

## RE

- Def. i)  $\Sigma, \lambda, \emptyset$  are RE  
 ii) If  $R_1$  &  $R_2$  then union of  $R_1$  &  $R_2$  is a R.G.  

$$R_1 + R_2$$
  
 iii) Concatenation of  $R_1$  &  $R_2$  i.e.  $R_1 R_2$   
 iv) Iteration on  $R$  i.e.  $R^*$   
 v) If  $R$  is a RE then  $(R)$  is also RE.  
 vi) RE w.r.t.  $\Sigma$  which can be obtained from above rules is also RE.

Ex. 1)  $L_1 = \text{set of all strings on } \Sigma = \{0, 1\}$   
 ending in 00.

$$R_1 = (0+1)^* 00$$

2)  $L_2 = \Sigma = \{0, 1\}$  starting with 0 &  
 ending with 1.  

$$R_2 = 0(0+1)^* L$$

3)  $L_3 = \{\lambda, 11, 1111, 11111, \dots\}$

$$R_3 = (11)^*$$

## Identities of REs

- 1)  $\emptyset + R = R$
- 2)  $\emptyset \cdot R = R \cdot \emptyset = \emptyset$
- 3)  $\lambda R = R \lambda = R$
- 4)  $\lambda^* = \lambda \text{ & } \emptyset^* = \lambda$
- 5)  $R + R = R$
- 6)  $R^* R^* = R^*$
- 7)  $R R^* = R^* R$
- 8)  $(R^*)^* = R^*$
- 9)  $\underbrace{\lambda + RR^*}_\lambda = R^* = \lambda + R^* R$
- 10)  $(RQ)^* R = R(QR)^*$
- 11)  $(P+Q)^* = (P^* Q^*)^* = (P^* + Q^*)^*$
- 12)  $(P+Q)R = PR + QR \text{ & }$   
 $R(P+Q) = RP + RQ$

## Arden's Theorem

Let  $P$  &  $Q$  be RE over  $\Sigma$ . If  $P$  does not contain  $\lambda$  then the following eq? in  $R$

viz.  $R = Q + RP$  has unique sol?

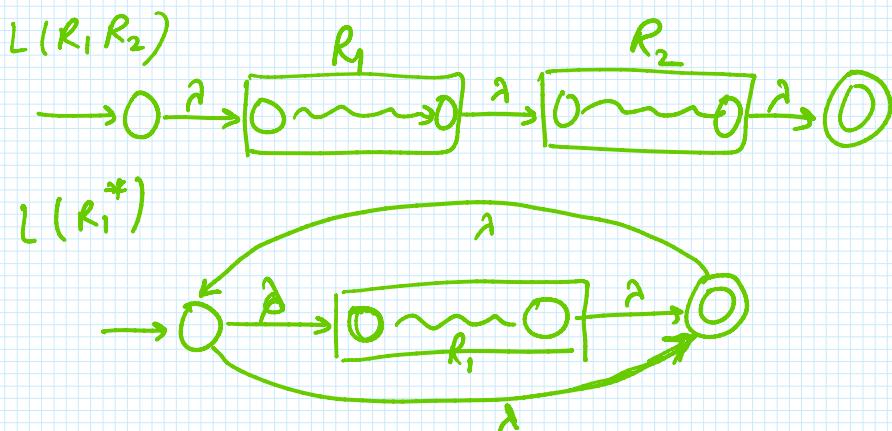
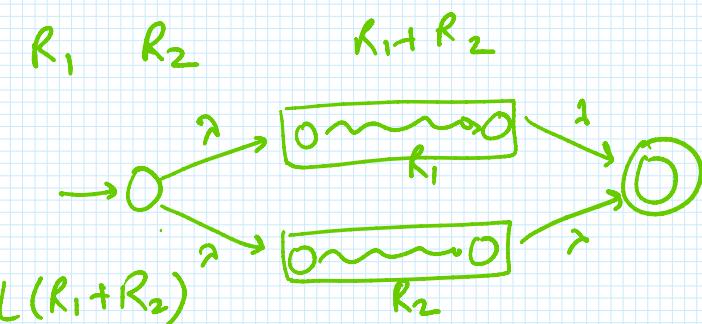
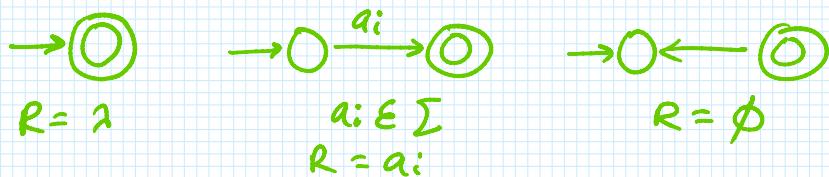
given by  $\underline{\underline{R = Q P^*}}$

