

Analysis

MOLECULAR COMMUNICATIONS: MODEL BASED AND DATA DRIVEN RECEIVER DESIGN AND OPTIMIZATION

Goal: To optimize the receiver design

Approach: 1) Model based (Analytical)
2) Data driven (ANN)

System Model

- Transmitter is of point size (small enough).
- The information particles follow the Brownian motion (diffuse randomly and independently of each other in medium)
- Receiver is spherical absorbing type.
- Assumptions are:
 - Temperature is constant
 - Viscosity remains unchanged
 - Diffusion co-efficient remains constant
 - No extra energy needed as particles diffuse freely.
- 3D unbounded diffusion channel
- As particles reach receiver in different time slots, it cause ISI.

Hitting rate of each particle will be

$$f_{hit}^{3D}(t) = \frac{a(d-a)}{d\sqrt{4\pi Dt^3}} e^{-\frac{(d-a)^2}{4Dt}}$$

→ On-off Keying (OOK) modulation scheme is taken

N_{Tx} = No. of particles transmitter releases

If $s_i = \begin{cases} 1 & \text{Transmitter emits particles} \\ 0 & \text{Transmitter doesn't emit particles} \end{cases}$
(signal)

So, hitting probability will be (after t seconds),

$$P_{hit}(t) = \int_0^t f_{hit}(t) dt = \frac{a}{d} \operatorname{erfc}\left(\frac{d-a}{\sqrt{4Dt}}\right)$$

where $\text{erfc}(y) = 1 - \text{erf}(y)$
 $\text{erf}(y) = \int_0^y \frac{2}{\sqrt{\pi}} e^{-x^2} dx$

Therefore, the hitting probability during the $(i-1)^{\text{th}}$ time slot is

$$P_{i-1} = \frac{a}{2} \left\{ \text{erfc}\left(\frac{d-v_i}{\sqrt{4D(i-1)T}}\right) - \text{erfc}\left(\frac{d-x}{\sqrt{4D(i-1)T}}\right) \right\}$$

The average number of received particles after emission from the transmitter will be C (for the j^{th} time slot),

$$C_j = N_{TX} P_j$$

The particles are generated at a time plot, so, it is a time series particle generation, so the number of received particles for a particular time slot will follow the Poisson random variable distribution.

Let the received particles for the i^{th} time slot be x_i ,

$$x_i \sim \text{Poisson}(I_i + s_i C_0)$$

where

$$I_i = \lambda_0 T + \sum_{j=1}^{\infty} s_{i-j} C_j$$

I_i is the sum of ISI and the background noise

λ_0 is the background noise power per unit time

So, probability of receiving particles x_i is,

$$P(x_i | I_i + s_i C_0) = \frac{e^{-(I_i + s_i C_0)} (I_i + s_i C_0)^{x_i}}{x_i!}$$

The, signal to noise ratio (SNR) will be

$$\text{SNR} = 10 \log_{10} \frac{C_0}{2\lambda_0 T}$$

Putting $C_0 = N_{TX} P_0$,

$$\text{SNR} = 10 \log_{10} \frac{N_{TX} P_0}{2\lambda_0 T}$$

$$\frac{\text{SNR}}{10} = \log_{10} \frac{N_{TX} P_0}{2\lambda_0 T}$$

$$\frac{N_{Tx} P_0}{2\lambda_0 T} = 10^{\frac{SNR}{10}}$$

$$N_{Tx} = \frac{2\lambda_0 T}{P_0} 10^{\frac{SNR}{10}}$$

For optimizing the BER, threshold based detector receiver design is proposed which uses the system model information.

The threshold based receivers consists of having a priori information of the previous bits and are of 3 types.

- 1) Zero-bit memory receiver
- 2) One-bit memory receiver
- 3) K-bit memory receiver

Zero-bit memory receiver

Let the demodulation threshold be τ .

