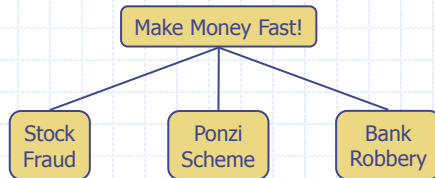
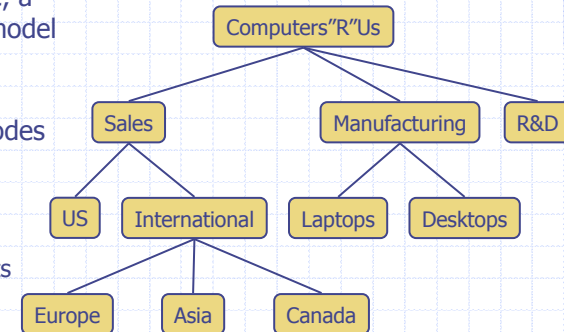


# Trees



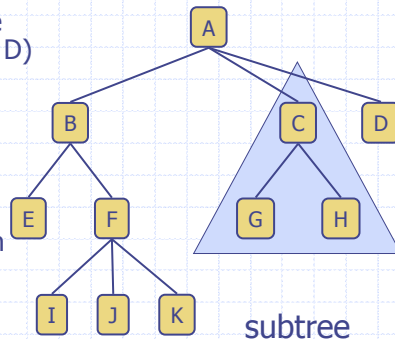
# What is a Tree

- In computer science, a tree is an abstract model of a hierarchical structure
- A tree consists of nodes with a parent-child relation
- Applications:
  - Organization charts
  - File systems
  - Programming environments



# Tree Terminology

- Root: node without parent (A)
- Internal node: node with at least one child (A, B, C, F)
- External node (a.k.a. leaf): node without children (E, I, J, K, G, H, D)
- Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- Depth of a node: number of ancestors
- Height of a tree: maximum depth of any node (3)
- Descendant of a node: child, grandchild, grand-grandchild, etc.
- Subtree: tree consisting of a node and its descendants



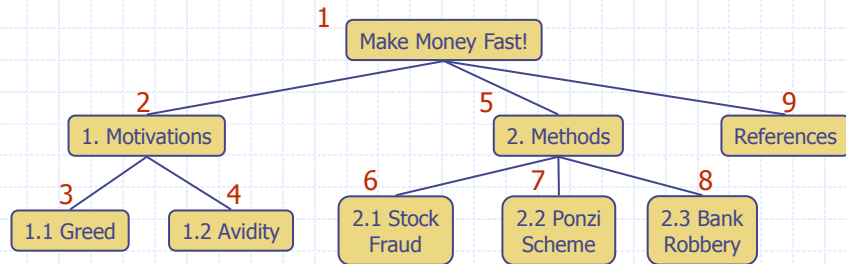
# Tree ADT

- We use positions to abstract nodes
- Generic methods:
  - integer **size**()
  - boolean **isEmpty**()
  - Iterator **iterator**()
  - Iterable **positions**()
- Accessor methods:
  - position **root**()
  - position **parent**(p)
  - Iterable **children**(p)
- ◆ Query methods:
  - boolean **isInternal**(p)
  - boolean **isExternal**(p)
  - boolean **isRoot**(p)
- ◆ Update method:
  - element **replace**(p, o)
- ◆ Additional update methods may be defined by data structures implementing the Tree ADT

# Preorder Traversal

- A traversal visits the nodes of a tree in a systematic manner
- In a preorder traversal, a node is visited before its descendants
- Application: print a structured document

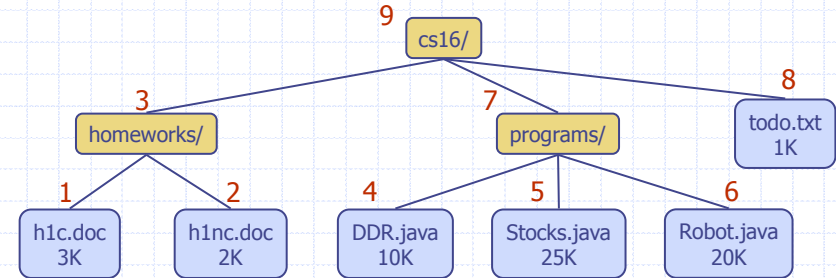
**Algorithm *preOrder*(v)**  
*visit(v)*  
**for each** child *w* of *v*  
     *preorder(w)*



