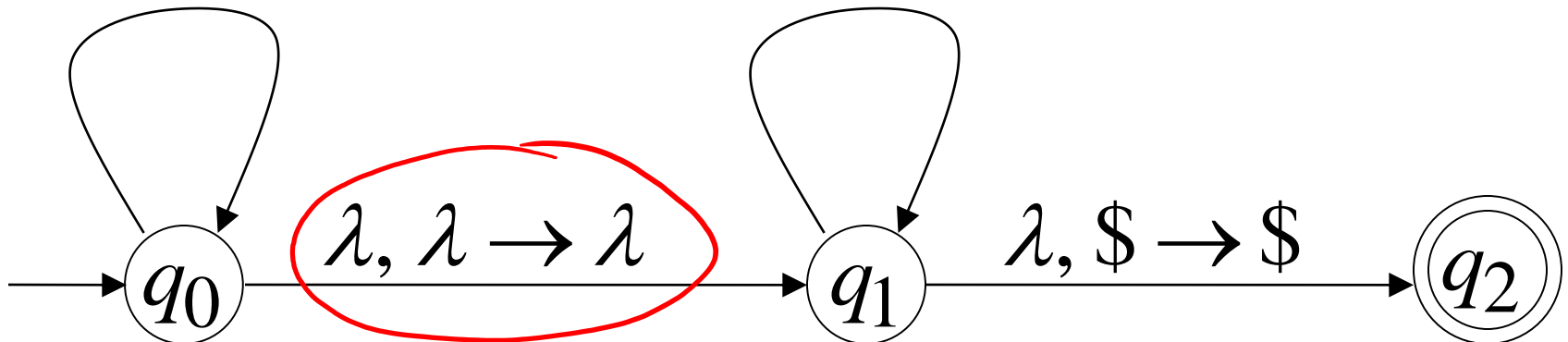
$$L(M) = \{ \underline{w} w^R \}$$

$a, \lambda \rightarrow a$

$a, a \rightarrow \lambda$

$b, \lambda \rightarrow b$

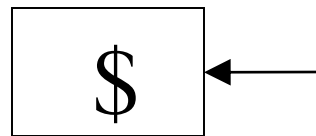
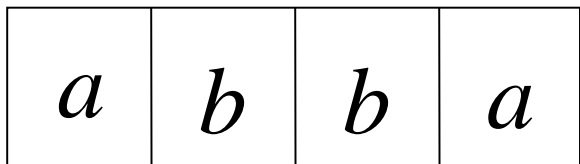
$b, b \rightarrow \lambda$



Execution Example:

Time 0

Input



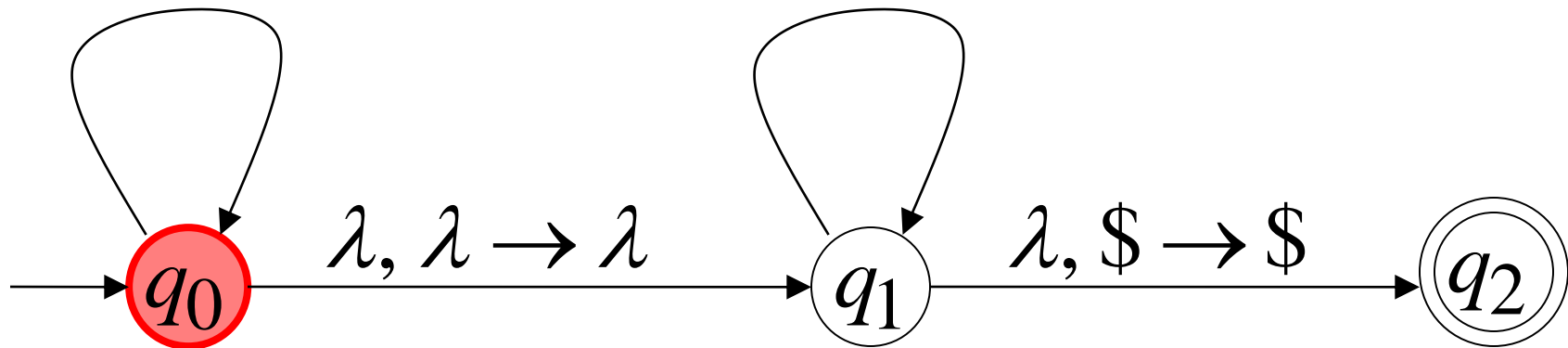
Stack

$a, \lambda \rightarrow a$

$a, a \rightarrow \lambda$

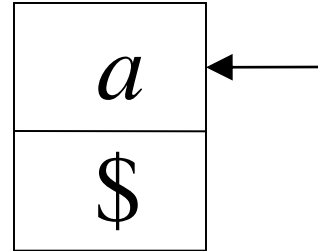
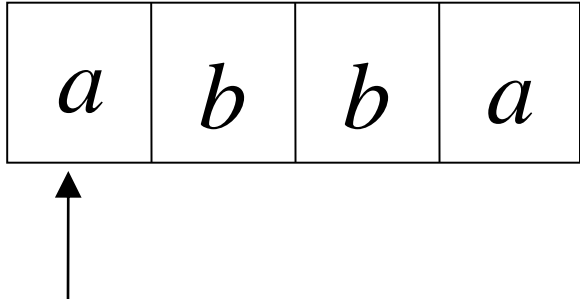
$b, \lambda \rightarrow b$

$b, b \rightarrow \lambda$



Time 1

Input



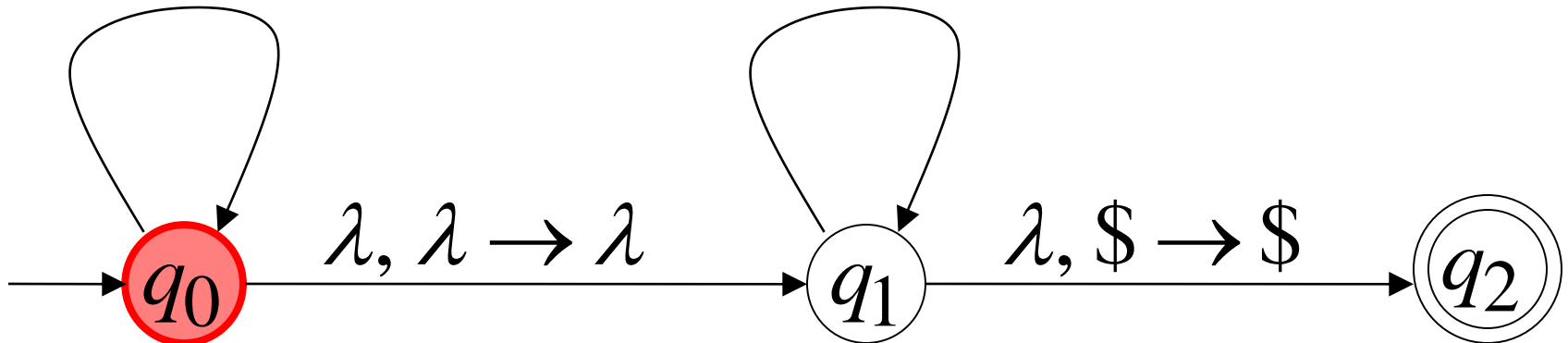
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

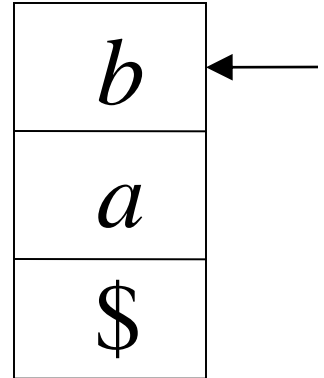
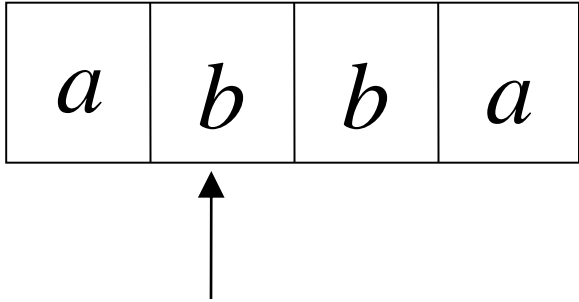
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 2

Input



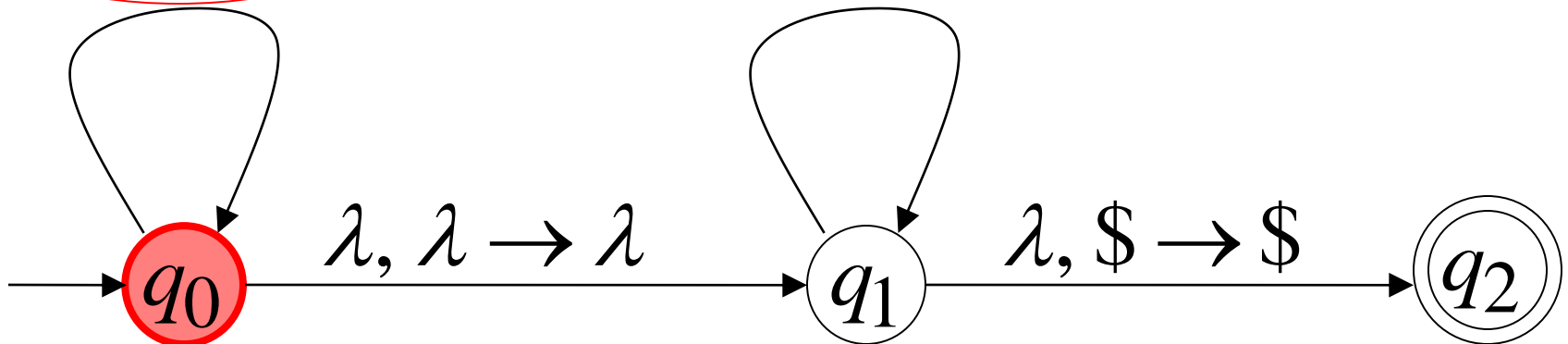
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

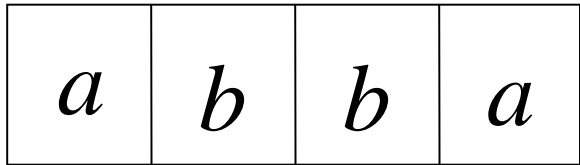
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$

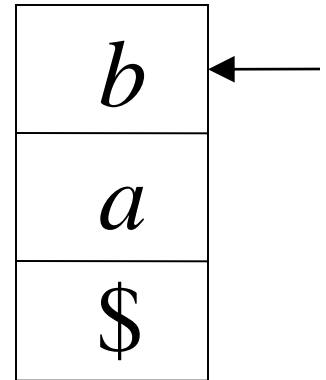


Time 3

Input



Guess the middle  
of string



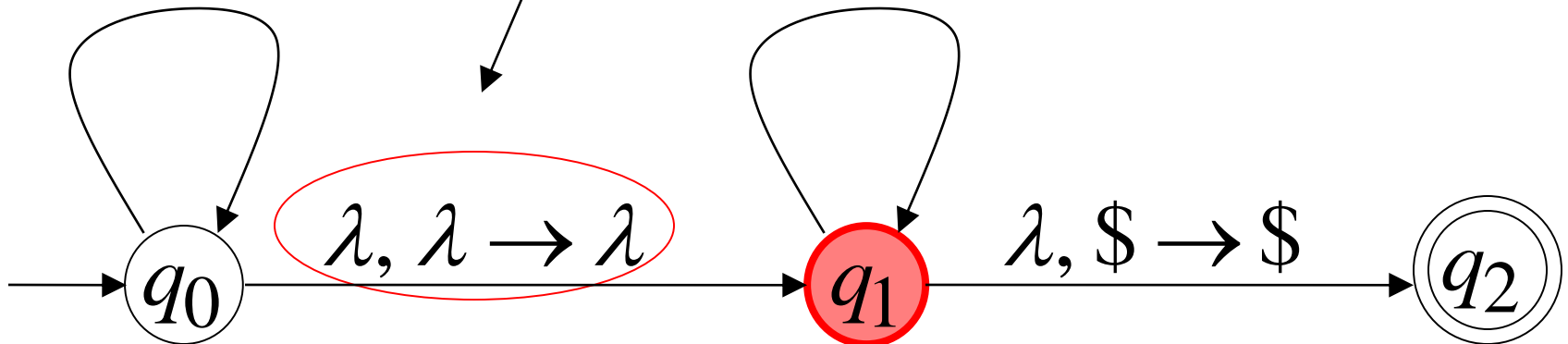
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$  ✓

$a, a \rightarrow \lambda$

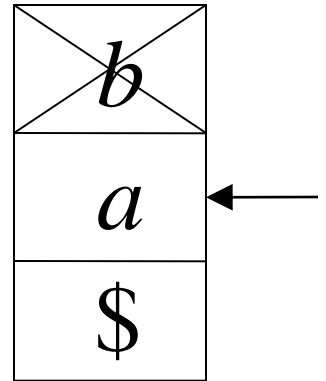
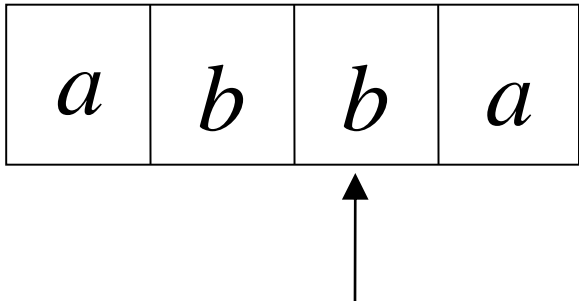
$b, b \rightarrow \lambda$





