

IT-Security 1

Chapter 5: Certificates and Public Key Infrastructures

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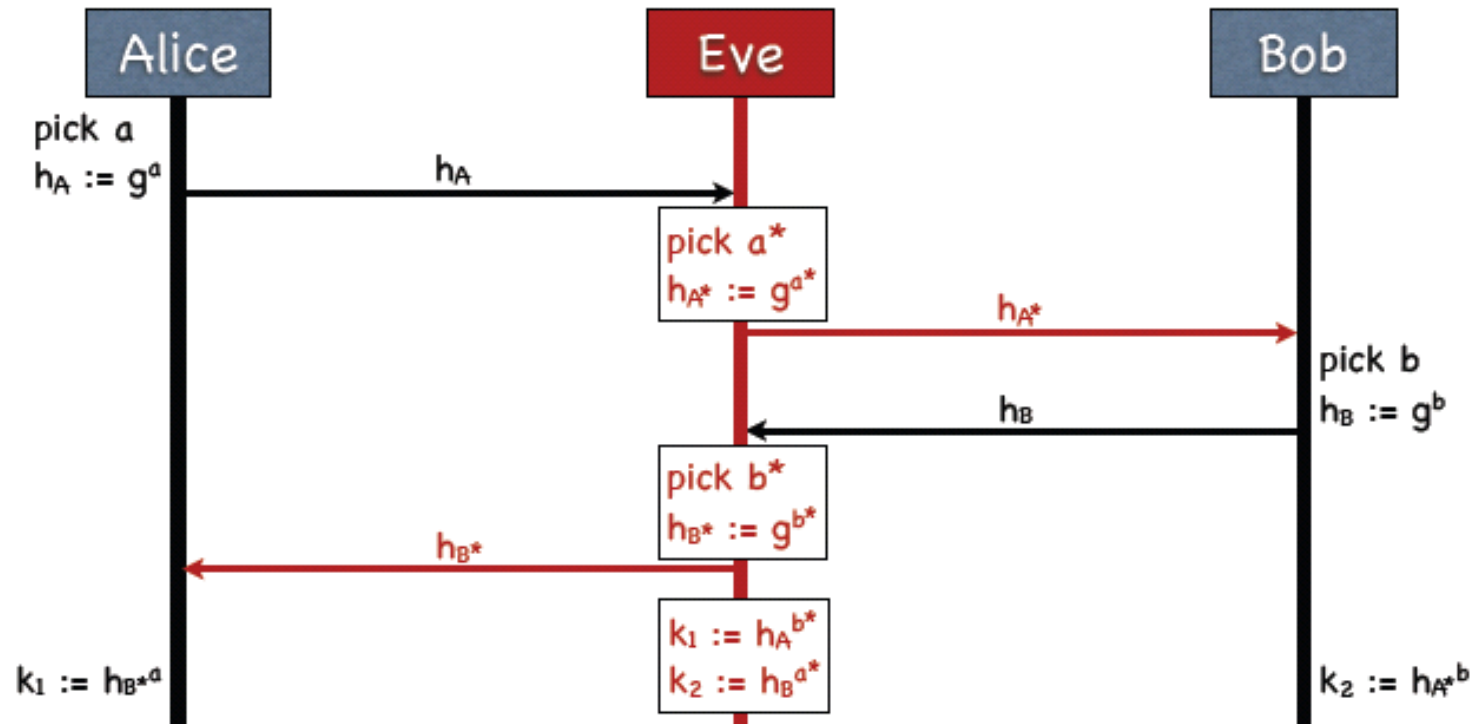
Chapter Overview

- General public key problem
- Trust models
- X.509 certificates
- Certificate revocation

General Public Key Problem

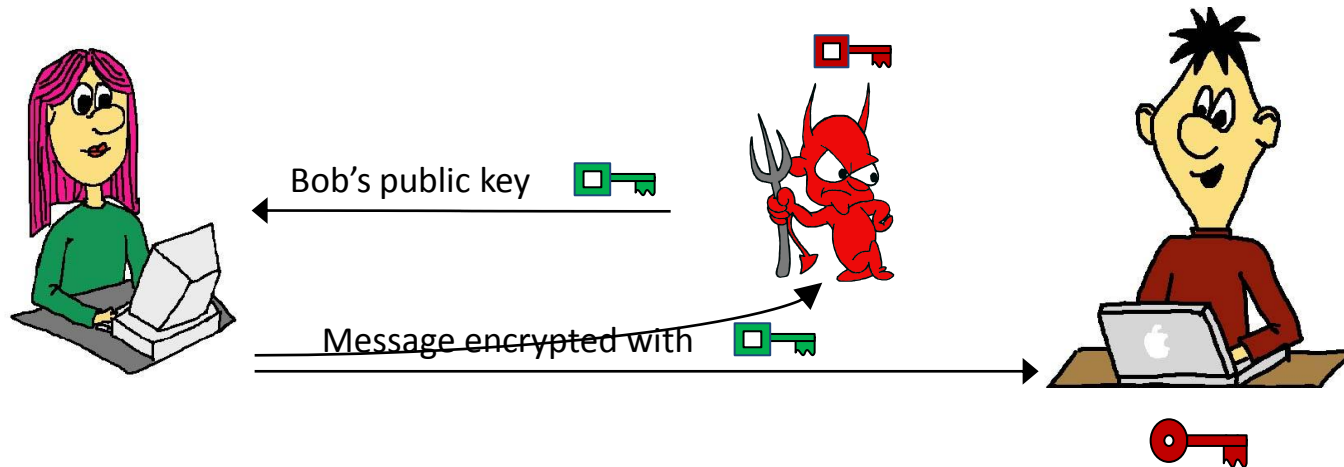
- Public key primitives require an authenticated way to provide / exchange public keys
- This holds for all asymmetric primitives:
 - Diffie-Hellman Key Exchange
 - Public key encryption
 - Digital signatures

