



# Search

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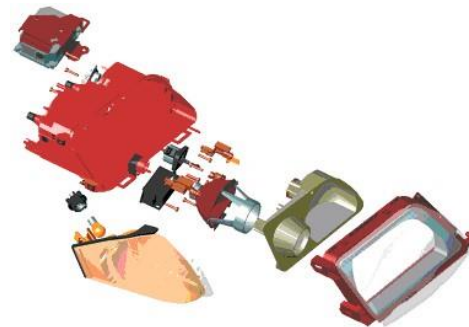
Russell chap 3, 4, 5

Luger chap. 3

# Search

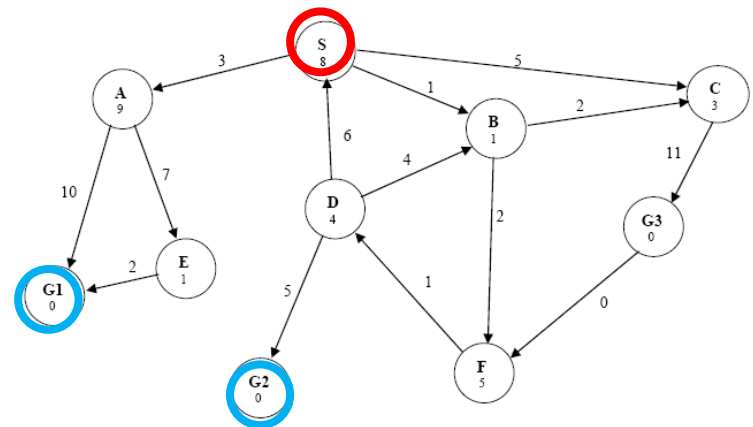
## ■ Problem solving

- Searching among alternative choices to **find a sequence of actions** by which a goal can be achieved
- Example
  - Route finding : finding a sequence of road
  - Assembly planning : finding a sequence of part assembly
  - Playing game : finding a sequence of move
  - Inference in FOPC : finding a sequence of applying M.P.



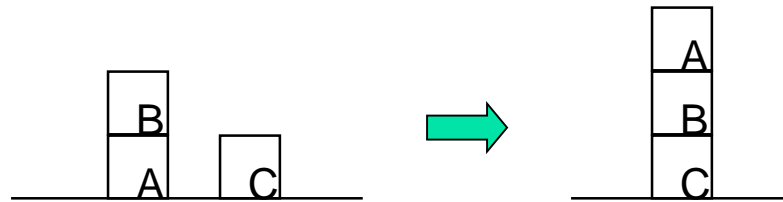
# State Space

- State space representation
  - Represent problem space as a *graph*
  - Nodes: states
  - Arcs: Operations to other states
- State space
  - [N, A, S, G]
    - N – nodes(states)
    - A – arcs(operations)
    - S – start state
    - G – goal state



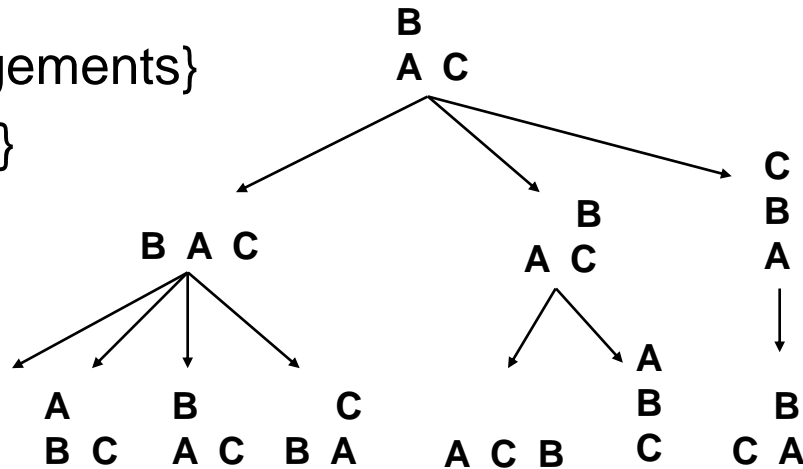
# State Space

- Planning in a blocks world



$N = \{\text{different arrangements}\}$

$A = \{\text{moving a block}\}$



# State Space

## ■ 8-Puzzle

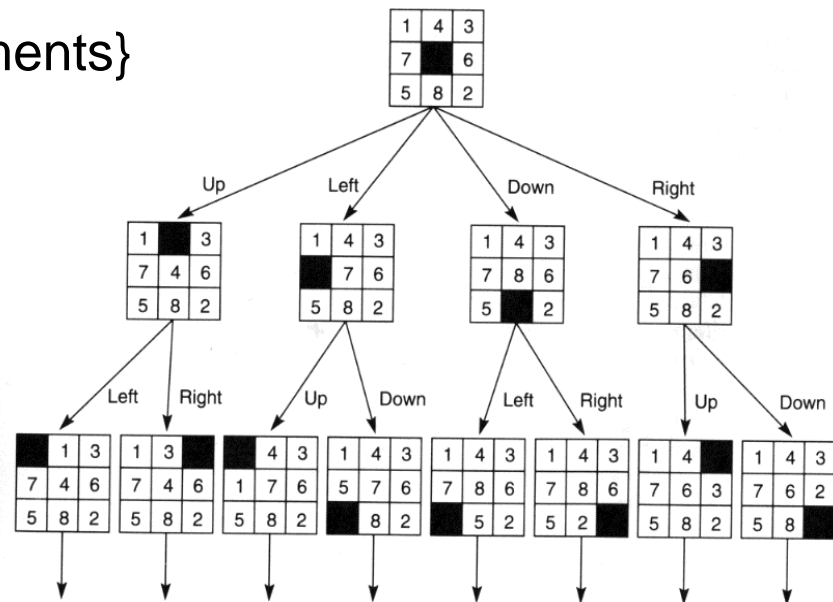
1	4	3
7		6
5	8	2



1	2	3
8		4
7	6	5

$N = \{\text{different arrangements}\}$

$A = \{\text{moving the blank}\}$



**Figure 3.6** State space of the 8-puzzle generated by "move blank" operations.



# State Space

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## ■ Inference

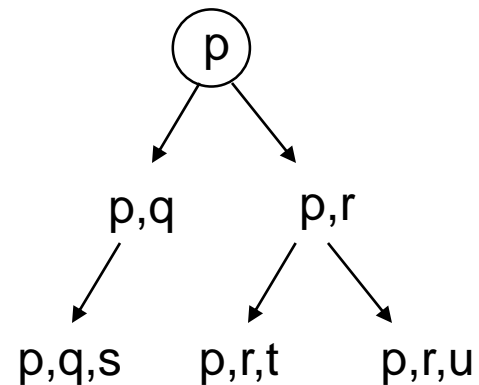
$p \Rightarrow q, p \Rightarrow r, q \Rightarrow s,$   
 $r \Rightarrow t, r \Rightarrow u, p$



$t ?$


$N = \{\text{true sentences}\}$

$A = \{\text{applying M.P.}\}$



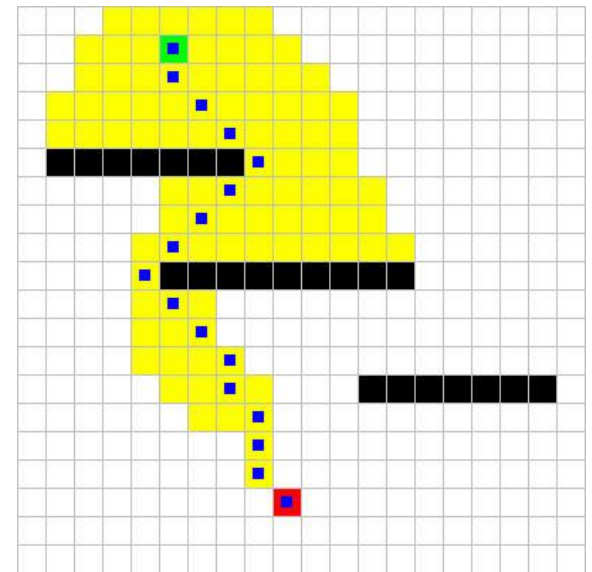
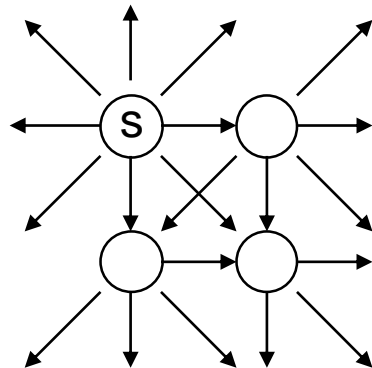
# State Space

