

# OverTheWire Krypton Walkthrough

## Introduction

This is a comprehensive walkthrough of the Krypton wargame offered by the OverTheWire platform. Krypton is a series of cybersecurity challenges that test the player's knowledge and understanding of cryptography. My walkthrough of it covers:

- The research I undertook to understand the concepts covered by each level,
- And the methodology I utilized to finally solve the problem.

## Connection

Each Krypton levels can be accessed via SSH, specifically the command:

'ssh kryptonX@krypton.labs.overthewire.org -p 2231', where X is the level you are attempting to access.

## Code Repository

All of the code I've written to complete the challenges can be found in this [GitHub repository](#). Most of the programs are quite simple, but if you'd be interested in reading about how they work, I've written comprehensive explanations for each program In the README file.

## Level 0 - Level 1

Link: <https://overthewire.org/wargames/krypton/krypton0.html>

## Level Info

Welcome to Krypton! The first level is easy. The following string encodes the password using Base64:

S1JZUFRPTKITR1JFQVQ=

Use this password to log in to krypton.labs.overthewire.org with username krypton1 using SSH on port 2231.

You can find the files for other levels in /krypton/

## Research

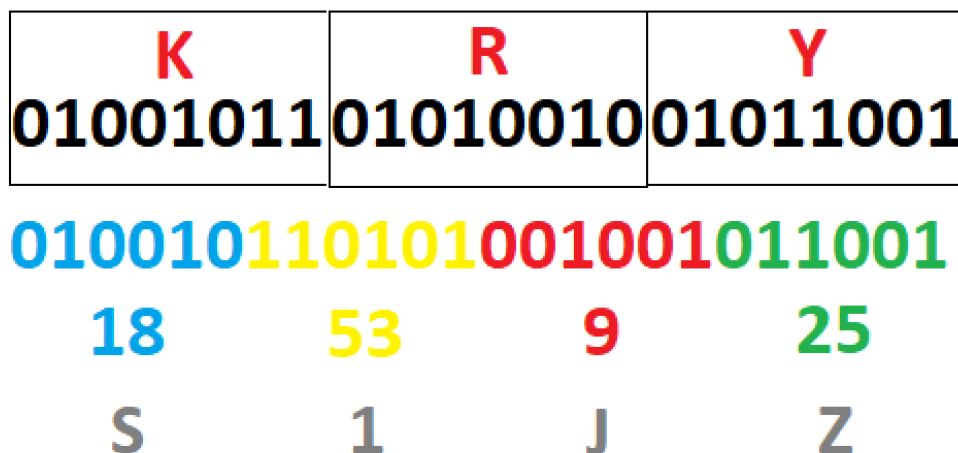
Base64 is known as a "binary-to-text" method of encryption. It represents binary data in blocks of 24 bits, split into four 6-bit long Base64 characters. The available Base64 characters are shown in the following table:

Table 1: The Base 64 Alphabet

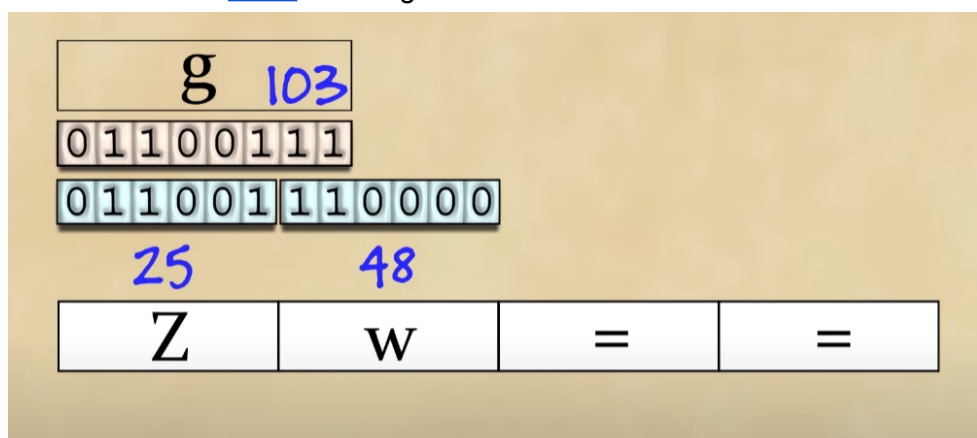
Value	Encoding	Value	Encoding	Value	Encoding	Value	Encoding
0	A	17	R	34	i	51	z
1	B	18	S	35	j	52	0
2	C	19	T	36	k	53	1
3	D	20	U	37	l	54	2
4	E	21	V	38	m	55	3
5	F	22	W	39	n	56	4
6	G	23	X	40	o	57	5
7	H	24	Y	41	p	58	6
8	I	25	Z	42	q	59	7
9	J	26	a	43	r	60	8
10	K	27	b	44	s	61	9
11	L	28	c	45	t	62	+
12	M	29	d	46	u	63	/
13	N	30	e	47	v		
14	O	31	f	48	w	(pad)	=
15	P	32	g	49	x		
16	Q	33	h	50	y		

Sourced from [RFC 4648](#)

The way that Base64 works is best shown using an example. Therefore, let's try encoding the characters KRY from Krypton.



In the above diagram, the characters K, R, and Y are converted into the 8-bit binary values in black. This can be extracted from an ASCII table. Afterwards, we line up these bits next to each other and split them up into the four color-coded 6-bit binary blocks. The binary values can then be converted to decimals. Using the Base64 table above, we can then reach the conclusion that the string 'KRY' is encrypted into 'S1JZ'. A string will not always be long enough to be perfectly divided into 6-bit sequences. Hence, Base64 encryption features the use of padding characters, namely the "=" character. Examine the following example by Youtuber 'schenken' in his [video](#) covering Base64:



The first 6 bits of 'g' are encoded into the letter 'Z'. The next 2 bits are padded with 0's to form the 6-bit sequence '110000', which is converted into 'w'. As there is nothing left to encode, Base64 pads the rest of the string with '=' characters. Hence, the letter 'g' on its own is encoded as 'Zw==' using Base64

## Method

Any Base64 decoder can be used to solve this problem. I used the Bash shell's built-in Base64 command like so:

```
ditto@Roblox-Hacking-Station:/mnt/c/Users/Kevin Ngo$ base64 -d -  
S1JZUFRPTk1TR1JFQVQ=  
KRYPTONISGREATditto@Roblox-Hacking-Station:/mnt/c/Users/Kevin Ngo$ |
```

The password for the next level is revealed to be **KRYPTONISGREAT**. I didn't actually know about the Bash shell's Base64 command, so I simply looked it up on Google and accessed the [documentation](#).

## Level 1 - Level 2

Link: <https://overthewire.org/wargames/krypton/krypton1.html>

### Level Info

The password for level 2 is in the file 'krypton2'. It is 'encrypted' using a simple rotation. It is also in non-standard ciphertext format. When using alpha characters for cipher text it is normal to group the letters into 5 letter clusters, regardless of word boundaries. This helps obfuscate any patterns. This file has kept the plain text word boundaries and carried them to the cipher text. Enjoy!

### Preamble

Before starting the level, we must connect to it using the password we obtained in the previous level and the connection instructions at the start of the document. Afterwards we'll be connected to the Krypton virtual machine, which is running a Linux distro. The level info above suggests that the password is stored in a file called "krypton2", and so we must navigate to it like so:

```
krypton1@bandit:~$ ls
krypton1@bandit:~$ cd /krypton
krypton1@bandit:/krypton$ ls
krypton1  krypton2  krypton3  krypton4  krypton5  krypton6
krypton1@bandit:/krypton$ cd krypton1
krypton1@bandit:/krypton/krypton1$ ls
krypton2  README
krypton1@bandit:/krypton/krypton1$
```

Examining the README file provides us with further information about the level, specifically that the password is encrypted using **ROT13**:

```
krypton1@bandit:/krypton/krypton1$ cat README
Welcome to Krypton!

This game is intended to give hands on experience with cryptography
and cryptanalysis. The levels progress from classic ciphers, to modern,
easy to harder.

Although there are excellent public tools, like cryptool, to perform
the simple analysis, we strongly encourage you to try and do these
without them for now. We will use them in later exercises.

** Please try these levels without cryptool first **

The first level is easy. The password for level 2 is in the file
'krypton2'. It is 'encrypted' using a simple rotation called ROT13.
It is also in non-standard ciphertext format. When using alpha characters for
cipher text it is normal to group the letters into 5 letter clusters,
regardless of word boundaries. This helps obfuscate any patterns.

This file has kept the plain text word boundaries and carried them to
the cipher text.

Enjoy!
krypton1@bandit:/krypton/krypton1$ |
```

Now, having a look at the krypton2 file gives us the encrypted password: **YRIRY GJB CNFFJBEQ EBGGRA**.

```
krypton1@bandit:/krypton/krypton1$ cat krypton2
YRIRY GJB CNFFJBEQ EBGGRA
krypton1@bandit:/krypton/krypton1$
```

From this initial reconnaissance, we've gathered that we need to understand how ROT13 works, and how to decrypt it to obtain the password. Therefore, I've conducted some research on it and have documented my findings below:

## Research

ROT13 is a substitution cipher where each plaintext letter is replaced with the letter that is 13 places after it in the alphabet. For example, the letter 'A' becomes 'N' because 'A' is the 1st letter in the alphabet, whereas 'N' is the 14th letter. Since 13 is half of the number of letters in the alphabet, the inverse would also be true. Encrypting 'N' using ROT13 would result in it being substituted for 'A' (14th letter + 13 = 27 = 26 + 1st letter in the alphabet).

A	B	C	D	E	F	G	H	I	J	K	L	M
↕	↕	↕	↕	↕	↕	↕	↕	↕	↕	↕	↕	↕
N	O	P	Q	R	S	T	U	V	W	X	Y	Z

ROT13

D	U	C	K
↕	↕	↕	↕
Q	H	P	X

Sourced from <https://www.geeksforgeeks.org/rot13-cipher/>

From this, it should be clear that decrypting ROT13 is extremely easy; we simply shift every letter forward by 13 places.

