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CHAPTER 2

Modeling with Linear Programming

Set 2.1a

- (a) $x_2 - x_1 \geq 1$ or $-x_1 + x_2 \geq 1$
- (b) $x_1 + 2x_2 \geq 3$ and $x_1 + 2x_2 \leq 6$
- (c) $x_2 \geq x_1$ or $x_1 - x_2 \leq 0$
- (d) $x_1 + x_2 \geq 3$
- (e) $\frac{x_2}{x_1 + x_2} \leq .5$ or $.5x_1 - .5x_2 \geq 0$

1

(a) $(x_1, x_2) = (1, 4)$

2

$$x_1, x_2 \geq 0$$

$$\begin{aligned} 6x_1 + 4x_2 &= 22 &< 24 \\ 1x_1 + 2x_2 &= 9 &\neq 6 \end{aligned}$$

infeasible

(b) $(x_1, x_2) = (2, 2)$

$$\begin{aligned} x_1, x_2 &\geq 0 \\ 6x_1 + 4x_2 &= 20 &< 24 \\ 1x_1 + 2x_2 &= 6 &= 6 \\ -1x_1 + 1x_2 &= 0 &< 1 \\ 1x_1 &= 2 &= 2 \end{aligned}$$

} feasible

$$Z = 5x_1 + 4x_2 = \$18$$

(c) $(x_1, x_2) = (3, 1.5)$

$$\begin{aligned} x_1, x_2 &\geq 0 \\ 6x_1 + 4x_2 &= 24 &= 24 \\ 1x_1 + 2x_2 &= 6 &= 6 \\ -1x_1 + 1x_2 &= -1.5 &< 1 \\ 1x_1 &= 1.5 &< 2 \end{aligned}$$

} feasible

$$Z = 5x_1 + 4x_2 = \$21$$

(d) $(x_1, x_2) = (2, 1)$

$$\begin{aligned} x_1, x_2 &\geq 0 \\ 6x_1 + 4x_2 &= 16 &< 24 \\ 1x_1 + 2x_2 &= 4 &< 6 \\ -1x_1 + 1x_2 &= -1 &< 1 \\ 1x_1 &= 1 &< 2 \end{aligned}$$

} feasible

$$Z = 5x_1 + 4x_2 = \$14$$

(e) $(x_1, x_2) = (2, -1)$
 $x_1 \geq 0, x_2 < 0$, infeasible

Conclusion: (c) gives the best feasible solution

3

$$(x_1, x_2) = (2, 2)$$

Let S_1 and S_2 be the unused daily amounts of M1 and M2.

$$\text{For M1: } S_1 = 24 - (6x_1 + 4x_2) = 4 \text{ tons/day}$$

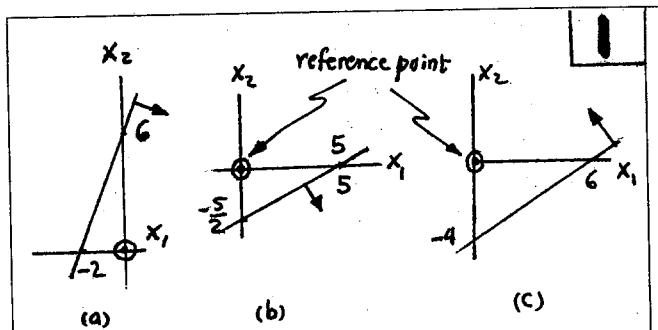
$$\begin{aligned} \text{For M2: } S_2 &= 6 - (x_1 + 2x_2) \\ &= 6 - (2 + 2x_2) = 0 \text{ tons/day} \end{aligned}$$

Quantity discount results in the **4**
following nonlinear objective function:

$$Z = \begin{cases} 5x_1 + 4x_2, & 0 \leq x_1 \leq 2 \\ 4.5x_1 + 4x_2, & x_1 > 2 \end{cases}$$

The situation cannot be treated as a linear program. Nonlinearity can be accounted for in this case using mixed integer programming (chapter 9).

Set 2.2a



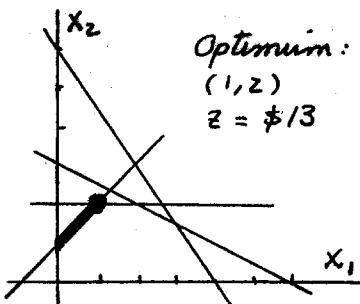
(a), (b)

(c)

(d)

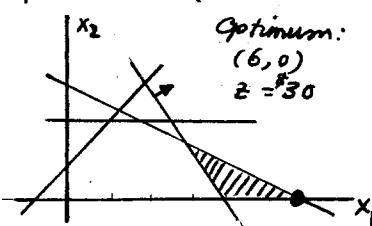
(e)

$$(c) -x_1 + x_2 = 1$$



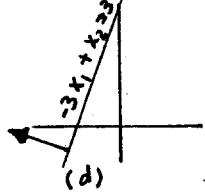
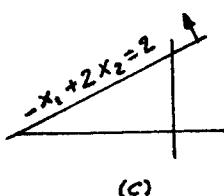
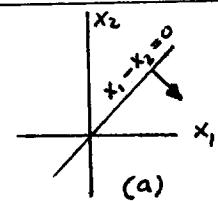
Optimum:
(1, 2)
 $Z = \$13$

$$(d) 6x_1 + 4x_2 \geq 24$$



Optimum:
(6, 0)
 $Z = \$30$

(e) No feasible space



x_1 = daily units of product 1
 x_2 = daily units of product 2

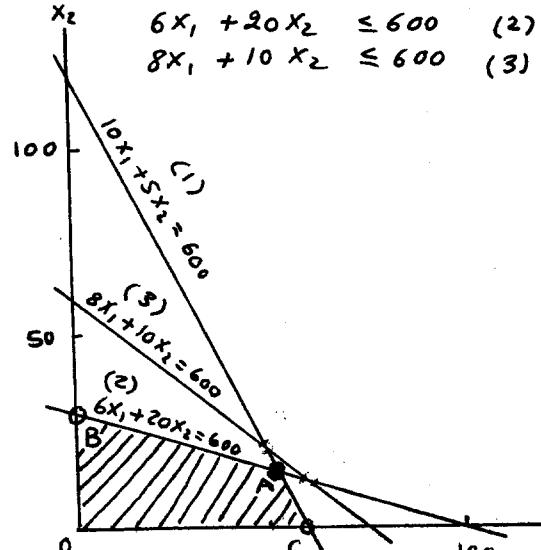
$$\text{Maximize } Z = 2x_1 + 3x_2$$

s.t.

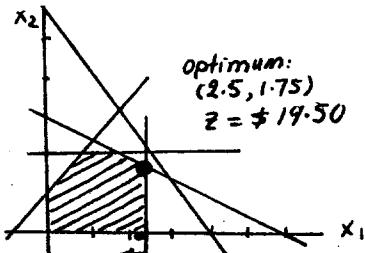
$$10x_1 + 5x_2 \leq 600 \quad (1)$$

$$6x_1 + 20x_2 \leq 600 \quad (2)$$

$$8x_1 + 10x_2 \leq 600 \quad (3)$$

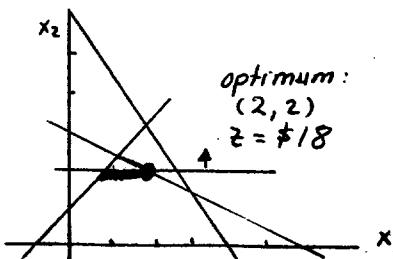


$$(a) x_1 \leq 2.5$$



Optimum:
(2.5, 1.75)
 $Z = \$19.50$

$$(b) x_2 \geq 2$$



Optimum:
(2, 2)
 $Z = \$18$

continued...

2-3

Optimum occurs at A:

$$x_1 = 52.94$$

$$x_2 = 14.12$$

$$Z = \$148.24$$

Set 2.2a

x_1 = number of units of A
 x_2 = number of units of B

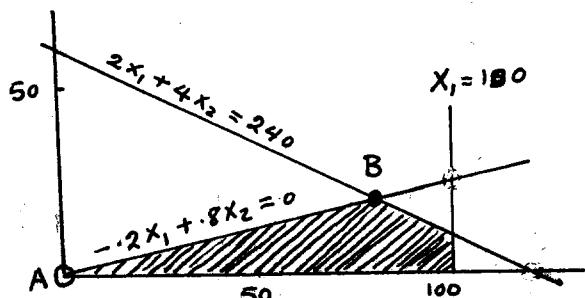
$$\text{Maximize } Z = 20x_1 + 50x_2$$

$$\frac{x_1}{x_1+x_2} \geq .8 \quad \text{or} \quad -2x_1 + 8x_2 \leq 0$$

$$x_1 \leq 100$$

$$2x_1 + 4x_2 \leq 240$$

$$x_1, x_2 \geq 0$$



Optimal occurs at B:

$$x_1 = 80 \text{ units}$$

$$x_2 = 20 \text{ units}$$

$$Z = \$2,600$$

5

x_1 = \$ invested in A

x_2 = \$ invested in B

$$\text{Maximize } Z = .05x_1 + .08x_2$$

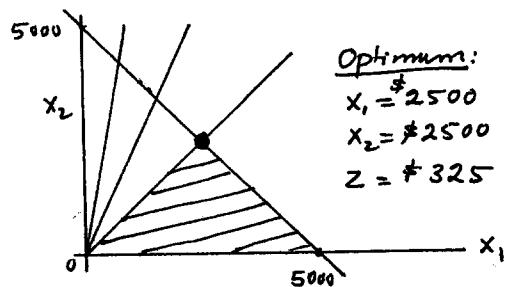
$$\text{s.t. } x_1 \geq .25(x_1 + x_2)$$

$$x_2 \leq .5(x_1 + x_2)$$

$$x_1 \geq .5x_2$$

$$x_1 + x_2 \leq 5000$$

$$x_1, x_2 \geq 0$$



7

Optimum:
 $x_1 = 2500$
 $x_2 = \$2500$
 $Z = \$325$

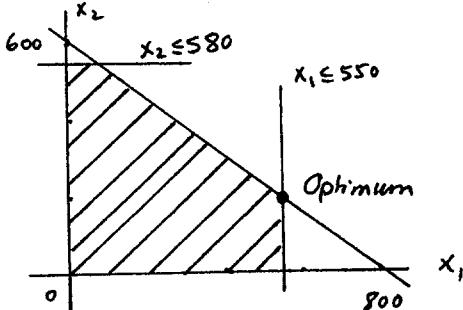
x_1 = number of sheets/day

x_2 = number of bars/day

$$\text{Maximize } Z = 40x_1 + 35x_2$$

$$\text{s.t. } \frac{x_1}{800} + \frac{x_2}{600} \leq 1$$

$$0 \leq x_1 \leq 550, \quad 0 \leq x_2 \leq 580$$



Optimum solution:

$$x_1 = 550 \text{ sheets}$$

$$x_2 = 187.13 \text{ bars}$$

$$Z = \$28,549.40$$

6

x_1 = number of practical courses

x_2 = number of humanistic courses

$$\text{Maximize } Z = 1500x_1 + 1000x_2$$

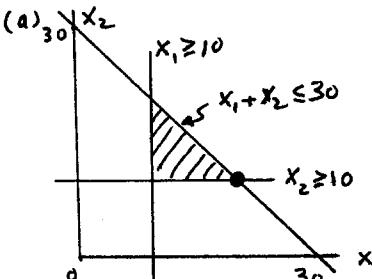
s.t.

$$x_1 + x_2 \leq 30$$

$$x_1 \geq 10$$

$$x_2 \geq 10$$

$$x_1, x_2 \geq 0$$



Optimum:
 $x_1 = 20$
 $x_2 = 10$
 $Z = \$40,000$

(b) Change $x_1 + x_2 \leq 30$ to $x_1 + x_2 \leq 31$

$$\text{Optimum } Z = \$41,500$$

$$\Delta Z = \$41,500 - \$40,000 = \$1500$$

Conclusion: Any additional course will be of the practical type.

Set 2.2a

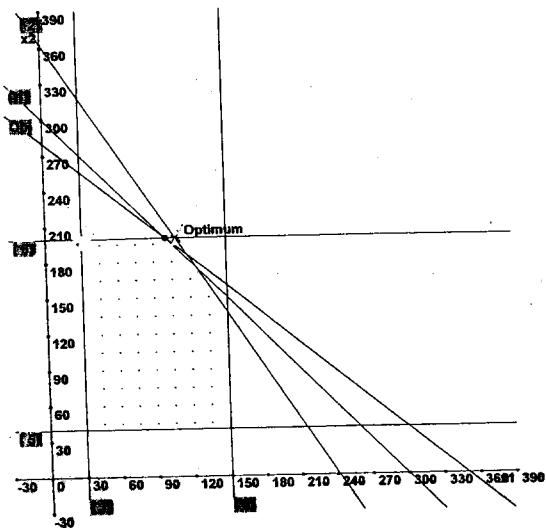
x_1 = units of solution A
 x_2 = units of solution B

$$\text{Maximize } Z = 8x_1 + 10x_2$$

subject to

$$\begin{aligned} .5x_1 + .5x_2 &\leq 150 \\ .6x_1 + .4x_2 &\leq 145 \\ x_1 &\geq 30 \\ x_1 &\leq 150 \\ x_2 &\geq 40 \\ x_2 &\leq 200 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Summary of Optimal Solution:
 Objective Value = 2800.00
 $x_1 = 100.00$
 $x_2 = 200.00$



x_1 = nbr. of grano boxes
 x_2 = nbr. of wheatie boxes

$$\text{Maximize } Z = x_1 + 1.35x_2$$

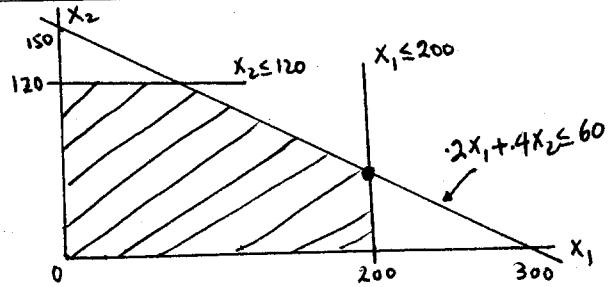
$$\text{s.t. } .2x_1 + .4x_2 \leq 60$$

$$x_1 \leq 200$$

$$x_2 \leq 120$$

$$x_1, x_2 \geq 0$$

9



Optimum: $x_1 = 200, x_2 = 50, Z = \267.50
Area allocation: 67% grano, 33% wheatie

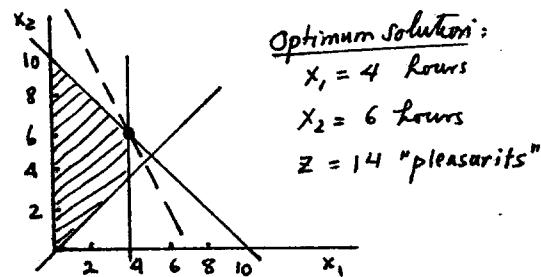
11

x_1 = play hours per day
 x_2 = work hours per day

$$\text{Maximize } Z = 2x_1 + x_2$$

s.t.

$$\begin{aligned} x_1 + x_2 &\leq 10 \\ x_1 - x_2 &\leq 0 \\ x_1 &\leq 4 \\ x_1, x_2 &\geq 0 \end{aligned}$$



Optimum solution:

$$x_1 = 4 \text{ hours}$$

$$x_2 = 6 \text{ hours}$$

$$Z = 14 \text{ "pleasurits"}$$

12

x_1 = Daily nbr. of type 1 hat

x_2 = Daily nbr. of type 2 hat

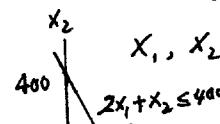
$$\text{Maximize } Z = 8x_1 + 5x_2$$

$$\text{s.t. } 2x_1 + x_2 \leq 400$$

$$x_1 \leq 150$$

$$x_2 \leq 200$$

$$x_1, x_2 \geq 0$$



Optimum:

$$x_1 = 100 \text{ type 1}$$

$$x_2 = 200 \text{ type 2}$$

$$Z = \$1800$$

10

continued...

continued...

Set 2.2a

$$x_1 = \text{radio minutes}$$

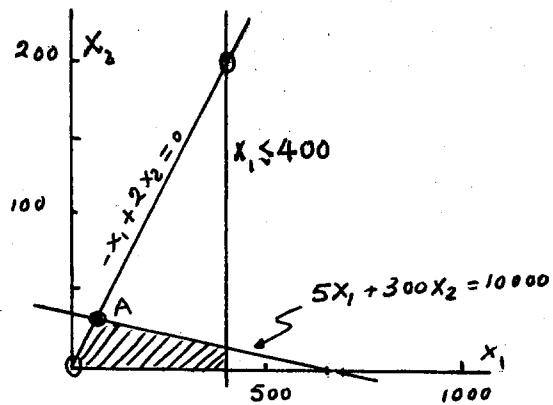
$$x_2 = \text{TV minutes}$$

$$\text{Maximize } Z = x_1 + 25x_2$$

$$\text{s.t. } 15x_1 + 300x_2 \leq 10,000$$

$$\frac{x_1}{x_2} \geq 2 \text{ or } -x_1 + 2x_2 \leq 0$$

$$x_1 \leq 400, x_1, x_2 \geq 0$$



Optimum occurs at A:

$$x_1 = 60.61 \text{ minutes}$$

$$x_2 = 30.3 \text{ minutes}$$

$$Z = 818.18$$

$$x_1 = \text{tons of C}_1 \text{ consumed per hour}$$

$$x_2 = \text{tons of C}_2 \text{ consumed per hour}$$

$$\text{Maximize } Z = 12000x_1 + 9000x_2$$

s.t.

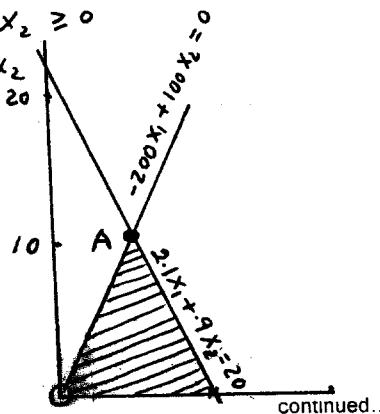
$$1800x_1 + 2100x_2 \leq 2000(x_1 + x_2)$$

or

$$-200x_1 + 100x_2 \leq 0$$

$$2.1x_1 + .9x_2 \leq 20$$

$$x_1, x_2 \geq 0$$



13

(a) Optimum occurs at A:

$$x_1 = 5.128 \text{ tons per hour}$$

$$x_2 = 10.256 \text{ tons per hour}$$

$$Z = 153,846 \text{ lb of Steam}$$

$$\text{Optimal ratio} = \frac{5.128}{10.256} = .5$$

$$(b) 2.1x_1 + .9x_2 \leq (20+1) = 21$$

$$\text{Optimum } Z = 161538 \text{ lb of Steam}$$

$$\Delta Z = 161538 - 153846 = 7692 \text{ lb}$$

15

$$x_1 = \text{Nbr. of radio commercials beyond the first}$$

$$x_2 = \text{Nbr. of TV ads beyond the first}$$

$$\text{Maximize } Z = 2000x_1 + 3000x_2 + 5000 + 2000$$

$$\text{s.t. } 300(x_1+1) + 2000(x_2+1) \leq 20,000$$

$$300(x_1+1) \leq .8 \times 20,000$$

$$2000(x_2+1) \leq .8 \times 20,000$$

$$x_1, x_2 \geq 0$$

or

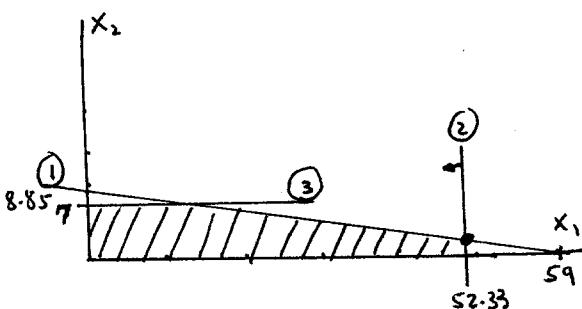
$$\text{Maximize } Z = 2000x_1 + 3000x_2 + 7000$$

$$\text{s.t. } 300x_1 + 2000x_2 \leq 17700 \quad (1)$$

$$300x_1 \leq 15700 \quad (2)$$

$$2000x_2 \leq 14000 \quad (3)$$

$$x_1, x_2 \geq 0$$



Optimum solution:

$$\text{Radio commercials} = 52.33 + 1 = 53.33$$

$$\text{TV ads} = 1 + 1 = 2$$

$$Z = 107666.67 + 7000 = 114666.67$$

continued...

2-6

Set 2.2a

x_1 = number of shirts per hour
 x_2 = number of blouses per hour

$$\text{Maximize } Z = 8x_1 + 12x_2$$

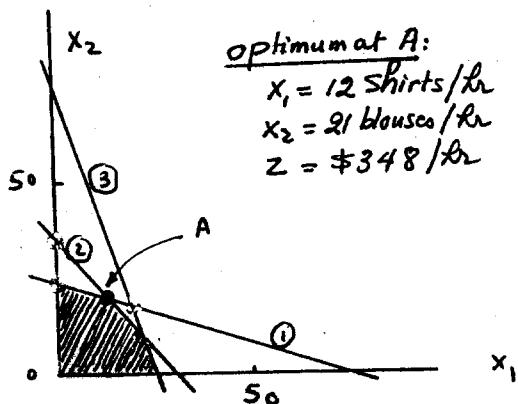
s.t.

$$20x_1 + 60x_2 \leq 25 \times 60 = 1500 \quad (1)$$

$$70x_1 + 60x_2 \leq 35 \times 60 = 2100 \quad (2)$$

$$12x_1 + 4x_2 \leq 5 \times 60 = 300 \quad (3)$$

$$x_1, x_2 \geq 0$$



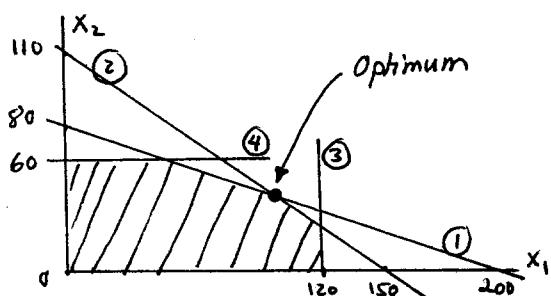
x_1 = Nbr. of desks per day
 x_2 = Nbr. of chairs per day

$$\text{Maximize } Z = 50x_1 + 100x_2$$

$$\frac{x_1}{200} + \frac{x_2}{80} \leq 1 \quad (1)$$

$$\frac{x_1}{150} + \frac{x_2}{110} \leq 1 \quad (2)$$

$$x_1 \leq 120, x_2 \leq 60 \quad (3,4)$$



Optimum:

$$x_1 = 90 \text{ desks}$$

$$x_2 = 44 \text{ chairs}$$

$$Z = \$8900$$

16

x_1 = number of HiFi 1 units
 x_2 = number of HiFi 2 units

18

Constraints:

$$6x_1 + 4x_2 \leq 480 \times 9 = 432$$

$$5x_1 + 5x_2 \leq 480 \times 8.6 = 412.8$$

$$4x_1 + 6x_2 \leq 480 \times 8.8 = 422.4$$

or

$$6x_1 + 4x_2 + s_1 = 432$$

$$5x_1 + 5x_2 + s_2 = 412.8$$

$$4x_1 + 6x_2 + s_3 = 422.4$$

Objective function:

$$\text{Minimize } S_1 + S_2 + S_3 = 1267.2 - 15x_1 - 15x_2$$

$$\text{Thus, } \min S_1 + S_2 + S_3 \equiv \max 15x_1 + 15x_2$$

$$\text{Maximize } Z = 15x_1 + 15x_2$$

s.t.

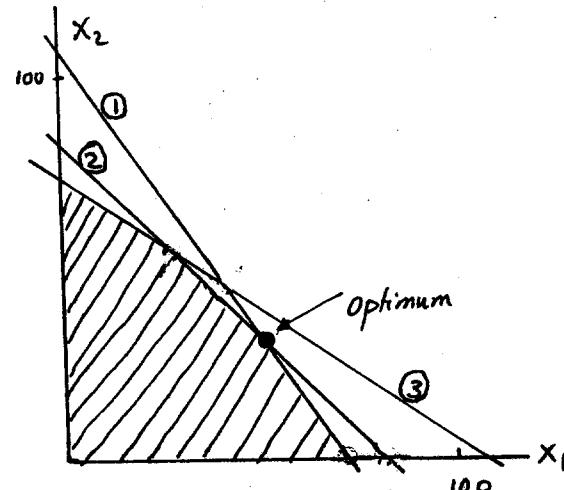
$$6x_1 + 4x_2 \leq 432 \quad (1)$$

$$5x_1 + 5x_2 \leq 412.8 \quad (2)$$

$$4x_1 + 6x_2 \leq 422.4 \quad (3)$$

$$x_1, x_2 \geq 0$$

17



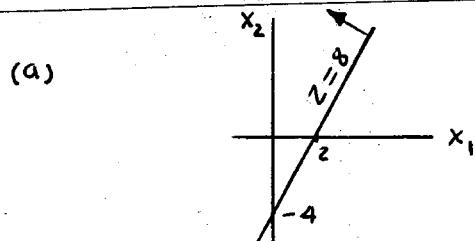
Optimum: (Problem has alternative optima)

$$x_1 = 50.88 \text{ units}$$

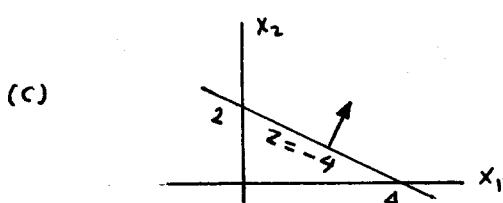
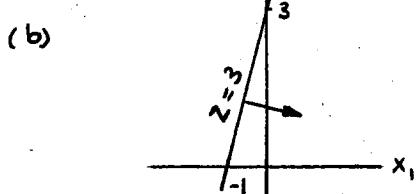
$$x_2 = 31.68 \text{ units}$$

$$Z = 1238.4 \text{ minutes}$$

Set 2.2b

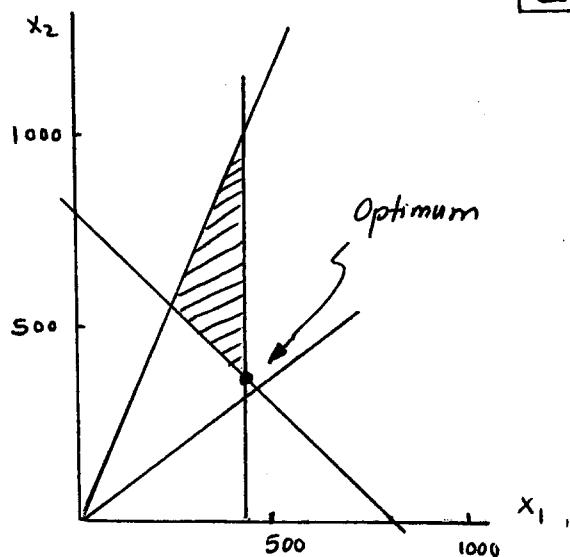


1



Additional constraint: $x_1 \leq 450$

2



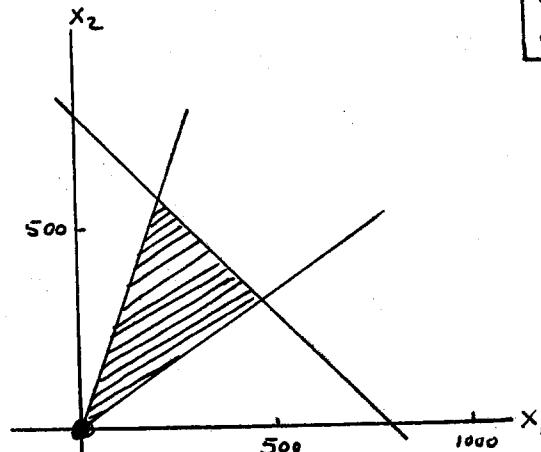
Optimum Solution:

$$x_1 = 450 \text{ lb}$$

$$x_2 = 350 \text{ lb}$$

$$Z = \$450$$

continued...



3

Optimum: $x_1 = 0, x_2 = 0, Z = 0$, which is nonsensical

4

x_1 = number of hours/week in store 1
 x_2 = number of hours/week in store 2

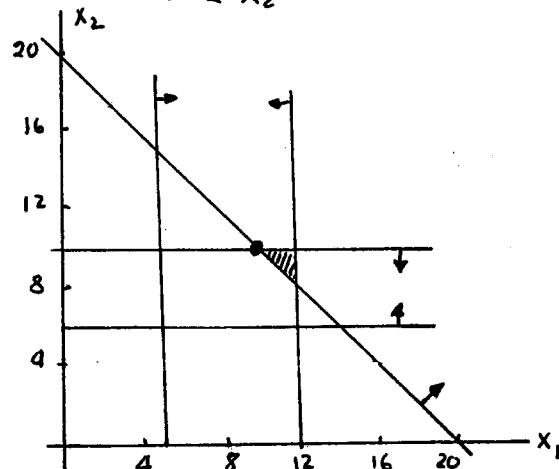
$$\text{Minimize } Z = 8x_1 + 6x_2$$

s.t.

$$x_1 + x_2 \geq 20$$

$$5 \leq x_1 \leq 12$$

$$6 \leq x_2 \leq 10$$



Optimum:

$$x_1 = 10 \text{ hours}$$

$$x_2 = 10 \text{ hours}$$

$$Z = 140 \text{ stress index}$$

continued...

Let

$$x_1 = 10^3 \text{ bbl/day from Iran}$$

$$x_2 = 10^3 \text{ bbl/day from Dubai}$$

$$\text{Refinery capacity} = x_1 + x_2 \quad 10^3 \text{ bbl/day}$$

$$\text{Minimize } Z = x_1 + x_2$$

Subject to

$$x_1 \geq .4(x_1 + x_2)$$

$$\text{or } -.6x_1 + .4x_2 \leq 0$$

$$.2x_1 + .1x_2 \geq 14$$

$$.25x_1 + .6x_2 \geq 30$$

$$.1x_1 + .15x_2 \geq 10$$

$$.15x_1 + .1x_2 \geq 8$$

$$x_1, x_2 \geq 0$$

Optimum solution from TORA:

5 Let

$$x_1 = 10^3 \text{ $ invested in blue chip stock}$$

$$x_2 = 10^3 \text{ $ invested in high-tech stocks}$$

$$\text{Minimize } Z = x_1 + x_2$$

Subject to

$$.1x_1 + .25x_2 \geq 10$$

$$.6x_1 - .4x_2 \geq 0$$

$$x_1, x_2 \geq 0$$

TORA optimum solution:

LINEAR PROGRAMMING - GRAPHICAL SOLUTION

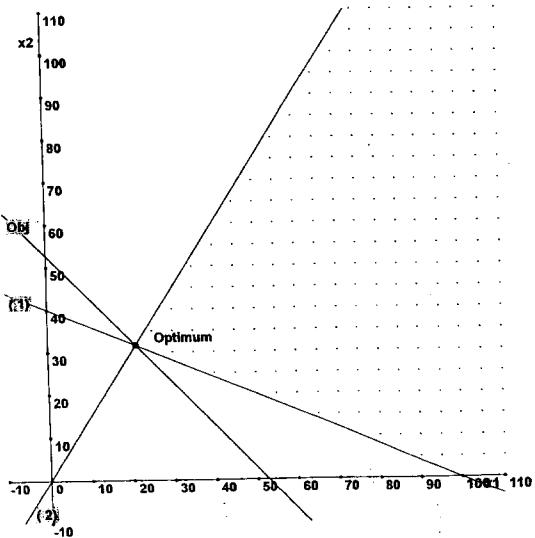
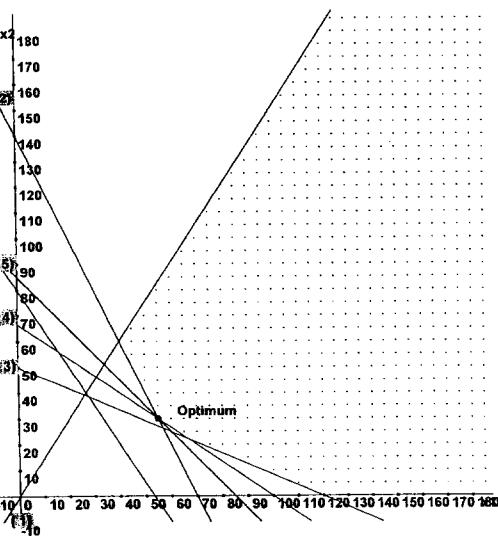
Title: diet problem

Summary of Optimal Solution:

Objective Value = 52.63

x1 = 21.05

x2 = 31.58



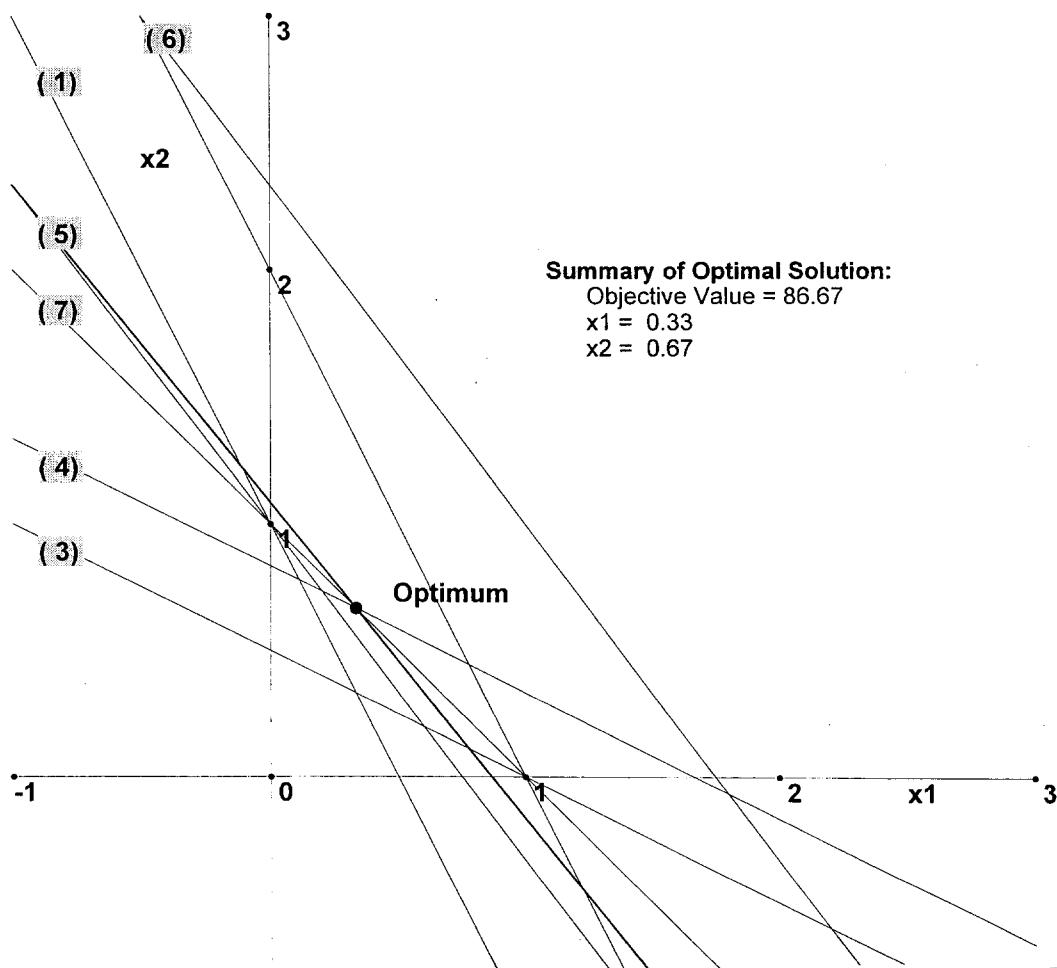
Set 2.2b

7

x_1 = Ratio of scrap A in alloy

x_2 = Ratio of scrap B in alloy

	x_1	x_2	
Minimize	100.00	80.00	
Subject to			
(1)	0.06	0.03	>= 0.03
(2)	0.06	0.03	<= 0.06
(3)	0.03	0.06	>= 0.03
(4)	0.03	0.06	<= 0.05
(5)	0.04	0.03	>= 0.03
(6)	0.04	0.03	<= 0.07
(7)	1.00	1.00	= 1.00



Set 2.3a

X_e = Nbr. of efficiency apartments
 X_d = Nbr. of duplexes
 X_s = Nbr. of single-family homes
 X_r = Retail space in ft²

$$\text{Maximize } Z = 600X_e + 750X_d + 1200X_s + 100X_r$$

$$\text{s.t. } X_e \leq 500, X_d \leq 300, X_s \leq 250$$

$$X_r \geq 10X_e + 15X_d + 18X_s$$

$$X_r \leq 10000$$

$$X_d \geq \frac{X_e + X_s}{2}$$

$$X_e, X_d, X_s, X_r \geq 0$$

Optimal Solution:

$$Z = 1,595,714.29$$

$$X_e = 207.14, X_d = 228.57$$

$$X_s = 250, X_r = 10,000$$

LP does not guarantee integer solution.
Use rounded solution or apply integer LP algorithm (Chapter 9).

x_i = Acquired portion of property i

Each site is represented by a separate LP.
The site that yields the smaller objective value is selected.

Site 1 LP:

$$\text{Minimize } Z = 25 + X_1 + 2.1X_2 + 2.35X_3 + 1.85X_4 + 2.95X_5$$

$$\text{s.t. } X_4 \geq .75, \text{ all } X_i \geq 0, i=1,2,\dots,5$$

$$20X_1 + 50X_2 + 50X_3 + 30X_4 + 60X_5 \geq 200$$

Optimum: $Z = 34.6625$ million \$
 $X_1 = .875, X_2 = X_3 = 1, X_4 = .75, X_5 = 1$

Site 2 LP:

$$\text{Minimize } Z = 27 + 2.8X_1 + 1.9X_2 + 2.8X_3 + 2.5X_4$$

$$\text{s.t. } X_3 \geq .5, X_1, X_2, X_3, X_4 \geq 0$$

$$80X_1 + 60X_2 + 50X_3 + 70X_4 \geq 200$$

Optimum: $Z = 34.35$ million \$
 $X_1 = X_2 = 1, X_3 = X_4 = .5$

Select Site 2.

X_{ij} = portion of project i completed in year j

3

$$\begin{aligned} \text{Maximize } Z &= .05(4X_4 + 3X_{12} + 2X_{13}) + \\ &+ .07(3X_{22} + 2X_{23} + X_{24}) + \\ &+ .15(4X_{31} + 3X_{32} + 2X_{33} + X_{34}) + \\ &+ .02(2X_{43} + X_{44}) \end{aligned}$$

s.t.

$$\sum_{j=1}^3 X_{1j} = 1, \quad \sum_{j=3}^4 X_{4j} = 1$$

$$.25 \leq \sum_{j=2}^5 X_{2j} \leq 1, \quad .25 \leq \sum_{j=1}^5 X_{3j} \leq 1$$

$$5X_{11} + 15X_{31} \leq 3$$

$$5X_{12} + 8X_{22} + 15X_{32} \leq 6$$

$$5X_{13} + 8X_{23} + 15X_{33} + 1.2X_{43} \leq 7$$

$$8X_{24} + 15X_{34} + 1.2X_{44} \leq 7$$

$$8X_{25} + 15X_{35} \leq 7$$

Optimum:

$$Z = \$523,750$$

$$X_{11} = .6, X_{12} = .4$$

$$X_{24} = .225, X_{25} = .025$$

$$X_{32} = .267, X_{33} = .387, X_{34} = .346$$

$$X_{43} = 1$$

x_l = Nbr. of low income units

4

x_m = Nbr. of middle income units

x_u = Nbr. of upper income units

x_p = Nbr. of public housing units

x_s = Nbr. of school rooms

x_r = Nbr. of retail units

x_c = Nbr. of condemned homes

$$\text{Maximize } Z = 7X_l + 12X_m + 20X_u + 5X_p + 15X_r$$

$$-10X_s - 7X_c$$

$$100 \leq X_l \leq 200, \quad 125 \leq X_m \leq 190$$

$$75 \leq X_u \leq 260, \quad 300 \leq X_p \leq 600$$

$$0 \leq X_s \leq 2.045$$

$$.05X_l + .07X_m + .03X_u + .025X_p +$$

$$.045X_s + .1X_c \leq .85(50 + .25X_c)$$

$$X_r \geq .023X_l + .034X_m + .046X_u +$$

$$.023X_p + .034X_s$$

continued...

Set 2.3a

$$25x_S \geq 1.3x_L + 1.2x_m + .5x_u + 1.4x_p$$

Optimum: $Z = 8290.30$ thousand \$

$$x_L = 100, x_m = 125, x_u = 227.04$$

$$x_p = 300, x_S = 32.54, x_n = 25$$

$$x_c = 0$$

x_1 = Nbr. of single-family homes

5

x_2 = Nbr. of double-family homes

x_3 = Nbr. of triple-family homes

x_4 = Nbr. of recreation areas

$$\text{Maximize } Z = 10,000x_1 + 12,000x_2 + 15,000x_3$$

s.t.

$$2x_1 + 3x_2 + 4x_3 + x_4 \leq .85 \times 800$$

$$\frac{x_1}{x_1 + x_2 + x_3} \geq .5 \text{ or } .5x_1 - .5x_2 - .5x_3 \geq 0$$

$$x_4 \geq \frac{x_1 + 2x_2 + 3x_3}{200} \text{ or } 200x_4 - x_1 - 2x_2 - 3x_3 \geq 0$$

$$1000x_1 + 1200x_2 + 1400x_3 + 800x_4 \geq 100,000$$

$$400x_1 + 600x_2 + 840x_3 + 450x_4 \leq 200,000$$

$$x_1, x_2, x_3, x_4 \geq 0$$

Optimum solution:

$$x_1 = 339.15 \text{ homes}$$

$$x_2 = 0$$

$$x_3 = 0$$

$$x_4 = 1.69 \text{ areas}$$

$$Z = \$339,521.20$$

New land use constraint:

6

$$2x_1 + 3x_2 + 4x_3 + x_4 \leq .85(800 + 100)$$

New Optimum Solution:

$$Z = \$381,5461.35$$

$$x_1 = 381.54 \text{ homes}$$

$$x_2 = x_3 = 0$$

$$x_4 = 1.91 \text{ areas}$$

$$\Delta Z = \$3815,461.35 - \$339,521.20$$

$$= \$423,940.35$$

$\Delta Z < \$450,000$, the purchasing cost of 100 acres. Hence, the purchase of the new acreage is not recommended.

Set 2.3b

The constraints remain unchanged, but the objective function is changed to

Maximize $Z = y - \text{commission}$

where

$$\begin{aligned} \text{Commission} &= .001(\text{all transactions in } \$) \\ &= .001[(X_{12} + X_{13} + X_{14} + X_{15}) + \\ &\quad \frac{1}{.769}(X_{21} + X_{23} + X_{24} + X_{25}) + \\ &\quad \frac{1}{.625}(X_{31} + X_{32} + X_{34} + X_{35}) + \\ &\quad \frac{1}{.105}(X_{41} + X_{42} + X_{43} + X_{45}) + \\ &\quad \frac{1}{.342}(X_{51} + X_{52} + X_{53} + X_{54})] \end{aligned}$$

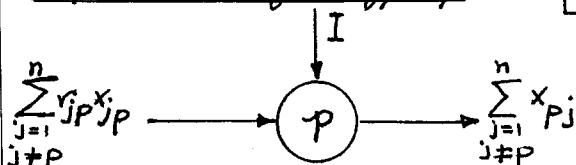
Optimum solution:

	Without	With
Z	5.09032	5.06211
y	5.09032	5.08986
Return	1.8064%	1.2421%

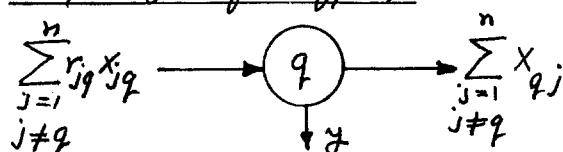
$$\begin{aligned} \text{Commission} &= 5.08986 - 5.06211 \\ &= \$27,750 \end{aligned}$$

or, .555% of the original investment of \$5 million

Input I in fund type p :



Output y in fund type q :

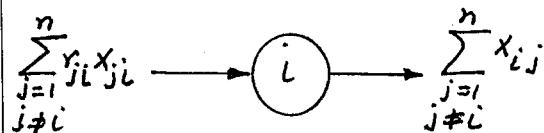


For specific p and q , the model below can be used to transform any fund to any other fund. In

continued...

the present problem, $p=1 (\$)$ and $q=2 (\text{€}), 3 (\text{£}), 4 (\text{¥}), \text{and } 5 (\text{KD})$.

General node i :



Maximize $Z = y$

s.t.

$$I + \sum_{\substack{j=1 \\ j \neq p}}^n r_{jp} x_{jp} = \sum_{\substack{j=1 \\ j \neq p}}^n x_{pj}$$

$$\sum_{\substack{j=1 \\ j \neq q}}^n r_{qj} x_{qj} = y + \sum_{\substack{j=1 \\ j \neq q}}^n x_{qj}$$

$$\sum_{\substack{j=1 \\ j \neq i}}^n r_{ij} x_{ij} = \sum_{\substack{j=1 \\ j \neq i}}^n x_{ij}, \quad i \neq p \text{ or } q$$

$$0 \leq x_{ij} \leq \text{Cap}_i, \quad \text{all } i \text{ and } j$$

Note: Solver or AMPL is ideal for solving this problem interactively. See files solver2.3b-2.xls and ampl2.3b-2.txt.

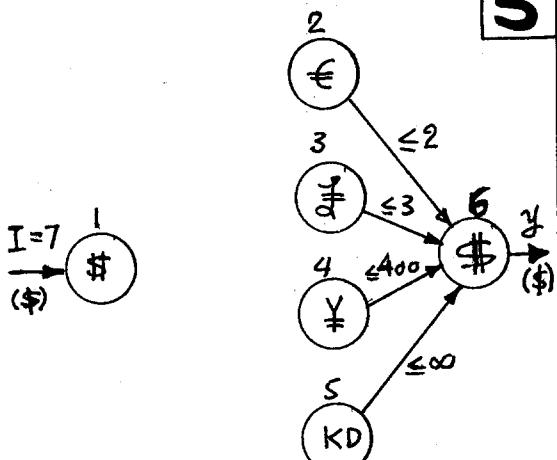
Results: (No commission)

p	q	Rate of return
\$	\$	1.8064%
\$	€	1.7966%
\$	£	1.8287%
\$	¥	2.8515%
\$	KD	1.0471%

Wide discrepancy in ¥ and KD currencies may be attributed to the fact that their exchange rates may not be consistent with the remaining rates. Nevertheless, the problem shows that there may be advantages in targeting accumulation in different currencies.

Set 2.3b

3



To formulate the objective function correctly, all output currencies are converted to a single currency (arbitrarily chosen to be \$). Thus

$$Y = r_{21}x_{26} + r_{31}x_{36} + r_{41}x_{46} + r_{51}x_{56}$$

Maximize $Z = Y$

s.t. $x_{26} \leq 2, x_{36} \leq 3, x_{46} \leq 400; x_{5j} \leq 3.5$
 $x_{ij} \leq 5, x_{2j} \leq 3, x_{4j} \leq 100, x_{5j} \leq 2.8, \text{ all } j$

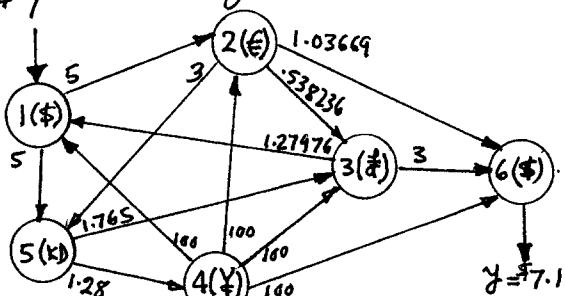
$$I + \sum_{i=2}^5 x_{ii} r_{ii} = \sum_{i=2}^5 x_{ij}$$

$$Y = r_{21}x_{26} + r_{31}x_{36} + r_{41}x_{46} + r_{51}x_{56}$$

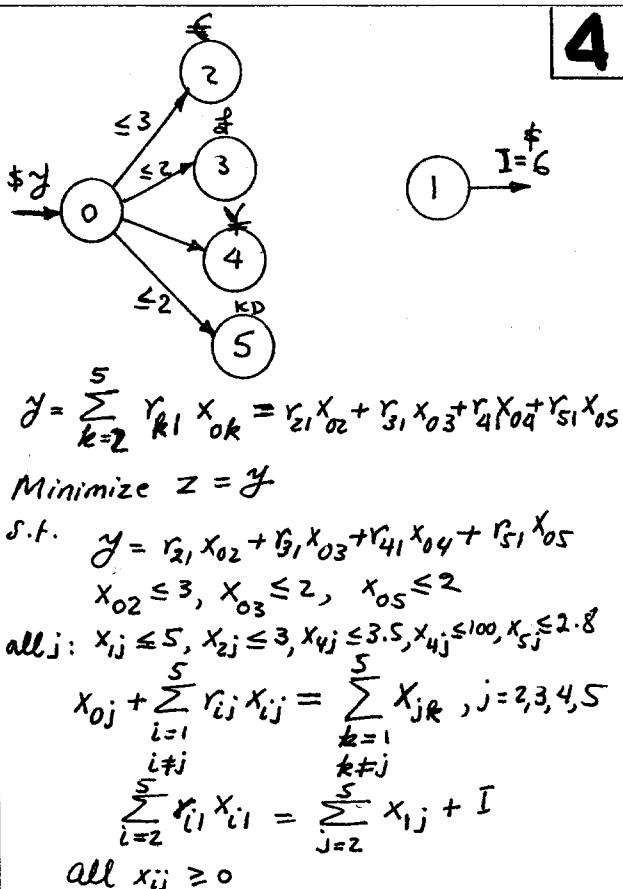
$$x_{i6} + \sum_{\substack{j=1 \\ j \neq i}}^5 x_{ij} = \sum_{\substack{j=1 \\ j \neq i}}^5 r_{ij}x_{jj}, i = 2, 3, 4, 5$$

all $x_{ij} \geq 0, i \neq j$

Solution: Total accumulation $Y = \$7.1$ million
Rate of return = 1.4354%



4



$$Y = \sum_{k=2}^5 r_{ki} x_{0k} = r_{21}x_{02} + r_{31}x_{03} + r_{41}x_{04} + r_{51}x_{05}$$

Minimize $Z = Y$

$$\text{s.t. } Y = r_{21}x_{02} + r_{31}x_{03} + r_{41}x_{04} + r_{51}x_{05}$$

$$x_{02} \leq 3, x_{03} \leq 2, x_{05} \leq 2$$

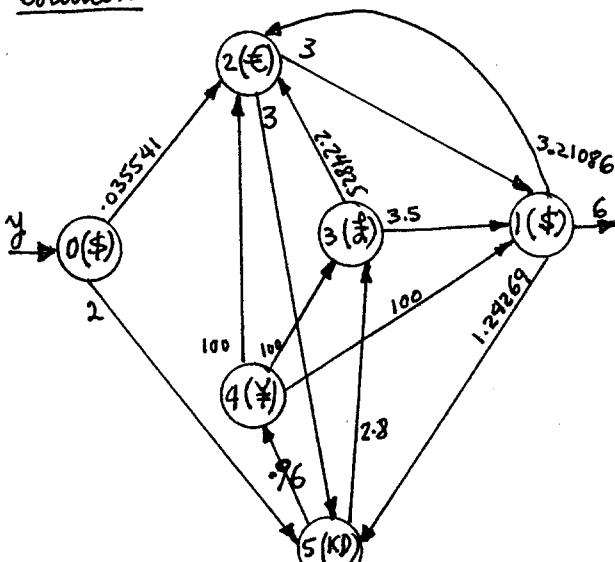
$$\text{all } j: x_{1j} \leq 5, x_{2j} \leq 3, x_{4j} \leq 3.5, x_{4j} \leq 100, x_{5j} \leq 2.8$$

$$x_{0j} + \sum_{\substack{i=1 \\ i \neq j}}^5 r_{ij} x_{ij} = \sum_{k=1}^5 x_{jk}, j = 2, 3, 4, 5$$

$$\sum_{i=2}^5 r_{i1} x_{i1} = \sum_{j=2}^5 x_{1j} + I$$

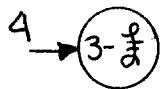
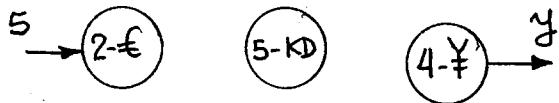
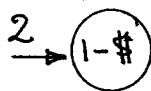
All $x_{ij} \geq 0$

Solution:



$$\text{Rate of return} = 1.7638\%$$

5



$$\text{Maximize } z = y$$

s.t.

$$y = r_{14}x_{14} + r_{24}x_{24} + r_{34}x_{34} + r_{54}x_{54}$$

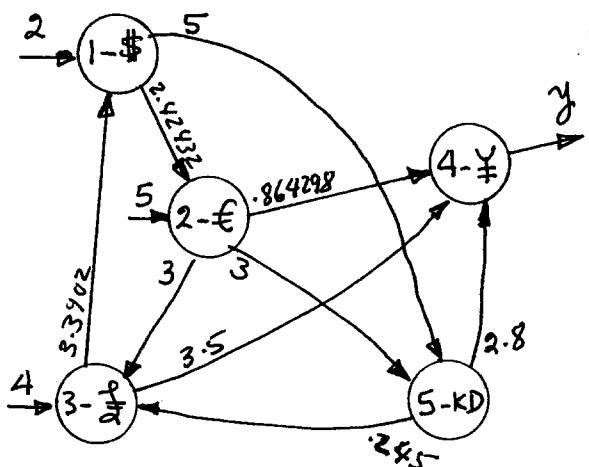
$$\sum_{\substack{i=1 \\ i \neq j}}^5 r_{ij}x_{ij} = \sum_{\substack{k=1 \\ k \neq j}}^5 x_{jk} - \begin{cases} 2, & j=1 \\ 5, & j=2 \\ 4, & j=3 \\ -y, & j=4 \\ 0, & j=5 \end{cases}$$

$$x_{ij} \leq C_i, \text{ for all } i \text{ and } j \quad \{i \neq j\}$$

$$x_{ij} \geq 0, \text{ for all } i \text{ and } j \quad \{i \neq j\}$$

Solution: $y = 1584.91 \text{ million } \frac{\alpha}{\beta}$

Rate of return = .8853%



Set 2.3c

(a) x_i = undertaken portion of project i

Maximize

$$Z = 32.4x_1 + 35.8x_2 + 17.75x_3 + 14.8x_4 + 18.2x_5 + 12.35x_6$$

Subject to

$$\begin{aligned} 10.5x_1 + 8.3x_2 + 10.2x_3 + 7.2x_4 + 12.3x_5 + 9.2x_6 &\leq 60 \\ 14.4x_1 + 12.6x_2 + 14.2x_3 + 10.5x_4 + 10.1x_5 + 7.8x_6 &\leq 70 \\ 2.2x_1 + 9.5x_2 + 5.6x_3 + 7.5x_4 + 8.3x_5 + 6.9x_6 &\leq 35 \\ 2.4x_1 + 3.1x_2 + 4.2x_3 + 5.0x_4 + 6.3x_5 + 5.1x_6 &\leq 20 \\ 0 \leq x_j &\leq 1, j=1,2,\dots,6 \end{aligned}$$

TOA optimum solution:

$$x_1 = x_2 = x_3 = x_4 = 1, x_5 = 0.84, x_6 = 0, Z = 116.06$$

(b) Add the constraint $x_2 \leq x_6$

TOA optimum solution:

$$x_1 = x_2 = x_3 = x_4 = x_5 = 1, x_6 = 0.03, Z = 113.68$$

(c) Let s_i be the unused funds at the end of year i and change the right-hand sides of constraints 2, 3, and 4 to $70+s_1$, $35+s_2$, and $20+s_3$, respectively.

TOA optimum solution:

$$x_1 = x_2 = x_3 = x_4 = x_5 = 1, x_6 = 0.71$$

$$Z = \$127.72 \text{ (thousand)}$$

The solution is interpreted as follows:

i	s_i	$s_i - s_{i-1}$	Decision
1	4.96	-	-
2	7.62	+2.66	Don't borrow from yr 1
3	4.62	-3.00	Borrow \$3 from year 2
4	0	-4.62	Borrow \$4.62 from yr 2

The effect of availing excess money for use in later years is that the first five projects are completed and 71% of project 6 is undertaken.

The total revenue increases from \$116,060 to 127,720.

(d) The slack s_i in period i is treated as an unrestricted variable.

TOA optimum solution: $Z = \$131.30$
 $s_1 = 2.3, s_2 = 4, s_3 = -5, s_4 = -6.1$

This means that additional funds are needed in years 3 and 4.

$$\begin{aligned} \text{Increase in return} &= 131.30 - 116.06 \\ &= \$15.24 \end{aligned}$$

Ignoring the time value of money, the amount borrowed $5+6.1-(2.3+4)$ $= \$8.4$. Thus, rate of return $= \frac{15.24-8.4}{8.4} \approx 81\%$

2

x_i = dollar investment in project i , $i=1,2,3,4$

y_j = dollar investment in bank in year j , $j=1,2,3,4,5$

Maximize $Z = y_5$

Subject to

$$\begin{aligned} x_1 + x_2 + x_3 + x_4 + y_1 &\leq 10,000 \\ .5x_1 + .6x_2 - x_3 + .4x_4 + 1.065y_1 - y_2 &= 0 \\ .3x_1 + .2x_2 + .8x_3 + .6x_4 + 1.065y_2 - y_3 &= 0 \\ 1.8x_1 + 1.5x_2 + 1.9x_3 + 1.8x_4 + 1.065y_3 - y_4 &= 0 \\ 1.2x_1 + 1.3x_2 + .8x_3 + .95x_4 + 1.065y_4 - y_5 &= 0 \end{aligned}$$

All variables ≥ 0

TOA optimal solution:

$$x_1 = 0, x_2 = \$10,000, x_3 = \$6000, x_4 = 0$$

$$y_1 = 0, y_2 = 0, y_3 = \$6800, y_4 = \$33,642$$

$Z = \$53,628.73$ at the start of year 5

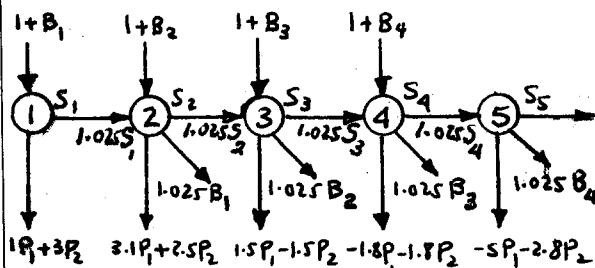
continued...

P_i = fraction undertaken of project

$i, i = 1, 2$

B_j = million dollars borrowed in quarter j , $j = 1, 2, 3, 4$

S_j = surplus million dollars at the start of quarter j , $j = 1, 2, 3, 4, 5$



(a) Maximize $Z = S_5$

subject to

$$\begin{aligned} P_1 + 3P_2 + S_1 - B_1 &= 1 \\ 3.1P_1 + 2.5P_2 - 1.02S_1 + S_2 + 1.025B_1 - B_2 &= 1 \\ 1.5P_1 - 1.5P_2 - 1.02S_2 + S_3 + 1.025B_2 - B_3 &= 1 \\ -1.8P_1 - 1.8P_2 - 1.02S_3 + S_4 + 1.025B_3 - B_4 &= 1 \\ -5P_1 - 2.8P_2 - 1.02S_4 + S_5 + 1.025B_4 &= 1 \\ 0 \leq P_1 \leq 1, \quad 0 \leq P_2 \leq 1 & \\ 0 \leq B_j \leq 1, \quad j = 1, 2, 3, 4 & \end{aligned}$$

Optimum Solution:

$$P_1 = .7113 \quad P_2 = 0$$

$Z = 5.8366$ million dollars

$B_1 = 0, \quad B_2 = .9104$ million dollars

$B_3 = 1$ million dollars, $B_4 = 0$

(b) $B_1 = 0, \quad S_1 = .2887$ million \$

$B_2 = .9104, \quad S_2 = 0$

$B_3 = 1, \quad S_3 = 0$

$B_4 = 0, \quad S_4 = 1.2553$

The solution shows that $B_i \cdot S_i = 0$, meaning that you can't borrow and also end up with surplus in any quarter. The result makes sense because the cost of borrowing (2.5%) is higher than the return on surplus funds (2%).

Assume that the investment program ends at the start of year 11. Thus, the 6-year bond option can be exercised in years 1, 2, 3, 4, and 5 only. Similarly, the 9-year bond can be used in years 1 and 2 only. Hence, from year 6 on, the only option available is insured savings at 7.5%.

Let

I_i = insured savings investments in year i , $i = 1, 2, \dots, 10$

G_i = 6-year bond investment in year i , $i = 1, 2, \dots, 5$

M_i = 9-year bond investment in year i , $i = 1, 2$

The objective is to maximize total accumulation at the end of year 10; that is,

$$\text{maximize } Z = 1.075 I_{10} + 1.079 G_5 + 1.085 M_2$$

The constraints represent the balance equation for each year's cash flow.

$$I_1 + .98G_1 + 1.02M_1 = 2$$

$$I_2 + .98G_2 + 1.02M_2 = 2 + 1.075I_1 + .079G_1 + .085M_1$$

$$I_3 + .98G_3 = 2.5 + 1.075I_2 + .079(G_1 + G_2) + .085(M_1 + M_2)$$

$$I_4 + .98G_4 = 2.5 + 1.075I_3 + .079(G_1 + G_2 + G_3) + .085(M_1 + M_2)$$

$$I_5 + .98G_5 = 3 + 1.075I_4 + .079(G_1 + G_2 + G_3 + G_4) + .085(M_1 + M_2)$$

$$I_6 = 3.5 + 1.075I_5$$

$$+ .079(G_1 + G_2 + G_3 + G_4 + G_5) + .085(M_1 + M_2)$$

continued...

