- Boolean Algebra:
- When no. of variables are less i.e. 1,2,3.
- It is preferred when output is 0 or 1.

- K-Map:
- When no. of variables are less i.e. 2,3,4 (upto 5 variables).
- It is preferred when output is 0, 1 or X.
- Tabulation Method:
- It is preferred when no. of variables are greater than 5.

- Operator order in Boolean expression:
- Parentheses
- NOT
- AND
- OR

• Example:

$$x + (y + z) = (x + y) + z$$

• Remember the operators order.

- Basic theorem based on Boolean postulates:
- NOT:
- (X')'= X
- AND:
- $\bullet$  X.X = X
- X.1 = X
- X.0 = 0
- X.X' = 0

- OR:
- $\bullet X+X=X$
- X+1 = X
- X+0 = X
- X+X'=1

- Problem:
- XY+XY'

- Solution:
- XY+XY'X(Y+Y')= X, [Since: Y+Y'= 1]

- Theorem: Identity
- 0(+)
- **1**(.)
- Theorem: Commutative
- $\bullet$  X+Y = Y+X
- $\bullet$  X.Y = Y.X

- Theorem: Distributive
- X(Y+Z)=(X+Y)+(Y+Z)
- $X + (Y.Z) = (X+Y) \cdot (Y+Z)$
- Theorem: Transposition
- $(X+Y) \cdot (X+Z) = (X+YZ)$

- Theorem: DeMorgan's
- (X+Y)'= X'. Y'
- (X.Y)' = X' + Y'

- Problem: Minimize the following Boolean function:
- $\bullet$  XY' + XYZ' + XY'Z'W

#### Solution:

• 
$$XY' + XYZ' + XY'Z'W = XYZ' + XY' (1 + Z'W)$$

• 
$$XYZ' + XY'$$
, [Since:  $1 + A = 1$ ]

• 
$$X (Y' + Z')$$
, [Since:  $Y + YZ' = Y' + Z'$ ]

Advantages of minimization:

No. of logic Gates: Less

Speed: High

Power dissipation: Lower

Complexity of logic circuits: Less

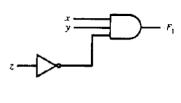
• Fan-in: Low

Cost: Low

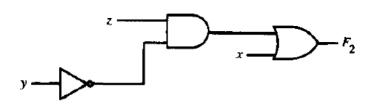
- Boolean Function:
- An algebraic expression formed by binary variables.
- An operators AND and OR, unary operator NOT.
- Parentheses, and equal sign.
- Example.
- F1 = xyz'
- variables: x, y, z
- Gates: 1 NOT gate; 1 AND Gate.
- Output: F1 can be either 0 or 1, for a given input.

F1 = xyz' when output is high / low?

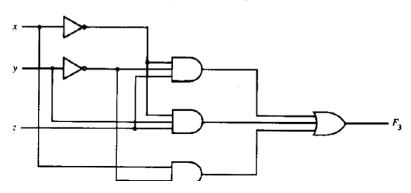
_	X	У	z	F <sub>1</sub>
	0	0	0	0
	0	0	1	0
	0	1	0	0
	0	1	1	0
	1	0	0	0
	1	0	1	0
	1	1	0	1
	1	1	1	0



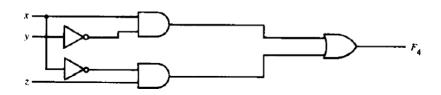
# F2 = x + y'z when output is high / low?



# F3 = x'y'z + x'yz + xy'when output is high / low?



# F4 = xy' + x'zwhen output is high / low?



- Problem:
- XY + XY'Z + XY'Z', find the minimum number of NAND Gates.

