

- Constructing arguments in propositional logic
- Normal forms

Enumeration

1.	Building Arguments	2
	• Our final topic on logic deals with constructing and validating arguments. We	
	start by giving examples of valid and non-valid arguments and define various	
	concepts that we will need to breakdown an argument.	

### 2. Inference Rules for Propositional Logic

• Breaking down arguments take effort. To simplify things we will collect some standard arguments which we will use, like lego bricks, when working with complicated arguments.

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## 3. Using the Rules of Inference to Build Valid Arguments

- In our final topic in logic, we will use the properties of logical operators to construct a valid argument.
- This is a relatively advanced topic and could be ignored until you are comfortable with the earlier topics in logic.

## Notation

## >Single-line vs Double-line Arrows >

For the purpose of this module the single line arrows (representing the IFTHEN and IFF connectives)

$$\rightarrow$$
 and  $\leftrightarrow$ 

mean the same thing as the corresponding double-line arrow

$$\Rightarrow$$
 and  $\Leftrightarrow$ 

I will use the double-lined arrows in places where I want to treat a complicated proposition as two smaller propositions. For example, I want to think of the proposition

$$(p \rightarrow q) \land \neg q \implies \neg p$$

in terms of the two proposition  $(p \rightarrow q) \land \neg q$  and  $\neg p$ .

## Motivation

Remember the Socrates example when we started Logic.

"All men are mortal.

Socrates is a man.

Therefore, Socrates is mortal."

Here we have two premises:

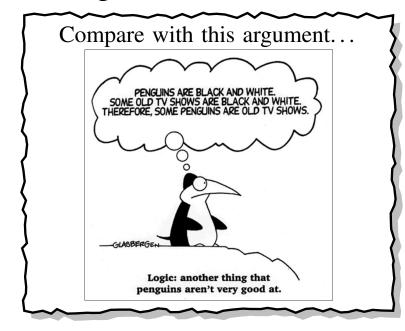
- All men are mortal
- Socrates is a man.

and the conclusion:

• Socrates is mortal.

**Q:** How do we get the conclusion from the premises?

**A:** We construct an argument, a sequence of propositions that follow from the rules of inference until we reach the conclusion.



## Arguments

## Definition 1 (Argument)

A argument in propositional logic is a sequence of propositions. All but the final proposition are called premises. The last statement is the conclusion. The argument is valid if the premises imply the conclusion.

• If the premises are  $p_1, p_2, \dots p_n$  and the conclusion is q then the argument is valid iff

$$(p_1 \wedge p_2 \wedge \cdots \wedge p_n) \rightarrow q$$

is a tautology.

- We could use truth tables to test if an argument is valid construct the above expression, then build the truth table and check the output column.
- Alternatively, we could sequently apply inference rules to arrive at the conclusion.
- Inference rules are simple arguments that will be used to construct more complex argument forms.

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## Detachment (Modus Ponens)

### Argument

$$\begin{array}{c} p \to q \\ \hline p \\ \therefore q \end{array}$$

Corresponding Tautology

$$(p \rightarrow q) \land p \implies q$$

#### Example

Let

p ="It is snowing."

q = "I will study discrete maths."

Then the argument is

"If it is snowing, then I will study discrete maths."

"It is snowing."

Therefore "I will study discrete maths."

# Indirect Reasoning (Modus Tollens)

#### Argument

$$\begin{array}{c}
p \to q \\
\neg q \\
\hline
\vdots \neg p
\end{array}$$

Corresponding Tautology

$$(p \rightarrow q) \land \neg q \implies \neg p$$

#### Example

Let

p ="It is snowing."

q = "I will study discrete maths."

Then the argument is

"If it is snowing, then I will study discrete maths."

"I will not study discrete maths."

Therefore "It is not snowing."

# Chain Rule (Hypothetical Syllogism)

#### Argument

$$\begin{array}{c}
p \to q \\
q \to r \\
\hline
\therefore p \to r
\end{array}$$

Corresponding Tautology

$$(p \rightarrow q) \land (q \rightarrow r) \implies (p \rightarrow r)$$

#### Example

#### Let

p ="It is snowing."

q = "I will study discrete maths."

r = "I will get an A."

### Then the argument is

"If it is snowing, then I will study discrete maths."

"If I will study discrete maths, then I will get an A."

Therefore "If it is snowing, then I will get an A."

# Chain Rule (Hypothetical Syllogism)

#### Argument

$$\begin{array}{c}
p \to q \\
q \to r \\
\hline
\therefore p \to r
\end{array}$$

Corresponding Tautology

$$(p \rightarrow q) \land (q \rightarrow r) \implies (p \rightarrow r)$$

#### Example

Let

p ="It is snowing."

q = "I will study discrete maths."

r = "I will get an A."

Then the argument is

"If it is snowing, then I will study discrete maths."

"If I will study discrete maths, then I will get an A."

Therefore "If it is snowing, then I will get an A."

# Disjunctive Simplification (Disjunctive Syllogism)

#### Argument

$$\begin{array}{ccc}
p \lor q & p \lor q \\
\hline
\neg p & \neg q \\
\hline
\therefore q & \vdots p
\end{array}$$

#### Corresponding Tautology

$$(p \lor q) \land (\neg p) \implies q$$
$$(p \lor q) \land (\neg q) \implies p$$

#### Example

Let

p = "I will study discrete maths."

q = "I will study programming."

