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I pledge my honor that I have abided by the Stevens Honor System.

4. Consider the following algorithm.

```
ALGORITHM Mystery(n)
//Input: A nonnegative integer n
S ← 0
for i ← 1 to n do
    S ← S + i * i
return S
```

a. What does this algorithm compute?

The sum of all perfect squares from 1 to n

b. What is its basic operation?

Reassignment

c. How many times is the basic operation executed?

The basic operation is executed n times

d. What is the efficiency class of this algorithm?

$\Theta(n)$

e. Suggest an improvement, or a better algorithm altogether, and indicate its efficiency class. If you cannot do it, try to prove that, in fact, it cannot be done.

1. Solve the following recurrence relations.

a. $x(n) = x(n - 1) + 5$ for $n > 1$, $x(1) = 0$

$x(n) = 5n - 5$

b. $x(n) = 3x(n - 1)$ for $n > 1$, $x(1) = 4$

$x(n) = 4 \cdot 3^n$

c. $x(n) = x(n - 1) + n$ for $n > 0$, $x(0) = 0$

$x(n) = n(n+1)/2$

d. $x(n) = x(n/2) + n$ for $n > 1$, $x(1) = 1$ (solve for $n = 2^k$)

$x(n) = 2n - 1$

e. $x(n) = x(n/3) + 1$ for $n > 1$, $x(1) = 1$ (solve for $n = 3^k$)

3. Consider the following recursive algorithm for computing the sum of the first n cubes: $S(n) = 1^3 + 2^3 + \dots + n^3$.

```
ALGORITHM S(n)
//Input: A positive integer n
//Output: The sum of the first n cubes
if n = 1 return 1
else return S(n - 1) + n * n * n
```

a. Set up and solve a recurrence relation for the number of times the algorithm's basic operation is executed.

$$\forall n: \{n > 1: S(n) = S(n-1) + n^3, S(1) = 1\}$$

$$S(n) = n^2(n+1)^2/2$$

b. How does this algorithm compare with the straightforward nonrecursive algorithm for computing this sum?

This algorithm has the same time complexity as the straightforward nonrecursive algorithm, $\Theta(n)$ if using a loop, but a time complexity of $\Theta(1)$ if just using the equation above