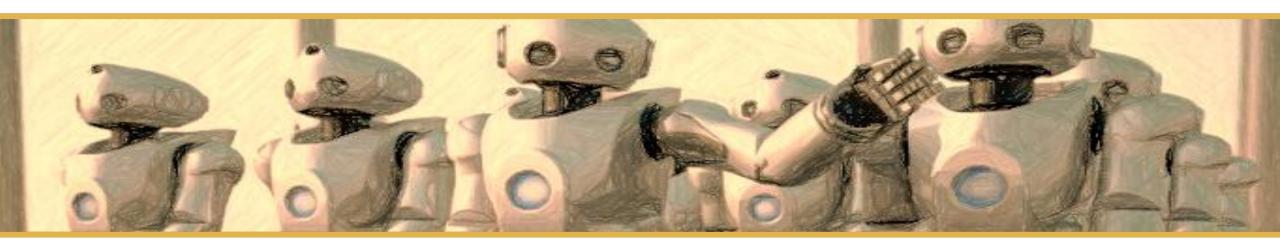


# Multiagent decision making and Coordination



#### **Outline**

- Introduction to multiagent coordination
- Coordination games
- Social conventions
- Social conventions with communication
- Coordination games + Social Conventions (with 3 agents)
- Roles
- Roles with communication
- Coordination Graphs



- What is coordination in multiagent systems?
  - Coordination is managing the interdependencies between activities
  - For example:
    - A non-sharable resource in the environment
    - Agents need to coordinate to use this resource

- What is coordination in multiagent systems?
  - Coordination is managing robustly and efficiently the interdependencies between activities
    - For example:
      - Two teammate soccer robots must coordinate their actions when deciding who should go for the ball

- Multiagent coordination is relevant for and multiagent cooperation
  - Cooperation is working together as a team to achieve a shared goal
  - But we should use coordination mechanisms within teamwork settings so that teams can efficiently achieve their goals

For example, in collaborative/cooperative agents, coordination mechanisms ensure that:

- Agents do not obstruct each other when taking actions
- Agents efficiently perform joint actions
- These actions serve the common goal of the team

Informally:

■ Coordination can be regarded as the process where every agent's individual decision leads to a good joint action for the group

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- We can model a coordination problem as a coordination game using the tools from game theory
  - We can use the normal-form representation
    - Actions sets and payoffs
  - Solve it using some solution concept
    - Nash equilibrium

- Consider the stag hunt game with two hunters:
  - If they both hunt hares, they each capture half of the hares in the range



If one hunts the stag and the other hunts hares, the stag hunter goes home empty-handed while the hare hunter captures all the hares



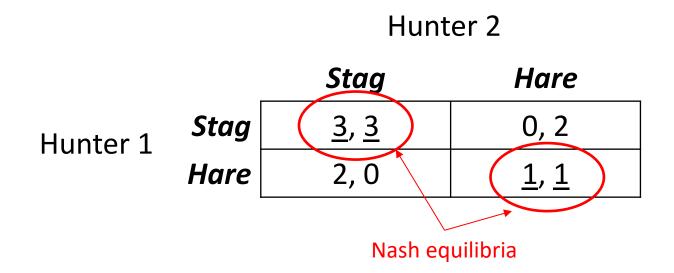
• Finally, if both hunt the stag, then each of their shares of the stag is greater than the value of all the hares

■ The payoff matrix of the **stag hunt game**:

Hunter 2

		Stag	Hare
Hunter 1	Stag	3, 3	0, 2
	Hare	2, 0	1, 1

What is the NE in the stag hunt game?



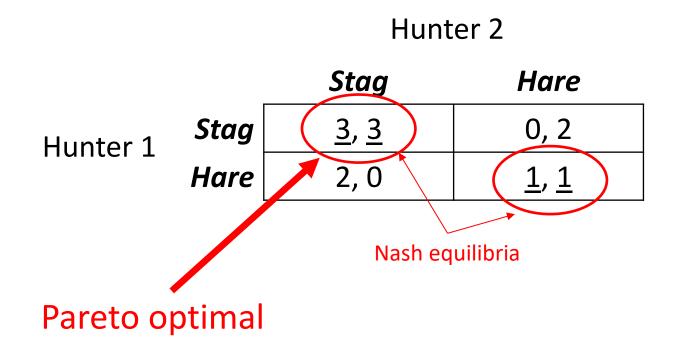
Which equilibrium should the hunters choose?

