

Isogenies of Oriented Elliptic Curves

Doron L Grossman-Naples (he/she/they)

University of Illinois, Urbana-Champaign

August 23rd, 2025



Some Preliminary Notes

Conventions/Terminology

- Stack = étale sheaf of ∞ -groupoids on $\mathcal{C}\text{Alg}_R$
- DM-stack (algebraic space) = spectral Deligne-Mumford stack (spectral algebraic space), not necessarily connective
- Formal DM-stack = formal filtered colimit of DM-stacks along closed immersions; called “honest” if actual DM-stack. Formal algebraic spaces defined similarly
- Isogeny = strict abelian variety map which is finite, flat, and locally almost of finite presentation



Some Preliminary Notes

Conventions/Terminology

- Stack = étale sheaf of ∞ -groupoids on $\mathcal{C}\text{Alg}_R$
- DM-stack (algebraic space) = spectral Deligne-Mumford stack (spectral algebraic space), not necessarily connective
- Formal DM-stack = formal filtered colimit of DM-stacks along closed immersions; called “honest” if actual DM-stack. Formal algebraic spaces defined similarly
- Isogeny = strict abelian variety map which is finite, flat, and locally almost of finite presentation

Related Work

Xuecai Ma and Yifei Zhu have a paper in the works ([MZ25]) which approaches this from a different perspective, defining level structures in terms of classical divisors. It isn't clear whether this is equivalent to my definition. A draft can be found on Professor Zhu's website.

Background: Classical \mathcal{M}_{ell}

Over \mathbb{C}

$$\left\{ \begin{array}{c} \mathcal{M}(\Gamma) \\ \downarrow \text{ét} \\ \mathcal{M}_{\text{ell}} \end{array} \right\} \longleftrightarrow \{ \text{Congruence subgroups } \Gamma \subset GL_2(\mathbb{Z}) \}$$



Background: Classical \mathcal{M}_{ell}

Over \mathbb{C}

$$\left\{ \begin{array}{c} \mathcal{M}(\Gamma) \\ \downarrow \text{ét} \\ \mathcal{M}_{\text{ell}} \end{array} \right\} \longleftrightarrow \{ \text{Congruence subgroups } \Gamma \subset GL_2(\mathbb{Z}) \}$$

$\mathcal{M}(\Gamma)$ = “moduli stack of elliptic curves with Γ -structure” (i.e. isogeny with prescribed kernel)



Background: Classical \mathcal{M}_{ell}

Over \mathbb{C}

$$\left\{ \begin{array}{c} \mathcal{M}(\Gamma) \\ \downarrow \text{ét} \\ \mathcal{M}_{\text{ell}} \end{array} \right\} \longleftrightarrow \{ \text{Congruence subgroups } \Gamma \subset GL_2(\mathbb{Z}) \}$$

$\mathcal{M}(\Gamma)$ = “moduli stack of elliptic curves with Γ -structure” (i.e. isogeny with prescribed kernel)

Over \mathbb{Z}

$\mathcal{M}(\Gamma)$ and the reduction map exist in the category of algebraic stacks, but



Background: Classical \mathcal{M}_{ell}

Over \mathbb{C}

$$\left\{ \begin{array}{c} \mathcal{M}(\Gamma) \\ \downarrow \text{ét} \\ \mathcal{M}_{\text{ell}} \end{array} \right\} \longleftrightarrow \{ \text{Congruence subgroups } \Gamma \subset GL_2(\mathbb{Z}) \}$$

$\mathcal{M}(\Gamma)$ = “moduli stack of elliptic curves with Γ -structure” (i.e. isogeny with prescribed kernel)

Over \mathbb{Z}

$\mathcal{M}(\Gamma)$ and the reduction map exist in the category of algebraic stacks, but

- $\mathcal{M}(\Gamma)$ may not be Deligne-Mumford (e.g. $\mathcal{M}_0(N)$)



