# Module 2 (Assembly Language Programming and Assemblers)

SIC/XE Programming, Basic functions of Assembler, Assembler output format – Header, Text and End Records. Assembler data structures, Two pass assembler algorithm.

## **Assemblers**

At one time, the computer programmer had at his disposal a basic machine that interpreted, through hardware, certain fundamental instructions. He would program this computer by writing a series of 1's and 0's (machine language), place them into the memory of the machine, and press a button, whereupon the computer would start to interpret them as instructions.

Programmers found it difficult to write or read programs in machine language. In their quest for a more convenient language they began to use a mnemonic (symbol) for each machine instructions, which they would subsequently translate into machine language. Such a mnemonic machine language is now called an assembly language. Programs known as Assemblers were written to automate the translation of assembly language into machine language. The input to an assembler program is called source program. The output is a machine language translation called object program.



## **Basic Assembler functions**

- Convert **symbolic operands** to their equivalent **machine addresses** (eg: RETADR to 1033)
- Convert **mnemonic operation codes** to their **machine language equivalents** (eg: STL to 14)
- Convert the **data constants** specified in the source program into their **internal machine representations** (eg: EOF to 454F46)

- Write the object program and the assembly listing
- Build the machine instruction into proper format

#### **Assembler Directive**

An assembler directive is a set of instructions used to instruct the assembler to perform certain actions during the assembly of the program. This statement neither represents machine instructions that are to be included in the object program nor indicate the storage allocation of constants or variables. It directs the assembler to take certain actions during the process of assembling a program. START, END, BYTE, WORD, RESW, RESB, BASE, EQU, ORG, LTORG are some of the assembler directives.

1. START : Specify name and starting address for the program

eg: COPY START 1000

2. END: Indicate the end of the source program and (optionally) specify the first executable instruction in the program

eg: END FIRST

There are four different ways of defining storage for data items in the SIC Assembler language:

- 1. BYTE: Generate **character or hexadecimal constant**, occupying as many bytes as needed to represent the constant
  - Eg: CHARZ BYTE C'Z'
- 2. WORD: Generate one-word integer constant
  - Eg: FIVE WORD 5
- 3. RESB: Reserve the indicated number of bytes for a data area
  - eg: C1 RESB
- 4. RESW: Reserve the indicated number of words for a data area
  - eg: ALPHA RESW 1

