4. Threads

ECE30021/ITP30002 Operating Systems

Agenda

- Overview
- Multithreading models
- Thread libraries
- Threading issues
- Operating system examples

Motivation

- Most modern applications are multithreaded
- Threads run within application
- Multiple tasks with the application can be implemented by separate threads
  - Update display, fetch data, spell checking, answer a network request
- Process creation is heavy-weight while thread creation is light-weight
- Can simplify code, increase efficiency
- Kernels are generally multithreaded

Overview

- Process: program in execution
  - Each process occupies resources required for execution
- Thread: a way for a program to split itself into two or more simultaneously running tasks
  - Smaller unit than process
  - Threads in a process share resources
- A thread is comprised of
  - Thread ID, (program counter, register set, stack, etc)

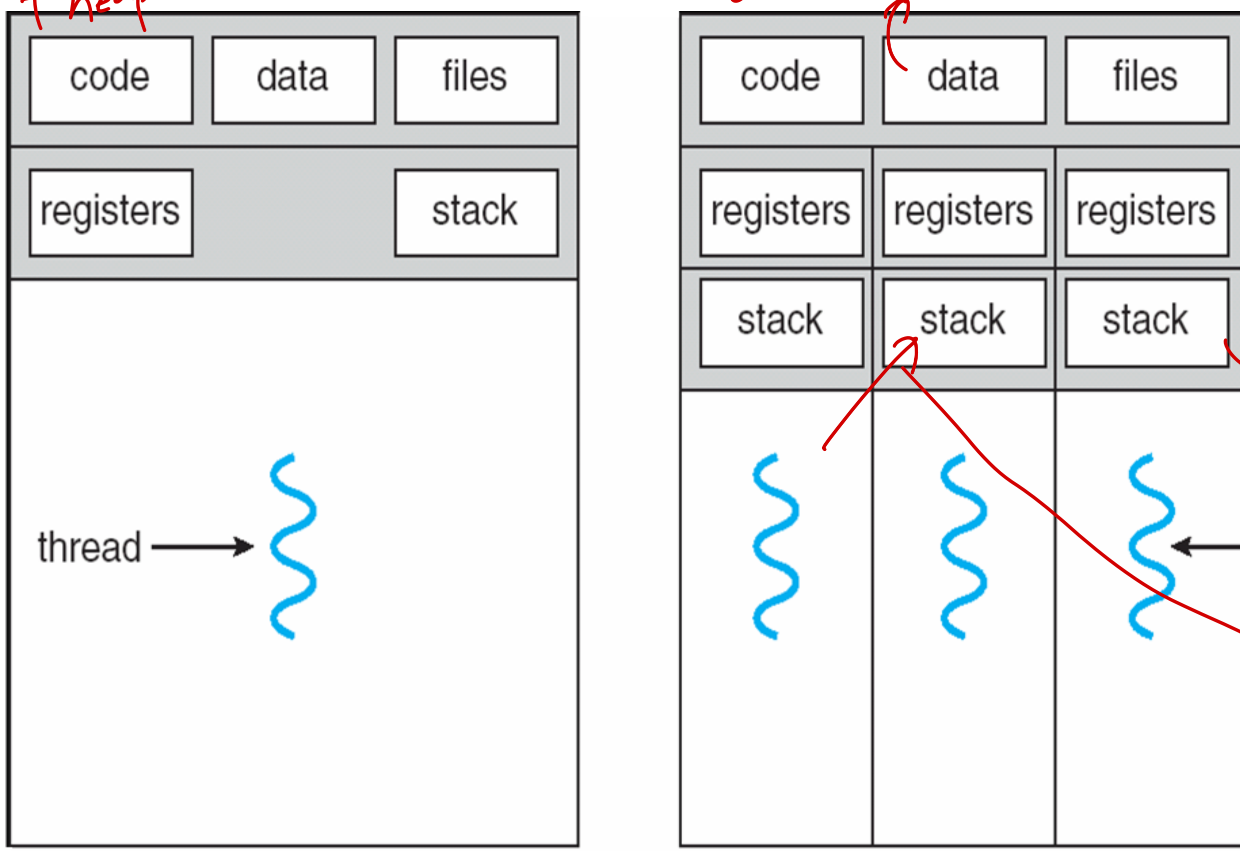
local variable
function call

단기기억 같은거

data memory
global variable
heap
stack

공통점: 프로그램 코드 실행에 직접적으로 필요함

Multithreaded Process



single-threaded process　　multithreaded process

+ heap

global variable 공유

공유자원

공유 X

local variable

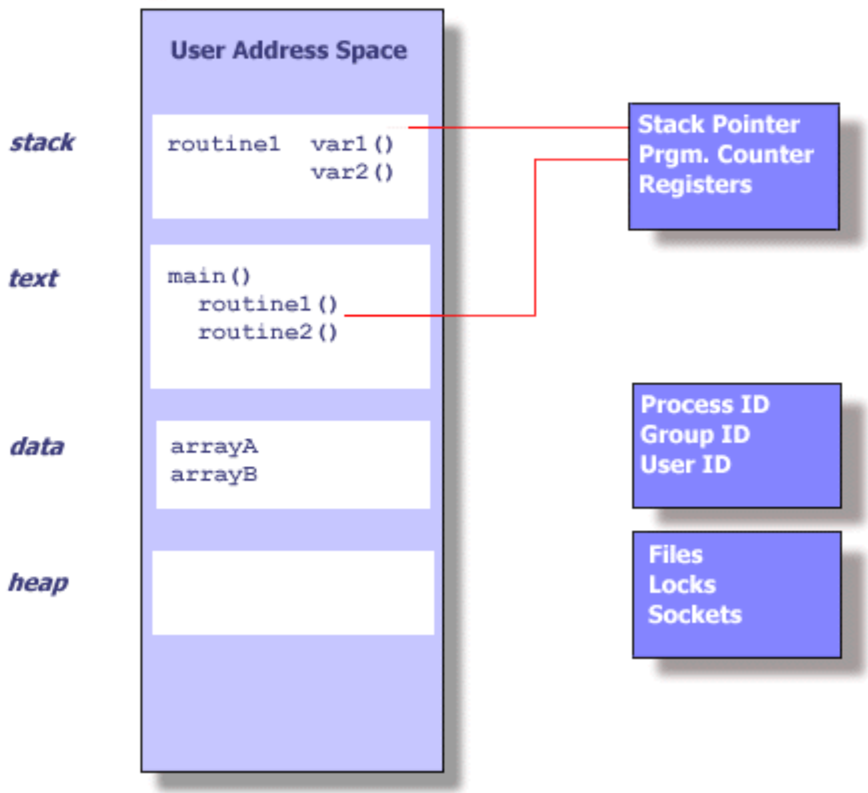
access 가능은 함

: 정확한 주소를 알아야 함

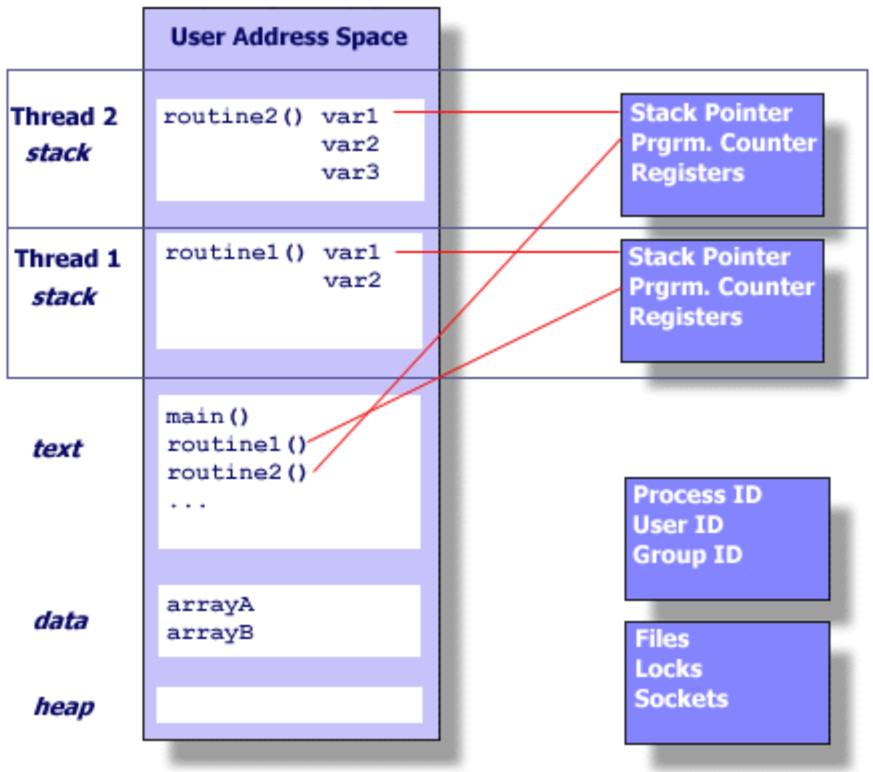
근데 위험함

local variable life time이 끝나서 access할 때 없을 수 있음

Process vs. Thread



Single-threaded process



Multi-threaded Process

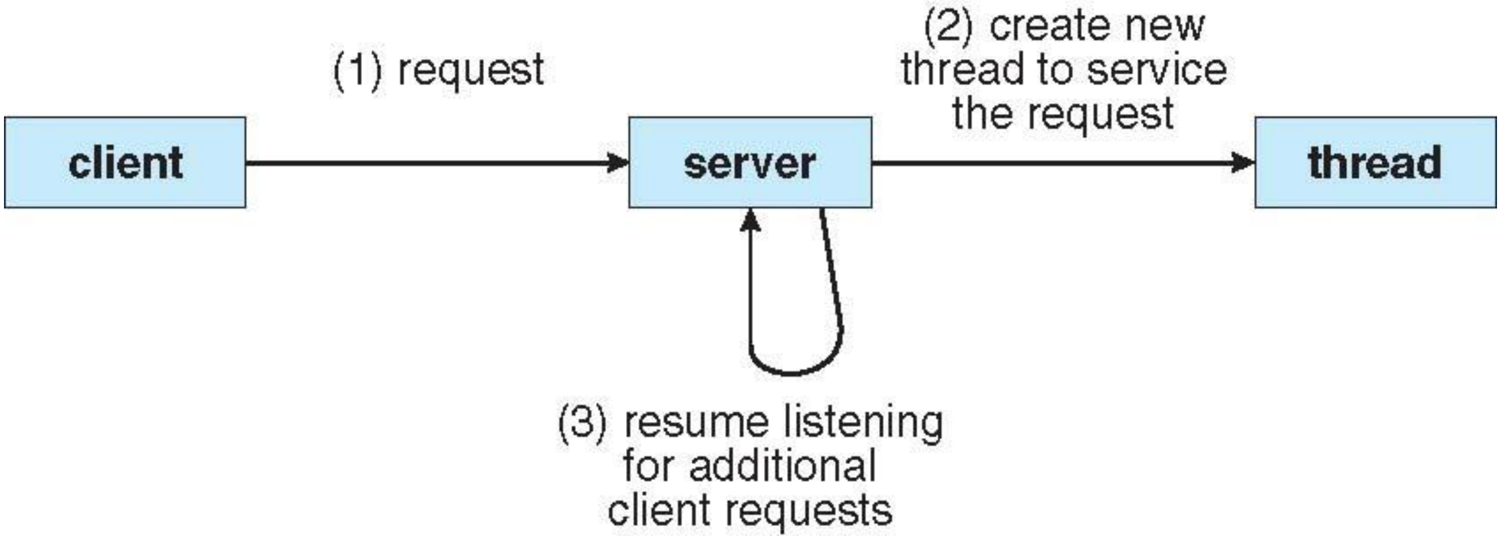
Thread Control Block (TCB)