Example: Diffie-Hellman MiM-Attack



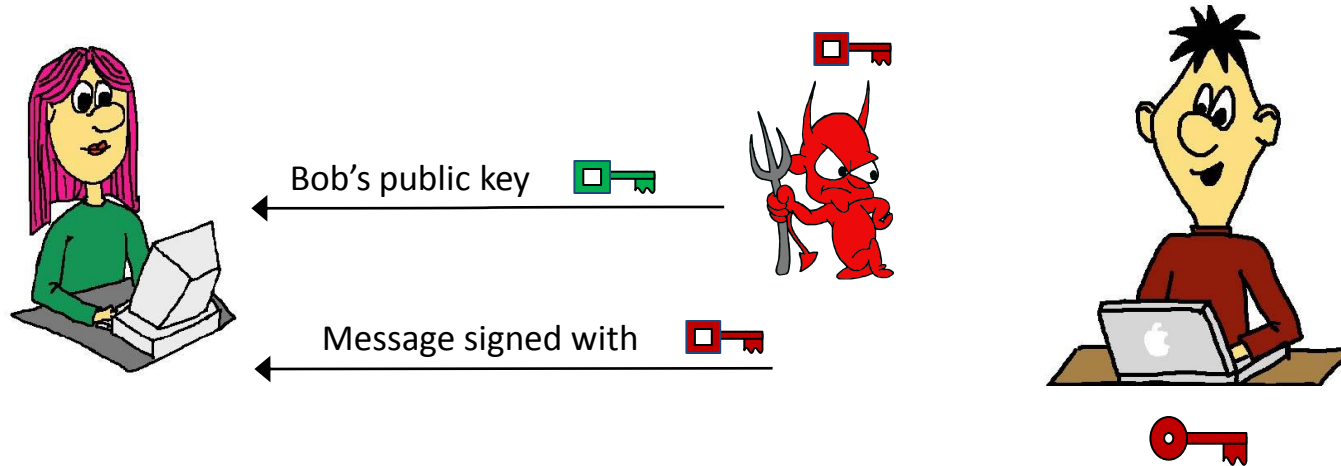
- Attack possible because Alice and Bob have no assurance of the authenticity of the public values received

Example: Public Key Encryption



- If an attacker can exchange Bob's public encryption key with his own, then he will be able to decipher confidential messages sent from Alice to Bob

Example: Public Key Signatures



- If an attacker can exchange Bob's public signature verification key with his own, the attacker can convince Alice that Bob has signed a certain message

Trusted Third Parties

- Solution: guarantee authenticity with the help of a trusted third party



- Assumption: Alice and Bob are in possession of an authentic copy of the public signature key of the trusted third party
- The trusted third party signs pairs of public keys and identifiers and thereby binds the keys to the identifiers
- I.e. the trusted third party issues certificates for the public keys of Alice and Bob
- Bob can verify the authenticity of Alice's public key by checking the signature of the trusted third party on Alice's certificate

- A trusted third party that issues certificates is called **Certification Authority**

Certificates

- A certificate binds the public key of a principal to some identifier (or attribute) of that principal
- A certificate is issued by some issuer, typically called a certification authority (CA)
- A certificate minimally comprises
 - The public key
 - An identifier of the principal
 - The issuer
 - The signature on the hash of the rest of the content of the certificate with the private key of the issuer
- Note:
 - We will discuss a real-world certificate format later on

Public Key Infrastructures

- Consists of the components to securely distribute public keys, these may include:
 - A certification authority that issues certificates for public keys
 - A repository for retrieving certificates
 - A method for revoking certificates
 - A method for evaluating “chains” of certificates
 - Sometimes a registration authority (RA) is used in addition to the certification authority
 - RA ensures the correctness of the binding between public key and principal and distributes the certificate (and the private key) to the principal
 - CA generates the actual certificate

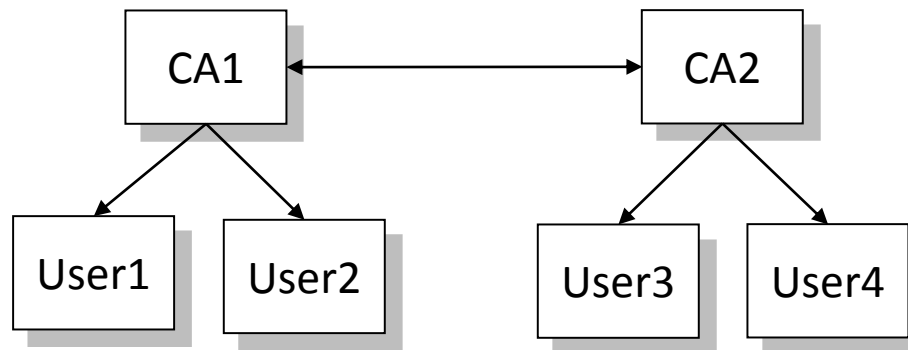
CAs and Trust

- The use of a CA requires
 - Alice and Bob to trust that the CA properly authenticates the principal before issuing the certificate
 - Alice and Bob to obtain the CA's public key in an authenticated manner
- Obviously a single world-wide CA is not realistic as
 - It would have to be trusted by any principal world-wide
 - Any principal would have to store the CA's public key
 - What if the CA's private key is broken and has to be exchanged?
 - Any principal would have to be able to obtain a certificate from that CA in a secure manner
 - ...

Some Ways to Use More Than one CA

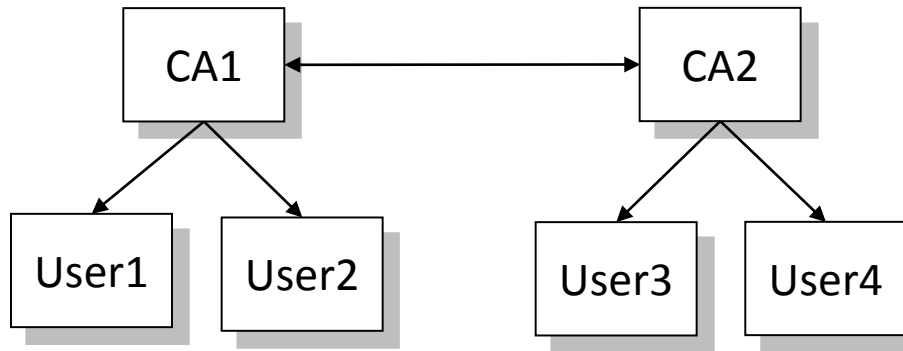
- Oligarchy: Several CAs are used simultaneously. All principles store the public keys of all CAs
- Cross certification: CA1 issues a certificate for CA2 and CA2 issues a certificate for CA1
- CA hierarchy: the CA's are hierarchically structured
 - A “root” CA issues certificates for the CAs on the second level of the hierarchy
 - The CAs on the second level of the hierarchy issue certificates for the CAs on the third level etc.
- Anarchy:
 - Everyone can sign everybody else's key
 - Everyone decides and configures whom they trust

Cross Certification



- How can User1 verify the certificate $\ll \text{User3} \gg_{\text{CA2}}$ of User3, assuming that User1 has an authentic copy of CA1's signature verification key?