Time 4

Input



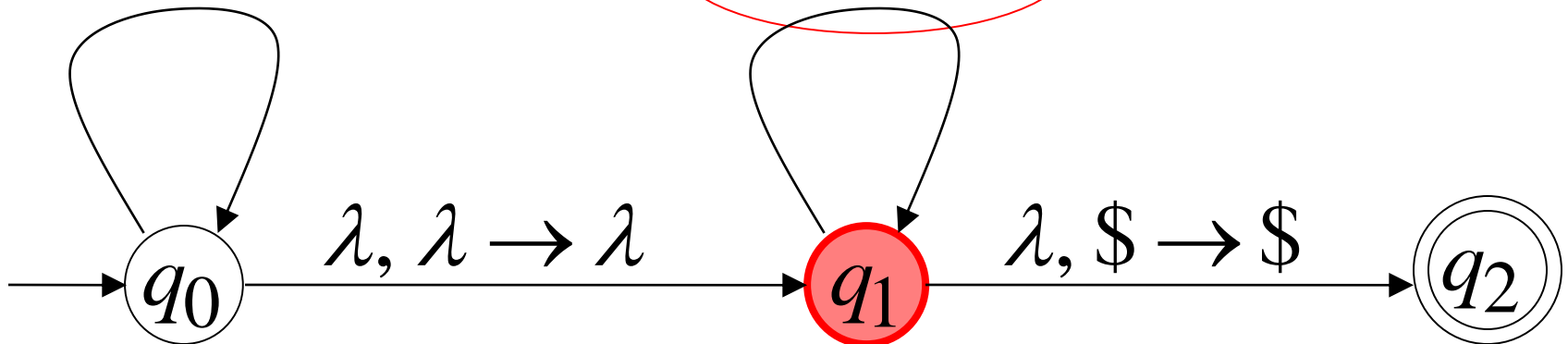
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

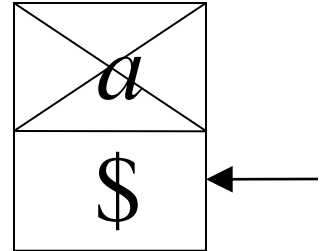
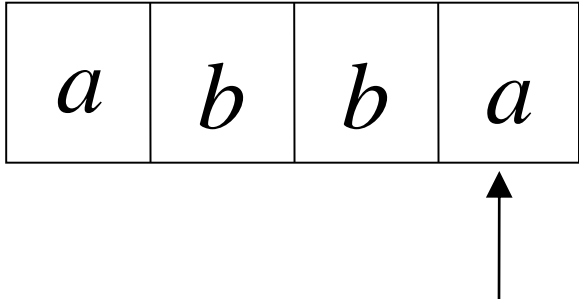
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 5

Input



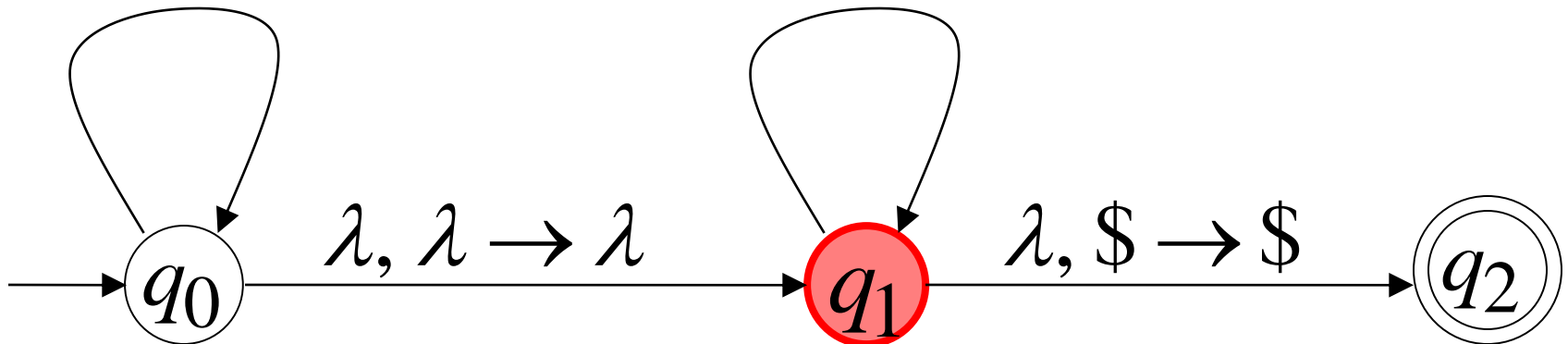
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

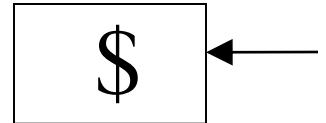
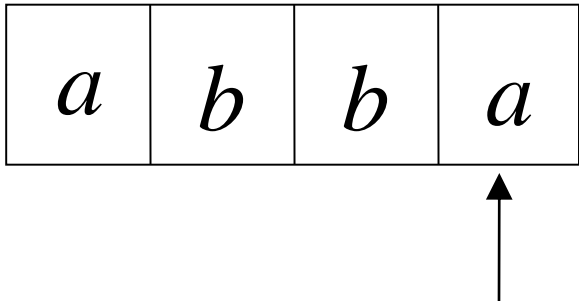
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 6

Input



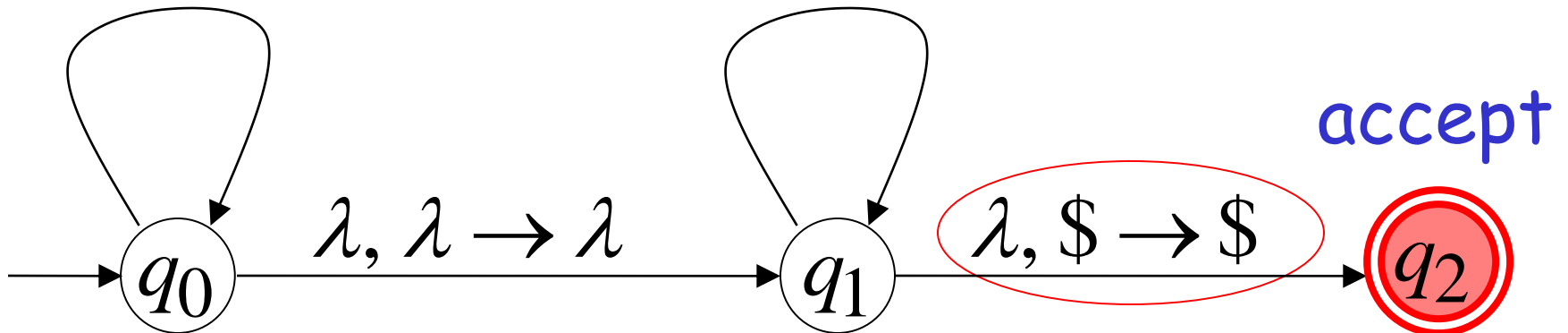
Stack

$a, \lambda \rightarrow a$

$a, a \rightarrow \lambda$

$b, \lambda \rightarrow b$

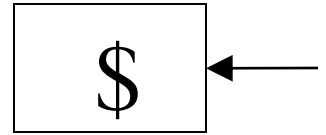
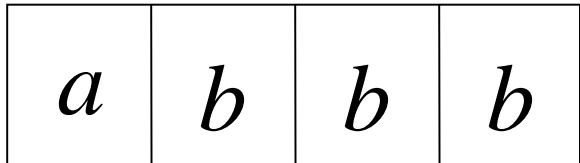
$b, b \rightarrow \lambda$



Rejection Example:

Time 0

Input



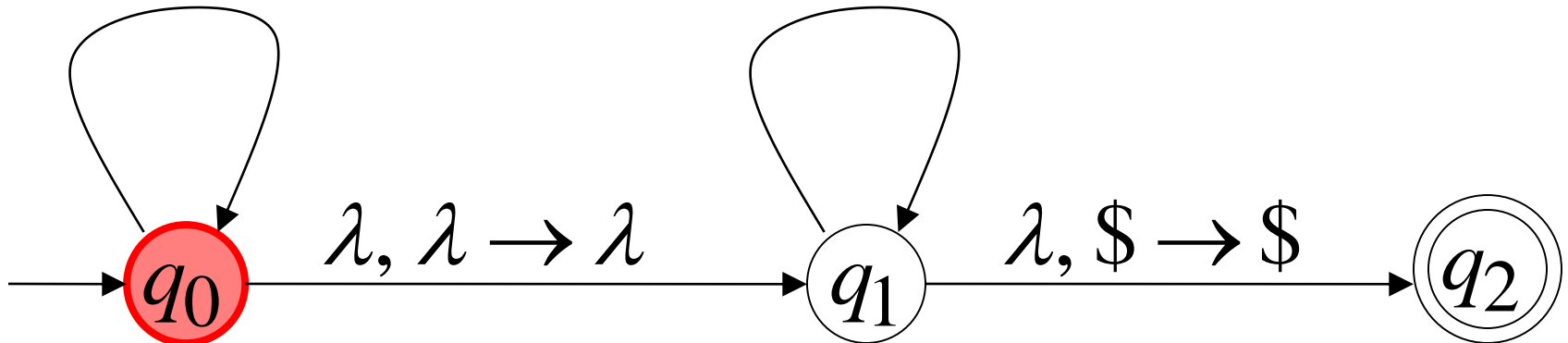
Stack

$a, \lambda \rightarrow a$

$a, a \rightarrow \lambda$

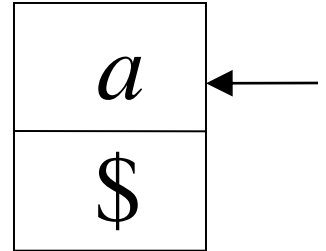
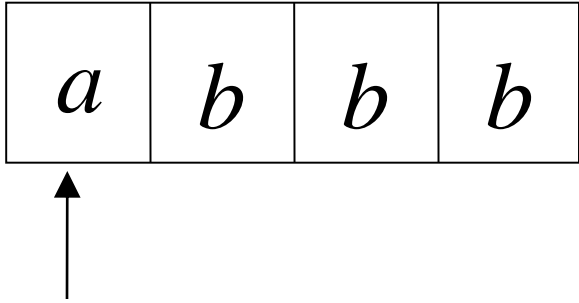
$b, \lambda \rightarrow b$

$b, b \rightarrow \lambda$



Time 1

Input



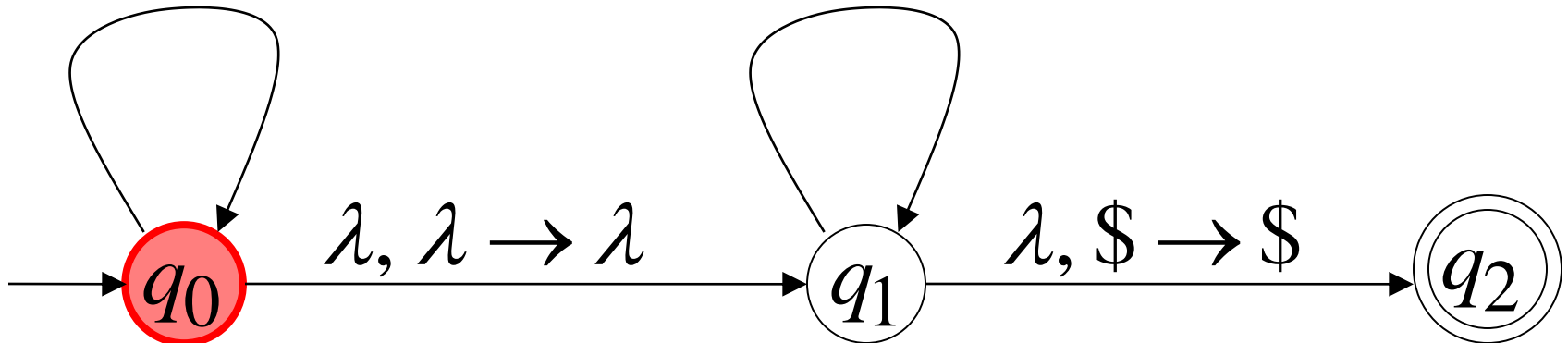
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

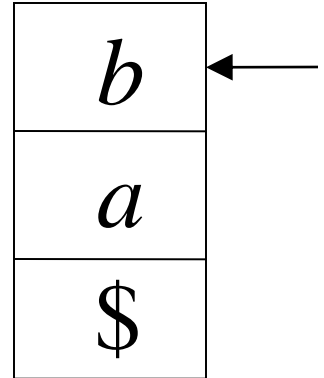
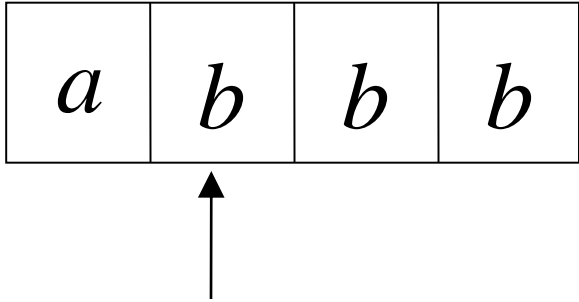
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 2

Input



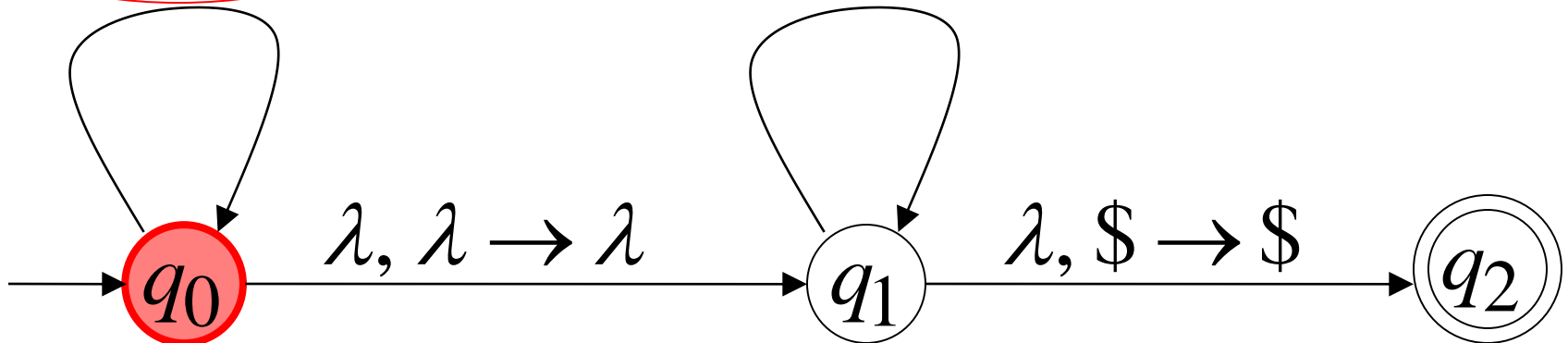
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

