# .\*VSTU.\*

# Team Reference Document

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## 1. Code Templates

```
1.1. Basic Configuration.
1.1.1. .vimrc.
set cin nu ts=2 sw=2 sts=2 mouse=a
function! Compile()
    :!q++ -std=qnu++11 -q % -o %<.exe
endfunction
function! Run()
    :!time ./%<.exe
endfunction
map <F4> :call Compile()<cr>
map <F5> :call Run()<cr>
map <C-A> qqVG"+y
1.1.2. stress and template.
// g++ -std=c++11 main.cpp -o main -D"_DEBUG_TEMICH_"
#include <algorithm>
#include <cmath>
#include <functional>
#include <iostream>
#include <map>
#include <queue>
#include <set>
#include <sstream>
#include <string>
#include <vector>
using namespace std;
using LL = long long;
using pii = pair<int, int>;
#define X first
#define Y second
struct Solver {
  void solve(istream& cin, ostream& cout) {
    int a. b:
    cin >> a >> b;
    cout << a + b << endl;
};
struct Brute {
  void solve(istream& cin, ostream& cout) {
    int a. b:
    cin >> a >> b;
    while (b--) ++a:
    cout << a << endl;</pre>
```

```
};
template <typename Solution>
struct SolutionStr {
  string solve(string input) {
    istringstream is(input);
    ostringstream os;
    Solution().solve(is, os);
    return os.str():
};
string gen_input(int it) {
  (void)it;
  return "10 20";
}
void stress() {
  for (int it = 0; it < 1000; ++it) {
    auto input = gen_input(it);
    auto brute_out = SolutionStr<Solver>().solve(input);
    auto sol_out = SolutionStr<Brute>().solve(input);
    if (sol_out != brute_out) {
      cerr << "WA #" << it << endl;
      cerr << "input: " << endl;</pre>
      cerr << input << endl;</pre>
      cerr << "expected: " << brute_out << endl;</pre>
      cerr << "got: " << sol_out << endl;</pre>
      exit(1);
  }
  cerr << "OK" << endl;</pre>
}
int main() {
  #ifdef _DEBUG_TEMICH_
  stress():
  #endif
  Solver().solve(cin, cout);
1.2. Vector.
struct Vec {
  explicit Vec(LL x = 0, LL y = 0) : x(x), y(y) {}
  Vec operator+(const Vec& o) const {
    return Vec(x + o.x, y + o.y); }
  Vec operator-(const Vec& o) const {
    return Vec(x - o.x, y - o.y); }
  Vec operator*(const LL p) const {
    return Vec(x * p, y * p); }
  double len() const { return sqrt(x * x + y * y); }
  LL cross(const Vec& o) const { return x * o.y - y * o.x; }
  LL dot(const Vec& o) const { return x * o.x + y * o.y; }
  static Vec read(istream& cin) {
```

```
LL x, y;
    cin >> x >> y;
    return Vec(x, y);
};
bool cmp(Vec a, Vec b) {
  return a.x < b.x \mid \mid (a.x == b.x \&\& a.y < b.y);
bool cw(Vec a, Vec b, Vec c) {
  return (b - a).cross(c - b) < 0;
bool ccw(Vec a, Vec b, Vec c) {
  return (b - a).cross(c - b) > 0:
void convex_hull(vector<Vec> & a) {
  if (a.size() == 1) return;
  sort(a.begin(), a.end(), &cmp);
  Vec p1 = a[0], p2 = a.back();
  vector<Vec> up, down;
  up.push_back(p1);
  down.push_back(p1);
  for (size_t i=1; i<a.size(); ++i) {</pre>
    if (i==a.size()-1 || cw(p1, a[i], p2)) {
      while (up.size()>=2
          && !cw(up[up.size()-2], up[up.size()-1], a[i]))
        up.pop_back();
      up.push_back (a[i]);
    if (i == a.size()-1 | | ccw(p1, a[i], p2)) {
      while (down.size()>=2
          && !ccw(down[down.size()-2],
            down[down.size()-1], a[i]))
        down.pop_back();
      down.push_back(a[i]);
   }
  a.clear();
  for (size_t i=0; i<up.size(); ++i)</pre>
    a.push_back(up[i]);
  for (size_t i=down.size()-2; i>0; --i)
    a.push_back(down[i]);
}
1.3. FFT.
struct Complex {
  long double re, im;
  explicit Complex(long double re = 0,
      long double im = 0) : re(re), im(im) {}
  Complex operator+(const Complex& o) const {
    return Complex(re + o.re, im + o.im); }
  Complex operator-(const Complex& o) const {
    return Complex(re - o.re, im - o.im); }
  Complex operator*(const Complex& o) const {
    return Complex(re * o.re - im * o.im. re * o.im + im * o.re)
};
```

```
const int MAX_SHIFT = 22;
const int MAX_N
                     = 1 << MAX_SHIFT;
const double Pi = acos(-1);
Complex roots[MAX_N / 2];
int bit_reverse[MAX_N];
void prep() {
  bit_reverse[0] = 0;
  for (int i = 1; i < MAX_N; ++i)
    bit_reverse[i] = (bit_reverse[i >> 1]
        | ((i & 1) << MAX_SHIFT)) >> 1;
  for (int i = 0; i + i < MAX_N; ++i) {
    double angle = 2 * i * Pi / MAX_N;
    roots[i] = Complex(cos(angle), sin(angle));
  }
}
Complex arr[MAX_N];
void fft(int k) {
  assert(k <= MAX_SHIFT);</pre>
  const int n = 1 \ll k;
  for (int i = 0; i < n; ++i) {
    int rv = bit_reverse[i] >> (MAX_SHIFT - k);
    if (rv < i) swap(arr[i], arr[rv]);</pre>
  }
  for (int bs = 2; bs <= n; bs *= 2) {
    const int hbs = bs / 2;
    const int factor = (MAX_N / 2) / hbs;
    for (int i = 0; i < n; i += bs) {
      for (int j = 0; j < hbs; ++j) {
        auto a = arr[i + j];
        auto b = arr[i + j + hbs] * roots[factor * j];
        arr[i + j] = a + b;
        arr[i + j + hbs] = a - b;
const int Base = 100;
void square(vector<int>& number) {
  int sz = number.size() * 2;
  int k = 1;
    int rsz = 2;
    while (rsz < sz) {</pre>
      rsz *= 2;
      ++k;
    }
```

```
sz = rsz;
assert(sz <= MAX_N);</pre>
for (int i = 0; i < sz; ++i)
 arr[i] = Complex(i < number.size() ? number[i] : 0);</pre>
fft(k):
for (int i = 0; i < sz; ++i)
 arr[i] = arr[i] * arr[i];
fft(k);
reverse(arr + 1, arr + sz);
number.resize(sz);
int cr = 0;
for (int i = 0; i < sz; ++i) {
 number[i] = cr + int(arr[i].re / sz + 0.5);
 cr = number[i] / Base;
 number[i] %= Base;
while (number.back() == 0) number.pop_back();
```