Background: Classical \mathcal{M}_{ell}

Over \mathbb{C}

$$\left\{ \begin{array}{c} \mathcal{M}(\Gamma) \\ \downarrow \text{ét} \\ \mathcal{M}_{\text{ell}} \end{array} \right\} \longleftrightarrow \{ \text{Congruence subgroups } \Gamma \subset GL_2(\mathbb{Z}) \}$$

$\mathcal{M}(\Gamma)$ = “moduli stack of elliptic curves with Γ -structure” (i.e. isogeny with prescribed kernel)

Over \mathbb{Z}

$\mathcal{M}(\Gamma)$ and the reduction map exist in the category of algebraic stacks, but

- $\mathcal{M}(\Gamma)$ may not be Deligne-Mumford (e.g. $\mathcal{M}_0(N)$)
- The reduction map is **never** étale (unless we invert the level)



Background: Classical \mathcal{M}_{ell}

Over \mathbb{C}

$$\left\{ \begin{array}{c} \mathcal{M}(\Gamma) \\ \downarrow \text{ét} \\ \mathcal{M}_{\text{ell}} \end{array} \right\} \longleftrightarrow \{ \text{Congruence subgroups } \Gamma \subset GL_2(\mathbb{Z}) \}$$

$\mathcal{M}(\Gamma)$ = “moduli stack of elliptic curves with Γ -structure” (i.e. isogeny with prescribed kernel)

Over \mathbb{Z}

$\mathcal{M}(\Gamma)$ and the reduction map exist in the category of algebraic stacks, but

- $\mathcal{M}(\Gamma)$ may not be Deligne-Mumford (e.g. $\mathcal{M}_0(N)$)
- The reduction map is **never** étale (unless we invert the level)

So we can't use the usual methods to lift to spectral AG.

Background: Classical \mathcal{M}_{ell}

Over \mathbb{C}

$$\left\{ \begin{array}{c} \mathcal{M}(\Gamma) \\ \downarrow \text{ét} \\ \mathcal{M}_{\text{ell}} \end{array} \right\} \longleftrightarrow \{ \text{Congruence subgroups } \Gamma \subset GL_2(\mathbb{Z}) \}$$

$\mathcal{M}(\Gamma)$ = “moduli stack of elliptic curves with Γ -structure” (i.e. isogeny with prescribed kernel)

Over \mathbb{Z}

$\mathcal{M}(\Gamma)$ and the reduction map exist in the category of algebraic stacks, but

- $\mathcal{M}(\Gamma)$ may not be Deligne-Mumford (e.g. $\mathcal{M}_0(N)$)
- The reduction map is **never** étale (unless we invert the level)

So we can't use the usual methods to lift to spectral AG.

Solution: Work with moduli interpretation directly!

The Moduli of Isogenies

Definition



The Moduli of Isogenies

Definition

The *moduli stack of isogenies* over R is the functor $\text{Isog} : \mathbf{CAlg}_R \rightarrow \mathcal{S}$ given by



The Moduli of Isogenies

Definition

The *moduli stack of isogenies* over R is the functor $\text{Isog} : \mathbf{CAlg}_R \rightarrow \mathcal{S}$ given by

$$A \mapsto \left\{ i : E \rightarrow E' \left| \begin{array}{l} E, E' \in \mathcal{M}_{\text{ell}}^{\text{or}}(A), \\ i \text{ isogeny.} \end{array} \right. \right\} \quad (1)$$



The Moduli of Isogenies

Definition

The *moduli stack of isogenies* over R is the functor $\text{Isog} : \mathbf{CAlg}_R \rightarrow \mathcal{S}$ given by

$$A \mapsto \left\{ i : E \rightarrow E' \left| \begin{array}{l} E, E' \in \mathcal{M}_{\text{ell}}^{\text{or}}(A), \\ i \text{ isogeny.} \end{array} \right. \right\} \quad (1)$$

(The isogenies are not required to preserve the orientation.)