- Route finding

start location  goal location

$N = \{\text{different locations}\}$

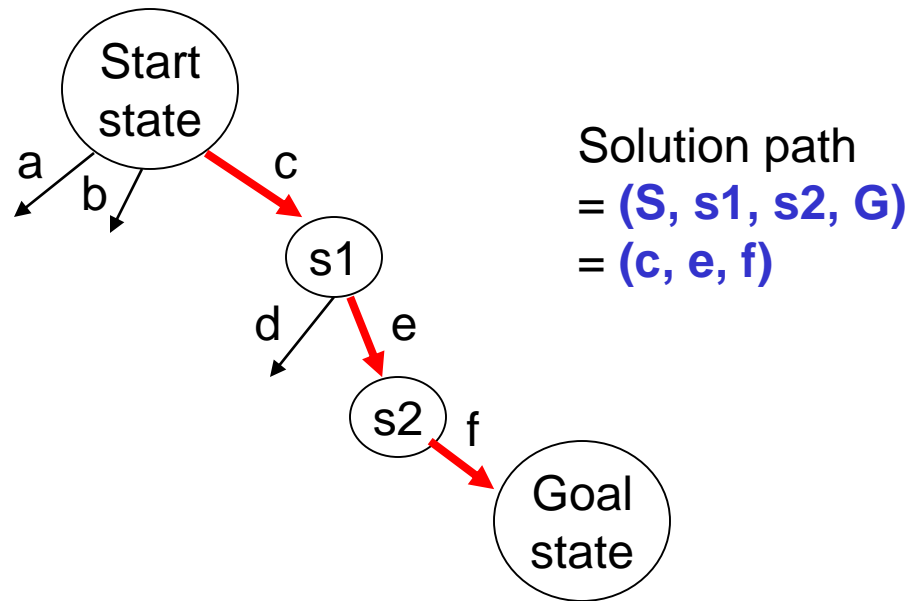
$A = \{\text{move to the adjacent location}\}$



# State Space Search

- Search

- Finding a solution path from S to G







# Depth-First Search

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- Goes deeper whenever possible

**open** = {START}, **closed** =  $\emptyset$

while (open  $\neq \emptyset$ )

    remove leftmost state from **open**  $\rightarrow X$

    if (X is GOAL) return (success)

    else

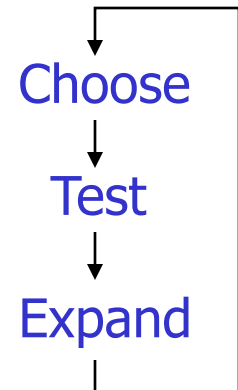
        generate children of X

        put X into **closed**

        eliminate child if it is in open or close

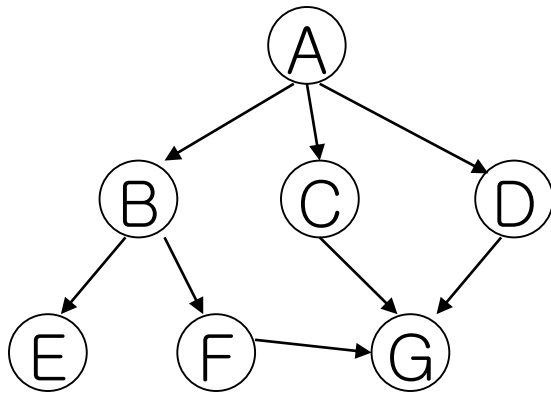
        put remaining children into left of **open**(stack)

return (fail)



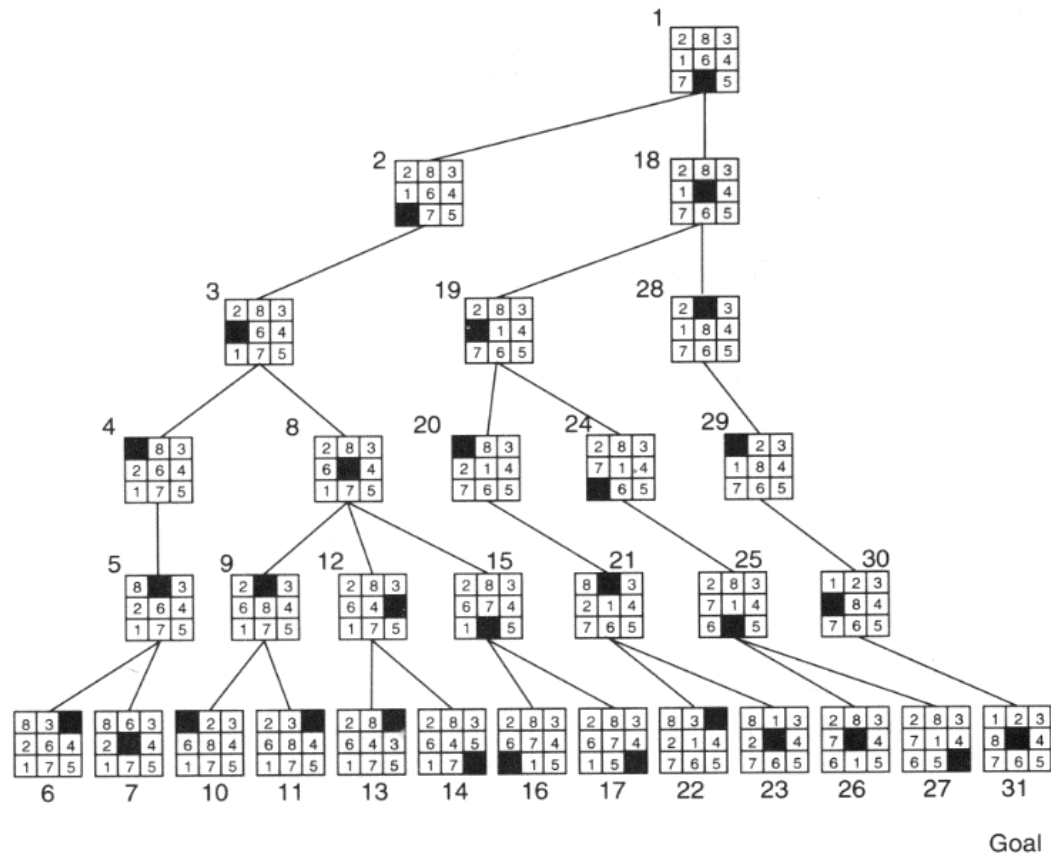
# Depth-First Search

## ■ Example



1. open = { }, closed = { }
2. X = , open = { }, closed = { }
3. X = , open = { }, closed = { }
4. X = , open = { }, closed = { }
5. X = , open = { }, closed = { }
- 6.

# Depth-First Search



**Figure 3.17** Depth-first search of the 8-puzzle with a depth bound of 5.



# Breadth-First Search

---

- Explore search space level by level

**open** = {START}, **closed** =  $\emptyset$

while (open  $\neq \emptyset$ )

    remove leftmost state from **open**  $\rightarrow X$

    if (X is GOAL) return (success)

    else

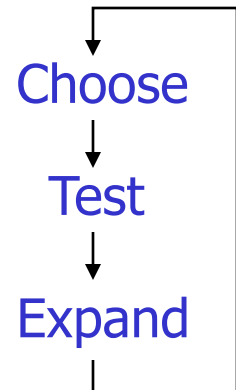
        generate children of X

        put X into **closed**

        eliminate child if it is in open or close

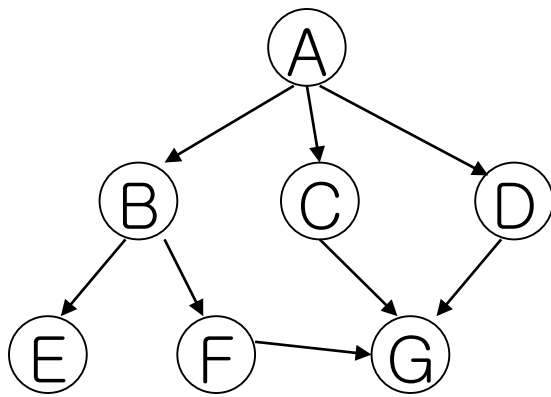
        put remaining children into right of **open**(queue)

return (fail)