Ex. Prove that

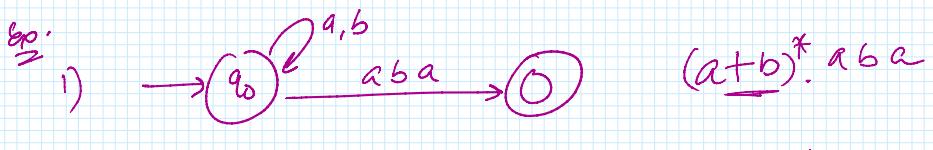
$$(1+00^*1) + (1+00^*1)(0+10^*1)^*(0+10^*1) \leftarrow \\ = 0^*1(0+10^*1)^*$$

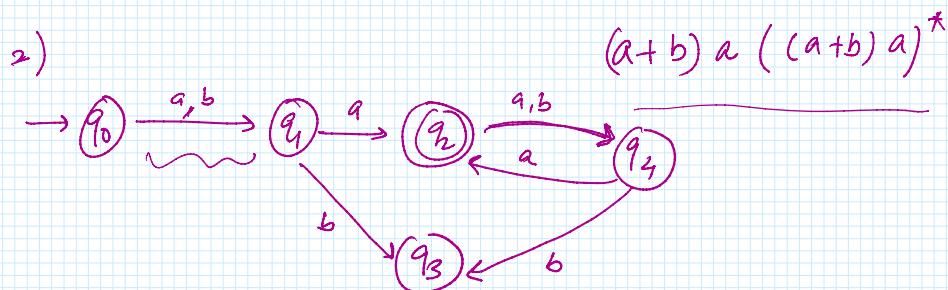
$$\begin{aligned}
 & (1+00^*) + (1+00^*)0(0+10^*)^* \\
 &= 0^* 1 (0+10^*)^* \\
 \stackrel{\text{LHS}}{=} & \frac{(1+00^*)}{R} \left[ 1 + \frac{(0+10^*)^*}{P^*} \frac{(0+10^*)}{P} \right] \\
 \equiv & (\lambda+00^*) 1 P^* \\
 \equiv & 0^* 1 ( \quad ) = R \stackrel{\text{RHS}}{=} \\
 & \text{FA } \& \text{ RE }
 \end{aligned}$$

Theorem: Every RE  $R$  can be recognised by a transition system and  $\exists$  a path from initial state to final state with path value  $\omega$ .



Theorem: Any set  $L$  accepted by a FA  $M$  is represented by a RE.



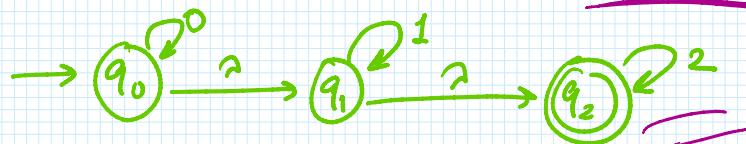


3)  $b a a + a b + a b b$  5)  $(a+b)^+$

4)  $(a+b)^* aa + (a+b)^*$

Transition system having null moves

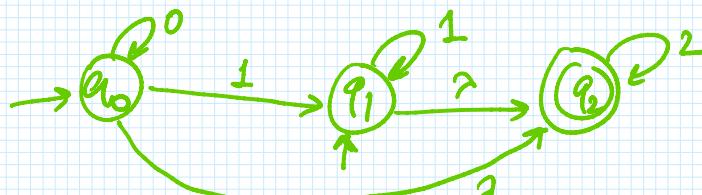
$\hat{\alpha}$ -move from  $v_1$  to  $v_2$   $v_1 \xrightarrow{\hat{\alpha}} v_2$



