

# Welcome!



**Pavel Naumov**

p.naumov@soton.ac.uk

Agents, Interaction, Complexity

Electronics and Computer Science

University of Southampton

Russia – (1994) – USA – (2021) – UK

PhD, Cornell University

# **COMP6207**

## **Algorithmic Game Theory**

### **Lecture 12      One-sided Matching With Initial Ownership**

Pavel Naumov  
p.naumov@soton.ac.uk  
Electronics and Computer Science  
University of Southampton

# Learning Outcomes

- By the end of this session, the students should be able to
  - *Describe* the one-sided matching problem
  - **Understand** the Top Trading Cycle mechanism for the *House Allocation* problem (*HA*)
  - **Understand** Pareto-optimality
  - *Compute* the outcome of TTC and YRMH-IGYT

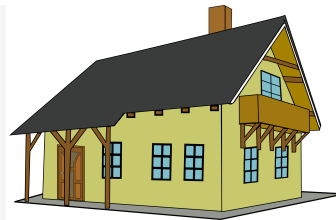
# One-sided Matching

- Bipartite structure
  - Individuals in one group has preferences over the individuals in the other group.
- Examples
  - *House allocation*
    - Each individual owns a house and is interested in improving its utility by swapping with others
  - Other allocation problems
    - No initial ownership

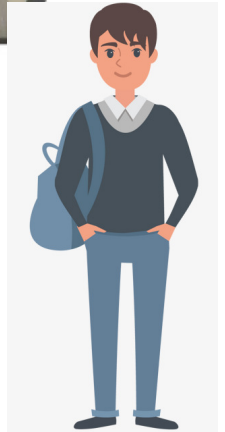
# House allocation



Can we benefit from swapping houses?



# Student rooms allocation



## **Random Serial Dictatorship with Squatting Rights.**

Used in undergrad housing in many universities.

1. Each existing tenant decides whether they want to participate in the housing lottery or keep the current house. Those who decide to keep their houses are assigned the current houses. All other houses become available for assignment in later steps.
2. An ordering of agents is decided. The ordering may be uniformly random or may favor some subgroup of agents (for example, seniors over juniors).
3. Serial dictatorship is applied to all available houses and agents (except for existing tenants already assigned their current houses).

What do you think of that mechanism?

**Problem:**

Existing tenants are not guaranteed to get at least as good a house as their current house: Individually irrational!

Some existing tenants may not want to enter the lottery even if they want to move.

This may result in loss of gains from trade, and the resulting matching may not be Pareto efficient.



# Desired properties

- Some good properties we want for the *student rooms allocation* mechanisms and *house allocation* mechanisms.
  - Individual rationality
  - Pareto efficiency
  - Strategy-proofness
- Can we find mechanisms with the above good properties?

# House Allocation Problem

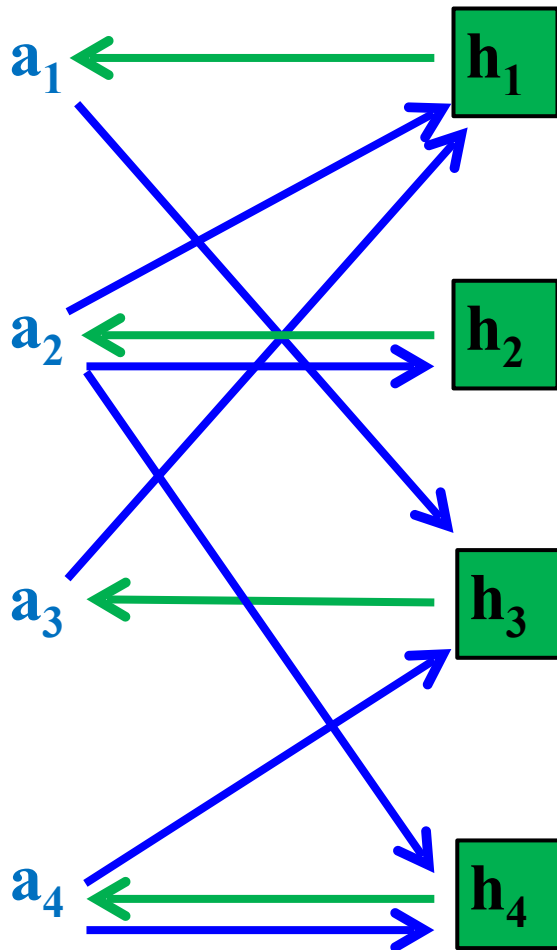
- A model for allocation of indivisible goods.
- A set of  $n$  agents, each **owns** a unique house and a strict preference ordering over all  $n$  houses.
- The objective is to reallocate the houses among the agents to improve the agents' utility.
- Top Trading Cycle (TTC) mechanism:
  - Introduced by David Gale and works in rounds.
  - Initially  $\#|\text{agents}| = \#|\text{houses}|$ , but let's check out this generalized version.