The Moduli of Isogenies

Definition

The *moduli stack of isogenies* over R is the functor $\text{Isog} : \mathbf{CAlg}_R \rightarrow \mathcal{S}$ given by

$$A \mapsto \left\{ i : E \rightarrow E' \left| \begin{array}{l} E, E' \in \mathcal{M}_{\text{ell}}^{\text{or}}(A), \\ i \text{ isogeny.} \end{array} \right. \right\} \quad (1)$$

(The isogenies are not required to preserve the orientation.)

Then $\mathcal{M}_{\text{ell}}^{\text{or}}(\Gamma)$ can be built from Isog .



The Moduli of Isogenies

Definition

The *moduli stack of isogenies* over R is the functor $\mathrm{Isog} : \mathcal{C}\mathrm{Alg}_R \rightarrow \mathcal{S}$ given by

$$A \mapsto \left\{ i : E \rightarrow E' \left| \begin{array}{l} E, E' \in \mathcal{M}_{\mathrm{ell}}^{\mathrm{or}}(A), \\ i \text{ isogeny.} \end{array} \right. \right\} \quad (1)$$

(The isogenies are not required to preserve the orientation.)

Then $\mathcal{M}_{\mathrm{ell}}^{\mathrm{or}}(\Gamma)$ can be built from Isog .

Main Theorem (GN)

Isog is a formal DM-stack.



The Moduli of Isogenies

Definition

The *moduli stack of isogenies* over R is the functor $\mathrm{Isog} : \mathcal{C}\mathrm{Alg}_R \rightarrow \mathcal{S}$ given by

$$A \mapsto \left\{ i : E \rightarrow E' \left| \begin{array}{l} E, E' \in \mathcal{M}_{\mathrm{ell}}^{\mathrm{or}}(A), \\ i \text{ isogeny.} \end{array} \right. \right\} \quad (1)$$

(The isogenies are not required to preserve the orientation.)

Then $\mathcal{M}_{\mathrm{ell}}^{\mathrm{or}}(\Gamma)$ can be built from Isog .

Main Theorem (GN)

Isog is a formal DM-stack.

Warning

It is not known whether Isog is an honest DM-stack.



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\mathbf{Ell}_{\text{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\mathbf{Ell}_{\mathbf{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.

Proof.



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\mathbf{Ell}_{\text{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.

Proof.

$$E \xrightarrow{i} E'$$



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\mathbf{Ell}_{\text{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.

Proof.

$$K \longrightarrow E \xrightarrow{i} E'$$



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\text{Ell}_{\text{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.

Proof.

$$\begin{array}{ccccc} K^{\circ} & & & & \\ \downarrow & & & & \\ K & \longrightarrow & E & \xrightarrow{i} & E' \\ \downarrow & & & & \\ K^{\text{ét}} & & & & \end{array}$$



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\mathbf{Ell}_{\text{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.

Proof.

$$\begin{array}{ccccc} K^\circ & \xlongequal{\quad} & K^\circ & & \\ \downarrow & & \downarrow & & \\ K & \longrightarrow & E & \xrightarrow{i} & E' \\ \downarrow & & \downarrow c & & \\ K^{\text{ét}} & \longrightarrow & E/K^\circ & & \end{array}$$



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\text{Ell}_{\text{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.

Proof.

$$\begin{array}{ccccc} K^\circ & \xlongequal{\quad} & K^\circ & & \\ \downarrow & & \downarrow & & \\ K & \longrightarrow & E & \xrightarrow{i} & E' \\ \downarrow & & \downarrow c & & \\ K^{\text{ét}} & \longrightarrow & E/K^\circ & \xrightarrow{e} & \widetilde{E}' \end{array}$$



Connected-Étale Factorization

Factorization Theorem (GN)

There is an orthogonal factorization system $(\mathcal{C}onn, \mathcal{E}t)$ on $\text{Ell}_{\text{Isog}}^{\text{or}}$ such that

- $\mathcal{C}onn$ is the class of connected isogenies, and
- $\mathcal{E}t$ is the class of étale isogenies.

