

EXPERIMENT - 10

Queue Using Linked List

AIM

Implementation of Queue using linked list.

Input:- Set of integers

Output:- Operation performed in queue.

ALGORITHM

1. Start
2. Create a structure 'node' containing datatype int to store the integers and struct node *next
3. struct node *head, *rear.
4. Choose from the options 1- Insert 2- Delete
3- Display 4- Exit.
 - 4.1. If (choice = 1) then
~~Insert()~~
 - 4.2. If (choice = 2) then
~~delete()~~
 - 4.3. If (choice = 3) then
~~display()~~
 - 4.4. If (choice = 4) then
~~exit.~~
 - 4.5. else .
~~Display "Invalid Choice"~~

5. Repeat Step 4 till option = 4.

6. Stop.

Insert()

1. struct node *ptr

2. Input value, item

3. ptr::data = item

4. if (head == NULL) then

 head = ptr

 rear = ptr

 head::next = NULL

 rear::next = NULL

5. else.

 rear::next = ptr

 rear = ptr

 rear::next = NULL.

Delete()

1. struct node *ptr.

2. If (head == NULL) then

 Display "Underflow"

3. else.

 ptr = head

 head = head::next.

 free(ptr)

Display()

1. struct node *ptr
2. ptr = head
3. if (head == NULL) then

 Display "Empty Queue"

else.

 while (ptr != NULL)

 Display ptr.data

 ptr = ptr.next

Result

Program ran successfully and output obtained.

P.L.

EXPERIMENT-11

POLYNOMIAL ADDITION Using LINKED LIST

Aim

Implementation of Polynomial addition using linked list.

Input:- 2 polynomials

Output: Sum of the two polynomials

ALGORITHM

1. Start
2. Declare a structure 'node' with datatype 'int' for the coefficient and exponent and struct node* next.
3. struct node* p1 = NULL, *p2 = NULL, *sum = NULL
4. int i = 1
5. Choose from the options 1 - Enter polynomial 1
2 - Enter polynomial 2 3 - Add Polynomial 4 - Exit.
5.1. If (option = 1) then
 - 5.1.1. Input the coefficient.
 - 5.1.2. Input the exponent.
 - 5.1.3. p1 = Create (p1, co, exp)
- 5.2. If (option = 2) then
 - 5.2.1. Input the coefficient.
 - 5.2.2. Input the Exponent.
 - 5.2.3. p2 = Create (p2, co, exp)

5.3. If (option = 3) then

5.3.1 sum = polyAdd (p1, p2, sum)

5.3.2. Display sum.

5.4. If (option = 4) then

i = 0

6. Repeat step 5 till i=1

7. stop.

Struct node* Create (struct node* head, int c0, int exp)

1. Struct node *temp, *ph

2. If (head == NULL) then

temp. c0 = c0

temp. exp = exp

temp. next = NULL

head = temp.

3. else.

temp = head

while (temp. next != NULL)

temp = temp. next

ph. c0 = c0

ph. exp = exp

ph. next = NULL

temp. next = ph

4. return head.

struct node * polyAdd (struct node * p1, struct node * p2,
 struct node * sum)

1. struct node * poly1 = p1, * poly2 = p2, * res
2. If (poly1 == NULL & and poly2 == NULL) then
 - 2.1. sum = poly1
 - 2.2. return sum
3. else if (poly1 == NULL and poly2 != NULL) then
 - 3.1. sum = poly2
 - 3.2. return sum.
4. while (poly1 != NULL and poly2 != NULL)
 - 4.1. If (sum == NULL) then
 res = sum
 - 4.2. else
 res = res. next.
 - 4.3. If (poly1. exp > poly2. exp) then
 res. co = poly1. co
 res. exp = poly1. exp
 poly1 = poly1. next.
 - 4.4. else if (poly1. exp < poly2. exp) then
 res. co = poly2. co
 res. exp = poly2. exp
 poly2 = poly2. next.
 - 4.5. else if (poly1. exp == poly2. exp) then
 res. co = poly1. co + poly2. co
 res. exp = poly1. exp
 poly1 = poly1. next
 poly2 = poly2. next.