# Top Trading Cycle mechanism (**TTC**)

## In each round:

1. Each **agent points** to her **most preferred house (possibly her own house)**; each **house points back** to its **owner**
2. This creates a directed graph; in this graph, identify cycles
  - Finite: cycle must exist
  - Strict preferences: each agent is in at most one cycle
3. Assign each agent in a cycle to the house she is pointing at and remove her from the mechanism with her assigned house.
4. If all houses are assigned or all agents are assigned or all preference lists are empty, stop.
5. Otherwise, Repeat (i.e. go to step 1)

# Example: TTC



$$a_1 : \mathbf{h}_3 > h_2 > \mathbf{h}_1$$

$$a_2 : h_1 > h_4 > \mathbf{h}_2$$

$$a_3 : \mathbf{h}_1 > h_4 > \mathbf{h}_3$$

$$a_4 : h_3 > \mathbf{h}_4$$

- $a_1$  is matched to  $\mathbf{h}_3$
- $a_3$  is matched to  $\mathbf{h}_1$
- $a_4$  is matched to  $\mathbf{h}_4$
- $a_2$  is matched to  $\mathbf{h}_2$

# Example: TTC

$$\mathbf{a}_1 : \mathbf{h}_3 > h_2 > \mathbf{h}_1$$

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$$\mathbf{a}_4 : h_3 > \mathbf{h}_4$$

- $\mathbf{a}_1$  is matched to  $\mathbf{h}_3$
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- $\mathbf{a}_4$  is matched to  $\mathbf{h}_4$
- $\mathbf{a}_2$  is matched to  $\mathbf{h}_2$

# Quiz

- What matching does TTC return?

$$\mathbf{a}_1 : h_2 > h_3 > \mathbf{h}_1$$

$$\mathbf{a}_2 : h_1 > \mathbf{h}_2$$

$$\mathbf{a}_3 : h_1 > h_4 > \mathbf{h}_3$$

$$\mathbf{a}_4 : h_2 > h_1 > h_3 > \mathbf{h}_4$$

# Desired Properties

- Reflection: what are the desired properties in this scenario?
  - Truthfulness
  - Pareto optimality
  - What is the matter with the agents' ownership?
    - Individual rationality
    - Core

# Incentive in TTC

## Theorem.

TTC is dominant-strategy truthful.

## Proof sketch.

For any agent matched in round  $k$  if truthful

- No change in her report can give her a house that was assigned in earlier rounds
  - (this step of the proof needs some more argument)
- No house assigned in a later round will make her better off
- So no benefit in reporting any other preferences but the truth.



# Pareto Optimality

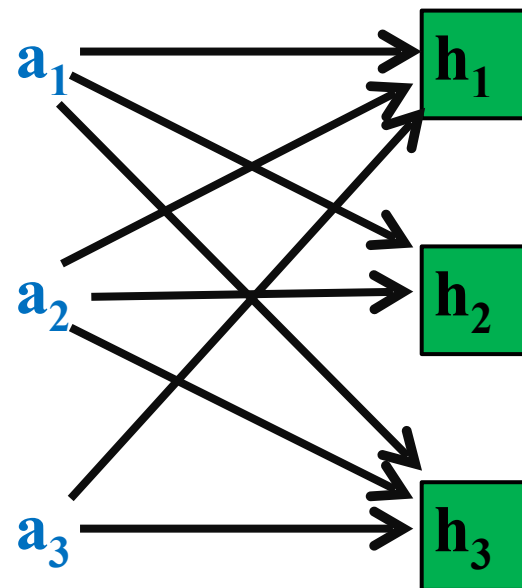
A matching  $\mathbf{M}_1$  is *Pareto optimal (PO)* if there is no other matching  $\mathbf{M}_2$  such that

- some agent prefers  $\mathbf{M}_2$  to  $\mathbf{M}_1$
- no agent prefers  $\mathbf{M}_1$  to  $\mathbf{M}_2$

$\mathbf{a}_1 : h_1 > h_2 > h_3$

$\mathbf{a}_2 : h_2 > h_1 > h_3$

$\mathbf{a}_3 : h_1 > h_3$



Pareto optimality is also known as the **Pareto efficiency**

# Pareto Optimality

A matching  $\mathbf{M}_1$  is *Pareto optimal (PO)* if there is no other matching  $\mathbf{M}_2$  such that

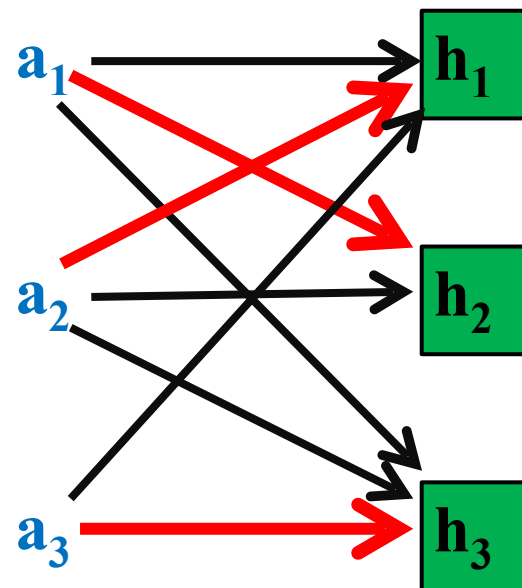
- some agent prefers  $\mathbf{M}_2$  to  $\mathbf{M}_1$
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$$\mathbf{a}_1 : h_1 > \mathbf{h}_2 > h_3$$

$$\mathbf{a}_2 : h_2 > \mathbf{h}_1 > h_3$$

$$\mathbf{a}_3 : h_1 > \mathbf{h}_3$$

$$\mathbf{M}_1 = \{(\mathbf{a}_1, \mathbf{h}_2), (\mathbf{a}_2, \mathbf{h}_1), (\mathbf{a}_3, \mathbf{h}_3)\}$$