- Thread Control Block (TCB) is a data structure in the operating system kernel which contains thread-specific information needed to manage it.

- Examples of information in TCB
  - Thread id
  - State of the thread (running, ready, waiting, start, done)
  - Stack pointer
  - Program counter
  - Thread's register values
  - Pointer to the process control block (PCB)

Why Thread?

- Process creation is expensive in time and resource
  - Ex) Web server accepting thousands of requests



Why Thread?

- Compared with single-threaded process
  - Scalability
    - Utilization of multiprocessor architectures
  - Responsiveness

- Compared with multiple processes
  - Resource sharing
  - Economy
    - Creating process is about 30 times slower than creating thread

Multicore Programming

- Multicore or multiprocessor systems putting pressure on programmers, challenges include:
  - Dividing activities
  - Balance
  - Data splitting
  - Data dependency
  - Testing and debugging - 특정 순서가 되면 버그가 발생하는 경우 디버깅이 어려움

- 추천 디버깅 방법: 1. 알고리즘 이해 2. 로그파일 작성 - thread 별로

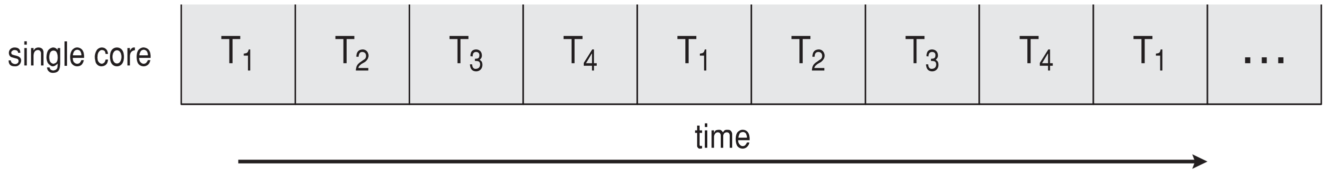
3. 애초에 처음부터 버그를 최소화하기

둘다 여러개 진행

Concurrency vs. Parallelism

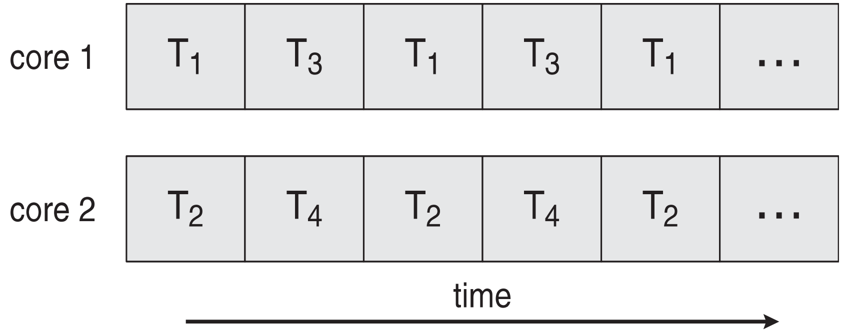
- Concurrency supports more than one task making progress
  - Single processor / core, scheduler providing concurrency

  Ex) Concurrent execution on single-core system:



- Parallelism implies a system can perform more than one task simultaneously

  Ex) Parallelism on a multi-core system:



Multicore Programming (Cont.)

