Homework 1

2-5

(a) Write down a recursion for $f^{(n)}(x)$ in terms of $f^{(n-1)}(\cdot)$ and the parameters of the problem.

For any participant, with n more spins, he/she can choose to spin or stop. When choosing **stopping**, he/she will take away the current accumulation of x. We can obtain $f^{(n)}(x) = x$. While choosing **spinning**, it is also easy to know that

$$f^{(n)}(x) = \sum_{i} f^{(n-1)}(x + r_i) \cdot p_i$$
, where $i = 1, \dots, m$.

$$f^{(n)}(x) = \max\{\sum_{i} f^{(n-1)}(x+r_i) \cdot p_i, x\}$$

(b) Solve the problem numerically for the case in which $m=3, p_i=1/4$, for $i=1,2,3, r_1=1, r_2=5, r_3=10$, and there are one or two more spins available.

At the beginning, the winning is 0, so must choose to spin.

$$f^{(2)}(0) = \begin{cases} f^{(1)}(1), \text{ probability } 1/4\\ f^{(1)}(5), \text{ probability } 1/4\\ f^{(1)}(10), \text{ probability } 1/4\\ 0 \quad \text{probability } 1/4 \end{cases}$$

$$f^{(1)}(x) = \begin{cases} x + 1, \text{ probability } 1/4\\ x + 5, \text{ probability } 1/4\\ x + 10, \text{ probability } 1/4\\ 0 \quad \text{ probability } 1/4\\ x \quad \text{ choose to stop} \end{cases}$$

So when spinning once, the excepted winnings are $1/4 \times (1 + 5 + 10) = 4$. And when spinning twice, $f^{(1)}(x) = 1/4 \times [(x+1) + (x+5) + (x+10)]$. $f^{(2)}(0) = 1/4 \times (f^{(1)}(1) + f^{(1)}(5) + f^{(1)}(10)) = 1/4 \times (4 + 3/4 + 4 + 15/4 + 4 + 30/4) = 6$

In particular, $f^{(0)}(x) = x, f^{(1)}(x) = \max\{3/4x + 4, x\}$. Let 3/4x + 4 = x, x = 16. When x > 16, the participant will take away the winnings; when x < 16, the participant will spin again. Consequently, the expected total winnings are 16 for this case.

2-6

(a) Write down a recursion for v(x) and use it to find an explicit expression for v(x) in terms of $x, \Phi(x), I(x)$, and α , where $I(x) := E(X - x)^+ = \int_x^{\infty} (\xi - x) \phi(\xi) d\xi$.

$$v(x,n) = v(x,n-1) \cdot \alpha \cdot \Phi(x) + \int_x^\infty \xi \phi(\xi) d\xi$$

Let
$$V = \int_{x}^{\infty} \xi \phi(\xi) d\xi$$
.

$$v(x,n) = v(x,1) \cdot [\alpha \cdot \Phi(x)]^{(n-1)} + V[1 + [\alpha \cdot \Phi(x)] + [\alpha \cdot \Phi(x)]^{2} + \cdots + [\alpha \cdot \Phi(x)]^{n-2}]$$

$$v(x,1) = \int_{x}^{\infty} \xi \phi(\xi) d\xi = V$$

Then we can obtain that $v(x) = V \frac{1 - (\alpha \cdot \Phi(x))^n}{1 - \alpha \cdot \Phi(x)}$, $\alpha \in (0, 1)$. Because this is an infinite period, let $n \to \infty$, so $v(x) = \frac{V}{1 - \alpha \cdot \Phi(x)}$. $I(x) = \int_x^{\infty} (\xi - x) \phi(\xi) d\xi = \int_x^{\infty} \xi \phi(\xi) d\xi - x(1 - \Phi(x))$

So
$$V=I(x)+x(1-\Phi(x)),$$
 substitute it in the $v(x),$ we get $v(x)=\frac{I(x)+x(1-\Phi(x))}{1-\alpha\cdot\Phi(x)}$

(b)
$$\alpha=0.99$$
 bids $\sim N(1000,200)$. $\mathbf{x}=1200$. So $v(1200)=\frac{I(1200)+1200(1-\Phi(1200))}{1-\alpha\Phi(1200)}$, $I(1200)=200I_N(1)$, $\Phi(1200)=\Phi_N(1)=0.841,I_N(1)=\phi_N(1)-1+\Phi_N(1)=0.083$. By calculating, $v(1200)=1238.9$

2-7

(a) Display Ω .

$$\Omega = \{(3,3,1), (2,3,1), (1,3,1), (0,3,1), (3,0,1),$$

$$(2,0,1), (1,0,1), (0,0,1), (1,1,1), (2,2,1),$$

$$(3,3,2), (2,3,2), (1,3,2), (0,3,2), (3,0,2),$$

$$(2,0,2), (1,0,2), (0,0,2), (1,1,2), (2,2,2)\}$$

Total: 20

$$(b)S_0 = \{(3,3,1)\}$$

$$S_1 = \{(2,3,2), (2,2,2), (1,3,2)\}$$

$$S_2 = S_0 \cup \{(2,3,1)\}$$

$$S_3 = S_1 \cup \{(0,3,2)\}$$

$$S_4 = S_2 \cup \{(1,3,1)\}$$

$$S_5 = S_3 \cup \{(1, 1, 2)\}$$

$$S_6 = S_4 \cup \{(2, 2, 1)\}$$

$$S_7 = S_5 \cup \{(2, 0, 2)\}$$

$$S_8 = S_6 \cup \{(3, 0, 1)\}$$

$$S_9 = S_7 \cup \{(1, 0, 2)\}$$

$$S_{10} = S_8 \cup \{(2, 0, 1), (1, 1, 1)\}$$

$$S_{11} = S_9 \cup \{(0, 0, 2)\}.$$

$$m = 11.$$

3-4

(a) Write down a recursion for $f^n(x|S)$.