# Pareto Optimality

A matching  $\mathbf{M}_1$  is *Pareto optimal (PO)* if there is no other matching  $\mathbf{M}_2$  such that

- some agent prefers  $\mathbf{M}_2$  to  $\mathbf{M}_1$
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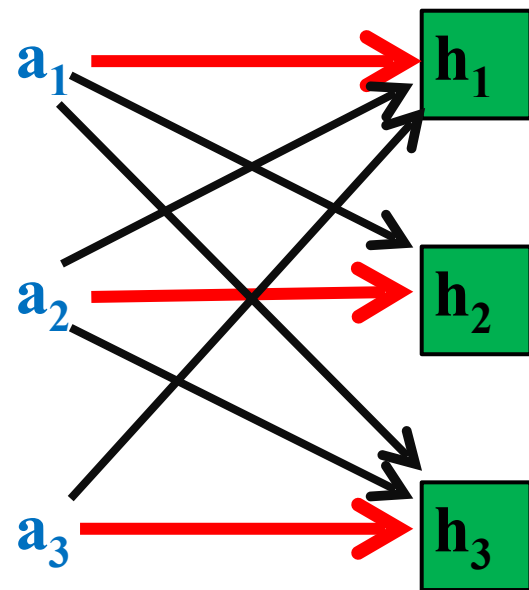
$$\mathbf{a}_1 : \mathbf{h}_1 > h_2 > h_3$$

$$\mathbf{a}_2 : \mathbf{h}_2 > h_1 > h_3$$

$$\mathbf{a}_3 : h_1 > \mathbf{h}_3$$

$$\mathbf{M}_1 = \{(\mathbf{a}_1, h_2), (\mathbf{a}_2, h_1), (\mathbf{a}_3, h_3)\}$$

$$\mathbf{M}_2 = \{(\mathbf{a}_1, \mathbf{h}_1), (\mathbf{a}_2, \mathbf{h}_2), (\mathbf{a}_3, \mathbf{h}_3)\}$$



$\mathbf{M}_2$  Pareto dominates  $\mathbf{M}_1$

# Quiz

Is the **red** matching Pareto optimal?

$$\mathbf{a}_1 : h_2 > \mathbf{h}_1$$

$$\mathbf{a}_2 : h_3 > \mathbf{h}_4 > h_2$$

$$\mathbf{a}_3 : h_4 > \mathbf{h}_3$$

$$\mathbf{a}_4 : h_3 > h_4$$

# Quiz

Is the **red** matching Pareto optimal?

$$\mathbf{a}_1 : \mathbf{h}_2 > h_1$$

$$\mathbf{a}_2 : h_3 > \mathbf{h}_4 > h_2$$

$$\mathbf{a}_3 : h_4 > \mathbf{h}_3$$

$$\mathbf{a}_4 : h_3 > h_4$$

# Quiz

Is the **red** matching Pareto optimal?

$$\mathbf{a}_1 : \mathbf{h}_2 > h_1$$

$$\mathbf{a}_2 : \mathbf{h}_3 > h_4 > h_2$$

$$\mathbf{a}_3 : h_4 > h_3$$

$$\mathbf{a}_4 : h_3 > \mathbf{h}_4$$

# Pareto optimality in TTC

## Theorem.

TTC allocation is always **Pareto optimal**.

## Proof sketch.

By induction:

- None of the agents matched in round 1 can improve their assignment (they are already getting their most favourite house)
- Fixing the assignment of those agents matched in round 1, none of the agents in round 2 can improve their assignment (they are already getting their most favourite among the remaining houses)
- And so on.

# The Core

- The agents own the houses  $\Rightarrow$  if a subset of agents can make all of its members better off by exchanging the houses within them, then the TTC allocation would be unstable as these agents would have an incentive to reallocate their houses.
- The allocation returned by TTC is such that no such subset of agents exists.
- Alternatively speaking, the TTC outcome is a core allocation.
- In addition, the TTC allocation is the only such assignment  $\Rightarrow$  it returns the unique core allocation.



# Student rooms allocation

- A set of existing students  $\{a_1, a_2, \dots, a_n\}$ , each occupies a room
- A set of newcomers  $\{a_{n+1}, a_{n+2}, \dots, a_{n+m}\}$
- A set of vacant rooms  $\{1, 2, \dots, m\}$
- A mechanism:
  - You request my house, I get your turn (YRMH-IGYT)

# You request my house, I get your turn (YRMH-IGYT)

- Fix a priority order of the agents
- Let the agent with the top priority receive her first-choice room, the second agent receive her top choice among the remaining goods and so on, until someone requests the room of an existing tenant.
- If the existing tenant whose room is requested has already received a room, then proceed the assignment to the next agent. *Otherwise, insert the existing tenant at the top of the priority order and proceed with the procedure.*
- If at any step a cycle forms, the cycle is formed by existing tenants  $(a_1, \dots, a_k)$  where  $a_1$  points to the house of agent  $a_2$ , who points to the house of  $a_3$ , and so on. In such a case assign these houses by letting them exchange, and then proceed with the algorithm.