This factorization system is natural with respect to change of base.

Proof.

$$\begin{array}{ccccc} K^\circ & \xlongequal{\quad} & K^\circ & \longrightarrow & 0 \\ \downarrow & & \downarrow & & \downarrow \\ K & \longrightarrow & E & \xrightarrow{i} & E' \\ \downarrow & & \downarrow c & & \parallel \\ K^{\text{ét}} & \longrightarrow & E/K^\circ & \xrightarrow{e} & \widetilde{E'} \end{array}$$



Digression: Why are the components isogenies?



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, classical version ([KM85])

Zariski-locally on the base, every morphism of classical elliptic curves is either 0 or an isogeny.



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \widetilde{R} & & & & \\ \downarrow & & & & \\ R & \longrightarrow & A & \xrightarrow{f} & B \end{array}$$



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \tilde{R} & \xrightarrow{\quad} & \tilde{A} & \xrightarrow{\quad} & \tilde{B} \\ \downarrow & & \downarrow & & \downarrow \\ R & \xrightarrow{\quad} & A & \xrightarrow{f} & B \end{array}$$



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \tilde{R} & \xrightarrow{\quad} & \tilde{A} & \xrightarrow{\quad} & \tilde{B} \\ \downarrow & & \downarrow & & \downarrow \\ R & \xrightarrow{\quad} & A & \xrightarrow{f} & B \end{array}$$

$$\begin{array}{c} \mathrm{Def}_{\tilde{A}}^{\tilde{A}}(B) \\ \downarrow \\ \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(B) \end{array}$$



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \tilde{R} & \xrightarrow{\quad} & \tilde{A} & \xrightarrow{\quad} & \tilde{B} \\ \downarrow & & \downarrow & & \downarrow \\ R & \xrightarrow{\quad} & A & \xrightarrow{f} & B \end{array}$$

$$\begin{array}{ccc} \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(f) & \longrightarrow & \mathrm{Def}_{\tilde{A}}^{\tilde{A}}(B) \\ \downarrow & \lrcorner & \downarrow \\ * & \longrightarrow & \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(B) \end{array}$$



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \tilde{R} & \xrightarrow{\quad} & \tilde{A} & \xrightarrow{\quad} & \tilde{B} \\ \downarrow & & \downarrow & & \downarrow \\ R & \xrightarrow{\quad} & A & \xrightarrow{f} & B \end{array}$$

f factors through R

$$\begin{array}{ccc} \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(f) & \longrightarrow & \mathrm{Def}_{\tilde{A}}^{\tilde{A}}(B) \\ \downarrow & \lrcorner & \downarrow \\ * & \longrightarrow & \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(B) \end{array}$$



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \tilde{R} & \xrightarrow{\quad} & \tilde{A} & \xrightarrow{\quad} & \tilde{B} \\ \downarrow & & \downarrow & & \downarrow \\ R & \xrightarrow{\quad} & A & \xrightarrow{f} & B \end{array}$$

$$\begin{array}{ccc} \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(f) & \longrightarrow & \mathrm{Def}_{\tilde{A}}^{\tilde{A}}(B) \\ \downarrow & \lrcorner & \downarrow \\ * & \longrightarrow & \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(B) \end{array}$$

f factors through $R \Rightarrow R$ is retract of A



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \tilde{R} & \xrightarrow{\quad} & \tilde{A} & \xrightarrow{\quad} & \tilde{B} \\ \downarrow & & \downarrow & & \downarrow \\ R & \xrightarrow{\quad} & A & \xrightarrow{f} & B \end{array}$$

$$\begin{array}{ccc} \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(f) & \longrightarrow & \mathrm{Def}_{\tilde{A}}^{\tilde{A}}(B) \\ \downarrow & \lrcorner & \downarrow \\ * & \longrightarrow & \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(B) \end{array}$$

f factors through $R \Rightarrow R$ is retract of A

$\Rightarrow B \otimes_A L_{A/R}$ vanishes



Digression: Why are the components isogenies?

