# Compile time adjoint in C++ 22<sup>nd</sup> EuroAD Workshop, Imperial College

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## Overview

- **1 Preliminaries:** resources, basic ideas, etc
- Mapping indices: separate tree structure and its data
- **3** Language extensions: functionality that would be very helpful

## **Preliminaries**

#### Zero or more

```
// no state, but instantiable
template < class ... T> struct list {};
// a more useful form...
template < class T, T...> struct seq {};
// _the_ operation
template < class L> struct front;
template < template < class ... > class L, class T1, class... T>
struct front<L<T1, T...> >
  using type = T1;
```

#### Simple C++11 metaprogramming

With variadic templates, parameter packs and template aliases

Peter Dimov. 26.05.2015

I was motivated to write this after I read Eric Niebler's thought-provoking Tiny Metaprogramming Library article. Thanks Eric.

#### C++11 changes the playing field

The wide acceptance of <u>Boost.MPL</u> made C++ metaprogramming seem a solved problem. Perhaps MPL wasn't ideal, but it was good enough to the point that there wasn't really a need to seek or produce alternatives.

C++11 changed the playing field. The addition of variadic templates with their associated parameter packs added a compiletime list of types structure directly into the language. Whereas before every metaprogramming library defined its own type list, and MPL defined several, in C++11, type lists are as easy as

```
// C++11
template<class... T> struct type_list {};
```

and there is hardly a reason to use anything else.

Template aliases are another game changer. Previously, "metafunctions", that is, templates that took one type and produced another, looked like

```
^{//\ C++\theta 3} template<class T> struct add_pointer { typedef T* type; };
```

and were used in the following manner:





MappedSurface Sensitivity of Force w.r.t. Normal Displacement -0.00050000 -0.00330000 -0.00010000 0.00010000 0.00030000

0.00050000

Compressible N-S adjoint: automatic at loop iteration level

# Differentiate a function losslessly

# Parity preserving transform

```
fn(A const &a, B const &b,
  R &r)
 auto c0 = 7;
 auto c1 = 9;
 auto t0 = a * b; // 1
 auto t1 = c0 + t0; // 2
 auto t2 = c1 + t0; // 3
 r = t1 / t2; // 4
```

write something like this ... to implement something like this

```
fn(A const &a, B const &b,
  R &r)
 auto c0 = 7;
 auto c1 = 9;
 auto t0 = a * b; // 1
 auto t1 = c0 + t0; // 2
 auto t2 = c1 + t0; // 3
// somehow add this...
 t1.d += (1/t2) * r.d; // 4
 t2.d = (t1/t2^2) * r.d; // 4
 t0.d += t2.d: // 3'
 t0.d += t1.d; // 2'
 a.d += b * t0.d; // 1'
 b.d += a * t0.d; // 1'
```

#### **Observations**

- Given a pure-functional algorithm, differentiate it
- 2 Implement the transpose of the chain of derivatives (the adjoint)
- The required 'extra' code is in the reversed sequence of the original and the data flow is reversed
- Eager and lazy evaluation: ctor-dtor pairs?

## Two hurdles

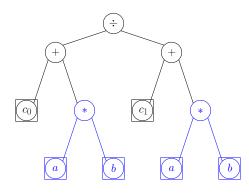
## 1. Dealing with duplicate nodes

Eager evaluation and capture by reference?

```
fn(A const &a, B const &b,
    R &r)
{
  auto c0 = 7;
  auto c1 = 9;

auto t0 = a * b;
  auto t1 = c0 + t0;
  auto t2 = c1 + t0;

r = t1 / t2;
}
```



# 2. Dealing with nested scoping

The *complete* tree, including **cm**, is needed for the transform

```
mul_dbl(A const &a, B const &b)
  auto cm = 2; // locally scoped
  return cm * a * b;
fn(A const &a, B const &b, R &r)
  auto c0 = 7;
  auto c1 = 9;
  auto t0 = mul_dbl(a, b);
  auto t1 = c0 + t0;
  auto t2 = c1 + t0;
  r = t1 / t2;
```

# 2. Dealing with nested scoping

The *complete* tree, including cm, is needed for the transform

```
fn(A const &a, B const &b,
  R &r)
  auto c0 = 7;
 auto c1 = 9;
  auto cm = 2;
  auto t0 = cm * a * b;
  auto t1 = c0 + t0;
  auto t2 = c1 + t0;
 r = t1 / t2;
```

```
fn(A const &a, B const &b,
  R &r)
  auto c0 = 7;
  auto c1 = 9;
 auto cm = 2;
auto t0 = cm * a * b;
 auto t1 = c0 + t0;
  auto t2 = c1 + t0:
  // the transform...
  t1.d += (1/t2) * r.d;
 t2.d = (t1/t2^2) * r.d;
  t0.d += t2.d;
  t0.d += t1.d;
  a.d += cm * b * t0.d;
  b.d += cm * a * t0.d;
```

## State of affairs

- Eager evaluation avoids duplicate branch evaluation but lazy evaluation will also be needed
- 'Capture by reference' to keep the tree small but cannot work with nested scoping
- Capture by value' is too inefficient the tree will get very large very quickly
- 4 Monolithic tree, supporting eager and lazy evaluation, of minimal size, and impartial to scoping is required

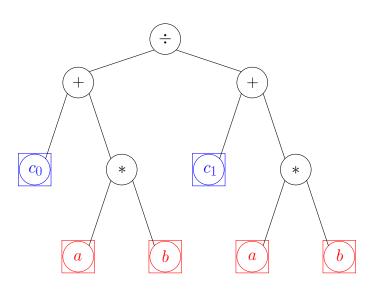
# Expression tree to type list

## Two kinds of data

```
auto fn(A const &a, B const &b)
{
  auto c0 = UQ(7);
  auto c1 = UQ(9);

auto t0 = a * b;
  auto t1 = c0 + t0;
  auto t2 = c1 + t0;
  return t1 / t2;
}
```

