Article: Hybrid Computing Using a Neural Network with Dynamic External Memory

Speaker:

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Prepare to me amazed...

- 1) What's the idea?
- 2) Does it even work?
- 3) Details (you'll need them!)

Neural Networks

Store weights

Learn from data

Learn some distributions

Weak with algorithms & data structures

What about... me?



Neural Networks

Computers

Store weights

Learn from data

Learn some distributions

Weak with algorithms & data structures

Processor and RAM

Passively store data

All data stored equally

Can process data structures

Differentiable Neural Computers

Store weights

Learn from data

Learn some distributions

Weak with algorithms & data structures

Processor and RAM Controller and Memory matrix

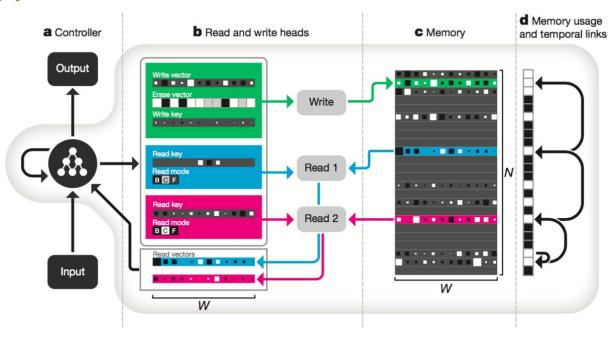
Passively store data

All data stored equally Data is stored with some distribution

Can process data structures

#NeuralTuringMachine

How?



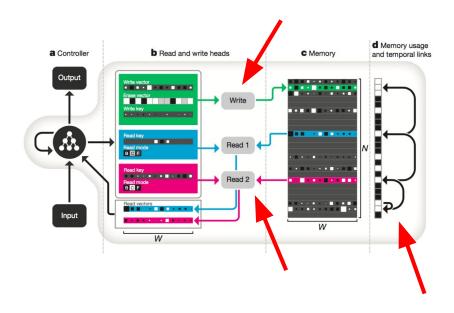
What's the idea?

Differentiable everything!

Yes, and memory as well....

3 forms of <differentiable> attention

- Content Lookup
- Temporal Linking
- Memory Allocation



*Hippocampus *CA3_synapses

3 forms of <differentiable> attention

Attention mechanism

Content Lookup

Temporal Linking

Memory Allocation

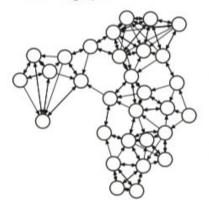
Computational considerations

Formation of associative data structures

Sequential retrieval of input sequences

Provides the write head with unused locations

a Random graph



b London Underground



LSTM: 37% accuracy after 2 million examples

Traversal

Shortest-path

Underground input:

(OxfordCircus, TottenhamCtRd, Central) (TottenhamCtRd, OxfordCircus, Central) (BakerSt, Marylebone, Circle) (BakerSt, Marylebone, Bakerloo) (BakerSt, OxfordCircus, Bakerloo)

(LeicesterSq, CharingCross, Northern) (TottenhamCtRd, LeicesterSq, Northern) (OxfordCircus, PiccadillyCircus, Bakerloo) (OxfordCircus, NottingHillGate, Central) (OxfordCircus, Euston, Victoria)

84 edges in total

Traversal question:

(BondSt, _, Central), (_, _, Circle), (_, _, Circle), (_, _, Circle), (_, _, Circle), (_, _, Jubilee), (_, _, Jubilee),

Answer:

(BondSt, NottingHillGate, Central) (NottingHillGate, GloucesterRd, Circle)

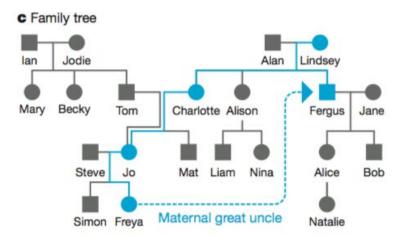
(Westminster, GreenPark, Jubilee) (GreenPark, BondSt, Jubilee)

Shortest-path question: (Moorgate, PiccadillyCircus, _)

