SWE3004 Operating Systems, Spring 2023

Project 3. Virtual memory

TA)
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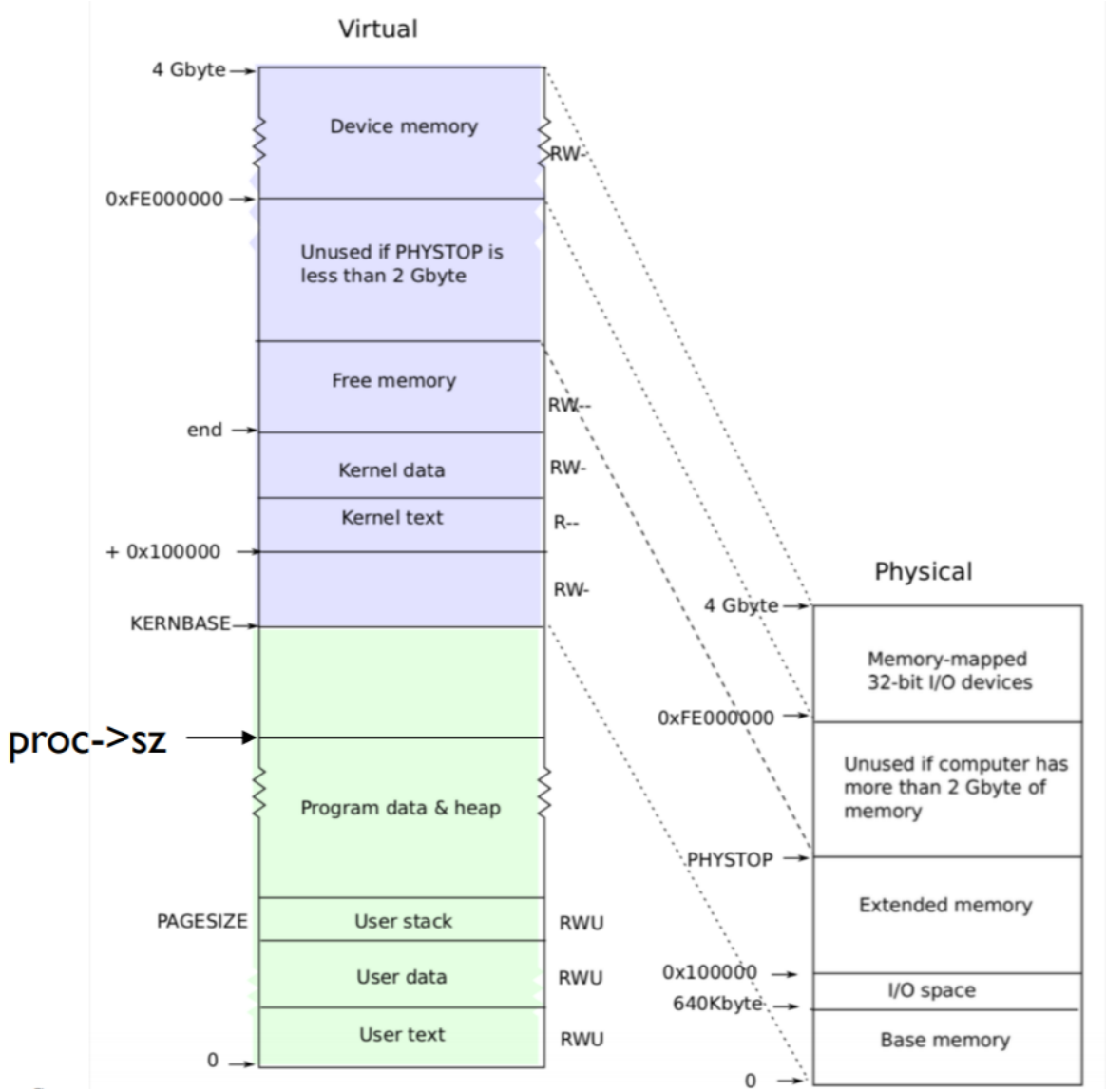
Project plan

- Total 6 projects
  0) ~~Booting xv6 operating system (10pt)~~
  1) ~~System call (15pt)~~
  2) ~~CPU scheduling (25pt)~~
  3) Virtual memory (25pt)
  4) Page replacement (25pt)
  5) File systems (0pt, optional)

xv6 Memory Layout



How Physical Memories Initialized in xv6

- main() of main.c

```
int
main(void)
{
  kinit1(end, P2V(4*1024*1024)); // phys page allocator
  kvmalloc();      // kernel page table
  mpinit();        // detect other processors
  lapicinit();     // interrupt controller
  seginit();       // segment descriptors
  picinit();       // disable pic
  ioapicinit();    // another interrupt controller
  consoleinit();   // console hardware
  uartinit();      // serial port
  pinit();         // process table
  tvinit();        // trap vectors
  binit();         // buffer cache
  fileinit();      // file table
  ideinit();       // disk
  startothers();   // start other processors
  kinit2(P2V(4*1024*1024), P2V(PHYSTOP)); // must come after startothers()
  userinit();      // first user process
  mpmain();        // finish this processor's setup
}
```

These two functions divide & manage physical memories with pages

How Physical Memories Initialized in xv6

```
void
kinit1(void *vstart, void *vend)
{
  initlock(&kmem.lock, "kmem");
  kmem.use_lock = 0;
  freerange(vstart, vend);
}

void
kinit2(void *vstart, void *vend)
{
  freerange(vstart, vend);
  kmem.use_lock = 1;
}

void
freerange(void *vstart, void *vend)
{
  char *p;
  p = (char*)PGROUNDUP((uint)vstart);
  for(; p + PGSIZE <= (char*)vend; p += PGSIZE)
    kfree(p);
}
```

```
void
kfree(char *v)
{
  struct run *r;

  if((uint)v % PGSIZE || v < end || V2P(v) >= PHYSTOP)
    panic("kfree");

  // Fill with junk to catch dangling refs.
  memset(v, 1, PGSIZE);

  if(kmem.use_lock)
    acquire(&kmem.lock);
  r = (struct run*)v;
  r->next = kmem.freelist;
  kmem.freelist = r;
  if(kmem.use_lock)
    release(&kmem.lock);
}
```

- kinit1() sets up for lock-less allocation in the first 4MB
- kinit2() arranges for more memory (until PHYSTOP) to be allocatable (224MB)

- freerange() kfree() with page size unit

- kfree() fills page with 1s, and put it into freelist (page pool)

How Physical Memories Initialized in xv6

- fork() creates a child with exactly the same memory contents as the parent

```
int
fork(void)
{
  int i, pid;
  struct proc *np;
  struct proc *curproc = myproc();

  // Allocate process.
  if((np = allocproc()) == 0){
    return -1;
  }

  // Copy process state from proc.
  if((np->pgdir = copyuvm(curproc->pgdir, curproc->sz)) == 0){
    kfree(np->kstack);
    np->kstack = 0;
    np->state = UNUSED;
    return -1;
  }
  np->sz = curproc->sz;
  np->parent = curproc;
  *np->tf = *curproc->tf;

  // Clear %eax so that fork returns 0 in the child.
  np->tf->eax = 0;

  for(i = 0; i < NOFILE; i++)
    if(curproc->ofile[i])
      np->ofile[i] = filedup(curproc->ofile[i]);
  np->cwd = idup(curproc->cwd);

  safestrcpy(np->name, curproc->name, sizeof(curproc->name));

  pid = np->pid;

  acquire(&ptable.lock);

  np->state = RUNNABLE;

  release(&ptable.lock);

  return pid;
}
```

- allocproc() allocates kernel stack
- copyuvm() copys parent's page table

```
pde_t*
copyuvm(pde_t *pgdir, uint sz)
{
  pde_t *d;
  pte_t *pte;
  uint pa, i, flags;
  char *mem;

  if((d = setupkvm()) == 0)
    return 0;
  for(i = 0; i < sz; i += PGSIZE){
    if((pte = walkpgdir(pgdir, (void *) i, 0)) == 0)
      panic("copyuvm: pte should exist");
    if(!(*pte & PTE_P))
      panic("copyuvm: page not present");
    pa = PTE_ADDR(*pte);
    flags = PTE_FLAGS(*pte);
    if((mem = kalloc()) == 0)
      goto bad;
    memmove(mem, (char*)P2V(pa), PGSIZE);
    if(mappages(d, (void*)i, PGSIZE, V2P(mem), flags) < 0) {
      kfree(mem);
      goto bad;
    }
  }
  return d;

bad:
  freevm(d);
  return 0;
}
```

