# **CS345A:** Design and Analysis of Algorithms Mock Quiz

Marks = 11 Date: 14 August 2023

NAME:			
ROLL No:			

Attempt any one of the following problems.

## 1. Two versions of Gayle and Shapley algorithm(6 marks)

Provide an instance of stable matching problem on a set of 3 men and 3 women such that the output of 'man proposing' version of Gayle Shapley is totally different from the output of 'woman proposing' version.

Note: You just need to provide only the following details in your answer.

- 1. The preference lists of 3 men and 3 women.
- 2. Output of the two versions of Gayle Shapley algorithm.

**Answer:** Let  $W_1, W_2$ , and  $W_3$  be three women and  $M_1, M_2$ , and  $M_3$  be three men.

The preference lists of women is the following.

	1	2	3
$W_1$	$M_1$	$M_2$	$M_3$
$W_2$	$M_2$	$M_1$	$M_3$
$W_3$	$M_3$	$M_1$	$M_2$

The preference lists of men is the following.

	1	2	3
$M_1$	$W_2$	$W_1$	$W_3$
$M_2$	$W_3$	$W_1$	$W_2$
$M_3$	$W_1$	$W_3$	$W_2$

The stable matching produced by the 'man proposing' version of Gayle Shapley is  $\{(W_2, M_1), (W_3, M_2), (W_1, M_3)\}$ .

The stable matching produced by the 'woman proposing' version is  $\{(W_1, M_1), (W_2, M_2), (W_3, M_3)\}.$ 

## 2. Non-dominated points in higher dimensions (11 marks)

You are given a set P of n points in 3-dimensions. A point  $q \in P$  is non-dominated if there is no point  $r \in P \setminus \{q\}$  that dominates q in each dimension. Assume without loss of generality that no two points in P have the same x-coordinates or y-coordinates or z-coordinates. Design an  $O(n \log n)$  time algorithm that computes all non-dominated points in P.

 $\mathit{Hint:}$  Process the points in the decreasing order of their z-coordinates. Under what conditions will ith point in this order be a non-dominated point? In order to achieve efficiency you might also like to make use of some well known data structure you learnt in ESO207 .

**Note:** You just need to describe the algorithm. There is no need to analyse its time complexity or prove its correctness.

#### Answer:

Let L be the list of points P sorted in the decreasing order of their z-coordinates.

Let  $p_i$  be the point at *i*th place in the list L.

Let  $P_i$  denote the set of first i points in the list L.

Let Proj(i) be the projection of points of set  $P_i$  on (x,y) plane, and let ND(i) be the set of non-dominated points of set Proj(i). We keep a height balanced Binary Search Tree T that, at the end of processing  $P_i$ , stores ND(i) using x-coordinate of points as the key. T initially stores point  $p_1$ . Note that  $p_1$  is surely a non-dominated point. We process points  $p_i$  with i > 1 in the increasing value of i using the following procedure.

### $\mathbf{Process}(p_i)$ :

```
q \leftarrow \operatorname{successor}(p_i,T); if q = null then
\operatorname{Print}(p_i \text{ is a non-dominated point});
\operatorname{Insert}(p_i,T); else
\operatorname{if}\ (y(q) < y(p_i)) \text{ then}
\operatorname{Print}(p_i \text{ is a non-dominated point});
q \leftarrow \operatorname{predecessor}(p_i,T);
\operatorname{while}\ (q \neq null \text{ and } (y(q) < y(p_i)) \text{ do}
\operatorname{Delete}\ q \text{ from } T;
q \leftarrow \operatorname{predecessor}(p_i,T);
\operatorname{end}\ \operatorname{while}
\operatorname{Insert}(p_i,T);
\operatorname{end}\ \operatorname{if}
```

**Note:** I have not added comments to the pseudocode above. For knowing proper comments, and explanation of the algorithm, please do attend the doubt clearing session at 12:00 noon on 19th August in RM101. You may write algorithm in English as well. However, it must be complete and unambiguous.