DATA STRUCTURE AND ALGORITHM

CLASS 11

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2017-03-06

Sorting

Sorting

Types of Sorting

- Internal Sorting The list is small enough to sort entirely in main memory
- External Sorting The list is too large to sort in main memory

Internal sorting algorithms

- Bubble sort
 - Selection sort
 - Insertion sort
 - Quick sort
 - Shell sort
 - Heap sort
 - Bucket sort
 - Radix sort



Bubble sort

```
repeat until no swaps
for i from 0 to n-2
if ith and i+1th elements out of order
swap them
```

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initial	15	4	8	3	50	9	20	unsorted array
1 st run	4	15	8	3	50	9	20	cmp 1 and 2, swap

Data Structure and Algorithm

initial	15	4	8	3	50	9	20	unsorted array
1 st run	4	15	8	3	50	9	20	cmp 1 and 2, swap
1 st run	4	8	15	3	50	9	20	cmp 2 and 3, swap

2017-03-06

initial	15	4	8	3	50	9	20	unsorted array
1 st run	4	15	8	3	50	9	20	cmp 1 and 2, swap
1 st run	4	8	15	3	50	9	20	cmp 2 and 3, swap
1 st run	4	8	3	15	50	9	20	cmp 3 and 4, swap

initial	15	4	8	3	50	9	20	unsorted array
1^{st} run	4	15	8	3	50	9	20	cmp 1 and 2, swap
1^{st} run	4	8	15	3	50	9	20	cmp 2 and 3, swap
1^{st} run	4	8	3	15	50	9	20	cmp 3 and 4, swap
1 st run	4	8	3	15	50	9	20	cmp 4 and 5, no swap

initial	15	4	8	3	50	9	20	unsorted array
1 st run	4	15	8	3	50	9	20	cmp 1 and 2, swap
1 st run	4	8	15	3	50	9	20	cmp 2 and 3, swap
1 st run	4	8	3	15	50	9	20	cmp 3 and 4, swap
1 st run	4	8	3	15	50	9	20	cmp 4 and 5, no swap
1 st run	4	8	3	15	9	50	20	cmp 5 and 6, swap

initial	15	4	8	3	50	9	20	unsorted array
1 st run	4	15	8	3	50	9	20	cmp 1 and 2, swap
1 st run	4	8	15	3	50	9	20	cmp 2 and 3, swap
1 st run	4	8	3	15	50	9	20	cmp 3 and 4, swap
1 st run	4	8	3	15	50	9	20	cmp 4 and 5, no swap
1 st run	4	8	3	15	9	50	20	cmp 5 and 6, swap
1 st run	4	8	3	15	9	20	50	cmp 6 and 7, swap

initial	15	4	8	3	50	9	20	unsorted array
1 st run	4	15	8	3	50	9	20	cmp 1 and 2, swap
1 st run	4	8	15	3	50	9	20	cmp 2 and 3, swap
1 st run	4	8	3	15	50	9	20	cmp 3 and 4, swap
1 st run	4	8	3	15	50	9	20	cmp 4 and 5, no swap
1 st run	4	8	3	15	9	50	20	cmp 5 and 6, swap
1 st run	4	8	3	15	9	20	50	cmp 6 and 7, swap
2 nd run	4	8	3	15	9	20	50	cmp 1 and 2, no swap
2 nd run	4	3	8	15	9	20	50	cmp 2 and 3, swap
2 nd run	4	3	8	15	9	20	50	cmp 3 and 4, no swap
2 nd run	4	3	8	15	9	20	50	cmp 4 and 5, no swap
2 nd run	4	3	8	9	15	2 0	50	cmp 5 and 6, swap
3 rd run	3	4	8	9	15	2 0	50	1, 2 swap
4^{th} run	3	4	8	9	15	2 0	50	no swap
5 th run	3	4	8	9	15	20	50	no swap
6 th run	3	4	8	9	15	20	50	no swap
Data Structure and Algorithm				S	ΙL	' '	2017-03	3-06 7 / 40

