REVIEW EXERCISE 05

Question 1. Given a knowledge base KB as follows, $\{P \to R, \neg S \to P, \neg S, R \to Q\}$. Consider the pseudo-code function PL-RESOLUTION given in the lecture to check whether **KB entails Q**.

Present your work to the table below, in which the first column contains KB $\wedge \neg \alpha$ in CNF, and every of the next columns includes new sentences added to KB after each loop. Note that

- Duplicated sentences are omitted from the table
- Circle the unit clauses that lead to the contradiction and hence the function ends successfully, if possible
- Process the clauses in order, that is first pair clause 1 with clause 2, 3, 4... then pair clause 2 with clause 3, 4,... and so on.

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CNF sentences	Loop 1	Loop 2	Loop 3	Loop 4	
1) ¬P∨R	6) R v S	10) R			
2) P ∨ S	7) ¬P∨Q	11) ¬P			
3) ¬S	8) P	12) S v Q			
4) Q∨¬R	9) ¬R	13) S			
5) ¬Q		14) Q			

Circle the correct option, IS or IS NOT.

Following the result of resolution, the sentence Q **IS** / **IS NOT** entailed by KB.

Question 2. Repeat Question 1. but this time you check whether **KB entails** \neg **Q**.

CNF sentences	Loop 1	Loop 2	Loop 3	Loop 4
1) ¬P∨R	6) R v S	9) R		
2) P ∨ S	7) ¬P∨Q	10) S v Q		
3) ¬S	8) P	11) Q		
4) Q∨¬R				
5) Q				

There is no new clause generated after Loop 2. There is no contradiction, either. Thus, **KB** does not entails $\neg \mathbf{Q}$.

Question 3. Are the above problems solved by using Forward chaining or Backward chaining? Give your reason.

NO, there exists a sentence that is not a definite clause, i.e $\neg S \Rightarrow P$

Question 4. Consider the following text. "Heather attended the meeting or Heather was not invited. If the boss wanted Heather at the meeting, then she was invited. Heather did not attend the meeting. If the boss did not want Heather there, and the boss did not invite her there, then she is going to be fired."

Use resolution to prove that **Heather is going to be fired**. *Hint: clauses in italic are good candidates for propositions.*

Let each of following propositions denote the facts represented in the corresponding clause.

- Proposition A represents for "Heather attended the meeting."
- Proposition I represents for "Heather was invited."
- Proposition W represents for "The boss wanted Heather at the meeting."
- Proposition F represents for "Heather is going to be fired."

Then the propositional KB in CNF will be

- 1) $A \vee \neg I$
- 2) $\neg W \vee I$
- 3) ¬A
- 4) $W \vee I \vee F$

Apply resolution to KB $\land \neg \alpha$

5) ¬F	Negation of conclusion		
6) <u>W</u> ∨ I	from sentences 4 and 5		
7) I	from sentences 2 and 6		
8) A	from sentences 1 and 7		
9) •	from sentences 3 and 8		

Conclusion: Therefore, Heather is going to be fired.

Question 5. Consider the following knowledge base of definte clauses.

1.
$$C \wedge D \rightarrow Y$$

5. B

2.
$$R \wedge Z \rightarrow C$$

6. $R \rightarrow D$

3.
$$\neg B \lor D$$

7. $D \rightarrow R$

Prove Y using backward chaining

Y requires two sub-goals, C and D (from 1)

D requires the sub-goal B (from 2). B is given (from 5). Thus **D** is satisfied.

C requires two sub-goals, R and Z (from 2). R requires D (from 6), which is proved. Z requires two sub-goals, D and R, which are both proved. Thus, **C is satisfied**.

Finally, Y can be proved because all required propositions can be obtained from KB.

Prove **Y** using forward chaining (only trigger a rule once for simplicity).

- 8. D is obtained from 3 and 5.
- 9. R is obtained from 7 and 8.
- 10. Z is obtained from 4 and 8-9 (rewrite 4 in implication form)
- 10'. D is obtained **again** from 6 and 9.
- 11. C is obtained from 2 and 9-10
- 12. Y is obtained from 1 and 5-11.

Finally, Y can be proved because it is generated from KB.