

- **Boolean Algebra:**
- **When no. of variables are less i.e. 1,2,3.**
- **It is preferred when output is 0 or 1.**

- **K-Map:**
- **When no. of variables are less i.e. 2,3,4 (upto 5 variables).**
- **It is preferred when output is 0, 1 or X.**
- **Tabulation Method:**
- **It is preferred when no. of variables are greater than 5.**

- **Operator order in Boolean expression:**
- **Parentheses**
- **NOT**
- **AND**
- **OR**

- **Example:**
- $x + (y + z) = (x + y) + z$
- $x \cdot (y \cdot z) = (x \cdot y) \cdot z$
- **Remember the operators order.**

- **Basic theorem based on Boolean postulates:**

- **NOT:**

- $(X')' = X$

- **AND:**

- $X.X = X$

- $X.1 = X$

- $X.0 = 0$

- $X.X' = 0$

- **OR:**
- **$X+X = X$**
- **$X+1 = X$**
- **$X+0 = X$**
- **$X+X' = 1$**

- $X + (X \cdot Y) = X$
- $X \cdot (X + Y) = X$

- **Problem:**
- **$XY + XY'$**

- **Solution:**
- **$XY + XY'$**
 $X(Y + Y') = X$, [Since: $Y + Y' = 1$]

- **Theorem: Identity**

- $0(+)$

- $1(.)$

- **Theorem: Commutative**

- $X+Y = Y+X$

- $X.Y = Y.X$

- **Theorem: Distributive**

- $X(Y+Z) = (X+Y) + (Y+Z)$

- $X + (Y.Z) = (X+Y) . (Y+Z)$

- **Theorem: Transposition**

- $(X+Y) . (X+Z) = (X+YZ)$

- **Theorem: DeMorgan's**
- $(X+Y)' = X' \cdot Y'$
- $(X.Y)' = X' + Y'$

- **Problem: Minimize the following Boolean function:**
- $XY' + XYZ' + XY'Z'W$

- **Solution:**

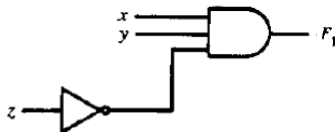
- $XY' + XYZ' + XY'Z'W = XYZ' + XY' (1 + Z'W)$
- $XYZ' + XY'$, [Since: $1 + A = 1$]
- $X (Y + YZ')$
- $X (Y' + Z')$, [Since: $Y + YZ' = Y' + Z'$]
- $XY' + XZ'$

- **Advantages of minimization:**
- **No. of logic Gates: Less**
- **Speed: High**
- **Power dissipation: Lower**
- **Complexity of logic circuits: Less**
- **Fan-in: Low**
- **Cost: Low**

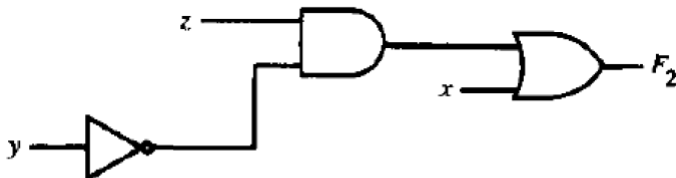
- **Boolean Function:**
- **An algebraic expression formed by binary variables.**
- **An operators AND and OR, unary operator NOT.**
- **Parentheses, and equal sign.**
- **Example.**
- $F1 = xyz'$
- variables: x, y, z
- Gates: 1 NOT gate; 1 AND Gate.
- Output: $F1$ can be either 0 or 1, for a given input.

$F_1 = xyz'$
when output is high / low ?

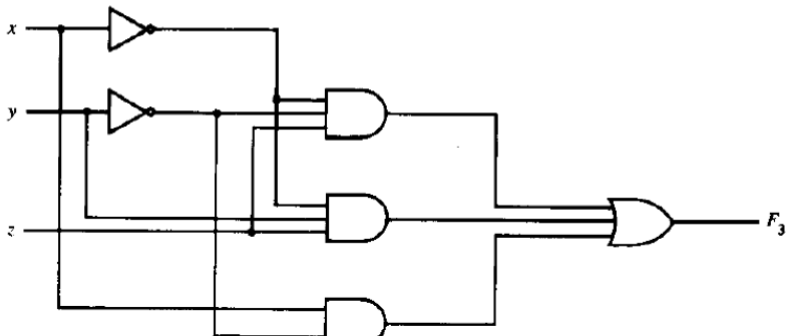
| x | y | z | F_1 |
|-----|-----|-----|-------|
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |



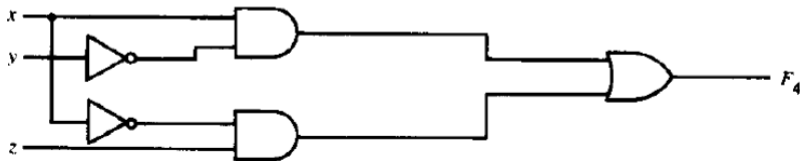
$F_2 = x + y'z$
when output is high / low ?



