


Logic

Discrete Mathematics

Topic 02 — Methods of Mathematical Proof

Lecture 02 — Proof Techniques

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Recurrence
Relations

Set Theory

Outline

- Proof by Cases, Contradiction, Construction, Induction ...

Enumeration

Outline

1. Proof by Contrapositive 2
 - We prove a statement by first switching to the original statement to its contrapositive.
2. Proof by Cases 5
 - We prove a statement by breaking it up into smaller and easier cases, which we prove separately.
3. Proof by Contradiction 12
 - We prove a statement using the process:
 - assume reverse of statement . . .
 - derive conclusions from assumption . . .
 - show conclusions are contradictory . . .
 - hence assumption must be **False**, so original statement is **True**.
4. Proof by Construction 16
 - We prove the existence of something by giving the instructions needed to construct it.
5. Proof by Induction 20
 - Special proof technique used to prove a family of statements,

Proof by Contrapositive

Proof by Contrapositive

In a **proof by contrapositive** argument you prove the contrapositive of the claim rather than the claim itself.

Proof by Contrapositive (Formal Structure)

Given claim

$$P \implies Q$$

the contrapositive (and logically equivalent claim) is

$$\neg Q \implies \neg P$$

- 1 Assume $\neg Q$.
- 2 Demonstrate that $\neg P$ must follow from $\neg Q$.

Please, please, . . . , pretty please don't confuse this with proof by contradiction (covered later).

Example

Example 1

If x^2 is odd then x must be odd.

(by contrapositive)*.

The contrapositive is

If x is even, then x^2 is even.

We assume x is even. Hence we can write $x = 2k$ for some integer k . Now

$$x^2 = (2k)^2 = 4k^2 = 2 \underbrace{(2k^2)}_{\substack{\text{integer} \\ \text{even integer}}}$$

Hence the contrapositive is true, and so is the original statement. □

*The above proof is certainly doable by a direct proof. However, a direct proof requires a cumbersome proof by cases approach.

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Proof by Cases

Proof by Cases

In a **Proof by cases** argument you

- List of of the possible cases and analyse each separately.
- Need to ensure that the cases are exhaustive — cover all possibilities

Proof by Cases (Formal Structure)

Given claim

$$P \implies Q$$

- 1 Show that there exist a number of distinct cases C_1, C_2, \dots such that whenever P is true then at least one of the cases must be true.
- 2 Then, for each case, C , in C_1, C_2, \dots ,
 - 1 Assume case C .
 - 2 Demonstrate that Q must follow from C .

Example 2

Example 2

In a cave you find three boxes. One contains gold, the other two are empty. Each box has imprinted on it a clue as to its contents; the clues are:



A: The gold is not here



B: The gold is not here



C: The gold is in box B

Only one message is true; the other two are false. Which box has the gold?

- Notice that I changed the question to “Which box has the gold?”. I could have left it as “Prove that the gold is in box A.” since, for this problem the two versions are equivalent.

Example 2

II

In a proof by cases, there are three cases based on where the gold is located. In each case we check the truth value of the three messages[†]

Gold is in box A

Gold is in box B

Gold is in box C

[†]You might complain that in the direct proof we did earlier building a truth table is really a proof by cases. You would be correct.

Example 2

In a proof by cases, there are three cases based on where the gold is located. In each case we check the truth value of the three messages[†]

Gold is in box A

A: *“The gold is not here”*

B: *“The gold is not here”*

C: *“The gold is in box B”*

Gold is in box B

A: *“The gold is not here”*

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Gold is in box A

A: *"The gold is not here"* **F**

B: *"The gold is not here"* **T**

C: *"The gold is in box B"* **F**

Gold is in box B

A: *"The gold is not here"*

B: *"The gold is not here"*

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A: “ <i>The gold is not here</i> ”	F	} Exactly one message true? ✓
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Gold is in box B

A: “*The gold is not here*”
B: “*The gold is not here*”
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A: "*The gold is not here*"
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So in order that exactly one message is true, the gold must be in box A.

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Every group of 6 minions includes a group of 3 minions who all know each other or a group of 3 minions who are mutual strangers.



Call one of the minions Bob. There are five others. Either Bob knows three of them, or he does not know three of them.

CASE 1: Bob knows three of the five others ...

Say that Bob knows three of the five others. Of those five minions either there exists two minions who know each other or no two know each other.

CASE 1.1: There exists two minions who know each other ...

Then those two and Bob form a mutually acquainted threesome.

CASE 1.2: No two of the five minions know each other ...