## Two kinds of data



```
template < std::size_t ID, typename T>
struct Unique
{
   T value;
};
```

```
// UQ
#define UQ(v) Unique<__COUNTER__, decltype(v)>{v}
```

## Tree type to list of types

## Generate tree with operator overloading

```
Binary<Div,
Binary<Add, C0,
Binary<Mul, A, B> >
Binary<Add, C1,
Binary<Mul, A, B> >
```

Group hierarchically and prune duplicates (at each binary node), then flatten (at the root node)

```
{A,
    B,
    Binary<Mul, A, B>,
    Binary<Add, C0, [...]>,
    Binary<Add, C1, [...]>,
    Binary<Div, [...], [...]>}
```

# Mapping nodes to data

#### Map CO and C1 to values

```
tuple<float, float> vars = {7.0, 9.0};
```

Indices of constants in 'L'

```
IC = \{3, 4\}
```

Map elements in 'L' to storage offsets (2 = null marker)

```
DC = \{2, 2, 2, 0, 1, 2\}
```

#### dual

```
// input output
//-----
DC_SIZE = 6 \
--- dual ---> DC = {2, 2, 2, 0, 1, 2}
IC = {3, 4} / 0 1 2 3 4 5
```

github.com/DominicJones/snippets/blob/master/Cxx/mp\_functions.cpp

# Mapping nodes to data

## Map A and A to addresses

```
tuple<float*, float*> args = {&a, &b};
```

Map offsets of left child nodes (6 = null marker)

```
IL_L = \{6, 6, 0, 6, 6, 3\}
```

Map offsets of right child nodes

```
IL_R = \{6, 6, 1, 2, 2, 4\}
```

# Evaluate operator list

```
tuple<float, float> vars = {7.0, 9.0};

tuple<float*, float*> args = {&a, &b};

DC = {2, 2, 2, 0, 1, 2}

IL_L = {6, 6, 0, 6, 6, 3}

IL_R = {6, 6, 1, 2, 2, 4}
```

## Iterate list to compute primal and adjoint

## Results

Case 1: nodes: 82, depth: 37, inputs: 2, constants: 12

Version	compilation	original	auto diff	manual diff
alt::tuple	2.2s	1x	1.25x	1.48x
std::tuple	4.5s	1x	1.25x	1.48x

**Case 2**: nodes: 331, depth: 25, inputs: 5, constants: 103

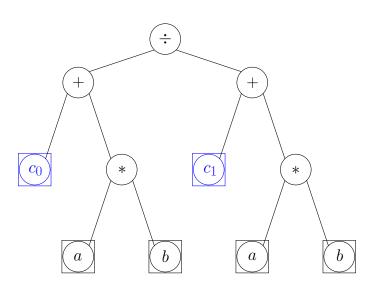
Version	compilation	original	auto diff	manual diff
alt::tuple	59s	1x	5.7x	1.9x
std::tuple	27m	1x	4.7x	1.9x

## Conclusion

- Manipulation the tree is obtained at immense effort
- Works in the range of exceptionally well (better than hand coded) to acceptably well (better than alternatives)
- Ompile-time features of the language are too limited to use this approach neatly (cf. preprocessor macros)
- Inlining gives up too readily; \_\_attribute\_\_((always\_inline)) used ubiquitously
- Only supports having one result
- On to obvious what is going on with std::tuple

# Float template parameter

# Distinguishing different constants



# **△** *Values as types* **△**

```
// C++ does not permit 'auto' to resolve as 'float'
template < auto V >
struct _float
{
   constexpr operator auto() const { return V; }
};
```

```
// type distinguished by value
auto constexpr c0 = -4.2_f;
static_assert(c0 == _float<-4.2>{});
```

# Values as types in D

```
// exactly what is wanted
struct _float(float v)
{
   static immutable auto value = v;
}
```

## "\_float" workaround

```
// exponent ignored...
template<auto H, auto L, auto E>
struct _float
{
   auto constexpr static value =
      (H + float(L) / multiplier<10, E, 1>::value);
   constexpr operator auto() const { return value; }
};
```

```
// and for operator+
template<auto H, auto L, auto E>
auto constexpr operator-(_float<H, L, E>)
{
   return _float<(-H), (-L), E>{};
}
```

## ... made palatable

```
// makes life easier
template < char...> struct mp_chars {};
```

```
// seamless conversion to literals
auto constexpr c0 = -4.2_f;
float t0 = 2 * c0;
```

# Reflect variable location

## The trusty preprocessor

```
auto fn(A const &a, B const &b)
{
  auto c0 = UQ(7); // auto = Unique < 724, float > {7}
  auto c1 = UQ(9); // auto = Unique < 725, float > {9}

auto t0 = a * b;
  auto t1 = c0 + t0;
  auto t2 = c1 + t0;
  return t1 / t2;
}
```

# A fundamental problem

Write a transform function to yield:

```
auto c0 = 7;
auto c1 = 9;
transform(c0 + c0) --> 2 * c0
transform(c1 + c0) --> c1 + c0
```

*Impossible!* ... despite it appearing so trivial

## Three kinds of reflection?

Reflect *location* of **x**, at [7, 4, main.cpp]?

## Reflect location in D

```
Terminal!double c0 = 7; // Terminal<double, "main.d", 724>
Terminal!double c1 = 9; // Terminal<double, "main.d", 725>
```

One definition per line could be run-time checked with a debug build which requires all Terminals be singletons.

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