## **Example of a SIC Assembler language program**

Line	Sour	ce statem	ent	
STATE OF	COPY	START	1000	COPY FILE PROM INPUT TO OUTPUT
10	FIRST	STL	REPADE	SAVE RETURN ADDRESS
15	CLOOP	JSUB	RDREET	READ INPUT RECORD
20		LDA	LHENCYTH	TEST FOR BOY (LENGTH = 0)
25		COMP	ZERO	THE RESERVE AND ASSESSMENT OF THE PARTY OF T
30		JIDQ	EMDEIL	EXIT IF EOF FOUND
35		JEOB	WRREC	WRITE OUTPUT RECORD
40		J	CLOOP	LOOP
45	ENDFIL	LDA	BOF	INSERT END OF FILE MARKER
50		STA	BUFFIRE	the wall telephone and built and I distant
55		LDA	THREE	SET LENGTH = 3
60		STA	LEONOTH	Ches where BUSH
65		JSUB	WRREC	WRITE EOF GET RETURN ADDRESS
70		LUL	RETADR	RETURN TO CALLER
75		RSUB	CONTRACTOR PROPERTY	RETURN TO CALLER
80	EOF	BYTE	C'EOF'	
85	THREE	WORD	3	
90	ZERO	WORD	0	2.1.1 A Simple SIC Assembler
95	RETADR	RESW	1	LENGTH OF RECORD
100	LENGTH	RESW	100	4096-BYTE BUFFER AREA
105	BUFFER	RESB	4096	4096-BILE BULLING STATES
110	a set of the second	1012 400000	Control March Street Co.	PROPER YANTS BITEPER
115	THE REAL PROPERTY.	SUBROU	FINE TO READ R	ECORD INTO BUFFER
120	SERVICE TO STATE OF THE PERSON NAMED IN COLUMN TWO IS NOT THE PERSON NAMED IN COLUMN TWO IS			CLEAR LOOP COUNTER
125	RDREC	LDK	ZERO	CLEAR A TO ZERO
130	September Sty Barri	LDA	ZERO	TEST INPUT DEVICE
135	RLOOP	TD	INPUT	LOOP UNTIL READY
140		JEQ	RLOOP	READ CHARACTER INTO REGISTER
145		RD	INPUT	TEST FOR END OF RECORD (X'00'
150		COMP	ZERO	FXIT LOOP IF EOR
155		JEQ	EXIT	STORE CHARACTER IN BUFFER
160		STCH	BUFFER, X	LOOP UNLESS MAX LENGTH
165		TIX	MAXLEN	
170		JLT	RLOOP	HAS BEEN REACHED
175	EXIT	STX	LENGTH	SAVE RECORD LENGTH
180		RSUB		RETURN TO CALLER
185	INPUT	BYTE	X'Fl'	CODE FOR INPUT DEVICE
190	MAXLEN	WORD	4096	
195				Committee to the committee of the commit
200		SUBROU	TINE TO WRITE	RECORD FROM BUFFER
205	MINDS NAMED			
210	WRREC	LDX	ZERO	CLEAR LOOP COUNTER
215	WLOOP	TD	OUTPUT	TEST OUTPUT DEVICE
220	WALCONE	JEO	WLOOP	LOOP UNTIL READY
225		LDCH	BUFFER, X	GET CHARACTER FROM BUFFER
2005		WD	OUTPUT	WRITE CHARACTER
230		100000	LENGTH	LOOP UNTIL ALL CHARACTERS
235		TIX	WLOOP	HAVE BEEN WRITTEN
240		JLT	WELCOF	RETURN TO CALLER
245		RSUB	****	CODE FOR OUTPUT DEVICE
250	OUTPUT	BYTE	X'05'	CODE FOR COLLEGE
255		END	FIRST	

Figure shows an assembler language program for the basic version of SIC. Line numbers are given only for reference and are not the part of the program. Then there are labels defined by the programmer. Then mnemonic instructions (opcode) eg: STL, JSUB. Indexing addressing is indicated by adding modifier "X" (line 160). Then comments are represented by "."

The program contains a main routine that reads records from an input device, identified by device code F 1 and copies them to an output device 05. Main routine calls subroutine RDREC to read a record into buffer and another subroutine WRREC to write the record from the buffer to the output device. Each subroutine must transfer the record one character at a time because the only I/O instructions available are RD and WD. The buffer is necessary: because the I/O rates for

two devices, such as a disk and a slow printing terminal, may be different. The end of each record is marked with a null character (hexadecimal 00).

If a record is longer than the length of a buffer (4096 bytes), only the first 4096 bytes are copied.(for simplicity, the program does not deal with the error recovery when a record containing 4096 bytes or more is read. The end of the file to be copied is indicated by a zero-length record. When the end of file is detected, the program writes EOF on the output device and terminates by executing an RSUB instruction. Assumed that the this program was called by the operating system using a JSUB instruction and thus the RSUB will return the control to the operating system

#### **Forward Reference**

Convert symbolic operands to their equivalent machine addresses (eg: RETADR to 1033). This cannot be achieved in the sequential processing of the source program, one line at a time. This poses a problem: Forward Reference

**Forward Reference** – a reference to label (RETADR) that is defined later in the program. If we attempt to translate the program line by line, we will unable to process this statement because we do not know the address that will be assigned to RETADR. Because of this, most assemblers make two passes over the source program.