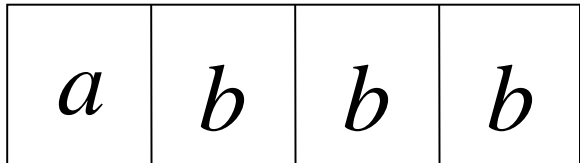
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$

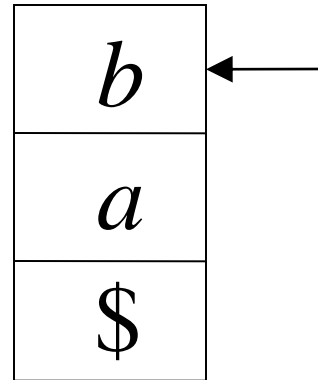


Time 3

Input



Guess the middle  
of string



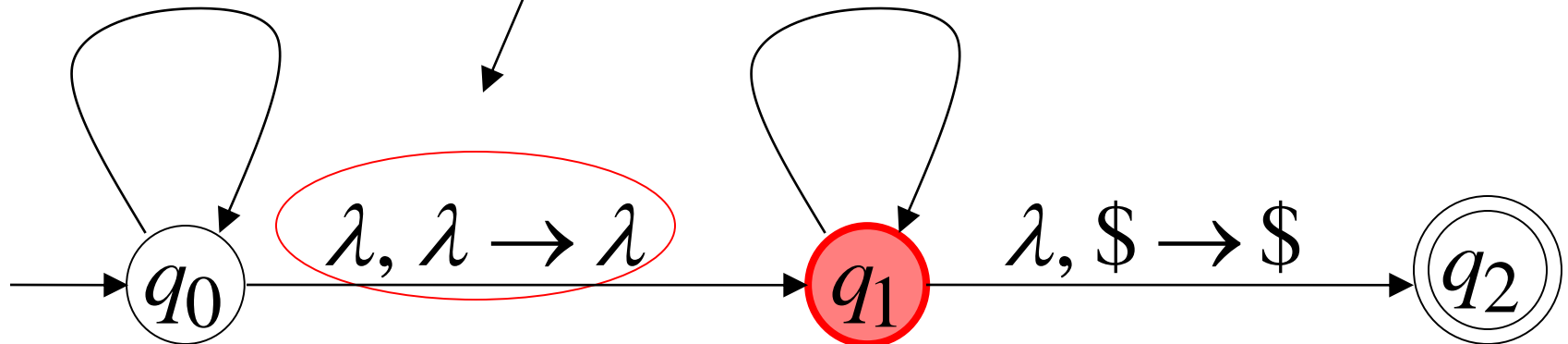
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

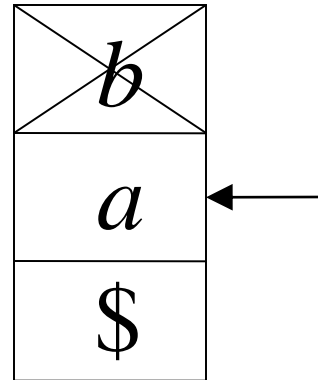
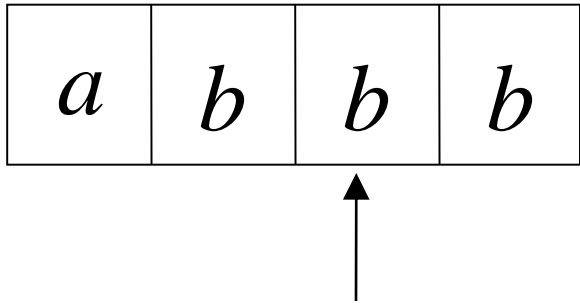
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 4

Input



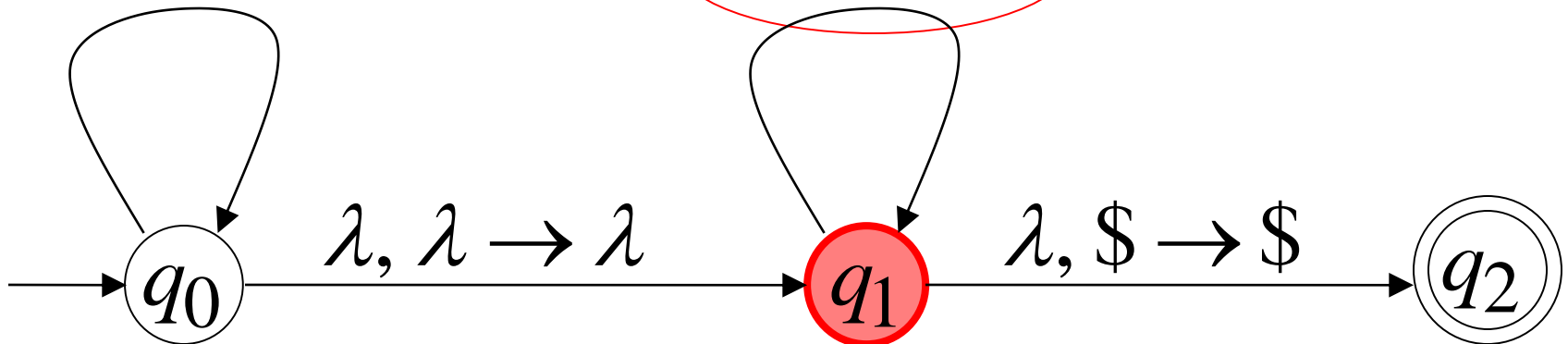
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$

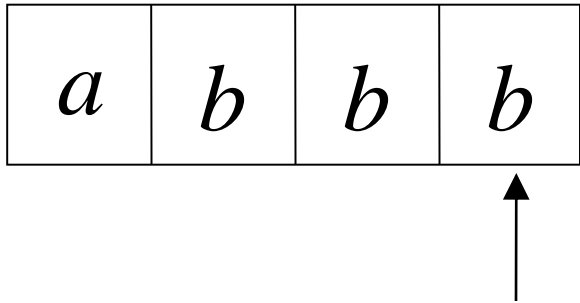




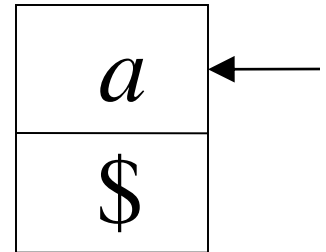
Time 5

Input

There is no possible transition.



Input is not consumed



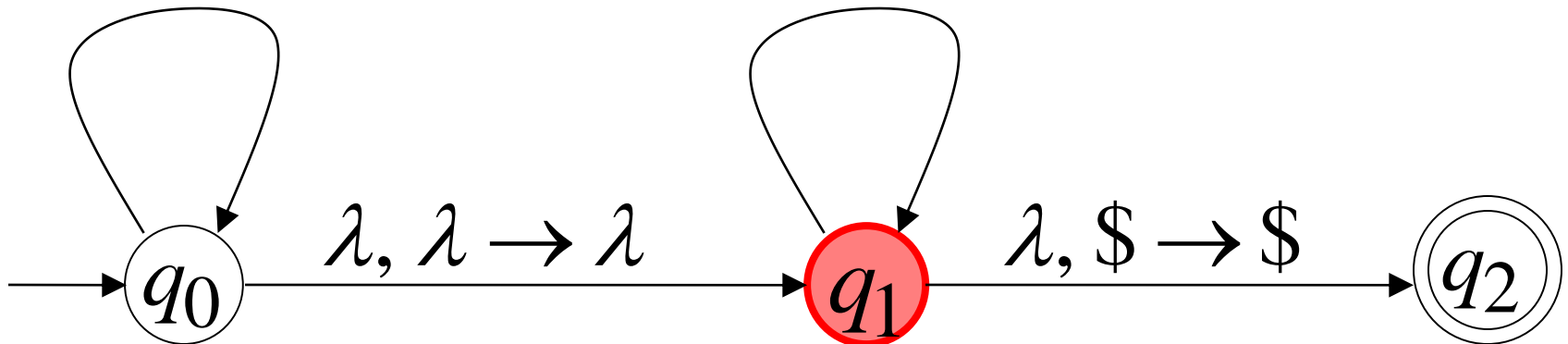
Stack

$a, \lambda \rightarrow a$

$a, a \rightarrow \lambda$

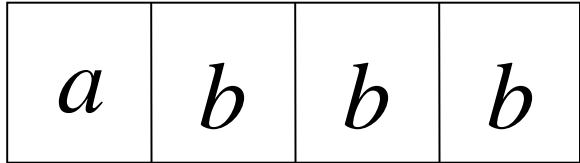
$b, \lambda \rightarrow b$

$b, b \rightarrow \lambda$

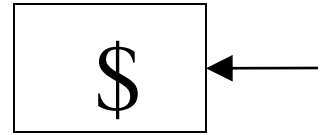


Another computation on same string:

Input



Time 0



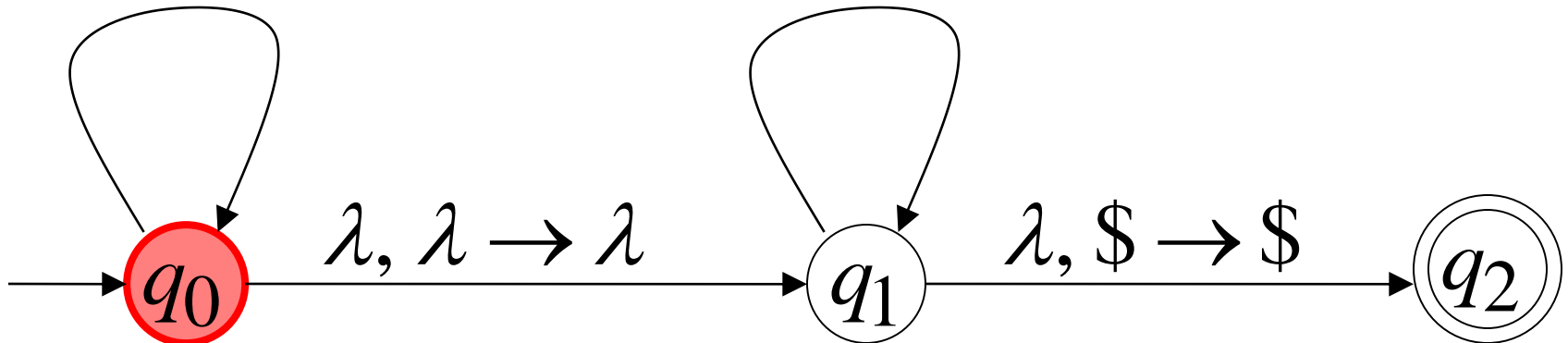
Stack

$a, \lambda \rightarrow a$

$a, a \rightarrow \lambda$

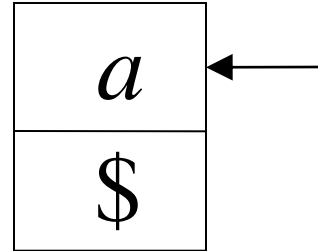
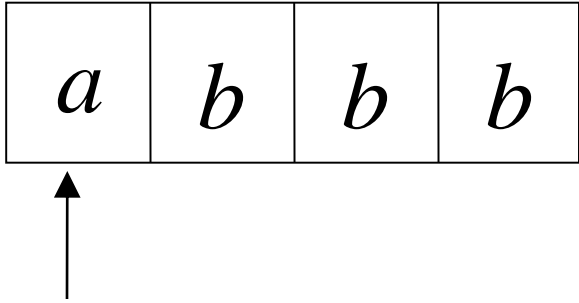
$b, \lambda \rightarrow b$

$b, b \rightarrow \lambda$



Time 1

Input



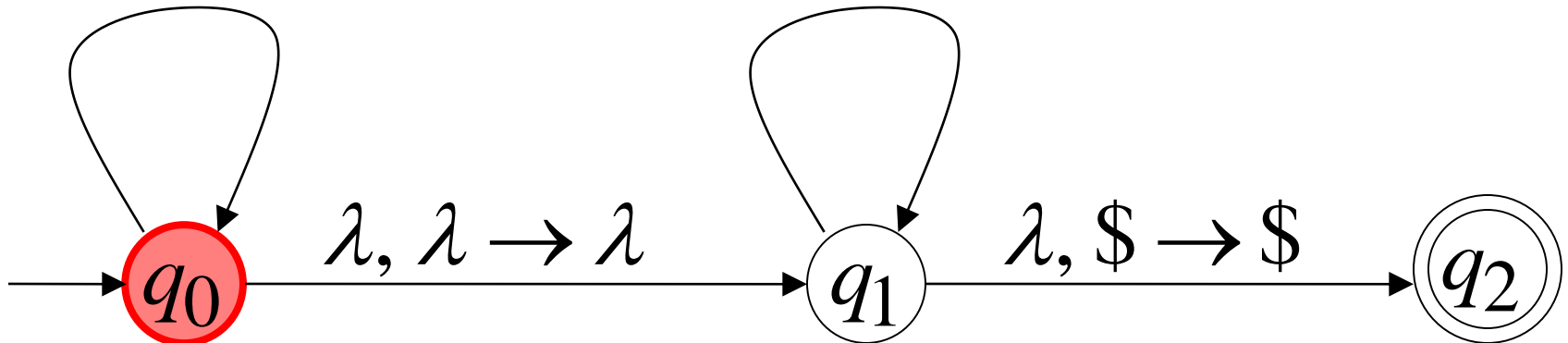
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

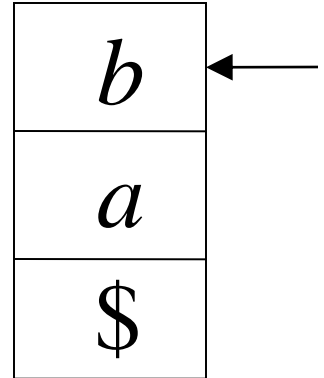
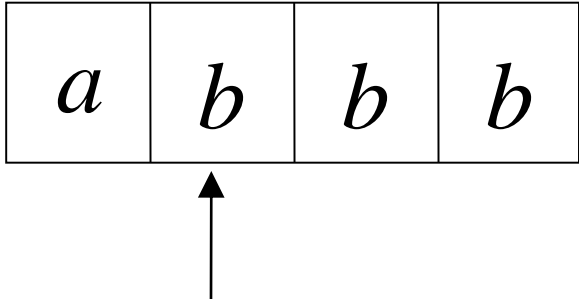
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 2

