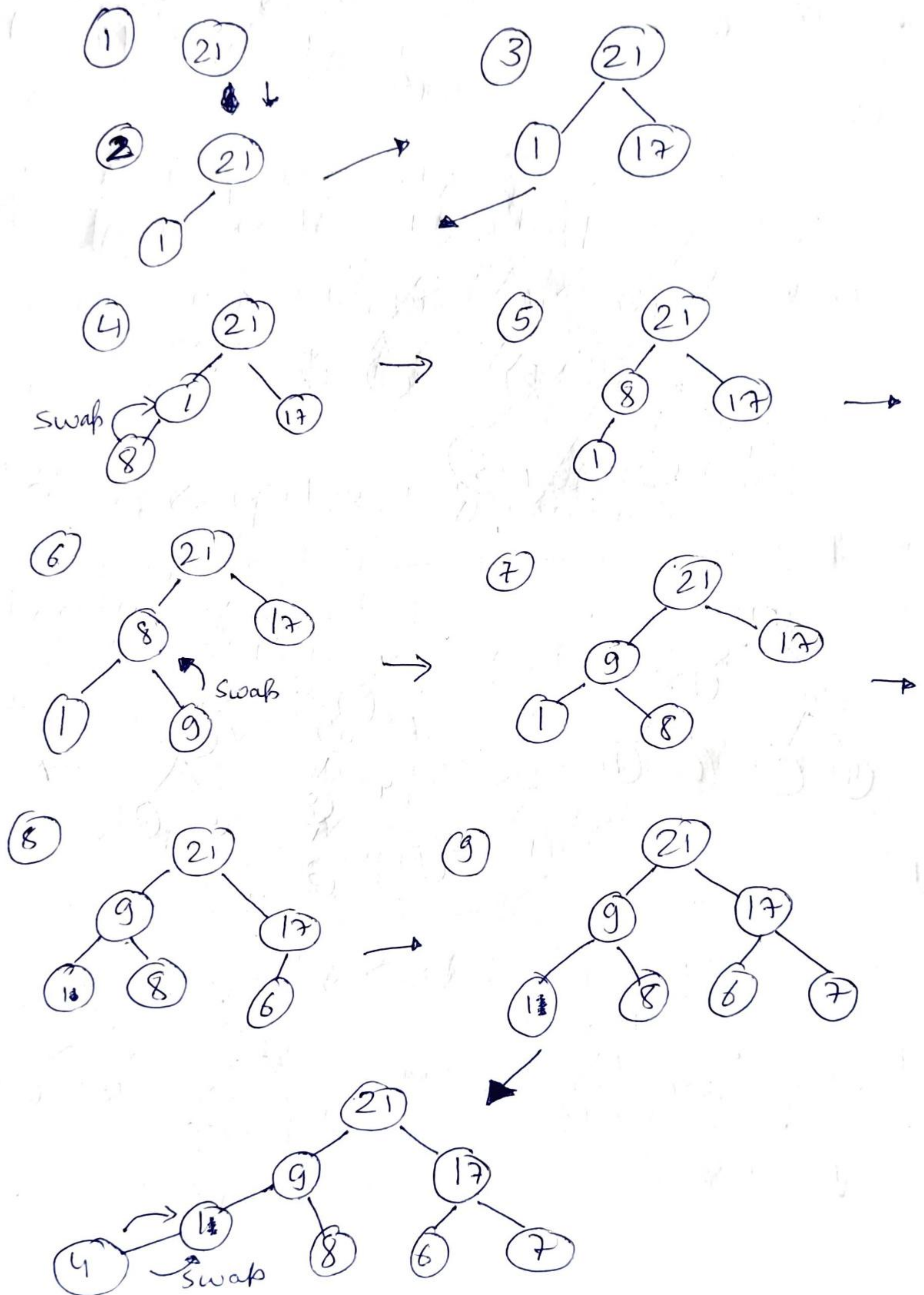


Ans 1) 21, 1, 17, 8, 9, 6, 7, 4, 3, 8, 5



① Max-Heap (A)
{

A-heap-size = A-length
for ($i = \text{floor}(A.\text{length}/2)$ to 1)
 max-heapify(A, i);
}

Time Complexity = $O(n \log n)$

Procedure

① Create a new node in the heap.

