

## MODULE II

More Number Theory and Some Algebra;  $Z_n$  ( $Z \bmod n$ )

As we already know  $(Z, +, \circ)$  has "algebraic" structure.

Indeed, both  $+$  and  $\circ$  are binary operations on  $Z$  that are associative and commutative; 0 is the identity of  $(Z, +)$ , 1 is the identity of  $(Z, \circ)$  and the operations are linked by the distributive law  $a \circ (b + c) = a \circ b + a \circ c$ .

We put these ideas in a more general context and begin a discussion of  $Z_n$ , the integers module  $n$ , in this algebraic setting.

Definition 1 A binary operation  $*$  on a non-empty set  $S$  is a function

$*$ :  $S \times S \rightarrow S$ . The operation  $*$  is said to be associative if

$$\forall a, b, c \in S \quad (a * b) * c = a * (b * c)$$

If  $*$  is associative then  $(S, *)$  is called a semigroup. An element  $e$  is called an identity provided  $a * e = e * a = a \quad \forall a \in S$ .

If an identity exists it is unique (indeed  $e_1 = e_1 * e_2 = e_2$ ) and  $(S, *)$  is referred to as a monoid. An element  $a$  is said to be invertible (or to have an inverse) if  $\exists b$  such that

$$a * b = b * a = e$$

If an inverse exists it is unique (Proof - Exercise 1) and is denoted by  $a^{-1}$ . If all elements of the monoid  $(S, *)$  are invertible then  $(S, *)$  is a group

(Exercise 2: If  $T = \{a \in S \mid a \text{ has an inverse}\}$  then  $(T, *)$  is a group).

If  $*$  is commutative, i.e.  $\forall a, b \in S \quad (a * b = b * a)$

$(S, *)$  is referred to as a commutative semigroup or commutative monoid or commutative group as the case may be. If the operation is denoted by  $+$  and  $(S, +)$  is a commutative group it is usually referred to as an abelian group.

If the non-empty set  $R$  is equipped with two binary operations  $+$  and  $\circ$ , the triple  $(R, +, \circ)$  is called a ring if

(i)  $(R, +)$  is an abelian group with identity 0 (inverses are denoted with minus signs, i.e. the inverse of  $a$  has the name  $-a$ )

(ii)  $(R, \circ)$  is a semigroup

and

(iii) (distributive laws)  $\forall a, b, c \in R$

$$a \circ (b + c) = a \circ b + a \circ c$$

and

$$(b + c) \circ a = b \circ a + c \circ a$$

Exercise 3. (Submit this one) If  $(R, +, \circ)$  is a ring then  $\forall a \in R$

$$a \circ 0 = 0 \circ a = 0$$

Thus if  $(R, \circ)$  has an identity, say 1, and  $|R| \geq 2$  then  $1 \neq 0$

Prove these contentions.

If  $|R| \geq 2$  and  $(R, \circ)$  is a monoid, then  $(R, +, \circ)$  is a ring with identity.

An element  $a \in R$ , where  $(R, +, \circ)$  is a ring with identity, is called a unit if it has a multiplicative inverse. The set of units  $U \subseteq R - \{0\}$  and  $(U, \circ)$  is a group. It is called the group of units of  $R$ .

Exercise 4. Prove  $0 \notin U$  and  $(U, \circ)$  is group

If  $(R, \circ)$  is a commutative semigroup then  $(R, +, \circ)$  is called a commutative ring.

A commutative ring with identity such that  $U = R - \{0\}$  is a field.

Remark1.  $(\mathbb{Z}, +, \circ)$  is a commutative ring with identity such that  $U = \{1, -1\}$ .

Definition 2. Let  $n \in \mathbb{Z}^+$ . We write  $a \equiv b \pmod{n}$  if and only if  $n \mid a-b$

Terminology: a is congruent to b modulo n

Observation Congruence modulo  $n$  is an equivalence relation. Each equivalence class contains a unique element from  $\{0, 1, \dots, n-1\}$  called the least residue of the class and is determined by the division algorithm, i.e. the least residue of the equivalence class containing  $a$  is given by  $r$  where  $a = qn + r$  and  $0 \leq r \leq n-1$ . Thus there is exactly one equivalence class for each  $0 \leq r \leq n-1$  and we use these numbers to represent the equivalence classes.

