
INTRODUCTION TO GENETIC PROGRAMMING

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OUTLINE

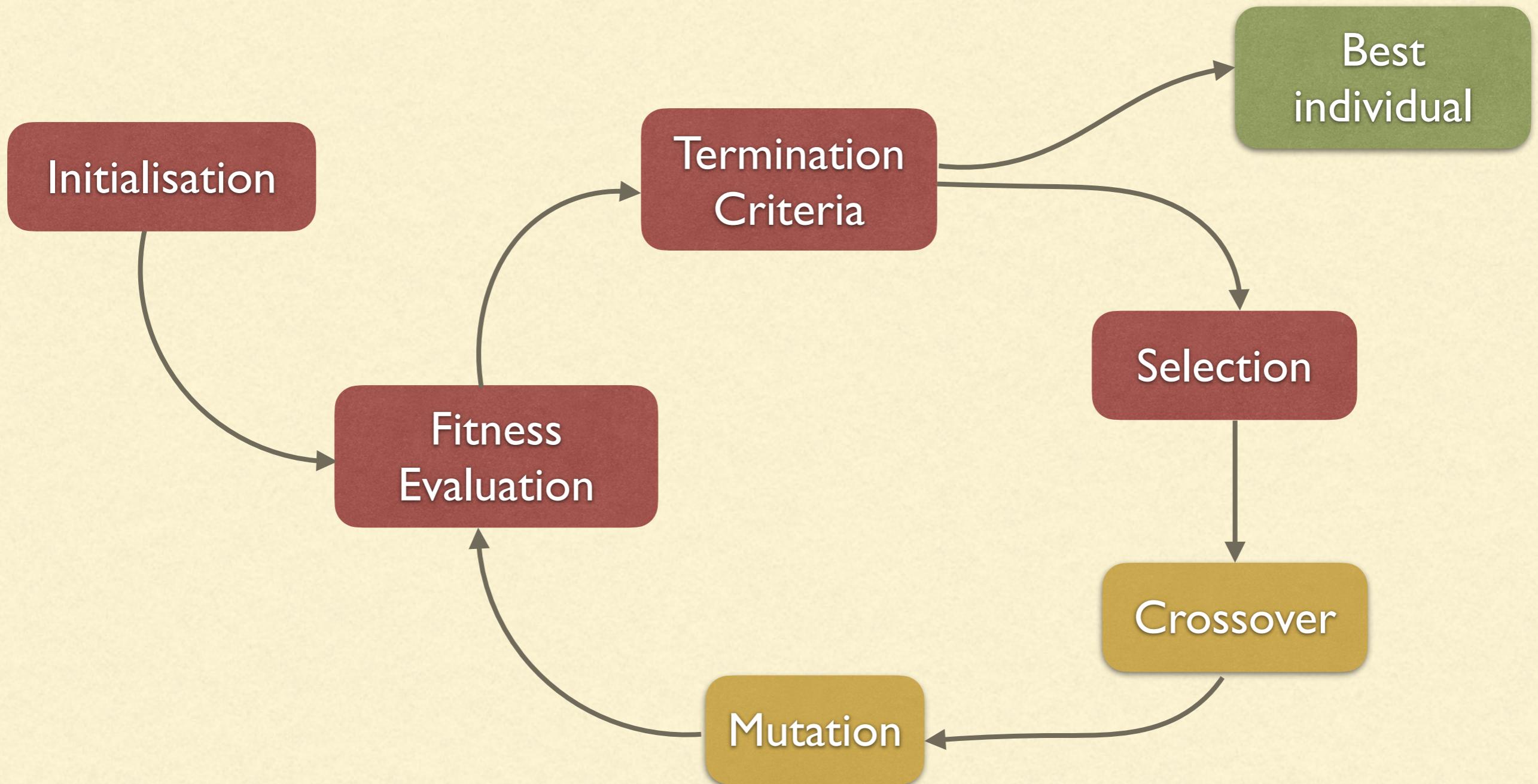
- How programs are represented: tree-based GP
- Initialisation, Crossover, and Mutation
- Overfitting and how to combat bloat
 - A PARTICULAR TYPE OF
OVERFITTING

WHAT'S GP?

- GP is a technique to *stochastically evolve* a *population* (multiset) of individuals encoding *computer programs*
- John Koza was one of the firsts to talk about evolving computer programs in the late 80s → TIMELINE FROM GENETIC ALGORITHM TO EVOLUTION STRATEGIES AND THEN GENETIC PROGRAMMING
- While the main “evolutionary cycle” is like the one for GA...
- ...GP has many peculiar properties, starting from the representation used.

EVOLUTION CYCLE

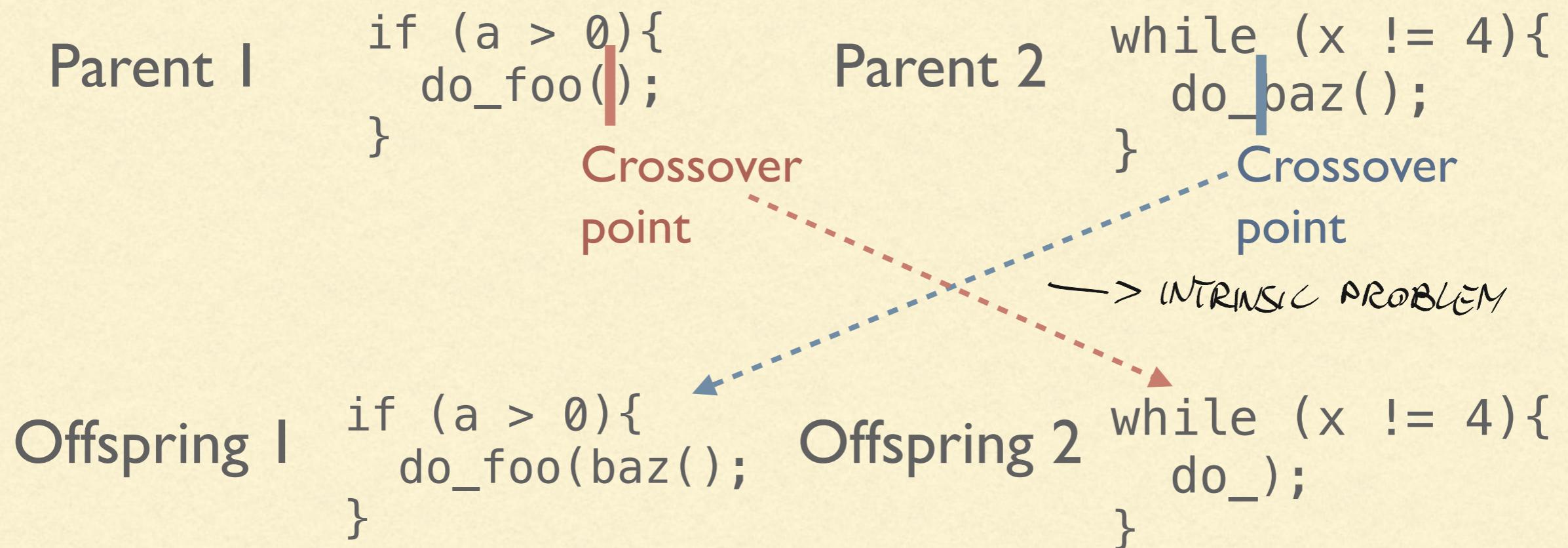
This is exactly like GA, but...