- **No coordination** in the stag hunt game:
  - The action to hunt a *Stag* is a **risk-taking strategy** 
    - If both hunters choose to hunt the Stag, then both get the highest payoff
    - If only one hunter chooses to hunt the Stag then he gets the lowest payoff
  - The action to hunt a *Hare* is a conservative strategy
    - The hunter will never get the highest payoff or lowest payoff
    - The hunter will **never go home empty-handed**

- With coordination in the stag hunt game:
  - What if both hunters enter into some kind of agreement (in advance) and both trust each other?
    - One agreement could be to choose the joint action that is Pareto optimal (or strictly Pareto efficient)

- **Definition:** A joint action a **Pareto dominates** joint action a' if for all  $i \in N$ ,  $u_i(a) \ge u_i(a')$ , and there exists some  $j \in N$  for which  $u_j(a) > u_j(a')$
- **Definition:** A joint action a is **Pareto optimal**, or strictly Pareto efficient, if there does not exist another joint action  $a' \in A$  that Pareto dominates a

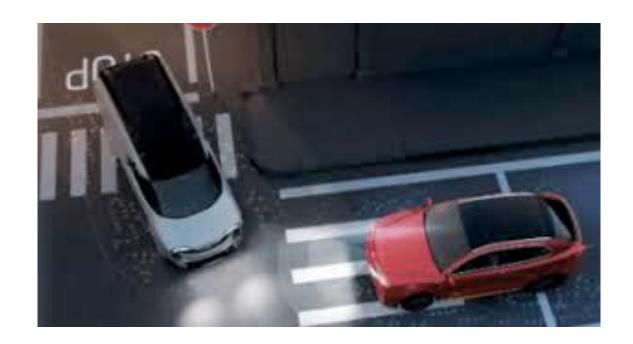
What is the NE in the stag hunt game?



- We can now formally define **coordination**:
  - The process in which a group of agents choose a single Pareto optimal Nash equilibrium in a game



Example of two autonomous vehicles at a crossroad



Example:

Each agent wants to cross first

But if they both cross they will crash

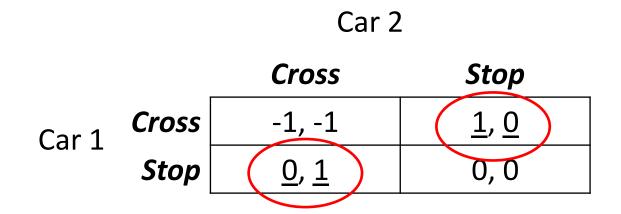
■ How can we represent this game with the normal-form representation and solve it?

■ The payoff matrix of the **car coordination game**:

r	•
	L

		Cross	Stop
Car 1	Cross	-1, -1	1, 0
	Stop	0, 1	0, 0

- How do we solve this coordination game?
  - Pareto optimal Nash equilibria



Which equilibrium should the agents choose?



■ Robot soccer example (within a **teamwork** or **collaborative** setting):

Two robots (of the same team) want to run to get the ball

But if they both try to get the ball at the same time, they will crash

• How can we represent this game with the normal-form representation and solve it?

Robot 2

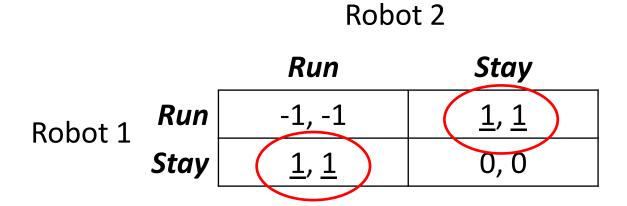
		Run	Stay
Robot 1	Run	-1, -1	1, 1
	Stay	1, 1	0, 0

• In the case of n collaborative agents, all agents in the team share the same payoff function:

$$u_1(a) = \dots = u_n(a) \equiv u(a)$$

■ These game are known as **pure coordination games** or **team games** 

- How do we solve this coordination game?
  - Pareto optimal Nash equilibria



Which equilibrium should the agents choose?

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- To solve a coordination problem:
  - A group of agents are faced with the problem of how to choose their actions in order to select the same Nash equilibrium
  - Unfortunately, there is not a single recipe that can be used for all games
  - Nevertheless, we can devise recipes that will instruct the agents on how to choose a single equilibrium for a given game

- A social convention (or social law) is such a recipe
  - It places constraints on the possible action choices of the agents
  - It can be regarded as a rule that dictates how the agents should choose their actions in a coordination game in order to reach an equilibrium
  - Moreover, given that the convention has been established and is common knowledge among agents, no agent can benefit from not abiding by it

- **Boutilier (1996)** proposed a **general convention** that achieves coordination in a large class of systems and **is very easy to implement**:
  - The convention assumes a unique ordering scheme of joint actions that is common knowledge among agents
  - In a particular game, each agent first computes all equilibria of the game, and then selects
    the first equilibrium according to this ordering scheme

Two agents who want to go to the movies together but not alone

Agent 2

		Thriller	Comedy
Agent 1	Thriller	1, 1	0, 0
	Comedy	0, 0	1, 1

■ Each agent first computes all equilibria of the game



- Then each agent selects the first equilibrium according to an ordering scheme:
  - Order the agents by 1 > 2
    - meaning that agent 1 has 'priority' over agent 2
  - Order the actions by *Thriller* > *Comedy*

What action should each agent choose?