# Postorder Traversal

- In a postorder traversal, a node is visited after its descendants
- Application: compute space used by files in a directory and its subdirectories

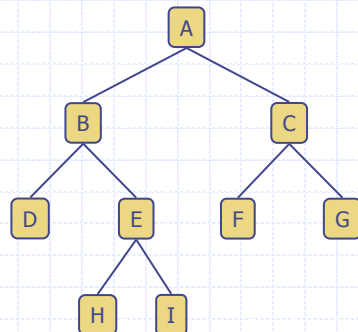
**Algorithm *postOrder*(v)**  
**for each** child *w* of *v*  
     *postOrder(w)*  
*visit(v)*



# Binary Trees

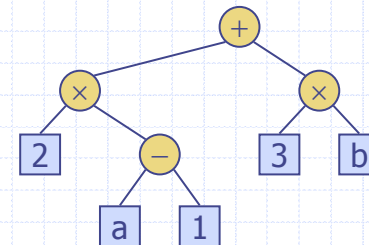
- A binary tree is a tree with the following properties:
  - Each internal node has at most two children (exactly two for **proper** binary trees)
  - The children of a node are an ordered pair
- We call the children of an internal node **left child** and **right child**
- Alternative recursive definition: a binary tree is either
  - a tree consisting of a single node, or
  - a tree whose root has an ordered pair of children, each of which is a binary tree

- Applications:
  - arithmetic expressions
  - decision processes
  - searching



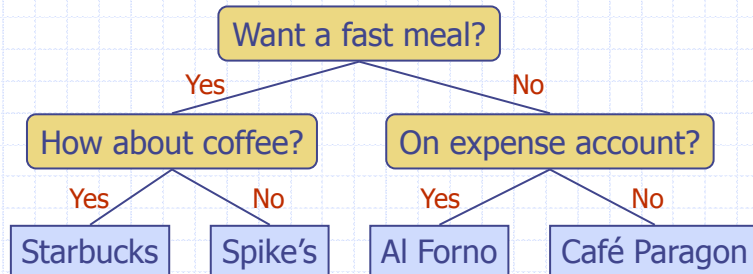
# Arithmetic Expression Tree

- Binary tree associated with an arithmetic expression
  - internal nodes: operators
  - external nodes: operands
- Example: arithmetic expression tree for the expression  $(2 \times (a - 1) + (3 \times b))$



# Decision Tree

- Binary tree associated with a decision process
  - internal nodes: questions with yes/no answer
  - external nodes: decisions
- Example: dining decision



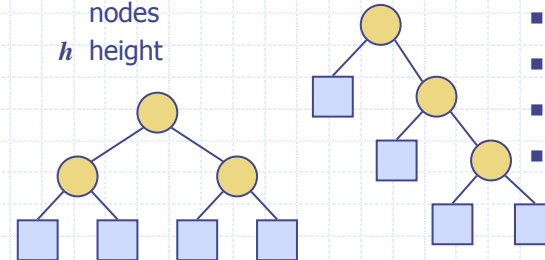
# Properties of Proper Binary Trees

## Notation

- $n$  number of nodes
- $e$  number of external nodes
- $i$  number of internal nodes
- $h$  height

## Properties:

- $e = i + 1$
- $n = 2e - 1$
- $h \leq i$
- $h \leq (n - 1)/2$
- $e \leq 2^h$
- $h \geq \log_2 e$
- $h \geq \log_2 (n + 1) - 1$



# BinaryTree ADT

- The BinaryTree ADT extends the Tree ADT, i.e., it inherits all the methods of the Tree ADT
- Update methods may be defined by data structures implementing the BinaryTree ADT
- Additional methods:
  - position **left**(p)
  - position **right**(p)
  - boolean **hasLeft**(p)
  - boolean **hasRight**(p)