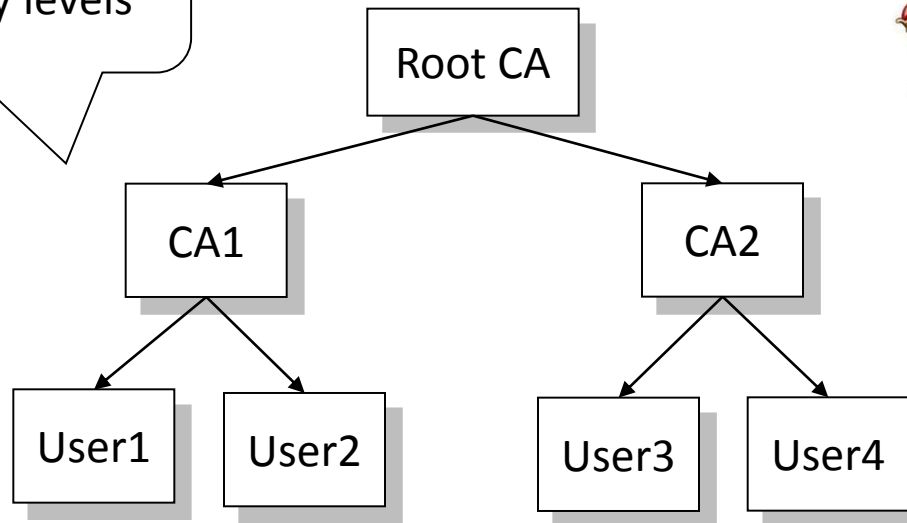
Cross Certification



- User3 presents User1 with the chain of certificates $\langle\langle\text{CA2}\rangle\rangle_{\text{CA1}}$ and $\langle\langle\text{User3}\rangle\rangle_{\text{CA2}}$
- User1 verifies CA1's signature on CA2's certificate, extracts CA2's key and uses it to verify the signature on $\langle\langle\text{User3}\rangle\rangle_{\text{CA2}}$

CA Hierarchies (Top down)

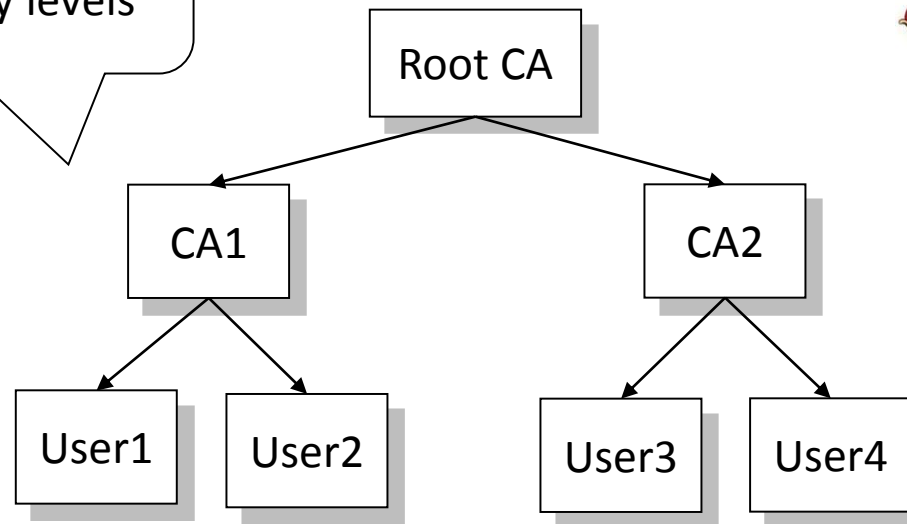
There can of course be more hierarchy levels



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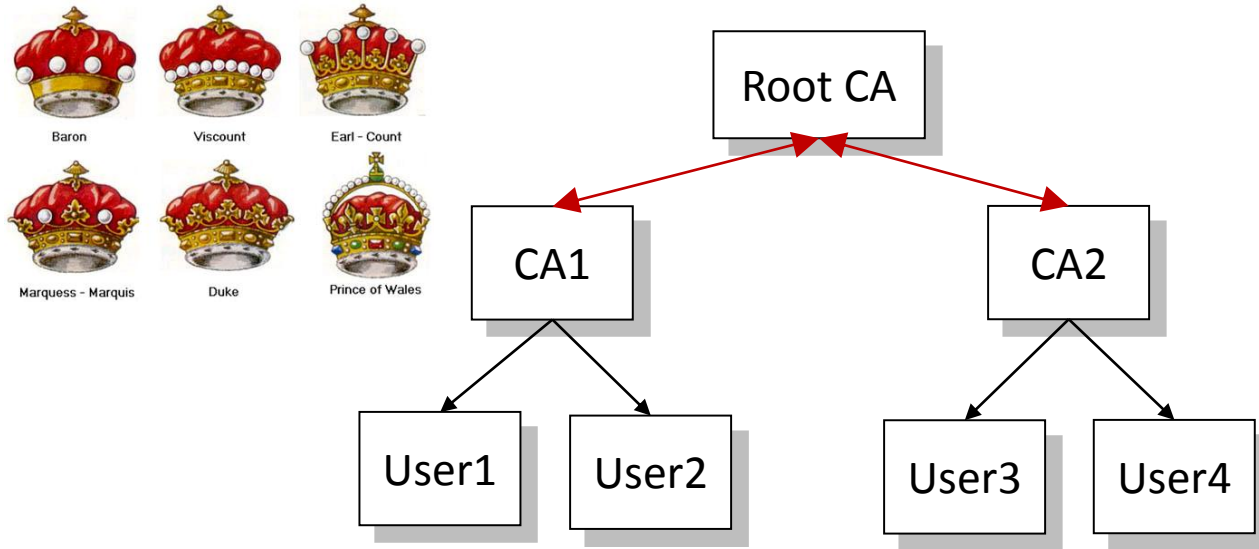
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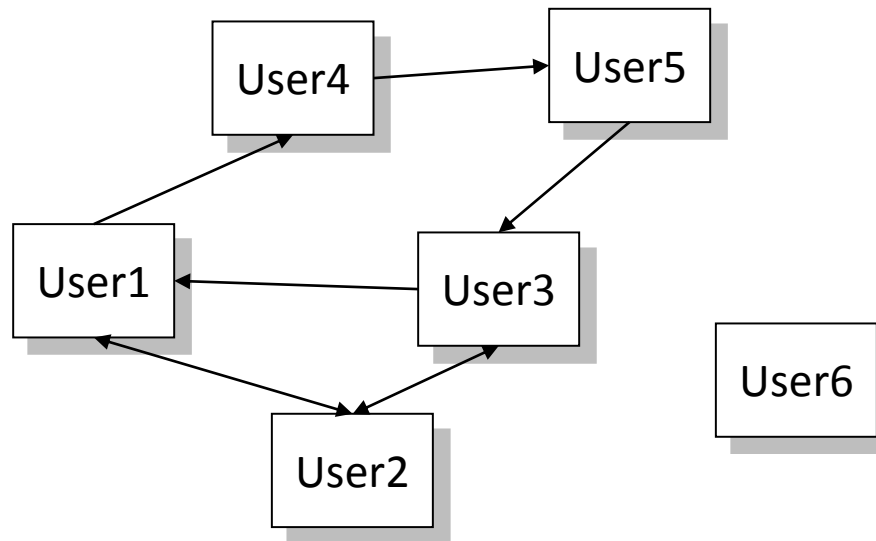
- User1 needs to have public key of the Root CA
- User 1 verifies the signature of the Root CA on CA2's certificate, extracts the public key of CA2 and uses it to verify the signature on User3's certificate

