Recursive compile time adjoint in C++ 23rd EuroAD Workshop

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Introduction

- The compile time differentiation of a C++ function which only calls built-in functions has been demonstrated in previous talks.
- This talk describes the approach for differentiating a function which may call any nested function.
- Siemen's Simcenter Star-CCM+ simulation software has an implementation of this approach, and is used to differentiate the Spalart Allmaras turbulence model, among other things.

What's new

Flat structure (old)

Nested structure (new)

Built-in functions

Hypotenuse

$$r = \sqrt{a^2 + b^2}$$

```
float a = 3;
float b = 4;
float r;

{
   auto d = a*a + b*b;
   r = sqrt(d);
}

std::cout << r << std::endl; // r = 5</pre>
```

Primal of Hypotenuse

$$r = \sqrt{a^2 + b^2}$$

```
float a = 3;
float b = 4;
float r;
auto constexpr mode = DrvMode::PRIMAL;
Drv<mode, float> a_{a}; // input
Drv<mode, float> b_{b}; // input
Drv<mode, float&> r_{r}; // output
 auto d = a_*a_+ + b_*b_;
 r_{-} = sqrt(d);
std::cout << r << std::endl; // r = 5
```

Tangents of Hypotenuse

 $\frac{dr}{da}$

```
float a = 3, a_drv = 1; // w.r.t. 'a'
float b = 4, b_drv = 0;
float r drv:
auto constexpr mode = DrvMode::TANGENT;
Drv<mode, float> a_{a, a_drv};
Drv<mode, float> b_{b, b_drv};
Drv<mode, float&> r_{r_drv};
 auto d = a_*a_+ + b_*b_;
 r_{-} = sqrt(d);
std::cout << r_drv << std::endl; // dr/da = 0.6
```

Tangents of Hypotenuse

 $\frac{dr}{db}$

```
float a = 3, a_drv = 0;
float b = 4, b_drv = 1; // w.r.t. 'b'
float r drv:
auto constexpr mode = DrvMode::TANGENT;
Drv<mode, float> a_{a, a_drv};
Drv<mode, float> b_{b, b_drv};
Drv<mode, float&> r_{r_drv};
 auto d = a_*a_+ + b_*b_;
 r_{-} = sqrt(d);
std::cout << r_drv << std::endl; // dr/db = 0.8
```

Adjoint of Hypotenuse

 $\begin{bmatrix} \frac{dr}{da} & \frac{dr}{db} \end{bmatrix}^T$

```
float a = 3, a_drv = 0;
float b = 4, b_drv = 0;
float r_drv = 1; // w.r.t. 'r'
auto constexpr mode = DrvMode::ADJOINT;
Drv<mode, float> a_{a, a_drv};
Drv<mode, float> b_{b, b_drv};
Drv<mode, float&> r_{r_drv};
 auto d = a * a + b * b:
 r_{-} = sqrt(d);
std::cout << a_drv << std::endl; // dr/da = 0.6
std::cout << b_drv << std::endl; // dr/db = 0.8
```

User defined functions

Defining a function

```
struct Discrimiant
  template < DrvMode mode > // 'float' could be templated, too
  static void
  evaluate(Drv<mode, float> const &a,
           Drv<mode, float> const &b,
           Drv<mode, float&> r)
    auto a2 = a * a;
    auto b2 = b * b;
    r = a2 + b2;
```

 All function arguments are treated as differentiable terms, even if in the function body some may be treated passively, e.g.

```
auto a2 = primal(a * a); // i.e. float a2 = a * a;
```

Calling it via a free function

Unified call syntax

```
Drv<mode, float> a_{a, a_drv};
Drv<mode, float> b_{b, b_drv};
Drv<mode, float&> r_{r_drv};

{
   auto d = discrimiant(a_, b_);
   r_ = sqrt(d);
}
```

- discrimiant can be used just like any built-in function, such as sqrt.
- a or b could be replaced with a or b, respectively, to treat either passively.