Definition 3.  $Z_n = \{0, 1, \dots, n-1\}$  the integers modulo n, is the collection of equivalence classes referred to above.

Also  $a +_n b \triangleq c$

where  $a + b \equiv c \pmod{n}$

and  $a \circ_n b \triangleq d$

where  $a \circ b \equiv d \pmod{n}$

Example 1.  $Z_{36} = \{0, 1, 2, \dots, 35\}$

$$24 +_{36} 25 = 13$$

$$9 \circ_{36} 8 = 0$$

Exercise 5. (Submit a and b)

a) Let  $a, b, c, d \in \mathbb{Z}$  and  $n \in \mathbb{Z}^+$ . Prove

i)  $a \equiv b \pmod{n} \Leftrightarrow b \equiv a \pmod{n} \Leftrightarrow a - b \equiv 0 \pmod{n}$

ii)  $a \equiv b \pmod{n}$  and  $b \equiv c \pmod{n} \Rightarrow a \equiv c \pmod{n}$

iii)  $a \equiv b \pmod{n}$  and  $c \equiv d \pmod{n} \Rightarrow a + c \equiv b + d \pmod{n}$   
and  $ac \equiv bd \pmod{n}$

iv)  $a \equiv b \pmod{n}$  and  $d|n \Rightarrow a \equiv b \pmod{d}$

v)  $a \equiv b \pmod{n}$  and  $c > 0 \Rightarrow ac \equiv bc \pmod{cn}$

Remark 2 i) and ii) verify symmetry and transitivity of the equivalence relation:

$a$  is related to  $b$  if and only if  $a \equiv b \pmod{n}$

iii) verifies that  $+_n$  and  $\circ_n$  are well-defined

b) Prove i)  $ax \equiv ay \pmod{n} \Leftrightarrow x \equiv y \pmod{\frac{n}{\gcd(a,n)}}$

ii)  $ax \equiv ay \pmod{n}$  and  $\gcd(a,n) = 1 \Rightarrow x \equiv y \pmod{n}$

iii)  $x \equiv y \pmod{n_i} \quad i = 1, \dots, k \Leftrightarrow x \equiv y \pmod{\text{lcm}(n_1, \dots, n_k)}$

iv) if  $n_1, \dots, n_k$  are pairwise relatively prime then  $x \equiv y \pmod{n_i} \quad i = 1, \dots, k$

$$\Leftrightarrow x \equiv y \pmod{\prod_{i=1}^k n_i}$$

c) Prove: if  $b \equiv c \pmod{n}$  then  $\gcd(b, n) = \gcd(c, n)$

d) Prove: if  $p$  is prime and  $a^2 \equiv b^2 \pmod{p}$  then  $p|(a+b)$  or  $p|(a-b)$

e) Suppose  $f(x)$  is a polynomial with integer coefficients.

Prove: If  $f(a) \equiv k \pmod{n}$  then  $f(a + tn) \equiv k \pmod{n} \quad \forall$  integers  $t$ .

Remark 3. Even though the notation used in the preceding discussion and exercise is rather standard it tends to obfuscate the essential nature of the operations. In actuality the operations are very simple to describe. Indeed, we leave it to the reader to show that  $+_n$  and  $\circ_n$  can be defined as follows: if  $a, b \in \mathbb{Z}$  and  $\bar{a}, \bar{b}$  represent the equivalence classes containing  $a$  and  $b$  respectively then

$$\bar{a} +_n \bar{b} = \overline{a+b}$$

and  $\bar{a} \circ_n \bar{b} = \overline{a \cdot b}$

In some of the discussions to follow we drop the subscript "n" on the operations and simply use  $+$  and  $\circ$  to denote modulo  $n$  arithmetic.

Proposition 1  $(\mathbb{Z}_n, +, \circ)$  is a commutative ring with identity. The elements 0 and 1

are the respective identities of  $(\mathbb{Z}, +)$  and  $(\mathbb{Z}, \circ)$ . The group of units is denoted by  $\mathbb{Z}_n^*$ .

Remark 5. We shall not prove this proposition as it is straight – forward and TEDIOUS.