Set 2.3c

$$\begin{aligned}
 I_7 &= 3.5 + 1.075 I_6 + 1.079 G_1 \\
 &\quad + .079 (G_2 + G_3 + G_4 + G_5) \\
 &\quad + .085 (M_1 + M_2) \\
 I_8 &= 4 + 1.075 I_7 + 1.079 G_2 \\
 &\quad + .079 (G_3 + G_4 + G_5) \\
 &\quad + .085 (M_1 + M_2) \\
 I_9 &= 4 + 1.075 I_8 + 1.079 G_3 \\
 &\quad + .079 (G_4 + G_5) \\
 &\quad + .085 (M_1 + M_2) \\
 I_{10} &= 5 + 1.075 I_9 + 1.079 G_4 \\
 &\quad + .079 G_5 + 1.085 M_1 + .085 M_2 \\
 \text{all variables } &\geq 0
 \end{aligned}$$

*** OPTIMUM SOLUTION SUMMARY ***				
Title: Problem 26a-14 Final iteration No: 14 Objective value (max) = 46.8500				
Variable	Value	Obj Coeff	Obj Val	Contrib
x1_11	0.0000	0.0000	0.0000	
x2_12	0.0000	0.0000	0.0000	
x3_13	0.0000	0.0000	0.0000	
x4_14	0.0000	0.0000	0.0000	
x5_15	0.0000	0.0000	0.0000	
x6_16	4.6331	0.0000	0.0000	
x7_17	9.6137	0.0000	0.0000	
x8_18	15.4678	0.0000	0.0000	
x9_19	24.6663	0.0000	0.0000	
x10_110	37.5201	1.0750	40.3341	
x11_61	0.0000	0.0000	0.0000	
x12_62	0.0000	0.0000	0.0000	
x13_63	2.9053	0.0000	0.0000	
x14_64	3.1395	0.0000	0.0000	
x15_65	3.9028	1.0790	4.2111	
x16_M1	1.9608	0.0000	0.0000	
x17_M2	2.1242	1.0850	2.3047	
Constraint	RHS	Slack(-)/Surplus(+)		
1 (=)	2.0000	0.0000		
2 (=)	2.0000	0.0000		
3 (=)	2.5000	0.0000		
4 (=)	2.5000	0.0000		
5 (=)	3.0000	0.0000		
6 (=)	3.5000	0.0000		
7 (=)	3.5000	0.0000		
8 (=)	4.0000	0.0000		
9 (=)	4.0000	0.0000		
10 (=)	5.0000	0.0000		

Year	Recommendation
1	Invest all in 9-yr bond
2	Invest all in 9-yr bond
3	Invest all in 6-yr bond
4	Invest all in 6-yr bond
5	Invest all in 6-yr bond
7	Invest all in insured savings
8	Invest all in insured savings
9	Invest all in insured savings
10	Invest all in insured savings

x_{iA} = amount invested in year i , plan A (\$1000)
5

x_{iB} = amount invested in year i , plan B (\$1000\$)

$$\text{Maximize } Z = 3x_{2B} + 1.7x_{3A}$$

subject to

$$x_{1A} + x_{1B} \leq 100$$

$$-1.7x_{1A} + x_{2A} + x_{2B} = 0$$

$$-3x_{1B} - 1.7x_{2A} + x_{3A} = 0$$

$$x_{iA}, x_{iB} \geq 0 \text{ for } i = 1, 2, 3$$

*** OPTIMUM SOLUTION SUMMARY ***

Title: Problem 2.6a-15
Final iteration No: 4
Objective value (max) = 510.0000
=> ALTERNATIVE solution detected at x2

Variable	Value	Obj Coeff	Obj Val	Contrib
x1_x1A	100.0000	0.0000	0.0000	
x2_x1B	0.0000	0.0000	0.0000	
x3_x2A	0.0000	0.0000	0.0000	
x4_x2B	170.0000	3.0000	510.0000	
x5_x3A	0.0000	1.7000	0.0000	

Constraint	RHS	Slack(-)/Surplus(+)
1 (<=)	100.0000	0.0000-
2 (<=)	0.0000	0.0000-
3 (<=)	0.0000	0.0000-

Optimum solution: Invest \$100,000 in A in yr 1 and \$170,000 in B in yr 2.

Alternative optimum: Invest \$100,000 in B in yr 1 and \$300,000 in A in yr 3.

x_i = dollars allocated to choice i ,
 $i = 1, 2, 3, 4$

6

y = minimum return

$$\text{Maximize } Z = \min \begin{cases} -3x_1 + 4x_2 - 7x_3 + 15x_4 \\ 5x_1 - 3x_2 + 9x_3 + 4x_4 \\ 3x_1 - 9x_2 + 10x_3 - 8x_4 \end{cases}$$

$$x_1 + x_2 + x_3 + x_4 \leq 500$$

$$x_1, x_2, x_3, x_4 \geq 0$$

The problem can be converted to a linear program as

continued...

Set 2.3c

Maximize $Z = y$

subject to

$$-3x_1 + 4x_2 - 7x_3 + 15x_4 \geq y$$

$$5x_1 - 3x_2 + 9x_3 + 4x_4 \geq y$$

$$3x_1 - 9x_2 + 10x_3 - 8x_4 \geq y$$

$$x_1 + x_2 + x_3 + x_4 \leq 500$$

$$x_1, x_2, x_3, x_4 \geq 0$$

if unrestricted

*** OPTIMUM SOLUTION SUMMARY ***

Title:

Final iteration No: 5

Objective value (max) = 1175.0000

Variable	Value	Obj Coeff	Obj Val Contrib
x1	0.0000	0.0000	0.0000
x2	0.0000	0.0000	0.0000
x3	287.5000	0.0000	0.0000
x4	212.5000	0.0000	0.0000
x5 y	1175.0000	1.0000	1175.0000

Constraint	RHS	Slack(-)/Surplus(+)
1 (>)	0.0000	0.0000+
2 (>)	0.0000	2262.5000+
3 (>)	0.0000	0.0000+
4 (<)	500.0000	0.0000-

Allocate \$287.50 to choice 3
and \$212.50 to choice 4. Return =
\$1175.00

$$i = \begin{cases} 1, \text{ regular savings} \\ 2, \text{ 3-month CD} \\ 3, \text{ 6-month CD} \end{cases}$$

7

x_{it} = Deposit in plan i at start of month t

$$t = \begin{cases} 1, 2, \dots, 12 & \text{if } i=1 \\ 1, 2, \dots, 10 & \text{if } i=2 \\ 1, 2, \dots, 7 & \text{if } i=3 \end{cases}$$

y_i = initial amount on hand to
ensure a feasible solution

r_i = interest rate for plan $i = 1, 2, 3$

$$\bar{J}_i = \begin{cases} 12, & i=1 \\ 10, & i=2 \\ 7, & i=3 \end{cases}$$

continued...

$$P_i = \begin{cases} 1, & i=1 \\ 3, & i=2 \\ 6, & i=3 \end{cases} \quad d_t = \$\text{demand for period } t$$

$$\text{Maximize } Z = \sum_{t=1}^{12} \sum_{i=1}^3 r_i x_{i,t} - y_i$$

$$t - P_i > 0$$

s.t.

$$y_i - x_{1i} - x_{2i} - x_{3i} \geq d_{i,3}$$

$$1000 + \sum_{i=1}^3 (1+r_i) x_{i,t} - \sum_{i=1}^3 x_{it} \geq d_t, t=2, \dots, 12$$

$$t - P_i > 0 \quad t \leq J_i$$

$$x_{it}, y_i \geq 0$$

Solution: (see file ampl2.3c-7.txt)

$$y_i = \$1200, Z = -1136.29$$

$$\text{Interest amount} = 1200 - 1136.29 = \$63.71$$

Deposits:

t	x_{1t}	x_{2t}	x_{3t}
1	0	0	0
2	0	200	0
3	286.48	313.53	0
4	0	587.43	0
5	314.37	289.30	0
6	0	734.69	0
7	0	98.20	0
8	0	294.60	0
9	0	848.16	0
10	0	0	0
11	0	0	0
12	0	0	0

Set 2.3d

X_{W1} = # wrenches/wk using regular time
 X_{W2} = # wrenches/wk using overtime
 X_{W3} = # wrenches/wk using subcontracting
 X_{C1} = # chisels/wk using regular time
 X_{C2} = # chisels/wk using overtime
 X_{C3} = # chisels/wk using subcontracting

$$\text{Minimize } Z = 2X_{W1} + 2.8X_{W2} + 3X_{W3} + 2.1X_{C1}$$

$$\text{Subject to} \quad + 3.2X_{C2} + 4.2X_{C3}$$

$$X_{W1} \leq 550, X_{W2} \leq 250$$

$$X_{C1} \leq 620, X_{C2} \leq 280$$

$$X_{C1} + X_{C2} + X_{C3} \geq 2$$

$$X_{W1} + X_{W2} + X_{W3}$$

or

$$2X_{W1} + 2X_{W2} + 2X_{W3} - X_{C1} - X_{C2} - X_{C3} \leq 0$$

$$X_{W1} + X_{W2} + X_{W3} \geq 1500$$

$$X_{C1} + X_{C2} + X_{C3} \geq 1200$$

$$\text{all variables} \geq 0$$

(a) Optimum from TORA:

$$X_{W1} = 550, X_{W2} = 250, X_{W3} = 700$$

$$X_{C1} = 620, X_{C2} = 280, X_{C3} = 2100$$

$$Z = \$14,918$$

(b) Increasing marginal cost ensures that regular time capacity is used before that of overtime, and that overtime capacity is used before that of subcontracting. If the marginal cost function is not monotonically increasing, additional constraints are needed to ensure that the capacity restriction is satisfied.

continued...

X_j = number of units produced of product j , $j = 1, 2, 3, 4$

2

Profit per unit:

$$\text{Product 1} = 75 - 2 \times 10 - 3 \times 5 - 7 \times 4 = \$12$$

$$\text{Product 2} = 70 - 3 \times 10 - 2 \times 5 - 3 \times 4 = \$18$$

$$\text{Product 3} = 55 - 4 \times 10 - 1 \times 5 - 2 \times 4 = \$2$$

$$\text{Product 4} = 45 - 2 \times 10 - 2 \times 5 - 1 \times 4 = \$11$$

$$\text{Maximize } Z = 12X_1 + 18X_2 + 2X_3 + 11X_4$$

s.t.

$$2X_1 + 3X_2 + 4X_3 + 2X_4 \leq 500$$

$$3X_1 + 2X_2 + X_3 + 2X_4 \leq 380$$

$$7X_1 + 3X_2 + 2X_3 + X_4 \leq 450$$

$$X_1, X_2, X_3, X_4 \geq 0$$

TORA Solution:

$$X_1 = 0, X_2 = 133.33, X_3 = 0, X_4 = 50$$

$$Z = \$2950$$

X_j = number of units of model j

3

$$\text{Maximize } Z = 30X_1 + 20X_2 + 50X_3$$

Subject to

$$(1) \quad 2X_1 + 3X_2 + 5X_3 \leq 4000$$

$$(2) \quad 4X_1 + 2X_2 + 7X_3 \leq 6000$$

$$(3) \quad X_1 + 5X_2 + \frac{1}{3}X_3 \leq 1500$$

$$(4) \quad \frac{X_1}{3} = \frac{X_2}{2}, \text{ or } 2X_1 - 3X_2 = 0$$

$$(5) \quad \frac{X_2}{2} = \frac{X_3}{5}, \text{ or } 5X_2 - 2X_3 = 0$$

$$X_1 \geq 200, X_2 \geq 200, X_3 \geq 150$$

*** OPTIMUM SOLUTION SUMMARY ***

Title: Problem 2.6a-12
Final iteration No: 4
Objective value (max) = 41081.0820

Variable	Value	Obj Coeff	Obj Val Contrib
x1	324.3243	30.0000	9729.7305
x2	216.2162	20.0000	4324.3242
x3	540.5405	50.0000	27027.0273
Constraint	RHS	Slack(-)/Surplus(+)	
1 (<)	4000.0000	0.0000	
2 (<)	6000.0000	486.4865	
3 (<)	1500.0000	887.3875	
4 (=)	0.0000	0.0000	
5 (=)	0.0000	0.0000	
LB-x1	200.0000	124.3243	
LB-x2	200.0000	16.2162	
LB-x3	150.0000	390.5405	

Set 2.3d

x_{ij} = Nbr. Cartons in month i from supplier j

I_i = End inventory in period i , $I_0 = 0$

C_{ij} = Price per unit of x_{ij}

h = Holding cost/unit/month

C = Supplier capacity/month

d_i = Demand for month i

$i = 1, 2, 3, j = 1, 2$

$$\text{Minimize } Z = \sum_{i=1}^3 \sum_{j=1}^2 C_{ij} x_{ij} + \frac{h}{2} \left(\sum_{i=1}^3 \left(\sum_{j=1}^2 x_{ij} + I_{i-1} + I_i \right) \right)$$

S.t. $x_{ij} \leq C$, all i and j

$$\sum_{j=1}^2 x_{ij} + I_{i-1} - I_i = d_i, \text{ all } i$$

Optimum solution:

i	x_{i1}	x_{i2}	I
1	400	100	0
2	400	400	200
3	200	0	0

Total cost = \$167,450.

X_i = Production amount in quarter i

I_i = End inventory for quarter i

$$\text{Minimize } Z = 20X_1 + 22X_2 + 24X_3 + 26X_4 + 3.5(I_1 + I_2 + I_3)$$

S.t.

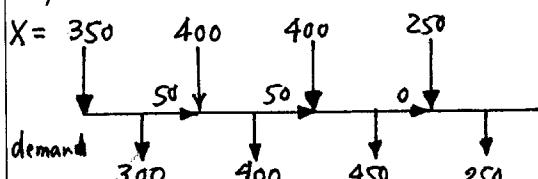
$$X_1 = 300 + I_1, \quad X_i \leq 400, i=1,2,3$$

$$I_1 + X_2 = 400 + I_2, \quad I_i \leq 100, i=1,2,3$$

$$I_2 + X_3 = 450 + I_3, \quad I_0 = I_4 = 0$$

$$I_3 + X_4 = 250$$

Optimum solution:



Total cost = \$32,250

4

x_{ij} = Qty of product i in month j ,
 $i = 1, 2, j = 1, 2, 3$

I_{ij} = End inventory of product i in month j

$$\text{Minimize } Z = 30(x_{11} + x_{12} + x_{13}) + 28(x_{21} + x_{22} + x_{23}) + .9(I_{11} + I_{12} + I_{13}) + .75(I_{21} + I_{22} + I_{23})$$

S.t.

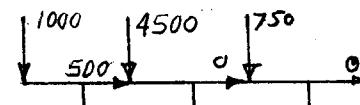
$$(x_{1j}/1.25) + x_{2j} \leq \begin{cases} 3000, & j=1 \\ 3500, & j=2 \\ 3000, & j=3 \end{cases}$$

$$I_{1,j-1} + x_{1j} - I_{1j} = \begin{cases} 500, & j=1 \\ 5000, & j=2 \\ 750, & j=3 \end{cases}$$

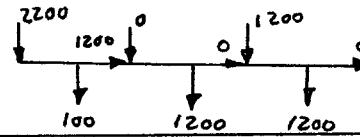
$$I_{2,j-1} + x_{2j} - I_{2j} = \begin{cases} 1000, & j=1 \\ 1200, & j=2 \\ 1200, & j=3 \end{cases}$$

Optimum solution: Cost = \$284,050

Product 1:



Product 2:



5

x_{ij} = Qty by operation i in month j ,
 $i = 1, 2, j = 1, 2, 3$

$$\text{Minimize } Z = .2 \sum_{j=1}^3 I_{1j} + .4 \sum_{j=1}^3 I_{2j} + 10x_{11} + 12x_{12} + 11x_{21} + 15x_{22} + 18x_{23}$$

$$+ 6x_{11} \leq 800, .6x_{12} \leq 700, .6x_{13} \leq 550$$

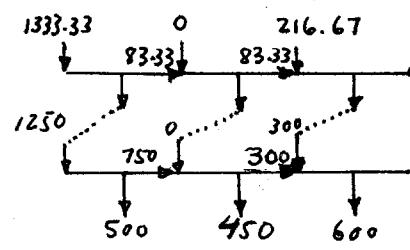
$$.8x_{21} \leq 1000, .8x_{22} \leq 850, .8x_{23} \leq 700$$

$$x_{1j} + I_{1,j-1} = x_{2j} + I_{1j} \quad \} j = 1, 2, 3$$

$$x_{2j} + I_{2,j-1} = I_{2j} + d_j \quad \} j = 1, 2, 3$$

$$I_{i0} = 0, i = 1, 2$$

Solution: Cost = \$39,720



6

7

I_{ij} = Ending inv. of op. i in month j

Set 2.3d

x_j = Units of product j , $j=1, 2$

8

y_i^- = Unused hours of machine i }
 y_i^+ = Overtime hours of machine i } $i=1, 2$

$$\text{Maximize } Z = 110x_1 + 118x_2 - 100(y_1^+ + y_2^+)$$

s.t.

$$\frac{x_1}{5} + \frac{x_2}{5} + y_1^- - y_1^+ = 8$$

$$\frac{x_1}{8} + \frac{x_2}{4} + y_2^- - y_2^+ = 8$$

$$y_1^+ \leq 4, \quad y_2^+ \leq 4$$

$$x_1, x_2, y_1^-, y_1^+, y_2^-, y_2^+ \geq 0$$

Solution:

$$\text{Revenue} = \$6,232$$

$$x_1 = 56, \quad y_1^+ = 4 \text{ hrs}$$

$$x_2 = 4, \quad y_2^+ = 0$$

$$y_1^-, y_2^- = 0$$

Set 2.3e

X_s = tons of strawberry / day

X_g = tons of grapes / day

X_a = tons of apples / day

X_A = cans of drink A / day
 X_B = cans of drink B / day } Each can holds one lb
 X_C = cans of drink C / day }

X_{SA} = 1b of strawberry used in drink A / day

X_{SB} = 1b of strawberry used in drink B / day

X_{gA} = 1b of grapes used in drink A / day

X_{gB} = 1b of grapes used in drink B / day

X_{gC} = 1b of grapes used in drink C / day

X_{aB} = 1b of apples used in drink B / day

X_{aC} = 1b of apples used in drink C / day

$$\text{Maximize } Z = 1.1X_A + 1.25X_B + 1.2X_C - 200X_S$$

$$\text{s.t. } -100X_g - 90X_a$$

$$X_S \leq 200, X_g \leq 100, X_a \leq 150$$

$$X_{SA} + X_{SB} = 1500X_S$$

$$X_{gA} + X_{gB} + X_{gC} = 1200X_g$$

$$X_{aB} + X_{aC} = 1000X_a$$

$$X_A = X_{SA} + X_{gA}$$

$$X_B = X_{SB} + X_{gB} + X_{aB}$$

$$X_C = X_{gC} + X_{aC}$$

$$X_{SA} = X_{gA},$$

$$X_{SB} = X_{gB}, X_{gB} = .5X_{aB}$$

$$3X_{gC} = 2X_{aC}$$

all variables ≥ 0

Optimum solution:

$$X_A = 90,000 \text{ cans}, X_B = 300,000 \text{ cans}, X_C = 0$$

<u>X_{ij}</u>	<u>j</u>		
<u>i</u>	<u>A</u>	<u>B</u>	<u>C</u>
S	45,000	75,000	0
g	45,000	75,000	0
a	0	150,000	0
	90,000	300,000	0

$$X_S = 80 \text{ tons}, X_g = 100 \text{ tons}, X_a = 150 \text{ tons}$$

$$Z = \$439,000/\text{day}$$

1 $X_S = 1b \text{ of screws per package}$

$X_b = 1b \text{ of bolts per package}$

$X_n = 1b \text{ of nuts per package}$

$X_w = 1b \text{ of washers per package}$

$$\text{Minimize } Z = 1.1X_S + 1.5X_b + \frac{70}{80}X_n + \frac{20}{30}X_w$$

$$\text{s.t. } Y = X_S + X_b + X_n + X_w$$

$$X_S \geq .1Y$$

$$X_b \geq .25Y, \frac{X_b}{50} \leq X_w, \frac{X_b}{10} \leq X_n$$

$$X_n \leq .15Y$$

$$X_w \leq .1Y$$

$$Y \geq 1$$

All variables are nonnegative

Optimum solution:

$$Y = 1, X_S = .5, X_b = .25, X_n = .15, X_w = .1$$

$$\text{Cost} = \$1.12$$

2 $X_{o,(A,B,C)} = 1b \text{ of oats in cereals A, B, C}$

$X_{r,(A,C)} = 1b \text{ of raisins in cereals A, C}$

$X_{c,(B,C)} = 1b \text{ of coconuts in cereals B, C}$

$X_{a,(A,B,C)} = 1b \text{ of almond in cereals A, B, C}$

$$Y_o = X_{oA} + X_{oB} + X_{oC}$$

$$Y_r = X_{rA} + X_{rC}$$

$$Y_c = X_{cB} + X_{cC}$$

$$Y_a = X_{aA} + X_{aB} + X_{aC}$$

$$W_A = X_{oA} + X_{rA} + X_{aA}$$

$$W_B = X_{oB} + X_{cB} + X_{aB}$$

$$W_C = X_{oC} + X_{rC} + X_{cC} + X_{aC}$$

$$\text{Maximize } Z = \frac{1}{5} (2W_A + 2.5W_B + 3W_C)$$

$$- \frac{1}{2000} (100Y_o + 120Y_r + 110Y_c + 200Y_a)$$

$$\text{s.t. } W_A \leq 500 \times 5 = 2500$$

$$W_B \leq 600 \times 5 = 3000$$

$$W_C \leq 500 \times 5 = 4000$$

continued...

Set 2.3e

$Y_0 \leq 5X_{2000} = 10,000$ $Y_r \leq 2X_{2000} = 4,000$ $Y_c \leq 1X_{2000} = 2,000$ $Y_a \leq 1X_{2000} = 2,000$ $X_{OA} = \frac{50}{5} X_{rA}, X_{OA} = \frac{50}{2} X_{aA}$ $X_{OB} = \frac{60}{2} X_{cB}, X_{OB} = \frac{60}{3} X_{aB}$ $X_{OC} = \frac{60}{3} X_{rC}, X_{OC} = \frac{60}{4} X_{cC}, X_{OC} = \frac{60}{2} X_{aC}$ <p>all variables are nonnegative.</p> <p><u>Optimum solution:</u> $Z = \\$5384.84/\text{day}$</p> $W_A = 2500 \text{ lb or } 500 \text{ boxes/day}$ $W_B = 3000 \text{ lb or } 600 \text{ boxes}$ $W_C = 5793.45 \text{ lb or } \approx 1158 \text{ boxes}$ $X_0 = 10,000 \text{ lb or } 5 \text{ tons/day}$ $X_r = 471.19 \text{ lb or } .236 \text{ ton}$ $X_c = 428.16 \text{ lb or } .214 \text{ ton}$ $X_a = 394.11 \text{ lb or } .197 \text{ ton}$
$\begin{aligned} X_{Ai} &= \text{bbl of gasoline A in fuel i} \\ X_{Bi} &= \text{bbl of gasoline B in fuel i} \\ X_{Ci} &= \text{bbl of gasoline C in fuel i} \\ X_{Di} &= \text{bbl of gasoline D in fuel i} \end{aligned} \quad i = 1, 2$ $Y_A = X_{A1} + X_{A2}$ $Y_B = X_{B1} + X_{B2}$ $Y_C = X_{C1} + X_{C2}$ $Y_D = X_{D1} + X_{D2}$ $F_1 = X_{A1} + X_{B1} + X_{C1} + X_{D1}$ $F_2 = X_{A2} + X_{B2} + X_{C2} + X_{D2}$ $\text{Maximize } Z = 200F_1 + 250F_2$ $- (120Y_A + 90Y_B + 100Y_C + 150Y_D)$

s.t.

$$X_{A1} = X_{B1}, X_{A1} = .5X_{C1}, X_{A1} = .25X_{D1}$$

$$X_{A2} = X_{B2}, X_{A2} = 2X_{C2}, X_{A2} = \frac{2}{3}X_{D2}$$

$$Y_A \leq 1000, Y_B \leq 1200, Y_C \leq 900, Y_D \leq 1500$$

$$F_1 \geq 200, F_2 \geq 400$$

Optimum solution: $Z = \$495,416.67$

$$Y_A = 958.33 \text{ bbl/day}$$

$$Y_B = 958.33 \text{ bbl/day}$$

$$Y_C = 516.67 \text{ bbl/day}$$

$$Y_D = 1500 \text{ bbl/day}$$

$$F_1 = 200 \text{ bbl/day}$$

$$F_2 = 3733.33 \text{ bbl/day}$$

5

$$A = \text{bbl of crude A / day}$$

$$B = \text{bbl of crude B / day}$$

$$R = \text{bbl of regular gasoline / day}$$

$$P = \text{bbl of premium gasoline / day}$$

$$J = \text{bbl of jet gasoline / day}$$

$$\begin{aligned} \text{Maximize } Z &= 50(R - R^+) + 70(P - P^+) \\ &+ 120(J - J^+) - (10R^- + 15P^- + 20J^-) \\ &- (2R^+ + 3P^+ + 4J^+) - (30A + 40B) \end{aligned}$$

s.t.

$$A \leq 2500, B \leq 3000$$

$$R = .2A + .25B, R + R^- - R^+ = 500$$

$$P = .1A + .3B, P + P^- - P^+ = 700$$

$$J = .25A + .1B, J + J^- - J^+ = 400$$

All variables ≥ 0

Optimum solution:

$$Z = \$21,852.94$$

$$A = 1176.47 \text{ bbl/day}$$

$$B = 1058.82 \text{ bbl/day}$$

$$R = 500 \text{ bbl/day}$$

$$P = 435.29 \text{ bbl/day}$$

$$J = 400 \text{ bbl/day}$$

continued...

Set 2.3e

$$NR = 661/\text{day of naphtha used in regular}$$

$$NP = 661/\text{day of naphtha used in premium}$$

$$NJ = 661/\text{day of naphtha used in jet}$$

$$LR = 661/\text{day of light used in regular}$$

$$LP = 661/\text{day of light used in premium}$$

$$LJ = 661/\text{day of light used in jet}$$

Using the other notation in Problem 5,

$$\begin{aligned} \text{Maximize } Z &= 50(R - R^+) + 70(P - P^+) + 12(J - J^+) \\ &\quad - (10R^+ + 15P^+ + 20J^+) - (2R^+ + 3P^+ + 4J^+) \\ &\quad - (30A + 40B) \end{aligned}$$

$$\text{s.t. } A \leq 2500, B \leq 3000$$

$$R + R^- - R^+ = 500$$

$$P + P^- - P^+ = 700$$

$$J + J^- - J^+ = 400$$

$$.35A + .45B = NR + NP + NJ$$

$$.6A + .5B = LR + LP + LJ$$

$$R = NR + LR$$

$$P = NP + LP$$

$$J = NJ + LJ$$

all variables are nonnegative

Optimum solution: $Z = \$71,473.68$

$$A = 1684.21, B = 0$$

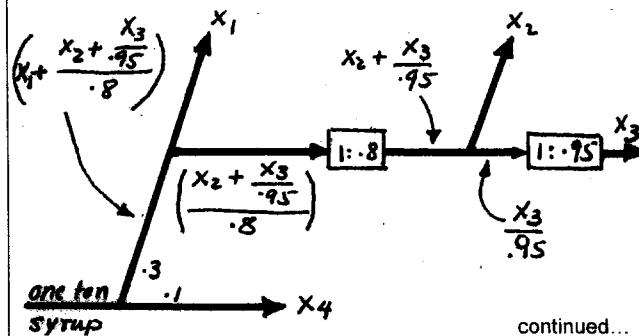
$$R = 500, P = 700, J = 400$$

$$x_1 = \text{tons of brown sugar per week}$$

$$x_2 = \text{tons of white sugar per week}$$

$$x_3 = \text{tons of powdered sugar per week}$$

$$x_4 = \text{tons of molasses per week}$$



continued...

2-25

6

$$\text{Maximize } Z = 150x_1 + 200x_2 + 280x_3 + 35x_4$$

s.t.

$$x_4 \leq 4000 \times 1$$

$$\text{or } x_4 \leq 400$$

$$x_1 + \frac{(x_2 + \frac{x_3}{.95})}{.8} \leq .3 \times 4000$$

$$\text{or } .76x_1 + .95x_2 + x_3 \leq 912$$

$$x_1 \geq 25, x_2 \geq 25$$

$$x_3 \geq 25, x_4 \geq 0$$

Optimum solution from TORA:

$$x_1 = 25 \text{ tons per week}$$

$$x_2 = 25 \text{ tons per week}$$

$$x_3 = 869.25 \text{ tons per week}$$

$$x_4 = 400 \text{ tons per week}$$

$$Z = \$222,677.50$$

8

$$A = 661/\text{hr of stock A}$$

$$B = 661/\text{hr of stock B}$$

$$Y_{Ai} = 661/\text{hr of A used in gasoline i} \quad i=1,2$$

$$Y_{Bi} = 661/\text{hr of B used in gasoline i} \quad i=1,2$$

$$\text{Maximize } Z = 7(Y_{A1} + Y_{B1}) + 10(Y_{A2} + Y_{B2})$$

s.t.

$$A = Y_{A1} + Y_{A2}, A \leq 450$$

$$B = Y_{B1} + Y_{B2}, B \leq 700$$

$$98Y_{A1} + 89Y_{B1} \geq 91(Y_{A1} + Y_{B1})$$

$$98Y_{A2} + 89Y_{B2} \geq 93(Y_{A2} + Y_{B2})$$

$$10Y_{A1} + 8Y_{B1} \leq 12(Y_{A1} + Y_{B1})$$

$$10Y_{A2} + 8Y_{B2} \leq 12(Y_{A2} + Y_{B2})$$

all variables are nonnegative

Optimum solution:

$$Z = \$10,675$$

$$A = 450 \text{ bbl/hr}$$

$$B = 700 \text{ bbl/hr}$$

$$\begin{aligned} \text{Gasoline 1 production} &= Y_{A1} + Y_{B1} \\ &= 61.11 + 213.89 = 275 \text{ bbl/hr} \end{aligned}$$

$$\begin{aligned} \text{Gasoline 2 production} &= Y_{A2} + Y_{B2} \\ &= 388.89 + 486.11 = 875 \text{ bbl/hr} \end{aligned}$$

Set 2.3e

S = tons of steel scrap / day
 A = tons of alum. scrap / day
 C = tons of Cast iron scrap / day
 A_b = tons of alum. briquettes / day
 S_b = tons silicon briquettes / day
 a = tons of alum. / day
 g = tons of graphite / day
 s = tons of silicon / day

aI = tons of alum. in ingot I / day
 aII = tons of alum. in ingot II / day
 gI = tons of graphite in ingot I / day
 gII = tons of graphite in ingot II / day
 SI = tons of silicon in ingot I / day
 SII = tons of silicon in ingot II / day
 I_1 = tons of ingot I / day
 I_2 = tons of ingot II / day.

Minimize $Z = 100S + 150A + 75C + 900A_b + 380S_b$
 s.t. $S \leq 1000$, $A \leq 500$, $C \leq 2500$

$$\begin{aligned} a &= .1S + .95A + A_b \\ g &= .05S + .01A + .15C \\ S &= .04S + .02A + .08C + S_b \end{aligned}$$

$$\begin{aligned} I_1 &= aI + gI + SI \\ I_2 &= aII + gII + SII \\ aI + gII &\leq 18, SI + SII \leq 8, gI + gII \leq 9 \\ .08I_1 &\leq aI \leq .108I_1 \\ .015I_1 &\leq gI \leq .03I_1 \\ .025I_1 &\leq SI < \infty \\ .062I_2 &\leq aII \leq .089I_2 \\ .041I_2 &\leq gII \leq \infty \\ .028I_2 &\leq SII \leq .041I_2 \\ I_1 &\geq 130, I_2 \geq 250 \end{aligned}$$

Optimum solution:

$$Z = \$117,435.65$$

$$S = 0, A = 38.2, C = 1489.41$$

$$A_b = S_b = 0$$

$$I_1 = 130, I_2 = 250$$

$$a = 36.29, g = 223.79, s = 119.92$$

9

10

x_{ij} = tons of ore i allocated to alloy k
 w_k = tons of alloy k produced

$$\begin{aligned} \text{Maximize } Z &= 200w_A + 300w_B \\ &- 30(x_{1A} + x_{1B}) \\ &- 40(x_{2A} + x_{2B}) \\ &- 50(x_{3A} + x_{3B}) \end{aligned}$$

Subject to

Specification constraints:

$$\begin{aligned} .2x_{1A} + .1x_{2A} + .05x_{3A} &\leq .8w_A \quad (1) \\ .1x_{1A} + .2x_{2A} + .05x_{3A} &\leq .3w_A \quad (2) \\ .3x_{1A} + .3x_{2A} + .2x_{3A} &\geq .5w_A \quad (3) \\ .1x_{1B} + .2x_{2B} + .05x_{3B} &\geq .4w_B \quad (4) \\ .1x_{1B} + .2x_{2B} + .05x_{3B} &\leq .6w_B \quad (5) \\ .3x_{1B} + .3x_{2B} + .7x_{3B} &\geq .3w_B \quad (6) \\ .3x_{1B} + .3x_{2B} + .2x_{3B} &\leq .7w_B \quad (7) \end{aligned}$$

Other constraints:

$$x_{1A} + x_{1B} \leq 1000$$

$$x_{2A} + x_{2B} \leq 2000$$

$$x_{3A} + x_{3B} \leq 3000$$

*** OPTIMUM SOLUTION SUMMARY ***

Title: Problem 26a-17
 Final iteration No: 12
 Objective value (max) = 400000.0000

Variable	Value	Obj Coeff	Obj Val Contrib
x1 wA	1799.9999	200.0000	359999.9688
x2 wB	1000.0001	300.0000	300000.0312
x3 x1A	1000.0000	-30.0000	-30000.0000
x4 x1B	0.0000	-30.0000	-0.0000
x5 x2A	0.0000	-40.0000	-0.0000
x6 x2B	2000.0001	-40.0000	-80000.0078
x7 x3A	3000.0000	-50.0000	-150000.0000
x8 x3B	0.0000	-50.0000	0.0000

Constraint	RHS	Slack(<)/Surplus(+)
1 (<)	0.0000	1090.0000-
2 (<)	0.0000	290.0000-
3 (>)	0.0000	0.0000+
4 (>)	0.0000	0.0000+
5 (<)	0.0000	200.0000-
6 (>)	0.0000	300.0002+
7 (<)	0.0000	100.0000-
8 (<)	1000.0000	0.0000-
9 (<)	2000.0000	0.0000-
10 (<)	3000.0000	0.0000-

Solution:

Produce 1800 tons of alloy A
 and 1000 tons of alloy B.

Set 2.3f

6

Let x_i = Nbr. starting on day i and lasting for 7 days

y_{ij} = Nbr. starting shift on day i and starting their 2 days off on day j, $i \neq j$

Thus, of the x_1 workers who start on Monday, y_{12} will take T and W off, y_{13} will take W and Th off, and so on, as the following table shows.

	x_1	x_2	x_3	x_4	x_5	x_6	x_7
1	start on Mon	y_{12}	$y_{12}+y_{13}$	$y_{13}+y_{14}$	$y_{14}+y_{15}$	$y_{15}+y_{16}$	y_{16}
2	y_{27}	Tue	y_{23}	$y_{23}+y_{24}$	$y_{24}+y_{25}$	$y_{25}+y_{26}$	$y_{26}+y_{27}$
3	$y_{31}+y_{37}$	y_{31}	Wed	y_{34}	$y_{34}+y_{35}$	$y_{35}+y_{36}$	$y_{36}+y_{37}$
4	$y_{41}+y_{47}$	$y_{41}+y_{42}$	y_{42}	Th	y_{45}	$y_{45}+y_{46}$	$y_{46}+y_{47}$
5	$y_{51}+y_{57}$	$y_{51}+y_{52}$	$y_{52}+y_{53}$	y_{53}	Fri	y_{56}	$y_{56}+y_{57}$
6	$y_{61}+y_{67}$	$y_{61}+y_{62}$	$y_{62}+y_{63}$	$y_{63}+y_{64}$	y_{64}	Sat	y_{67}
7	y_{71}	$y_{71}+y_{72}$	$y_{72}+y_{73}$	$y_{73}+y_{74}$	$y_{74}+y_{75}$	y_{75}	Su

Minimize $Z = x_1+x_2+x_3+x_4+x_5+x_6+x_7$

Each employee has 2 days off: $x_i = \sum\{j \text{ in } 1..7, j \neq i\} y_{ij}$

Mon (1) constraint: $s - (y_{27} + y_{31} + y_{37} + y_{41} + y_{47} + y_{51} + y_{57} + y_{61} + y_{67} + y_{71}) \geq 12$

Tue (2) constraint: $s - (y_{12} + y_{31} + y_{41} + y_{42} + y_{51} + y_{52} + y_{61} + y_{62} + y_{71} + y_{72}) \geq 18$

Wed (3) constraint: $s - (y_{12} + y_{13} + y_{23} + y_{42} + y_{52} + y_{53} + y_{62} + y_{63} + y_{72} + y_{73}) \geq 20$

Th (4) constraint: $s - (y_{13} + y_{14} + y_{23} + y_{24} + y_{24} + y_{53} + y_{63} + y_{64} + y_{73} + y_{74}) \geq 28$

Fri (5) constraint: $s - (y_{14} + y_{15} + y_{24} + y_{25} + y_{34} + y_{35} + y_{45} + y_{64} + y_{74} + y_{75}) \geq 32$

Sat(6) constraint: $s - (y_{15} + y_{16} + y_{25} + y_{26} + y_{35} + y_{36} + y_{45} + y_{46} + y_{56} + y_{75}) \geq 40$

Sun(7) constraint: $s - (y_{16} + y_{26} + y_{27} + y_{36} + y_{37} + y_{46} + y_{47} + y_{56} + y_{57} + y_{67}) \geq 40$

continued

Set 2.3f

Solution: 42 employees

Starting		Nbr off						
On	Nbr	M	Tu	Wed	Th	Fri	Sat	Sun
M	16		16	16				
Tu	8				8	8		
Wed	8	8	8					
Th	0							
Fri	6			6	6			
Sat	2	2					2	
Sun	2				2	2		
Nbr off		10	24	22	14	10	2	2
Nbr at work		32	18	20	28	32	40	40
Surplus above minimum		20	0	0	0	0	0	0

Set 2.3g

Setting		Number produced	Surplus rolls	1
1	3			
5'	0	200	50	
7'	1	200	0	
9'	1	300	0	
Loss/ft ²	4	1		
No. rolls	200	100		

Trim loss area =

$$L(200 \times 4 + 100 \times 1 + 50 \times 5) = 1150 \text{ ft}^2$$

(b) 15' standard roll:

Setting				
	1	2	3	4
5'	3	1	1	0
7'	0	1	0	2
9'	0	0	1	0
trim loss				
perf ft	0	3	1	1

$$(c) x_1 + x_2 + 2x_5 \geq 120$$

New solution calls for decreasing the number of standard 20'-rolls by 30

$$(d) x_1 + x_3 + 2x_6 \geq 240$$

New solution calls for increasing the number of standard 20'-rolls by 50

x_i = Space (in^2) allocated to cereal i

2

$$\text{Maximize } Z = 1.1x_1 + 1.3x_2 + 1.08x_3 + 1.25x_4 + 1.2x_5$$

s.t.

$$16x_1 + 24x_2 + 18x_3 + 22x_4 + 20x_5 \leq 5000$$

$$x_1 \leq 100, x_2 \leq 85, x_3 \leq 140, x_4 \leq 80, x_5 \leq 90$$

$$x_i \geq 0 \text{ for all } i = 1, 2, \dots, 5$$

Solution:

$$Z = \$314/\text{day}$$

$$x_1 = 100, x_3 = 140, x_5 = 44$$

$$x_2 = x_4 = 0$$

x_i = Nbr. of ads for issue i , $i = 1, 2, 3, 4$

3

$$\text{Minimize } Z = S_1^- + S_2^- + S_3^- + S_4^-$$

s.t.

$$(-30,000 + 60,000 + 30,000)x_1 + S_1^- - S_1^+ = .51 \times 400,000$$

$$(80,000 + 30,000 - 45,000)x_2 + S_2^- - S_2^+ = .51 \times 400,000$$

$$(40,000 + 10,000)x_3 + S_3^- - S_3^+ = .51 \times 400,000$$

$$(90,000 - 25,000)x_4 + S_4^- - S_4^+ = .51 \times 400,000$$

$$1500(x_1 + x_2 + x_3 + x_4) \leq 100,000$$

$$x_1, x_2, x_3, x_4 \geq 0$$

Solution:

$$x_1 = 3.4, x_2 = 3.14, x_3 = 4.08, x_4 = 3.14$$

x_{ij} = Units of part j produced by department i , $i = 1, 2, 3, j = 1, 2$

4

$$\text{Maximize } Z = \min \{x_{11} + x_{21}, x_{12} + x_{22}, x_{13} + x_{23}\}$$

or

$$\text{Maximize } Z = y$$

s.t.

$$y \leq x_{11} + x_{21}$$

$$y \leq x_{12} + x_{22}$$

$$y \leq x_{13} + x_{23}$$

$$\frac{x_{11}}{8} + \frac{x_{12}}{5} + \frac{x_{13}}{10} \leq 100$$

$$\frac{x_{21}}{6} + \frac{x_{22}}{12} + \frac{x_{23}}{4} \leq 80$$

$$\text{all } x_{ij} \geq 0$$

Solution:

$$\text{Nbr. of assembly units} = y = 556.2 \approx 557$$

$$x_{11} = 354.78, x_{21} = 0$$

$$x_{12} = 556.52, x_{22} = 201.74$$

$$x_{13} = 556.52, x_{23} = 0$$

x_i = tons of coal i , $i = 1, 2, 3$

5

$$\text{Minimize } Z = 30x_1 + 35x_2 + 33x_3$$

s.t.

$$2500x_1 + 1500x_2 + 1600x_3 \leq 2000(x_1 + x_2 + x_3)$$

$$x_1 \leq 30, x_2 \leq 30, x_3 \leq 30$$

$$x_1 + x_2 + x_3 \geq 50$$

Solution: $Z = \$1361.11$

$$x_1 = 22.22 \text{ tons}, x_2 = 0, x_3 = 27.78 \text{ tons.}$$

Set 2.3g

t_i = Green time in secs for highway i ,
 $i = 1, 2, 3$

$$\text{Maximize } Z = 3\left(\frac{500}{3600}\right)t_1 + 4\left(\frac{600}{3600}\right)t_2 + 5\left(\frac{400}{3600}\right)t_3$$

s.t.

$$\left(\frac{500}{3600}\right)t_1 + \left(\frac{600}{3600}\right)t_2 + \left(\frac{400}{3600}\right)t_3 \leq \frac{510}{3600} (2.2 \times 60 - 3 \times 10)$$

$$t_1 + t_2 + t_3 + 3 \times 10 \leq 2.2 \times 60, t_1 \geq 25, t_2 \geq 25, t_3 \geq 25$$

Solution: $Z = \$58.04/\text{hr}$

$$t_1 = 25, t_2 = 43.6, t_3 = 33.4 \text{ sec}$$

6

Cost (\$) per cubic yd:

	(5) A2	(6) A4
(1) A1	$.2 + 2x \cdot .15 = .50$	$.20 + 7x \cdot .15 = 1.25$
(2) A3	$.20 + 2x \cdot .15 = .50$	$.20 + 3x \cdot .15 = .65$
(3) P1	$1.70 + 3x \cdot .15 = 2.15$	$1.70 + 8x \cdot .15 = 2.90$
(4) P3	$2.10 + 7x \cdot .15 = 3.15$	$2.10 + 2x \cdot .15 = 2.40$

Using the code $A1 \equiv 1, A3 \equiv 2, P1 \equiv 3, P3 \equiv 4, A2 \equiv 5, A4 \equiv 6$, let

$x_{ij} = 10^3 \text{ Yd}^3$ from source i to destination j
 $i = 1, 2, 3, 4, j = 5, 6$

$$\text{Minimize } Z = 1000(.5x_{15} + 1.25x_{16} + .5x_{25} + .65x_{26} + 2.15x_{35} + 2.9x_{36} + 3.15x_{45} + 2.4x_{46})$$

s.t.

$$x_{15} + x_{16} \leq 1760 \quad x_{35} + x_{36} \leq 20,000$$

$$x_{25} + x_{26} \leq 1760 \quad x_{45} + x_{46} \leq 15,000$$

$$x_{15} + x_{25} + x_{35} + x_{45} \geq 3520$$

$$x_{16} + x_{26} + x_{36} + x_{46} \geq 3520$$

Solution:

$$A1 \rightarrow A2: x_{15} = 1760 \text{ (1000 Cu Yd)}$$

$$A1 \rightarrow A4: x_{16} = 0$$

$$A3 \rightarrow A2: x_{25} = 0$$

$$A3 \rightarrow A4: x_{26} = 1760$$

$$P1 \rightarrow A2: x_{35} = 1760$$

$$P1 \rightarrow A4: x_{36} = 0$$

$$P2 \rightarrow A2: x_{45} = 0$$

$$P2 \rightarrow A4: x_{46} = 1760$$

$$\text{Cost} = \$10,032,000$$

8



$$A1 = 2 \times 1760 \times 10 \times 50 = 1760 \text{ (thousand) Yd}^3$$

$$A2 = 3520, A3 = 1760, A4 = 3520$$

Distances (center to center) in miles:

	A2	A4
A1	2	7
A3	2	3
P1	3	8
P2	7	2

continued...

9

x_{ij} = Blue regulars on front i in defense line $j, i =$

y_{ij} = Blue reserves on front i in defense line j .

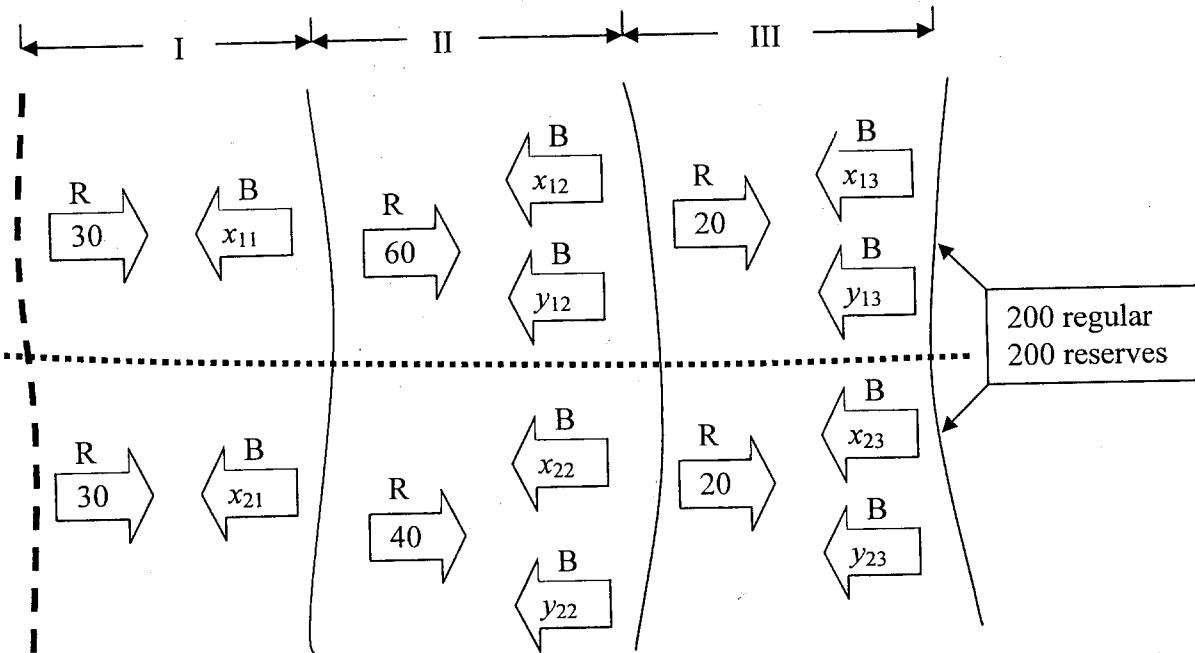
t_{ij} = Delay days on front i in defense line j :

$$\text{Maximize } Z = \min \{t_{11} + t_{12} + t_{13}, t_{21} + t_{22} + t_{23}\}$$

or

continued...

Set 2.3g



$$\text{Maximize } Z = T$$

s.t.

$$T \leq t_{11} + t_{12} + t_{13}$$

$$T \leq t_{21} + t_{22} + t_{23}$$

$$x_{11} + x_{12} + x_{13} + x_{21} + x_{22} + x_{23} \leq 200$$

$$y_{12} + y_{13} + y_{22} + y_{23} \leq 200$$

$$t_{11} = .5 + 8.8 \frac{x_{11}}{30}$$

$$t_{12} = .75 + 7.9 \frac{x_{12} + y_{12}}{60}$$

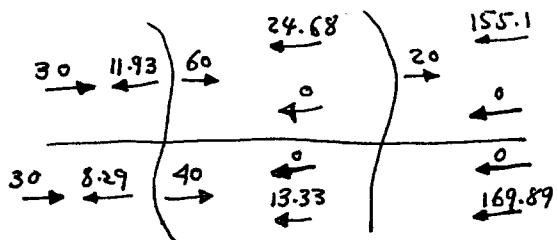
$$t_{13} = .55 + 10.2 \frac{x_{13} + y_{13}}{20}$$

$$t_{21} = 1.1 + 10.5 \frac{x_{21}}{30}$$

$$t_{22} = 1.3 + 8.1 \frac{x_{22} + y_{22}}{40}$$

$$t_{23} = 1.5 + 9.2 \frac{x_{23} + y_{23}}{20}$$

Solution: Battle duration = 87.65 days



continued...

2-32

x_i = Efficiency of plant i

10

$$\text{Minimize } Z = .2(500)x_1 + .25(3000)x_2 + .15(6000)x_3 + .18(1000)x_4$$

s.t.

$$500(1-x_1) \leq .00085 \times 215,000$$

$$.94(500)(1-x_1) + 3000(1-x_2) \leq .0009 \times 220,000$$

$$.94^2(500)(1-x_1) + .94(3000)(1-x_2) + 6000(1-x_3) \leq .0008 \times 200,000$$

$$.94^3(500)(1-x_1) + .94^2(3000)(1-x_2) + .94(6000)(1-x_3) + 1000(1-x_4) \leq .0008 \times 210,000$$

$$0 \leq x_1 \leq .99$$

$$0 \leq x_2 \leq .99$$

$$0 \leq x_3 \leq .99$$

$$0 \leq x_4 \leq .99$$

Solution

Cost per hour = \$1891.41

Plant 1 efficiency = .99

Plant 2 efficiency = .9661

Plant 3 efficiency = .99

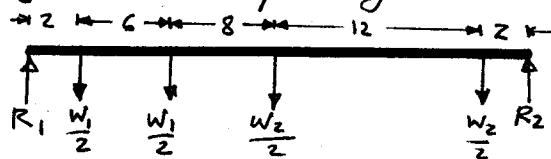
Plant 4 efficiency = .9824

Set 2.3g

w_i = Capacity of yoke i (Kips)

R_1 = Reaction in Kips at left end

R_2 = Reaction in Kips at right end



$$\text{Maximize } Z = w_1 + w_2$$

s.t.

$$R_1 + R_2 = w_1 + w_2$$

$$2\left(\frac{w_1}{2}\right) + 8\left(\frac{w_1}{2}\right) + 16\left(\frac{w_2}{2}\right) + 28\left(\frac{w_2}{2}\right) \\ = 30R_2$$

$$R_1 \leq 25, \quad R_2 \leq 25$$

$$\frac{w_1}{2} \leq 20, \quad \frac{w_2}{2} \leq 20$$

Solution:

$$w_1 = 20.59 \text{ Kips}$$

$$w_2 = 29.41 \text{ Kips}$$

x_{ij} = Nbr. of aircraft of type i allocated to route j
($i = 1, 2, 3, 4$, $j = 1, 2, 3, 4$)

s_j = Nbr. of passengers not served on route j , $j = 1, 2, 3, 4$

$$\begin{aligned} \text{Minimize } Z &= 1000(3x_{11}) + 1100(2x_{12}) \\ &+ 1200(2x_{13}) + 1500(x_{14}) \\ &+ 800(4x_{21}) + 900(3x_{22}) \\ &+ 1000(3x_{23}) + 1000(2x_{24}) \\ &+ 600(5x_{31}) + 800(5x_{32}) \\ &+ 800(4x_{33}) + 900(2x_{34}) \\ &+ 40S_1 + 50S_2 + 45S_3 + 70S_4 \end{aligned}$$

Subject to

$$\sum_{j=1}^4 x_{1j} \leq 5, \quad \sum_{j=1}^4 x_{2j} \leq 8, \quad \sum_{j=1}^4 x_{3j} \leq 10$$

$$50(3x_{11}) + 30(4x_{21}) + 20(5x_{31}) + S_1 = 1000$$

$$50(2x_{12}) + 30(3x_{22}) + 20(5x_{32}) + S_2 = 2000$$

$$50(2x_{13}) + 30(3x_{23}) + 20(4x_{33}) + S_3 = 900$$

$$50(x_{14}) + 30(2x_{24}) + 20(2x_{34}) + S_4 = 1200$$

All x_{ij} and $s_j \geq 0$

continued...

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*** OPTIMUM SOLUTION SUMMARY ***

Title: Problem 26a-16
Final iteration No: 16
Objective value (min) = 221900.0000
=> ALTERNATIVE solution detected at x13

Variable	Value	Obj Coeff	Obj Val Contrib
x1 x11	5.0000	3000.0000	14999.9990
x2 x12	0.0000	2200.0000	0.0000
x3 x13	0.0000	2400.0000	0.0000
x4 x14	0.0000	1500.0000	0.0000
x5 x21	0.0000	3200.0000	0.0000
x6 x22	0.0000	2700.0000	0.0000
x7 x23	0.0000	3000.0000	0.0000
x8 x24	8.0000	2000.0000	15999.9990
x9 x31	2.5000	3000.0000	7500.0015
x10 x32	7.5000	4000.0000	29999.9980
x11 x33	0.0000	3200.0000	0.0000
x12 x34	0.0000	1800.0000	0.0000
x13 s1	0.0000	40.0000	0.0000
x14 s2	1250.0000	50.0000	62500.0000
x15 s3	899.9998	45.0000	40499.9922
x16 s4	720.0001	70.0000	50400.0078

Constraint	RHS	Slack(-)/Surplus(+)
1 (<)	5.0000	0.0000-
2 (<)	8.0000	0.0000-
3 (=)	10.0000	0.0000-
4 (=)	1000.0000	0.0000
5 (=)	2000.0000	0.0000
6 (=)	900.0000	0.0000
7 (=)	1200.0000	0.0000

Solution:

Aircraft Type	Route	Nbr. Aircraft
1	1	5
2	4	8
3	1	2.5
3	2	7.5

Fractional solution must be rounded.

Cost = \$ 221,900

12

2-33

CHAPTER 3

The Simplex Method and Sensitivity Analysis

Set 3.1a

$$(x_1, x_2) = (3, 1)$$

$$M1: S_1 = 24 - (6x_3 + 4x_1) = 2 \text{ tons/day}$$

$$M2: S_2 = 6 - (1x_3 + 2x_1) = 1 \text{ ton/day}$$

$$S_1 = x_1 + x_2 - 800$$

$$= 500 + 600 - 800 = 300 \text{ lb}$$

$$10x_1 - 3x_2 \geq -5 \equiv -10x_1 + 3x_2 \leq 5$$

$$\text{Thus, } -10x_1 + 3x_2 + S_1 = 5 \quad ①$$

$$\text{Also, } 10x_1 - 3x_2 \geq -5 \equiv 10x_1 - 3x_2 - S_2 = -5$$

$$\text{Thus, } -10x_1 + 3x_2 + S_2 = 5 \quad ②$$

① and ② are the same

x_{ij} = number of units of product **4**
 i manufactured on machine j

LP model

$$\text{Maximize } Z = 10(x_{11} + x_{12}) + 15(x_{21} + x_{22})$$

Subject to

$$|(x_{11} + x_{21}) - (x_{12} + x_{22})| \leq 5$$

$$x_{11} + x_{21} \leq 200$$

$$x_{12} + x_{22} \leq 250$$

$$x_{ij} \geq 0 \text{ for all } i \neq j$$

Equation form:

$$|(x_{11} + x_{21}) - (x_{12} + x_{22})| \leq 5$$

to

$$x_{11} + x_{21} - x_{12} - x_{22} \leq 5$$

$$x_{11} + x_{21} - x_{12} - x_{22} \geq -5$$

$$\text{Maximize } Z = 10x_{11} + 10x_{12} + 15x_{21} + 15x_{22}$$

Subject to

$$x_{11} + x_{21} - x_{12} - x_{22} + S_1 = 5$$

$$-x_{11} - x_{21} + x_{12} + x_{22} + S_2 = 5$$

$$x_{11} + x_{21} + S_3 = 200$$

$$x_{12} + x_{22} + S_4 = 250$$

$$x_{ij} \geq 0 \text{ for all } i \text{ and } j$$

$$S_i \geq 0 \text{ for all } i$$

continued...

$$y = \max \{|x_1 - x_2 + 3x_3|, |-x_1 + 3x_2 - x_3|\}$$

Hence

$$|x_1 - x_2 + 3x_3| \leq y$$

$$|-x_1 + 3x_2 - x_3| \leq y$$

LP model:

$$\text{minimize } Z = y$$

Subject to

$$x_1 - x_2 + 3x_3 \leq y$$

$$x_1 - x_2 + 3x_3 \geq -y$$

$$-x_1 + 3x_2 - x_3 \leq y$$

$$-x_1 + 3x_2 - x_3 \geq -y$$

$$x_1, x_2, x_3, y \geq 0$$

Equation form:

$$\text{Minimize } Z = y$$

Subject to

$$-y + x_1 - x_2 + 3x_3 + S_1 = 0$$

$$-y - x_1 + x_2 - 3x_3 + S_2 = 0$$

$$-y - x_1 + 3x_2 - x_3 + S_3 = 0$$

$$-y + x_1 - 3x_2 + x_3 + S_4 = 0$$

$$x_1, x_2, x_3, y, S_1, S_2, S_3, S_4 \geq 0$$

$$\sum_{j=1}^n a_{ij} x_j = b_i \Leftrightarrow \begin{cases} \sum_{j=1}^n a_{ij} x_j \leq b_i & ① \\ \sum_{j=1}^n a_{ij} x_j \geq b_i & ② \end{cases} \quad 6$$

From ②, for $i = 1, 2, \dots, m$, we have

$$\sum_{j=1}^n a_{ij} x_j \geq b_i \Leftrightarrow \sum_{i=1}^m \left(\sum_{j=1}^n a_{ij} x_j \right) \geq \sum_{i=1}^m b_i$$

$$\Leftrightarrow \sum_{j=1}^n \left(\sum_{i=1}^m a_{ij} \right) x_j \geq \sum_{i=1}^m b_i$$

Thus, ① and ② are equivalent to

$$\sum_{j=1}^n a_{ij} x_j \leq b_i, \quad i = 1, 2, \dots, m$$

$$\sum_{j=1}^n \left(\sum_{i=1}^m a_{ij} \right) x_j \geq \sum_{i=1}^m b_i$$

Set 3.1b

$$X_1 = \text{Nbr. } \frac{1}{4} - \text{lb / day}$$

$$X_2 = \text{Nbr. cheeseburgers/day}$$

$$\text{Maximize } Z = .2X_1 + .15X_2 - .25X_3^+$$

s.t.

$$.25X_1 + .2X_2 + X_3^- - X_3^+ = 200$$

$$X_1 + X_2 \leq 900$$

$$\underline{\text{Solution: }} Z = \$173.35$$

$$X_1 = 900, X_2 = 0, X_3^+ = 25 \text{ lb}$$

$$(a) X_j = \# \text{ units of product } j \text{ per day}, j=1,2$$

$$X_3^+ = \text{unused minutes of machine time / day}$$

$$X_3^- = \text{machine overtime per day in minutes}$$

$$\text{Maximize } Z = 6X_1 + 7.5X_2 - .5X_3^-$$

Subject to

$$10X_1 + 12X_2 + X_3^+ - X_3^- = 2500$$

$$150 \leq X_1 \leq 200$$

$$X_2 \leq 45$$

$$X_1, X_2 \geq 0$$

$$X_3^+, X_3^- \geq 0$$

TORA optimum solution:

$$X_1 = 200 \text{ units/day}$$

$$X_2 = 45 \text{ units/day}$$

$$X_3^- = \text{overtime minutes} \\ = 40 \text{ minutes/day}$$

$$Z = \$1517.50$$

(b) Overtime at \$1.50/min yields $X_3^- = 0$, which means no overtime is needed

$$X_j = \# \text{ of units of products } 1, 2, \text{ and } 3$$

$$\text{Maximize } Z = 2X_1 + 5X_2 + 3X_3 - 15X_4^+ - 10X_5^+$$

Subject to

$$2X_1 + X_2 + 2X_3 + X_4^+ - X_4^- = 80$$

$$X_1 + X_2 + 2X_3 + X_5^+ - X_5^- = 65$$

$$\text{all variables } \geq 0$$

$$\underline{\text{Solution: }} Z = \$325$$

$$X_2 = 65 \text{ units}, X_4^- = 15$$

$$\text{All other variables} = 0$$

continued...

4

Formulation 1:

$$\text{Maximize } Z = -2X_1 + 3X_2^+ - 3X_2^- - 2X_3^+ + 2X_3^-$$

Subject to

$$4X_1 - X_2^+ + X_2^- - 5X_3^+ + 5X_3^- = 10$$

$$2X_1 + 3X_2^+ - 3X_2^- + 2X_3^+ - 2X_3^- = 12$$

all variables ≥ 0

Optimum solution:

$$X_1 = 0$$

$$X_2^+ = 6.15 \quad \} \Rightarrow X_2 = 6.15$$

$$X_2^- = 0$$

$$X_3^+ = 0 \quad \} \Rightarrow X_3 = -3.23$$

$$X_3^- = 3.23$$

$$Z = 24.92$$

Formulation 2:

$$\text{Maximize } Z = -2X_1 + 3X_2^+ - 2X_3^+ - w$$

Subject to

$$4X_1 - X_2^+ - 5X_3^+ + 6w = 10$$

$$2X_1 + 3X_2^+ + 2X_3^+ - 5w = 12$$

all variables ≥ 0

Optimum solution:

$$X_1 = 0$$

$$X_2^+ = 9.38 \quad \} \Rightarrow X_2 = 9.38 - 3.23 = 6.15$$

$$w = 3.23$$

$$X_3^+ = 0 \quad \} \Rightarrow X_3 = 0 - 3.23 = -3.23$$

$$w = 3.23$$

$$Z = 24.92$$

continued...