## Method

There are a plethora of websites out there that will easily decipher a ROT13 cipher, such as this [one](#). However, to challenge myself and improve my programming skills, I decided to write a [Python program](#) to accomplish this:

```
1  #!/usr/bin/env python3
2
3  lower = {
4      "a" : 1, "b" : 2, "c" : 3, "d" : 4, "e" : 5, "f" : 6, "g" : 7, "h":8,
5      "i" : 9, "j" : 10, "k" : 11, "l" : 12, "m" : 13, "n" : 14, "o" : 15,
6      "p" : 16, "q" : 17, "r" : 18, "s" : 19, "t" : 20, "u" : 21, "v" : 22,
7      "w" : 23, "x" : 24, "y" : 25, "z" : 26
8  }
9
10 l_key_list = list(lower.keys())
11
12 upper = {
13     "A" : 1, "B" : 2, "C" : 3, "D" : 4, "E" : 5, "F" : 6, "G" : 7, "H":8,
14     "I" : 9, "J" : 10, "K" : 11, "L" : 12, "M" : 13, "N" : 14, "O" : 15,
15     "P" : 16, "Q" : 17, "R" : 18, "S" : 19, "T" : 20, "U" : 21, "V" : 22,
16     "W" : 23, "X" : 24, "Y" : 25, "Z" : 26
17 }
18
19 u_key_list = list(upper.keys())
20
21 SHIFT = 13
22 LETTERS = 26
23
24 message = input("Enter a string:")
25 decoded = ''
26
27
28 for character in message:
29     # calculate (charpos + 13) mod 26
30     if character in lower:
31         new_pos = (lower[character] + SHIFT) % LETTERS
32         decoded = decoded + l_key_list[new_pos - 1]
33     elif character in upper:
34         new_pos = (upper[character] + SHIFT) % LETTERS
35         decoded = decoded + u_key_list[new_pos - 1]
36     else:
37         decoded = decoded + character
38
39 print(decoded)
```

The above code works using two dictionaries, one to store lower case letters along with their indexes, and the other to store upper case letters. The user inputs a message via STDIN, then the program loops through each character in the string. Each character is shifted forward by 13 characters using the operation:  $(characterIndex + 13) \bmod 26$ . The shifted character is concatenated to the “decoded” string and printed after all characters have been looped over. Due to the limitations of ROT13, any character that isn’t in the alphabet is not shifted.

Now, let’s input the cipher **YRIRY GJB CNFFJBEQ EBGGRA** into the program:

```
/python.exe" f:/Code/Krypton/rot13.py
Enter a string:YRIRY GJB CNFFJBEQ EBGGRA
LEVEL TWO PASSWORD ROTTEN
```

Hence, the password for this level is **ROTTEN**.

## **Level 2 - Level 3**

Link: <https://overthewire.org/wargames/krypton/krypton2.html>

### **Level Info**

ROT13 is a simple substitution cipher.

Substitution ciphers are a simple replacement algorithm. In this example of a substitution cipher, we will explore a ‘monoalphabetic’ cipher. Monoalphabetic means, literally, “one alphabet” and you will see why.

This level contains an old form of cipher called a ‘Caesar Cipher’. A Caesar cipher shifts the alphabet by a set number. For example:

plain: a b c d e f g h i j k ...

cipher: G H I J K L M N O P Q ...

In this example, the letter ‘a’ in plaintext is replaced by a ‘G’ in the ciphertext so, for example, the plaintext ‘bad’ becomes ‘HGJ’ in ciphertext.

The password for level 3 is in the file krypton3. It is in 5 letter group ciphertext. **It is encrypted with a Caesar Cipher**. Without any further information, this cipher text may be difficult to break. You do not have direct access to the key, however you do have access to a program that will encrypt anything you wish to give it using the key. If you think logically, this is completely easy.

One shot can solve it!

Have fun.

### **Additional Information:**

The encrypt binary will look for the keyfile in your current working directory. Therefore, it might be best to create a working directory in /tmp and in there a link to the keyfile. As the encrypt binary runs setuid krypton3, you also need to give krypton3 access to your working directory.

### **Here is an example:**

```
krypton2@melinda:~$ mkdir -p
```

```
/tmp/tmp.Wf2OnCpCDQ
```

```
krypton2@melinda:~$ cd /tmp/tmp.Wf2OnCpCDQ
```

```
krypton2@melinda:/tmp/tmp.Wf2OnCpCDQ$ ln -s /krypton/krypton2/keyfile.dat
krypton2@melinda:/tmp/tmp.Wf2OnCpCDQ$ ls
keyfile.dat
krypton2@melinda:/tmp/tmp.Wf2OnCpCDQ$ chmod 777 .
krypton2@melinda:/tmp/tmp.Wf2OnCpCDQ$ /krypton/krypton2/encrypt /etc/issue
krypton2@melinda:/tmp/tmp.Wf2OnCpCDQ$ ls
ciphertext keyfile.dat
```

## Research

Our earlier research on ROT13 applies here, since it's actually a type of Caesar Cipher. The level info also provides sufficient information on the cipher. However, unlike ROT13, unless the shift is 13 letters, we cannot decrypt the cipher by simply running the encryption algorithm again on the ciphertext. Instead, we must perform a **left shift** to reverse the encryption's right shift.. Recall the ROT13 program I wrote earlier, where shifting a letter was performed with the following calculation:

*.Encrypted = (characterIndex + N) mod 26, where N = shift*

A left shift is performed with a similar algorithm, simply changing the addition to a subtraction.

*.Decrypted = (characterIndex - N) mod 26, where N = shift*

## Method

We are provided a program that will encrypt anything with the same key used to encrypt the password. In terms of a Caesar Cipher, the key would simply refer to the number of positions each letter is shifted. Using this, we can find the key:

1. Follow the example provided in the level info to create a temporary working directory.
2. Create a file that simply contains the letter 'A':

```
krypton2@bandit:/tmp/tmp.xDNBQOSI0c$ echo 'A' > a.txt
krypton2@bandit:/tmp/tmp.xDNBQOSI0c$ ls
a.txt  keyfile.dat
krypton2@bandit:/tmp/tmp.xDNBQOSI0c$ cat a.txt
A
krypton2@bandit:/tmp/tmp.xDNBQOSI0c$ |
```

3. Encrypt a.txt and open the produced ciphertext:

```
krypton2@bandit:/tmp/tmp.xDNBQOSI0c$ /krypton/krypton2/encrypt a.txt
krypton2@bandit:/tmp/tmp.xDNBQOSI0c$ ls
a.txt  ciphertext  keyfile.dat
krypton2@bandit:/tmp/tmp.xDNBQOSI0c$ cat ciphertext
Mkrypton2@bandit:/tmp/tmp.xDNBQOSI0c$
```

4. Encrypting 'A' shifts it to 'M'. **Therefore, we can deduce that the Caesar Cipher shifts letters by 12 positions.**

I've written a simple program to decipher Caesar Ciphers, it's essentially the same as the ROT13 program with modified shift operations and allows for user-inputted shift amounts:

```

1  #!/usr/bin/env python3
2
3  lower = {
4      "a" : 1, "b" : 2, "c" : 3, "d" : 4, "e" : 5, "f" : 6, "g" : 7, "h":8,
5      "i" : 9, "j" : 10, "k" : 11, "l" : 12, "m" : 13, "n" : 14, "o" : 15,
6      "p" : 16, "q" : 17, "r" : 18, "s" : 19, "t" : 20, "u" : 21, "v" : 22,
7      "w" : 23, "x" : 24, "y" : 25, "z" : 26
8  }
9
10 l_key_list = list(lower.keys())
11
12 upper = {
13     "A" : 1, "B" : 2, "C" : 3, "D" : 4, "E" : 5, "F" : 6, "G" : 7, "H":8,
14     "I" : 9, "J" : 10, "K" : 11, "L" : 12, "M" : 13, "N" : 14, "O" : 15,
15     "P" : 16, "Q" : 17, "R" : 18, "S" : 19, "T" : 20, "U" : 21, "V" : 22,
16     "W" : 23, "X" : 24, "Y" : 25, "Z" : 26
17 }
18
19 u_key_list = list(upper.keys())
20
21 LETTERS = 26
22
23 message = input("Enter the ciphertext: ")
24 shift = int(input("Enter the shift amount: "))
25 decoded = ''
26
27
28 for character in message:
29     # calculate (charpos - 13) mod 26
30     if character in lower:
31         new_pos = (lower[character] - shift) % LETTERS
32         decoded = decoded + l_key_list[new_pos - 1]
33     elif character in upper:
34         new_pos = (upper[character] - shift) % LETTERS
35         decoded = decoded + u_key_list[new_pos - 1]
36     else:
37         decoded = decoded + character
38
39 print(decoded)

```

Now, feeding it the ciphertext **OMQEMDUEQMEK**, and the shift amount of 12, we get the password **'CAESARISEASY'**.

```

Enter the ciphertext: OMQEMDUEQMEK
Enter the shift amount: 12
CAESARISEASY

```

Now that we've completed the level using this more long winded route, I wanted to point out that it's also extremely easy to complete via simple brute forcing. In the English language, there are only 25 possible shifts for the Caesar Cipher (since there are 26 letters in the alphabet). Therefore, you could ignore the process of using the encryption program and just try all shifts until it works. This is made easy using readily accessible websites such as [Cryptii](#).

## Level 3 - Level 4

Link: <https://overthewire.org/wargames/krypton/krypton3.html>

### Level Info

Well done. You've moved past an easy substitution cipher.

The main weakness of a simple substitution cipher is repeated use of a simple key. In the previous exercise you were able to introduce arbitrary plaintext to expose the key. In this example, the cipher mechanism is not available to you, the attacker.

However, you have been lucky. You have intercepted more than one message. The password to the next level is found in the file 'krypton4'. You have also found 3 other files. (found1, found2, found3)

You know the following important details:

- The message plaintexts are in American English (\*\* very important)
- They were produced from the same key (\*\* even better!)

Enjoy.

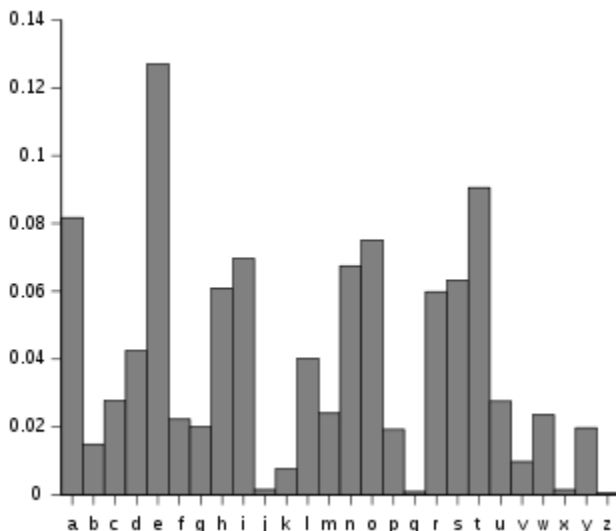
### Preamble

This challenge is slightly more difficult than the previous ones because the cipher mechanism is not provided to us. Therefore to complete it we must learn about and apply cryptanalysis techniques, namely **Frequency Analysis** along with trial and error.