작업은 나눌수 있어야 병렬가능
(병렬처리)

- Types of parallelism

  학생 100명
  TA 4명

  - Data parallelism – distributes subsets of the same data across multiple cores, same operation on each thread

    TA 1명당 25명씩
    채점

  - Task parallelism – distributing threads across cores, each thread performing unique operation

    TA에게 분업하기 채점
    1명은 과목조사
    1 회차 1~2
    1. 회차 2~3

- As # of threads grows, so does architectural support for threading
  - CPUs have cores as well as hardware threads
  - Consider Oracle SPARC T4 with 8 cores, and 8 hardware threads per core

암산

Amdahl's Law

- Identifies performance gains from adding additional cores to an application that has both serial and parallel components
  - S is serial portion, N processing cores

$$speedup \le \frac{1}{S + \frac{(1-S)}{N}}$$

Ex) Application is 75% parallel / 25% serial, moving from 1 to 2 cores results in speedup of 1.6 times

  - As N approaches infinity, speedup approaches 1 / S
  - Serial portion of an application has disproportionate effect on performance gained by adding additional cores.

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- Thread libraries
- Threading issues
- Operating system examples

Types of Threads

- User thread: thread supported by thread library in user level

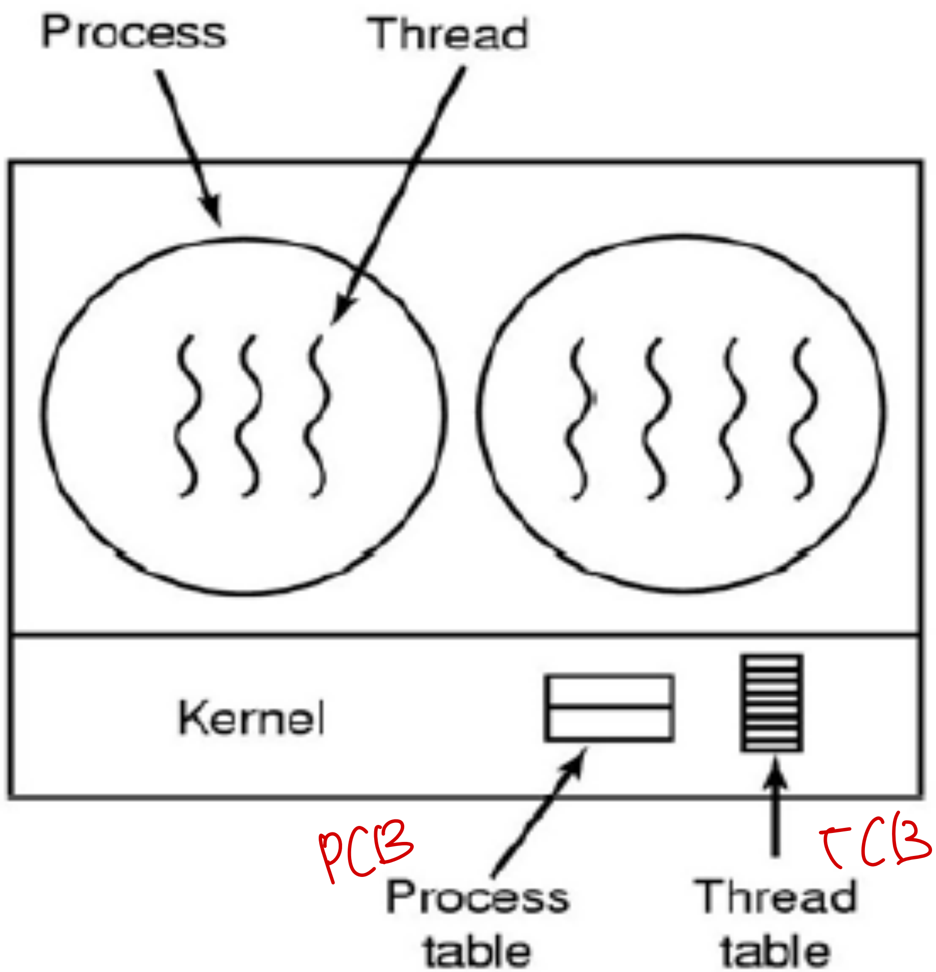
  user level library

  - Created by library function call (not system call)
  - Kernel is not concerned in user thread.
  - Switching of user thread is faster than kernel thread.

kernel scheduler가 관

- Kernel thread: thread supported by kernel
  - Created and managed by kernel
  - Scheduled by kernel
  - Cheaper than process user thread 보다 무겁고 느리다.
  - More expensive than user thread

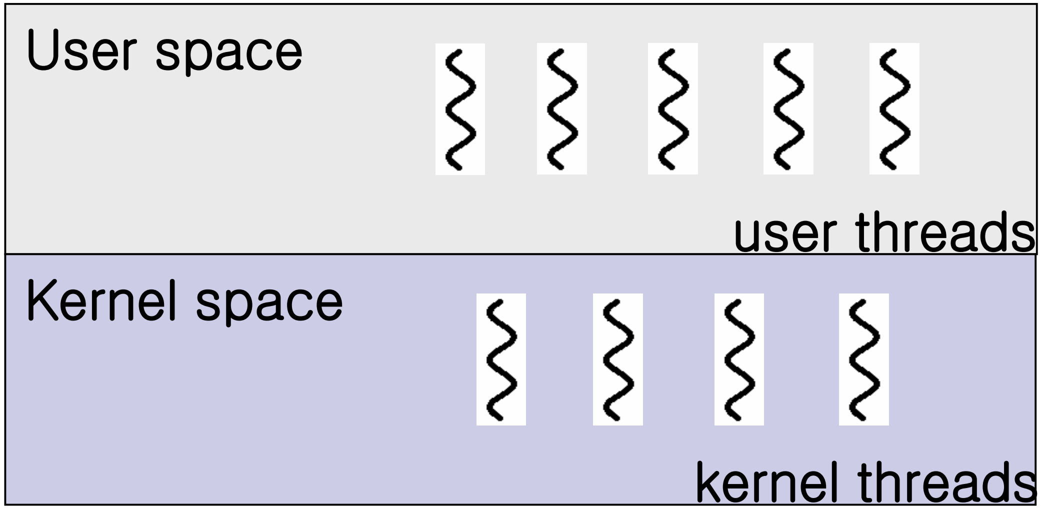
Kernel Thread



여기서
linked list가 쓰이게 되는 이유
: 생성과 소멸이 자주 이루어져서
linked list가 좋다.

