

# Homework 1

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## 2-5

(a) Write down a recursion for  $f^{(n)}(x)$  in terms of  $f^{(n-1)}(\cdot)$  and the parameters of the problem.

For any participant, with  $n$  more spins, he/she can choose to spin or stop. When choosing **stopping**, he/she will take away the current accumulation of  $x$ . We can obtain  $f^{(n)}(x) = x$ . While choosing **spinning**, it is also easy to know that

$$f^{(n)}(x) = \sum_i f^{(n-1)}(x + r_i) \cdot p_i, \text{ where } i = 1, \dots, m.$$

$$f^{(n)}(x) = \max\{\sum_i f^{(n-1)}(x + r_i) \cdot p_i, x\}$$

(b) Solve the problem numerically for the case in which  $m = 3, p_i = 1/4$ , for  $i = 1, 2, 3, r_1 = 1, r_2 = 5, r_3 = 10$ , and there are one or two more spins available.

At the beginning, the winning is 0, so must choose to spin.

$$f^{(2)}(0) = \begin{cases} f^{(1)}(1), \text{ probability } 1/4 \\ f^{(1)}(5), \text{ probability } 1/4 \\ f^{(1)}(10), \text{ probability } 1/4 \\ 0 \quad \text{probability } 1/4 \end{cases}$$

$$f^{(1)}(x) = \begin{cases} x + 1, \text{probability } 1/4 \\ x + 5, \text{probability } 1/4 \\ x + 10, \text{probability } 1/4 \\ 0 \quad \text{probability } 1/4 \\ x \quad \text{choose to stop} \end{cases}$$

So when spinning once, the expected winnings are  $1/4 \times (1 + 5 + 10) = 4$ . And when spinning twice,  $f^{(1)}(x) = 1/4 \times [(x + 1) + (x + 5) + (x + 10)]$ .  $f^{(2)}(0) = 1/4 \times (f^{(1)}(1) + f^{(1)}(5) + f^{(1)}(10)) = 1/4 \times (4 + 3/4 + 4 + 15/4 + 4 + 30/4) = 6$

In particular,  $f^{(0)}(x) = x, f^{(1)}(x) = \max\{3/4x + 4, x\}$ . Let  $3/4x + 4 = x, x = 16$ . When  $x > 16$ , the participant will take away the winnings; when  $x < 16$ , the participant will spin again. Consequently, the expected total winnings are 16 for this case.

## 2-6

(a) Write down a recursion for  $v(x)$  and use it to find an explicit expression for  $v(x)$  in terms of  $x, \Phi(x), I(x)$ , and  $\alpha$ , where  $I(x) := E(X - x)^+ = \int_x^\infty (\xi - x)\phi(\xi)d\xi$ .

$$v(x, n) = v(x, n - 1) \cdot \alpha \cdot \Phi(x) + \int_x^\infty \xi \phi(\xi) d\xi$$

$$\text{Let } V = \int_x^\infty \xi \phi(\xi) d\xi.$$

$$v(x, n) = v(x, 1) \cdot [\alpha \cdot \Phi(x)]^{(n-1)} + V[1 + [\alpha \cdot \Phi(x)] + [\alpha \cdot \Phi(x)]^2 + \dots + [\alpha \cdot \Phi(x)]^{n-2}]$$

$$v(x, 1) = \int_x^\infty \xi \phi(\xi) d\xi = V$$

Then we can obtain that  $v(x) = V \frac{1-(\alpha \cdot \Phi(x))^n}{1-\alpha \cdot \Phi(x)}$ ,  $\alpha \in (0, 1)$ . Because this is an infinite period, let  $n \rightarrow \infty$ , so  $v(x) = \frac{V}{1-\alpha \cdot \Phi(x)}$ .

$$I(x) = \int_x^\infty (\xi - x) \phi(\xi) d\xi = \int_x^\infty \xi \phi(\xi) d\xi - x(1 - \Phi(x))$$

So  $V = I(x) + x(1 - \Phi(x))$ , substitute it in the  $v(x)$ , we get  $v(x) = \frac{I(x) + x(1 - \Phi(x))}{1 - \alpha \cdot \Phi(x)}$

(b)  $\alpha = 0.99$  bids  $\sim N(1000, 200)$ .  $x = 1200$ . So  $v(1200) = \frac{I(1200) + 1200(1 - \Phi(1200))}{1 - \alpha \Phi(1200)}$ ,  $I(1200) = 200I_N(1)$ ,  $\Phi(1200) = \Phi_N(1) = 0.841$ ,  $I_N(1) = \phi_N(1) - 1 + \Phi_N(1) = 0.083$ .

By calculating,  $v(1200) = 1238.9$

## 2-7

(a) Display  $\Omega$ .

$$\Omega = \{(3, 3, 1), (2, 3, 2), (2, 2, 2), (1, 3, 2), (2, 3, 1), (0, 3, 2), (1, 3, 1), (1, 1, 2), (2, 2, 1), (2, 0, 2), (3, 0, 1), (1, 0, 2), (2, 0, 1), (1, 1, 1), (0, 0, 2)\}$$

(b)  $S_0 = \{(3, 3, 1)\}$ ,  $S_1 = \{(2, 3, 2), (2, 2, 2), (1, 3, 2)\}$ ,

$$S_2 = \{(2, 3, 1)\}$$
,  $S_3 = \{(0, 3, 2)\}$ ,

$$S_4 = \{(1, 3, 1)\}$$
,  $S_5 = \{(1, 1, 2)\}$ ,

$$S_6 = \{(2, 2, 1)\}$$
,  $S_7 = \{(2, 0, 2)\}$ ,

$$S_8 = \{(3, 0, 1)\}$$
,  $S_9 = \{(1, 0, 2)\}$ ,

$$S_{10} = \{(2, 0, 1), (1, 1, 1)\}$$
,  $S_{11} = \{(0, 0, 2)\}$ .