PROGRAM REPRESENTATION

WE PRACTICALLY HAVE A STRING OF TEXT

How can we represent programs that evolve?

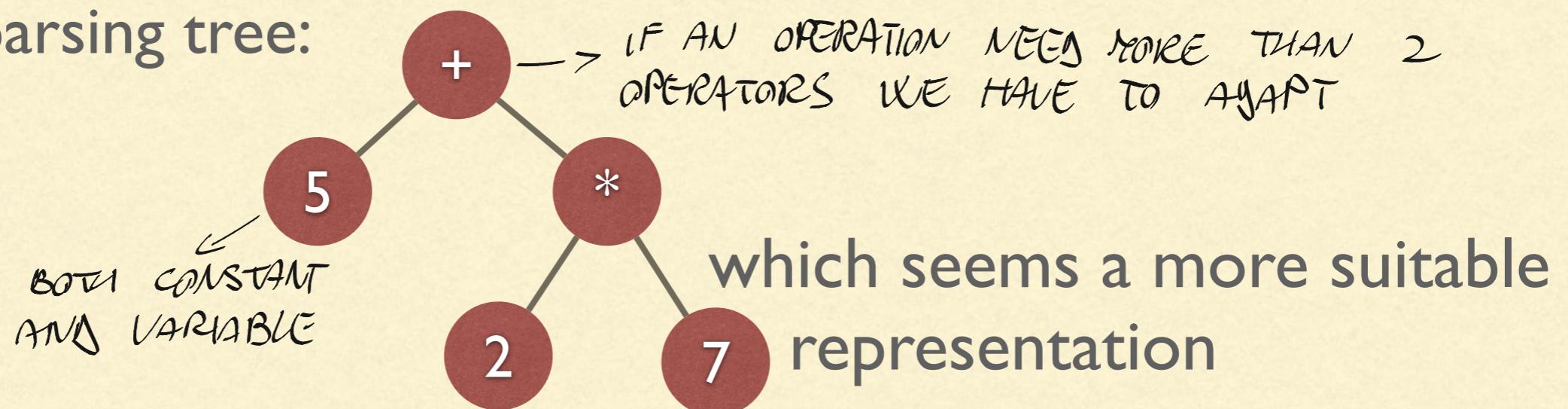


WE RISK TO MAKE A VERY BIG ERROR
A STRING CAN NOT BE USED

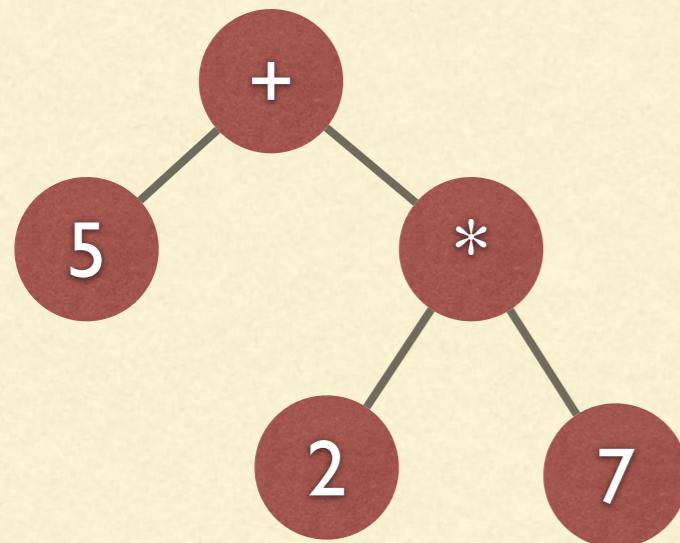
⇒ WE NEEDS ANOTHER INTERPRETATION

PROGRAM REPRESENTATION

- We must select a representation where crossover and mutation can easily work...
- ...which means that “a string” is usually not the best choice
- If we start with an expression like “ $5 + (7 * 2)$ ” we can represent it as its parsing tree:



PROGRAM REPRESENTATION



Crossover and mutation can be written as operations on subtrees

Evaluation of a tree is quite simple

ONCE WE CHOOSE A REPRESENTATION WHAT WE ACTUALLY NEED:

What do we need to encode a program as a tree?

What are the possible internal nodes and leaves?

TERMINAL SET

The terminal set contains all the possible leaves, for example:

Constants

{0,1,2,3,4,...}

{true, false}

{ e , π , -1 , ...}

Input variables

{ x_0, x_1, \dots }

FUNCTION SET

The function set contains the possible inner nodes, for example:

Arithmetical operations $\{ +, -, \times, \div \}$

Trigonometric functions $\{ \sin, \cos, \tan \}$

Boolean operators $\{ \wedge, \vee, \neg \}$

Choice/conditional $\{ \text{if } \dots \text{ then } \dots \text{ else } \dots \}$

CLOSURE

- The primitives sets (functionals and terminals) must respect the property of closure:
- Intuitively, we must be able to “mix” the primitives without problems. *⇒ THE OUTPUT OF EVERY FUNCTION MUST BE USABLE FROM EVERY OTHER FUNCTION.*
- Type consistency: the terminals and the output of any functional must be valid inputs to all functionals
IT MEANS THAT IN WHATEVER PART OF THE TREE IS, IT CAN BE USED BY ANY OTHER FUNCTION
- Evaluation safety: primitives that can fail at runtime (e.g., division) should be “protected” to avoid runtime failures.
EX: %0 INSTEAD OF RETURN FAILURE YOU CAN RETURN A DEFAULT VALUE

SUFFICIENCY

- To find a solution we must be able to represent it, which means that the primitives must be sufficient to write a solution
- For example: with real constants, variables, and $\{+, -, *\}$ we can represent any polynomial...
- ...which is not very useful if the function that we must fit is an exponential
- Usually we cannot assure sufficiency, but we might still obtain solutions that are good approximations

