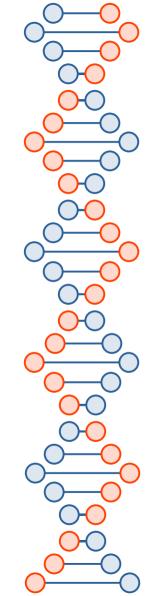


#### The "Shiva runtime environment"



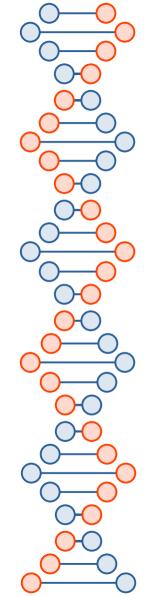
Advancing the programmability of the Linux process address space

Ryan "elfmaster" O'Neill: ryan@bitlackeys.org



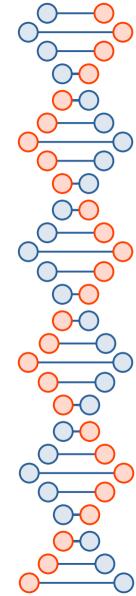
# What is Shiva exactly?

- A weird machine: programmable runtime engine
- Custom dynamic linker:
  - Loads ELF micro-programs into the address space
- Shiva provides a rich API for program transformation, tracing, hooking, and debugging.
- Think "LKM's (Loadable kernel modules) for userland processes".
- A programmable debugging engine that does NOT require the ptrace system call



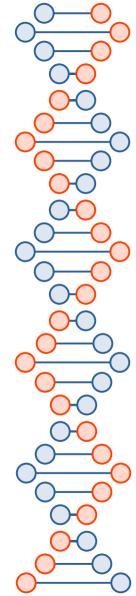
#### What can Shiva be used for?

- Building debuggers and tracers
- Designing process Security modules
   Anti-exploitation, Sandboxing, Process hardening
- In-memory fuzzing harnesses
- Software profiling
- Virus detection engines
- Malware unpackers
- Hot-patching code and data
- Userland-Rootkit detection & disinfection
- Extremely fast instrumentation, hooking, and process injection of all types (Without the use of SYS\_ptrace).



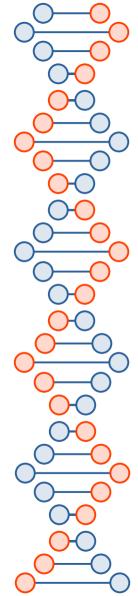
# From a blackhat perspective

- Shiva can be used to crack software
- Design powerful in-process rootkits
- Perform Modular virus infection: "Preloading the linker for fun and profit": https://tmpout.sh/2/6.html
- Create ELF binary packers



#### How does Shiva run?

- Shiva loads and links "Shiva modules" at runtime
- Shiva loads the module "/opt/shiva/modules" directory
  - 1. Shiva can be invoked directly as an executable
    - Shiva works as a userland-exec to load and map the target executable into its own process address space (See grugqs work on ul\_exec: https://grugq.github.io/docs/ul\_exec.txt)
    - i.e. ./shiva /bin/test\_prog
  - 2. Shiva can be invoked indirectly as a runtime interpreter (i.e. A program declares shiva as the dynamic linker)
    - The kernel loads the target executable and loads Shiva as the interpreter
    - Good for writing modules that are linked at every single execution
    - i.e. /bin/test\_proq



# Shiva ran directly on program ./test

- Shiva is loading a module: pltgot\_hook.o
- This module hooks the PLT function 'puts' and modifies the string output

```
elfmaster@esoteric-labs:~/git/shiva$ ./test abc

Hello World

Hello World

Hello World

Hello World

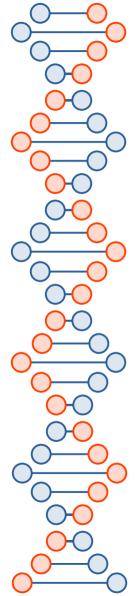
abc

elfmaster@esoteric-labs:~/git/shiva$ ./shiva test abc

hijacked your string: 'Hello World'

hijacked your string: 'abc'

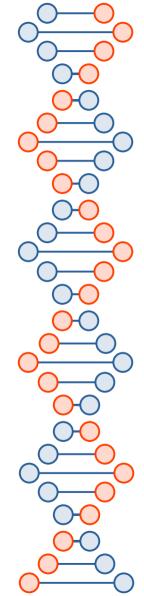
elfmaster@esoteric-labs:~/git/shiva$
```



