



HA NOI UNIVERSITY OF SCIENCE AND TECHNOLOGY
SCHOOL OF INFORMATION AND COMMUNICATION TECHNOLOGY

C Programming Basic

Sorting - part 2

CONTENT

- Data generation
- Implement merge sort, quick sort, heap sort algorithms
- Experiments

Exercise 1: data generation

- Profile of each staff has following information
 - Full name
 - Date of birth
- Write a program that generates n profiles randomly and store to a file profile-n.txt under the format:
 - Line $2i-1$ and $2i$ ($i = 1, \dots, n$) respectively write the full name and date of birth of the profile i . Full name has the format <last_name> <middle_name> <first_name> and the date of birth has the format YYYY-MM-DD
 - Line $2n+1$: write the character # to indicate the end of the data

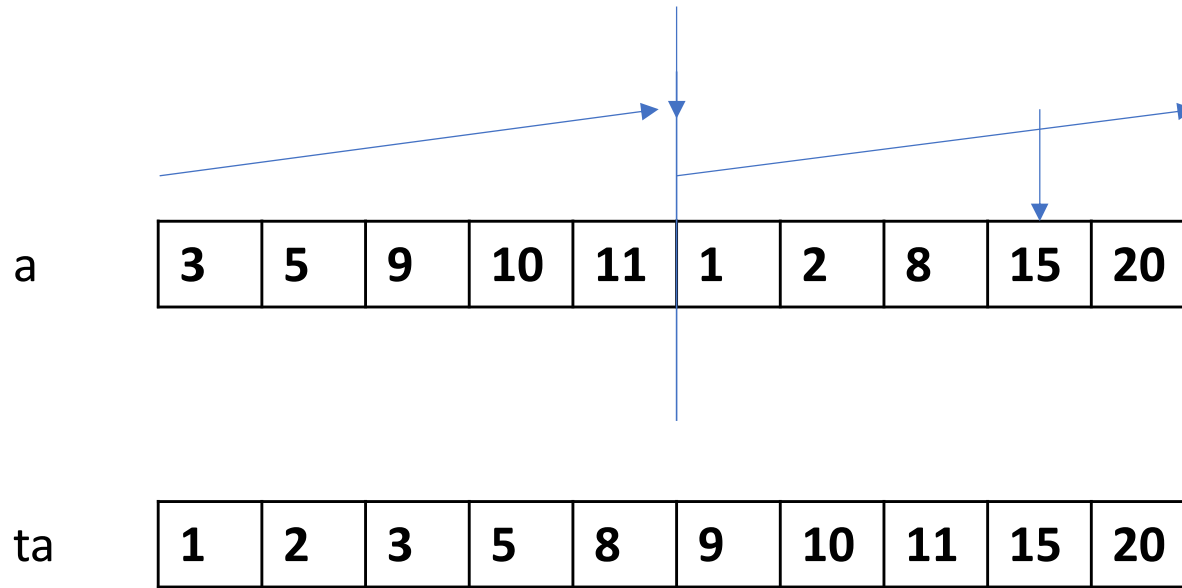
Profile-5.txt
Bui Hai An
1980-02-30
Pham Viet Anh
1986-10-08
Do Duc Bang
1990-04-24
Dang Van Cuong
1987-08-17
Pham Viet Anh
1986-05-20
#

Merge sort

- Problem: Sorting n elements of an array $S[0] \dots S[n-1]$
- Merge-sort on an input sequence \mathbf{S} with n elements consists of three steps:
 - **Divide**: partition \mathbf{S} into two sequences \mathbf{S}_1 and \mathbf{S}_2 of about $n/2$ elements each
 - **Conquer**: recursively sort \mathbf{S}_1 and \mathbf{S}_2 by using merge sort
 - **Combine**: merge \mathbf{S}_1 and \mathbf{S}_2 into a unique sorted sequence

```
void mergeSort(int A[], int L, int R) {  
    if(L < R){  
        int M = (L+R)/2;  
        mergeSort(A,L,M);  
        mergeSort(A,M+1,R);  
        merge(A,L,M,R);  
    }  
}
```

Merge sort



Merge sort

- Use auxiliary array
- Running time
 - Worst case $O(n \log n)$
 - Best case: $O(n \log n)$

```
void merge(int A[], int L, int M, int R) {  
    // tron 2 day da sap A[L..M] va A[M+1..R]  
    int i = L; int j = M+1;  
    for(int k = L; k <= R; k++){  
        if(i > M){ TA[k] = A[j]; j++;}  
        else if(j > R){TA[k] = A[i]; i++;}  
        else{  
            if(A[i] < A[j]){  
                TA[k] = A[i]; i++;  
            }  
            else {  
                TA[k] = A[j]; j++;  
            }  
        }  
    }  
    for(int k = L; k <= R; k++) A[k] = TA[k];  
}
```