Project 3: Implement Simple mmap()

- **Implement three system calls and page fault handler on xv6**
- **What your code should handle**

1. mmap() syscall
2. Page fault handler
3. munmap() syscall
4. freemem() syscall

1. mmap() system call on xv6

- Simple mmap() synopsis

```
uint mmap(uint addr, int length, int prot, int flags, int fd, int offset)
```

1. ***addr*** is always page-aligned
   - MMAPBASE + addr is the start address of mapping
   - MMAPBASE of each process's virtual address is 0x40000000
2. ***length*** is also a multiple of page size
   - MMAPBASE + addr + length is the end address of mapping
3. ***prot*** can be **PROT_READ** or **PROT_READ|PROT_WRITE**
   - *prot* should be match with file's open flag

1. mmap() system call on xv6

- Simple mmap() synopsis

```
uint mmap(uint addr, int length, int prot, int flags, int fd, int offset)
```

4. *flags* can be given with the combinations
   1) **If MAP_ANONYMOUS** is given, it is anonymous mapping
   2) **If MAP_ANONYMOUS** is not given, it is file mapping
   3) If **MAP_POPULATE** is given, allocate physical page & make page table for whole mapping area.
   4) If **MAP_POPULATE** is not given, just record its mapping area.

      If page fault occurs to according area (access to mapping area's virtual address), allocate physical page & make page table to according page
   5) Other flags will not be used

1. mmap() system call on xv6

- Simple mmap() synopsis

```
uint mmap(uint addr, int lenth, int prot, int flags, int fd, int offset)
```

5. *fd* is given for file mappings, if not, it should be -1
6. *offset* is given for file mappings, if not, it should be 0

Return

Succeed: return the start address of mapping area

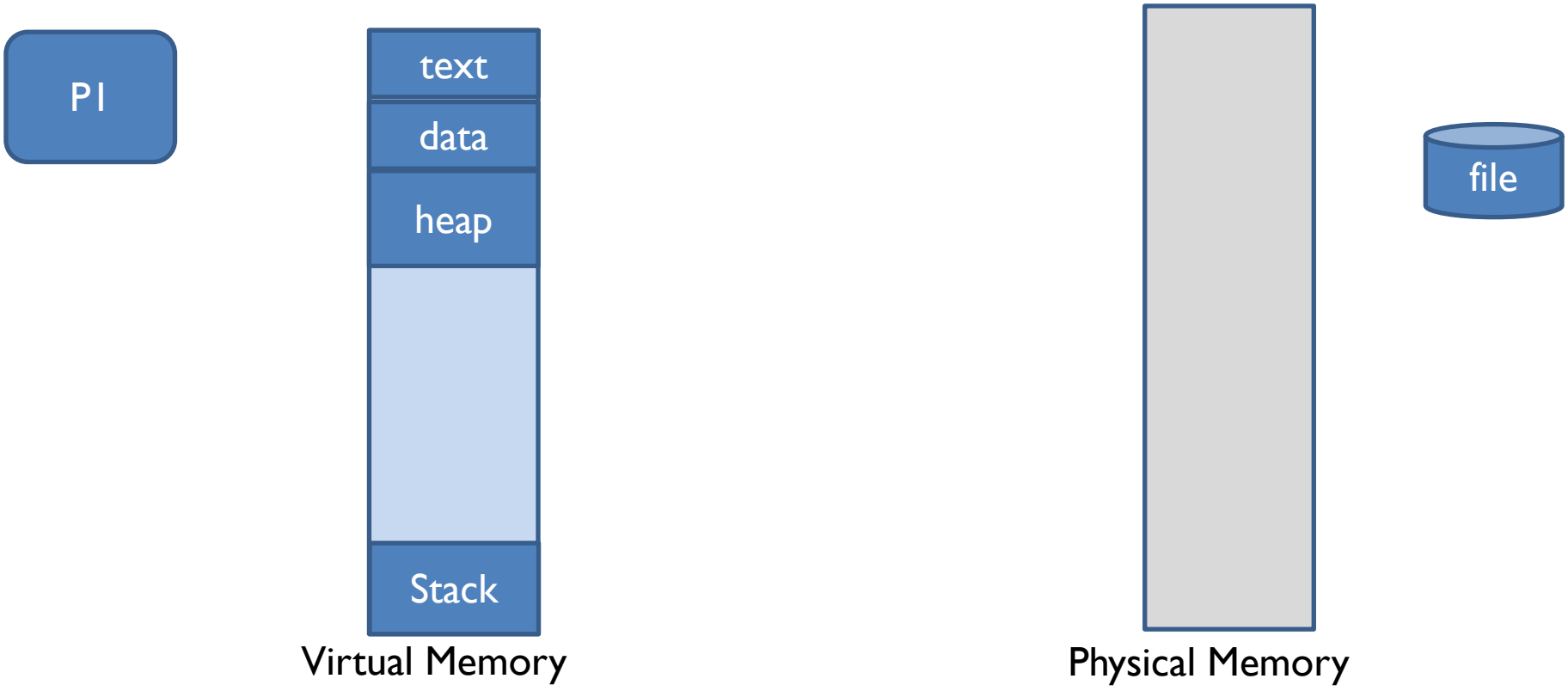
Failed: return 0

- It's not anonymous, but when the fd is -1
- The protection of the file and the prot of the parameter are different
- The situation in which the mapping area is overlapped is not considered
- If additional errors occur, we will let you know by writing notification

How file mmap() Works

1) Private file mapping **with MAP_POPULATE**

- mmap(0, 8192, PROT_READ, MAP_POPULATE, fd, 4096)
  - mmap 2pages



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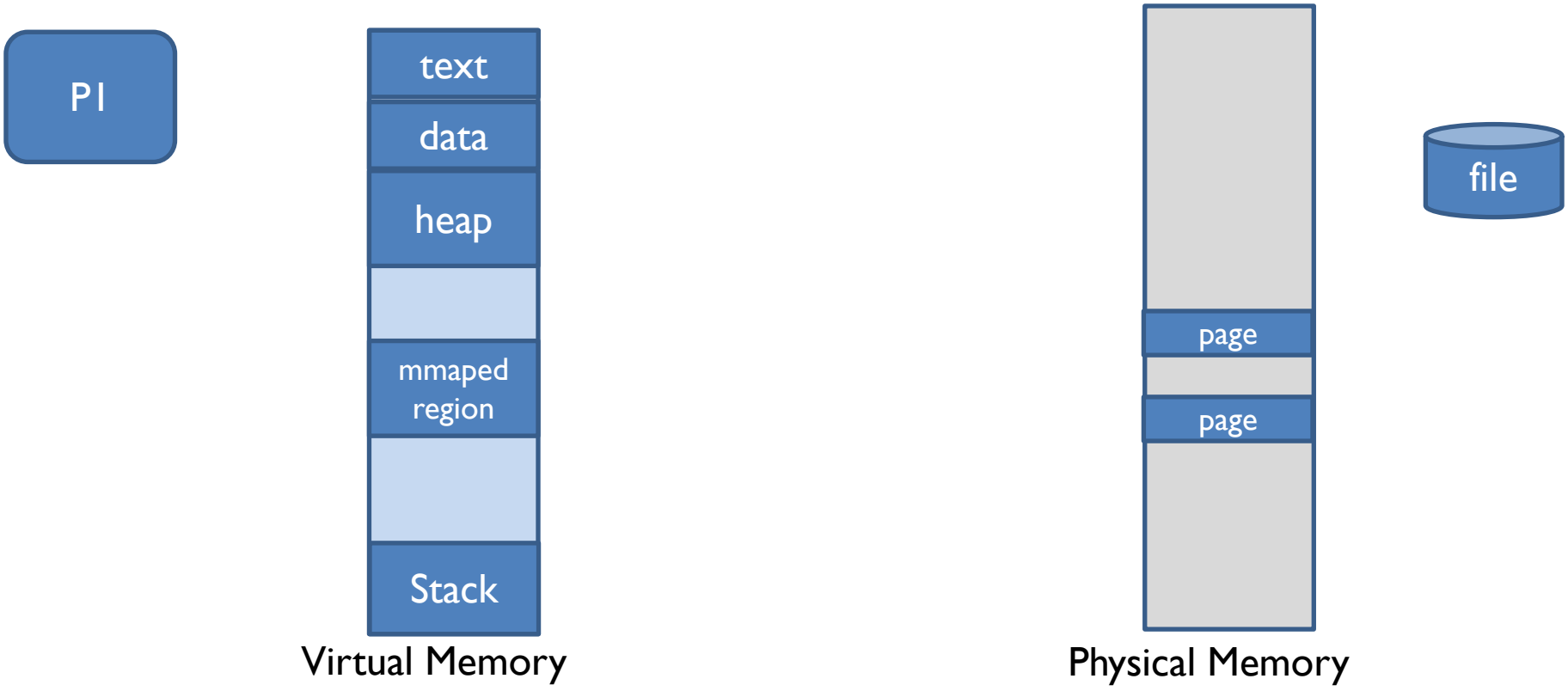