CA Hierarchies (Bottom up)



- The lower level CAs also sign the upper level CAs
- Advantage: User1 only has to know CA1's public key in advance
- Disadvantage: User1 has to check the signatures on more certificate

Anarchy: e.g. used in PGP



- Obviously hard to find chains if used by a large group
- Left to the choice of user's whom to trust

Certificate Revocation

- Certificates are typically issued for a certain validity period
- However, the private key may be compromised before the validity period of the corresponding certificate ends
- It is often required to be able to revoke a certificate before it expires
- Consequence: the verifier of a certificate needs to check if
 - the signature(s) on the (chain of) certificate(s) is/are correct
 - the certificate(s) is/are still valid or are already expired
 - the certificate(s) has/have not been revoked e.g. by fetching the appropriate Certificate Revocation Lists

Certificate Revocation with CRLs

- Certificate Revocation List (CRL) is issued periodically by the CA
- CRL is a signed list of all revoked certificates which validity period has not expired yet
- CRL is signed with the private key of the CA

X.509 Certificates and CRLs

- ITU-T standard for public key certificates and certificate revocation lists
- Profile for use on the Internet defined in RFC 5280
- Used e.g. for
 - SSL/TLS in web browsers
 - IPsec
 - In various authentication protocols

X.509 Certificates

- Signed Content:
 - Version – The version number
 - Serial Number – An issuer-unique serial number
 - Signature – The signature algorithm identifier (includes hash function used)
 - Issuer – Name of the issuer
 - Validity – from: / to: validity period
 - Subject – Name of the subject
 - SubjectPublicKeyInfo – Public key
 - IssuerUniqueID – Unique identifier of the issuer (optional)
 - SubjectUniqueID – Unique identifier of the subject (optional)
 - Extensions
- Additional unsigned content
 - SignatureAlgorithm – Algorithm identifier (includes hash function used)
 - SignatureValue – Signature on the hash of the content

Examples for Extensions in X.509

- **Key Usage** – defines the purpose of the key (e.g. signature verification or encryption)
- **Authority Key Identifier** – e.g. issuer name, serial number of certificate corresponding to the key with which the certificate was signed
- **Subject Alternative Name** – may be used in addition or instead of the Subject field. May include e.g. an email address, a DNS name, an IP address etc.
- **Name Constraints** – may be used in CA certificates only. Describes constraints on the subject names for which a CA can issue certificates

X.509 CRLs

- Signed content
 - Version – if present must be v2
 - Signature – Signature algorithm identifier (includes hash function)
 - Issuer - Name
 - thisUPdate - Time
 - nextUpdate – Time (optional)
 - revokedCertificates – List of revoked certificates
 - userCertificate – Certificate serial number
 - revocationDate - Time
 - crlEntryExtensions – optional extensions
 - CrlExtensions – optional
- SignatureAlgorithm (includes hash function)
- SignatureValue

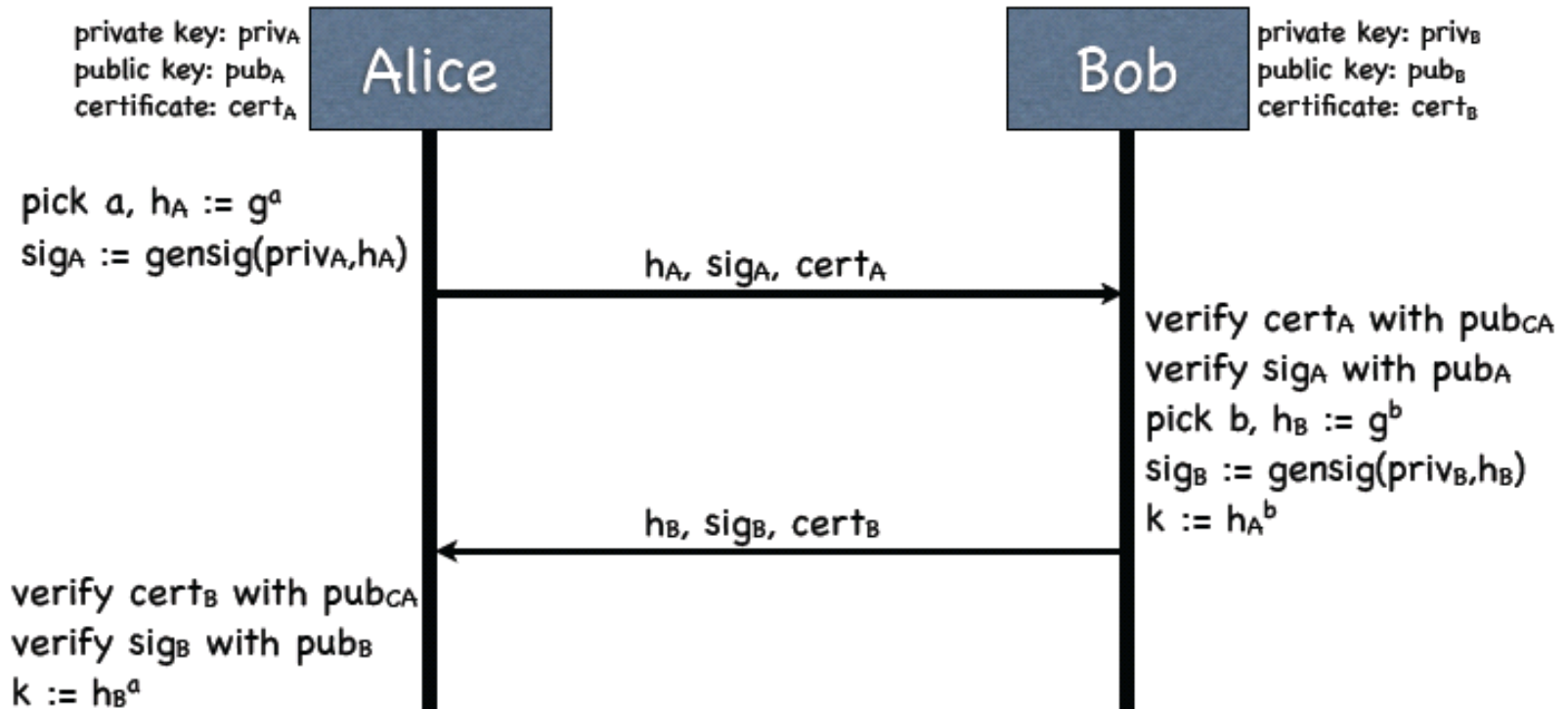
Online Certificate Status Protocol

- The disadvantage of CRLs is that revocation is only as timely as the period used for issuing CRLs
- The OCSP allows applications to explicitly query if certificates are still valid
 - OCSP client sends status request to OCSP responder
 - Includes list of serial numbers of certificates
 - OCSP responder replies with the status for all certificates in the list
 - The status is one of: good, revoked, unknown
 - The response is signed by the responder
- OCSP only defines the content of messages, not their format as OCSP can be carried over LDAP, HTTP, SMTP, etc.

