

# SEARCHING & SORTING

## [LEVEL-3]

Date 20/09/23...

# Questions on search space :-

Q1. We have 2 numbers as input and we have to divide them using Binary Search. Return the integer value of quotient.

Ex:- i/p  $\rightarrow \frac{29}{7}$  o/p  $\rightarrow 4$

Divisor  $\overline{)$  Dividend  $\overline{)$  Quotient

$\hookrightarrow$  Quotient \* Divisor + Remainder = Dividend Remainder

$\text{Quotient} * \text{Divisor} \leq \text{Dividend}$

Search space for possible answers will be -

0  $\longleftarrow$   $\longrightarrow$  Dividend

Ex:- i/p  $\rightarrow \frac{29}{7} \rightarrow$  dividend

0  $\longleftarrow$   $\xrightarrow{\text{mid}}$  29  
                    14

$S=0, e=29, \text{mid}=14$

$\hookrightarrow$  Is this a possible ans?

Quotient x divisor  $\leq$  Dividend

$$14 \times 7 = 98 \neq 29$$

$98 > 29 \rightarrow$  false

$\hookrightarrow$  Left me jao  $\rightarrow e = \text{mid} - 1$

$$e = 14 - 1 = 13$$

0  $\longleftarrow$   $\xrightarrow{\text{mid}}$  13  
                    6

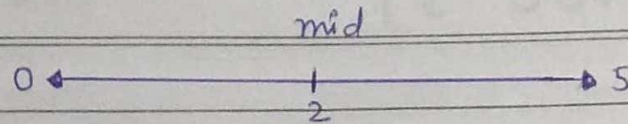
$S=0, e=13, \text{mid} = \frac{0+13}{2} = 6$

$Q \times \text{Divisor} \leq \text{dividend}$

$$6 \times 7 = 42$$

$42 > 29 \rightarrow$  left  $\rightarrow e = \text{mid} - 1 = 6 - 1 = 5$

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$$s=0, e=5, \text{mid} = \frac{0+5}{2} = 2$$

$$Q \times \text{divisor} \leq \text{dividend}$$

$$2 \times 7 \leq 29$$

$$14 \leq 29 \rightarrow \text{Valid ans}$$

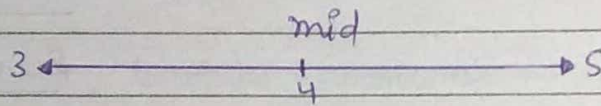
$$\text{ans} = 2$$

↳ ans store

↳ Right

$$\rightarrow s = \text{mid} + 1$$

$$s = 2 + 1 = 3$$



$$s=3, e=5, \text{mid} = \frac{3+5}{2} = 4$$

$$Q \times \text{divisor} \leq \text{dividend}$$

$$4 \times 7 \leq 29$$

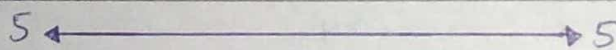
$$28 \leq 29 \rightarrow \text{Valid Ans}$$

$$\text{ans} = 4$$

↳ store ans

↳ Right  $\rightarrow s = \text{mid} + 1$

$$s = 4 + 1 = 5$$



$$s=5, e=5, \text{mid}=5$$

$$Q \times \text{divisor} \leq \text{dividend}$$

$$5 \times 7 \leq 29$$

$$35 \leq 29 \rightarrow \text{False}$$

↳ Left  $\rightarrow e = \text{mid} - 1$

$$e = 5 - 1 = 4$$

$$s=5, e=4$$

$$s > e$$

↳ STOPPING

CONDITION



## Conditions -

- Quotient  $\times$  divisor == dividend
  - ↳ Q is final Ans
- Quotient  $\times$  divisor < dividend
  - ↳ ans store
  - ↳ right
- Quotient  $\times$  divisor > dividend
  - ↳ left me jao

## Code:-

```

int getQuotient (int divisor , int dividend) {
    int s = 0; int e = dividend;
    int mid = s + (e - s) / 2;
    int ans = -1;
    while (s <= e) {
        if (mid * divisor == dividend) {
            return mid;
        }
        else if (mid * divisor < dividend) {
            // store ans
            // right
            ans = mid;
            s = mid + 1;
        }
        else {
            // left
            e = mid - 1;
        }
        // update mid
        mid = s + (e - s) / 2;
    }
    return ans;
}

