

**Instruction Set for the problems below, unless otherwise specified, assume that the instruction set includes: (i) R-type instructions (ADD, SUB, AND, ORR), (ii) LDUR, (iii) STUR, (iv) CBZ, and (v) B.**

**Problem 1 (15 points)** When silicon chips are fabricated, defects in materials (e.g., silicon) and manufacturing errors can result in defective circuits. A very common defect is for one signal wire to get broken and always register a logical 0. This is often called a stuck-at-0 fault. Answer the following questions based on **Section 4.4** of the textbook (see slide 19 in Chapter 4 on Canvas).

- (a) [5] Which instructions fail to operate correctly if the MemToReg wire is stuck at 0?
  - (i) Write data in the Register block will always get an address instead of the data stored in the address. STUR will fail.
- (b) [5] Which instructions fail to operate correctly if the ALUSrc wire is stuck at 0?
  - (i) Sign Extension will fail, so adding, subtracting, multiplying, and dividing will fail.
- (c) [5] Which instructions fail to operate correctly if the Reg2Write wire is stuck at 0?
  - (i) STUR and LDUR will fail

**Problem 2 (15 points)** Consider the addition of a multiplier to the CPU shown in Figure 4.23 (slide 19 in Chapter 4 on Canvas). This addition will add 300 ps to the latency of the ALU, but will reduce the number of instructions by 10% (because there will no longer be a need to emulate the multiply instruction). Answer the following questions based on **Section 4.4** of the textbook and the following table for the latencies of the stages of the pipeline.

IF	ID	EX	MEM	WB
250 ps	350 ps	150 ps	300 ps	200 ps

- (a) [5] What is the clock cycle time with and without this improvement?
  - (i) Without: 1500 ps; With: 2250 ps
- (b) [5] What is the speedup achieved by adding this improvement?
  - (i) The speedup is achieved by slowing down the clock rate to lower the total number of instructions,
- (c) [5] What is the slowest the new ALU can be and still result in improved performance?
  - (i) 450 ps

**Problem 3 (30 points)** In this exercise, we examine how pipelining affects the clock cycle time of the processor. Questions in this problem assume that individual stages of the datapath have the latencies shown in Problem 2 above. Also, assume that instructions executed by the processor are broken down as follows:

ALU/Logic	Jump/Branch	LDUR	STUR
45%	20%	20%	15%

Answer the following questions based on **Section 4.5** of the textbook.

- (a) [5] What is the clock cycle time in a pipelined and non-pipelined processor?
  - (i) Pipelined: 400 ps; Non-Pipelined: 1200 ps
- (b) [10] What is the total latency of an LDUR instruction in a pipelined and non-pipelined processor?
  - (i) Total Instructions = I, CPI = 1  
 Pipelined:  $1 \cdot (0.2 \cdot I) \cdot 400 \text{ ps} = 80I \text{ ps}$ ; Non-Pipelined:  $1 \cdot (0.2 \cdot I) \cdot 1200 \text{ ps} = 250I \text{ ps}$
- (c) [10] If we can split one stage of the pipelined datapath into two new stages, each with half the latency of the original stage, which stage would you split and what is the new clock cycle time of the processor?

- (d) [5] Assuming there are no stalls or hazards, what is the utilization of the data memory?

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LOOP:LDUR X10, [X1, #0]
      LDUR X11, [X1, #8]
      SUB X13, X11, X10
      ADD X12, X12, X11
      SUBI X1, X1, #162
      CBNZ X12, LOOP

```

IF ED EX MEM WB

IF   ED   EX   MEM   WB

IF    ED    EX    MEM    WB

IF    ED    EX    MEM   WB

IF    ED    EX    MEM    WB

IF    ED    EX   MEM   WB

IF    ED   EX   MEM   WB

IF ED EX MEM WB

```
ADD X5, X2, X1
LDUR X3, [X5, #4]
LDUR X2, [X2, #0]
ORR X3, X5, X3
STUR X3, [X5, #0]
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(a) [5] If there is no forwarding or hazard detection, insert NOPs to ensure correct execution.

- Nop**

**Nop**

**LDUR X3, [X5, #4]**

**LDUR X2, [X2, #0]**

**ORR X3, X5, X3**

**Nop**

**Nop**

**STUR X3, [X5, #0]**

- (b) [10] Now, change and/or rearrange the code to minimize the number of NOPs needed. You can assume register X7 can be used to hold temporary values in your modified code.

(i) **ADD X5, X2, X1**  
**LDUR X2, [X2, #0]**  
**LDUR X7, [X5, #4]**  
**Nop**  
**Nop**  
**ORR X3, X5, X7**  
**STUR X3, [X5, #0]**

**Problem 6 (15 points)** This problem examines the accuracy of various branch predictors for the following repeating pattern of branch outcomes: T, NT, T, NT, NT. Answer the following questions based on **Section 4.8** of the textbook.

- (a) [5] What is the accuracy of always-taken and always-not-taken predictors for this sequence of branch outcomes?
- (i) Always Taken:  $\frac{2}{5} = 40\%$ ; Always Not-Taken:  $\frac{3}{5} = 60\%$
- (b) [10] What is the accuracy of the 2-bit predictor for this pattern, assuming that the predictor starts off in the bottom left state from Figure 4.62 (predict not taken)? Show the state of the predictor at each step.
- (i) Assuming that the bottom left is encoded 11, bottom right 10, top right 01, and top left 00:  
Predictor: 11 (start, miss), 10 (hit), 01 (hit), 00 (miss), 01 (miss)  
 $\frac{2}{5} = 40\%$  accuracy.