Input



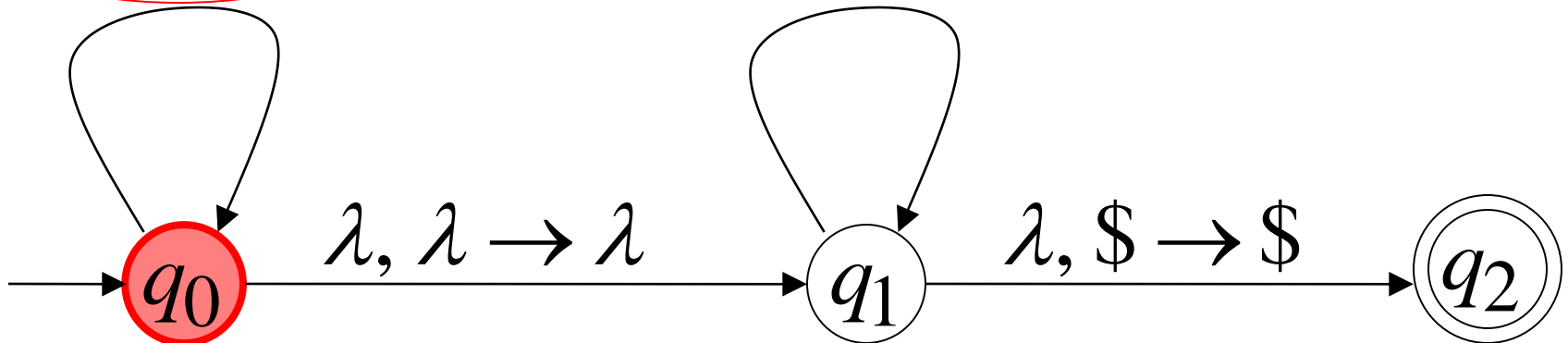
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

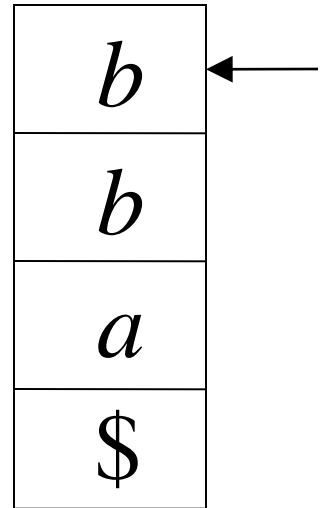
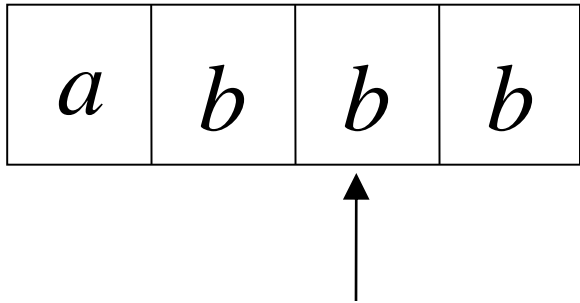
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 3

Input



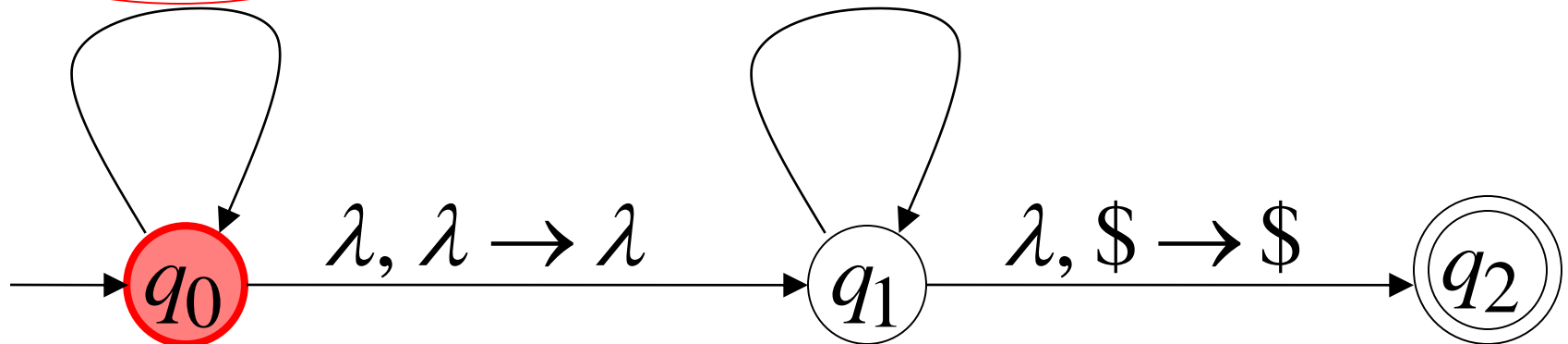
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

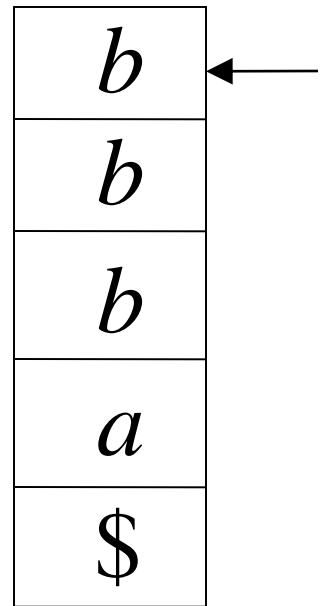
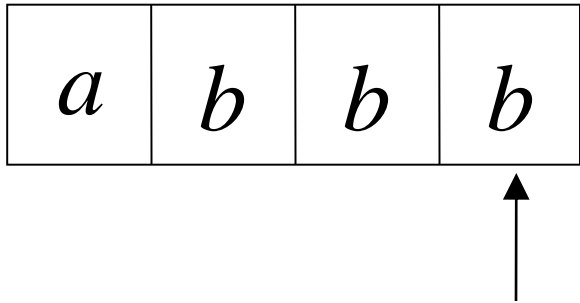
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



Time 4

Input



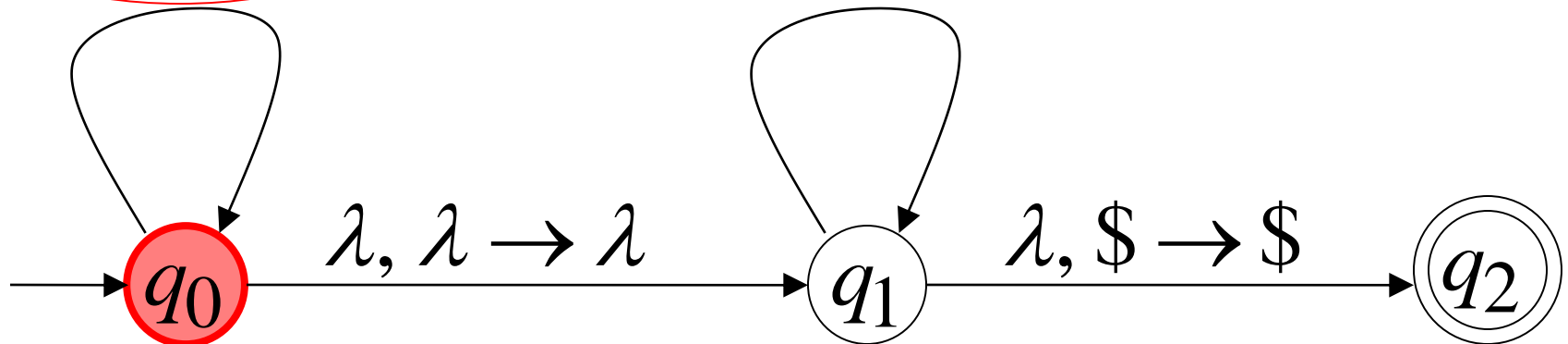
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

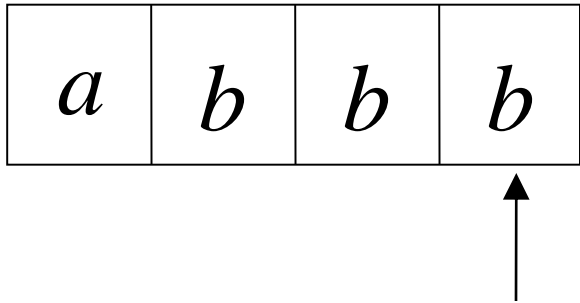
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$

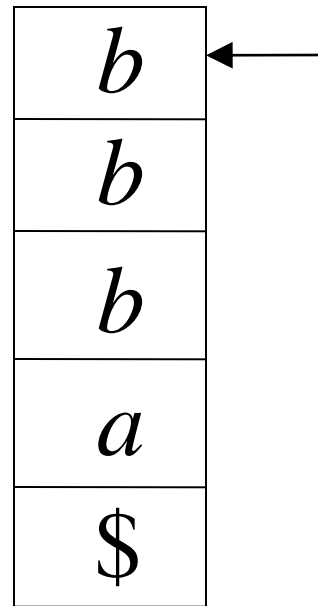


Time 5

Input



No final state  
is reached



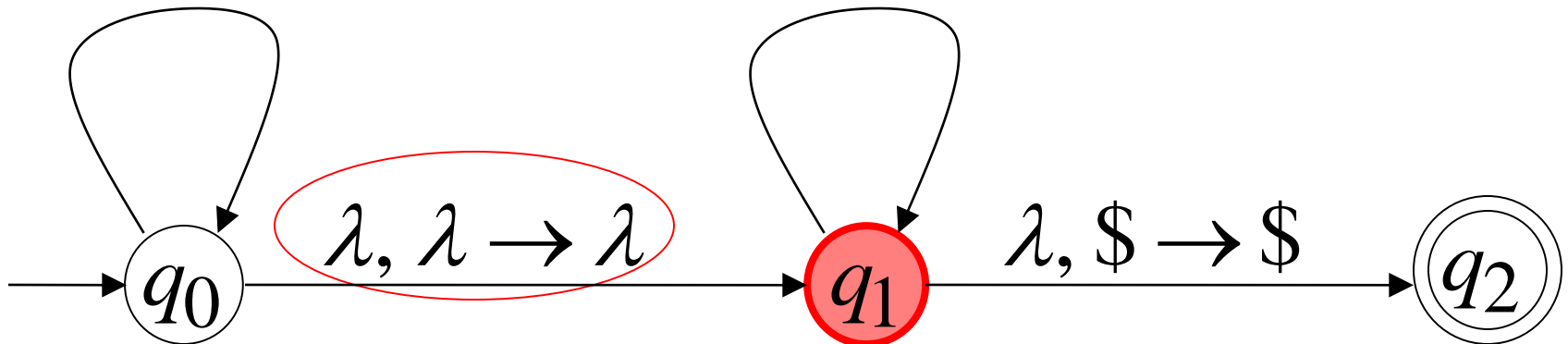
Stack

$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

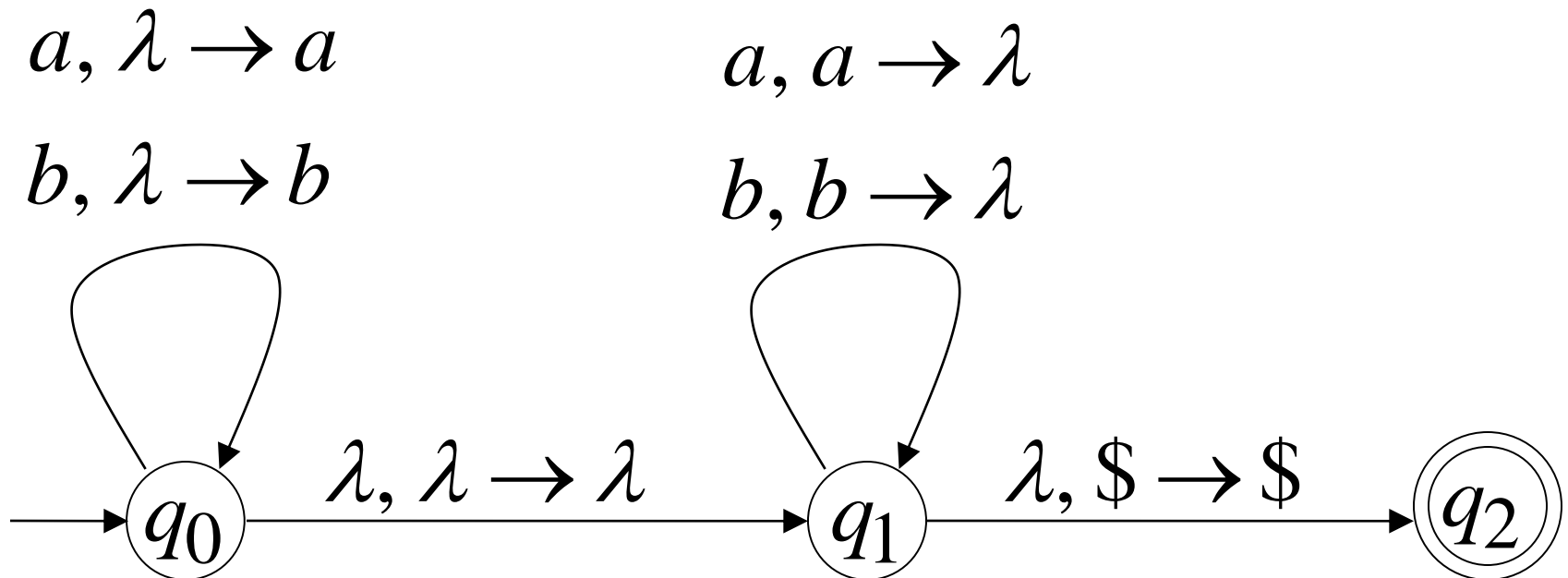
$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$



There is no computation  
that accepts string  $abbb$

$$abbb \notin L(M)$$





A string is rejected if there is  
NO computation such that:

All the input is consumed

**AND**

The last state is a final state

At the end of the computation,  
we do not care about the stack contents

In other words, a string is rejected  
if in every computation with this string:

The input cannot be consumed

OR

The input is consumed and the last  
state is not a final state

OR

The stack head moves below the  
bottom of the stack

# Another NPDA example

NPDA  $M$

$$L(M) = \{a^n b^m : n \geq m - 1\}$$

# Another NPDA example

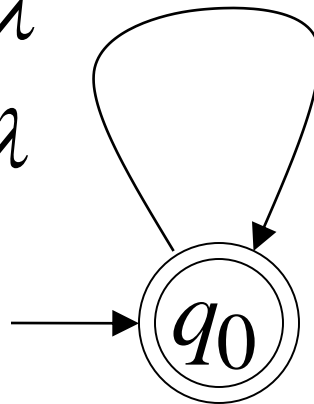
NPDA  $M$

$$L(M) = \{a^n b^m : n \geq m - 1\}$$

$$a, \lambda \rightarrow a$$

$$b, a \rightarrow \lambda$$

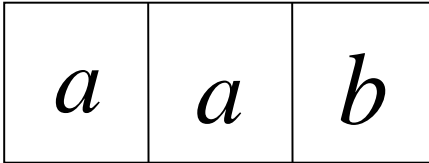
$$b, \$ \rightarrow \lambda$$



Execution Example:

Time 0

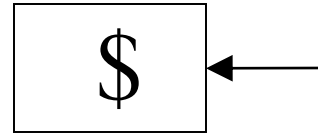
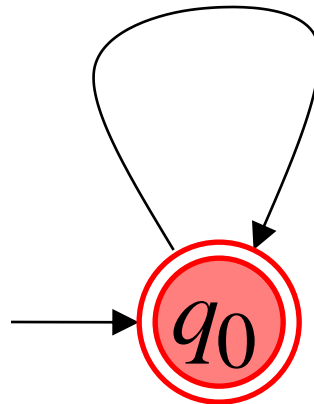
Input



$a, \lambda \rightarrow a$

$b, a \rightarrow \lambda$

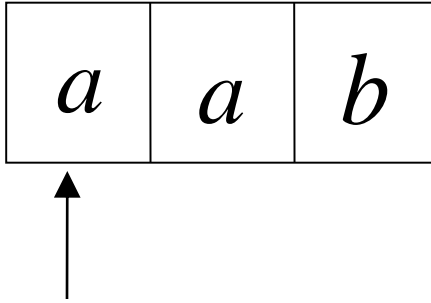
$b, \$ \rightarrow \lambda$



Stack

Time 1

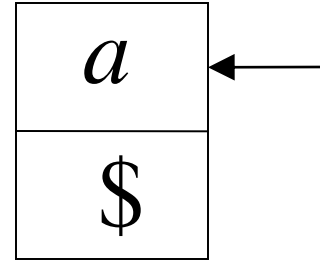
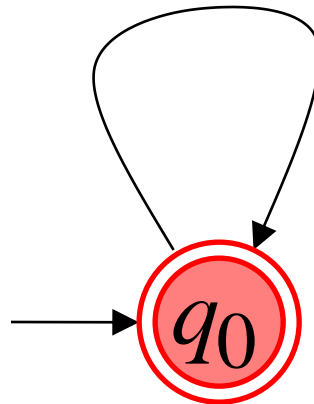
Input



$$a, \lambda \rightarrow a$$

$$b, a \rightarrow \lambda$$

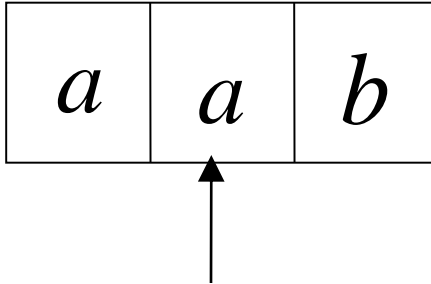
$$b, \$ \rightarrow \lambda$$



Stack

Time 2

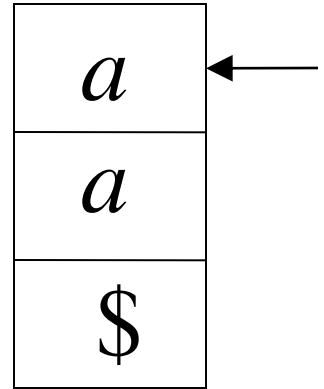
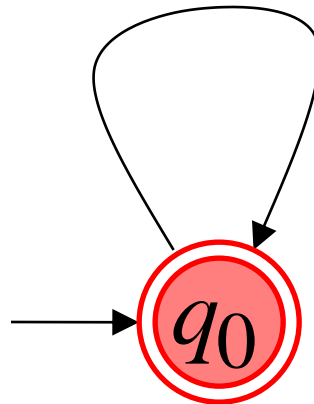
Input



$$a, \lambda \rightarrow a$$

$$b, a \rightarrow \lambda$$

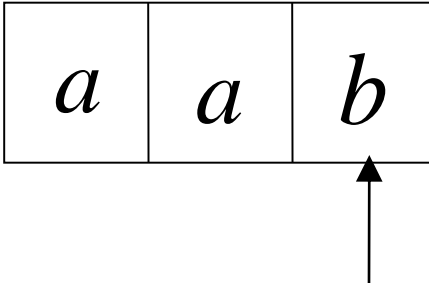
$$b, \$ \rightarrow \lambda$$



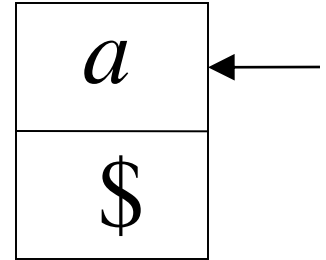
Stack

Time 3

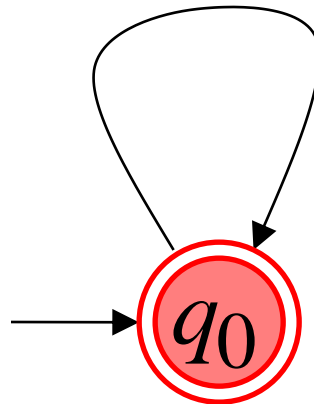
Input



$a, \lambda \rightarrow a$   
 $b, a \rightarrow \lambda$   
 $b, \$ \rightarrow \lambda$



Stack



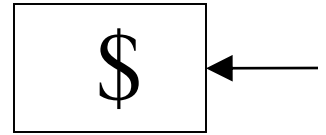
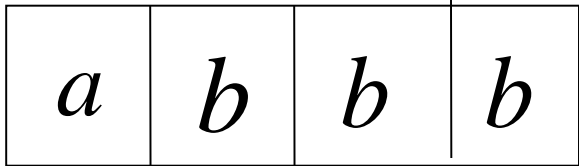
accept



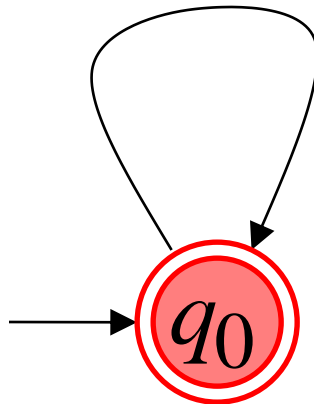
Rejection example:

Time 0

Input

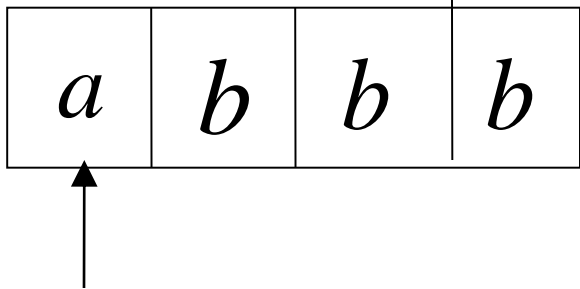


Stack



Time 1

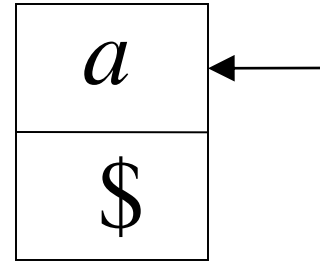
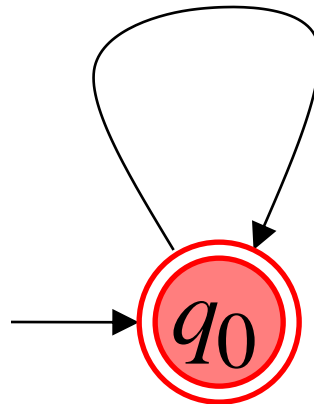
Input



$$a, \lambda \rightarrow a$$

$$b, a \rightarrow \lambda$$

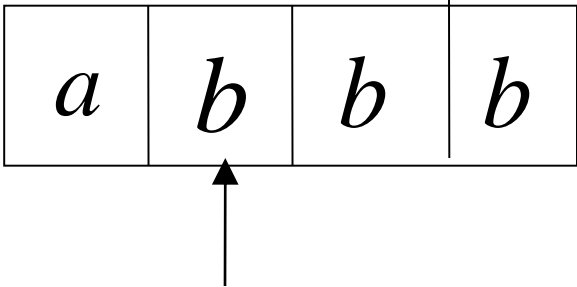
$$b, \$ \rightarrow \lambda$$



Stack

Time 2

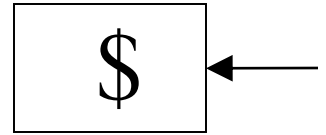
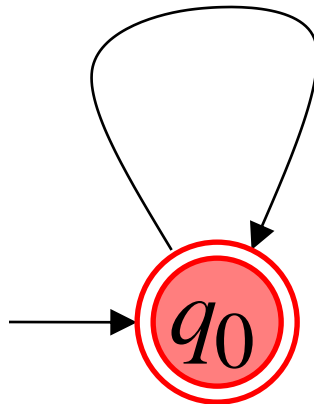
Input



$a, \lambda \rightarrow a$

$b, a \rightarrow \lambda$

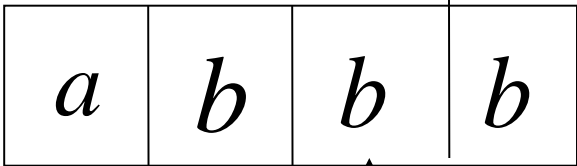
$b, \$ \rightarrow \lambda$



Stack

Time 3

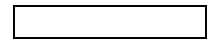
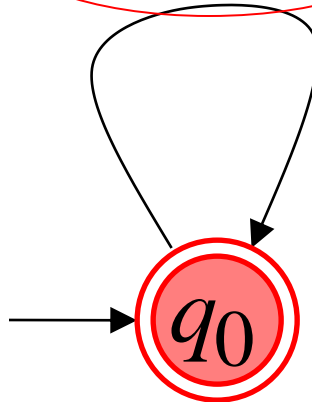
Input



$a, \lambda \rightarrow a$

$b, a \rightarrow \lambda$

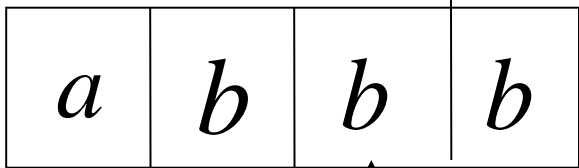
$b, \$ \rightarrow \lambda$



Stack

Time 4

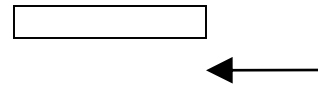
Input



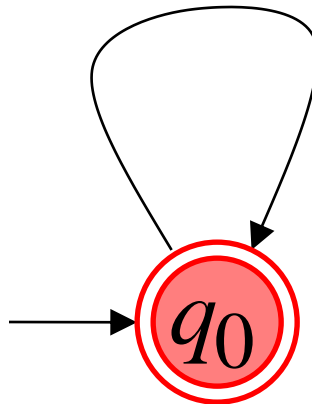
$a, \lambda \rightarrow a$

$b, a \rightarrow \lambda$

$b, \$ \rightarrow \lambda$

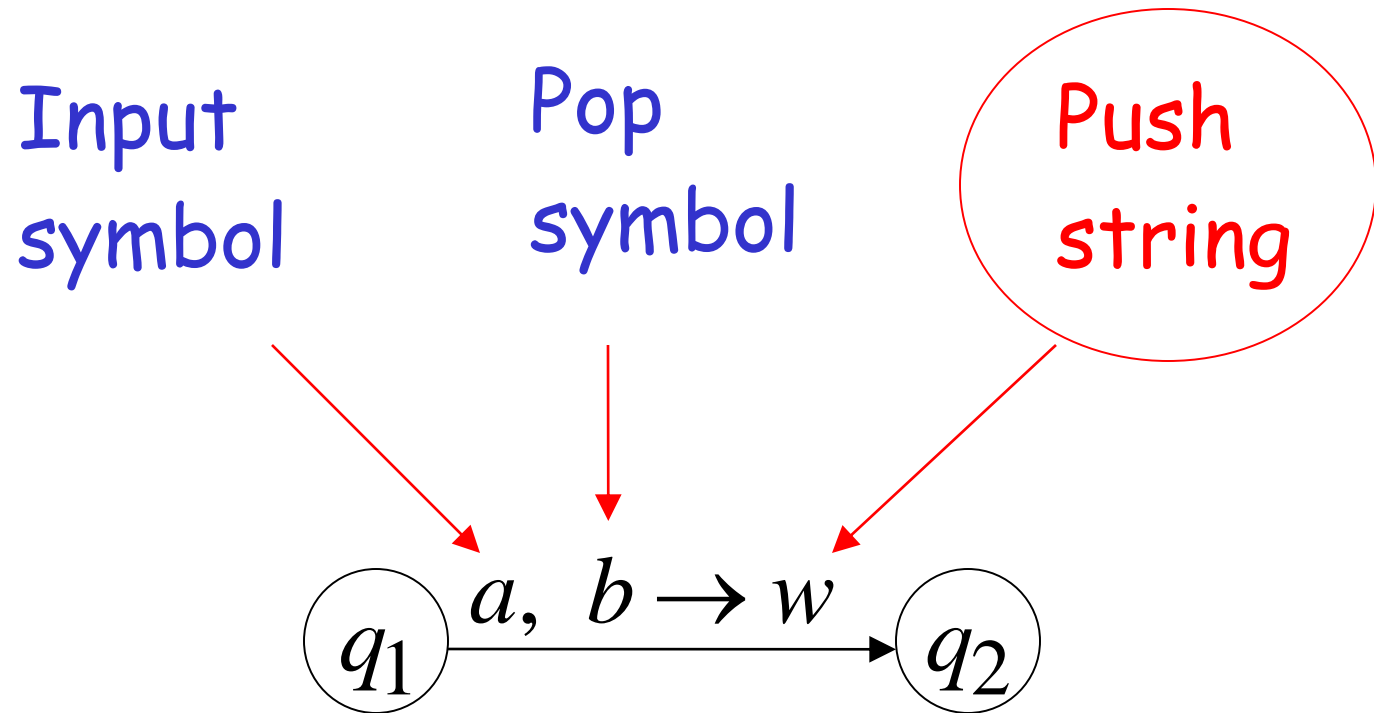


Stack

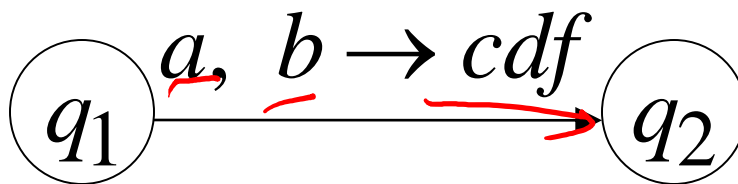


Halt and Reject

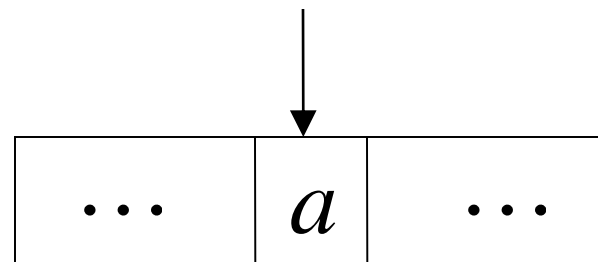
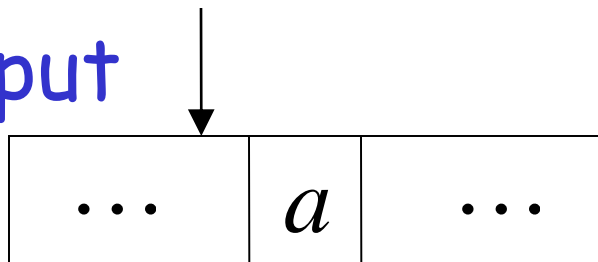
# Pushing Strings