```

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```
int main() {
    int dividend = 29;
    int divisor = 7;
    int ans = getQuotient(abs(divisor), abs(dividend));
    // now we need to decide the sign either +ve or -ve
    if ((divisor > 0 && dividend < 0) || (divisor < 0 && dividend > 0)) {
        ans = 0 - ans;
    }
    cout << "final ans is : " << ans << endl;
}
```

Output:-  
Final ans is : 4

Imp

Q2. Binary Search on nearly sorted array.

Sorted Array	10	20	30	40	50	60	70
	0	1	2	3	4	5	6
Can be (i-1)	-1	0✓	1	2	3✓	4	5✓
(i)	0	1	2✓	3	4	5	6✓
Found at (i+1)	1✓	2	3	4✓	5	6✓	7

If a no. is at  $i^{\text{th}}$  index  
in sorted  
array

Then it can be  
present at any  
position among  
them in nearly  
sorted array

- $(i-1)$  index
- $(i)$  index
- $(i+1)$  index

Nearly Sorted Array	20	10	30	50	40	70	60
	0	1	2	3	4	5	6



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## Binary Search in Normal sorted array

Conditions

→ if (arr[mid] == target)  
return mid

→ if (target > arr[mid])  
→ Right

else  
→ Left

## Binary Search in Nearly Sorted array

But here  
ans can be  
at  
(mid-1)  
(mid)  
(mid+1)

if (arr[mid] == target)  
return mid;

if (arr[mid-1] == target)  
return mid-1;

if (arr[mid+1] == target)  
return mid+1;

if (target > arr[mid])  
→ Right → **catch** → Yaha  
else  
→ Left → **catch** → galtri  
Karung  
mai

Dry Run:-

	0	1	2	3	4	5	6
arr	20	10	30	50	40	70	60
s	↑						↑
e							↑

target = 70

s = 0, e = 6, mid = 3

arr[mid-1] == 70 → 30 == 70 → F

arr[mid] == 70 → 50 == 70 → F

arr[mid+1] == 70 → 40 == 70 → F

if (target > arr[mid])

70 > 50 → Right → s = mid + 1

s = mid + 2 (catch)

To skip the already  
checked number  
and same is applicable to  
Left wala case. ∴ e = mid - 2  
and not use mid - 1 here.

	5	6
arr	70	60
s	↑	
e		↑

s = 5, e = 6, mid = 5

arr[mid-1] Not exist

arr[mid] == 70 → 70 == 70 → True

return mid

So, s = 3 + 2 = 5

Spiral



Code:-

```

int searchNearlySorted (int arr[], int n, int target) {
    int s=0; int e= n-1, int mid = s+(e-s)/2;
    while ( s <= e) {
        if ( arr[mid-1] == target)
            return mid-1;
        if ( arr[mid] == target)
            return mid;
        if ( arr[mid+1] == target)
            return mid+1;

        if ( target > arr[mid]) {
            // right
            s = mid + 2; // Yaha catch hai s=mid+2
        }
        else {
            // left
            e = mid - 2; // Yaha catch hai e=mid-2
        }
        mid = s + (e-s)/2;
    }
    return -1;
}

```

```

int main() {

```

```

    int arr[] = {10, 20, 30, 50, 40, 70, 60}; int n = 7; int target = 10;
    int targetIndex = searchNearlySorted (arr, n, target);
    if (targetIndex == -1)
        cout << "Target Not found";
    else

```