Set 3.2a

(a)

Equation form:

$$\text{Maximize } Z = 2x_1 + 3x_2$$

Subject to

$$x_1 + 3x_2 + x_3 = 6$$

$$3x_1 + 2x_2 + x_4 = 6$$

$$x_1, x_2, x_3, x_4 \geq 0$$

(b) Basic (x_1, x_2) (Point B):

$$x_1 + 3x_2 = 6$$

$$3x_1 + 2x_2 = 6$$

$$\text{Solution: } (x_1, x_2) = \left(\frac{6}{7}, \frac{12}{7}\right), Z = 6 \frac{6}{7}$$

Basic (x_1, x_3) (Point E):

$$x_1 + x_3 = 6$$

$$3x_1 = 6$$

$$\text{Solution: } (x_1, x_3) = (2, 4), Z = 4$$

Basic (x_1, x_4) (Point C):

$$x_1 = 6$$

$$3x_1 + x_4 = 6$$

$$\text{Solution: } (x_1, x_4) = (6, -12)$$

Unique but infeasible

Basic (x_2, x_3) (Point A):

$$3x_2 + x_3 = 6$$

$$2x_2 = 6$$

$$\text{Solution: } (x_2, x_3) = (3, -3)$$

Unique but infeasible

Basic (x_2, x_4) (Point D):

$$3x_2 = 6$$

$$2x_2 + x_4 = 6$$

$$\text{Solution: } (x_2, x_4) = (2, 2), Z = 6$$

Basic (x_3, x_4) (Point F):

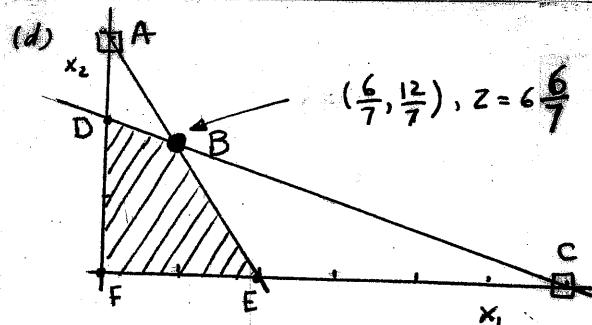
$$x_3 = 6$$

$$x_4 = 6$$

$$\text{Solution: } (x_3, x_4) = (6, 6), Z = 0$$

(c) Optimum solution occurs at B:

$$(x_1, x_2) = \left(\frac{6}{7}, \frac{12}{7}\right) \text{ with } Z = 6 \frac{6}{7}$$



(e) From the graph in (d), we have

$$A: x_2 = 3, x_3 = -3$$

$$C: x_1 = 6, x_4 = -12$$

2

(a) Maximize $Z = 2x_1 - 4x_2 + 5x_3 - 6x_4$

Subject to

$$x_1 + 4x_2 - 2x_3 + 8x_4 + x_5 = 2$$

$$-x_1 + 2x_2 + 3x_3 + 4x_4 + x_6 = 1$$

$$x_1, x_2, x_3, x_4, x_5, x_6 \geq 0$$

Combination	Solution	Status	Z
x_1, x_2	$0, 1/2$	Feasible	-2
x_1, x_3	$8, 3$	Feasible	31
x_1, x_4	$0, 1/4$	Feasible	-3/2
x_1, x_5	$-1, 3$	Infeasible	-
x_1, x_6	$2, 3$	Feasible	4
x_2, x_3	$1/2, 0$	Feasible	-2
x_2, x_4	$1/2, 0$	Feasible	-2
x_2, x_5	$1/2, 0$	Feasible	-2
x_2, x_6	$1/2, 0$	Feasible	-2
x_3, x_4	$0, 1/4$	Feasible	-3/2
x_3, x_5	$1/3, 8/3$	Feasible	5/3
x_3, x_6	$-1, 4$	Infeasible	-
x_4, x_5	$1/4, 0$	Feasible	-3/2
x_4, x_6	$1/4, 0$	Feasible	-3/2
x_5, x_6	$2, 1$	Feasible	0

Optimum Solution:

$$x_1 = 8, x_2 = 0, x_3 = 3, x_4 = 0$$

$$Z = 31$$

continued...

Set 3.2a

(b) Minimize $Z = x_1 + 2x_2 - 3x_3 - 2x_4$
Subject to

$$x_1 + 2x_2 - 3x_3 + x_4 = 4$$

$$x_1 + 2x_2 + x_3 + 2x_4 = 4$$

$$x_1, x_2, x_3, x_4 \geq 0$$

Combination Solution Status Z

x_1, x_2	infinity of solutions	-
x_1, x_3	4, 0	Feasible 4
x_1, x_4	4, 0	Feasible 4
x_2, x_3	2, 0	Feasible 4
x_2, x_4	2, 0	Feasible 4
x_3, x_4	$-\frac{4}{7}, \frac{16}{7}$	Infeasible -

Alternative optima:

x_1	x_2	x_3	x_4	Z
4	0	0	0	4
0	2	0	0	4

maximize $Z = x_1 + x_2$
Subject to

$$x_1 + 2x_2 + x_3 = 6$$

$$2x_1 + x_2 - x_4 = 16$$

$$x_1, x_2, x_3, x_4 \geq 0$$

Combination Solution Status

x_1, x_2	$2\frac{6}{3}, -4\frac{1}{3}$	Infeasible
x_1, x_3	8, -2	Infeasible
x_1, x_4	6, -4	Infeasible
x_2, x_3	16, -26	Infeasible
x_2, x_4	3, -13	Infeasible
x_3, x_4	6, -16	Infeasible

3

Maximize $Z = 2x_1 + 3x_2 - 3x_2^+ - 5x_3^-$

4

Subject to

$$-6x_1 + 7x_2^- - 7x_2^+ - 9x_3^- - x_4 = 4$$

$$x_1 + x_2^- - x_2^+ + 4x_3^- = 10$$

$$x_1, x_2^-, x_2^+, x_3^-, x_4 \geq 0$$

(x_2^-, x_2^+) :

$$7x_2^- - 7x_2^+ = 4$$

$$x_2^- - x_2^+ = 10$$

Since $(7x_2^- - 7x_2^+)$ and $(x_2^- - x_2^+)$ are dependent, it is impossible for x_2^- and x_2^+ to be basic simultaneously. This means that at least x_2^- and x_2^+ must be nonbasic at zero level; thus making it impossible for x_2^- and x_2^+ to assume positive values simultaneously in any basic solution.

maximize $Z = x_1 + 3x_2$

5

Subject to

$$x_1 + x_2 + x_3 = 2$$

$$-x_1 + x_2 + x_4 = 4$$

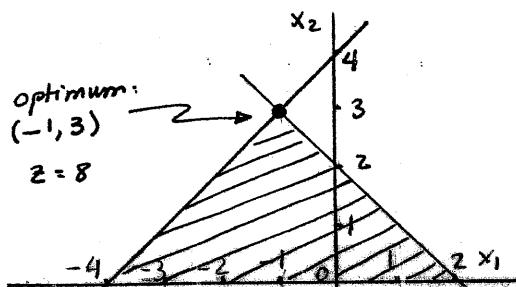
x_i unrestricted

$$x_2, x_3 \geq 0$$

Combination	Solution	Status	Z
x_1, x_2	-1, 3	Feasible	8
x_1, x_3	-4, 6	Feasible	-4
x_1, x_4	2, 6	Feasible	2
x_2, x_3	4, -2	Infeasible	-
x_2, x_4	2, 2	Feasible	6
x_3, x_4	2, 4	Feasible	0

Optimum: $x_1 = -1, x_2 = 3, Z = 8$

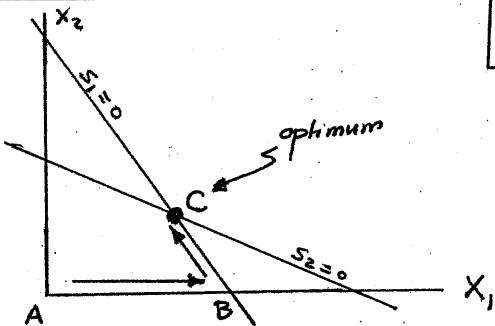
(c)



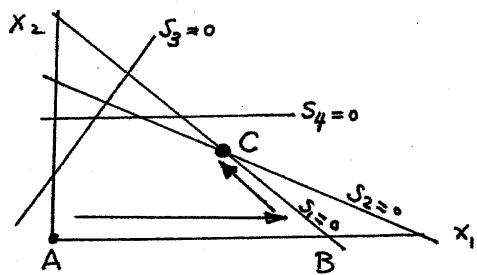
continued...

3-5

Set 3.3a



Extreme point	Basic	Nonbasic
A	S_1, S_2	X_1, X_2
B	X_1, S_2	X_2, S_1
C	X_1, X_2	S_1, S_2



Extreme point	Basic	Nonbasic
A	S_1, S_2, S_3, S_4	X_1, X_2
B	X_1, S_2, S_3, S_4	S_1, X_2
C	X_1, X_2, S_3, S_4	S_1, S_2

- (a) (A, B) adjacent, hence can be on a simplex path. Remaining pairs cannot be on a simplex path because they are not adjacent.
- (b) (i) Yes, because connects adjacent extreme points
(ii) No, because C and I are not adjacent.
(iii) No, because the path returns to a previous extreme point.

Extreme Point	Basic	Nonbasic
A	S_1, S_2, S_3, S_4	X_1, X_2, X_3
B	S_1, X_1, S_3, S_4	S_2, X_2, S_3
C	X_1, S_2, S_3, S_4	S_1, X_1, X_3
D	S_1, S_2, X_3, S_4	X_1, X_2, S_3
E	X_1, X_2, S_3, S_4	S_1, S_2, X_3
F	X_1, S_2, X_3, S_4	X_1, S_1, S_3
G	S_1, X_1, X_3, S_4	S_2, X_2, S_3
H	S_1, X_1, X_2, X_3	S_2, S_3, S_4
I	X_1, X_2, X_3, S_3	S_1, S_2, S_4
J	X_1, S_2, X_2, X_3	S_1, S_3, S_4

- (a) x_3 enters at value 1
 $Z = 0 + 3 \times 1 = 3$
- (b) x_1 enters at value 1
 $Z = 0 + 5 \times 1 = 5$
- (c) x_2 enters at value 1
 $Z = 0 + 7 \times 1 = 7$
- (d) Tie broken arbitrarily between x_1, x_2 , and x_3 . Entering value = 1
 $Z = 0 + 1 \times 1 = 1$

3

Set 3.3b

								1
Basic	Z	X ₁	X ₂	S ₁	S ₂	S ₃	S ₄	Sol
Z	1	-5	-4	0	0	0	0	0
S ₁	0	6	4	1	0	0	0	24
S ₂	0	1	2	0	1	0	0	6
S ₃	0	-1	1	0	0	1	0	1
S ₄	0	0	1	0	0	0	1	2
Z	1	0	6	0	5	0	0	30
S ₁	0	0	-8	1	-6	0	0	-12
X ₁	0	1	2	0	1	0	0	6
S ₃	0	0	3	0	1	1	0	7
S ₄	0	0	1	0	0	0	1	2

								2	
(a)	Basic	x ₁	x ₂	x ₃	x ₄	sx ₅	sx ₆	sx ₇	Solution
Z	-2.00	-1.00	3.00	-5.00	0.00	0.00	0.00	0.00	0.00
1)sx ₅	1.00	2.00	2.00	4.00	1.00	0.00	0.00	40.00	
2)sx ₆	2.00	-1.00	1.00	2.00	0.00	1.00	0.00	8.00	
3)sx ₇	4.00	-2.00	1.00	-1.00	0.00	0.00	1.00	10.00	
Z	3.00	-3.50	5.50	0.00	0.00	2.50	0.00	20.00	
1)sx ₅	-3.00	4.00	0.00	0.00	1.00	-2.00	0.00	24.00	
2)sx ₄	1.00	-0.50	0.50	1.00	0.00	0.50	0.00	4.00	
3)sx ₇	5.00	-2.50	1.50	0.00	0.00	0.50	1.00	14.00	
Z	0.38	0.00	5.50	0.00	0.88	0.75	0.00	41.00	
1)sx ₂	-0.75	1.00	0.00	0.00	0.25	-0.50	0.00	6.00	
2)sx ₄	0.62	0.00	0.50	1.00	0.12	0.25	0.00	7.00	
3)sx ₇	3.12	0.00	1.50	0.00	0.62	-0.75	1.00	29.00	

								3
Basic	Ratios							
Basic	X ₁	X ₂	X ₃	X ₄				
X ₅	4/1	4/2	--					4/5
X ₆	8/5	--	--					8/6
X ₇	3/2	3/3	--					3/3
X ₈	--	--	0/1	--				
Value	1.5	1	0	0.8				
Leaving var	X ₇	X ₇	X ₈	X ₅				

(a) Nonbasic x_1 will improve solution.								
<u>Basic x_i-ratios</u>								
$x_2 \quad (4/5) \Rightarrow x_2 \text{ leaves}, x_i = \frac{4}{5}$								

$X_1 = \frac{4}{5} = .8$	$X_3 = 8 - 6 \times .8 = 3.6$	$X_4 = 3 - 3 \times .8 = .6$
$X_2 = 0$	$Z = .8 \times 1 = .8$	

(b) x_1 remains nonbasic at zero. Current solution, $X_2 = 4$, $X_3 = 8$, $X_4 = 3$, $Z = 0$ is optimum								
--	--	--	--	--	--	--	--	--

Basic solutions consist of one variable each. Thus,								
$X_1 = 90/1 = 90$	$Z = 5 \times 90 = 450$							
$X_2 = 90/3 = 30$	$Z = -6 \times 30 = -180$							
$X_3 = 90/5 = 18$	$Z = 3 \times 18 = 54$							
$X_4 = 90/6 = 15$	$Z = -5 \times 15 = -75$							
$X_5 = 90/3 = 30$	$Z = 12 \times 30 = 360$							

Optimum solution:								
$X_1 = 90, X_2 = X_3 = X_4 = X_5 = 0, Z = 450$								

(a) Basic: $(X_8, X_3, X_1) = (12, 6, 0)$, $Z = 620$								
Nonbasic: $(X_2, X_4, X_5, X_6, X_7) = (0, 0, 0, 0, 0)$								
(b) X_2, X_5, X_6 will improve solution.								

<u>X_2 enters: $X_2 = \min(\frac{12}{3}, \frac{6}{1}, -)$ = 4. Thus,</u>								
<u>X_8 leaves, $\Delta Z = 4 \times 5 = 20$</u>								

continued...

Set 3.3b

x_5 enters: $x_5 = \min(-, \frac{6}{1}, \frac{0}{6}) = 0$. Thus,
 $\Delta Z = 1 \times 0 = 0$ (x_5 leaves)

x_6 enters: $x_6 = \min(-, -, -)$. Thus, no
leaving variable and x_6 can
be increased to ∞ . $\Delta Z = +\infty$

(C) x_4 can improve solution.

x_4 enters: $x_4 = \min(-, \frac{6}{3}, -) = 2$. Thus,
 x_3 leaves. $\Delta Z = -4 \times 2 = -8$

(d) As shown in (b), x_5 cannot
change Z because it enters the
solution at level zero. x_7 cannot change
 Z either because its objective equation
coefficient = 0. $\Delta Z = 0 \times \min(\frac{12}{5}, \frac{6}{3}, -) = 0$

(a) Maximize $Z = 3x_1 + 6x_2$: 7

x_2 is the first entering variable.
Resulting path is $A \rightarrow G \rightarrow F \rightarrow E$.

(b) Maximize $Z = 4x_1 + x_2$:

Entering variable $x_1 = \min$ (intercept with
 x_1 -axis)

$$x_1 = \min(2, 3, 5) = 2 \text{ at } B$$

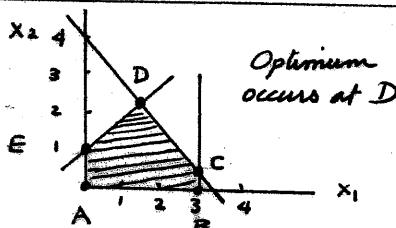
$$\Delta Z = 4 \times 2 = 8$$

(c) Maximize $Z = x_1 + 4x_2$:

Entering variable $x_2 = \min$ (intercept with
 x_2 -axis)

$$x_2 = \min(1, 2, 4) = 1$$

$$\Delta Z = 4 \times 1 = 4$$



8

(a) x_1 will enter first and the iterations
will follow the path $A \rightarrow B \rightarrow C \rightarrow D$

(b) x_2 enters first and the iterations will
follow the path $A \rightarrow E \rightarrow D$

(c) The most-negative criterion requires
more iterations (4 vs. 3). This criterion
is only a heuristic, and although it does
not guarantee the smallest number of

iterations, computational experience
demonstrates that, on the average,
the most-negative criterion is
more efficient.

(d) Iterations are identical, with
the exception of the objective row, which
should appear with an opposite sign

Optimum tableau:

Basic	x_1	x_2	s_1	s_2	s_3	s_4	
Z	0	0	$\frac{3}{4}$	$\frac{1}{2}$	0	0	21
x_1	1	0	$\frac{1}{4}$	$-\frac{1}{2}$	0	0	3
x_2	0	1	$-\frac{1}{8}$	$\frac{3}{4}$	0	0	$\frac{3}{2}$
s_3	0	0	$\frac{3}{8}$	$-\frac{5}{4}$	1	0	$\frac{5}{2}$
s_4	0	0	$\frac{1}{8}$	$-\frac{3}{4}$	0	1	$\frac{1}{2}$

9

If s_5 enters, its value = $\min\{\frac{3}{1/4}, -\frac{5/2}{3/4}, \frac{1/2}{1/8}\} = 4$
New $Z = 21 - 3/4 \times 4 = 18$

If s_2 enters, its value = $\min\{-\frac{3/2}{3/4}, -\frac{1}{2}\} = 2$
New $Z = 21 - 1/2 \times 2 = 20$. The second best Z is
associated with s_2 entering the basic solution

Not easily extendable because the third
best solution may not be an adjacent
corner point of the current optimum point.

10

x_1 = number of purses per day

x_2 = number of bags per day

x_3 = number of backpacks per day

$$\text{Maximize } Z = 24x_1 + 22x_2 + 45x_3$$

Subject to

$$2x_1 + x_2 + 3x_3 \leq 42$$

$$2x_1 + x_2 + 2x_3 \leq 40$$

$$x_1 + 5x_2 + x_3 \leq 45$$

$$x_1, x_2, x_3 \geq 0$$

TORA's optimum solution:

$$x_1 = 0, x_2 = 36, x_3 = 2, Z = \$882$$

Status of resources:

Resource	slack	Status
Leather	0	scarce
Sewing	0	scarce
Finishing	25	abundant

continued...

From TORA Iterations module, **12**

Click **All Iterations**, then go to the optimal iteration and click any of the associated nonbasic variables (X_4, Sx_6, Sx_7, Sx_8). Now, click **Next Iteration** to produce the new iteration in which the selected variable becomes basic. The associated value of Z will deteriorate.

To determine the next-best solution, follow the procedure in Problem 1. First, let X_4 enter the basic solution and record the associated value of Z . Next, click **View/Modify Input Data** and re-solve the problem to produce the same optimum tableau that was used before X_4 was entered into the basic solution. Now, enter Sx_6 into the basic solution and record the associated value of Z . Repeat the procedure for Sx_7 and Sx_8 . You will get the following results:

Entering Variable	Z
X_4	2.63
Sx_6	1.00
Sx_7	6.40
Sx_8	1.90

The next-best solution is associated with entering Sx_7 into the basic solution. Associated values of the variables are

$$X_1 = 1.6$$

$$X_2 = 0$$

$$X_3 = 1.6$$

$$X_4 = 0$$

$$Z = 6.40$$

Set 3.4a

Iteration	Basic	x_1	x_2	x_3	R_1	R_2	x_4	Solution
0 (starting)	z	$-4 + 7M$	$-1 + 4M$	$-M$	0	0	0	$9M$
	R_1	3	1	0	1	0	0	3
	R_2	4	3	-1	0	1	0	6
	x_4	1	2	0	0	0	1	4
1	z	0	$\frac{1+5M}{3}$	$-M$	$\frac{4-7M}{3}$	0	0	$4+2M$
	x_1	1	$\frac{1}{3}$	0	$\frac{1}{3}$	0	0	1
	R_2	0	$\frac{5}{3}$	-1	$-\frac{4}{3}$	1	0	2
	x_4	0	$\frac{5}{3}$	0	$-\frac{1}{3}$	0	1	3
2	z	0	0	$\frac{1}{5}$	$\frac{8}{5} - M$	$-1/5 - M$	0	$18/5$
	x_1	1	0	$\frac{1}{5}$	$\frac{3}{5}$	$-\frac{1}{5}$	0	$\frac{3}{5}$
	x_2	0	1	$-\frac{3}{5}$	$-\frac{4}{5}$	$\frac{3}{5}$	0	$\frac{6}{5}$
	x_4	0	0	1	1	-1	1	1
3	z	0	0	0	$\frac{7}{5} - M$	$-M$	$-\frac{1}{5}$	$17/5$
	x_1	1	0	0	$\frac{2}{5}$	0	$-\frac{1}{5}$	$\frac{2}{5}$
	x_2	0	1	0	$-\frac{1}{5}$	0	$\frac{3}{5}$	$\frac{9}{5}$
	x_3	0	0	1	1	-1	1	1
(optimum)								

M = 1:

Optimum solution: $x_1 = 0, x_2 = 2, x_3 R_4 = 1$
 $Z = 3$

Solution is infeasible because xR_4 is positive. The reason $M=1$ produces an infeasible solution is that it does not play the role of a penalty relative to the objective coefficients of the real variables, x_1 and x_2 . Using $M=1$ makes xR_4 more attractive than x_1 from the standpoint of minimization.

M = 10:

Optimum solution: $x_1 = -4, x_2 = 1.8, Z = 3.4$

The solution is feasible because it does not include artificials at positive level. $M=10$ is relatively much larger than the objective coefficients of x_1 and x_2 , and hence properly plays the role of a penalty.

M = 1000:

It produces the optimum solution as with $M=10$. The conclusion is that it suffices to select M reasonably larger than the objective coefficients of the real variables. Actually, $M=1000$ is an "over kill" in this case, and selecting such huge values could result in adverse round-off error.

(a) Minimize $Z = 4x_1 + x_2 + M(R_1 + R_2 + R_3)$
 subject to

$$\begin{aligned} 3x_1 + x_2 + R_1 &= 3 \\ 4x_1 + 3x_2 - S_2 + R_2 &= 6 \\ x_1 + 2x_2 - S_3 + R_3 &= 4 \\ x_1, x_2, S_2, S_3, R_1, R_2, R_3 &\geq 0 \end{aligned}$$

3

Basic	x_1	x_2	S_2	S_3	R_1	R_2	R_3	
Z	-4	-1			(-M)	(-M)	(-M)	0
R_1	3	1			(1)			3
R_2	4	3	-1			(1)		6
R_3	1	2		-1			(1)	4
Z	$-4+8M$	$-1+6M$	$-M$	$-M$	0	0	0	$10M$
R_1	3	1				1		3
R_2	4	3	-1			1		6
R_3	1	2		-1			1	4

(b) Minimize $Z = 4x_1 + x_2 + M R_1$
 subject to

$$\begin{aligned} 3x_1 + x_2 + R_1 &= 3 \\ 4x_1 + 3x_2 + S_2 &= 6 \\ x_1 + 2x_2 + S_3 &= 4 \end{aligned}$$

Basic	x_1	x_2	R_1	S_2	S_3	
Z	-4	-1	(-M)			0
R_1	3	1	(1)			3
S_2	4	3		1		6
R_3	1	2			1	4
Z	$-4+3M$	$-1+M$		0	0	$3M$
R_1	3	1		1		3
S_2	4	3		1		6
R_3	1	2			1	4

(c) Minimize $Z = 4x_1 + x_2 + M(R_1 + R_2)$

subject to

$$\begin{aligned} 3x_1 + x_2 + R_1 &= 3 \\ 4x_1 + 3x_2 + R_2 &= 6 \\ x_1 + 2x_2 + S_3 &= 4 \end{aligned}$$

Basic	x_1	x_2	R_1	R_2	S_3	
Z	-4	-1	(-M)	(-M)	0	0
R_1	3	1	(1)			3
R_2	4	3		(1)		6
S_3	1	2			1	4
Z	$-4+7M$	$-1+4M$	0	0	0	$9M$
R_1	3	1	1			3
R_2	4	3		1		6
S_3	1	2			1	4

continued...

(d) Maximize $Z = 4x_1 + x_2 - M(R_1 + R_2)$
subject to

$$\begin{aligned} 3x_1 + x_2 + R_1 &= 3 \\ 4x_1 + 3x_2 - S_2 + R_2 &= 6 \\ x_1 + 2x_2 + S_3 &= 4 \end{aligned}$$

Basic	x_1	x_2	S_2	R_1	R_2	S_3	
Z	-4	-1	0	(M)	(M)	0	0
R_1	3	1		1			3
R_2	4	3	-1		1		6
S_3	1	2				1	4
Z	-4-7M	-1-4M	M	0	0	0	-9M
R_1	3	1		1			3
R_2	4	3	-1		1		6
S_3	1	2				1	4

(a) Maximize $Z = 5x_1 + 6x_2 - M(R_1)$
subject to

$$\begin{aligned} -2x_1 + 3x_2 + (R_1) &= 3 \quad (1) \\ x_1 + 2x_2 + S_3 &= 5 \quad (3) \\ 6x_1 + 7x_2 + S_4 &= 3 \quad (4) \end{aligned}$$

$$Z - (5-2M)x_1 - (6+3M)x_2 = -3M$$

(b) Maximize $Z = 2x_1 - 7x_2 - M(R_1 + R_2 + R_5)$
subject to

$$\begin{aligned} -2x_1 + 3x_2 + R_1 &= 3 \quad (1) \\ 4x_1 + 5x_2 - S_2 + R_2 &= 10 \quad (2) \\ 6x_1 + 7x_2 + S_4 &= 3 \quad (4) \\ 4x_1 + 8x_2 - S_5 + R_5 &= 5 \quad (5) \end{aligned}$$

$$Z - (2+6M)x_1 - (-7+16M)x_2 + MS_2 + MS_5 = -18M$$

(c) Minimize $Z = 3x_1 + 6x_2 + MR_S$

subject to

$$\begin{aligned} x_1 + 2x_2 + S_1 &= 5 \quad (3) \\ 6x_1 + 7x_2 + S_2 &= 3 \quad (4) \\ 4x_1 + 8x_2 - S_5 + R_S &= 5 \quad (5) \end{aligned}$$

$$Z - (3-4M)x_1 - (6-8M)x_2 - MS_5 = 5M$$

(d) Minimize $Z = 4x_1 + 6x_2 + M(R_1 + R_2 + R_5)$

subject to

$$\begin{aligned} -2x_1 + 3x_2 + R_1 &= 3 \quad (1) \\ 4x_1 + 5x_2 - S_2 + R_2 &= 10 \quad (2) \\ 4x_1 + 8x_2 - S_5 + R_5 &= 5 \quad (5) \end{aligned}$$

$$Z - (4-6M) - (6-16M)x_2 - MS_2 - MS_5 = 18M$$

(e) Minimize $Z = 3x_1 + 2x_2 + M(R_1 + R_5)$

subject to

$$\begin{aligned} -2x_1 + 3x_2 + R_1 &= 3 \quad (1) \\ 4x_1 + 8x_2 - S_5 + R_5 &= 5 \quad (5) \end{aligned}$$

$$Z - (3-2M)x_1 - (2-11M)x_2 - MS_5 = 8M$$

(a)

Basic	x_1	x_2	x_3	S_2	R_1	R_2	
$\bar{3}$	-2	-3	5	-M	0	0	-17M
0	R_1	1	1	1	0	1	0
R_2	2	-5	1	-1	0	1	10
$\bar{3}$	0	-8	6	-1	1	1	16
I	R_1	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	1	$-\frac{1}{2}$	2
x_1	1	$-\frac{5}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$	0	$\frac{1}{2}$	5
$\bar{3}$	0	0	$50/7$	$1/7$	$16/7$	$-1/7$	$\frac{102}{7}$
II	x_2	0	1	$1/7$	$1/7$	$3/7$	$-1/7$
x_1	1	0	$6/7$	$-1/7$	$5/7$	$1/7$	$\frac{45}{7}$

(b)

Basic	x_1	x_2	x_3	S_2	R_1	R_2	$S_0/2$
$\bar{3}$	-2	-3	5	-M	0	0	-17M
0	R_1	1	1	1	0	1	7
R_2	2	-5	1	-1	0	1	10
$\bar{3}$	0	-8	6	-1	1	1	15
I	R_1	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	1	$-\frac{1}{2}$	2
x_1	1	$-\frac{5}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$	0	$\frac{1}{2}$	5
$\bar{3}$	0	0	$50/7$	$1/7$	$16/7$	$-1/7$	$\frac{102}{7}$
II	x_2	0	1	$1/7$	$1/7$	$3/7$	$-1/7$
x_1	1	0	$6/7$	$-1/7$	$5/7$	$1/7$	$\frac{45}{7}$
$\bar{3}$	0	-50	0	-7	-12	7	-14
III	x_3	0	7	1	1	2	-1
x_1	1	-6	0	-1	-1	1	3

continued...

continued...

Set 3.4a

(c)

	x_1	x_2	x_3	s_1	R_1	R_2	Sols
0	3	-1	-2	-1	0	m	m
0	R_1	1	1	1	0	1	0
0	R_2	2	-5	1	-1	0	1
I	3	-1	-2	-1	m	0	0
I	3	-3m	+4m	-2m	m	0	-17m
I	R_1	1	1	1	0	1	0
I	R_2	2	-5	1	-1	0	1
II	3	0	$-\frac{9}{2}$	$-\frac{1}{2}$	$-\frac{1}{2}$	$\frac{1}{2}$	$\frac{5}{2}$
II	R_1	0	$\frac{7}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	1	$-\frac{1}{2}$
II	x_1	1	$-\frac{5}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$	0	$\frac{1}{2}$
III	3	0	0	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}m$	$-\frac{1}{7}m$
III	x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{2}{7}$	$-\frac{1}{7}$
III	x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{5}{7}$	$\frac{45}{7}$

(d)

	x_1	x_2	x_3	s_1	R_1	R_2	Sols
0	3	-4	8	-3	0	-m	-m
0	R_1	1	1	1	0	1	0
0	R_2	2	-5	1	-1	0	1
I	3	-4	8	-3	m	0	6
I	3	+3m	-4m	+2m	-m	0	17m
I	R_1	1	1	1	0	1	0
I	R_2	2	-5	1	-1	0	1
II	3	0	-2	-1	-2	2	$\frac{20}{7}$
II	3	0	$\frac{7}{2}m/2$	$\frac{1}{2}m/2$	$\frac{1}{2}m/2$	0	$-\frac{3}{2}m/2$
II	R_1	0	$\frac{7}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	1	$-\frac{1}{2}$
II	x_1	1	$-\frac{5}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$	0	$\frac{1}{2}$
III	3	0	0	$-\frac{5}{7}$	$-\frac{12}{7}$	$\frac{4}{7}m$	$-\frac{1}{7}m$
III	x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$	$-\frac{1}{7}$
III	x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$-\frac{5}{7}$	$\frac{45}{7}$

In the first iteration, we must substitute out the starting solution variables, x_3 and x_4 , in the Z-equation, exactly as we do with the artificial variables.

6

Basic	x_1	x_2	x_3	x_4	Solution	
0	Z	-2	-4	<u>(-4)</u>	<u>(3)</u>	-
0	x_3	1	1	<u>1</u>	0	4
0	x_4	1	4	0	<u>1</u>	8
I	Z	-1	-12	0	0	-8
I	x_3	1	1	1	0	4
I	x_4	1	<u>4</u>	0	1	8
II	Z	2	0	0	3	16
II	x_3	$\frac{3}{4}$	0	1	$-\frac{1}{4}$	2
II	x_2	$\frac{1}{4}$	1	0	$\frac{1}{4}$	2

After adding surplus S_1 and S_2 , substitute out x_3 in the Z-equation.

7

Basic	x_1	x_2	s_1	s_2	x_3	x_4	Solution
0	Z	-3	-2	0	0	<u>(-3)</u>	0
0	x_3	1	4	-1	0	<u>1</u>	0
0	x_4	2	1	0	-1	0	1
I	Z	0	10	-3	0	0	0
I	x_3	1	4	-1	0	1	0
I	x_4	2	1	0	-1	0	1
II	Z	$-\frac{5}{2}$	0	$-\frac{1}{2}$	0	$-\frac{5}{2}$	0
II	x_2	$\frac{1}{4}$	1	$-\frac{1}{4}$	0	$\frac{1}{4}$	0
II	x_4	$\frac{7}{4}$	0	$\frac{1}{4}$	-1	$-\frac{1}{4}$	1

Both x_3 and R (the starting solution variables) must be substituted out in the Z-equation.

8

Basic	x_1	x_2	x_3	R	Solution	
0	3	-1	-5	<u>(-3)</u>	<u>(M)</u>	-
0	x_3	1	2	<u>1</u>	0	3
0	R	2	-1	0	<u>1</u>	4
I	3	$2-2m$	$1+m$	0	0	$9-4m$
I	x_3	1	2	1	0	3
I	R	<u>2</u>	-1	0	1	4
II	3	0	2	0	$-1+m$	5
II	x_3	0	$\frac{5}{2}$	1	$-\frac{1}{2}$	1
II	x_1	1	$-\frac{1}{2}$	0	$\frac{1}{2}$	2

Maximize $Z = 2x_1 + 5x_2 - MR_1$
subject to

$$\begin{aligned} 3x_1 + 2x_2 - S_1 + R_1 &= 6 \\ 2x_1 + x_2 &+ S_2 = 2 \\ x_1, x_2, S_1, R_1, S_2 &\geq 0 \end{aligned}$$

9

Basic	x_1	x_2	S_1	R_1	S_2	
Z	-2	-5	0	M	0	-
R_1	3	2	-1	1	0	6
S_2	2	1	0	0	1	2
Z	$-2-3M$	$-5-2M$	M	0	0	$-6M$
R_1	3	2	-1	1	0	6
S_2	2	1	0	0	1	2
Z	0	$-4-M/2$	M	0	$1+3M/2$	$+3M$
R_1	0	$1/2$	-1	1	$-3/2$	3
x_1	1	$1/2$	0	0	$1/2$	1
Z	$8+M$	0	M	0	$5+2M$	$-2M$
R_1	-1	0	-1	1	-2	2
x_2	2	1	0	0	1	2

The Z-row shows that the solution is optimal (all nonbasic coefficients in the Z-row are ≥ 0). However, the solution is infeasible because the artificial variable R_1 assumes a positive value. Having a positive value for the artificial variable R_1 is the same as regarding the constraint $3x_1 + 2x_2 \geq 6$ as $3x_1 + 2x_2 \leq 6$, which violates the constraints of the original model.

Set 3.4b

In Phase I, we always minimize the sum of the artificial variables because the sum represents a measure of infeasibility in the problem

- (a) Minimize $r = R_1$
 (b) Minimize $r = R_1 + R_2 + R_5$
 (c) Minimize $r = R_5$
 (d) Minimize $r = R_1 + R_2 + R_5$
 (e) Minimize $r = R_1 + R_5$

(a) Phase I:

Basic	x_1	x_2	x_3	s_2	R_1	R_2	Sol ^{1a}
R_1	0	0	0	0	-1	-1	0
R_2	1	1	1	0	1	0	7
R_2	2	-5	1	-1	0	1	10
R_1	3	-4	2	-1	0	0	17
R_1	1	1	1	0	1	0	7
R_2	2	-5	1	-1	0	1	10
R_1	0	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	0	$-\frac{3}{2}$	2
R_1	0	$\frac{7}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	1	$-\frac{1}{2}$	2
x_1	1	$-\frac{5}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$	0	$\frac{1}{2}$	5
R_1	0	0	0	0	-1	-1	0
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{2}{7}$	$-\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{5}{7}$	$\frac{1}{7}$	$\frac{45}{7}$

2

(c) Phase I is the same as in (a)

Phase II:

Basic	x_1	x_2	x_3	s_2	Sol ^{1a}
Z	-1	-2	-1	0	0
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$
Z	0	0	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{53}{7}$
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$

(d) Phase I is the same as in (a)

Phase II:

Basic	x_1	x_2	x_3	x_4	Sol ^{1a}
Z	-4	8	-3	0	0
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$
Z	0	0	$-\frac{5}{7}$	$-\frac{12}{7}$	$\frac{21}{7}$
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$

Minimize $r = R_1$

Subject to

$$\begin{aligned} 3x_1 + 2x_2 - s_1 + R_1 &= 6 \\ 2x_1 + x_2 + s_2 &= 2 \\ x_1, x_2, s_1, R_1, s_2 &\geq 0 \end{aligned}$$

Solution of Phase I by TORA yields $r = 2$, which indicates that the problem has no feasible space

Basic	x_1	x_2	x_3	s	Sol ^{1a}
Z	-2	-3	5	0	0
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$
Z	0	0	$\frac{50}{7}$	$\frac{1}{7}$	$\frac{102}{7}$
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$

(b) Phase I is the same as in (a)

Basic	x_1	x_2	x_3	s_2	Sol ^{1a}
Z	-2	-3	5	0	0
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$
Z	0	0	$\frac{50}{7}$	$\frac{1}{7}$	$\frac{102}{7}$
x_2	0	1	$\frac{1}{7}$	$\frac{1}{7}$	$\frac{4}{7}$
x_1	1	0	$\frac{6}{7}$	$-\frac{1}{7}$	$\frac{45}{7}$
Z	0	-50	0	-7	-14
x_3	0	7	1	1	4
x_4	1	-6	0	-1	3

continued...

Minimize $Z = R_2$

Subject to

$$\begin{aligned} 2x_1 + x_2 + x_3 + s_1 &= 2 \\ 3x_1 + 4x_2 + 2x_3 - s_2 + R_2 &= 8 \\ x_1, x_2, x_3, s_1, s_2, R_2 &\geq 0 \end{aligned}$$

Phase I Optimal solution:

Basic	x_1	x_2	x_3	s_1	s_2	R_2	Sol ^{1a}
r	-5	0	-2	-1	-4	0	0
x_2	2	1	1	0	1	0	2
R_2	-5	0	-2	-1	-4	1	0

$R_2 = 0$ is basic in the Phase I Solution
continued...

3-14

Set 3.4b

6

(b)

Phase I (continued): R2 leaves, x₁ enters (also x₃, s₂, and s₁ are candidates for the entering variable).

	x ₁	x ₂	x ₃	s ₂	s ₁	R ₂	Sol.
r	-5	0	-2	-1	-4	0	0
x ₂	2	1	1	0	1	0	2
R ₂	-5	0	-2	-1	-4	1	0
r	0	0	0	0	0	-1	
x ₂	0	1	1/5	-2/5	-3/5	2/5	2
x ₁	1	0	2/5	1/5	4/5	-1/5	0

Drop R₂-column.

Phase II:

	x ₁	x ₂	x ₃	s ₂	s ₁	Sol.
z	-2	-2	-4	0	0	0
x ₂	0	1	1/5	-2/5	-3/5	2
x ₁	1	0	2/5	1/5	4/5	0
z	0	0	-14/5	-2/5	2/5	4
x ₂	0	1	1/5	-2/5	-3/5	2
x ₁	1	0	2/5	1/5	4/5	0
z	7	0	0	1	6	4
x ₂	-1/2	1	0	-1/2	-1	2
x ₃	5/2	0	1	1/2	2	0

Optimum solution:

$$x_1 = 0, x_2 = 2, x_3 = 0, z = 4$$

Phase I:

	x ₁	x ₂	x ₃	R ₁	R ₂	R ₃	Sol.
r	-10	0	-4	-8	0	0	0
x ₂	2	1	1	1	0	0	2
R ₂	-5	0	-2	-3	1	0	0
R ₃	-5	0	-2	-4	0	1	0
r	0	0	1	-2	-2	0	0
x ₂	0	1	1/5	-1/5	2/5	0	2
x ₁	1	0	2/5	3/5	-1/5	0	0
R ₃	0	0	0	-1	-1	1	0

Remove R₁- and R₂ columns, which gives

	x ₁	x ₂	x ₃	R ₃	Sol.
r	0	0	1	0	0
x ₂	0	1	1/5	0	2
x ₁	1	0	2/5	0	0
R ₃	0	0	0	1	0

The R₃-row is R₃ = 0, which is redundant. Hence the R₃-row and R₃-column can be dropped from the tableau with no consequences.

Phase II:

	x ₁	x ₂	x ₃	Sol.
z	-3	-2	-3	0
x ₂	0	1	1/5	2
x ₁	1	0	2/5	0
z	0	0	-7/5	4
x ₂	0	1	1/5	2
x ₁	1	0	2/5	0
z	7/2	0	0	4
x ₂	-1/2	1	0	2
x ₁	5/2	0	1	0

Optimum solution:

$$x_1 = 0, x_2 = 2, x_3 = 0, z = 4$$

Set 3.4b

If x_1, x_3, x_4 , or x_5 assume a positive value, the value of the objective function at the end of Phase I must necessarily become positive. This follows because these variables have nonzero Z-row coefficients in the optimal Phase I tableau. A positive objective value at the end of Phase I means that Phase I solution is infeasible. Since Phase II uses the same constraints as in Phase I, it follows that Phase II must have $x_1 = x_3 = x_4 = x_5 = 0$ as well.

7

Phase II:

Basic	x_2	R	S _{0/Z}
Z	(-2)	0	0
x_2	1	0	2
R	0	1	0
Z	0	0	4
x_2	1	0	2
R	0	1	0

Optimum solution:

$$x_1 = 0 \quad x_2 = 2 \quad x_3 = x_4 = x_5 = 0 \\ Z = 4$$

$$\begin{array}{lcl} -5x_1 + 6x_2 - 2x_3 + x_4 & = -5 \\ x_1 - 3x_2 - 5x_3 + x_5 & = -8 \\ 2x_1 + 5x_2 - 4x_3 + x_6 & = 9 \end{array}$$

8

x_1	x_2	x_3	x_4	x_5	x_6	R
0	0	0	0	0	0	-1
-5	6	-2	1	0	0	-1
1	-3	-5	0	1	0	-1
2	5	-4	0	0	1	0
-1	3	5	0	-1	0	0
-6	9	3	1	-1	0	0
-1	3	5	0	-1	0	1
2	5	-4	0	0	1	0

Phase I problem:

$$\text{minimize } r = R$$

subject to

$$\begin{array}{lcl} -6x_1 + 9x_2 + 3x_3 + x_4 - x_5 & = 3 \\ -x_1 + 3x_2 + 5x_3 - x_5 + R & = 8 \\ 2x_1 + 5x_2 - 4x_3 + x_6 & = 9 \end{array}$$

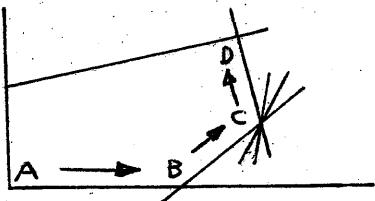
all variables ≥ 0

The logic of the procedure is as follows:

In the R-column, enter -1 for any constraint with negative RHS and 0 for all other constraints.

Next, use the R-column as a pivot column and select the pivot element as the one corresponding to the most negative RHS. This procedure will always require one artificial variable regardless of the number of constraints.

(a)

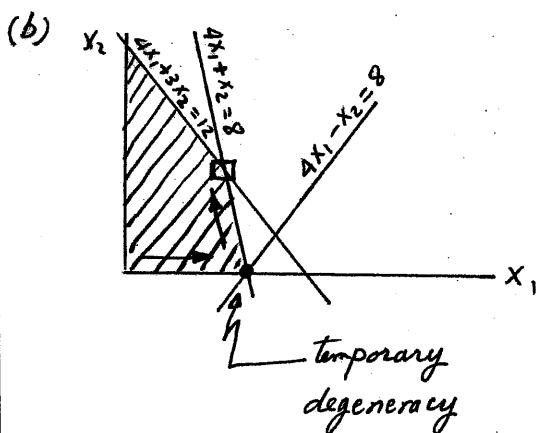


1

(b) $A: 1, B: 1, C: \binom{3}{2} = 3, D: 1$

(a) From TORA, iterations 2 and 3 are degenerate. Degeneracy is removed in iteration 4.

2



(a) Four iterations

3

(b) Three iterations: In iteration 2, degeneracy is removed because basic $SX_5 = 0$ corresponds to a negative constraint coefficient in the entering variable column (x_2).

(c) In part (a), solution encounters 2 degenerate basic solution at the same corner point. In part (b), only one basic solution was encountered.

Set 3.5b

Basic	x_1	x_2	x_3	s_1	s_2	s_3	Solution
\bar{Z}	-1	-2	-3	0	0	0	0
s_1	1	2	$\boxed{3}$	1	0	0	10
s_2	1	1	0	0	1	0	5
s_3	1	0	0	0	0	1	1
\bar{Z}	0	$\boxed{0}$	0	1	0	0	10
x_3	$\frac{1}{3}$	$\frac{2}{3}$	1	$\frac{1}{3}$	0	0	$\frac{10}{3}$
s_2	1	$\boxed{1}$	0	0	1	0	5
s_3	1	0	0	0	0	1	1
\bar{Z}	$\boxed{0}$	0	0	1	0	0	10
x_3	$-\frac{1}{3}$	0	1	$\frac{1}{3}$	$-\frac{2}{3}$	0	0
x_2	1	1	0	0	1	0	5
s_3	$\boxed{1}$	0	0	0	0	1	1
\bar{Z}	0	0	0	1	$\boxed{0}$	0	10
x_3	0	0	1	$\frac{1}{3}$	$-\frac{2}{3}$	$\frac{1}{3}$	$\frac{1}{3}$
x_2	0	1	0	0	$\boxed{1}$	-1	4
x_1	1	0	0	0	0	1	1
\bar{Z}	0	0	0	1	0	0	10
x_3	0	$\frac{2}{3}$	1	$\frac{1}{3}$	0	$-\frac{1}{3}$	3
s_2	0	1	0	0	1	-1	4
x_1	1	0	0	0	0	1	1

Three alternative basic optima:

$$(x_1, x_2, x_3) = \begin{cases} (0, 0, 10/3) \\ (0, 5, 0) \\ (1, 4, 1/3) \end{cases}$$

The associated nonbasic alternative optima are

$$\hat{x}_1 = \lambda_3$$

$$\hat{x}_2 = 5\lambda_2 + 4\lambda_3$$

$$\hat{x}_3 = 10/3\lambda_1 + 1/3\lambda_3$$

where

$$\lambda_1 + \lambda_2 + \lambda_3 = 1$$

$$0 \leq \lambda_i \leq 1, i=1,2,3$$

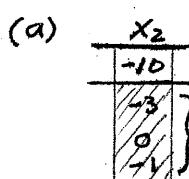
Basic	x_1	x_2	x_3	s_1	s_2	s_3	2
\bar{Z}	-2	1	3	0	0	0	0
s_1	1	-1	$\boxed{5}$	1	0	10	
s_2	2	-1	3	0	1	40	
\bar{Z}	$-\frac{7}{5}$	$\frac{2}{5}$	0	$\frac{3}{5}$	0	6	
x_3	$\boxed{1/5}$	$-\frac{1}{5}$	1	$\frac{1}{5}$	0	2	
s_2	$\frac{7}{5}$	$-\frac{2}{5}$	0	$-\frac{3}{5}$	1	34	
\bar{Z}	0	-1	7	2	0	20	
x_1	1	-1	5	1	0	10	
s_2	0	$\boxed{1}$	-7	-2	1	20	
\bar{Z}	0	0	$\boxed{0}$	$\boxed{0}$	1	40	
x_1	1	0	-2	-1	0	30	
x_2	0	1	-7	-2	1	20	

x_3 and s_1 can yield alternative optima. However, because all their constraint coefficients are negative (in general, ≤ 0), none can yield an alternative (corner point) basic solution.

Basic	x_1	x_2	x_3	s_1	s_2	s_3	3
\bar{Z}	-3	-1	0	0	0	0	0
s_1	1	2	0	1	0	0	5
s_2	$\boxed{1}$	1	-1	0	1	0	2
s_3	7	3	-5	0	0	1	20
\bar{Z}	0	2	-3	0	3	0	6
s_1	0	1	$\boxed{1}$	1	-1	0	3
x_1	1	1	-1	0	1	0	2
s_3	0	-4	2	0	-7	1	6
\bar{Z}	0	5	0	3	$\boxed{0}$	0	15
x_3	0	1	1	1	-1	0	3
x_1	1	2	0	1	0	0	5
s_3	0	-6	0	-2	-5	1	0

The optimum solution is degenerate because s_3 is basic and equal to zero. Also, it has alternative nonbasic solutions because s_2 has a zero coefficient in the Z -row and all its constraint coefficients are ≤ 0 .

Basic	x_1	x_2	s_1	s_2	
Z	-2	-1	0	0	0
s_1	1	-1	1	0	10
s_2	2	0	0	1	40
Z	0	-3	2	0	20
x_1	1	-1	1	0	10
s_2	0	2	-2	1	20
Z	0	0	-1	$\frac{3}{2}$	50
x_1	1	0	0	$\frac{1}{2}$	20
x_2	0	1	-1	$\frac{1}{2}$	10

unbounded \rightarrow 

2

\Rightarrow Solution space unbounded
in the direction of x_2

(b) Objective value is unbounded
because each unit increase
in x_2 increases Z by 10

If, at any iteration, all the
constraint coefficients of a
variable are ≤ 0 , then the
solution space is unbounded in
the direction of that variable.

3

A more "foolproof" way of
accomplishing this task is to solve
a sequence of LPs in which the
objective function is

Maximize $Z = x_j$, $j=1, 2, \dots, n$
Subject to the constraints of the
problem. For the unbounded
variables, $Z = \infty$.

Set 3.5d

x_1 = number of units of T1
 x_2 = number of units of T2
 x_3 = number of units of T3

Constraints:

$$3x_1 + 5x_2 + 6x_3 \leq 1000$$

$$5x_1 + 3x_2 + 4x_3 \leq 1200$$

$$x_1 + x_2 + x_3 \geq 500$$

$$x_1, x_2, x_3 \geq 0$$

We can use Phase I to see whether the problem has a feasible solution; that is,

minimize $r = R_3$

subject to

$$3x_1 + 5x_2 + 6x_3 + S_1 = 1000$$

$$5x_1 + 3x_2 + 4x_3 + S_2 = 1200$$

$$x_1 + x_2 + x_3 - S_3 + R_3 = 500$$

$$x_1, x_2, x_3, S_1, S_2, S_3, R_3 \geq 0$$

Optimum solution from TORA:

$$R_3 = r = 225 \text{ units}$$

This is interpreted as a deficiency of 225 units. The most that can be produced is $500 - 225 = 275$ units

Basic	x_1	x_2	x_3	S_1	S_2	R_3	Sol'n
Z	-3	-2	-3				
	-3M	-4M	-2M	M	0	0	-8M
S_1	2	1	0	1	0	0	2
R_1	3	4	2	-1	0	1	8
Z	-1	-1	2				
	+5M	0	+2M	M	+4M	0	4
x_2	2	1	0	1	0	0	2
R_1	-5	0	-2	-1	-4	1	0

Because $R_1 = 0$ in the optimal tableau, the problem has a feasible solution. The optimum solution is

$$x_1 = 0, x_2 = 2, Z = 4$$

Note that in the first iteration, R_1 could have been used as the leaving variable, in which case it would not be basic in the optimum iteration.

Set 3.6a

X_1 = Nbr. units of product A
 X_2 = Nbr. units of product B

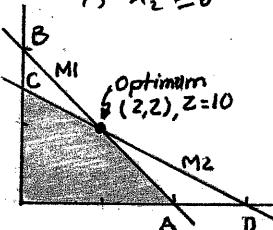
$$\text{Maximize } Z = 2X_1 + 3X_2$$

S.t.

$$2X_1 + 2X_2 \leq 8 \quad (\text{M1})$$

$$3X_1 + 6X_2 \leq 18 \quad (\text{M2})$$

$$X_1, X_2 \geq 0$$



	M1	M2	Z
A = (4, 0)	12	8	
B = (0, 4)	24	12	
C = (0, 3)	6	9	
D = (2, 2)	12	12	

$$(a) M1 \text{ at } C = 2(0) + 2(3) = 6$$

$$M1 \text{ at } D = 2(6) + 2(0) = 12$$

$$Z \text{ at } C = 2(0) + 3(3) = 9$$

$$Z \text{ at } D = 2(6) + 3(0) = 12$$

$$\text{Dual price} = \frac{12-9}{12-6} = .50/\text{unit}$$

$$\text{Allowable range} = (6 \leq M1 \leq 12)$$

$$M2 \text{ at } A = 3(4) + 6(0) = 12$$

$$M2 \text{ at } B = 3(0) + 6(4) = 24$$

$$Z \text{ at } A = 2(4) + 3(0) = 8$$

$$Z \text{ at } B = 2(0) + 3(4) = 12$$

$$\text{Dual price} = \frac{12-8}{24-12} = .33/\text{unit}$$

$$\text{Range: } 12 \leq M2 \leq 24$$

$$(b) \text{ Dual price} = .50/\text{unit valid in the range } 6 \leq M1 \leq 12$$

$$\text{Increase in revenue} = .5 \times 4 = \$2.00$$

$$\text{Increase in cost} = .3 \times 4 = \$1.20$$

Cost < Revenue - purchase recommended

$$(c) \text{ Dual price} = .33/\text{unit valid in the range } 12 \leq M2 \leq 24$$

$$\text{Purchase price/unit} < .33$$

$$(d) \text{ Dual price} = .33/\text{unit valid in the range } 12 \leq M2 \leq 24. M2 \text{ is increased from 18 to 23 units}$$

$$\text{Increase in revenue} = 5 \times .33 = \$1.65$$

$$\text{New optimum revenue} = 10 + 1.65 = \$11.65$$

2

X_1 = daily number of type 1 rats
 X_2 = daily number of type 2 rats

$$\text{Maximize } Z = 8X_1 + 5X_2$$

$$2X_1 + X_2 \leq 400$$

$$X_1 \leq 150$$

$$X_2 \leq 200$$

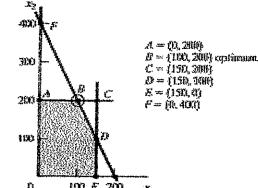
$$X_1, X_2 \geq 0$$

(a) Optimum occurs at B:

$$X_1 = 100 \text{ type 1 rats}$$

$$X_2 = 200 \text{ type 2 rats}$$

$$Z = \$1800$$



$$(b) A = (0, 200), C = (150, 200)$$

capacity

$$A \quad 2 \times 0 + 1 \times 200 = 200 \quad 8 \times 0 + 5 \times 200 = 1000$$

$$C \quad 2 \times 150 + 1 \times 200 = 500 \quad 8 \times 150 + 5 \times 200 = 2200$$

$$\text{worth/capacity unit} = \frac{2200 - 1000}{500 - 200}$$

$$= \$4 \text{ per type 2 rat}$$

$$\text{Range: } (200, 500)$$

$$(c) \text{ Dual price} = 0 \text{ in the range } (100, \infty)$$

Thus, change from $X_1 \leq 150$ to $X_1 \leq 120$

has no effect on optimum Z.

$$(d) \text{ Let } d = \text{demand limit for type 2 rat}$$

$$d \quad Z$$

$$D(150, 100) \quad 100 \quad 8(150) + 5(100) = \$1700$$

$$F(0, 400) \quad 400 \quad 8(0) + 5(400) = \$2000$$

$$\text{Dual price} = \frac{2000 - 1700}{400 - 100} = \$1.00$$

$$\text{Range } (100, 400)$$

$$\text{Maximum increase in demand limit for type 2 rat} = 400 - 200 = 200 \text{ rats}$$

Set 3.6b

(a) $\frac{3}{6} \leq \frac{C_A}{C_B} \leq \frac{2}{2}$, or
 $0.5 \leq \frac{C_A}{C_B} \leq 1$ or $1 \leq \frac{C_B}{C_A} \leq 2$

(b) Maximize $Z = 2x_A + 3x_B$

$C_B = 3$: $3x_A \leq C_A \leq 3x_1$
 $1.5 \leq C_A \leq 3$

$C_A = 2$: $2x_A \leq C_B \leq 2x_2$
 $1 \leq C_B \leq 4$

(c) $\frac{C_A}{C_B} = \frac{5}{4} = 1.25$, which falls outside the range $0.5 \leq \frac{C_A}{C_B} \leq 1$. Optimum solution changes and must be computed anew.
New solution: $x_A = 4$, $x_B = 0$, $Z = \$20$.

(d) Case 1: $Z = 5x_A + 3x_B$

$C_A = 5$ falls outside the range $(1.5, 3)$, hence the optimum changes. New Optimum is $x_A = 4$, $x_B = 0$, $Z = \$20$.

Case 2: $Z = 2x_A + 4x_B$

$C_B = 4$ falls in the range $(1, 4)$, hence optimum is unchanged at $x_A = x_B = 2$,
 $Z = 2(2) + 4(2) = \$12$

(a) $\frac{1}{2} \leq \frac{C_A}{C_B} \leq \frac{6}{4}$, or

$0.5 \leq \frac{C_A}{C_B} \leq 1.5$ or $\frac{2}{3} \leq \frac{C_B}{C_A} \leq 2$

(b) Given $C_A = 5$, then

$5\left(\frac{2}{3}\right) \leq C_B \leq 5(2)$, or $\frac{10}{3} \leq C_B \leq 10$

(c) $\frac{C_A}{C_B} = \frac{5}{3} = 1.67$, which falls

outside the range $0.5 \leq \frac{C_A}{C_B} \leq 1.5$.

Hence the solution changes

1

(a) $0 \leq \frac{C_A}{C_B} \leq \frac{2}{1}$, or

$0 \leq \frac{C_B}{C_A} \leq 2$

(b) $\frac{C_A}{C_B} = 1$, which falls in the range $0 \leq \frac{C_B}{C_A} \leq 2$. Hence, the solution is unchanged.

3

2

Set 3.6c

Feasibility conditions:

$$X_2 = 100 + \frac{1}{2}D_1 - \frac{1}{4}D_2$$

$$X_3 = 230 + \frac{1}{2}D_2$$

$$X_6 = 20 - 2D_1 + D_2 + D_3$$

$$(a) D_1 = 438 - 430 = 8 \text{ min}$$

$$D_2 = 500 - 460 = 40$$

$$D_3 = 410 - 420 = -10$$

$$X_2 = 100 + \frac{1}{2}(8) - \frac{1}{4}(40) = 94 > 0$$

$$X_3 = 230 + \frac{1}{2}(40) = 250 > 0$$

$$X_6 = 20 - 2(8) + 40 - 10 = 34 > 0$$

Dual prices:

$$\text{Resource 1} = \$1/\text{min}, -200 \leq D_1 \leq 10$$

$$2 = \$2/\text{min}, -20 \leq D_2 \leq 400$$

$$3 = \$0/\text{min}, -20 \leq D_3 < \infty$$

$$\text{New profit} = 1350 + D_1 + 2D_2 + 0D_3 \\ = 1350 + 8 + 2 \times 40 = \$1438$$

$$(b) D_1 = 460 - 430 = 30 \text{ min}$$

$$D_2 = 440 - 460 = -20$$

$$D_3 = 380 - 420 = -40$$

$$X_2 = 100 + \frac{1}{2}(30) - \frac{1}{4}(-20) = 120 > 0$$

$$X_3 = 230 + \frac{1}{2}(-20) = 220 > 0$$

$$X_6 = 20 - 2(30) - 20 - 40 = -100 < 0$$

$$(a) Overtime cost \frac{50}{60} = \$0.83/\text{min}$$

Revenue (dual price) for operation 1 is \\$1/min.

Cost < Revenue \Rightarrow advantageous

$$(b) \text{Dual price for operation 2} = \$2/\text{min} \text{ valid in the range } -20 \leq D_2 \leq 400$$

$$D_2 = 120 \text{ minutes}$$

$$\text{Revenue increase} = 120 \times 2 = \$240$$

$$\text{Cost increase} = 2 (\$55) = \$110$$

Revenue > Cost \Rightarrow accept.

(c) No, resource 3 is already abundant. This is the reason its dual price = 0

(d) Dual price for operation 1 is \\$1/min, valid in the range $-200 \leq D_1 \leq 10$

continued...

3-23

$$D_1 = 440 - 430 = 10 \text{ min}$$

$$\text{Cost} = \frac{10}{60} \times 40 = \$6.67$$

$$\text{New revenue} = 1350 + 1 \times 10 = \$1360$$

$$\text{Net revenue} = 1360 - \$6.67 = \$1353.33$$

$$(e) \text{Dual price} = \$2/\text{min}, -20 \leq D_2 \leq 400$$

$$D_2 = - \text{ min}$$

$$\text{Decrease in cost} = \frac{15}{60} \times 30 = \$7.50$$

$$\text{Lost revenue} = 15 \times \$2.00 = \$30.00$$

Lost revenue > Decrease in cost

Not recommended.

X_j = units of product $i = 1, 2, 3$

$$\text{Maximize } Z = 20X_1 + 50X_2 + 35X_3$$

s.t.

$$-.5X_1 + .5X_2 + .5X_3 \leq 0$$

$$X_1 \leq 75$$

$$2X_1 + 4X_2 + 3X_3 \leq 240$$

$$X_1, X_2, X_3 \geq 0$$

$$(a) \text{Solution: } Z = \$2800$$

$$X_1 = X_2 = 40, X_3 = 0$$

	X_1	X_2	X_3	S_1	S_2	S_3	
Z	0	0	$10/3$	$20/3$	0	$35/3$	2800
X_2	0	0	$5/6$	$2/3$	0	$1/6$	40
S_2	1	0	$1/6$	$4/3$	1	$-1/6$	35
X_1	0	1	$-1/6$	$-4/3$	0	$1/6$	40

$$(b) Z + 10/3X_3 + 20/3S_1 + 0S_2 + 35/3S_3 = 2800$$

$$\text{Dual price for raw material} = \$35/3 / 16$$

$$\left. \begin{aligned} X_2 &= 40 + D_3/6 \\ S_2 &= 35 - D_3/6 \end{aligned} \right\} \Rightarrow -240 \leq D_3 \leq 210$$

$D_3 = 120$ falls in the range (-240, 210)

New solution:

$$X_1 = 40 + \frac{120}{6} = 60 \text{ units}$$

$$X_2 = 40 + \frac{120}{6} = 60 \text{ units}$$

$$X_3 = 0$$

$$\text{New revenue} = 2800 + (35/3)(120) \\ = \$4200$$

continued...

Set 3.6c

(e) Dual price = 0, $-35 \leq D_2 < \infty$
 $\pm 10\% \text{ of } 75 = \pm 7.5$ or
 Change has no effect on the solution

X_j = units of product j , $j = 1, 2, 3$,

Maximize $Z = 4.5X_1 + 5X_2 + 4X_3$
 s.t.

$$\begin{aligned} 10X_1 + 5X_2 + 6X_3 &\leq 600 \\ 6X_1 + 8X_2 + 9X_3 &\leq 600 \\ 8X_1 + 10X_2 + 12X_3 &\leq 600 \\ X_1, X_2, X_3 &\geq 0 \end{aligned}$$

(a) Solution: $Z = \$325$

$$X_1 = 50, X_2 = 20, X_3 = 0$$

(b) Optimum tableau

X_1	X_2	X_3	S_1	S_2	S_3	
Z	0	0	2	.083	0	.458 325
X_1	1	0	0	.167	0	-.083 50
S_2	0	0	-.6	.067	1	-.833 140
X_2	0	1	1.2	-.133	0	.167 20

$$Z + 2X_3 + .083S_1 + 0S_2 + .458S_3 = 325$$

Dual prices:

Process 1: \$.083/min

2: \$0/min

3: \$.458/min

Process 3 > Process 1

(c) Process 1: $60X_1 \cdot .083 = \$4.98$

2: 0

3: $60X_1 \cdot .458 = \$27.48$

X_1 = Nbr. of practical courses

X_2 = Nbr. of humanistic courses

Maximize $Z = 1500X_1 + 1000X_2$

$$X_1 + X_2 + S_1 = 30 \quad (1)$$

$$X_1 - S_2 = 10 \quad (2)$$

$$X_2 - S_3 = 10 \quad (3)$$

$$X_1, X_2, S_1, S_2, S_3 \geq 0$$

(a) Solution:

$$Z = \$40,000$$

$$X_1 = 20 \text{ courses}$$

$$X_2 = 10 \text{ courses}$$

5

continued...

(b) From TORA,
 $Z + 1500S_1 + 0S_2 + 500S_3 = 40,000$
 S_1 is a slack, S_2 and S_3 are surplus

Dual prices:

constraint 1: \$1500/course
 constraint 2: \$0/min limit course
 constraint 3: -\$500/min limit course

Dual price for constraint 1 equals the revenue per practical course. Hence, an additional course must necessarily be of the practical type.

