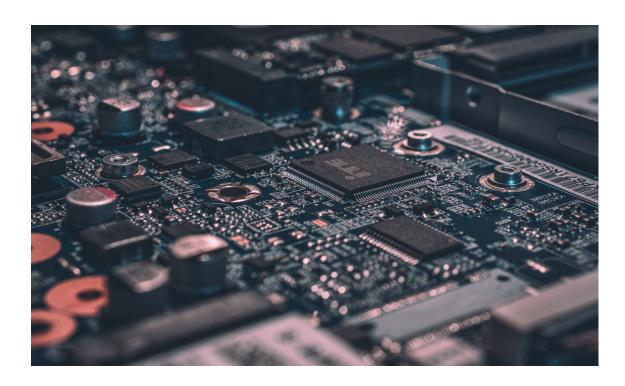
DIC - Project 2

Zephyr RTOS Crypto Processor

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1 General

1.1 Abstract

Simply put, the task is to implement a cryptoprocessor in Zephyr RTOS. The processor should successfully complete the tests given in the specification. The virtual serial interface is used to address the cryptoprocessor. This is decrypted using AES-128 (Advanced-Encryption-Standard-128-bit). All this is implemented on the operating-system-version 2.4.0. In this project, a 64-bit Linux microcontroller board ("native_posix_64") is used. The native_posix_64 board implements a virtual serial interface called "UART_0". At the same time, a Crypo API called "CRYPTO_TC" is implemented using libtinycrypt. Both implemented drivers are to be used.

2 Prior Knowledge and Component Information

2.1 Specifications

MAIN Thread

The main thread of the programm should start all other threads and then only sporadically give a life message.

UART IN Thread

The native_posix_64 UART driver does not support IRQ, so reading is blocking. Therefore a thread is needed that reads and as soon as it has a character it sends it to a "Message_Queue". The processing threads consume this character from there. When the character "." has been read, it should be echoed back to the serial port. If the processing thread is busy, the message "BUSY" is output via the serial interface.

UART OUT Thread

This thread ensures that multiple threads can output messages to the serial port without interfering with each other. To do this, it reads the messages to be sent from a queue and sends it character by character via the serial interface before it sends another before fetching another message from its queue.

PROCESSING Thread

This thread performs the crypto operations. These can take a very long time. The processing thread can store exactly one message in its message queue, This is because one is in processing, one waiting before the system has to generate "BUSY" messages.

2.2 Serial Protocol

Serial connection possible (= alive)

Whenever a '.' is received, a '.' is immediately returned.

Cryptoprocessor available (= avail)

When a 'P' is received, the processing thread responds with "PROCESSING AVAIL".

Load key (= key)

A 'K' followed by 16 bytes of AES-128 key and any character (discarded) loads a new key into the cryptoprocessor.

Load Initial Vecotr (= iv)

An 'I' followed by 16 bytes of AES-128 key and an arbitrary character (discarded) loads a new key into the cryptoprocessor

Decrypt in CBC Mode

A 'D' followed by a byte length (ciphertext, must be multiple of cipherblocksize) followed by correspondingly long ciphertext and any character (discarded). The cryptoprocessor decrypts with the key/iv in aes128-cbc and returns the plaintext on the serial interface. "XERROR" if errors have occurred (e.g.: length).

Further Information

All further information regarding the start of the test suite and the start of the crypto processor can be found in the Markdown file (readme.md) provided by the supervisor.