Let the estimate of symbol s_i at i th time slot be \hat{s}_i

The demodulation rule will be:

$$\hat{s}_i = \begin{cases} 0 & r_i \leq \tau \\ 1 & r_i > \tau \end{cases}$$

The sub-optimal threshold value can be obtained from

$$P_{app}(r_i | s_i) = \frac{e^{-\lambda |s_i|} (\lambda |s_i|)^{r_i}}{r_i!}$$

$$\text{where } \lambda |s_i| = \cos \theta_i + \frac{\sum_{j=1}^L C_j}{2} + \lambda_0 T$$

(L = slot length)

By using the equality $P_{app}(r_i | s_i = 0) = P_{app}(r_i | s_i = 1)$, we get,

$$\frac{e^{\lambda |s_i=0|} (\lambda |s_i=0|)^{r_i}}{r_i!} = \frac{e^{\lambda |s_i=1|} (\lambda |s_i=1|)^{r_i}}{r_i!}$$

Here $r_i = \tau$

$$\frac{e^{\lambda |s_i=0|} (\lambda |s_i=0|)^{\tau}}{\tau!} = \frac{e^{\lambda |s_i=1|} (\lambda |s_i=1|)^{\tau}}{\tau!}$$

$$e^{\lambda |s_i=0|} (\lambda |s_i=0|)^z = e^{\lambda |s_i=1|} (\lambda |s_i=1|)^z$$

$$e^{-\frac{\sum G_j}{2} + \lambda_0 T} \left(\frac{\sum G_j}{2} + \lambda_0 T \right)^z = e^{C_0 + \frac{\sum G_j}{2} \lambda_0 T} \left(C_0 + \frac{\sum G_j}{2} + \lambda_0 T \right)^z$$

$$e^{C_0} = \left(\frac{C_0 + \frac{\sum G_j}{2} + \lambda_0 T}{\frac{\sum G_j}{2} + \lambda_0 T} \right)^z$$

$$z = \frac{C_0}{\ln \left(1 + \frac{C_0}{\frac{\sum G_j}{2} + \lambda_0 T} \right)}$$

This will be the sub optimal threshold value.

The optimal threshold value will be as follows:

$$(\tau^*, P_e^*) = \arg \min_{\tau} P_e(\tau)$$

where $P_e(\tau)$ is the BER (function of τ)

$$P_e(\tau) = \frac{1}{2^L} \sum_{s_{i-1}} P_e(s_{i-1}, \tau)$$

$$\begin{aligned} \text{where, } P_e(s_{i-1}, \tau) &= \frac{1}{2} \left[Q\left(\lambda_0 T + \sum_{j=1}^L s_{i-j} G_j, \sqrt{\tau}\right) + 1 - Q\left(\lambda_0 T + \sum_{j=1}^L s_{i-j} G_j + C_0, \sqrt{\tau}\right) \right] \\ &= \frac{1}{2} \left[P(a_i \geq \tau | s_i = 0, s_{i-1}) + P(a_i < \tau | s_i = 1, s_{i-1}) \right] \end{aligned}$$

Proof of $P_e(s_{i-1}, \tau)$ is :

$$\begin{aligned} P(a_i \geq \tau | s_i = 0, s_{i-1}) &= P(a_i \geq \tau | \lambda_0 T + \sum_{j=1}^L s_{i-j} G_j) \\ &= \sum_{k=\lceil \tau \rceil}^{\infty} \frac{e^{-(\lambda_0 T + \sum_{j=1}^L s_{i-j} G_j)} (\lambda_0 T + \sum_{j=1}^L s_{i-j} G_j)^k}{k!} \\ &= Q\left(\lambda_0 T + \sum_{j=1}^L s_{i-j} G_j, \sqrt{\tau}\right) \end{aligned}$$

$$\begin{aligned} P(a_i < \tau | s_i = 1, s_{i-1}) &= P(a_i < \tau | \lambda_0 T + \sum_{j=1}^L s_{i-j} G_j + C_0) \\ &= \sum_{k=0}^{\lceil \tau \rceil - 1} \frac{e^{-(\lambda_0 T + \sum_{j=1}^L s_{i-j} G_j + C_0)} (\lambda_0 T + \sum_{j=1}^L s_{i-j} G_j + C_0)^k}{k!} \end{aligned}$$

