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1  |----- MODULE PaxosStore -----|
   | Specification of the consensus protocol in PaxosStore. |
   | See [PaxosStore@VLDB2017](https://www.vldb.org/pvldb/vol10/p1730-lin.pdf) by Tencent. |
8  | EXTENDS Integers, FiniteSets |
9  |-----|
10 |  $Max(m, n) \triangleq \text{IF } m > n \text{ THEN } m \text{ ELSE } n$  |
11 |  $Injective(f) \triangleq \forall a, b \in \text{DOMAIN } f : (a \neq b) \Rightarrow (f[a] \neq f[b])$  |
12 |-----|
13 | CONSTANTS |
14 |   Participant, the set of participants |
15 |   Value the set of possible input values for Participant to propose |
17 |  $None \triangleq \text{CHOOSE } b : b \notin \text{Value}$  |
18 |  $NP \triangleq \text{Cardinality}(\text{Participant})$  number of  $p \in \text{Participants}$  |
20 |  $Quorum \triangleq \{Q \in \text{SUBSET } Participant : \text{Cardinality}(Q) * 2 = NP + 1\}$  |
21 | ASSUME QuorumAssumption  $\triangleq$  |
22 |    $\wedge \forall Q \in Quorum : Q \subseteq Participant$  |
23 |    $\wedge \forall Q1, Q2 \in Quorum : Q1 \cap Q2 \neq \{\}$  |
25 |  $Ballot \triangleq Nat$  |
27 |  $PIndex \triangleq \text{CHOOSE } f \in [Participant \rightarrow 1 \dots NP] : Injective(f)$  TODO: (1) symmetry set? (2) model |
28 |  $Bals(p) \triangleq \{b \in Ballot : b \% NP = PIndex[p] - 1\}$  allocate ballots for each  $p \in Participant$  |
29 |-----|
30 |  $State \triangleq [maxBal : Ballot \cup \{-1\},$  |
31 |    $maxVVal : Ballot \cup \{-1\}, maxVVal : Value \cup \{None\}]$  |
33 |  $InitState \triangleq [maxBal \mapsto -1, maxVVal \mapsto -1, maxVVal \mapsto None]$  |
   | For simplicity, in this specification, we choose to send the complete state of a participant each |
   | time. When receiving such a message, the participant processes only the “partial” state it needs. |
39 |  $Message \triangleq [type : \{\text{“Prepare”, “Accept”, “ACK”}\},$  |
40 |    $from : Participant, to : \text{SUBSET } Participant,$  TODO: remove “to” |
41 |    $state : [Participant \rightarrow State]]$  |
42 |-----|
43 | VARIABLES |
44 |   state,  $state[p][q]$ : the state of  $q \in Participant$  from the view of  $p \in Participant$  |
45 |   msgs the set of messages that have been sent |
47 |  $vars \triangleq \langle state, msgs \rangle$  |
49 |  $TypeOK \triangleq$  |
50 |    $\wedge state \in [Participant \rightarrow [Participant \rightarrow State]]$  |
51 |    $\wedge msgs \subseteq Message$  |
53 |  $Send(m) \triangleq msgs' = msgs \cup \{m\}$  |
54 |-----|

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55 Init  $\triangleq$ 
56    $\wedge \text{state} = [p \in \text{Participant} \mapsto [q \in \text{Participant} \mapsto \text{InitState}]]$ 
57    $\wedge \text{msgs} = \{\}$ 

p  $\in$  Participant starts the prepare phase by issuing a ballot b  $\in$  Ballot.


61 Prepare(p, b)  $\triangleq$ 
62    $\wedge \text{state}[p][p].\text{maxBal} < b$ 
63    $\wedge b \in \text{Bals}(p)$ 
64    $\wedge \text{state}' = [\text{state} \text{ EXCEPT } ![p][p].\text{maxBal} = b]$ 
65    $\wedge \text{Send}([type \mapsto \text{"Prepare"}, from \mapsto p, to \mapsto \text{Participant}, state \mapsto \text{state}'[p]])$ 

q  $\in$  Participant updates its own state state[q] according to the actual state pp of p  $\in$  Participant extracted from a message m  $\in$  Message it receives. This is called by OnMessage(q).



Note: pp is m.state[p]; it may not be equal to state[p][p] at the time UpdateState is called.


74 UpdateState(q, p, pp)  $\triangleq$ 
75   state' = [state EXCEPT
76      $![q][p].\text{maxBal} = \text{Max}(@, pp.\text{maxBal}),$ 
77      $![q][p].\text{maxVVal} = \text{Max}(@, pp.\text{maxVVal}),$ 
78      $![q][p].\text{maxVVal} = \text{IF } \text{state}[q][p].\text{maxVVal} < pp.\text{maxVVal}$ 
79       THEN pp.maxVVal ELSE @,
80      $![q][q].\text{maxBal} = \text{Max}(@, pp.\text{maxBal}),$ 
81      $![q][q].\text{maxVVal} = \text{IF } \text{state}[q][q].\text{maxBal} \leq pp.\text{maxVVal}$ 
82       THEN pp.maxVVal ELSE @, make promise
83      $![q][q].\text{maxVVal} = \text{IF } \text{state}[q][q].\text{maxBal} \leq pp.\text{maxVVal}$  TODO: write-once?
84     THEN pp.maxVVal ELSE @] accept


q  $\in$  Participant receives and processes a message in Message.


88 OnMessage(q)  $\triangleq$ 
89    $\exists m \in \text{msgs} :$ 
90      $\wedge m.type = \text{"ACK"} \Rightarrow m.to = \{q\}$ 
91      $\wedge \text{LET } p \triangleq m.from$ 
92     IN UpdateState(q, p, m.state[p])
93      $\wedge \text{IF } \vee m.state[q].\text{maxBal} < \text{state}'[q][q].\text{maxBal}$  TODO: delete "if"?
94        $\vee m.state[q].\text{maxVVal} < \text{state}'[q][q].\text{maxVVal}$ 
95       THEN Send([type  $\mapsto$  "ACK", from  $\mapsto$  q, to  $\mapsto$  {m.from}, state  $\mapsto$  state'[q])
96       ELSE UNCHANGED msgs

p  $\in$  Participant starts the accept phase by issuing the ballot b  $\in$  Ballot with value v  $\in$  Value.


101 Accept(p, b, v)  $\triangleq$ 
102    $\wedge \neg \exists m \in \text{msgs} :$  TODO: delete it? to allow repeating Phase2a?
103     m.type = "Accept"  $\wedge m.state[p].\text{maxBal} = b$ 
104    $\wedge b \in \text{Bals}(p)$  TODO: delete it? to break "client-restricted config"?
105    $\wedge \exists Q \in \text{Quorum} : \forall q \in Q : \text{state}[p][q].\text{maxBal} = b$ 
106    $\wedge \forall q \in \text{Participant} : \text{state}[p][q].\text{maxVVal} = -1$  free to pick its own value
107    $\vee \exists q \in \text{Participant} : v$  is the value with the highest maxVVal
108      $\wedge \text{state}[p][q].\text{maxVVal} = v$ 
109      $\wedge \forall r \in \text{Participant} : \text{state}[p][q].\text{maxVVal} \geq \text{state}[p][r].\text{maxVVal}$ 

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