2.3 Zephyr (real time operating system)

2.3.1 General infromation

"The Zephyr OS is based on a small-footprint kernel designed for use on resource-constrained and embedded systems: from simple embedded environmental sensors and LED wearables to sophisticated embedded controllers, smart watches, and IoT wireless applications." $^{\rm 1}$

2.3.2 Features

"Zephyr intends to provide all components needed to develop resource-constrained and embedded or microcontroller-based applications. This includes, but is not limited to:

- A small kernel
- A flexible configuration and build system for compile-time definition of required resources and modules
- A set of protocol stacks (IPv4 and IPv6, OMA LWM2M, MQTT, 802.15.4, Bluetooth Low Energy, CAN)
- A virtual file system interface with several flash file systems for non-volatile storage
- Management and device firmware update mechanisms" ²

2.3.3 West

"The Zephyr project includes a swiss-army knife command line tool named west1. West is developed in its own repository. West's built-in commands provide a multiple repository management system with features inspired by Google's Repo tool and Git submodules. West is also "pluggable": you can write your own west extension commands which add additional features to west. Zephyr uses this to provide conveniences for building applications, flashing and debugging them, and more." ³

2.3.4 Ninja

Ninja is a tool designed for speed optimisation and contains only the most necessary functions to describe arbitrary dependency graphs. Due to the lack of syntax, it is not possible to express complex decisions. Instead, Ninja is intended to be used with a separate program that generates the input files. Autotools mean that a lot of factors have to be taken into account for the build time. Ninja build files allow Ninja to evaluate incremental builds quickly.

 $^{^{1}} Source:\ 17.04.2021,\ https://docs.zephyrproject.org/latest/introduction/index.html$

²Source: 17.04.2021, https://en.wikipedia.org/wiki/Zephyr_(operating_system)

³Source: 17.04.2021, https://docs.zephyrproject.org/latest/guides/west/index.html

2.3.5 Kconfig

"KConfig is a selection-based configuration system originally developed for the Linux kernel. It is commonly used to select build-time options and enable or disable features."

Through a curses-based or graphical menu interface, most users interact with KConfig. This is usually invoked by running "menuconfig" or "guiconfig". In this interface, the user selects the desired options and functions and saves a configuration file, which is then used as input for the build process. The Zephyr kernel and subsystems can be configured at build time to suit specific application and platform requirements. Configuration is done through Kconfig, the same configuration system used by the Linux kernel. The main goal is to provide configuration support without changing the source code.

2.3.6 Devicetree

The Zephyr RTOS has a devicetree, a data structure that simplifies the hardware description. As soon as a project is compiled here, a ".dts" file is automatically and immediately created which corresponds to the devicetree. However, this also speeds up the process, as the data structure of all components is simply transferred during booting and thus everything does not have to be described in detail.

2.4 Threads

Threads enable parallelism in programmes that would otherwise be executed sequentially. The thread itself is executed sequentially. Threads are, so to speak, processes that are executed line by line. However, it is possible to connect threads in parallel and ensure that the entire programme is no longer executed sequentially but in parallel. A symbolic structure of a programme with threads looks like this:

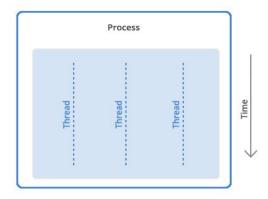


Figure 1: Process Structure (17.04.2021,

https://www.backblaze.com/blog/whats-the-diff-programs-processes-and-threads/)

⁴Source: 17.04.2021, https://docs.legato.io/19 11/toolsKconfig.html

2.5 Message Queues

"A message queue is a kernel object that implements a simple message queue, allowing threads and ISRs to asynchronously send and receive fixed-size data items. Any number of message queues can be defined (limited only by available RAM). Each message queue is referenced by its memory address. A message queue has the following key properties:

- A ring buffer of data items that have been sent but not yet received.
- $\bullet\,$ A data item size, measured in bytes.
- A maximum quantity of data items that can be queued in the ring buffer."⁵

⁵Source: 17.04.2021, https://docs.zephyrproject.org/latest/reference/kernel/data_passing/message_queues.html

3 Implementation

3.1 Github - Repo

 $https://github.com/Djordje-Stojanovic/FSST/tree/main/Stojanovic_DIC_PROJECT_2$

3.2 Code

Disclaimer: The basis on which I built the code was Patrick Wintners implementation idea. He helped me with the structure.