5. while ($\text{poly1} \neq \text{NULL}$)

$\text{res} = \text{res}.\text{next}$

$\text{res}.\text{co} = \text{poly1}.\text{co}$

$\text{res}.\text{exp} = \text{poly1}.\text{exp}$

$\text{poly1} = \text{poly1}.\text{next}$

6. while ($\text{poly2} \neq \text{NULL}$)

$\text{res} = \text{res}.\text{next}$

$\text{res}.\text{co} = \text{poly2}.\text{co}$

$\text{res}.\text{exp} = \text{poly2}.\text{exp}$

$\text{poly2} = \text{poly2}.\text{next}$

7. $\text{res}.\text{next} = \text{NULL}$

8. return sum

void display (struct node * head)

1. struct node * temp = head

2. while ($\text{temp} \neq \text{NULL}$)

Display $\text{temp}.\text{co}$, $\text{temp}.\text{exp}$

$\text{temp} = \text{temp}.\text{next}$.

Result

Program ran successfully and sum obtained

S.P.

EXPERIMENT - 12

SPARSE ADDITION

AIM

Implementation of Addition of Sparse Matrix

Input:- 2 sparse matrix

Output :- Sum of 2 sparse matrix

ALGORITHM

1. Start
2. Input 2 matrix of size $m \times n$

3. Convert sparse matrix to Triplet form
Triplet()

4. Addition of 2 matrix.

Add();

5. Display the sum,

6. Stop.

Triplet()

1. $a[0][0] = m$

2. $a[0][1] = n$

3. $k = 1$

4. for $i = 0$ to $m-1$

 for $j = 0$ to $n-1$

 if $[A[i][j]] \neq 0$

$a[k][0] = i$

$a[k][1] = j$

$a[k][2] = A[i][j]$

$K = K + 1$

5. $a[0][2] = K - 1$

6. Return $a[][]$

Add()

1. Input 2 matrix a, b in tripld form.

2. If $a[0][0] \neq b[0][0]$ || $a[0][1] \neq b[0][1]$ then
 Display "Addition not possible"

3. $\sigma = a[0][0]$

4. $C = a[0][1]$

5. $S[0][0] = \sigma$

6. $S[0][1] = C$

7. $t_1 = a[0][2]$

8. $t_2 = b[0][2]$

9. $i = j = 1; k = 1$

10. while ($i < t_1$) +4 ($j < t_2$)

 if ($a[i][0] < b[j][0]$)

$S[k][0] = a[i][0]$

$S[k][1] = a[i][1]$

$S[k][2] = a[i][2]$

$i = i + 1; k = k + 1$

 else if ($a[i][0] > b[j][0]$) then

$S[k][0] = b[j][0]$

$S[k][1] = b[j][1]$

$S[k][2] = b[j][2]$

$k = k + 1; j = j + 1$

else if ($a[i][j] < b[j][i]$) then

$$S[k][0] = a[i][0]$$

$$S[k][1] = a[i][1]$$

$$S[k][2] = a[i][2]$$

$i = i + 1$, $k = k + 1$;

else if ($a[i][j] > b[j][i]$) then

$$S[k][0] = b[j][0]$$

$$S[k][1] = b[j][1]$$

$$S[k][2] = b[j][2]$$

$j = j + 1$; $k = k + 1$;

else.

$$S[k][0] = a[i][0]$$

$$S[k][1] = a[i][1]$$

$$\cancel{S[k][2] = a[i][2]}$$

$k = k + 1$; $i = i + 1$;

12. while ($i < t_2$)

$$\cancel{S[k][0] = b[j][0]}$$

$$\cancel{S[k][1] = b[j][1]}$$

$$\cancel{S[k][2] = b[j][2]}$$

$k = k + 1$; $j = j + 1$;

13. $S[0][2] = k - 1$

EXPERIMENT- 13

TRANSPOSE OF SPARSE MATRIX

AIM.

Implementation of transpose of sparse matrix.

Input :- A sparse matrix

Output:- Transpose of matrix.

ALGORITHM.