### Research

According to the [OWASP Foundation](https://owasp.org/), "Cryptanalysis is a process of finding weaknesses in cryptographic algorithms and using these weaknesses to decipher the ciphertext without knowing the secret key (instance deduction)". We will be learning how to conduct some basic cryptanalysis to determine the cipher mechanism.

A common cryptanalysis technique is "frequency analysis". It involves counting the frequencies of letters in the ciphertext, and comparing them to the expected frequencies of letters in the corresponding alphabet. All written languages have certain patterns; for example, 'E' is the most common letter in the English alphabet. Below is a graph depicting a typical distribution of letters in written English:



Source: [https://en.wikipedia.org/wiki/Frequency\\_analysis](https://en.wikipedia.org/wiki/Frequency_analysis)



It's also possible to do frequency analysis on pairs of letters (bigrams), triplets (trigrams) and other N-letter combinations (N-grams).

Using this information, it is possible to make certain deductions about the cipher. For example, if the frequency of letters were consistently shifted then it could be concluded that a Caesar Cipher was used as well as the number of shifts. For example, suppose that the frequency analysis of a ciphertext showed that 'H' was the most common letter, 'W' the second most common, 'D' the third most common etc. Therefore, we could conclude that a Caesar Cipher with a 2 letter right shift was used.

## Method

Firstly, we'll perform a frequency analysis on the ciphertext to obtain information about its language patterns. Using this information, we'll brute force the solution.

I wrote a simple [Python program](#) to perform single letter frequency analysis. It prints out each character sorted in order of the number of times they appear in the ciphertext.

```
1  #!/usr/bin/env python3
2
3  frequencies = {}
4
5  ciphertext = input("Enter the Ciphertext: ")
6  ciphertext = ciphertext.replace(" ", "")
7  ciphertext = ciphertext.upper()
8
9  for character in ciphertext:
10     if character.isalpha():
11         frequencies[character] = frequencies.get(character, 0) + 1
12
13  for character, frequency in sorted(frequencies.items(), key = lambda x:x[1], reverse = True):
14     print(f"{character}: {frequency} ")
```

Examining the krypton4 file, we see that the ciphertext is **KSVVW BGSJD SVSIS VXBMN YQUUK BNWCU ANMJS**. Furthermore, there are a number of other files that also contain 5-letter group ciphers, encrypted using the same key.

Inserting the ciphertext into the cryptanalysis program, we obtain the following result:

```
Enter the Ciphertext: KSVVW BGSJD SVSIS VXBMN YQUUK BNWCU ANMJS
S: 6
V: 4
B: 3
N: 3
U: 3
K: 2
W: 2
J: 2
M: 2
G: 1
D: 1
I: 1
```

The letter 's' appears most commonly. However the frequency of the other letters are too close to each other to make any conclusions. Therefore, I also analysed the found files. Their results of the first two are pasted below:

found1:

```
Enter the Ciphertext: CGZNL YJBEN QYDLQ ZQSUQ NZCYD SNQVU BFGBK GQUQZ QSUQN UZCYD SNJDS UDCXJ ZCYDS NZQSU QNUZB WSBNZ QSUQN
JCBGZ CYDSN CGKDC ZDSQZ DVSJJ SNGGJ DSYVQ CGJSO JCUNS YVQZS WALQV SJJSN UBT SX COSWG MTASN BXYBU CJCBG UWBKG JDSQV YDQAS JXBN
M WCMUZ QSUQN KDBMU SWCJJ BZBTT MGCZQ JSKCT DDCUE SGNQ VUJDS SGNL YJCBG UJSYV SNXBN TSWAL QZQSU QNZCY DSNCU BXJSG CGZBN YB
GCG JDSNB JULUJ STQIK CJDQV VUCGE VSQVY DQASJ UMAUJ CJMJC BGZCY DSNUJ DSZQS UQNZC YDSNC USQUC VLANB FSGQG WCGYN QZJZC SBXXS
QGWDC USQNV LYVQL UKSNS TQCGV LZBTS WCSUQ GNDUJ JBNCS UESGN SUDSN QCUSW JBIDS YSQFB XUBYD CUJZC QJCBG QGWQN JCUNJ LALJD SSGW
Q GWTQZ ASJDZ BGUCW SNSWU BSBXJ JDSXC GSUJS OQTYV SUGGJ DSSGE VCUVJ QGEMQ ESCGD CUVQJ JYDQU SSKN BJSJN QECZB TSWCS UQVUB FG
JCB GUBXI QNLGQ EHMVQ CJLQG WQZZM NQZLW MNCGE DCUVC XSJCT SQGWC GJKBB XDCUX BNTSN JDSQJ NCZQV ZBVVS QEMSU YMAVC UDSWJ DSXCN
EGCUS WQUUD QFSUY SQNSU
S: 155
C: 107
Q: 106
J: 102
U: 100
B: 87
G: 81
N: 74
D: 69
Z: 57
V: 56
W: 47
Y: 42
T: 32
X: 29
```

found2:

```
Enter the Ciphertext: QVJDB MEDGB QJJSQ WQGSZ NSZBN WUXBN JDSYS NCBMU MNICI STBUJ ACBEN QYDSN UQENS SJDQJ UQFES UYSQN SKQUS WMZQJ SWQJJ DSFCG EUGSK
UZDBB VCGUJ NQJXB NQXN SSUZD BBVZD QNJSN SWCGQ ABMJQ HMQNJ SNBXQ TCYSX NBTDC UDBTS ENQTT QNUZD BBVUI QNCNW CGHMV VCJLW MNCGE JDSV CPQAS JDQGS NQAM
J JDSZM NNCZM VMTKQ UWCZJ QJSHA LVQKJ DNBME DBMDS GEVQG WQGNJ DSUZD BBVKB MWDQ ISYNB ICWSW QGCGJ SGUCI SSWMZ QJCBG CGVQJ CGENQ TTQNJ GWJDS ZVQUU CZ
UQJ JDSQE SBXUD QFSUY SQNST QNNCS WJDSL SQNBV WQGSZ DQJQJ KQLJD SZBGU CUJBN LZBMN JBXJD SWCBZ SUSBX KBNZS UJSNC UUMSW QTQNN QCESV CZSGZ SBGG BISTAS
NJKBB XDQJD QKQLU GSCED ABMNU YBUBS WABGW UJDSG SOJWQ LQUUM NSJLJ DQJJD SNSKS NSGBC TYSWC TSGUJ JBIDS TQNNC QESJD SZBMY VSTQL DQISQ NNQGE SWJDS ZSNS
T BGLCG UBTSJ QJJSU CGZSJ DSKBN ZSUJS NZDQG ZSVVB NQVVB KSWJD STQNN QESA QGGUJ BASNS QNBGZ SCGUJ SQNBX JDSMU MQVJD NSSJC TSUQG GSUYN SEGQ ZLZBM VW
DQI SASSG JDSNS QUBGX BNJDC UUCOT BGJDU QXJSN JDSQJ NNCQE SUDSE QISAC NJDJB QWQME DJSNU MUQGG QKDBK QUAQY JCUSW BGTQL JKCGU UBGDQ TGSJQ GWNQM EDJSN
RMWJ DXBVV BKSWQ VTBUT JKBL SNUVQ JSNQG WKSNS AQYJC USWBG XSANM QNLQJ TGSJW CSWBX MGFGB KGZQM USUQJ JDSQE SBXQG WKQUA MNCNW BQWME MUJQX JSNJD SACH
J DBXJD SJKCG UJDSN SQNSX SKDCU JBNCZ QVJNQ ZSUBX UQFES UYSQN SMGJC VDSUJ TSGJC BGSWQ UYQNJ BXJDS VBGWB GJDSQ JNSUZ SGSCG ASZQM USBXJ DCUEQ YUZDB VQ
NUN SXSNJ BJDSL SQNUA SJKSS GQGNQ UUDQF SUYSQ NSUVB UJLSQ NUACB ENQYD SNUQJ JSTYJ CGEJB QZZBM GJXBN JDCUY SNCBW DQISN SYBNJ SWTQG LQYBZ NLYDQ VUJBN
CSUGC ZDBVQ UNBKS UQFES UYSQN SUXCN UJACB ENQYD SNNSZ BMGJS WQUJN QJXBN WVS ES GWJQJ JUDQF SUYSQ NSXVS WJDSJ BKGXB NVBGW BGJBS UZQYS YNBUS ZMJC B GXBN
W SSNVB QZDCG EQGBJ DSNCS EDJSS GJDSZ GJNML UJBNL DQUUD QFSUY SQNSU JQNJC GEDCU JDSQJ NCZQV ZQNS NTCGW CGEJD SDBNU SUBXJ DSQJN SYQJN BGUCG VBGWB GR
BDG QMANS LNSYB NJSNJ DQJUD QFSUY SQNSD QWASS GQZBM GJNLU ZDBBV TQJJS NUBTS JKSGJ CSJZD SGJMN LUZDB VQNUJ QISUM EESUJ SWJQJ JUDQF SUYSQ NSTQL DQISA
SSGST YVBLS WQUQU ZDBBV TQJJS NALQV SOQGW SNDBE DJBGB XVQZ QUDCN SQZQJ DBVCZ VQGNB KGSNK DBGQT SWQZS NJQCG KCWVC QTUDQ FSUDQ XJSCG DCUKC WGBS ICWS
G ZSUMA UJQJG CQJJS UMZDU JBNCB UBIDS NJDQG DSQNU QLZBV VSZJS WQXJS NDCUW SQJD
S: 243
Q: 186
J: 158
N: 135
U: 130
B: 129
D: 119
G: 111
C: 86
W: 66
Z: 59
V: 53
M: 45
T: 37
E: 34
X: 33
Y: 33
K: 30
L: 27
```

The results consistently show that ‘S’ is the most common letter in the ciphers. Therefore, it is likely that the letter ‘E’ has been substituted for ‘S’. However the results for the other characters are still inconsistent. The encrypted password’s second-most common letter is ‘V’, while the ‘found’ files’ most common letters are either C or Q. These inconsistencies rule out the possibility of the cipher being a simple shift substitution cipher (such as a Caesar cipher). This is because we’d expect to see essentially the same frequency distribution as the usual English alphabet but shifted to different letters.