3.2.1 Main

```
//this is the main function
       void main(void)
               // Initializations
               uart_dev = device_get_binding(DT_LABEL(UART_DEVICE));
          // If wart is not found
               if(!uart_dev)
11
               //print that it couldnt be found
                       printk("UART could not be found\n");
           //else print that UART was found
              printk("UART has been found\n");
17
               //Configuration of the UART
19
               uartconf.baudrate = 9600; //initialization of a fixed baudrate
               \verb|vartconf.parity| = \verb|VART_CFG_PARITY_NONE|; //initialization| of a fixed parity|
21
               \verb| uartconf.stop\_bits = \verb| UART\_CFG\_STOP\_BITS\_1; // initialization of stop bits| \\
               {\tt uartconf.data\_bits} = {\tt UART\_CFG\_DATA\_BITS\_8;} \ // initialization \ of \ data \ bits
23
               {\tt uartconf.flow\_ctrl} = {\tt UART\_CFG\_FLOW\_CTRL\_NONE}; \ // initialization \ of \ flow \ control \ (none)
25
           //If UART couldnt be configured
26
               if(!uart_configure(uart_dev, &uartconf))
27
28
                //Print that it could not be configured
29
                        printk("Configuration of UART failed\n");
30
                        return;
31
              }
32
           //else print that it could been configured
33
               printk("UART configured\n");
34
35
               for(;;)
36
           {
                        printk("\nmain is waiting for death\n");
37
38
                        k_msleep(10*1000); // 10s time sleep
39
40
41
```

Listing 1: Stojanovic DIC Project2.c - Main

3.2.2 Initializations

```
//Definition of constant variables:
     //For the UART:
     #define UART_DEVICE DT_NODELABEL(uart0) //Nodelable according to Zephyr Documentation is used for
                                    // "Get a node identifier for a node label."
     #define sizeOfStack 1024 //Reporesents the Stack Size for the Thread
     #define preference 0 //Represents the preference of the thread
10
12
     //Further Initialization:
13
     enum {init, avail}; //initialize enum
14
     int state = init; //initialize state
     const struct device *uart_dev; //initialize the struct "device"
16
     struct uart_config uartconf; ///initialize the struct "uart_config"
17
18
     //Initialization for Wart
19
     void uartInput(void *, void *, void *); //Input of the UART
20
     void uartOutput(void*, void*, void*); //Output of the UART
21
     void uartProcess(void*, void*, void*); //Process of the UART
22
23
24
25
26
     //Thread Initialization:
27
     28
                  uartInput, NULL, NULL, NULL,
29
                  preference, 0, 0);
30
     31
                 uartOutput, NULL, NULL, NULL,
32
                  preference, 0, 0);
33
     34
                  uartProcess, NULL, NULL, NULL,
35
                  preference, 0, 0);
36
37
38
39
     // Message Queue Initialization:
40
     K_MSGQ_DEFINE(uartMessageQueue, 100*sizeof(char), 10, 1); //message Queue inizialization
41
42
     K_MSGQ_DEFINE(uartProcessMessageQueue, 100*sizeof(char), 10, 1); //message Queue Process initialization
43
```