(c) From TORA,

$$\begin{cases} S_2 = 10 + D_1 \geq 0 \\ X_1 = 20 + D_1 \geq 0 \\ X_2 = 10 \end{cases} \quad \begin{cases} -10 \leq D_1 < \infty \\ -10 \leq D_1 \leq 10 \\ X_2 = 10 \end{cases}$$

Thus, the dual price of \$1500 for constraint 1 is valid for any number of courses $\geq 30 - 10 = 20$.

(d) Dual price = -\$500. To determine the range where it applies, we have from TORA

$$\begin{cases} S_3 = 10 - D_3 \geq 0 \\ X_1 = 20 - D_3 \geq 0 \\ X_2 = 10 + D_3 \geq 0 \end{cases} \quad \begin{cases} -10 \leq D_3 \leq 10 \\ -10 \leq D_3 \leq 10 \\ X_2 = 10 + D_3 \geq 0 \end{cases}$$

A unit increase in lower limit on humanistic course offering (i.e. from 10 to 11) decreases revenue by \$500

6

X_1 = Radio minutes

X_2 = TV minutes

X_3 = Newspaper ads

Maximize $Z = X_1 + 50X_2 + 5X_3$

s.t.

$$15X_1 + 300X_2 + 50X_3 \leq 10000 \quad (1)$$

$$X_3 \geq 5 \quad (2)$$

$$X_1 \leq 400 \quad (3)$$

$$-X_1 + 2X_2 \leq 0 \quad (4)$$

$$X_1, X_2, X_3 \geq 0$$

Solution: $Z = 1561.36$

$$X_1 = 59.09 \text{ min}, X_2 = 29.55 \text{ min}, X_3 = 5 \text{ ads}$$

continued...

Set 2.3c

(b) S_1, S_3, S_4 = slacks associated with constraints 1, 3, and 4
 S_2 = surplus associated with constraint 2

From TORA's optimum tableau:

$$\begin{aligned} Z + 2.879S_2 + .158S_1 + 0S_2 + 1.364S_3 &= 1561.36 \\ 59.091 + .006D_1 - .303D_2 &\quad -.909D_4 \geq 0 \\ 5 &+ D_2 \\ 340.909 - .006D_1 + .303D_2 + D_3 &+.909D_4 \geq 0 \\ 29.545 + .003D_1 - .152D_2 &+.045D_4 \geq 0 \end{aligned}$$

Constraint	Dual Price	RHS Range
1.	.158	(250, 66250)
2	-2.879*	(0, 2000)
3	0	(59.09, ∞)
4	1.3636	(-375, 65)

* Negative because S_2 is a surplus variable
 + These results are taken from TORA output. They differ from those computed from the given D_i -conditions because of roundoff error

Conclusions:

- Increasing the lower limit on the number of newspaper ads is not advantageous because the associated dual price is negative ($= -2.879$)
- Increasing the upper limit on radio minutes is not warranted because its dual price is zero (the current limit is already abundant).
- Dual price $= .158$ /budget \$ valid in the range $250 \leq \$ \leq 66250$.

50% budget increase = \$5000, or budget will be increased to 15,000.

Increase in $Z = .158 \times 5000 = 790$

(a) X_1 = Nbr. shirts / week
 X_2 = Nbr. blouses / week

Maximize $Z = 8X_1 + 12X_2$

s.t.

$$\begin{aligned} 20X_1 + 60X_2 &\leq 25 \times 60 \times 40 = 60,000 \\ 70X_1 + 60X_2 &\leq 35 \times 60 \times 40 = 84,000 \\ 12X_1 + 4X_2 &\leq 5 \times 60 \times 40 = 12,000 \\ X_1, X_2 &\geq 0 \end{aligned}$$

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Solution: $Z = \$13920/\text{week}$

$X_1 = 480$ shirts, $X_2 = 840$ blouses

(b) Let S_1, S_2 , and S_3 be the slack variables associated with the cutting, sewing, and packaging constraints. From the optimum TORA tableau, we have

$$Z + .12S_1 + .08S_2 + 0S_3 = 13920$$

Dept. Worth/hr (Dual price)

Cutting	\$.12 /	= \$7.20/hr
Sewing	\$.08/min	= \$4.80/hr
Packaging	\$ 0/hr	

(c) Break-even wages are \$7.20/hr for cutting and \$4.80 for sewing

(a) X_1 = units of solution A

X_2 = units of solution B

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Maximize $Z = 8X_1 + 10X_2$

s.t. $.5X_1 + .5X_2 \leq 150$ (1)

$.6X_1 + .4X_2 \leq 145$ (2)

$30 \leq X_1 \leq 150$ (3)

$40 \leq X_2 \leq 200$ (4)

Solution: $Z = \$2800$

$X_1 = 100$ units, $X_2 = 200$ units

(b) Define

S_1, S_2, S_3, S_4 = slacks in constraints 1, 2, 3, 4

S_5, S_6 = surplus variables associated with the lower bounds of constraints 3 and 4.

From TORA's optimum tableau:

$$Z + 16S_1 + 0S_2 + 0S_3 + 2S_4 + 0S_5 + 0S_6 = 2800$$

Conditions:

$$S_1 = 70 + 2D_1 - D_2 - D_5 \geq 0$$

$$S_2 = 5 - 1.2D_1 + D_2 + 2D_4 \geq 0$$

$$S_3 = 50 - 2D_1 + D_3 + D_4 \geq 0$$

$$X_1 = 100 + 2D_1 - D_4 \geq 0$$

$$X_2 = 200 + D_4 \geq 0$$

$$S_4 = 160 + D_4 - D_6 \geq 0$$

continued...

continued...

Set 3.6c

Constraint	Dual price	RHS-range
1	16	(115, 154.17)
2	0	(140, ∞)
3 (upper)	0	(100, ∞)
3 (lower)	0	(- ∞ , 100)
4 (upper)	2	(175, 270)
4 (lower)	0	(- ∞ , 200)

Increase in raw material 1 and in the upper bound on solution B is advantageous because their dual prices (16 and 2) are positive.

(c) Increase in revenue/unit = \$16

Increase in cost/unit = \$20

Not recommended!

(d) Dual price for raw material 2 is zero because it is abundant. No increase is warranted.

$$X_1 = \text{Nbr. } D_1 G_i - 1$$

$$X_2 = \text{Nbr. } D_2 G_i - 2$$

S_i = Idle minutes for station i , $i=1, 2, 3$

The objective is to minimize $S_1 + S_2 + S_3$.

To express the objective function in terms of X_1 and X_2 , consider

$$6X_1 + 4X_2 + S_1 = .9 \times 480 = 432$$

$$5X_1 + 4X_2 + S_2 = .86 \times 480 = 412.8$$

$$4X_1 + 6X_2 + S_3 = .88 \times 480 = 422.4$$

Thus, $S_1 + S_2 + S_3 = 1267.2 - 15X_1 - 14X_2$

(a)

Maximize $Z = 15X_1 + 14X_2$

s.t.

$$6X_1 + 4X_2 + S_1 = 432$$

$$5X_1 + 4X_2 + S_2 = 412.8$$

$$4X_1 + 6X_2 + S_3 = 422.4$$

$$X_1, X_2, S_1, S_2, S_3 \geq 0$$

Z represents the total used time in the three stations in minutes.

Solution: $Z = 1241.28$ minutes
 $X_1 = 45.12$ units, $X_2 = 40.32$ units

Utilization = $\frac{1241.28}{1267.20} \times 100 = 97.95\%$

continued...

(b) From TORA,

$$Z + 1.7S_1 + 0S_2 + 1.2S_3 = 1241.28$$

Conditions:

$$X_1 = .3D_1 - .2D_3 + 45.12 \geq 0$$

$$S_2 = -.7D_1 + D_2 - .2D_3 + 25.92 \geq 0$$

$$X_2 = -.2D_1 + .3D_3 + 40.32 \geq 0$$

Station	Dual Price	RHS Range
1	1.7	281.6, 469.03
2	0	386.88, ∞
3	1.2	288, 552

1% decrease in maintenance time is equivalent to $D_1 = D_2 = D_3 = 4.8$ minutes. This is equivalent to having Daily minutes

station	Daily minutes
1	436.8
2	417.6
3	427.2

All three daily minutes fall within the allowable ranges. Thus

station	Increase in utilized time/day
1	$4.8 \times 1.7 = 8.16$ minutes
2	$4.8 \times 0 = 0$
3	$4.8 \times 1.2 = 5.76$

$$(c) D_1 = .9(600 - 480) = 108 \text{ min}$$

$$D_2 = .86(600 - 480) = 103.2$$

$$D_3 = .88(600 - 480) = 105.6$$

From the conditions in (b)

$$X_1 = .3 \times 108 - .2 \times 105.6 + 45.12 = 56.4$$

$$S_2 = -.7 \times 108 + 103.2 - .2 \times 105.6 + 25.92 = 32.4$$

$$X_2 = -.2 \times 108 + .3 \times 105.6 + 40.32 = 50.4$$

Solution is feasible. Hence dual prices remain applicable and the net utilization is increased by $.1.7 \times 108 + 0 \times 103.2 + 1.2 \times 105.6 = 310.32$ minutes. Because station 2 has zero dual price, its capacity need not be increased. The associated cost thus equals $1.5(600 - 480) + 0 + 1.5(600 - 480) = \360 .

The proposal can be improved by recommending that station 2 time remain unchanged.

Set 3.6c

$$\begin{aligned} X_1 &= \text{Nbr. purses/day} \\ X_2 &= \text{Nbr. bags/day} \\ X_3 &= \text{Nbr. backpacks/day} \end{aligned}$$

$$\begin{aligned} \text{Maximize } Z &= 24X_1 + 22X_2 + 45X_3 \\ \text{s.t.} \end{aligned}$$

$$\begin{aligned} 2X_1 + X_2 + 3X_3 &\leq 42 \\ 2X_1 + X_2 + 2X_3 &\leq 40 \\ X_1 + .5X_2 + X_3 &\leq 45 \\ X_1, X_2, X_3 &\geq 0 \end{aligned}$$

Solution: $Z = \$882, X_1 = 0, X_2 = 2, X_3 = 36$

Letting S_1, S_2, S_3 be the slacks in constraints 1, 2, and 3, we get

$$Z + 20X_1 + S_1 + 21S_2 + 0S_3 = 882$$

Conditions:

$$\begin{aligned} X_3 &= 2 + D_1 - D_2 \geq 0 \\ X_2 &= 36 - 2D_1 + 3D_2 \geq 0 \\ S_3 &= 25 - .5D_2 + D_3 \geq 0 \end{aligned}$$

Resource	Dual price	RHS Range
Leather	1	(40, 60)
Sewing	21	(28, 42)
Finishing	0	(20, ∞)

(a) Available leather = 45 ft² falls in the RHS range. Solution remains feasible.

$$D_1 = 45 - 42 = 3. \text{ New solution:}$$

$$\begin{aligned} X_1 &= 0 \\ X_2 &= 36 - 2X_3 = 30 \\ X_3 &= 2 + 3 = 5 \end{aligned}$$

$$Z = 882 + 1 \times D_1 = 882 + 1 \times 3 = \$885$$

(b) Available leather = 41 ft² falls in the RHS range and the solution remains feasible. $D_1 = 41 - 42 = -1$

$$X_2 = 36 - (2 \times -1) = 38$$

$$X_3 = 2 - 1 = 1$$

$$Z = 882 + (-1) = \$881$$

(c) Sewing hours = 38 falls within the RHS range. $D_2 = 38 - 40 = -2$. Dual price = 21

$$X_2 = 36 + 3 \times -2 = 30$$

$$X_3 = 2 - (-2) = 4$$

$$Z = 882 + (21 \times -2) = \$840$$

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(d) Sewing hours = 46 hours falls outside the RHS range. Thus, the current optimum basic solution is infeasible. To obtain the new solution, either solve the problem anew or use the algorithms in chapter 4.

(e) Finishing hours = 15, which falls outside the RHS range. Hence, resolve the problem.

(f) Sewing hours = 50, which falls in the RHS range. $D_3 = 50 - 45 = 5$. Solution remains unchanged because dual price is zero and D_3 does not appear in the expression for X_2 or X_3 .

(g) Dual price = \$21/hr, which is higher than the cost of an additional worker per hour. Hiring is recommended.

$$X_1 = \text{Nbr. model 1 units}$$

$$X_2 = \text{Nbr. model 2 units}$$

$$\text{Maximize } Z = 3X_1 + 4X_2$$

s.t.

$$2X_1 + 3X_2 \leq 1200$$

$$2X_1 + X_2 \leq 1000$$

$$4X_2 \leq 800$$

$$X_1, X_2 \geq 0$$

Solution: $Z = \$1750$

$$X_1 = 450, X_2 = 100$$

(a) $S_1 = 0 \Rightarrow$ Resistors scarce

$S_2 = 0 \Rightarrow$ Capacitors scarce

$S_3 = 400 \Rightarrow$ chips abundant

$$(b) Z + \frac{5}{4}S_1 + \frac{1}{4}S_2 = 1750$$

Resource	Dual price
Resistors	\$1.25/resistor
Capacitors	\$0.25/capacitor
Chips	\$0/chip

(c) Conditions:

$$X_1 = 450 - \frac{1}{4}D_1 + \frac{3}{4}D_2 \geq 0$$

$$S_3 = 400 - 2D_1 + 2D_2 + D_3 \geq 0$$

$$X_2 = 100 + \frac{1}{2}D_1 - \frac{1}{2}D_2 \geq 0$$

Feasibility ranges:

$$\begin{cases} 450 - \frac{1}{4}D_1 \geq 0 \\ 400 - 2D_1 \geq 0 \\ 100 + \frac{1}{2}D_1 \geq 0 \end{cases} \Rightarrow -200 \leq D_1 \leq 200$$

continued...

continued...

Set 3.6c

$$\begin{aligned} 450 + .75D_2 &\geq 0 \\ 400 + 2D_2 &\geq 0 \\ 100 - .5D_2 &\geq 0 \end{aligned} \Rightarrow -200 \leq D_2 \leq 200$$

$$400 + D_3 \geq 0 \Rightarrow -400 \leq D_3 < \infty$$

(d) $D_1 = 1300 - 1200 = 100$ in the allowable range $-200 \leq D_1 \leq 200$.

$$\Delta Z = 100 \times 1.25 = \$125$$

$$X_1 = 450 - .25 \times 100 = 425$$

$$X_2 = 100 + .5 \times 100 = 150$$

$$\text{New } Z = 1750 + \Delta Z = \$1875$$

(e) $D_3 = 350 - 800 = -450$, which falls outside allowable range $-400 \leq D_3$.

Thus, basic solution and dual price change and the problem must be solved anew.

(f) $-200 \leq D_2 \leq 200$, dual price = .25.

Thus,

$$-200 \times .25 \leq \Delta Z \leq 200 \times .25$$

$$-50 \leq \Delta Z \leq 50$$

$$\$1700 \leq Z \leq \$1800$$

$$450 - .75 \times 200 \leq X_1 \leq 450 + .75 \times 200$$

$$100 - \frac{1}{2}(-200) \leq X_2 \leq 100 - \frac{1}{2}(+200)$$

(g) Cost of purchasing 500 additional resistors = $500 \times .40 = \$200$

$$D_1 = 500 \text{ resistors}$$

Dual price of \$1.25 is valid in $-200 \leq D_1 \leq 200$. Thus, for the first 200 resistors alone, H/D/C will get an additional revenue of $200 \times 1.25 = \$250$, which is more than the cost of all 500 resistors. Accept.

From Example 3.6-2, we have for the TOYCO model

$$\begin{aligned} -200 \leq D_1 &\leq 10 \\ -20 \leq D_2 &\leq 400 \\ -20 \leq D_3 &< \infty \end{aligned}$$

(g) $D_1 = 8, D_2 = 40, D_3 = -10$

All $D_i, i=1, 2, 3$ fall within the feasibility ranges. Thus

continued...

$$r_1 = \frac{8}{10}, r_2 = \frac{40}{400}, r_3 = \frac{-10}{-20}$$

$$r_1 + r_2 + r_3 = -8 + .1 + .5 = 1.4 > 1$$

Hence, no conclusion can be made about the feasibility of the new RHS (438, 500, 410). Problem 1(a) shows that these new values do produce a feasible solution.

(b) $D_1 = 30, D_2 = -20, D_3 = -40$. Because D_1 and D_3 fall outside the given feasibility ranges, the 100% rule cannot be applied in this case.

(a) From TORA,

$$x_1 = 2 + \frac{2}{3}D_1 + \frac{1}{3}D_2 \geq 0$$

$$x_2 = 2 - \frac{1}{3}D_1 + \frac{2}{3}D_2 \geq 0$$

Feasibility ranges:

$$-3 \leq D_1 \leq 6$$

$$-3 \leq D_2 \leq 6$$

(b) $D_1 = D_2 = \Delta > 0$. Thus

$$x_1 = 2 + \Delta/3 > 0 \quad \text{for all } \Delta > 0$$

$$x_2 = 2 + \Delta/3 > 0$$

100% rule for $0 < \Delta \leq 3$:

$$r_1 = r_2 = \frac{\Delta}{6} \leq \frac{3}{6} \Rightarrow r_1 + r_2 \leq 1, \text{ which confirms feasibility for } 0 < \Delta \leq 3$$

100% rule for $3 < \Delta \leq 6$:

$$r_1 = r_2 = \frac{\Delta}{6} \Rightarrow \frac{3}{6} \leq r_1, r_2 \leq \frac{6}{6}$$

$r_1 + r_2 \geq 1 \Rightarrow$ cannot confirm feasibility.

100% rule for $\Delta > 6$:

Δ is outside $-3 \leq D_1, D_2 \leq 6$. Thus, the rule is not applicable.

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Set 3.6d

From Section 3.6.3, we have the following optimality conditions for the TOYCO model:

$$x_1: 4 - \frac{1}{4}d_2 + \frac{3}{2}d_3 - d_1 \geq 0$$

$$x_4: 1 + \frac{1}{2}d_2 \geq 0$$

$$x_5: 2 - \frac{1}{4}d_2 + \frac{1}{2}d_3 \geq 0$$

$$(i) Z = 2x_1 + x_2 + 4x_3$$

$$d_1 = 2-3 = -1, d_2 = 1-2 = -1, d_3 = 4-5 = -1$$

$$x_1: 4 - \frac{1}{4}(-1) + \frac{3}{2}(-1) - (-1) = 3.75 > 0$$

$$x_4: 1 + \frac{1}{2}(-1) = -0.5 > 0$$

$$x_5: 2 - \frac{1}{4}(-1) + \frac{1}{2}(-1) = 1.75 > 0$$

Conclusion: Solution is unchanged

$$(ii) Z = 3x_1 + 6x_2 + x_3$$

$$d_1 = 3-3 = 0, d_2 = 6-2 = 4, d_3 = 1-5 = -4$$

$$x_1: 4 - \frac{1}{4}(4) + \frac{3}{2}(4) - (0) = -3 < 0$$

Conclusion: Solution changes

$$(iii) Z = 8x_1 + 3x_2 + 9x_3$$

$$d_1 = 8-3 = 5, d_2 = 3-2 = 1, d_3 = 9-5 = 4$$

$$x_1: 4 - \frac{1}{4}(1) + \frac{3}{2}(4) - (5) = 4.75 > 0$$

$$x_4: 1 + \frac{1}{2}(1) = 1.5 > 0$$

$$x_5: 2 - \frac{1}{4}(1) + \frac{1}{2}(4) = 3.75 > 0$$

Conclusion: Solution is unchanged

x_1 = Nbr. cans of A1

x_2 = Nbr. cans of A2

x_3 = Nbr. cans of BK

$$\text{Maximize } Z = 80x_1 + 70x_2 + 60x_3$$

$$\text{s.t. } x_1 + x_2 + x_3 \leq 500 \quad \leftarrow S_1$$

$$x_1 \geq 100 \quad \leftarrow S_2$$

$$4x_1 - 2x_2 - 2x_3 \leq 0 \quad \leftarrow S_3$$

$$x_1, x_2, x_3 \geq 0$$

TORA optimum tableau:

Basic	x_1	x_2	x_3	S_1	S_2	S_3	Solution
Z	0	0	10	73.33	0	1.67	36666.67
x_2	0	1	1	-0.67	0	-0.17	333.33
x_1	1	0	0	0.33	0	0.17	166.67
S_2	0	0	0	0.33	1	0.17	66.67

continued...

(a) $Z = \$366.67$

$$x_1 = 166.67, x_2 = 333.33, x_3 = 0$$

(b) Reduced cost for $x_3 = 10$ cents. Price should be increased by more than 10 cents/can

(c) $d_1 = d_2 = d_3 = -5$ cents

From the optimum tableau, reduced costs:

$$x_3: 10 + d_2 - d_3 = 10 - 5 - (-5) = 10 > 0$$

$$S_1: 73.33 + .67d_2 + .33d_3 = 73.33 + .67(-5) + .33(-5) = 68.33 > 0$$

$$S_3: 1.67 - .17d_2 + .17d_3 = 1.67 - .17(-5) + .17(-5) = 1.67 > 0$$

Conclusion: Solution is unchanged.

(d) Available carpenter hours in a 10-day period = $4 \times 10 \times 8 = 320$

x_1 = Nbr. chairs assembled in 10 days

x_2 = Nbr. tables assembled in 10 days

$$\text{Maximize } Z = 50x_1 + 135x_2$$

s.t.

$$5x_1 + 2x_2 \leq 320$$

$$4 \leq \frac{x_1}{x_2} \leq 6 \Rightarrow \begin{cases} x_1 - 4x_2 \geq 0 \\ x_1 - 6x_2 \leq 0 \end{cases}$$

$$x_1, x_2 \geq 0$$

$$\text{Solution: } Z = \$27,840, x_1 = 384, x_2 = 64$$

(e) Optimum tableau:

	x_1	x_2	S_1	S_2	S_3	Solution
Z	0	0	87	0	6.5	27840
x_2	0	1	-2	0	-1	64
x_1	1	0	1.2	0	.4	384
S_2	0	0	.4	1	.8	128

Optimality conditions:

$$S_1: 87 + 1.2d_1 + 2d_2 \geq 0$$

$$S_3: 6.5 + 4d_1 - 1d_2 \geq 0$$

For $d_1 = -5, d_2 = -13.5$:

$$S_1: 87 + 1.2(-5) + 2(-13.5) = 78.3 > 0$$

$$S_3: 6.5 + 4(-5) - 1(-13.5) = 5.85 > 0$$

Solution remains the same

$$(f) d_1 = 25 - 50 = -25, d_2 = 120 - 135 = -15$$

$$S_1: 87 + 1.2(-25) + 2(-15) = 58.5 > 0$$

$$S_3: 6.5 + 4(-25) - 1(-15) = -2 < 0$$

Solution changes

Set 3.6d

(a) x_1 = Amt. of personal loan (\$) x_2 = Amt. of car loan (\$)	4
Maximize $Z = .14(x_1 - .03x_1) + .12(x_2 - .02x_2)$ $= .1058x_1 + .0976x_2$	
S.t.	
$x_1 + x_2 \leq 200,000$	
$\frac{x_2}{x_1} \geq 2$ or $2x_1 - x_2 \leq 0$	
$x_1, x_2 \geq 0$	
Solution: $Z = \$20,067$ $x_1 = \$66,667, x_2 = \$133,333$	
Rate of return = $\frac{20,067}{200,000} \times 100 = 10.03\%$	

(b) Optimum tableau:

x_1	x_2	s_1	s_2	Solution
0	0	.1003	.0027	200666.67
0	1	.6667	-.3333	133333.33
1	0	.3333	.3333	66666.67

Optimality conditions:

$$S_1: .1003 + .333d_1 + .6667d_2 \geq 0$$

$$S_2: .0027 + .3333d_1 - .3333d_2 \geq 0$$

$$\text{New } x_1 \text{-objective coefficient} = .14(1 - .04) - .04 \\ = .0944$$

$$\text{New } x_2 \text{-objective coefficient} = .12(1 - .03) - .03 \\ = .0864$$

$$d_1 = .0944 - .1058 = -.0114$$

$$d_2 = .0864 - .0976 = -.0112$$

$$S_1: .1003 + .3333(-.0114) + .6667(-.0112) \\ = .08907 > 0$$

$$S_2: .0027 + .3333(-.0114) - .3333(-.0112) \\ = .00267 > 0$$

Solution does not change

(a) x_i = Nbr of units of motor $i, i=1, 2, 3, 4$	5
Maximize $Z = 60x_1 + 40x_2 + 25x_3 + 30x_4$	
S.t.	
$8x_1 + 5x_2 + 4x_3 + 6x_4 \leq 8000$	
$x_1 \leq 500, x_2 \leq 500, x_3 \leq 800, x_4 \leq 750$	
$x_1, x_2, x_3, x_4 \geq 0$	
Solution: $Z = \$59,375, x_1 = 500, x_2 = 500, x_3 = 375$	
$x_4 = 0$	

continued...

(b) Optimality conditions (from TORA):

$$x_4: 7.5 + 1.5d_3 - d_4 \geq 0$$

$$S_1: 6.25 + .25d_3 \geq 0$$

$$S_2: 10 - 2d_3 + d_1 \geq 0$$

$$S_3: 8.75 - 1.25d_3 + d_2 \geq 0$$

$$\text{From } S_3, 8.75 + d_2 \geq 0 \Rightarrow -8.75 \leq d_2 < \infty$$

Thus, price of type 2 motor can be reduced by at most \$8.75 without causing a solution change.

$$(c) d_1 = -15, d_2 = -10, d_3 = -6.25, d_4 = -7.5$$

Solution remains the same because

$$x_4: 7.5 + 1.5(-6.25) - (-7.5) = 5.625 > 0$$

$$S_1: 6.25 + .25(-6.25) = 4.6875 > 0$$

$$S_2: 10 - 2(-6.25) + (-15) = 7.5 > 0$$

$$S_3: 8.75 - 1.25(-6.25) + (-10) = 6.5625 > 0$$

(d) Reduced cost for $x_4 = 7.5$. Increase price of type 4 motor by more than \$7.50.

$$(a) x_1 = Cases of juice/day$$

$$x_2 = Cases of sauce/day$$

$$x_3 = Cases of pasta/day$$

$$\text{Maximize } Z = 21x_1 + 9x_2 + 12x_3$$

S.t.

$$(1 \times 24)x_1 + (\frac{1}{2} \times 24)x_2 + (\frac{3}{4} \times 24)x_3 \leq 60,000$$

$$x_1 \leq 2000, x_2 \leq 5000, x_3 \leq 6000$$

$$x_1, x_2, x_3 \geq 0$$

$$\text{Solution: } Z = \$51,000$$

$$x_1 = 2000, x_2 = 1000, x_3 = 0$$

(b) From TORA, optimality conditions given d_2 :

$$x_3: 1.5 + 1.5d_2 \geq 0 \Rightarrow d_2 \geq -1$$

$$S_1: .75 + .083d_2 \geq 0 \Rightarrow d_2 \geq -9$$

$$S_2: 3 - 2d_2 \geq 0 \Rightarrow d_2 \leq 1.5$$

Thus, $-1 \leq d_2 \leq 1.5$, or

$$9-1 \leq \text{price/case of sauce} \leq 9+1.5$$

Solution mix remains the same if the price per case of sauce remains between \$8 and \$10.50.

Set 3.6d

(a) $X_1 = \text{Nbr. regular cabinets/day}$
 $X_2 = \text{Nbr. deluxe cabinets/day}$
 Maximize $Z = 100X_1 + 140X_2$
 s.t.

$$5X_1 + X_2 \leq 180$$

$$X_1 \leq 200$$

$$X_2 \leq 150$$

$$X_1, X_2 \geq 0$$

Solution: $Z = \$31,200$
 $X_1 = 200$ regular
 $X_2 = 80$ deluxe

(b) From TORA, optimality conditions:

$$S_1: 140 + d_2 \geq 0$$

$$S_2: 30 + d_1 - .5d_2 \geq 0$$

$$d_1 = 80 - 100 = -20$$

$$d_2 = 80 - 140 = -60$$

$$S_1: 140 + (-60) = 80 > 0$$

$$S_2: 30 + (-20) - .5(-60) = 40 > 0$$

Solution remains the same

(a) For the original TOYCO model,
 TORA gives (also see Section 3.6.3)

$$-\infty < d_1 \leq 4, -2 \leq d_2 \leq 8, -8/3 \leq d_3 < \infty$$

(ii) Original $Z = 3X_1 + 2X_2 + 5X_3$

$$\text{New } Z = 3X_1 + 6X_2 + X_3$$

i	d_i	u_i	v_i	r_i
1	0	4	0/4 = 0	
2	4	8	4/8 = 1/2	
3	-4	-8/3	-4/-8/3 = 3/2	

$$r_1 + r_2 + r_3 = 0 + 1/2 + 3/2 = 2 > 1$$

The 100% rule is nonconclusive in this case. The solution in Problem 1 (ii) shows that the solution will change.

(iii) Original $Z = 3X_1 + 2X_2 + 5X_3$

$$\text{New } Z = 8X_1 + 3X_2 + 9X_3$$

i	d_i	u_i	v_i	r_i
1	5	4	5/4	
2	1	8	1/8	
3	4	0	4/0 = 0	

$$r_1 + r_2 + r_3 = \frac{5}{4} + \frac{1}{8} = \frac{11}{8} > 1$$

continued...

7

The 100% rule is nonconclusive. Yet Problem 1 (iii) shows that the solution remains unchanged.

The two cases demonstrate that the 100% rule is too weak to be effective in decision making, and that it is more reliable to utilize the simultaneous optimality conditions given in Section 3.6.3.

$$(b) -30 \leq d_1 < \infty, -140 \leq d_2 \leq 60$$

$$\text{New } Z = 80X_1 + 80X_2$$

$$\text{Original } Z = 100X_1 + 140X_2$$

i	d_i	u_i	v_i	r_i
1	-20	-30	0	-20/-30 = 2/3
2	-60	-140	60	-60/-140 = 3/7

$$r_1 + r_2 = \frac{2}{3} + \frac{3}{7} = \frac{23}{21} > 1$$

The 100% rule is nonconclusive. Yet, Problem 7(b) shows that the solution remains unchanged.

8

(a) For the original TOYCO model,
 TORA gives (also see Section 3.6.3)

$$-\infty < d_1 \leq 4, -2 \leq d_2 \leq 8, -8/3 \leq d_3 < \infty$$

(ii) Original $Z = 3X_1 + 2X_2 + 5X_3$

$$\text{New } Z = 3X_1 + 6X_2 + X_3$$

i	d_i	u_i	v_i	r_i
1	0	4	0/4 = 0	
2	4	8	4/8 = 1/2	
3	-4	-8/3	-4/-8/3 = 3/2	

$$r_1 + r_2 + r_3 = 0 + 1/2 + 3/2 = 2 > 1$$

The 100% rule is nonconclusive in this case. The solution in Problem 1 (ii) shows that the solution will change.

(iii) Original $Z = 3X_1 + 2X_2 + 5X_3$

$$\text{New } Z = 8X_1 + 3X_2 + 9X_3$$

i	d_i	u_i	v_i	r_i
1	5	4	5/4	
2	1	8	1/8	
3	4	0	4/0 = 0	

$$r_1 + r_2 + r_3 = \frac{5}{4} + \frac{1}{8} = \frac{11}{8} > 1$$

3-31

Set 3.6e

See file solver3.6e-1.xls in ch3Files

Dual prices for years 1, 2, 3, and 4 are 0, 0, 0, 2.89. Thus, for year 4, one (thousand) additional dollars increases Z by \$2.89 thousand. It is worthwhile to increase the funding for year 4.

See file tora3.6e-2.txt

Constraint	Dual Price	Range
1	5.36	(0, ∞)
2	-3.73	(-∞, 6000)
3	-1.13	(-∞, 6800)
4	-1.07	(-∞, 33642)
5	-1.00	(-∞, 53628.73)

2

(a) Constraint 1: $x_1 + x_2 + x_3 + y_1 \leq 10,000$

Dual price = \$5.36 / invested \$

Rate of return = 536%

(b) Constraint 2: \$1000 spent on pleasure

$$.5x_1 + .6x_2 - x_3 + .4x_4 + 1.065y_2 - y_1 = 1000$$

Dual price = -3.73 / pleasure \$

Range = (-∞, 6000)

Spending \$1000 at end of year 1 reduces total return by \$3730.

See file tora3.6e-3.txt in ch3Files

Quarter	Dual price	Range
1	1.2488	.6647, 2.5806
2	1.2443	.6580, 2.6122
3	1.1945	-.2646, 1.1245
4	1.0200	-.2553, 00
5	1.0000	-4.8366, 00

3

(a) An additional \$ available at the start of quarter 1 is worth \$1.2488 at the end of 4 quarters. Similarly, an additional dollar at the start of periods 2, 3, and 4 is worth \$1.2443, \$1.1945, and \$1.02, respectively. The dual price for quarter 4 (= \$1.02) shows that all we can do with the money then is to invest it at 2% for the quarter.

We can use the dual price to determine

continued...

the rate of return for each quarter — namely,

quarter 1:

$$1.2488 = 1.2243(1+i_1) \Rightarrow i_1 = .02$$

quarter 2:

$$1.2243 = 1.1945(1+i_2) \Rightarrow i_2 = .025$$

quarter 3:

$$1.1945 = 1.02(1+i_3) \Rightarrow i_3 = .171$$

quarter 4:

$$1.02 = 1.0(1+i_4) \Rightarrow i_4 = .02$$

(b) The dual price associated with the upper bound on B_3 (UB-X10) is \$.149. It represents the networth per dollar borrowed in period 3. Also, an extra dollar in period 3 is worth \$.11945 at the end of the horizon. However, if that dollar is borrowed, it must be repaid as \$1.025 in the next quarter. The repayment is equivalent to forgoing making 2% in interest. Thus, the networth of borrowing in period 3 is

$$1.1945 - 1.025 \times 1.02 = .149$$

This result is consistent with the dual price for the upper bound on B_3 .

Constraint	Current RHS	Min RHS	Max RHS	Dual Price
1 (=)	2.0000	0.0000	infinity	2.1756
2 (=)	2.0000	-0.1667	infinity	2.0173
3 (=)	2.5000	-0.3472	infinity	1.8647
4 (=)	2.5000	-0.5767	infinity	1.7296
5 (=)	3.0000	-0.8248	infinity	1.6044
6 (=)	3.5000	-1.1331	infinity	1.4356
7 (=)	3.5000	-6.1137	infinity	1.3355
8 (=)	4.0000	-11.4678	infinity	1.2423
9 (=)	6.0000	-20.6663	infinity	1.1556
10 (=)	9.0000	-32.5201	infinity	1.0759

4

The dual price provides the worth per additional \$ at the end of year 10.

Annual rate of return:

$$\text{Period 1: } 2.1756 = 2.0173(1+i_1) \Rightarrow i_1 = .0785$$

$$\text{Period 2: } 2.0173 = 1.8647(1+i_2) \Rightarrow i_2 = .0818$$

$$\text{Period 3: } 1.8647 = 1.7296(1+i_3) \Rightarrow i_3 = .0781$$

$$\text{Period 4: } 1.7296 = 1.6044(1+i_4) \Rightarrow i_4 = .0780$$

etc...

Set 3.6e

See file *tora3.6e-5.txt* in ch3files

The dual price for constraint 1

$$x_{1A} + x_{1B} \leq 100,000$$

is \$5.10. Thus, each invested \$ is worth \$5.10 at the end of the investment horizon. Range (0, ∞)

5

See file *tora3.6e-9.txt* in ch3files

(a) Constraint $2x_1 + 3x_2 + 5x_3 \leq 4000$

corresponds to raw material A. Its dual price is \$10.27/lb. For a purchase price of \$12/lb, acquisition of additional raw material A is not recommended.

9

(b) Constraint $4x_1 + 2x_2 + 7x_3 \leq 6000$

is associated with raw material B. Its dual price is \$0/lb. Resource B is already abundant. Thus, no additional purchase is recommended.

Dual price for the constraint

6

$$x_1 + x_2 + x_3 + x_4 \leq 500$$

is \$2.35 per \$ invested, range (0, ∞)

The gambler should bet the largest amount possible.

See file *tora3.6e-7.txt* in ch3files

7

For, $x_{W1} + x_{W2} + x_{W3} \geq 1500$, the dual price is \$11.4, range (800, ∞)

One extra wrench automatically requires the production of two chisels, thus leading to the following changes:

Cost of one wrench using subcont. = \$3.00
Cost of 2 chisels using subcont. = $2 \times \$4.20$
total = \$11.40

$x_{W1} \leq 550$, dual price = -\$1, range (- ∞ , 1250). If regular time capacity for wrenches is increased by 1 unit, one less wrench will be produced by subcontractor, which saves $\$3 - \$2 = \$1$.

Similar interpretations can be given for the remaining dual prices

(a) See file *tora3.6e-10.txt*

10

Constraint	Dual price
1	0
2	0
3	-400
4	-750
5	0
6	0
7	0

Constraints 3 and 4 have negative dual price. These correspond respectively to the third specification for alloy A and the first specification for alloy B. Changes in these specifications affects profit adversely.

(b) For the ore constraints, the dual prices are \$90, \$110, and \$30 per additional ton of ores 1, 2, and 3, respectively. These are the maximum prices the company should pay.

See file *tora3.6e-8.txt* in ch3files

8

Machine Capacity Dual price Range

1	500	2	(253.33, 570)
2	380	12	(333.33, 750)

The company should pay less than \$2/hr for machine 1 and less than \$12/hr for machine 2.

CHAPTER 5

Transportation Model and its Variants

Set 5.1a

- (a) False
 (b) True
 (c) True

(a) $\sum a_i = 25, \sum b_j = 31$

Add a dummy source whose supply amount is $31 - 25 = 6$ units

(b) $\sum a_i = 74, \sum b_j = 65$

Add a dummy destination whose demand amount is $74 - 65 = 9$ units

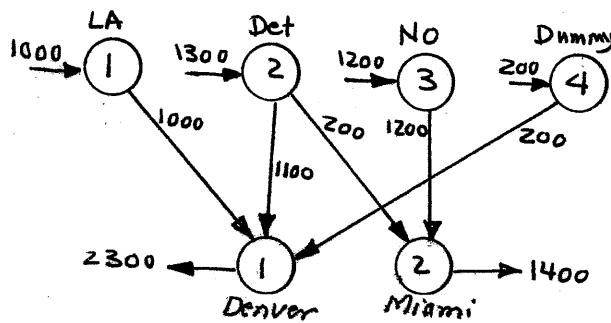
Denver will be 150 cars short.

Similarly, Miami will be 50 cars short of satisfying its demand

Assign a very high cost M
 to the route from Detroit to Dummy

	Den	Miami	
	1	2	
LA 1	80	M	
	1000		1000
	100	108	
Det 2	1100	200	1300
	102	68	
NO 3		1200	1200
	200	300	
Dummy 4	200		200
	2300	1400	

Optimum solution from TORA

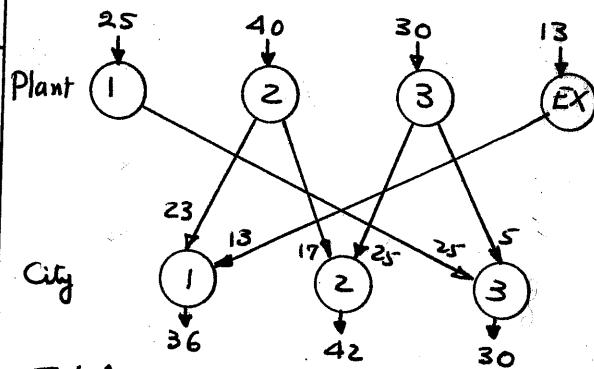


Denver is 200 cars short, Cost = \$33,200

5-2

	City			
	1	2	3	
Plant 1	600	700	400	25
Plant 2	320	300	350	40
Plant 3	500	480	450	30
Excess plant 4	1000	1600	M	13
	36	42	30	

(b) $M = \$10,000$ in TORA

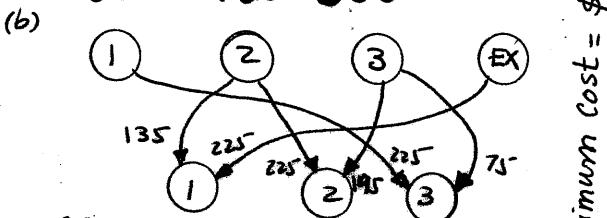


Total cost = \$ 49,710

(c) City 1 excess cost = $13 \times 1000 = \$13,000$

Assume units in 100,000 kWh

	1	2	3	
Plant 1	60	70	40	225
2	32	30	35	360
3	50	48	45	270
EX	100	100	M	225
	360	420	300	



(c) City 1 excess cost = \$22,500

Optimum cost = \$55,305*

Set 5.1a

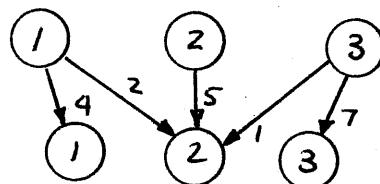
Unit transportation cost in
thousand \$ per million
ton miles

$$\text{gallons} = \left(\frac{104}{1000} \times 10^6 \times \text{mileage} \right) \times \frac{1}{100} \times \frac{1}{1000}$$

$$= \frac{\text{mileage}}{1000000}$$

¹⁰ Distribution Area

Ref.	1	2	3
4	12	18	M
5	30	10	8
6	20	25	12
7	1	7	8



Total cost = \$243,000

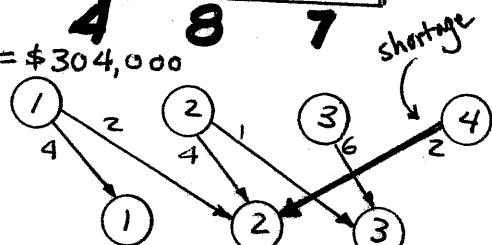
Unit cost in thousand \$ from
Dummy source to distribution
areas 2 or 3 6

$$= \frac{5\phi}{100} \times \frac{10^6}{10^3} = 50 \text{ thousand \$ / million gal}$$

Distribution area

		2	3	
Ref.	1	12	18	M
2	30	4	10	1 8
3	20		25	12
Dummy	M	2	50	50

Cost = \$304,000



8

Unit costs in thousand \$ per million gallons:

from refinery 1 to Dummy

$$= \frac{\$1.50}{100} \times \frac{10^6}{10^3} = 15$$

from refinery 2 to Dummy

$$= \frac{\$2.20}{100} \times \frac{10^6}{10^3} = 22$$

	1	2	3	Dummy
Ref. 1	4	2	M	15
2	30	5	10	22
3	20	1	25	12
	4	8	4	3

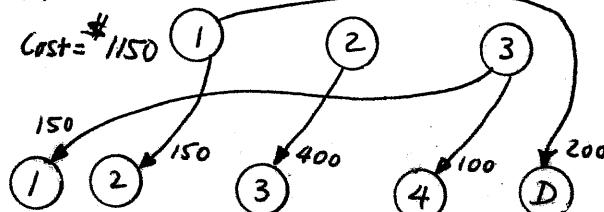
Refinery 3 diverts 3 million gallons for use within.

Total cost = \$20,700

$$\begin{aligned} \text{(a) Total supply} &= 150 + 200 + 250 = 600 \text{ crates} \\ \text{Total demand} &= 150 + 150 + 400 + 100 = 800 \text{ crates} \\ (\text{Potential overtime supply}) &= 800 - 600 = 200 \\ (\text{by each of orchards 1 \& 2}) \end{aligned}$$

	1	2	3	4	Dummy
Arch 1	1 (150)	2 400	3 400	2 100	0 200
2	2 150	4 150	1 400	2 100	0 200
3	1 150	3 150	5 400	3 100	M 200
	150	150	400	100	200

(6)



Problem has alternative optima.

(c) Orchard 1 = 0 overtime costs

Orchad 2 = 200 overtime system

Set 5.1a

Supply/demand quantities are expressed in truck loads, determined by dividing the number of cars by 18 and rounding the result up, if necessary. For example, supply amount at center 1 is $\frac{400}{18} = 22.22$ or 23 truck loads.

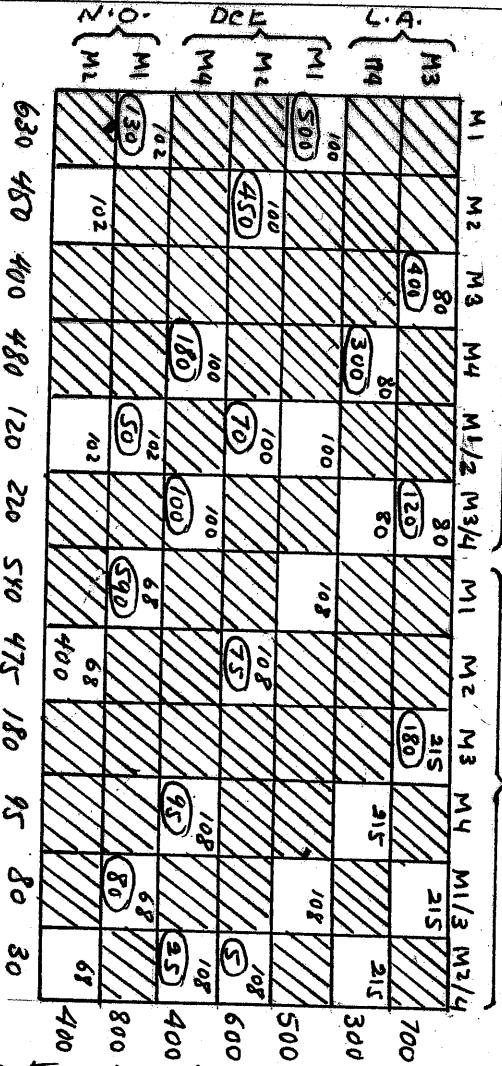
¹⁸ Expressing unit transportation costs in \$1000 per truck load, we get

	1	2	3	4	5	
1	2.5	3.75	5	3.5	8.75	23
2	1.25	1.75	1.5	1.625	2	12
3	1	2.25	2.5	3.75	3.25	9
	6	12	9	9	8	

(b) Alternative solution exists
Cost = \$ 92,500

12

13



Optimum Solutions:

LA - Denver M4	=	300	CARS
Det. - Denver M1	=	500	CARS
Det. - Denver M2	=	450	CARS
Det. - Denver M1/M2	=	70	CARS
Det. - Miami M2	=	75	CARS
Det. - Miami M2/4	=	5	CARS
Det. - Denver M4	=	180	CARS
Det. - Denver M3/4	=	100	CARS
Det. - Miami M4	=	95	CARS
Det. - Miami M2/4	=	25	CARS
N.O. - Denver M1	=	130	CARS
N.O. - Denver M1/2	=	50	CARS
N.O. - Miami M1	=	540	CARS
N.O. - Miami M1/3	=	80	CARS
N.O. - Miami M2	=	400	CARS
Total Cost =		\$343,620	

5-3A

Set 5.2a

	1	2	3	4		1	Sharpening Service				
1	40 50	40.4	40.7	41.4	50	Day	New	Overnite	2-day	3-day	Disposal
2	42 50	40	40.3	41	180	Mon	24	0	6	18	0
3	44 70	42	40	40.7	280	Tue	12	12	0	0	0
4	46 270	44	42	40	270	Wed	2	14	0	0	0
	100	200	180	300		Thu	0	0	20	0	0
						Fri	0	14	0	0	4
						Sat	0	2	0	0	12
						Sun	0	0	0	0	22

	1	2	3	4		1	3
	50	180	280	270		50	12
	100	200	180	300		12	12
	50	50	130	70		12	12
	200	200	180	180		12	12
	300	300	300	270		12	12

Cost = \$31,461

	Least-cost starting solution. (Problem has alternative optima.)						2		
	Mon	Tue	Wed	Thu	Fri	Sat	Sun	Disposal	
New	12 24	12 12	12 2	12	12	12	12	0	124
Mon	6	6	3 6	1 18	1	1	1	0	24
Tue	6	6	3	1	1	1	0		12
Wed	6	6	3	1	1	0			14
Thu	6	6	3	0					20
Fri				6	6	0			18
Sat				14			6	0	14
Sun				20			14	22	22
	24	12	14	20	18	14	22	124	

	Day	New	Overnite	2-day	Disposal
Mon	24	12	12	0	
Tue	0	6	6	0	
Wed	8	8	6	0	
Thu	0	12	8	0	
Fri	0	8	0	10	
Sat	0	14	0	0	
Sun	0	0	0	22	

The given optimum solution is interpreted as summarized below.

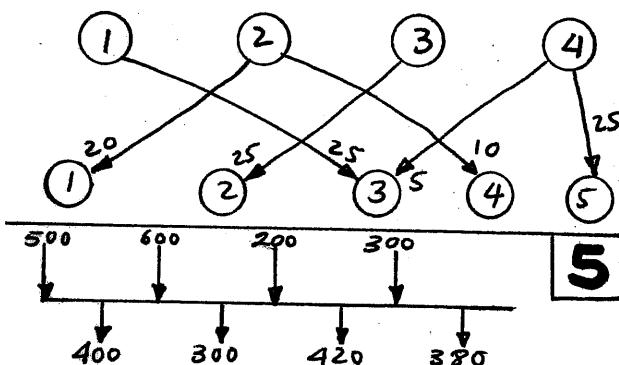
Total cost = \$840

continued...

Set 5.2a

	Task					
	1	2	3	4	5	4
Machine	10	2	3	15	9	25
1			(25)			
2	5	10	15	(10)	4	30
3	15	5	14	7	15	20
4	20	15	13	M	(25)	30
	20	20	30	10	25	

Total cost = \$560



C: \$100 \$140 \$120 \$150
h: \$3 \$3 \$3 \$3

	1	2	3	4	Surplus
1	100	103	106	M	0
2	M	140	143	146	0
3	M	M	120	123	0
4	M	M	M	150	0
	400	300	420	380	100

Cost = \$190,040, Alternative solution exists

Period	Capacity	Actual Prod.	Delivery
1	500	500	400 for 1 100 for 2
2	600	600	200 for 2 220 for 3 180 for 4
3	200	200	200 for 3
4	300	200	200 for 4

	1	2	3	4	5	Surplus	
R ₁	100	180					180
O ₁	150	20	154	158	162	166	90
R ₂		96	150	80	104	108	230
O ₂		144	148	152	156	115	115
R ₃			116	220	120	124	430
O ₃			174	178	182	186	215
R ₄				102	106	0	300
O ₄				250	50	150	150
R ₅					153	157	300
O ₅					106	0	150
					300	159	150
					200	150	860

Cost = \$137,720

Alternative solution exists.

Period	Production schedule
1	Regular - 180 engines Overtime - 20 engines
2	Regular: 230 engines
3	Regular 270 engines
4	Regular 300 engines
5	Regular 300 engines

Set 5.2a

New	1	2	3	4	5	6	Disposal
	200	210	220.5	231.53	243.1	255.26	0
1	200	180	140				878
	120	121.5	35	36.5	38	0	1398
2		120	121.5	35	36.5	0	
		148	32				
3		120	121.5	35	0		
		10	290				
4		120	121.5	0			
		198					
5		120	0				
		230					
6			0				
	200	180	300	198	230	290	1398

Cost = \$ 170,698

Alternative solution exists

Month	New	overhaul			Disposal
		1-day	3-day		
1	200	12	188	0	
2	180	148	32	0	
3	140	10	290	0	
4	0	198	0	0	
5	0	0	0	230	
6	0	0	0	290	

7

(a) Use negative cost values

		Bidder			
Loc		1	2	3	4
1		-520	M	-650	-180
2		-210	-510	M	-430
3		-570	-495	-240	-710
Dummy		30	10	20	0
		30	30	30	30

8

(b) Bidder 1 = 0 acre

Bidder 2 = 20 acres (location 1)

Bidder 3 = 10 acres (location 2)

Bidder 4 = 30 acres (location 3)

Set 5.3a

(a)

Northwest:

Cost = \$42

0	2	1
5	1	0
2	4	3

6

7

7

5 5 10

Least-cost:

Cost = \$37

0	2	1
5	1	0
2	4	3

6

7

5 5 10

Vogel:

Cost = \$37

0	2	1	Penalty	Penalty
5	1	0	1	1
2	5	1	1	4
2	4	3	1	1

6

7

7

Penalty

(2) 1 2 ← Step 1

Penalty

- 1 2 ← Step 2

(b)

Northwest:

Cost = \$94

1	2	6
7		
3	4	2

7

12

11

10 10 10

Least-Cost:

Cost = \$61

1	2	6
10	4	2
3	10	5

7

12

11

10 10 10

continued...

VAM:

Cost = \$40

1	2	6	Penalties
7	0	4	1 1 1
2	10	2	2 4 -
1	10	5	2 2 2

Penalties

{ 1

2

2

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Set 5.3b

(i)

<u>u</u>	<u>v</u>	0	2	6	
0		(5) 0	1 2	5	6
-1		2	1	5	9
-3		4	5		
-5		2	4	(5) 3	5
		5	5	10	

<u>u</u>	<u>v</u>	0	-3	1	
0		(5) 0	-5	1	6
4		2	5	4	9
7		2	4	(5) 3	5
		5	5	10	

<u>u</u>	<u>v</u>	0	-1	1	
0		(1) 0	2	(5) 1	6
2		(4) 2	(5) 1	5	9
2		2	4	(5) 3	5
		5	5	10	

Cost = \$33

Alternative solution exists

(ii)

<u>u</u>	<u>v</u>	0	4	2	
0		(7) 0	1 4	0	8
-1		2	3	4	5
-3		(5)	-3		
-2		1	2	(6) 0	6
		7	6	6	

Problem has alternative optima. Cost = \$19

Note: If x_{23} were selected as the zero in place of x_{32} , solution would require one more iteration.

<u>u</u>	<u>v</u>	M	M-3	M-5	
0		(4) M	3	5	4
7-M		1	M-6	M-10	7
11-M		1	8	(19) 6	19
		5	6	19	

<u>u</u>	<u>v</u>	6	3	1	
0		6-M	(4) 3	-4	4
1		(5) 7	(2) 4	9	7
5		1	8	(19) 6	19
		10	6	19	

<u>u</u>	<u>v</u>	6	3	11	
0		6-M	(4) 3	5	4
1		(5) 7	(2) 4	9	7
-5		1	8	(19) 6	19
		10	6	19	

<u>u</u>	<u>v</u>	0	-3	5	
0		-M	-6	(4) 5	4
1		(1) -	(6) 3	9	7
1		(4) +	-10	(15) -	19
		5	6	19	

<u>u</u>	<u>v</u>	0	0	5	
-M		-3	(4) 5	4	
-3		7	(6) 4	(1) 9	7
(5)		1	8	(14) 6	19
		5	6	19	

Cost = \$142

continued...

Set 5.3b

(c)

Method

	(i)	Nbr. of iterations	(ii)	(iii)
NW	3	4	5	
Least cost	2	2	2	
Vogel	2	1	1	

Least-cost starting solution:

u v

2	1	3	
5	(10)	1	7
-3		-5	
(70) 6	-1	4	(10) 6
(5) 3	+ (10)	2	5
5		-2	
-3	-2		(40)

75 20 50

u v

3	1	2	
5	(10)	1	7
-2		-4	
(60) 6	(10)	4	(10) 6
(15) 3	3	2	5
0	-1	-2	
-1	5	3	2
-3	-3		(40)

75 20 50

Destination 3 will be 40 units short. Optimum cost = \$ 595

Least-cost starting solution:

u v

2	1	1	
5	10	1	7
-3		-6	
30 6		4	50 6
5 3	10	2	5
-2	40 0	-1	M

75 20 50

3	1	2	
5	10	1	7
-2		-4	
20 6	10	4	50 6
15 0		2	5
40 0	-2	-M	M

75 20 50

Total cost = \$ 515. Dest. 1 is 40 units short.

Vogel method:

1	2	3	5
3	4	5	M
2	3	3	20 3
1	1	2	2

1	2	3	5
3	4	5	
2	3	3	
1	1	2	

3	4	5	
2	3	0 3	
1	1	2	

20 3	20 4		
10 2	3		
1	1		

Set 5.3b

$u \backslash v$	0	1	1	1	
0	-1	-1	20	1	20
3	20	20	4	5	M
2	10	2	3	3	20
	30	20	20	20	30

Cost = \$240 - Alternative solution exists

$u \backslash v$	2	5	10	
-2	(15)	c_{12}	c_{13}	15
3	(5)	(25)	c_{23}	30
5	c_{31}	(5)	(80)	c_{33}
	20	30	80	

(a) $c_{ij} = u_i + v_j$ for basic x_{ij}

Thus,

$$c_{11} = 2 - 2 = 0$$

$$c_{21} = 3 + 2 = 5$$

$$c_{22} = 3 + 5 = 8$$

$$c_{32} = 5 + 5 = 10$$

$$c_{33} = 5 + 10 = 15$$

$$\text{Cost} = 15 \times 0 + 5 \times 5 + 25 \times 8 + 5 \times 10 \\ + 80 \times 15 = \$1475$$

(b) $u_i + v_j - c_{ij} \leq 0$ for nonbasic x_{ij}

$$-2 + 5 - c_{12} \leq 0 \Rightarrow c_{12} \geq 3$$

$$-2 + 10 - c_{13} \leq 0 \Rightarrow c_{13} \geq 8$$

$$3 + 10 - c_{23} \leq 0 \Rightarrow c_{23} \geq 13$$

$$5 + 2 - c_{31} \leq 0 \Rightarrow c_{31} \geq 7$$

5

Problems 6 and 7 on
next page

continued...

continued...

Set 5.3b

(a) For basic x_{ij} , $c_{ij} = u_i + v_j$.

u	2	2	5	
1	$c_{11}=3$ (10)	$1+2\theta$	$1+3\theta$	10
-1	$2+\theta$	$c_{21}=1$ (20)	$c_{23}=4$ (20)	40

6

$$10 \quad 20 \quad 20 \\ \text{Cost} = 3 \times 10 + 1 \times 20 + 4 \times 20 = \$130$$

(b) For nonbasic x_{ij} : $u_i + v_j - c_{ij} \leq 0$ to satisfy optimality. Hence

$$2+1-(1+2\theta) \leq 0 \Rightarrow \theta \geq 1$$

$$5+1-(1+3\theta) \leq 0 \Rightarrow \theta \geq 5/3$$

$$2-1-(2+\theta) \leq 0 \Rightarrow \theta \geq -1$$

Take $\theta = \frac{5}{3}$ to yield $x_{13} = 0$ as the zero basic variable.

7

$$\begin{array}{ccccccc} & x_{11} & x_{12} & x_{13} & x_{21} & x_{22} & x_{23} \\ \text{Min } Z = & 1 & 1 & 2 & 6 & 5 & 1 \end{array}$$

s.t.

$$\begin{array}{ccccc} 1 & 1 & 1 & & \geq 5 \\ & 1 & 1 & 1 & \geq 6 \\ 1 & & 1 & & \geq 2 \\ 1 & & & 1 & \geq 7 \\ 1 & & & 1 & \geq 1 \end{array}$$

$x_{ij} \geq 0$ for all i and j

Optimum LP solution using TORA:

$$Z = 15, x_{11} = 2, x_{12} = 3, x_{23} = 6$$

If we replace the first two constraints with equations, we get the optimum solution:

$$Z = 27, x_{11} = 2, x_{12} = 3,$$

$$x_{22} = 4, x_{23} = 2$$

The new solution is worse!

Set 5.3c

U_1	U_2	U_3	V_1	V_2	V_3	V_4	
Max 15	25	10	5	15	15	15	

s.t.

$$\begin{array}{l} | \\ | \quad \quad \quad \leq 10 \\ | \quad \quad \quad \leq 2 \\ | \quad \quad \quad \leq 20 \\ | \quad \quad \quad \leq 11 \\ | \quad \quad \quad \leq 12 \\ | \quad \quad \quad \leq 7 \\ | \quad \quad \quad \leq 9 \\ | \quad \quad \quad \leq 20 \\ | \quad \quad \quad \leq 4 \\ | \quad \quad \quad \leq 14 \\ | \quad \quad \quad \leq 16 \\ | \quad \quad \quad \leq 18 \end{array}$$

From Table 5-25:

$$\begin{aligned} U_1 &= 0, U_2 = 5, U_3 = 7 \\ V_1 &= -3, V_2 = 2, V_3 = 4, V_4 = 11 \end{aligned}$$

$$\begin{aligned} \text{Optimum } W &= 15x_0 + 25x_5 + 10x_7 \\ &\quad + 5x_{-3} + 15x_2 + \\ &\quad 15x_4 + 15x_{11} \\ &= \$435 \end{aligned}$$

Minimize $Z = \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij}$

subject to

$$\sum_{j=1}^n x_{ij} = a_i, \quad i=1, 2, \dots, m$$

$$\sum_{i=1}^m x_{ij} = b_j, \quad j=1, 2, \dots, n$$

Next, consider

$$\begin{aligned} Z' &= \sum_{i=1}^m \sum_{j=1}^n (c_{ij} + K) x_{ij} \\ &= \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} + K \sum_{i=1}^m \left(\sum_{j=1}^n x_{ij} \right) \\ &= \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} + K \sum_{i=1}^m a_i \end{aligned}$$

continued...

$$\begin{aligned} &= \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} + K, \quad K \text{ is a constant} \\ &= Z + K \end{aligned}$$

This result shows that optimization using Z and Z' yield the same optimum values of x_{ij} .

To show why the dual values associated with a given primal basic solution are not unique, note that, for any constant K ,

$$(\text{Dual Values}) = \begin{pmatrix} \text{Original basic} \\ \text{obj. coefficients} \\ + \\ K \end{pmatrix} \times \text{Inverse}$$

This means that even though the optimal primal solution is unique for all K , there are infinity of dual values, each corresponding to a given value of K .

The conclusion is that an arbitrary value assigned to one of the dual variables (e.g., $U_1 = 0$) implies a specific value for the constant K .

2

Set 5.4a

(a-i)

3	8	2	10	3	2
8	7	2	9	7	2
6	4	2	7	5	2
8	4	2	3	5	2
9	10	6	9	10	6

Row
min

1

0	7	0	0	5	Optimum:
4	0	4	4	5	1-1
5	1	4	6	0	2-2
0	4	3	0	0	3-5
6	4	0	1	4	4-4

Cost = \$11

1	6	0	8	1
6	5	0	7	5
4	2	0	5	3
6	2	0	1	3
3	4	0	3	4

Col min → 1 2 0 1 1

Assignment:

0	4	2	7	0
3	1	0	4	2
3	0	2	4	2
5	0	2	0	2
0	0	0	0	1

Cost = \$21

(a-ii)

3	9	2	2	7	2
6	1	5	6	6	1
9	4	7	10	3	3
2	5	4	2	1	1
9	6	2	4	6	2

1	7	0	1	5
5	0	4	5	5
6	1	4	7	0
1	4	3	1	0
7	4	0	2	4

Col min 1 0 0 1 0

continued...

5	5	M	2
7	4	2	3
9	3	5	M
7	2	6	7

3	3	M-2	0
5	2	0	1
6	0	2	M-3
5	0	4	5

(All entries are divided by 10 for convenience)

0	3	M-2	0
2	2	0	1
3	0	2	M-3
2	0	4	5

0	5	M-2	0
2	4	0	1
1	0	0	M-5
0	0	4	5

Optimum: 1-4, 2-3, 3-2, 4-1
Cost = \$140

worker	Job				
	1	2	3	4	5
1	50	50	M	20	0
2	70	40	20	30	0
3	90	30	50	M	0
4	70	20	60	70	0
5	60	45	30	80	0

Job 5 is dummy

worker	Job				
	1	2	3	4	5
1	0	30	M-20	0	0
2	20	20	0	10	0
3	40	10	30	M-20	0
4	20	0	40	50	0
5	10	25	10	60	0

worker	Job				
	1	2	3	4	5
1	0	30	M-20	0	10
2	20	20	0	10	10
3	30	0	20	M-30	0
4	20	0	40	50	10
5	0	15	0	50	0

Optimum:

1-4
2-3
3-5
4-2
5-1

Worker 3 is assigned to dummy job 5.
Thus, worker 5 must replace worker 3.