#### • PASS 1:

 Scan the source program for label definitions and assign addresses (such as the Loc column)

## • PASS 2:

Performs the actual translation

5	COPY	START	1000	COPY FILE FROM INPUT TO OUTPUT
10	FIRST	STL	RETADR	SAVE RETURN ADDRESS
15	CLOOP	JSUB	RDREC	READ INPUT RECORD
20		LDA	LENGTH	TEST FOR EOF (LENGTH = 0)
25		COMP	ZERO	
30		JEQ /	ENDFIL	EXIT IF EOF FOUND
35		JSUB /	WRREC	WRITE OUTPUT RECORD
40		J /	CLOOP	LOOP
45	ENDFIL	LDA/	EOF	INSERT END OF FILE MARKER
50		STA	BUFFER	
55		LDA	THREE	SET LENGTH = 3
60	HALL MAN TO SERVICE STATE OF THE PARTY OF TH	STA	LENGTH	
65	Forward	JSUB	WRREC	WRITE EOF
70	reference	LDL	RETADR	GET RETURN ADDRESS
75	1	RSUB		RETURN TO CALLER
80	EOF /	BYTE	C'EOF'	
85	THREE /	WORD	3	
90	ZERO 🗲	WORD	0	
95	RETADR	RESW	1	
100	LENGTH	RESW	1	LENGTH OF RECORD
105	BUFFER	RESB	4096	4096-BYTE BUFFER AREA

## **Assembler Output Format**

Finally, the assembler must write the generated object code onto some output device. The object program will later be loaded into memory for execution. The simple object program format uses 3 types of records: Header, Text and End.

- **Header record** contains the program name, starting address and length.
- **Text records** contain the translated (ie., machine code) instructions and data of the program, together with an indication of the addresses where these are to be loaded.
- **End record** marks the end of the object program and specifies the address in the program where execution is to begin.

## Header record

Col. 1 H
Col. 2~7 Program name

Col. 8~13 Starting address of object program (hex)

Col. 14-19 Length of object program in bytes (hex)

Text record

Col. 1 T

Col. 2~7 Starting address for object code in this record (hex)
Col. 8~9 Length of object code in this record in bytes (hex)
Col. 10~69 Object code, represented in hex (2 col. per byte)

End record

Col.1 E

Col.2~7 Address of first executable instruction in object program (hex)

Line	Loc	Source statement			Object code
5	1000	COPY	START	1000	
10	1000	FIRST	STL	RETADR	141033
15	1003	CLOOP	JSUB	RDREC	482039
20	1006		LDA	/ LENGTH	001036
25	1009		COMP	ZERO	281030
30	100C		JEQ /	ENDFIL	301015
35	100F		JSUB /	WRREC	482061
40	1012		J /	CLOOP	3C1003
45	1015	ENDFIL	LDA/	EOF	00102A
50	1018		STA	BUFFER	0C1039
55	101B		LDA	THREE	00102D
60	101E		STA	LENGTH	0C1036
65	1021		JSUB	WRREC	482061
70	1024		LDL	RETADR	081033
75	1027		/ RSUB		4C0000
80	102A	EOF	BYTE	C'EOF'	454F46
85	102D	THREE /	WORD	3	000003
90	1030	ZERO ¥	WORD	3	000000
95	1033	RETADR	RESW	1	
100	1036	LENGTH	RESW	1	
105	1039	BUFFER	RESB	4096	
110					

115			SUBROUT	TINE TO READ REC	CORD INTO BUFFER
120					
125	2039	RDREC	LDX	ZERO	041030
130	203C		LDA	ZERO	001030
135	203F	RLOOP	TD	INPUT	E0205D
140	2042		JEQ	RLOOP	30203F
145	2045		RD	INPUT	D8205D
150	2048		COMP	ZERO	281030
155	204B		JEQ	EXIT	302057
160	204E		STCH	BUFFER, X	549039
165	2051		TIX	MAXLEN	2C205E
170	2054		JLT	RLOOP	38203F
175	2057	EXIT	STX	LENGTH	101036
180	205A		RSUB		4C0000
185	205D	INPUT	BYTE	X'F1'	F1
190	205E	MAXLEN	WORD	4096	001000
105					