Following the ordering scheme:

Agent 2

		Thriller	Comedy
Agent 1	Thriller	$\left(\underline{1},\underline{1}\right)$	0, 0
	Comedy	0, 0	<u>1</u> , <u>1</u>

• Hence, the first equilibrium in the resulting ordering of joint actions is:
(Thriller, Thriller)

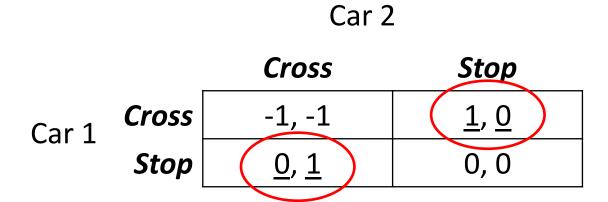
- And this will be the unanimous choice of the agents
- Now each agent can choose his individual action

• Recall the **car coordination game**:

•	1	r	,
•	. a		

		Cross	Stop
Car 1	Cross	-1, -1	1, 0
	Stop	0, 1	0, 0

■ Each agent first computes all equilibria of the game



## **Social Conventions**

- Then each agent selects the first equilibrium according to an ordering scheme:
  - The driver coming from the right will always have priority over the other driver
    - Note that agents need to have more information about the environment and/or other agents (and not just the payoffs and action sets)
    - Let us assume that Car 1 sees Car 2 coming from the right
      - Then Car 2 has 'priority' over Car 1
  - Order the actions by Cross > Stop
- What action should each agent choose?

## **Social Conventions**

Following the ordering scheme:

 Cross
 Stop

 Car 1
  $\frac{1}{5}$  

 Car 1
  $\frac{1}{5}$  

 Car 1
  $\frac{1}{5}$ 
 $\frac{0}{5}$   $\frac{1}{5}$ 
 $\frac{0}{5}$   $\frac{1}{5}$ 
 $\frac{1}{5}$   $\frac{1}{5}$ </td

#### **Social Conventions**

Hence, the first equilibrium in the resulting ordering of joint actions is:
(Stop, Cross)

- And this will be the unanimous choice of the agents
- Now each agent can choose his individual action

#### **Exercise**

#### Battle of the sexes:

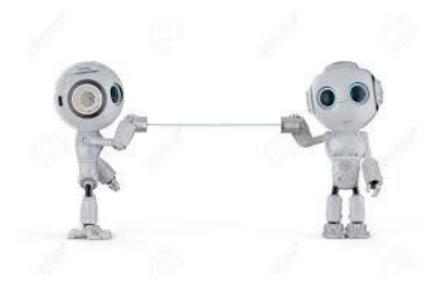
- A man and woman want to get together for an evening of entertainment, but they have no means of communication
- They can either go to the ballet or the fight
  - The man prefers going to the fight
  - The woman prefers going to the ballet
  - But they both prefer being together than being alone
- Create an ordering scheme and solve the coordination problem

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- When **communication** is available:
  - We only need to impose an ordering of the agents
    - And this ordering is common knowledge



- We first assume an ordering i = 1, ..., n of agents
- Coordination can now be achieved by the following algorithm:
  - 1. Each agent i (except agent 1) waits until all previous agents 1, ..., i-1 in the ordering have broadcast their chosen actions
  - 2. agent i computes its component  $a_i^*$  of an equilibrium that is consistent with the choices of the previous agents
  - 3. agent i broadcasts  $a_i^*$  to all agents that have not chosen an action yet

Note that:

■ The fixed ordering of the agents together with the wait/send primitives induce a synchronized sequential execution order of the coordination algorithm

■ Recall the **car coordination game**:

	$\mathbf{a}$	r	•
Ι.	<b>a</b>		

		Cross	Stop
Car 1	Cross	-1, -1	1, 0
Cai I	Stop	0, 1	0, 0

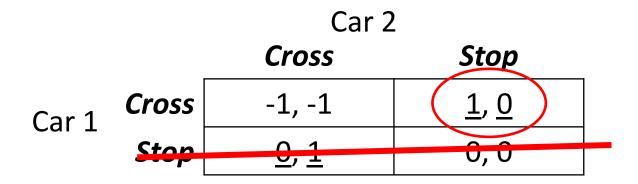
Let us assume the following ordering of agents:

• Car 1 chooses action  $a_1^*$  = *Cross* of an equilibrium

		Car 2	•
		Cross	Stop
Car 1	Cross	-1, -1	<u>1</u> , <u>0</u>
	Stop	<u>0</u> , <u>1</u>	0, 0

- Car 1 broadcasts  $a_1^*$  = Cross to all agents that have not chosen an action yet
  - i.e., sends a message to Car 2

- Car 2 waits until all previous agents in the ordering have broadcast their chosen actions
  - Thus, it waits a message from Car 1:  $a_1^*$  = *Cross*
- Car 2 computes its component  $a_2^*$  of an equilibrium that is consistent with the choices of the previous agents
  - Thus,  $a_2^*$  = *Stop*



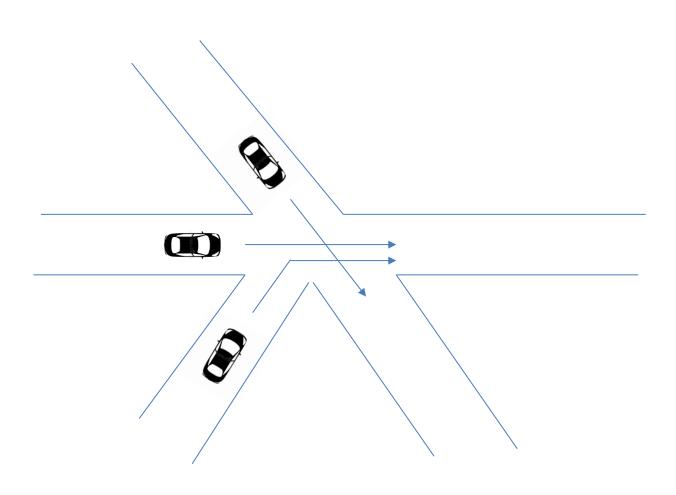
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Example of THREE autonomous vehicles at a crossroad





Example:

Each agent/car wants to cross first

But if they two cross they will crash

■ How can we represent this game with the normal-form representation and solve it?