- Solution:
- XY + XY'Z + XY'Z'
- XY + XY' (Z + Z')
- XY + XY'
- X (Y + Y')
- X

#### Zero NAND Gate.

- Problem:
- (X+Y+Z)(X+Y'+Z)(X+Y+Z')

- Solution:
- Take A = X+Y
- (A+Z)(X+Y'+Z)(A+Z')
- (A+ZZ')(X+Y'+Z)
- A(Y + Y' + Z)
- (X+Y)(X+Y'+Z)
- X + YZ

- Problem:
- (X+Y)(X+Y')(X'+Y)(X'+Y')

- Solution:
- (X+YY')(X'+YY')
- X.X'
- 0

- Redundant term: Not needed
- Example:
- XY+X'Z+YZ
- YZ is redundant term
- Solution:
- XY+X'Z+YZ(X+X')
- XY+X'Z+XYZ+X'YZ)
- XY(1+Z)+X'Z(1+Y)
- XY+X'Z)

- Redundancy Theorem:
- Three variables.
- Each variable comes twice.
- Out of three variables one variable is complement.
- The term which is complement is taken.
- Example:
- AB+A'C+BC
- AB+A'C

- Complement of a Boolean function:
- Use of DeMorgan's theorem
- Use of Duality

## Complement by DeMorgan's theorem

Example: 
$$F1 = x'yz' + x'y'z$$
  
Find  $F1'$ 

$$F1' = (x'yz' + x'y'z)'$$
= (x'yz')' (x'y'z)'
= (x + y' + z) (x + y + z')

Check: truth-table

- Problem:
- Find F2' for F2 = x (y'z' + yz)

- Solution:
- F2' = (x (y'z' + yz))'
- F2' = x' + (y'z' + yz)'
- F2' = x' + (y'z')' (yz)'
- F2' = x' + (y + z) (y' + z')

## Complement by Duality

Each AND => OR
 Dual of Boolean function

 Each OR => AND

Complement each literal.

## Again do F1' using Duality

$$F1 = x'yz' + x'y'z$$

$$F1' = (x' + y + z') (x' + y' + z)$$
 (dual of F1)  
=  $(x + y' + z) (x + y + z')$  (complementing each literal)

Look back to the previous answer of F1'.

## Now let us do F2' by Duality

$$F2 = x (y'z' + yz)$$

$$F2' = x + (y' + z') (y + z)$$
 (dual of F2)  
=  $x' + (y + z) (y' + z')$  (complement of literals)

#### • Minterms or Standard products:

- Two binary variables expression can have following four AND operations or product terms:
- x.y
- x'.y
- x.y'
- x'.y'

#### Maxterms or Standard sums:

- Two binary variables expression can have following four OR operations or product terms:
- x+y
- x'+y
- x+y'
- x'+y'

x .			М	interms	Maxterms	
	У	Z	Term	Designation	Term	Designation
0	0	0	x'y'z'	$m_0$	x + y + z	$M_0$
0	0	1	x'y'z	$m_1$	x + y + z'	$M_1$
0	1	0	x'yz'	$m_2$	x + y' + z	$M_2$
ö	1	1	x'yz	$m_3$	x + y' + z'	$M_3$
1	0	0	xy'z'	$m_{\perp}$	x' + y + z	$M_4$
1	0	l	xy'z	$m_5$	x' + y + z'	$M_5$
1	1	0	xyz'	$m_6$	x' + y' + z	$M_6$
1	ì	1	xvz	$m_7$	x' + y' + z'	$M_7$

### • Expressing a Boolean function from its Truth-table:

У	Z	Function $f_1$	Function f <sub>2</sub>
0	0	0	0
0	1	1	0
1	0	0	0
l	1	0	1
0	0	1	0
0	1	0	1
1	0	0	1
1	1	1	1
	y 0 0 1 1 0 0 1 1 1	y z 0 0 0 1 1 0 1 1 0 0 0 1 1 1 1 1 1 1 1 1	y     z     Function f1       0     0     0       0     1     1       1     0     0       1     1     0       0     0     1       0     1     0       1     0     0       1     1     0       1     1     1

$$F1 = x'y'z + xy'z' + xyz = m1 + m4 + m7$$

$$F2 = x'yz + xy'z + xyz' + xyz = m3 + m5 + m6 + m7$$