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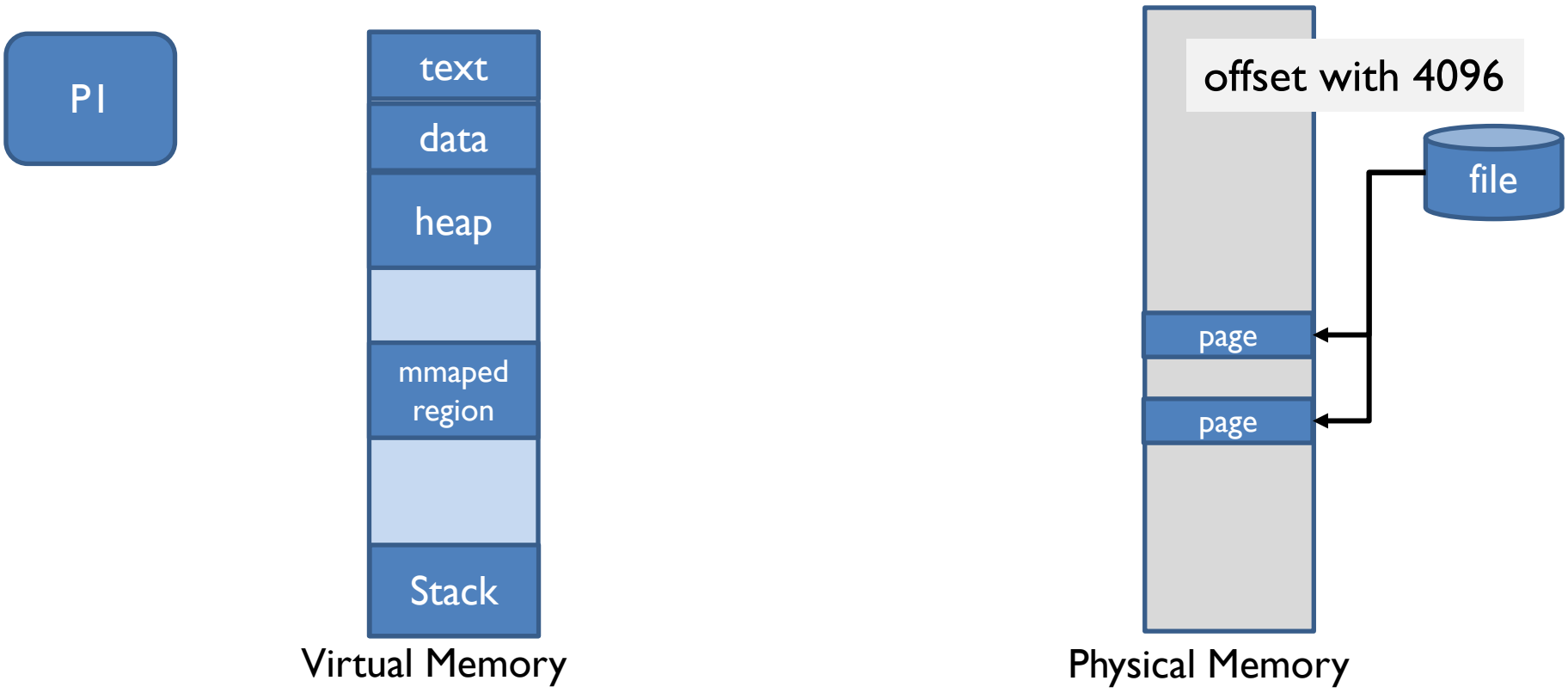

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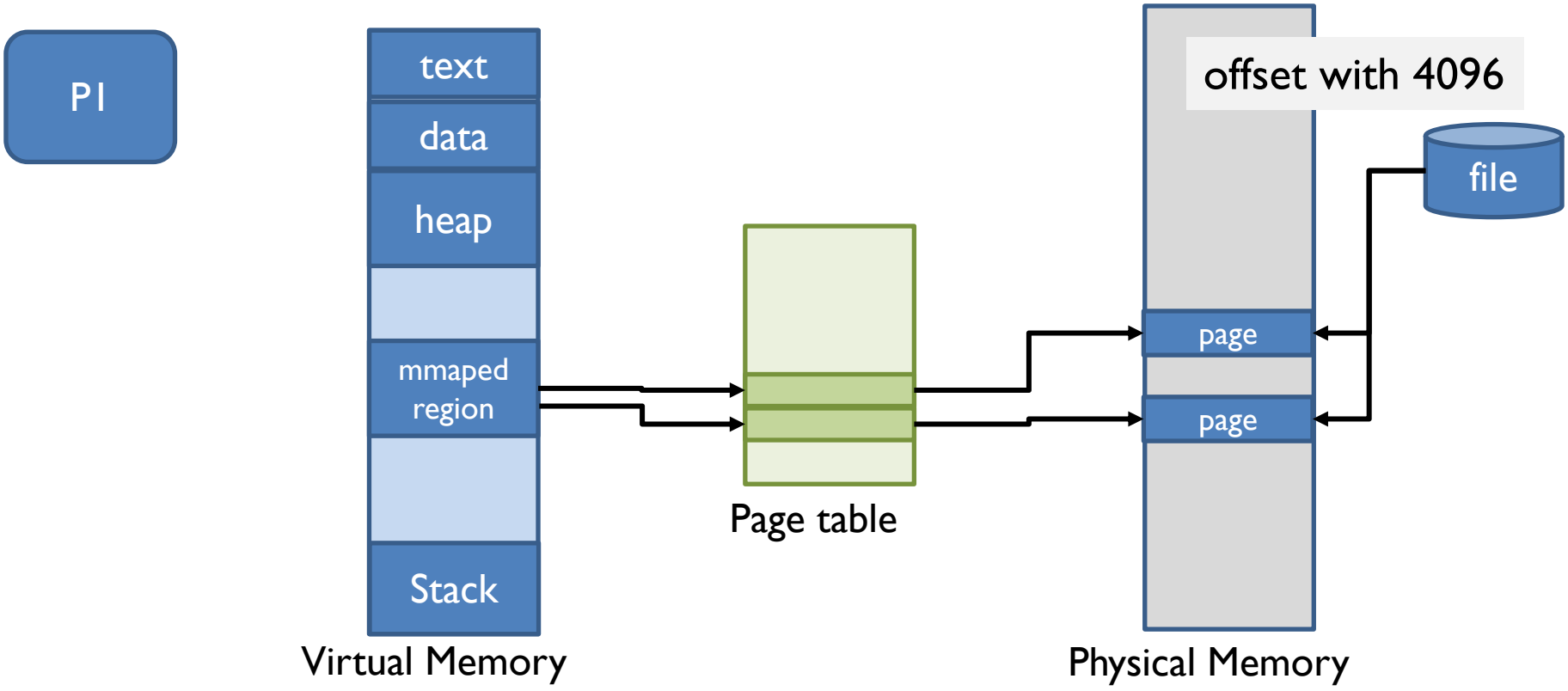