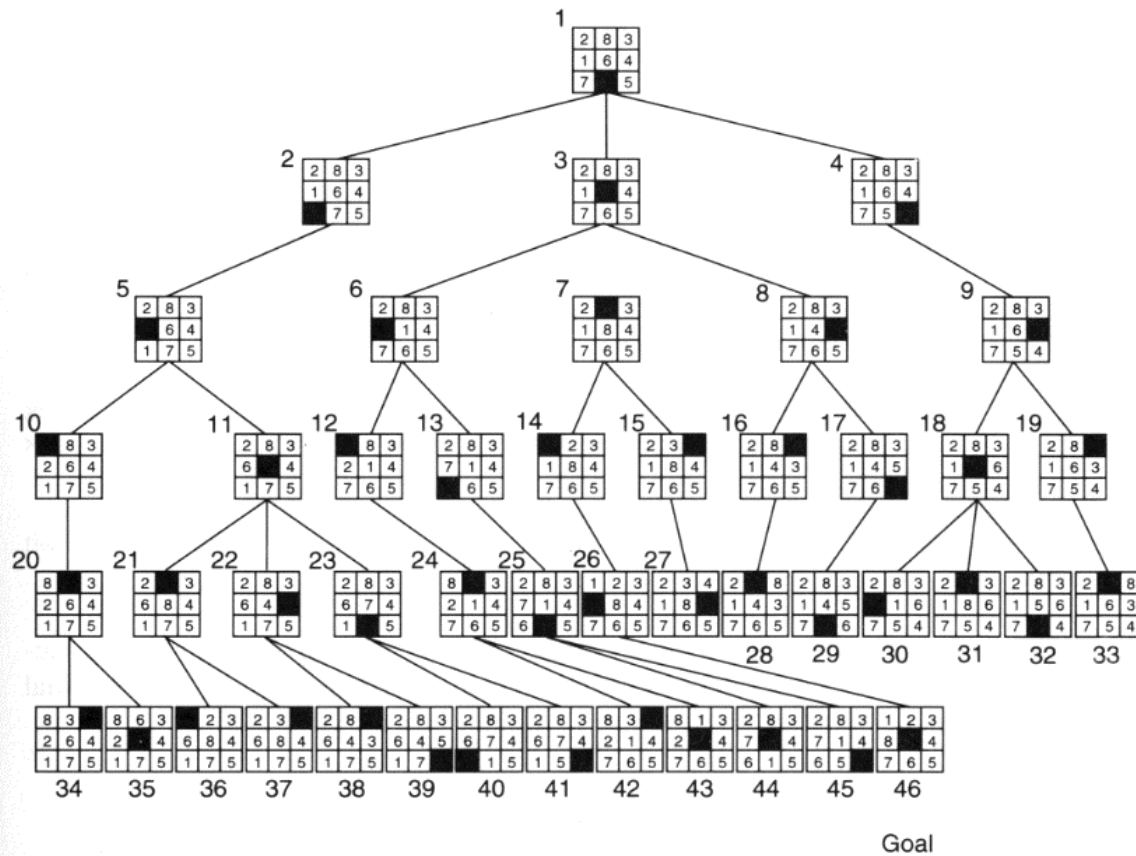
# Breadth-First Search

## ■ Example



1. open = { }, closed = { }
2. X = , open = { }, closed = { }
3. X = , open = { }, closed = { }
4. X = , open = { }, closed = { }
5. X = , open = { }, closed = { }
6. X = , open = { }, closed = { }
7. X = , open = { }, closed = { }
- 8.

# Breadth-First Search



**Figure 3.15** Breadth-first search of the 8-puzzle, showing order in which states were removed from open.



# Solution Path

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- Store parent information with the states
- Example
  - $\text{closed} = \{ (A,0), (B,A), (C,A), (D,A), (E,B), (F,B) \}$
  - $\text{Goal} = (G, C)$

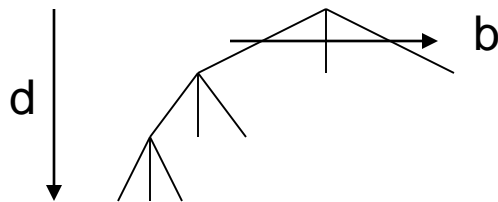
→  $G \leftarrow C \leftarrow A$

# DFS vs. BFS

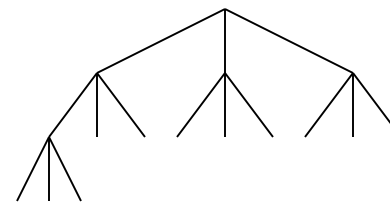
## ■ Let

- $b$ : Average branching factor
- $d$ : Search depth

	OPEN	Time	Space	Optimal
DFS	stack	$O(b^d)$	$O(b \cdot d)$	No
BFS	queue	$O(b^d)$	$O(b^d)$	Yes



■ DFS:  $b + b + \dots + b$



BFS:  $b + b^2 + \dots + b^{d-1}$





# Heuristics

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- Uninformed search (DFS, BFS)
  - Impractical for large state space
    - 15-puzzle:  $\sim 10^{12}$ , Chess:  $\sim 10^{120}$ , Baduk:  $\sim 10^{360}$
  - Inefficient in terms of time and space
- Informed search
  - Use *heuristics*
    - Rules for choosing branches in a state space
    - Guide search to states that are most likely lead to the goal



*Informed search, heuristic search*

# Heuristics

- Example
  - State to search next?

<table><tr><td>2</td><td>8</td><td>3</td></tr><tr><td>1</td><td>6</td><td>4</td></tr><tr><td></td><td>7</td><td>5</td></tr></table>	2	8	3	1	6	4		7	5	5	6	0
2	8	3										
1	6	4										
	7	5										
<table><tr><td>2</td><td>8</td><td>3</td></tr><tr><td>1</td><td></td><td>4</td></tr><tr><td>7</td><td>6</td><td>5</td></tr></table>	2	8	3	1		4	7	6	5	3	4	0
2	8	3										
1		4										
7	6	5										
<table><tr><td>2</td><td>8</td><td>3</td></tr><tr><td>1</td><td>6</td><td>4</td></tr><tr><td>7</td><td>5</td><td></td></tr></table>	2	8	3	1	6	4	7	5		5	6	0
2	8	3										
1	6	4										
7	5											
	Tiles out of place	Sum of distances out of place	2 x the number of direct tile reversals									

1	2	3
8	<b>7</b>	4
7	6	5

Goal

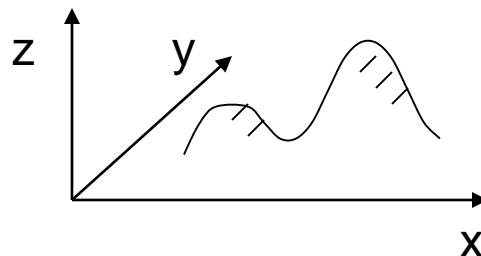
# Hill Climbing

## ■ Algorithm

1. Expand current state
2. Select best child state (use heuristics)
3. Stop if current state is best, or goto 1

## ■ Problem

- It may stuck at local maximum
  - Find  $\max z = f(x, y)$  (S: (1,1), G: unknown)

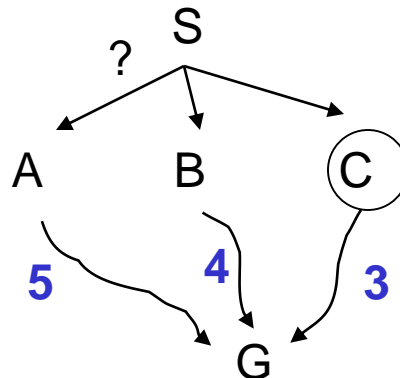


# Best-First Search

- Best-first

- Use heuristic value  $f(n)$  to evaluate each state
- Choose most promising state from OPEN

→ If heuristic value estimate the distance (cost) to the goal, choose smallest one





# Best-First Search

**open** = {START}, **closed** =  $\emptyset$

while (**open**  $\neq \emptyset$ )

    remove leftmost state from **open**  $\rightarrow X$

    if (X is GOAL) return (success)

    else

        generate and evaluate children of X

        put X into **closed**

        for each child C

            if C is in **open**      update path

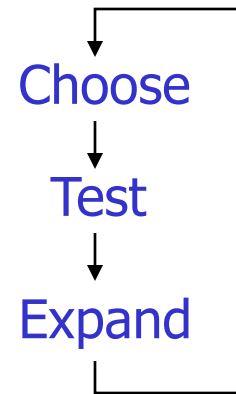
            if C is in **closed** and reached by shorter path

                remove C from **closed**, put into **open**

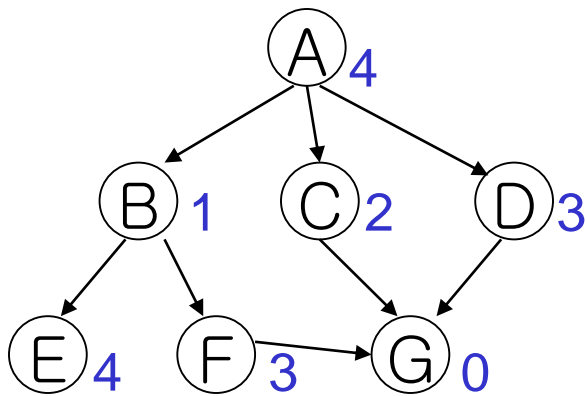
            else                  put C into **open**

        reorder **open**(priority queue)

return (fail)