The variadic expression node

Unary, binary, . . . and variadic

- Allied with variadic templates, perfect forwarding and std::tuple, unary and binary expression nodes can be generalised into an *N*-ary form.
- In principle, all built-in functions, such as + * / and sin cos tan pow, could return an *N*-ary node.
- A mechanism to distinguish between built-in and user defined functions is still required. This is done with an operator binding, such as Bind_evaluate<0P>.

Perfect forwarding

```
template < class OP, class E0>
struct Unary { Unary(E0 &&e0) {} };
class Sqrt;
template < class E0>
auto sqrt(E0 &&e0)
 return Unary < Sqrt, E0 > (std::forward < E0 > (e0));
float const a0{1};
auto a1 = sqrt(a0);  // Unary<Sqrt, float const &>
float b0{1};
auto b1 = sqrt(b0); // Unary<Sqrt, float &>
auto c1 = sqrt(float{1}); // Unary<Sqrt, float>
```

Base of all nodes

```
template < DrvMode m, typename R, typename E>
struct DrvExpression
{
   static constexpr auto mode = m;
   using Result = R;
   using Expression = E;
};
```

Variadic node*

```
template < DrvMode m,
                         // mode
        typename R, // result
        typename OP, // operator
        typename... EE> // children
struct DryVariadicNode
  : DrvExpression<m, R, DrvVariadicNode>
 DrvVariadicNode(EE &&... ee) : _ee(ee...) {}
 template<int I> auto const &node() const {
   return std::get<I>(_ee);
 private: std::tuple<EE...> _ee;
```

Variadic node

- The *N*-ary node is minimal, providing nothing more than access to its children.
- Children will be associated to the node by value, for in-place expressions or literals, or by reference. This difference is distinguished by perfect forwarding.
- The values and references are held in the std::tuple member.

Evaluating the primal*

```
template < DrvMode m,
         typename OP, typename R, typename... EE,
         typename AA, int... I>
auto
evaluatePrimal(
  DrvVariadicNode<m, R, OP, EE...> const &expr,
  AA const &aa,
                                        // primal values: a. b
  std::index_sequence<I...>)
  auto constexpr lm = DrvMode::PRIMAL; // local mode
  auto r_pri = R(0);
  OP::evaluate(
    Drv<lm, decltype(DrvVariadicNode::node<I>())::Result>
      {std::get<I>(aa)}...,
                                    // inputs
    Drv < lm, R&> {r_pri});
                                      // output
  return r_pri;
```

Evaluating the primal

- All work relating to differentiable expressions is done by free functions.
- The expression mode, m, is distinguished from the local mode, lm, in order to generate the correct instance of the user defined function.
- The arguments for OP::evaluate are constructed in-place, generating a new (and independent) differentiation context inside the user defined function.

Evaluating the adjoint*

```
template < DrvMode m,
         typename OP, typename R, typename... EE,
         typename AA, typename RHS, int... I>
void
evaluateAdjoint(
 DrvVariadicNode<m, R, OP, EE...> const &expr,
 AA const &aa.
                           // primal values: a, b
                            // adjoint refs: a_drv, b_drv
 AA &aa_drv,
 RHS const &rhs,
                            // seed value: r_drv
  std::index_sequence<I...>)
 OP::evaluate(
   Drv<m, decltype(DrvVariadicNode::node<I>())::Result>
      {std::get<I>(aa), std::get<I>(aa_drv)}...,
   Drv < m, R\& > \{rhs\});
```

Evaluating the adjoint

- Here, OP::evaluate is compiled in adjoint mode, as deduced from the arguments.
- Despite the fact that OP::evaluate will have been evaluated in primal mode already, the values of any intermediate results cannot be captured and used here because there is no monolithic expression tree 'view'.
- The work presented in EuroAD 2019 did have such a view, and so avoided duplicate primal evaluations. But compilation time proved to be too expensive.

Summary

- The variadic template expression node, combined with perfect forwarding, facilitates generic user defined functions.
- User defined virtual functions can be supported, and so too functions with multiple results.
- Despite being difficult to implement, the added functionality of variadic nodes has proved to be a major advancement in deploying automatic differentiation in Star-CCM+.
- The tool is fully constexpr qualified, and with aggressive compiler optimisations performs very well: 2.82x on the harmonic test (331 nodes, 5 inputs); cf. EuroAD 2019.

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