Listing 2: Stojanovic DIC Project2.c - Initialization

3.2.3 Input-Function

```
2
      //this function is responsible for the uart input
3
       //there are 3 possible cases which will be stated bellow
       void uartInput(void *firstPointer, void *secondPointer, void *thirdPointer)
               // Initializations
               ARG_UNUSED(firstPointer);
9
               ARG_UNUSED(secondPointer);
10
               ARG_UNUSED(thirdPointer);
11
               unsigned char input;
12
13
           //next thing is an infinite loop
14
              for(;;)
15
           {
16
                       switch(state)
17
18
                   //if state is init:
19
                               case init:
                                       if(!uart_poll_in(uart_dev, &input))
20
21
                       { //if data is recieved, the following is printed.
                                              printk("Data has been recieved: %c\n", input);
22
23
                                               switch(input)
24
                               //if "." is recieved we should print "." back
25
26
                                                       case '.':
                                                               \label{eq:k_msgq_put} $$k_msgq_put(\&uartMessageQueue, ".\n", K_FOREVER);$
27
28
                                                               break;
                               //else if we got "P" then we should change the state to avail
29
30
                                                       case 'P':
                                                               printk("Sate is beeing changed to avail\n");
31
32
                                                               state = avail:
33
                                                               break;
34
                               //if we dont get P or . then it should break
35
                                                       default: break;
36
37
38
                   //if state is avail:
39
                              case avail: break; //the function is breaked
40
                               default: break;
41
42
                       k_msleep(1); //According to Zephyr documentation:
                            // "This routine puts the current thread to sleep for -duration- milliseconds."
43
44
               return;
46
```

Listing 3: Stojanovic_DIC_Project2.c - Input-Function

3.2.4 Output-Function

```
2
      //this function is responsible for the uart output
3
       void uartOutput(void *firstPointer, void *secondPointer, void *thirdPointer)
               // Initializations
               ARG_UNUSED(firstPointer);
              ARG_UNUSED(secondPointer);
9
               ARG_UNUSED(thirdPointer);
10
               unsigned char *output=malloc(100*sizeof(char));
11
12
          //next thing is an infinite loop
13
              for(;;)
14
15
               //resets\ the\ memmory
16
                       memset(output, 0, strlen(output));
               //Scan the message queue, if it is 0 then
17
18
                       19
                   //send data
20
21
                               printk("Sending the data: <%s>\n", output);
22
23
                               for(int i=0; i<strlen(output); i++)</pre>
24
                   {
25
                                       uart_poll_out(uart_dev, *(output+i));
26
                       //prints out the sent data
                                       printk("Data \ that \ has \ been \ sent: <\%x>\n", *(output+i));
27
                               7
28
29
30
                       \label{eq:k_msleep(1); //According to Zephyr documentation:} k\_msleep(1); //According to Zephyr documentation:
31
                            //\ {\it "This\ routine\ puts\ the\ current\ thread\ to\ sleep\ for\ -duration-\ milliseconds."}
32
33
               return:
34
35
```

Listing 4: Stojanovic DIC Project2.c - Output-Function

3.2.5 Process-Function

```
//this function is responsible for the uart process in general
 3
       void uartProcess(void *firstPointer, void *secondPointer, void *thirdPointer)
                // Initializations
               ARG_UNUSED(firstPointer);
               ARG_UNUSED(secondPointer);
               ARG_UNUSED(thirdPointer);
10
           //next thing is an infinite loop
              for(;;)
12
13
                        switch(state)
14
                    //here the process should break if it is in state init
16
                                case init:
18
                    /\!/if\ it\ is\ in\ state\ avail\ it\ should\ make\ sure\ to\ tell\ that\ processing\ is\ now\ avaliable
19
                                case avail:
20
                                         \verb|k_msgq_put(\&uartMessageQueue, "PROCESSING AVAILABLE \n", K_FOREVER);|
21
                        //telling that its gonna change state
22
                                        printk("Changing state the state back to init\n");
23
24
                                         state = init;
25
                                         break;
^{26}
                                default: break;
27
                        {\tt k\_msleep(1);} \ /\!/{\tt According} \ to \ {\tt Zephyr} \ documentation:
28
29
                              //\ {\it "This\ routine\ puts\ the\ current\ thread\ to\ sleep\ for\ -duration-\ milliseconds."}
30
                }
31
                return;
32
33
```

Listing 5: Stojanovic_DIC_Project2.c - Process-Function

4 Test of Function

4.1 Test 00 connection alive

This test has been passed sucessfuly:

Figure 2: Test: 00_connection_alive

$4.2 \ Test_01_availibility$

This test has been passed successfuly:

Figure 3: Test: 01_availibility