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1  ┌────────────────── MODULE UniversalPaxosStore ───────────────────┐
    Specification of the consensus protocol in PaxosStore.
    See [PaxosStore@VLDB2017](https://www.vldb.org/pvldb/vol10/p1730-lin.pdf) by Tencent.
    In this version (adopted from "PaxosStore.tla"):
    - Client-restricted config (Ballot)
      - Message types (i.e., "Prepare", "Accept", "ACK") are deleted. No state flags (such as "Prepare",
        "Wait-Prepare", "Accept", "Wait-Accept" are needed.
15 EXTENDS Integers, FiniteSets
16 ┌────────────────────────────────────────────────────────────────────────┐
17  $Max(m, n) \triangleq$  IF  $m > n$  THEN  $m$  ELSE  $n$ 
18  $Injective(f) \triangleq \forall a, b \in \text{DOMAIN } f : (a \neq b) \Rightarrow (f[a] \neq f[b])$ 
19 ┌────────────────────────────────────────────────────────────────────────┐
20 CONSTANTS
21   Participant, the set of participants
22   Value         the set of possible input values for Participant to propose
23
24   None  $\triangleq$  CHOOSE  $b : b \notin \text{Value}$ 
25    $NP \triangleq \text{Cardinality}(\text{Participant})$  number of  $p \in \text{Participants}$ 
26
27    $Quorum \triangleq \{Q \in \text{SUBSET } Participant : \text{Cardinality}(Q) * 2 = NP + 1\}$ 
28   ASSUME  $QuorumAssumption \triangleq$ 
29      $\wedge \forall Q \in Quorum : Q \subseteq Participant$ 
30      $\wedge \forall Q1, Q2 \in Quorum : Q1 \cap Q2 \neq \{\}$ 
31
32    $Ballot \triangleq Nat$ 
33
34    $PIndex \triangleq$  CHOOSE  $f \in [Participant \rightarrow 1..NP] : Injective(f)$ 
35    $Bals(p) \triangleq \{b \in Ballot : b \% NP = PIndex[p] - 1\}$  allocate ballots for each  $p \in Participant$ 
36 ┌────────────────────────────────────────────────────────────────────────┐
37    $State \triangleq [maxBal : Ballot \cup \{-1\},$ 
38      $maxVVal : Value \cup \{None\}]$ 
39
40    $InitState \triangleq [maxBal \mapsto -1, maxVVal \mapsto -1, maxVVal \mapsto None]$ 
    For simplicity, in this specification, we choose to send the complete state of a participant each
    time. When receiving such a message, the participant processes only the "partial" state it needs.
46    $Message \triangleq [from : Participant, to : \text{SUBSET } Participant, state : [Participant \rightarrow State]]$ 
47 ┌────────────────────────────────────────────────────────────────────────┐
48 VARIABLES
49   state,  $state[p][q]$ : the state of  $q \in Participant$  from the view of  $p \in Participant$ 
50   msgs   the set of messages that have been sent
51
52    $vars \triangleq \langle state, msgs \rangle$ 
53
54    $TypeOK \triangleq$ 
55      $\wedge state \in [Participant \rightarrow [Participant \rightarrow State]]$ 
56      $\wedge msgs \subseteq Message$ 

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58  $Send(m) \triangleq msgs' = msgs \cup \{m\}$ 
59
60  $Init \triangleq$ 
61    $\wedge state = [p \in Participant \mapsto [q \in Participant \mapsto InitState]]$ 
62    $\wedge msgs = \{\}$ 
63    $p \in Participant$  starts the prepare phase by issuing a ballot  $b \in Ballot$ .
64
65  $Prepare(p, b) \triangleq$ 
66    $\wedge state[p][p].maxBal < b$ 
67    $\wedge b \in Bals(p)$ 
68    $\wedge state' = [state \text{ EXCEPT } ![p][p].maxBal = b]$ 
69    $\wedge Send([from \mapsto p, to \mapsto Participant, state \mapsto state'[p]])$ 
70    $q \in Participant$  updates its own state  $state[q]$  according to the actual state  $pp$  of  $p \in Participant$ 
    extracted from a message  $m \in Message$  it receives. This is called by  $OnMessage(q)$ .
    Note:  $pp$  is  $m.state[p]$ ; it may not be equal to  $state[p][p]$  at the time  $UpdateState$  is called.
71
72  $UpdateState(q, p, pp) \triangleq$ 
73    $state' = [state \text{ EXCEPT}$ 
74      $![q][p].maxBal = Max(@, pp.maxBal),$ 
75      $![q][p].maxVVal = Max(@, pp.maxVVal),$ 
76      $![q][p].maxVVal = \text{IF } state[q][p].maxVVal < pp.maxVVal$ 
77        $\text{THEN } pp.maxVVal \text{ ELSE } @,$ 
78      $![q][q].maxBal = Max(@, pp.maxBal),$ 
79      $![q][q].maxVVal = \text{IF } state[q][q].maxVVal \leq pp.maxVVal$ 
80        $\text{THEN } pp.maxVVal \text{ ELSE } @, \text{ make promise}$ 
81      $![q][q].maxVVal = \text{IF } state[q][q].maxVVal \leq pp.maxVVal$ 
82        $\text{THEN } pp.maxVVal \text{ ELSE } @] \text{ accept}$ 
83    $q \in Participant$  receives and processes a message in  $Message$ .
84
85  $OnMessage(q) \triangleq$ 
86    $\exists m \in msgs :$ 
87      $\wedge q \in m.to$ 
88      $\wedge \text{LET } p \triangleq m.from$ 
89      $\text{IN } UpdateState(q, p, m.state[p])$ 
90      $\wedge \text{IF } \vee m.state[q].maxBal < state'[q][q].maxBal$ 
91        $\vee m.state[q].maxVVal < state'[q][q].maxVVal$ 
92      $\text{THEN } Send([from \mapsto q, to \mapsto \{m.from\}, state \mapsto state'[q]])$ 
93      $\text{ELSE UNCHANGED } msgs$ 
94    $p \in Participant$  starts the accept phase by issuing the ballot  $b \in Ballot$  with value  $v \in Value$ .
95
96  $Accept(p, b, v) \triangleq$ 
97    $\wedge b \in Bals(p)$ 
98    $\wedge \exists Q \in Quorum : \forall q \in Q : state[p][q].maxBal = b$ 
99    $\wedge \forall q \in Participant : state[p][q].maxVVal = -1 \text{ free to pick its own value}$ 
100    $\vee \exists q \in Participant : v \text{ is the value with the highest } maxVVal$ 
101    $\wedge state[p][q].maxVVal = v$ 
102    $\wedge \forall r \in Participant : state[p][q].maxVVal \geq state[p][r].maxVVal$ 

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