

Formal Verification of Distributed Algorithms using Distributed PlusCal

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Erasmus Mundus Joint Masters Degree MSc in Advanced Systems Dependability







Motivations Contributions

Formal Methods for Distributed Algorithms

- Distributed Algorithms are prone to deadlocks and race conditions.
- Formal verification methods have been employed successfully to model the system and its properties and then verify its correctness.

Modeling Distributed Algorithms

- TLA+¹ is a formal language used to describe algorithms, it is used to specify and verify complicated algorithms concisely.
- TLA+ relies on mathematical logic, PlusCal²-was designed as an algorithm language with a more familiar syntax that can be translated into TLA+ specifications.
- PlusCal² lacks primitives for communication between processes which makes modeling distributed algorithms difficult.

References

[1]Lamport, Leslie. (2000). Specifying Concurrent Systems with TLA+.

TLA+. [2] Lamport, Leslie. (2009). The PlusCal Algorithm Language. 5684. 36-60. 10.1007/978-3-642-03466-4 2.

Contributions

- We propose Distributed PlusCal, an extension of PlusCal that introduces constructs aid in modeling distributed algorithms.
- We provide a backward compatible translator that translates from Distributed PlusCal and PlusCal to TLA+.

Distributed PlusCal

- Distributed PlusCal is a PlusCal extension that offers primitives that aid in modeling distributed algorithms such as
 - Sub-processes that take the general form process(...) { * sub-process

```
...
} { \* sub-process
...
}
```

- Communication channels :
 - Unordered channels with the syntax
 channel (identifier)[(dimension1,...dimensionN)]
 - FIFO channels with the syntax
 fifo (identifier)[(dimension1,...dimensionN)]
 - Channels support the following operations
 - send
 - 2. broadcast
 - 3. multicast
 - 4. receive
 - 5. clear

```
--algorithm TwoPhaseCommit {
\* unordered channels
channels agt[Agent], coord;
  fair process (a \in Agent)
  variable aState = "unknown"; { \* sub-process
    a1: if (aState = "unknown") {
        with(st \in {"accept", "refuse"})
          aState := st;
          send(coord, [type |-> st, agent |-> self]);};};
    a2: await(aState \in {"commit", "abort"})
 } { \* sub-process
    a3:await (aState # "unknown");
       receive(agt[self], aState);
  fair process (c = Coord)
 variables cState = "unknown",
            commits = {},
            msg = {}; { \ \ }*sub-process
   c1: await(cState \in {"commit", "abort"});
      broadcast(agt, [ag \in Agent|-> cState]);
  } { \* sub-process
     c2:while (cState \notin {"abort", "commit"}) {
        receive(coord, msg);
            if (msg.type = "refuse") {
                cState := "abort";
            else if (msg.type = "accept") {
                commits := commits \cup {msg.agent};
                 if (commits = Agent) {
                    cState := "commit";}
```