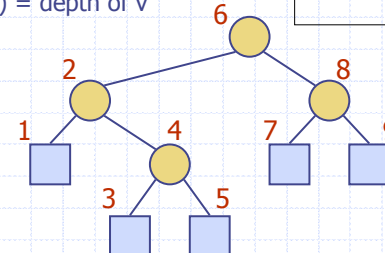
# Inorder Traversal

- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Application: draw a binary tree
  - $x(v)$  = inorder rank of  $v$
  - $y(v)$  = depth of  $v$

## Algorithm **inOrder**( $v$ )

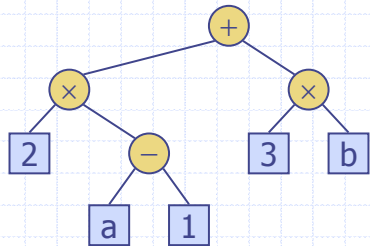
```

if hasLeft( $v$ )
    inOrder(left( $v$ ))
visit( $v$ )
if hasRight( $v$ )
    inOrder(right( $v$ ))
    
```



# Print Arithmetic Expressions

- Specialization of an inorder traversal
  - print operand or operator when visiting node
  - print "(" before traversing left subtree
  - print ")" after traversing right subtree



## Algorithm *printExpression(v)*

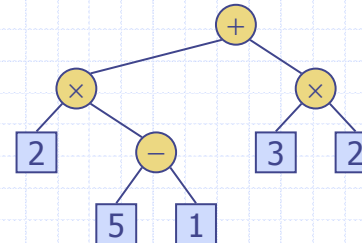
```

if hasLeft (v)
    print("(")
    inOrder (left(v))
    print(v.element ())
if hasRight (v)
    inOrder (right(v))
    print(")")
    
```

$((2 \times (a - 1)) + (3 \times b))$

# Evaluate Arithmetic Expressions

- Specialization of a postorder traversal
  - recursive method returning the value of a subtree
  - when visiting an internal node, combine the values of the subtrees



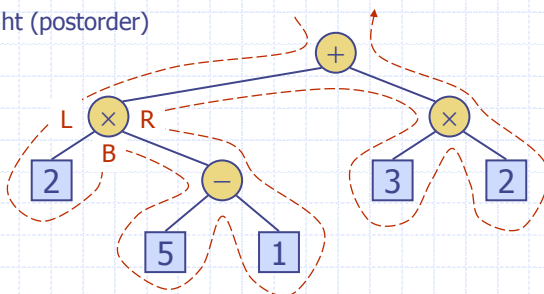
## Algorithm *evalExpr(v)*

```

if isExternal (v)
    return v.element ()
else
    x ← evalExpr(leftChild (v))
    y ← evalExpr(rightChild (v))
    ◇ ← operator stored at v
    return x ◇ y
    
```

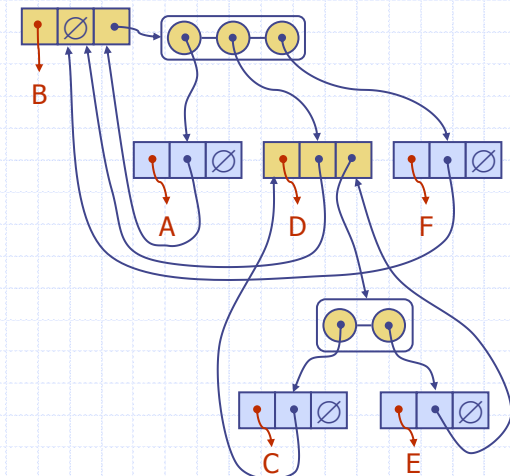
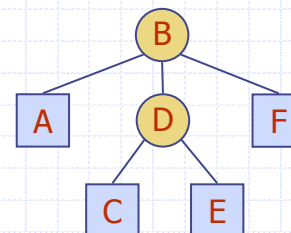
# Euler Tour Traversal

- Generic traversal of a binary tree
- Includes a special cases the preorder, postorder and inorder traversals
- Walk around the tree and visit each node three times:
  - on the left (preorder)
  - from below (inorder)
  - on the right (postorder)