Set 5.4a

Add a "dummy" operator with zero assignment cost to each job (including the fifth). The optimal solution will show the replacement by indicating which of the current jobs (1 thru 4) is assigned to the dummy operator. If the dummy operator is assigned to the new job, then the new job must assume lower priority to the current four jobs. (All assignment cost are divided by 10 for convenience.)

	Job				
Operator	1	2	3	4	5
1	5	5	M	2	2
2	7	4	2	3	1
3	9	3	5	M	2
4	7	2	6	7	8
5	0	0	0	0	0

← Dummy

3	3	M-3	0	0
6	3	1	2	0
7	1	3	M-2	0
5	0	4	5	0
0	0	0	0	0

Optimum:

2	2	M-4	0	0
5	2	0	2	0
6	0	2	M-2	0
5	0	4	6	7
0	0	0	1	0

(S-1)

Since dummy operator is assigned to job 1, new job 5 has higher priority over job 1.

Define the following two sets:

Set 1: (DA,3), (DA,10), (DA,17), (DA,25)

continued

5

Set 2: (AT,7), (AT,12), (AT,21), (AT,28).

The idea is to match one element from set 1 with another element from set 2. The matching automatically decides the date and location for the purchase of each ticket. For example, consider the following assignment:

- (DA,3) - (AT,21)
- (DA,10) - (AT,7)
- (DA,17) - (AT,28)
- (DA,25) - (AT,12)

This assignment can be interpreted as follows:

Ticket 1: June 3 DA → AT
June 21 AT → DA

Ticket 2: June 7 AT → DA
June 10 DA → AT

Ticket 3: June 17 DA → AT
June 28 AT → DA

Ticket 4: June 12 AT → DA
June 25 DA → AT

The complete assignment model is given below

	A,7	A,12	A,21	A,28
D,3	400	300	300	(280)
D,10	(300)	400	300	300
D,17	300	(300)	400	300
D,25	300	300	(300)	400

Optimum:

- (D,3) - (A,28) (A,21) - (D,25)
- (A,7) - (D,10) (A,12) - (D,17)

Problem has alternative optima.

Set 5.4a

Distance matrix in meters:

	candidate areas				
	a	b	c	d	
existing centers	1	50	50	95	45
	2	30	30	55	65
	3	70	50	25	55
	4	100	60	55	25

A measure of the optimal assignment of new centers to candidate locations must reflect both distance and frequency of trips; that is

	existing				candidate			
	1	2	3	4	a	b	c	d
I	10	7	0	11	50	50	95	45
II	2	1	8	4	30	30	55	65
new III	4	9	6	0	70	50	25	55
IV	3	5	2	7	100	60	55	25

	a	b	c	d
I	1810	1370	1940	(1180)
II	1090	770	(665)	695
III	(890)	770	1025	1095
IV	1140	(820)	995	745

TORA optimum assignment:

- I - d
- II - c
- III - a
- IV - b

6

The ranking of the projects by the different teams can use the following numeric score

7

1: Highest preference

10: Lowest preference

A tie in preference between two or more projects is indicated by assigning the projects the same score. For example, the scores

Project 1 2 3 4 5 6 7 8 9 10

Score 9 9 8 7 3 5 4 1 2 6

indicate that project 8 is the most preferred and projects 1 and 2 tie for the least preferred status.

For the development of the model, we use the following numeric designations for the projects

Project nbr.	Project name
1	Boeing-F15
2	Boeing-F18
3	Boeing Simulation
4	Cargil
5	Cobb-Vantress
6	ConAgra
7	Cooper
8	DaySpring(layout)
9	DaySpring(Materials)
10	JB Hunt
11	Raytheon
12	Tyson South
13	Tyson East
14	WAL-MART
15	Yellow

continued...

Set 5.4a

The following is a typical summary of preference scores submitted by the 11 teams:

	1*	2	3	Team	4	5	6*	7*	8*	9*	10	11
1	-	①	2	2	1	-	-	-	-	2	15	
2	-	1	3	①	2	-	-	1	-	10	12	
3	1	2	5	3	2	13	5	1	4	15	①	
4	②	3	6	4	10	5	14	2	1	4	14	
5	3	5	4	5	9	4	12	3	3	13	13	
6	3	4	2	5	9	8	12	①	2	1	13	
7	4	6	①	12	8	9	10	2	5	2	5	
8	5	6	7	14	7	9	10	4	6	3	15	
9	7	8	9	14	7	1	①	15	1	15	1	
10	7	9	12	15	6	3	9	5	4	7	5	
11	-	9	13	6	5	-	-	7	-	6	7	
12	13	10	14	7	4	②	8	9	15	4	9	
13	14	11	1	8	3	13	7	8	①	8	9	
14	15	12	5	9	①	14	7	6	2	9	10	
15	15	13	7	10	2	15	6	1	3	①	11	

* Team does not meet citizenship requirements

② Project requiring US citizenship

The problem is modeled as an assignment model. Entries — are replaced by M, a very large value. The model is unbalanced. Thus, 4 artificial teams must be added to balance the model. In the end four projects will not be assigned.

TOA Solution:

Project	Team	Score
1	2	1
2	4	1
3	11	1

continued...

Project	Team	Score
4	1	1
5	None	—
6	8	1
7	3	1
8	None	—
9	7	1
10	None	—
11	None	—
12	6	2
13	10	1
14	5	1
15	10	1

Total score 13

$$\text{Average score} = \frac{13}{11} = 1.18$$

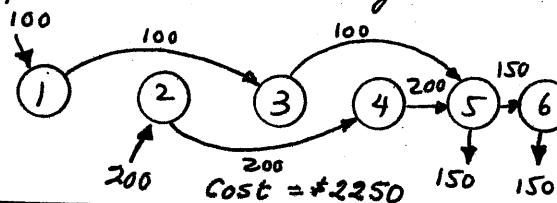
The average score is close to 1, meaning that all preferences are well met.

Set 5.5a

	1	4	M	M	6
1	(100)				100
2	3	2	M	M	
3	(200)	1	(100)	6	M
4	3	0	5	8	
5	M	M	(150)	(150)	1
	B	B	150+B	150	

Let $B = 300$ units

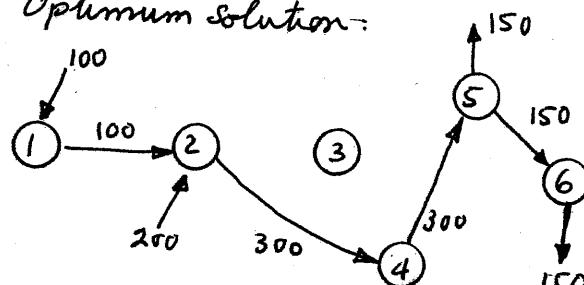
Optimum solution using TORA:



$B = 300$ units

	1	2	3	4	5	6	7
1	(100)						100
2	0	3	(300)	2	M	M	
3	M	0	1	6	M		
4	M	3	0	(300)	5	8	
5	M	M	M	(150)	(150)	1	
	B	B	B	150+B	150		

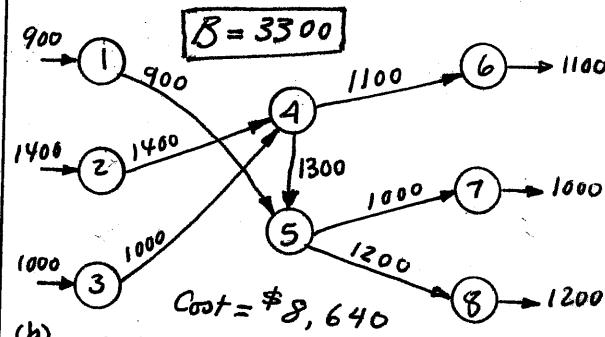
Optimum solution:



Cost = \$2,350

	4	5	6	7	8	9
1		(900)	M	M	M	900
2	8	4.3	M	M	M	1400
3	2	4.6	M	M	M	1000
4	0	1.5	.2	4.5	6	B
5	900	(1300)	(1100)	2.1	1.9	B
	M	0	3	(1000)	(1200)	
	B	B	1100	1000	1200	

$B = 3300$

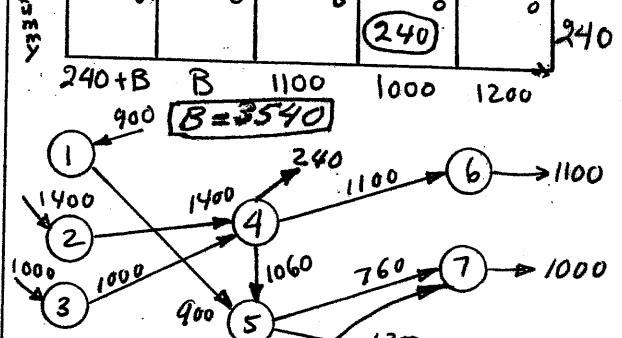


2

	1	2	3	4	5	6	7
1	(100)						100
2	0	3	(300)	2	M	M	
3	M	0	1	6	M		
4	M	3	0	(300)	5	8	
5	M	M	M	(150)	(150)	1	
	B	B	B	150+B	150		

2

	4	5	6	7	8	9
1		(900)	M	M	M	900
2	8	4.3	M	M	M	1400
3	2	4.6	M	M	M	1000
4	0	1.5	.2	4.5	6	B
5	1300	(1060)	(1100)	2.1	1.9	B
	M	0	3	(760)	(1200)	
				(240)	0	
	240+B	B	1100	1000	1200	



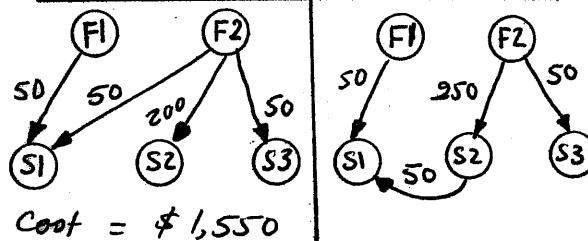
5-17

Set 5.5a

	F1	F2	S1	S2	S3	Dum.
F1	0 (500)	6 (50)	7 (50)	8 (50)	9 (50)	0 (150)
F2	6 (500)	0 (50)	5 (200)	4 (50)	3 (50)	0 (300)
S1	7 (500)	2 (500)	0 (500)	5 (500)	1 (500)	0 (500)
S2	1 (500)	5 (500)	1 (500)	0 (500)	4 (500)	0 (500)
S3	8 (500)	9 (500)	7 (500)	6 (500)	0 (500)	0 (500)
	B B	B B	100 + B	200 + B	50 + B	150

$$B = 500$$

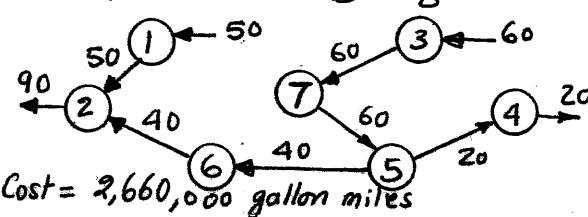
Alternative optima using TORA:
#1 #2



$$\text{Cost} = \$1,550$$

Assume that units of supply and demand are in thousand gallons. $B = 110$

	2	4	5	6	7	
1	20 (50)	m	m	m	3 (50)	50
3	m	30	m	m	9 (60)	60
5	m	2	0	4	10 (40)	B
6	8 (40)	m	4	0 (70)	m	B
7	40	m	10 (60)	m	0 (50)	B
	90	20	B	B	B	



$$\text{Cost} = 2,660,000 \text{ gallon miles}$$

	2	3	4	5	6	7	
1	5 (1)	3 (1)	m	m	m	m	1
2	0 (1)	4	1	7	m	m	0+1
3	6	0	5	1	2 (1)	m	0+1
4	m	m	0 (1)	9	m	4	0+1
5	m	m	2	0	5	8	0+1
6	m	3	m	7	0	3 (1)	0+1
	0+1	0+1	0+1	0+1	0+1	1	

Optimum route using TORA:

$$1 \rightarrow 3 \rightarrow 6 \rightarrow 7$$

$$\text{Distance} = 3 + 2 + 3 = 8$$

minimize Z.

$$x_{13} x_{14} x_{23} x_{24} x_{34} x_{35} x_{36} x_{46} x_{47} x_{56} x_{67}$$

$$z = 3 \quad 4 \quad 2 \quad 5 \quad 7 \quad 8 \quad 6 \quad 4 \quad 9 \quad 5 \quad 3$$

①	1	1									= 1000
②			1	1							= 1200
③	-1	-1	1	1	1						= 0
④	-1	-1	-1	-1		1	1				= 0
⑤					-1			1			= -800
⑥					-1	-1	-1	1			= -900
⑦					-1	-1	-1	-1	-1	-1	= -500

Each node yields a constraint. The special characteristics of the model show that each column has exactly +1 and -1, with the remainder of the elements equal to zero.

Set 5.5a

8

x_{ij} = number of laborers hired at the start of period i
and terminated at the start of period j .

Define nodes 1, 2, 3, 4, and 5 to correspond to the five months of the horizon. Node 6 is added to allow defining the variables x_{i6} that terminate at the end of the five-month planning horizon. The associated LP is defined below.

	x_{12}	x_{13}	x_{14}	x_{15}	x_{16}	x_{23}	x_{24}	x_{25}	x_{26}	x_{34}	x_{35}	x_{36}	x_{45}	x_{46}	x_{56}	
	100	130	180	220	250	100	130	180	220	100	130	180	100	130	100	min
(1)	1	1	1	1	1											≥ 100
(2)		1	1	1	1	1	1	1								≥ 120
(3)			1	1	1		1	1	1	1	1					≥ 80
(4)				1	1			1	1	1	1	1	1			≥ 170
(5)						1			1		1		1	1	1	≥ 50

Let s_1, s_2, s_3, s_4 , and s_5 be the surplus variables associated with constraints 1, 2, 3, 4, and 5, respectively. The LP after adding the surplus variables thus appears as

	x_{12}	x_{13}	x_{14}	x_{15}	x_{16}	x_{23}	x_{24}	x_{25}	x_{26}	x_{34}	x_{35}	x_{36}	x_{45}	x_{46}	x_{56}	s_1	s_2	s_3	s_4	s_5	
	100	130	180	220	250	100	120	180	220	100	130	180	100	130	100		min				
1	1	1	1	1	1											-1					100
	1	1	1	1	1	1	1	1								-1					120
		1	1	1		1	1	1	1	1	1					-1					80
			1	1			1	1	1	1	1	1	1			-1					170
						1			1	1	1	1	1	1		-1					50

Next, perform the following transformation:

1. Leave equation (1) unchanged.
2. Replace equation (2) with $(2) - (1)$.
3. Replace equation (3) with $(3) - (2)$.
4. Replace equation (4) with $(4) - (3)$.
5. Replace equation (5) with $(5) - (4)$.
6. Add a new equation that equals $-(5)$.

These transformations lead to the following LP

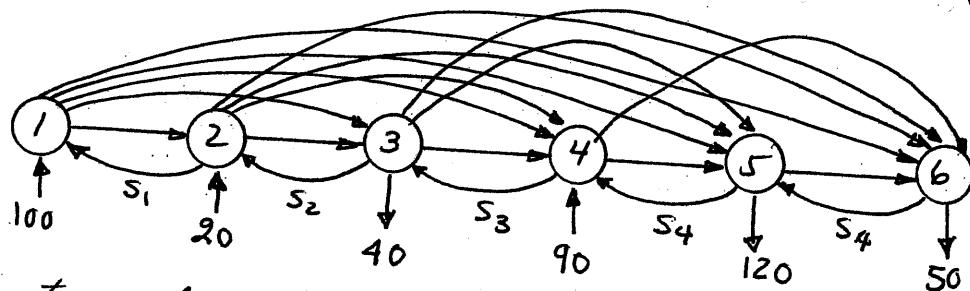
	x_{12}	x_{13}	x_{14}	x_{15}	x_{16}	x_{23}	x_{24}	x_{25}	x_{26}	x_{34}	x_{35}	x_{36}	x_{45}	x_{46}	x_{56}	s_1	s_2	s_3	s_4	s_5	
	100	130	180	220	250	100	130	180	220	100	130	180	100	130	100		min				
1	1	1	1	1	1											-1					100
	-1					1	1	1	1							1	-1				20
		-1				-1				1	1	1				1	-1				-40
			-1				-1		-1	1	1	1				1	-1				90
				-1				-1	-1	-1	-1	-1	1	1		1	-1				-120
					-1				-1	-1	-1	-1	-1	-1		1	-1				-50

continued...

Set 5.5a

The last LP has the structure of a transshipment model (See Problem 7). Let
 $S_1 = x_{21}$ $S_3 = x_{43}$ $S_5 = x_{65}$
 $S_2 = x_{32}$ $S_4 = x_{54}$

Then the LP above can be translated as a network as follows:

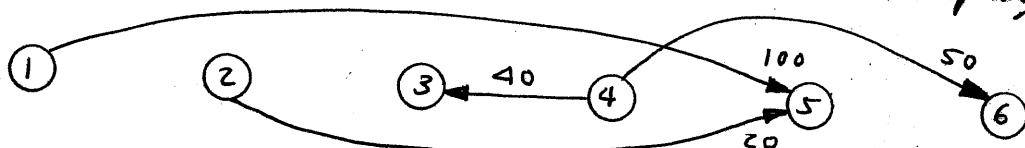


The transshipment model thus appears as

	1	2	3	4	5	6	
1	0	100	130	180	220	250	$100 + B$
2	0	0	100	130	180	220	$20 + B$
3	M	0	0	100	130	180	B
4	M	M	0	0	100	130	$90 + B$
5	M	M	M	0	0	100	B
6	M	M	M	M	0	0	B
	B	B	$40 + B$	B	$120 + B$	$50 + B$	

$$B = 550$$

The optimum solution from TORA is (Problem has alternative optima)



This solution can be interpreted as follows

1. Hire 100 laborers at the start of period 1 and terminate them at the start of period 5.
2. Hire 20 workers at the start of period 2 and terminate them at the start of period 5.
3. Hire 50 workers at the start of period 4 and terminate them at the start of period 6.

The solution satisfies the labor requirements exactly, except for period 3 where there is a surplus of 40 workers ($x_{43} = 40$).

CHAPTER 11

Deterministic Inventory Models

Set 11.3a

$$y^* = \sqrt{\frac{2KD}{h}}, t_0 = \frac{y^*}{D}, TCU(y^*) = \sqrt{2KDh}$$

a) $y^* = \sqrt{\frac{2 \times 100 \times 30}{.05}} = 346.4 \text{ units}$
 $t_0 = \frac{346.4}{30} = 11.55 \text{ days}$
 $TCU(y^*) = \frac{100 \times 30}{346.4} + \frac{.05 \times 346.4}{2} = \17.32

Policy: Order 346.4 units whenever inventory drops to 207.2 units
 Effective lead time = 6.91 days

b) $y^* = \sqrt{\frac{2 \times 50 \times 30}{.05}} \approx 245 \text{ units}$
 $t_0 = \frac{245}{30} = 8.16 \text{ days}$
 $L_e = 5.51 \text{ days}$
 $TCU(y^*) = \frac{50 \times 30}{245} + \frac{.05 \times 245}{2} = \12.25

Policy: Order 245 units whenever inventory drops to 165.15 units

c) $y^* = \sqrt{\frac{2 \times 100 \times 40}{.01}} = 894.4 \text{ units}$
 $t_0 = \frac{894.4}{40} = 22.36 \text{ days}$

$L_e = 7.64 \text{ days}$
 $TCU(y^*) = \frac{100 \times 40}{894.4} + \frac{.01 \times 894.4}{2} = \8.94

Policy: Order 894.4 units whenever inventory drops to 305.57 units.

d) $y^* = \sqrt{\frac{2 \times 100 \times 20}{.04}} = 316.23 \text{ units}$
 $t_0 = \frac{316.23}{20} = 15.81 \text{ days}$
 $L_e = 14.19 \text{ days}$
 $TCU(y^*) = \frac{100 \times 20}{316.23} + \frac{.04 \times 316.23}{2} = \12.65

Policy: Order 316.23 units whenever inventory drops to 283.8 units.

$D = 300 \text{ lb/wk}, K = \$20, h = \$0.03/\text{lb/day}$ 2

(a) $TC/\text{wk} = \frac{KD}{y} + \frac{hy}{2}$
 $= \frac{20 \times 300}{300} + \frac{7 \times .03 \times 300}{2} = \51.50

(b) $y^* = \sqrt{\frac{2 \times 20 \times 300}{.03 \times 7}} = 239 \text{ lb}$

$t_0 = \frac{239}{300/7} = .8 \text{ wk}$
 $TC/\text{wk} = \sqrt{2 \times 20 \times 300 \times .03 \times 7} = \50.20

continued...

$L_e = 0 \text{ days}$

Policy: Order 239 lb whenever inventory drops to zero level.

c) Cost difference = $51.50 - 50.20 = \$1.30$

2) $h = \frac{.35}{7} = \$0.05/\text{unit/day}$

$D = 50 \text{ units/day}, K = \20

$y^* = \sqrt{\frac{2 \times 20 \times 50}{.05}} = 200 \text{ units}$

$t_0 = \frac{200}{50} = 4 \text{ days}$

$L = 7 \text{ days}, L_e = 3 \text{ days}$

$R = 3 \times 50 = 150 \text{ units}$

Policy: Order 200 units whenever inventory drops to 150 units.

b) Optimum number of orders = $\frac{365}{4} \approx 91 \text{ orders}$

(a) Policy 1: $D = \frac{R}{L_e} = \frac{50}{10} = 5 \text{ units/day}$ 4

Cost/day = $\frac{KD}{y} + \frac{hy}{2}$
 $= \frac{20 \times 5}{150} + \frac{.02 \times 150}{2} = \2.17

Policy 2: $D = \frac{75}{15} = 5 \text{ units/day}$

Cost/day = $\frac{20 \times 5}{200} + \frac{.02 \times 200}{2} = \2.50

choose policy 1.

(b) $K = \$20, D = 5 \text{ units/day}$
 $h = \$0.02, L = 22 \text{ days}$

$y^* = \sqrt{\frac{2 \times 20 \times 5}{.02}} = 100 \text{ units}$

$t_0 = \frac{100}{5} = 20 \text{ days}$

$L_e = 22 - 20 = 2 \text{ days}$

Reorder level = $2 \times 5 = 10 \text{ units}$

Order 100 units whenever the level drops to 10 units

Cost/day = $\frac{20 \times 5}{100} + \frac{.02 \times 100}{2} = \2.00

Set 11.3a

$$D = 5 \text{ units/day}$$

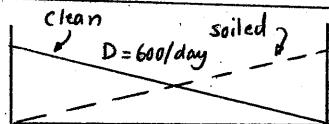
$$h = \$0.10/\text{day}$$

$$K = \$100$$

$$yt = \sqrt{\frac{2x5 \times 100}{h}} = 100 \text{ pallets}$$

$$t_0 = \frac{100}{5} = 20 \text{ days}$$

Pick up 100 pallets every 20 days.



5

$$\begin{aligned} TC/\text{day} &= \frac{K}{y/D} + \frac{h_1 y}{2} + \frac{h_2 y}{2} + .6D \\ &= \frac{KD}{y} + (h_1 + h_2) \frac{y}{2} + .6D \end{aligned}$$

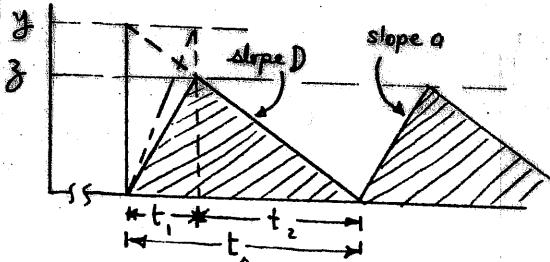
$$y^* = \sqrt{\frac{2KD}{(h_1 + h_2)}} = \sqrt{\frac{2 \times 81 \times 600}{(.01 + .02)}} = 1800 \text{ towels}$$

$$t_0 = 1800/600 = 3 \text{ days}$$

$$\text{Cost/day} = \frac{81 \times 600}{1800} + \frac{.03 \times 1800}{2} = \$54$$

Optimal policy: Pick up soiled towels and deliver an equal batch of 1800 towels every 3 days

6



a) From the geometry of the figure,

$$z = t_1(a-D) = \frac{y}{a}(a-D) = y\left(1 - \frac{D}{a}\right)$$

$$\begin{aligned} b) TCU(y) &= \frac{K + (\frac{3}{2})t_0 * h}{t_0} \\ &= \frac{KD}{y} + \frac{h}{2}\left(1 - \frac{D}{a}\right)y \end{aligned}$$

c) $\frac{\partial TCU(y)}{\partial y} = 0$ gives

$$-\frac{KD}{y^2} + \frac{h}{2}\left(1 - \frac{D}{a}\right) = 0$$

$$y^* = \sqrt{\frac{2KD}{h\left(1 - \frac{D}{a}\right)}}$$

$$(d) \lim_{a \rightarrow \infty} \sqrt{\frac{2KD}{h\left(1 - \frac{D}{a}\right)}} = \sqrt{\frac{2KD}{h}}$$

7

The basic assumption is that the employee will deposit sufficient funds in Europe to take advantage of the higher interest rate and periodically send lump sums to the US to take care of the obligations. This problem in the context of an application of the simple economic lot size formula with no shortages. The idea is that it may be more economical to hold funds longer in European banks to take advantage of their considerably higher interest rate. The cost of wiring funds from overseas ($= \$50$) may be regarded as the "setup" cost and the lost interest per dollar per year ($= .065 - .015 = .05$) can be treated as the "holding" cost. Using this information, the economic lot size formula will yield

$$\text{Deposit amount} = \sqrt{\frac{2KD}{h}} = \sqrt{\frac{2 \times 50 \times 12000}{.05}} = \$4899$$

$$\begin{aligned} \text{Time between deposits} &= t_0 = \frac{4899}{12000} = .408 \text{ year} \\ &= 4.9 \text{ months} \end{aligned}$$

Optimal policy: Send \$4899 ($\approx \5000) every 4.9 (≈ 5) months to the US. The first installment occurs at the start of the year

Alternative 1: Produce

$$\begin{aligned} y^* &= \sqrt{\frac{2KD}{h\left(1 - \frac{D}{a}\right)}} \\ &= \sqrt{\frac{2 \times 20 \times \frac{26000}{365}}{.02\left(1 - \frac{26000/365}{100}\right)}} = 703.7 \text{ units} \end{aligned}$$

Total cost / day

$$\begin{aligned} &= \frac{KD}{y^*} + \frac{h}{2}\left(1 - \frac{D}{a}\right)y^* \\ &= \frac{200 \times \frac{2600}{365}}{703.7} + \frac{.02}{2}\left(1 - \frac{26000}{100 \times 365}\right) \times 703.7 \end{aligned}$$

$$= \$4.05 \text{ per day}$$

9

continued...

Set 11.3a

alternative 2: Buy

$$y^* = \sqrt{\frac{2KD}{h}}$$

$$= \sqrt{\frac{2 \times 15 \times \frac{26000}{365}}{0.02}}$$

$$= 326.87 \text{ units}$$

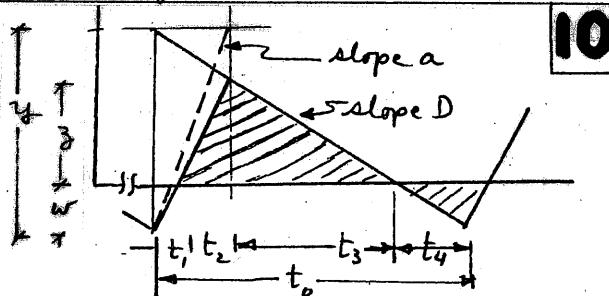
Total cost /day

$$= \frac{KD}{y^*} + \frac{h}{2} y^*$$

$$= \frac{15 \times \frac{26000}{365}}{326.87} + \frac{0.02}{2} \times 326.87$$

$$= \$6.54/\text{day}$$

The company should produce its own.



$$\bar{z} = y(1 - \frac{D}{a}) - w$$

$$TCU(y, w) = \left[K + \frac{h \{ y(1 - \frac{D}{a}) - w \}^2 + pw^2}{2D(1 - D/a)} \right] / t_0$$

$$= \frac{KD}{y} + \frac{h \{ y(1 - \frac{D}{a}) - w \}^2 + pw^2}{2y(1 - D/a)}$$

Partial derivatives = 0 give

$$-\frac{KD}{y^2} + h \left(\frac{1}{2} \left(1 - \frac{D}{a} \right) - \frac{w^2}{2y^2(1 - D/a)} \right) - \frac{pw^2}{2y^2(1 - \frac{D}{a})} = 0$$

$$h \left(\frac{w}{y(1 - \frac{D}{a})} - 1 \right) + \frac{pw}{y(1 - D/a)} = 0$$

This gives,

$$y^* = \sqrt{\frac{2KD(p+h)}{ph(1 - D/a)}}, \quad w^* = \sqrt{\frac{2KDh(1 - \frac{D}{a})}{p(p+h)}}$$

Set 11.3b

EOQ before quantity discount = 1800
towels per Problem 6, Set 11.2a.

$$\begin{aligned} \text{Total cost/day given batches of 1800 towels} \\ = DC_1 + \frac{KD}{y} + \frac{(h_1+h_2)y}{2} \\ = 600 \times 6 + \frac{81 \times 600}{1800} + \frac{.03 \times 1800}{2} = \$414 \end{aligned}$$

$$\begin{aligned} \text{Total cost/day given batches of 2500 towels} \\ = DC_2 + \frac{KD}{y} + \frac{(h_1+h_2)y}{2} \\ = 600 \times 5 + \frac{81 \times 600}{2500} + \frac{.03 \times 2500}{2} = \$356.94 \end{aligned}$$

Take advantage of price discount.

$$y_m = \sqrt{\frac{2KD}{h}} = \sqrt{\frac{2 \times 100 \times 30}{.05}} = 346.41$$

$y = 500$ units

Because $y_m < q$, we need to compute Q .

$$\begin{aligned} TCU_1(y_m) &= DC_1 + \frac{KD}{y_m} + \frac{hy_m}{2} \\ &= 30 \times 10 + \frac{100 \times 30}{346.41} + \frac{.05 \times 346.41}{2} \\ &= 317.32 \end{aligned}$$

The equation for computing Q is

$$Q^2 + \left(\frac{2(8 \times 30 - 317.32)}{.05} \right) Q + \frac{2 \times 100 \times 30}{.05} = 0$$

$$\text{or } Q^2 - 3092.82Q + 120000 = 0$$

This yields $Q = 3053.52$ units

$$\text{Because } y_m < q < Q \Rightarrow y^* = q = 500$$

$$t_o = \frac{500}{30} = 16.67 \text{ days} \Rightarrow L_C = 4.33$$

Order 500 units when inventory drops to 130.

$$y_m = \sqrt{\frac{2KD}{h}} = \sqrt{\frac{2 \times 50 \times 20}{.3}} = 81.65 \text{ units}$$

Because $q > y_m$, we need to compute Q .

$$\begin{aligned} TCU_1(y_m) &= 20 \times 25 + \frac{50 \times 20}{81.65} + \frac{.3 \times 81.65}{2} \\ &= 524.49 \end{aligned}$$

Q -equation:

$$Q^2 + \left(\frac{2(22.5 \times 20 - 524.49)}{.3} \right) Q + \frac{2 \times 50 \times 20}{.3} = 0$$

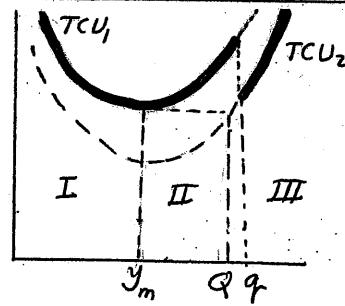
$$Q^2 - 496.63Q + 6666.67 = 0$$

continued...

Thus, $Q = 482.83$

Because $y_m < q < Q \Rightarrow y^* = 150$

Order 150 units when inventory drops to 0



4

From the preceding figure, the discount is not advantageous if

$$TCU_1(y_m) \leq TCU_2(q)$$

or

$$DC_1 + \frac{KD}{y_m} + \frac{hy_m}{2} \leq DC_2 + \frac{KD}{q} + \frac{hq}{2}$$

or

$$\begin{aligned} 20C_1 + \frac{50 \times 20}{81.65} + \frac{.3 \times 81.65}{2} \\ \leq 20C_2 + \frac{50 \times 20}{150} + \frac{.3 \times 150}{2} \end{aligned}$$

Thus, the condition reduces to

$$C_1 - C_2 \leq -23359$$

Let $d = \text{discount factor } (< 1)$.

$$\text{Then } C_2 = (1-d)C_1, \quad 0 < d < 1$$

Given $C_1 = 25$, we have

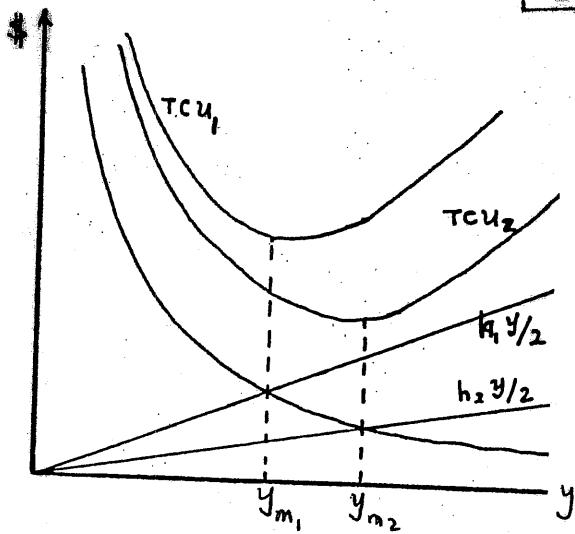
$$25d \leq -233588$$

$$\text{or } d \leq -0.09344$$

Thus, no advantage of the % discount is $\leq .9344\% (\approx 1\%)$

Set 11.3b

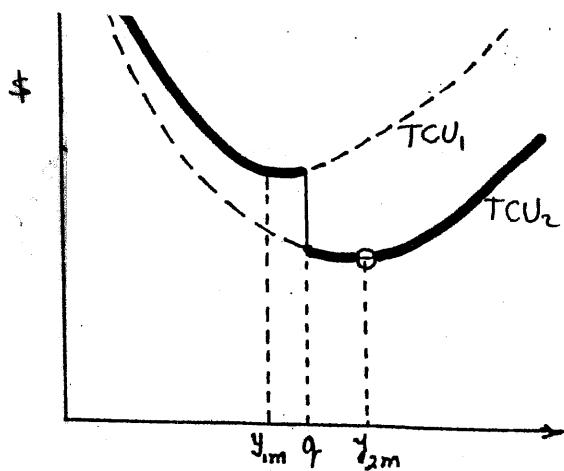
5



$$TCU_1(y) = \frac{KD}{y} + \frac{h_1 y}{2}$$

$$TCU_2(y) = \frac{KD}{y} + \frac{h_2 y}{2}$$

Case 1: $q < y_{2m}$



Solution:

$$y^* = y_{2m}$$

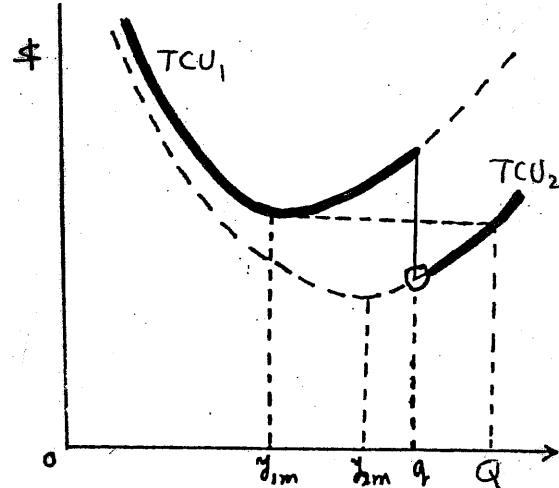
$$TCU(y^*) = TCU_2(y_{2m})$$

continued...

Case 2: $y_{2m} < q \le Q$

The value of Q is determined from the equation:

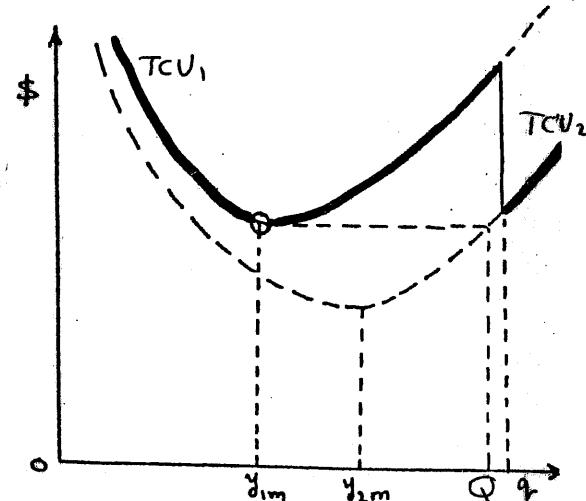
$$TCU_1(y_{1m}) = TCU_2(Q)$$



Solution: $y^* = q$

$$TCU(y^*) = TCU_2(q)$$

Case 3: $y_{2m} < Q < q$



Solution: $y^* = y_{1m}$, $TCU(y^*) = TCU_1(y_{1m})$

$$TCU(y^*) = \begin{cases} TCU_2(y_{2m}), & q < y_{2m} \\ TCU_2(q), & y_{2m} < q \le Q \\ TCU_1(y_{1m}), & y_{2m} < Q < q \end{cases}$$

Set 11.3c

See file ampl11.3c-1.txt.

AMPL model will not converge unless
 $K_i D_i / y_i$ is replaced with $K_i D_i / (y_i + \epsilon)$, where $\epsilon > 0$
and very small.

1

SOLUTION:

Total cost = 568.11

$y_1 = 4.42$

$y_2 = 6.87$

$y_3 = 4.12$

$y_4 = 7.20$

$y_5 = 5.80$

See file ampl11.3c-2.txt.

New constraint:

$$(1/2)(y_1 + y_2 + y_3) \leq 25$$

2

SOLUTION:

Total cost = 10.42

$y_1 = 10.83$

$y_2 = 16.85$

$y_3 = 22.32$

See file ampl11.3c-3.txt.

New constraint:

Average inventory for item i = $y_i / 2$.

$$(1/2)(100y_1 + 55y_2 + 100y_3) \leq 1000$$

3

SOLUTION:

Total cost = 14.31

$y_1 = 5.58$

$y_2 = 7.90$

$y_3 = 10.07$

See file ampl11.3c-4.txt.

AMPL model will not converge unless

$K_i D_i / y_i$ is replaced with $K_i D_i / (y_i + \epsilon)$, where $\epsilon > 0$
and very small.

4

New constraint:

$$365(10/y_1 + 20/y_2 + 5/y_3 + 10/y_4) \leq 150$$

SOLUTION:

Total cost = 54.71

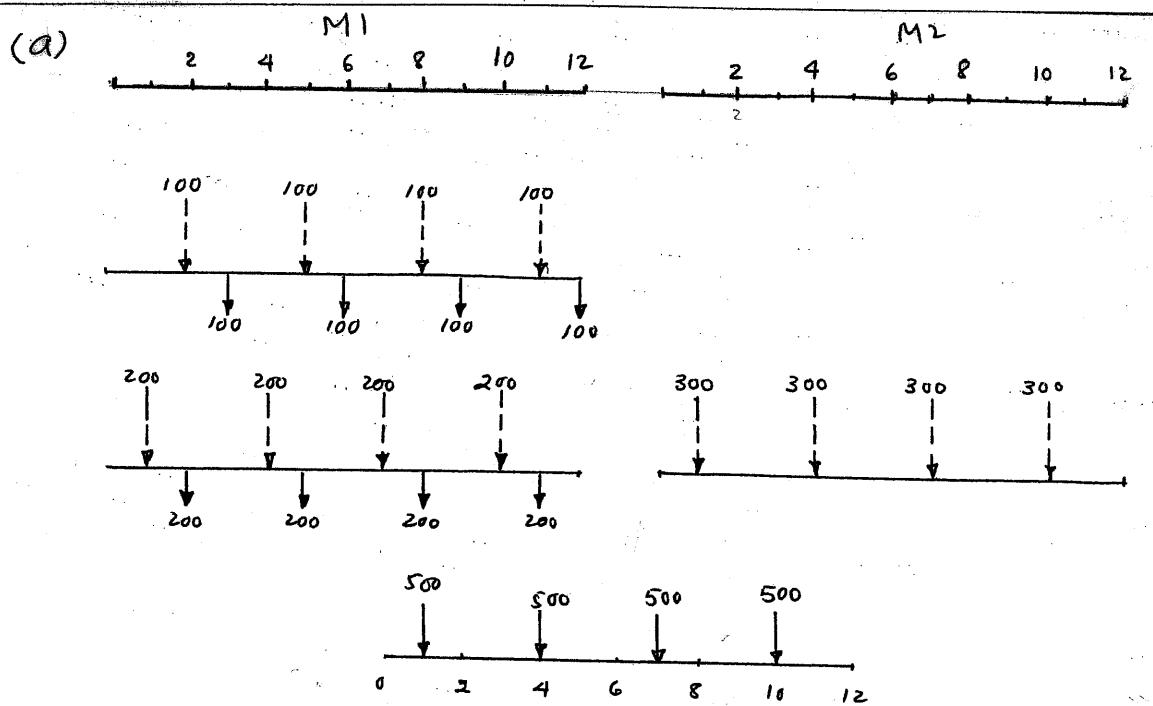
$y_1 = 155.30$

$y_2 = 118.81$

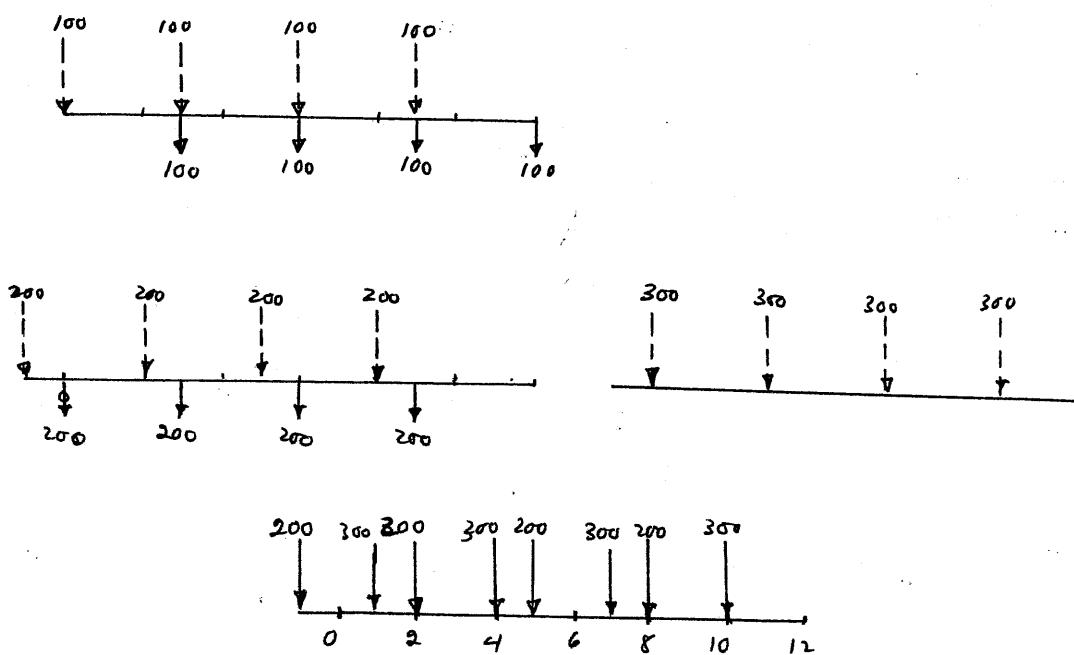
$y_3 = 74.36$

$y_4 = 90.09$

Set 11.4a



(b)



	1	2	3	4	Surplus		
R_1	90	5	5.1	5.25	5.37	0	
O_1	10	75	30	7.6	7.75	7.87	0
R_2			3	3.15	3.27	0	
O_2		100	4.5	4.65	4.77	0	
R_3			4	4.12	4.12	0	
O_3			6	6.12	6.12	0	
R_4			110	1.5	1.5	0	
O_4			50	20	20		
	100	190	210	160	20		

	1	2	3	4	Surplus		
R_1	100	4	4.5	5.	5.5	6.	0
O_1	50	6	6.5	7	7.5	8.	0
S	3	20	7	7.5	8	8.5	9
R_2		40	4	4.5	5.	6	0
O_2		60	6	6.5	7.	7.5	0
S		80	7	7.5	8	8.5	0
R_3			90	4	4.5	5	0
O_3			60	20	6.5	7	0
S			70	7	7.5	8	0
R_4				60	4	4.5	0
O_4				50	6	6.5	0
				50	7	7.5	20
					20	0	100

(a)

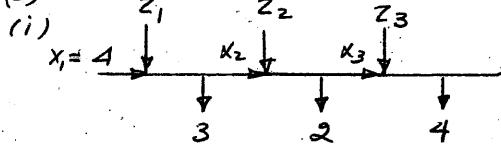
	1	2	3	4	Surplus		
I	11	1	1.3	1.65	1.85	0	
II	2	1	2.3	3	2.65	2.85	0
III	5	5.3	5.65	5.85	5.85	0	
I		3	2	2.35	2.55	0	
II		11	4.35	1	4.55	0	
III		6	6.35	5	6.55	10(5)	0
I			3	2	2.2	0	
II			5	8	5.2	0	
III			7	7.2	4	0	
IV			10	10.2	10	0	
I				3	3	0	
II				8	4	0	
III				4	5	0	
IV				7	10	0	
	11	4	17	29	39		

(b) Additional 10 units are produced as shown by the circled entries in period 4. The problem has alternative solutions.

Set 11.4c

(a) No, because inventory should not be held needlessly at the end of planning horizon

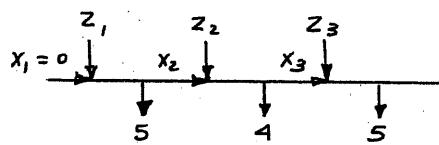
(b)



$$0 \leq z_1 \leq 5, 1 \leq z_2 \leq 5, 0 \leq z_3 \leq 4$$

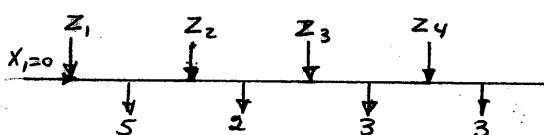
$$x_1 = 4, 1 \leq x_2 \leq 6, 0 \leq x_3 \leq 4$$

(ii)



$$5 \leq z_1 \leq 14, 0 \leq z_2 \leq 9, 0 \leq z_3 \leq 5$$

$$x_1 = 0, 0 \leq x_2 \leq 9, 0 \leq x_3 \leq 5$$



Stage 1: $f_1(x_2) = \min\{K_1 + C_1(z_1) + h_1 x_2\}$

$$z_1 = D_1 + x_2$$

where $C_i(z_i) = \begin{cases} 1z_i, & 0 \leq z_i \leq 6 \\ 2z_i, & z_i \geq 7 \end{cases}, i=1,2,\dots,4$

x_1	z_1	$K_1 = 5, h_1 = 1$	f_1	$Opt. Sol.$
0	5	10	10	5
1	12	12	12	6
2	15	15	15	7
3	18	18	18	8
4	21	21	21	9
5	24	24	24	10
6	27	27	27	11
7	30	30	30	12
8	33	33	33	13

Stage 2:

$$f_2(x_3) = \min_{0 \leq z_2 \leq D_2 + x_3} \{K_2 + C_2(z_2) + h_2 x_3 + f_1(x_3 + D_2 - z_2)\}$$

$$0 \leq z_2 \leq 8, 0 \leq x_3 \leq 6, D_2 = 2$$

$$K_2 = 7, h_2 = 1$$

x_3	z_2	f_2	$Opt. Sol.$
0	15	15	0
1	19	19	0
2	23	23	0, 4
3	27	27	5
4	31	31	6
5	35	35	6
6	39	39	7, 8

Stage 3: $0 \leq z_3 \leq 6, 0 \leq x_4 \leq 3, D_3 = 3$

x_4	z_3	f_3	$Opt. Sol.$
0	25	25	0
1	28	28	0
2	32	32	5
3	36	36	6

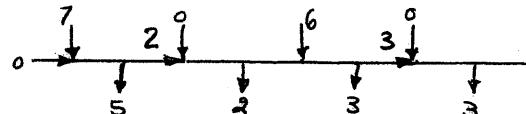
Stage 4: $0 \leq z_4 \leq 3, x_5 = 0, D_4 = 3$

x_5	z_4	f_4	$Opt. Sol.$
0	33	33	6

Solution:

$$(x_5 = 0) \rightarrow z_4 = 0 \rightarrow (x_4 = 3) \rightarrow z_3 = 6 \rightarrow (x_3 = 0) \rightarrow$$

$$z_2 = 0 \rightarrow (x_2 = 2) \rightarrow z_1 = 7$$



$$\text{Total cost} = \$33$$

continued...

Set 11.4c

$$f_i(x_2) = \min_{0 \leq z_i \leq D_i + x_2} \left\{ C_i(z_i) + K_i + h_i \left(\frac{x_i + z_i + x_2}{2} \right) \right\}$$

$$= \min_{0 \leq z_i \leq D_i + x_2} \left\{ K_i + C_i(z_i) + h_i \left(x_2 + \frac{D_i}{2} \right) \right\}$$

$$f_i(x_{i+1}) = \min_{0 \leq z_i \leq D_i + x_{i+1}} \left\{ K_i + C_i(z_i) + h_i \left(x_{i+1} + \frac{D_i}{2} \right) \right.$$

$$\left. + f_{i+1} \left(x_{i+1} + D_{i+1} - z_i \right) \right\}$$

3

Stage 1: $D_1 = 3$

								Opt. Sol.	
x_1	$z_1 = 2$	3	4	5	6	7	8	f_1	z_1
1	99	100	111	115	129	193	151	99	2

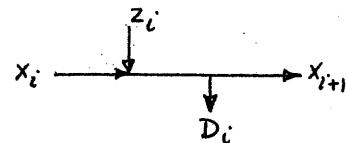
Solution:

$$(x_1 = 1) \rightarrow z_1 = 2 \rightarrow (x_2 = 0) \rightarrow z_2 = 3 \rightarrow$$

$$(x_3 = 1) \rightarrow z_3 = 3$$

Cost = \$99

5



$$f_n(x_n) = \min_{z_n + x_n = D_n} \{ K_n + C_n(z_n) \}$$

4

$$f_i(x_i) = \min_{D_i \leq x_i + z_i \leq D_i + \dots + D_n} \{ K_i + C_i(z_i) + h_i(x_i + z_i - D_i) \}$$

$$+ f_{i+1}(x_{i+1} + z_{i+1} - D_{i+1}) \}$$

Stage 3: $D_3 = 4, 0 \leq x_3 \leq 4$

x_3						Opt. Sol.	
	$z_3 = 0$	1	2	3	4	f_3	z_3
0				56	56	4	
1			36		36	3	
2		26			26	2	
3			16		16	1	
4	0				0	0	

Stage 2: $D_2 = 2$

x_2							Opt. Sol.		
	$z_2 = 0$	1	2	3	4	5	6	f_2	z_2
0		83	76	89	102	109	76	3	
1	73	66	69	82	89		66	2	
2	56	56	59	62	69		56	0, 1	
3	39	49	52	49			34	0	
4	32	42	39				32	0	
5	25	29					25	0	
6	12						12	0	

continued...

11-11

$$\begin{aligned} \text{Average inventory} &= \frac{x_i + z_i + x_{i+1}}{2} \\ &= \frac{x_i + z_i + x_i + z_i - D_i}{2} \\ &= x_i + z_i - \frac{D_i}{2} \end{aligned}$$

Replace $h_i(x_i + z_i - D_i)$ with $h_i(x_i + z_i - \frac{D_i}{2})$ in the backward formulation of problem 4.

Set 11.4d

Period 1:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model											
Number of periods, N		4		Current period		1		Optimum Solution Summary			
Period		1	2	3	4			x	f	z	
1	0	2	2	2							
2	0	114	114	105	70						
3	1	1	1	1							
4	0	22	90	67							
Optimum solution computed											
Period 1		1	0	22	112	179					
Period 2		2	0	142	322	466					
x ₁ = 0		0	114	114	114	114	114	0	112	434	112
x ₂ = 22		22	112	112	112	112	112	0	179	635	179
x ₃ = 112		112	112	112	112	112	112	112	112	112	112
x ₄ = 179		179	179	179	179	179	179	179	179	179	179

Stage 1: $D_1 = 150, X_1 = 50$

Opt. Sol.													
X_2	25	100	200	220	260	330	420	550	730	870	920	f_1	Z_1
0	700		1400									1400	1400
100		1400		1540								1400	1400
120			1540		1820							1540	120
160				1820		2310						1820	160
230					2310		2940					2310	230
320						2940		3850				2940	920
450							3850		5110			3850	550
630								5110		6090		5110	730
770									6090		6490	670	770
920										6490	6700	920	

2

Period 2:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model											
Number of periods, N		4		Current period		2		Optimum Solution Summary			
Period		1	2	3	4			x	f	z	
1	0	2	2	2							
2	0	114	114	105	70						
3	1	1	1	1							
4	0	22	90	67							
Optimum solution computed											
Period 1		1	0	22	112	179					
Period 2		2	0	142	322	466					
x ₁ = 0		0	114	114	114	114	114	0	112	434	112
x ₂ = 22		22	112	112	112	112	112	0	179	635	179
x ₃ = 112		112	112	112	112	112	112	112	112	112	112
x ₄ = 179		179	179	179	179	179	179	179	179	179	179

Period 3:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model											
Number of periods, N		4		Current period		3		Optimum Solution Summary			
Period		1	2	3	4			x	f	z	
1	0	2	2	2							
2	0	114	114	105	70						
3	1	1	1	1							
4	0	22	90	67							
Optimum solution computed											
Period 1		1	0	22	112	179					
Period 2		2	0	142	322	466					
x ₁ = 0		0	114	114	114	114	114	0	112	434	112
x ₂ = 22		22	112	112	112	112	112	0	179	635	179
x ₃ = 112		112	112	112	112	112	112	112	112	112	112
x ₄ = 179		179	179	179	179	179	179	179	179	179	179

Stage 2: $D_2 = 100$

Opt. Sol.												
X_3	25	100	120	160	220	320	450	630	770	870	f_2	Z_2
0	1400	1400									1400	1400
20	1560		1540								1540	120
60	1880			1820							1820	160
120	2440				2310						2310	230
220	3160					2940					2940	320
350	4200						3850				3850	450
530	5690							5110			5110	630
670	6760								6090		6090	770
770	7160									6490	6700	770

Stage 3: $D_3 = 20$

Opt. Sol.											
X_4	25	40	60	130	220	350	530	670	720	f_3	Z_3
0	1540	1580								1540	0
90	1900		1820							1820	60
120	2530			2240						2240	130
200	3340				2780					2780	220
380	4510					3560				3560	350
510	6130						4640			4640	530
650	7390							5480		5480	670
780	8240								5780	5780	720

Stage 4: $D_4 = 40$

Opt. Sol.										
X_5	25	40	110	200	330	510	650	750	f_4	Z_4
0	1820	1900							1820	0
70	2310		2250						2250	110
160	2940			2700					2700	200
220	3850				3350				3350	330
370	5110					4250			4250	510
610	6090						4950		4950	650
660	6440							5200	5200	700
780	6380								6380	0

continued...

Set 11.4d

3

Stage 6: $D_6 = 90$

							Opt. Sol.	
x_7	26	90	220	400	540	590	f_6	Z_6
0	2880	3110					2880	0
130	4180		4600				4180	0
310	5980			6580			6980	0
450	7380				8120		7380	0
500	7880					8670	7880	0

Stage 7: $D_7 = 130$

							Opt. Sol.	
x_8	27	130	310	450	500	f_7	Z_7	
0	4180	3700		4600			3700	130
180	6160						4600	310
320	7700				5300		5300	450
370	8250					5550	5580	500

Stage 8: $D_8 = 180$

							Opt. Sol.
x_9	28	180	320	370	f_8	Z_8	
0	4600	4720				4600	0
140	5860		5840			5840	220
190	6310			6240		6240	370

Stage 9: $D_9 = 140$

							Opt. Sol.
x_{10}	29=0	140	190	f_9	Z_9		
0	5840	5180			5180	140	
50	6340		5380		5380	190	

Stage 10: $D_{10} = 50$

							Opt. Sol.
x_{11}	$z_{10} = 0$	50		f_{10}	Z_{10}		
0	5380		5780		5380	0	

Solution:

Period	Order Amount
1	100
2	120
3	0
4	200
5	0
6	0
7	310
8	0
9	190
10	0

Minimum cost = \$5380

Period 1:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model						
Number of periods, N		5	Current period	1		
Period	1	2	3	4	5	
$x_{1 \rightarrow 5}$	10	10	10	10	10	
$x_{1 \rightarrow 4}$	80	70	60	60	60	
$x_{1 \rightarrow 3}$	1	1	1	1	1	
$x_{1 \rightarrow 2}$	50	70	100	30	60	
$x_{1 \rightarrow 1}$	0	0	0	0	0	
$C1(z_1)$	50	120	220	250	310	
$x_{2 \rightarrow 5}$	0	50	1111111111111111111111			
$x_{2 \rightarrow 4}$	70	1111111111111111111111				
$x_{2 \rightarrow 3}$	170	1111111111111111111111				
$x_{2 \rightarrow 2}$	200	1111111111111111111111				
$x_{2 \rightarrow 1}$	260	1111111111111111111111				
$C2(z_2)$	0	170	240	270	310	
$x_{3 \rightarrow 5}$	0	1350	1111111111111111111111			
$x_{3 \rightarrow 4}$	250	1111111111111111111111				
$x_{3 \rightarrow 3}$	2910	1111111111111111111111				
$x_{3 \rightarrow 2}$	2780	1111111111111111111111				
$x_{3 \rightarrow 1}$	3440	1111111111111111111111				
$C3(z_3)$	0	1350	200	270	310	
$x_{4 \rightarrow 5}$	0	1350	1111111111111111111111			
$x_{4 \rightarrow 4}$	30	2910	1111111111111111111111			
$x_{4 \rightarrow 3}$	2450	1111111111111111111111				
$x_{4 \rightarrow 2}$	2780	1111111111111111111111				
$x_{4 \rightarrow 1}$	3440	1111111111111111111111				
$C4(z_4)$	0	1350	120	30	2740	100
$x_{5 \rightarrow 5}$	0	1060	1350	1950		
$x_{5 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{5 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{5 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{5 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C5(z_5)$	0	1350	120	30	2740	100
$x_{6 \rightarrow 5}$	0	1060	1350	1950		
$x_{6 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{6 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{6 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{6 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C6(z_6)$	0	1350	120	30	2740	100
$x_{7 \rightarrow 5}$	0	1060	1350	1950		
$x_{7 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{7 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{7 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{7 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C7(z_7)$	0	1350	120	30	2740	100
$x_{8 \rightarrow 5}$	0	1060	1350	1950		
$x_{8 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{8 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{8 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{8 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C8(z_8)$	0	1350	120	30	2740	100
$x_{9 \rightarrow 5}$	0	1060	1350	1950		
$x_{9 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{9 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{9 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{9 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C9(z_9)$	0	1350	120	30	2740	100
$x_{10 \rightarrow 5}$	0	1060	1350	1950		
$x_{10 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{10 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{10 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{10 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C10(z_{10})$	0	1350	120	30	2740	100

Period 2:

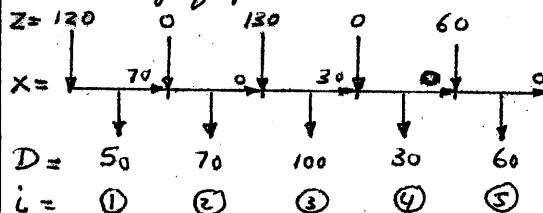
Wagner-Whitin (Forward) Dynamic Programming Inventory Model						
Number of periods, N		5	Current period	2		
Period	1	2	3	4	5	
$x_{1 \rightarrow 5}$	10	10	10	10	10	
$x_{1 \rightarrow 4}$	80	70	60	60	60	
$x_{1 \rightarrow 3}$	1	1	1	1	1	
$x_{1 \rightarrow 2}$	50	70	100	30	60	
$x_{1 \rightarrow 1}$	0	0	0	0	0	
$C1(z_1)$	50	120	220	250	310	
$x_{2 \rightarrow 5}$	0	50	1111111111111111111111			
$x_{2 \rightarrow 4}$	70	1111111111111111111111				
$x_{2 \rightarrow 3}$	170	1111111111111111111111				
$x_{2 \rightarrow 2}$	200	1111111111111111111111				
$x_{2 \rightarrow 1}$	260	1111111111111111111111				
$C2(z_2)$	0	170	240	270	310	
$x_{3 \rightarrow 5}$	0	1350	1111111111111111111111			
$x_{3 \rightarrow 4}$	250	1111111111111111111111				
$x_{3 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{3 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{3 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C3(z_3)$	0	1350	120	30	2740	100
$x_{4 \rightarrow 5}$	0	1060	1350	1950		
$x_{4 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{4 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{4 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{4 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C4(z_4)$	0	1350	120	30	2740	100
$x_{5 \rightarrow 5}$	0	1060	1350	1950		
$x_{5 \rightarrow 4}$	2450	2410	1111111111111111111111			
$x_{5 \rightarrow 3}$	2910	2740	1111111111111111111111			
$x_{5 \rightarrow 2}$	3440	3460	1111111111111111111111			
$x_{5 \rightarrow 1}$	3440	3460	1111111111111111111111			
$C5(z_5)$	0	1350	120	30	2740	100
$x_{6 \rightarrow $						

Set 11.4d

Period 5:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model						
Number of periods, N	6	Current period, t	5	6	7	8
Period	1	2	3	4	5	6
$c_{1,1}$	10	10	10	10	10	
$c_{1,2}$	10	10	10	10	10	
$c_{1,3}$	10	70	60	60	60	
$c_{1,4}$	1	1	1	1	1	
$c_{1,5}$	50	70	100	30	60	
$c_{1,6}$						
$x_{1,1}$	0	50	50	0	240	100
$x_{1,2}$	0	70	240	120	30	2740
$x_{1,3}$	5	25	170	240	20	3400
$x_{1,4}$	3400	60	200	2700	260	period 4
$x_{1,5}$	3400	340	310	0	2740	0
$x_{1,6}$						
$f_{1,1}$	0	1350	0	period 5		
$f_{1,2}$	0	100	2450	170	0	3400
$f_{1,3}$	130	2700	200			
$f_{1,4}$	190	3400	260			
$f_{1,5}$						
$f_{1,6}$						
$z_{1,1}$	120	0	130	0	60	
$z_{1,2}$	0	70	0	30	0	
$z_{1,3}$	70	0	0	30	0	
$z_{1,4}$	0	0	30	0	0	
$z_{1,5}$	0	0	0	0	0	
$z_{1,6}$						
i_1	①	②	③	④	⑤	

Summary of optimum solution:



Cost = \$3400

Period 3:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model						
Number of periods, N	6	Current period, t	3	4	5	6
Period	1	2	3	4	5	6
$c_{1,1}$	10	10	10	10	10	
$c_{1,2}$	10	10	10	10	10	
$c_{1,3}$	10	70	60	60	60	
$c_{1,4}$	1	1	1	1	1	
$c_{1,5}$	50	70	100	30	60	
$c_{1,6}$						
$x_{1,1}$	0	50	50	0	240	100
$x_{1,2}$	0	70	240	120	30	2740
$x_{1,3}$	5	25	170	240	20	3400
$x_{1,4}$	3400	60	200	2700	260	period 4
$x_{1,5}$	3400	340	310	0	2740	0
$x_{1,6}$						
$f_{1,1}$	0	1350	0	period 5		
$f_{1,2}$	0	100	2450	170	0	3400
$f_{1,3}$	130	2700	200			
$f_{1,4}$	190	3400	260			
$f_{1,5}$						
$f_{1,6}$						
$z_{1,1}$	10	15	7	20	13	25
$z_{1,2}$	0	15	7	20	13	25
$z_{1,3}$	15	7	20	13	25	
$z_{1,4}$	0	0	0	0	0	
$z_{1,5}$	0	0	0	0	0	
$z_{1,6}$						
i_1	①	②	③	④	⑤	

Period 4:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model						
Number of periods, N	6	Current period, t	4	5	6	7
Period	1	2	3	4	5	6
$c_{1,1}$	10	10	10	10	10	
$c_{1,2}$	10	10	10	10	10	
$c_{1,3}$	10	70	60	60	60	
$c_{1,4}$	1	1	1	1	1	
$c_{1,5}$	50	70	100	30	60	
$c_{1,6}$						
$x_{1,1}$	0	50	50	0	240	100
$x_{1,2}$	0	70	240	120	30	2740
$x_{1,3}$	5	25	170	240	20	3400
$x_{1,4}$	3400	60	200	2700	260	period 4
$x_{1,5}$	3400	340	310	0	2740	0
$x_{1,6}$						
$f_{1,1}$	0	1350	0	period 5		
$f_{1,2}$	0	100	2450	170	0	3400
$f_{1,3}$	130	2700	200			
$f_{1,4}$	190	3400	260			
$f_{1,5}$						
$f_{1,6}$						
$z_{1,1}$	10	15	7	20	13	25
$z_{1,2}$	0	15	7	20	13	25
$z_{1,3}$	15	7	20	13	25	
$z_{1,4}$	0	0	0	0	0	
$z_{1,5}$	0	0	0	0	0	
$z_{1,6}$						
i_1	①	②	③	④	⑤	

Period 5:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model						
Number of periods, N	6	Current period, t	5	6	7	8
Period	1	2	3	4	5	6
$c_{1,1}$	10	10	10	10	10	
$c_{1,2}$	10	10	10	10	10	
$c_{1,3}$	10	70	60	60	60	
$c_{1,4}$	1	1	1	1	1	
$c_{1,5}$	50	70	100	30	60	
$c_{1,6}$						
$x_{1,1}$	0	50	50	0	240	100
$x_{1,2}$	0	70	240	120	30	2740
$x_{1,3}$	5	25	170	240	20	3400
$x_{1,4}$	3400	60	200	2700	260	period 4
$x_{1,5}$	3400	340	310	0	2740	0
$x_{1,6}$						
$f_{1,1}$	0	1350	0	period 5		
$f_{1,2}$	0	100	2450	170	0	3400
$f_{1,3}$	130	2700	200			
$f_{1,4}$	190	3400	260			
$f_{1,5}$						
$f_{1,6}$						
$z_{1,1}$	10	15	7	20	13	25
$z_{1,2}$	0	15	7	20	13	25
$z_{1,3}$	15	7	20	13	25	
$z_{1,4}$	0	0	0	0	0	
$z_{1,5}$	0	0	0	0	0	
$z_{1,6}$						
i_1	①	②	③	④	⑤	

Period 6:

Wagner-Whitin (Forward) Dynamic Programming Inventory Model						
Number of periods, N	6	Current period, t	6	7	8	9
Period	1	2	3	4	5	6
$c_{1,1}$	10	10	10	10	10	
$c_{1,2}$	10	10	10	10	10	
$c_{1,3}$	10	70	60	60	60	
$c_{1,4}$	1	1	1	1	1	
$c_{1,5}$	50	70	100	30	60	
$c_{1,6}$						
$x_{1,1}$	0	50	50	0	240	100
$x_{1,2}$	0	70	240	120	30	2740
$x_{1,3}$	5	25	170	240	20	3400
$x_{1,4}$	3400	60	200	2700	260	period 4
$x_{1,5}$	3400	340	310	0	2740	0
$x_{1,6}$						
$f_{1,1}$	0	1350	0	period 5		
$f_{1,2}$	0	100	2450	170	0	3400
$f_{1,3}$	130	2700	200			
$f_{1,4}$	190	3400	260			
$f_{1,5}$						
$f_{1,6}$						
$z_{1,1}$	10	15	7	20	13	25
$z_{1,2}$	0	15	7	20	13	25
$z_{1,3}$	15	7	20	13	25	
$z_{1,4}$	0	0	0	0	0	
$z_{1,5}$	0	0	0	0	0	
$z_{1,6}$						
i_1	①	②	③	④	⑤	

continued...