# An example

Let  $A_E = \{a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9\}$  be the set of existing tenants

$A_N = \{a_{10}, a_{11}, a_{12}, a_{13}, a_{14}, a_{15}, a_{16}\}$  be the set of newcomers

$H_0 = \{h_1, h_2, h_3, h_4, h_5, h_6, h_7, h_8, h_9\}$  be the set of occupied houses

$H_V = \{h_{10}, h_{11}, h_{12}, h_{13}, h_{14}, h_{15}, h_{16}\}$  be the set of vacant houses

Let the preference profile  $P$  be given as

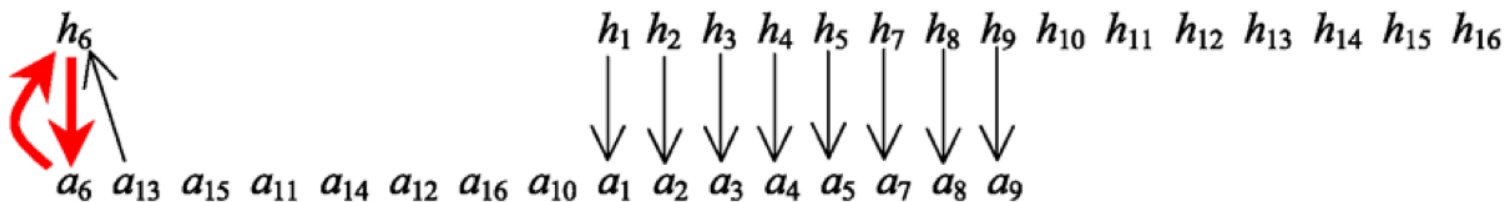
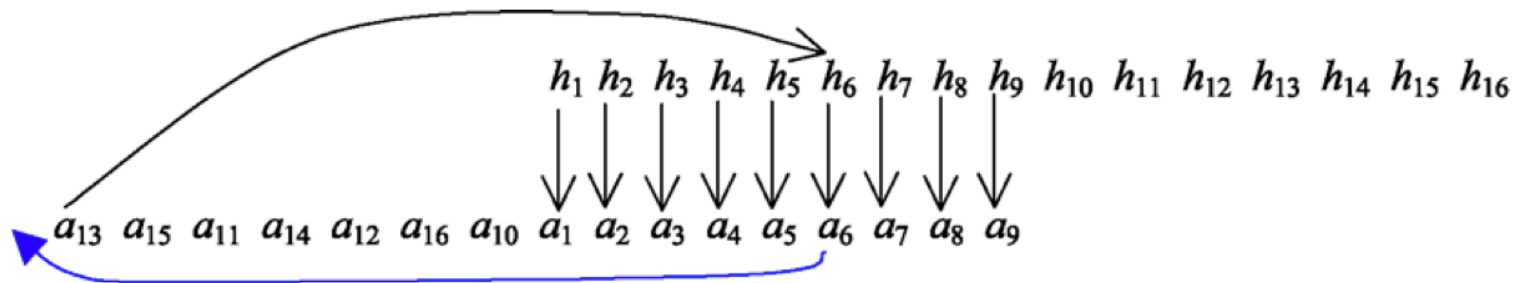
$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					

Priority order  $>$ :  $a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

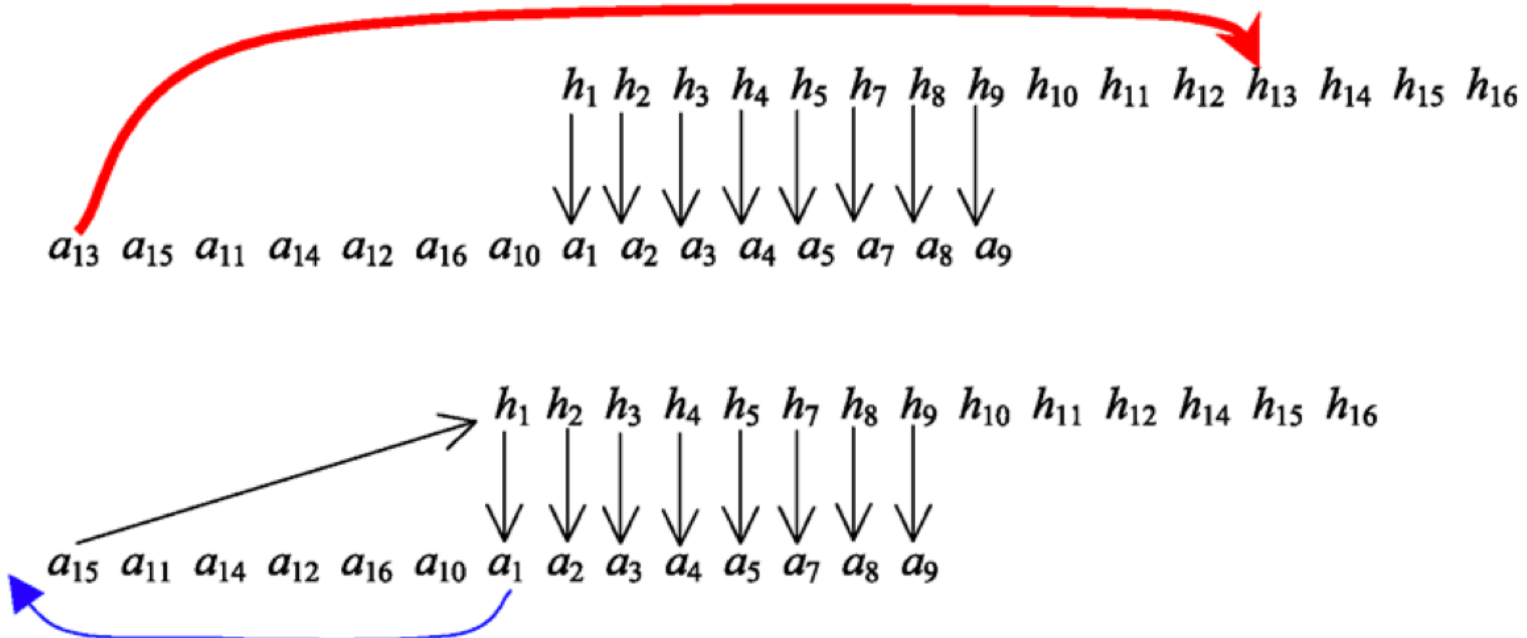
$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					