Answer:

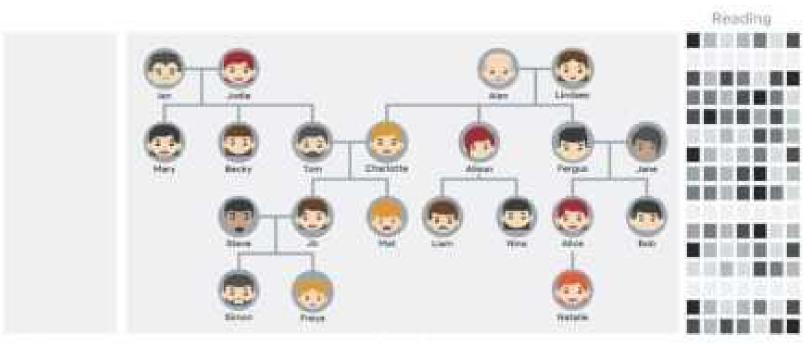
(Moorgate, Bank, Northern)
(Bank, Holborn, Central)
(Holborn, LeicesterSq, Piccadilly)
(LeicesterSq, PiccadillyCircus, Piccadilly)

DNC: 98.8% accuracy after 1 million examples



Family tree input: Inference question: (Charlotte, Alan, Father) (Freya, _, MaternalGreatUncle) (Simon, Steve, Father) (Steve, Simon, Son1) (Nina, Alison, Mother) (Lindsey, Fergus, Son1) (Bob, Jane, Mother) Answer: (Natalie, Alice, Mother) (Freya, Fergus, MaternalGreatUncle) (Mary, Ian, Father) (Jane, Alice, Daughter1) (Mat, Charlotte, Mother) 54 edges in total

DNC: 81.8% accuracy on four-step relations



Input Graph Memory

Sheep are afraid of wolves.

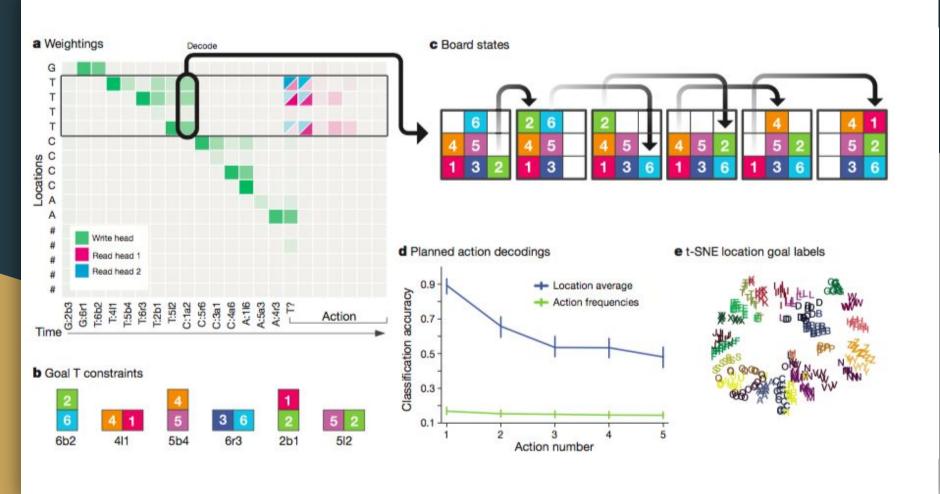
Gertrude is a sheep.

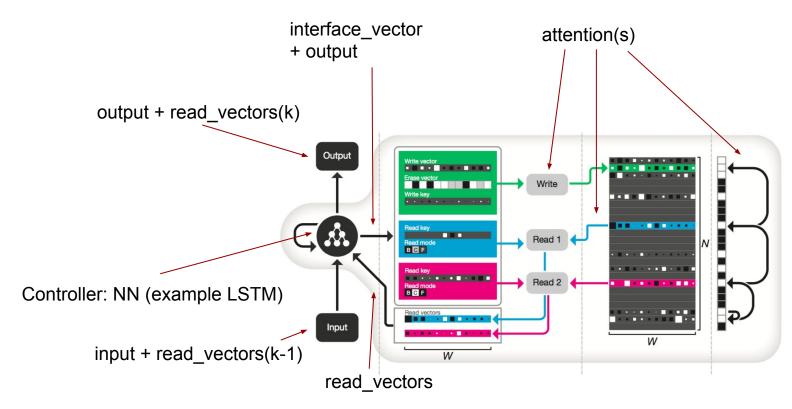
Mice are afraid of cats.