Multithreading Models

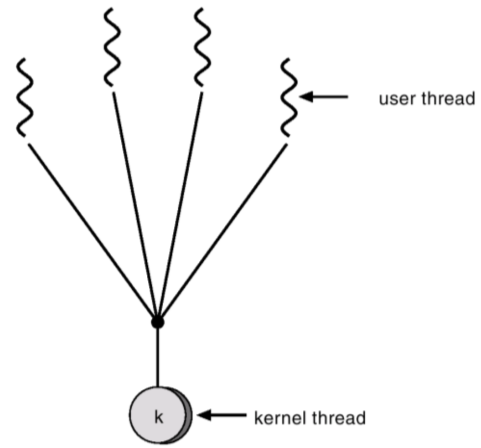
- Major issue: correspondence between user treads and kernel threads



Multithreading Models

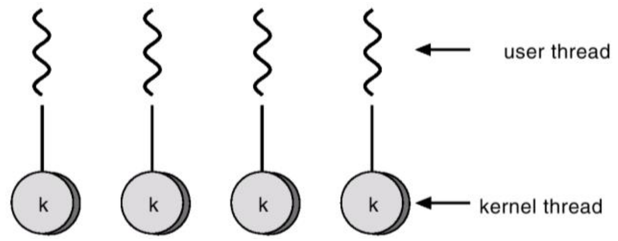
- Many-to-one model
  - Many user threads are mapped to single kernel thread
  - Threads are managed by user-level thread library

  Ex) Green threads,
  GNU Portable Threads



- One-to-one model
  - Each user thread is mapped to a kernel thread
  - Provides more concurrency
  - Problem: overhead

  Ex) Linux, Windows, Solaris

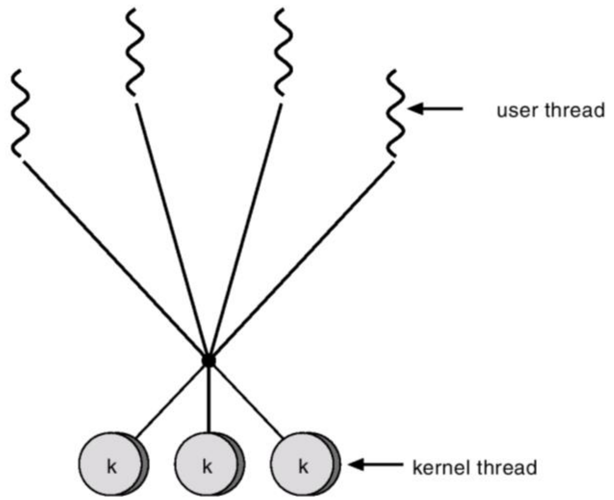


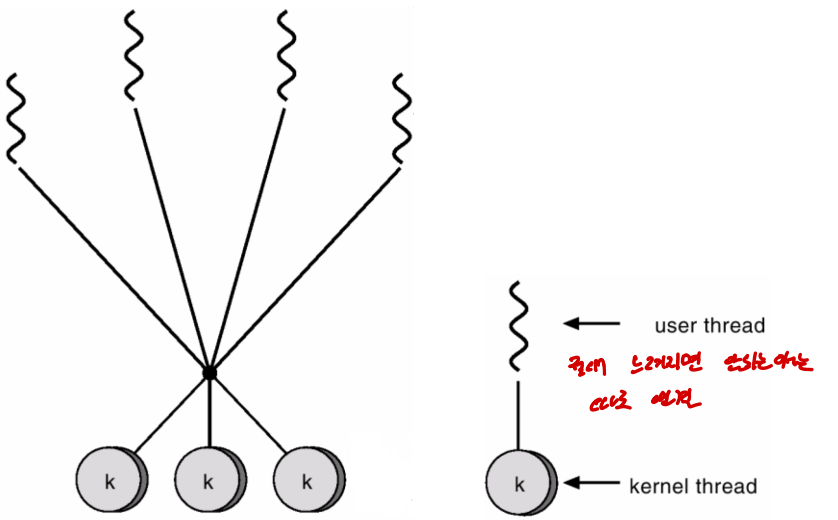
Multithreading Models

- Many-to-Many model
  - Multiplex many user level threads to smaller or equal number of kernel threads
  - Compromise between n:1 model and 1:1 model

- Two-level model
  - Variation of N:M model
  - Basically N:M model
  - A user thread can be bound to a kernel thread

  Ex) IRIX, HP-UX, Tru64, Solaris (<=8)





Scheduler Activation

- Communication between the kernel and the thread library in many-to-many model and two-level model.
  - Scheduler activation is one scheme for communication between user thread library and kernel

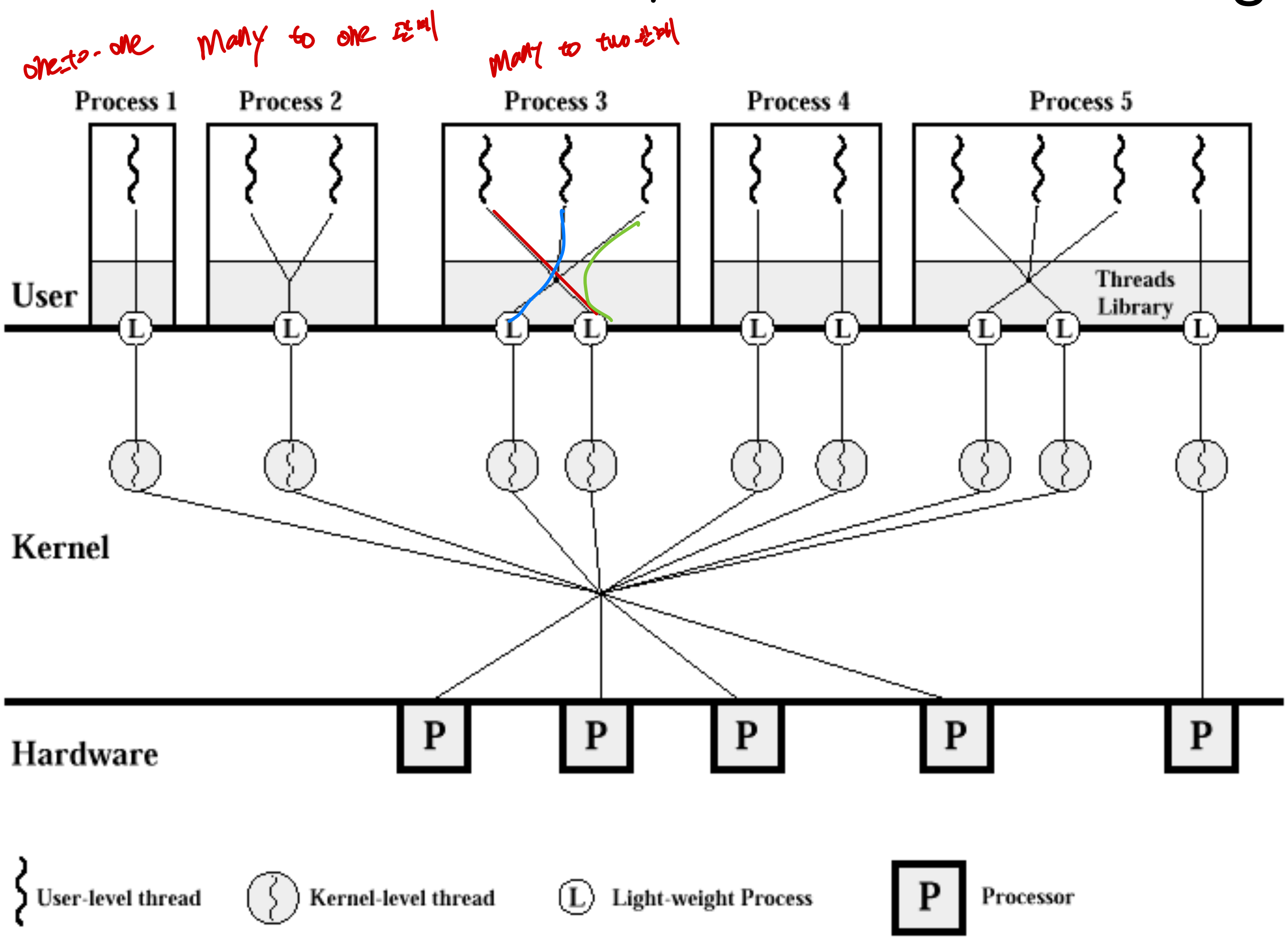
- In many-to-many model and two-level model, user threads are connected with kernel threads through LWP.

kernel thread와 1대1 관계, 환경에따라서 kernel thread로 되기도 함

- Lightweight process (LWP)
  - A data structure connecting user thread to kernel thread.
  - Basically, a LWP corresponds to a kernel thread, but there are some exceptions.
  - To the user-level thread library, a LWP appears to be a virtual processor.