# Shiva interpreter path

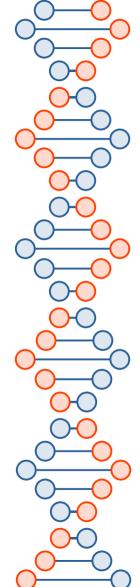
 The INTERP segment contains the path to 'shiva' instead of "Id-linux.so".

```
Elf file type is DYN (Position-Independent Executable file)
Entry point 0x1080
There are 14 program headers, starting at offset 64
Program Headers:
          Offset
                      VirtAddr
                                  PhysAddr
 Type
          FileSiz
                      MemSiz
                                  Flags Align
 PHDR
          0x0000000000000310 0x000000000000310
                                       0x8
 INTERP
          0x1
    [Requesting program interpreter: /home/elfmaster/git/shiva/shiva]
          LOAD
```

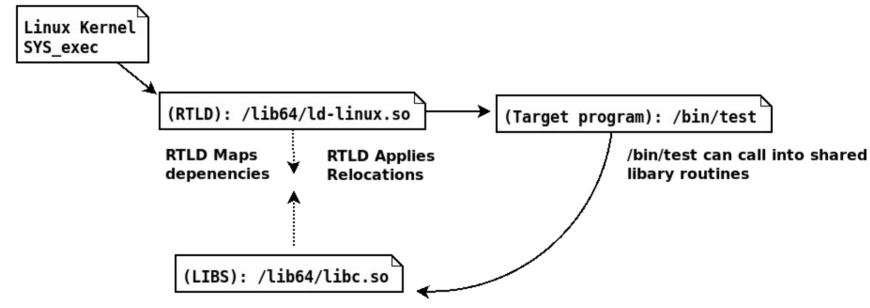


# Shiva running as the interpreter for ./test

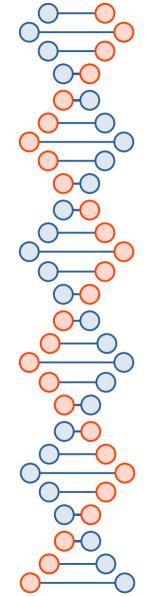
```
lfmaster@esoteric-labs:~/git/shiva$ ./test abc
hijacked your string: 'Hello World'
hijacked your string: 'abc'
elfmaster@esoteric-labs:~/git/shiva$
```



# Process execution flow of a normal dynamically linked ELF

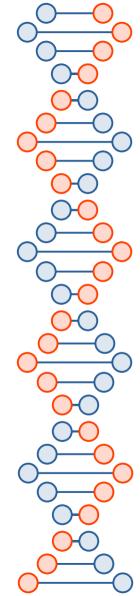


#### Process execution flow with Shiva as an interpreter. Linux Kernel SYS exec (Shiva): /bin/shiva (RTLD): /lib64/ld-linux.so (Target program): /bin/test Shiva maps the **RTLD Applies RTLD Maps** /bin/test dan call into shared Shiva maps modules RTLD "ld-linux.so" depenencies Relocations libary routines into the address into memory. space. /bin/test / is set with breakpoints/hooks trigger code within the Shiva module 'mod.o' Shiva API (LIBS): /lib64/libc.so (MODULES): /opt/shiva/modules/mod.o



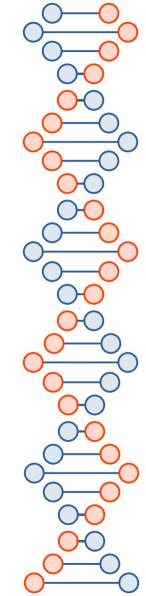
# Understanding Shiva modules

- Shiva modules are loaded from: "/opt/shiva/modules"
- Similar to LKM (Loadable kernel module):
  - Shiva modules are ELF relocatable object
  - Shiva modules must have an initialization routine int shakti\_main(shiva\_ctx\_t \*ctx)