Proposition 2.  $Z_n^* = \{a \in Z_n \mid \gcd(a, n) = 1\}$

Proof Suppose  $a \in Z_n^*$  so that  $a \cdot a^{-1} = 1$  i.e.  $aa^{-1} \equiv 1 \pmod{n}$

Thus  $aa^{-1} = qn + 1$

so if  $k \mid a, n$  it follows that  $k \mid 1$ , i.e.  $\gcd(a, n) = 1$

Conversely, suppose  $\gcd(a, n) = 1$ . Then  $\exists x, y \in Z$  such that

$$ax + ny = 1$$

or  $ax \equiv 1 + (-y)n$ .

Hence  $ax \equiv 1 \pmod{n}$ .

Choose  $b \in \{0, 1, \dots, n-1\}$  such that  $b \equiv x \pmod{n}$

Then  $ab \equiv ax \equiv 1 \pmod{n}$  i.e.  $a^{-1} \equiv b \pmod{n}$

Corollary 2.1  $|Z_n^*| = \phi(n)$

Corollary 2.2 If  $p$  is prime then  $Z_p^* = Z_p - \{0\}$ , so  $Z_p$  is a field. Furthermore, if

$Z_n$  is a field then  $n$  is prime.

Exercise 6 (Submit b.)

a) Determine  $Z_{30}^*$

b) Prove the following generalization of the previous proposition:

Let  $d = \gcd(a, n)$ . Then  $\exists x \in Z_n$ , a solution of  $ax \equiv b \pmod{n}$ ,

if and only if  $d \mid b$ . In this case there are exactly  $d$  solutions in  $Z_n$  and they are all

congruent modulo  $n/d$ .

Example 2. In  $Z_9$ ,  $Z_9^* = \{1, 2, 4, 5, 7, 8\}$  where

$$1^{-1} = 1$$

$$2^{-1} = 5 \text{ (because } 2 \cdot 5 = 10 \equiv 1 \pmod{9}\text{)}$$

$$4^{-1} = 7 \text{ (because } 4 \cdot 7 = 28 \equiv 1 \pmod{9}\text{)}$$

$$8^{-1} = 8 \text{ (because } 8 \cdot 8 = 64 \equiv 1 \pmod{9}\text{)}$$

Consider the equation

$$3x \equiv b \pmod{9}$$

If  $b = 0$  then solutions are 0, 3, 6 (which are congruent mod 3)

If  $b = 3$  the solutions are 1, 4, 7 (which are congruent mod 3)

If  $b = 6$  the solutions are 2, 5, 8 (which are congruent mod 3)

No solutions exist otherwise, i.e. for  $b \in Z_9$ ,  $b \neq 0, 3, 6$ .

In the proof of the multiplicativity of the Euler-phi- function  $\varphi$  we established a special case of our next theorem, the important Chinese Remainder Theorem. We prove the theorem by using the so-called Gauss algorithm for computing the solution. It is actually an explicit formula for the solution which becomes intuitively appealing upon some inspection.

Theorem 1. Let  $n_1, n_2, \dots, n_k \in \mathbb{Z}^+$  be pairwise relatively prime. Then the system of simultaneous congruences

$$\begin{aligned} x &\equiv a_1 \pmod{n_1} \\ x &\equiv a_2 \pmod{n_2} \\ &\cdot \\ &\cdot \\ &\cdot \\ x &\equiv a_k \pmod{n_k} \end{aligned}$$

has a UNIQUE solution modulo  $n = n_1 n_2 \dots n_k$  (i.e. in  $\mathbb{Z}_n$ )

Proof First we prove existence. Let  $\hat{x} = \sum_{i=1}^k a_i \left( \frac{n}{n_i} \right) M_i$

where  $M_i \equiv \left( \frac{n}{n_i} \right)^{-1} \pmod{n_i}$ ,  $i = 1, \dots, k$ . Observe that  $\left( \frac{n}{n_i} \right)^{-1}$  exists

in  $\mathbb{Z}_{n_i}$  because  $n_i$  and  $\frac{n}{n_i}$  are relatively prime. Also

$$a_i \frac{n}{n_i} M_i \equiv a_i \pmod{n_i}$$

because  $\frac{n}{n_i} M_i \equiv 1 \pmod{n_i}$ . Moreover

$$a_j \left( \frac{n}{n_j} \right) M_j \equiv 0 \pmod{n_i} \text{ for } j \neq i$$

because  $n_i \mid \left( \frac{n}{n_j} \right)$ . Thus

$$\begin{aligned} \hat{x} &= \sum_{j=1}^k a_j \left( \frac{n}{n_j} \right) M_j \\ &\equiv a_i \pmod{n_i} \text{ for each } i = 1, 2, \dots, k. \end{aligned}$$