Then any three of the five minions are a mutually unacquainted threesome.

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Example 3

II

CASE 2: *Bob does not know three of the five others ...*

CASE 2.1: No two of the five minions know each other ...

Then those two and Bob form a mutually unacquainted threesome.

CASE 2.2: All pairs within the five minions know each other ...

Then any three of the five minions are a mutually acquainted threesome.

We have covered all possibilities, and in every instance come up either with a mutually acquainted threesome or a mutually unacquainted threesome.

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Examples

- a) Prove that for any integer n , the number $(n^3 - n)$ is even.
- b) Prove that every prime number greater than 3 is either one more or one less than a multiple of 6. Hint. Prove the contrapositive by cases.
- c) Let a, b, c, d be integers. If $a > c$ and $b > c$, then $\max(a, b) - c$ is always positive.

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Proof by Contradiction

Proof by Contradiction

In a **proof by contradiction** argument you:

- Assume the negative of the claim
 - So a universal claim will become an existence claim, and an existence claim will become a universal claim.
- Then show that the assumption leads to a contradiction.

Proof by Contradiction (Formal Structure)

Given claim

$$P \implies Q$$

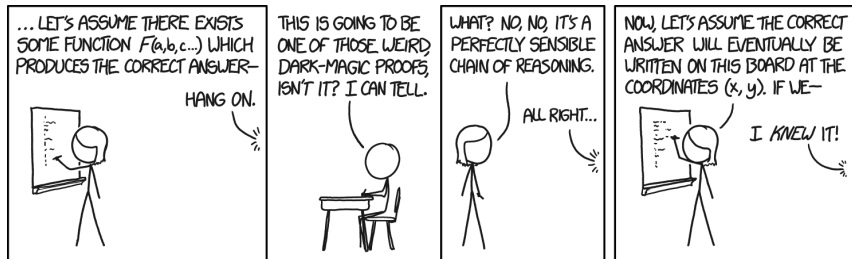
Show that the negative, i.e. $P \implies \neg Q$, leads to a contradiction, by

- 1 Assume P .
- 2 Assume $\neg Q$.
- 3 Use P and $\neg Q$ to demonstrate a contradiction.

Proof by Contradiction

Proofs by contradiction can be tricky, you

- Need to be very clear as to what statement you are assuming in order to generate a contradiction.
- In particular, take case when the statement involves a qualifier.



Examples

- a) Prove that a triangle cannot have more than one right angle.
- b) Prove that the $\sqrt{2}$ is irrational.[§]
- c) Prove that $\log_2(3)$ is irrational.
- d) Let n be an integer. If $3n + 2$ is odd, then n is odd.
- e) Prove that there are an infinite number of primes.[¶]
- f) There are no integers x and y such that $x^2 = 4y + 2$.
- g) The Pigeonhole Principle: If more than n pigeons fly into n pigeon holes, then at least one pigeon hole will contain at least two pigeons. Prove this.

[§]“irrational”= “not rational”. A **rational** number is a number that can be expressed as quotient of two integers p and q which don't have a common factor.

[¶]A **prime** is an integer greater than one with exactly two divisors.

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Proof by Construction

Proof by Construction

In a **proof by construction** argument you:

- Are dealing with an existence claim.
 - Prove existence of an object by actually constructing it.
 - The proof usually involves stating the steps (algorithm) needed to construct the required object.
-
- This type of proof is more powerful than just an existence proof — this not only proves existence but also create an example.
 - Very common in geometry, and graph theory

Example 4

Example 4

If a and b are real numbers and $a \neq 0$, then there exists a unique r such that

$$ar + b = 0$$

Here we have two claims: existence and uniqueness.

Existence (by construction).

Construction

$$\begin{aligned} ar + b &= 0 \\ \implies ar &= -b \\ \xRightarrow{a \neq 0} r &= -b/a \end{aligned}$$

Verify

$$ar + b = a \left(\frac{-b}{a} \right) + b = -b + b = 0$$

Construction: $r = -b/a$



Uniqueness (by contradiction).

Assume there are two values, r and s , with $r \neq s$ satisfying the equation. Then

$$\begin{aligned} ar + b &= 0 = as + b \\ \implies ar + b &= as + b \\ \implies ar &= as \\ \xRightarrow{a \neq 0} r &= s \end{aligned}$$

Contradiction! $\implies r = s$ is unique.



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Examples

- a) Prove that x^n can be computed using only $\log_2(n)$ multiplications when n is a power of 2.

This is a special case of the Montgomery algorithm for computing large integer power quickly — a big deal in cryptography!

