

GTU Department of Computer Engineering  
CSE 222 - Homework 8 - Report

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# Complexity Analysis

## BFS Algorithm

It is defined in terms of the number of nodes ( $V$ ) and edges ( $E$ ) in the graph.

**Time Complexity:  $O(V + E)$**

The algorithm visits each vertex once and moves through all adjacent vertices via the edges. Here  $V$  represents the number of vertices and  $E$  represents the number of edges in the graph.

It takes  $O(V)$  time to visit each vertex

It takes  $O(E)$  time to check each edge

So the total time complexity of the BFS algorithm is  $O(V + E)$ .

**Space Complexity:  $O(V)$**

The algorithm maintains a queue data structure, the cluster visited, and a hashmap.

In the worst-case scenario, all vertices are added to the queue and the visited cluster, so this takes  $O(V)$  space. Each entry in the HashMap prev takes up fixed space and in the worst case it takes up  $O(V)$  space as it will grab all the vertices. Therefore, the total area complexity is  $O(V)$ .

## Dijkstra Algorithm

*Checking Initialization and Starting Point:* The complexity here is fixed, i.e.  $O(1)$ , because we are just creating a few data structures and performing a check on the start and end nodes.

*Main Loop:* This loop runs until the priority queue is empty. Each operation inside the loop has a different complexity:

*Querying from the queue:* This operation takes  $O(\log V)$  time because polling on a binary stack takes logarithmic time.

*Checking and adding visited cluster:* This takes  $O(1)$  time.

*Checking Neighbors:* Checking neighbors is actually going around all the edges of the node and collectively takes  $O(E)$  time as it is done for each node.

*Updating and queuing distances:* These operations take  $O(\log V)$  time because they include operations in the priority queue.

*Path reconfiguration:* This process takes  $O(V)$  time in the worst case scenario when the path includes all nodes.

*Writing the path to a file:* This takes  $O(V)$  time as we iterate over every node in the path.

Summing all this up, **the time complexity of Dijkstras Algorithm is  $O((V + E) \log V)$** , assuming the priority queue uses a binary heap data structure. The  $E$  term results from visiting all edges of each vertex, and the  $V$  term originates from operations in the priority queue.

space complexity:

The visited cluster, distance map, and previous node map can each store all vertices, so each takes up  $O(V)$  space. The queue can also grab all the vertices in the worst-case scenario, which takes up  $O(V)$  space.

The final path list holds at most all vertices, which is  $O(V)$ .

Hence, **the space complexity of Dijkstra's Algorithm is  $O(V)$** .

## Running time performance

```
long endTime = System.nanoTime();  
duration = endTime - startTime;  
System.out.println("BFS: duration: "+duration);  
return path;
```

Working times were measured with time variables added to the code.

```
<terminated> Main (0) [Java Application] C:\Users\Dogukan\AppData\Local\Temp\p  
Dijkstra: duration: 1104589900  
BFS: duration: 382005000
```

(nanosecond)

As it can be seen, it provides a faster result in accordance with the theoretical calculation.

## Outputs:

**Green Dijkstra** and **Blue BFS**