WE HAVE TO MAKE PEACE WITH IT

INITIALISATION METHODS

INITIALISATION

- For binary strings (traditional GA), a random generation is easy: select each bit independently with a uniform probability
- Finite trees are infinite in number...
- ...and to generate them we must recall that:
 - ↳ WHICH MEANS THAT WE NEED A SPECIFIC FUNCTION
- “*Random numbers should not be generated with a method chosen at random.*” — Donald Knuth
 - WE COULD END UP WITH ALL TREES WITH THE SAME EXPRESSION OR IN THE SAME SPACE THAT COMPUTE THE SAME FUNCTION
- ...also holds for randomly generating trees.

$$(x \cdot (x \cdot (x \cdot (x \cdot x))))$$

GROW

Up to a maximum depth, select randomly across all primitives
Once the maximum depth is reached select a terminal

$$\mathcal{F} = \{ +, *, -, /\}$$

$$\mathcal{T} = \{x, y, 1\}$$

NEEDS TO BE LARGE ENOUGH TO REPRESENT THE GOOD SOLUTION BUT NOT TOO LARGE TO END UP WITH ENORMOUS TREE.
BETWEEN 5 AND 20 SHOULD BE GOOD

EX: IF THE FUNCTION IS GRADE 5 WE NEED AT LEAST A DEPTH OF 3

$$\text{max_depth} = 2$$



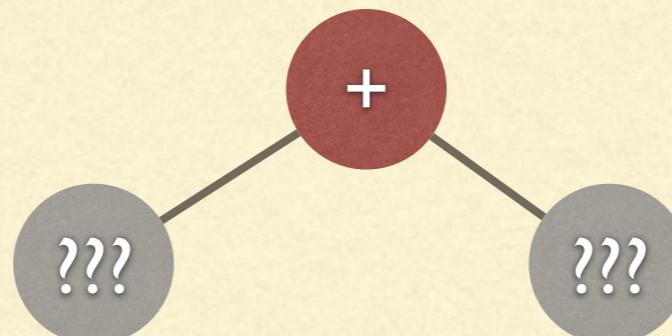
GROW

Up to a maximum depth, select randomly across all primitives
Once the maximum depth is reached select a terminal

ALWAYS START WITH
FUNCTIONS

$$\mathcal{F} = \{ +, *, -, / \}$$

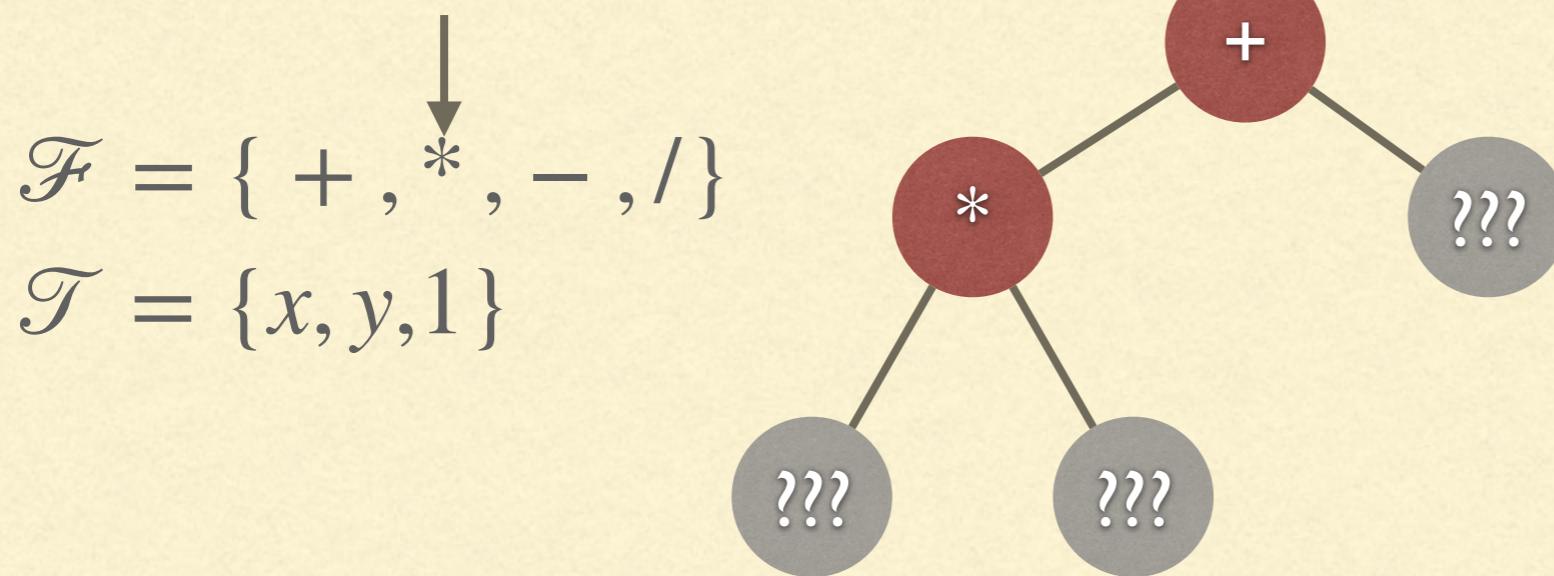
$$\mathcal{T} = \{x, y, 1\}$$



$$\text{max_depth} = 2$$

GROW

Up to a maximum depth, select randomly across all primitives
Once the maximum depth is reached select a terminal



`max_depth = 2`

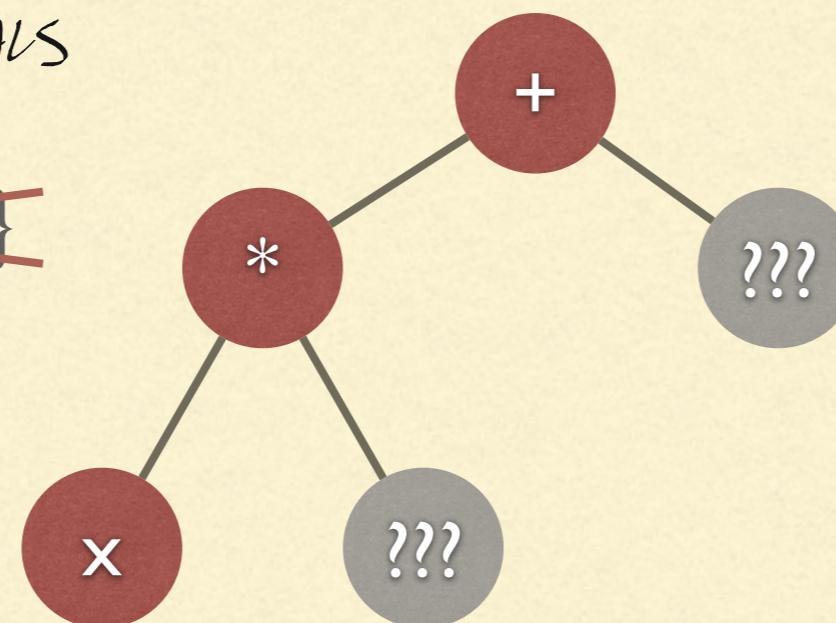
GROW

Up to a maximum depth, select randomly across all primitives
Once the maximum depth is reached select a terminal