195	STATE STATE OF STREET	STATE OF THE STATE	1 1 2 1 1 5 W		AND THE RESIDENCE OF THE PARTY OF
200			SUBROU	TIME TO WRITE R	ECORD FROM BUFFER
205		a Partition in			
210	2061	WRREC	LDX	ZERO	041030
215	2064	WLOOP	TD	OUTPUT	E02079
220	2067		JEQ	WLOOP	302064
225	206A		LDCH	BUFFER, X	509039
230	206D		WD	OUTPUT	DC2079
235	2070		TIX	LENGTH	2C1036
240	2073		JLT	WLOOP	382064
245	2076		RSUB		4C0000
250	2079	OUTPUT	BYTE	X'05'	05
255	. (timisat	ELEGITIVE TO	END	FIRST	

Figure below shows the sample object program generated for the above given simple SIC assembly program.

To avoid confusions, we have used he term column rather than byte to refer to positions within object program records. This is not meant to imply the use of any particular medium for the object program. "^" used to separate fields visually and is not present in the actual object program. Note there is no object code corresponding to the addresses  $1033-2038 \rightarrow$  this storage is simply reserved by the loader for use by the program during execution.

## **Passes of Assembler**

A Pass is defined as the processing activity of every single statement in the source code to perform a set of language processing functions. Pass can also be defined as the activity of scanning the assembly language programming.

• **Single pass Assembler:** The assembler scans the entire source program (assembly language program) once and convert into an object code.

• **Multi-pass Assembler:** The translation of assembly language program into object code requiring many passes.

The breaking of the entire assembly process into passes makes design simpler and enables better control over the subtasks and intermediate operations.

## **Functions of Two Passes of Assembler**

- PASS 1 (Define symbols)
  - Assign addresses to all statements in the program
  - Save the values (addresses) assigned to all labels for use in Pass 2
  - Perform some processing of assembler directives. (This includes processing that
    affects address assignment, such as determining the length of data areas defined
    by BYTE, RESW, etc.)
- PASS 2 (Assemble instructions and generate object program)
  - Assemble instructions (translating operation codes and looking up addresses)
  - Generate data values defined by BYTE, WORD, etc.
  - Perform processing of assembler directives not done during Pass 1
  - Write the object program and assembly listing

#### **Assembler Data Structures**

Simple Assembler uses two major internal data structures:

- Operation Code Table (OPTAB)
- Symbol Table (SYMTAB)

Also need a variable Location Counter (LOCCTR).

#### **OPTAB**

OPTAB is used to look up mnemonic operation codes and translate them to machine language equivalents. This must contain at least mnemonic operation code and its machine language equivalent. In more complex assemblers, this table also contains information about instruction format and length. During Pass 1 OPTAB is used to look up and validate operation codes in the source program. In Pass 2, it is used to translate the operation codes to machine language.

In case of SIC/XE machine that has instruction of different length. We must search OPTAB in the first pass to find the instruction length for incrementing LOCCTR. In second pass, the information from OPTAB tell us which instruction format to use in assembling the instruction, and any peculiarities of the object code instructions (typically most real assemblers).

OPTAB is usually organised as a hash table, with mnemonic operation code as the key. This information in OPTAB is predefined when the assembler itself is written, rather than being loaded into the table at the execution time. This hash table organisation provides fast retrieval with a minimum of searching. OPTAB is static table – entries are not normally added to or deleted from it.

#### **SYMTAB**

SYMTAB is used to store values(address) assigned to labels. SYMTAB includes the name and value(address) for each label in the source program, together with flags to indicate error conditions(eg: a symbol defined in two different places). The table may contain other information about the data area or instruction labelled (eg: it's type or length). During Pass 1, the labels are entered into SYMTAB as they are encountered in the source program, along with their assigned addresses (from LOCCTR). During Pass 2, symbols used as operands are looked up in SYMTAB to obtain the addresses to be inserted in the assembled instructions

SYMTAB is usually organised as hash table for efficiency in insertion and retrieval. Entries are rarely deleted from this table. Programmers often select many labels that have similar characteristics (eg: label start or end with the same characters, like LOOP1, LOOP2, LOOPA,...or are of same length like A, X, Y, Z). Hashing function selected should perform well with such non random keys. Care should be taken in the selection of hashing function because the SYMTAB is used throughout the assembly. Good option is the selection of hash function which divides the entire key by a prime table length.

