

Computing



Tutor: Kevin O'Brien

Tutorial: Maths for Computing



Online Tutorial 4

Chapter 7: Sequence, Series and Proof by Induction

Covered heavily in the last onsite tutorial. Leave until end of this class.

Chapter 8 : Trees

Emphasis of Today's Class



Trees

 A lot of concepts and definitions follows from Chapter 5: Introduction to Graph Theory

Syllabus

- Properties of Trees
- Rooted Trees and Binary Trees
- Binary Search Trees



1) Characteristics of a Tree

A tree is a connected graph that contains no cycles. A tree has no loops and no multiple edges. All trees are simple graphs.

2) Path Graphs

A tree that contains only vertices of degree one or two is called a *path graph*. The length of a path graph is the number of edges in it.

3) Number of Edges

(Theorem 3.3) Let T be a tree with n vertices. Then T has n − 1 edges.



4) Spanning Subgraphs

The graph H is a **subgraph** of a graph G if H's vertices are a subset of the G's vertex, its edges are a subset of the edge set of G, and each edge of H has the same end-vertices in G and H.

H is called a **spanning subgraph** of G if the vertices of H are the same as the vertices of G.

5) Spanning Trees

If H is a spanning subgraph which is also a tree, then H is said to be a spanning tree of G. (G does not need to be a tree)



Spanning Trees (Figure)



Trees: Properties of Trees 2008 Zone A Q9

Question 9

(a) A graph with 5 vertices: a, b, c, d, e has the following adjacency list:

a:b,e

b:a, c, d

c:b, d

d:b, c, e

e:d, a.

- Draw this graph, G.
- (ii) Draw a spanning tree of G.
- (iii) Draw all the non-isomorphic spanning trees of G and call this set S.
- (iv) How many non-isomorphic trees can be created by adding a new vertex and edge to the trees in S.
 [6]



Trees: Properties of Trees 2008 Zone A Q9

a:b, e

 $b:a,\ c,\ d$

c:b, d

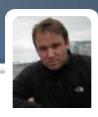
d:b, c, e

 $e:d,\ a.$



Trees: Properties of Trees 2008 Zone A Q9

- Part IV
- Examiner's Commentaries: Then it is a question of adding another vertex and edge to each of these in all possible places and finally eliminating the isomorphic ones to do part (iv).
- Simpler Exercise (2006 Q9)
 - (b) (i) Draw the 3 non-isomorphic trees on 5 vertices.
 - (ii)Draw, on a separate diagram, all the non-isomorphic trees on 6 vertices, by adding a vertex to copies of the trees you have drawn or otherwise.
 [6]



Trees: Properties of Trees 2006 Zone A Q9



Question 7 (a) (i) What properties must a graph have in order for it to be a tree?

- (ii) Say, with reason, whether or not it is possible to construct a tree with degree sequence 4, 3, 3, 1, 1.
- (iii) Say, with reason, whether it is possible to construct a tree with degree sequence 4, 3, 2, 2, 1.
- (iv) What properties must a graph have in order for it to be a binary tree?
 [5]





Trees: Properties of Trees (2005)

(c) Let G be the simple graph with vertex set V(G) = {a, b, c, d, e} and adjacency matrix

$$\mathbf{A} = \begin{bmatrix} a & b & c & d & e \\ a & 0 & 1 & 0 & 0 & 0 \\ b & 1 & 0 & 1 & 0 & 1 \\ c & 0 & 1 & 0 & 1 & 0 \\ d & 0 & 0 & 1 & 0 & 1 \\ e & 0 & 1 & 0 & 1 & 0 \end{bmatrix}.$$

- (i) What do the numbers on the leading diagonal of this matrix tell you about the graph?
- (ii) Say how the number of edges in G is related to the entries in the adjacency matrix A and calculate this number.
- (iii) Draw G.
- (iv) Find a spanning tree T₁ for G and give its degree sequence.
- (v) Find a spanning tree T₂ for G which is **not** isomorphic to T₁ and give a reason why it is not isomorphic. [7]







Trees: Rooted Trees and Binary Trees

Terminology (Page 37)

- Root
- Nodes
- Key
- · Children and Parents
- Ancestors and Descendants
- Height



Binary Search Tree

A Binary Search Tree is a binary tree in symmetric order

Symmetric order means that:

- every node has a key (or number)
- every node's key is
 - larger than all keys in its left subtree
 - smaller than all keys in its right subtree

The root *r* is the record

$$\# [(1+N)/2].$$



Binary Search Tree

If the first record in the subtree is #a and the last record is #b, then the **root of the subtree is**

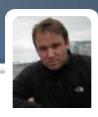
$$\# \lfloor (a+b)/2 \rfloor$$
.



Trees: Binary Search Trees

The height h of a binary search tree with N records stored at internal nodes is

$$h = \lceil \log_2 (N+1) \rceil.$$



Trees: Binary Search Trees (2004)

- (b) A binary search tree is designed to store an ordered list of 3000 records at its internal nodes.
 - (i) Find which record is stored at the root (level 0) of the tree and at each of the nodes at level 1.
 - (ii) What is the height of the tree?
 - (iii) What is the maximum number of comparisons needed in order to find an existing record in the tree?
 [5]



Trees: Binary Search Trees



Trees: Binary Search Trees (2006)

Question 9

- (a) A binary search tree is designed to store an ordered list of 10000 records numbered 1,2,3,...10000 at its internal nodes.
 - (i) Draw levels 0, 1 and 2 of this tree showing which number record is stored at the root and at each of the nodes at level 1 and 2, making it clear which records are at each level.
 - (ii) What is the maximum number of comparisons that would have to be made in order to locate an existing record from the list of 10000?
 [4]



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Maths for Computing





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