Bubble sort: the code (codes/bubble.c) I

```
// Unoptimized and Optimized implementation of bubble sort
   #include <stdio.h>
   void swap(int *x, int *y);
   void bubble_unoptimized(int list[], int size);
   void bubble(int list[], int size);
   void printArray(int list[], int size);
8
   int main(){
       int list[] = \{30, 28, 18, 29, 48, 40\};
      // get the number of items in the list
11
       int n = sizeof(list)/sizeof(list[o]);
       printf("initial list: \t");
       printArray(list, n);
15
```

Bubble sort: the code (codes/bubble.c) II

```
// bubble_unoptimized(list, n);
   bubble(list, n);
   printf("sorted list: \t");
   printArray(list, n);
   return o;
void swap(int *x, int *y){ // swap the two value
   int temp = *x;
   *x = *y;
   *y = temp;
```

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23 24 25

26

27

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Bubble sort: the code (codes/bubble.c) III

```
32
    void printArray(int list[], int size){
       for (int i = 0; i < size; i++)
           printf("%d ", list[i]);
       printf("\n");
38
39
   // unoptimized bubble sort with worst case of O(n^2)
   void bubble unoptimized(int list[], int size){
       int i, j;
       for (i = 0; i < size-1; i++)
          for (j = 0; j < size-i-1; j++) // skip sorted items in the list
              if (list[j] > list[j+1])
                  swap(\&list[i], \&list[i+1]);
```

Bubble sort: the code (codes/bubble.c) IV

```
// optimzed bubble sort
   void bubble(int list[], int size){
       int i, j;
       int swapped = 1;
      // loop until no two elements were swapped by inner loop
       for (i = size - 1; swapped > 0; i - -)
          swapped = 0;
          for (j = 0; j < i; j++)
              if (list[i] > list[i+1]){
57
                 swap(\&list[i], \&list[i+1]);
                 swapped = 1;
```

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Bubble sort: Time complexity

loop	number	of	com-
	parison		
1	n-1		
2	n-2		
3	n-3		
•	•		
•	•		
n-1	1		

Time complexity: Worst case $O(n^2)$

SELECTION SORT

- $_{1}$ for i from 0 to n-1
- find smallest element between ith and n-1th
- 3 swap smallest with ith element

init	ial	15	4	8	3	50	9	20	unsorted array
1^{st}	run	3	4	8	15	50	9	20	midx 15, min 3, swap

initial	15	4	8	3	50	9	20	unsorted array
1^{st} run	3	4	8	15	50	9	20	midx 15, min 3, swap
2 nd run	3	4	8	15	50	9	20	midx 4, min 4

initial	15	4	8	3	50	9	20	unsorted array
1 st run	3	4	8	15	50	9	20	midx 15, min 3, swap
2 nd run	3	4	8	15	50	9	20	midx 4, min 4
3 rd run	3	4	8	15	50	9	20	midx 8, min 8

initial	15	4	8	3	50	9	20	unsorted array
1 st run	3	4	8	15	50	9	20	midx 15, min 3, swap
2 nd run	3	4	8	15	50	9	20	midx 4, min 4
3 rd run	3	4	8	15	50	9	20	midx 8, min 8
4 th run	3	4	8	9	50	15	20	midx 15, min 9, swap

initial	15	4	8	3	50	9	20	unsorted array
1 st run	3	4	8	15	50	9	20	midx 15, min 3, swap
2 nd run	3	4	8	15	50	9	20	midx 4, min 4
3 rd run	3	4	8	15	50	9	20	midx 8, min 8
4 th run	3	4	8	9	50	15	20	midx 15, min 9, swap
4 th run	3	4	8	9	15	50	20	midx 50, min 15, swap

initial	15	4	8	3	50	9	20	unsorted array
1 st run	3	4	8	15	50	9	20	midx 15, min 3, swap
2 nd run	3	4	8	15	50	9	20	midx 4, min 4
3 rd run	3	4	8	15	50	9	20	midx 8, min 8
4 th run	3	4	8	9	50	15	20	midx 15, min 9, swap
4 th run	3	4	8	9	15	50	20	midx 50, min 15, swap
5 th run	3	4	8	9	15	20	50	midx 50, min 20, swap