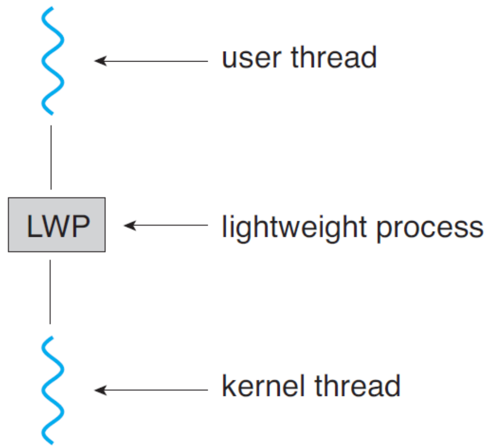
Scheduler Activation and LWP

- Connection between user/kernel threads through LWP



Scheduler Activation

- Kernel provides a set of virtual processors(LWP's).
- User level thread library schedules user threads onto virtual processors.
- If a kernel thread is blocked or unblocked, kernel notices it to thread library(upcall).
- Upcall handler schedules properly.
  - If a kernel thread is blocked, assign the LWP to another thread
  - If a kernel thread is unblocked, assign an LWP to it.



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Thread Libraries

- Thread library: set of API's to create and manage threads
  - User level library
  - Kernel level library

Examples)

- POSIX Pthreads: a specification with various implementations.
  - LinuxThreads
  - NPTL (Native POSIX Thread Library)
  - GNU Portable Threads
  - Open source Pthreads for win32
  - Etc.
- Win32 threads
- Java threads

POSIX Pthreads

- Reading assignment: Search the Internet for the following functions, and study them.

- API functions

  포인터 thread ID

  - int pthread_create(pthread_t *thread, const pthread_attr_t *attr, void * (*start_routine)(void *), void *arg);

    function pointer

    - Creates thread

  - int pthread_attr_init(pthread_attr_t *attr);
    - Initializes *attr* by default values

  - int pthread_join(pthread_t th, void **thread_return);
    - Waits for a thread *th*.

  exit는 process가 종료

  - void pthread_exit(void *retval);
    - Terminates thread

Example

55

```
#include <pthread.h>
#include <stdio.h>

int sum = 0; /* this data is shared by the thread(s) */
void *runner(void *param); /* the thread */

int main (int argc, char *argv[])
{
      pthread_t tid = 0; /* the thread identifier */
      pthread_attr_t attr; /* set of thread attributes */
      if (argc < 2 ) {
               fprintf(stderr,"usage: a.out <integer>\n");   10
               exit(0);
      }
      if (atoi(argv[1]) < 0) {
               fprintf(stderr,"%d must be <= 0\n",atoi(argv[1]));
               exit(0);
      }
      /* get the default attributes */
      pthread_attr_init(&attr);   NULL도 가
      /* create the thread */
      pthread_create(&tid,&attr,runner,argv[1]);   10
      /* now wait for the thread to exit */
      pthread_join(tid,NULL);
      printf("sum  = %d\n",sum);
      return 0;
}

/* The thread will begin control in this function */
void *runner(void *param)
{
       int upper = atoi(param);   ask to integer
       int i = 0;
       sum = 0;
       if (upper > 0) {
              for (i = 1; i <= upper ; i++)
                            sum += i;
       }
       return NULL;
}
```

Windows Threads

- Create

```
HANDLE WINAPI CreateThread(
        LPSECURITY_ATTRIBUTES lpThreadAttributes,
        SIZE_T dwStackSize,
        LPTHREAD_START_ROUTINE lpStartAddress,
        LPVOID lpParameter,
        DWORD dwCreationFlags,
        LPDWORD lpThreadId
);
```

  See http://msdn.microsoft.com/library/default.asp?url=/library/en-us/dllproc/base/createthread.asp

- Wait

```
DWORD WINAPI WaitForSingleObject(
        HANDLE hHandle,
        DWORD dwMilliseconds                    // time-out interval
);
```

- Close (deallocate) handle

```
BOOL CloseHandle(LPDWORD lpThreadId);
```

Windows Multithreaded C Program

```c
#include <windows.h>
#include <stdio.h>
DWORD Sum; /* data is shared by the thread(s) */

/* the thread runs in this separate function */
DWORD WINAPI Summation(LPVOID Param)
{
   DWORD Upper = *(DWORD*)Param;
   for (DWORD i = 0; i <= Upper; i++)
      Sum += i;
   return 0;
}

int main(int argc, char *argv[])
{
   DWORD ThreadId;
   HANDLE ThreadHandle;
   int Param;

   if (argc != 2) {
      fprintf(stderr,"An integer parameter is required\n");
      return -1;
   }
   Param = atoi(argv[1]);
   if (Param < 0) {
      fprintf(stderr,"An integer >= 0 is required\n");
      return -1;
   }

   /* create the thread */
   ThreadHandle = CreateThread(
     NULL, /* default security attributes */
     0, /* default stack size */
     Summation, /* thread function */
     &Param, /* parameter to thread function */
     0, /* default creation flags */
     &ThreadId); /* returns the thread identifier */

   if (ThreadHandle != NULL) {
      /* now wait for the thread to finish */
     WaitForSingleObject(ThreadHandle,INFINITE);

     /* close the thread handle */
     CloseHandle(ThreadHandle);

     printf("sum = %d\n",Sum);
   }
}
```

Java Threads

- Java threads are managed by the JVM
  - Typically, implemented using the threads model provided by underlying OS
- Java threads may be created by:
  - Extending Thread class
    - For detail, search 'Java Thread class' from Internet
  - Implementing the Runnable interface implement는 여러개 할수 있다.

```
public interface Runnable
{
    public abstract void run();
}
```

Java Thread using Thread Class

- Extending Thread class

```
class PrimeThread extends Thread {
    long minPrime;
    PrimeThread(long minPrime) {
        this.minPrime = minPrime;
    }

    public void run() {
        // compute primes larger than minPrime
         . . .
    }
}
```

- Launching thread

```
PrimeThread p = new PrimeThread(143);
p.start();
```

Java Thread using Running Interface

- Extending Thread class

```
class PrimeRun implements Runnable {
    long minPrime;
    PrimeRun(long minPrime) {
        this.minPrime = minPrime;
    }

    public void run() {
        // compute primes larger than minPrime
         . . .
    }
}
```

- Launching thread

```
PrimeRun p = new PrimeRun(143);
new Thread(p).start();
```