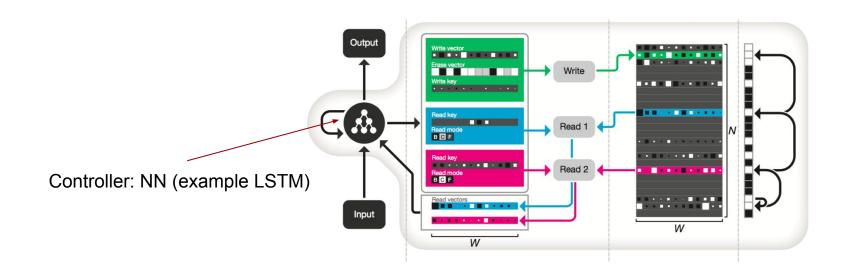
LSTM: 7.5% mean test error

What is Gertrude afraid of?

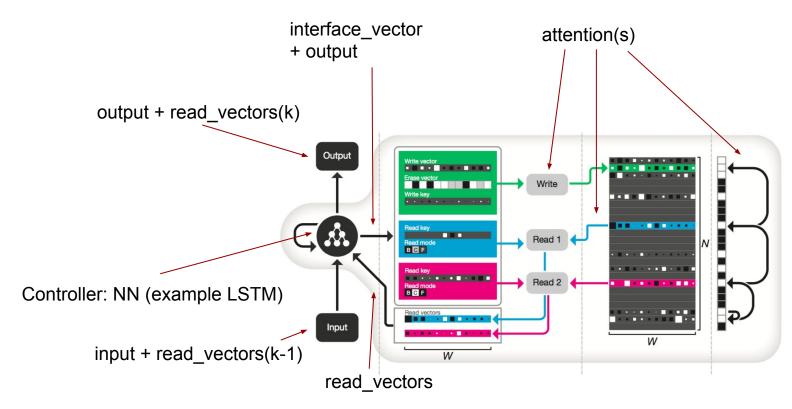
DNC: 3.8% mean test error







input
$$\begin{aligned} & \boldsymbol{i}_t^l = \sigma(\boldsymbol{W}_{\boldsymbol{i}}^l[\boldsymbol{\chi}_t; \boldsymbol{h}_{t-1}^l; \boldsymbol{h}_t^{l-1}] + \boldsymbol{b}_{\boldsymbol{i}}^l) \\ & \text{forget} & \boldsymbol{f}_t^l = \sigma(\boldsymbol{W}_{\boldsymbol{f}}^l[\boldsymbol{\chi}_t; \boldsymbol{h}_{t-1}^l; \boldsymbol{h}_t^{l-1}] + \boldsymbol{b}_{\boldsymbol{f}}^l) \\ & \text{state} & \boldsymbol{s}_t^l = \boldsymbol{f}_t^l \boldsymbol{s}_{t-1}^l + \boldsymbol{i}_t^l \mathrm{tanh}(\boldsymbol{W}_{\boldsymbol{s}}^l[\boldsymbol{\chi}_t; \boldsymbol{h}_{t-1}^l; \boldsymbol{h}_t^{l-1}] + \boldsymbol{b}_{\boldsymbol{s}}^l) \\ & \text{output gate activation} & \boldsymbol{o}_t^l = \sigma(\boldsymbol{W}_{\boldsymbol{o}}^l[\boldsymbol{\chi}_t; \boldsymbol{h}_{t-1}^l; \boldsymbol{h}_t^{l-1}] + \boldsymbol{b}_{\boldsymbol{o}}^l) \\ & \text{hidden} & \boldsymbol{h}_t^l = \boldsymbol{o}_t^l \mathrm{tanh}(\boldsymbol{s}_t^l) \end{aligned}$$