```

        cout << "Target found at index : " << targetIndex;
    }

```

Output:-  
Target found at index: 0



**Imp Hard**

Q Find the odd occurring element.

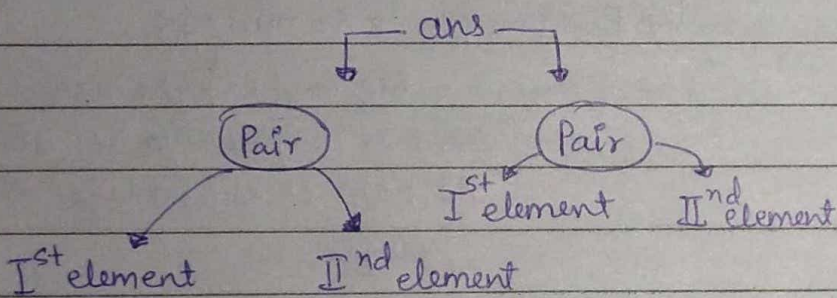
- ↳ All elements are occurring even no. of times except one
  - ↳ Even no. of elements are in pairs and pairs are not repeated.
  - ↳ Ek baar me koi bhi no. 2 se jada baar nahi aa skta
- find that element which occurs odd times.

Left								Right						
1	1	5	5	2	2	3	3	2	4	4	6	6	7	7
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
E	O	E	O	E	O	E	O	E	O	E	O	E	O	E

Types of B.S Questions

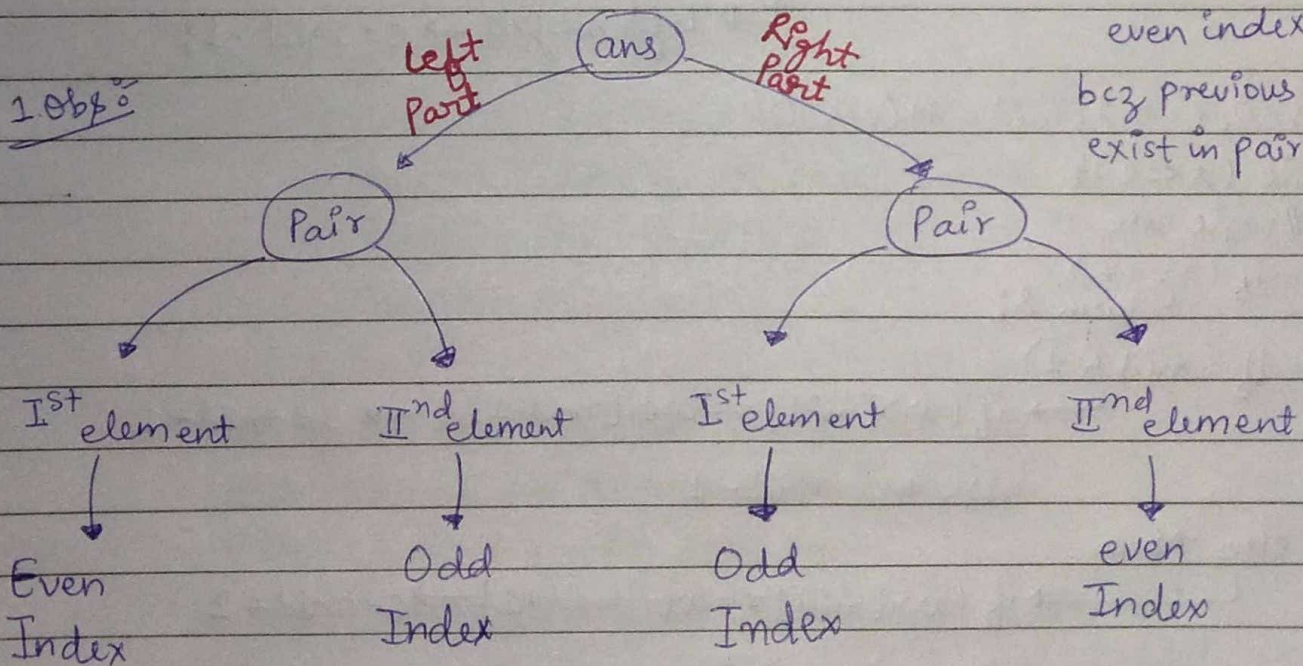
- ① Classical
  - ② Search space
  - ③ Predicate function
  - ④ Index pe observation
- Karke logi banana

Observation:-



2 obs:- single element  
 $s \rightarrow 2 \leftarrow e$   
 if  $(s=e)$   
 return s

1 Obs:-



3 obs:- We get ans at even index always bcz previous elements exist in pair of 2



logic

if (mid % 2 == 0)  $\Rightarrow$  Even index

abhi ans k left me hai

 $\rightarrow$  if (arr[mid] == arr[mid+1]) $\rightarrow$  Right me jao $\rightarrow$  s = mid + 2already +1 check  
koiya that's why +2else  $\rightarrow$  Ya toh ans k right me hai/  
ya ans pehi hai $\rightarrow$  e = mid;

why?

ans lost na hojye  
or left bhichle jaye that's why.

{Peak in mountain}

if (mid % 2 == 1)

 $\rightarrow$  ans k left me hai abhi $\rightarrow$  if (arr[mid] == arr[mid-1]) $\rightarrow$  Right me jao  $\rightarrow$  s = mid + 1;Bcz mid k piche wale  
element se compare kiya hai isliye  
+1 koke ek element aage  
jyenge.

else

 $\rightarrow$  left me jao  $\rightarrow$  e = mid - 1;

int s = 0, e = n - 1, mid = s + (e - s) / 2

while (s &lt;= e) {

// single case

if (s == e)  
return s;

if (mid &amp; 1)

 $\rightarrow$  if (arr[mid] == arr[mid-1])  $\rightarrow$  s = mid + 1else  $\rightarrow$  e = mid - 1else  $\rightarrow$  Even $\rightarrow$  if (arr[mid] == arr[mid+1])  $\rightarrow$  s = mid + 2else  $\rightarrow$  e = mid;  $\rightarrow$  ans lost na ho



Code:-

```

int findOddOccurringElement ( int arr[], int n) {
    int s = 0; int e = n-1; int mid = s + (e-s)/2;
    while (s <= e) {
        if (s == e) {
            return s;
        }

        // check mid → even or odd
        if (mid & 1) { // mid & 1 → True → odd no.
            if (mid + 1 < n && arr[mid] == arr[mid+1]) {
                // left me jao
                e = mid - 1;
            } else {
                // ans ke left side me hai to
                // right me jao
                s = mid + 1;
            }
        }

        // even no
        if (mid + 1 < n && arr[mid] == arr[mid+1]) {
            // mid+1 already checked so move towards
            // right using mid+2
            s = mid + 2;
        } else {
            // Ya toh mai right part par Khada hu
            // Ya toh mai ans ke uppr Khada hu
            // That's why e = mid krna hai
            // Kyunki e = mid - 1 se ans lost ho skta hai
            e = mid;
        }
    }
}

```



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```
mid = s + (e - s) / 2;
```

```
}
```

```
return -1;
```

```
}
```

```
int main() {
```

```
int arr[] = { 20, 20, 5, 5, 3, 3, 1 };
```

```
int n = 7;
```

```
int ans = findOddOccurringElement(arr, n);
```

```
cout << "Final ans is: " << arr[ans];
```

```
}
```

Output:-

Final ans is 1