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1  |----- MODULE OneValePerBallotPaxosStore -----|
  | Specification of the consensus protocol in PaxosStore.
  | See [PaxosStore@VLDB2017](https://www.vldb.org/pvldb/vol10/p1730-lin.pdf) by Tencent.
  | In this version (adapted from “UniversalPaxosStore.tla”):
  | - Adding
12 | EXTENDS Integers, FiniteSets
13 |-----|
14 |  $Max(m, n) \triangleq \text{IF } m > n \text{ THEN } m \text{ ELSE } n$ 
15 |-----|
16 | CONSTANTS
17 |   Participant, the set of participants
18 |   Value         the set of possible input values for Participant to propose
20 |  $None \triangleq \text{CHOOSE } b : b \notin \textit{Value}$ 
22 |  $Quorum \triangleq \{Q \in \text{SUBSET } \textit{Participant} : \text{Cardinality}(Q) * 2 = \text{Cardinality}(\textit{Participant}) + 1\}$ 
23 | ASSUME  $QuorumAssumption \triangleq$ 
24 |    $\wedge \forall Q \in Quorum : Q \subseteq \textit{Participant}$ 
25 |    $\wedge \forall Q1, Q2 \in Quorum : Q1 \cap Q2 \neq \{\}$ 
27 |  $Ballot \triangleq \textit{Nat}$ 
28 |-----|
29 |  $State \triangleq [maxBal : Ballot \cup \{-1\},$ 
30 |    $maxVVal : Ballot \cup \{-1\}, maxVVal : Value \cup \{None\}]$ 
32 |  $InitState \triangleq [maxBal \mapsto -1, maxVVal \mapsto -1, maxVVal \mapsto None]$ 
  | For simplicity, in this specification, we choose to send the complete state of a participant each
  | time. When receiving such a message, the participant processes only the “partial” state it needs.
38 |  $Message \triangleq [from : \textit{Participant}, to : \text{SUBSET } \textit{Participant}, state : [\textit{Participant} \rightarrow \textit{State}]]$ 
39 |-----|
40 | VARIABLES
41 |   state,  $state[p][q]$ : the state of  $q \in \textit{Participant}$  from the view of  $p \in \textit{Participant}$ 
42 |   bals, the subset of Ballot consisting of  $b \in \textit{Ballot}$  such that  $\text{Accept}(\_, b, \_)$ 
43 |   msgs the set of messages that have been sent
45 |  $vars \triangleq \langle state, bals, msgs \rangle$ 
47 |  $TypeOK \triangleq$ 
48 |    $\wedge state \in [\textit{Participant} \rightarrow [\textit{Participant} \rightarrow \textit{State}]]$ 
49 |    $\wedge bals \subseteq \textit{Ballot}$ 
50 |    $\wedge msgs \subseteq \textit{Message}$ 
52 |  $Send(m) \triangleq msgs' = msgs \cup \{m\}$ 
53 |-----|
54 |  $Init \triangleq$ 
55 |    $\wedge state = [p \in \textit{Participant} \mapsto [q \in \textit{Participant} \mapsto \textit{InitState}]]$ 

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56  $\wedge \text{bals} = \{\}$   
57  $\wedge \text{msgs} = \{\}$

$p \in \text{Participant}$  starts the prepare phase by issuing a ballot  $b \in \text{Ballot}$ .

61  $\text{Prepare}(p, b) \triangleq$   
62  $\wedge \text{state}[p][p].\text{maxBal} < b$   
63  $\wedge \text{state}' = [\text{state} \text{ EXCEPT } ![p][p].\text{maxBal} = b]$   
64  $\wedge \text{Send}([from \mapsto p, to \mapsto \text{Participant}, state \mapsto \text{state}'[p]])$   
65  $\wedge \text{UNCHANGED } \text{bals}$

$q \in \text{Participant}$  updates its own state  $\text{state}[q]$  according to the actual state  $pp$  of  $p \in \text{Participant}$  extracted from a message  $m \in \text{Message}$  it receives. This is called by  $\text{OnMessage}(q)$ .

Note:  $pp$  is  $m.\text{state}[p]$ ; it may not be equal to  $\text{state}[p][p]$  at the time  $\text{UpdateState}$  is called.

74  $\text{UpdateState}(q, p, pp) \triangleq$   
75  $\text{state}' = [\text{state} \text{ EXCEPT}$   
76  $![q][p].\text{maxBal} = \text{Max}(@, pp.\text{maxBal}),$   
77  $![q][p].\text{maxVVal} = \text{Max}(@, pp.\text{maxVVal}),$   
78  $![q][p].\text{maxVVal} = \text{IF } \text{state}[q][p].\text{maxVVal} < pp.\text{maxVVal}$   
79  $\text{THEN } pp.\text{maxVVal} \text{ ELSE } @,$   
80  $![q][q].\text{maxBal} = \text{Max}(@, pp.\text{maxBal}),$   
81  $![q][q].\text{maxVVal} = \text{IF } \text{state}[q][q].\text{maxBal} \leq pp.\text{maxVVal}$   
82  $\text{THEN } pp.\text{maxVVal} \text{ ELSE } @, \text{ make promise}$   
83  $![q][q].\text{maxVVal} = \text{IF } \text{state}[q][q].\text{maxBal} \leq pp.\text{maxVVal}$   
84  $\text{THEN } pp.\text{maxVVal} \text{ ELSE } @] \text{ accept}$

$q \in \text{Participant}$  receives and processes a message in  $\text{Message}$ .

88  $\text{OnMessage}(q) \triangleq$   
89  $\wedge \exists m \in \text{msgs} :$   
90  $\wedge q \in m.to$   
91  $\wedge \text{LET } p \triangleq m.from$   
92  $\text{IN } \text{UpdateState}(q, p, m.\text{state}[p])$   
93  $\wedge \text{IF } \vee m.\text{state}[q].\text{maxBal} < \text{state}'[q][q].\text{maxBal}$   
94  $\vee m.\text{state}[q].\text{maxVVal} < \text{state}'[q][q].\text{maxVVal}$   
95  $\text{THEN } \text{Send}([from \mapsto q, to \mapsto \{m.from\}, state \mapsto \text{state}'[q]])$   
96  $\text{ELSE } \text{UNCHANGED } \text{msgs}$   
97  $\wedge \text{UNCHANGED } \text{bals}$

$p \in \text{Participant}$  starts the accept phase by issuing the ballot  $b \in \text{Ballot}$  with value  $v \in \text{Value}$ .

102  $\text{Accept}(p, b, v) \triangleq$   
103  $\wedge b \notin \text{bals}$   
104  $\wedge \exists Q \in \text{Quorum} : \forall q \in Q : \text{state}[p][q].\text{maxBal} = b$   
105  $\wedge \forall q \in \text{Participant} : \text{state}[p][q].\text{maxVVal} = -1 \text{ free to pick its own value}$   
106  $\vee \exists q \in \text{Participant} : v \text{ is the value with the highest } \text{maxVVal}$   
107  $\wedge \text{state}[p][q].\text{maxVVal} = v$   
108  $\wedge \forall r \in \text{Participant} : \text{state}[p][q].\text{maxVVal} \geq \text{state}[p][r].\text{maxVVal}$   
109  $\wedge \text{state}' = [\text{state} \text{ EXCEPT } ![p][p].\text{maxVVal} = b,$   
110  $![p][p].\text{maxVVal} = v]$   
111  $\wedge \text{Send}([from \mapsto p, to \mapsto \text{Participant}, state \mapsto \text{state}'[p]])$