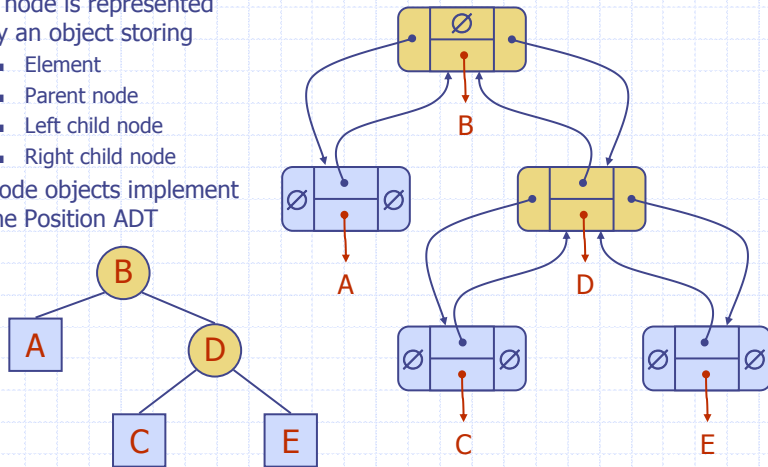
# Linked Structure for Trees

- A node is represented by an object storing
  - Element
  - Parent node
  - Sequence of children nodes
- Node objects implement the Position ADT



## Linked Structure for Binary Trees

- A node is represented by an object storing
  - Element
  - Parent node
  - Left child node
  - Right child node
- Node objects implement the Position ADT

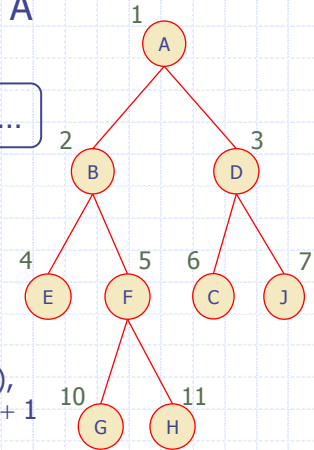


## Array-Based Representation of Binary Trees

- Nodes are stored in an array A



- Node  $v$  is stored at  $A[\text{rank}(v)]$ 
  - $\text{rank}(\text{root}) = 1$
  - if node is the left child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node}))$
  - if node is the right child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node})) + 1$



## Template Method Pattern

- Generic algorithm
- Implemented by abstract Java class
- Visit methods redefined by subclasses
- Template method `eulerTour`
  - Recursively called on left and right children
  - A `TourResult` object with fields `left`, `right` and `out` keeps track of the output of the recursive calls to `eulerTour`

```
public abstract class EulerTour<E, R> {
    protected BinaryTree<E> tree;
    public abstract R execute(BinaryTree<E> T);
    protected void init(BinaryTree<E> T) { tree = T; }
    protected R eulerTour(Position<E> v) {
        TourResult<R> r = new TourResult<R>();
        visitLeft(v, r);
        if (tree.hasLeft(p))
            { r.left=eulerTour(tree.left(v)); }
        visitBelow(v, r);
        if (tree.hasRight(p))
            { r.right=eulerTour(tree.right(v)); }
        return r.out;
    }
    protected void visitLeft(Position<E> v, TourResult<R> r) {}
    protected void visitBelow(Position<E> v, TourResult<R> r) {}
    protected void visitRight(Position<E> v, TourResult<R> r) {}
}
```

## Specializations of EulerTour

- Specialization of class `EulerTour` to evaluate arithmetic expressions
- Assumptions
  - Nodes store `ExpressionTerm` objects with method `getValue`
  - `ExpressionVariable` objects at external nodes
  - `ExpressionOperator` objects at internal nodes with method `setOperands(Integer, Integer)`

```
public class EvaluateExpressionTour
    extends EulerTour<ExpressionTerm, Integer> {
    public Integer execute
        (BinaryTree<ExpressionTerm> T) {
        init(T);
        return eulerTour(tree.root());
    }
    protected void visitRight
        (Position<ExpressionTerm> v,
         TourResult<Integer> r) {
        ExpressionTerm term = v.element();
        if (tree.isInternal(v)) {
            ExpressionOperator op = (ExpressionOperator) term;
            op.setOperands(r.left, r.right);
            r.out = term.getValue();
        }
    }
}
```