Having stumbled upon a dead end, I decided to look towards the trigram analysis of the found files. Since my program isn’t capable of doing the analysis, I used the [CrypTool](#) website. Via this, we see that the most common trigram is JDS:



## Graphical Frequency Analysis

N-gram analysis that works interactively and graphically

Analysis

Description

Background

Language frequency

Text to analyze

Restrict to English alphabet

GWTQZ ASJDZ BGUCW SNSWU BTSBX JDSXC GSUJS OQTYV SUCGJ DSSGE VCUDV QGEMQ ESCGD CUVQU JYDQU SDSKN BJSJN QECZB TSWCS UQVUB  
FGBKG QUNBT QGZSU QGWZB VVQAB NQJSW KCJDB JDSNY VQLKN CEDJU TQGLB XDCUY VQLUK SNSYM AVCUD SWCGS WCJCB GUBXI QNLCG EHMV  
CJLQG WQZM NQZLW MNCGE DCUVC XSJCT SQGWC GJKBB XDCUX BNTSN JDSQJ NCZQV ZBVVS QEMSU YMAVC UDSWJ DSXCN UJXBV CBQZB VVSZJ  
SWSWC JCBGB XDCUN NQTQJ CZKBN FUJDQ JCGZV MWSWQ VVAMJ JKBX JDSYV QLUGB KNSZB EGCUS WQUUD

Text length: 1529

N-gram:

-

3

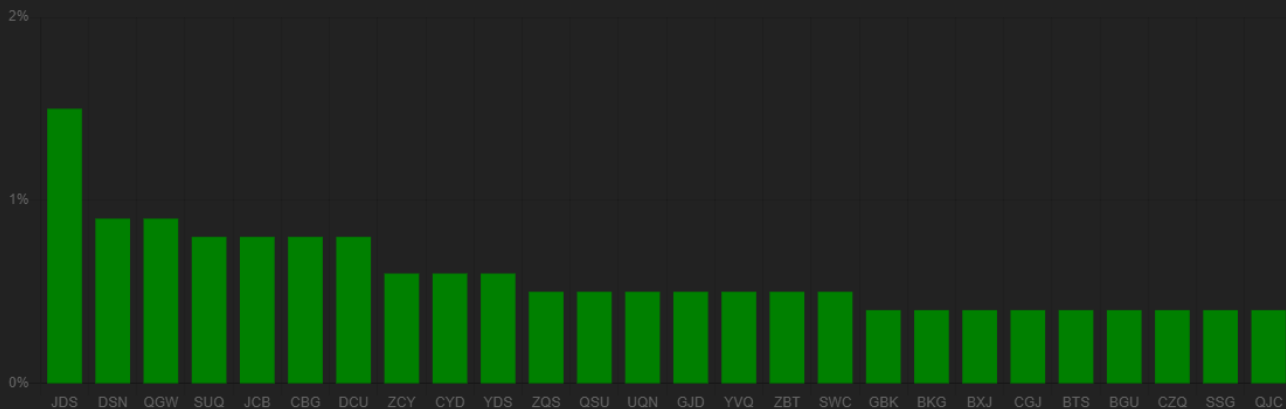
+

Number of bars:

26

Sort by frequency

Sort alphabetically



Now, the most common trigram in the English language is 'THE'. Therefore, we can deduce that '**THE**' is mapped to '**JDS**'.

Enough information has been obtained such that we can now brute force the solution. I'll highlight a few of my steps below:

- ☐ Use the 'tr' command to undo the 'THE' -> 'JDS' mapping:

```
ditto@Roblox-Hacking-Station:~$ echo 'KSVVW BGSJD SVSIS VXBMN YQUUK BNWCU ANMJS' | tr 'JDS' 'THE'
KEVVW BGETH EVEIE VXBMN YQUUK BNWCU ANMTE
```

- ☐ Seeing as 'THE' is a word, let's try rearranging the ciphertext a bit: 'KEVVW BGE **THE** VEIE VXBMN YQUUK BNWCU ANMTE'
- ☐ Consider KEVV, if we assume it is a 4 letter word with its 2nd letter being E, then it's possible that 'L' was mapped to 'V'. Undoing this, the ciphertext is now: '**KELL** W BGE **THE** LEIE LXBMN YQUUK BNWCU ANMTE'.
- ☐ Substitute 'K' for 'W' since it looks like the first word could be 'well': '**WELL** WBGE **THE** LEIE LXBMN YQUUW BNWCU ANMTE'.
- ☐ Continue this process, and we finally reach the conclusion that the plaintext is '**WELL DONE THE LEVEL FOUR PASSWORD IS BRUTE**'.

## **Level 4 - Level 5**

Link: <https://overthewire.org/wargames/krypton/krypton4.html>

### **Level Info**

Good job!

You more than likely used some form of FA and some common sense to solve that one.

So far we have worked with simple substitution ciphers. They have also been 'monoalphabetic', meaning using a fixed key, and giving a one to one mapping of plaintext (P) to ciphertext (C). Another type of substitution cipher is referred to as 'polyalphabetic', where one character of P may map to many, or all, possible ciphertext characters.

An example of a polyalphabetic cipher is called a Vigenère Cipher. It works like this:

If we use the key(K) 'GOLD', and P = PROCEED MEETING AS AGREED, then "add" P to K, we get C. When adding, if we exceed 25, then we roll to 0 (modulo 26).

P P R O C E E D M E E T I N G A S A G R E E D \  
K G O L D G O L D G O L D G O L D G O \  
becomes:

P 15 17 14 2 4 4 3 12 4 4 19 8 13 6 0 18 0 6 17 4 4 3\  
K 6 14 11 3 6 14 11 3 6 14 11 3 6 14 11 3 6 14 11 3 6 14\  
C 21 5 25 5 10 18 14 15 10 18 4 11 19 20 11 21 6 20 2 8 10 17\  
So, we get a ciphertext of:

VFZFK SOPKS ELTUL VGUCH KR

This level is a Vigenère Cipher. You have intercepted two longer, english language messages (American English). You also have a key piece of information. You know the key length!

For this exercise, the key length is 6. The password to level five is in the usual place, encrypted with the 6 letter key.

Have fun!

### **Research**

Vignere ciphers are a type of polyalphabetic substitution ciphers, where each alphabetic character in the plaintext is encrypted using its position and the position of the corresponding letter in the key. This is perfectly demonstrated in the level info, where the plaintext 'PROCEED MEETING AS AGREED', is encrypted using the key 'GOLD'.

When counting the position of letters in the context of Vignere Ciphers, it is important to note that A is the 0th letter, not the 1st.