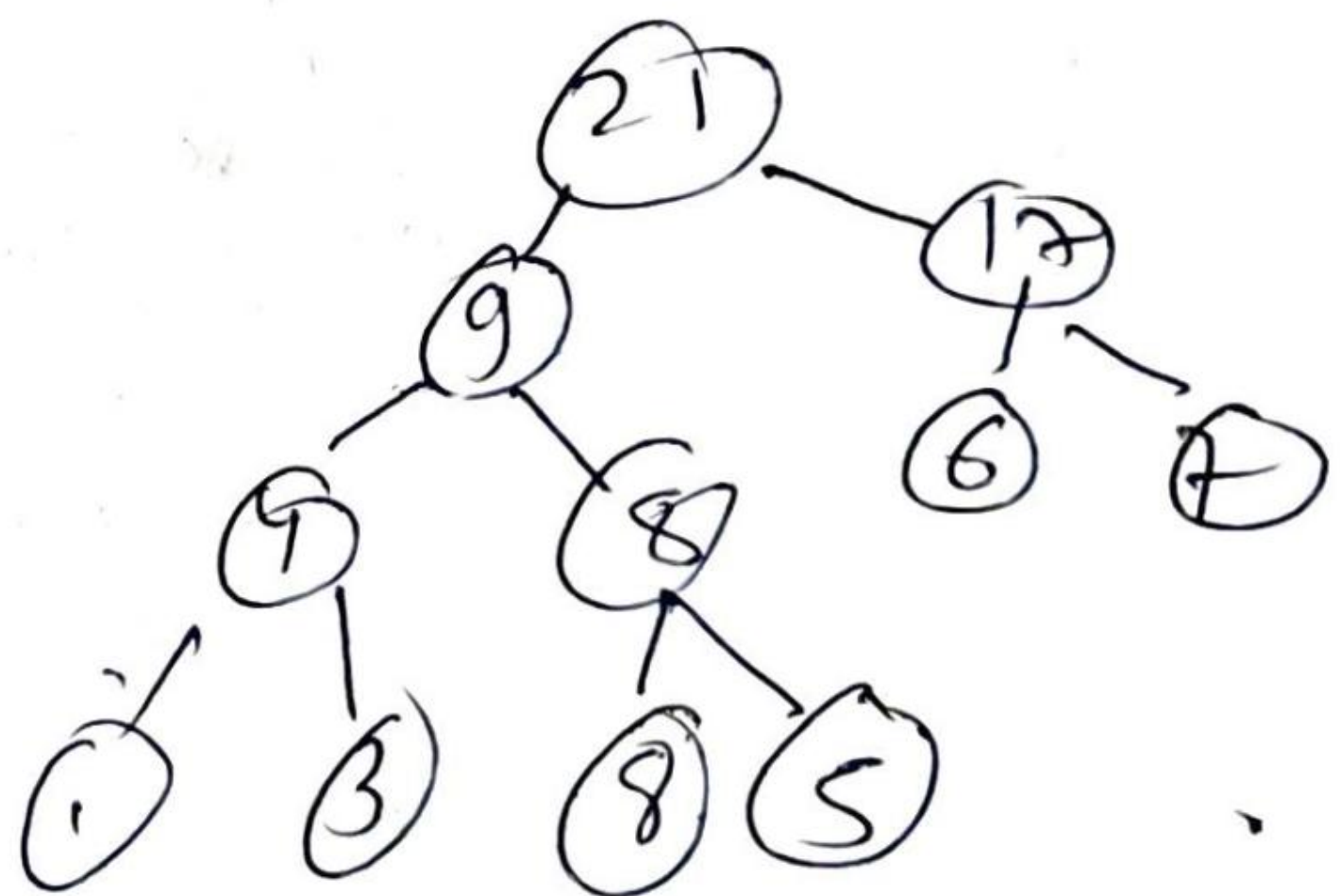
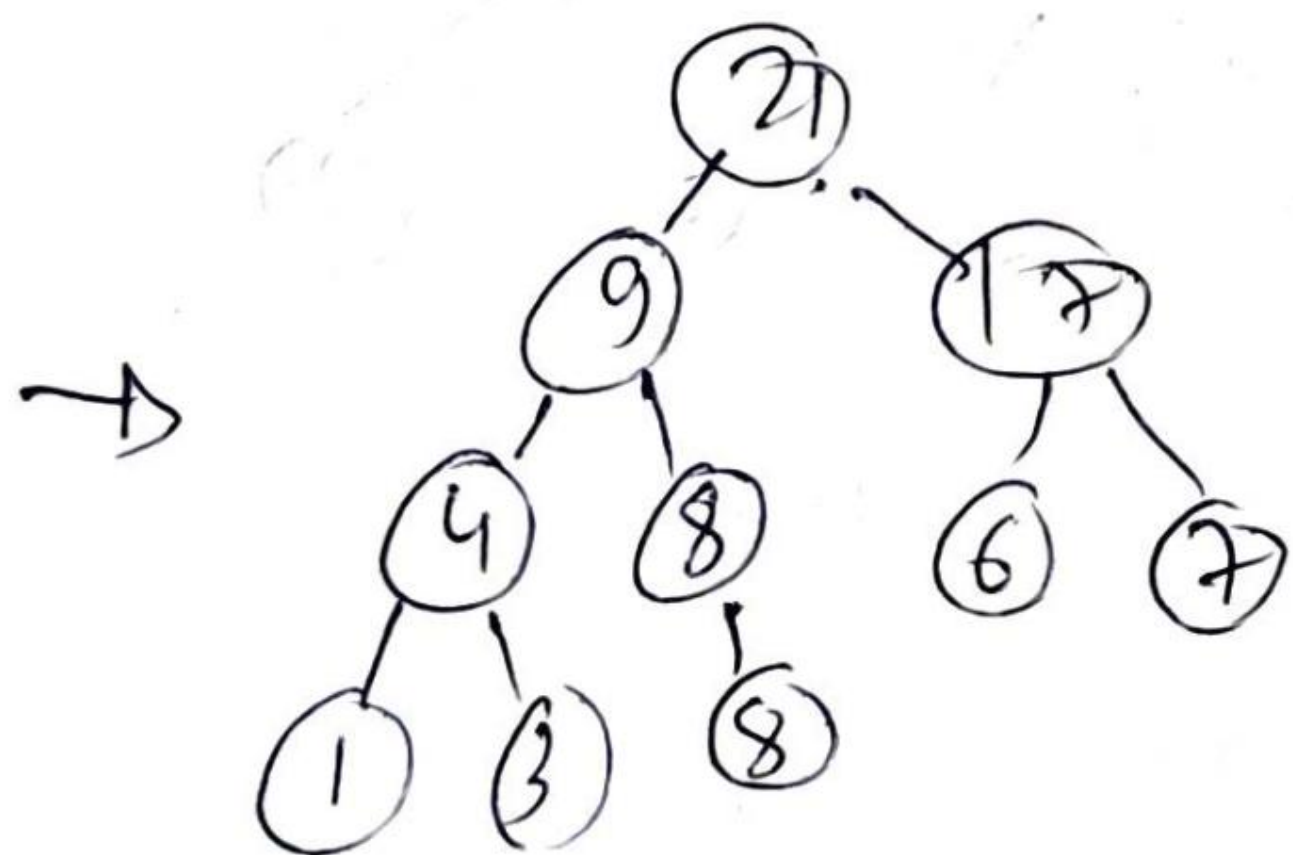
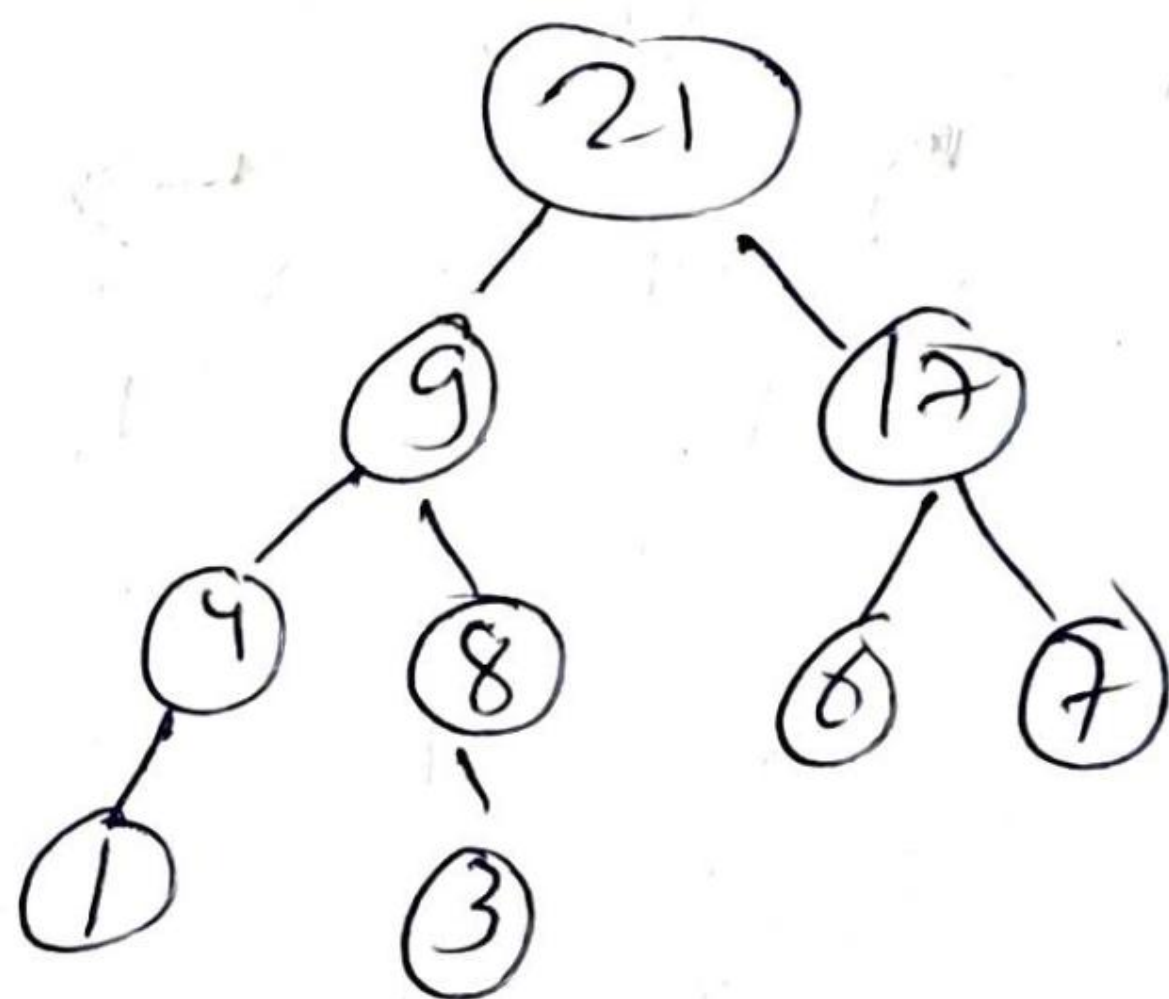
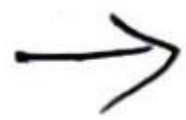
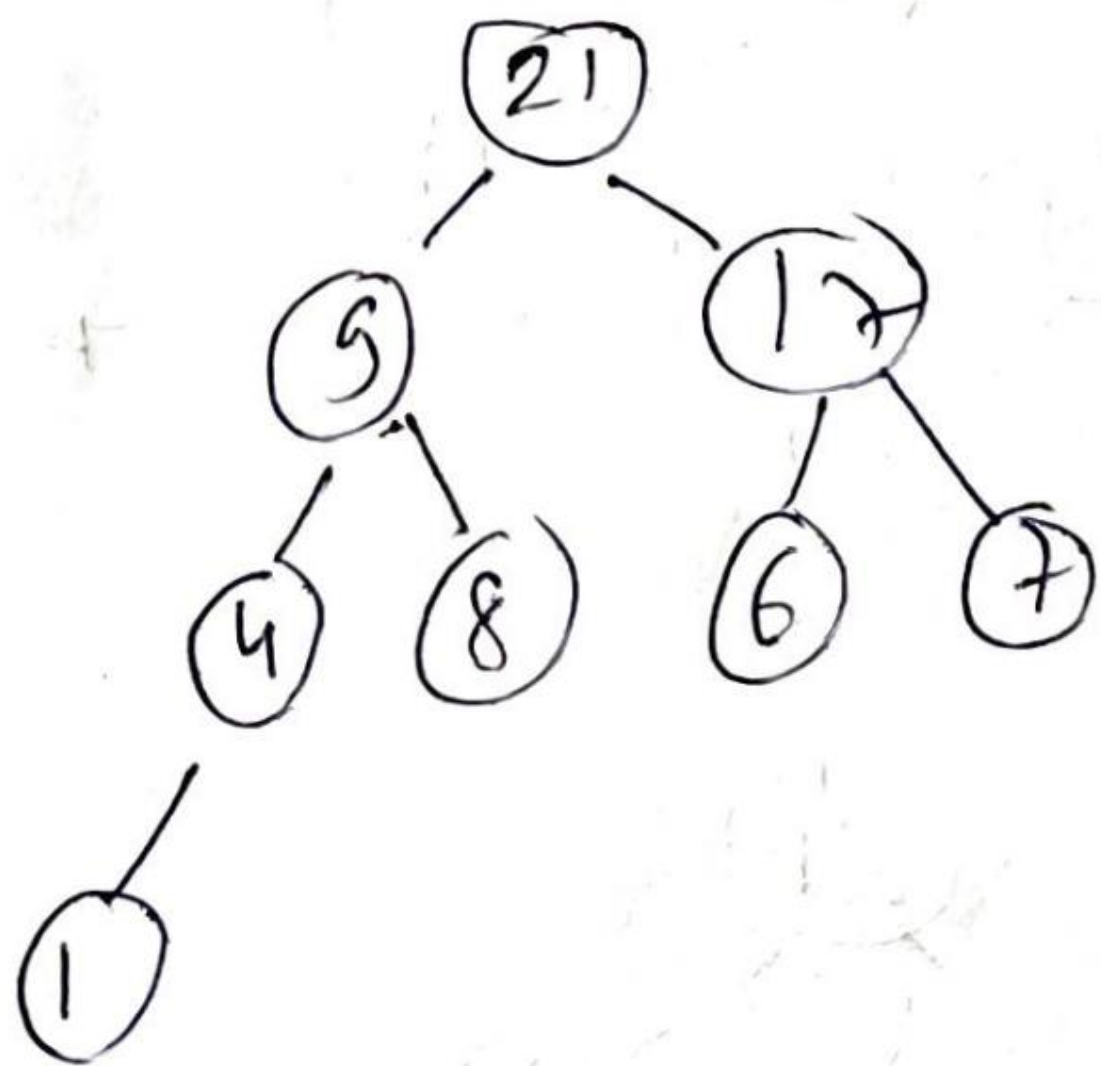
② Assign value to node

③ Compare values of child node with parent node

If parent $<$ child then swap

④ If parent $>$ child, then normally we increase i with first left and then right.

⑤ Repeat steps ③ ④ until heap property holds.



②① PARTITION (A, p, r)

$x = A[r]$

$i = p - 1$

for $j = p$ to $r - 1$

if $A[j] \leq x$

$i = i + 1$

exchange $A[i]$ with $A[j]$

exchange $A[i+1]$ with $A[r]$

return $i+1$

when all are equal then loop terminates
 $k = r - 1$, so value of $q = k + 1 = r$.

• PARTITION'(A, p, r) // modified

$x = A[r]$

$k = p - 1$

$l = p - 1$

for $j = p$ to $r - 1$

if $A[j] < x$

$k = k + 1$

exchange $A[j]$ with $A[k]$

$l = l + 1$

~~else~~

else if $A[j] = x$

$l = l + 1$

exchange $A[j]$ with $A[l]$

exchange $A[l + 1]$ with $A[x]$

return $(A, [(l + k) / 2] + 1)$

- (b) The basic idea of the greedy approach is to calculate the ratio value weight of each item and sort the item on basis of this ratio. Then take the item with highest ratio & add them until, the condition reaches of weight or other condition.