Let  $x \in \mathbb{Z}_n$  such that  $\hat{x} \equiv x \pmod{n}$

Because  $n_i \mid n$ ,

$$x \equiv a_i \pmod{n_i} \text{ for each } i = 1, \dots, k$$

so setting  $x = \hat{x} \pmod{n}$  we obtain a solution in  $\mathbb{Z}_n$ .

As for uniqueness suppose  $x_1$  and  $x_2$  are solutions and  $x_1, x_2 \in \mathbb{Z}_n$ .

Then, assuming  $x_1 \geq x_2$

$$x_1 - x_2 \equiv a_i - a_i = 0 \pmod{n_i}$$

Thus  $n_i \mid x_1 - x_2$   $i = 1, \dots, k$

But then  $n = n_1 \dots n_k \mid x_1 - x_2$  because  $n_1, n_2, \dots, n_k$  are pairwise relatively prime.

Finally then  $x_1 - x_2 = 0$  because  $0 \leq x_2 \leq x_1 \leq n-1$ .

Example 3. Consider  $n_1 = 7, n_2 = 13, n_3 = 15$  and

$$x \equiv 3 \pmod{7}$$

$$x \equiv 7 \pmod{13}$$

$$x \equiv 13 \pmod{15}$$

Then  $n = (7)(13)(15) = 1365$

$$\begin{aligned} \frac{n}{n_1} &= 195, \quad M_1 = (195)^{-1} \pmod{7} \\ &= (6)^{-1} \pmod{7} = 6 \end{aligned}$$

$$\begin{aligned} \frac{n}{n_2} &= 105, \quad M_2 = (105)^{-1} \pmod{13} \\ &= (1)^{-1} \pmod{13} = 1 \end{aligned}$$

$$\begin{aligned} \frac{n}{n_3} &= 91, \quad M_3 = (91)^{-1} \pmod{15} \\ &= (1)^{-1} \pmod{15} = 1 \end{aligned}$$

Thus

$$\begin{aligned} \hat{x} &= (3)(195)(6) + (7)(105)(1) + (13)(91)(1) \\ &= 1560 + 735 + 1183 \\ &= 5428 \end{aligned}$$

so  $\underline{\underline{x}} = \underline{\underline{1333}}$

Check:

$$\begin{aligned} 1333 &= (190)(7) + \underline{3} \\ 1333 &= (102)(13) + \underline{7} \\ 1333 &= (88)(15) + \underline{13} \end{aligned}$$

Exercise 7. Solve

$$\begin{aligned}x &\equiv 3 \pmod{5} \\x &\equiv 3 \pmod{7} \\x &\equiv 5 \pmod{12}\end{aligned}$$

for the unique solution in  $Z_{420}$ .

Remark 6. Of course this explicit expression for the solution provides for a convenient algorithmic computational procedure given algorithms for computing the terms in the expressions. The only computational procedure that deserves mention is that of computing inverses modulo  $n$ .

### Modulo $n$ Inverse Algorithm

Input:  $a \in Z_n$

Output:  $a^{-1} \bmod n$ , provided it exists

1. Apply the extended Euclidean algorithm to find  $(d, x, y)$  such that  
 $d = \gcd(x, y)$  and  $d = a x + n y$
2. If  $d > 1$  then  $a^{-1} \bmod n$  doesn't exist. Otherwise apply the Division Algorithm and return  $(x \bmod n)$ .

Exercise 8. Write an algorithm for the solution of the system of simultaneous congruences given pairwise relatively prime  $n_1, n_2, \dots, n_k$  and elements  $a_i \in Z_{n_i}$   $i = 1, \dots, k$ . What is its complexity?

Remark 7. We noted that  $\varphi(n) = |Z_n^*|$  so that  $Z_n^*$  is a FINITE group. Before examining  $Z_n^*$  in more detail we discuss some generalities regarding finite groups.