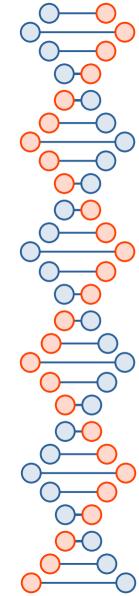
### Shiva Module: context struct

- The context struct "shiva\_ctx\_t"
- Passed into the modules init function
- Contains:
  - Register state
  - Process-state
  - Thread-state
  - Control flow data
  - Signal data
  - Shiva-Trace API specific data structures
  - Full access to **elfobj\_t** of target executable



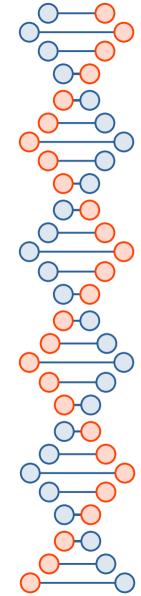
# Shiva modules: Linking

- Shiva modules use the Shiva trace API for program instrumentation
- Modules are written in C and have access to:
  - . musl-libc
  - . libelfmaster
  - Libudis86
  - Shiva-Trace API
- Shiva loader handles loading and linking of modules
- The module is an **ELF ET\_REL** (relocatable object file)
- Sections are mapped appropriately into memory segments: Text, and Data



# Shiva modules execute in-process

- Shiva modules are executed in-process.
- In-process means:
  - The Shiva runtime engine is in the same process as the target program
  - The module also executes within the same process address space
- Instrumentation and hooking is extremely fast
- Module has an *initialization function*
- Module init function sets hooks and breakpoints within the target program
- Handler functions are callbacks for hooks and breakpoints

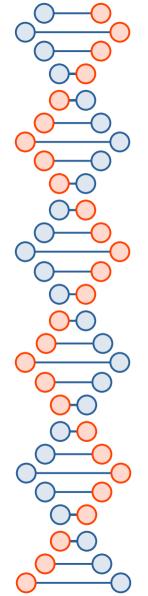


# Understanding Shiva modules: **gcc code models**

• Shiva modules are currently compiled with GCC and must be compiled with a large code model (i.e.

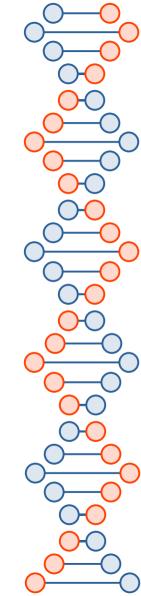
gcc -c -mcmodel=large)

- A large code model is necessary to handle module relocation offsets that exceed 4GB
- On 64bit system: The Shiva executable, the Shiva module, and the target executable will be mapped to addresses that are vastly far apart
- In some instances we can and must use a small code model for modules (More on this later)



A security module. Prevent command injection via system function

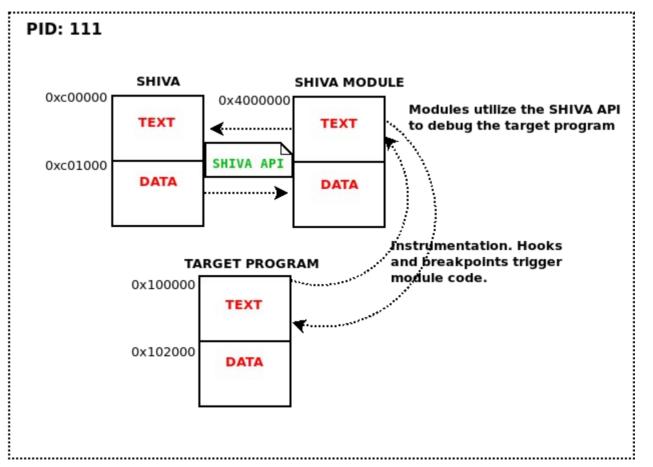
```
#include "../shiva.h"
int n system(const char *s)
       printf("s: %s\n", s);
       if (strstr(s, ";") != NULL || strstr(s, "|") != NULL) {
               printf("Detected possible OS command injection attack '%s'\n", s);
               abort();
       return system(s);
shakti main(shiva ctx t *ctx)
       bool res:
       shiva error t error;
       res = shiva trace(ctx, 0, SHIVA TRACE OP ATTACH,
           NULL, NULL, 0, &error);
       if (res == false) {
               printf("shiva trace failed: %s\n", shiva error msg(&error));
       res = shiva trace register handler(ctx, (void *)&n system,
           SHIVA TRACE BP PLTGOT, &error);
       if (res == false) {
               printf("shiva register handler failed: %s\n",
                   shiva error msg(&error));
       res = shiva trace set breakpoint(ctx, (void *)&n system,
           0, "system", &error);
       if (res == false) {
               printf("shiva trace set breakpoint failed: %s\n", shiva error msg(&error));
       return 0;
```