Joint actions ( $a$ )	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	1	0	0
(cross, cross, stop)	-1	-1	0
(cross, stop, cross)	-1	0	-1
(cross, cross, cross)	-1	-1	-1
(stop, stop, stop)	0	0	0
(stop, cross, stop)	0	1	0
(stop, stop, cross)	0	0	1
(stop, cross, cross)	0	-1	-1

Let us calculate the best response functions:

■ 
$$B_1(a_2 = stop, a_3 = stop) = ?$$

■ 
$$B_1(a_2 = cross, a_3 = stop) = ?$$

■ 
$$B_1(a_2 = stop, a_3 = cross) = ?$$

$$\blacksquare B_1(a_2 = cross \ a_3 = cross) = ?$$

■ 
$$B_2(a_1 = stop, a_3 = stop) = ?$$

■ 
$$B_2(a_1 = cross, a_3 = stop) = ?$$

• 
$$B_2(a_1 = stop, a_3 = cross) = ?$$

■ 
$$B_2(a_1 = cross, a_3 = cross) = ?$$

• 
$$B_3(a_1 = stop, a_2 = stop) = ?$$

• 
$$B_3(a_1 = cross, a_2 = stop) = ?$$

• 
$$B_3(a_1 = stop, a_2 = cross) = ?$$

• 
$$B_3(a_1 = cross \ a_2 = cross) = ?$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	(1)	0	0
(stop, stop, stop)	0	0	0

	stop, stop
cross	1,0,0
stop	0, 0, 0

$$B_1(a_2 = stop, a_3 = stop) = cross$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, cross, stop)	-1	-1	0
(stop, cross, stop)	(0)	1	0

	cross, stop
cross	-1, -1, 0
stop	0,1,0

$$B_1(a_2 = cross, a_3 = stop) = stop$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, cross)	-1	0	-1
(stop, stop, cross)	$\left(\begin{array}{c}0\end{array}\right)$	0	1

	stop, cross
cross	-1, 0, -1
stop	0,0,1

$$B_1(a_2 = stop, a_3 = cross) = stop$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, cross, cross)	-1	-1	-1
(stop, cross, cross)	(0)	-1	-1

	cross, cross
cross	-1, -1, -1
stop	0, 1, 1

$$B_1(a_2 = cross, a_3 = cross) = stop$$

Joint actions ( $a$ )	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	1	0	0
(cross, cross, stop)	-1	-1	0
(cross, stop, cross)	-1	0	-1
(cross, cross, cross)	-1	-1	-1
(stop, stop, stop)	0	0	0
(stop, cross, stop)	0	1	0
(stop, stop, cross)	0	0	1
(stop, cross, cross)	0	-1	-1

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(stop, stop, stop)	0	0	0
(stop, cross, stop)	0	(1)	0

	stop, stop
stop	0, 0, 0
cross	0, 1,0

$$B_2(a_1 = stop, a_3 = stop) = cross$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	1	( 0 )	0
(cross, cross, stop)	-1	-1	0

	cross, stop
stop	(, 0, 0
cross	-1, -1, 0

$$B_2(a_1 = cross, a_3 = stop) = stop$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(stop, stop, cross)	0	( 0 )	1
(stop, cross, cross)	0	-1	-1

	stop, cross
stop	0, 0, 1
cross	0, -1, -1

$$B_2(a_1 = stop, a_3 = cross) = stop$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, cross)	-1	( 0 )	-1
(cross, cross, cross)	-1	-1	-1

	cross, cross
stop	-(, 0, 1
cross	-1, -1, -1

$$B_2(a_1 = cross, a_3 = cross) = stop$$

Joint actions ( $a$ )	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	1	0	0
(cross, cross, stop)	-1	-1	0
(cross, stop, cross)	-1	0	-1
(cross, cross, cross)	-1	-1	-1
(stop, stop, stop)	0	0	0
(stop, cross, stop)	0	1	0
(stop, stop, cross)	0	0	1
(stop, cross, cross)	0	-1	-1