- b) Prove that the sum of the first n positive integers equals $n(n+1)/2$

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Proof by Induction

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A **proof by induction** argument, can be applied when Q , the conclusion in the claim $P \Rightarrow Q$, can be represented as a sequence of related claims, Q_1, Q_2, Q_3, \dots . Then we show, that

- the first claim is true, and
- if any claim is true, then the next claim must also be true.

Proof by Induction (Formal Structure)

Given family of claims, where integer n is $1, 2, 3, 4, \dots$,

$$P \Rightarrow Q_n$$

- 1 Assume P , (now we need to prove that all of Q_1, Q_2, Q_3, \dots , are true)
- 2 Prove Q_1 . (the basic/initial step)
- 3 Prove $Q_k \Rightarrow Q_{k+1}$ for arbitrary integer k . (the inductive step)^{||}

^{||}Instead of attacking the problem directly, we only explain how to get a proof for Q_{k+1} when given a proof for Q_k .

Example 5

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Suppose an ATM has only twenty euro and fifty euro bills. You can type in the amount you want, and it will figure out how to divide things up into the proper number of twenty and fifty euro bills.

Prove that the ATM can generate any multiple of 10 euro amount ≥ 40 .

(by induction).

First we define the proposition (or family of propositions)

$$Q_n : \text{ATM can output } 10n \text{ euro} = 20a + 50b \quad \text{where } a \text{ and } b \text{ are nonnegative integers}$$

We want to prove the sequence

$$Q_4, Q_5, Q_6, \dots$$

Note: In this example we did not start at one (and our stride was 10). □

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Q_n : $10n = 20a + 50b$ where a and b are nonnegative integers

the basis step, Q_4

$$10(4) \stackrel{?}{=} 20a + 50b \quad \text{True} \quad a = 2, b = 0$$

i.e., ATM can output forty euro by outputting two twenty euro and no fifty euro notes.

the inductive step

Assume Q_k . It is equivalent to assuming

$$10k = 20a + 50b \quad \text{for some non-negative integers } a \text{ and } b$$

i.e., ATM can output $10k$ euro by outputting only twenty euro and fifty euro notes.

and now we want to prove Q_{k+1} , ie,

$$10(k+1) = 20A + 50B \quad \text{for some non-negative integers } A \text{ and } B$$

i.e., ATM can output $10(k+1)$ euro by outputting only twenty euro and fifty euro notes.

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i.e., ATM can output $10(k+1)$ euro by outputting only twenty euro and fifty euro notes.

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To prove Q_{k+1} we have two cases:

CASE 1: *The ATM used at least one fifty when outputting 10k euro.*

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CASE 1: *The ATM used at least one fifty when outputting $10k$ euro.*

Hence $b > 0$, To get ten more euro out we replace one fifty by three twenties, i.e.,

$$10k = 20a + 50b \implies 10k + 10 = 10(k + 1) = 20(a + 3) + \underbrace{50(b - 1)}_{\text{OK, since } b > 0}$$

CASE 1: *The ATM used no fifty euro notes when outputting $10k$ euro.*

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CASE 1: *The ATM used no fifty euro notes when outputting $10k$ euro.*

Hence $a \geq 2$, since $10k \geq 40$. To get ten more euro out we replace two twenties by one fifty, i.e.,

$$10k = 20a + 50(0) \implies 10k + 10 = 10(k + 1) = \underbrace{20(a - 2)}_{\text{OK, since } a > 2} + 50(1)$$

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To prove Q_{k+1} we have two cases:

CASE 1: *The ATM used at least one fifty when outputting $10k$ euro.*

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$$10k = 20a + 50b \implies 10k + 10 = 10(k + 1) = 20(a + 3) + \underbrace{50(b - 1)}_{\text{OK, since } b > 0}$$

CASE 1: *The ATM used no fifty euro notes when outputting $10k$ euro.*

Hence $a \geq 2$, since $10k \geq 40$. To get ten more euro out we replace two twenties by one fifty, i.e.,

$$10k = 20a + 50(0) \implies 10k + 10 = 10(k + 1) = \underbrace{20(a - 2)}_{\text{OK, since } a > 2} + 50(1)$$

We have proven the two steps required in an induction argument, hence we can conclude the sequence of claims are true.

Examples

- a) For all integers n , prove that $n^2 + 5n + 6$ is even.
- b) Prove that the sum of the first n positive integers equals $n(n + 1)/2$
- c) Prove for integer $n \geq 4$, that $3^n > 2n^2 + 3n$.