Elliptic Rigidity Theorem, spectral version (GN)

Zariski-locally on the base, every morphism of strict elliptic curves is either 0 or an isogeny.

Proof sketch.

Main idea: 0 map deforms uniquely through square-zero extensions
 \Rightarrow through Postnikov tower.

$$\begin{array}{ccccc} \tilde{R} & \xrightarrow{\quad} & \tilde{A} & \xrightarrow{\quad} & \tilde{B} \\ \downarrow & & \downarrow & & \downarrow \\ R & \xrightarrow{\quad} & A & \xrightarrow{f} & B \end{array}$$

$$\begin{array}{ccc} \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(f) & \longrightarrow & \mathrm{Def}_{\tilde{A}}^{\tilde{A}}(B) \\ \downarrow & \lrcorner & \downarrow \\ * & \longrightarrow & \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(B) \end{array}$$

f factors through $R \Rightarrow R$ is retract of A

$\Rightarrow B \otimes_A L_{A/R}$ vanishes

$\Rightarrow \mathrm{Def}_{\tilde{R}}^{\tilde{R}}(f) = 0. \quad \square$



Corollary

We have a pullback of functors

$$\begin{array}{ccc} \mathrm{Isog} & \longrightarrow & \mathrm{Isog}^{\mathrm{ét}} \\ \downarrow & & \downarrow s \\ \mathrm{Isog}^{\mathrm{conn}} & \xrightarrow{t} & \mathcal{M}_{\mathrm{ell}}^{\mathrm{or}} \end{array}$$



Corollary

We have a pullback of functors

$$\begin{array}{ccc} \mathrm{Isog} & \longrightarrow & \mathrm{Isog}^{\mathrm{\acute{e}t}} \\ \downarrow & & \downarrow s \\ \mathrm{Isog}^{\mathrm{conn}} & \xrightarrow{t} & \mathcal{M}_{\mathrm{ell}}^{\mathrm{or}}. \end{array}$$

Just need to show that $\mathrm{Isog}^{\mathrm{\acute{e}t}}, \mathrm{Isog}^{\mathrm{conn}}$ are formal DM-stacks.



Theorem (GN)

$\mathrm{Isog}^{\mathrm{\acute{e}t}}$ is a DM-stack.



Theorem (GN)

$\mathrm{Isog}^{\mathrm{\acute{e}t}}$ is a DM-stack.

Proof sketch.



Theorem (GN)

$\text{Isog}^{\text{ét}}$ is a DM-stack.

Proof sketch.

- 1 [KM85]: $\{(E, K) \mid E \text{ elliptic curve}, K \subset E \text{ finite}\} \rightarrow (\mathcal{M}_{\text{ell}}^{\text{or}})^{\heartsuit}$
relative scheme.



Theorem (GN)

$\text{Isog}^{\text{ét}}$ is a DM-stack.

Proof sketch.

- 1 [KM85]: $\{(E, K) \mid E \text{ elliptic curve}, K \subset E \text{ finite}\} \rightarrow (\mathcal{M}_{\text{ell}}^{\text{or}})^{\heartsuit}$
relative scheme.
- 2 $\{(E, K) \mid E \text{ elliptic curve}, K \subset E \text{ finite étale}\}$ open substack.



Theorem (GN)

$\mathrm{Isog}^{\mathrm{\acute{e}t}}$ is a DM-stack.

Proof sketch.

- 1 [KM85]: $\{(E, K) \mid E \text{ elliptic curve}, K \subset E \text{ finite}\} \rightarrow (\mathcal{M}_{\mathrm{ell}}^{\mathrm{or}})^{\heartsuit}$
relative scheme.
- 2 $\{(E, K) \mid E \text{ elliptic curve}, K \subset E \text{ finite \acute{e}tale}\}$ open substack.
- 3 Leverage \acute{e}taleness and use ([Lur18c], Theorem 18.1.0.2) to lift from classical to spectral. \square



Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.