$$= 1 - \sum_{k=\lceil \tau \rceil}^{\infty} \frac{e^{-(\lambda_0 T + \sum_{j=1}^L s_{i-j} G + C_0)} (\lambda_0 T + \sum_{j=1}^L s_{i-j} G + C_0)^k}{k!}$$

$$= 1 - Q(\lambda_0 T + \sum_{j=1}^L s_{i-j} G + C_0, \lceil \tau \rceil)$$

On putting this values in the equation, we get BER.

One bit receiver design

It has more a priori information compared to zero-bit receiver.

The estimate of symbol s_i at i th time slot be \bar{s}_i

The demodulation rule will be:

$$\bar{s}_i = \begin{cases} 0, & r_i \leq \tau | s_{i-1} \\ 1, & r_i > \tau | s_{i-1} \end{cases}$$

where $\tau | s_{i-1}$ denotes the threshold for i th symbol when previously transmitted symbol is s_{i-1}

The optimal threshold can be calculated as follows:

$$\tau^* | s_{i-1} = \arg \min_{\tau} P_e(\tau, s_{i-1})$$

$$\text{where, } P_e(\tau, s_{i-1}) = \frac{1}{2^{L-1}} \sum_{s_{i-2} \dots s_{i-L}} P_e(s_{i-1}, \tau)$$

So, P_e can be calculated as follows:

$$P_e = \frac{m+n}{2}$$

$$P_e = \frac{1}{2} [P(\bar{s}_i = 1 | s_i = 0) + P(\bar{s}_i = 0 | s_i = 1)]$$

$$m = \frac{1}{2^L} \sum_{s_{i-1}} \sum_{\bar{s}_{i-1}} Q(\lambda | s_{i-1}, s_i = 0, \lceil \tau \rceil \bar{s}_i) \psi(s_{i-1}, \bar{s}_{i-1}, m, n)$$

$$n = \frac{1}{2^L} \sum_{s_{i-1}} \sum_{\bar{s}_{i-1}} Q(1 - Q(\lambda | s_{i-1}, s_i = 1, \lceil \tau \rceil \bar{s}_i)) \psi(s_{i-1}, \bar{s}_{i-1}, m, n)$$

where

$$\psi(s_{i-1}, \bar{s}_{i-1}, m, n) = \begin{cases} m & (s_{i-1} = 0, \bar{s}_{i-1} = 1) \\ 1-m & (s_{i-1} = 0, \bar{s}_{i-1} = 0) \\ n & (s_{i-1} = 1, \bar{s}_{i-1} = 0) \\ 1-n & (s_{i-1} = 1, \bar{s}_{i-1} = 1) \end{cases}$$

Proof of P_e is:

$$P_e = \frac{1}{2} [P(\bar{s}_i = 1 | s_i = 0) + P(\bar{s}_i = 0 | s_i = 1)]$$

$$\begin{aligned} P(\bar{s}_i = 1 | s_i = 0) &= \sum_{s_{i-1}, \bar{s}_{i-1}} P(\bar{s}_i = 1 | s_i = 0, s_{i-1}, \bar{s}_{i-1}) P(\bar{s}_{i-1}, s_{i-1}, \dots, s_{i-L}) \\ &= \sum_{s_{i-1}, \bar{s}_{i-1}} P(x_i \geq \tau | s_{i-1} | s_i = 0, s_{i-1}, \bar{s}_{i-1}) P(\bar{s}_{i-1} | s_{i-1}) P(s_{i-1}) P(s_{i-2}, \dots, s_{i-L}) \\ &= \frac{1}{2^L} \sum_{s_{i-1}, \bar{s}_{i-1}} P(\bar{s}_{i-1} | s_{i-1}) Q(\lambda | s_{i-1}, s_i = 0, \tau | \bar{s}_{i-1}) \\ &= \frac{1}{2^L} \sum_{s_{i-1}, \bar{s}_{i-1}} Q(\lambda | s_{i-1}, s_i = 0, \tau | \bar{s}_{i-1}) \psi(s_{i-1}, \bar{s}_{i-1}, m, n) \end{aligned}$$