interface_vector + output Output Write vector Erase vector Write key Write Read 1 Read mode Read key Read mode B C F Read 2 Read vectors Input W

interface_vector

Interface parameters. Before being used to parameterize the memory interactions, the interface vector ξ_t is subdivided as follows:

$$\boldsymbol{\xi}_{t} = \left[\boldsymbol{k}_{t}^{\mathrm{r},1}; \ldots; \boldsymbol{k}_{t}^{\mathrm{r},R}; \hat{\boldsymbol{\beta}}_{t}^{\mathrm{r},1}; \ldots; \hat{\boldsymbol{\beta}}_{t}^{\mathrm{r},R}; \boldsymbol{k}_{t}^{\mathrm{w}}; \hat{\boldsymbol{\beta}}_{t}^{\mathrm{w}}; \hat{\boldsymbol{e}}_{t}; \boldsymbol{v}_{t}; \hat{\boldsymbol{f}}_{t}^{1}; \ldots; \hat{\boldsymbol{f}}_{t}^{R}; \hat{\boldsymbol{g}}_{t}^{\mathrm{a}}; \hat{\boldsymbol{g}}_{t}^{\mathrm{w}}; \hat{\boldsymbol{\pi}}_{t}^{1}; \ldots; \hat{\boldsymbol{\pi}}_{t}^{R}\right]$$

- the write key $\mathbf{k}_t^{\mathrm{w}} \in \mathbb{R}^W$;
- the write strength $\beta_t^{\text{w}} = \text{oneplus}(\hat{\beta}_t^{\text{w}}) \in [1,\infty)$;
- the erase vector $\mathbf{e}_t = \sigma(\hat{\mathbf{e}}_t) \in [0,1]^W$;
- the write vector $\mathbf{v}_t \in \mathbb{R}^W$;

- R free gates $\{f_t^i = \sigma(\hat{f}_t^i) \in [0,1]; 1 \le i \le R\};$
- the allocation gate $g_t^a = \sigma(\hat{g}_t^a) \in [0,1]$;
- the write gate $g_t^w = \sigma(\hat{g}_t^w) \in [0,1]$; and
- $R \text{ read modes } \{\pi_t^i = \operatorname{softmax}(\hat{\pi}_t^i) \in S_3; 1 \le i \le R\}.$

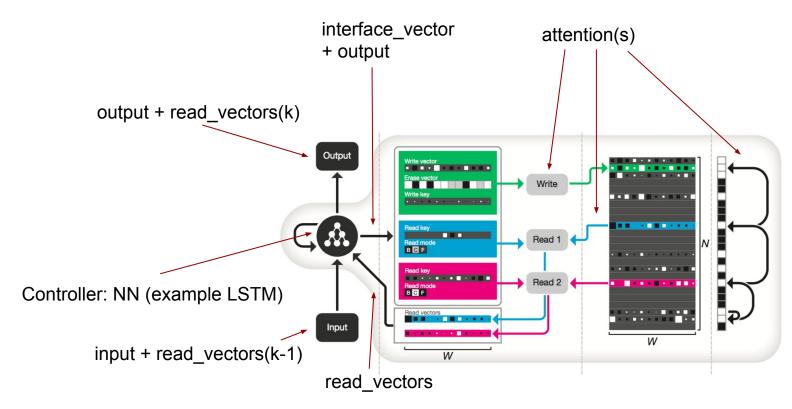
interface_vector

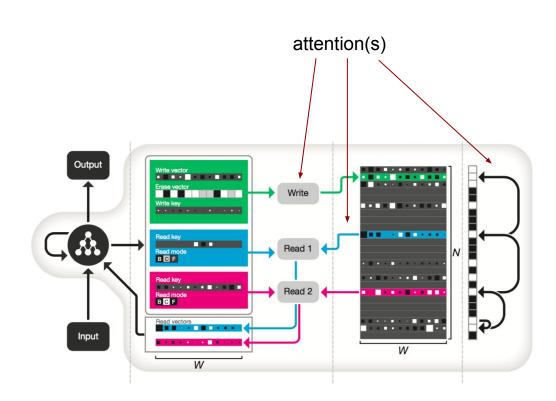
Interface parameters. Before being used to parameterize the memory interactions, the interface vector ξ_t is subdivided as follows:

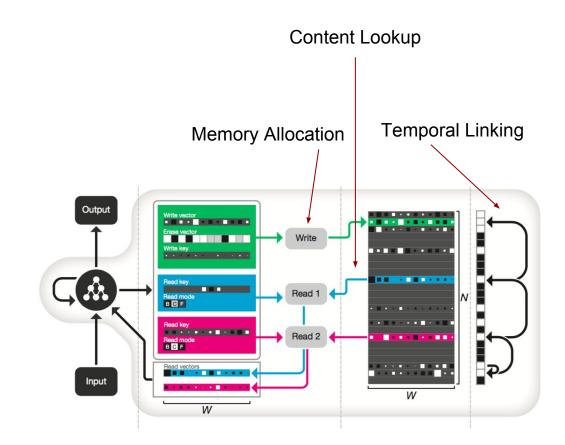
$$\boldsymbol{\xi}_{t} = \left[\boldsymbol{k}_{t}^{\mathrm{r},1}; \ldots; \boldsymbol{k}_{t}^{\mathrm{r},R}; \hat{\boldsymbol{\beta}}_{t}^{\mathrm{r},1}; \ldots; \hat{\boldsymbol{\beta}}_{t}^{\mathrm{r},R}; \boldsymbol{k}_{t}^{\mathrm{w}}; \hat{\boldsymbol{\beta}}_{t}^{\mathrm{w}}; \hat{\boldsymbol{e}}_{t}; \boldsymbol{v}_{t}; \hat{\boldsymbol{f}}_{t}^{1}; \ldots; \hat{\boldsymbol{f}}_{t}^{R}; \hat{\boldsymbol{g}}_{t}^{\mathrm{a}}; \hat{\boldsymbol{g}}_{t}^{\mathrm{w}}; \hat{\boldsymbol{\pi}}_{t}^{1}; \ldots; \hat{\boldsymbol{\pi}}_{t}^{R}\right]$$

$$\boldsymbol{v}_t = W_{\boldsymbol{y}}[\boldsymbol{h}_t^1; \dots; \boldsymbol{h}_t^L]$$

$$\boldsymbol{\xi}_t = W_{\boldsymbol{\xi}}[\boldsymbol{h}_t^1; \dots; \boldsymbol{h}_t^L]$$







Content lookup (attention 1)

$$C(M, \mathbf{k}, \beta)[i] = \frac{\exp\{D(\mathbf{k}, M[i, \cdot])\beta\}}{\sum_{j} \exp\{D(\mathbf{k}, M[j, \cdot])\beta\}}$$

where $k \in \mathbb{R}^W$ is a lookup key, $\beta \in [1, \infty)$ is a scalar representing key strength and \mathcal{D} is the cosine similarity:

$$\mathcal{D}(u,v) = \frac{u \cdot v}{|u||v|}$$

Memory allocation (attention 2)

$$\psi_t = \prod_{i=1}^R \left(\mathbf{1} - f_t^i \mathbf{w}_{t-1}^{r,i} \right)$$

$$\mathbf{u}_t = (\mathbf{u}_{t-1} + \mathbf{w}_{t-1}^{w} - \mathbf{u}_{t-1} \circ \mathbf{w}_{t-1}^{w}) \circ \psi_t$$

$$\mathbf{a}_t[\phi_t[j]] = (1 - \mathbf{u}_t[\phi_t[j]]) \prod_{i=1}^{j-1} \mathbf{u}_t[\phi_t[i]]$$