How file mmap() Works

2) Private file mapping **without MAP_POPULATE**

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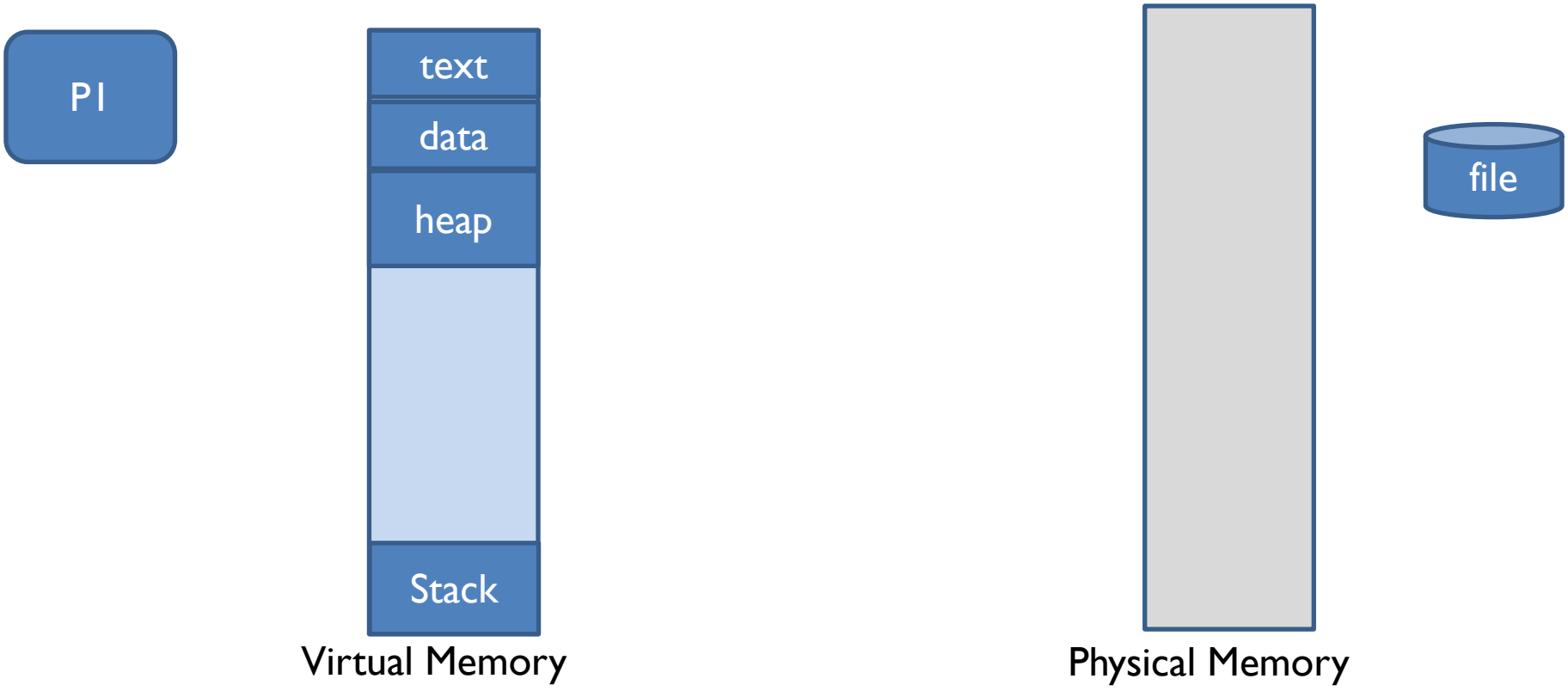

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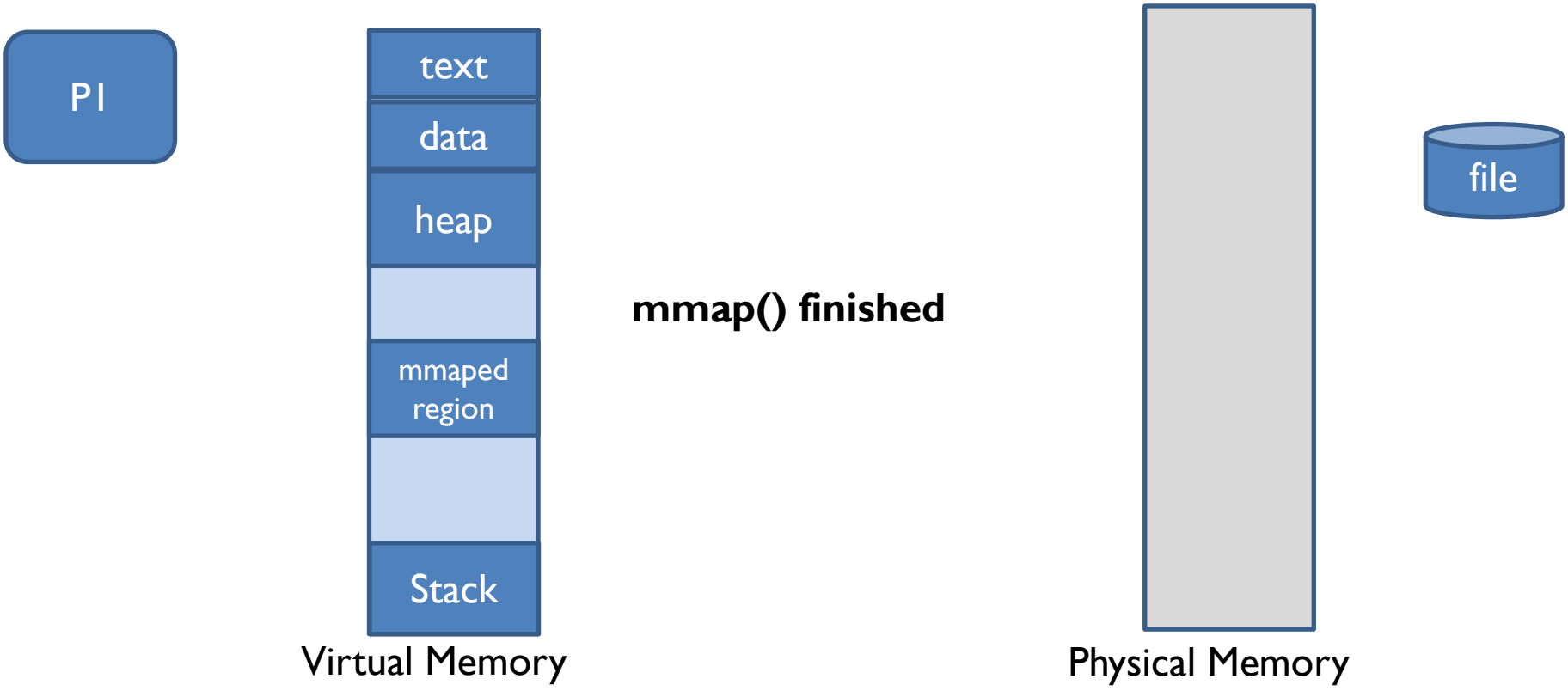