Additional PKI-related Protocols

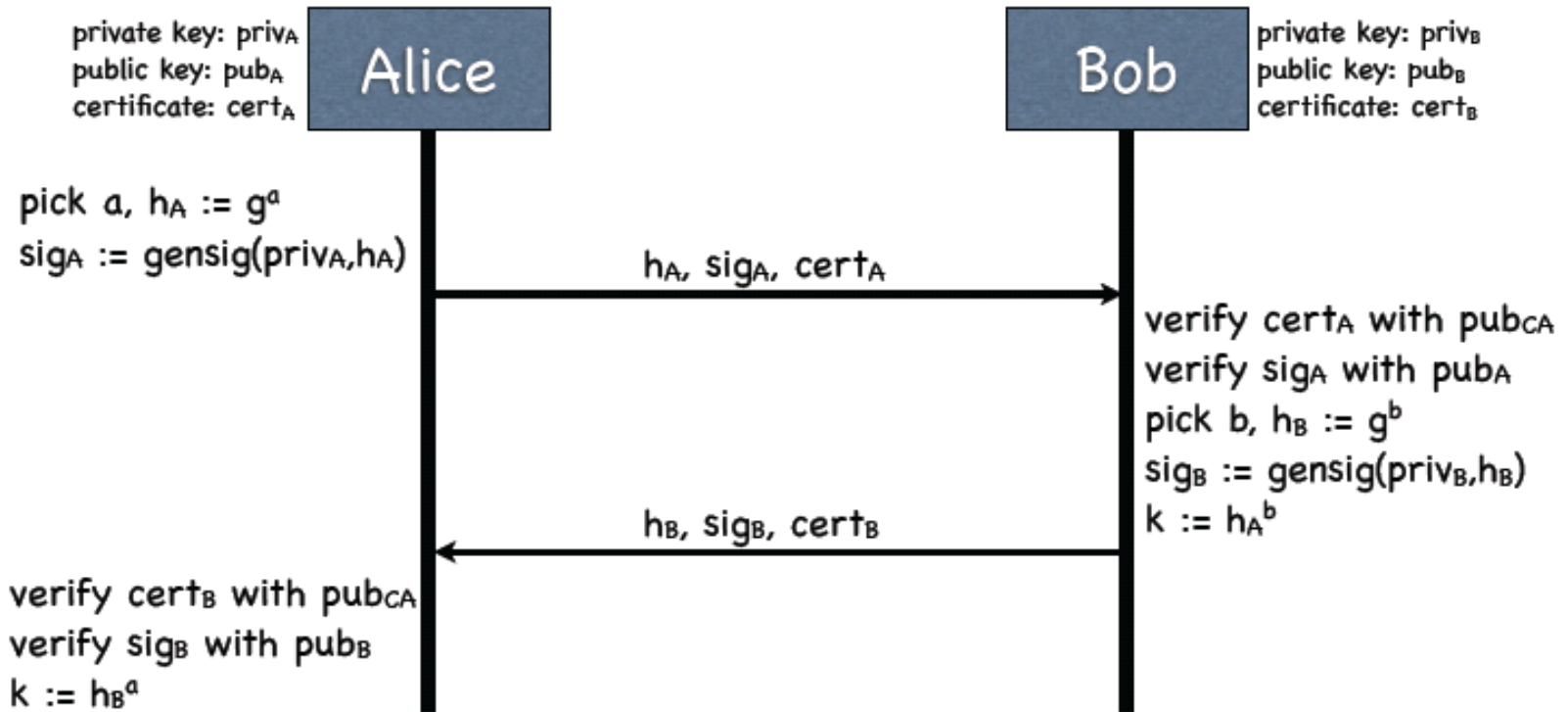
- Lightweight Directory Access Protocol
 - Protocol that specifies how to retrieve certificates and CRLs from a repository storing that information
 - Specifies message formats for
 - reading entries, searching for entries (all clients)
 - adding/deleting entries in the repository (CA only)
- Certificate Management Protocol
 - Specifies message formats for certificate creation, initial distribution, and management
- Certificate Request Message Format
 - Specifies message formats for requesting a certificate for a particular public key from a CA

