# Words Worth Finding: Solving Word Search 2



By: Donal Moloney, Zirui Huang, Berker Ustura

# Lost in a Sea of Letters: Project Overview

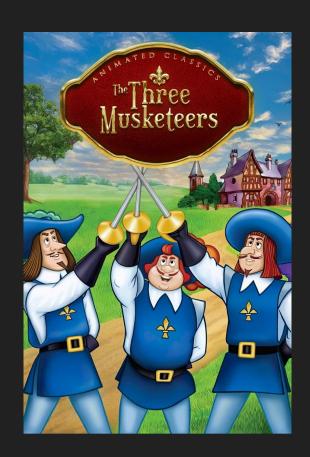


#### **Project Goal**

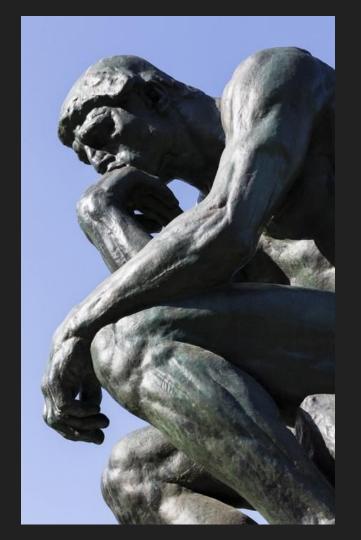
- Implement three approaches to solve the Word Search II Leetcode Hard Problem
- Compare the performance of each.

#### The three approaches:

- 1. Backtracking with Trie
- 2. DFS with Trie
- 3. DFS with Set



### Finding Needles in a Haystack: **Problem Description**



- Objective: Find all the words that can be constructed from adjacent cells.
- Constraints:
- Words are formed from adjacent cells on the board
- Each letter on the board is only used once in a single word.
- Expected Input:
  - 2D matrix of chars representing the board
  - Strings array representing list of words.
- Expected Output: Array of strings representing words found on the board.

#### Example:

Given Board:

[['c','a','t','s'],

['o','d','g','e'],

['e','a','t','i'],

['f','e','e','l']]

and Given Words:

["cat", "dog", "eat", "feel"]

Expected output: ["cat", "dog", "eat", "feel"].

# Putting the Pieces Together: Implementation Details

## **Solution 1:**Backtracking With Trie



#### Approach:

Backtracking with Trie data structure

#### Algorithm:

 Starting from each cell, explore all possible paths & use a Trie to efficiently check if the path is a valid word.

#### Time Complexity:

- O(M \* (4\*3^(L-1)))
- M (number of cells in the board) and L (maximum length of a word in the list of words).

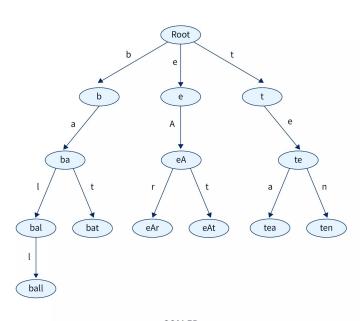
#### **Space Complexity**:

O(N), N total number of letters in the word list.

#### Key Features:

- Trie allows us to efficiently check if a path is a valid word
- Backtracking helps explore all possible paths without revisiting same cell.

## Solution 2: DFS with Trie



Topics

#### Approach:

Depth-first search (DFS) with Trie data structure

#### Algorithm:

• At each cell, explore all possible paths using DFS & use a Trie to check if path is a valid word.

#### Time Complexity:

- O(M(4⋅3^(L−1)))
- M (number of cells in the board), and L (max length of a word in the list of words).

#### **Space Complexity**:

- O(N + M)
- N is total number of letters in the word list, M is the number of cells in the board.

#### **Key Features**:

 Approach can simplify the code compared to backtracking; Trie allows checking if a path is a valid word.

# Solution 3: DFS With Set

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#### Approach:

DFS with SET data structure

#### Algorithm:

 At each cell, explore all possible paths using DFS with a Set to store the words list found on board.