AND END WITH TERMINALS

$$\mathcal{F} = \{ +, *, -, / \}$$

$$\mathcal{T} = \{ x, y, 1 \}$$



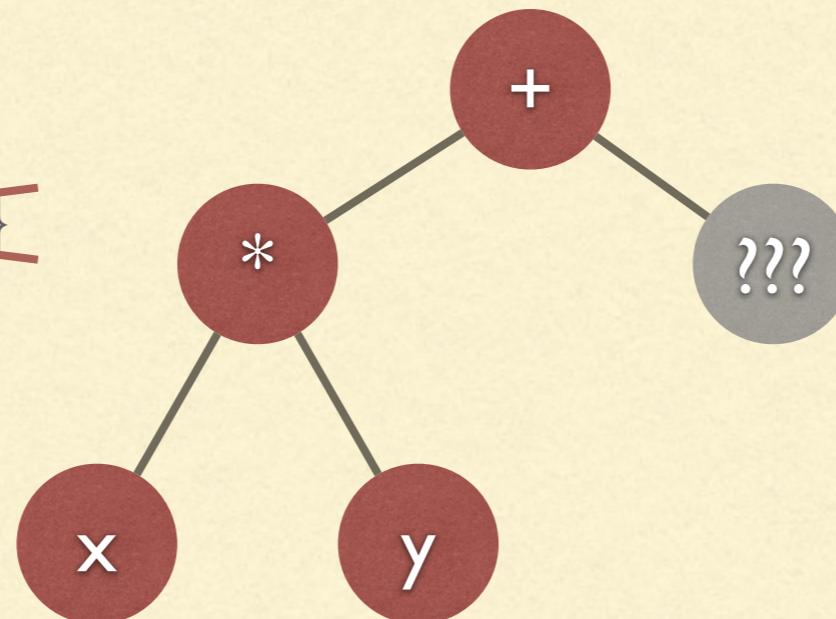
$$\text{max_depth} = 2$$

GROW

Up to a maximum depth, select randomly across all primitives
Once the maximum depth is reached select a terminal

$$\mathcal{F} = \{ +, *, -, / \}$$

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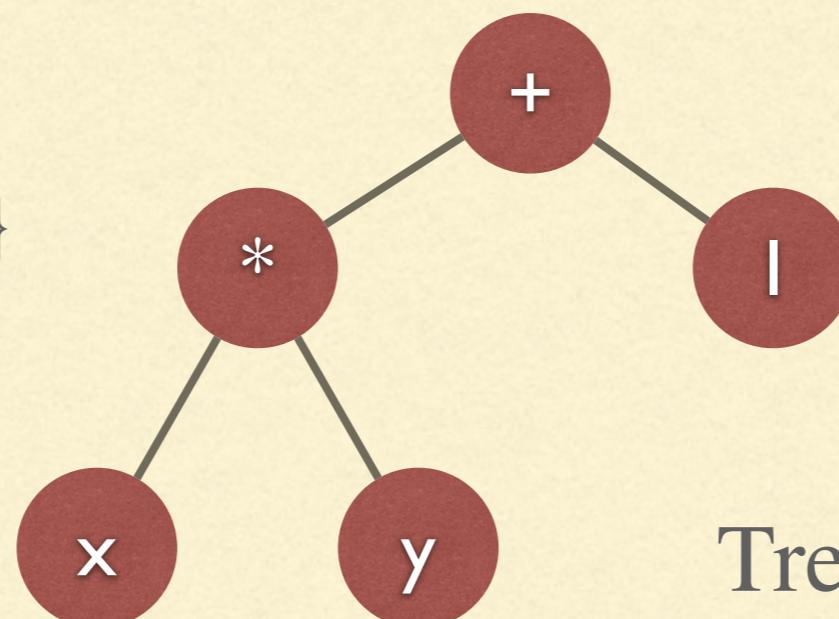
max_depth = 2

GROW

Up to a maximum depth, select randomly across all primitives
Once the maximum depth is reached select a terminal

$$\mathcal{F} = \{ +, *, -, / \}$$

$$\mathcal{T} = \{x, y, 1\}$$



$$\text{Tree} = xy + 1$$

$$\text{max_depth} = 2$$

THIS COULD BE
A SIMPLE BUT
FUNCTIONING
INITIALIZATION

FULL

- Like the “grow” method, but only functional symbols are selected before reaching the maximum depth
- This means that terminals appears only in the last level of the tree
- Different distribution of trees w.r.t. the “grow” method
(obviously)
THERE ARE TREES THAT CANNOT GENERATE BY FULL THAT GROW CAN AND VICE VERSA.
- Generally bigger trees

RAMPED HALF AND HALF

- Randomly select between “grow” and “full”...
- ...with a random maximum depth between a minimum and a maximum
- Generally trees with a better “variability” among them (non all representing similar functions)