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(stop, stop, stop)	0	0	0
(stop, stop, cross)	0	0	(1)

	stop, stop
stop	0, 0, 0
cross	0, 0, 1

$$B_3(a_1 = stop, a_2 = stop) = cross$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	1	0	( 0 )
(cross, stop, cross)	-1	0	-1

	cross, stop
stop	1, 0, 0
cross	-1, 0, -1

$$B_3(a_1 = cross, a_2 = stop) = stop$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(stop, cross, stop)	0	1	( 0 )
(stop, cross, cross)	0	-1	-1

	stop, cross	
stop	0, 1, 0	
cross	0, -1, -1	

$$B_3(a_1 = stop, a_2 = cross) = stop$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, cross, stop)	-1	-1	( 0 )
(cross, cross, cross)	-1	-1	-1

	cross, cross
stop	-1, -(1, 0)
cross	-1, -1, -1

$$B_3(a_1 = cross, a_2 = cross) = stop$$

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	1	0	0
(cross, cross, stop)	-1	-1	0
(cross, stop, cross)	-1	0	-1
(cross, cross, cross)	-1	-1	-1
(stop, stop, stop)	0	0	0
(stop, cross, stop)	0	1	0
(stop, stop, cross)	0	0	1
(stop, cross, cross)	0	-1	-1

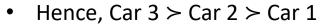
## **Coordination Games & Social Conventions**

- Then each agent selects the first equilibrium according to an ordering scheme:
  - The driver coming from the right will always have priority over the other driver
  - Order the actions by Cross > Stop

What action should each agent choose?

## **Coordination Games & Social Conventions**

- Car 1 sees Car 2 and Car 3 coming from the right
- Car 2 sees Car 3 coming from the right



1

2



# **Coordination Games & Social Conventions**

Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
(cross, stop, stop)	1	0	0
(cross, cross, stop)	-1	-1	0
(cross, stop, cross)	-1	0	-1
(cross, cross, cross)	-1	-1	-1
(stop, stop, stop)	0	0	0
(stop, cross, stop)	0	1	0
(stop, stop, cross)	0	0	1
(stop, cross, cross)	0	-1	-1

Now each agent can choose his individual action

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- Coordination by social conventions relies on the following assumption:
  - An agent can compute all equilibria in a game before choosing a single one
- However, computing equilibria can be expensive when the agents' action sets are large

- Hence, it makes sense to start trying to reduce the size of the action sets
- Such a reduction can lead to:
  - computational advantages in terms of speed
  - simplify the equilibrium selection problem

■ A natural way to reduce the agents' action sets is to assign roles to the agents

■ Formally:

Given a particular state, a role can be viewed as a masking operator on the agent's action
 set

In practical terms:

• if an agent is assigned a role at a particular state, then some of the agent's actions are deactivated at this state

For instance, a **robot soccer agent** that is currently with a **defender role cannot attempt to** *Score* 



- **Example**: two agents who want to go to the movies together but not alone
- if Agent 2 is assigned a role that forbids him to select the action Thriller
  - Because Agent 2 is under 12



So how can we assign a role to an agent?

- Suppose that:
  - there are n available roles (not necessarily distinct)
  - the state is fully observable to the agents

- Suppose that the following facts are common knowledge among agents:
  - There is a **fixed ordering**  $\{1,2,...,n\}$  **of the roles** 
    - In other words, role 1 must be assigned first, followed by role 2, etc.

- Suppose that the following facts are common knowledge among agents:
  - For each role there is a function that assigns to each agent a 'potential'
    - The potential reflects how appropriate that agent is for the specific role, given the current state
  - For example, in robot soccer scenario:
    - the potential of a robot for the role attacker can be given by its negative Euclidean distance to the ball

- Suppose that the following facts are common knowledge among agents:
  - Each agent can be assigned only one role

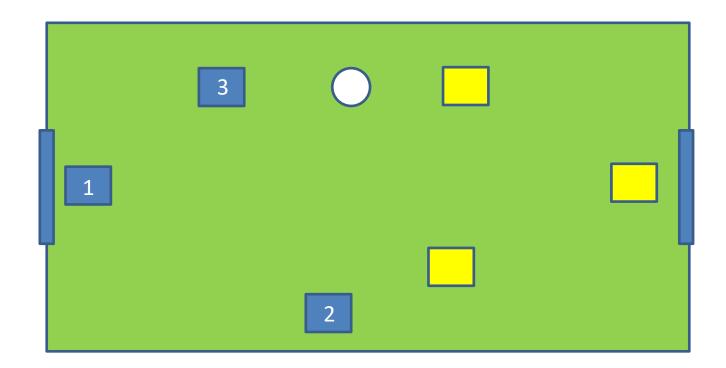
- Then role assignment can be carried out, for instance, by a greedy algorithm:
  - each role (starting from role 1) is assigned to the agent that has the highest potential for that role
  - and so on until all agents have been assigned a role

■ Example: Robot soccer



**3** robots: 1, 2, and 3

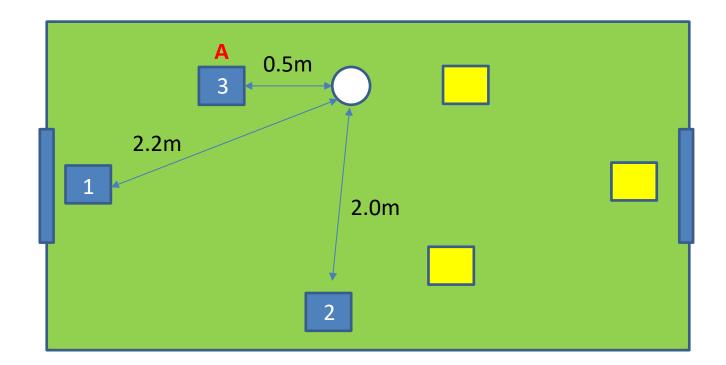
• Action sets:  $A_1$ ,  $A_2$ ,  $A_3 = \{run, stay\}$ 