#### **Time Complexity**:

- O(M \* 4^L)
- M (number of cells in the board) and L (max length of a word in list of words).

#### **Space Complexity**:

O(N), N total number of letters in the word list.

#### **Key Features**:

- DFS / Set can simplify the code compared to backtracking or using a Trie
- Set allows the store of list of words found on the board without requiring a Trie.

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#### **Backtracking with Trie Pros / Cons**

#### **Pros:**

- Can be faster than the DFS approaches for smaller boards and lists of words.
- Trie allows for efficient checking if path forms valid word.
- Adaptable to solve various word search problems.

#### Cons:

- May not find all possible words on the board.
- Can be slow & memory intensive for large boards & words.
- Needs careful management of the Trie to allow correct update when backtracking.

#### **Comparison with DFS with Trie:**

- DFS with Trie may be slower for smaller boards and lists of words.
- DFS with Trie may be easier to understand
- DFS with Trie less memory intensive (doesn't need to store multiple paths).

#### **Comparison with DFS with Set:**

- DFS with Set slower for smaller boards & lists of words.
- DFS with Set may be simpler to implement.
- DFS with Set may be more memory-efficient.
- DFS with Set may not find all possible words on the board.

#### **DFS with Trie Pros / Cons**

#### Pros:

- Simpler implementation than backtracking.
- Efficient checking if path forms a valid word.
- Adaptable to solve various word search problems.

#### Cons:

- May not find all possible words on the board.
- Can be slow & use a lot of memory for larger boards / lists of words.
- During DFS traversal careful management needed to ensure Trie updates correctly.

#### **Comparison with Backtracking with Trie:**

- DFS with Trie may be easier to understand and implement.
- Backtracking with Trie may be faster for smaller boards and lists of words.
- DFS with Trie may require less memory.

#### **Comparison with DFS with Set:**

- DFS with Trie may be faster for larger boards and lists of words.
- DFS with Trie may have better prefix search capabilities.
- DFS with Trie may be harder implement.
- DFS with Set may not find all possible words on board.

#### **DFS & Set Pros / Cons**

#### **Pros:**

- Simple & easy to implement compared to other.
- Uses less memory than the Trie approach.
- Adaptable to solve various word search problems.

#### Cons:

- May not find all possible words on board.
- Can be slow for larger boards / word lists.
- Set doesn't allow efficient prefix search.
- Unsuitable for frequently updated word lists.

#### **Comparison with Backtracking with Trie:**

- DFS with Set may be simpler to implement.
- DFS with Set may be slower for smaller boards & lists of words.
- DFS with Set may be more memory-efficient.

#### **Comparison with DFS with Trie:**

- DFS with Set may be simpler to implement.
- DFS with Trie may be faster for larger boards and lists of words.
- DFS with Trie may have better prefix search capabilities.

#### **Real Time of Each Algorithm**

We ran two separate tests for each algorithm, using a short and long input. The times are listed in the following table (runtime measurements is in ns):

|                      | Short  | Long         |
|----------------------|--------|--------------|
| Backtracking w/ Trie | 35500  | 42600        |
| DFS w/ Trie          | 15700  | 32800        |
| DFS w/ Set           | 208500 | 292569774800 |

#### **Final Considerations**

- Each solution has strengths & weaknesses.
- <u>Backtracking with Trie (BTT)</u>: good for finding all possible words on the board & can be highly efficient for small to medium-sized boards / lists of words.
- DFS with Trie (DFST): is simpler to implement than BTT & can be very efficient for small to medium-sized boards / lists of words.
- <u>DFS with Set (DFSS)</u>: easiest to implement; good for finding a subset of words on the board, may not explore all possible paths & slow for larger boards or words list.
- The best option depends on: board size, word size, freq. of updates to words etc.
- <u>Understanding</u> the pros & cons helps choose the correct solution for the job.