Exercise 2

- Write a program that reads a sequence of n profiles from data files generated from the exercise 1 above, sort the sequence in non-decreasing order (prioritize full name, then date of birth) by the merge sort algorithm
- Input (profile-n.txt)
 - Line $2i-1$ and $2i$ ($i = 1, \dots, n$) respectively write the fullname and date of birth of the profile i . Fullname has the format <last_name> <middle_name> <first_name> and the date of birth has the format YYYY-MM-DD
 - Last line: write the character # to indicate the end of the data
- Result (sorted-profile-n.txt)
 - Write the sorted sequence under the format:
 - Line $2i-1$ and $2i$ ($i = 1, \dots, n$) respectively write the fullname and date of birth of the profile i . Fullname has the format <last_name> <middle_name> <first_name> and the date of birth has the format YYYY-MM-DD
 - Last line: write the character # to indicate the end of the data

Quick sort

The quick sort algorithm is described recursively as following (similar to merge sort):

1. Base case. If the array has only one element, then the array is sorted already, return it without doing anything.

2. Divide:

- Select an element in the array, and call it as the pivot ***p***.
- Divide the array into 2 subarrays: Left subarray (*L*) consists of elements \leq the pivot, right subarray (*R*) consists of elements \geq the pivot. This operation is called “**Partition**”.

3. Conquer: recursively call QuickSort for 2 subarrays $L = A[p \dots q]$ and $R = A[q+1 \dots r]$.

4. Combine: The sorted array is ***L p R***.

In contrast to Merge Sort, in Quick Sort: division operation is complicate, but the Partition operation is simple.

Quick sort

```
void quickSort(int A[], int L, int R) {  
    if(L < R){  
        int index = (L + R)/2;  
        index = partition(A, L, R, index);  
        if(L < index)  
            quickSort(A, L, index-1);  
        if(index < R)  
            quickSort(A, index+1, R);  
    }  
}
```

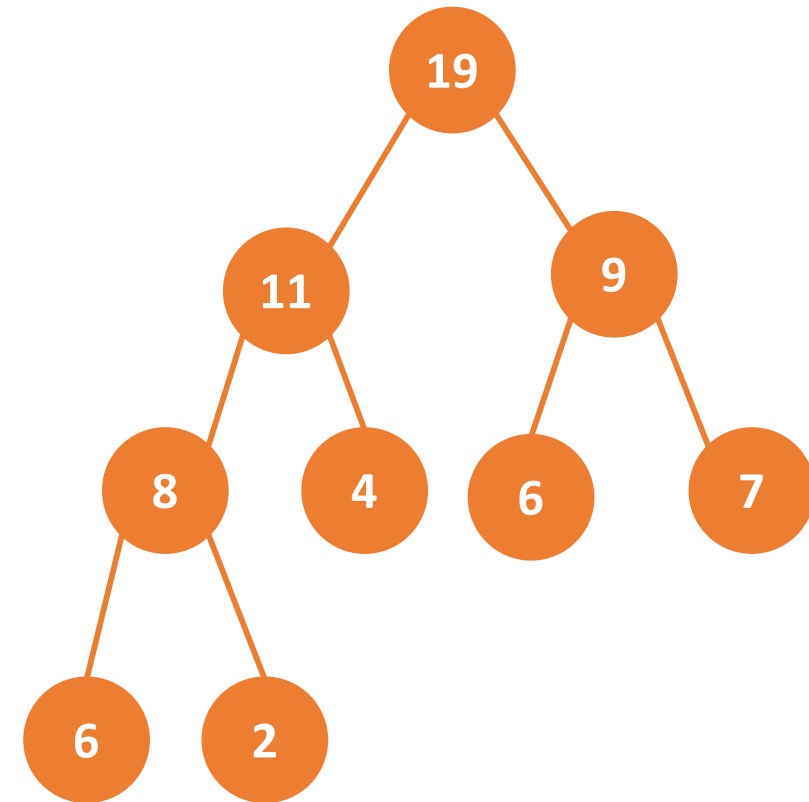
```
int partition(int A[], int L, int R, int  
            indexPivot) {  
    int pivot = A[indexPivot];  
    swap(A[indexPivot], A[R]);  
    int storeIndex = L;  
    for(int i = L; i <= R-1; i++){  
        if(A[i] < pivot){  
            swap(A[storeIndex], A[i]);  
            storeIndex++;  
        }  
    }  
    swap(A[storeIndex], A[R]);  
    return storeIndex;  
}
```