Memory allocation (attention 2) Free gates (from interface vector) Retention vector $\boldsymbol{\psi}_{t} = \prod \left(\mathbf{1} - f_{t}^{i} \mathbf{w}_{t-1}^{r,i} \right)$ Weights (learnable) Element-wise product Memory usage vector $\mathbf{u}_{t} = (\mathbf{u}_{t-1} + \mathbf{w}_{t-1}^{\mathrm{w}} - \mathbf{u}_{t-1} \circ \mathbf{w}_{t-1}^{\mathrm{w}}) \circ \mathbf{\psi}_{t}$ Free list $\phi_t[1]$ is the index of the least used location $\mathcal{A}_t[\dot{\boldsymbol{\phi}}_t[j]] = (1 - \boldsymbol{u}_t[\boldsymbol{\phi}_t[j]]) \prod^t \boldsymbol{u}_t[\boldsymbol{\phi}_t[i]]$ Allocation weighting

Write weighting. The controller can write to newly allocated locations, or to locations addressed by content, or it can choose not to write at all. First, a write content weighting $c_t^w \in S_N$ is constructed using the write key k_t^w and write strength β_t^w :

$$\boldsymbol{c}_{t}^{\mathrm{w}} = \mathcal{C}(M_{t-1}, \boldsymbol{k}_{t}^{\mathrm{w}}, \boldsymbol{\beta}_{t}^{\mathrm{w}})$$

 c_t^{w} is interpolated with the allocation weighting a_t defined in equation (1) to determine a write weighting $w_t^{\text{w}} \in \Delta_N$:

$$\boldsymbol{w}_{t}^{W} = \boldsymbol{g}_{t}^{W} \left[\boldsymbol{g}_{t}^{a} \boldsymbol{a}_{t} + (1 - \boldsymbol{g}_{t}^{a}) \boldsymbol{c}_{t}^{W} \right]$$
 (2)

where $g_t^a \in [0,1]$ is the allocation gate governing the interpolation and $g_t^w \in [0,1]$ is the write gate. If the write gate is 0, then nothing is written, regardless of the other write parameters; it can therefore be used to protect the memory from unnecessary modifications.

Temporal linking (attention 3)

$$\mathbf{p}_0 = \mathbf{0}$$

$$\mathbf{p}_t = \left(1 - \sum_i \mathbf{w}_t^{\mathrm{w}}[i]\right) \mathbf{p}_{t-1} + \mathbf{w}_t^{\mathrm{w}}$$
Precedence weights (writing order)

$$L_0[i,j]{=}0 \quad orall \ i,j$$
 — Temporal linking matrix $L_t[i,i]{=}0 \quad orall \ i$

$$L_t[i,j] = (1 - \mathbf{w}_t^{w}[i] - \mathbf{w}_t^{w}[j])L_{t-1}[i,j] + \mathbf{w}_t^{w}[i]\mathbf{p}_{t-1}[j]$$

Read weighting. Each read head *i* computes a content weighting $\mathbf{c}_t^{\mathbf{r},i} \in \Delta_N$ using a read key $\mathbf{k}_t^{\mathbf{r},i} \in \mathbb{R}^W$:

$$\boldsymbol{c}_t^{\mathrm{r},i} = \mathcal{C}(M_t, \boldsymbol{k}_t^{\mathrm{r},i}, \boldsymbol{\beta}_t^{\mathrm{r},i})$$

Each read head also receives a read mode vector $\boldsymbol{\pi}_t^i \in \mathcal{S}_3$, which interpolates among the backward weighting \boldsymbol{b}_t^i , the forward weighting \boldsymbol{f}_t^i and the content read weighting $\boldsymbol{c}_t^{r,i}$, thereby determining the read weighting $\boldsymbol{w}_t^{r,i} \in \mathcal{S}_3$:

$$\mathbf{w}_{t}^{\mathrm{r},i} = \mathbf{\pi}_{t}^{i}[1]\mathbf{b}_{t}^{i} + \mathbf{\pi}_{t}^{i}[2]\mathbf{c}_{t}^{\mathrm{r},i} + \mathbf{\pi}_{t}^{i}[3]\mathbf{f}_{t}^{i}$$

If $\pi_t^i[2]$ dominates the read mode, then the weighting reverts to content lookup using $k_t^{r,i}$. If $\pi_t^i[3]$ dominates, then the read head iterates through memory locations in the order they were written, ignoring the read key. If $\pi_t^i[1]$ dominates, then the read head iterates in the reverse order.