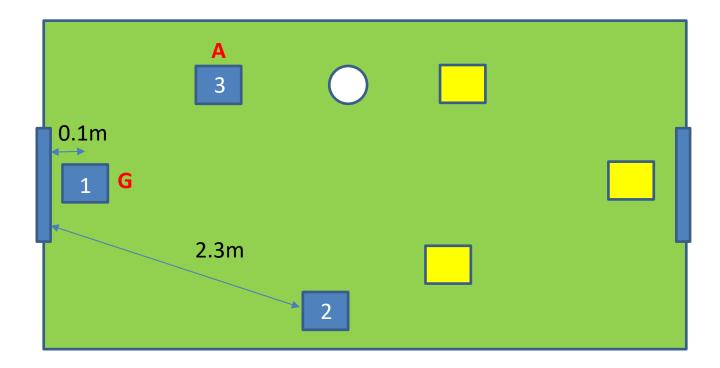
Joint actions (a)	$u_1(a)$	$u_2(a)$	$u_3(a)$
( <u>run</u> , <u>run</u> , <u>run</u> )	-1	-1	-1
(run, run, <u>stay</u> )	-1	-1	-1
(run, stay, run)	-1	-1	-1
( <u>run</u> , <u>stay</u> , <u>stay</u> )	1	1	1
(stay, run, run)	-1	-1	-1
(stay, run, stay)	1	1	1
(stay, stay, run)	1	1	1
(stay, stay, stay)	0	0	0

- Fixed ordering of the roles = {Attacker, Goalkeeper, Defender}
- Potential for role Attacker: negative Euclidean distance to the ball
- Potential for role Goalkeeper: negative Euclidean distance to its own team's goal
- Potential for role Defender: positive Euclidean distance to the ball

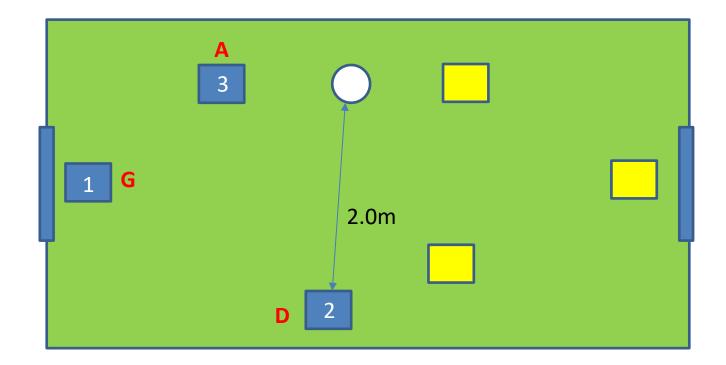
- Assigning the Attacker role (A)
  - Potential is the negative Euclidean distance to the ball



- Assigning the Goalkeeper role (G)
  - Potential is the negative Euclidean distance to its own team's goal



- Assigning the **Defender** role (**D**)
  - Potential is the positive Euclidean distance to the ball



- Hence, if an agent is assigned a role at a particular state then some of the agent's actions
   are deactivated at this state
  - The role **Defender** and **Goalkeeper** cannot choose the action run (so they can only choose stay)
  - lacksquare The role **Attacker** can either choose the action run or stay

Joint actions ( $a$ )	$u_1(a)$	$u_2(a)$	$u_3(a)$
(stay, stay, run)	1	1	1
(stay, stay, stay)	0	0	0

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### Roles with communication

- When **communication** is **not** available:
  - Each agent can run the previous algorithm identically and in parallel
    - Assuming that each agent can compute the potential of each other agent

### Roles with communication

- When communication is available:
  - An agent only needs to compute its own potentials for the set of roles
  - Then broadcast them to the rest of the agents.
  - Next it can wait for all other potentials to arrive in order to compute the assignment of roles to agents as the previous algorithm

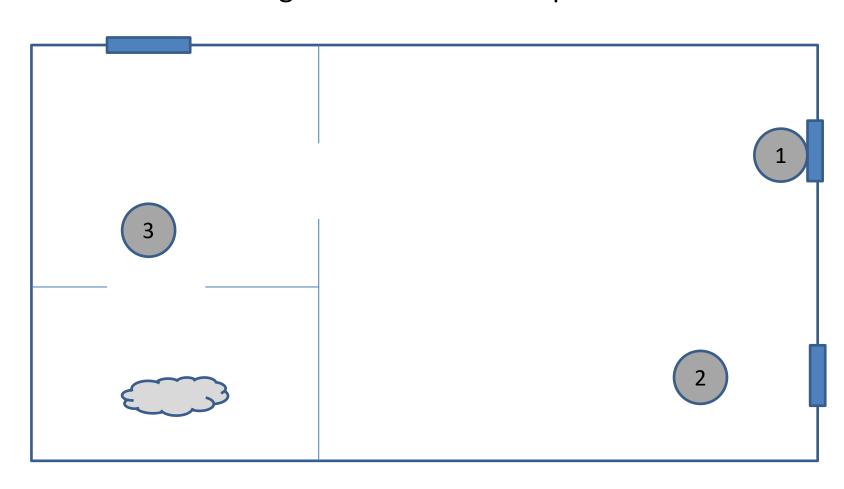
## Roles with communication

```
For each agent i in parallel
    I = \{\}.
    For each role j = 1, \ldots, n
         Compute the potential r_{ij} of agent i for role j.
         Broadcast r_{ij} to all agents.
    End
    Wait until all r_{i'j}, for j = 1, ..., n, are received.
    For each role j = 1, \ldots, n
         Assign role j to agent i^* = \arg \max_{i' \notin I} \{r_{i'j}\}.
         Add i^* to I.
    End
End
```