Java Multithreaded Program

```
class Sum
{
  private int sum;

  public int getSum() {
   return sum;
  }

  public void setSum(int sum) {
   this.sum = sum;
  }
}

class Summation implements Runnable
{
  private int upper;
  private Sum sumValue;

  public Summation(int upper, Sum sumValue) {
   this.upper = upper;
   this.sumValue = sumValue;
  }

  public void run() {
   int sum = 0;
   for (int i = 0; i <= upper; i++)
      sum += i;
   sumValue.setSum(sum);
  }
}
```

Java Multithreaded Program (Cont.)

```
public class Driver
{
   public static void main(String[] args) {
     if (args.length > 0) {
       if (Integer.parseInt(args[0]) < 0)
        System.err.println(args[0] + " must be >= 0.");
       else {
        Sum sumObject = new Sum();
        int upper = Integer.parseInt(args[0]);
        Thread thrd = new Thread(new Summation(upper, sumObject));
        thrd.start();
        try {
           thrd.join();
           System.out.println
                   ("The sum of "+upper+" is "+sumObject.getSum());
       } catch (InterruptedException ie) { }
       }
     }
     else
       System.err.println("Usage: Summation <integer value>"); }
}
```

Implicit Threading

- Creation and management of threads done by compilers and run-time libraries rather than programmers

- Three methods explored
  - Thread Pools
  - OpenMP
  - Grand Central Dispatch

- Other methods include Microsoft Threading Building Blocks (TBB), java.util.concurrent package

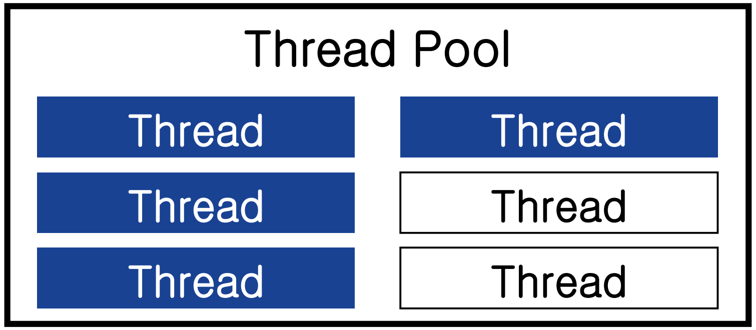
Thread Pools

Thread의 생성과 해제가 잦은경우

(overhead를 줄일 수 있다.

thread 전체 갯수를 조절하기 쉽다. → 안정성↑

- Create a number of threads in a pool where they await work
- Advantages
  - Usually slightly faster to service a request with an existing thread than create a new thread
  - Allows the number of threads in the application(s) to be bound to the size of the pool
  - Separating task to be performed from mechanics of creating task allows different strategies for running task
    - i.e., Tasks could be scheduled to run periodically



Windows Thread Pool (New API)

- See https://docs.microsoft.com/en-us/windows/win32/procthread/thread-pools
- Creating/Closing Thread pool
  - PTP_POOL pool = CreateThreadpool(NULL);
  - void CloseThreadpool(PTP_POOL ptpp);
- Setting maximum/minimum number of threads
  - void SetThreadpoolThreadMaximum(PTP_POOL ptpp, DWORD cthrdMost);
  - BOOL SetThreadpoolThreadMinimum(PTP_POOL ptpp, DWORD cthrdMic);
- Create and associate a callback environment to a thread pool
  - TP_CALLBACK_ENVIRON CallBackEnviron;
  - InitializeThreadpoolEnvironment(&CallBackEnviron);
  - SetThreadpoolCallbackPool(&CallBackEnviron, pool);
  - https://docs.microsoft.com/en-us/archive/msdn-magazine/2011/september/windows-with-c-the-thread-pool-environment
- Create a cleanup group
  - PTP_CLEANUP_GROUP cleanupgroup = CreateThreadpoolCleanupGroup();
  - SetThreadpoolCallbackCleanupGroup(&CallBackEnviron, cleanupgroup, NULL);
  - https://docs.microsoft.com/en-us/archive/msdn-magazine/2011/october/windows-with-c-thread-pool-cancellation-and-cleanup

Windows Thread Pool (New API)

- See https://docs.microsoft.com/en-us/windows/win32/procthread/thread-pools

- Create work
  - PTP_WORK work = CreateThreadpoolWork(workcallback, NULL, &CallBackEnviron);

- Submit a work to the thread pool
  - SubmitThreadpoolWork(work);

- Close all members of clearnup group to finish
  - CloseThreadpoolCleanupGroupMembers(cleanupgroup, FALSE, NULL);

컴파일러와 함께 사용되는 Tool

OpenMP

- Provides support for parallel programming in shared-memory environments
  - Set of compiler directives and an API for C, C++, FORTRAN
  - Identifies parallel regions – blocks of code that can run in parallel

- Create as many threads as there are cores or H/W threads

```
#pragma omp parallel                    // each runs the statement
    printf("Hello, World!₩n");
```

컴파일러에게 주는 특수 명령

CPU 코어수만큼 병렬적으로 수행

- Run for loop in parallel

```
#pragma omp parallel for                // unroll loop over cores
    for(i=0;i<N;i++) {
        c[i] = a[i] + b[i];
    }
```

For 문을 나눠서 실행

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- Overview
- Multithreading models
- Thread libraries
- <u>Threading issues</u>
- Operating system examples

Threading Issues

- fork() and exec()
- Cancellation
- Signal handling
- Thread-local storage
- Scheduler activation (already covered)

fork() and exec()

- fork() on multithreaded process
  - Duplicates all threads in the process?
  - Duplicates only corresponding thread?
  - → UNIX supports two versions of fork
    - fork(), fork1()

      프로세스 다 복사 스레드만 복사

- exec() on multithreaded process
  - Replace entire process

Cancellation

- Thread cancellation
  - Terminating a thread (target thread) before it has completed

```
pthread_t tid;

/* create the thread */
pthread_create(&tid, 0, worker, NULL);

. . .

/* cancel the thread */
pthread_cancel(tid);
```

너 죽어 X 죽여주시겠습니까

- Problem with thread cancellation
  - A thread share the resource with other threads

  cf. A process has its own resource.

  → A thread can be cancelled while it updates data shared with other threads

Cancellation

- Setting thread cancellation type
  - pthread_t pthread_setcanceltype(int type, int *oldtype);
    - Asynchronous cancellation (PTHREAD_CANCEL_ASYNCHRONOUS)
      - Immediate termination
    - Deferred cancellation (PTHREAD_CANCEL_DEFERRED)
      - Target thread checks periodically whether it should terminate.