Using Certificates to Protect Diffie-Hellman – First Try



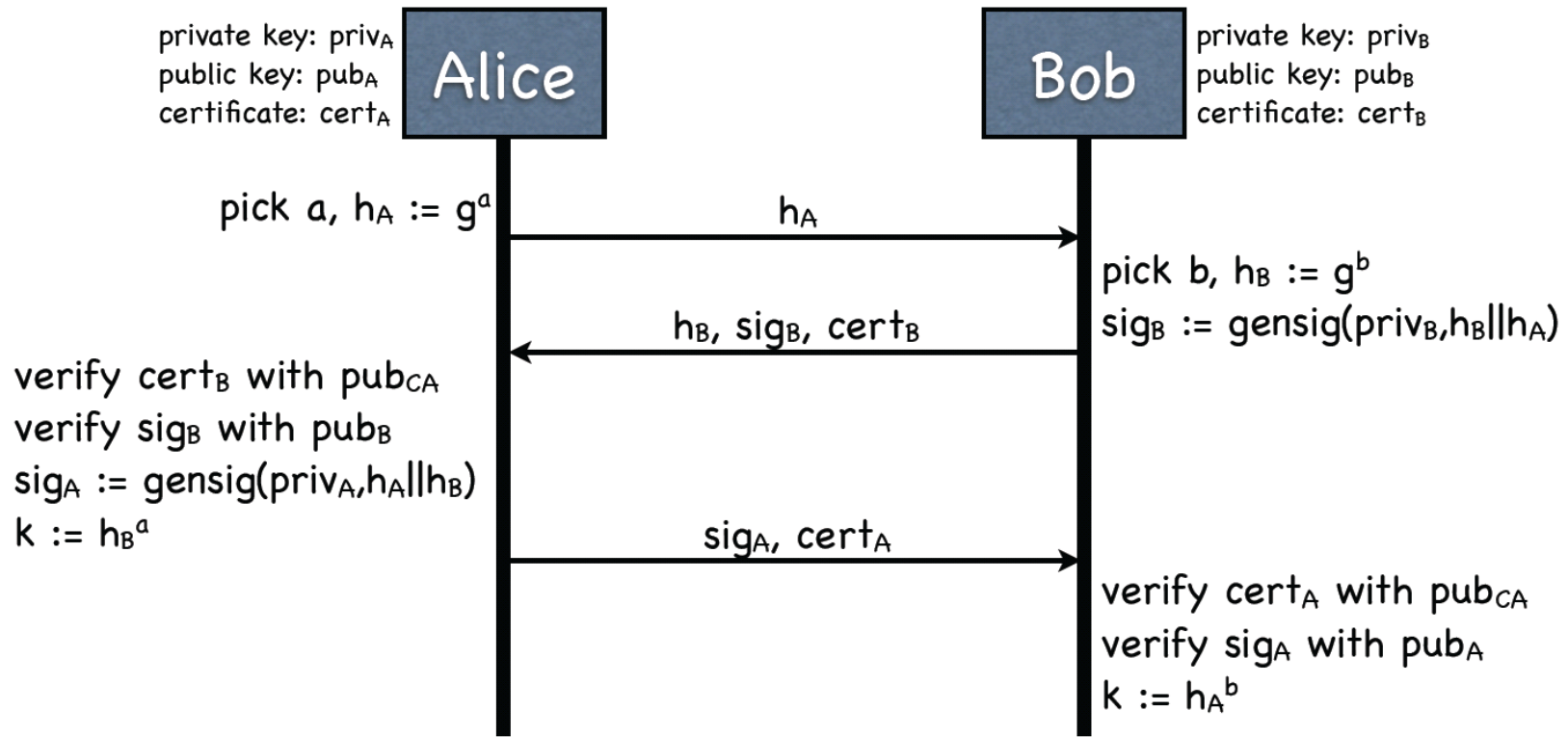
■ Problem?

Using Certificates to Protect Diffie-Hellman – First Try



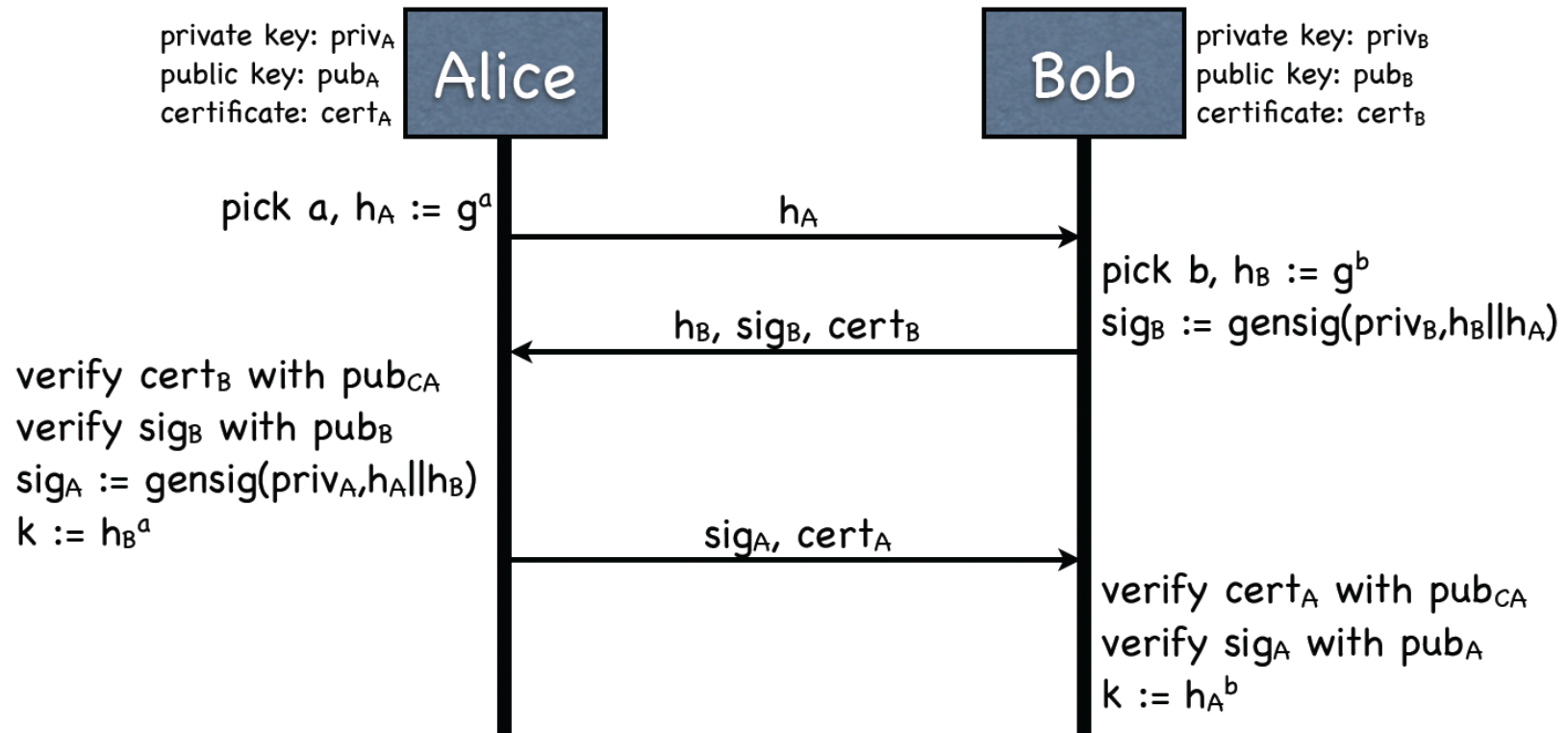
- The signatures are not bound to THIS particular run of the protocol
- An attacker could record e.g. Alice's message and replay it to Bob