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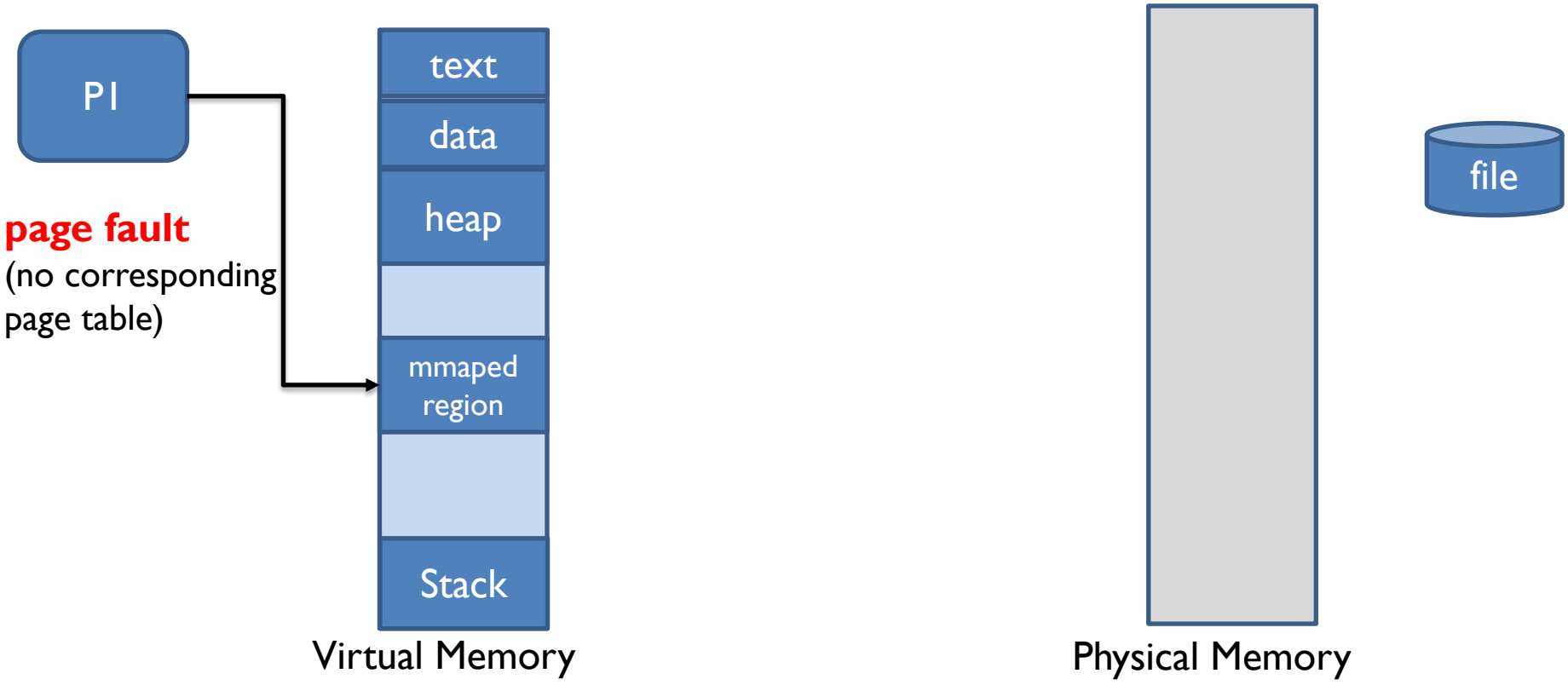

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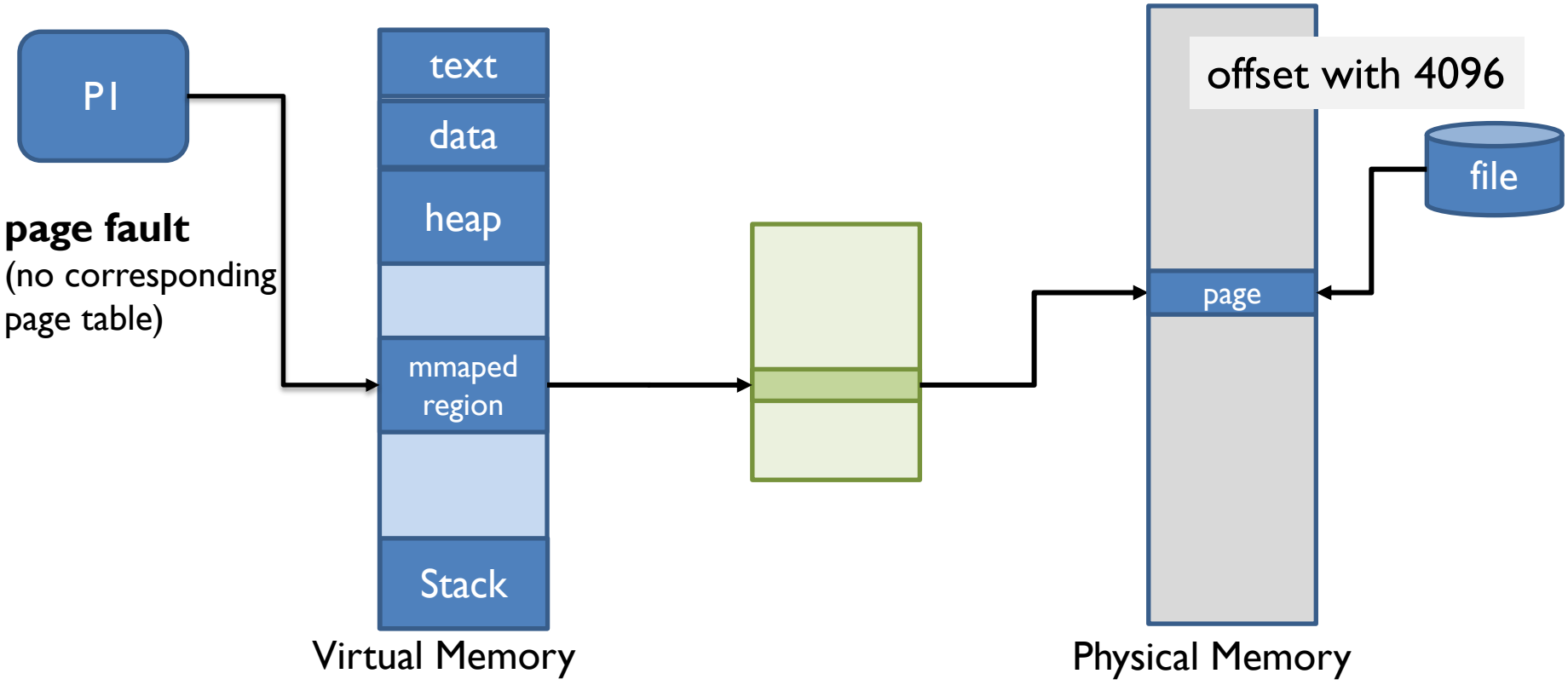
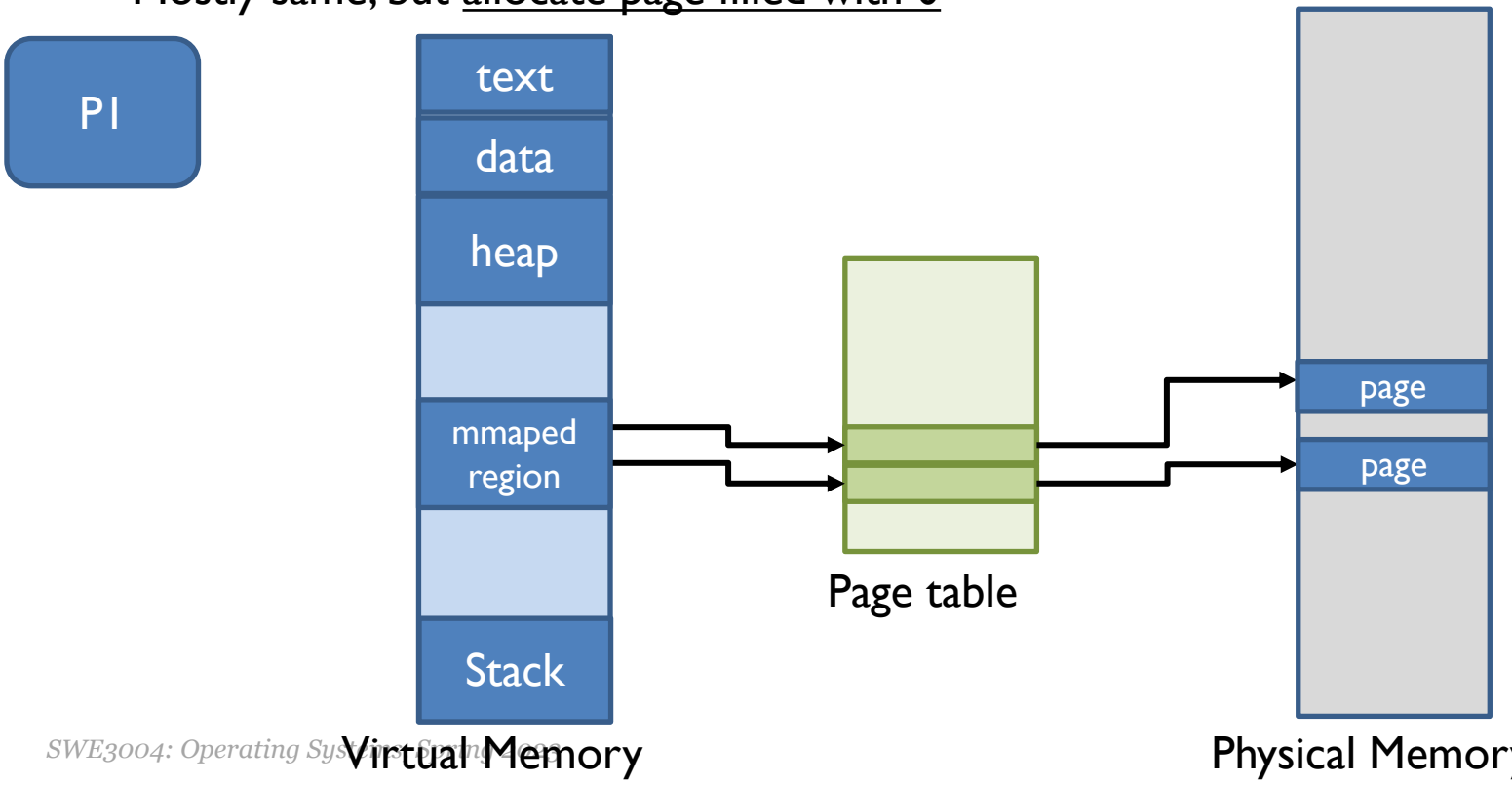