## Method

Examining the krypton5 file, we see that the password's ciphertext is **HCIKV RJOX**. The two intercepted messages are:

```
krypton4@bandit:/krypton/krypton4$ cat found1
YYICS JIZIB AGYYX RIEWV IXAFN JOOVQ QVHDL CRKLB SSLYX RIQYI IOXQT WXRIC RVVVKP BHZXI YLYZP DLCDI IKGFJ UXRIP TFQGL CWVXR
IEZRV NMYSF JDLCL RXOWJ NMIXN FNJSP JGHVV ERJTT OOHMR VMBWN JTXKG JJJXY TSYKL OQZFT OSRFN JKBIY YSSHE LIKLO RFJGS VMRJC
CYTCS VHDLC LRXOJ MWFYB JPNVR NWUMZ GRVMF UPOEB XKSDL CBZGU IBBZX MLMKK LOACX KECOC IUSBS RMPXR IPJZW XSPTR HQRQB VVOHR
MVKEE PIZEK SDYVI QERJJ VYSLJ VZOVU NJLOW RTXSD LYYNE ILMBK LORVW VAOXM KZRNL CWZRA YGWWH DLCLZ VVXFF KASPJ GVIKW WWTVT
MCIKL OQYSW SBAFJ EWRII SFACC MZRVO MLYYI MSSSK VISDY YIGML PZICW FJNMV PDNEH ISSFE HWEIJ PSEEJ QYIBW JFMIC TCWYE ZWLTK
WKMBY YICGY WVGBS UKFVG IKJRR DSBJJ XBSWM VVYLR MRXSW BNWJO VCSKW KMBYY IQYYW UMKRM KKLOK YVWVX SMSVL KWCAV VNIQY ISIIB
MVVLI DTIIC SGRSX EVYQC CDLMZ XLDWF JNSEP BRROO WJFMI CSDDF YKQWM VLKWM KKLOV CXKFE XRFBI MEPJW SBWFJ ZWGMA PVHNR BKZIB
GCFEH WEWSF XKPJT NCCYR TUICX PTPLO VIJVT DSRMV AOWRB YIBIR MVWER QJKWK RBDYF MELS F XPEGQ KSPML IYIBX FJXPX ELPVH RMKFE
HLEBJ YMWKM TUFII YSUXE VLJUX YAYWU XRIUJ JXGEJ PZRQS TJIJS IJIJS PWMKK KBEQX USDXC IYIBI YSUXR IPJNM DLBFZ WSIQF EHLR
YVVMY NXUSB SRMPW DMJQN SBIRM VTBIR YWPSW IIIIC WQMV LKHNZ SXMLY YIZEJ FTILY RSFAD SFJIW EVNWZ WOFWJ WSERB NKAHW LTCXS
KCMXV OILGL XZYPJ NLSXC YYIBF ZGFRK VMZEH DSRTJ ROGIM RHKQP TCSCX GYJKB ICSTS VSPFE HGEQF JARMR JRWNS PTKLI WBWVW CXFJV
QOYVQ UGSXW BRWCS MSCIP XDFIF OLGUS ECXFJ PENZY STINX FJXVY YLISI MEKJI SEKFJ IEXHF NCPSI PKFVD LCWVA OVCSF JKVXK ESBLM
ZJTCM LYVMC GMZEX BCMKK LOACX KEXHR MVKBS SSUAK WSSKM VPCIZ RDLCF WXOVL TFRDL CXLRC LMSVL YXGSK LOMPK RGOWD TIXRI PJNIB
ILTKV OIQYF SPJCW KLOQQ MRHOW MYIED FCKFV ORGLY XNSPT KLIEL IKSDS YSUXR IJNFR GIPJK MBIBF EHVEW IFAXY NTEXR IEWRW CELIW
IVPYX CIOTU NKLDL CBFNS QYSRR NXFJJ GKVCH ISGOC JGMXX UFKGR krypton4@bandit:/krypton/krypton4$
```

```
krypton4@bandit:/krypton/krypton4$ cat found2
YYIIA CWVSL PGLVH DSAFD TYRYV YEDRG LYXER BJIEV EPLVX BICNE XRIDT IICXD TIXRI PJNIB ILTYS EWCXE IKVRM VXBIC RRHOE ETFHD
LGHBB YZCZW RQXMU RSDIA YKLOQ DMFQD LCIVA KRBYI IDMLB FSNQY STLYT NJUEQ VCFKT SPCTW AYSBB ZXRLG XRBOE LIUSB SRMPF EMJYR
WZPCS UMNJG WVXRE IBRVW IBMVU KRBRR HOLCW WIOPI JJWVS LJCCC LCFEH DSRTT XOXFJ CECXM KKLOM PGIIK HYSUR YAMVH HSHLT KOXSU
BYEDX FJPAY YJIUS PSPGI IKODF JXSJW TLASW FXRNM XFJCM YRGBZ PVKMN EXYXF JWSBI QYRRN OGQCE NICWV SBCMZ PSEGY SISKW RNKFI
XFJWM BIQNE GOCMZ IKXWR JJEI B QTGIM YJNRV DLYYP SETPJ WIBGM TBINJ MTUEX HRMVR ISSBZ PVLVA VEFIP DXSYH ZWVEU JYXKH YRRUC
IKWCI FRDFC LXINX FJMKX AMTUQ KRQYX SEPBH VVDEG SCCGI CUZJI SSPZP VIBFG SYVBJ VVKRB YYIXQ WORAC AMZCH BYQYR KKMLG LXDLC
QZSXA CSKEG EWNEX YXFJW SBIQY RRNMJ ZEHRM QTNRC YNUVU KRBSF SXICA VVURC BNLKX GYNEC JMWYI NMBSK QORRN FRXSY SUXRI QHRVO
GPTNJ YYLIR XBICK LPVSD SLXCE LIWMM PCUIS BSRMP WLEQP VXGMR MKLOQ QTKLK XQMVV YYJIE SDFCM LRQVW KFKVP MSXXS QCXYI DLMZX
LDXFN JAKWT JICUM LIRRN XFTLK RDXZC SPXFG JGKVC HISGF SYJLO PYZXL OHFJR VDMJD RDXLC FNOGE PINEI MLBYM MLRMV TYPSP IIXXS
WVTSJ IJUJY XFJEY DMFJN TKHBJ ULKRB XNIBI QTTPE QQDRR NXFJE YDMUJ IICSQ RRPVX FFKLO HPTGT OHYQD SCXYX DEXCY XYIZY RNEXR
IZFJO OXZZK XRIQH RVQGP TNHSH LTKQS RBMFA VLLZ XDSMP YMWXM KZPVX FJSEC OCYWS BMRJE ELPIC YMWXM PVIZE UFPJB SKYYI PMPJR
WRIDJ RVOHY XGEBU KNXLD KCYZR DSFNJ WDVYB RNRFS WELSQ SUJSR IJGK KKMTU HSWRF EGOEU FPJBS KYYIP PYRVW KRBT E PIGYR VROEP
YFGYZ CWUSB SRMPA SXFII CVIYA VWGLC SJLOP YDUSG RRTJP OINYY ICIIJ GXRIP AVVIW LZSEX HUFIQ KRBXY ICPCU KWYYL ICCER RNCQY
VLNEK GLCSZ XGEQI RCVME MKXRI ENIPL ERMVH RPKR GOMLF CDMXJ JIMZT JNEKL VMTBE XHQT F RKJRJ IXRIW FCPCX YWKIN XMBRV NXFJV
QOYVQ UGSXW YVMCA YXSLJ IYSVZ ORRLK PNEWK FVCLC YIEFI JJIWD LCDYE NLYWU PIFCJ EAKPI NEKKR FTLVG LCSKL OCQFN FOJMW VXRIK
FXVOE RIZXM LRMRX VMXJ INXFJ ISKH Y SUHSZ GIVHD LCKFV OWRFJ JKVYX KLOCA TLPNW CJFRO MRMVV CMBJZ XGEQF MIBCU NUINM RHYEX
HUMVR DLCDT VOTRZ GXYXF JVHQI YSUYI SIJUM XMMNK XRIWH FVYHQ JMVDA YXRPC STJIC NICUR RNXFJ IIGIP JDEXC ZNXNK KEJUV YGIXR
XDLCG FXDSK YYICM BJJA OVCXFW DICUK LKXLT EIJJR MVQMS SQUGV MKGUS GRYSU JYVVR FQORR NKWOI KJUXR ERYVI SVHTL VXIWR LWDIL
INLXX QMRPV ACIFE COCIU SBSRM PHOWN FZVSR EQPMR ETJEX DLCKR MXXCX KMNIY XRMNX FJKMX AMTUQ KRYSU XRIJN FRCLM TBSW QMRKQ
CKFEI KRBQF SUIBY YSEKF YWYVF SYKLO WAFII VMVBJ ESHUJ TEXRM YWPX FFKMC GCWKE SRLJZ XRIPH RRGIA QZQLH MBEMX XMYM CKPJR
XNMRH YXRI P JWSBI GKNIM ELSFX TYKUF ZOYGV NIWYQ YJXYT UMVVO ACFII SXFNE OSGMZ CHTYK UFZOV GYJES HRMVG YAYWU PIPGT EEPXC
WDIKW SWZRQ XFJUM CXYST IMEPJ WYVPW NELSW KNEHD LCSNI KVCFC PBHEM KEXWU JIINX FJJGK VCHIS GJMWP SEGVS TEBVW ZJEVP MAVVY
RWTLV LEAPF ROERF KMWU JCPSP JYICS XQFZH DLQZ SXAFT NMVPE TWMBW RNMVU PBJTP KVCIK LOWAF IIMVM BWSBM DDFYP SSSUX RERDF
YMSSQ URYXH ZDTYZ CWKLO KSQWH YVMYY CGSSQ UFOOG QCINS PYYID MLBFS NQYSS ENPWI VRDIB TEXRI PTTOC FCQFA LYRNW MQQMS PSEVZ
FTOSX UNCPX SRRRX DIPXF QEGFK FVDLC KRPVA MZCHX SRMLV DQCFK EVPkrypton4@bandit:/krypton/krypton4$
```

We know that the key length is 6 letters long from the level info. Therefore, it is possible to deduce that every 6 letters in the cipher, we'll see the same shift. For example, have a look at the example Vignere Cipher in the level info:

```
P R O C E E D M E E T I
G O L D G O L D G O L D
V F Z F K S O P K S E L
```

In this case, every 4 letters from the start of the key, the plaintext is shifted by the letter 'G'. Shifting over one position, every 4 letters the plaintext is shifted by 'O'. This applies for every letter in the key.

In other words, a simple Caesar cipher shift occurs every  $n$  letters in a Vignere cipher, where  $n$  refers to the length of the key.

Therefore, if we have the length of the key, we can form individual strings compiled from the letters in the ciphertext that have been shifted by the same letter in the key. In the example, the start of two of these strings would be: ["VKK", "FSS"], as highlighted above. Performing a basic frequency analysis on one of these individual strings will theoretically yield the letter that it was shifted by. Hence, performing frequency analysis on all of the strings will yield the key.

Now, I've written a [Python program](#) to perform this task for us. Here it is below:



```

1  #!/usr/bin/env python3
2
3  key_length = input("Enter the key length: ")
4  key_length = int(key_length)
5  ciphertext = input("Enter the ciphertext: ")
6  ciphertext = ciphertext.replace(" ", "")
7
8  string_list = []
9
10 for step in range(key_length):
11     curr_string = ''
12     for pos in range(step, len(ciphertext), key_length):
13         curr_string = curr_string + ciphertext[pos]
14     string_list.append(curr_string)
15
16 print(string_list)

```

I inserted a substring from the found1 file into the program to demonstrate it:

```

python.exe" f:/Code/Krypton/split2.py
Enter the key length: 6
Enter the ciphertext: YYICS JIZIB AGYYX RIEW IXAFN JOOVQ QVHDL CRKLB SSLYX RIQYI I
OXQT WXRIC RVVKP BHZXI YLYZP DLCDI IKGFJ UXRIIP TFQGL CWVXR
['YIYWNQRLYTRHYDJTW', 'YZYVJVKYIWVZZIUUFV', 'IIXIOHLXIXVPIXQX', 'CBRXODBRORKIDKRGR'
, 'SAIAVLSIXIPYLGIL', 'JGEFQCSQQCBLCFCPC']
PS F:\Code\Krypton>

```

The first string in the array is made of letters in the ciphertext that are shifted by the first letter in the key, the second string made of letters shifted by the second letter in the key, etc.