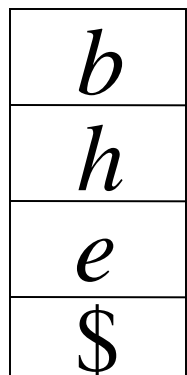
Example:



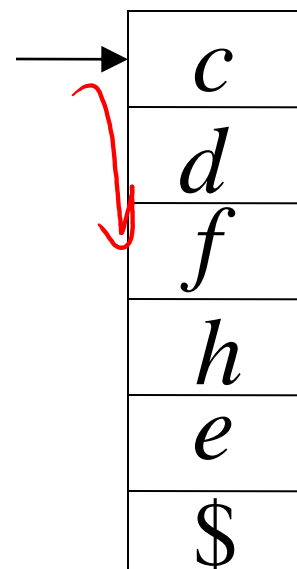
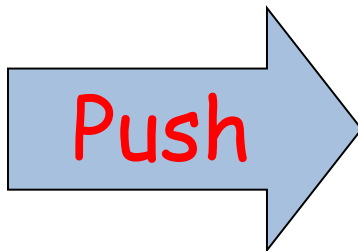
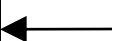
input



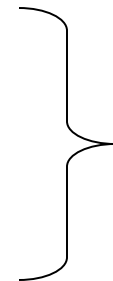
stack



top



pushed  
string



# Another NPDA example

NPDA  $M$

$$L(M) = \{w : n_a = n_b\}$$



# Another NPDA example

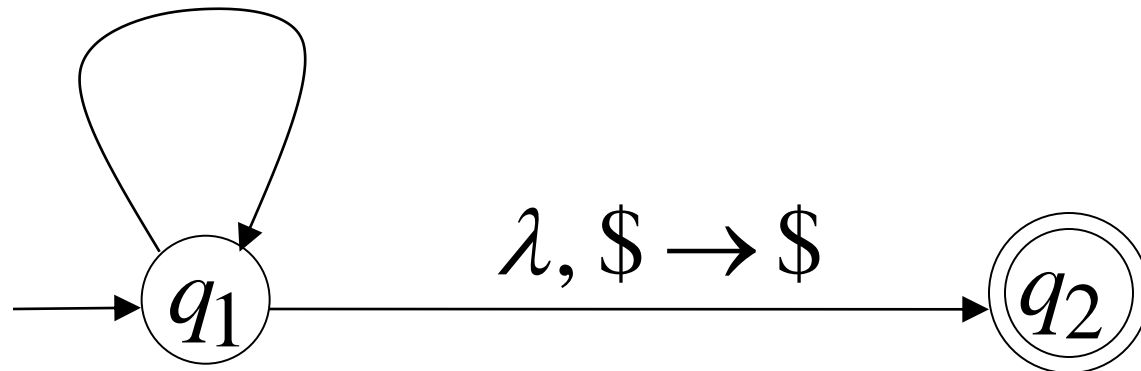
NPDA  $M$

$$L(M) = \{w : n_a = n_b\}$$

$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$

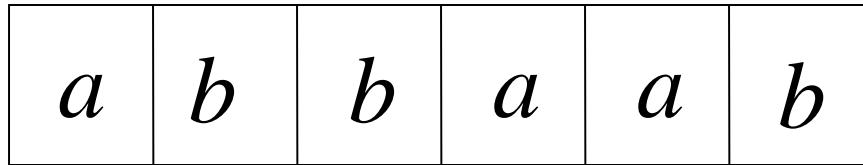
$a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$



# Execution Example:

Time 0

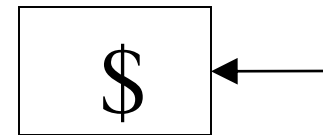
Input



$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$

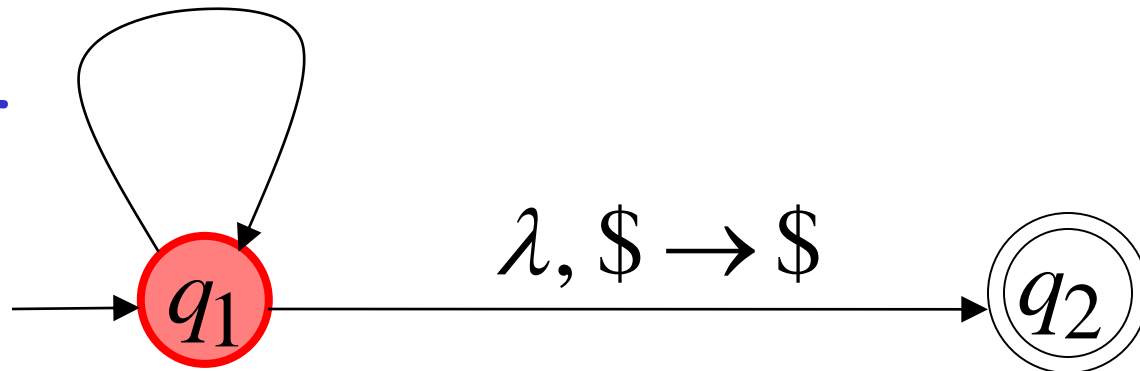
$a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$



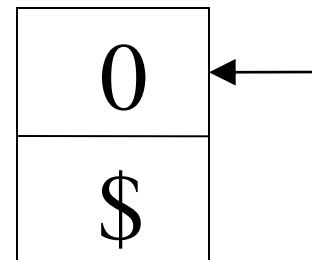
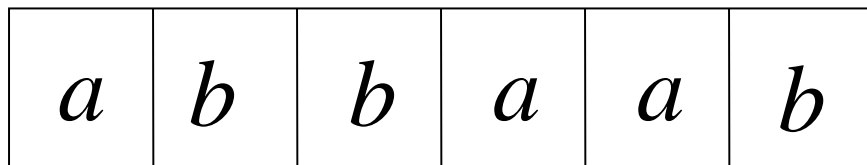
Stack

current  
state



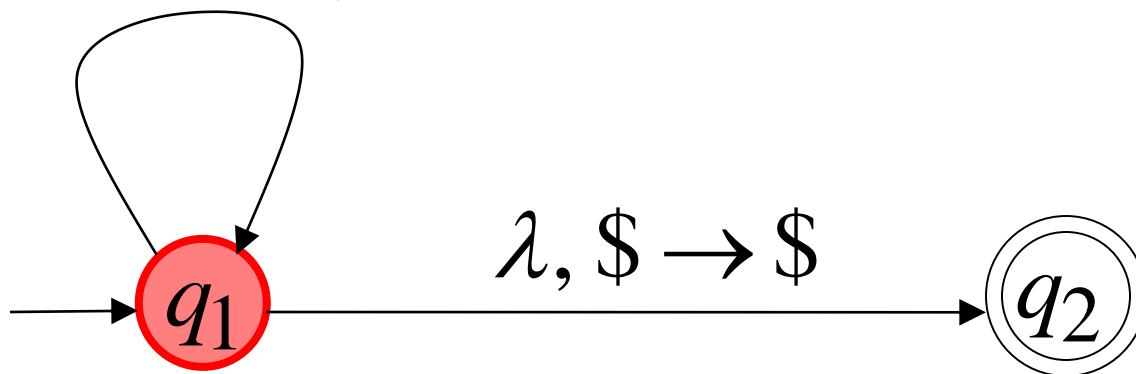
Time 1

Input



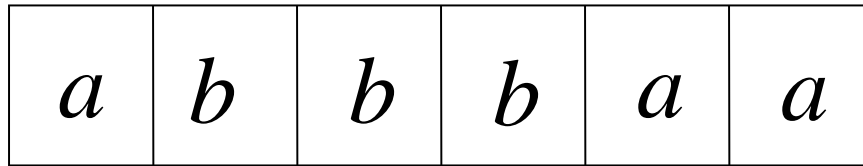
Stack

$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$   
 $a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$   
 $a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$



Time 3

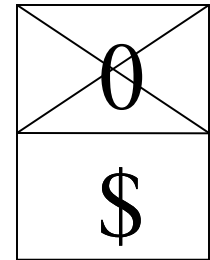
Input



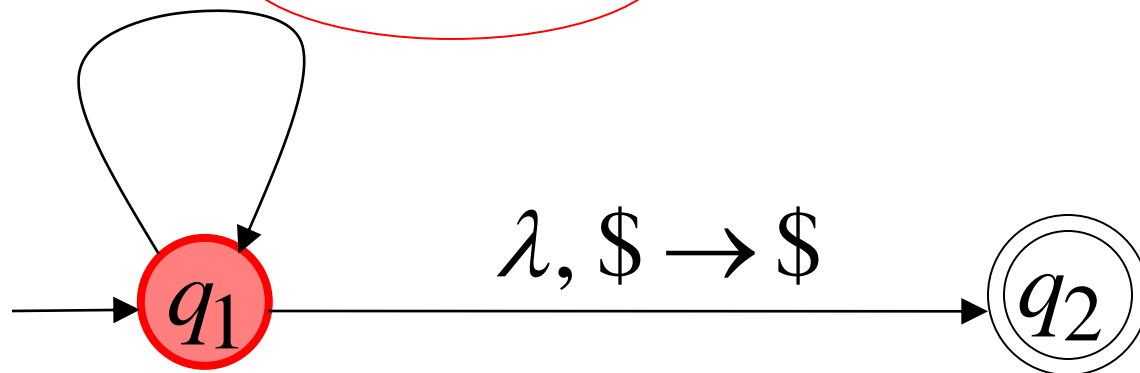
$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$

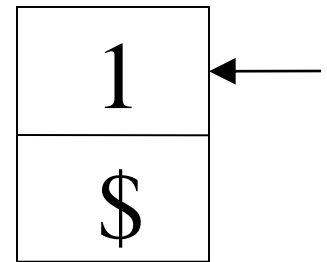
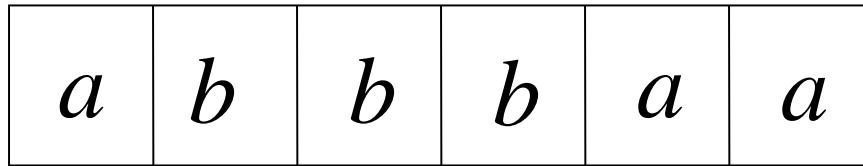


Stack



Time 4

Input

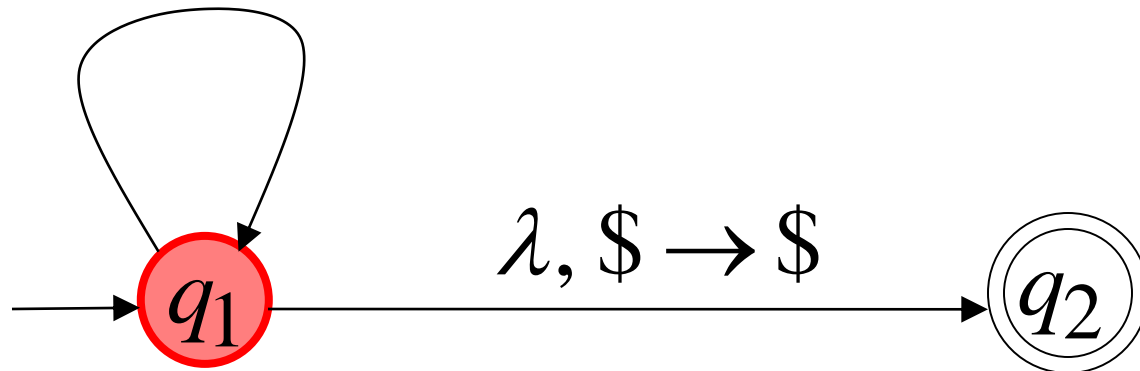


Stack

$a, \$ \rightarrow 0\$$       $b, \$ \rightarrow 1\$$

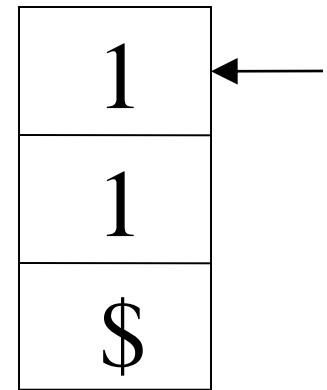
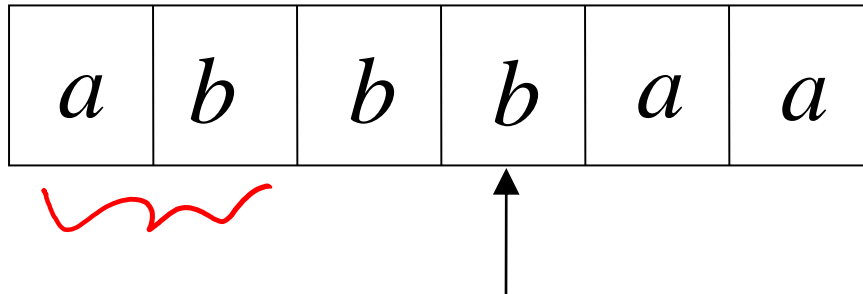
$a, 0 \rightarrow 00$       $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$       $b, 0 \rightarrow \lambda$