Authenticated Diffie Hellman – Second Try



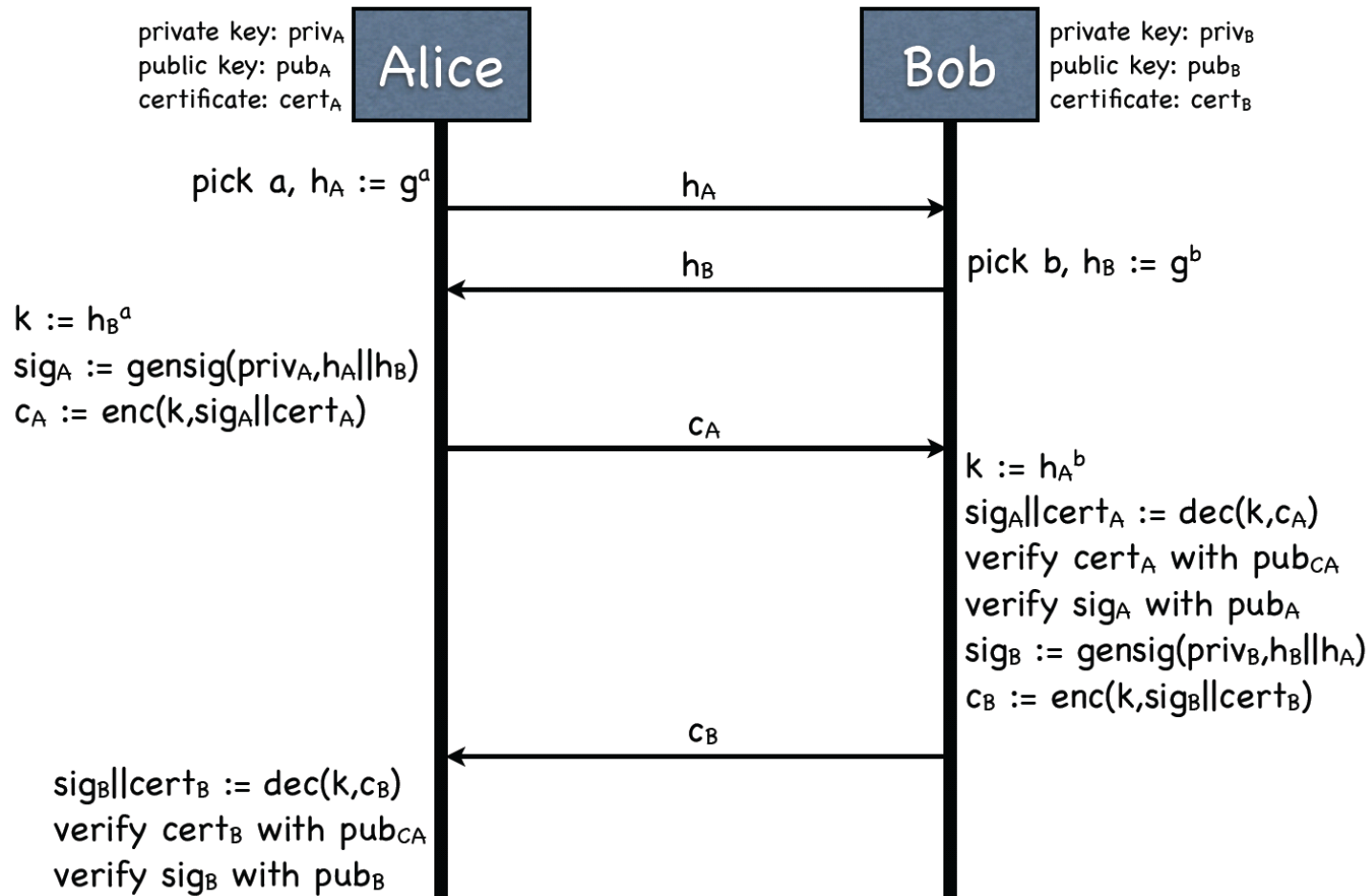
- Both parties sign the public DH values received from each other
- Both parties use public key certificates to exchange their public signature keys
- Now replay is not possible because Alice can check if Bob correctly received her fresh public DH value and vice versa

Authenticated Diffie Hellman – Second Try



- Problem: Attacker Eve can make Alice believe that she is exchanging a key with Eve and not Bob
 - Eve just has to exchange Bob's public key and signature with her own

Secure Authentic Diffie-Hellman



- Bob and Alice confirm that they both agreed upon the correct key by using it to encrypt their certificate and their signature

So...

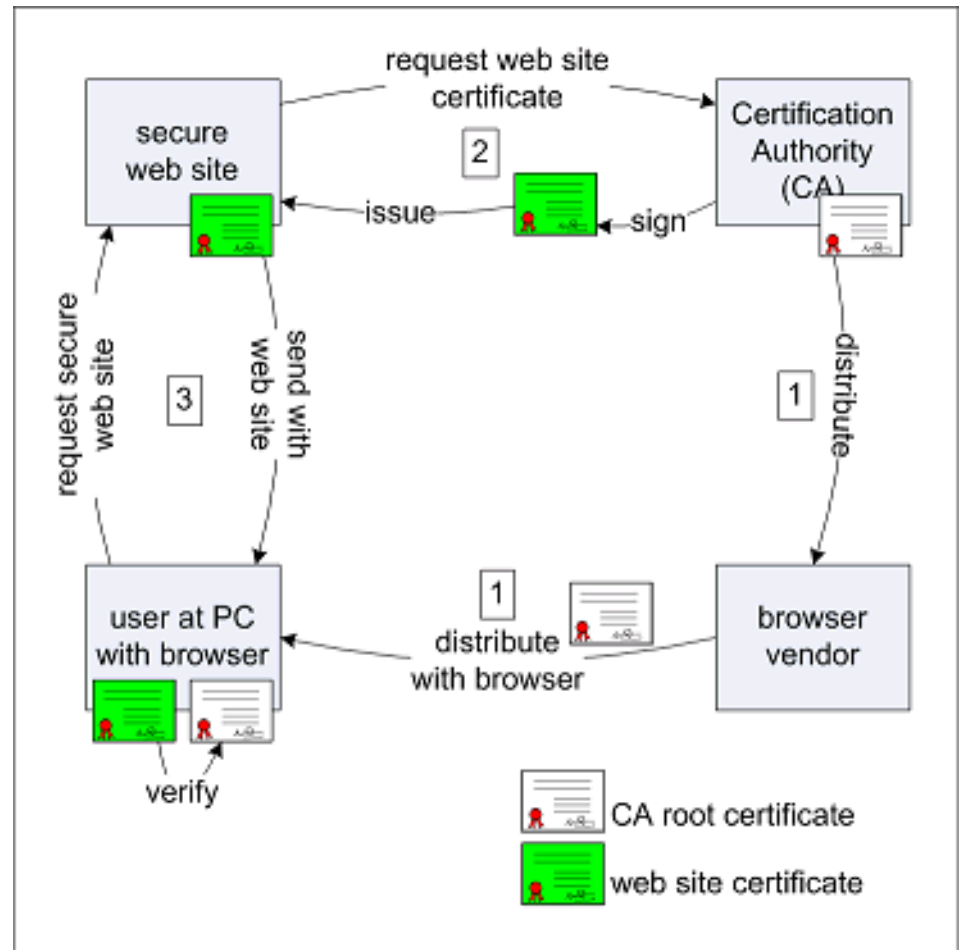
- ... certificates can be used to distribute public keys in an authenticated way
- The content of a certificate is hashed and then signed with the private key of a certification authority
- In Chapter 3 we discussed hash functions and saw that a tool is available online that finds collisions in MD5
- What are the consequences of the MD5 attack for faking certificates?