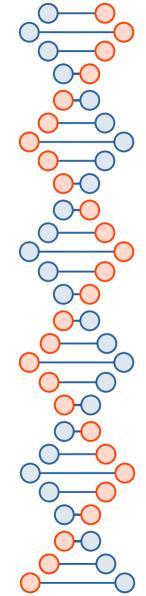


# The Shiva API (shiva\_trace)

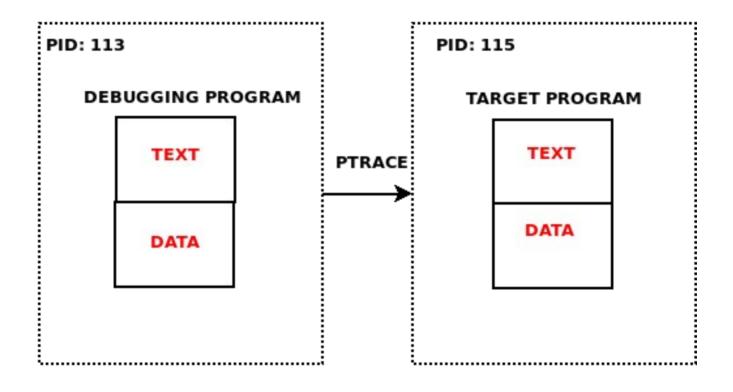
- Shiva modules are written in C
- The Shiva\_trace API can: access, read, transform and instrument the process image.
- Rich API for process transformation
- Excellent for writing debuggers, tracers, security modules, etc.
- In very early development phase

# Shiva debugging model





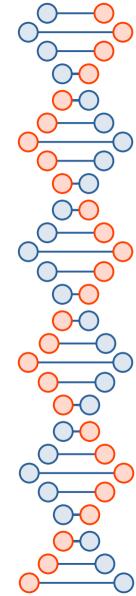
# PTRACE Debugging model





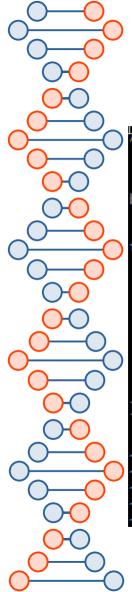
#### Shiva-Trace vs. PTRACE

- The Shiva-Trace API does not rely on the PTRACE system call
- Shiva-Trace is effective against hardened binaries and anti-debugging
- Shiva-Trace is implemented as an "in-process" programmable runtime engine.
- PTRACE on the other hand, executes from a separate process context.
- PTRACE: Slow. And is useless against many Linux anti-debugging techniques.
- PTRACE operations send a SIGSTOP to the target process.
- PTRACE is a system call and requires a context switch, whereas Shiva\_trace is implemented purely in userland and less expensive.



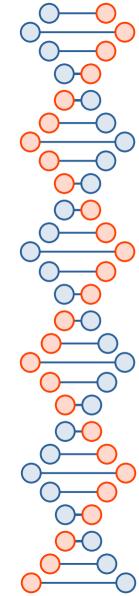
# Shiva Trace p3.

- shiva\_trace can cleverly instrument the target program at runtime
- shiva\_trace uses the terminology "Hooks" and "Breakpoints" interchangeably.
- Many breakpoints are really just various types of hooks
- shiva\_trace can defeat anti-debugging techniques
- Powerful for debugging hostile binaries