Now, inputting all of found1's ciphertext into the program gives us:

```

P XDFIF OLGSI ECFJ PENZY STINX FJXVY YLISI MEKJI SEKFJ IEXHF NCPST PKFVD LCWA OVCSE JKVXK ESBM ZJICM LYYMC GMZEX BCMKK LO
ACX KEXHR MKVBS SSUAK WSSKM VPCIZ RDLCE WXOVL TFRDL CXLRC LMSVL YXGSK LOMPK RGOWD TIXRI PJNIB ILTKV OIQYF SPJCW KLOQQ MRHOW
MYVED FCKFV ORGLY XNSPT KLIEL IKSDS YSUXR IJNFR GIPJK MBIBF EHVEW IFAXY NTEXR IEWRW CELIW
['YIYWNQRLYTRHYDJTWZSLNNHTMJJYFNVIJJSLSLWNMFXXBBKXIMJTBMIYJJNTYBWKWWLFGWI
SJSZSYSPJNFJQFWTYWKJMMNSYWKYSAYMTSQZJRFDKMXFJJPKFSTTTJMBMJDSQIJPJFJT
SJWJPIKXISJFFYXMQMYIMZYFSJWJNTWGJYGZTMTYSFFJTWJQBSFSJSJIJJNKWSXZYKX
MSSIFTXSSKTJTYWMYKLTISNJFITWI', 'YZYVJVKY
IWVZZIUUFVRFVRFVJVTJKTJSGKVRVZUKZKKUPZRVVZYJXJXNKVZZVZKWKWIEFRYKYZNEEPYMYKYVFRJVRWYKUKVWVIVIRCNRMFVKKBWZVEFNUPVWYKFFKYP
VEYUJUJZIKUYUNZEVUPNPIVKYTFIZWKCXLYFEJRCJTEJRWVURCIUPTXSIICFVFEJYEKKVUKZWFLVKRINKFKRYFYKKUFKEFERW', 'IIXIOHLXIXVPIXQXV
JXISVOMXLOKLSCHXYRGPXGLESXWVKEIRZLSELARRHVAITLSWAVIVIIIMHHSIIEWIGVRXVXJWIMLWLVSVIXCLSOIYLLFISWHTXICILTAIWMYXSIXHHMFXXXRJ
SKSIXMHVSWSTWILSIIAWWSASVXLIHRHRSKSHAWLWQGWIFEEIVISEPVAJSIMXLEKAMRXXRLLGXIVSLHEVXLSXRMHAXW', 'CBRXODBRORKIDKRGRNDONPEOBKYOS
BHOVYDOBNDODUMOCBRXKXOEYQOODIOONADVSKVOBRCOMIGCVIWEBCZKCBGDBYSOKQXKNILCEDDEOCKKOEMBGBKWKYCODOBEKMPBRRLWIEYRGQSPBDBRDSLMBD
BBSCXKZLDEOEKXOZSBKDKOCBVGNRNICOSCPOCNMYEXSDOKBCCBOXBKVDODCYOORBOPOODONIDRGBVXRC', 'SAIAVLSIXIPYLGILIMLWXJRHMGTOQRIERMTLJJWE
LILAOSTSRHESSESWLLRLXLYLXPWMQAICMSSMWPSEEWMTMGSISSLWMYRKSIIISVLWPWSWWXEWMRGEPYXVSWIRREEMXEMEKIVAIESIWEXIILIYYSMIIPWHMEYSW
RWKIYXWVSGPXISEMSWVXSXLXZYXKHLVLMGCAHSWPLVLLXMMWIIJQWFRSESIIEYIE', 'JGEFQCSQQCBLCFCPCCEYCFJGJRNJZSFYLFRCMPUMBCEBMCCRPPQRP
DRLURYMYMCGCFJWCYFIMLSDLFDSIJJCLBYUKBWRBCBYMYMCQBDGYMFBJDQMCPRFABCWJRPJRRRQBLGLFLKBMVLYUTJTMQCYPBQRNRJRIQNLJRFNFBLCPCZMRIQ
GCPQRPBFYWMDFYFLKFFPCCKMLMCRSSCLCMGPDPLQCMGPLYJPBWNEL']
PS F:\Code\Krypton>

```

You may need to zoom into the image to see the output

Let's perform a frequency analysis on each string using the frequency analysis program, like we did in the previous level.

### First String:

YIYWNQRLYTRHYDJTWZSLNNHTMJJYFNVIJJSLSLWNMFXXBBKXIMJTBMIYJJNTYBWKWWLFGWI  
SJSZSYSPJNFJQFWTYWKJMMNSYWKYSAYMTSQZJRFDKMXFJJPKFSTTTJMBMJDSQIJPJFJT  
SJWJPIKXISJFFYXMQMYIMZYFSJWJNTWGJYGZTMTYSFFJTWJQBSFSJSJIJJNKWSXZYKX  
MSSIFTXSSKTJTYWMYKLTISNJFITWI

The single letter frequency analysis is:

```
Enter the Ciphertext: YIYWNQRLYTRHYDJTWZSLNNHTMJYFNYIJJSLWNMFXBKXIMJTBMIYJJNTYBWKWWLFGWISJZSYSPJNFJQFWTYWKJJMMNSYWKYSAYMTSQ
ZJRFDKXKFJJPKFSTTTJMBMDJSQIJPFJTSJWJPJIKXISJFFYXMQMYIMZYFSJWJNTWJYGTMTYSFFJTWJQBSFSJSJIJJNKWSXZYKXMSIFTXSSTJTYWJMYKLTISNJF
ITWI
J: 35
S: 23
Y: 22
T: 20
W: 17
F: 17
M: 16
I: 14
N: 11
K: 11
X: 8
Z: 7
Q: 6
B: 6
L: 5
P: 4
R: 3
D: 3
G: 3
H: 2
A: 1
```

The most common letter in the string is 'J'. In typical plaintext English, the most common letter is 'E'. Therefore, we can deduce that E has been shifted to J in the ciphertext. E is the 4th letter in the alphabet, whereas J is the 9th letter (remembering that A is the 0th letter) . Therefore, a Caesar shift of 5 positions has been performed on this string; The 5th letter in the alphabet is F, so we can conclude that **every letter in the first string has been shifted by F and so the first letter of the Vignere cipher key is F**.

We now repeat this process for the other 5 strings in our list. I'll show the process for the next string and move on.

### Second String:

**YZYVJVKYIWVZZIUFRFRMJVTVTJKTJSKGCVRFVZUKZZKKUPZRVVZYJVJXNKVZZVZKVVKW  
EFRYKYZNEEPYMYKYVFRJVRWKYUKVVVIVIRCXNRMFVKKBWZVZEFNUPVVVYKFFKYVVEYU  
UUUJZIJKUYUNZEVUPNPVIVKYTFIZWKXCXLYFEJRCJTEJRKVVURCIUPTXSIICFVFEJYEKKVU  
KZWFLVKRINKFKRYFYKKUFKEFERW**

```
Enter the Ciphertext: YZYVJVKYIWVZZIUFRFRMJVTVTJKTJSKGCVRFVZUKZZKKUPZRVVZYJVJXNKVZZVZKVVKW
EFRYKYZNEEPYMYKYVFRJVRWKYUKVVVIVIRCXNRMFVKKBWZVZEFNUPVVVYKFFKYVVEYU
UUUJZIJKUYUNZEVUPNPVIVKYTFIZWKXCXLYFEJRCJTEJRKVVURCIUPTXSIICFVFEJYEKKVU
KZWFLVKRINKFKRYFYKKUFKEFERW
V: 35
K: 30
Y: 19
Z: 18
F: 18
U: 16
R: 15
J: 13
E: 12
I: 11
N: 8
W: 7
P: 7
T: 6
C: 6
X: 4
M: 3
S: 2
L: 2
G: 1
B: 1
```

V is the most common letter. E is the 4th letter in the alphabet, whereas V is the 21st letter.  
 $21 - 4 = 17$ , and hence a Caesar Shift of 17 positions was performed on the string. **The 17th letter is R and so we can conclude that the second letter of the key is R.**

After analyzing all of the other strings, **the key is discovered to be FREKEY**. We can now easily decipher the password by subtracting the index of each letter in the ciphertext with the corresponding letter in the key (similar to what we just did):

**Ciphertext: H C I K V R J O X**

**Key: F R E K E Y F R E**

**Password: C L E A R T E X T**

**H - F = 7 - 5**

**= 2**

**= C**

**C - R = 2 - 17**

**= -15**

**= 26 - 15**

**= 11**

**= L**

**I - E = 8 - 4**

**= 4**

**= E**

Note that what we've done above is actually equivalent to performing a modulo 26 operation. For example:

$(7 - 5) \bmod 26 = 2$

$(2 - 17) \bmod 26 = 11$

## **Level 5 - Level 6**

Link: <https://overthewire.org/wargames/krypton/krypton5.html>

### **Level Info**

FA can break a known key length as well. Lets try one last polyalphabetic cipher, but this time the key length is unknown. Note: the text is written in American English

Enjoy.

### **Research**

A method of deducing the key length of a polyalphabetic substitution cipher is the Kasiski Examination. It involves searching for repeated strings in the ciphertext, then finding the distance between the beginning of each of the found strings. The factors of the distance value are then considered; the key length is one of these factors. An analyst will narrow down the factors based on logic and trial and error.

Kasiski Examination works based on the assumption that if a repeated string occurs in the plaintext, and the distance between said string is a multiple of the key length, then they'll be encrypted in the exact same way.

Consider the following example from the [Wikipedia entry](#) on Kasiski Examination:



Key: ABCDABCDABCDABCDABCDABCDABCD  
Plaintext: *cryptoisshortfor cryptography*  
Ciphertext: CSASTPKVSIQUTGQUCSASTPIUAQJB

In the ciphertext, we can see that the strings CSASTP are repeated. By counting each letter from the start of the first CSASTP string, we deduce that their distance is 16. The factors of 16 are 16, 8, 4, 2, 1. The key length is likely not 16 as this seems too long. On the other hand, the key length also likely isn't 1 or 2 as these are too short. Hence, we only need to try 8 and 4 as the key length.

If the ciphertext is longer, and multiple instances of different repeated strings are found, we can count the distance between each of these strings. The greatest common factor of the distances is likely to be the correct key length. An example of this is demonstrated in the method below.