Identifying $\text{Isog}^{\text{conn}}$

Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch.



Identifying $\text{Isog}^{\text{conn}}$

Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch.

$$\{E \xrightarrow{\text{conn}} E'\}$$



Identifying $\mathrm{Isog}^{\mathrm{conn}}$

Theorem (GN)

$\mathrm{Isog}^{\mathrm{conn}}$ is a formal algebraic space.

Proof sketch.

$$\{E \xrightarrow{\mathrm{conn}} E'\}$$



$\{K \subset E \text{ closed, proper, connected} \quad \quad \quad \}$



Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch.

$$\begin{array}{c} \{E \xrightarrow{\text{conn}} E'\} \\ \updownarrow \\ \{K \subset E \text{ closed, proper, connected} \} \\ \updownarrow \\ \{K \subset \widehat{\mathbb{G}}_R^Q \text{ honest subgroup} \} \end{array}$$



Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch.

$$\{E \xrightarrow{\text{conn}} E'\}$$

$$\updownarrow$$

$$\{K \subset E \text{ closed, proper, connected AND equiv } \widehat{E/K} \simeq \widehat{\mathbb{G}}_R^Q\}$$

$$\updownarrow$$

$$\{K \subset \widehat{\mathbb{G}}_R^Q \text{ honest subgroup AND equiv } \widehat{\mathbb{G}}_R^Q/K \simeq \widehat{\mathbb{G}}_R^Q\}$$



Identifying $\text{Isog}^{\text{conn}}$

Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch.

$$\{E \xrightarrow{\text{conn}} E'\}$$

$$\updownarrow$$

$$\{K \subset E \text{ closed, proper, connected AND equiv } \widehat{E/K} \simeq \widehat{\mathbb{G}}_R^Q\}$$

$$\updownarrow$$

$$\{K \subset \widehat{\mathbb{G}}_R^Q \text{ honest subgroup AND equiv } \widehat{\mathbb{G}}_R^Q/K \simeq \widehat{\mathbb{G}}_R^Q\}$$

$$\Rightarrow \text{Isog}^{\text{conn}} \simeq \mathcal{M}_{\text{ell}}^{\text{or}} \times \text{QuilIsog}$$



Identifying $\text{Isog}^{\text{conn}}$

Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

$$\begin{array}{c} \text{QuilIsog} \\ \downarrow \\ \text{Sub}^h(\widehat{\mathbb{G}}_R^Q) \end{array}$$



Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

$$\begin{array}{ccc} & & \text{QuilIsog} \\ & & \downarrow \\ * & \xrightarrow{K} & \text{Sub}^h(\widehat{\mathbb{G}}_R^Q) \end{array}$$



Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

$$\begin{array}{ccc} \text{OrDat}(\widehat{\mathbb{G}}_R^Q/K) & \longrightarrow & \text{QuilIsog} \\ \downarrow & \lrcorner & \downarrow \\ * & \xrightarrow{K} & \text{Sub}^h(\widehat{\mathbb{G}}_R^Q) \end{array}$$



Theorem (GN)

$\mathrm{Isog}^{\mathrm{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

$$\begin{array}{ccc} \mathrm{OrDat}(\widehat{\mathbb{G}}_R^Q/K) & \longrightarrow & \mathrm{QuilIsog} \\ \downarrow & \lrcorner & \downarrow \\ * & \xrightarrow{K} & \mathrm{Sub}^h(\widehat{\mathbb{G}}_R^Q) \end{array}$$

[Lur18b]: $\mathrm{OrDat}(\widehat{\mathbb{G}}_R^Q/K)$ is an affine DM-stack.