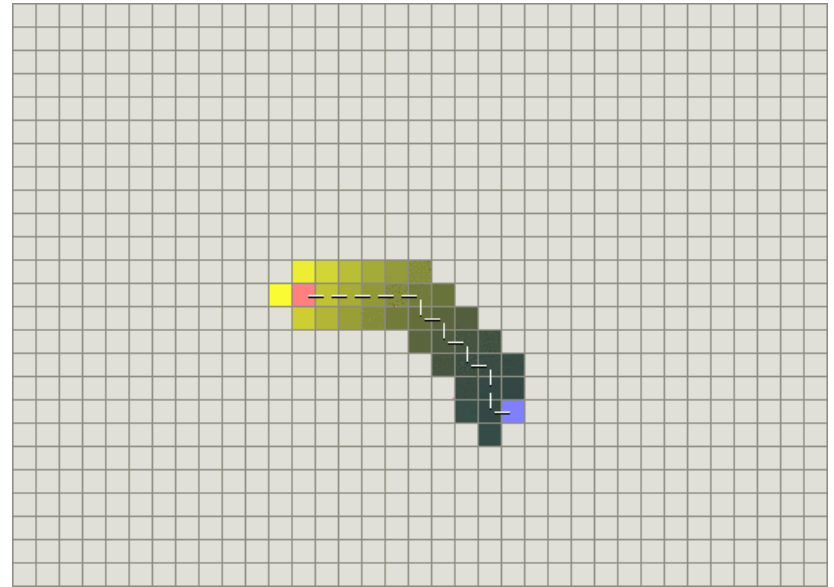
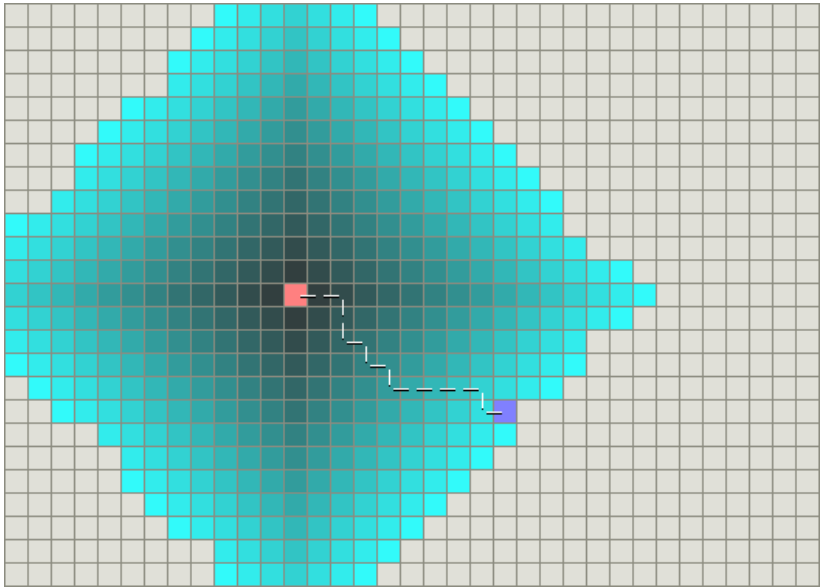


# Best-First Search



1. open = { },  
closed = { }
2. X = , open = { },  
closed = { }
3. X = , open = { },  
closed = { }
4. X = , open = { },  
closed = { }
- 5.

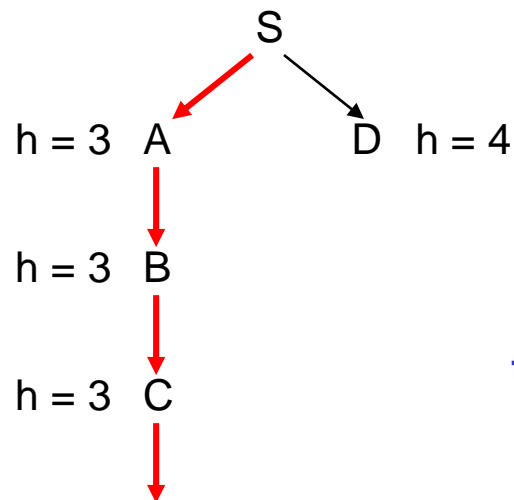
# BFS vs. Best-First



# Evaluation Function

- Heuristics are fallible

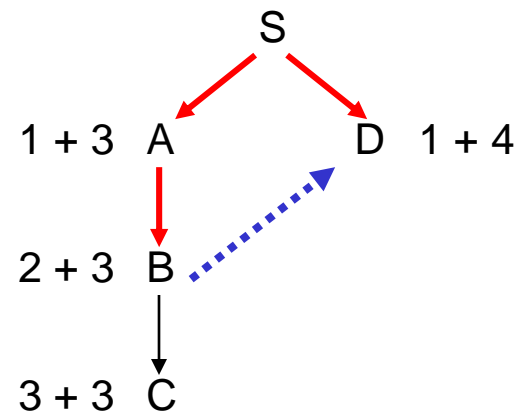
- We want to find good path (not only fast search)  
→ use (heuristic value to G + actual path length from S)



