CPEN 333 Final Project : Multithreaded Game (Alternative Implementation)

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1 Redesign

Our main redesign of the program structure was to change the model of Inter-Process Communication (IPC) from message passing to shared-memory. This task was reduced to producer-consumer problems for each resource, in which the producer must wait for the consumer to yield the resource and vice-versa, in order for the program to be thread-safe.

Having removed the use of the Queue, QueueHandler classes, we aimed to achieve synchronization with semaphores for the 4 shared resources (i.e. "game_over", "score", "prey", "move") to indicate when they have been produced. All, except for "game_over" are protected via critical section using a dict of mutexes.

Please see the comment at the top of the alternative.py file for a description of the implementation. This file will describe the benefits, disadvantages, and challenges of this.

2 Benefits

Since we were able to eliminate the need for 2 classes, this simplified the relationships between classes, as shown in Figure 1.

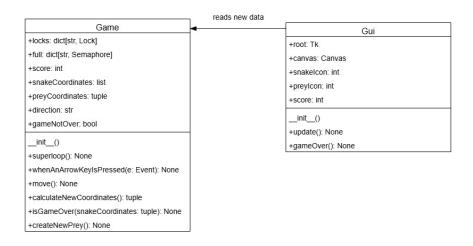


Figure 1: Alternative Implementation - UML Relationship

2.1 Reader-Writer Abstraction

If we abstract this to be a reader-writer synchronization problem, a single writer / many readers can access access the resource at all times. We need to ensure there isn't a race condition between them. This is a problem we have addressed before (e.g. Lab 5).

We have designed this program to have the following constraints:

- {Readers} $\in \{0, 1\}$ where {Writers} $\in \{0\}$
- {Writers} $\in \{0, 1\}$ where {Readers} $\in \{0\}$
- {Readers} \in {0} where {Writers} \in {1}

2.1.1 Producer Behaviour

The game instance acts as both a reader and a writer in different places. It reads the current value(s) of the shared resource and copies it for processing. It also writes the new data to the shared resource. When it is behaving as a reader, we don't need to acquire the mutex or release the semaphore. Race conditions become a possibility for cases where the shared resources are being modified.

3 Disadvantages

3.1 Additional Complexity

We have to ensure that each of the shared resources are accessed effectively. More locks means that there are more critical sections to manage among threads. We needed to ensure that critical sections were kept short and to the point to be able to navigate this added complexity.

In addition, we have the full semaphores to indicate when a new value is available. This condition must be evaluated by the consumer once it successfully acquires a mutex.

3.2 Storage

Along with the mutexes and semaphores, we also need to store the game.preyCoordinates data field for the gui instance. This wasn't necessary with the queue implementation.

4 Challenges

We initially spawned two threads for the game.superloop() and gui.update() methods. This was highly problematic since the program functioned as intended for initial gameplay. Critical sections seemed to be thread-safe, both theoretically and in practice. However, through extensive testing we observed this consistently caused the program to crash when the snake grew at a score $\sim 25-26$.

We realized why the original program design had the QueueHandler.queueHandler() method schedule itself with the Tk.after(100) method. From referring to the Python documentation, we realized that Tkinter widget updates must be done by the main thread.

(See The Python Software Foundation. (n.d.). Tkinter - Python interface to TCL/TK. Python Documentation. https://docs.python.org/3/library/tkinter.html#threading-model)

4.1 Solution

Instead of having the gui.update() method execute loop with spinlocks, we wrote it to be a conditional statement which schedules itself using the Tk.after(100) method. We implemented this with non-blocking semaphore acquisition as well, such that the a critical section should only be entered by the consumer if new data is available.