Selection sort: the Code (codes/selection.c)

```
// selection sort
  void selection(int list[], int size){
    int i, j, midx;
    for (i = 0; i < size-1; i++){}
       // find the minimum element in unsorted list
       midx= i:
       for (j = i+1; j < size; j++)
          if (list[j] < list[midx])</pre>
             midx = j;
       // swap the minimum element with the first element
       swap(&list[midx], &list[i]);
time complexity: O(n^2)
```

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Data Structure and Algorithm



for i from 1 to n-1
call 0th through i-1th elements the ''sorted side''
remove ith element
insert it into sorted side in order

	idx	1	2	3	4	5	
		36	28	30	17	23	initial state
_	1	28	36	30	17	23	28 is smaller than 36, move 28 to the sorted side

idx	1	2	3	4	5	
	36	28	30	17	23	initial state
1	28	36	30	17	23	28 is smaller than 36, move 28 to the sorted side
2	28	30	36	17	23	30 is smaller than 36, move it to the sorted side

idx	1	2	3	4	5	
	36	28	30	17	23	initial state
1	28	36	30	17	23	28 is smaller than 36, move 28 to the sorted side
2	28	30	36	17	23	30 is smaller than 36, move it to the sorted side
3	17	28	30	36	23	17 is smaller than 36, 30, and 28, move it to the
						sorted side

idx	1	2	3	4	5	
	36	28	30	17	23	initial state
1	28	36	30	17	23	28 is smaller than 36, move 28 to the sorted side
2	28	30	36	17	23	30 is smaller than 36, move it to the sorted side
3	17	28	30	36	23	17 is smaller than 36, 30, and 28, move it to the
						sorted side
4	17	23	28	30	36	23 is smaller than 36, 30, and 28, move it to the
						sorted side

Insertion sort: the Code (codes/insertion.c)

```
// insertion sort
   void insertion(int list[], int size){
      int i, j, key;
31
     for (i = 1; i < size; i++){
        kev = list[i];
        // if item is greater than the key
35
        // move the item in list to one position ahead
        for (j = i-1; j \ge 0 \&\& list[j] > key; j--)
37
           list[j+1] = list[j];
        list[j+1] = key;
39
```

time complexity: worst case $O(n^2)$, best case O(n)



Shell sort

Shell sort is based on insertion sort algorithm to reduce the number of shifts. For example, if the list in decreasing order and we want to sort in increasing order, the far left record has to be moved to the far left, one by one.

QUICK SORT

Divide and Conquer

- two phase
- split and control

using recursion

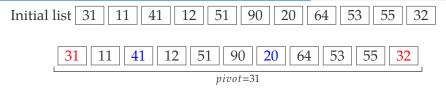
stack is needed

Best average time

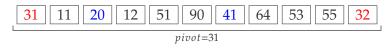
 \bigcirc $O(n \cdot log_2n)$

Quick sort: The principle

- 1. Pick a "**pivot**" element
- 2. "partition" the array into three parts
 - 1st part: all elements in this is less than the pivot
 - 2nd part: the pivot itself (only one element)
 - \circ 3rd part: all elements in this part is greater than or equal to the pivot
- 3. Recursively apply quick sort to the first and the third part
 - swap the out of order entries from the first and the thir part



- examine the first, and last item of the full list that is 31 (first, pivot), and 32 (last)
- \odot search forward until we find an item larger than the pivot 41 > 31
- \odot search backward until we find an item smaller than the pivot 20 < 31
- swap the two items