11-14

Set 11.4e

$i = 1, K_1 = \$250$:

Period, t	D_t	$TC(1, t)$	$TCU(1, t)$
1	60	250	$250/1 = 250$
2	70	$250 + 1 \times 70 = 320$	$320/2 = 160^*$
3	80	$320 + 2 \times 80 = 480$	$480/3 = 160^*$
4	90	$480 + 3 \times 90 = 750$	$750/4 = 187.50$

Produce $60 + 70 + 80 = 210$ for 1, 2, and 3

$i = 4, K_4 = \$300$:

Period, t	D_t	$TC(4, t)$	$TCU(4, t)$
4	90	300	$300/1 = 300$
5	85	$300 + 1 \times 85 = 385$	$385/2 = 192.5$
6	80	$385 + 2 \times 80 = 545$	$545/3 = 181.67$
7	75	$545 + 3 \times 75 = 770$	$770/4 = 192.5$

Produce $90 + 85 + 80 = 255$ for 4, 5, and 6

$i = 7, K_7 = \$250$:

Period, t	D_t	$TC(7, t)$	$TCU(7, t)$
7	75	250	$250/1 = 250$
8	70	$250 + 1 \times 70 = 320$	$320/2 = 160$
9	65	$320 + 2 \times 65 = 450$	$450/3 = 150$
10	60	$450 + 3 \times 60 = 630$	$630/4 = 157.50$

Produce $75 + 70 + 65 = 210$ for 7, 8, and 9

$i = 10, K_{10} = \$250$:

Period, t	D_t	$TC(10, t)$	$TCU(10, t)$
10	60	250	$250/1 = 250$
11	55	$250 + 1 \times 55 = 305$	$305/2 = 152.50$
12	50	$305 + 2 \times 50 = 405$	$405/3 = 135$

Produce $60 + 55 + 50 = 165$ for 10, 11, and 12

$i = 1, K = 200$:

t	D_t	$TC(1, t)$	$TCU(1, t)$
1	100	200	$200/1 = 200$
2	120	$200 + 1 \times 120 = 344$	$344/2 = 172$
3	50	$344 + 2 \times 50 = 464$	$464/3 = 154.67$
4	70	$464 + 3 \times 70 = 716$	$716/4 = 179$

$i = 4, K = \$200$:

t	D_t	$TC(4, t)$	$TCU(4, t)$
4	70	200	$200/1 = 200$
5	90	$200 + 1 \times 90 = 308$	$308/2 = 154$
6	105	$308 + 2 \times 105 = 560$	$560/3 = 186.67$

$i = 6, K = \$200$:

t	D_t	$TC(6, t)$	$TCU(6, t)$
6	105	200	$200/1 = 200$
7	115	$200 + 1.2 \times 115 = 338$	$338/2 = 169$
8	95	$338 + 2 \times 1.2 \times 95 = 566$	$566/3 = 188.67$

$i = 8, K = \$200$:

t	D_t	$TC(8, t)$	$TCU(8, t)$
8	95	200	$200/1 = 200$
9	80	$200 + 1.2 \times 80 = 296$	$296/2 = 148$
10	85	$296 + 2 \times 1.2 \times 85 = 500$	$500/3 = 166.67$

$i = 10, K = \$200$:

t	D_t	$TC(10, t)$	$TCU(10, t)$
10	85	200	$200/1 = 200$
11	100	$200 + 1.2 \times 100 = 320$	$320/2 = 160$
12	110	$320 + 2 \times 1.2 \times 110 = 584$	$584/3 = 194.67$

Schedule:

Produce	for periods
270	1, 2, and 3
160	4, and 5
220	6 and 7
175	8 and 9
185	10 and 11
110	12

2

Continued...

Chapter 15

Queuing Systems

Set 15.1a

(a) Efficiency = $100 - 29 = 71\%$

(b) For average waiting time ≤ 3 minutes, at least 5 cashiers are needed

For efficiency $\geq 90\%$, the associated idleness percentage is $\leq 10\%$. The corresponding number of cashiers is at most 2.

Conclusion:

The two conditions cannot be satisfied simultaneously.
At least one of the two conditions must be relaxed.

1

$C_A = \$18 \text{ per hour}$

$C_B = \$25 \text{ per hour}$

Length of queue A = 4 jobs

Length of queue B = $.7 \times 4 = 2.8$ jobs

Cost of A = $\$18 + 4 \times \$10 = \$58 \text{ per hour}$

Cost of B = $\$25 + 2.8 \times \$10 = \$53 \text{ per hour}$

2

Decision:

Select Model B.

Set 15.2a

3

Situation	Customer	Server
a	Plane	Runway
b	Passenger	Taxi
c	Machinist	Clerk at tool crib
d	Letter	Clerk
e	Student	Registrar's office
f	Cases	Judge
g	Shopper	Cashier
h	Car	Parking space

2

Situation	Calling Source	Customers arrival
a	∞	Individual
b	∞	Individual
c	∞	Individual
d	∞	Bulk
e	∞	Individual
f	∞	Individual
g	∞	Individual
h	∞	Individual

Situation	Interarrival time	Service time
a	Probabilistic	Time to clear runway
b	Probabilistic	Ride time
c	Probabilistic	Time to receive tool
d	Deterministic	Time to process letter
e	Probabilistic	Time to process registr.
f	Probabilistic	Trial time
g	Probabilistic	Check-out time
h	Probabilistic	Parking time.

Situation	Queue Capacity	Queue Discipline
a	∞	FIFO
b	∞	FIFO
c	∞	FIFO
d	∞	Random
e	∞	FIFO
f	∞	FIFO
g	∞	FIFO
h	0	None

#	Queueing situation	Customer
1	Arrival of orders	Orders
2	Processing (single machine)	Rush orders
3	Processing (single machine)	Regular jobs
4	Processing (Prod. line)	Rush jobs
5	Processing (Prod. line)	Regular jobs
6	Receipt of completed jobs	Completed orders
7	Tool crib	Tools
8	Machine breakdown	machines

#	Servers	Discipline	Service time	Queue length	Source
1	Foreman	Priority	Sorting time	∞	∞
2	Machine	FIFO	Prod. time	∞	∞
3	machine	FIFO	Prod. time	∞	∞
4	Prod. line	FIFO	Prod. time	∞	∞
5	Prod. line	FIFO	Prod. time	∞	∞
6	Shipping facilities	FIFO	Loading time	finite	∞
7	Tool crib	Priority	Exchange time	finite	finite
8	Repair persons	Priority	Repair time	finite	finite

(a) T. (b) T. (c) T.

4

- (a) None.
- (b) None.
- (c) None.
- (d) None.
- (e) Jockey or balk
- (f) None
- (g) Jockey
- (h) None

5

Set 15.3a

(a) Av. interarrival time (in time units)

$$= \frac{1}{\text{arrival rate } \lambda} \quad |$$

(b) Let \bar{I} = av. interarrival time

(i) $\lambda = \frac{60}{10} = 6 \text{ arrivals/hr}$
 $\bar{I} = 10 \text{ minutes} = \frac{1}{6} \text{ hour}$

(ii) $\lambda = \frac{60}{3} = 20 \text{ arrivals/hr}$
 $\bar{I} = \frac{6}{2} = 3 \text{ minutes} = \frac{1}{20} \text{ hr}$

(iii) $\lambda = \frac{10}{30} \times 60 = 20 \text{ arrivals/hr}$
 $\bar{I} = \frac{30}{10} = 3 \text{ minutes} = \frac{1}{20} \text{ hour}$

(iv) $\lambda = 1/5 = 2 \text{ arrivals/hour}$
 $\bar{I} = .5 \text{ hour}$

(c) Let \bar{S} = av. service time

(i) $\mu = \frac{60}{12} = 5 \text{ services/hour}$
 $\bar{S} = 12 \text{ minutes} = .2 \text{ hour}$

(ii) $\mu = \frac{60}{7.5} = 8 \text{ services/hr}$
 $\bar{S} = 7.5 \text{ min} = .125 \text{ hr}$

(iii) $\mu = \frac{5}{30} \times 60 = 10 \text{ services/hr}$
 $\bar{S} = \frac{30}{5} = 6 \text{ min} = \frac{1}{10} \text{ hr}$

(iv) $\mu = \frac{1}{.3} = 3.33 \text{ services/hr}$
 $\bar{S} = .3 \text{ hour}$

(a) $\lambda_{\text{hour}} = .2 \text{ failures/hr}$

$\lambda_{\text{week}} = .2 \times 24 \times 7 = 33.6 \text{ failures/week}$

(b) $P\{\text{at least one failure in 2 hours}\}$

$= P\{\text{time betw. failures} \leq 2\}$

$= P\{t \leq 2\} = 1 - e^{-2 \times 2} \approx .33$

(c) $P\{t > 3 \text{ hrs}\} = 1 - P\{t \leq 3\} = e^{-3 \times 2} \approx .55$

(d) $P\{t \leq 1 \text{ hour}\} = 1 - e^{-2 \times 1} = .18$

$\lambda = \frac{1}{.05} = 20 \text{ arrivals/hr}$

(a) $f(t) = \lambda e^{-\lambda t}$
 $= 20 e^{-20t}, t > 0$

(b) $P\{t > \frac{15}{60}\} = P\{t > .25\}$
 $= e^{-20 \times .25}$
 $= .00674$

(c) $P\{t \leq \frac{3}{60}\} = P\{t \leq .05\}$
 $= 1 - e^{-20 \times .05} = .632$

$P\{t > \frac{5}{60}\} = e^{-20 \times \frac{5}{60}} = .189$

(d) $t = 45 - 10 = 35 \text{ minutes}$

Av. # of arrivals in 35 min.
 $= 20 \times \frac{35}{60} = 11.67 \text{ arrivals}$

$\lambda = \frac{1}{6} \text{ arrivals/hr}$

$P\{t \geq 1\} = e^{-\frac{1}{6} \times 1} = .846$

$P\{t \leq .5\} = 1 - e^{-\frac{1}{6} \times .5}$
 $= 1 - e^{-1/12} = .08$

(a) $\lambda = \frac{60}{10} = 6 \text{ arrivals/hr}$

(b) $P\{t \geq \frac{15}{60}\} = e^{-6 \times \frac{15}{60}} = .223$

(c) $P\{t \leq \frac{20}{60}\} = 1 - e^{-6 \times \frac{20}{60}} = .865$

(a) $P\{t \leq \frac{2}{60}\} = 1 - e^{-35(2/60)} = .6886$

(b) $P\{\frac{2}{60} \leq t \leq \frac{3}{60}\}$
 $= P\{t \leq \frac{3}{60}\} - P\{t \leq \frac{2}{60}\}$
 $= (1 - e^{-35 \times 3/60}) - (1 - e^{-35 \times 2/60})$
 $= e^{-70/60} - e^{-105/60} = .1376$

(c) $P\{t > \frac{3}{60}\} = e^{-35(3/60)} = .1738$

Set 15.3a

$$\lambda = \frac{60}{1.5} = 40 \text{ arrivals/hr}$$

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Jim's Payoff	-2¢	+2¢	
Prob.	$P\{t \geq 1\}$	$P\{t \leq 1\}$	

$$P\{t \geq 1\} = e^{-40(1/60)} = .5134$$

$$P\{t \leq 1\} = 1 - .5134 = .4866$$

$$\begin{aligned} \text{Jim's exp. payoff / arriving customer} \\ = -2 \times .5134 + 2 \times .4866 \end{aligned}$$

$$= -.0536 \text{ cent}$$

Jim's exp. payoff / 8 hours

$$= -.0536(8\lambda)$$

$$= -.0536 \times 8 \times 40$$

$$\approx -17.15 \text{ cent}$$

Conclusion: Jim will pay Ann an average of 17 cents every 8 hrs

Jim's payoff	2	0	-2
Probability	$P\{t \leq 1\}$	$P\{1 \leq t \leq 1.5\}$	$P\{t \geq 1.5\}$
$P\{t \leq 1\} = .4866$			
$P\{t \geq 1.5\} = e^{-40(1.5/60)}$			
$= .3679$			
	2	0	-2
	.4866	.1455	.3679

Jim's expected payoff / 8 hours

$$= [2 \times .4866 + 0 \times .1455 - 2 \times .3679] \times 40 \times 8$$

$$\approx 76 \text{ cents}$$

2¢	3¢	-5¢	-6¢
$t \leq 1$	$1 \leq t \leq 1.5$	$1.5 \leq t \leq 2$	$t \geq 2$

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$$\lambda = 40 \text{ arrivals/hr}$$

$$-40/60$$

$$P\{t \leq 1\} = 1 - e^{-40(1/60)} = .4866$$

$$P\{1 \leq t \leq 1.5\} = e^{-40(1.5/60)} - e^{-40(1/60)}$$

$$= .1455$$

$$P\{1.5 \leq t \leq 2\} = e^{-40(2/60)} - e^{-40(1.5/60)}$$

$$= .1043$$

$$P\{t \geq 2\} = e^{-40(2/60)} = .2636$$

Jim's exp. payoff / 8 hours

$$= 8 \times 40 (2 \times .4866 + 3 \times .1455 - 5 \times .1043 - 6 \times .2636)$$

$$\approx -2.22 \text{ cents}$$

Jim pays Ann an average of \$2.22 / 8 hours.

$$(a) \lambda = \frac{60}{6} = 10 \text{ customers/hr}$$

$$P\{t \leq 4 \text{ min}\} = 1 - e^{-10(4/60)} = .4866$$

(b)

$$\% \text{ discount} = \begin{cases} 10\%, & \text{if } t \leq 4 \\ 6\%, & \text{if } 4 < t \leq 5 \\ 2\%, & \text{if } t > 5 \end{cases}$$

$$P\{t \leq 4\} = .4866$$

$$P\{4 < t \leq 5\} = e^{-10(4/60)} - e^{-10(5/60)}$$

$$= .0788$$

$$P\{t > 5\} = e^{-10(5/60)} =$$

$$= .4346$$

Expected % discount

$$= 10 \times .4866 + 6 \times .0788 + 2 \times .4346$$

$$= 6.208\%$$

Set 15.3a

$$\lambda = \frac{365 \times 24}{9000} = .973 \text{ failure / yr}$$

$$P\{t \leq 1\} = 1 - e^{-0.973 \times 1} = .622$$

Lack-of-memory property applies.

(a) The waiting time for the green bus is exponential with mean 10 minutes:

$$f(t) = .1 e^{-0.1t}, t \geq 0$$

(b) The waiting time for the red bus is exponential with mean 7 minutes:

$$f(t) = \frac{1}{7} e^{-t/7}, t \geq 0$$

$$E\{t\} = \int_0^\infty t \lambda e^{-\lambda t} dt$$

$$= - \int_0^\infty t d e^{-\lambda t}$$

$$= - \left(t e^{-\lambda t} - \int_0^\infty e^{-\lambda t} dt \right)$$

$$= - \left(t e^{-\lambda t} - \frac{1}{\lambda} e^{-\lambda t} \right) \Big|_0^\infty$$

$$= \frac{1}{\lambda}$$

$$E\{t^2\} = \lambda \int_0^\infty t^2 e^{-\lambda t} dt$$

$$= - \int_0^\infty t^2 d e^{-\lambda t}$$

$$= - \left[t^2 e^{-\lambda t} - \int_0^\infty 2t e^{-\lambda t} dt \right]$$

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$$= - \left[t^2 e^{-\lambda t} - \frac{2}{\lambda} \int_0^\infty t \lambda e^{-\lambda t} dt \right] \Big|_0^\infty$$

$$= + \frac{2}{\lambda^2}$$

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$$\text{Var}\{t\} = E\{t^2\} - E\{t\}^2$$

$$= \frac{2}{\lambda^2} - \left(\frac{1}{\lambda}\right)^2$$

$$= \frac{1}{\lambda^2}$$

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continued...

Set 15.4a

TORA input = $(5, 0, 0, \infty, \infty)$

$$\begin{aligned} P_{n \geq 5}(t=1 \text{ hr}) &= 1 - [P_0(1) + \dots + P_4(1)] \\ &= 1 - e^{-5} (1 + 5 + \frac{5^2}{2!} + \frac{5^3}{3!} + \frac{5^4}{4!}) \\ &= 1 - .44049 = .55951 \end{aligned}$$

$$\lambda = 1 \text{ trip/month}$$

(a) $\lambda t = 3$: TORA input = $(3, 0, 0, \infty, \infty)$

$$P_0(3) = \frac{(1 \times 3)^0 e^{-1 \times 3}}{0!} = .049787$$

(b) $\lambda t = 12$: TORA input = $(12, 0, 0, \infty, \infty)$

$$\begin{aligned} P_{n \leq 8}(t=12) &= P_0(12) + \dots + P_8(12) \\ &= \frac{12^0 e^{-12}}{0!} + \frac{12^1 e^{-12}}{1!} + \dots + \frac{12^8 e^{-12}}{8!} \\ &= .15503 \end{aligned}$$

$$(c) P_0(t) = \frac{1^0 e^{-1}}{0!} = e^{-1} = .3679$$

TORA input = $(1, 0, 0, \infty, \infty)$

$$\lambda = 2 \text{ arrivals/minute}$$

(a) $\lambda t = 2 \times 5 = 10 \text{ arrivals}$

$$(b) \lambda t = 2 \times 5 = 1$$

TORA input = $(1, 0, 0, \infty, \infty)$

$$P_0(t=.5) = e^{-2 \times .5} = .3679$$

$$(c) 1 - P_0(t=.5) = 1 - .3679 = .6321$$

$$(d) \lambda t = 2 \times 3 = 6 \text{ arrivals}$$

TORA input = $(6, 0, 0, \infty, \infty)$

$$P_0(t=3) = \frac{(2 \times 3)^0 e^{-2 \times 3}}{0!} = .00248$$

$$\lambda = 1/5 = .2 \text{ arrival/min}$$

$$(a) P_2(t=7) = \frac{(.2 \times 7)^2 e^{-2 \times 7}}{2!} = .24167$$

TORA input = $(1.4, 0, 0, \infty, \infty)$

$$(b) P_1(t=5) = \frac{(.2 \times 5)^1 e^{-2 \times 5}}{1!} = .36788$$

$$\lambda = 25 \text{ books per day}$$

$$(a) \lambda t = 25 \times 30 = 750 \text{ books} = 7.5 \text{ shelves}$$

$$(b) 10 \text{ bookcases} = 10 \times 5 \times 100 = 5000 \text{ books}$$

$$\begin{aligned} P_{n \geq 5000}(t=30) &= 1 - [P_0(30) + \dots + P_{5000}(30)] \\ &\approx 0 \end{aligned}$$

$$(a) \lambda_{\text{green}} = .1 \text{ stop/min}, \lambda_{\text{red}} = 1/7 \text{ stop/min}$$

$$\lambda_{\text{combined}} = .1 + \frac{1}{7} = .24286 \text{ stop/min}$$

$$P_2(S) = \frac{(.24286 \times 5)^2}{2!} e^{-24286 \times 5} = .219$$

The two buses could be 2G, 2R or 1G and 1R.

$$(b) P\{t \leq 2\} = 1 - e^{-243 \times 2} = .3849$$

$$E\{n|t\} = \sum_{n=1}^{\infty} n \frac{(\lambda t)^n e^{-\lambda t}}{n!}$$

$$= \lambda t + e^{-\lambda t} \sum_{n=1}^{\infty} \frac{(\lambda t)^{n-1}}{(n-1)!}$$

$$= \lambda t + e^{-\lambda t} e^{\lambda t} = \lambda t$$

$$E\{n^2|t\} = \sum_{n=0}^{\infty} n^2 \frac{(\lambda t)^n e^{-\lambda t}}{n!}$$

$$= \sum_{n=1}^{\infty} n^2 \frac{(\lambda t)^n e^{-\lambda t}}{n!}$$

$$= \lambda t + e^{-\lambda t} \sum_{n=1}^{\infty} \frac{n(\lambda t)^{n-1}}{(n-1)!}$$

$$= \lambda t e^{-\lambda t} \frac{\partial}{\partial \lambda t} \left(\lambda t \sum_{n=1}^{\infty} \frac{(\lambda t)^{n-1}}{(n-1)!} \right)$$

$$= \lambda t e^{-\lambda t} \frac{d}{d \lambda t} (\lambda t + e^{\lambda t})$$

$$= \lambda t e^{-\lambda t} (\lambda t e^{\lambda t} + e^{\lambda t})$$

$$= (\lambda t)^2 + \lambda t$$

Thus,

$$\text{var}\{n|t\} = (\lambda t)^2 + \lambda t - (\lambda t)^2$$

$$= \lambda t$$

Set 15.4a

$$p'_0(t) = -\lambda p_0(t) \quad (1)$$

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$$p'_n(t) = -\lambda p_n(t) + \lambda p_{n-1}(t) \quad (2)$$

From (1)

$$dp_0(t) = -\lambda p_0(t) dt$$

which yields

$$p_0(t) = A e^{-\lambda t}$$

$$\text{Because } p_0(0) = 1 \Rightarrow A = 1, p_0(t) = e^{-\lambda t}$$

For $n=1$:

$$\begin{aligned} p'_1(t) &= -\lambda p_1(t) + \lambda p_0(t) \\ &= -\lambda p_1(t) + \lambda e^{-\lambda t} \end{aligned}$$

or

$$p'_1(t) + \lambda p_1(t) = \lambda e^{-\lambda t}$$

This yields the solution:

$$\begin{aligned} p_1(t) &= e^{\int \lambda dt} \left\{ \int \lambda e^{-\lambda t} e^{-\int \lambda dt} dt + C \right\} \\ &= \lambda t e^{-\lambda t} + C \end{aligned}$$

$$\text{Because } p_1(0) = 0, C = 0, \text{ and}$$

$$p_1(t) = \frac{\lambda t e^{-\lambda t}}{1!}$$

Induction proof:

Given

$$p_i(t) = \frac{(\lambda t)^i e^{-\lambda t}}{i!}$$

then

$$p'_{i+1}(t) + \lambda p_{i+1}(t) = \lambda \frac{(\lambda t)^{i+1} e^{-\lambda t}}{i+1!}$$

The solution is

$$\begin{aligned} p_{i+1}(t) &= e^{\int \lambda dt} \left\{ \frac{\int \lambda (\lambda t)^i e^{-\lambda t} e^{\int \lambda dt} dt}{i!} + C \right\} \\ &= \frac{e^{-\lambda t} (\lambda t)^{i+1}}{(i+1)!} + C \end{aligned}$$

Because $p_{i+1}(0) = 0, C = 0$, and

$$p_{i+1}(t) = \frac{e^{-\lambda t} (\lambda t)^{i+1}}{(i+1)!}$$

continued...

Set 15.4b

$M = 3 \text{ dozens/day}$, $N = 18$
 TORA input data = $(0, Mt, 1, 18, 18)$

$$(a) \mu = 3 \times 3 = 9$$

$$P_0(t=3) = .00532 \text{ (from TORA)}$$

$$(b) Mt = 3 \times 2 = 6$$

$$\sum_{n=0}^{18} n P_n(2) = 11.955$$

(c) This part can be solved using Poisson or exponential distributions:

$$\text{Poisson: } Mt = 3 \times 1 = 3$$

$$\text{Probability} = P_0(1) + P_1(1) + \dots + P_{17}(1) \\ = .9502 \text{ (from TORA)}$$

Exponential: mean = $1/3$ day

$$P\{\text{purchasing at least one dozen in 1 day}\} \\ = P\{\text{time between purchases} \leq 1\} \\ = 1 - e^{-3 \times 1} = .9502$$

$$(d) \text{Exponential: } P\{t \leq .5\} = 1 - e^{-3 \times .5} = .7769$$

$$\text{Poisson: } P_0(.5) + P_1(.5) + \dots + P_{17}(.5) = .7769$$

$$(e) P_0(1) = 0 \quad (Mt = 3 \times 1 = 3)$$

$$N = 40, M = 10 \text{ calls/hr}$$

TORA input $(0, Mt, 1, 40, 40)$

$$(a) P_{n>0}(t=4) = 1 - P_0(4) \\ = 1 - .521 = .479$$

$$(b) E\{n|t=4\} = \sum_{n=0}^{40} n P_n(4) \approx 2.5 \text{ blocks} \\ \approx 25 \text{ tickets}$$

$$N = 48, M = \frac{4 \times 10}{8} = 5 \text{ calls/hr}$$

$$Mt = 5 \times 4 = 20 \text{ calls}$$

$$P_0(4) \approx .000005 \text{ (from TORA)}$$

$$N = 48, Mt = 5 \times 8 = 40, P_0(8) = .11958$$

$$M = 1/1 = 1 \text{ withdrawal/week}$$

$$N = 5, Mt = 4$$

$$P_0(4) = .37116$$

2

$N = 80 \text{ items}, M = 5 \text{ items/day}$

6

$$(a) Mt = 5 \times 2 = 10 \text{ items}$$

$$P_0(2) = .1251$$

$$(b) Mt = 5 \times 4 = 20 \text{ items}$$

$$P_0(4) = .00001$$

$$(c) Mt = 5 \times 4 = 20 \text{ items}$$

$$E\{n|4 \text{ days}\} = \sum_{n=0}^{80} n P_n(4) \approx 60 \text{ items}$$

$$\text{Av. # of withdrawals} = 80 - 60 \\ = 20 \text{ items}$$

$$M = 1/1 = 1 \text{ breakdown/day}$$

$$N = 10, Mt = 1 \times 2 = 2$$

7

$$\text{From TORA, } P_0(2) = .00005$$

$$(a) N = 25, M = 3/day$$

$$t = 6 \text{ days}, Mt = 18$$

$$\text{Av. stock remaining after 6 days} \\ = E\{n|t=6\} = 7.11$$

$$\text{Av. order size} = 25 - 7.11$$

$$\approx 18 \text{ items}$$

$$(b) t = 4, Mt = 3 \times 4 = 12$$

$$P_0(4) = .00069$$

$$(c) t = 6, Mt = 3 \times 6 = 18$$

$$P_{n \leq 14}(6) = P_0(6) + \dots + P_{14}(6) = .9696$$

3

4

$$P\{\text{time betw. departures} > T\}$$

$$= P\{\text{no departures during } T\}$$

$$= P\{N \text{ left after time } T\}$$

$$= P_N(T)$$

$$P\{t > T\} = P_N(T) = \frac{(MT)^0 e^{-MT}}{0!} \\ = e^{-MT}$$

9

Set 15.4b

10

$$p'_N(t) = -\mu p_N(t) \quad (1)$$

$$p'_n(t) = -\mu p_n(t) + \mu p_{n+1}(t), \quad 0 \leq n < N \quad (2)$$

From (1), we get

$$p_N(t) = C e^{-\mu t}$$

Given $p_N(0) = 1$, then $C = 1$ and

$$p_N(t) = e^{-\mu t}$$

Next, consider (2) for $n = N-1$

$$\begin{aligned} p'_{N-1}(t) &= -\mu p_{N-1}(t) + \mu p_N(t) \\ &= -\mu p_{N-1}(t) + \mu e^{-\mu t} \end{aligned}$$

$$\begin{aligned} \text{Thus, } p_{N-1}(t) &= e^{-\int \mu dt} \left\{ \int \mu e^{-\mu t} e^{\int \mu dt} dt + C \right\} \\ &= e^{\mu t} \mu t + C \end{aligned}$$

Because $p_{N-1}(0) = 0$, $C = 0$ and $p_{N-1}(t) = (\mu t) e^{-\mu t}$

Induction proof:

$$\text{Given } p_{n+1}(t) = \frac{(\mu t)^{N-n-1} e^{-\mu t}}{(N-n-1)!}, \text{ then}$$

$$p'_n(t) = -\mu p_n(t) + \frac{\mu (\mu t)^{N-n-1} e^{-\mu t}}{(N-n-1)!}$$

Solution gives

$$p_n(t) = \frac{(\mu t)^{N-n} e^{-\mu t}}{(N-n)!}$$

Set 15.5a

$$(a) P\{0 \text{ counter open}\} = P_0 = \frac{1}{55}$$

$$P\{1 \text{ counter open}\} = P_1 + P_2 + P_3 \\ = \frac{1}{55} (2+8+8) = \frac{14}{55}$$

$$P\{2 \text{ counters open}\} = P_4 + P_5 + P_6 \\ = \frac{1}{55} (8+8+8) = \frac{24}{55}$$

$$P\{3 \text{ counters open}\} = P_7 + P_8 + \dots \\ = 1 - (P_0 + \dots + P_6) \\ = 1 - (\frac{1}{55} + \frac{14}{55} + \frac{24}{55}) = \frac{16}{55}$$

$$(b) \text{Av. # busy counters} \\ = 0 \times \frac{1}{55} + 1 \times \frac{14}{55} + 2 \times \frac{24}{55} + 3 \times \frac{16}{55} \\ = 2 \text{ counters}$$

$$(c) \text{Av. # idle counters} = 3 - 2 = 1$$

$$\lambda = 1/5 = .2 \text{ arrival/min} \\ = 12 \text{ arrivals/hr}$$

$$(a) M_n = \begin{cases} 5 \text{ customers/hr}, & n=0,1,2 \\ 10 \text{ customers/hr}, & n=3,4 \\ 15 \text{ customers/hr}, & n=5,6 \\ 20 \text{ customers/hr}, & n \geq 7 \end{cases}$$

$$P_0 = \frac{12}{5} P_0 = 2.4 P_0$$

$$P_1 = \left(\frac{12}{5}\right)^2 P_0 = 5.76 P_0$$

$$P_2 = \left(\frac{12}{5}\right)^3 \left(\frac{12}{10}\right) P_0 = 6.912 P_0$$

$$P_3 = \left(\frac{12}{5}\right)^4 \left(\frac{12}{10}\right)^2 P_0 = 8.2944 P_0$$

$$P_4 = \left(\frac{12}{5}\right)^5 \left(\frac{12}{10}\right)^3 \left(\frac{12}{15}\right) P_0 = 6.63552 P_0$$

$$P_5 = \left(\frac{12}{5}\right)^6 \left(\frac{12}{10}\right)^4 \left(\frac{12}{15}\right)^2 P_0 = 5.308416 P_0$$

$$P_{n \geq 7} = \left(\frac{12}{5}\right)^7 \left(\frac{12}{10}\right)^5 \left(\frac{12}{15}\right)^3 \left(\frac{12}{20}\right)^2 P_0 = 5.308416(6)^{-6} P_0$$

$$\text{From } \sum_{n=0}^{\infty} P_n = 1, \text{ we get } P_0 = .002587$$

$$P_1 = .05421, P_2 = .13010, P_3 = .15612$$

$$P_4 = .18735, P_5 = .14988, P_6 = .1199$$

$$P_{n \geq 7} = .1199 (6)^{-6}$$

$$(b) P_{n \geq 7} = 1 - (P_0 + P_1 + \dots + P_6) = .8$$

Continued...

$$(c) P\{0 \text{ counter}\} = P_0 = .002587$$

$$P\{1 \text{ counter}\} = P_1 + P_2 = .18431$$

$$P\{2 \text{ counters}\} = P_3 + P_4 = .34347$$

$$P\{3 \text{ counters}\} = P_5 + P_6 = .26978$$

$$P\{4 \text{ counters}\} = P_7 + P_8 + \dots = .199853$$

Av. # idle counters

$$= 4 - (1 \times .18431 + 2 \times .34347 + 3 \times .26978 \\ + 4 \times .199853) \approx 1.52$$

$$M_n = \begin{cases} 5^n, & n=1,2 \\ 15, & n=3,4, \dots \end{cases}$$

$$P_1 = \left(\frac{10}{5}\right) P_0 = 2 P_0$$

$$P_2 = \left(\frac{10}{5}\right) \left(\frac{10}{10}\right) P_0 = 2 P_0$$

$$P_{n \geq 3} = \left(\frac{10}{5}\right) \left(\frac{10}{10}\right) \left(\frac{10}{15}\right)^{n-2} P_0 = 2 \left(\frac{2}{3}\right)^{n-2} P_0$$

Thus,

$$P_0 + 2P_0 + 2P_0 + \left[2 \left(\frac{2}{3}\right) + 2 \left(\frac{2}{3}\right)^2 + \dots\right] P_0 = 1$$

which gives $P_0 = .1111$

(a) Prob that 3 counters are in use

$$= P_{n \geq 3} = 1 - (P_0 + P_1 + P_2) \\ = 1 - (.1111 + .2222 + .2222) \\ = .4445$$

$$(b) P_{n \leq 2} = P_0 + P_1 + P_2 = .5555$$

$$\lambda_n = \begin{cases} 12 \text{ cars/hr}, & n=0,1,\dots,10 \\ 0 & n \geq 11 \end{cases}$$

$$M_n = 60/6 = 10 \text{ cars/hr}$$

$$P_n = \left(\frac{12}{10}\right)^n P_0, \quad n=1,2,\dots,10 \\ = 0, \quad n \geq 11$$

$$P_0 \left(1 + 1.2 + 1.2^2 + \dots + 1.2^{10}\right) = P_0 \frac{1 - 1.2^{11}}{1 - 1.2}$$

$$\text{Thus, } P_0 = .0311$$

Continued...

Set 15.5a

$$(a) P_{10} = \left(\frac{12}{10}\right)^{10} P_0 = .19259$$

$$(b) P_{n \geq 1} = 1 - P_0 = 1 - .0311 = .9689$$

(c) Av. length of the line

$$= 0P_0 + 1P_1 + \dots + 10P_{10}$$

$$= 1x.03732 + 2x.04479 + 3x.05375 + 4x.0645 + 5x.0774 + 6x.09288 + 7x.11145 + 8x.13374 + 9x.16049 + 10x.19259 = 6.71071$$

$$\lambda_n = 6 \text{ arrivals/hr, } n=0, 1, \dots, 8$$

$$= 5 \text{ arrivals/hr, } n=9, 10, \dots, 11, 12$$

$$M_n = n/5 = 2n/\text{hr}, n=1, 2, 3, 4$$

$$= 10/\text{hr}, n \geq 5$$

$$P_1 = \frac{6}{2} P_0 = 3P_0$$

$$P_2 = \frac{6}{2} \cdot \frac{6}{4} P_0 = 4.5 P_0$$

$$P_3 = \frac{6}{2} \cdot \frac{6}{4} \cdot \frac{6}{6} P_0 = 4.5 P_0$$

$$P_4 = \frac{6}{2} \cdot \frac{6}{4} \cdot \frac{6}{6} \cdot \frac{6}{8} P_0 = 3.375 P_0$$

$$P_5 = \frac{6}{2} \cdot \frac{6}{4} \cdot \frac{6}{6} \cdot \frac{6}{8} \cdot \frac{6}{10} P_0 = 2.025 P_0$$

$$P_6 = \frac{6}{2} \cdot \frac{6}{4} \cdot \frac{6}{6} \cdot \frac{6}{8} \cdot \frac{6}{10} P_0 = 1.215 P_0$$

$$P_7 = \frac{6}{2} \cdot \frac{6}{4} \cdot \frac{6}{6} \cdot \frac{6}{8} \cdot \frac{6}{10} \cdot \frac{6}{10} P_0 = .729 P_0$$

$$P_8 = \frac{6}{2} \cdot \frac{6}{4} \cdot \frac{6}{6} \cdot \frac{6}{8} \cdot \frac{(6)}{10}^4 P_0 = .4374 P_0$$

$$P_{n \geq 9} = \frac{6}{2} \cdot \frac{6}{4} \cdot \frac{6}{6} \cdot \frac{6}{8} \cdot \frac{(6)}{10}^4 \cdot \left(\frac{5}{10}\right)^{n-8} P_0 = .4374 \cdot (5)^{n-8} P_0$$

$$\text{From } \sum_{n=0}^{12} P_n = 1, \text{ we get } P_0 = .0495$$

$$(a) P_{12} = .4374 \times .5^4 \times .0495 = .00135$$

$$(b) P_{n \geq 5} = 1 - (P_0 + P_1 + \dots + P_4) = .2385$$

$$(c) \text{Av. # busy tables} = 0P_0 + 1P_1 + 2P_2 + 3P_3 + 4P_4 + 5P_{n \geq 5} = 2.9768$$

4 continued

$$(d) 1P_6 + 2P_7 + \dots + 7P_{12}$$

5 continued

$$= 1x.0602 + 2x.0361 + 3x.0217 + 4x.0108 + 5x.0054 + 6x.0027 + 7x.00135 \\ = .2935 \text{ pair}$$

$$\lambda = 4 \text{ customers/hr}$$

6

$$\lambda_n = \begin{cases} 4, & n=0, 1, \dots, 4 \\ 0, & n \geq 5 \end{cases}$$

$$M_n = \frac{60}{15} = 4 \text{ customers/hr}$$

$$(a) P_1 = \frac{4}{4} P_0$$

$$P_2 = \left(\frac{4}{4}\right)^2 P_0$$

$$P_3 = \left(\frac{4}{4}\right)^3 P_0$$

$$P_4 = \left(\frac{4}{4}\right)^4 P_0$$

$$P_0 + P_1 + \dots + P_4 = 1 \Rightarrow P_0 = 1/5$$

$$P_0 = P_1 = P_2 = P_3 = P_4 = 1/5$$

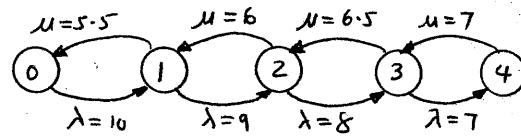
(b) expected # in shop =

$$0P_0 + 1P_1 + 2P_2 + 3P_3 + 4P_4$$

$$= \frac{1}{5}(1+2+3+4) = 2$$

$$(c) P_4 = .2$$

7



$$(a) 5.5 P_1 = 10 P_0$$

$$10P_0 + 6P_2 = (5.5 + 9)P_1$$

$$9P_1 + 6.5P_3 = (6 + 8)P_2$$

$$8P_2 + 7P_4 = (6.5 + 7)P_3$$

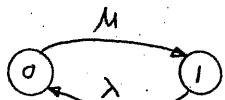
$$(b) P_1 = 1.82 P_0 \Rightarrow P_2 = 2.727 P_0$$

$$P_3 = 3.3566 P_0, P_4 = 3.3566 P_0$$

$$P_0 + P_1 + \dots + P_4 = 1 \Rightarrow P_0 = .088882$$

$$P_1 = .1614, P_2 = .2422, P_3 = .2981, P_4 = .33566$$

8



$$(a) \mu P_1 = \lambda P_0$$

$$P_1 = \frac{\lambda}{\mu} P_0$$

$$(b) P_0 + \frac{\lambda}{\mu} P_0 = 1$$

$$P_0 = \frac{1}{1+\rho}, \quad \rho = \lambda/\mu$$

$$P_1 = \frac{\rho}{1+\rho}$$

$$(c) L_s = 0P_0 + 1P_1 = \frac{\rho}{1+\rho}$$

$$(d) \lambda_{eff} = \lambda P_0 = \frac{\lambda}{1+\rho}$$

$$(e) W_f = \frac{L_s}{\lambda_{eff}} - \frac{1}{\mu} \\ = \frac{\rho/(1+\rho)}{\lambda/(1+\rho)} - \frac{1}{\mu} = 0$$

9

$$\lambda_{n-1} P_{n-1} + M_{n+1} P_{n+1} =$$

$$\lambda_{n-1} \left(\frac{\lambda_0}{M_1} \cdot \frac{\lambda_1}{M_2} \cdots \frac{\lambda_{n-2}}{M_{n-1}} \right) +$$

$$M_{n+1} \left(\frac{\lambda_0}{M_1} \cdot \frac{\lambda_1}{M_2} \cdots \frac{\lambda_n}{M_{n+1}} \right)$$

$$= M_n \left(\frac{\lambda_0}{M_1} \frac{\lambda_1}{M_2} \cdots \frac{\lambda_{n-1}}{M_n} \right) +$$

$$\lambda_n \left(\frac{\lambda_0}{M_1} \frac{\lambda_1}{M_2} \cdots \frac{\lambda_{n-1}}{M_n} \right)$$

$$= M_n P_n + \lambda_n P_n$$

$$= (M_n + \lambda_n) P_n$$

Set 15.6a

$$\begin{aligned}
 (a) L_q &= \sum_{n=6}^8 (n-5) P_n \\
 &= 1P_6 + 2P_7 + 3P_8 \\
 &= 1 \times 0.05847 + 2 \times 0.03508 + 3 \times 0.02105 \\
 &= .19177
 \end{aligned}$$

$$\begin{aligned}
 (b) W_q &= \frac{L_q}{\lambda_{\text{eff}}} \\
 &= \frac{.19177}{5.8737} = .03265 \text{ hour} \\
 W_s &= W_q + \frac{1}{\mu} \\
 &= .03265 + \frac{1}{2} = .53265 \text{ hour}
 \end{aligned}$$

$$\begin{aligned}
 (c) \lambda_{\text{lost}} &= \lambda P_8 \\
 &= 6 \times 0.02105 = .1263 \text{ car/hr}
 \end{aligned}$$

$$\text{Number lost}/8 \text{ hrs} = .1263 \times 8 = 1.01 \text{ cars}$$

$$\begin{aligned}
 (d) \text{Average number of empty spaces} \\
 &= C - (L_s - L_q) \\
 &= C - \sum_{n=0}^8 n P_n + \sum_{n=c+1}^8 (n-C) P_n \\
 &= (C \sum_{n=0}^8 P_n - C \sum_{n=c+1}^8 P_n) \\
 &\quad - \left(\sum_{n=0}^8 n P_n - \sum_{n=c+1}^8 n P_n \right) \\
 &= C \sum_{n=0}^c P_n - \sum_{n=0}^c n P_n \\
 &= \sum_{n=0}^{c-1} (C-n) P_n
 \end{aligned}$$

$$(e) \lambda_n = 6 \text{ cars/hr}, n=0, 1, \dots, 6$$

$$M_n = \begin{cases} \left(\frac{4}{3}\right)n, & n=1, 2, \dots, 6 \\ 8, & n=7, 8, 9, 10 \end{cases}$$

$$P_n = \left(\frac{6}{413}\right)^n \frac{1}{n!} P_0, n=0, 1, \dots, 6$$

2

continued...

$$P_n = \frac{\left(\frac{6}{413}\right)^n}{6!} P_0, n=7, 8, 9, 10$$

2 continued

$$\begin{aligned}
 P_0 \left(1 + \frac{9/2}{1!} + \frac{(9/2)^2}{2!} + \frac{(9/2)^3}{3!} + \frac{(9/2)^4}{4!} + \frac{(9/2)^5}{5!} + \frac{(9/2)^6}{6!} \right. \\
 \left. + \frac{(9/2)^7}{6!6^1} + \frac{(9/2)^8}{6!6^2} + \frac{(9/2)^9}{6!6^3} + \frac{(9/2)^{10}}{6!6^4} \right) = 1
 \end{aligned}$$

$$\text{Thus, } P_0 = .0004$$

n	P_n	n	P_n
1	.0004	6	.10027
2	.001141	7	.12534
3	.002852	8	.15667
4	.005348	9	.19584
5	.008022	10	.24480

$$\begin{aligned}
 (b) \lambda_{\text{eff}} &= \lambda (1 - P_{10}) = 10 (1 - 0.2448) \\
 &= 7.552 \text{ cars/hr}
 \end{aligned}$$

$$\begin{aligned}
 (c) L_s &= 0P_0 + 1P_1 + 2P_2 + \dots + 10P_{10} \\
 &= 7.6941 \text{ cars}
 \end{aligned}$$

$$(d) W_s = \frac{L_s}{\lambda_{\text{eff}}} = \frac{7.6941}{7.552} = 1.0155 \text{ cars}$$

$$W_q = 1.0155 - \frac{1}{4/3} = .2655$$

$$\begin{aligned}
 (e) L_q &= \lambda_{\text{eff}} W_q \\
 &= .2655 \times 7.552 \\
 &= 2.005 \text{ cars}
 \end{aligned}$$

$$\text{Average number of occupied spaces} = L_s - L_q$$

$$= 7.6941 - 2.005$$

$$= 5.6891 \text{ spaces}$$

Set 15.6b

(a) % utilization = $100(1-P_0)$

$$= 100 \cdot \frac{\lambda}{\mu}$$

$$= 100 \left(\frac{4}{6}\right) = 66.67\%$$

(b) $P_{n \geq 1} = 1 - P_0 = \frac{\lambda}{\mu} = \frac{4}{6} = .6667$

(c) $P_{n \leq 7} = P_0 + P_1 + \dots + P_7$

$$= 1 - \left(\frac{\lambda}{\mu}\right)^8 = 1 - \left(\frac{4}{6}\right)^8 = .961$$

(d) $P_0 + P_1 + \dots + P_K \geq .99$

From Figure 17-6, $K = 11$

Also, we can determine K from

$$1 - P^{K+1} \geq .99$$

$$(K+1) \geq \frac{\ln .01}{\ln (4/6)} = 11.$$

or $K \geq 11.350 - 1 = 10.358$

Thus, $K \geq 11$

Note that the desired number of parking spaces is almost doubled (from 5 to 11) to accommodate the increase in the acceptance percentage from 90% to 99%.

$\lambda = 1/5 = .2 \text{ job/day}$

$\mu = 1/4 = .25 \text{ job/day}$

From the TOR A output on the next column,

(a) $P_0 = .2$

(b) Av. income/month = \$50 μt

$$= 50 \times .25 \times 30 \\ = \$375$$

(c) Av. number of jobs awaiting completion = $L_q = 3.2 \text{ jobs}$

cost = $3.2 \times \$40 = \128

Continued...

Lambda = 0.20000	Mu = 0.25000
Lambda eff = 0.20000	Rho/c = 0.80000
Ls = 4.00000	Lq = 3.20000
Ws = 20.00000	Wq = 16.00000

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.20000	0.20000	23	0.00118	0.99528
1	0.16000	0.38000	24	0.00094	0.99622
2	0.12800	0.48800	25	0.00075	0.99693
3	0.10240	0.59040	26	0.00060	0.99758
4	0.08192	0.67232	27	0.00048	0.99807
5	0.06554	0.73786	28	0.00039	0.99845
6	0.05243	0.79028	29	0.00031	0.99876
7	0.04194	0.83223	30	0.00025	0.99901
8	0.03355	0.86578	31	0.00020	0.99921
9	0.02684	0.89263	32	0.00016	0.99937
10	0.02147	0.91410	33	0.00013	0.99949
11	0.01718	0.93128	34	0.00010	0.99957
12	0.01374	0.94502	35	0.00008	0.99965
13	0.01100	0.95602	36	0.00006	0.99974
14	0.00880	0.96482	37	0.00005	0.99979
15	0.00704	0.97185	38	0.00004	0.99983
16	0.00563	0.97748	39	0.00003	0.99987
17	0.00450	0.98199	40	0.00003	0.99989
18	0.00360	0.98559	41	0.00002	0.99991
19	0.00288	0.98847	42	0.00002	0.99993
20	0.00231	0.99078	43	0.00001	0.99995
21	0.00184	0.99262	44	0.00001	0.99996
22	0.00148	0.99410			

$\lambda = 1/4 = .25 \text{ case/wk}$

$\mu = 1/1.5 = .66667 \text{ case/wk}$

M/M/c/GD/N/K Queueing Model

Input Data		Output Results	
$\lambda =$	0.25	$\lambda_{eff} =$	0.66667
$c =$	1	∞	∞
Sys. Lim. / N =		Source limit / K =	∞
$\lambda_{eff} =$	0.2500	$\lambda_{eff} =$	0.3750
$L_s =$	0.6000	$L_q =$	0.2250
$W_s =$	2.4000	$W_q =$	0.9000
n	Pn	CPn	1-CPn
0	0.625002	0.625002	0.374998
1	0.234375	0.859376	0.140624
2	0.087890	0.947266	0.052734
3	0.032959	0.980225	0.019775
4	0.012359	0.992584	0.007416
5	0.004635	0.997219	0.002781
6	0.001738	0.998957	0.001043
7	0.000652	0.999609	0.000391
8	0.000244	0.999853	0.000147
9	0.000092	0.999945	0.000055
10	0.000034	0.999979	0.000021
11	0.000013	0.999992	0.000008
12	0.000005	0.999997	0.000003
13	0.000002	0.999999	0.000001
14	0.000001	1.000000	0.000000

(a) $L_q = .225 \text{ case}$

(b) $1 - P_0 = 1 - .625 = .375 \text{ or } 37.5\%$

(c) $W_s = 2.4 \text{ weeks}$

Present situation :

$\lambda = 90 \text{ cars/hr}$

$\mu = \frac{3600}{38} = 94.7368 \text{ cars/hr}$

New situation:

$\lambda = 90 \text{ cars per hour}$

$\mu = \frac{3600}{30} = 120 \text{ cars per hour}$

Continued...

Set 15.6b

Comparative Analysis

Scenario	c	Lambda	Mu	Lambda eff	p0	Ls	Lq	Ws	Wq
1	1	90.00000	94.73680	90.00000	0.05000	19.00017	18.00017	0.21111	0.20056
2	1	90.00000	120.00000	90.00000	0.25000	3.00000	2.25000	0.03333	0.02500

$$L_s (\text{present}) = 19 \text{ cars}$$

$$\% \text{ of idle time (new)} = p_0 (\text{new}) \times 100 \\ = 100 \times .25 = 25\%$$

The device can be justified based on the number of waiting customers, L_q , in the present system, but not on the basis of % idle time in the new one.

$$(b) 1 - CP_2 = 1 - .4213 = .5787$$

$$(c) W_q = .417 \text{ hour}$$

(d) Let $N = \text{spaces (including car being served)}$

$$CP_{N-1} \geq .9$$

Because $CP_1 = .88784$ and $CP_2 = .90654$, $N-1 \geq 12 \Rightarrow N \geq 13$.

In general, $L_s < L_q + 1$. The reason is that $p > 0$, usually. Consider

$$L_q = \sum_{n=1}^{\infty} (n-1) p_n \\ = \sum_{n=1}^{\infty} n p_n - \sum_{n=1}^{\infty} p_n \\ = L_s - (1-p)$$

The closer p is to zero, the more likely $L_s \approx L_q + 1$ will hold.

7

Scenario 1-(MM/1):(GD/infinity/infinity)

5

Lambda =	0.40000	Mu =	0.66667
Lambda eff =	0.40000	Rho/c =	0.60000
Ls =	1.49998	Lq =	0.89998
Ws =	3.74995	Wq =	2.24996

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.40000	0.40000	11	0.00145	0.99782
1	0.24000	0.64000	12	0.00087	0.99869
2	0.14400	0.78400	13	0.00052	0.99922
3	0.08640	0.87040	14	0.00031	0.99953
4	0.05184	0.92224	15	0.00019	0.99972
5	0.03110	0.95335	16	0.00011	0.99983
6	0.01866	0.97201	17	0.00007	0.99990
7	0.01210	0.98320	18	0.00004	0.99994
8	0.00672	0.98992	19	0.00002	0.99996
9	0.00403	0.99395	20	0.00001	0.99998
10	0.00242	0.99637			

$$(a) p_0 = .4$$

$$(b) L_q = .9 \text{ car}$$

$$(c) W_q = 2.25 \text{ minutes}$$

$$(d) P_{n \geq 11} = 1 - CP_{10} = 1 - .99637 = .0036$$

Scenario 1-(MM/1):(GD/infinity/infinity)

6

Lambda =	10.00000	Mu =	12.00000
Lambda eff =	10.00000	Rho/c =	0.83333
Ls =	5.00000	Lq =	4.16667
Ws =	0.50000	Wq =	0.41667

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.16667	0.16667	27	0.00121	0.99393
1	0.13889	0.30556	28	0.00101	0.99494
2	0.11574	0.42130	29	0.00084	0.99579
3	0.09565	0.51775	30	0.00070	0.99649
4	0.08030	0.59812	31	0.00059	0.99707
5	0.06569	0.66510	32	0.00049	0.99756
6	0.05582	0.72092	33	0.00041	0.99797
7	0.04651	0.76743	34	0.00034	0.99831
8	0.03876	0.80619	35	0.00028	0.99859
9	0.03230	0.83849	36	0.00024	0.99882
10	0.02692	0.86541	37	0.00020	0.99902
11	0.02243	0.88784	38	0.00016	0.99918
12	0.01869	0.90654	39	0.00014	0.99932
13	0.01558	0.92211	40	0.00011	0.99943
14	0.01298	0.93509	41	0.00009	0.99953
15	0.01082	0.94591	42	0.00008	0.99961
16	0.00901	0.95493	43	0.00007	0.99967
17	0.00751	0.96244	44	0.00005	0.99973
18	0.00626	0.96870	45	0.00005	0.99977
19	0.00522	0.97392	46	0.00004	0.99981
20	0.00435	0.97826	47	0.00003	0.99984
21	0.00362	0.98189	48	0.00003	0.99987
22	0.00302	0.98491	49	0.00002	0.99989
23	0.00252	0.98742	50	0.00002	0.99991
24	0.00210	0.98952	51	0.00002	0.99992
25	0.00175	0.99126	52	0.00001	0.99994
26	0.00146	0.99272	53	0.00001	0.99995

$$(a) p_0 + p_1 + p_2 = .4213$$

continued...

15-16

$$(b) 1 - CP_2 = 1 - .4213 = .5787$$

$$(c) W_q = .417 \text{ hour}$$

(d) Let $N = \text{spaces (including car being served)}$

$$CP_{N-1} \geq .9$$

Because $CP_1 = .88784$ and $CP_2 = .90654$, $N-1 \geq 12 \Rightarrow N \geq 13$.

7

$$L_q = \sum_{n=1}^{\infty} (n-1) p_n \\ = \sum_{n=1}^{\infty} n p_n - \sum_{n=1}^{\infty} p_n \\ = L_s - (1-p)$$

The closer p is to zero, the more likely $L_s \approx L_q + 1$ will hold.

8

Consider

$$L_q = \sum_{n=1}^{\infty} (n-1) p_n \\ = \sum_{n=1}^{\infty} (n-1)(1-p)^{n-1} \\ = (1-p) \sum_{n=1}^{\infty} \frac{d}{dp} (1-p)^{n-1} \\ = (1-p) \sum_{n=1}^{\infty} \frac{d}{dp} \sum_{n=0}^{\infty} p^n \\ = (1-p) \sum_{n=1}^{\infty} \frac{p^2}{(1-p)^2} \left(\frac{1}{1-p} \right) \\ = p^2 (1-p) \frac{1}{(1-p)^2} \\ = \frac{p^2}{1-p}$$

$$(a) P\{j \text{ in queue} | j \geq 1\}$$

$$= P\{n \text{ in system} | n \geq 2\}$$

$$= \frac{P_n}{\sum_{j=2}^{\infty} P_j}$$

Thus,

$$\text{expected number} = \sum_{n=2}^{\infty} (n-1) \frac{P_n}{\sum_{j=2}^{\infty} P_j}$$

$$= \frac{\sum_{n=2}^{\infty} np_n - \sum_{n=2}^{\infty} P_n}{\sum_{n=2}^{\infty} P_n}$$

$$= \frac{\sum_{n=1}^{\infty} np_n - P_1}{\sum_{n=2}^{\infty} P_n} - 1$$

$$= \frac{\frac{P}{1-P} - P(1-P)}{1 - [(1-P) + P(1-P)]} - 1$$

$$= \frac{1}{1-P}$$

(b) Exp. number in queue given the system is not empty

$$= \sum_{n=1}^{\infty} (n-1) \left(\frac{P_n}{\sum_{j=1}^{\infty} P_j} \right)$$

$$= \frac{\sum_{n=1}^{\infty} np_n - \sum_{n=1}^{\infty} P_n}{\sum_{j=1}^{\infty} P_j}$$

$$= \frac{\left(\frac{P}{1-P} \right) - P}{P}$$

$$= \frac{P}{1-P}$$

Thus,

Exp. waiting time in queue for those who must wait

$$= \frac{P/(1-P)}{\lambda}$$

$$= \frac{1}{M-\lambda}$$

Continued...

Set 15.6c

1

$$w(\tau) = (\mu - \lambda) e^{-(\mu - \lambda)\tau}, \tau > 0$$

$$\begin{aligned} \lambda &= 1/4 = .25/\text{hr} \\ \mu &= 1/1.5 = .667/\text{hr} \end{aligned} \quad \left. \begin{aligned} (\mu - \lambda) &= .417 \\ f &= \lambda/\mu = \frac{1.5}{4} = .375 \end{aligned} \right\}$$

$$w(\tau) = .417 e^{-0.417\tau}, \tau > 0$$

$$P\{\tau > 1\} = e^{-0.417 \times 1} = .659$$

(a) Standard deviation = $\frac{1}{\mu - \lambda} = \frac{1}{0.417} = 2.4$ 2

(b) $w(\tau) = (\mu - \lambda) e^{-(\mu - \lambda)\tau}, \tau > 0$

$$P\left\{\frac{1}{2(\mu - \lambda)} \leq \tau \leq \frac{3}{2(\mu - \lambda)}\right\}$$

$$= (1 - e^{-1.5}) - (1 - e^{-0.5})$$

$$= e^{-0.5} - e^{-1.5}$$

$$= .3834$$

$W_s \leq 10$ minutes, $\lambda = 4/\text{hr}$ 3

$$\frac{1}{(\mu - \lambda)} \leq \frac{10}{60} \text{ hr}$$

or $\mu - \lambda \geq 6$

or $\mu \geq 6 + \lambda = 10/\text{hr}$

$P\{\tau > \frac{10}{60}\} \leq .1$, or

$$e^{-\frac{1}{6}(\mu - 4)} \leq .1$$

$\mu - 4 \geq 13.8$

$\mu \geq 17.8/\text{hr}$

$P\{\tau > 5\} = e^{-(\mu - \lambda)t} = e^{-0.267 \times 5} = .2636$ 5

where $\lambda = .4/\text{min}$, $\mu = .667/\text{min}$

Exp. # customers in a 12-hr day
 $= \lambda \times 12 \times 60 = .4 \times 12 \times 60 = 288$ cust.