$$f(n)=h(n)$$



**D is better  
than C !**



$$f(n)=g(n)+h(n)$$



# Evaluation Function

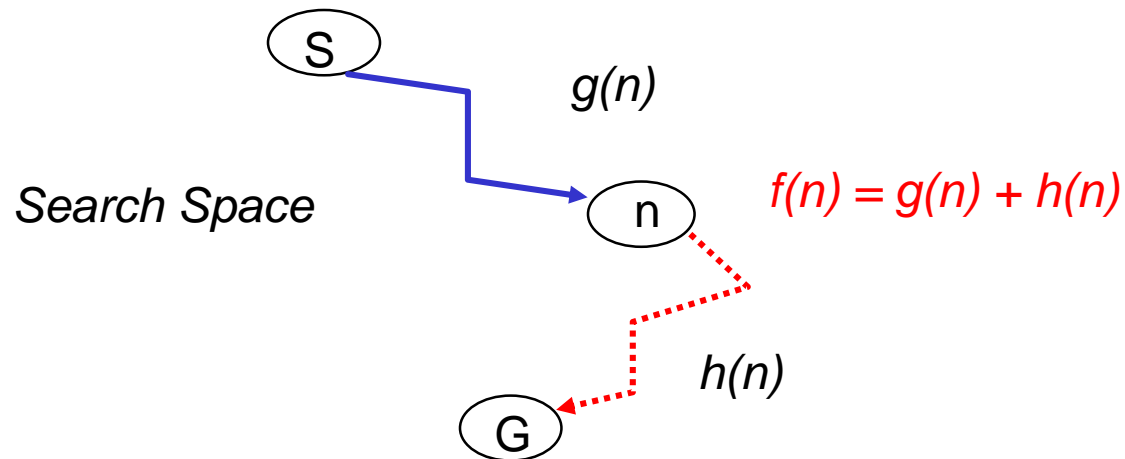
- Evaluation function

$$f(n) = g(n) + h(n)$$

$g(n)$  : actual cost (distance) from S to n

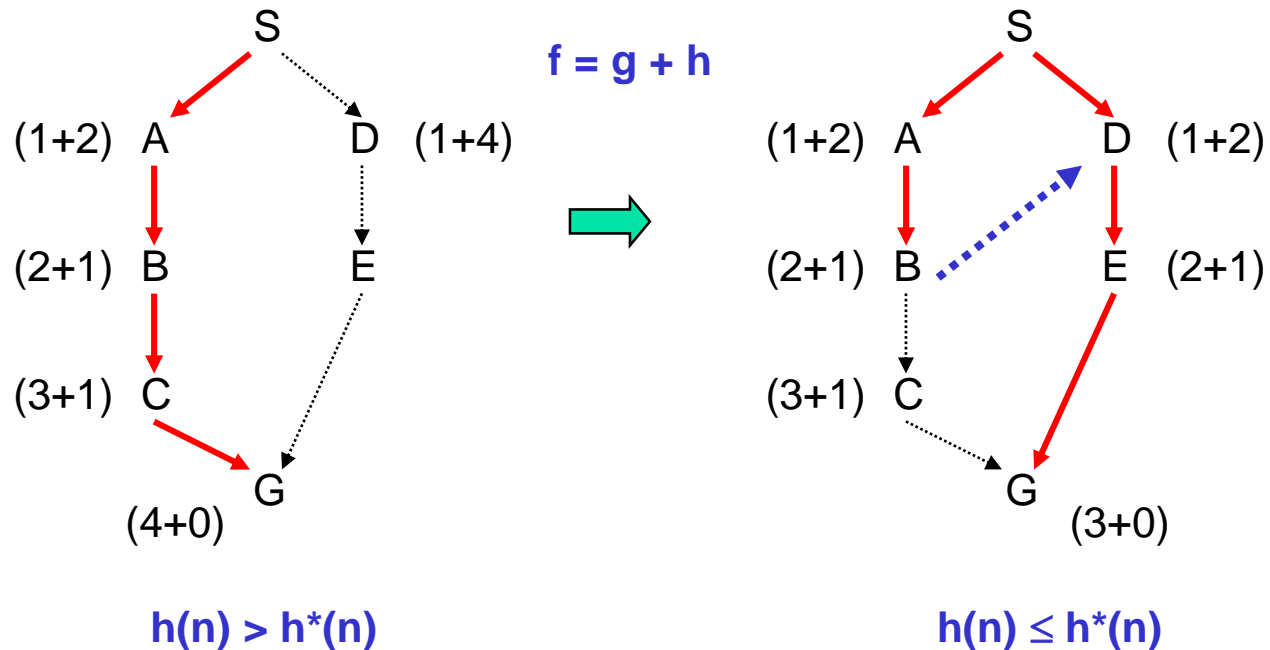
$h(n)$  : estimated cost (distance) from n to G

- →  **$f(n)$**  is the estimated cost (distance) *from S to G*



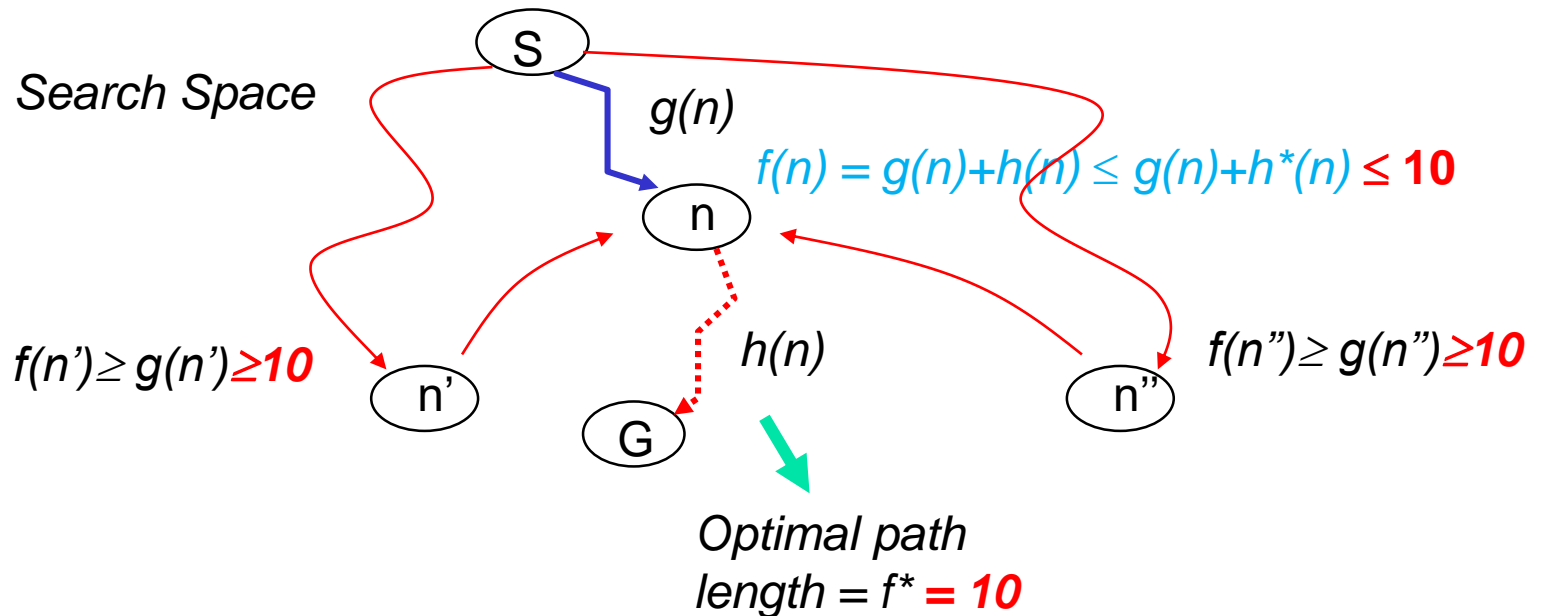
# Under Estimation

- Heuristics are fallible
  - Optimal path may not be found  
→ use heuristic value smaller than real value ( $h^*(n)$ )



# Under Estimation

- $h(n) \leq h^*(n)$  guarantees optimal path
  - Example: optimal path length =  $f^* = 10$





# A\* Search

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- A search algorithm is admissible
  - If it is **guaranteed to find a minimal path** to a solution whenever such a path exist
- A\* algorithm
  - A best-first search with

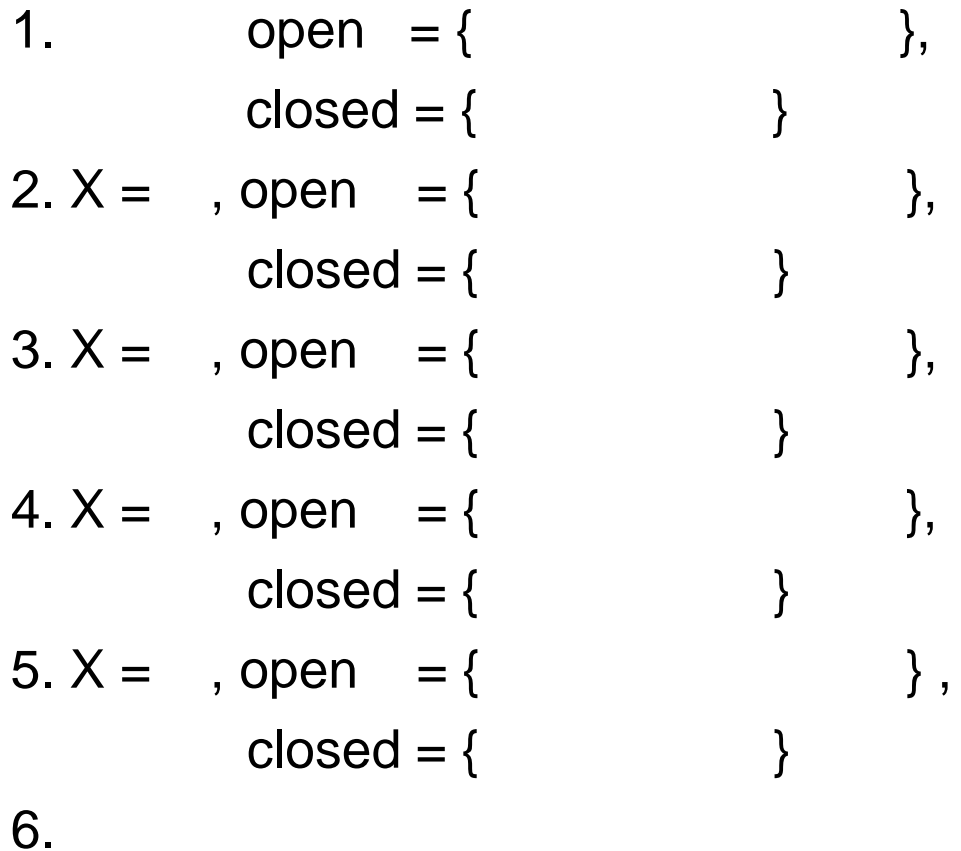
$$f(n) = g(n) + h(n)$$

where  **$h(n) \leq h^*(n)$**



*Admissible !*

( $h^*(n)$ ) : actual cost (distance) from n to G)



# Admissible Heuristics

## ■ BFS

- $f(n) = g(n)$

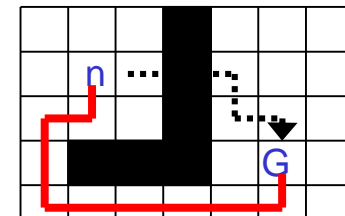
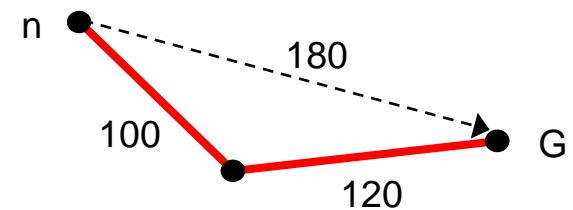
$\rightarrow h(n) = 0 < h^*(n)$

$\therefore$  BFS is an A\* algorithm !