Time 5

Input

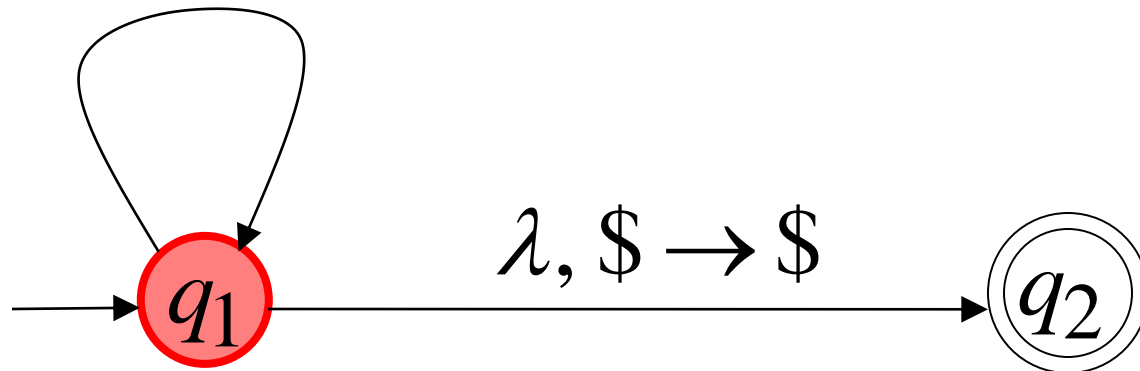


Stack

$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$

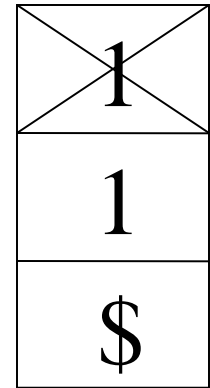
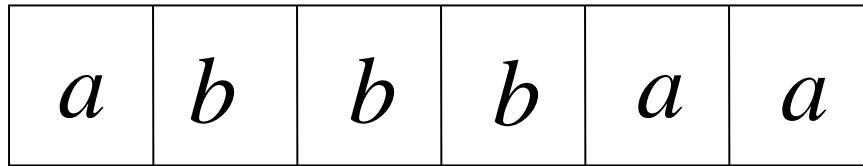
$a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$



Time 6

Input

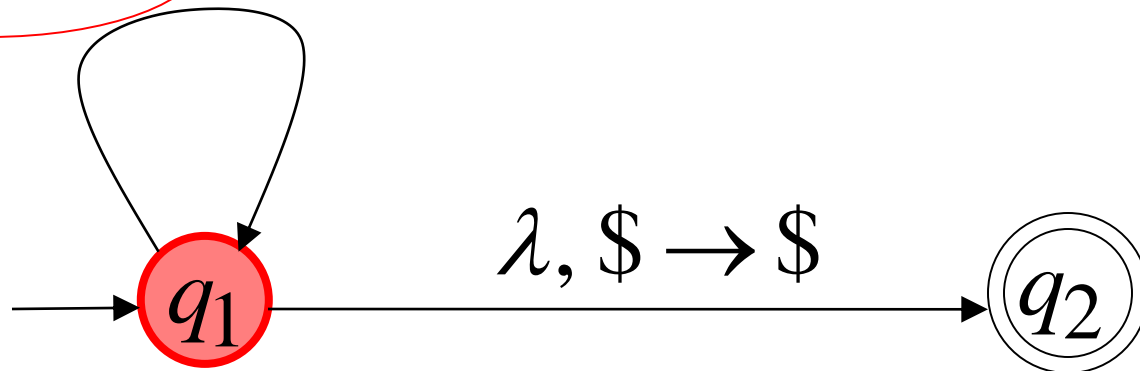


Stack

$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$

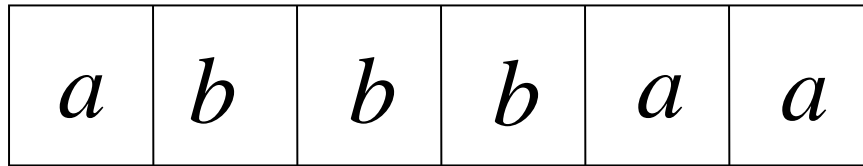
$a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$



Time 7

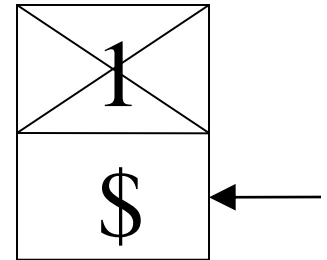
Input



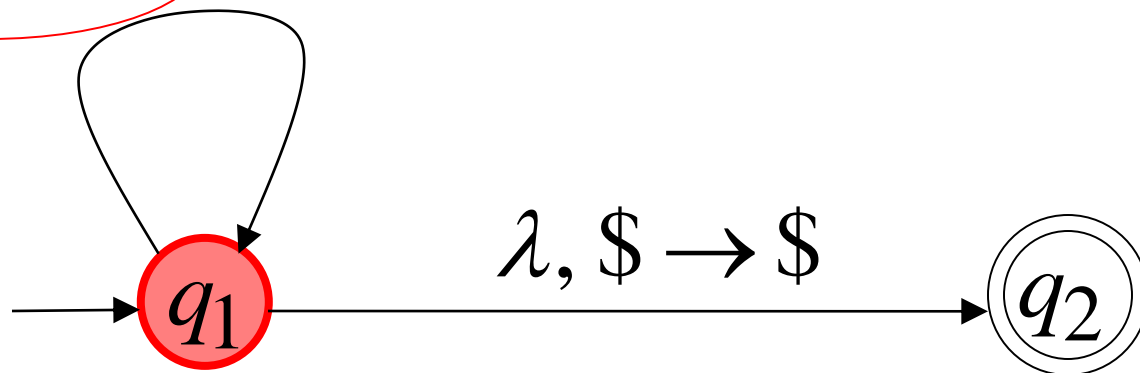
$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$



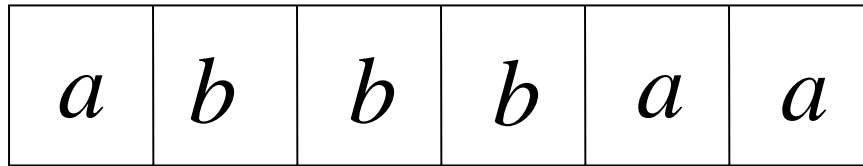
Stack





Time 8

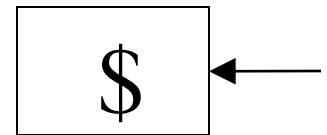
Input



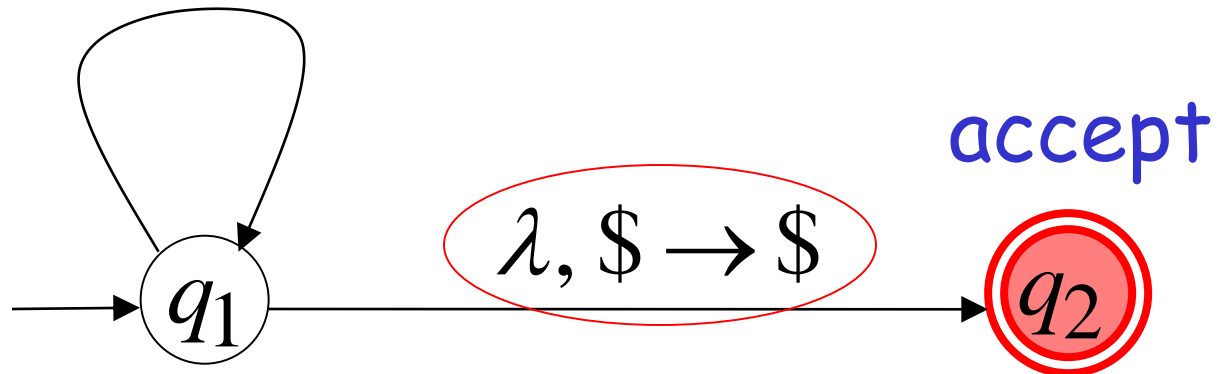
$a, \$ \rightarrow 0\$$        $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$        $b, 1 \rightarrow 11$

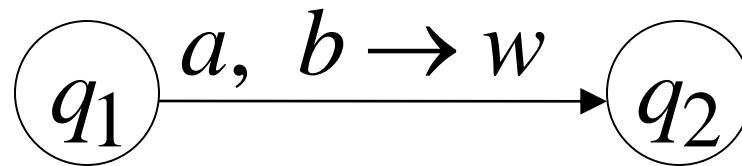
$a, 1 \rightarrow \lambda$        $b, 0 \rightarrow \lambda$



Stack

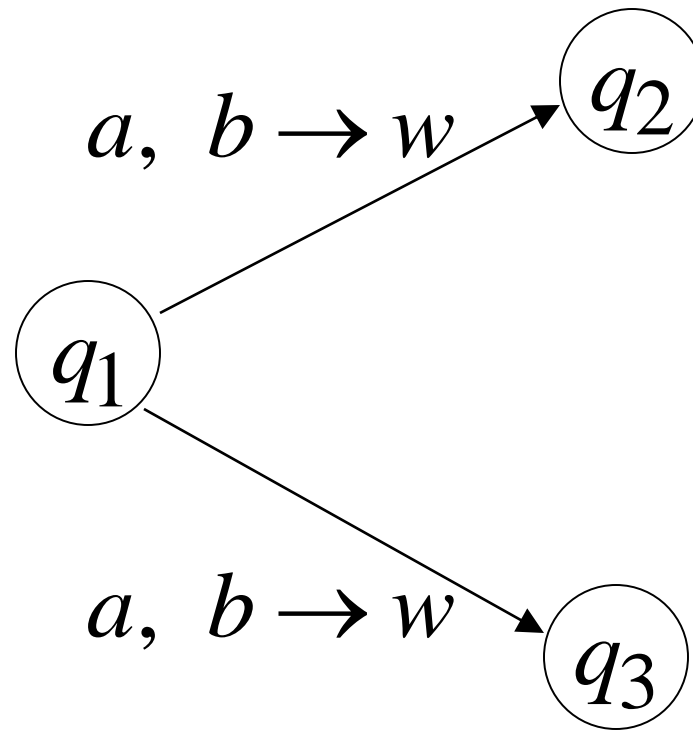


# Formalities for NPDAs



Transition function:

$$\delta(\underline{q_1}, \underline{a}, \underline{b}) = \{(\underline{q_2}, \underline{w})\}$$



Transition function:

$$\delta(q_1, a, b) = \{(q_2, w), (q_3, w)\}$$

# Formal Definition

## Non-Deterministic Pushdown Automaton NPDA

$$M = (Q, \Sigma, \Gamma, \delta, q_0, z, F)$$

States

Input  
alphabet

Stack

alphabet

Transition

Initial

Stack

start

symbol

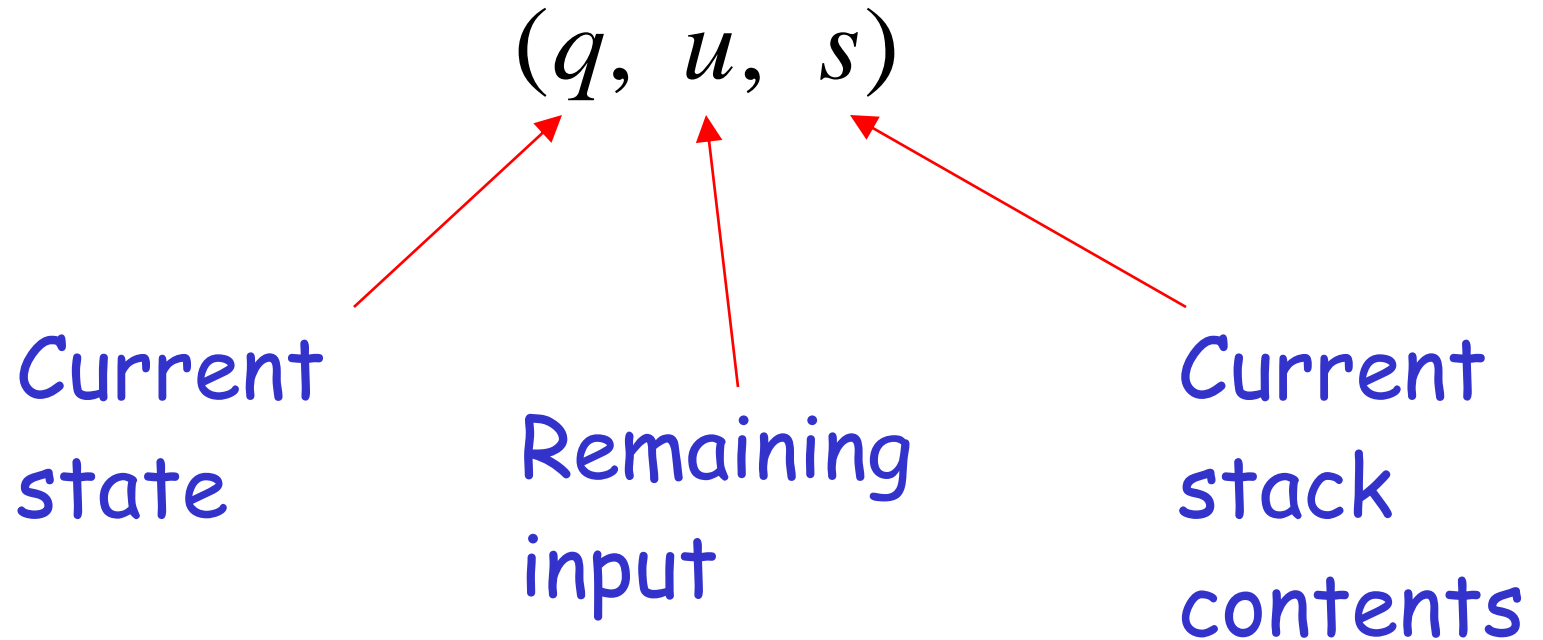
Final  
states

function

state

$\delta: Q \times (\Sigma \cup \{\lambda\}) \times \Gamma \rightarrow \text{finite subsets of } Q \times \Gamma^*$