## Method

Examining the found1 file returns the following ciphertext.

```
SXULW GNXIO WRZJG OFLCM RHEFZ ALGSP DXBLM PWIQT XJGLA RI/RI BLPPC HMYMG CTZDL CLKRU YMYSJ TWUTX ZCMRH EFZAL OTMNL BLULV  
MCCMG CTZDL CPTBI AVPML NVRJN SSXVT XJGLA RIQPE FUGVP PGRLG OMDKW RSIFK IZYRM QHNXD UOWQT XJGLA RIQAV VTZVP LMAIV ZPHCX  
FPAVT MLBSD OIFVT PBACS EQKOL BCRSM AMULP SPYPF CXOKH LZUO GNLID ZVRAL DOACC INREN YMLRH VXXJD XMSIN BXUGI UPVRG ESQSG  
YKQOK LMXRS IBZAL BAYJM AYAVB XRSIC KKPYP ULWFU YHBPG VIGNX WBIQP RGVXY SSBEL NZLVW IMQMG YGVSX GPWGG NARSP TXVKL PXWGD  
XRJHU SXQHI VTZYO GCTZR JYVBK MZHBX YVBIT TPTVM OOWSA IERTA SZCOI TXXYL JAZQC GKPCS LZRYE MOOVC HIEKT RSREH MGNIS KVEPN  
NCTUN EQFIR TPPDL YAPNO GMLGC ZRGNX ARVMY IBLXU QPYH GNXYO ACCIN QBUQA GELNR TYQIH LANTW HAYCP RJOMO KJYTV SGVLY RRSIG  
NKVXI MQJEG GJOML MSGNV VERRC MRYBA GEQNP RGLBL XFLRP XRZDE JESGN XSYVB DSSZA LCXYE ICXXZ OVTPW BLEVK ZCDEA JYPCL CDXUG  
MARM L RWVTZ LXIPL PJKKL CIREP RJYVB ITPVV ZPHCX FPCRG KVPSS CPBXW VXIRS SHYTU NWCGI ANNUN VCOEA JLLFI LECSO OLCGT CMGAT  
SBITH PNZBV XWUPV RIHUM IBPHG UXUQP YYHNZ MOKXD LZBAK LNTCC MBJTZ KXRSI FSKZC SSEL P UMARE BCIPK GAVCY EXNOG LNLCC JVBXH  
XHRHI AZBLD LZWIF YXKLM PELQG RVPAP ZQNVK VZLCE MPVKP FERPM AZALV MDPKH GKXCL YOLRX TSNIB ELRYN IVMKP ECVXH BELNI OETUX  
SSYGV TZARE RLVEG GNOQC YXFCX YOQYO ISUKA RIQHE YRHDS REFTB LEVXH MYEAJ PLCXK TRFXZ YOZCY XUKVV MOJLR RMAVC XFLHO KXUVE  
GOSAR RHBSS YHQS LXSJ INXLH PXCCV NVIPX KMFV ZLTOW QLKRY TZDLC DTVXB ACSDE LVYOL BCWPE ERTZD TYDXF AILBR YEYEG ESIHC  
QMPX UDM LZ VVMBU KPGE EGIWO HMFV NXPBW KPVRS XZCEE PWVTM OOIYC XURRV BHCCS SKOLX XQSEQ RTAOP WNSZK MVDLC PRTRB ZRGPZ  
AAGGK ZIMAP RLKVV EAZRT XXZCS DMVVZ BZRWS MNRIM ZSRX IEOVH GLGNL FZKH KCESE KEHDI FLZRV KVFIB XSEKB TZSPE EAZMV DLCSY  
ZGGYK GCELN TUIIG MXQHT BJXG ZRFEX ABIAP MIKWA RVMFK UGGFY JRSIP NBUI LDSSZ ALMSA VPNTX IBSMO krypton5@bandit:/krypton  
/krypton5$
```

As circled in red and blue, there are two repeating words: **MGCTZDLC** and **TXJGLARI**. Both of these words are 8 letters long, and have a **distance of 45 characters** in between them, excluding whitespaces. The factors of 45 are: 45, 15, 9, 5, 3, and 1. 45 and 15 are likely too long for the key length, whereas 1 and 3 are likely too short. Hence, by method of elimination, **the key length is likely 9 characters long**. To confirm this, let's examine the ciphertext in the found2 file:

```
GLCYX UKFHS PEZXF AVJOW QQYYR RAYHM GIEOG ARIAZ YEYXV PXFPJ BXUY SLELR NXNNH PLARX TADLC CSLGE NOSPR IUUML VSNPR RJMOO  
GMLGU JHVBE QSMFI NZDSK HEFNX KSHGE AVZAZ YQCQP BAKPC LMQGR XXTYR WQSEG FHSPH ZYETX FVPMX PBTWV XMLHM AZXYG EQLRN IAPDZ  
CXIAZ MVMSL RVNZN SKXCL RNJOL XXSCS HYMYK ZCWPR XNWYR ZJXUG MASQC ELRXX DKWMY PLUGL KHTPR GAKVE WRCEI KESOV JPJGH XJYRE  
CEGAE HDIBQ SEZAL DAMZX UKKZR EBMIR TLLDH MHRNZ MOOMP CIFVX JDMTP VBGWZ SHCOI FZBUK XGZRF ZALWM JOIJE BUCMB PSSZA LMSYN  
LJOMO SXQOE ZVTUN HGCXL YMYKA GEWQO LHQIC LFYKL TOPJL RQOMZ YFQNY EOMFG EQCEG NXYVM IPEYG KNOVB ZKXKG UOPKC PBXKF DLCAE  
FYXUQ IPDLN QBUQL GXWRR VYEXM QMGOG JREGY WBLLA BEULX NTZSO SDDL N MZFGV YATRX YSKTN TRTNT AKRBX YQJRS OKQHE FXSTAR IPWMX  
KTSKV EPVFU KAYJB ZKGNX YOAGW POKTW KGIPX GUVHV EGDXB SHYBS UOVNC XYIIQ DMEY ARIUP EGNXY RSJOW NTWAR IUTRQ YXACX MWIEG  
USOJY TVGNX ASHCH MYRL L BZCAV RZMFV MAPPL GMHLS SEXJU BUDLC LJGKK UYSLD MEHXX CMPTW UGESX SRRSG UULNX GWPAO ZODFS EMJGG  
AKFCO VBUFH XHYME EHXYK RBELR TUYOE IQEFZ LPBCC DWVXM OKXUL CFOKP PCMTF YKTZO WFZAP UGJYV BRIAZ ELWEL DZNRB ZOELD LBZPH  
DIPES PUGJY VBAYY RHMPK CYXYK FHXWZ ZSGYB UMSLN SEJRV EAGWP SOGKK JGYIF KTJYE JQMEK LPBJC EGUHT YLIPE SPUGJ YVBDX VXTIY  
YREL R XXUYA DZVPU GJYVB ELRIH UMSP O FRJVO KQZPV OKBUQ EJHEL YTZCM EYIQZ HZHEQ DIAMX YLCRS IZGBS KRBAE FYXUQ IPDFL ZALWE  
GWFRO GNKPU LCFNX HFMJJ AEGIW OHSAJ EUFOO EBESS UHADL CCSBS AHNXF PSQJB UDIPP WGLHY DLCPW GGUSS WFXIA ZHMDL CCSLG ENOSP  
RIGNT AKPRS SHMAI EXMYI XOGKY JKLRJ GLZOI LESTU BUDSG EERYD PXHQL RQBTY SIRT I FUYTO RALQR UNAYJ GEGBT LLAYC XYXET UYXFP  
VQXTD OVYYH GCHWY VRPVF GKKCI TPNVR FHS HQ LRQZA LVELO PNJRD OVCLP YRHPD IPTRT HRRMG GOIAZ TAFEP TSHYI VSRRD SSZAL BSYOF  
RZPLO RRSIP UGJYV BLRQZ ALMSD QIRXH VWAFF RNMXU DPCXE AUYSZ BRJJB XFHVP WOVRY LLNML LFEUP UCYGE SSIEV DLCDT EKMAI ACWPJ  
UKULY RGIEE PLVPI PTGCB ARPYC KRYJB KVCNY SLLHX HJLVT KYSKT QESGN XWYGI PXFVT ZCIBL PBTZV XLGDA NEMVR MQMVR GDMKW Rkrypt  
on5@bandit:/krypton/krypton5$
```

As circled in red, the word **SPUGJYVB** is repeated. It is 8 letters long, and there is a distance of 81 characters in between them. The factors of 81 are: 81, 27, 9, 3 and 1. **9 is the greatest common divisor of the two distances, 81 and 45**. Hence, we have obtained evidence to further reinforce the proposition that the key length is 9.

From here on, the steps to decipher the password ciphertext, which is **BELOS Z**, are identical to the method utilized in the previous level. I'll briefly illustrate the steps here