## ■ Route-finding

- $h(n) = \text{straight line distance to } G$   
 $\leq h^*(n)$

- $h(n) = \text{Manhattan distance to } G$   
 $\leq h^*(n)$



# Admissible Heuristics

## ■ 8-puzzle

- $h_1(n) = \# \text{ of tiles in wrong position} \leq h^*(n)$
- $h_2(n) = \text{Sum of Manhattan distance} \leq h^*(n)$

<table><tr><td>2</td><td>8</td><td>3</td></tr><tr><td>1</td><td>6</td><td>4</td></tr><tr><td></td><td>7</td><td>5</td></tr></table>	2	8	3	1	6	4		7	5	5	6	0
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2	8	3										
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7	5											
	Tiles out of place	Sum of distances out of place	2 x the number of direct tile reversals									

1	2	3
8	4	4
7	6	5

Goal



# Admissible Heuristics

---

- Informedness

- Accuracy of  $h(n)$

- $0 < h_1(n) < h_2(n) \leq h^*(n)$

- ➡  $h_2$  is more accurate

- ➡  $h_2$  expands less nodes

- Example

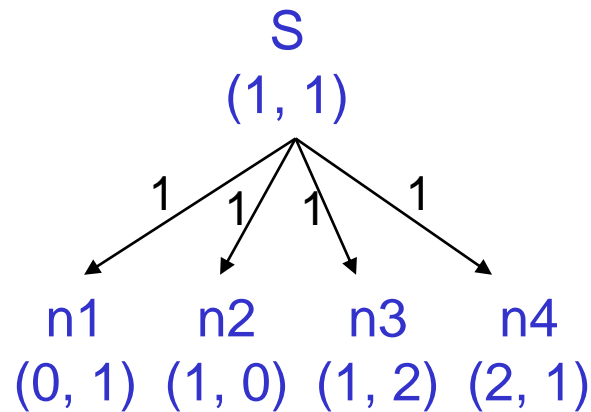
- $h_1(n) = \# \text{ of tiles in wrong position} < h_2(n) = \text{Manhattan distance}$   
→  $h_2(n)$  is better !