# Instantaneous Description

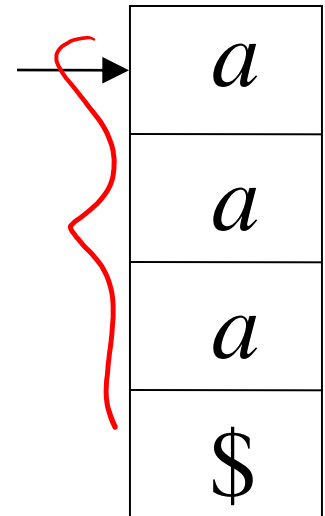
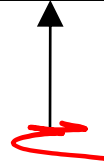
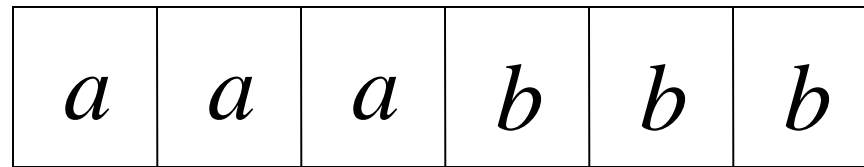


# Example: Instantaneous Description

$(q_1, \underline{b}\underline{b}\underline{b}, aaa\$)$

Time 4:

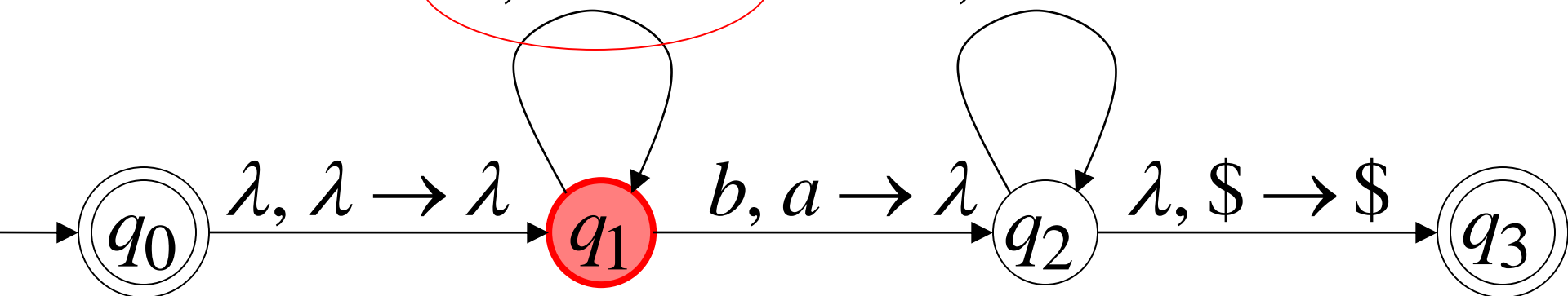
Input



Stack

$a, \lambda \rightarrow a$

$b, a \rightarrow \lambda$

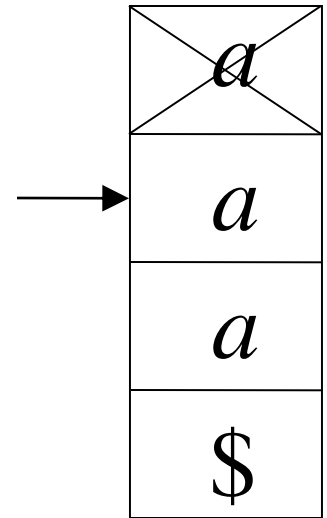
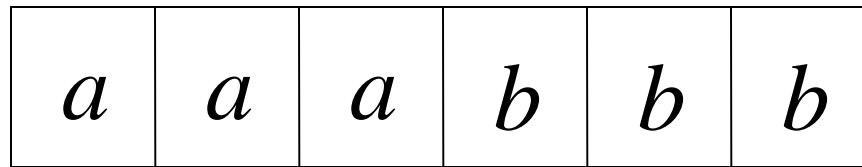


# Example: Instantaneous Description

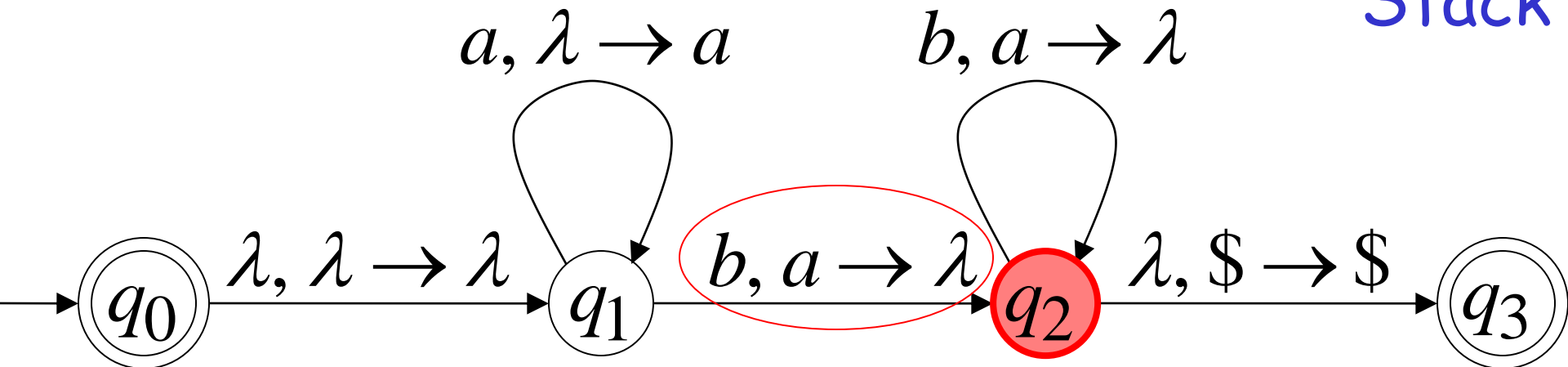
$(q_2, bb, aa\$)$

Time 5:

Input



Stack





We write:

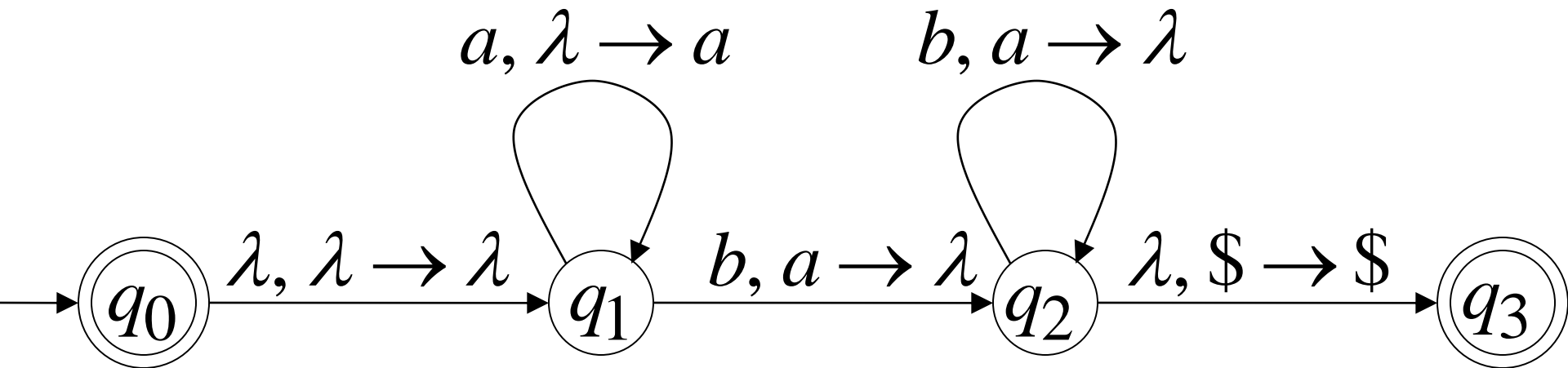
$$(q_1, bbb, aaa\$) \underset{=}{\succ} (q_2, bb, aa\$)$$

Time 4

Time 5

## A computation:

$$\begin{aligned} &(\underline{q_0}, \underline{aaabbbb}, \underline{\lambda}) \succ (q_1, aaabbbb, \$) \succ \\ &(q_1, aabbbb, a\$) \succ (q_1, abbbb, aa\$) \succ (q_1, bbb, aaa\$) \succ \\ &(q_2, bb, aa\$) \succ (q_2, b, a\$) \succ (q_2, \lambda, \$) \succ (\underline{q_3}, \lambda, \$) \end{aligned}$$



$$\begin{aligned}
 &(q_0, aaabbb, \$) \succ (q_1, aaabbb, \$) \succ \\
 &(q_1, aabbbb, a\$) \succ (q_1, abbbb, aa\$) \succ (q_1, bbb, aaa\$) \succ \\
 &(q_2, bb, aa\$) \succ (q_2, b, a\$) \succ (q_2, \lambda, \$) \succ (q_3, \lambda, \$)
 \end{aligned}$$

For convenience we write:

$$\underbrace{(q_0, aaabbb, \$)} \overset{*}{\succ} \underbrace{(q_3, \lambda, \$)}$$

# Formal Definition

Language  $L(M)$  of NPDA  $M$

$$L(M) = \{w : (q_0, w, s) \stackrel{*}{\succ} (q_f, \lambda, s')\}$$

Initial state

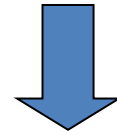


Final state



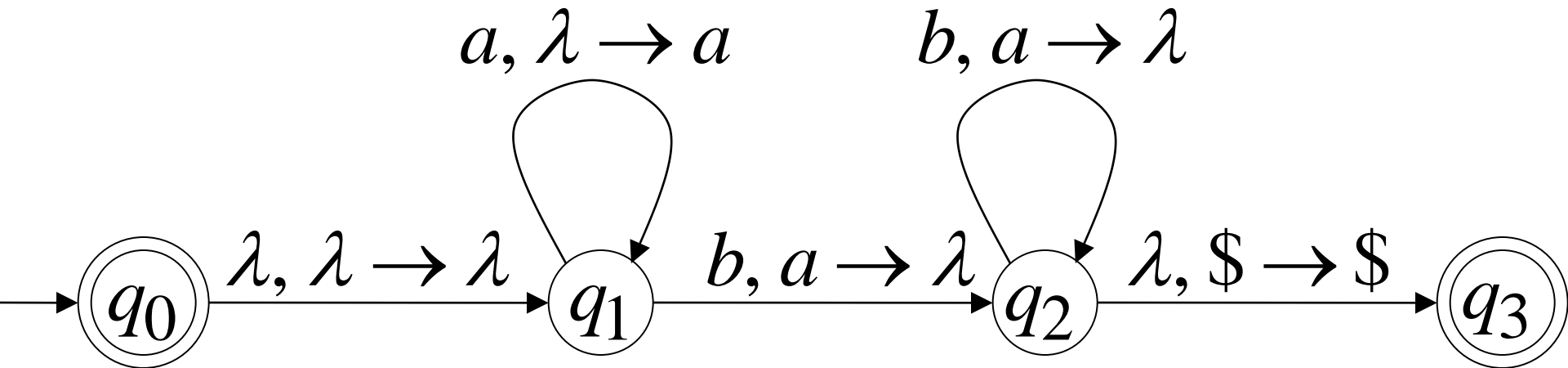
Example:

$$(q_0, aaabbb, \$) \stackrel{*}{\succ} (q_3, \lambda, \$)$$



$$aaabbb \in L(M)$$

NPDA  $M$ :

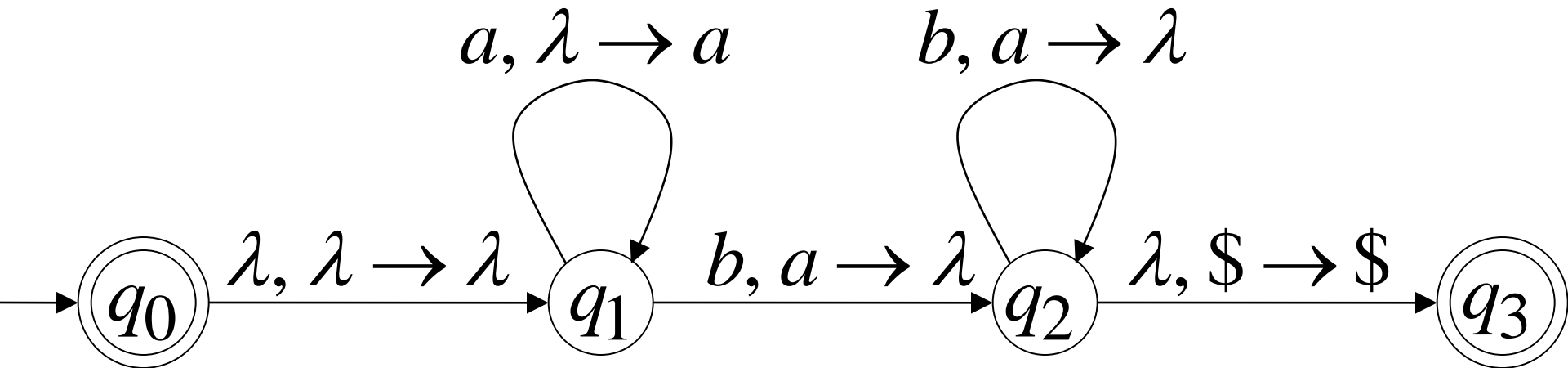


$$(q_0, a^n b^n, \$) \stackrel{*}{\succ} (q_3, \lambda, \$)$$



$$a^n b^n \in L(M)$$

NPDA  $M$ :



Therefore:  $L(M) = \{a^n b^n : n \geq 0\}$

NPDA  $M$ :

