

# Formal Verification of Distributed Algorithms using Distributed PlusCal

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#### **Motivations** Contributions

## Formal Methods for Distributed Algorithms

- Distributed Algorithms are prone to deadlocks and race conditions.
- Formal verification methods have been employed successfully to model the system and its properties and then verify its correctness.

### **Modeling Distributed Algorithms**

- TLA+<sup>1</sup> is a formal language used to describe algorithms, it is used to specify and verify complicated algorithms concisely.
- TLA+ relies on mathematical logic, PlusCal<sup>2</sup>-was designed as an algorithm language with a more familiar syntax that can be translated into TLA+ specifications.
- PlusCal<sup>2</sup> lacks primitives for communication between processes which makes modeling distributed algorithms difficult.

### References

[1]Lamport, Leslie. (2000). Specifying Concurrent Systems with TLA+.

[2] Lamport, Leslie. (2009). The PlusCal Algorithm Language. 5684. 36-60. 10.1007/978-3-642-03466-4 2.

- We propose Distributed PlusCal, an extension of PlusCal that introduces constructs that can overcome some of the limitation of PlusCal.
- Distributed PlusCal offers primitives that aid in modeling distributed algorithms.
- We provide a backward compatible translator that translates from Distributed PlusCal and PlusCal to TLA+.

#### Distributed PlusCal

- Distributed PlusCal is a PlusCal extension that offers primitives that aid in modeling distributed algorithms such as
  - Sub-processes
  - Communication channels:
    - Unordered channels with the syntax

channel (identifier)[(dimension1,...dimensionN)]

FIFO channels with the syntax

fifo (identifier)[(dimension1,...dimensionN)]

- Channels support the following operations
- 1. send

```
2. broadcast
3. multicast
4. receive
5. clear
```

```
--algorithm TwoPhaseCommit {
\* unordered channels
channels agt[Agent], coord;
  fair process (a \in Agent)
  variable aState = "unknown"; { \* sub-process
    a1: if (aState = "unknown") {
        with(st \in {"accept", "refuse"})
          aState := st;
          send(coord, [type |-> st, agent |-> self]);};};
    a2: await(aState \in {"commit", "abort"})
  } { \* sub-process
    a3:await (aState # "unknown");
       receive(agt[self], aState);
  fair process (c = Coord)
 variables cState = "unknown",
            commits = {},
            msg = {}; { \ \ }*sub-process
   c1: await(cState \in {"commit", "abort"});
      broadcast(agt, [ag \in Agent|-> cState]);
  } { \* sub-process
     c2:while (cState \notin {"abort", "commit"}) {
        receive(coord, msg);
            if (msg.type = "refuse") {
                cState := "abort";
            else if (msg.type = "accept") {
                commits := commits \cup {msg.agent};
                 if (commits = Agent) {
                    cState := "commit";}
```