Definition(s) 4. Let  $(G, *)$  be a finite group. The order of  $G$  is just  $|G|$ . If  $H \subseteq G$  and the restriction of  $*$  to  $H$  renders  $(H, *)$  a group then  $H$  is called a subgroup of  $G$ , denoted by  $H \subseteq_g G$  (This is the case even if  $|G| = \infty$ ).

Proposition 3. Let  $(G, *)$  be a group.

- i)  $H \subseteq_g G$  if and only if  $\forall a, b \in H \Rightarrow a b^{-1} \in H$
- ii)  $|G| < \infty$  and  $H \subseteq_g G \Rightarrow |H| \mid |G|$
- iii)  $|G| < \infty$  and  $a \in G \Rightarrow \exists n \in \mathbb{Z}^+$  such that  $a^n = e$

The smallest such  $n$  is called the order of  $a$ , denoted  $\text{ord}(a)$

and  $\langle a \rangle = \{a^k \mid k \geq 1\} \subseteq_g G$  such that  $|\langle a \rangle| = \text{ord}(a)$ .

Thus  $\text{ord}(a) \mid |G|$  and so  $a^{|G|} = e$  as well. Furthermore, if  $a^k = e$  then  $\text{ord}(a) \mid k$ .

- iv) if  $a, b \in G$ , a finite group, such that  $a b = b a$  and

$$\gcd(\text{ord } a, \text{ord } b) = 1$$

then

$$\text{ord}(a b) = \text{ord}(a) \cdot \text{ord}(b)$$

In fact, if  $a_1, \dots, a_k \in G$  commute in pairs and  $\text{ord } a_1, \dots, \text{ord } a_k$  are pairwise relatively prime then  $\text{ord}(a_1 \dots a_k) = (\text{ord } a_1) \dots (\text{ord } a_k)$

Proof (i) ( $\Rightarrow$ ): Let  $e_H$  denote the identity of  $H$ . Then  $e_H = e_H \cdot e_H$ . But

$e_H^{-1}$  exists in  $G$  so

$$e = e_H^{-1} * e_H = e_H^{-1} * (e_H * e_H) = (e_H^{-1} * e_H) * e_H = e_H, \text{ i.e. the}$$

identity of  $G$  is also the identity of  $H$ . But then uniqueness of inverse implies that the inverse of an element of  $H$  in  $G$  is the inverse in  $H$ .

Hence  $a, b \in H \Rightarrow a, b^{-1} \in H \Rightarrow a * b^{-1} \in H$ .

( $\Leftarrow$ ): First choose  $a \in H$  so  $a, a^{-1} \in H$  and therefore,  $e = a * a^{-1} \in H$ .

Hence  $b \in H$  implies  $e, b \in H$  which implies  $b^{-1} = e * b^{-1} \in H$ .



Next  $a, b \in H$  implies  $a, b^{-1} \in H$  which implies  $a * b = a * (b^{-1})^{-1} \in H$

Since  $*$  is associative on  $H$  we have that  $H \subseteq_g G$ .

(ii) Define  $a R b$  if and only if  $aH = bH$  (where  $aH \triangleq \{a * h \mid h \in H\}$ ).

It is the case that  $R$  is an equivalence relation on  $G$ ; we only verify that  $aH = \{a * h \mid h \in H\}$  is the equivalence class containing  $a$ :

- $a \in aH$  because  $a = a * e$  and  $e \in H$
- $bR a$  if and only if  $b \in aH$  - indeed if  $b \in aH$

i.e.  $b = a * h$  for some  $h$ , then  $b * \hat{h} = a * (h * \hat{h}) \in aH \forall \hat{h}$  so  $bH \subseteq aH$ .

Likewise  $a = b * h^{-1} \in bH$  so  $aH \subseteq bH$  and so  $aR b$ . Conversely, if  $aH = bH$  then  $b = b * e \in bH = aH$ .

Next observe that  $|H| = |bH| \forall b \in G$

since  $H \rightarrow bH$  is easily seen to be  

$$h \rightarrow bh$$

a bijection. Thus  $G$ , being the pairwise disjoint union of equal size equivalence classes of common size  $|H|$ , has order which is divisible by  $|H|$ .