$F_3 = x'y'z + x'yz + xy'$
when output is high / low ?



$F_4 = xy' + x'z$
when output is high / low ?



- **Problem:**
- **$XY + XY'Z + XY'Z'$, find the minimum number of NAND Gates.**

- **Solution:**

- $XY + XY'Z + XY'Z'$

- $XY + XY' (Z + Z')$

- $XY + XY'$

- $X (Y + Y')$

- X

Zero NAND Gate.

- **Problem:**
- $(X+Y+Z)(X+Y'+Z)(X+Y+Z')$

- **Solution:**

- **Take $A = X+Y$**

- $(A+Z)(X+Y'+Z)(A+Z')$

- $(A+ZZ')(X+Y'+Z)$

- $A(Y + Y' + Z)$

- $(X+Y)(X+Y'+Z)$

- $X + YZ$

- **Problem:**
- $(X+Y)(X+Y')(X'+Y)(X'+Y')$

- **Solution:**
- $(X + YY')(X' + YY')$
- $X.X'$
- 0

- **Redundant term: Not needed**

- **Example:**

- $XY + X'Z + YZ$
- YZ is redundant term

- **Solution:**

- $XY + X'Z + YZ(X + X')$
- $XY + X'Z + XYZ + X'YZ$
- $XY(1 + Z) + X'Z(1 + Y)$
- $XY + X'Z$

- **Redundancy Theorem:**
- **Three variables.**
- **Each variable comes twice.**
- **Out of three variables one variable is complement.**
- **The term which is complement is taken.**
- **Example:**
 - $AB + A'C + BC$
 - $AB + A'C$

- **Complement of a Boolean function:**
- **Use of DeMorgan's theorem**
- **Use of Duality**

Complement by DeMorgan's theorem

Example: $F1 = x'yz' + x'y'z$

Find $F1'$

$$\begin{aligned} F1' &= (x'yz' + x'y'z)' \\ &= (x'yz')' (x'y'z)' \\ &= (x + y' + z) (x + y + z') \end{aligned}$$

Check: truth-table

- **Problem:**
- **Find $F2'$ for $F2 = x (y'z' + yz)$**

- **Solution:**

- $F2' = (x (y'z' + yz))'$

- $F2' = x' + (y'z' + yz)'$

- $F2' = x' + (y'z')' (yz)'$

- $F2' = x' + (y + z) (y' + z')$

Complement by Duality

- Each AND \Rightarrow OR
 - Each OR \Rightarrow AND
- Dual of Boolean function
- Complement each literal.

Again do $F1'$ using Duality

$$F1 = x'yz' + x'y'z$$

$$\begin{aligned} F1' &= (x' + y + z') (x' + y' + z) \quad (\text{dual of } F1) \\ &= (x + y' + z) (x + y + z') \quad (\text{complementing each literal}) \end{aligned}$$

Look back to the previous answer of $F1'$.

Now let us do F2' by Duality

$$F2 = x (y'z' + yz)$$

$$\begin{aligned} F2' &= x + (y' + z') (y + z) && \text{(dual of F2)} \\ &= x' + (y + z) (y' + z') && \text{(complement of literals)} \end{aligned}$$

- **Minterms or Standard products:**
- Two binary variables expression can have following four AND operations or product terms:
 - $x.y$
 - $x'.y$
 - $x.y'$
 - $x'.y'$

- **Maxterms or Standard sums:**
- Two binary variables expression can have following four OR operations or product terms:
 - $x+y$
 - $x'+y$
 - $x+y'$
 - $x'+y'$

| | | | Minterms | | Maxterms | |
|-----|-----|-----|----------|-------------|----------------|-------------|
| x | y | z | Term | Designation | Term | Designation |
| 0 | 0 | 0 | $x'y'z'$ | m_0 | $x + y + z$ | M_0 |
| 0 | 0 | 1 | $x'y'z$ | m_1 | $x + y + z'$ | M_1 |
| 0 | 1 | 0 | $x'yz'$ | m_2 | $x + y' + z$ | M_2 |
| 0 | 1 | 1 | $x'yz$ | m_3 | $x + y' + z'$ | M_3 |
| 1 | 0 | 0 | $xy'z'$ | m_4 | $x' + y + z$ | M_4 |
| 1 | 0 | 1 | $xy'z$ | m_5 | $x' + y + z'$ | M_5 |
| 1 | 1 | 0 | xyz' | m_6 | $x' + y' + z$ | M_6 |
| 1 | 1 | 1 | xyz | m_7 | $x' + y' + z'$ | M_7 |

- Expressing a Boolean function from its Truth-table:

| x | y | z | Function f_1 | Function f_2 |
|-----|-----|-----|----------------|----------------|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |

$$F1 = x'y'z + xy'z' + xyz = m1 + m4 + m7$$

$$F2 = x'yz + xy'z + xyz' + xyz = m3 + m5 + m6 + m7$$