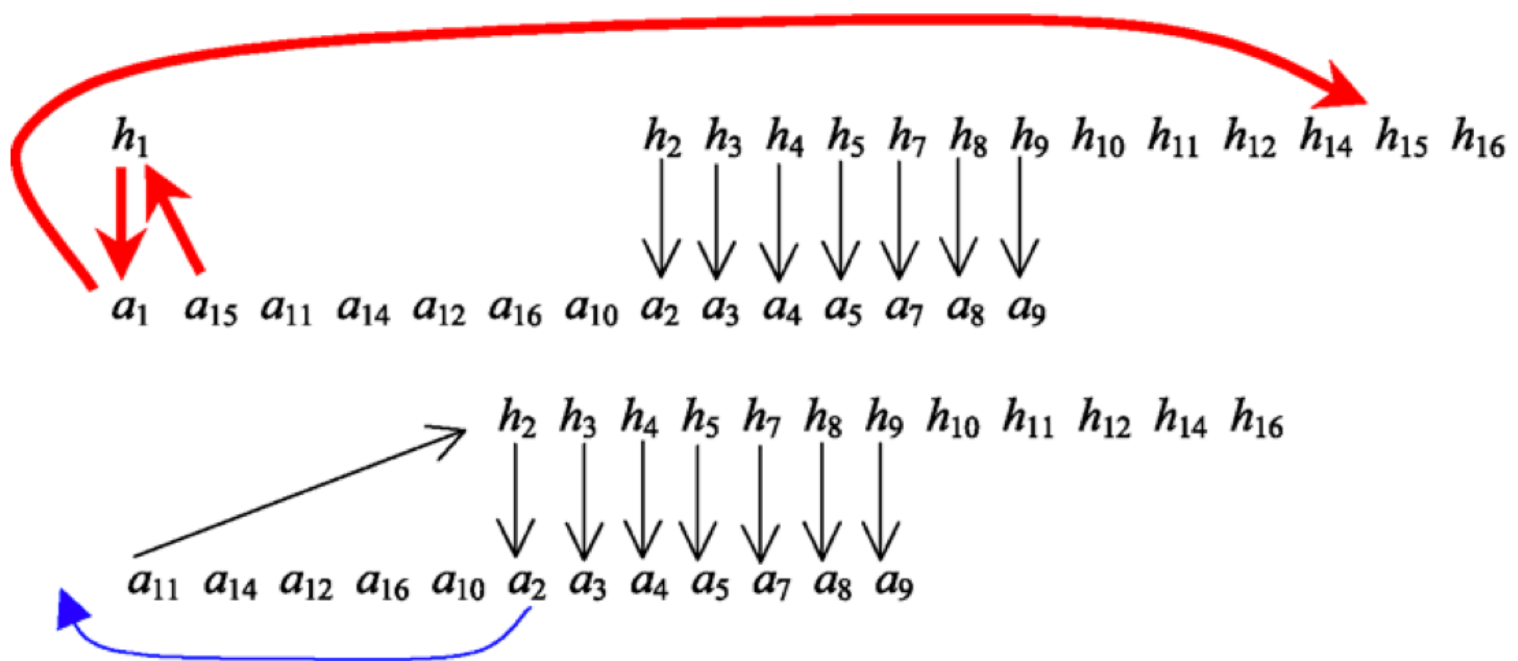
Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
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Priority order  $\succ$ :

$$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$$
$$\begin{array}{c} \overbrace{a_1 \ a_2 \ a_3 \ a_4 \ a_5 \ a_6 \ a_7 \ a_8 \ a_9}^{A_E} \\ \hline h_{15} \ h_3 \ h_1 \ h_2 \ h_9 \ h_6 \ h_6 \ h_6 \ h_{11} \\ \vdots \quad h_4 \ h_3 \quad \vdots \quad \vdots \quad \vdots \quad h_7 \ h_{12} \quad \vdots \\ \quad \quad \vdots \quad \vdots \quad \quad \quad \vdots \quad \vdots \\ \end{array} \quad \begin{array}{c} \overbrace{a_{10} \ a_{11} \ a_{12} \ a_{13} \ a_{14} \ a_{15} \ a_{16}}^{A_N} \\ \hline h_7 \ h_2 \ h_4 \ h_6 \ h_8 \ h_1 \ h_5 \\ h_3 \ h_4 \ h_{14} \ h_{13} \quad \vdots \quad \vdots \quad \vdots \\ h_{12} \ h_{16} \quad \vdots \quad \vdots \\ h_{10} \quad \vdots \end{array}$$


Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					

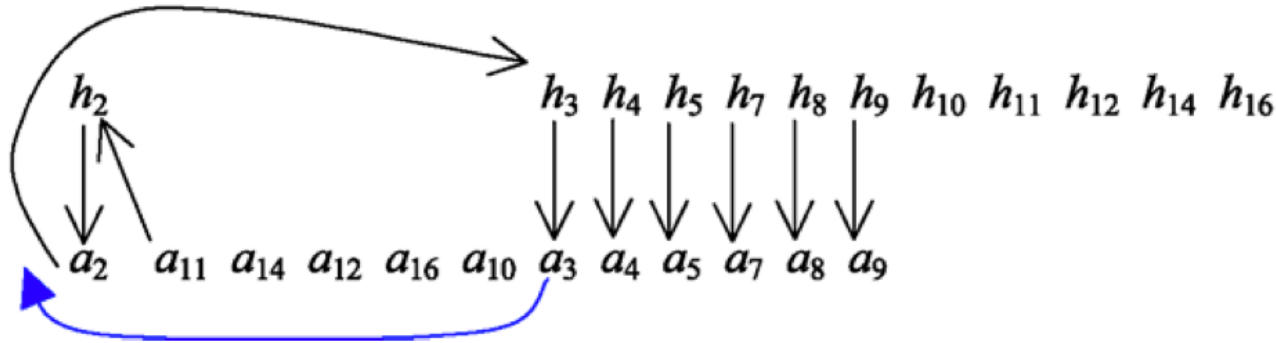
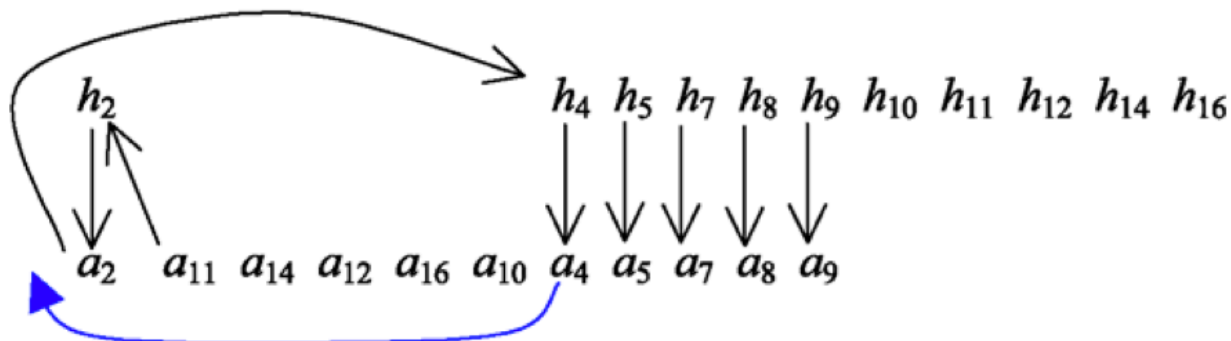
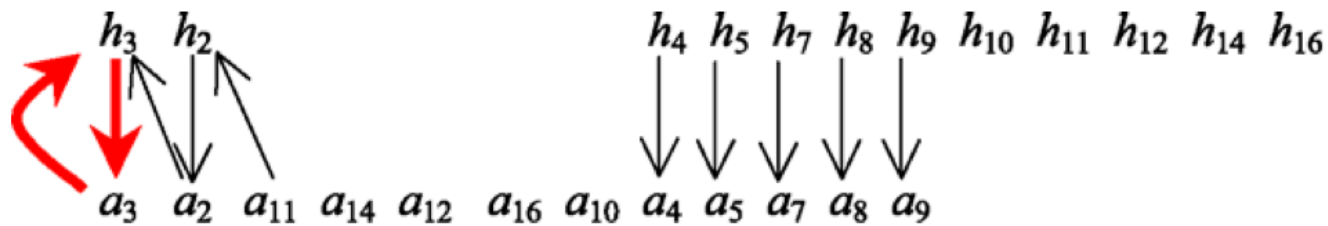


Fig. 1. The sequence of first seven events under the YRMH-IGYT algorithm.

Priority order  $\succ$ :

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$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					

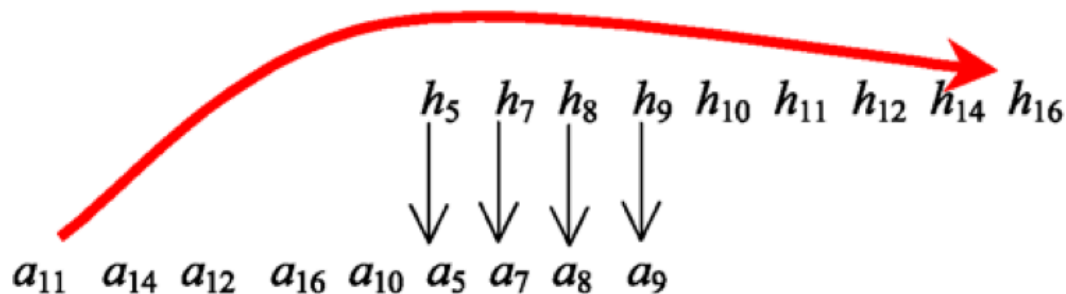
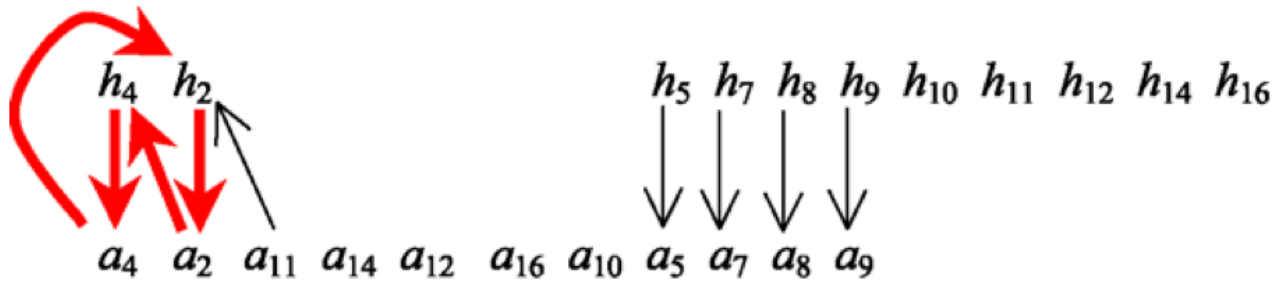