(iii) Since  $(a) = \{a^k \mid k \geq 1\} \subseteq G$ ,  $(a)$  is finite.

Thus  $\exists 1 \leq \ell < k$  such that  $a^k = a^\ell$ . But  $a^{-\ell} = (a^{-1})^\ell$  so  $a^k * (a^{-1})^\ell = (a^\ell) * (a^{-1})^\ell = e$ .

Hence  $a^{k-\ell} = e$  and  $k - \ell > 0$ .

Let  $m$  be the smallest positive integer such that  $a^m = e$

By the division algorithm  $(a) = \{e, a, \dots, a^{m-1}\}$  so, it being clear that  $(a) \subseteq_g G$ ,

$$m = |(a)| \mid |G|$$

Of course if  $a^k = e$  then, by the Division Algorithm,  $k = qm + r$  such that  $0 \leq r \leq m-1$  and

$$e = a^k = (a^m)^q a^r = a^r$$

But  $r < m$  forces  $r = 0$  and so  $m \mid k$ .

(iv) Realize to begin with that, by induction,  $(a b)^t = a^t b^t$

$\forall t \geq 1$ . Thus

$$\begin{aligned} e &= (a b)^{\text{ord}(a b) \text{ ord}(b)} \\ &= a^{\text{ord}(a b) \text{ ord}(b)} b^{\text{ord}(b) \text{ ord}(a b)} \\ &= a^{\text{ord}(a b) \cdot \text{ord}(b)} \end{aligned}$$

and so, by (iii)

$$\text{ord}(a) \mid \text{ord}(a b) \cdot \text{ord}(b)$$

But  $\gcd(\text{ord}(a), \text{ord}(b)) = 1$  forces

$$\text{ord}(a) \mid \text{ord}(a b)$$

Likewise

$$\text{ord}(b) \mid \text{ord}(ab)$$

and because  $\gcd(\text{ord}(a), \text{ord}(b)) = 1$ , it follows that  $(\text{ord}(a)) (\text{ord}(b)) \mid \text{ord}(ab)$ .

Of course

$$\begin{aligned} &(a b)^{\text{ord}(a) \text{ ord}(b)} \\ &= a^{\text{ord}(a) \cdot \text{ord}(b)} \cdot b^{\text{ord}(b) \cdot \text{ord}(a)} \\ &= e \cdot e = e \end{aligned}$$

and, again by (iii),  $\text{ord}(a b) \mid \text{ord}(a) \cdot \text{ord}(b)$ .

The result for arbitrary  $k$  follows by induction

Exercise 9. Prove (iv) for arbitrary  $k \geq 2$ .

Exercise 10. (Submit c)

a) Let  $m$  be a negative integer. Define  $a^m \triangleq (a^{-1})^{|m|}$

for  $a \in G$  (a group). Prove:  $\forall m, n \in \mathbb{Z} \quad \forall a \in G \quad a^{m+n} = a^m * a^n$

and  $(a^m)^n = a^{mn}$ .

Recall:  $a^k \triangleq a \cdot a^{k-1} \quad \forall k \geq 1$  and  $a^0 \triangleq e$

b) Prove:  $\text{ord}(a^{-1}) = \text{ord}(a)$  for each  $a \in G$  (a finite group)

c) Determine the orders of the elements of  $Z_{30}^*$ .

Corollary 3.1 a) (Euler's theorem) If  $a \in Z_n^*$ ,  $n \geq 2$ , then

$$a^{\varphi(n)} \equiv 1 \pmod{n}$$

b) If  $a \in Z_n^*$ ,  $\text{ord}(a) = m$  and  $a^k \equiv 1 \pmod{n}$

then  $m \mid k$ . In particular  $m \mid \varphi(n)$ .

Proof We need only note that  $|Z_n^*| = \varphi(n)$  and 1 is the identity of  $Z_n^*$ .