Exercise 3

- Write a program that reads a sequence of n profiles from data files generated from the exercise 1 above, sort the sequence in non-decreasing order (prioritize full name, then date of birth) by the quick sort algorithm
- Input (profile-n.txt)
 - Line $2i-1$ and $2i$ ($i = 1, \dots, n$) respectively write the fullname and date of birth of the profile i . Fullname has the format <last_name> <middle_name> <first_name> and the date of birth has the format YYYY-MM-DD
 - Last line: write the character # to indicate the end of the data
- Result (sorted-profile-n.txt)
 - Write the sorted sequence under the format:
 - Line $2i-1$ and $2i$ ($i = 1, \dots, n$) respectively write the fullname and date of birth of the profile i . Fullname has the format <last_name> <middle_name> <first_name> and the date of birth has the format YYYY-MM-DD
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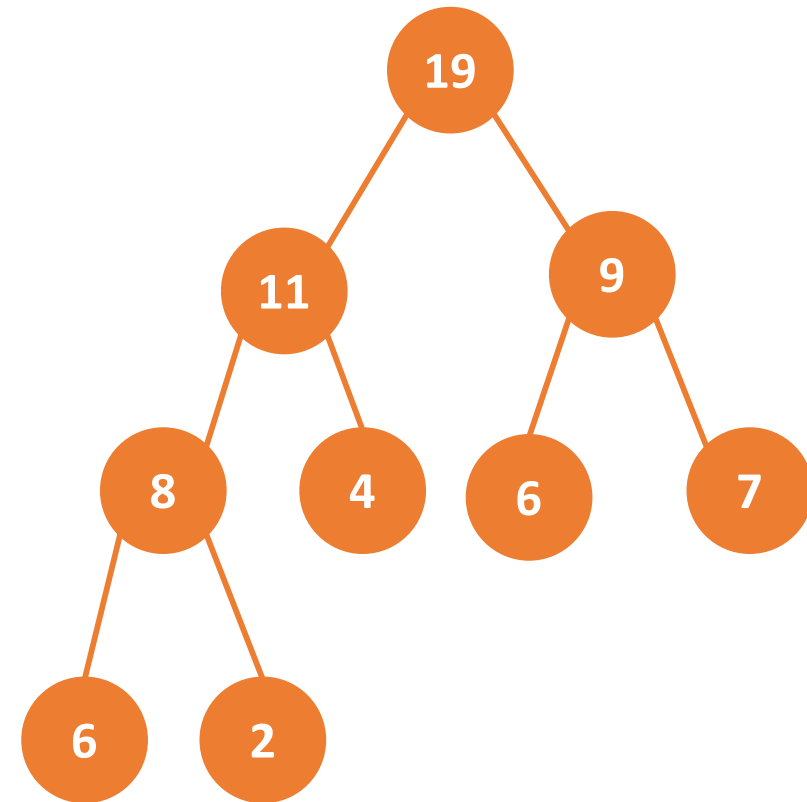
Heap sort

- Heap structure (max-heap)
 - Complete tree
 - Key of each node is greater or equal to the keys of its children (max-heap property)
- Map the sequence $A[1 \dots N]$ to a complete tree
 - Root is $A[1]$
 - $A[2i]$ and $A[2i+1]$ are respectively the left child and the right child of $A[i]$
 - The height of the tree is $\log N + 1$



Heap sort

- Heapify
 - Status:
 - max-heap property at $A[i]$ is destroyed
 - max-heap property at children of $A[i]$ is satisfied
 - Adjust the tree to recover the max-heap property at the root $A[i]$



Heap sort

- Running time: $O(\log N)$

```
void heapify(int A[], int i, int N) {  
    int L = 2*i;  
    int R = 2*i+1;  
    int max = i;  
    if(L <= N && A[L] > A[i])  
        max = L;  
    if(R <= N && A[R] > A[max])  
        max = R;  
    if(max != i){  
        swap(A[i], A[max]);  
        heapify(A,max,N);  
    }  
}
```

Heap sort

- Heap sort
 - Build max-heap
 - Swap $A[1]$ and $A[N]$
 - Heapify at $A[1]$ until $A[1..N-1]$
 - Swap $A[1]$ and $A[N-1]$
 - Heapify at $A[1]$ until $A[1..N-2]$
 - ...
- Run time: $O(N \log N)$

```
void buildHeap(int A[], int N) {
    for(int i = N/2; i >= 1; i--)
        heapify(A,i,N);
}

void heapSort(int A[], int N) {
    // index tu 1 -> N
    buildHeap(A,N);
    for(int i = N; i > 1; i--) {
        swap(A[1], A[i]);
        heapify(A, 1, i-1);
    }
}
```

Exercise 4

- Write a program that reads a sequence of n profiles from data files generated from the exercise 1 above, sort the sequence in non-decreasing order (prioritize full name, then date of birth) by the heap sort algorithm
- Input (profile-n.txt)
 - Line $2i-1$ and $2i$ ($i = 1, \dots, n$) respectively write the fullname and date of birth of the profile i . Fullname has the format <last_name> <middle_name> <first_name> and the date of birth has the format YYYY-MM-DD
 - Last line: write the character # to indicate the end of the data
- Result (sorted-profile-n.txt)
 - Write the sorted sequence under the format:
 - Line $2i-1$ and $2i$ ($i = 1, \dots, n$) respectively write the fullname and date of birth of the profile i . Fullname has the format <last_name> <middle_name> <first_name> and the date of birth has the format YYYY-MM-DD
 - Last line: write the character # to indicate the end of the data