Exp. cost = $288 \times .2636 \times .5 = \37.95

6

Let $w_{n+1}(t|n)$ = conditional pdf for waiting in queue given there are n customers ahead

= n -fold convolution of the exponential pdf

$$= \frac{\mu(\mu t)^{n-1} e^{-\mu t}}{(n-1)!}$$

$w(t)$ = absolute pdf of waiting time in queue

$g(t, n)$ = joint pdf of t and n

$$= w_{n+1}(t|n) p_n$$

$$= \frac{\mu(\mu t)^{n-1} e^{-\mu t}}{(n-1)!} f^n(1-p)$$

(a) For $t > 0$

$$w(t) = \sum_{n=1}^{\infty} g(t, n)$$

$$= \frac{\mu p e^{-\mu t} (1-p)}{e^{-\mu p t}} \sum_{n=1}^{\infty} \frac{(\mu p t)^{n-1} e^{\mu p t}}{(n-1)!}$$

$$= \mu p (1-p) e^{\mu(1-p)t}, t > 0$$

For $t = 0$, $w(0) = p_0 = (1-p)$

$$w(t) = \begin{cases} 1-p, & t=0 \\ \mu p (1-p) e^{\mu(1-p)t}, & t>0 \end{cases}$$

(b) $W_q = E\{t\}$

$$= \int_0^\infty t w(t) dt$$

$$= 0 w(0) + \int_0^\infty t w(t) dt$$

$$= \int_{0^+}^\infty \mu p t (1-p) e^{\mu(1-p)t} dt$$

$$= \frac{p}{\mu(1-p)}$$

Set 15.6d

(a) $p_0 = .3654$

(b) $W_q = .207$ hour

(c) Average number of empty spaces = $4 - L_q$

$$= 4 - .788$$

= 3.212 spaces

(d) $\frac{p_0}{5} = .04812$

(e) $W_s \leq 10$ minutes

Title: 17.6d-1

Comparative Analysis

Scenario	c	Lambda	Mu	Lambda eff	p0	Ls	Lq	Ws	Wq
1	1	4.00000	6.00000	3.80752	0.36541	1.42256	0.78797	0.37362	0.20695
2	1	6.00000	9.00000	3.91779	0.44003	1.1691	0.56034	0.28659	0.14413
3	1	4.00000	6.00000	3.83752	0.36541	1.42256	0.78797	0.37362	0.20695
4	1	4.00000	6.00000	3.95118	0.56987	0.75340	0.31257	0.19520	0.10684
5	1	4.00000	10.00000	3.97532	0.60247	0.64199	0.24446	0.16149	0.06149

	M (cars/hr)	W_s (hrs)	W_s (min)
6		.3736	22.4
7		.287	17.16
8		.23	13.80
9		.19	11.40
10		.16	9.60

Desired service rate = 10 cars/hr
 Thus, the service time must be reduced from $\frac{60}{10} = 6$ minutes to $\frac{60}{15} = 4$ minutes. a 40% reduction

m = number of parking spaces

An arriving car will not find a space

if there are $m+1$ cars in the system. Thus,
 find m such that $P_{m+1} \leq .01$

TO RA input = (4, 6, 1, $m+1$, ∞)

m	$N=m+1$	P_N
4	5	.04812
5	6	.0311
6	7	.0203
7	8	.01335
8	9	.009

Select the number of parking spaces $m \geq 8$

Continued...

m = number of seats.

The $N = m+1$, and

$$\lambda_{eff} = \frac{\lambda}{N} = \frac{\lambda}{m+1} \text{ customers/hr}$$

TO RA input = (6, 5, 1, N , ∞)

Title: 17.6d-3

Comparative Analysis

Scenario	c	Lambda	Mu	Lambda eff	p0	Ls	Lq	Ws	Wq
1	1	6.00000	6.00000	3.62337	0.27473	1.12088	0.36560	0.30909	0.10909
2	1	6.00000	9.00000	4.22310	0.27473	0.91207	0.42418	0.22418	0.09091
3	1	6.00000	5.00000	4.32310	0.13438	2.72578	0.13438	0.27258	0.06149
4	1	6.00000	5.00000	4.49647	0.10071	3.02117	2.12188	0.67190	0.47190
5	1	6.00000	5.00000	4.51288	0.07742	3.70584	2.78728	0.80423	0.50423

$$m \quad N = m+1 \quad \lambda_{eff} (\text{customers/hr})$$

$$1 \quad 2 \quad 3.63 \\ 2 \quad 3 \quad 4.07$$

Use two seats or less

$\lambda = 10$ generators per hour

$$M = \frac{60}{15} = 4 \text{ generators per hour}$$

$$N = 7 + 1 = 8$$

Title: 17.6d-4

Scenario 1—(M/M/1):(GD/8/Infinity)

Lambda = 10.00000	Mu = 4.00000
Lambda eff = 3.99843	Rho/c = 2.50000
Ls = 7.33569	Lq = 6.33609
Ws = 1.83454	Wq = 1.58464

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.00039	0.00039	5	0.03841	0.06375
1	0.00098	0.00138	6	0.09603	0.15978
2	0.00246	0.00383	7	0.24006	0.39984
3	0.00615	0.00998	8	0.60016	1.00000
4	0.01536	0.02534			

$$(a) P_8 \approx .6$$

$$(b) Lq = 6.34 \text{ generators}$$

(c) Let C = belt capacity. Thus,
 $N = C + 1$. The assembly department is kept in operation so long as at least one empty space remains on the belt; that is,

$$P\{\text{empty space on belt}\} = P_0 + P_1 + \dots + P_C$$

$$= \frac{1-p}{1-p^{C+2}} \sum_{n=0}^C p^n$$

$$= \frac{1-p}{1-p^{C+2}} \cdot \frac{1-p^{C+1}}{1-p}$$

$$= \frac{1-p^{C+1}}{1-p^{C+2}}$$

Continued...

Continued...

Set 15.6d

$$\lim_{c \rightarrow \infty} \frac{1 - p^{c+1}}{1 - p^{c+2}} = \lim_{c \rightarrow \infty} \frac{-(c+1)p^c}{-(c+2)p^{c+1}}$$

$$= \lim_{c \rightarrow \infty} \frac{c+1}{(c+2)p}$$

$$= \lim_{c \rightarrow \infty} \frac{(1 + 1/c)}{(1 + 2/c)} \frac{1}{p}$$

$$= \frac{1}{p}$$

In the present example, $p = 10/4$ and $1/p = .4$. Thus,

$$\lim_{c \rightarrow \infty} (p_0 + p_1 + \dots + p_c) = 1/p = .4$$

This result means that regardless of how large the belt is, the probability of finding an empty space cannot exceed .4. Thus, achieving a 95% utilization for the assembly dept. is impossible.

The result makes sense because the arrival rate λ ($= 10/\text{hr}$) is 2½ times larger than the service rate ($= 4$). The only way we can accomplish the desired result is to reduce λ and/or increase M .

(a) $P_{50} \approx .00002$

(b) $P\{\text{w} \leq 48 \text{ or more in restaurant}\}$

$$= P_{48} + P_{49} + P_{50}$$

$$= 1 - (P_0 + P_1 + \dots + P_{47})$$

$$= 1 - .99993$$

$$= .00007$$

5

Title: 17-6d-5
Scenario 1-(M/M/1):(GD/50/infinity)

TORA input = (10, 12, 1, 50, ∞)

Lambda = 10.00000 Mu = 12.00000
Lambda eff = 9.99982 Rho/c = 0.83333
Ls = 4.99533 Lq = 4.16201
Ws = 0.49954 Wq = 0.41621

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.16668	0.16668	26	0.00146	0.99281
1	0.13890	0.30558	27	0.00121	0.99402
2	0.11575	0.42133	28	0.00101	0.99504
3	0.09646	0.51779	29	0.00084	0.99584
4	0.08038	0.59818	30	0.00070	0.99558
5	0.06599	0.66516	31	0.00059	0.99717
6	0.05592	0.72098	32	0.00049	0.99765
7	0.04652	0.76759	33	0.00041	0.99806
8	0.03876	0.80627	34	0.00034	0.99840
9	0.03230	0.83257	35	0.00028	0.99868
10	0.02691	0.85549	36	0.00024	0.99892
11	0.02243	0.88792	37	0.00020	0.99911
12	0.01869	0.90662	38	0.00016	0.99928
13	0.01558	0.92220	39	0.00014	0.99941
14	0.01298	0.93518	40	0.00011	0.99952
15	0.01082	0.94600	41	0.00009	0.99962
16	0.00902	0.95501	42	0.00008	0.99970
17	0.00751	0.96253	43	0.00007	0.99976
18	0.00626	0.96879	44	0.00006	0.99982
19	0.00522	0.97401	45	0.00005	0.99986
20	0.00435	0.97835	46	0.00004	0.99990
21	0.00362	0.98198	47	0.00003	0.99993
22	0.00302	0.98500	48	0.00003	0.99995
23	0.00252	0.98751	49	0.00002	0.99998
24	0.00210	0.98961	50	0.00002	1.00000
25	0.00175	0.99136			

6

TORA input = (20, 7.5, 1, 15, ∞)

Title: 17-6d-6
Scenario 1-(M/M/1):(GD/15/infinity)

Lambda = 20.00000 Mu = 7.50000
Lambda eff = 7.50000 Rho/c = 2.66667
Ls = 14.40000 Lq = 13.40000
Ws = 1.92000 Wq = 1.78667

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.00000	0.00000	8	0.00065	0.00104
1	0.00000	0.00000	9	0.00174	0.00278
2	0.00000	0.00000	10	0.00463	0.00742
3	0.00000	0.00001	11	0.01236	0.01678
4	0.00001	0.00002	12	0.03296	0.05273
5	0.00003	0.00005	13	0.08789	0.14062
6	0.00009	0.00015	14	0.23438	0.37500
7	0.00024	0.00039	15	0.62500	1.00000

- (a) $P_0 \approx 0$
(b) $P_{n \leq 14} = P_0 + \dots + P_{14} = .375$
(c) $W_s = 1.92 \text{ hours}$

7

(a) $P_{n \leq 4} = P_0 + P_1 + \dots + P_4$
= .962

(b) $\lambda_{lost} = \lambda P_5$
= $5 \times .038 = .19 \text{ cust./hr}$

(c) $L_s = 0 \times .399 + 1 \times .249 + 2 \times .156$
+ $3 \times .097 + 4 \times .061$
+ $5 \times .038$
= 1.286

continued...

continued...

$$(d) W_q = W_s - \frac{1}{\mu}$$

$$\lambda_{eff} = 5(1 - 0.38) = 4.81 \text{ cust/hr}$$

$$W_s = \frac{L_s}{\lambda_{eff}}$$

$$= \frac{1.286}{4.81}$$

$$= .2675 \text{ hour}$$

$$W_q = .2675 - \frac{1}{8}$$

$$= .1424 \text{ hour}$$

$$P_n = \frac{(1-p)p^n}{1-p^{N+1}}$$

8

$$\begin{aligned} \lim_{p \rightarrow 1} p_n &= \lim_{p \rightarrow 1} \frac{p^n - p^{n+1}}{1 - p^{N+1}} \\ &= \lim_{p \rightarrow 1} \frac{n p^{n-1} - (n+1)p^n}{-(N+1)p^N} \\ &= \frac{1}{N+1} \end{aligned}$$

Thus,

$$\begin{aligned} L_s &= \sum_{n=0}^N n p_n \\ &= \frac{1}{N+1} \sum_{n=0}^N n \\ &= \frac{N(N+1)}{2(N+1)} = \frac{N}{2} \end{aligned}$$

$$W_s = W_q + \frac{1}{\mu}$$

$$\lambda_{eff} W_s = \lambda_{eff} W_q + \frac{\lambda_{eff}}{\mu}$$

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Thus,

$$L_s = L_q + \frac{\lambda_{eff}}{\mu}$$

or

$$\lambda_{eff} = \mu(L_s - L_q)$$

Set 15.6e

TOA input = (8, 5, 2, 00, 00)

Title: 17.6e-1
Scenario 1-(M/M/2):(GD/infinity/infinity)

Lambda = 8.00000 Mu = 5.00000
Lambda eff = 8.00000 Rho/c = 0.80000
Ls = 4.44444 Lq = 2.84444
Ws = 0.55556 Wq = 0.35556

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.11111	0.11111	23	0.00131	0.99475
1	0.17778	0.28889	24	0.00105	0.99580
2	0.14222	0.43111	25	0.00084	0.99524
3	0.11378	0.54489	26	0.00067	0.99731
4	0.09102	0.63591	27	0.00054	0.99785
5	0.07282	0.70873	28	0.00043	0.99828
6	0.05825	0.76698	29	0.00034	0.99862
7	0.04660	0.81359	30	0.00028	0.99890
8	0.03728	0.85087	31	0.00022	0.99912
9	0.02983	0.88707	32	0.00018	0.99930
10	0.02386	0.90456	33	0.00014	0.99944
11	0.01909	0.92365	34	0.00011	0.99955
12	0.01527	0.93892	35	0.00009	0.99954
13	0.01222	0.95113	36	0.00007	0.99971
14	0.00977	0.96091	37	0.00006	0.99977
15	0.00782	0.96873	38	0.00005	0.99982
16	0.00625	0.97498	39	0.00004	0.99985
17	0.00500	0.97998	40	0.00003	0.99988
18	0.00400	0.98399	41	0.00002	0.99991
19	0.00320	0.98719	42	0.00002	0.99992
20	0.00256	0.98975	43	0.00002	0.99994
21	0.00205	0.99180	44	0.00001	0.99995
22	0.00164	0.99344			

TOA input = (16, 5, 4, 00, 00)

Title: 17.6e-1
Scenario 2-(M/M/4):(GD/infinity/infinity)

Lambda = 16.00000 Mu = 5.00000
Lambda eff = 16.00000 Rho/c = 0.80000
Ls = 5.58573 Lq = 2.38573
Ws = 0.34911 Wq = 0.14911

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.02730	0.02730	24	0.00138	0.99450
1	0.08737	0.11467	25	0.00110	0.99560
2	0.13979	0.25446	26	0.00086	0.99648
3	0.14911	0.40357	27	0.00070	0.99718
4	0.11929	0.52285	28	0.00056	0.99779
5	0.09543	0.61828	29	0.00045	0.99820
6	0.07634	0.69463	30	0.00036	0.99856
7	0.06107	0.75570	31	0.00029	0.99885
8	0.04886	0.80458	32	0.00023	0.99908
9	0.03939	0.84585	33	0.00018	0.99926
10	0.03127	0.87492	34	0.00015	0.99941
11	0.02502	0.89994	35	0.00012	0.99953
12	0.02001	0.91995	36	0.00009	0.99962
13	0.01601	0.93596	37	0.00008	0.99970
14	0.01281	0.94877	38	0.00006	0.99976
15	0.01025	0.95901	39	0.00005	0.99981
16	0.00820	0.96721	40	0.00004	0.99985
17	0.00656	0.97377	41	0.00003	0.99988
18	0.00525	0.97901	42	0.00002	0.99990
19	0.00420	0.98321	43	0.00002	0.99992
20	0.00336	0.98657	44	0.00002	0.99994
21	0.00269	0.98926	45	0.00001	0.99995
22	0.00215	0.99140	46	0.00001	0.99996
23	0.00172	0.99312			

(a) C=2:

$$P\{\text{all servers are busy}\} = \left(p_{n \geq 2} \right)^2 \\ = (1 - .29)^2 \\ = .504$$

C=4:

$$P\{\text{all servers are busy}\} = 1 - P_{n \leq 3} \\ = 1 - .404 \\ = .596$$

C=4 yields a higher probability that all servers are busy.

continued...

(b)

Title: 17.6e-1
Comparative Analysis

Scenario	c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
1	4	16.00000	5.00000	16.00000	0.02730	5.58573	2.38573	0.34911	0.14911
2	5	16.00000	5.00000	16.00000	0.02715	3.71299	0.51299	0.23208	0.03208
3	6	16.00000	5.00000	16.00000	0.02777	3.34328	0.14520	0.20908	0.00903

for C = 5, Wq = .032 hour ≈ 2 min

C = 4, Wq = .149 hour ≈ 9 min

Select C = 5

C = 2: $\lambda = 8 \text{ calls/hr}$

$$M = \frac{60}{14.5} = 4.1379 \text{ calls/hr}$$

C = 4: $\lambda = 16 \text{ calls/hr}$

$$M = 4.1379 \text{ calls per hour}$$

$$\text{utilization} = \lambda/M = .967$$

Title: 17.6e-1
Comparative Analysis

Scenario	c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
1	2	16.00000	8.00000	16.00000	0.01685	28.49905	27.58471	3.65728	3.44559

$Wq = \begin{cases} 3.446 \text{ hours for } C=2 \\ 1.681 \text{ hours for } C=4 \end{cases}$

Consolidation reduces the waiting time by more than 51%.

(a) $\lambda = \frac{60}{5} = 12 \text{ per hour}$

$$M = 10 \text{ per hour}$$

$$C > \frac{\lambda}{M} = 1.2 \Rightarrow C \geq 2$$

(b) $\lambda = \frac{60}{2} = 30 \text{ per hour}$

$$M = \frac{60}{6} = 10 \text{ per hour}$$

$$C > \frac{\lambda}{M} = \frac{30}{10} = 3 \Rightarrow C \geq 4$$

(c) $\lambda = 30 \text{ per hour}, M = 40 \text{ per hr.}$

$$C > \frac{30}{40} = .75 \Rightarrow C \geq 1$$

$\lambda = 45 \text{ customers/hr}$

$$M = \frac{60}{5} = 12 \text{ customers/hr}$$

$$C > \frac{45}{12} \text{ or } C \geq 4$$

Desired Wq ≤ 30 seconds = .0083 hr

Comparative Analysis

Scenario	c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
1	4	45.00000	12.00000	45.00000	0.00355	18.72545	12.57545	0.37193	0.28154
2	5	45.00000	12.00000	45.00000	0.01555	5.12537	3.35537	0.11412	0.03879
3	6	45.00000	12.00000	45.00000	0.02208	4.12903	0.37903	0.03176	0.00842
4	7	45.00000	12.00000	45.00000	0.02309	3.86873	0.11873	0.08597	0.00264

Select C ≥ 7.

2

4

Set 15.6e

TOA input: (20, 12, 3, ∞, ∞)

Title: 15.6e-5
Scenario 1- (M/M/3):(GD/infinity/infinity)

Lambda = 20.00000	Mu = 12.00000
Lambda eff = 20.00000	Rho/c = 0.55556
Ls = 2.04137	Lq = 0.37470
Ws = 0.10207	Wq = 0.01874

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.17266	0.17266	10	0.00218	0.99728
1	0.28777	0.46043	11	0.00121	0.99849
2	0.23981	0.70024	12	0.00067	0.99916
3	0.13323	0.83347	13	0.00037	0.99953
4	0.07401	0.90748	14	0.00021	0.99974
5	0.04112	0.94860	15	0.00012	0.99986
6	0.02284	0.97144	16	0.00006	0.99992
7	0.01269	0.98414	17	0.00004	0.99996
8	0.00705	0.99119	18	0.00002	0.99998
9	0.00392	0.99510	19	0.00001	0.99999

$m = \text{size of waiting room.}$

$$P_0 + P_1 + \dots + P_{m+2} \geq .999 \Rightarrow m \geq 10$$

$$C=2, \lambda_{\text{windows}} = .8 \times \frac{60}{3} = 16 \text{ /hr}$$

$$\mu = \frac{60}{5} = 12 \text{ per hour}$$

Title: 15.6e-6
Scenario 1- (M/M/2):(GD/infinity/infinity)

Lambda = 16.00000	Mu = 12.00000
Lambda eff = 16.00000	Rho/c = 0.66667
Ls = 2.40000	Lq = 1.06667
Ws = 0.15000	Wq = 0.06667

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.20000	0.20000	14	0.00137	0.99726
1	0.26667	0.46667	15	0.00091	0.99817
2	0.17778	0.64444	16	0.00061	0.99878
3	0.11852	0.76296	17	0.00041	0.99919
4	0.07901	0.84198	18	0.00027	0.99946
5	0.05267	0.89465	19	0.00018	0.99964
6	0.03512	0.92977	20	0.00012	0.99976
7	0.02341	0.95318	21	0.00008	0.99984
8	0.01611	0.96879	22	0.00005	0.99989
9	0.01040	0.97919	23	0.00004	0.99993
10	0.00694	0.98613	24	0.00002	0.99995
11	0.00462	0.99075	25	0.00002	0.99997
12	0.00308	0.99383	26	0.00001	0.99998
13	0.00206	0.99589			

$$(a) P_{n \geq 2} = 1 - (P_0 + P_1)$$

$$= 1 - .46667$$

$$= .5333$$

$$(b) P_0 = .2$$

$$(c) L_q = 1.067$$

(d) NO, because $\lambda > \mu$. The minimum number of windows should $\geq \frac{\lambda}{\mu} = \frac{16}{12} = 1.33$
Number of windows ≥ 2

5

$$\lambda = 25 \times \frac{60}{15} = 100 \text{ jobs/hour}$$

$$\mu = \frac{60}{2} = 30 \text{ jobs/hour}, C = 4$$

7

Title: 15.6e-7
Scenario 1- (M/M/4):(GD/infinity/infinity)

Lambda = 100.00000	Mu = 30.00000
Lambda eff = 100.00000	Rho/c = 0.83333
Ls = 6.62194	Lq = 3.28861
Ws = 0.06622	Wq = 0.03289

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.02131	0.02131	28	0.00138	0.99811
1	0.07163	0.09234	29	0.00115	0.99826
2	0.11839	0.21073	30	0.00096	0.99851
3	0.13154	0.34228	31	0.00080	0.99880
4	0.09202	0.45190	32	0.00066	0.99901
5	0.09135	0.54325	33	0.00055	0.99923
6	0.07613	0.61937	34	0.00046	0.99945
7	0.06344	0.68281	35	0.00038	0.99968
8	0.05286	0.73568	36	0.00032	0.99980
9	0.04045	0.77973	37	0.00027	0.99989
10	0.03671	0.81644	38	0.00022	0.99993
11	0.03059	0.84703	39	0.00019	0.99997
12	0.02549	0.87253	40	0.00015	0.99999
13	0.02125	0.89377	41	0.00013	0.99999
14	0.01770	0.91148	42	0.00011	0.99999
15	0.01475	0.92623	43	0.00009	0.99999
16	0.01229	0.93853	44	0.00007	0.99993
17	0.01025	0.94877	45	0.00006	0.99996
18	0.00854	0.95731	46	0.00005	0.99999
19	0.00711	0.96443	47	0.00004	0.99998
20	0.00593	0.97035	48	0.00004	0.99999
21	0.00494	0.97530	49	0.00003	0.99995
22	0.00412	0.97941	50	0.00002	0.99998
23	0.00343	0.98294	51	0.00002	0.99999
24	0.00286	0.98570	52	0.00002	0.99999
25	0.00238	0.98869	53	0.00001	0.99999
26	0.00199	0.99007	54	0.00001	0.99994
27	0.00165	0.99173	55	0.00001	0.99995

$$(a) P_{n \geq 4} = 1 - C P_3$$

$$= 1 - .34228 = .65772$$

$$(b) W_S = .06622 \text{ hour}$$

$$(c) L_q = 3.29 \text{ jobs}$$

$$(d) P_0 = .021 \Rightarrow 2.1\% \text{ idleness}$$

$$(e) \text{Av. # of idle computers} = 4 - (L_S - L_q)$$

$$= 4 - (6.62 - 3.29) = .67$$

$$\lambda = 15 + 10 + 20 = 45 \text{ customers/hour}$$

$$\mu = \frac{60}{6} = 10 \text{ customers/hour}$$

$$C > 45/10 = 4.5 \Rightarrow C \geq 5$$

8

Title: 15.6e-8
Comparative Analysis

Scenario	c	Lambda	Mu	L's eff	p0	Ls	Lq	Ws	Wq
1	5	45.00000	10.00000	45.00000	0.00496	11.36244	6.86244	0.25250	0.16250
2	6	45.00000	10.00000	45.00000	0.00464	5.76465	1.28646	0.16250	0.11251
3	7	45.00000	10.00000	45.00000	0.01046	4.89100	0.39100	0.10369	0.00969

$$(a) W_S \leq 15/60 = .25 \text{ hour} \Rightarrow C \geq 6$$

$$(b) \% \text{ idle} = \frac{C - (L_S - L_q)}{C} \times 100$$

C	Ls	Lq	C - (Ls - Lq)	% idle
5	11.362	6.862	.5	10%
6	5.765	1.265	1.5	25%

Select C = 5

C	5	6	7
P0	.00496	.00414	.01046

Select C ≤ 6

Set 15.6e

- Limited space inside a bank or a grocery store.
- Multiple queues appear to offer more courteous service.

9

For C parallel servers:

10

$$L_q = \frac{\rho}{C-\rho}, \text{ provided } \frac{\rho}{C} \rightarrow 1$$

Thus,

$$W_{q_c} = \frac{1}{\lambda_c} \frac{\rho}{C-\rho} = \frac{1}{(C\mu - \lambda_c)}$$

For a single server

$$W_{q_1} = \frac{\lambda_1}{\mu(\mu - \lambda_1)}$$

Because $\lambda_c = C\lambda_1$, we have

$$\begin{aligned} \frac{W_{q_c}}{W_{q_1}} &= \left(\frac{\frac{1}{C(C\mu - \lambda_1)}}{\frac{\lambda_1}{\mu(\mu - \lambda_1)}} \right) = \frac{1}{C(\frac{\lambda_1}{\mu})} \\ &= \frac{1}{C(\frac{\lambda c / \mu}{c})} \\ &= \frac{1}{C(S/c)} \end{aligned}$$

$$\lim_{\frac{\rho}{c} \rightarrow 1} \frac{W_{q_c}}{W_{q_1}} = \frac{1}{C}$$

Determination of p_0 involves the finite series sum

$$\sum_{n=c}^{\infty} \left(\frac{\rho}{c}\right)^{n-c} = \sum_{j=0}^{\infty} \left(\frac{\lambda}{\mu c}\right)^j$$

The series will diverge if $\lambda \geq \mu c$. The condition requires that customers be serviced at a rate faster than the rate at which they arrive at the facility. Else, the queue will build up to infinity.

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$$\begin{aligned} L_q &= \sum_{n=c}^{\infty} (n-c) p_n \\ &= \sum_{n=c}^{\infty} n p_n - c \sum_{n=c}^{\infty} p_n + \sum_{n=0}^{c-1} n p_n - \\ &\quad \sum_{n=0}^{c-1} n p_n + c \sum_{n=0}^{c-1} p_n - c \sum_{n=0}^{c-1} p_n \\ &= \sum_{n=0}^{\infty} n p_n - c \sum_{n=0}^{\infty} p_n + \sum_{n=0}^{c-1} (c-n) p_n \\ &= L_s - c + (\text{number of idle servers}) \\ &= L_s - \bar{C} \end{aligned}$$

12

Now, by definition

$$L_s = L_q + \frac{\lambda_{\text{eff}}}{\mu}$$

$$\text{It follows that } \bar{C} = \frac{\lambda_{\text{eff}}}{\mu}$$

13

$$p_n = \begin{cases} \frac{\lambda^n}{n! \mu^n} p_0, & n \leq c \\ \frac{\lambda^n}{c! c^{n-c} \mu^n} p_0, & n \geq c \end{cases}$$

for $c = 1$,

$$p_n = \begin{cases} \frac{\lambda}{\mu} p_0, & n=1 \\ \left(\frac{\lambda}{\mu}\right)^n p_0, & n \geq 1 \end{cases}$$

Thus,

$$p_n = \left(\frac{\lambda}{\mu}\right)^n p_0, \quad n=1, 2, \dots$$

14

$$\begin{aligned} L_q &= p_0 \frac{1}{c!} \sum_{n=c+1}^{\infty} (n-c) \frac{(\lambda/\mu)^n}{c^{n-c}} \\ &= p_0 \frac{(\lambda/\mu)^c}{c!} \sum_{n=c+1}^{\infty} (n-c) \left(\frac{\lambda}{\mu c}\right)^{n-c} \\ &= p_0 \frac{(\lambda/\mu)^c}{c!} \sum_{j=1}^{\infty} j \left(\frac{\lambda}{\mu c}\right)^j \\ &= p_0 \frac{(\lambda/\mu)^c}{c!} \frac{\lambda}{\mu c} \frac{d}{d(\lambda/\mu c)} \sum_{j=0}^{\infty} \left(\frac{\lambda}{\mu c}\right)^j \\ &= p_0 \frac{(\lambda/\mu)^c}{c!} \left\{ \frac{\lambda/\mu c}{(1 - \lambda/\mu c)^2} \right\} \\ &= p_0 \frac{\rho/c}{(1 - \rho/c)^2} = \frac{\rho}{(C-\rho)^2} p_0 \end{aligned}$$

15-24

(a) $P\{\text{a customer is waiting}\}$

$$\begin{aligned}
 &= P\{\text{at least } c+1 \text{ in system}\} \\
 &= \sum_{n=c+1}^{\infty} p_n \\
 &= \sum_{n=c}^{\infty} p_n - p_c \\
 &= p_0 \frac{\rho^c}{c!} \frac{1}{1-\frac{\rho}{c}} - p_c \\
 &= p_c \left\{ \frac{1}{1-\frac{\rho}{c}} - 1 \right\} \\
 &= p_c \left(\frac{\rho}{c-\rho} \right)
 \end{aligned}$$

(b) Expected number in queue given the queue is not empty

$$\begin{aligned}
 &= \sum_{i=c+1}^{\infty} (i-c) \frac{p_i}{\sum_{j=c+1}^{\infty} p_j} \\
 &= \frac{L_q}{\sum_{j=c+1}^{\infty} p_j} = \frac{L_q}{p_c \left(\frac{\rho}{c-\rho} \right)}
 \end{aligned}$$

$$\begin{aligned}
 \text{Now, } L_q &= \frac{p_0}{c!} \sum_{n=c+1}^{\infty} (n-c) \frac{\rho^n}{c^{n-c}} \\
 &= p_0 \frac{\rho^c}{c!} \sum_{j=1}^{\infty} j \left(\frac{\rho}{c}\right)^j \\
 &= p_0 \frac{\rho^c}{c!} \left(\frac{\rho/c}{(1-\rho/c)^2}\right), \frac{\rho}{c} < 1 \\
 &= p_c \left\{ \frac{c\rho}{(c-\rho)^2} \right\}, \frac{\rho}{c} < 1
 \end{aligned}$$

Substitution for L_q yield the desired result.(c) Exp. waiting time for those who must wait = Exp. waiting time given there are c in the system.

$$\begin{aligned}
 &= \frac{1}{\lambda} \sum_{i=c+1}^{\infty} (i-c) \frac{p_i}{\sum_{n=0}^{\infty} p_n} \\
 &= \frac{L_q/\lambda}{p_c/(1-\rho/c)} = \frac{1}{\mu(c-\rho)}
 \end{aligned}$$

15

16

First convert the c -channel case into an equivalent single channel. Let the customer just arriving be the j th in queue. Because there are c channels in parallel, the service time, t_i , of each of the other $j-1$ customers and the (one) customer in service are determined as follows: let t_1, t_2, \dots, t_c be the actual service times in the c channels. Then,

$$\begin{aligned}
 P\{t > T\} &= P\{\min_{1 \leq i \leq c} t_i > T\} \\
 &= (e^{-\mu c T})^c = e^{-\mu c T}
 \end{aligned}$$

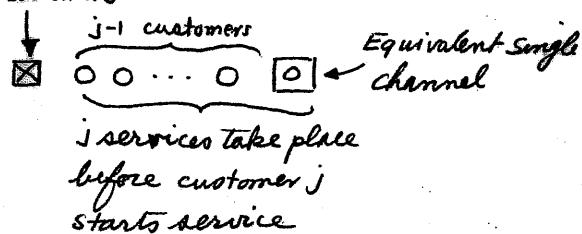
This is true because if $\min_i t_i > T$, then every t_i must be $> T$.

Now,

$$\begin{aligned}
 F_t(T) &= 1 - P\{t > T\} \\
 &= 1 - e^{-\mu c T}, \quad T > 0
 \end{aligned}$$

Thus,

$$f(t) = \frac{\partial F_t(T)}{\partial T} = \mu c e^{-\mu c T}, \quad T > 0$$

which is exponential with mean $\frac{1}{\mu c}$.The c channels can be converted into an equivalent single channel as customer j 

Before customer j starts service, j other customers each with a service time T must be processed first.

Continued...

Set 15.6e

The assumption here is that all c channels are busy. If there are any idle servers, arriving customer j will have zero waiting time in queue and the special case is treated separately.

Let τ be the waiting time in queue given there are j other customer yet to be serviced. Then

$$\tau = T_1' + T_2' + \dots + T_j'$$

where T_1', T_2', \dots, T_j' are exponential with mean $1/\mu c$. T_i' represents the remaining service time for the customer already in service. The lack of memory property indicate that T_i' is also exponential with mean $1/\mu c$. Thus,

$$W_q(\tau|j) = \frac{\mu c (\mu c \tau)^{j-1} e^{-\mu c \tau}}{(j-1)!}, \tau > 0$$

Let $W_q(\tau)$ be the absolute pdf, then

$$W_q(\tau) = \sum_{j=1}^{\infty} W_q(\tau|j) q_j.$$

where

$$q_j = \begin{cases} \sum_{k=0}^{c-1} p_k, & j=0 \\ p_{c+j-1}, & j>0 \end{cases}$$

Hence, for $\tau > 0$

$$\begin{aligned} W_q(\tau) &= \sum_{j=1}^{\infty} \frac{\mu c (\mu c \tau)^{j-1} e^{-\mu c \tau}}{(j-1)!} \frac{\rho^{c+j-1}}{c! c^{j-1}} p_0 \\ &= \frac{\rho^c \mu c e^{-\mu c \tau}}{c!} p_0 \sum_{j=0}^{\infty} \frac{(\rho \mu c \tau/c)^j}{j!} \\ &= \frac{\rho^c \mu c e^{-\mu c \tau}}{c!} p_0 e^{-\lambda \tau} \\ &= \frac{\rho^c \mu c e^{-\mu c(c-\rho)\tau}}{(c-1)!} p_0 \end{aligned}$$

continued...

For $\tau = 0$, the corresponding probability is $\sum_{k=0}^{c-1} p_k$, or

$$\begin{aligned} 1 - \sum_{k=c}^{\infty} p_k &= 1 - \sum_{j=0}^{\infty} p_{c+j} \\ &= 1 - \sum_{j=0}^{\infty} \frac{\rho^{c+j}}{c! c^j} p_0 \\ &= 1 - \frac{\rho^c}{c!} \left(\frac{p_0}{1 - \frac{\rho}{c}} \right) \\ &= 1 - \left\{ \frac{\rho^c p_0}{(c-1)! (c-\rho)} \right\} \end{aligned}$$

Hence,

$$W_q(\tau) = \begin{cases} 1 - \frac{\rho^c p_0}{(c-1)! (c-\rho)}, & \tau = 0 \\ \frac{\rho^c e^{-\mu(c-\rho)\tau}}{(c-1)!} p_0, & \tau > 0 \end{cases}$$

17

$$\begin{aligned} P\{\tau > y\} &= \int_y^{\infty} W_q(\tau) d\tau \\ &= \frac{c \mu \rho^c p_0}{c!} \int_y^{\infty} e^{-(c\mu - \lambda)\tau} d\tau \\ &= \frac{\rho^c c \mu}{c! (c\mu - \lambda)} e^{-(c\mu - \lambda)y} p_0 \\ &= \frac{\rho^c p_0}{c! (1 - \frac{\rho}{c})} e^{-(c\mu - \lambda)y} \\ &= P\{\tau > 0\} e^{-(c\mu - \lambda)y} \end{aligned}$$

where $P\{\tau > 0\} = 1 - P\{\tau = 0\}$

From Problem 16, the waiting time in the system is computed as

$$T = T_1 + T_2 + \dots + T_j + t_j$$

where

t_j = actual service time for customer j .

t_j is exponential with mean $1/\mu$

Thus, T is the convolution of the waiting time in queue and the actual service time of customer j . This means that $w(T)$ is the convolution of $w_q(\tau)$ and $g(t)$; that is,

$$w(T) = w_q(\tau) * g(t)$$

where

$$g(t) = \mu e^{-\mu t}, \quad t > 0$$

$$w(T) = w_q(0)g(T)$$

$$+ \int_{0^+}^T w_q(\tau)g(T-\tau)d\tau$$

$$= \left(1 - \frac{\rho^c p_0}{(c-1)! (c-p)}\right) \mu e^{-\mu T}$$

$$+ P_0 \int_{0^+}^T \frac{\mu p^c e^{-\mu(c-p)\tau}}{(c-1)!} \mu e^{-\mu(T-\tau)} d\tau$$

$$= \left(1 - \frac{\rho^c p_0}{(c-1)! (c-p)}\right) \mu e^{-\mu T}$$

$$+ \frac{\mu p^c e^{-\mu T}}{(c-1)! (c-1-p)} P_0 \left\{ 1 - e^{-\mu(c-1-p)T} \right\}$$

$$= \mu e^{-\mu T} \frac{\rho^c p_0 N e^{-\mu T}}{(c-1)! (c-1-p)} \frac{(c-p-1)}{(c-p)}$$

$$+ \frac{\mu p^c e^{-\mu T} P_0}{(c-1)! (c-1-p)} - \frac{\mu p^c e^{-\mu T} e^{-\mu(c-1-p)T}}{(c-1)! (c-1-p)}$$

Continued...

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$$= \mu e^{-\mu T} + \frac{\rho^c p_0 N e^{-\mu T}}{(c-1)! (c-1-p)} \left\{ \frac{1}{c-p} - e^{-\mu(c-1-p)T} \right\}$$

$T > 0$

Set 15.6f

(a) $C - (L_s - L_q) = 4 - (4.24 - 1.54)$
 $= 1.3 \text{ cars}$

(b) $P_q = .04468$

(c) Title: 6t-1
Comparative Analysis

Scenario	c	Lambda	Mu	U/ds eff	p ₀	L _s	L _q	W _s	W _q
1	4	16.00000	5.00000	15.2815	0.03121	4.2398	1.1542	0.27481	0.07481
2	4	16.00000	5.00000	15.2869	0.03238	4.0283	0.57450	0.26387	0.06387
3	4	16.00000	5.00000	15.2854	0.03393	3.7847	0.77903	0.25184	0.05184
4	4	16.00000	5.00000	15.2849	0.03513	3.51216	0.87078	0.23881	0.04881
5	4	16.00000	5.00000	14.24151	0.03931	3.02050	0.35719	0.22508	0.02508

m = length of waiting list

N = m + 4

m	N	W _q (hr)	W _q (min)
6	10	.075	4.5
5	9	.064	3.83
4	8	.052	3.12
3	7	.039	2.33
2	6	.025	1.5

Select m ≤ 3

C = 2, λ = 20/hr, N = 5

μ = 60/6 = 10/hr

(a) P₅ = .1818 or 18.18%

(b) P₁ = .1818 or 18.18%

(c) % utilization = 100 $\left(\frac{L_s - L_q}{C}\right)$
 $= \frac{2.727 - 1.091}{2} \times 100$
 $= 81.8\%$

(d) Probability = P₂ + P₃ + P₄ = .54546

(e) P_N ≤ .1

N	5	...	8	9	10
P _N	.1818		.1176	.1053	.0952

N ≥ 10 spaces (including the pumps)

(f) P₀ ≤ .05

N	5	...	8	9	10
P _N	.0909		.0588	.0526	.0476

N ≥ 10

λ = 60/10 = 6/hr

μ = 60/30 = 2/hr, N = 18

3

(a) # idle mechanics

$$= C - (L_s - L_q)$$

$$= 3 - (9.54 - 6.71) = .17$$

(b) P₁₈ = .0559

$$\lambda_{lost} = .0559 \times 6 = .3354 \text{ job/hr}$$

lost jobs in 10 hrs = 3.354 jobs

(c) P_{n≤17} = P₀ + P₁ + ... + P₁₇
 $= .9441$

(d) P_{n≤2} = P₀ + P₁ + P₂ = .10559

(e) L_q = 6.7081 mower

(f) $\frac{L_s - L_q}{C} = \frac{9.54 - 6.71}{3} = .944$

N = 40, C = 30, λ = 20/hr

μ = 60/60 = 1/hr

4

(a) P₄₀ = .00014

(b) P₃₀ + P₃₁ + ... + P₃₉ = P_{n≤39} - P_{n≤29}
 $= .99986 - .97533$
 $= .02453$

(c) P₂₉ = .01248

(d) L_s - L_q = 20.043 - 0.96 ≈ 20 spaces

(e) L_q = .046

continued...

continued...

(f) If there are 30 cars or more in the lot, the student will not make it to class. Thus,

$P\{\text{not finding a parking space}\}$

$$= P_{30} + P_{31} + \dots + P_{40} = 1 - P_{n \leq 29}$$

$$= 1 - .97533 = .02467$$

No. of students who cannot park during an 8-hr period = $20 \times .02467 \times 8$
 ≈ 4 students

$$1 = P_0 \left\{ \sum_{n=0}^{c-1} \frac{p^n}{n!} + \frac{p^c}{c!} \sum_{n=c}^N \left(\frac{p}{c}\right)^{n-c} \right\} \quad 5$$

$$= P_0 \left\{ \sum_{n=0}^{c-1} \frac{p^n}{n!} + \frac{p^c}{c!} \frac{1 - (\frac{p}{c})^{N-c+1}}{1 - \frac{p}{c}} \right\}$$

$$P_0 = \left\{ \sum_{n=0}^{c-1} \frac{p^n}{n!} + \frac{p^c}{c!} \left(\frac{1 - (\frac{p}{c})^{N-c+1}}{1 - \frac{p}{c}} \right) \right\}$$

$$\bar{C} = L_s - L_q$$

$$= \lambda_{\text{eff}} (W_s - W_q)$$

$$= \lambda_{\text{eff}} \left(\frac{1}{\mu} \right)$$

$$1 = \frac{P_0}{c!} \sum_{n=c}^N \frac{p^n}{c^{n-c}} + P_0 \sum_{n=0}^{c-1} \frac{p^n}{n!} \quad 7$$

$$= \frac{P_0 p^c}{c!} \sum_{n=0}^{N-c} \left(\frac{p}{c}\right)^n + P_0 \sum_{n=0}^{c-1} \frac{p^n}{n!}$$

$$= \frac{P_0 p^c}{c!} (N-c+1) + P_0 \sum_{n=0}^{c-1} \frac{p^n}{n!}$$

Thus,

$$P_0 = \left\{ \sum_{n=0}^{c-1} \frac{p^n}{n!} + \frac{p^c}{c!} (N-c+1) \right\}^{-1}$$

$$L_q = \sum_{n=c}^N (n-c) P_n$$

$$= \sum_{j=0}^{N-c} j P_{j+c}$$

$$= \frac{p}{c!} \frac{p^c}{c} \sum_{j=0}^{N-c} j \left(\frac{p}{c}\right)^{j-1} P_0$$

$$= \frac{p^c}{c!} \sum_{j=0}^{N-c} j P_0 \quad (\text{because } \frac{p}{c} = 1)$$

$$= \frac{p^c}{c!} \frac{(N-c)(N-c+1)}{2} P_0$$

$$= \frac{p^c (N-c)(N-c+1)}{2c!} P_0$$

8

$$\lambda_n = \begin{cases} \lambda, & n=0, 1, 2, \dots, c-1 \\ 0, & n=c \end{cases}$$

$$M_n = nM, \quad n=0, 1, \dots, c$$

Thus,

$$P_n = \frac{p^n}{n!} P_0, \quad n=0, 1, 2, \dots, c$$

$$\sum_{n=0}^c P_n = \sum_{n=0}^c \frac{p^n}{n!} P_0 = 1$$

$$P_0 = \left\{ \sum_{n=0}^c \frac{p^n}{n!} \right\}^{-1}$$

continued...

Set 15.6g

(a) $P_0 = 0$

1

(b) $P_{n \geq 10} = 1 - P_{n \leq 9} = 1$

(c) $P_{n \leq 40} - P_{n \leq 29} = .7771 - .13787$
 $= .63923$

(d) $L_S = 36$

Net annual equity

$$= \$1000 \times 36 \{ .1(1-.3) + .9(1+.15) \}$$

$$= \$39,780$$

$\lambda = \frac{100}{8} = 12.5 / \text{hr}$

2

$M = \frac{60}{30} = 2 / \text{hr}$

(a) $L_S = 6.25 \approx 7 \text{ seats}$

(b) $P_{n \geq 8} = 1 - (P_0 + P_1 + \dots + P_7)$
 $= 1 - .7089 = .2911$

(c) $P_0 = .00193$

$P = .1$

3

c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
2	1.00000	10.00000	1.00000	0.90476	0.10025	0.00025	0.10025	0.00025
4	1.00000	10.00000	1.00000	0.90484	0.10000	0.00000	0.10000	0.00000
10	1.00000	10.00000	1.00000	0.90484	0.10000	0.00000	0.10000	0.00000
20	1.00000	10.00000	1.00000	0.90484	0.10000	0.00000	0.10000	0.00000
50	1.00000	10.00000	1.00000	0.90484	0.10000	0.00000	0.10000	0.00000
9999	1.00000	10.00000	1.00000	0.90484	0.10000	0.00000	0.10000	0.00000

$P = .9$

4

c	λ	M	λ_{eff}	P_0	L_s	L_q	W_s	W_q
10	9.00000	1.00000	9.00000	0.00007	15.01858	6.01858	1.66873	0.66873
15	9.00000	1.00000	9.00000	0.00012	9.07235	0.07235	1.00004	0.00004
25	9.00000	1.00000	9.00000	0.00012	9.00000	0.00000	1.00000	0.00000
50	9.00000	1.00000	9.00000	0.00012	9.00000	0.00000	1.00000	0.00000
9999	9.00000	1.00000	9.00000	0.00012	9.00000	0.00000	1.00000	0.00000

- For very small P_0 , $(M/M/\infty) : (GD/\alpha_{\text{eff}})$ provides reliable estimates for $(M/M/c) : (GD/\alpha/\alpha)$.
- For large P_0 , $(M/M/\infty)$ gives reliable estimates only if c is large

Set 15.6h

$$(a) R=1: \lambda_{eff} = \lambda(22-L_S) \\ = .5(22 - 12.004) \\ = 4.998$$

$$R=4: \lambda_{eff} = .5(22 - 2.1) = 9.95$$

$$(b) \text{No. of idle repair persons} \\ = 4 - (L_S - L_q) \\ = 4 - (2.1 - .11) = 2.01$$

$$(c) P_0 = .10779$$

$$(d) R = 3:$$

$$P\{2 \text{ or } 3 \text{ are idle}\} = P_0 + P_1 \\ = .34492$$

Title: 6h-1
Scenario 3-(MM/3):(GD/22/22)

Λ = 0.50000	M_u = 5.00000
Λ_{eff} = 9.76696	Rho/c = 0.03333
L_s = 2.46596	L_q = 0.51257
W_s = 0.25248	W_q = 0.05248

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.10779	0.10779	8	0.00953	0.99244
1	0.23713	0.34492	9	0.00445	0.99689
2	0.24899	0.59380	10	0.00193	0.99881
3	0.16599	0.75980	11	0.00077	0.99959
4	0.10513	0.86502	12	0.00028	0.99987
5	0.06308	0.92810	13	0.00009	0.99996
6	0.03574	0.96384	14	0.00003	0.99999
7	0.01906	0.98291			

Productivity of repair persons

= Av. # busy repair persons

R	Repair prod.	Shop prod.
1	100%	45.44%
2	88.2%	80.15%
3	65.1%	88.7%
4	49.7%	90.45%

$R=2$ yield 80.15% shop productivity and also maintain repair productivity at 88.2%

Increasing R , in effect, increases the number of machines that remain operative, and hence the chance of additional breakdowns. Stated differently, if all machines remain broken, there will be no new calls for repair service, and $\lambda_{eff} = 0$.

3

$$\lambda = \frac{60}{45} = 1.33 \text{ machines/hr}$$

$$M = \frac{60}{8} = 7.5 \text{ machines/hr}$$

$$R=1, K=5$$

Title: 6h-4
Scenario 1-(MM/1):(GD/5/5)

Λ = 1.33333	M_u = 7.50000
Λ_{eff} = 4.99939	Rho/c = 0.17778
L_s = 1.25045	L_q = 0.58386
W_s = 0.25012	W_q = 0.11679

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.33341	0.33341	3	0.11240	0.95293
1	0.29637	0.62978	4	0.03996	0.99290
2	0.21075	0.84053	5	0.00710	1.00000

$$(a) L_s = 1.25 \text{ machines}$$

$$(b) P_0 = .33341$$

$$(c) W_s = .25 \text{ hour}$$

$$\lambda = 60/45 = 1.33/\text{hr}$$

$$M = 60/20 = 3/\text{hr}$$

$$R=4, K=4$$

Title: 6h-5
Scenario 1-(MM/4):(GD/4/4)

Λ = 1.33333	M_u = 3.00000
Λ_{eff} = 3.69230	Rho/c = 0.11111
L_s = 1.23077	L_q = 0.00000
W_s = 0.33333	W_q = 0.00000

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.22972	0.22972	3	0.08067	0.99104
1	0.40839	0.63811	4	0.00896	1.00000
2	0.27226	0.91037			

$$(a) L_s = 1.23 \text{ workers}$$

$$(b) P_0 = .22972$$

Set 15.6h

$$\lambda = \frac{60}{30} = 2 \text{ calls/hr/baby}$$

6

$$\mu = \frac{60}{120} = .5 \text{ /hr}$$

$$R = 5, K = 5$$

Title: 6h-6
Scenario 1- (M/M/5); (GD/5/5)

Lambda = 2.00000	Mu = 0.50000
Lambda eff = 2.00000	Rho/c = 0.80000
Ls = 4.00000	Lq = 0.00000
Ws = 2.00000	Wq = 0.00000

n	Probability, pn	Cumulative, Pn	n	Probability, pn	Cumulative, Pn
0	0.00032	0.00032	3	0.20480	0.26272
1	0.00640	0.00672	4	0.40960	0.67232
2	0.05120	0.05792	5	0.32768	1.00000

(a) No. "awake" babies

$$= 5 - L_S = 5 - 4 = 1 \text{ baby}$$

$$(b) p_5 = .32768$$

$$(c) P_{n \leq 2} = p_0 + p_1 + p_2 = .05792$$

$$\bar{R} = L_S - L_Q$$

8

$$= \lambda_{\text{eff}} (W_S - W_Q)$$

$$= \lambda_{\text{eff}} \left(\frac{1}{\mu} \right)$$

$$\text{Hence } \lambda_{\text{eff}} = \mu \bar{R}$$

$$P_n = \begin{cases} C_n^k \rho^n n! p_0, & n=0,1 \\ C_n^k n! \rho^n p_0, & n=1,2,\dots,K \\ \frac{K!}{(K-n)!} \rho^n p_0, & n=0,1,2,\dots,K \end{cases}$$

$$L_S = \sum_{n=0}^K n p_n = p_0 K! \sum_{n=0}^K \frac{n \rho^n}{(K-n)!}$$

$$= K - \left(\frac{1-p_0}{\rho} \right)$$

7

$$P_n = \begin{cases} \frac{K\lambda}{\mu} \frac{(K-1)\lambda}{2\mu} \dots \frac{(K-n)\lambda}{n\mu} p_0, & 0 \leq n \leq R \\ \frac{K\lambda}{\mu} \frac{(K-1)\lambda}{2\mu} \dots \frac{(K-R)\lambda}{R\mu} \dots \frac{K-n}{R\mu} p_0, & R \leq n \leq K \end{cases}$$

Thus,

$$P_n = \begin{cases} \frac{K(K-1)\dots(K-n)}{1 \times 2 \times \dots \times n} \left(\frac{\lambda}{\mu}\right)^n p_0, & 0 \leq n \leq R \\ \frac{C_n^k n!}{R! R^{n-R}} \left(\frac{\lambda}{\mu}\right)^n p_0, & R \leq n \leq K \end{cases}$$

$$= \begin{cases} C_n^k \rho^n p_0, & 0 \leq n \leq R \\ C_n^k \frac{n! \rho^n}{R! R^{n-R}} p_0, & R \leq n \leq K \end{cases}$$

Set 15.7a

$$\begin{aligned}\% \text{idle} &= \frac{1 - (L_s - L_q)}{1} \times 100 \\ &= [1 - (L_s - L_q)] \times 100 \\ &= (1 - 1.333 + .667) \times 100 \\ &= 33.3\%\end{aligned}$$

1 $\lambda = .3 \text{ job/day}$

3

Service time distribution:

$$f(t) = .5, \quad 2 \leq t \leq 4 \text{ days}$$

$$E\{t\} = 3 \text{ days}$$

$$\text{Var}\{t\} = \frac{4}{12} = .333 \text{ days}^2$$

$$(a) L_q = 4.2 \text{ homes}$$

$$(b) W_s = 17 \text{ days}$$

$$(c) E\{t\} = 1.5, \text{ Var}\{t\} = \frac{1}{12} = .0833$$

$$L_q = .191 \text{ home}$$

$$W_s = 2.14 \text{ days}$$

2 (a) $E\{t\} = 14 \text{ min}$

$$\text{Var}\{t\} = \frac{(20-8)^2}{12} = 12 \text{ min}^2$$

$$\lambda = 4/\text{hr} = .0667/\text{min}$$

$$L_s = 7.867 \text{ cars}$$

$$W_s = 118 \text{ min} = 1.967 \text{ hours}$$

$$L_q = 6.933 \text{ cars}$$

$$W_q = 104 \text{ min} = 1.733 \text{ hours}$$

(b) $E\{t\} = 12 \text{ min}$

$$\text{Var}\{t\} = 9 \text{ min}^2$$

$$\lambda = .0667/\text{min}$$

$$L_s = 2.5 \text{ cars}$$

$$W_s = 37.5 \text{ min} = .625 \text{ hours}$$

$$L_q = 1.7 \text{ cars}$$

$$W_q = 25.5 \text{ min} = .425 \text{ hr}$$

(c) $E\{t\} = 4x.2 + 8x.6 + 15x.2 = 8.6 \text{ min}$

$$\text{Var}\{t\} = (4-8.6)^2(.2) + (8-8.6)^2(.6)$$

$$+ (15-8.6)^2(.2) = 12.64 \text{ min}^2$$

$$L_s = 1.0244 \text{ cars}$$

$$W_s = 15.3657 \text{ min} = .256 \text{ hr}$$

$$L_q = .451 \text{ car}$$

$$W_q = 6.765 \text{ min} = .113 \text{ hr}$$

4 $\lambda = \frac{30}{8 \times 60} = .0625 \text{ prescr./min}$

$$E\{t\} = 12 + 3 = 15 \text{ min}$$

$$\text{Var}\{t\} = 9 + \frac{(4-2)^2}{12} = 9.333 \text{ min}^2$$

(a) $P_0 = .0625$

(b) $L_q = 7.3 \text{ prescriptions}$

(c) $W_s = 132.17 \text{ min} = 2.2 \text{ hours}$

5 $\lambda = 1/45 / \text{min} = .0222 / \text{min}$

$$E\{t\} = 28 + 4.5 = 32.5 \text{ min}$$

$$\text{Var}\{t\} = \frac{(6-3)^2}{12} = .75$$

(a) $L_q = .9395 \text{ item}$

(b) $P_0 = .278$

(c) $W_s = 74.78 \text{ min}$

6 $L_s = \lambda E\{t\} + \frac{\lambda^2 (E^2\{t\} + \text{Var}\{t\})}{2(1-\lambda E\{t\})}$

$$= \lambda E\{t\} + \frac{(\lambda E\{t\})^2}{2(1-\lambda E\{t\})}$$

$$= P + \frac{P^2}{2(1-P)}$$

Set 15.7a

7

$$\begin{aligned}
 L_S &= \frac{m\lambda}{\mu} + \frac{\lambda^2}{2(1-\frac{m\lambda}{\mu})} \left(\frac{m^2}{\mu^2} + \frac{n}{\mu^2} \right) \\
 &= mp + \frac{m^2 p^2 + mp^2}{2(1-mp)} \\
 &= mp + \frac{m(m+1)p^2}{2(1-mp)}
 \end{aligned}$$

8

$$\begin{aligned}
 E\{t\} &= \frac{1}{\mu}, \quad \text{Var}\{t\} = \frac{1}{\mu^2} \\
 L_S &= \frac{\lambda}{\mu} + \frac{\lambda^2 \left(\frac{1}{\mu^2} + \frac{1}{\mu^2} \right)}{2(1-\lambda/\mu)} \\
 &= p + \frac{p^2}{1-p} \\
 &= \frac{p}{1-p}
 \end{aligned}$$

9

(a) Because each server receives every c^{th} customer and the interarrival time at the channel is exponential with mean $1/\lambda$, the interarrival time at each server is the convolution of c exponential distributions each with mean $\frac{1}{\lambda}$. This means that the interarrival time is gamma with mean c/λ and variance c/λ^2 .

(b) The interarrival time at the i^{th} server is exponential with mean $\frac{1}{\alpha_i \lambda}$. This means that the arrivals at server i is Poisson with mean $\alpha_i \lambda$, $i=1, 2, \dots, c$

$$(a) M_2 = \frac{24}{(\frac{1000}{36}) \times \frac{1}{60}} = 5.184 \text{ jobs/day}$$

$$M_3 = \frac{24}{(\frac{1000}{50}) \times \frac{1}{60}} = 7.2 \text{ jobs/day}$$

$$M_4 = \frac{24}{(\frac{1000}{66}) \times \frac{1}{60}} = 9.5 \text{ jobs/day}$$

$$(b) ETC_i = 24 C_{ii} + 80 Lq_i$$

i	λ_i	M_i	Lq_i	C_{ii}	ETC_i
1	4	4.32	11.57	\$15	\$1285.60
2	4	5.18	2.62	20	689.60
3	4	7.20	.69	24	631.20
4	4	9.50	.31	27	672.80

Select model 3.

$$\lambda = 3 \text{ /hr}$$

$$M_1 = 5 \text{ /hr}, \quad C_1 = \$15$$

$$M_2 = 8 \text{ /hr}, \quad C_2 = \$20$$

Cost/Broken machine = \$50/hr

(M/M/1) : (GD/10/10):

$$\lambda = 3, \mu = 5 \Rightarrow L_S = 8.33$$

(M/M/1) : (GD/10/10):

$$\lambda = 3, \mu = 8 \Rightarrow L_{S_2} = 7.33$$

$$TC_1 = 50L_{S_1} + 15 = 50 \times 8.33 + 15 \\ = \$431.50/\text{hr}$$

$$TC_2 = 50L_{S_2} + 20 = 50 \times 7.33 + 20 \\ = \$386.50/\text{hr}$$

Here second repair person.

$$\lambda = 10/\text{hr} = .167/\text{min}$$

Scanner A:

Service time distribution:

$$f_A(t) = \frac{1}{(\frac{35}{10}) - (\frac{25}{10})} = 1, 2.5 \leq t \leq 3.5$$

Continued...

$$E_A\{t\} = 3 \text{ min}$$

$$\text{Var}_A\{t\} = \frac{1}{12} \text{ min}^2$$

Scanner B:

$$f_B(t) = \frac{1}{\frac{35}{15} - \frac{25}{15}} = 1.5, \quad \frac{5}{3} \leq t \leq \frac{7}{3}$$

$$E_B\{t\} = 2 \text{ min}$$

$$\text{Var}_B\{t\} = \frac{(2/3)^2}{12} = \frac{1}{27} \text{ min}^2$$

From Excel file PKFormula.xls,

$$L_{SA} = .755 \text{ customer}$$

$$L_{SB} = .419 \text{ customer}$$

$$TC_A = .2L_{SA} + C_A \\ = (-.2 \times .755 + \frac{10}{10 \times 60}) \times 60 = \$10.06/\text{hr}$$

$$TC_B = .2L_{SB} + C_B \\ = (-.2 \times .419 + \frac{15}{10 \times 60}) \times 60 = \$6.53/\text{hr}$$

Select scanner B

(a) $M = \text{number of filled orders/hr}$

$\lambda = \text{number of requested orders/hr}$

$C_1 = \text{cost/unit increase in production rate}$

$C_2 = \text{cost of waiting/unit waiting time/cust.}$

$TC(M) = \text{Total cost/unit waiting time}$
given μ

$$= C_1 \mu + C_2 L_S$$

$$= C_1 \mu + C_2 \frac{\lambda}{M-\lambda}$$

$$\frac{\partial TC(M)}{\partial M} = C_1 - C_2 \frac{\lambda}{(M-\lambda)^2} = 0$$

$$M = \lambda + \sqrt{\frac{C_2}{C_1} \lambda}$$

$$(c) \lambda = 3, C_1 = .1 \times 500 = \$50, C_2 = \$100$$

$$M = 3 + \sqrt{\frac{100}{50} \times 3} = 5.45 \text{ orders/hr}$$

Optimum production rate

$$= 500 \times 5.45 \approx 2725 \text{ pieces/hr}$$

Set 15.9a

$$\lambda = 80 \text{ jobs/wk}$$

$$C_1 = \$250/\text{wk} \quad C_2 = \$500/\text{job/wk}$$

$$M = \lambda + \sqrt{\frac{C_2}{C_1}} \lambda$$

$$= 80 + \sqrt{\frac{500}{250} \times 80} = 92.65 \text{ jobs/wk}$$

5

$$\lambda = 25 \text{ groups/hr}$$

$$\underline{\text{Model A: }} M = 26/\text{hr}, N = 20$$

$$\text{Operating cost } C_A = \$12000/\text{month}$$

$$\text{From TORA: } P_{20} = .03128$$

$$L_q = 7.65 \text{ groups}$$

6

$$\text{Cost/hr} = \text{operating cost/hr} + \text{waiting cost/hr} + \text{cost of lost customers/hr}$$

$$= \frac{C_A}{30 \times 10} + 10L_q + \lambda P_N \times 15$$

$$= \frac{12000}{30 \times 10} + 10 \times 7.65 + 25 \times .03128 \times 15$$

$$= \$128.23/\text{hr}$$

$$\underline{\text{Model B: }} M = 29/\text{hr}, N = 30$$

$$C_B = \$16000/\text{month}$$

$$\text{From TORA: } P_{30} = .0016$$

$$L_q = 5.07 \text{ groups}$$

$$\text{Cost/hr} = \frac{\$16000}{30 \times 10} + 10 \times 5.07 + 25 \times .0016 \times 15$$

$$= \$104.63$$

Select model B

Let

$$C_3 = \text{cost/unit time/additional capacity unit.}$$

The cost model in Problem 6 is modified by adding the term $C_3 N$ to the cost equation.

7

P_0 is the probability of running out of stock. Thus,

8

$$\begin{aligned} \text{Cost of lost sales per hour} &= C_1 \lambda P_0 \\ E\{\text{cost}\}/\text{unit time} &= E\{\text{lost sales cost}\}/\text{unit time} \\ &\quad + E\{\text{holding cost}\}/\text{unit time} \\ &= C_1 \lambda P_0 + C_2 L_S \end{aligned}$$

For (M/M/1): (GD/oo/o)

$$P_0 = (1 - \rho)$$

$$L_S = \frac{\rho}{1 - \rho}$$

Thus,

$$E\{\text{cost}\}/\text{unit time} = C_1 \lambda (1 - \rho) + C_2 \frac{\rho}{1 - \rho}$$

$$\frac{\partial E\{\text{cost}\}}{\partial \rho} = -C_1 \lambda + \frac{C_2}{(1 - \rho)^2} = 0$$

Thus,

$$\rho = 1 \pm \sqrt{\frac{C_1 \lambda}{C_2}}$$

Under steady state, ρ must be less than 1. Thus,

$$\rho = 1 - \sqrt{\frac{C_1 \lambda}{C_2}}$$

The solution requires $\sqrt{\frac{C_1 \lambda}{C_2}} < 1$ in order for ρ not to assume a negative value. Note that

$$\rho = \frac{\lambda}{M}, \text{ where } M \text{ is a constant.}$$

This means that M is the actual optimization variable.