1. Start

2. Input the sparse matrix, A with size $m \times n$

3. Convert to triplet form, g

Triplet()

4. Transpose()

5. Display Transpose of sparse matrix.

Triplet()

1. $a[0][0] = m$

2. $a[0][1] = n$

3. $k = 1$

4. for $i = 0$ to $m - 1$

 for $j = 0$ to $n - 1$

 if ($A[i][j] \neq 0$)

$a[k][0] = i$

$a[k][1] = j$

$a[k][2] = A[i][j]$; $k = k + 1$

5. $a[0][2] = k - 1$

6. Return $a[0][0]$

Transpose (2)

1. $T[0][0] = a[0][1]$

2. $T[0][1] = a[0][0]$

3. $T[0][2] = a[0][2]$

4. $t = a[0][2]$

5. $n = a[0][0]$

6. $k = 1$

7. For $i = 0$ to $n - 1$, $i++$

For $j = 1$ to t , $+t+$,

if ($a[j][1] == i$)

$T[k][0] = a[j][1]$

$T[k][1] = a[j][0]$

$T[k][2] = a[j][2]$

$k = k + 1$

Result

Program ran successfully and transpose of matrix obtained.

S.H.

EXPERIMENT-14

INFIX TO POSTFIX

AIM

Convert an infix expression to postfix and evaluate it.

Input:- Expression in infix form.

Output:- Postfix form and result.

ALGORITHM

1. Scan the I/p expression from left to right.
2. If the scanned character is an operand, then place it to o/p
3. else.
 - 3.1. if the precedence of the scanned operator is greater than the precedence of stack 'Top' operator, or if the stack is empty, then push the operator.
 - 3.2. else
 - i) repeatedly pop from stack & add to o/p , each operator which has the same or higher precedence than the scanned operator.
 - ii) Push the scanned operator onto stack

4) If the scanned character is opening bracket, push it to stack.

5) If the scanned character is closing bracket, pop the contents from the stack & O/P it until "c" is found. Discard "c" & "

6) Repeat steps 2 to 5 until the expression is scanned completely.

7) If the stack is not empty, pop the contents from the stack & add it to O/P.

8) Display the O/P

9) To evaluate the postfix expression push the operand.

10) pop the operand according to the operators present and find the result.

RESULT

Program ran successfully and postfix expression and result obtained.

✓

EXPERIMENT - 15

BINARY TREE TRAVERSAL

AIM

Write a program to implement traversal in a binary tree (recursive traversal)

Input:- A binary tree

Output:- post order, preorder and inorder traversal results.

ALGORITHM

1. Start
2. Define a structure 'node' having datatype "int" for data and struct node *LC & *RC
3. Input data of different nodes of the tree from user
4. Display the options, get input and call corresponding function.
5. Stop.

Create (ptr, item).

1. If (ptr ≠ NULL) then
 - 1.1. ptr·data = item
 - 1.2. Display "Node ptr has LST (Y/N)?"
 - 1.3. If (option == 'Y') then
 - 1.3.1. Allocate memory and let lcptr point to it.
 - 1.3.2. ptr·LC = lcptr

1.3.3. Create (lcptr, l.data)

1.4. else

ptr.LC = NULL.

1.5. Display "Node ptr has RST(Y/N)?"

1.6. If (option == 'Y') then

1.6.1. Allocate memory and let rcptr point to it

1.6.2. ptr.RC = rcptr

1.6.3. Create (rcptr, r.data)

1.7. else

ptr.RC = NULL.

Inorder (root)

1. Inorder (root, LC)

2. visit root

3. Inorder (root, RC)

Preorder (root)

1. visit (root)

2. Preorder (root-LC)

3. preorder (root.RC)

Postorder (root)

1. Postorder (root.LC)

2. Postorder (root.RC)

3. visit root

Result

~~Success~~ Program ran successfully and traversed through tree performed.

EXPERIMENT-16

BINARY TREE OPERATIONS

AIM

Perform Binary tree operations like search, insertion and deletion.

Input :- A binary tree

Output :- Operations on binary tree.

ALGORITHM

1. Start
2. Create a new node 'new'
3. Input the value of root node.
4. Create a binary tree.
5. Display Menu 1-Insert 2-Search 3-Delete
4-Display 5-Exit.
6. do
 - 6.1. Display "Enter the choice"
 - 6.2. Read the choice
 - 6.3. switch (choice)
 - case 1 : Insert()
break;
 - case 2 : Search()
break;
 - case 3 : Delete()
break;

case 4 :- Display the tree.
break

case 5 :- Exit(0)
break

default :- display " Enter a valid choice!"
while (choice != 4)

7. Stop.

Search (ptr, key)

1. If ptr-data ≠ key
2. i) if ptr.LC ≠ NULL

 search (ptr.LC, key)

 ii) else

 return 0

 iii) If ptr.RC ≠ NULL

 search (ptr.RC, key)

 iv) else

 return 0

2. else

 Display element found at index ptr.

