

PROGRAM 5

Implement 0/1 Knapsack problem using dynamic programming.

ALGORITHM

→ Algorithm:

Function Knapsack (weight [], value [], n, w):

// Create a 2D table dp[n+1][w+1]

create array dp[0...n][0...w]

// Initialize base case: zero items or zero capacity

for i from 0 to n:

for w from 0 to w:

If $i == 0$ or $w == 0$:

dp[i][w] = 0

// Fill the DP table

for i from 1 to n:

for w from 1 to w:

If weight[i-1] <= w:

// Either take the item or don't

dp[i][w] = max (value[i-1] + dp[i-1]

(w - weight[i-1]), dp[i-1](w))

Else:

// cannot include item i-1

dp[i][w] = dp[i-1][w]

Return dp[n][w] // This is the maximum value achieved

CODE

```
#include <stdio.h>
```

```
int max(int a, int b) {  
    return (a > b) ? a : b;  
}
```

```
// Function to solve knapsack problem using DP
```

```
int knapsack(int W, int wt[], int val[], int n) {  
    int dp[n + 1][W + 1];
```

```
    // Build table dp[][] in bottom-up manner
```

```
    for (int i = 0; i <= n; i++) {  
        for (int w = 0; w <= W; w++) {  
            if (i == 0 || w == 0)  
                dp[i][w] = 0;  
            else if (wt[i - 1] <= w)  
                dp[i][w] = max(val[i - 1] + dp[i - 1][w - wt[i - 1]], dp[i - 1][w]);  
            else  
                dp[i][w] = dp[i - 1][w];  
        }  
    }
```

```
    return dp[n][W];  
}
```

```
int main() {  
    int n, W;  
  
    printf("Enter number of items: ");
```

```

scanf("%d", &n);

int wt[n], val[n];

printf("Enter weights of items:\n");
for (int i = 0; i < n; i++)
    scanf("%d", &wt[i]);

printf("Enter values of items:\n");
for (int i = 0; i < n; i++)
    scanf("%d", &val[i]);

printf("Enter capacity of knapsack: ");
scanf("%d", &W);

int maxValue = knapsack(W, wt, val, n);

printf("Maximum value that can be obtained = %d\n", maxValue);

return 0;
}

```

OUTPUT

```

Enter number of items: 4
Enter weights of items:
2 3 4 5
Enter values of items:
3 4 5 6
Enter capacity of knapsack: 5
Maximum value that can be obtained = 7

```

TRACING

Tracing:

Item	wt	value
1	2	3
2	3	4
3	4	5
4	5	6

capacity = 5

S1: Initialize to 0, max val = 0

S2: Add $i = 1$ ($w = 2, v = 3$)
 $cap \geq 2$, take $i = 1$ & max val = 3

S3: Add $i = 2$ ($w = 3, v = 4$)
 * check if we can take P item 2 alone $\rightarrow v = 4$
 * check if we can take $i_1 + i_2$ together $2 + 3 = 5 \leq capacity$
 * value if both take $3 + 4 = 7 \rightarrow > 4 \& 3$.
 * max val = 7.

S4: Add $i = 3$ ($w = 4, v = 5$)
 * i_1 alone $\rightarrow v = 5$

* $i_1 + i_3 \rightarrow wt = 2 + 4 = 6 > capacity - no$
 * $i_2 + i_3 \rightarrow wt = 3 + 4 = 7 > cap \rightarrow no$
 So max value stays the same 7.

S5: Add $i = 4$ ($w = 5, v = 6$)
 * $i = 4$ alone $\rightarrow v = 6$.
 $i_1 + i_4 \rightarrow wt = 2 + 5 = 7 < capacity - no$
 $i_2 + i_4 \rightarrow wt = 3 + 5 = 8 < capacity - no$
 max value remains same

So Final maximum value = 7

Items chosen 1 & 2 ($wt\ 2 + 3 = 5$).
 (val $3 + 4 = 7$)

W	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	3	3	3	3
2	0	0	3	4	4	7
3	0	0	3	4	5	7
4	0	0	3	4	5	7

Backtracking

* $i=4, w=5$

* $dp[w][5] = dp[3][5] = 7$

4 not included.

$w=5 \quad i=3$

* $dp[3][5] = 7$ } equal

$dp[2][5] = 7$

3 not included

$w=5 \quad i=2$

* $dp[2][5] = 7$ } no equal

$dp[1][5] = 3$

2 included

$w = 5 - \text{weight}(2-1) = 5 - 3 = 2$

$i=1$

* $dp[1][2] = 3 \neq dp[0][2] = 0$

Item 1 was included.