Insert the ciphertext of one of the found files into the [split](#) program:

```

Enter the key length: 9
Enter the ciphertext: SXULW GNXIO WRZJG OFLCM RHEFZ ALGSP DXBLM PWIQT XJGLA RIYRI BLPPC HMXMG CTZDL CLKRU YMYSJ TWUTX ZC
MRH EFZAL OTMNL BLULV MCQMG CTZDL CPTBI AVPM L NVRJN SSWWT XJGLA RIQPE FUGVP PGRLG OMDKW RSIFK TZYRM QHNXD UOWQT XJGLA RI
QAV VTZVP LMAIV ZPHCX FPAVT MLBSD OIFVT PBACS EQKOL BCRSM AMULP SPYPF CXOKH LZUO GNLID ZVRAL DOACC INREN YMLRH VXXJD XM
SIN BXUGI UPVRG ESQSG YKQOK LMXRS IBZAL BAYJM AYAVB XRSIC KKPYP ULWU YHBPV VIGNX WBIQP RGVXY SSBEL NZLWV IMQMG YGVSU GP
WGG NARSP TXVKL PXWGD XRJHU SXQMI VTZYO GCTZR JYVBK MZHBX YVBIT TPVTM OOWSA IERTA SZCOI TXXLY JAZQC GKPCS LZRYE MOOVC HI
EKT RSREH MGNTS KVEPN NCTUN EOFIR TPPDL YAPNO GMLGC ZRGNX ARVMY IBLXU QPYH GNXYO ACCIN QBUQA GELNR TYQIH LANTW HAYCP RJ
OMO KJYTV SGVLY RRSIG NKVXI MQJEG GJOML MSGNV VERRC MRYBA GEQNP RGKLB XFLRP XRZDE JESGN XSYYB DSSZA LCXYE ICXXZ OVTPW BL
EVK ZCDEA JYPC L CDXUG MARM L RIWVTZ LXIPL PJKKL CIREP RJYVB ITPV ZPHCX FPCRG KVPSS CPBXW VXIRS SHYTU NWCIG ANNUN VCOEA JL
LFI LEC SO OLCTG CMGAT SBITP PNZBV XWUPV RIHUM IBPHG UXUQP YYH NZ MOKXD LZBAK LNTCC MBJ TZ KXRS M FSKZC SSEL P UMARE BCIPK GA
VCY EXNOG LNLCC JVBX HXRHI AZBLD LZWIF YXKLM PELQG RVP AF ZQNVK VZLCE MPVKP FERPM AZALV MDPKH GKKCL YOLRX TSNIB ELRYN IV
MKP ECVX H BELNI OETUX SSGV TZARE RLVEG GNOQC YXFCX YOQYO ISUKA RIQHE YRHS REFTB LEVXH MYE AJ PLCXK TRFZX YOZCY XUKW MO
JLR RMAVC XFLHO KXUVE GOSAR RHBSS YHQUS LXSDJ INXLH PXCCV NVIPX KMFV ZLTOW QLKRY TZDLC DTVXB ACSDE LVVOL BCWPE ERTZD TY
DXF AILBR YEYEG ESIHC QMPOX UMLZ VVMBU KPGE C EGIWO HMFVG NXPBW KPVRS XZCEE PWVTM OOIYC XURRV BHCCS SKOLX XQSEQ RTAOP WN
SZK MVDLC PRTRB ZRGPZ AAGGK ZIMAP RLK W EAZRT XXZCS DMVZ BZRWS MNRIM ZSRVX IEOVH GLGNL FZKH KCESE KEHDI FLZRV KVFIB XS
EKB TZSPE EAZMV DLCSY ZGGYK GCELN TTUIG MXQHT BJXG ZRFEX ABIAP MIKWA RVMFK UGGFY JRSIP NBJUI LDSSZ ALMSA VPNTX IBSMO
[ 'SOCGWRCDYCOVDPSPKYPORPCBBPKLOYDGMBBYBBSWSRSOBBOSYOSKNDKRQOQAKYXOEGBDYCOKCRXCBCPXNNFCBXMQKNKCRVNXDLPFZFCNVBXRROORSX
COMCVBXPXODCDBIDKOBKOBXODRZWCERSGXDFZDKIKBRYUMB', 'XWMSIIHLSMTLMXIPWRWILXSACSHIAMXISXAXHPISISWSXGKIWZJSVRVELGVPAAIYJRIM
REXEVXVZLMIITXSIWVITWIPXTXSECLHLMALIMLIMESEQQIRHXZOXESSXKWLSCTRHMPHWEITHXPLGIESWRLKITS LGGXIVJISS', 'URRPQYMCJRMCC L WQGRMQ
QMFD CRPLDCLMUGRYRUGQBMGPQCMTSCALCEE OYCMYCGHCYRMLRQFJB YTCCLPRTFSRCCLGTUB YDCRSBYCRZPFCDYBKL SRCYQEMKCJFGSDCMQCDWYCLGMKEY
CQWC PMADSYGCFBPCCMGAMRLAM', 'LZHDTRXLTHNQPNTPRSQTAAPSSPZZCRSPYSJSLVPEQPTDMT ZTAOZZHHPFAZYCELP TSQMCNLEDEPDRLP PCSGOECPP
PYLCESECEHWEZEPPOEPNYLYOHFYTYLLOYJCF LDEPDEQZEFPPCCSNPZAZMMXNEXSEXPFS DVO', 'WJEXXIMKWELMTVXELIHXVIAIEMXYVIHVKIMIWIRLM
WXXIZHP IIRIMNIPRIHILARVIJSMPRSSIWEIXWPPVCPSEICMPVHHZMMLIXJII LQMMLLEIGVXIETERXRHSHIVXKTL EXYMVXVWXSES RAPRVNLSZSEYLQRMKI
SP', 'GGFB JBG RUFBGBRJFGFNJVVVFQAFURNVNRQBACFGGNGGVVRVBVETCYEGNRNGBGNNNJSGEGRRPGSCBAUVJRVRBHAASGNRGNBBFPNNVAFQNP AHRRCOV EF
SYBAFUROAQNNVRVVEFEPVEGRVUSQZTARTVREFEREAZNHFIUPS N', 'NOZLGLCUTZLCIJGUOKXGTZTVKMCOARXBGOZYKUNVZYGKJTJXTRXGEKNCTONLNQRTOG
NGNYGXNZXLJGTKJZGXYNJOAZIUZAJ SUKOBZYGWVZGXVYETGCURLJZKMKRUXVZYXYRAGOMGNSTRKRRLGXZIOZKVKZGTTEKGNZT', 'XFAMLP TYXAUTANLGMT
DLZPMTOUXGLEXXEKA AKYXXLGNLHZYYMTXKMTTPGXXXBTWMMVKGVBKRXAXEYMKZYPKWTLNLOTBHMKTGMGGXBXRRKAKTNTXZGXKHEPVAXRSLILTBOTIEXBIDX
MROTMBGXKBMMVKEKBMTBXWGBAX', 'ILLPAPZMZLLZVSAVDZUAVHLLP LOND NJUSLLVPHWYVAPUYVVOALPORSUPMAUYVYHOLVJVALZSLZVPALLVHVUULLSV
UJOLZZAALHLKWLPLKSIHUNYADVLYVUHLHPTZALZSUJWPZOV LAVZKVZZZHHIVTVYUJAAFJLT'] ]

```

Perform a frequency analysis on each of the resultant strings. For example, inputting the second string into the frequency analysis program shows that I is the most common letter, meaning that E has been Caesar shifted to I. This is a shift of 4 positions, a shift by the letter E in other words. Hence, the second letter of the key is E.

After analysing all 9 strings, we find that the key is supposedly XEYLENGTH. This can obviously be corrected to **KEYLENGTH**. I can't clearly explain why the first letter turned out to be X, but luckily it was an easy correction.

Now, the ciphertext, **BELOS Z**, can easily be deciphered:

**Ciphertext = B E L O S Z**

**Key = KEYLENGTH**

**Password = RANDOM**

**B - K = 1 - 10**

**= - 9**

**= 26 - 9**

**= 17**

**= R**

**E - E = 4 - 4**

**= 0**

**= A**

**L - Y = 11 - 24**

**= -13**

**= 26 - 13**

**= 13**

**= N**

## Level 6 - Level 7

Link: <https://overthewire.org/wargames/krypton/krypton6.html>

## Level Info

Examine the link since the contents are really long.

## Preamble

The final level involves cracking a **stream cipher**. Similar to the second level, we are given a program, `encrypt6`, that will encrypt any plaintext using the same key that was used to encrypt the password.

## Research

In modern cryptography, there are two types of ciphers: block ciphers and stream ciphers. The substitution ciphers we dealt with in previous levels share similarities with modern block ciphers which work by grouping plaintext into blocks of specific sizes (128 bits for example) and encrypting them all at once. Stream ciphers on the other hand, encrypt the plaintext one bit at a time. Stream ciphers consist of three main parts:

- Plaintext
- Keystream: A pseudorandom sequence of bits produced by inputting a specific key into a pseudorandom number generator.
- Ciphertext: The ciphertext is produced by performing XOR operations on the plaintext bit-by-bit with the keystream.

The keystream being pseudorandom is an important detail. Pseudorandom number generators create statistically random digit strings, but these strings are not truly random. This is because they are determined via an initial value, known as the “seed value”. In this case, the seed value is the key. Therefore, if the key used to generate the keystream is obtained, then the keystream itself can easily be found.

XOR operations are a type of Linear-feedback Shift Registers (LFSR) that when used on binary bits, work like so:

## Boolean Math: XOR ( $\oplus$ )

X	Y	$X \oplus Y$
0	0	0
0	1	1
1	0	1
1	1	0

Source: <https://www.secureideas.com/blog/2020/10/boolean-math-xor-logic-cissp-domain-3.htm>

## Method

Let us observe how the `encrypt6` function encrypts a string consisting of only A's:

```
krypton6@bandit:/krypton/krypton6$ python3 -c "print('A'*200)" > /tmp/plainA.txt
krypton6@bandit:/krypton/krypton6$ cat /tmp/plainA.txt
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
krypton6@bandit:/krypton/krypton6$ ./encrypt6 /tmp/plainA.txt /tmp/ciperA.txt
krypton6@bandit:/krypton/krypton6$ cat /tmp/ciperA.txt
EICTDGYIYZKTHNSIRFXYPFUEOCKRNEICTDGYIYZKTHNSIRFXYPFUEOCKRNEICTDGYIYZKTHNSIRFXYPFUEOCKRN
EICTDGYIYZKTHNSIRFXYPFUEOCKRNEICTDGYIYZKTHNSIRFXYPFUEOCKRNEICTDGYIYZKTHNSIRFXykrypton6@bandit:/krypton/krypton6$
```

The number of A's to test is arbitrary, I simply chose 200 because I felt that a longer plaintext would better illustrate potential patterns. This assumption was correct; the ciphertext repeats several times. To be precise,

the word “**EICTDGYIYZKTHNSIRFXYCPFUEOCKRNEICTDGYIYZKTHNSIRFXYCPFUEOCKRN**” is repeated. It is my suspicion that this is simply the key.

Let's test this theory by observing the program's behaviour on 200 B's instead:

```
PNUKLYLWRQKGKBekrypton6@bandit:/krypton$ python3 -c "print('B'*100)" > /tmp/plainB.txt  
krypton6@bandit:/krypton/krypton6$ python3 -c "print('B'*200)" > /tmp/plainB.txt  
krypton6@bandit:/krypton/krypton6$ ./encrypt6 /tmp/plainB.txt /tmp/cipherB.txt  
krypton6@bandit:/krypton/krypton6$ cat /tmp/cipherB.txt  
FJDUEHZJJALUIOTJSYGYZDQGVPDL SO FJDUEHZJJALUIOTJSYGYZDQGVPDL SO FJDUEHZJJALUIOTJSYGYZDQGVPDL SO FJDUEHZJJALUIOTJSYGYZDQGVPDL SO  
FJDUEHZJJALUIOTJSYGYZDQGVPDL SO FJDUEHZJJALUIOTJSYGYZDQGVPDL SO FJDUEHZJJALUIOTJSYGYZDQGVPDL SO FJDUEHZJJALUIOTJSYGYZ
```

We get repetitions of the word

“FJDUEHJZJZALUIOTJSGYZDQGVFPDLSOFJDUEHJZJZALUIOTJSGYZDQGVFPDLSO” instead. Notice that this is simply the first repeated word with each letter shifted forward by one position. Therefore, we can confirm that plainB.txt has been encrypted with the first word. From the start of the string, if B is encrypted by E, we get F (B is Caesar shifted to the right by E), if B is encrypted by I we get J etc.

Therefore, the ciphertext, **PNUKLYLWRQKGKBE**, can be decrypted easily using the same method used in the past two levels.

**Ciphertext = P N U K L Y L W R Q K G K B E**

Key = E I C T D G Y I Y Z K T H N S

Password = LFSR IS NOT RANDOM

$$P - E = 15 - 4$$

**= 11**

$$= L$$
$$N - I = 13 - 8$$

**= 5**

= F

$$U - C = 20 - 2$$

**= 18**

= S

**K - T = 10 - 19**

$$= -9$$
$$= 26 - 9$$

**= 17**

$$= R$$

The level info made this level much more difficult than it needed to be. I spent a while going in circles, trying to make sense of hexdump output to detect patterns in it and messing with online pseudorandom generators. I'm sure it is possible to complete this level using hexdump and other similar tools but I believe the method I used here is the most straightforward.

## Conclusion

With that, we have completed all of Krypton, congratulations! I believe that Krypton is a great resource for beginners to start their journey into cryptography. Through both self-research and practice, players learn the mechanisms of classical substitution ciphers and the cryptanalysis techniques used to crack them. Most importantly, the challenges force players to adopt an attacker's mindset which is something that all security people must learn to understand.

Thank you for following my guide, I hope that it was easy to understand and that you learnt something from it.

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