= m.

$$\text{where } \lambda | s_{i-1}, s_i = 0 = \lambda_0 T + \sum_{j=1}^L s_{i-j} G_j$$

$$\begin{aligned} P(\bar{s}_i = 0 | s_i = 1) &= \sum_{s_{i-1}, \bar{s}_{i-1}} P(\bar{s}_i = 0 | s_i = 1, s_{i-1}, \bar{s}_{i-1}) P(\bar{s}_{i-1}, s_{i-1}, \dots, s_{i-L}) \\ &= \sum_{s_{i-1}, \bar{s}_{i-1}} P(x_i < \tau | \bar{s}_{i-1} | s_i = 1, s_{i-1}, \bar{s}_{i-1}) P(\bar{s}_{i-1} | s_{i-1}) P(s_{i-1}) P(s_{i-2}, \dots, s_{i-L}) \\ &= \frac{1}{2^L} \sum_{s_{i-1}, \bar{s}_{i-1}} P(\bar{s}_{i-1} | s_{i-1}) (1 - Q(\lambda | s_{i-1}, s_i = 1, \tau | \bar{s}_{i-1})) \end{aligned}$$

$$= \frac{1}{2^L} \sum_{s_{i-1}, \bar{s}_{i-1}} (1 - Q(\lambda | s_{i-1}, s_i = 1, \tau | \bar{s}_{i-1})) \psi(s_{i-1}, \bar{s}_{i-1}, m, n)$$

$$\text{where } \lambda | s_{i-1}, s_i = 1 = \lambda_0 T + \sum_{j=1}^L s_{i-j} G_j + C_0$$

By putting these values in the equation, we get BER.

K-bit memory receiver

It has more a priori information as compared to one-bit receiver. The estimate of symbol s_i at i th time slot be \bar{s}_i .

The demodulation rule will be

$$\bar{s}_i = \begin{cases} 0 & x_i \leq \tau | \bar{s}_{i-1}, \dots, \bar{s}_{i-K} \\ 1 & x_i > \tau | \bar{s}_{i-1}, \dots, \bar{s}_{i-K} \end{cases}$$

where $\bar{s}_{i-1}, \dots, \bar{s}_{i-K}$ are the K previously transmitted symbols.

The optimal threshold can be calculated as follows.

$$z^* | s_{i-1}, \dots, s_{i-k} = \arg \min_z P_e(z, s_{i-1}, \dots, s_{i-k})$$

where,

$$P_e(z, s_{i-1}, \dots, s_{i-k}) = \frac{1}{2^{L-k}} \sum_{s_{i-k+1}, \dots, s_{i-L}} P_e(s_{i-1}, z)$$

Here, $P_e = \frac{m+n}{2}$

where, $m = \frac{1}{2^L} \sum_{s_{i-1}} \sum_{\bar{s}_{i-1}, \dots, \bar{s}_{i-k}} Q(\lambda | s_{i-1}, s_i=0, \Gamma z | \bar{s}_{i-1}, \dots, \bar{s}_{i-k}) \prod_{j=1}^k \psi(s_{i-j}, \bar{s}_{i-j}, m, n)$

$$n = \frac{1}{2^L} \sum_{s_{i-1}} \sum_{\bar{s}_{i-1}, \dots, \bar{s}_{i-k}} (1-Q(\lambda | s_{i-1}, s_i=1, \Gamma z | \bar{s}_{i-1}, \dots, \bar{s}_{i-k})) \prod_{j=1}^k \psi(s_{i-j}, \bar{s}_{i-j}, m, n)$$