Priority order  $\succ$ :

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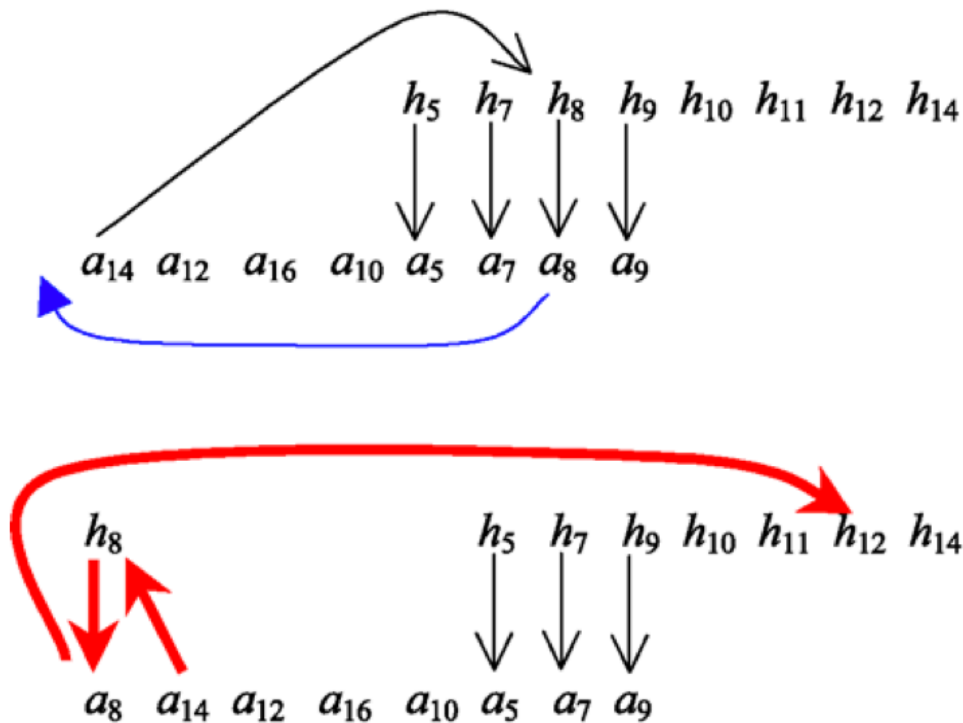
$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					



Priority order  $>$ :

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$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					



Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					

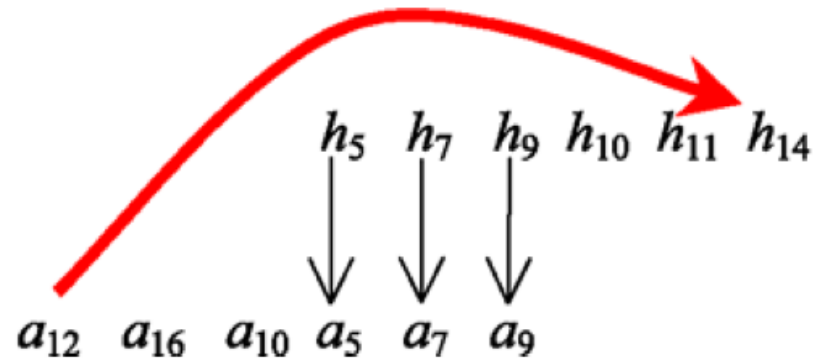
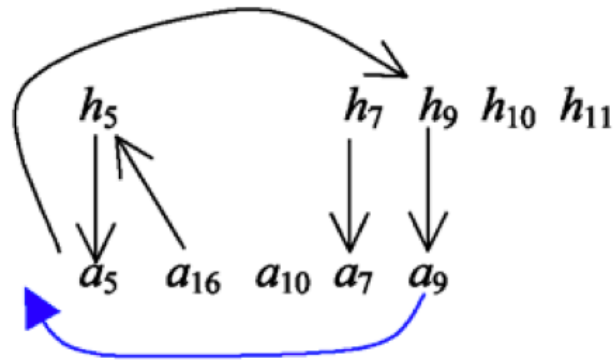
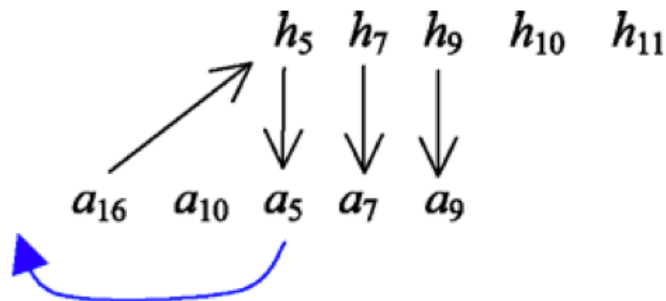


Fig. 2. The sequence of second seven events under the YRMH-IGYT algorithm.

Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

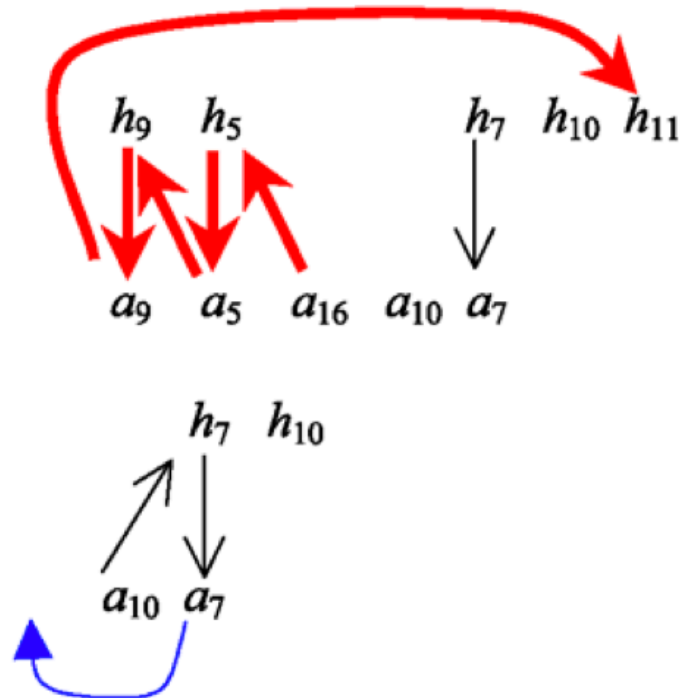
$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					



Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

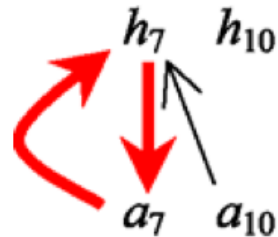
$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					



Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

$$\overbrace{\begin{array}{ccccccccc} a_1 & a_2 & a_3 & a_4 & a_5 & a_6 & a_7 & a_8 & a_9 \\ \hline h_{15} & h_3 & h_1 & h_2 & h_9 & h_6 & h_6 & h_6 & h_{11} \\ \vdots & h_4 & h_3 & \vdots & \vdots & \vdots & h_7 & h_{12} & \vdots \\ & \vdots & \vdots & & & & \vdots & \vdots & \end{array}}^{A_E}$$

$$\overbrace{\begin{array}{ccccccc} a_{10} & a_{11} & a_{12} & a_{13} & a_{14} & a_{15} & a_{16} \\ \hline h_7 & h_2 & h_4 & h_6 & h_8 & h_1 & h_5 \\ h_3 & h_4 & h_{14} & h_{13} & \vdots & \vdots & \vdots \\ h_{12} & h_{16} & \vdots & \vdots & & & \\ h_{10} & \vdots & & & & & \end{array}}^{A_N}$$


Priority order  $\succ$ :

$a_{13}, a_{15}, a_{11}, a_{14}, a_{12}, a_{16}, a_{10}, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$

$A_E$									$A_N$						
$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$a_{16}$
$h_{15}$	$h_3$	$h_1$	$h_2$	$h_9$	$h_6$	$h_6$	$h_6$	$h_{11}$	$h_7$	$h_2$	$h_4$	$h_6$	$h_8$	$h_1$	$h_5$
$\vdots$	$h_4$	$h_3$	$\vdots$	$\vdots$	$\vdots$	$h_7$	$h_{12}$	$\vdots$	$h_3$	$h_4$	$h_{14}$	$h_{13}$	$\vdots$	$\vdots$	$\vdots$
	$\vdots$	$\vdots$				$\vdots$	$\vdots$		$h_{12}$	$h_{16}$	$\vdots$	$\vdots$			
									$h_{10}$	$\vdots$					

The outcome of the algorithm is

$$\begin{bmatrix} a_1 & a_2 & a_3 & a_4 & a_5 & a_6 & a_7 & a_8 & a_9 & a_{10} & a_{11} & a_{12} & a_{13} & a_{14} & a_{15} & a_{16} \\ h_{15} & h_4 & h_3 & h_2 & h_9 & h_6 & h_7 & h_{12} & h_{11} & h_{10} & h_{16} & h_{14} & h_{13} & h_8 & h_1 & h_5 \end{bmatrix}.$$

# Properties

Theorem (Theorem 3 in [Abdulkadiroğlu and Sönmez \(1999\)](#)): For a given ordering  $f$ , the YRMH-IGYT algorithm yields the same outcome as the top trading cycles algorithm.

We can think of YRMH-IGYT as a variant of Galei's TTC in which all vacant houses (and houses whose initial owners are already assigned houses) point to the highest priority agents rather than the owners of the houses. So we sometimes call the mechanism TTC as well.

The YRMH-IGYT mechanism is Pareto efficient, strategyproof, and makes no existing tenant worse off.



# Summary

- One-sided matching problem
  - When agents own items, i.e., so-called House Allocation problem
    - Top Trading Cycle mechanism (TTC)
- Extra reading
  - House allocation
    - Chapter 10, Section 3, AGT book

# Preview

- When agents do not own items
  - Probabilistic Serial mechanism (PS)
  - Random Priority (RP)