Theorem (GN)

$\mathrm{Isog}^{\mathrm{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

$$\begin{array}{ccc} \mathrm{OrDat}(\widehat{\mathbb{G}}_R^Q/K) & \longrightarrow & \mathrm{QuilIsog} \\ \downarrow & \lrcorner & \downarrow \\ * & \xrightarrow{K} & \mathrm{Sub}^h(\widehat{\mathbb{G}}_R^Q) \end{array}$$

[Lur18b]: $\mathrm{OrDat}(\widehat{\mathbb{G}}_R^Q/K)$ is an affine DM-stack.

\Rightarrow Enough to show that $\mathrm{Sub}^h(\widehat{\mathbb{G}}_R^Q)$ is formal algebraic space.



Identifying $\text{Isog}^{\text{conn}}$

Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

We have a retract:



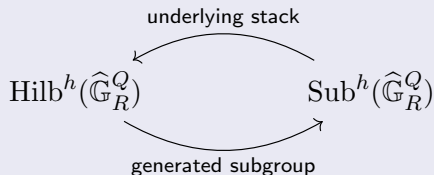
Identifying $\text{Isog}^{\text{conn}}$

Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

We have a retract:

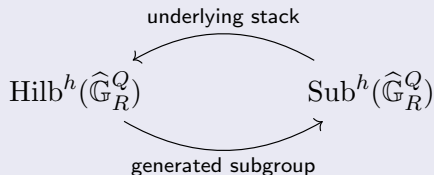


Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

We have a retract:



[Lur04]: Hilb of separated algebraic space is algebraic space.

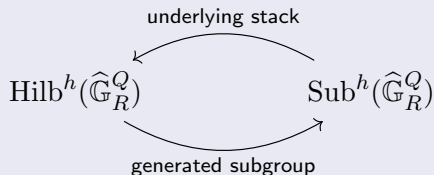


Theorem (GN)

$\text{Isog}^{\text{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

We have a retract:



[Lur04]: Hilb of separated algebraic space is algebraic space.
 $\Rightarrow \text{Hilb}^h(\widehat{\mathbb{G}}_R^Q)$ is formal algebraic space.

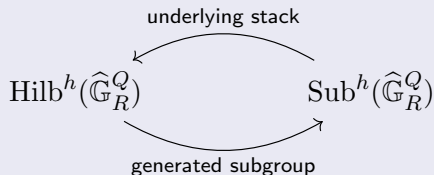


Theorem (GN)

$\mathrm{Isog}^{\mathrm{conn}}$ is a formal algebraic space.

Proof sketch (ctd).

We have a retract:



[Lur04]: Hilb of separated algebraic space is algebraic space.

$\Rightarrow \mathrm{Hilb}^h(\widehat{\mathbb{G}}_R^Q)$ is formal algebraic space.

$\Rightarrow \mathrm{Sub}^h(\widehat{\mathbb{G}}_R^Q)$ is formal algebraic space. □



Thank you!

References

- [KM85] Nicholas M. Katz and Barry Mazur. *Arithmetic Moduli of Elliptic Curves*. 108. Princeton University Press, 1985.
- [Lur04] Jacob Lurie. “Derived Algebraic Geometry”. PhD thesis. Massachusetts Institute of Technology, 2004. URL: <http://oastats.mit.edu/handle/1721.1/30144>.
- [Lur18a] Jacob Lurie. *Elliptic Cohomology*. 2018. URL: <https://www.math.ias.edu/~lurie/papers/Elliptic-I.pdf>. Pre-published.
- [Lur18b] Jacob Lurie. *Elliptic Cohomology II: Orientations*. Apr. 2018. URL: <https://www.math.ias.edu/~lurie/papers/Elliptic-II.pdf>. Pre-published.
- [Lur18c] Jacob Lurie. *Spectral Algebraic Geometry*. 2018. URL: <https://www.math.ias.edu/~lurie/papers/SAG-rootfile.pdf>.
- [MZ25] Xuecai Ma and Yifei Zhu. *Spectral Moduli Problems for Level Structures and an Integral Jacquet-Langlands Dual of Morava E-theory*. 2025. URL: <https://yifeizhu.github.io/sagreal.pdf>. Pre-published.