Remark 8. A special case of Corollary 3.1a) yields

Fermat's theorem: if  $p$  is prime and  $\gcd(a, p) = 1$  then  $a^{p-1} \equiv 1 \pmod{p}$ . Indeed, by the Division Algorithm,  $\exists r \in \mathbb{Z}_p$  such that  $a \equiv r \pmod{p}$  and  $r \neq 0$ . But, as previously noted in Corollary 2.2,  $\mathbb{Z}_p - \{0\}$ , is a group, i.e.

$$\mathbb{Z}_p - \{0\} = \mathbb{Z}_p^* \text{ and so}$$

$$r^{p-1} \equiv 1 \pmod{p}$$

$$\text{Thus} \quad a^{p-1} \equiv 1 \pmod{p}$$

$$\text{since} \quad a^{p-1} \equiv r^{p-1} \pmod{p}.$$

Also note that  $a^p \equiv a \pmod{p} \quad \forall a$ .

Another important result on congruences which generalizes Euler's theorem is given next.

Proposition 4. If  $n$  is a product of distinct primes and  $r, s > 0$  then

$$\forall a \in \mathbb{Z} \quad (r \equiv s \pmod{\varphi(n)}) \Rightarrow a^r \equiv a^s \pmod{n}$$

In particular if  $n = p$  a prime, then  $\forall a (r \equiv s \pmod{p-1})$

$$\Rightarrow a^r \equiv a^s \pmod{p}.$$

Proof First observe that we may assume  $a \in \mathbb{Z}_n$  and that  $a \neq 0$ . Indeed if  $a \notin \mathbb{Z}_n$  then  $\exists b \in \mathbb{Z}_n$  such that  $a \equiv b \pmod{n}$ . But then  $a^k \equiv b^k \pmod{n} \quad \forall k$ . If  $a = 0$  the result is trivially true.

Now suppose  $a \in \mathbb{Z}_n^*$  so that  $a^{\varphi(n)} \equiv 1 \pmod{n}$ .

$$\text{But} \quad r = q \varphi(n) + s$$

$$\text{so} \quad a^r = (a^{\varphi(n)})^q \cdot a^s$$

$$\text{Since} \quad a^{\varphi(n)} \equiv 1 \pmod{n}, \quad a^r \equiv a^s \pmod{n}$$

Suppose  $a \notin \mathbb{Z}_n^*$  so that  $\gcd(a, n) = d > 1$

Now  $n = n'd$ ,  $a = a'd$  where  $\gcd(n', a') = 1$ . Also, since  $n$  is a product of distinct primes  $\gcd(n', d) = 1$  so  $\varphi(n) = \varphi(n') \varphi(d)$ . As  $r \equiv s \pmod{\varphi(n)}$  we may write  $r - s = \alpha \varphi(n') \varphi(d)$  for some  $\alpha \in \mathbb{Z}$ . Consider

$$\begin{aligned} a^{r-s} &= (a')^{r-s} (d)^{r-s} \\ &= (a')^{\varphi(n') \varphi(d) \alpha} (d)^{\varphi(n') \varphi(d) \alpha} \end{aligned}$$

But  $\gcd(a', n') = 1 \Rightarrow (a')^{\varphi(n')} \equiv 1 \pmod{n'}$  which, in turn, forces

$$(a')^{\varphi(n')\varphi(d)\alpha} \equiv 1 \pmod{n'}$$

Likewise  $\gcd(d, n') = 1 \Rightarrow (d)^{\varphi(n')} \equiv 1 \pmod{n'}$

and so

$$(d)^{\varphi(n')\varphi(d)\alpha} \equiv 1 \pmod{n'}$$

Thus

$$a^{r-s} \equiv 1 \pmod{n'}$$

so

$$a^r \equiv a^s \pmod{n'}$$

follows.

As  $r, s > 0$  we see that  $d \mid a \Rightarrow d \mid a^r - a^s$

Now

$$n' \mid a^r - a^s \quad \text{and} \quad d \mid a^r - a^s$$

together with  $\gcd(n', d) = 1$  forces  $n = n' d \mid a^r - a^s$  i.e.  $a^r \equiv a^s \pmod{n}$ .

Remark 9. If  $r > 0, s = 0$  (or  $r = 0, s > 0$ ) the result is not true.

