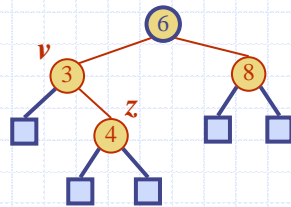
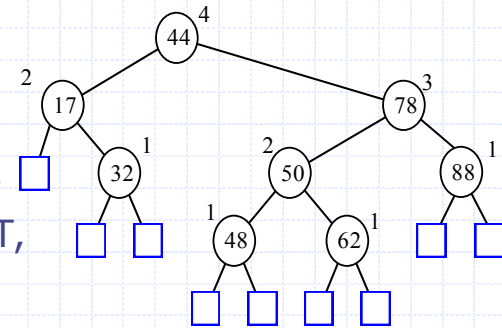


# AVL Trees



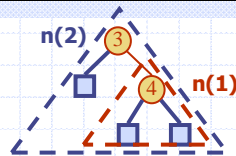
# AVL Tree Definition

- ◆ AVL trees are balanced
- ◆ An AVL Tree is a **binary search tree** such that for every internal node  $v$  of  $T$ , the **heights of the children of  $v$**  can differ by at most 1



An example of an AVL tree where the heights are shown next to the nodes:

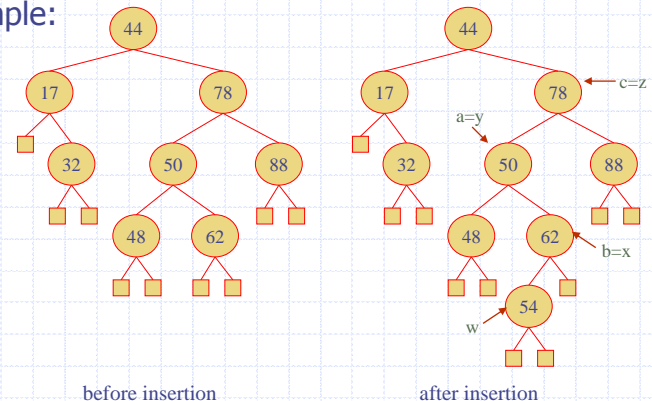
# Height of an AVL Tree



- ◆ **Fact:** The **height** of an AVL tree storing  $n$  keys is  $O(\log n)$ .
- ◆ **Proof:** Let us bound  $n(h)$ : the minimum number of internal nodes of an AVL tree of height  $h$ .
- ◆ We easily see that  $n(1) = 1$  and  $n(2) = 2$
- ◆ For  $n > 2$ , an AVL tree of height  $h$  contains the root node, one AVL subtree of height  $n-1$  and another of height  $n-2$ .
- ◆ That is,  $n(h) = 1 + n(h-1) + n(h-2)$
- ◆ Knowing  $n(h-1) > n(h-2)$ , we get  $n(h) > 2n(h-2)$ . So  
 $n(h) > 2n(h-2)$ ,  $n(h) > 4n(h-4)$ ,  $n(h) > 8n(h-6)$ , ... (by induction),  
 $n(h) > 2^{h/2}n(h-2i)$
- ◆ Solving the base case we get:  $n(h) > 2^{h/2-1}$
- ◆ Taking logarithms:  $h < 2\log n(h) + 2$
- ◆ Thus the height of an AVL tree is  $O(\log n)$

# Insertion

- ◆ Insertion is as in a binary search tree
- ◆ Always done by expanding an external node.
- ◆ Example:



before insertion

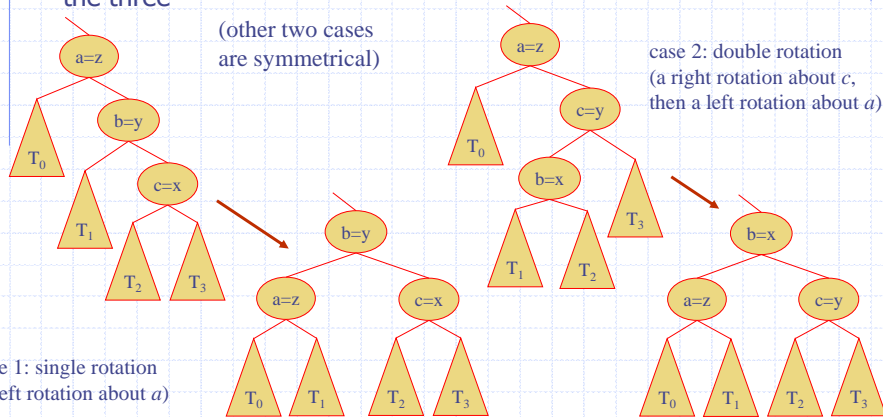
after insertion

# Trinode Restructuring

- ◆ let  $(a, b, c)$  be an inorder listing of  $x, y, z$
- ◆ perform the rotations needed to make  $b$  the topmost node of the three

(other two cases are symmetrical)

case 2: double rotation  
(a right rotation about  $c$ ,  
then a left rotation about  $a$ )



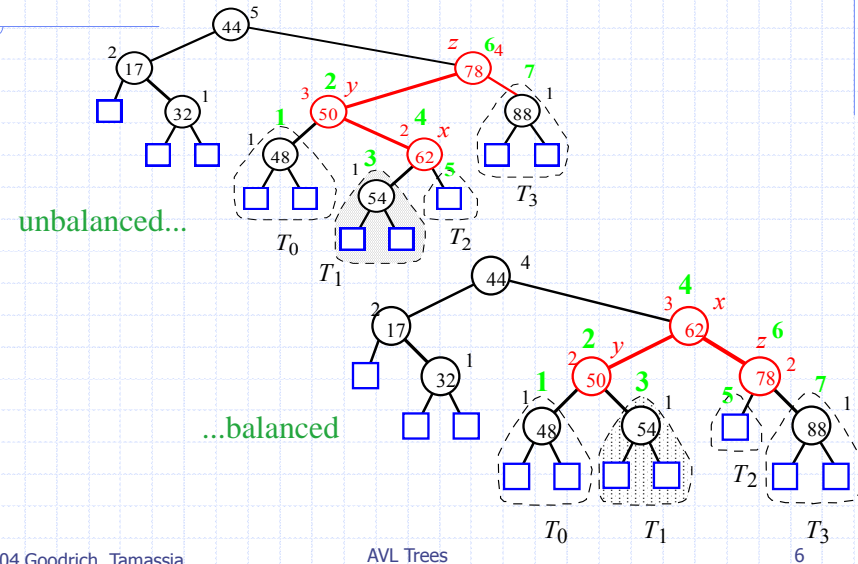
case 1: single rotation  
(a left rotation about  $a$ )

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AVL Trees

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# Insertion Example, continued



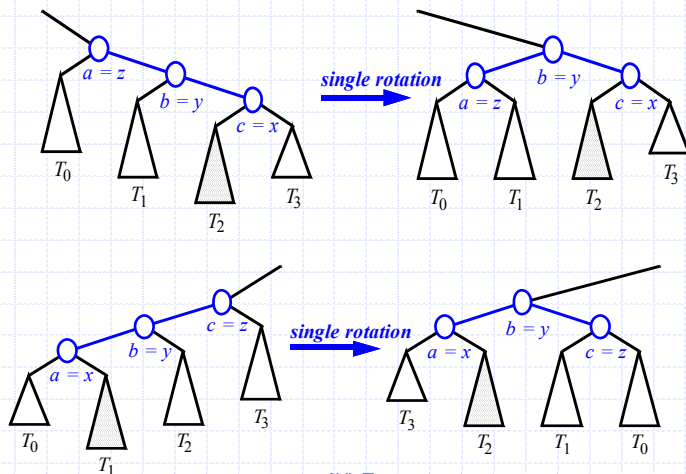
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AVL Trees

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# Restructuring (as Single Rotations)

- ◆ Single Rotations:



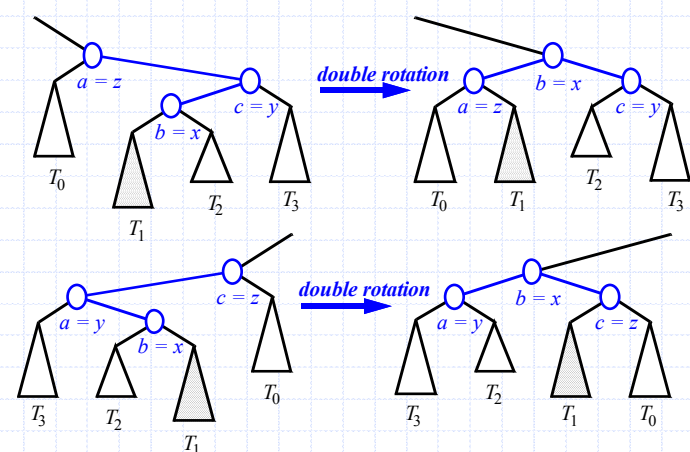
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AVL Trees

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# Restructuring (as Double Rotations)

- ◆ double rotations:



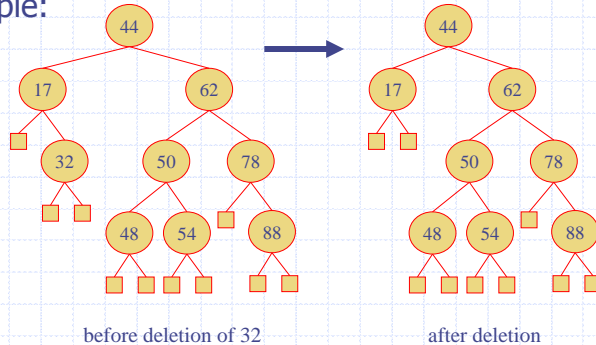
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AVL Trees

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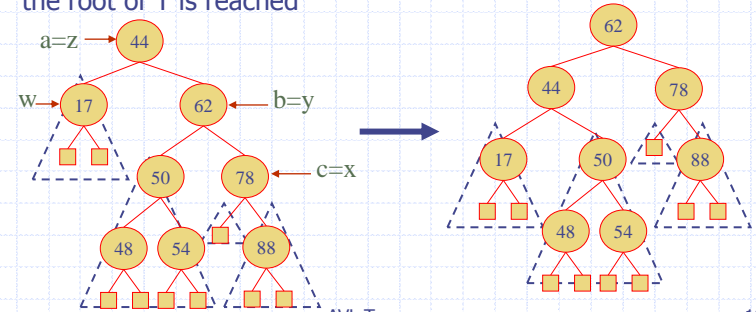
# Removal

- ◆ Removal begins as in a binary search tree, which means the node removed will become an empty external node. Its parent,  $w$ , may cause an imbalance.
- ◆ Example:

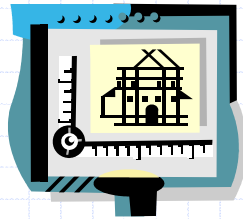


# Rebalancing after a Removal

- ◆ Let  $z$  be the **first unbalanced** node encountered while travelling up the tree from  $w$ . Also, let  $y$  be the child of  $z$  with the larger height, and let  $x$  be the child of  $y$  with the larger height
- ◆ We perform **restructure**( $x$ ) to restore balance at  $z$
- ◆ As this restructuring may upset the balance of another node higher in the tree, we must continue checking for balance until the root of  $T$  is reached



# AVL Tree Performance



- ◆ a single restructure takes  $O(1)$  time
  - using a linked-structure binary tree
- ◆ **get** takes  $O(\log n)$  time
  - height of tree is  $O(\log n)$ , no restructures needed
- ◆ **put** takes  $O(\log n)$  time
  - initial find is  $O(\log n)$
  - Restructuring up the tree, maintaining heights is  $O(\log n)$
- ◆ **remove** takes  $O(\log n)$  time
  - initial find is  $O(\log n)$
  - Restructuring up the tree, maintaining heights is  $O(\log n)$