# A\* Search - Example

- Route finding in grid space

- $h(n)$  : Manhattan distance



$$g(n1) = 1$$

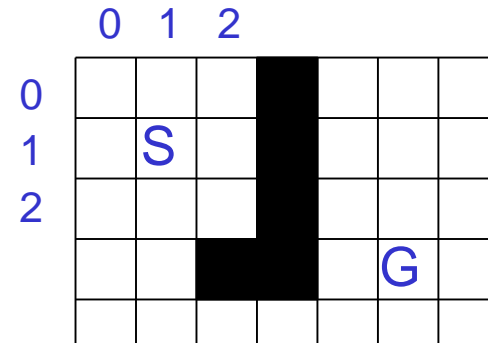
$$h(n1) = 7$$

$$f(n1) = 1 + 7 = 8$$

$$g(n4) = 1$$

$$h(n4) = 5$$

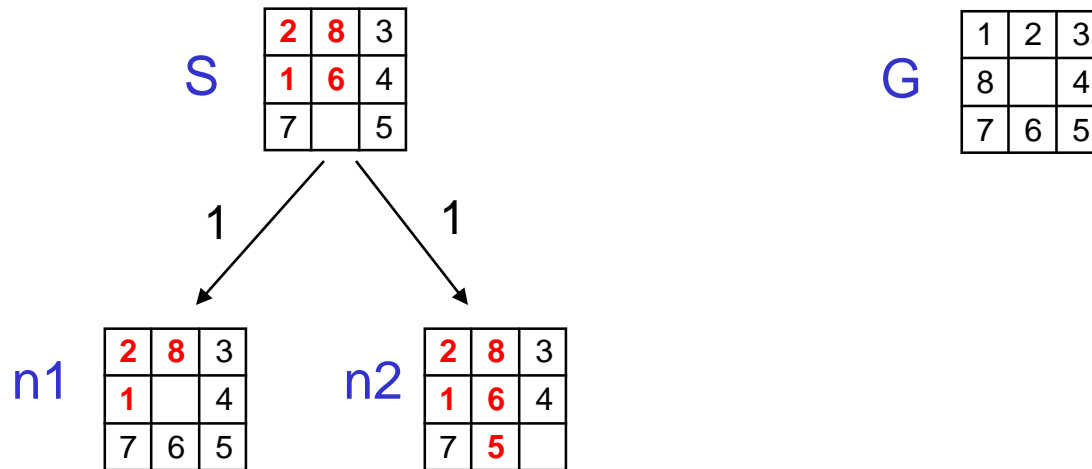
$$f(n4) = 1 + 5 = 6$$



# A\* Search - Example

- 8-puzzle

- $h(n)$  : # of tiles out-of-place



$$g(n1) = 1$$

$$h(n1) = 3$$

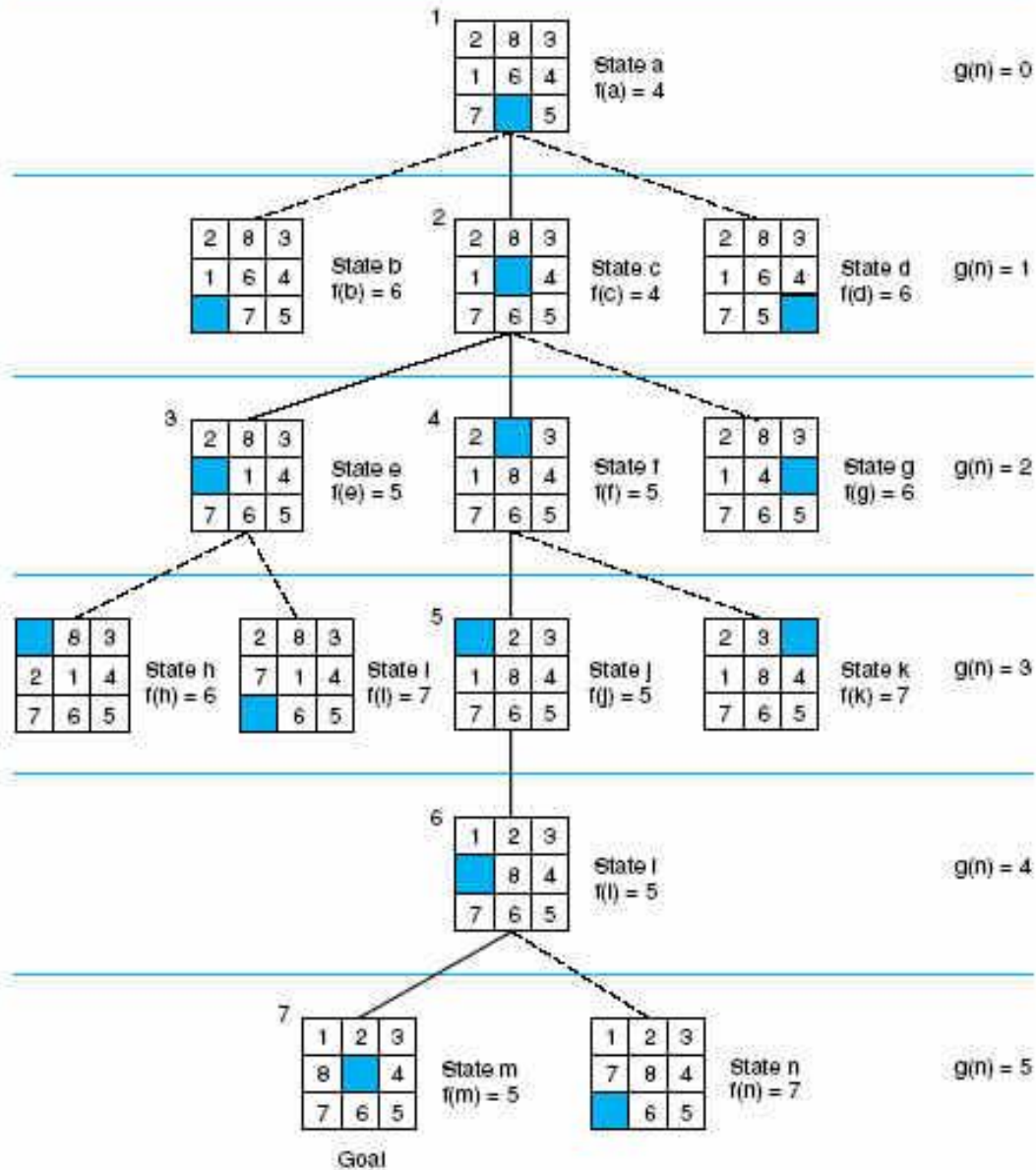
$$f(n1) = 1 + 3 = 4$$

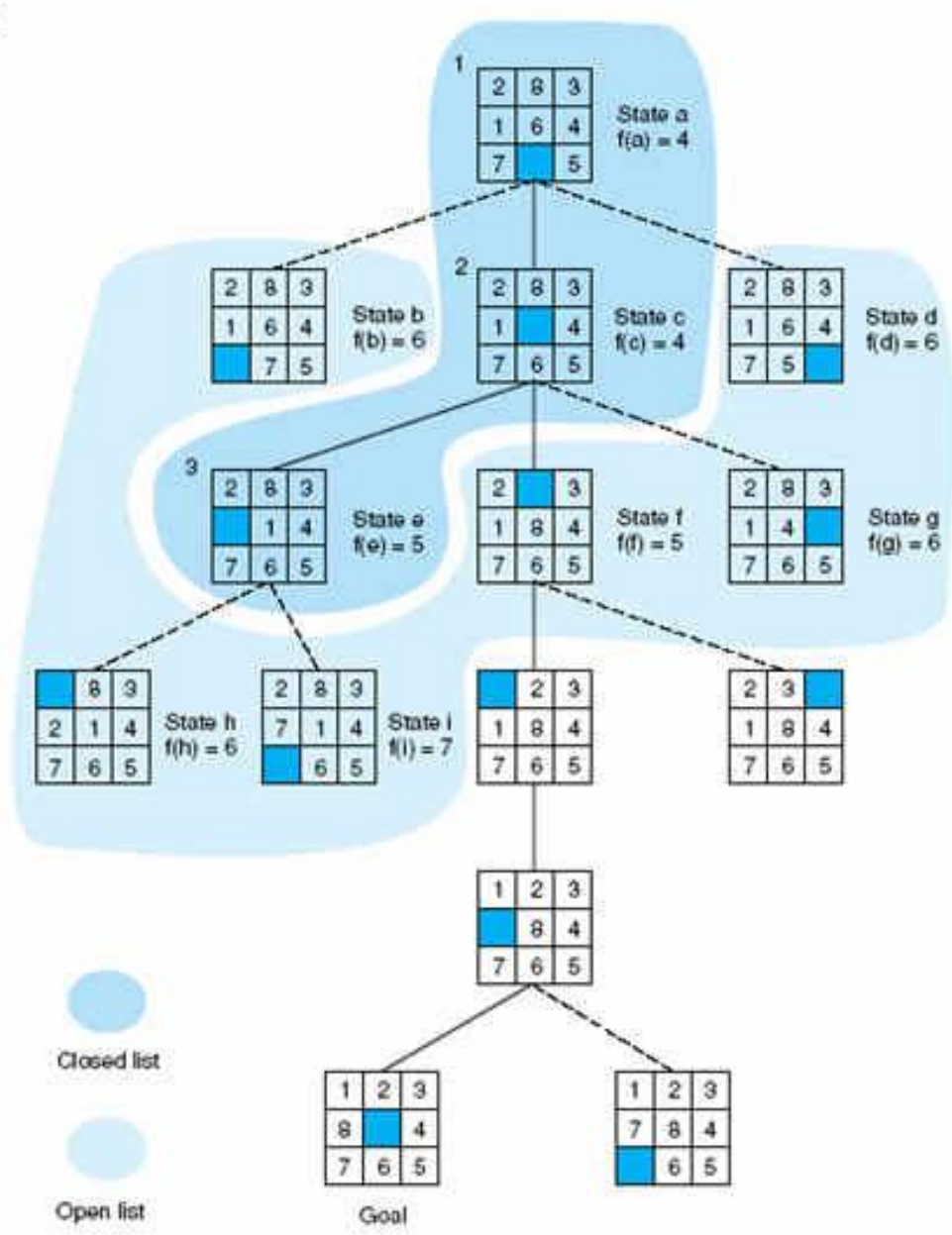
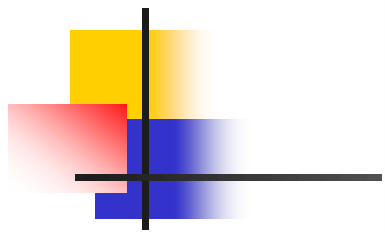
$$g(n2) = 1$$

$$h(n2) = 5$$

$$f(n2) = 1 + 5 = 6$$

Level of search  
 $g(n) =$

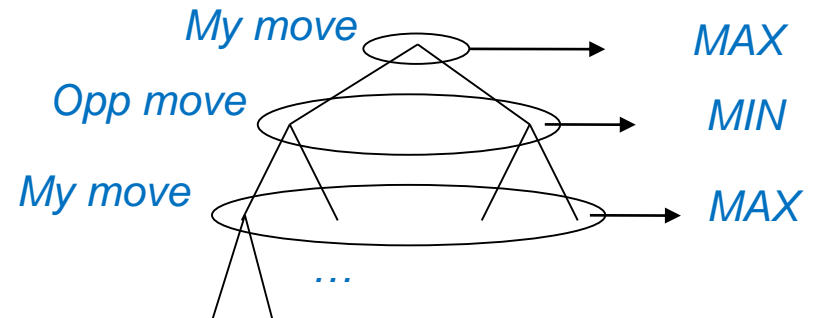






# Using Heuristics in Game

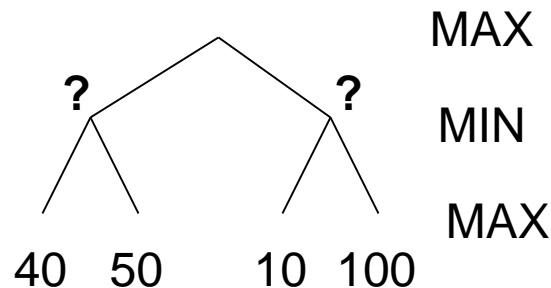
- Search in 2-persons game
  - Chess, Baduk, ...
- Assumption
  - I select a state that **MAXimize** my winning probability
  - **Opponent** select a state that **MINimize** my winning probability
- Divide states
  - MAX - my move
  - MIN - opponent's move



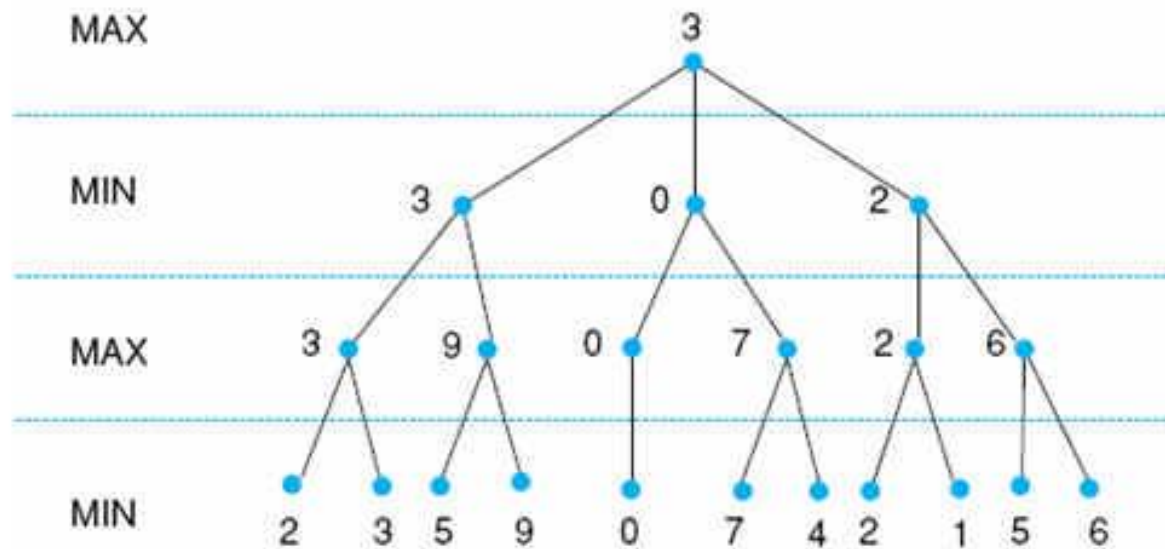
# Minimax Procedure

## ■ Evaluation

1. Search to certain depth
2. Evaluate states (higher value means better state)
3. In MAX node, value of node = max of children
4. In MIN node, value of node = min of children

