

Tutorial - 2

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Q1 What is the time complexity of below code and how?

```
void fun(int n)
{
    int j=1, i=0;
    while(i < n)
    {
        i+=j;
        j++;
    }
}
```

↳ $\left. \begin{array}{ll} j=1 & i=0 \\ j=2 & i=1+2 \\ j=3 & i=1+2+3 \end{array} \right\} m\text{-level}$

for(i)

$$\therefore 1+2+3+\dots + < n$$

$$\therefore 1+2+3+m < n$$

$$\frac{m(m+1)}{2} < n$$

$$m \approx \sqrt{n}$$

By summation method

$$\sum_{i=1}^{\infty} 1 \Rightarrow 1+1+\dots + \sqrt{n} \text{ times}$$

$$\boxed{T(n) = \sqrt{n}} \quad \text{--- Ans}$$

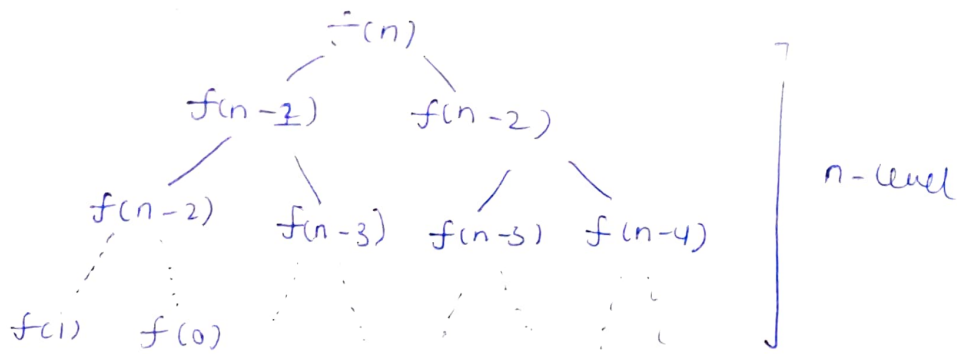
Q2 write recurrence relation for function that prints Fibonacci series. Solve it to get the time complexity. What will be the space complexity and why?

→ For Fibonacci series

$$f(n) = f(n-1) + f(n-2)$$

By forming a tree

$$\begin{array}{l} f(0) = 0 \\ f(1) = 1 \end{array}$$



∴ At every function call we get = Function calls
 ∴ for n levels
 we have = $2 \times 2 \dots n$ times

$$\therefore T(n) = 2^n$$

MAXIMUM SPACE

Considering Recursive

Stack:

no. of calls maximum = n

For each call we have space complexity $O(1)$

$$\therefore T(n) = O(n)$$

without considering Recursive Stack:
 each call we have time complexity $O(1)$

$$\therefore T(n) = O(1)$$

Q3 Write program which have complexity:
 $n(\log n)$, n^3 , $\log(\log n)$

1) $n \log n \rightarrow$ Quick Sort

```
void quicksort (int arr[], int low, int high)
```

```
{ if (low < high)
```

```
{ int pi = partition(arr, low, high);
```

```
  quicksort (arr, low, pi-1);
```

```
  quicksort (arr, pi+1, high);
```

```
} }
```

int partition(int arr[], int low, int high)

```
{ int pivot = arr[high];  
  int i = (low-1);  
  for (int j = low; j <= high-1; j++)  
  { if (arr[j] < pivot)  
    { i++;  
      swap(arr[i], arr[j]);  
    }  
  }  
  swap(arr[i+1], arr[high]);  
  return (i+1);  
}
```

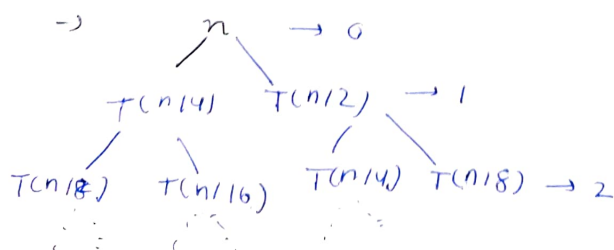
2 $n^3 \rightarrow$ Multiplication of 2 square number matrix

```
for (int i = 0; i < n; i++)  
{  
  for (j = 0; j < n; j++)  
  {  
    for (k = 0; k < n; k++)  
    { arr[i][j] += arr[i][k] * arr[k][j];  
    }  
  }  
}
```

3) $\log(\log n)$

```
for (i = 2; i < n; i = i * 1)  
{ count++;  
}
```

Q4 Solve the following recurrence relation
 $T(n) = T(n/4) + T(n/2) + cn^2$



At level

$$0 \rightarrow cn^2$$

$$1 \rightarrow \frac{n^2}{4^2} + \frac{n^2}{2^2} = \frac{5n^2}{16}$$

$$2 \rightarrow \frac{n^2}{8^2} + \frac{n^2}{16^2} + \frac{n^2}{4^2} + \frac{n^2}{8^2} = \left(\frac{5}{16}\right)^2 n^2$$

$$\text{max level} = \frac{n}{2^k} = 1$$

$$= k = \log_2 n$$

$$T(n) = c(n^2 + 5/16)n^2 + (5/16)^2 n^2 + \dots + (5/16)^{\log n} n^2$$

$$T(n) = cn^2 \left[1 + \left(\frac{5}{16}\right) + \left(\frac{5}{16}\right)^2 + \dots + \left(\frac{5}{16}\right)^{\log n} \right]$$

$$T(n) = cn^2 \times 1 \times \left(\frac{1 - (5/16)^{\log n}}{1 - (5/16)} \right)$$

$$T(n) = cn^2 \times \frac{11}{5} \times \left(1 - \left(\frac{5}{16}\right)^{\log n} \right)$$

$$T(n) = O(n^2)$$

$$O(n^2) \quad \underline{\text{Ans}}$$

Q5 > what is the time complexity of following fun()?

int fun(int n)

```
{
  for(int i=1; i<=n; i++)
  {
    for(int j=1; j<=n; j+=i)
    {
      Some task O(1)
    }
  }
}
```

c) 8^{2n} , $\log(n)$, $n(\log_6(n))$, $n(\log_2(n))$, $\log(n!)$, $n!$, $\log_8(n)$,
 96 , $8n^2$, $7n^3$, $5n$

$$\rightarrow 96 < \log_8 n < \log_2 n < 5n < n \log_6(n) < n \log_2 n < \log(n!) < 8n^2 < 7n^3 < n! < 8^{2n}$$

①