Example 4. Let  $n = 6$  so  $\varphi(6) = 2$ . Set  $r = 2$  so  $r \equiv 0 \pmod{2}$ . Note for  $a = 4 \notin \mathbb{Z}_6^*$

$$a^r = 4^2 = 16 \equiv 4 \pmod{6} \neq 1 \pmod{6}.$$

Of course, it is the case that

$$r \equiv 0 \pmod{\varphi(n)} \Rightarrow a^r \equiv 1 \pmod{n} \text{ for } a \in \underline{\underline{\underline{\mathbb{Z}_n^*}}}$$

Remark 10. If one wants to raise an integer  $a$  to a power modulo  $n$  and the power is not congruent to 0 modulo  $\varphi(n)$  then the computation may be done with a power which is congruent to the original power modulo  $\varphi(n)$  (Here  $n$  is a product of distinct primes).

Example 5. Consider  $n = (2)(3)(5) = 30$ ,  $a = 63$  and  $r = 10$ . We desire the value of  $63^{10} \pmod{30}$ .

First  $63 \equiv 3 \pmod{30}$

so  $63^{10} \equiv 3^{10} \pmod{30}$

Now  $\phi(30) = 8$

and  $10 \equiv 2 \pmod{8}$

so  $3^{10} \equiv 3^2 \pmod{30}$

and, finally,  $63^{10} \equiv 9 \pmod{30}$

On the other hand consider  $r = 8 (= 0 \pmod{\phi(30)})$

Then  $63^8 \equiv 6561 \equiv 21 \pmod{30} \not\equiv 1 \pmod{30}$

Finally consider  $a = 7 \in \mathbb{Z}_{30}^*$  and  $r = 8$ ; then

$$7^8 = 1443001 \equiv 1 \pmod{30}.$$

Exercise 11. Let  $n = 4 = 2^2$  and  $a = 2$  and  $r = 3, s = 1$

Calculate:  $a^s$  and  $a^r \pmod{4}$ . Does  $a^r = a^s \pmod{n}$ ? What does this

say about the advisability of dropping the condition that  $n$  be a product of distinct primes in the previous result?

An algorithm for modular exponentiation is given next; but first an observation:

Observation If  $k = \sum_{i=0}^t k_i 2^i$  ( $k_i = 0, 1$ ) is the binary representation of  $k$  then

$$a^k = (a^{2^0})^{k_0} (a^{2^1})^{k_1} \dots (a^{2^t})^{k_t}$$

Algorithm (Repeated Square -and-Multiply)

Input:  $a \in \mathbb{Z}_n$  and  $k = \sum_{i=0}^t k_i 2^i$

Output:  $a^k \pmod{n}$

1. Set  $b \leftarrow 1$ ; if  $k = 0$  return  $(b)$
2. Set  $A \leftarrow a$
3. If  $k_0 = 1$  then set  $b \leftarrow a$
4. For  $i = 1, \dots, t$  do the following:
  - 4.1 Set  $A \leftarrow A^2 \pmod{n}$
  - 4.2 If  $k_i = 1$  set  $b \leftarrow A \cdot b \pmod{n}$
5. Return  $(b)$

Example 6. Evaluate  $4^{35} \bmod 30$

Input:  $a = 4$ ,  $k = 1 \cdot 2^0 + 1 \cdot 2^1 + 0 \cdot 2^2 + 0 \cdot 2^3 + 0 \cdot 2^4 + 1 \cdot 2^5$

Steps 1.  $b = 1$

2.  $A = 4$

3. Since  $k_0 = 1$ ,  $b = 4$

4's. i)  $i = 1$  4.1  $A = 16 \bmod 30 = 16$

4.2  $b = (16)4 \bmod 30 = 4$

ii)  $i = 2$  4.1  $A = 256 \bmod 30 = 16$

iii)  $i = 3$  4.1  $A = 256 \bmod 30 = 16$

iv)  $i = 4$  4.1  $A = 256 \bmod 30 = 16$

v)  $i = 5$  4.1  $A = 256 \bmod 30 = 16$

4.2  $b = (16)(4) \bmod 30 = 4$

5.  $b = 4$

Check:  $n = 30$   $\phi(n) = 8$ ,  $35 = 3 \bmod 8$  so  $4^{35} = 4^3 \bmod 30 = 4 \bmod 30$ .

Exercise 12. (Submit a) a) evaluate  $5^{75} \bmod 35$  using the algorithm. Can you check this using the previous theorem? If so do so.

b) If  $k < n$  then the algorithmic complexity is  $O((\lg n)^3)$

Prove this.