How anonymous mmap() Works

3) Private anonymous mapping **with MAP_POPULATE**

- mmap(0, 8192, PROT_READ,

  MAP_POPULATE|MAP_ANONYMOUS, **-1**, 0)
  - mmap 2pages
- Mostly same, but <u>allocate page filled with 0</u>



Implementation detail of mmap()

- Parameters will be defined at param.h
  - PROT_READ 0x1
  - PROT_WRITE 0x2

  - MAP_ANONYMOUS 0x1
  - MAP_POPULATE 0x2

Implementation detail of mmap()

- struct mmap_area {

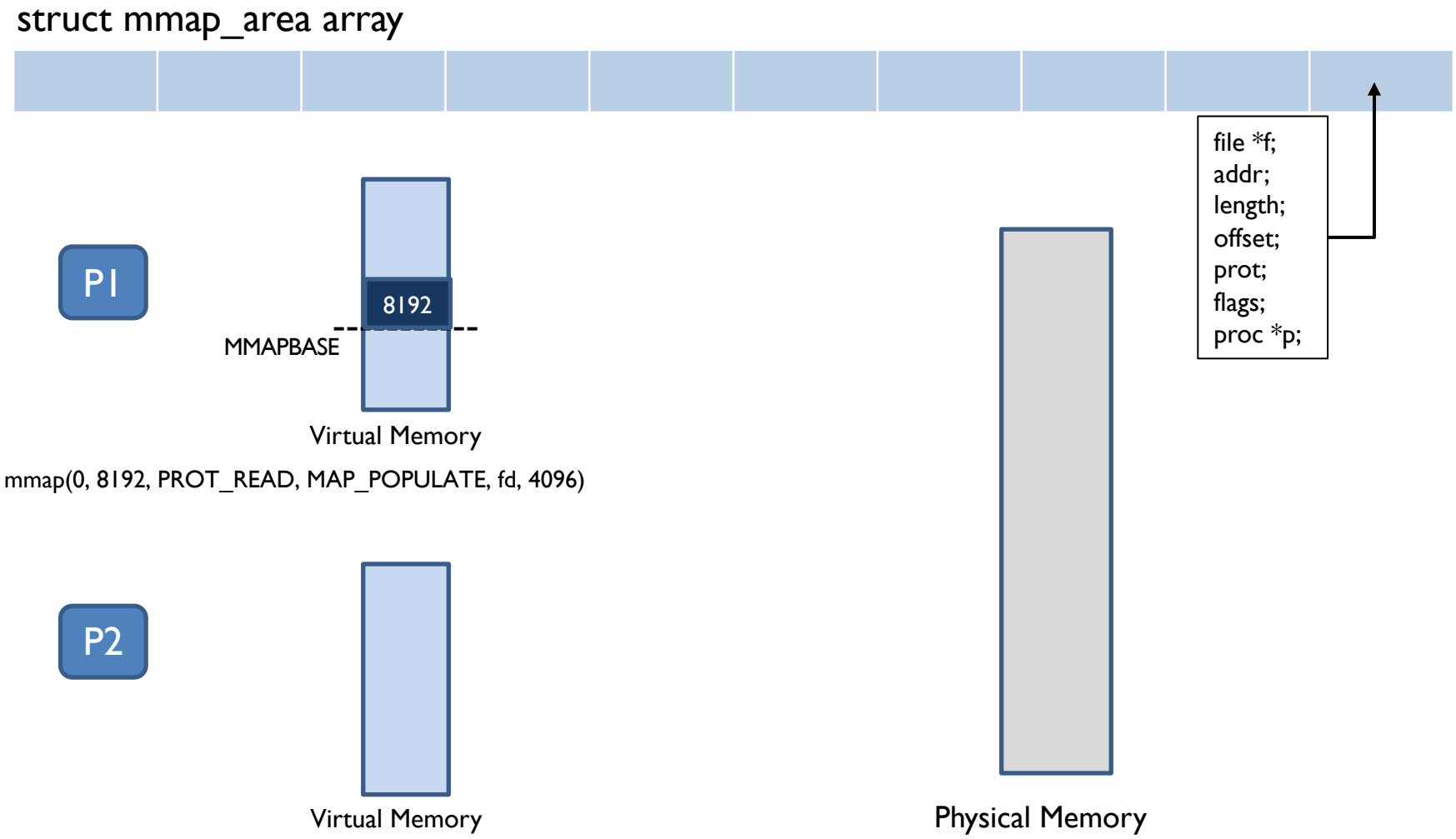
```
strict file *f;
uint addr;
int length;
int offset;
int prot;
int flags;
struct proc *p // the process with this mmap_area
```

}

- Manage all mmap areas created by each mmap() call in one mmap_area array.
- Maximum number of mmap_area array is 64.

mmap_area Array Example

In the case of populate...



mmap_area Array Example

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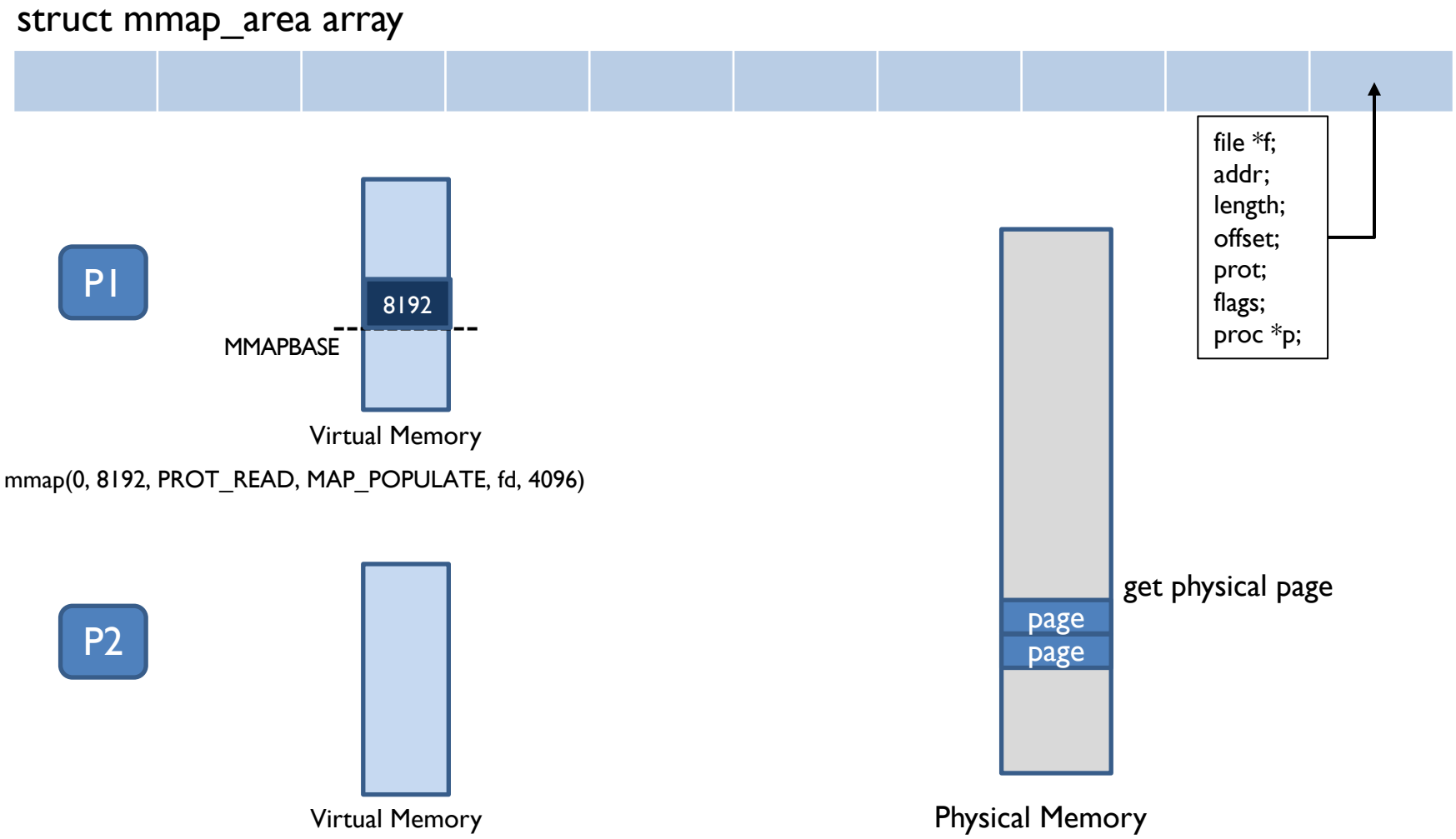