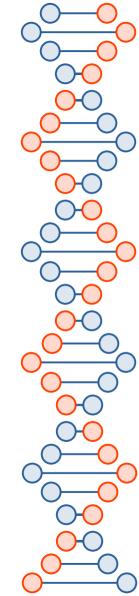
## Shiva API basics: shiva\_trace()

```
shiva trace lets the module attach to the process, read and write to the process,
 * get and set the register state.
bool shiva trace(shiva ctx t *, pid t, shiva trace op t, void *, void *, size t, shiva error t *);
typedef enum shiva trace op {
        SHIVA TRACE OP \overrightarrow{CONT} = 0,
        SHIVA TRACE OP ATTACH,
        SHIVA TRACE OP POKE,
        SHIVA TRACE OP PEEK,
        SHIVA TRACE OP GETREGS,
        SHIVA TRACE OP SETREGS,
        SHIVA TRACE OP SETFPREGS,
        SHIVA TRACE OP GETSIGINFO,
        SHIVA TRACE OP SETSIGINFO
 shiva trace op t;
```



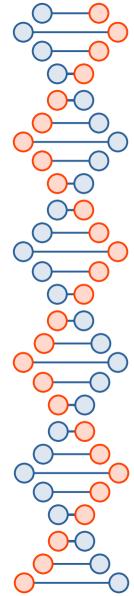
# Shiva-Trace API - Breakpoints

- Breakpoints are registered to a handler function
- Similar to Linux Kprobes, but for userland
- Multiple breakpoints can be set for a single handler
- Some breakpoints are hooks
- Some breakpoints are signal driven
- A complex series of in-process trampolines



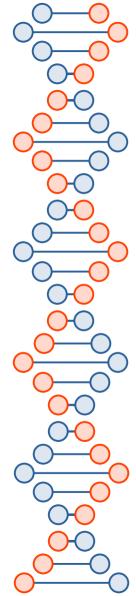
# Shiva breakpoints and signals

- We do not use PTRACE to catch signals
- Our debugging module exists within the same PID as the program it's debugging
- Signal driven breakpoints (i.e. int3) register a handler internally with sigaction()
- Shiva internally hooks **signal()** and **sigaction()** in the target program to prevent the target from overwriting breakpoint handlers



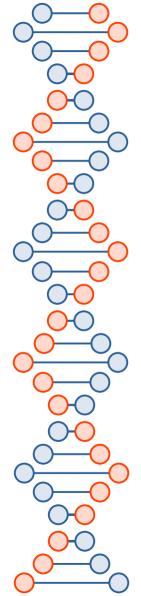
# A look at all Shiva breakpoint types

```
typedef enum shiva trace bp type {
    SHIVA TRACE BP JMP = 0,
    SHIVA TRACE BP CALL,
    SHIVA TRACE BP INT3,
    SHIVA_TRACE_BP_SEGV,
    SHIVA TRACE BP SIGILL,
    SHIVA TRACE BP TRAMPOLINE,
    SHIVA_TRACE_BP_PLTGOT
} shiva trace bp type t;
```



# Setting hooks and breakpoints

Breakpoints/hooks are registered to a handler function



#### Basic handler function



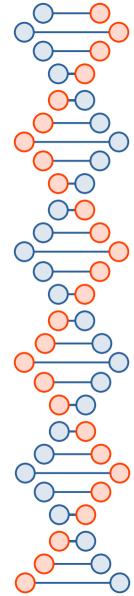
# SHIVA\_TRACE\_BP\_JMP

- Breakpoint type: Control flow hook
- Set on 5 byte jmp instructions
- The jmp offset is replaced by an offset to your registered handler function
- When done, the handler jumps back to the original target.
- **SHIVA\_TRACE\_LONGJMP\_RETURN** macro: rewinds the stack and jmp's to the original offset
- Defeats breakpoint detection with anti-debugging



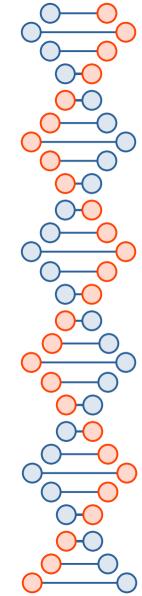
# SHIVA\_TRACE\_BP\_CALL

- Breakpoint type: Control flow hook
- Set on 5 byte immediate call instructions
- Original call target is replaced with offset to registered handler
- The registered handler returns after calling the original function (If desired)
- Extremely fast way to build a function tracer: (See modules/examples/func\_tracer.c)
- Only works when executing Shiva directly (Not as an interpreter) due to required small-code model.
- Defeats breakpoint detection



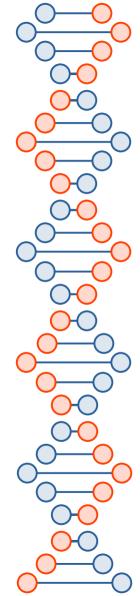
### SHIVA\_TRACE\_BP\_TRAMPOLINE

- Breakpoint type: Control flow hook
- This type of breakpoint inserts a trampoline on the entry point of the function
  - movq \$target\_addr, %rax jmp \*%rax
- Excellent for hooking any function locally or globally within the address space
- Trampoline function restores the original code bytes before calling the original function
- Not currently thread-safe. Memcpy's that place and restore hooks are not atomic operations
- Works well in a large code model whereas call hooks do not.