$w = 2 - w(1-1) = 2 - 2 = 0$

\Rightarrow Total weight

$= 2 + 3 = 5$

Total value

$= 3 + 4 = 7$

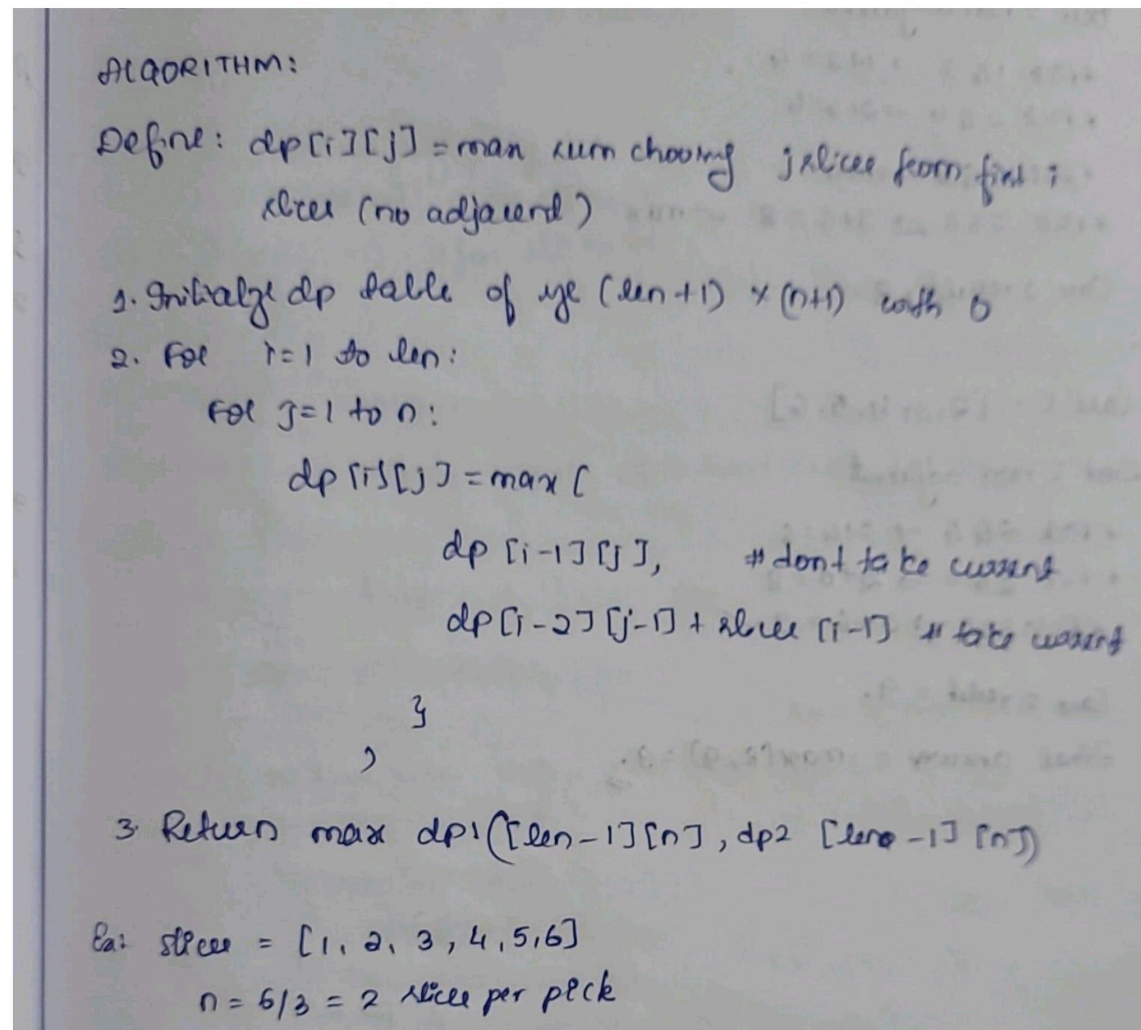
Shi
17/8/25

6/6

LEETCODE 4

PIZZA WITH 3n SLICES

ALGORITHM



CODE

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
#include <string.h>
```

```
int max(int a, int b) {  
    return (a > b) ? a : b;  
}
```

```

// Function to compute the max sum using DP (like 0/1 Knapsack)
int calculate(int* slices, int start, int end, int n) {
    int len = end - start + 1;
    int dp[len + 2][n + 1];
    memset(dp, 0, sizeof(dp));

    for (int i = len - 1; i >= 0; i--) {
        for (int j = 1; j <= n; j++) {
            int take = slices[start + i] + dp[i + 2][j - 1]; // pick slice
            int skip = dp[i + 1][j]; // skip slice
            dp[i][j] = max(take, skip);
        }
    }
    return dp[0][n];
}

int maxSizeSlices(int* slices, int slicesSize) {
    int n = slicesSize / 3;
    int max1 = calculate(slices, 0, slicesSize - 2, n); // exclude last
    int max2 = calculate(slices, 1, slicesSize - 1, n); // exclude first
    return max(max1, max2);
}

```

OUTPUT

Accepted Runtime: 0 ms

• Case 1 • Case 2

Input

```
slices =  
[1,2,3,4,5,6]
```

Output

```
10
```

Expected

```
10
```

♥ Contribute a testcase

Accepted Runtime: 0 ms

• Case 1 • Case 2

Input

```
slices =  
[8,9,8,6,1,1]
```

Output

```
16
```

Expected

```
16
```

♥ Contribute a testcase

TRACING

Tracing:
Case 1: Use slice [1,2,3,4,5] (exclude last)
Case 2: Use slice [2,3,4,5,6] (exclude first).

Case 1: [1, 2, 3, 4, 5]

Best 2 non adjacent

* Pick 1 & 3 $\rightarrow 1+3=4$

* Pick 2 & 4 $\rightarrow 2+4=6$

* Pick 1 & 4 $\rightarrow 1+4=5$

* Pick 3 & 5 $\rightarrow 3+5=8$ ✓ max.

Case 1 result = 8.

Case 2: [2, 3, 4, 5, 6]

Best 2 non adjacent

* Pick 2 & 4 $\rightarrow 2+4=6$

* Pick 2 & 5 $\rightarrow 2+5=7$

* Pick 3 & 5 $\rightarrow 3+5=8$ ✓

Case 2 result = 9.

Final answer = $\max(8, 9) = 9$.