Then the argument is

"I will study discrete maths or I will study programming."

"I will not study discrete maths."

Therefore "I will study programming."

# Disjunctive Addition

#### Argument

$$\frac{p}{\therefore p \vee q}$$

Corresponding Tautology

$$p \implies (p \lor q)$$

#### Example

Let

p = "I will study discrete maths."

q = "I will get high."

Then the argument is

"I will study discrete maths."

Therefore "I will study discrete maths or I will get high."

# Conjunctive Simplification

#### Argument

$$\begin{array}{c} p \wedge q \\ \therefore p \end{array} \quad \begin{array}{c} p \wedge q \\ \vdots q \end{array}$$

#### Corresponding Tautology

$$(p \wedge q) \implies p$$

$$(p \land q) \implies p$$

$$(p \land q) \implies q$$

#### Example

Let

p = "I will study discrete maths."

q = "I will get high."

Then the argument is

"I will study discrete maths and I will get high."

Therefore "I will study discrete maths."

## Resolution

#### Argument

$$\begin{array}{c}
\neg p \lor r \\
p \lor q \\
\hline
\therefore q \lor r
\end{array}$$

Corresponding Tautology

$$(\neg p \lor r) \land (p \lor q) \implies (q \lor r)$$

#### Example

#### Let

p = "I will study discrete maths."

p = "I will study programming."

p = "I will study databases."

### Then the argument is

"I will not study discrete maths or I will study programming."

"I will study discrete maths or I will study databases."

Therefore "I will study programming or I will study databases."

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• In our final topic in logic, we will use the properties of logical operators to

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construct a valid argument.

able with the earlier topics in logic.

A valid argument is a sequence of statements. Each statement is either a premise or follows from previous statements by rules of inference. The last statement is called conclusion.

## Example 2

From the single proposition

$$p \land (p \rightarrow q)$$

show that q is a conclusion.

### Method 1

Construct argument using inference rules ...

	Step	Reason
1)	$p \land (p \rightarrow q)$	Premise
2)	p	Conjunctive Simplification from (1)
3)	$p \rightarrow q$	Conjunctive Simplification from (1)
•	q	Detachment (Modus Ponens) from (2) and (3)

### >Method 2>

Construct an expression of the form

(premise 1) 
$$\land$$
 (premise 2)  $\land \cdots \land$  (premise  $n$ )  $\implies$  (conclusion)

and verify that the expression is a tautology (using a truth table).

p	q	$(p \rightarrow q)$	$p \wedge (p \mathop{\rightarrow} q)$	$p \land (p \rightarrow q) \Rightarrow q$
F	F	T	F	$\mathbf{T}$
$\mathbf{F}$	T	${f T}$	${f F}$	${f T}$
$\mathbf{T}$	${f F}$	${f F}$	${f F}$	${f T}$
T	T	$\mathbf{T}$	${f T}$	$oldsymbol{T}$

## Example 3

### With these hypotheses:

- "It is not sunny this afternoon and it is colder than yesterday."
- We will go swimming only if it is sunny."
- "If we do not go swimming, then we will take a canoe trip."
- "If we take a canoe trip, then we will be home by sunset."

Using the inference rules, construct a valid argument for the conclusion:

• "We will be home by sunset."

### General procedure ...

- STEP 1 Choose propositional variables.
- (STEP 2) Translation into propositional logic.
- STEP 3 Construct the valid argument (OR verify related tautology using truth table.)

### STEP 1) Choose propositional variables.

- s = "It is Sunny this afternoon."
- c = "It is Colder than yesterday."
- w = "We will go sWimming"
- t = "We will take a canoe Trip."
- h = "We will be Home by sunset."

### STEP 2 Translation into propositional logic.

#### Premises ...

- ① "It is not sunny this afternoon and it is colder than yesterday."  $\neg s \land c$
- "We will go swimming only if it is sunny."  $w \rightarrow s$
- "If we do not go swimming, then we will take a canoe trip."  $\neg w \rightarrow t$
- "If we take a canoe trip, then we will be home by sunset."  $t \rightarrow h$

### and conclusion

• "We will be home by sunset."

h

STEP 3 Construct the valid argument

(Note the truth table here would have 32 rows)

We have

Premises Conclusion

(a) (b) (c) (d) 

$$\neg s \land c \quad w \rightarrow s \quad \neg w \rightarrow t \quad t \rightarrow h \quad h$$

And our argument is ...

	Step	Reason
1)	$\neg s \wedge c$	Premise (a)
2)	$\neg s$	Conjunctive Simplification from (1)
3)	$w \rightarrow s$	Premise (b)
4)	$\neg w$	Indirect Reasoning (Modus Tollens) from (2) and (3)
5)	$\neg w \rightarrow t$	Premise (c)
6)	t	Detachment (Modus Ponens) from (4) and (5)
7)	$t \rightarrow h$	Premise (d)
• •	h	Detachment (Modus Ponens) from (6) and (7)