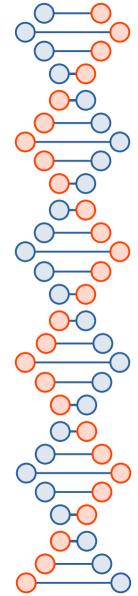
# SHIVA\_TRACE\_BP\_PLTGOT

- Hook any shared library function within the PLT/GOT
- Modifies the .got.plt[entry] of the given symbol
- Was tricky to implement. The RTLD wants to patch over our GOT patch. (i.e. with DT\_BINDNOW)
- Very low overhead, and thread safe
- Handler is usually a hook-replacement function for the function being hooked
- A single handler can also be used for multiple PLTGOT hooks. (i.e. Any PLT call triggers the same handler)



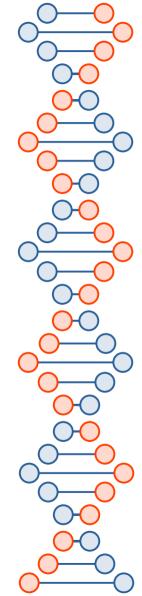
# SHIVA\_TRACE\_BP\_INT3

- Breakpoint type: Signal driven
- Traditional int3 breakpoint uses Oxcc opcode
- Causes kernel to deliver SIGTRAP
- The breakpoint handler is essentially a signal handler
- Uses struct ucontext to get register state in this case, as any signal handler would
- If the target program already had a **SIGTRAP** handler: ours replaces it, and optionally can invoke the original handler



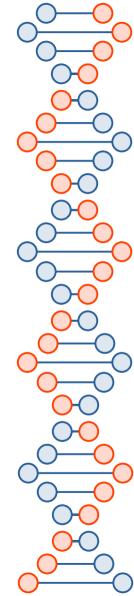
# SHIVA\_TRACE\_BP\_SIGILL

- Breakpoint type: Signal driven
- This breakpoint inserts a two byte ud2 (undefined instruction) to trigger SIGILL. "\x0f\x0b\"
- Breakpoint is registered to your signal handler function
- Can be helpful against hardened binaries that perform anti-debugging on standard breakpoints



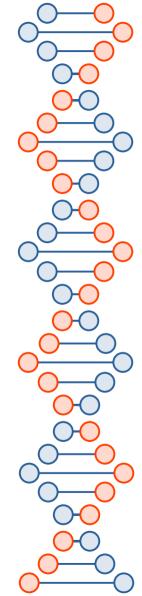
# SHIVA\_TRACE\_BP\_SEGV

- Breakpoint type: Signal driven
- Breakpoint code:
  - push <invalid\_offset>(%rip)
- Shiva\_trace API inserts an instruction with an invalid offset or address
- Causes an invalid memory dereference, and thus a **sigsegv**
- The handler function is a signal handler that catches the SEGV, and executes your custom code



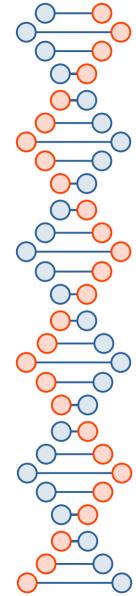
# Reading and writing to process memory

- The shiva\_trace() API call
- SHIVA\_TRACE\_OP\_PEEK
  - Allows the module to read anywhere in process memory
- SHIVA\_TRACE\_OP\_POKE
  - Allows the module to write anywhere in process memory (Except to the Shiva code itself)