        Ex) pthread_testcancel() 취소 되어 확인
      - Safer than asynchronous cancellation

- Enabling/disabling thread cancelation
  - pthread_t pthread_setcancelstate(int state, int *oldstate);
    - state: PTHREAD_CANCEL_DISABLE or PTHREAD_CANCEL_ENABLE
    - PTHREAD_CANCEL_DISABLE: cancellation remains pending until thread enables it

Mode	State	Type
Off	Disabled	–
Deferred	Enabled	Deferred
Asynchronous	Enabled	Asynchronous

Signal Handling

- Signal: mechanism provided by UNIX to notify a process a particular event has occurred
  - A signal can be generated by various sources.
  - The signal is delivered to a process.
  - The process handles it.
    - Default signal handler (kernel)
    - User-defined signal handler

- Types of signal
  - Synchronous: signal from same process

    Ex) illegal memory access, division by 0
  - Asynchronous: signal from external sources

    Ex) <Ctrl>-C

Signal

- A signal is a software interrupt delivered to a process.
  - The operating system uses signals to report exceptional situations to an executing program
- A signal is a limited form of inter-process communication used in Unix and other POSIX-compliant operating systems.
  - Essentially it is an asynchronous notification sent to a process in order to notify it of an event that occurred

Examples of UNIX Signals

- Signals in UNIX (in signal.h)

```
#define SIGHUP     1  /* hangup */
#define SIGINT     2  /* interrupt */
#define SIGQUIT    3  /* quit */
#define SIGILL     4  /* illegal instruction (not reset when caught) */
#define SIGTRAP    5  /* trace trap (not reset when caught) */
#define SIGIOT     6  /* IOT instruction */
#define SIGABRT    6  /* used by abort, replace SIGIOT in the future */
#define SIGEMT     7  /* EMT instruction */
#define SIGFPE     8  /* floating point exception */
#define SIGKILL    9  /* kill (cannot be caught or ignored) */
#define SIGBUS     10 /* bus error */
#define SIGSEGV    11 /* segmentation violation */
#define SIGSYS     12 /* bad argument to system call */
#define SIGPIPE    13 /* write on a pipe with no one to read it */
#define SIGALRM    14 /* alarm clock */
#define SIGTERM    15 /* software termination signal from kill */
```

Installing Signal Handler

- **Defining signal handler**

```
typedef void (*sighandler_t)(int);
sighandler_t signal(int signum,
                    sighandler_t handler);
```

- **Example**

```
#include <signal.h>
#include <unistd.h>

void sig_handler(int signo);

// press CTRL-₩ to terminate this program
int main()
{
    int i = 0;
    signal(SIGINT, (void *)sig_handler);

    while(1){
        printf("%d₩n", i++);
        sleep(1);
    }
    return 1;
}

void sig_handler(int signo)
{
    printf("SIGINT was received!₩n");
}
```

Signal Handling

- Question: To what thread the signal should be delivered?

- Possible options
  - To the thread to which the signal applies
  - To every thread in the process
  - To certain threads in the process
  - Assign a specific thread to receive all signals

  ➔ depend on type of signal

- Another scheme: specify a thread to deliver the signal

  Ex) pthread_kill(tid, signal) in POSIX

  signal 보낼때 사용

Thread-Local Storage global 변수를 thread 별로 할당

- In a process, all threads share global variables
- Sometimes, thread-local storage(TLS) is required
  - Many OS's support thread specific data

Ex) A key is assigned to each thread.

```
+------------------------------------+
| Process                            |
|                                    |
|     Global_data                    |
|     +-------+-------+-------+      |
|     | data1 | data2 | data3 |      |
|     +-------+-------+-------+      |
|     ^       ^       ^              |
|     |       |       |              |
|     |       |       |              |
|    key1    key2    key3            |
|     |       |       |              |
|   +---+   +---+   +---+            |
|   |   |   |   |   |   |            |
|   | A |   | B |   | C |            |
|   +---+   +---+   +---+            |
+------------------------------------+
```

Thread-Local Storage in pthread

- Thread-local storage (TLS) allows each thread to have its own copy of data

  Ex) __thread int tls; global 변수 // on pthread

  - Each thread has its own 'int tls' variable

  - Different from local variables
    - Local variables visible only during single function invocation
    - TLS visible across function invocations

  - Similar to static data
    - TLS is unique to each thread

- Useful when you do not have control over the thread creation process (i.e., when using a thread pool)

Thread-Local Storage in pthread

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>

#define THREADS 3

__thread int tls;
int global;

void *func(void *arg)
{
    int num = (int) arg;
    tls = num;
    global = num;
    sleep(1);
    printf("Thread = %d tls = %d global = %d₩n",
            num, tls, global);
}
```

```
int main()
{
    int ret;
    pthread_t thread[THREADS];
    int num;

    for(num = 0; num < THREADS; num++){
        ret = pthread_create(&thread[num], NULL, &func,
                                    (void *)num);
        if(ret){
            printf("error pthread_create₩n");
            exit(1);
        }
    }

    for(num = 0; num < THREADS; num++){
        ret = pthread_join(thread[num], NULL);
        if(ret){
                printf("error pthread_join₩n");
                exit(1);
        }
    }