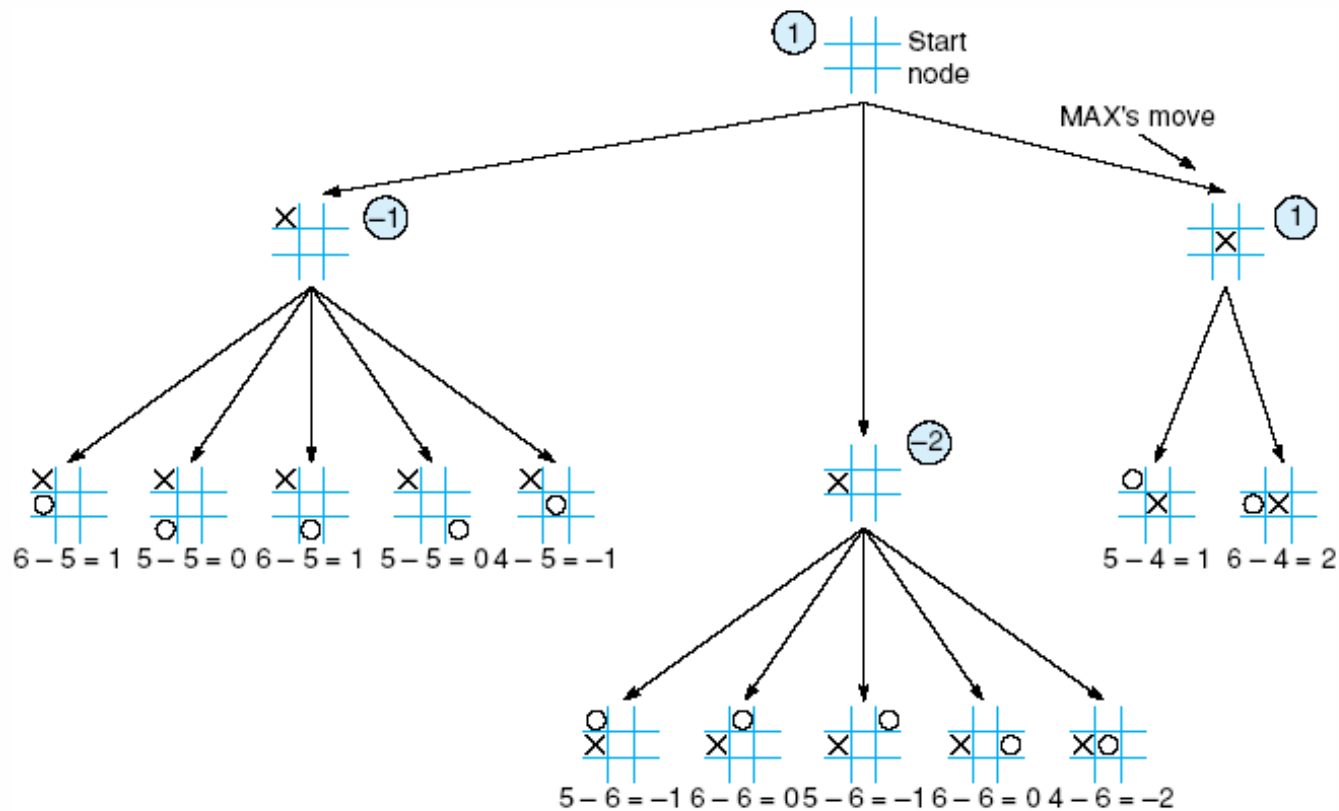


# Minimax Procedure





# Example: TicTacToe

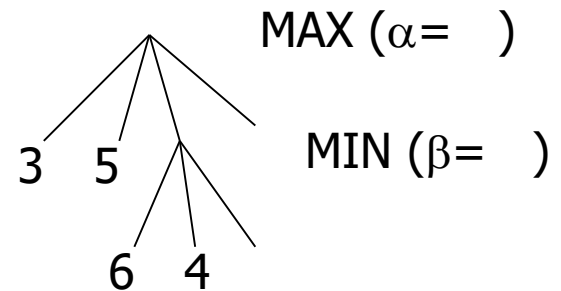


# Alpha-Beta Pruning

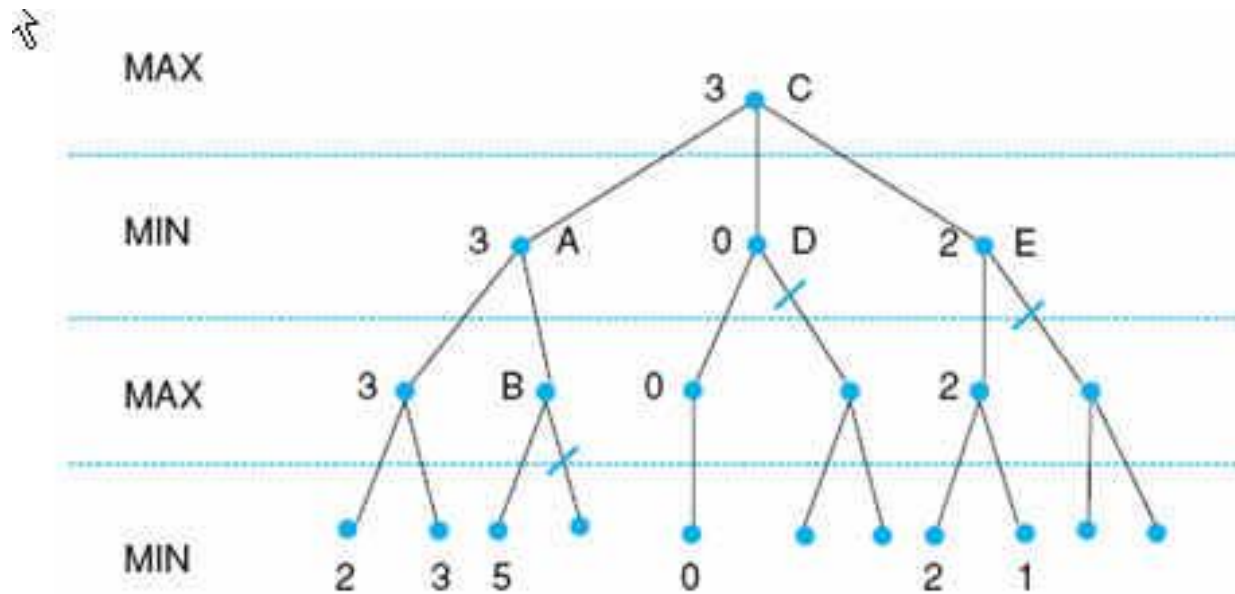
- Reduce the search space
- Algorithm
  - Search the game tree like DFS with depth bound
  - $\alpha$  - current maximum value of MAX node
  - $\beta$  - current minimum value of MIN node
  - If  $\beta \leq \alpha$  of parent, stop search
  - If  $\alpha \geq \beta$  of parent, stop search

*$\alpha$  always increase*

*$\beta$  always decrease*



# Alpha-Beta Pruning



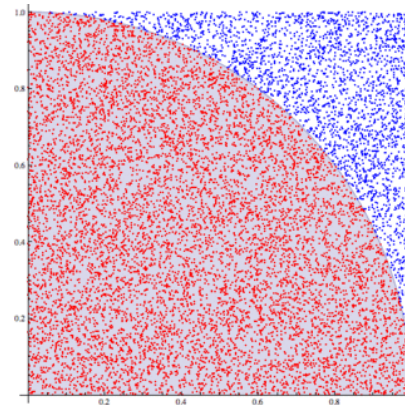
# MCTS

## ■ Monte Carlo Method

- Algorithms rely on repeated *random sampling* to obtain numerical results
- Example: computing  $\pi$ 
  - Generate random  $(x, y)$  in  $[0, 1] \times [0, 1]$
  - Count the number of points inside the circle of radius 1

$$\frac{\# \text{ of inside points}}{\# \text{ of total points}} \approx$$

$$\frac{\text{area of } \frac{1}{4} \text{ circle}}{\text{area of square}} = \frac{\frac{1}{4}\pi}{1}$$



# MCTS

## ■ MCTS: Monte Carlo Tree Search

- A heuristic search algorithm for decision processes in game
  1. The game is played out to the end by selecting random moves
  2. Playout is then used to weight the nodes in the game tree→ better nodes are more likely to be chosen in future play