EPHEMERAL CONSTANTS

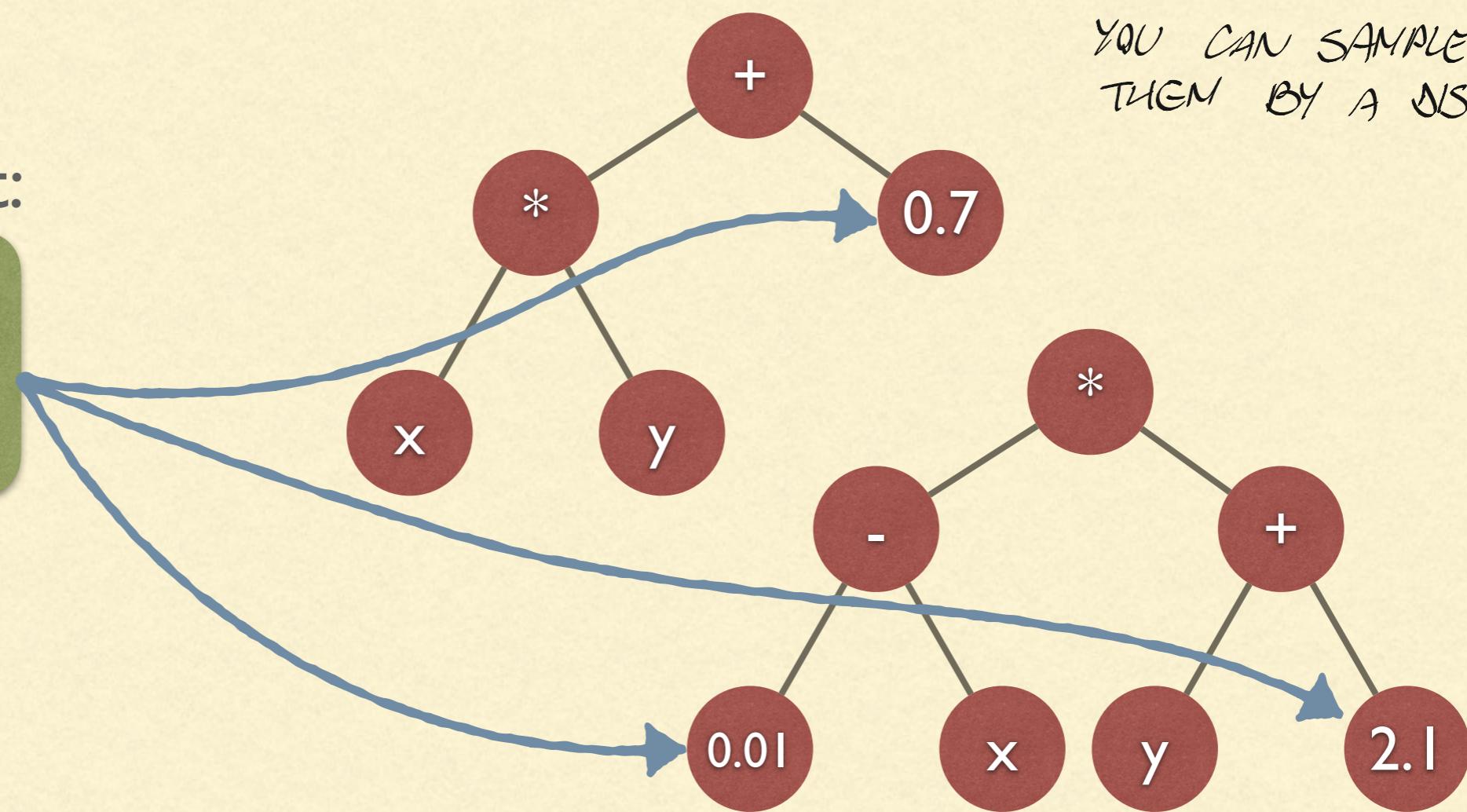
ONE ADDITIONAL STEP IN INITIALIZATION

The terminals might include constant. But how to chose them?
An alternative to the choice is the use of ephemeral constants

In the
terminal set:

A random
constant
in $[0,3]$

YOU CAN SAMPLE
THEM BY A DISTRIBUTION



CROSSOVER

HOMOLOGOUS CROSSOVER

Lets move back to GA

$x = \boxed{0} \boxed{1} \boxed{0} \boxed{1} \boxed{1} \boxed{0}$

Can we obtain a different individual
by crossing x with itself?

$x = \boxed{0} \boxed{1} \boxed{0} \boxed{1} \boxed{1} \boxed{0}$

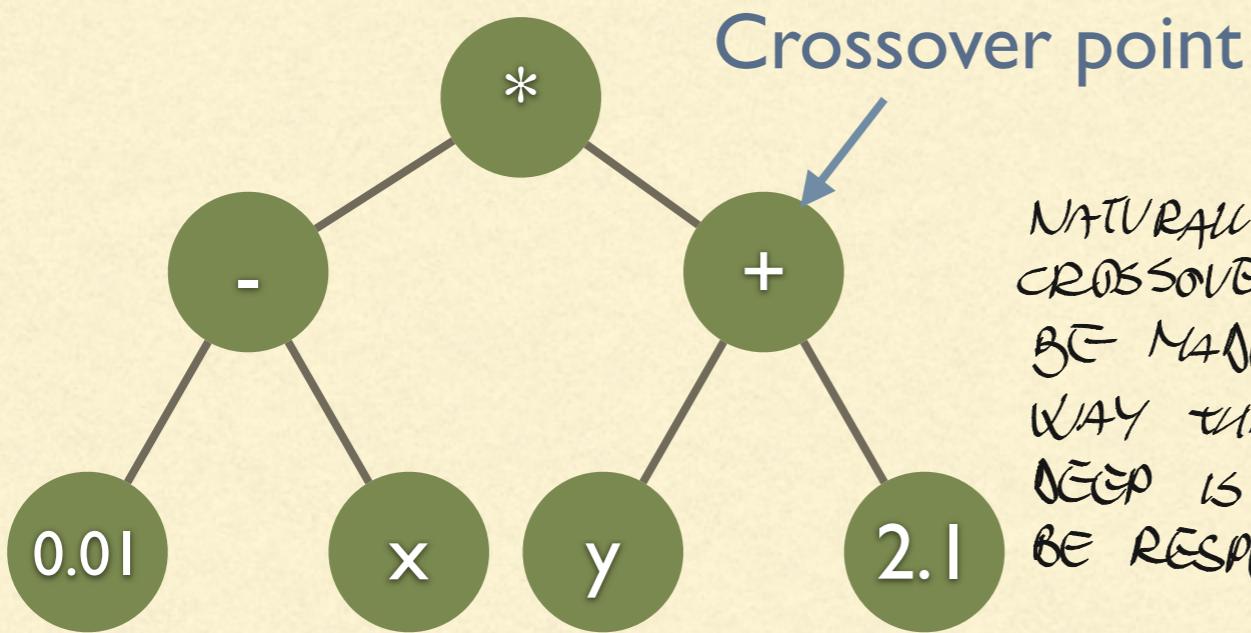
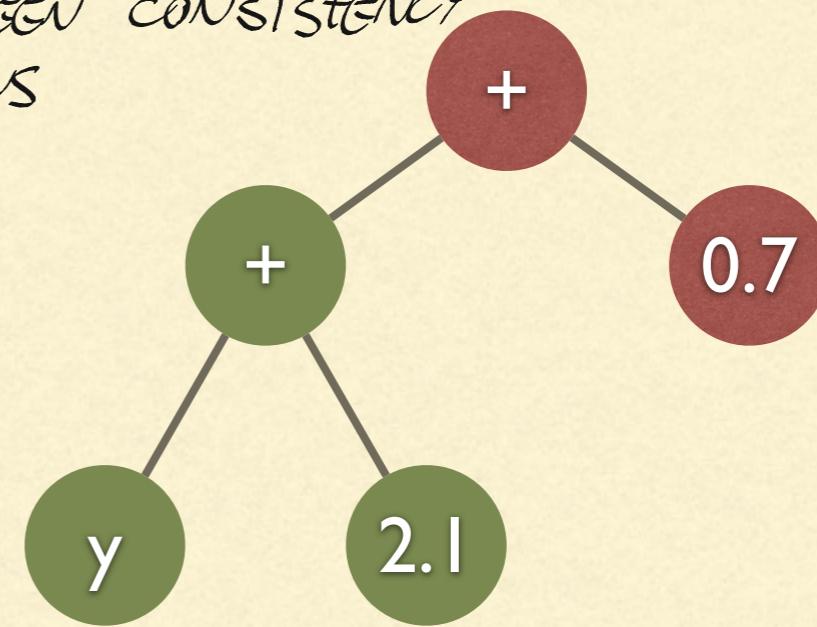
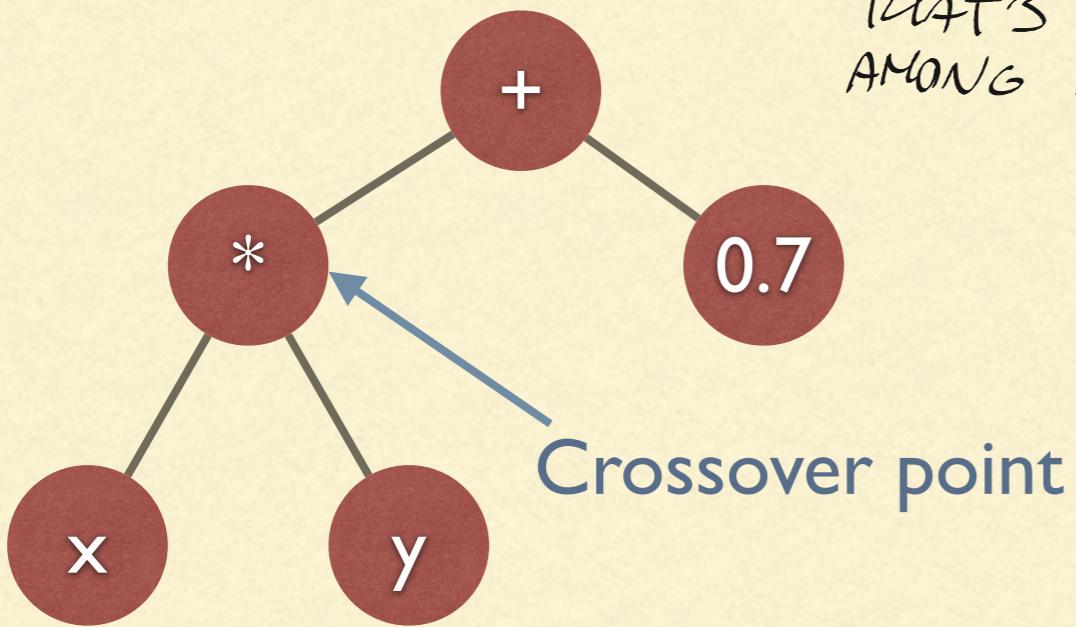
NO