Insert (item)

1. If search (ptr, key)

 i) Input choice whether to insert at LC or RC

 ii) if choice = LC

 if ptr.LC = NULL

 ptr.LL = item

else

iii) else display "LC exists"

if ptr.RC = NULL

ptr.RC = item

else

display "RL exists"

2. else

display "parent not found."

Delete (Item)

1. ptr = root

2. parent = search(ptr, Item)

3. if parent ≠ NULL

i) ptr1 = parent.LC

ii) ptr2 = parent.RC

iii) if ptr.data = item

if (ptr1.LC = NULL & ptr1.RC = NULL)

parent.LC = NULL

else

display "Node is not leaf node"

else if (ptr2.LC = NULL & ptr2.RC = NULL)

parent.RC = NULL

else

display "Node not leaf node"

~~RESULT~~

Program run successfully and operations
done on binary tree.

EXPERIMENT - 17

HEAP SORT

AIM

Implementation of heap sort.

Input :- Array (binary tree) to be sorted.

Output :- sorted tree (array representation)

ALGORITHM

1. Start

2. Input the array representation of a binary tree

3. Call Heap sort (a[], n)

4. Stop.

Heap sort (a[], n)

1. for (i = n/2 - 1) to 0
~~heapsort (a, n, i)~~

~~heapsort (a, n, i)~~

2. for (i = n - 1) to 0

~~2.1. swap (a[0], a[i])~~

~~2.2. heapsort (a, i, 0)~~

3. ~~Display~~ the sorted tree.

4. Stop

heapsort (a[], n, i)

1. largest = i

2. left = $2^*i + 1$

$$3) sc = 2^*i + 2$$

4) if (lc < n & a[lc] > a[largest])
 largest = lc

5) if (rc < n & a[rc] > a[largest])
 largest = rc

6) if (largest != i)

 6.1) swap (a[i] & a[largest])

 6.2) heapify (a, n, largest)

7) return.

RESULT

Program ran successfully & sorted tree
obtained.

J.P.
1.

EXPERIMENT-18

MERGE SORT

AIM

Implementation of Merge Sort

Input:- Array representation of a tree

Output:- Sorted array.

ALGORITHM

Merge sort (a[], beg, end)

1. if (beg < end)

(i) mid = $\lfloor (\text{beg} + \text{end}) / 2 \rfloor$

(ii) Merge sort (a, beg, mid)

(iii) Merge sort (a, mid+1, end)

iv) merge (a, beg, mid, end)

merge (a[], beg, mid, end)

1. i = beg, j = mid + 1, k = beg

2. while (i <= mid and j <= end)

(i) if (a[i] <= a[j])

b[k] = a[i]

k++, i++

(ii) else

b[k] = a[j]

k++, j++

3) while ($i \leq mid$)

(i) $b[k] = a[i]$

(ii) $k++, i++$

4) while ($j \leq end$)

(i) $b[k] = a[j]$

(ii) $k++, j++$

5) for ($i = beg$ to end)

(i) $a[i] = b[i]$

(ii) $i++$

RESULT

Program ran successfully and sorted array obtained.

✓
SAHIL

EXPERIMENT-19

Quick SORT

AIM

Implementation of quick sort

Input:- Array representation of tree

Output:- Sorted array.

ALGORITHM

Quick sort ($a[]$, lb , ub)

i) if ($lb < ub$)

(i) loc = partition (a , lb , ub)

(ii) Quick sort (a , lb , $loc-1$)

(iii) Quick sort (a , $loc+1$, ub)

Partition ($a[]$, lb , ub).

1) pivot = $a[lb]$, start = lb , end = ub

2) while ($start < end$)

(i) while ($a[start] \leq pivot$)

 start++

(ii) while ($a[end] \geq pivot$)

 end--

(iii) if ($start < end$)

 swap ($a[start]$, $a[end]$)

3) swap ($a[lb]$, $a[end]$)

4) return end