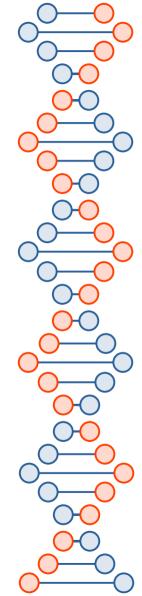
# Getting and setting register state

- The shiva trace API call
- SHIVA\_TRACE\_OP\_GETREGS
  - Get the current register values
- SHIVA\_TRACE\_OP\_SETREGS
  - Set the register values



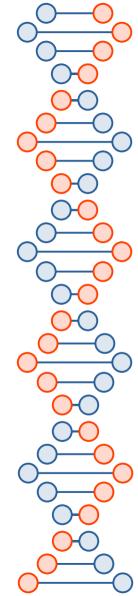
## In-process debugging tools with Shiva

- When running Shiva directly it can be used for creating powerful debugging tools
- "In-process" debugging technology is extremely fast
- Less context switching, more in-process trampolines
- modules/examples/func\_tracer.c is a module for function tracing
- Sets a **SHIVA\_TRACE\_BP\_CALL** hook on every immediate call
- Registered handler function simply prints each function name being called



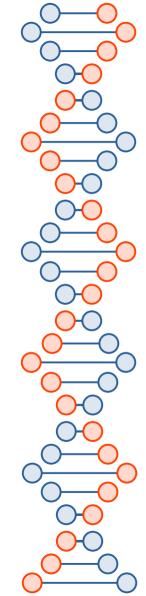
#### Shiva module: func\_tracer.o

- func tracer.c
- Only 84 lines of code
- Much faster than tracers that are written using **PTRACE**
- Benchmarks demonstrating performance:
- Command ./test:
  - Real 0m0.052s
- Tracing ./test with Shiva:
  - Real 0m0.155s
- Tracing ./test with /bin/ltrace
  - Real 0m3.582



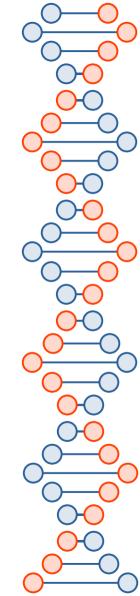
### Control flow integrity: ret2PLT protection

- modules/examples/plt\_cfi.c
- Only 109 lines of code
- Forward edge / Backward edge CFI
- Sets a SHIVA\_TRACE\_BP\_PLTGOT hook on every PLT call
- shiva\_trace\_set\_breakpoint() installs
   hooks/breakpoints and stores CFI data in breakpoint struct
- Every PLT call is redirected to our modules handler:
   int plt\_handler(void \*arg)



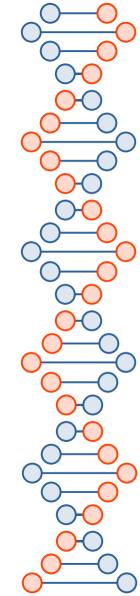
### Ret2plt handler function

- int plt\_handler(void \*arg)
- Checks the return address
- Prevents arbitrary calls to the PLT
- A ret2PLT attack will fail against this



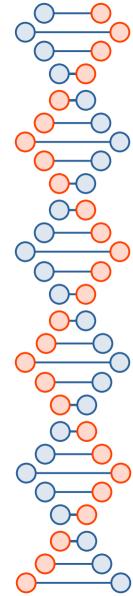
## Mitigate OS command injection

- modules/examples/os\_inject.o
- Only 48 lines of code
- SHIVA\_TRACE\_BP\_PLTGOT hook on system()
   function
- Replaced with secure function n\_system()
- n\_system() sanitizes the command string



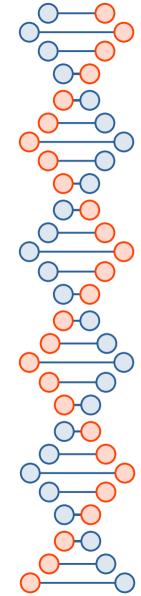
#### Crack software

- ./crackme
- Shiva module to crack serial number validator
- modules/examples/crackme\_bypass.o
- Redirects: int check\_serial(char \*serial)
- To a modified version of the function
- Uses a **SHIVA\_TRACE\_BP\_TRAMPOLINE** breakpoint



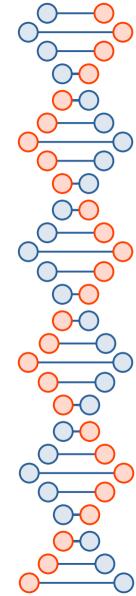
## Shiva for debugging hardened binaries