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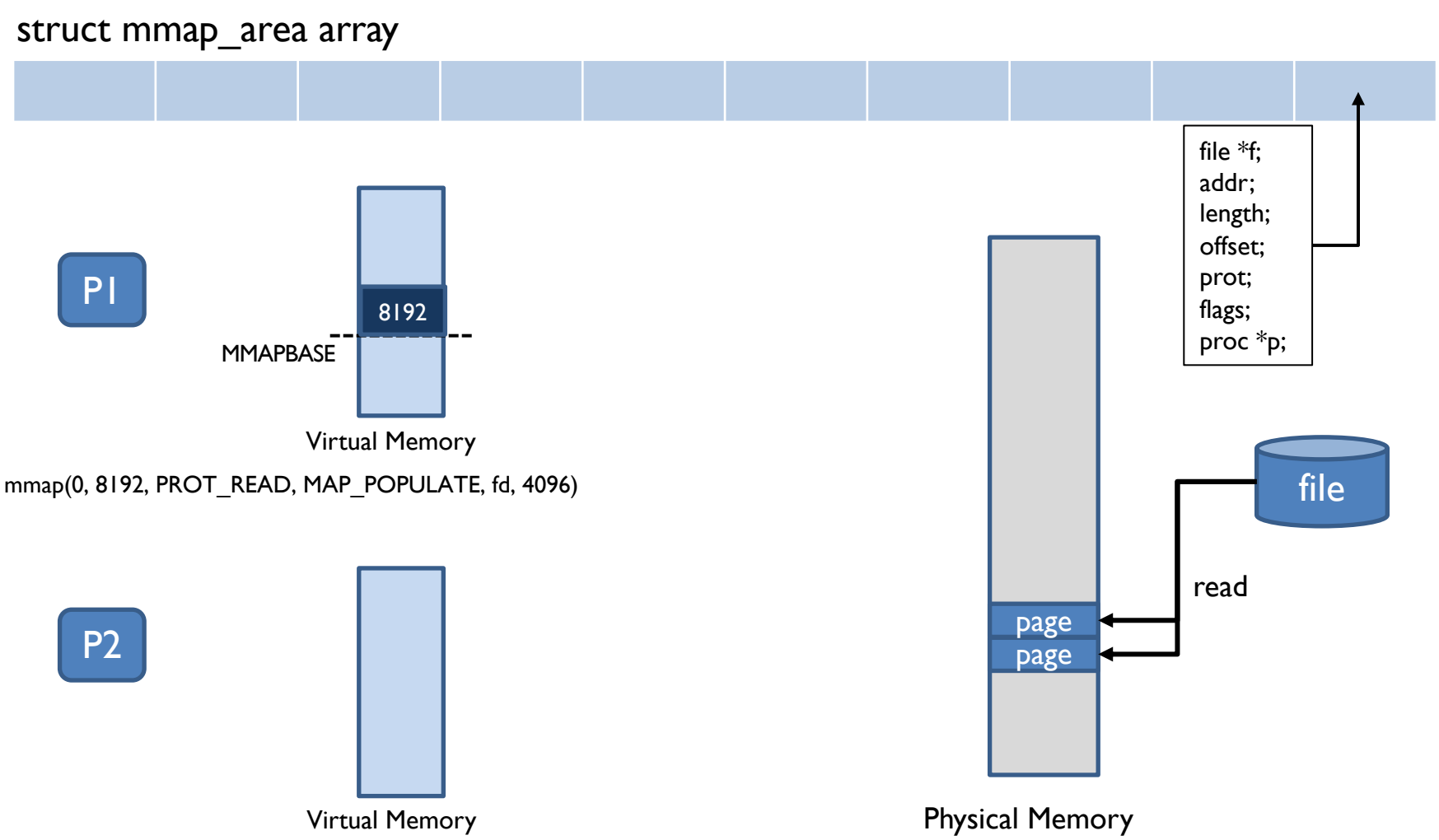

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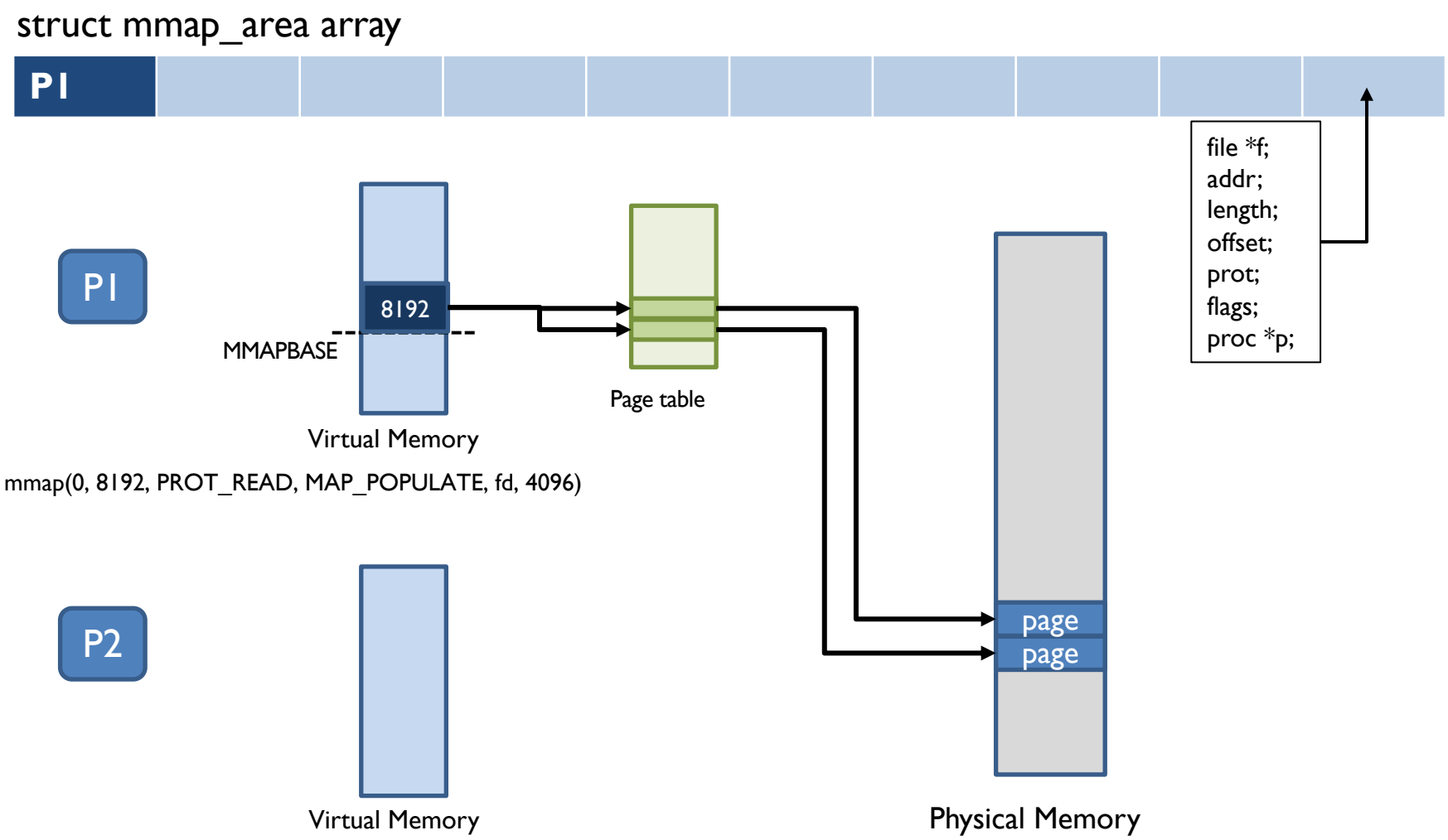