The classical GA crossovers are homologous

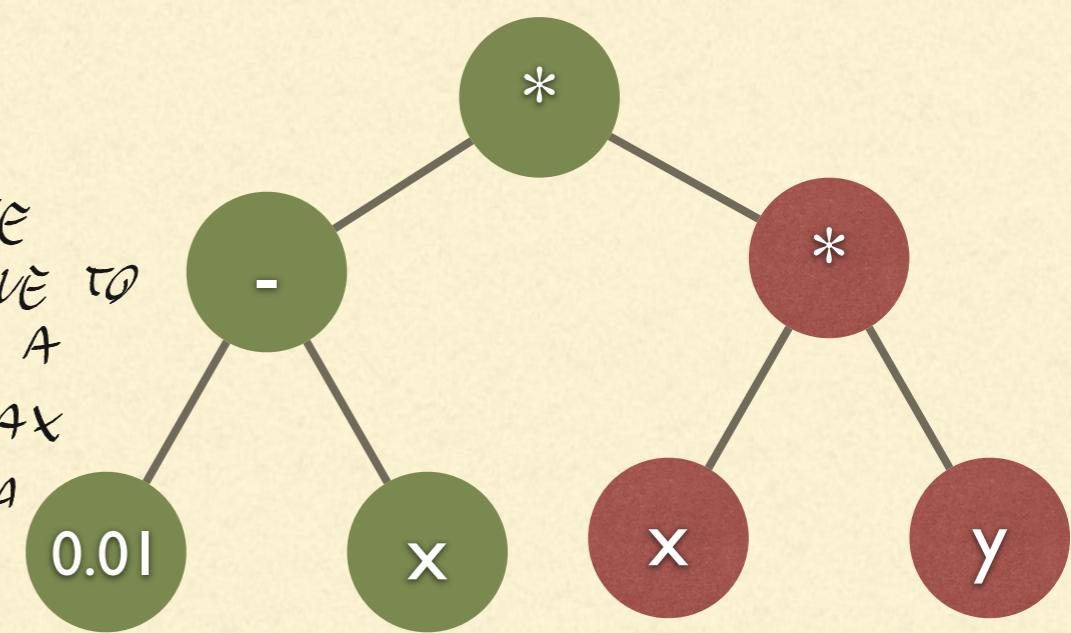
This is not usually the case for GP

SUBTREE CROSSOVER

WE CAN SELECT RANDOM POINT
THAT'S WHY WE NEED CONSISTENCY
AMONG ALL FUNCTIONS



NATURALLY THE
CROSSOVER HAVE TO
BE MADE IN A
WAY THAT MAX
DEPTH IS GUNNA
BE RESPECTED

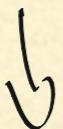


SUBTREE CROSSOVER

- In many cases trees have a maximum depth during evolution
- Which means that subtree crossover must not exceed it
- Subtree crossover is non-homologous...
- ...but there exist crossovers for GP that are homologous

WE ARE GOING TO
SEE THEM LATER

MUTATION(S)