$$f^{(n)}(x|S) = \begin{cases} \sum_{i} f^{(n-1)}(x + r_i|S) \cdot p_i, & x < S. \\ x, & x \ge S. \end{cases}$$

(b) Being indifferent means $\sum_i f^{(n-1)}(S+r_i|S) \cdot p_i = S = f^{(n)}(S|S) \Rightarrow \sum_i^m (S^*+r_i)p_i = S^* \Rightarrow S^* = \frac{\mu}{p}.$

3-5

(a) Write a recursion for f_n in terms of f_{n-1} .

$$f_n = f_{n-1} \cdot \alpha \cdot \Phi(\alpha f_{n-1}) + \int_{\alpha f_{n-1}}^{\infty} \xi \phi(\xi) d\xi = I(\alpha f_{n-1}) + \alpha f_{n-1}$$

(b) Calculate f_n .

$$I_N(-0.05) = \phi_N(-0.05) + 0.05 \times (1 - \Phi_N(-0.05)) = 0.398 - 1 + 0.48 = 0.4244$$

$$I(990) = 200I_N(-0.05) = 200 \times (0.4244) = 84.88$$

 $f_2 = \alpha \mu + I(\alpha \mu) = 0.99 * 1000 + 84.88 = 1074.88$
 $f_3 = \alpha f_2 + I(\alpha f_2) = 1115.92$

$$f_4 = \alpha f_3 + I(\alpha f_3) = 1142.87$$

(c) Let $f_n = f_{n-1}$, we can obtain that $f_n = \alpha f_n \Phi(\alpha f_n) + I(\alpha f_n) + \alpha f_n (1 - \Phi(\alpha f_n))$. Obviously, we know that f_n is the solution of $x = \frac{I(\alpha x)}{1-\alpha}$.

3-6

(a) Because there is a discount factor which is less than 1, when you wait it for a long period, your profit will decrease as time went by. So, it is obvious that we don't need to wait for all bids to be received.

$$f_n(x) = \begin{cases} \alpha E(f_{n-1}(X) - f_{n-1}(x))^+ + f_{n-1}(x), \text{ not accept this period.} \\ I(x) + x \quad , \text{ accept.} \end{cases}$$

$$f_n(x) = \max\{\alpha(E(f_{n-1}(X) - f_{n-1}(x))^+ + f_{n-1}(x)), I(x) + x\}$$

4-4

The original formulation is
$$v(x,S) = \begin{cases} c & \text{if } x = 0 \\ px + qv(x-1,S) & \text{if } 1 \leq x \leq S \\ v(S,S) & \text{if } S < x \end{cases}$$

for this case, we can continue drive until we find an unoccupied space. So define a function F(x), f(x,i), where x < 0.

$$F(x) = pf(x,0) + qf(x,1) \text{ and}$$

$$f(x,i) = \begin{cases} -x & \text{if } i = 0\\ F(x-1) & \text{if } i = 1 \end{cases}$$

i=0 means the space isn't occupied, only choose to park here.

Now, we can obtain
$$F(x) = p(-x) + qF(x-1) = p(\sum_{i=-x}^{n-x-1} iq^{i+x}) = \frac{q[1+(n-1)q^n - nq^{n-1}]}{(1-q)^2} - x\frac{1-q^n}{1-q}, x < 0, \text{ let } n = \infty, F(-1) = \frac{q}{p^2} + \frac{1}{p}.$$

Let v(0,S) = F(-1), we can obtain the formulation for this question:

$$v(x,S) = \begin{cases} v(0,S) = \frac{q}{p^2} + \frac{1}{p} & \text{if } x = 0 \\ px + qv(x - 1, S) & \text{if } 1 \le x \le S \\ v(S,S) & \text{if } S < x \end{cases}$$
Consequently, $v(S,S) = p \sum_{i=0}^{S} q^i(S-i) + q^S \frac{1}{(1-q)^2} = S - \frac{q(1-q^S)}{p} + q^S (\frac{q}{p^2} + \frac{1}{p}).$

4-5

$$f^{(n)}(x|S) = \begin{cases} \sum_{i} f^{(n-1)}(x + r_i|S) \cdot p_i, \ x < S. \\ x, \ x \ge S. \end{cases}$$

Define $G^{n}(x) = -x + \sum_{i=1}^{m} f^{(n-1)}(x + r_{i})p_{i}$.

$$G^{(n)}(x) = \begin{cases} -x + \sum_{i} (G^{(n-1)}(x) + x) \cdot p_i, & x < S. \\ -x + \sum_{i} p_i \cdot x, & x \ge S. \end{cases}$$

Use mathematical induction to prove $G^{(n)}(x)$ is decreasing in x for each n.

1.
$$G'^{(i+1)}(x) = -1 + \sum_{i=1}^{m} p_i < 0.$$

If
$$f^{i}(x|S) = 0$$
, then $G'^{(i+1)}(x) = -1 < 0$.

2. Suppose $G'^{(n)}(x) < 0$, then prove $G'^{(n+1)}(x) < 0$.

Omitted.