

## Chapter 6: Digraphs and Relations

Chapter 1 D ' h d R l t' Summary Digraphs; Relations, using digraphs to illustrate relations, equivalence relations, partial orders, relations and cartesian products.

### Digraphs

Some problems require that we direct the edges of a graph, from one endpoint to the other, in order to express the relationship between the vertices. A graph in which every edge has a direction assigned to it is called a **digraph** (an abbreviation of directed graph).

The directed edges are often called **arcs**.

**Example 1.1** Suppose we want to model a plan for traffic flow in part of a city, where some streets allow traffic flow in only one direction.

For this model, a digraph would be appropriate. We would represent each road junction by a vertex; a one-way street from junction  $u$  to junction  $v$ , by an arc directed from  $u$  to  $v$  and a street allowing traffic flow in both directions between junctions  $x$  and  $y$ , by two arcs, one directed from  $x$  to  $y$  and the other from  $y$  to  $x$ .

An example is given in Figure 1.17 where the direction of each arc is indicated by an arrow. In a digraph we define directed paths and directed cycles as in graphs, but with the added condition that arcs must be used in their correct direction.

We say a digraph  $D$  is strongly connected if for every pair of vertices  $x$  and  $y$  of  $D$ , there is a directed path from  $x$  to  $y$  and a directed path from  $y$  to  $x$  in  $D$ . Example 1.2 The digraph in Figure 1.1 contains a directed path  $P = v_1 v_2 v_3$ ; and

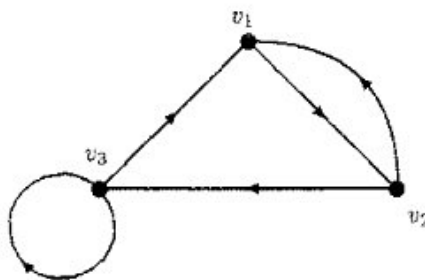


Figure 1:

a directed cycle  $C : v_1 v_2 v_3 v_1$ . It is strongly connected. Digraphs can be used to illustrate a variety of "one-way" relationships such as predator-prey models in ecology, the scheduling of tasks in a manufacturing process, family trees, flow diagrams for a computer program and many others.

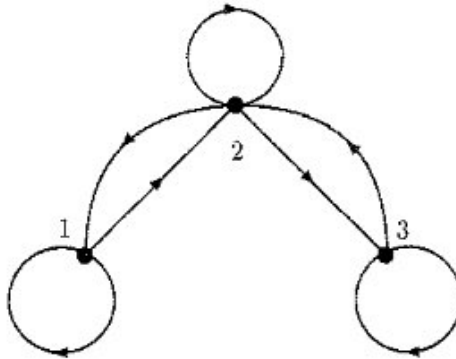


Figure 2:

## Relations

Let  $S$  be a set. A relation  $R$ , on  $S$  is a rule which compares any two elements  $x, y \in S$  and tells us either that  $x$  is related to  $y$  or that  $x$  is not related to  $y$ .

We write  $xRy$  to mean  $x$  is related to  $y$  under the relation  $R$ . We are already familiar with many examples of relations defined in society: for example we could let  $S$  be the set of all people in London and say that two people  $x$  and  $y$  are related if  $x$  is a parent of  $y$ . We have also already seen several examples of relations in mathematics, for example " $=$ ", " $<$ ", and " $\leq$ " are all relations on the set of integers.

A different example is the following:

**Example 1.3** Let  $S = \{2, 3\}$ . We define a relation  $R$  on  $S$  by saying that  $x$  is related to  $y$  if  $|xy| = 1$ ; for all  $x, y \in S$ . Thus  $2R2$  but  $1$  is not related to  $3$ .

## Using digraphs to illustrate relations

Given a relation  $R$  on a set  $S$ ; we can model  $R$  by defining the digraph  $D$  with  $V(D) = S$  in which, for any two vertices  $x$  and  $y$  there is an arc in  $D$  from  $x$  to  $y$  if and only if  $xRy$ . We shall call  $D$  the relationship digraph corresponding to  $R$ .

**Example 1.4** Figure 1.2 gives the relationship digraph corresponding to the relation defined in Example 1.3.

Conversely, given a digraph  $D$ , we can define a relation  $R$  on the set  $S = V(D)$  by saying  $xRy$  if and only if there is an arc in  $D$  from  $x$  to  $y$ .

**Example 1.5** The digraph given in Figure 1.1 defines the relation  $R$  on the set  $S = \{v, w, x, y, z\}$  given by  $zRx, yRx, vRy, wRy, wRx, wRz, xRy, xRz, yRz$ .

## Equivalence Relations

We can see from Example 1.5 that the definition of a relation on a set may be rather arbitrary. Many relations which occur in practical situations however have a more “*regular structure*”.

**Definition 1.6** Let  $\mathcal{R}$  be a relation defined on a set  $S$ . We say that  $\mathcal{R}$  is:

- reflexive if for all  $x \in S$ , we have  $x\mathcal{R}x$ .
- symmetric if for all  $x, y \in S$  such that  $x\mathcal{R}y$ , we have  $y\mathcal{R}x$ .
- transitive if for all  $x, y, z \in S$  such that  $x\mathcal{R}y$  and  $y\mathcal{R}z$ , we have  $x\mathcal{R}z$ .

In terms of the relationship digraph  $D$  of  $\mathcal{R}$ , it can be seen that:

- $\mathcal{R}$  is reflexive if all vertices of  $D$  are in a directed loop.
- $\mathcal{R}$  is symmetric if all arcs  $xy$  in  $D$  are in a directed cycle of length two.
- $\mathcal{R}$  is transitive if for all directed paths of length two  $P : xyz$  of  $D$ , we have

an arc  $xz$ ; and for all directed cycles of length two  $C : xyx$  of  $D$ , we have a loop  $xx$  and a loop  $yy$ .

Using Figure 1.2 we deduce that the relation defined in Example 1.3 is reflexive and symmetric but not transitive. Similarly, using Figure 1.1 we deduce that the relation defined in Example 1.5 is not reflexive, symmetric or transitive.

**Example 1.7** The relation  $\leq$  defined on  $\mathbb{Z}$  is reflexive and transitive but not symmetric. The relation  $<$  defined on  $\mathbb{Z}$  is transitive but not reflexive or symmetric.