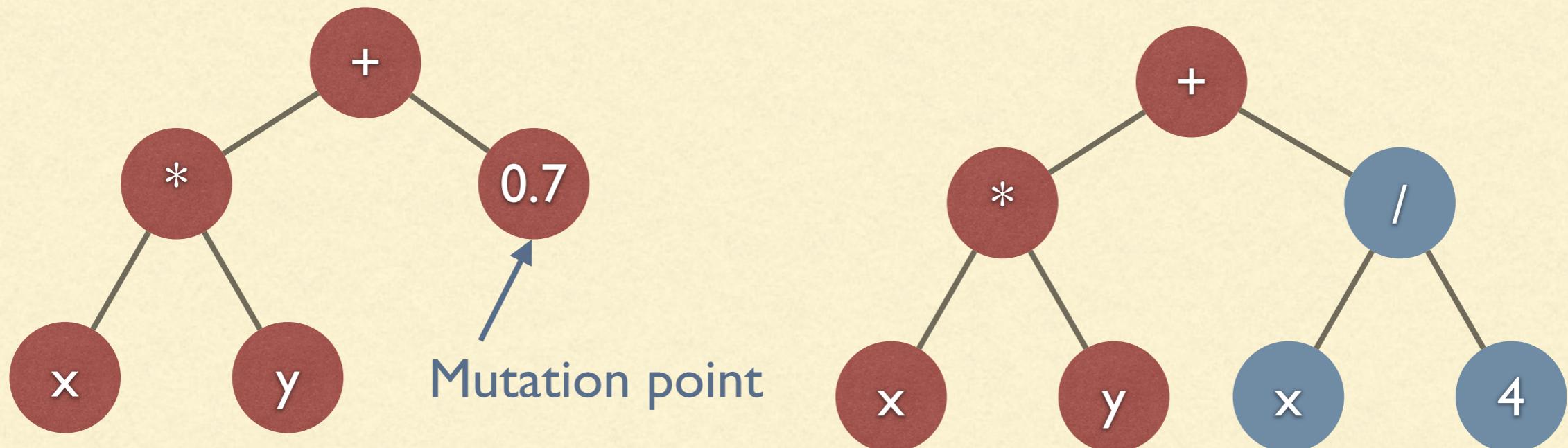


THEY CAN BE DEFINED IN
MANY WAY

SUBTREE MUTATION

Replacement of a randomly selected subtree
with a new random subtree

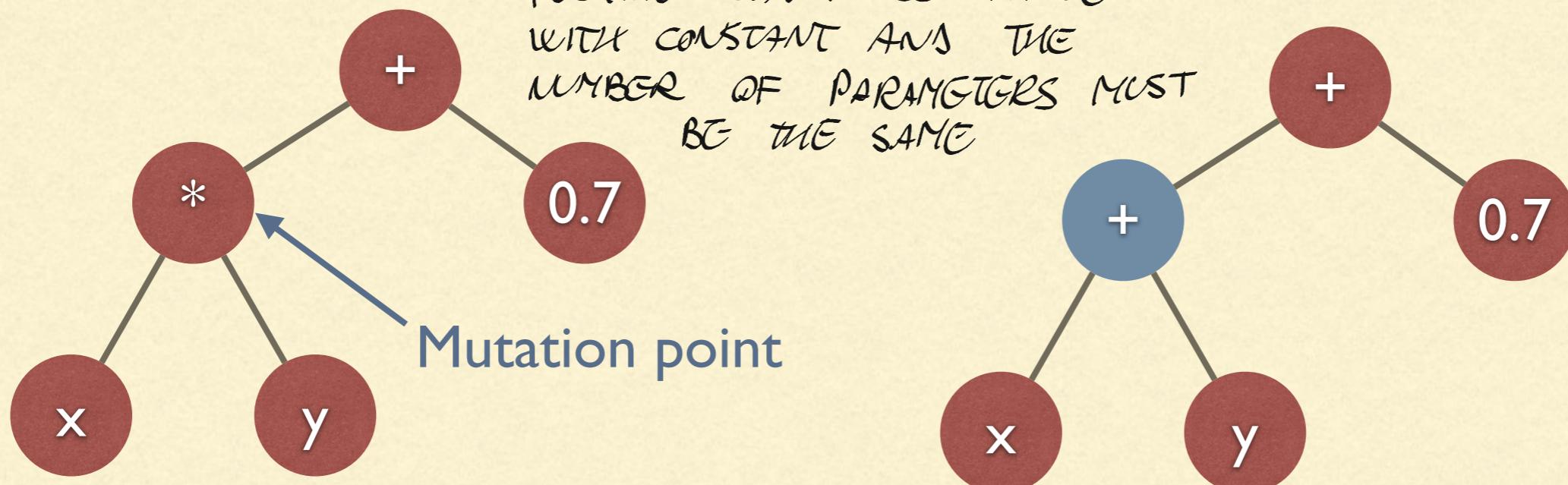
THE TREE CAN EITHER GROW OR REDUCE IN SIZE



POINT MUTATION

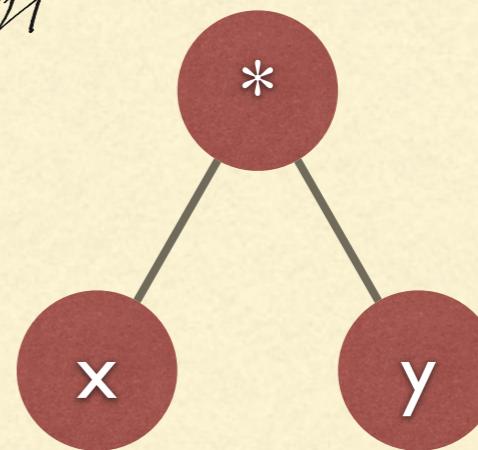
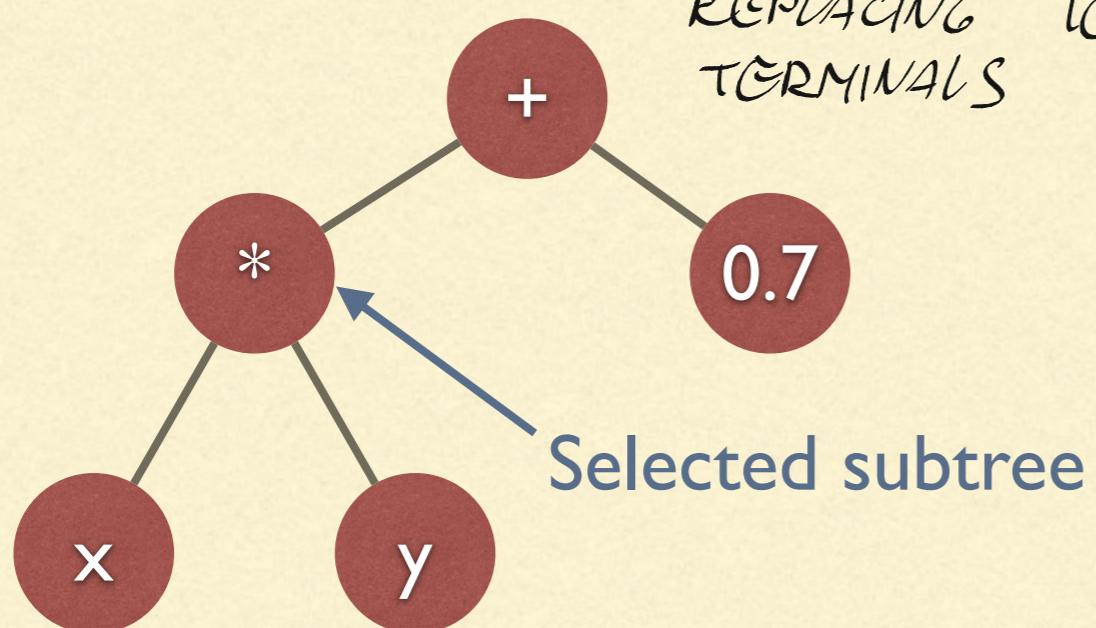
Replacement of a randomly selected node
with a compatible randomly selected node