#### 2. Misc

## 2.1. Debugging Tips.

- Stack overflow? Recursive DFS on tree that is actually a long path?
- Floating-point numbers
  - Getting NaN? Make sure acos etc. are not getting values out of their range (perhaps 1+eps).
  - Rounding negative numbers?
  - Outputting in scientific notation?
- Wrong Answer?
  - Read the problem statement again!
  - Are multiple test cases being handled correctly? Try repeating the same test case many times.
  - Integer overflow?
  - Think very carefully about boundaries of all input parameters
  - Try out possible edge cases:
    - \*  $n = 0, n = -1, n = 1, n = 2^{31} 1$  or  $n = -2^{31}$
    - \* List is empty, or contains a single element
    - \* n is even, n is odd
    - \* Graph is empty, or contains a single vertex
    - \* Graph is a multigraph (loops or multiple edges)
    - \* Polygon is concave or non-simple
  - Is initial condition wrong for small cases?
  - Are you sure the algorithm is correct?
  - Explain your solution to someone.
  - Are you using any functions that you don't completely understand? Maybe STL functions?
  - Maybe you (or someone else) should rewrite the solution?
  - Can the input line be empty?
- Run-Time Error?
  - Is it actually Memory Limit Exceeded?

#### 2.2. Solution Ideas.

- Dynamic Programming
  - Parsing CFGs: CYK Algorithm
  - Drop a parameter, recover from others
  - Swap answer and a parameter
  - When grouping: try splitting in two
  - $-2^k$  trick
  - When optimizing
    - \* Convex hull optimization
      - $\cdot \operatorname{dp}[i] = \min_{i < i} \{\operatorname{dp}[j] + b[j] \times a[i]\}$
      - $b[j] \geq b[j+1]$
      - · optionally  $a[i] \leq a[i+1]$
      - $O(n^2)$  to O(n)
    - \* Divide and conquer optimization
      - $dp[i][j] = \min_{k < j} \{dp[i-1][k] + C[k][j]\}$
      - $A[i][j] \leq A[i][j+1]$
      - ·  $O(kn^2)$  to  $O(kn\log n)$
      - · sufficient:  $C[a][c] + C[b][d] \le C[a][d] + C[b][c], a \le$  $b \le c \le d \text{ (QI)}$
    - \* Knuth optimization
      - $dp[i][j] = \min_{i < k < j} \{dp[i][k] + dp[k][j] + C[i][j]\}$
      - $A[i][j-1] \le A[i][j] \le A[i+1][j]$
      - $O(n^3)$  to  $O(n^2)$
      - · sufficient: QI and C[b][c] < C[a][d], a < b < c < d