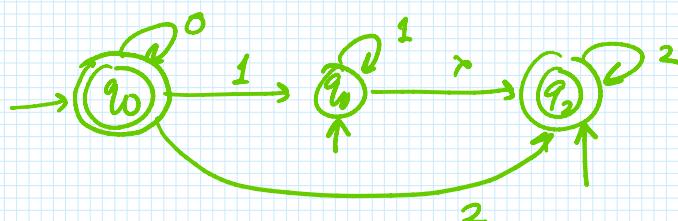
Steps

- Find all edges starting from  $v_2$
- Duplicate all these edges starting from  $v_1$  without changing the edge labels.
- If  $v_1$  is an initial state, make  $v_2$  also as initial state
- If  $v_2$  is final state, then make  $v_1$  as final state.

(i)  $\hat{\alpha}$ -move  $q_0 \xrightarrow{\hat{\alpha}} q_1$

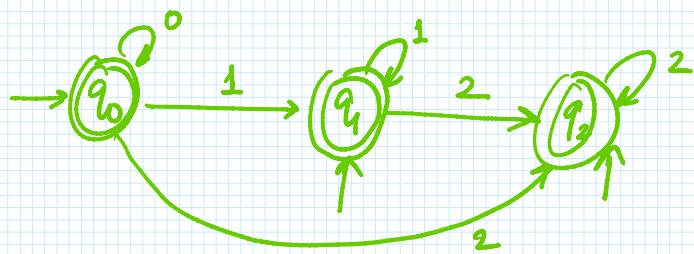


(ii)  $\hat{\alpha}$ -move  $q_0 \xrightarrow{\hat{\alpha}} q_2$   
 $v_1$        $v_2$



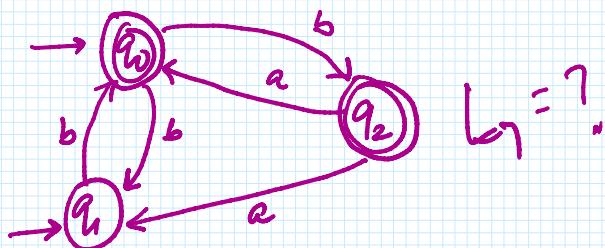
(iii)  $\hat{\alpha}$ -move  $q_1 \xrightarrow{\hat{\alpha}} q_2$

(ii) 2 move  $q_1 \xrightarrow{2} q_2$

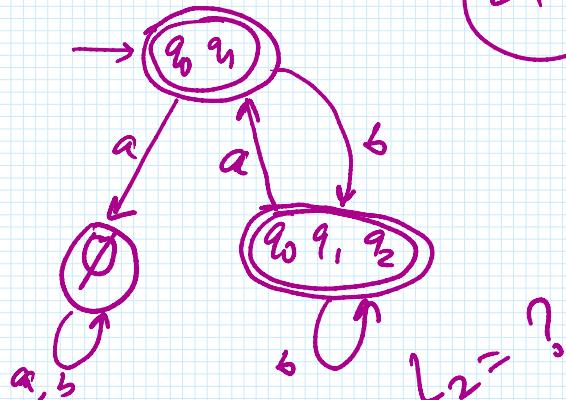


Conversion of non-deterministic system  
to deterministic system

<u>q<sub>0</sub></u>	a	b	q <sub>1</sub> , q <sub>2</sub>
q <sub>1</sub>			q <sub>0</sub>
q <sub>2</sub>			q <sub>0</sub> , q <sub>1</sub>



<u>q<sub>0</sub>, q<sub>1</sub></u>	a	b	q <sub>0</sub> , q <sub>1</sub> , q <sub>2</sub>
q <sub>0</sub> , q <sub>1</sub> , q <sub>2</sub>	q <sub>0</sub> , q <sub>1</sub>	q <sub>0</sub> , q <sub>1</sub> , q <sub>2</sub>	φ
φ	φ	φ	φ



bab