$$m = 11.$$

## 3-4

(a) Write down a recursion for  $f^n(x|S)$ .

$$f^{(n)}(x|S) = \sum_i f^{(n-1)}(x + r_i|S) \cdot p_i.$$

$$f^{(n)}(x|S) = \begin{cases} \sum_i f^{(n-1)}(x + r_i|S) \cdot p_i, & \text{spin or } x < S. \\ x & , \text{stop.} \end{cases}$$

(b) Being indifferent means  $\sum_i f^{(n-1)}(S + r_i|S) \cdot p_i = S = f^{(n)}(S|S) \Rightarrow \sum_i^m (S^* + r_i)p_i = S^* \Rightarrow S^* = \frac{\mu}{p}$ .

### 3-5

(a) Write a recursion for  $f_n$  in terms of  $f_{n-1}$

$$f_n = f_{n-1} \cdot \alpha \cdot \Phi(\alpha f_{n-1}) + \int_{\alpha f_{n-1}}^{\infty} \xi \phi(\xi) d\xi$$

$$\text{Let } V = \int_{\alpha x}^{\infty} \xi \phi(\xi) d\xi = I(\alpha x) + \alpha x(1 - \Phi(\alpha x)).$$

(b) Calculate  $f_n$ .  $I_N(-0.05) = \phi_N(-0.05) - 1 + \Phi_N(-0.05) = 0.398 - 1 + 0.48 = -0.122$ ,

$$V = I(990) + 990(1 - \Phi(990)) = 200I_N(-0.05) + 990(1 - \Phi_N(-0.05)) = 200 \times (-0.12) + 990(1 - 0.48) = 490.5$$

$$f_2 = \alpha \mu \Phi(\alpha \mu) + \int_{\alpha \mu}^{\infty} \xi \phi(\xi) d\xi = 0.48 * 0.99 * 1000 + 490.5 = 965.7$$

$$f_3 = \alpha \Phi(\alpha f_2) f_2 + \int_{\alpha f_2}^{\infty} \xi \phi(\xi) d\xi = 916.5$$

$$f_4 = \alpha \Phi(\alpha x) f_3 + \int_{\alpha x}^{\infty} \xi \phi(\xi) d\xi = 843.3$$

(c) Let  $f_n = f_{n-1}$ , we can obtain that  $f_n = \alpha f_n \Phi(\alpha f_n) + I(\alpha f_n) + \alpha f_n(1 - \Phi(\alpha f_n))$ . Obviously, we can obtain that  $f_n = \frac{I(\alpha f_n)}{1 - \alpha}$ .

### 3-6

(a) Because there is a discount factor which is less than 1, when you wait it for a long period, your profit will be very low as time went by. So, it is obvious that we don't need to wait for all bids to be received.

(b)

$$f_n(x) = \begin{cases} \alpha f_{n-1}(x/\alpha), \text{ not accept this period.} \\ \alpha^t x, \text{ accept.} \end{cases}$$

$$f_n(x) = \alpha f_{n-1}(x/\alpha) \int_{x/\alpha}^{\infty} \phi(\xi) d\xi + \alpha^t x \Phi(x/\alpha)$$

To be honest, I'm confused by this question....

#### 4-4

$$\text{The original formulation is } v(x, S) = \begin{cases} c & \text{if } x = 0 \\ px + qv(x-1, S) & \text{if } 1 \leq x \leq S \\ v(S, S) & \text{if } S < x \end{cases},$$

for this case, we can continue drive until we find an unoccupied space. So define a function  $F(x)$ ,  $f(x, i)$ , where  $x < 0$ .

$$F(x) = pf(x, 0) + qf(x, 1) \text{ and}$$

$$f(x, i) = \begin{cases} -x & \text{if } i = 0 \\ F(x-1) & \text{if } i = 1 \end{cases}$$

$i = 0$  means the space isn't occupied, only choose to park here.

Now, we can obtain  $F(x) = p(-x) + qF(x-1) = p(\sum_{i=-x}^{n-x-1} iq^{i+x}) = \frac{q[1+(n-1)q^n - nq^{n-1}]}{(1-q)^2} - x\frac{1-q^n}{1-q}$ ,  $x < 0$ , let  $n = \infty$ ,  $F(-1) = \frac{1}{(1-q)^2}$

Let  $v(0, S) = F(-1)$ , we can obtain the formulation for this question:

$$v(x, S) = \begin{cases} v(0, S) = \frac{1}{(1-q)^2} & \text{if } x = 0 \\ px + qv(x-1, S) & \text{if } 1 \leq x \leq S \\ v(S, S) & \text{if } S < x \end{cases}$$

Consequently,  $v(S, S) = p \sum_{i=0}^S q^i (S-i) + q^S \frac{1}{(1-q)^2} = S - \frac{q(1-q^S)}{p} + q^S \frac{1}{(1-q)^2}$ .

#### 4-5

Define  $G^m(x) = -x + \sum_i^m f^{n-1}(x+r_i)p_i$ .

The derivative of  $G^n(x)$  equals  $(-1 + \sum_i^m f'^{n-1}(x + r_i)p_i)$ , because  $f^{n-1}(x + r_i)$  is a linear function of  $x$ , and the coefficient of  $x$  is 1. So  $\sum_i^m f'^{n-1}(x + r_i)p_i = \sum_i^m p_i = 1 - p < 1$ . So the derivative of  $G^n(x)$  is less than 0. Therefore,  $G^n(x) < G^n(x + 1)$ . According to Exercise 3.4, the decision rule we decided is indeed optimal.