Set 15.9b

$$C_1 = \$20, C_2 = \$45,$$

$$\lambda = 17.5/\text{hr}, \mu = 10/\text{hr}$$

Title: 15.9b-1 (M/M/c/GD/infinity/infinity)
Comparative Analysis

Scenario	c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
1	2	17.50000	10.00000	17.50000	0.00000	7.49597	5.71667	0.45667	0.32667
2	3	17.50000	10.00000	17.50000	0.15654	1.47172	0.45712	0.22222	0.00000
3	4	17.50000	10.00000	17.50000	0.17038	1.84206	0.09206	0.10528	0.00628
4	5	17.50000	10.00000	17.50000	0.17314	1.76982	0.01862	0.10112	0.00112

$$ETC(c) = 20c + 45L_s$$

C	$L_s(c)$	$ETC(c)$
2	7.467	$20 \times 2 + 45 \times 7.467 = \376.07
→ 3	2.217	$20 \times 3 + 45 \times 2.217 = \159.77
4	1.842	$20 \times 4 + 45 \times 1.842 = \162.89
5	1.770	$20 \times 5 + 45 \times 1.770 = \179.65

Use three clerks

$$\text{Cost/hr} = C_1 L_s + C_2 c$$

$$C_1 = \$30, C_2 = \$18$$

$$(M/M/c):(GD/10/10): \lambda = 1/20 = 0.05/\text{hr}$$

$$\mu = 1/3 = 0.333/\text{hr}$$

Title: 15.9b-2
Comparative Analysis

Scenario	c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
1	2	0.05000	0.33300	0.41603	0.21439	1.67942	0.43010	4.03663	1.03683
2	3	0.05000	0.33300	0.43167	0.24269	1.36246	0.06554	3.15478	0.15178

$$(\text{Cost/hr for } c=2) = 30 \times 1.68 + 18 \times 2 = \$86.40$$

$$(\text{Cost/hr for } c=3) = 30 \times 1.36 + 18 \times 3 = \$94.80$$

(a) No, because the cost is higher

(b) Schedule loss/breakdown = $C_1 W_s$

$$C=2: W_s = 4.037 \text{ hours}$$

$$\text{Schedule loss} = 30 \times 4.037 = \$121.11$$

$$C=3: W_s = 3.155 \text{ hours}$$

$$\text{Schedule loss} = 30 \times 3.155 = \$94.65$$

The problem is similar to the machine repair model. The executives are the "machines" and the WATS line is the "server".

Arrival rate/executive = 2 calls/day

$$\text{Service rate} = \frac{480}{6}$$

$$= 80 \text{ calls/day}$$

Continued...

TOVA input:

$$R=1: (2, 80, 1, 100, 100)$$

$$R=2: (2, 80, 2, 100, 100)$$

Title: 15.9b-3
Comparative Analysis

Scenario	c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
1	2	2.00000	80.00000	80.00000	0.00200	59.99881	59.99881	0.74999	0.73249
2	2	2.00000	80.00000	129.28000	0.00200	20.33920	18.36920	0.12782	0.11532

(a) No WATS:

$$\begin{aligned} \text{Cost/month} &= (2 \text{ calls}/8 \text{ hrs/exec}) \times \\ &(100 \text{ exec}) \times (6 \text{ min/call}) \times \\ &(50 \text{¢/min}) \times (200 \text{ hrs/month}) \\ &= \$15,000/\text{month} \end{aligned}$$

One WATS Line: $L_q = 59$

$\text{Cost/month} = \text{Cost of WATS line} +$

$$\begin{aligned} C_1 L_q \\ = \$2000/\text{month} + 59 \left(\frac{14}{100} \times 60 \times 200 \right) \\ = \$9080 \end{aligned}$$

$$\text{Savings} = 15,000 - 9080$$

$$= \$5920/\text{month}$$

(b) Two WATS lines: $L_q = 18.4$

$$\begin{aligned} \text{Cost/month} &= 2 \times 2000 + \\ &18.4 \left(\frac{14}{100} \times 200 \times 60 \right) \\ &= \$6200 \end{aligned}$$

Additional savings

$$= 9080 - 6200 = \$2880$$

Lease a second WATS line

Set 15.9b

Rate of breakdown/machine, λ

$$= \frac{57.8}{8 \times 20} = .36125/\text{hr}$$

$$\mu = \frac{60}{6} = 10/\text{hr}$$

TORA model: $(M/M/3):(GD/20/20)$

W_s = lost time per breakdown

λ = number of breakdowns/hr/mach

lost time per mach/hr = λW_s

From TORA, $W_s = .10118$ hr

Lost revenue/machine/hr

$$= 25 \times (.36125 \times .10118) \times \$2 \\ = \$1.83$$

Lost revenue for all machines

$$= 20 \times 1.83 = \$36.50$$

Cost of 3 repairpersons/hr

$$= 3 \times 20 = \$60.$$

$$TC(c) = C_1 + C_2 L_s(c)$$

$$TC(c-1) = (c-1) C_1 + C_2 L_s(c-1)$$

$$TC(c+1) = (c+1) C_1 + C_2 L_s(c+1)$$

$$TC(c-1) - TC(c) \\ = -C_1 + C_2 \{L_s(c-1) - L_s(c)\}$$

$$TC(c+1) - TC(c) \\ = C_1 - C_2 \{L_s(c) - L_s(c+1)\}$$

At a minimum point, we must have

$$TC(c-1) \geq TC(c)$$

$$TC(c+1) \geq TC(c)$$

Thus,

$$L_s(c-1) - L_s(c) \geq \frac{C_1}{C_2}$$

$$L_s(c) - L_s(c+1) \leq \frac{C_1}{C_2}$$

4

or

$$L_s(c) - L_s(c+1) \leq \frac{C_1}{C_2} \leq L_s(c-1) - L_s(c)$$

$$\frac{C_1}{C_2} = \frac{12}{50} = .24$$

C	$L_s(c)$	$L_s(c) - L_s(c+1)$	
2	7.467	-	
3	2.217	5.25	
4	1.842	.375	$\leftarrow \frac{C_1}{C_2} = .24$
5	1.764	.078	

$$C^* = 4$$

5

continued...

$$\lambda = 1/7 = .1428 \text{ breakdown/hr}$$

$M = .95$ repair per hour

TORA model: $(M/M/R) : (G.D/10/10)$

Comparative Analysis

Scenario	c	Lambda	Mu	L'da eff	p0	Ls	Lq	Ws	Wq
1	1	0.14280	0.25000	0.25000	0.00001	8.24932	7.24934	32.99772	28.99772
2	1	0.14280	0.25000	0.25000	0.00001	8.24932	7.24934	32.99772	28.99772
3	3	0.14280	0.25000	0.71071	0.00548	5.02302	2.18017	7.06758	3.06758
4	4	0.14280	0.25000	0.83618	0.00949	4.14443	0.79972	4.95641	0.95641
5	5	0.14280	0.25000	0.89109	0.01168	3.81048	0.20000	4.25887	0.25887
6	6	0.14280	0.25000	0.90407	0.01081	3.66859	0.05272	4.05931	0.05931
7	7	0.14280	0.25000	0.90807	0.01089	3.64098	0.00867	4.00955	0.00955
8	8	0.14280	0.25000	0.90878	0.01091	3.63662	0.00091	4.00100	0.00100

(a) From TORA's output

$$L_s < 4 \Rightarrow R \geq 5$$

(b) From TORA's output

$$W_q < 1 \Rightarrow R \geq 4$$

$$C_1 = \$12$$

2

C	L_s
2	7.467
3	2.217
4	1.842

$$2.217 - 1.842 \leq \frac{12}{C_2} \leq 7.467 - 2.217$$

$$.375 \leq \frac{12}{C_2} \leq 5.25$$

or

$$\$2.29 \leq C_2 \leq \$32$$

Chapter 16

Simulation Modeling

Set 16.1a

R1	R2	X	Y	$(X-1)^2 + (Y-2)^2 \leq 1$	in?
0.0589	0.6733	3.411	3.733	22.46021	1
0.4799	0.9486	0.799	6.486	20.164597	1
0.6139	0.5933	2.139	2.933	2.16781	1
0.9341	0.1782	5.341	-1.218	29.199805	0
0.3473	0.5644	-0.527	2.644	2.746465	1
0.3529	0.3646	-0.471	0.646	3.997157	1
0.7676	0.8931	3.676	5.931	22.613737	1
0.3919	0.7876	-0.061	4.876	9.439937	1
0.5199	0.6358	1.199	3.358	1.883765	1
0.7472	0.8954	3.472	5.954	21.7449	1
				Total=	9
				Area estimate=	90

Exact area = π cm². Estimate from Figure 18-2 = 78.56 cm² for a sample size of n=30,000. Current estimate = 90 cm², which is unreliable because the sample size is too small.

1

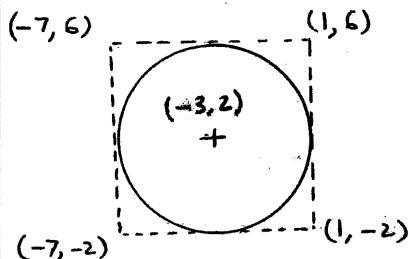
3

$$(a) X = -7 + 8R_1$$

$$Y = -2 + 8R_2$$

$$f(x) = \frac{1}{8}, \quad -7 \leq x \leq 1$$

$$f(y) = \frac{1}{8}, \quad -2 \leq y \leq 6$$



(b)

Monte Carlo Estimation of the Area of a Circle	
Input Data	
No. Replications, N =	10
Sample size, n =	100,000 Steps = 1
Radius, r =	4
Center, cx =	-3
Center, cy =	2
Output results	
Exact area =	50.265
Press to Execute Monte Carlo:	
Monte Carlo Calculations:	
n=100000	
Replication 1	50.223
Replication 2	50.378
Replication 3	50.113
Replication 4	50.260
Replication 5	50.244
Replication 6	50.330
Replication 7	50.327
Replication 8	50.252
Replication 9	50.236
Replication 10	50.467
Mean =	50.283
Std. Deviation =	0.099
95% lower conf. limit =	50.212
95% upper conf. limit =	50.354

2

R1	R2	X	Y	in?
0.0589	0.6733	4.123	2.6932	No
0.4799	0.9486	3.3593	3.7944	Yes
0.6139	0.5933	4.2973	2.3732	Yes
0.9341	0.1782	6.5387	.7128	No
0.3473	0.5644	2.4311	2.9576	Yes
0.3529	0.3646	2.4703	1.4584	Yes
0.7676	0.8931	5.3732	3.5724	No
0.3919	0.7876	2.7433	3.1504	Yes
0.5199	0.6358	3.6393	2.5432	No
0.7472	0.8954	5.2304	3.5816	No

points in = 5

$$\text{Area estimate} = \frac{5}{10} \times (4 \times 7) = 14 \text{ miles}^2$$

(4) $P\{H\} = .5 \quad P\{T\} = .5$
 If $0 \leq R \leq .5$, Juri gets \$10
 $.5 < R \leq 1$, Jan gets \$10

4

R	Jan's pay	R	Jan's pay
0.0589	-10	0.3529	-10
0.6733	10	0.3646	-10
0.4799	-10	0.7676	10
0.9486	10	0.8931	10
0.6139	10	0.3919	-10
0.5933	10	0.7876	10
0.9341	10	0.5199	10
0.1782	-10	0.6358	-10
0.3473	-10	0.7472	10
0.5644	10	0.8954	10

$$\bar{X}_1 = \$2 \quad \bar{X}_2 = \$4$$

continued...

Set 16.1a

<u>R</u>	Jan's pay	<u>R</u>	Jan's pay	<u>R</u>	Jan's pay
.5861	10	.3455	-10	.7900	10
.1281	-10	.4871	-10	.7698	10
.2867	-10	.8111	10	.2871	-10
.8216	10	.8912	10	.9534	10
.3866	-10	.4291	-10	.1394	-10
.7125	10	.2302	-10	.9025	10
.2108	-10	.5423	10	.1605	-10
.3575	-10	.4208	-10	.3567	-10
.2926	-10	.6975	10	.3670	-10
.8261	10	.5954	10	.5513	10

$$\bar{X}_3 = -\$2$$

$$\bar{X}_4 = \$0$$

$$\bar{X}_5 = \$0$$

(b) Av. Jan's pay based on 5 rep'ts.

$$= 2 + 4 - 2 + 0 + 0$$

$$= \$.8$$

$$S = \sqrt{\frac{(2 - .8)^2 + (4 - .8)^2 + (-2 - .8)^2 + 2(0 - .8)^2}{5-1}}$$

$$= \sqrt{\frac{80.8}{4}} = 2.28$$

Confidence interval:

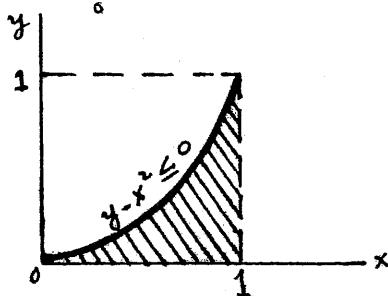
$$.8 - \frac{2.28}{\sqrt{5}} t_{.025, 4} \leq \mu \leq .8 + \frac{2.28}{\sqrt{5}} t_{.025, 4}$$

Given $t_{.025, 4} = 2.776$, the 95% confidence interval is

$$-2.03 \leq \mu \leq 3.63$$

(c) Theoretical Jan's payoff = \$0.

Estimate $\int_0^1 x^2 dx$



Continued...

5

(a)

Let $x=R1$ and $y=R2$.

Experiment: If $R2 < R1^2$, count point "in".

Estimate of integral = $(1x1)(\text{Points "in"})/5$

(b)

	R1	R2	1=in, 0=out
--	----	----	-------------

Rep 1	0.0589	0.6733	0
	0.4799	0.9486	0
	0.6139	0.5933	0
	0.9341	0.1782	1
	0.3473	0.5644	0

Integral estimate = 0.2

Rep 2	0.3529	0.3646	0
	0.7676	0.8931	0
	0.3919	0.7876	0
	0.5199	0.6358	0
	0.7472	0.8954	0

Integral estimate = 0

Rep 3	0.5869	0.1281	1
	0.2867	0.8216	0
	0.8261	0.3866	1
	0.7125	0.2108	1
	0.3575	0.2926	0

Integral estimate = 0.6

Rep 4	0.3455	0.4871	0
	0.8111	0.8912	0
	0.4291	0.2302	0
	0.5954	0.5423	0
	0.4208	0.6975	0

Integral estimate = 0

overall integral estimate = 0.2

Std. Deviation = 0.244949

95% lower confidence limit = -0.189714

95% upper confidence limit = 0.5485706

Exact integral value = 0.3333

The given estimate is not "good" when compared with the exact value because sample size ($n = 5$) is too small.

6

$$T = (6, 1), (5, 2), (4, 3), (3, 4), (2, 5), (1, 6)$$

$$II \equiv (6, 5), (5, 6)$$

Monte Carlo experiment:

<u>R</u>	outcome
$0 \leq R \leq \frac{1}{6}$	1
$\frac{1}{6} < R \leq \frac{1}{3}$	2
$\frac{1}{3} < R \leq \frac{1}{2}$	3
$\frac{1}{2} < R \leq \frac{2}{3}$	4
$\frac{2}{3} < R \leq \frac{5}{6}$	5
$\frac{5}{6} < R \leq 1$	6

$0 \leq R \leq .167$	1
$.167 < R \leq .333$	2
$.333 < R \leq .5$	3
$.5 < R \leq .667$	4
$.667 < R \leq .833$	5
$.833 < R \leq 1$	6

Continued...

16-3

Set 16.1a

R_1	R_2	Sum	Payoff
.0589	.6733	1+5 = 6 point	
.4799	.9486	3+6 = 9	
.6139	.5933	4+4 = 8	
.9341	.1782	6+2 = 8	
.3473	.5644	3+4 = 7 \rightarrow	-\$10
.3529	.3646	3+3 = 6 point	
.7676	.8931	5+6 = 11	
.3919	.7876	3+5 = 8	
.5199	.6358	4+4 = 8	
.7472	.8954	5+6 = 11	
.5869	.1281	4+1 = 5	
.2867	.8216	2+5 = 7 \rightarrow	-\$10
.8261	.3866	5+3 = 8 point	
.7125	.2108	5+2 = 7 \rightarrow	-\$10
.3575	.2926	3+2 = 5 point	
.3455	.4871	3+3 = 6	
.8111	.8912	5+6 = 11	
.4291	.2302	3+2 = 5 \rightarrow	\$10
.5954	.5423	4+4 = 8 point	
.4208	.6975	3+5 = 8 \rightarrow	\$10

Lead time:

$$0 \leq R \leq .5, \quad L = 1 \text{ day} \\ .5 < R \leq 1, \quad L = 2 \text{ days}$$

Demand/day:

$$0 \leq R \leq .2, \quad d = 0 \text{ unit} \\ .2 < R \leq .9, \quad d = 1 \text{ unit} \\ .9 < R \leq 1, \quad d = 2 \text{ units}$$

Let $p(d, L)$ be the joint pdf of demand and lead time. The procedure calls for constructing a frequency table of demand and lead time.

The maximum demand during lead time is $2 \times 2 = 4$ units, so that the demand $d = 0, 1, 2, 3, 4$. We will use the random numbers in Table 16-1 in the following manner: First use a random number to generate a lead time. If $L=1$ day, use one

random number to generate 7 continued
the demand in that day. If $L=2$ days, use two random numbers to generate the demands for the two days. For example, $R = .058962$ yields $L=1$. Next, $R = .6733$ gives $d=1$. Thus, we update the frequency table by increasing the frequency of the entry $(d=1, L=1)$ by one. The frequency table using the first two columns of R in Table 16-1 is

		0	1	2	3	4	
		1	1	1111	11	0	0
L		2	11	0	1111	1111	0

		0	1	2	3	4	
		1	1	7	2	0	0
L		2	2	0	7	4	0

$$\text{Total } n = 23$$

Relative frequency table: $P(L)$

		0	1	2	3	4	
L		1	1/23	7/23	2/23	0	0
L		2	2/23	0	7/23	4/23	0

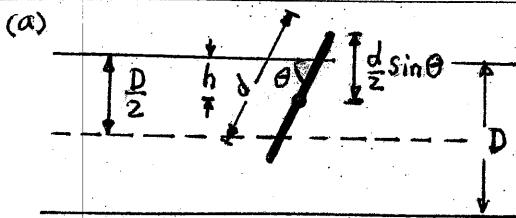
$$p(d) = 3/23 \quad 7/23 \quad 9/23 \quad 4/23 \quad 0$$

Notice that

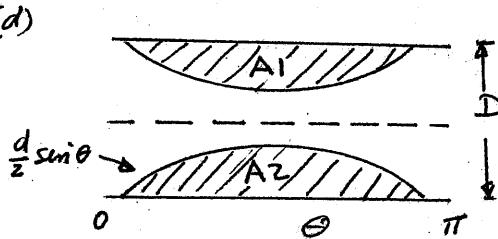
$$P(d) = \sum_L p(d, L)$$

$$P(L) = \sum_d p(d, L)$$

Set 16.1a



8



From graph, needle will touch line or cross it if

$$h \leq \frac{d}{2} \sin \theta$$

(b) Generate $h = R_1 \times D/2$

$$\theta = \pi \times R_2$$

If $h \leq \frac{d}{2} \sin \theta$, needle touches. Else it doesn't.

$$\text{Probability estimate} = \frac{\# \text{ touches}}{\text{sample size}}$$

(c)

A	B	C	D	E
D=	20	d=	10	
(RAND()*.5+.5 RAND()*.PI() SES1*.5*SIN(C4)) IF(B4<=D4,1,0)				
h	theta	$\frac{d}{2} \sin(\theta)/2$	1=touch, 0=else	

Rep 1	8.396953573	1.3165558	4.839272983	0
	7.107859045	2.9048959	1.172463622	0
	0.27542965	0.8440783	3.736795168	1
	1.267504547	2.8354706	1.506816139	1
	9.237262421	0.7436482	3.38488765	0
	2.495379696	2.9719552	0.844125326	0
	4.253169953	2.8396976	1.486650397	0
	8.516662244	1.4161445	4.940326141	0
	4.224254495	0.7887632	3.547410981	0
	3.690266876	3.0811598	0.301979787	0
	Estimate of probability=			0.2
Rep 2	0.712918949	1.5238102	4.994481772	1
	9.381794079	2.5979258	2.586388239	0
	1.360072144	2.0189288	4.506289193	1
	8.477675064	1.9724771	4.60202594	0
	0.99443686	1.300734	4.81877136	1
	5.170438974	1.4568612	4.967582038	0
	5.056822846	1.6844549	4.967739087	0
	5.864264693	0.0683356	0.341412027	0
	6.87137267	2.6283793	2.454895584	0
	1.092023022	2.6522347	2.350296303	1
	Estimate of probability=			0.4
Rep 3	9.712756211	1.694489	4.961799031	0
	6.686447356	1.2243834	4.702983326	0
	6.436673778	2.4581589	3.157296664	0
	1.324134345	2.2441568	3.908652279	1
	1.775706228	2.255079	3.874363448	1
	0.090587765	2.7080167	2.100592855	1
	4.979938633	2.5138689	2.936520016	0
	8.678634219	2.7348178	1.978247037	0
	2.179672677	1.8339609	4.827857959	1
	9.640572895	1.2431615	4.734030551	0
	Estimate of probability=			0.4
Rep 4	8.227016322	2.6999829	2.136976805	0
	8.757368267	2.1537385	4.174233356	0
	4.203914479	0.1860064	0.92467824	0
	6.098369885	2.1672345	4.13670754	0
	4.960185836	0.7841548	3.531135292	0
	3.899078191	1.8047989	4.863730557	1
	5.840727605	0.727722	3.325852126	0
	6.645324046	0.498725	2.391531067	0
	5.361422671	0.89898	3.913462424	0
	3.223016816	1.6715052	4.974665749	1
	Estimate of probability=			0.2
	Mean value =			0.3
	Std. Deviation =			0.1155
	95% LCL =			0.1163
	95% UCL =			0.4837

(c) From (c),

$$\tilde{p} = .3$$

Thus,

$$\frac{2d}{\pi D} = .3$$

$$\text{or } \pi \approx \frac{2d}{.3D}$$

$$\approx \frac{2 \times 10}{.3 \times 20}$$

$$\approx 3.33$$

Set 16.2a

(a) Discrete

1

(b) Continuous

(c) Discrete

In discrete simulation, there
are two main events: arrivals and
departures. An arrival event may
experience delay before starting
service. When service has been
completed, customer leaves the
facility.

2

The description of the discrete
simulation situation by arrival
and departure events is the
reason discrete simulation
is associated with queues.

Events:

1

 A_1 = rush job arrives A_2 = regular job arrives D_1 = rush job departs D_2 = regular job departs A_0 = job arrives at carousel

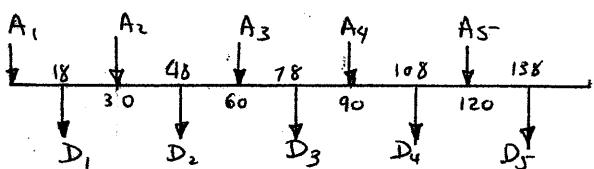
2

 A_1 = job arrives at station 1 A_2 = job arrives at station 2 A_3 = job arrives at station 3 D_1 = job departs station 1 D_2 = job departs station 2 D_3 = job departs station 3 A_1 = car enters lane 1

3

 A_2 = car enters lane 2 A_3 = car goes elsewhere D_1 = car departs lane 1 D_2 = car departs lane 2.

4



Set 16.3b

$$t = -\frac{1}{\lambda} \ln(1-R)$$

$\lambda = 4 \text{ customers/hr}$

Customer R t(hrs) Arrival time

Customer	R	t(hrs)	Arrival time
1	-	-	0
2	.0589	.015	$0 + .015 = .015$
3	.6733	.280	$.015 + .28 = .295$
4	.4799	.163	$.295 + .163 = .458$
A ₁ A ₂		A ₃	A ₄
↓		↓	↓
0 .015		.295	.458

$$f(t) = \frac{1}{b-a}, \quad a \leq t \leq b$$

$$F(t) = \int_0^t \frac{1}{b-a} dx = \frac{t-a}{b-a}, \quad a \leq t \leq b$$

$$R = \frac{t-a}{b-a}$$

$$t = a + (b-a)R$$

$$f_1(t_1) = .5 e^{-5t_1}, \quad \lambda = 1/2 \text{ arrival/hr}$$

$$f_2(t) = \frac{1}{.9}, \quad 1.1 < t < 2$$

$$R = .0589, \quad a_1 = -2 \ln(1-.0589) = .12 \text{ hr}$$

$$R = .6733, \quad d_1 = 1.1 + .9 \times .6733 = 1.71 \text{ hrs}$$

$$R = .4799, \quad a_2 = -2 \ln(1-.4799) = 1.31 \text{ hrs}$$

$$R = .9486, \quad a_3 = -2 \ln(1-.9486) = 5.94 \text{ hrs}$$

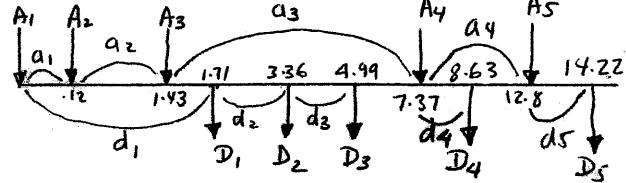
$$R = .6139, \quad d_2 = 1.1 + .9 \times .6139 = 1.65 \text{ hrs}$$

$$R = .5933, \quad d_3 = 1.1 + .9 \times .5933 = 1.63 \text{ hrs}$$

$$R = .9341, \quad a_4 = -2 \ln(1-.9341) = 5.44 \text{ hrs}$$

$$R = .1782, \quad d_4 = 1.1 + .9 \times .1782 = 1.26 \text{ hrs}$$

$$R = .3473, \quad d_5 = 1.1 + .9 \times .3473 = 1.41 \text{ hrs}$$



- (a) $0 \leq R < .2, \quad d = 0$
 $.2 \leq R < .5, \quad d = 1$
 $.5 \leq R < .9, \quad d = 2$
 $.9 \leq R \leq 1., \quad d = 3$

4

Day	R	Demand d	Stock level
0	-	-	5
1	.0589	0	5
2	.6733	2	3
3	.4799	1	2

Replenish stock on day 3

Repair/.2, Package/.8:

5

$0 \leq R < .2, \quad \text{goto Repair}$

$.2 \leq R \leq 1., \quad \text{goto Package}$

Package/.8, Repair/.2:

$0 \leq R < .8, \quad \text{goto Package}$

$.8 \leq R \leq 1, \quad \text{goto Repair}$

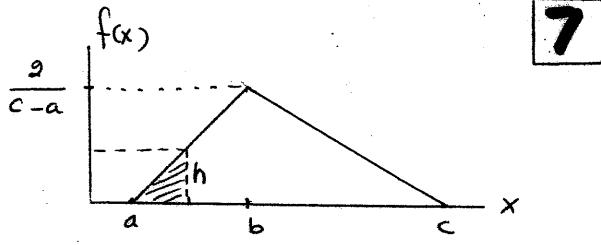
Example: $R = .1$ leads to Repair in the first case and to Package in the second case

$0 \leq R < .5 : H$

$.5 \leq R \leq 1. : T$

n R outcome Payoff

1	.0589	H	\$2
—	—	—	—
1	.6733	T	0
2	.4799	H	$2^2 = \$4$



continued...

Set 16.3b

(a) $F(x) = \begin{cases} \frac{(x-a)^2}{(b-a)(c-a)}, & a \leq x \leq b \\ 1 - \frac{(c-x)^2}{(c-b)(c-a)}, & b \leq x \leq c \end{cases}$ [7 continued]

For $R = \frac{(x-a)^2}{(b-a)(c-a)}$,

$$x = a + \sqrt{R(b-a)(c-a)}, \quad 0 \leq R \leq \frac{b-a}{c-a}$$

For $R = 1 - \frac{(c-x)^2}{(c-b)(c-a)}$,

$$x = c - \sqrt{(c-b)(c-a)(1-R)}, \quad \frac{b-a}{c-a} \leq R \leq 1$$

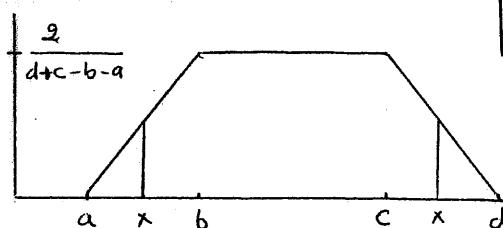
(b) $a = 1, b = 3, c = 7$

$$\frac{b-a}{c-a} = \frac{3-1}{7-1} = -0.333$$

Thus,

$$X = \begin{cases} 1 + \sqrt{(3-1)(7-1)R} \\ = 1 + \sqrt{12R}, \quad 0 \leq R \leq 0.333 \\ 7 - \sqrt{(7-3)(7-1)(1-R)} \\ = 7 - \sqrt{24(1-R)}, \quad -0.333 \leq R \leq 1 \end{cases}$$

R	X
.0589	1.84
.6733	4.20
.4799	3.47
.9486	5.89
.6139	3.96



8

(a) $F(x) = \begin{cases} \frac{(x-a)^2}{(b-a)(d+c-b-a)}, & a \leq x \leq b \\ \frac{1}{(b-a)(d+c-b-a)} + \frac{2(x-a)}{(d+c-b-a)}, & b \leq x \leq c \\ 1 - \frac{(d-x)^2}{(d-c)(d+c-b-a)}, & c \leq x \leq d \end{cases}$

Continued...

$$R = \frac{(x-a)^2}{(b-a)(d+c-b-a)} \text{ gives}$$

$$x = a + \sqrt{(b-a)(d+c-b-a)R}, \quad 0 \leq R \leq \frac{b-a}{(d+c-b-a)}$$

$$R = \frac{1}{(b-a)(d+c-b-a)} + \frac{2(x-b)}{(d+c-b-a)} \text{ gives}$$

$$x = \frac{1}{2} \left(R - \frac{1}{(b-a)(d+c-b-a)} \right) (d+c-b-a),$$

$$\frac{b-a}{d+c-b-a} \leq R \leq 1 - \frac{d-c}{(d+c-b-a)}$$

$$R = 1 - \frac{(d-x)^2}{(d-c)(d+c-b-a)}$$

$$x = d - \sqrt{(d-c)(d+c-b-a)(1-R)},$$

$$1 - \frac{d-c}{(d+c-b-a)} \leq R \leq 1$$

(b) $a = 1, b = 2, c = 4, d = 6$

$$1 + \sqrt{(2-1)(6+4-2-1)R} = 1 + \sqrt{7R}, \quad 0 \leq R \leq .143$$

$$2 + \frac{6+4-2-1}{2}(R - \frac{1}{(2-1)(6+4-2-1)})$$

$$= 2 + 3.5(R - .143),$$

$$.143 \leq R \leq .714$$

$$6 - \sqrt{(6-4)(6+4-2-1)(1-R)} \\ = 6 - \sqrt{14(1-R)} \\ .714 \leq R \leq 1$$

R	X
.0589	1.64
.6733	3.86
.4799	3.18
.9486	5.15
.6139	3.65

$f(x) = pq^x, \quad x = 0, 1, 2, \dots$

$$(p+q) = 1$$

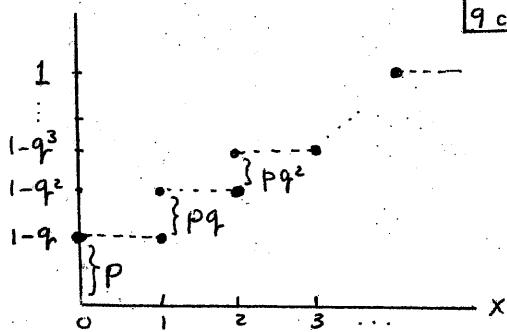
$$F(x) = p \sum_{t=0}^x q^t \\ = 1 - q^{x+1}, \quad x = 0, 1, 2, \dots$$

9

Continued...

16-9

Set 16.3b



9 continued

Sampling procedure:

if $0 \leq R \leq p$, then $x = 0$.

For $p < R \leq 1$, we have

$$1 - q^n \leq R \leq 1 - q^{n+1}$$

or $n \leq \frac{\ln(1-R)}{\ln q} \leq n+1$

Thus, for $p \leq R \leq 1$, compute

$$x = \left[\frac{\ln(1-R)}{\ln q} \right]$$

where $[a]$ is the largest integer less than or equal to a .

For $p = .6$, $q = .4$, we have

R	$\frac{\ln(1-R)}{\ln q}$	x
.0589	—	0
.6733	1.22	1
.4799	—	0
.9486	3.24	3
.6139	1.03	1

$$f(x) = \alpha \beta^{-\alpha} x^{\alpha-1} e^{-(x/\beta)^{\alpha}}, \quad x > 0$$

10

$$= \frac{\alpha}{\beta} \left(\frac{x}{\beta} \right)^{\alpha-1} e^{-\left(\frac{x}{\beta}\right)^{\alpha}}, \quad x > 0$$

$$F(x) = 1 - e^{-\left(\frac{x}{\beta}\right)^{\alpha}}, \quad x > 0$$

Thus, $R = 1 - e^{-\left(\frac{x}{\beta}\right)^{\alpha}}$

or $x = \beta \left[-\ln(1-R) \right]^{1/\alpha}$

Set 16.3c

1

$$y = -\frac{1}{5} \ln \{ (0.0589x \cdot 6733x \cdot 4799x \cdot 9486x) \}$$

$$= .803 \text{ hour}$$

2

$$\lambda = 5 \text{ events/hr}, t = 1$$

$$e^{-5x_1} = e^{-5} = .00673$$

i $R_1 R_2 \dots R_4$

1	.0589
2	.0589x.6733 = .0397
3	.0397x.4799 = .0190
4	.0190x.9486 = .0181
5	.0181x.6139 = .0111
6	.0111x.5933 = .00656
7	.00656x.9341 = .00614

Hence $n = 6$

$\mu = 8, \sigma = 1, N(8, 1)$

2

3

Convolution method:

$$x = R_1 + R_2 + \dots + R_6 = 6.1094$$

$$y = 8 + 1(6.1094 - 6) = 8.1094$$

Box-Miller method:

$$x = \sqrt{-2 \ln R_1} \cos(2\pi R_2)$$

$$= \sqrt{-2 \ln 0.0589} \cos(2\pi x.6733)$$

$$\approx -1.103$$

$$y = 8 + 1(-1.103) = 6.897$$

4

$$\lambda = 6 \text{ / day } m = 5$$

$$y = -\frac{1}{6} \ln (0.0589x.6733x.4799x \\ \cdot 9486x.6139) = .751 \text{ hour}$$

$N(27, 3): \mu = 27, \sigma = 3$

5

Given R_1 and R_2 , we have

$$x_1 = \sqrt{-2 \ln R_1} \cos(2\pi R_2)$$

$$x_2 = \sqrt{-2 \ln R_1} \sin(2\pi R_2)$$

$$y_1 = \mu + \sigma x_1$$

$$y_2 = \mu + \sigma x_2$$

Continued...

J	K	L	M	N	O
Mean = 27	Std. Dev. = 3				
R1	R2	x1	x2	y1	y2
5 0.0589	0.6733	-1.1030306	-2.108827	23.69091	20.67352
6 0.4799	0.9486	1.149111	-0.384576	30.44733	25.84627
7 0.6139	0.5933	-0.8229152	-0.546495	24.53125	25.36051
			mean y=	25.09163	
			Sy	3.197533	

Formulas:

$$L5 = \text{SQRT}(-2 * \text{LN}(J5)) * \text{COS}(2 * \text{PI}() * K5)$$

$$M4 = \text{SQRT}(-2 * \text{LN}(J5)) * \text{SIN}(2 * \text{PI}() * K5)$$

$$N4 = \$K\$1 + L4 * \$M\$1$$

$$O4 = \$K\$1 + M4 * \$M\$1$$

$$x_i = 10 + (20 - 10) R_i$$

$$= 10 + 10 R_i, i = 1, 2, 3, 4$$

$$t = x_1 + x_2 + x_3 + x_4$$

$$= 40 + 10(R_1 + R_2 + R_3 + R_4)$$

R ₁	R ₂	R ₃	R ₄	t (sec)	$\sum t$	7
1 0.0589	0.6733	4.799	9.486	61.61	61.60	
2 0.6139	0.5933	9.341	1.782	63.20	124.81	
3 3.4793	7.676	8.931	3.919	64.00	188.81	
4 7.876	15.199	6.358	7.472	66.91	255.72	
5 8.954	5.869	1.128	2.867	58.94	314.69	

The number of mice that exit the maze in 300 seconds is 4

Let x_1, x_2, \dots, x_n be n successive random deviates obtained from the geometric distribution as given in Problem 9, Set 18.3b. Then

$$x_i = \left[\frac{\ln R_i}{\ln(1-p)} \right], i = 1, 2, \dots, n$$

Because the negative binomial is the convolution of r independent geometric random variables, it follows that a random negative binomial sample can be determined as

$$X = \sum_{i=1}^n \left[\frac{\ln R_i}{\ln(1-p)} \right]$$

Note that $[a]$ represents the largest integer $\leq a$

Set 16.3d

Step 1: $R = .6139$
 $X = .6139$

Step 2: $R = .5933$

Step 3: $\frac{f(.6139)}{g(.6139)} = .948 > .5933$

Step 1: $R = .9341$, $X = .9341$ Reject X

Step 2: $R = .1782$

Step 3: $\frac{f(.9341)}{g(.9341)} = \frac{.3693}{1.5} = .246 > .1782$

Reject X

Step 1: $R = .3473$, $X = .3473$

Step 2: $R = .5644$

Step 3: $\frac{f(.3473)}{g(.3473)} = \frac{.9067}{1.5} = .6044 > .5644$

Reject X

Step 1: $R = .3529$, $X = .3529$

Step 2: $R = .3646$

Step 3: $\frac{f(.3529)}{g(.3529)} = \frac{.913}{1.5} = .6086 > .3646$

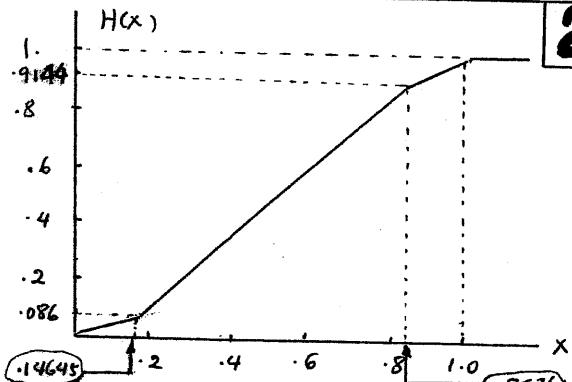
Reject X

Step 1: $R = .7676$, $X = .7676$

Step 2: $R = .8931$

Step 3: $\frac{f(.7676)}{g(.7676)} = \frac{.7135}{1.5} = .4756 < .8931$

Accept $X = .7676$



Step 1: $R = .4799$, $X = .4831$

Step 2: $R = .9486$

Step 3: $\frac{f(.4831)}{g(.4831)} = \frac{.9988}{1.5} = .6659 > .9486$

Reject X

Step 1: $R = .6139$, $X = .5974$

Step 2: $R = .5933$

continued...

Step 3: $\frac{f(.5974)}{g(.5974)} = \frac{.962}{1.5} = .6413$ 2 continued
 reject X

Step 1: $R = .9341$, $X = .8804$

Step 2: $R = .1782$

Step 3: $\frac{f(.8804)}{g(.8804)} = \frac{.842}{1.5} = .5613 > .1782$

Reject X

Step 1: $R = .3529$, $X = .375$

Step 2: $R = .3646$

Step 3: $\frac{f(.375)}{g(.375)} = \frac{.937}{1.5} = .6247 > .3646$

Reject X

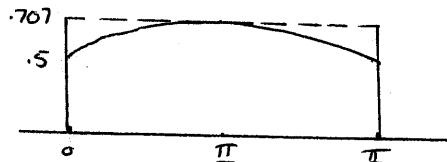
Step 1: $R = .7676$, $X = .7286$

Step 2: $R = .8931$

Step 3: $\frac{f(.7286)}{g(.7286)} = \frac{.791}{1.5} = .5273 < .8931$

Accept X

3



$f(x) = \frac{\sin(x) + \cos(x)}{2}$ $0 \leq x \leq \frac{\pi}{2}$

$\max f(x) = .707 \text{ at } x = \frac{\pi}{4}$

$g(x) = .707 \quad 0 \leq x \leq \pi/2$

$h(x) = \frac{g(x)}{\text{area under } g(x)}$

$= \frac{.707}{.707 \times \frac{\pi}{2}} = .637 \quad 0 \leq x \leq \frac{\pi}{2}$

$\int_{12}^{20} \frac{K_1}{t} dt = K_1 \ln \frac{20}{12} = 1$

Thus, $K_1 = 1.96$

$\int_{18}^{22} \frac{K_2}{t^2} dt = K_2 \left(\frac{1}{18} - \frac{1}{22} \right) = 1$

Thus, $K_2 = 99$

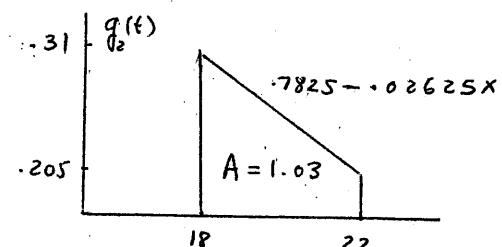
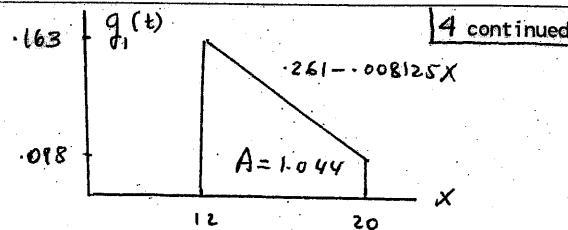
$f_1(t) = \frac{1.96}{t}, \quad 12 \leq t \leq 20$

$f_2(t) = \frac{99}{t^2}, \quad 18 \leq t \leq 22$

continued...

4

Set 16.3d



$$h_1(t) = \frac{.261 - .008125t}{1.044}$$

$$= .25 - .007783t$$

$$H_1(t) = .025t - .00778 \left. \frac{t^2}{2} \right|_{12}$$

$$= .25t - .003892t^2 - 2.44$$

$$h_2(t) = \frac{.7825 - .02625t}{1.03}$$

$$= .76 - .0255t$$

$$H_2(t) = .76t - .01275t^2 - 9.55$$

Sample computations from $H_2(t)$:

Step 1: $R_1 = .0589$

$$.76t - .01275t^2 - 9.55 = .0589$$

$$t^2 - 59.6t + 753.64 = 0$$

$$t = \frac{59.6 \pm \sqrt{(-59.6)^2 - 4 \times 753.64}}{2}$$

$$= 18.2$$

Step 2: $R = .6733$

$$\text{Step 3: } \frac{f_2(18.2)}{g_2(18.2)} = \frac{\left(\frac{99}{18.21^2}\right)}{.7825 - .02625 \times 18.21}$$

$$= .98 > .6733$$

Reject t .

continued...

Set 16.4a

1

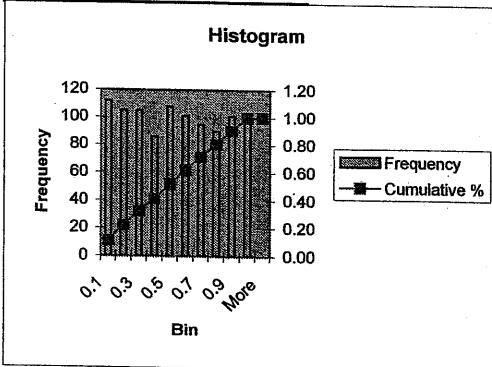
Multiplicative Congruential Method	
Input data	
$a_1 =$	17
$c_1 =$	111
$m_1 =$	7
$m_1^2 =$	103
How many numbers?	50
Output results	
Press to Generate Sequence	
Generated random numbers:	
1	0.23301
2	0.03883
3	0.73786
4	0.62136
5	0.64078
6	0.97087
7	0.58252
8	0.98058
9	0.74757
10	0.78641
11	0.44660
12	0.66990
13	0.46602
14	0.00000
15	0.07767
16	0.39806
17	0.84466
18	0.43689
19	0.50485
20	0.66019
21	0.30097
22	0.19417
23	0.37864
24	0.51456
25	0.82524
26	0.10680
27	0.89320
28	0.26214
29	0.53398
30	0.15534
31	0.71845
32	0.29126
33	0.02913
34	0.57282
35	0.81553
36	0.94175
37	0.08738
38	0.56311
39	0.65049
40	0.13592
41	0.38835
42	0.67961
43	0.63107
44	0.80583
45	0.77670
46	0.28155
47	0.86408
48	0.76699
49	0.11650
50	0.05825

2

R=RAND()	Bin
0.813455	0.1
0.21757	0.2
0.937991	0.3
0.840823	0.4
0.19536	0.5
0.681599	0.6
0.829291	0.7
0.377723	0.8
0.149187	0.9
0.965781	1
0.808752	
0.957601	
0.502469	
0.620944	
0.992405	
0.97218	
0.051905	
0.144368	
0.129308	
0.676603	
0.140868	
0.486705	
0.12415	
0.821802	
0.954853	
0.301267	
0.827929	
0.917179	
0.07369	
0.462159	
0.333902	
0.390604	
0.723163	
0.041401	
0.805603	
0.556012	

Bin	Frequency	umulative %
0.1	112	0.11
0.2	105	0.22
0.3	105	0.32
0.4	86	0.41
0.5	108	0.52
0.6	101	0.62
0.7	95	0.71
0.8	90	0.80
0.9	101	0.90
1	97	1.00
More	0	1.00

Sample
Size = 1000



Set 16.5a

$C = 2$ barbers

$$f_1(t) = .1 e^{-0.1t}, \quad t > 0$$

$$f_2(t) = \frac{1}{15}, \quad 15 \leq t \leq 30$$

$$t_1 = -12 \ln R$$

$$t_2 = 15 + 15R$$

A_1 at $T=0$:

$$T(A_1) = 0 + (-10 \ln .0589) = 28.3$$

$$T(D_2) = 0 + (15 + 15 \times .6733) = 25.1$$

Barber 1 busy

D_2 at $T=25.1$:

Barber 1 idle

A_2 at $T=28.3$:

$$T(A_2) = 28.3 - 10 \ln .4799 = 35.6$$

$$T(D_2) = 28.3 + (15 + 15 \times .9486) = 57.5$$

Barber 1 busy $A_3 \ D_2$

A_3 at $T=35.6$:

$$T(A_3) = 35.6 - 10 \ln .6139 = 40.5$$

$$T(D_3) = 35.6 + (15 + 15 \times .5933) = 59.5$$

Barber 2 busy $A_4 \ D_2 \ D_3$

A_4 at $T=40.5$:

$$T(A_4) = 40.5 - 10 \ln .9341 = 41.2$$

A_4 waits in queue

 $A_5 \ D_2 \ D_3$
 A_4 ← queue

A_5 at $T=41.2$:

$$T(A_5) = 41.2 - 10 \ln .1782 = 58.4$$

A_5 waits in queue

 $D_2 \ A_6 \ D_3$
 $A_4 \ A_5$ ← queue

D_2 at $T=57.5$:

Barber 1 idle

Take A_4 out of queue

$$T(D_4) = 57.5 + 15 + 15 \times .3473 = 77.7$$

Barber 1 busy

 $A_6 \ D_3 \ D_4$
 A_5 ← queue

A_6 at $T=58.4$:

$$T(A_7) = 58.4 - 10 \ln .5644 = 64.1$$

Put A_6 in queue $D_3 \ A_7 \ D_4$

D_3 at $T=59.5$:

 $A_5 \ A_6$ ← queue

Barber 2 idle

Take A_5 out of queue

$$T(D_5) = 59.5 + 15 + 15 \times .3529 = 79.8$$

Barber 2 busy

 $A_7 \ D_4 \ D_5$
 A_6 ← queue

A_7 at $T=64.1$:

$$T(A_8) = 64.1 - 10 \ln .3646 = 74.2$$

Put A_7 in queue

 $A_8 \ D_4 \ D_5$
 $A_6 \ A_7$ ← queue

A_8 at $T=74.2$:

$$T(A_9) = 74.2 + (-10 \ln .7676)$$

$$= 76.8$$

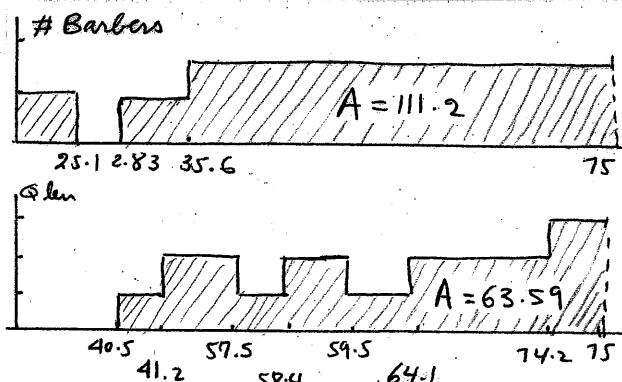
Place A_8 in queue.

 $A_9 \ D_4 \ D_5$
 $A_6 \ A_7 \ A_8$ ← queue

continued...

continued...

Set 16.5a



$$\text{Av. facility utilization} = \frac{111.2}{75}$$

$$= 1.48 \text{ barbers}$$

$$\text{Av. queue length} = \frac{63.59}{75} = .8 \text{ customer}$$

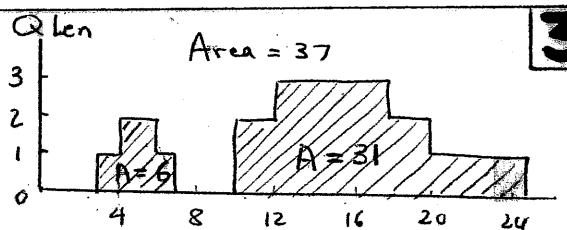
$$\text{Av. waiting time in queue} = \frac{63.59}{8}$$

$$= 7.95 \text{ min}$$

$$\text{Av. waiting time for those who must wait} = \frac{63.59}{5} = 12.72 \text{ min}$$

- (a) Observation.
- (b) Time.
- (c) Observation.
- (d) Observation.
- (e) Observation.
- (f) Time.

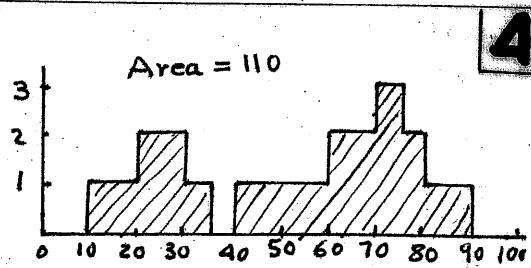
2



$$(a) \bar{Q} = \frac{37}{25} = 1.48 \text{ customers}$$

$$(b) \text{Number of waiting customers} = 5$$

$$\bar{W} = \frac{37}{5} = 7.4 \text{ hours}$$



(a) Average utilization

$$= \frac{110}{100} = 1.1 \text{ barber}$$

(b) Average idle time

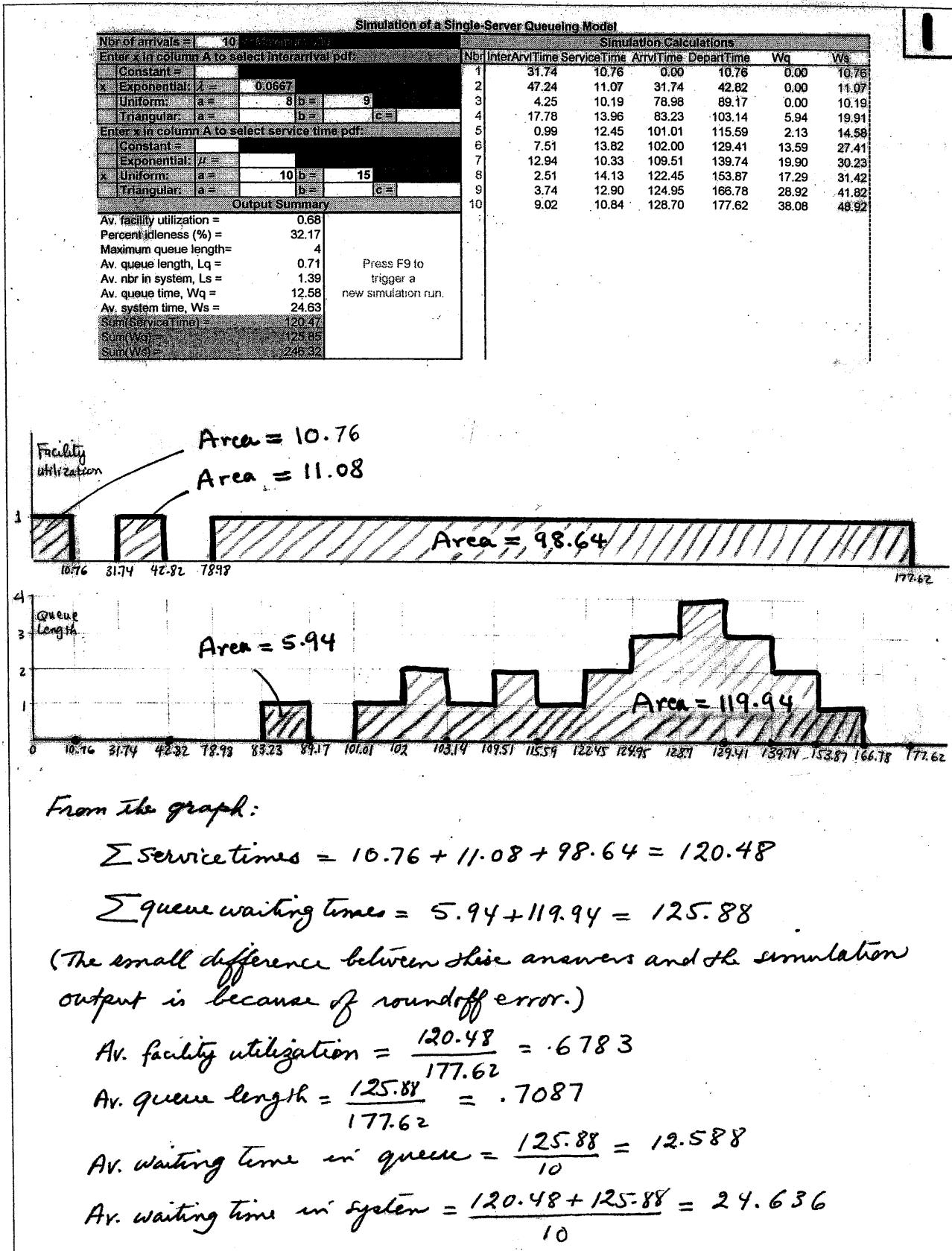
$$= \frac{10 + (40 - 35) + (100 - 90)}{3}$$

$$= \frac{25}{3}$$

$$= 8.33 \text{ minutes}$$

3

Set 16.5b



Set 16.5b

Nbr. of arrivals	500	<<Maximum 5000
Enter x in column A to select distribution		
Constant		
x	Exponential:	/
	Uniform:	a =
	Triangular:	a =
Enter x in column A to select service time		
Constant		
x	Exponential:	/
	Uniform:	a =
	Triangular:	a =
Output Summary		
Av. facility utilization	=	0.66
Percent idleness (%)	=	33.84
Maximum queue length	=	0
(1) Av. queue length, Lq	=	1.42
Av. nbr in system, Ls	=	2.08
Av. queue time, Wq	=	0.37
Av. system time, Ws	=	0.54
Av. facility utilization	=	0.61
Percent idleness (%)	=	38.65
Maximum queue length	=	0
(2) Av. queue length, Lq	=	0.91
Av. nbr in system, Ls	=	1.52
Av. queue time, Wq	=	0.24
Av. system time, Ws	=	0.40
Av. facility utilization	=	0.65
Percent idleness (%)	=	35.11
Maximum queue length	=	0
(3) Av. queue length, Lq	=	0.91
Av. nbr in system, Ls	=	1.56
Av. queue time, Wq	=	0.22
Av. system time, Ws	=	0.38
Av. facility utilization	=	0.68
Percent idleness (%)	=	31.70
Maximum queue length	=	0
(4) Av. queue length, Lq	=	1.35
Av. nbr in system, Ls	=	2.03
Av. queue time, Wq	=	0.32
Av. system time, Ws	=	0.48
Av. facility utilization	=	0.60
Percent idleness (%)	=	39.83
Maximum queue length	=	0
(5) Av. queue length, Lq	=	1.14
Av. nbr in system, Ls	=	1.74
Av. queue time, Wq	=	0.30
Av. system time, Ws	=	0.46

Summary:

	Utiliz	Lq.	4s	Wq	Wa
Mean	.64	1.146	1.786	.29	.452
Std. Dev.	.0339	.2388	.2598	.0608	.0642

95% confidence limits:

$$t_{4,025} = 2.776$$

$$UCL = \bar{x} + \frac{2.776s}{\sqrt{n}} = \bar{x} + 1.245s$$

$$LCL = \bar{x} - 1.245s$$

	utiliz	Lq	Ls	Wq	Ws
LCL	.598	.850	1.464	.215	.372
UCL	.682	1.442	2.108	.365	.531

Poisson queue output

Scenario 1 – (M/M/1): (GD/infinity/infinity)

Lambda = 4.00000 Mu = 6.00000
Lambda eff = 4.00000 Rho/c = 0.66667

Ls = 2.00000 Lq = 1.33333
Ws = 0.50000 Wq = 0.33333

S=	200	<<Maximum 500
Column A to select Integrator value:		
rat:	11.5	
rat:	a =	
rat:	b =	
rat:	c =	
Column A to select Service time (sec):		
rat:	0	
rat:	d =	
rat:	e =	
rat:	f =	

Av. facility utilization =	0.96
Percent idleness (%) =	4.20
Maximum queue length =	2
Av. queue length, L_q =	0.12
Av. nbr in system, L_s =	1.08
Av. queue time, W_q =	1.36
Av. system time, W_s =	12.38

16-18

	Av. facility utilization =	0.96
	Percent idleness (%) =	3.85
	Maximum queue length=	2
(2)	Av. queue length, Lq =	0.12
	Av. nbr in system, Ls =	1.08
	Av. queue time, Wq =	1.33
	Av. system time, Ws =	12.39
	Av. facility utilization =	0.97
	Percent idleness (%) =	2.98
	Maximum queue length=	2
(3)	Av. queue length, Lq =	0.19
	Av. nbr in system, Ls =	1.16
	Av. queue time, Wq =	2.14
	Av. system time, Ws =	13.33
	Av. facility utilization =	0.96
	Percent idleness (%) =	3.58
	Maximum queue length=	2
(4)	Av. queue length, Lq =	0.16
	Av. nbr in system, Ls =	1.13
	Av. queue time, Wq =	1.88
	Av. system time, Ws =	12.97
	Av. facility utilization =	0.97
	Percent idleness (%) =	3.39
	Maximum queue length=	2
(5)	Av. queue length, Lq =	0.17
	Av. nbr in system, Ls =	1.14
	Av. queue time, Wq =	2.00
	Av. system time, Ws =	13.12

utilization:

$$\text{mean} = \frac{.96 + .96 + .97 + .96 + .97}{5}$$

$$= .964$$

$$\text{st.dev.} = .0311$$

Set 16.6a

$$W_1 = \frac{14}{3} = 4.67 \text{ (time units)}$$

$$W_2 = \frac{10}{4} = 2.5$$

$$W_3 = \frac{11}{3} = 3.67$$

$$W_4 = \frac{6}{3} = 2$$

$$W_5 = \frac{15}{4} = 3.75$$

$$\bar{W} = \frac{4.67 + 2.5 + 3.67 + 2 + 3.75}{5} = 3.32 \text{ time units}$$

Discard observations during the transient period (0, 100)

$$W_1 = \frac{12 + 30 + 10 + 14 + 16}{5} = 16.4 \text{ time units}$$

$$W_2 = \frac{15 + 17 + 20 + 22}{4} = 18.5$$

$$W_3 = \frac{10 + 20 + 30 + 15 + 25 + 31}{6} = 21.83$$

$$W_4 = \frac{15 + 17 + 20 + 14 + 13}{5} = 15.8$$

$$W_5 = \frac{25 + 30 + 15}{3} = 23.33$$

$$\bar{W} = 19.17 \quad S = 3.3$$

Confidence interval

$$\bar{W} \pm t_{.025, 4} \frac{S}{\sqrt{n}}$$

$$= 19.17 \pm 2.776 \frac{3.3}{\sqrt{5}}$$

or

$$15.07 \leq \mu \leq 23.27$$

Batch	a_i	b_i	y_i
1	6	7	.869
2	10	7	1.369
3	6	9	.584

$$\bar{a} = 7.33 \quad \bar{b} = 7.67 \quad \bar{y} = .941$$

$$Sy = .397$$

continued...

16-20

$$y_i = \frac{3x7.33 - (3-1)(3x7.33 - a_i)}{7.67 - 3x7.67 - b_i}$$

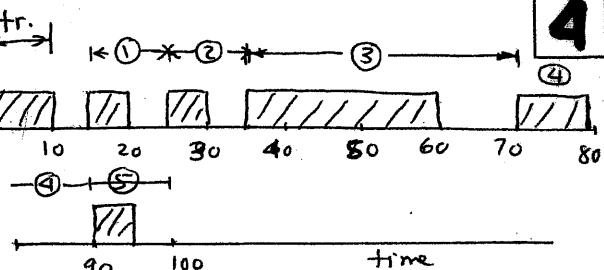
$$= 2.867 - \frac{43.98 - 2a_i}{23.01 - b_i}$$

95% confidence interval:

$$.941 - 2.776 \frac{.397}{\sqrt{3}} \leq \mu \leq .941 + 2.776 \frac{.397}{\sqrt{3}}$$

$$.305 \leq \mu \leq 1.577$$

3 continued



(a) Start points are 15, 25, 35, 70, 90

(b)

Batch	a_i	b_i	y_i
1	5	10	.54
2	5	10	.54
3	25	35	.94
4	10	20	.45
5	5	10	.54

$$\bar{a} = 10 \quad \bar{b} = 17 \quad \bar{y} = .602$$

$$Sy = .193$$

$$y_i = \frac{5x10 - 4(5x10 - a_i)}{17 - 5x17 - b_i}$$

$$= 2.94 - \frac{200 - 4a_i}{85 - b_i}$$

$$.602 - 2.776 \frac{.193}{\sqrt{5}} \leq \mu \leq .602 + 2.776 \frac{.193}{\sqrt{5}}$$

$$\text{or} \quad .36 \leq \mu \leq .84$$

$$(9) \quad t = \frac{90}{5} = 18$$

i	1	2	3	4	5
A	8	13	14	10	5
a_i	.44	.72	.78	.56	.28

Mean = .556, Std. Dev. = .2042