Typical Way of Obtaining Certificates

- User creates private key
- User creates a Certificate Signing Request
 - Containing
 - User identity
 - Domain name
 - Public key
- CA processes the Certificate Signing Request
 - Includes
 - Validating user identity
 - Validating domain ownership
 - Signs and returns the certificate
- User installs private key and certificate e.g. on a web server

Distribution of Web Site Certificates

- CA root certificates are distributed via browser vendors
- Web servers can request certificates from a CA
- Users can verify certificate presented by web site with the pre-installed CA certificate

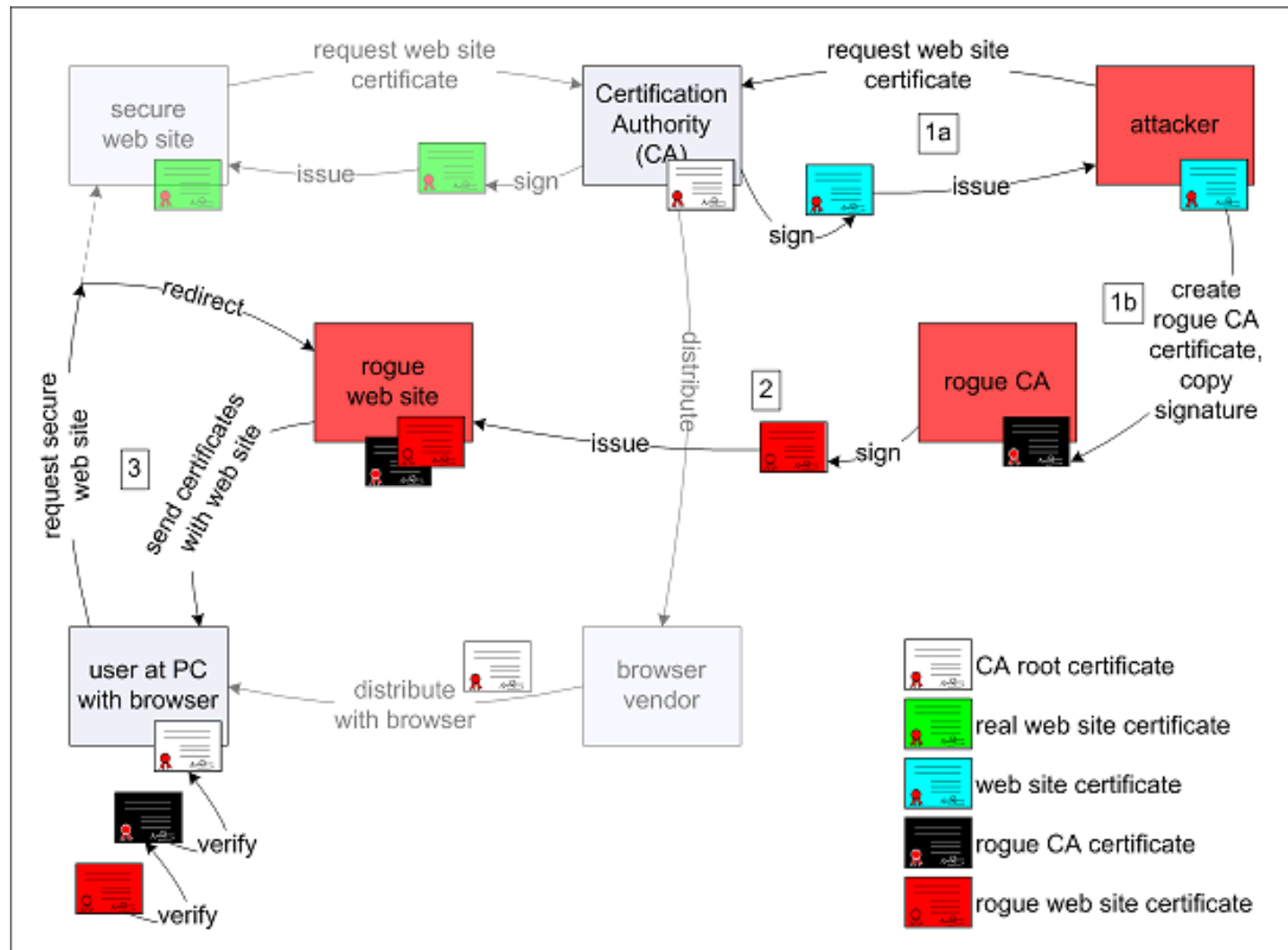


End of 2008

- A group of researcher showed as a proof of concept that
 - They made a regular, browser pre-installed root CA hash, sign and issue a web server certificate to them
 - They constructed a second certificate with the same MD5 hash
- Even worse
 - The second certificate could be constructed to be a CA certificate, i.e. it could further be used to sign other certificates



Overview Of Attack and Potential Consequences



MD5 Collision Attack

- Given any P, P' a pair C, C' can be computed such that

$$\text{MD5}(P \parallel C) = \text{MD5}(P' \parallel C')$$

Approach to Find Collisions

set by the CA, need to be predicted by attacker	serial number	chosen prefix (difference)	serial number
	validity period		validity period
	real subject name		rogue subject name
	real cert RSA key	collision bits (computed)	rogue cert RSA key
	X.509 extensions		X.509 extensions*
	signature	identical bytes (copied from real cert) except for CA=FALSE set to CA = TRUE	signature

The Authors State:

- “With optimizations the attack might be done for \$2000 on Amazon EC2 in 1 day”
 - EC2 = Elastic Compute Cloud
- “We want to prevent malicious entities from repeating the attack:
 - We are not releasing our collision finding implementation or improved methods until we feel it’s safe
 - We’ve talked to the affected CAs: they will switch to SHA-1 very, very soon”

Vulnerable CAs in 2008

- Of 30 000 website certificates
 - 9 000 were signed with MD5
 - 97% of them were issued by RapidSSL
- CAs that still used MD5 in 2008
 - RapidSSL
 - FreeSSL
 - TrustCenter
 - RSA Data Security
 - Thawte
 - Verisign.co.jp
- All of them moved to other hash functions by March 2009

References and Further Reading

- Kaufmann et al., Chapter 15
- RFC 5280: X.509 Certificates and CRLs
- RFC 2560: Online Certificate Revocation Protocol
- RFC 2559: X.509 Public Key Infrastructure Operational Protocols - LDAPv2
- RFC 4210: Certificate Management Protocol (CMP)
- Consequences of MD5 attacks
 - <http://www.win.tue.nl/hashclash/rogue-ca/#secPres>
 - Figures taken from there!