## **Final remarks**

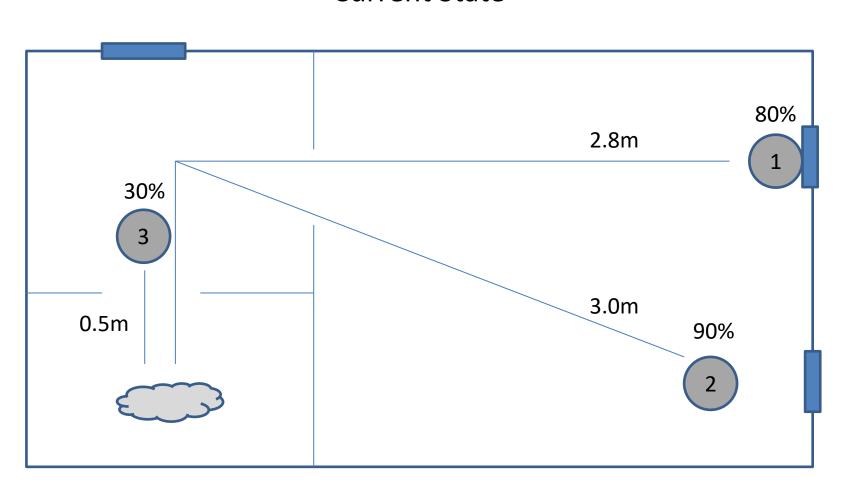
- In the communication-free algorithm
  - The greedy algorithm has a complexity of  $O(n^2)$
- In the communication-based algorithm
  - Each agent needs to compute O(n) potentials (its own)
  - However, this is added to the total number  $O(n^2)$  of potentials that need to be broadcast and processed by the agents

Cleaning robot coordination problem



- 3 cleaning robots: 1, 2, and 3
- Action sets:  $A_1$ ,  $A_2$ ,  $A_3 = \{clean, charge, stay\}$
- Ordering of roles = {Need Charging (NC), Go Clean (GC), Wait for Work (WfW)}
- Potential for NC: negative of the battery level
- Potential for GC: negative distance to dirt
- Potential for WfW: positive distance to dirt

#### **Current State**



- Deactivated actions for each role:
  - Need Charging (NC): cannot clean and cannot stay
  - Go Clean (GC): cannot stay and cannot charge
  - Wait for Work (WfW): cannot clean

 Use the communication-based algorithm to assign the roles and solve the coordination problem

- Role assignment
  - Need Charging (NC): 3
  - Go Clean (GC): 1
  - Wait for Work (WfW): 3

Can we improve this?

- Change the ordering?
- Role admissibility / activation conditions
  - Add a function that tests if the role is needed in the current state

NC.available(A)

If A.battery\_level < 15%

## **Outline**

- Introduction to multiagent coordination
- Coordination games
- Social conventions
- Social conventions with communication
- Coordination games + Social Conventions (with 3 agents)
- Roles
- Roles with communication
- Coordination Graphs



- Roles can help us solve coordination game by reducing the agents' action sets prior to computing the equilibria
- However, computing equilibria in a subgame can still be a difficult task when the number of agents is large
- Recall that the number of joint actions is exponential in the number of agents:

 $m^n$ 

where m is the number of the agents' actions and n is the number of agents

- We also need a method that reduces the number of agents in a coordination game
- Coordination graph (Guestrin et al., 2002) is a framework for solving large-scale coordination problems

- A coordination graph decomposes a coordination game into several smaller subgames that are easier to solve
  - In roles, we reduce the action sets and form a single subgame
  - In **coordination graphs**, we form **various subgames**, each involving a smaller number of agents

- Main assumption in coordination graphs:
  - the global payoff function u(a) can be written as a linear combination of k local payoff functions  $f_i$ , for  $i=1,\ldots,k$ , each involving fewer agents
- For instance, a coordination problem with n=4 agents and k=3 local payoff functions, each involving two agents, could have the following decomposition:

$$u(a) = f_1(a_1, a_2) + f_2(a_1, a_3) + f_3(a_3, a_4)$$

Example of the decomposition of the global payoff function:

$$u(a) = f_1(a_1, a_2) + f_2(a_1, a_3) + f_3(a_3, a_4)$$

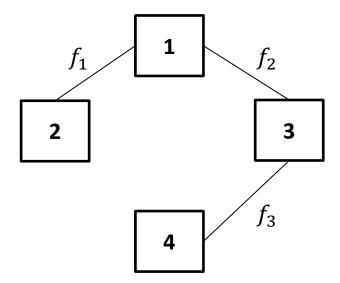
Note:

- $f_1(a_1, a_2)$  involves only agents 1 and 2, with their actions  $a_1$  and  $a_2$
- $f_2(a_1, a_3)$  involves only agents 1 and 3, with their actions  $a_1$  and  $a_3$
- $f_3(a_3, a_4)$  involves only agents 3 and 4, with their actions  $a_3$  and  $a_4$

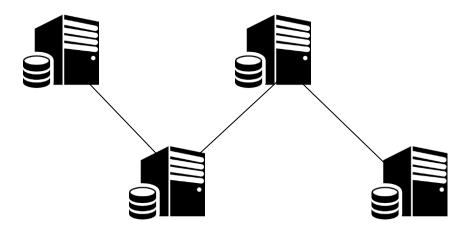
Example of the decomposition of the global payoff function:

$$u(a) = f_1(a_1, a_2) + f_2(a_1, a_3) + f_3(a_3, a_4)$$

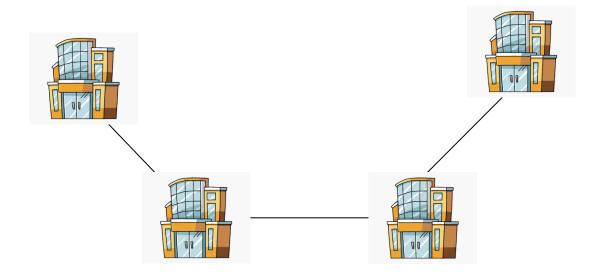
■ The decomposition can be represented with the following graph:



- Practical examples that could use coordination graphs:
  - Servers in a computer network that need to coordinate in order to optimize the overall network traffic



- **Practical examples** that could use coordination graphs:
  - Within a corporation, offices in nearby cities may need to coordinate in order to maximize global sales

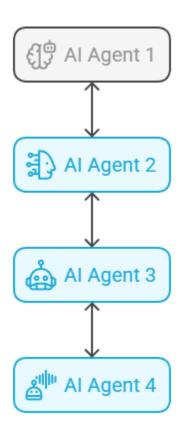


- Practical examples that could use coordination graphs:
  - In a soccer team, nearby players may need to coordinate in order to improve team performance

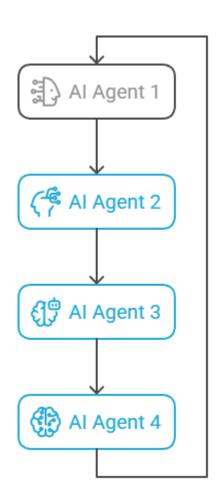




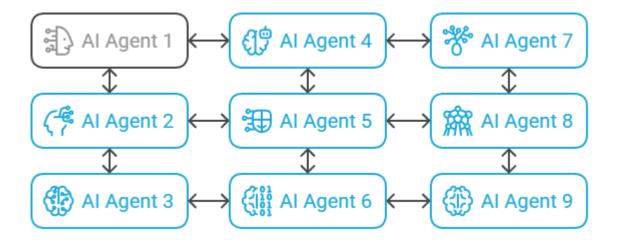
Linear coordination



Cycle/ring coordination



Grid coordination

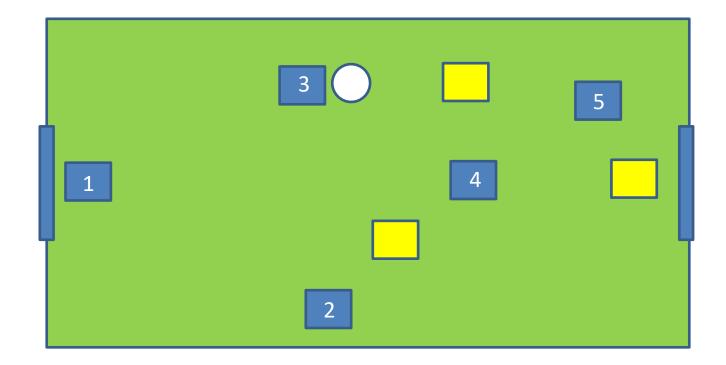


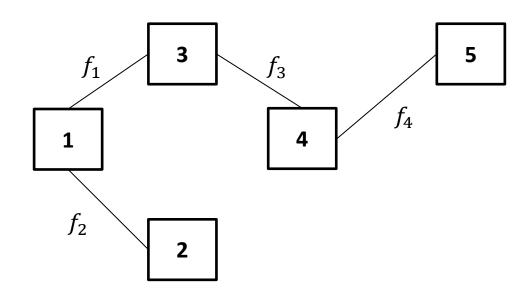
■ Example: Robot soccer



■ 5 robots: 1, 2, 3, 4 and 5

• Action sets:  $A_1$ ,  $A_2$ ,  $A_3$ ,  $A_4$ ,  $A_5 = \{stay, pass\}$ 





$$u(a) = f_1(a_1, a_3) + f_2(a_1, a_2) + f_3(a_3, a_4) + f_4(a_4, a_5)$$

# **Thank You**



rui.prada@tecnico.ulisboa.pt