    return 0;
}
```

Agenda

- Overview
- Multithreading models
- Thread libraries
- Threading issues
- Operating system examples skip

Windows XP Threads

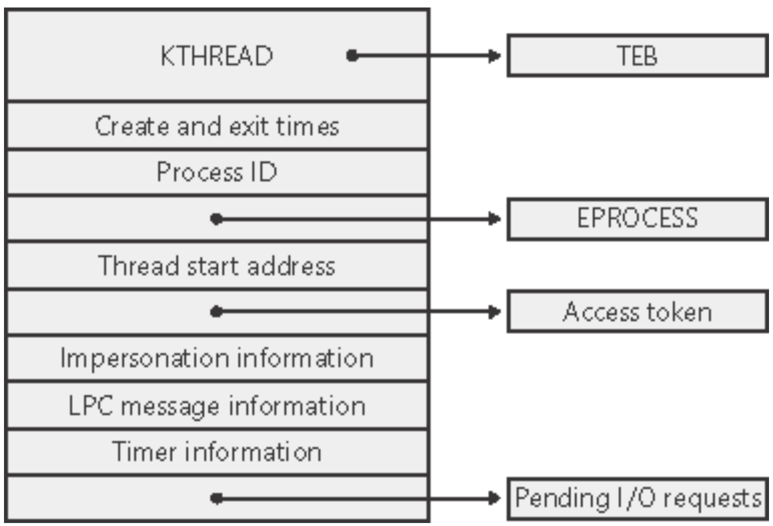
- Implements the one-to-one mapping, kernel-level
  - Additionally, many-to-many model is supported by fiber library
- Each thread contains
  - A thread id, register set, separate user and kernel stacks, private data storage area (for run-time libraries and DLLs)

  Cf. The register set, stacks, and private storage area are known as the context of the thread.

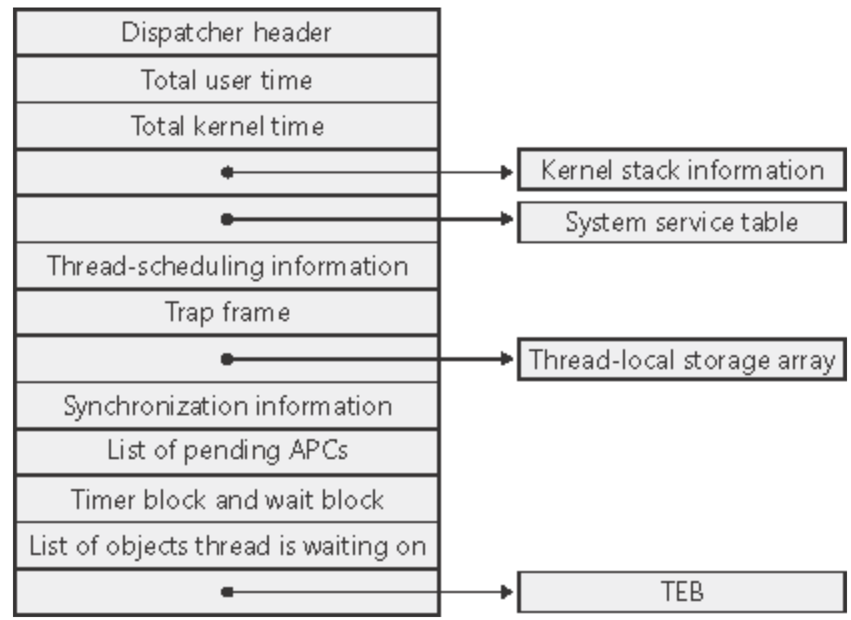
Internal Data Structures

- The primary data structures of a thread include:
  - ETHREAD (executive thread block)
  - KTHREAD (kernel thread block)
    - Information for thread scheduling and synchronization
  - TEB (thread environment block)
    - Context information for the image loader and various Windows DLL

Layout of ETHREAD and KTHREAD

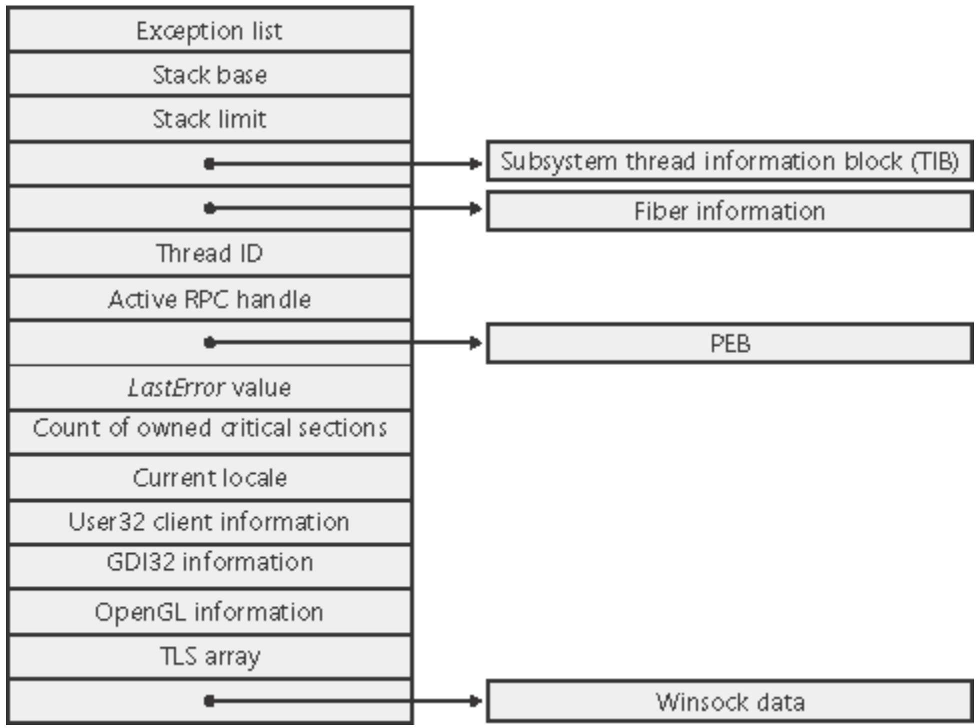


ETHREAD

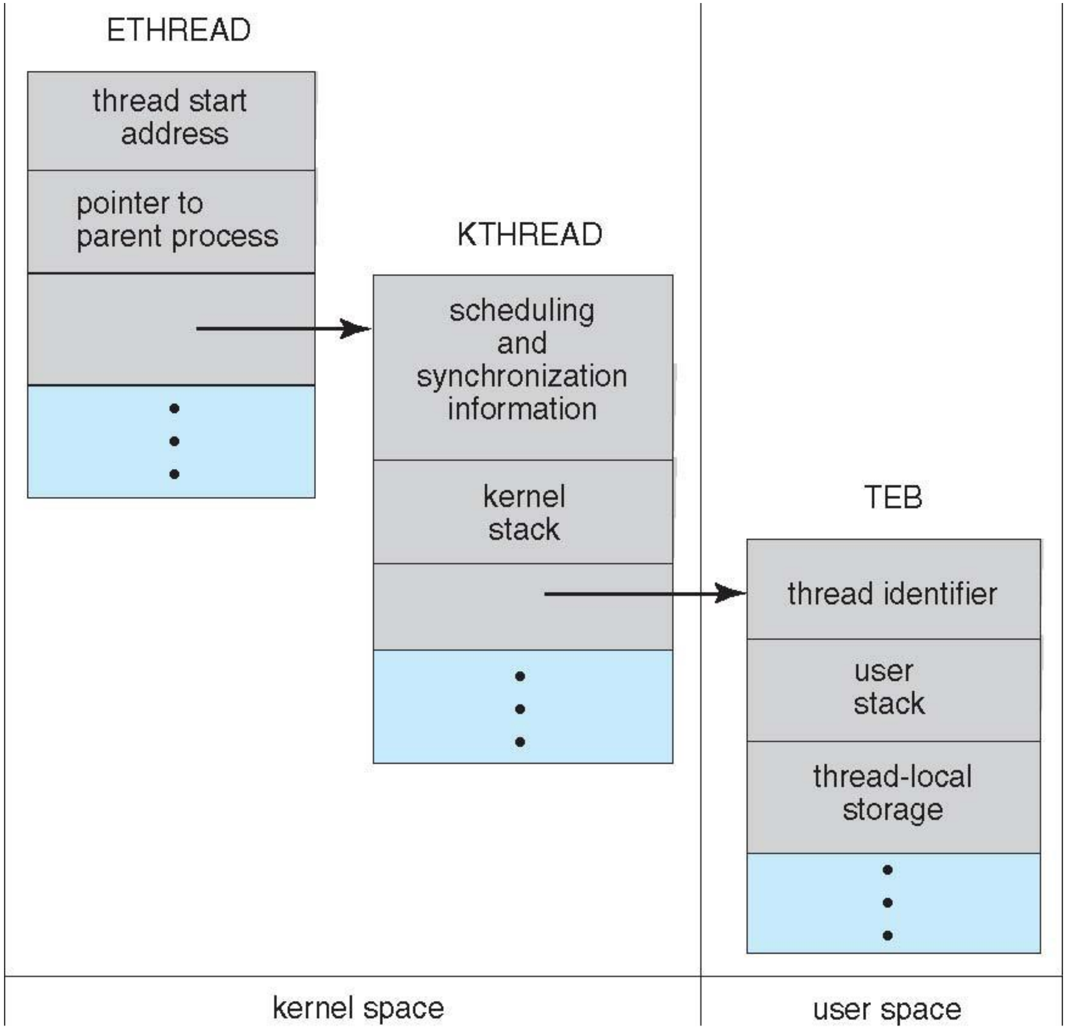


KTHREAD

Layout of TEB



Windows XP Threads



Linux Threads

- First implementation: LinuxThreads
  - Thread is created by clone() system call
  - However, no difference between process and thread
    - Term '**task**' is used more frequently than 'process' or 'thread'
  - User can control how much the resource is shared between parent and child.

Linux Threads

- Problems of *LinuxThreads*
  - It did not scale well,
    - There is a limitation on the number of threads, generally between 1024 and 8192.
    - The overhead of creating and destroying processes is relatively high.
  - The manager thread resulted in some fragility and another scaling bottleneck.
  - Some required POSIX semantics were not possible. Each thread had its own PID.
- NPTL (Native POSIX Thread Library)
  - Better than LinuxThreads in performance and compatibility