- There is a long history of anti-debugging techniques
- PTRACE system call is easily foiled
- PTRACE fails on PaX hardened binaries
  - mprotect restrictions prevent ptrace (PTRACE\_POKETEXT, ...);
- Shiva maps the hostile/hardened program into it's own address space
- In-process debugging can be very effective for Reverse engineering tough problems



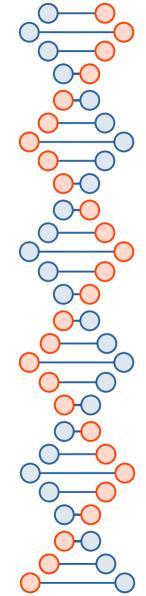
## Tracing a hardened binary

- modules/examples/func\_tracer.o
- Capable of tracing binaries with anti-debugging
- test\_antidebug.c: Simple anti-debugging techniques
  - ptrace(PTRACE\_TRACEME, 0, 0, 0, 0);
  - prctl(PR\_SET\_DUMPABLE, 0, 0, 0, 0);



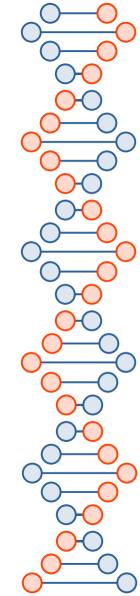
### Shiva works on stripped ELF binaries

- Corrupted meta-data
- Missing symbol tables
- Missing section header table
- Shiva uses libelfmaster: https://github.com/elfmaster/libelfmaster
- Forensic reconstruction of stripped binaries:
  - Symbol table reconstruction
  - Section header table reconstruction



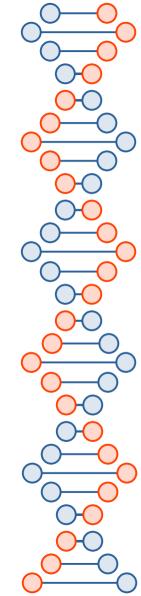
## Tracing a stripped binary

- Itrace fails to get any tracing info:
  - Requires section headers and symbol tables
- Our Shiva module for function tracing succeeds
  - Forensically reconstructs section and symbol data



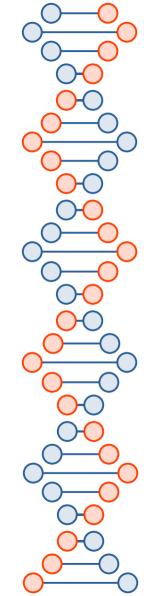
#### **Current limitations**

- Support for multi-threaded applications
- Shiva doesn't work on ELF's with **-fcf-protection**
- Shiva can load only one module at a time
- Shiva only works with **x86\_64 ET\_DYN** binaries
  - Will add support for **x86\_64 ET\_EXEC**
- Running Shiva as an interpreter requires a large code model for all modules
  - Some Shiva features require a small code model (i.e. SHIVA\_TRACE\_BP\_CALL breakpoints



# Future plans

- Open source the technology
- Continue to develop and expand the usability and functionality
- shiva-gcc wrapper for modules
- Design a suite of security modules for process hardening
- Handle loading multiple modules
- Handle threads



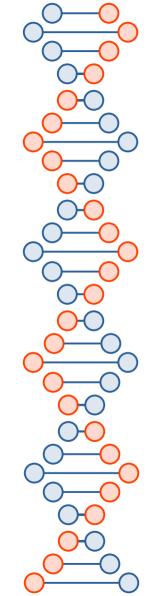
## A note on handling threads

- How do we handle threads without using SYS\_ptrace?
- TLS resolver hooks
- pthread\_create hooks
- SYS\_clone hooks
- Shared memory segments



#### Related works

- Linux kprobe instrumentation
  - http://phrack.org/issues/67/6.html
- Preloading the linker for fun and profit
  - https://tmpout.sh/2/6.html
- Embedded ELF debugging (by Mayhem & ERESI)
  - http://phrack.org/issues/63/9.html
- Userland exec implementation (By the Grugq: )
  - https://grugq.github.io/docs/ul\_exec.txt



## Questions?

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- Bitlackeys.org
- https://github.com/elfmaster/shiva