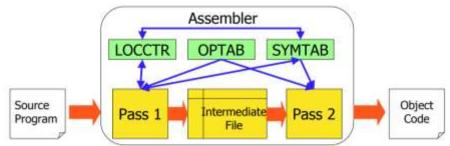
10

## **LOCCTR**

LOCCTR is a variable that is used to help in the assignment of addresses. LOCCTR is initialized to the beginning address specified in the START statement. After each source statement is processed, the length of the assembled instruction or data area to be generated is added to LOCCTR. Thus whenever we reach a label in the source program, the current value of LOCCTR gives the address to be associated with that label.

## **Assembler Algorithm**

Both passes of the assembler reads the original source program as input. However, there is certain information (such as location counter values and error flags for the statements) that can or should be communicated between the two passes. Pass 1 usually writes an *intermediate file* that contains each source statements together with its assigned address, error indicators etc. This file is used as input to Pass 2. Means this working copy of the source program(intermediate file) can also be used to retain the results of certain operations that may be performed during Pass 1 (such as scanning the operand field for symbols and addressing flags), so these need not be performed again during Pass 2. Similarly, pointers into OPTAB and SYMTAB may be retained for each operation code and symbol used.



The intermediate file include each source statement, assigned address and error indicator

Algorithm explains the logic flow of two passes of assembler. Apply the algorithm to source program (assembly language) to generate object program. For simplicity, we assume that source lines are written in the fixed format with fields:

#### LABEL OPCODE OPERAND

If one of these fields contains a character string that represents a number, we denote its numeric value with a prefix #. (eg: #(OPERAND))

```
Pass 1:
begin
   read first input line
   if OPCODE = 'START' then
      begin
          save #[OPERAND] as starting address
          initialize LOCCTR to starting address
         write line to intermediate file
          read next input line
      end {if START}
  else
      initialize LOCCTR to 0
  while OPCODE # 'END' do
      begin
         if this is not a comment line then
             begin
                if there is a symbol in the LABEL field then
                    begin
                       search SYMTAB for LABEL
                       if found then
                           set error flag (duplicate symbol)
                           insert (LABEL, LOCCTR) into SYMTAB
                    end {if symbol}
                search OPTAB for OPCODE
                if found then
                    add 3 (instruction length) to LOCCTR
                else if OPCODE = 'WORD' then
                    add 3 to LOCCTR
                else if OPCODE = 'RESW' then
                    add 3 * #[OPERAND] to LOCCTR
                else if OPCODE = 'RESB' then
                    add #[OPERAND] to LOCCTR
                else if OPCODE = 'BYTE' then
                   begin
                       find length of constant in bytes
                       add length to LOCCTR
                   end (if BYTE)
                else
                   set error flag (invalid operation code)
             end {if not a comment}
         write line to intermediate file
         read next input line
     end {while not END}
  write last line to intermediate file
  save (LOCCTR - starting address) as program length
end {Pass 1}
```

## Pass 2: begin read first input line (from intermediate file) if OPCODE = 'START' then begin write listing line read next input line end (if START) write Header record to object program initialize first Text record while OPCODE # 'END' do begin if this is not a comment line then begin search OPTAB for OPCODE if found then begin if there is a symbol in OPERAND field then begin search SYMTAB for OPERAND if found then store symbol value as operand address else begin store 0 as operand address set error flag (undefined symbol) end end {if symbol} else store 0 as operand address assemble the object code instruction end {if opcode found} else if OPCODE = 'BYTE' or 'WORD' then convert constant to object code if object code will not fit into the current Text record then begin write Text record to object program initialize new Text record end add object code to Text record end {if not comment} write listing line read next input line end {while not END} write last Text record to object program write End record to object program write last listing line end {Pass 2}