- orepeat the process until the left and right indices cross each other
- then swap the item with the pivot

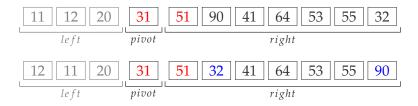


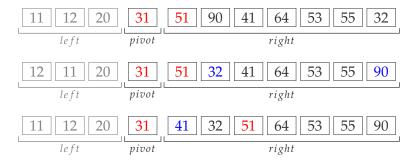
repeat the quick sort for the left and right part

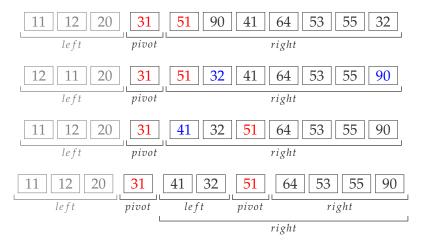
repeat the quick sort for the left and right part

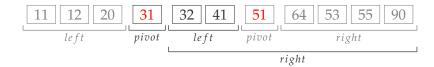
repeat the quick sort for the left and right part

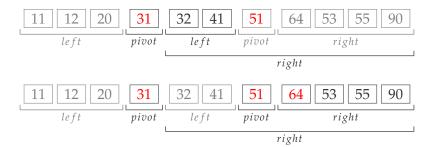


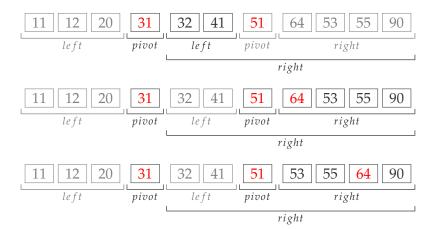


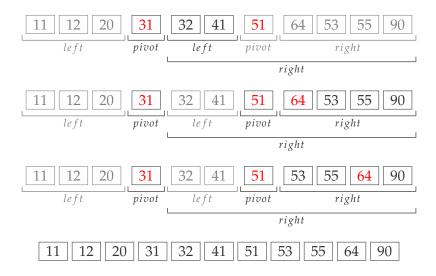












Quick Sort: the Code

```
void quicksort(int list[], int left, int right){
if (left < right){
   int pivot = partition(list, left, right);
   quicksort(list, left, pivot-1);
   quicksort(list, pivot+1, right);
}
</pre>
```

Quick Sort: the Code

```
int partition(int list[], int left, int right){
42
      int pivot, temp;
43
      int low, high;
44
      low = left;
      high = right +1;
      pivot = list[left];
47
      do{
          do
             low++:
50
          while(low <= right && list[low] < pivot);</pre>
51
          do
52
             high--;
53
          while (high >= left && list[high] > pivot);
54
          if (low < high)
55
             swap(list[low], list[high], temp);
      } while (low < high);</pre>
57
swap(list[left], list[high], temp);
Data Structure and Algorithm, list[high], temp);
                                                            32 / 40
       return high:
59
```

Quick Sort: time complexity

Average case: $O(n \cdot log_2 n)$, when elements can be split into "equal size" Let T(n) be average time to sort n records

$$T(n) = c \cdot n + 2 \cdot T(n/2)$$

$$\leq c \cdot n + 2 \cdot (c \cdot n/2 + 2 \cdot T(n/4))$$

$$\leq 2 \cdot c \cdot n + 4 \cdot T(n/4))$$

$$\cdots$$

$$\leq c \cdot n \cdot \log_2 n + n \cdot T(1))$$

$$O(n \cdot \log_2 n)$$

Worst case: $O(n^2)$

Other approaches:

- http://me.dt.in.th/page/Quicksort/
- $\bigcirc \ \ https://ece.uwaterloo.ca/\sim cmoreno/ece250/quick-sort-complete-example.pdf$



BUCKET SORT

RADIX SORT