- Randomized
- Optimizations
  - Use bitset (/64)
  - Switch order of loops (cache locality)
- Process queries offline
  - Mo's algorithm
- Square-root decomposition
- Precomputation
- Efficient simulation
  - Mo's algorithm
  - Sqrt decomposition
  - Store  $2^k$  jump pointers
- Data structure techniques
  - Sqrt buckets
  - Store  $2^k$  jump pointers
  - $-2^k$  merging trick
- Counting
  - Inclusion-exclusion principle
  - Generating functions
- Graphs
  - Can we model the problem as a graph?
  - Can we use any properties of the graph?
  - Strongly connected components
  - Cycles (or odd cycles)
  - Bipartite (no odd cycles)
    - \* Bipartite matching
    - \* Hall's marriage theorem
    - \* Stable Marriage
  - Cut vertex/bridge
  - Biconnected components
  - Degrees of vertices (odd/even)
  - Trees
    - \* Heavy-light decomposition
    - \* Centroid decomposition
    - \* Least common ancestor
    - \* Centers of the tree
  - Eulerian path/circuit
  - Chinese postman problem
  - Topological sort
  - (Min-Cost) Max Flow
  - Min Cut
    - \* Maximum Density Subgraph
  - Huffman Coding
  - Min-Cost Arborescence

  - Steiner Tree
  - Kirchoff's matrix tree theorem
  - Prüfer sequences
  - Lovász Toggle
  - Look at the DFS tree (which has no cross-edges)
  - Is the graph a DFA or NFA?
    - \* Is it the Synchronizing word problem?
- Mathematics
  - Is the function multiplicative?
  - Look for a pattern
  - Permutations
    - \* Consider the cycles of the permutation

- Functions
  - \* Sum of piecewise-linear functions is a piecewise-linear
  - \* Sum of convex (concave) functions is convex (concave)
- Modular arithmetic
  - \* Chinese Remainder Theorem
  - \* Linear Congruence
- Sieve
- System of linear equations
- Values too big to represent?
  - \* Compute using the logarithm
  - \* Divide everything by some large value
- Linear programming
  - \* Is the dual problem easier to solve?
- Can the problem be modeled as a different combinatorial problem? Does that simplify calculations?
- Logic
  - 2-SAT
  - XOR-SAT (Gauss elimination or Bipartite matching)
- Meet in the middle
- Only work with the smaller half  $(\log(n))$
- Strings
  - Trie (maybe over something weird, like bits)
  - Suffix array
  - Suffix automaton (+DP?)
  - Aho-Corasick
  - eerTree
  - Work with S + S
- Hashing
- Euler tour, tree to array
- Segment trees
  - Lazy propagation
  - Persistent
  - Implicit
  - Segment tree of X
- Geometry
  - Minkowski sum (of convex sets)
  - Rotating calibers
  - Sweep line (horizontally or vertically?)
  - Sweep angle
  - Convex hull
- Fix a parameter (possibly the answer)
- Are there few distinct values?
- Binary search
- Sliding Window (+ Monotonic Queue)
- Computing a Convolution? Fast Fourier Transform • Computing a 2D Convolution? FFT on each row, and then on each column
- Exact Cover (+ Algorithm X)
- Cycle-Finding
- What is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?
- Look at the complement problem
  - Minimize something instead of maximizing

• Greedy

- $\bullet$  Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/Bucket sort

## PRACTICE CONTEST CHECKLIST

- How many operations per second? Compare to local machine.
- What is the stack size?
- How to use printf/scanf with long long/long double?
- Are \_\_int128 and \_\_float128 available?
- Does MLE give RTE or MLE as a verdict? What about stack overflow?
- What is RAND\_MAX?
- How does the judge handle extra spaces (or missing newlines) in the output?
- Look at documentation for programming languages.
- Try different programming languages: C++, Java and Python.
- Try the submit script.
- Try local programs: i?python[23], factor.
- Try submitting with assert(false) and assert(true).
- Return-value from main.
- Look for directory with sample test cases.
- Make sure printing works.
- Remove this page from the notebook.