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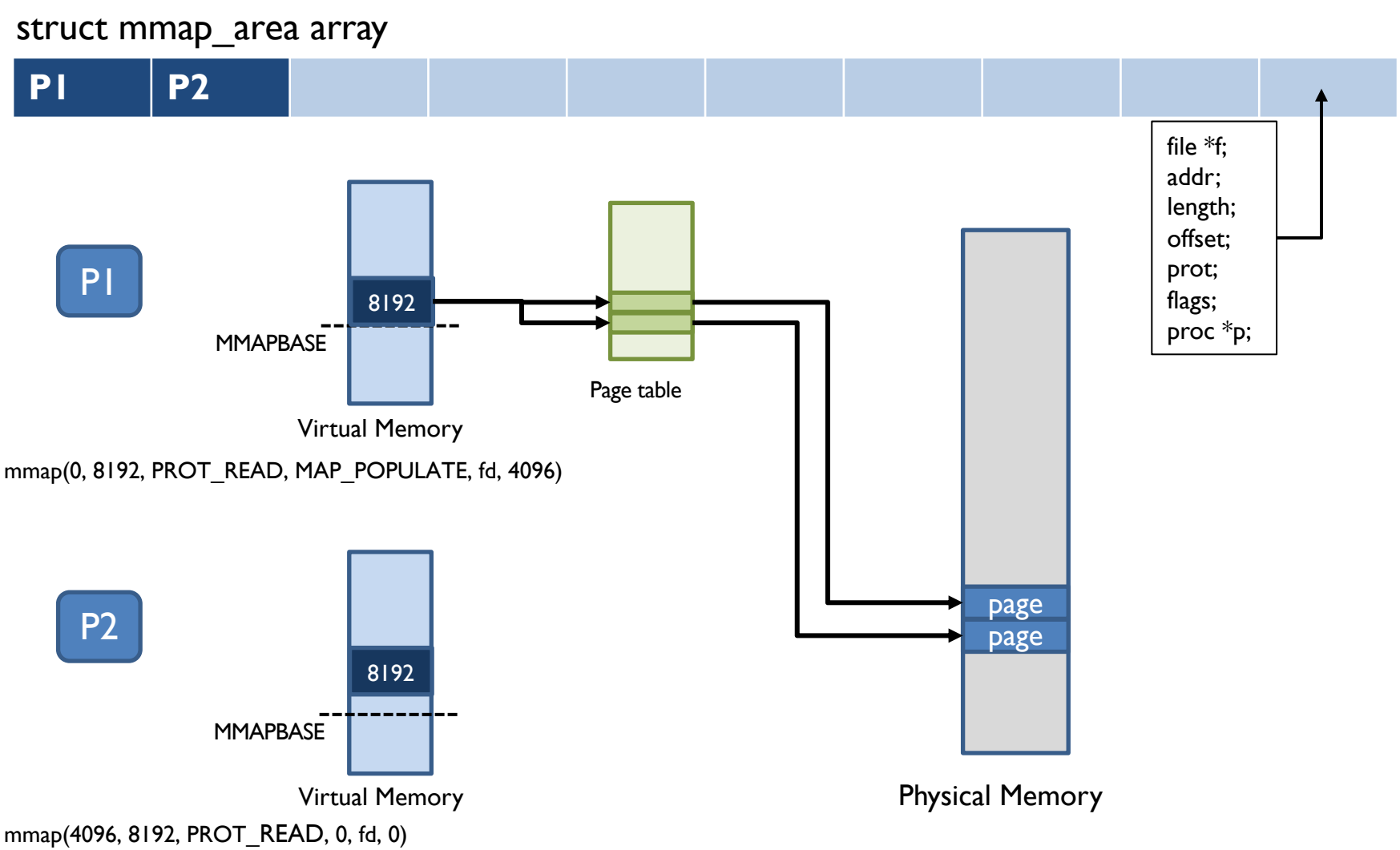

2. Page Fault Handler on xv6

- **Page fault handler is for dealing with access on mapping region with physical page & page table is not allocated**
- **Succeed:** Physical pages and page table entries are created normally, and the process works without any problems
- **Failed:** The process is terminated

1. When an access occurs (read/write), catch according page fault (interrupt 14, T_PGFLT) in ***traps.h***
2. In page fault handler, determine fault address by reading CR2 register(using rcr2()) & access was read or write

   read: tf->err&2 == 0 / write: tf->err&2 == 1
3. Find according mapping region in mmap_area

   If faulted address has no corresponding mmap_area, return -1
4. If fault was write while mmap_area is write prohibited, then return -1
5. For only one page according to faulted address
   1. Allocate new physical page
   2. Fill new page with 0
   3. If it is file mapping, read file into the physical page with offset
   4. If it is anonymous mapping, just left the page which is filled with 0s
   5. Make page table & fill it properly (if it was PROT_WRITE, PTE_W should be 1 in PTE value)

3. munmap() system call on xv6

- **munmap(*addr*)**
- **Unmaps corresponding mapping area**
- **Return value: 1 (succeed), -1 (failed)**

1. *addr* will be always given with the start address of mapping region, which is page aligned
2. munmap() should remove corresponding mmap_area structure
   - If there is no mmap_area of process starting with the address, return -1
3. If physical page is allocated & page table is constructed, should free physical page & page table
   - When freeing the physical page should fill with 1 and put it back to freelist
4. If physical page is not allocated (page fault has not been occurred on that address), just remove mmap_area structure.
5. Notice) In one mmap_area, situation of some of pages are allocated and some are not can happen.

4. freemem() system call on xv6

- syscall to return current number of free memory pages

1. When kernel frees (put page into free list), freemem should be increase
2. When kernel allocates (takes page from free list and give it to process), freemem should decrease

Submission

- This project is to implement three system calls and page fault handler
  - mmap() syscall
  - Page fault handler
  - munmap() syscall
  - freemem() syscall

- Use the *submit* & *check-submission* binary file in Ji Server
  - Make clean
  - $ ~swe3004/bin/submit pa3 xv6-public
  - you can submit several times, and the submission history can be checked through check-submission
    - Only the last submission will be graded

Submission

- Report
  - Submit to iCampus
  - Modified code and explanation
  - Free length
  - pa3_2019310392.pdf

Submission

- PLEASE DO NOT COPY
  - We will run inspection program on all the submissions
  - Any unannounced penalty can be given to **both students**
    - 0 points / negative points / F grade …

- Due date:5/10(Wed.), 23:59:59 PM
  - **-25**% per day for delayed submission

Questions

- If you have questions, please ask on i-campus discussion section
  - Please use the discussion board
  - We don't reply i-campus messages and e-mails
- You can also visit Corporate Collaboration Center #85533
  - Please e-mail TA before visiting
- Reading xv6 commentary will help you a lot
  - http://csl.skku.edu/uploads/SSE3044S20/book-rev11.pdf

Appendix. Hint

- File structures corresponding to fd are contained in proc->ofile[fd]
  - File structure can be used to get file data and file protection
- Page table entry can be created using mappages as in copyuvm
- At fork time, if the parent process has mmap areas the child process will also have mmap areas at the same address, so this needs to be processed
- Page fault invokes the trap function in trap.c after similar processing of system calls in Project 1.
  - Here, you can utilize rcr2() and tf->err