THIS IS MORE STRICT.
FUNCTION CAN'T BE CHANGE
WITH CONSTANT AND THE
NUMBER OF PARAMETERS MUST
BE THE SAME



HOIST MUTATION

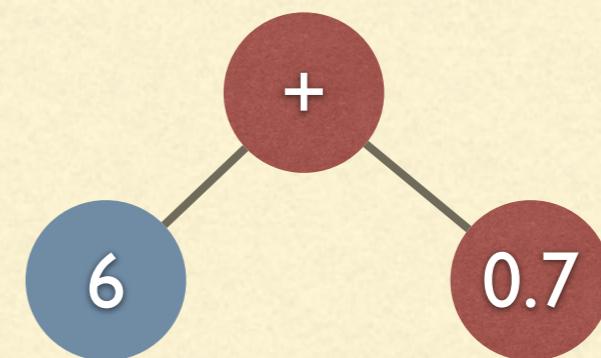
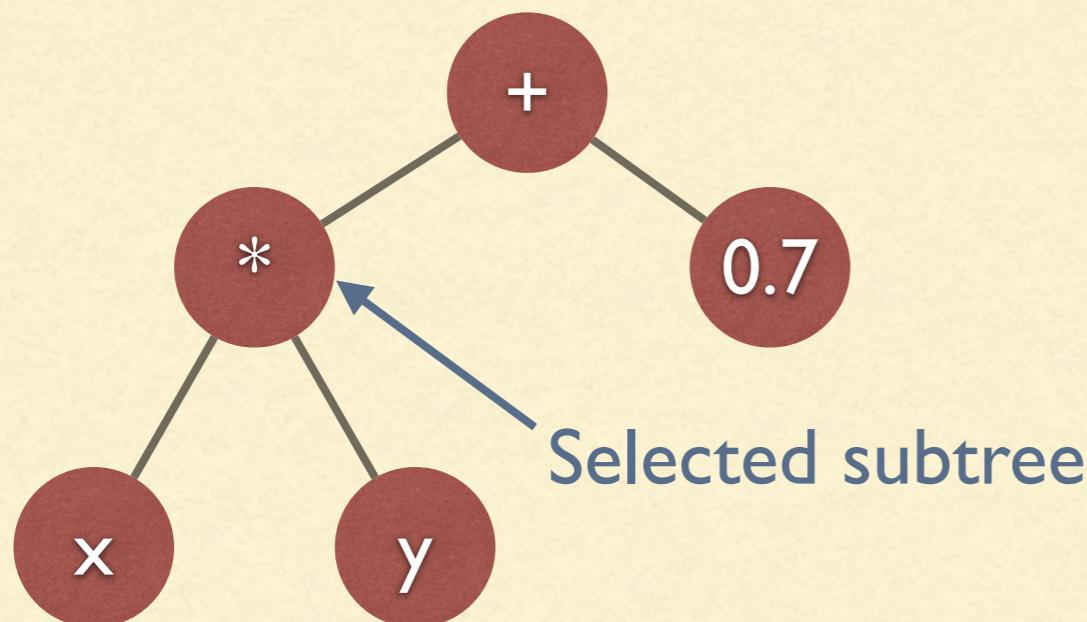
Replacement of the entire tree
with one of its subtree



Used to reduce program size

SHRINK MUTATION

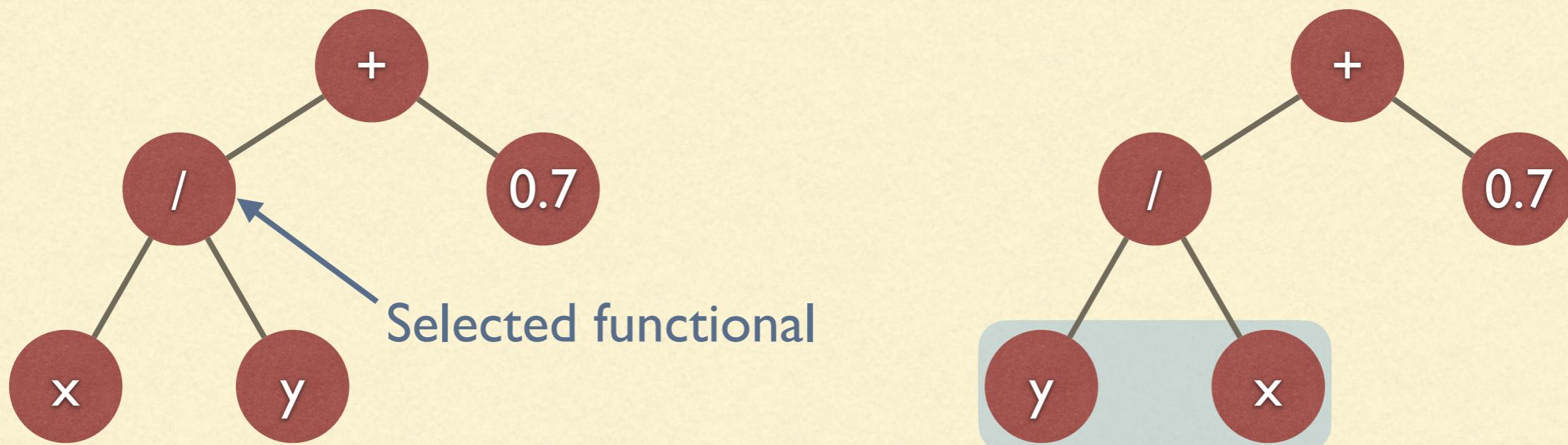
Replacement a randomly selected subtree
with a randomly selected terminal



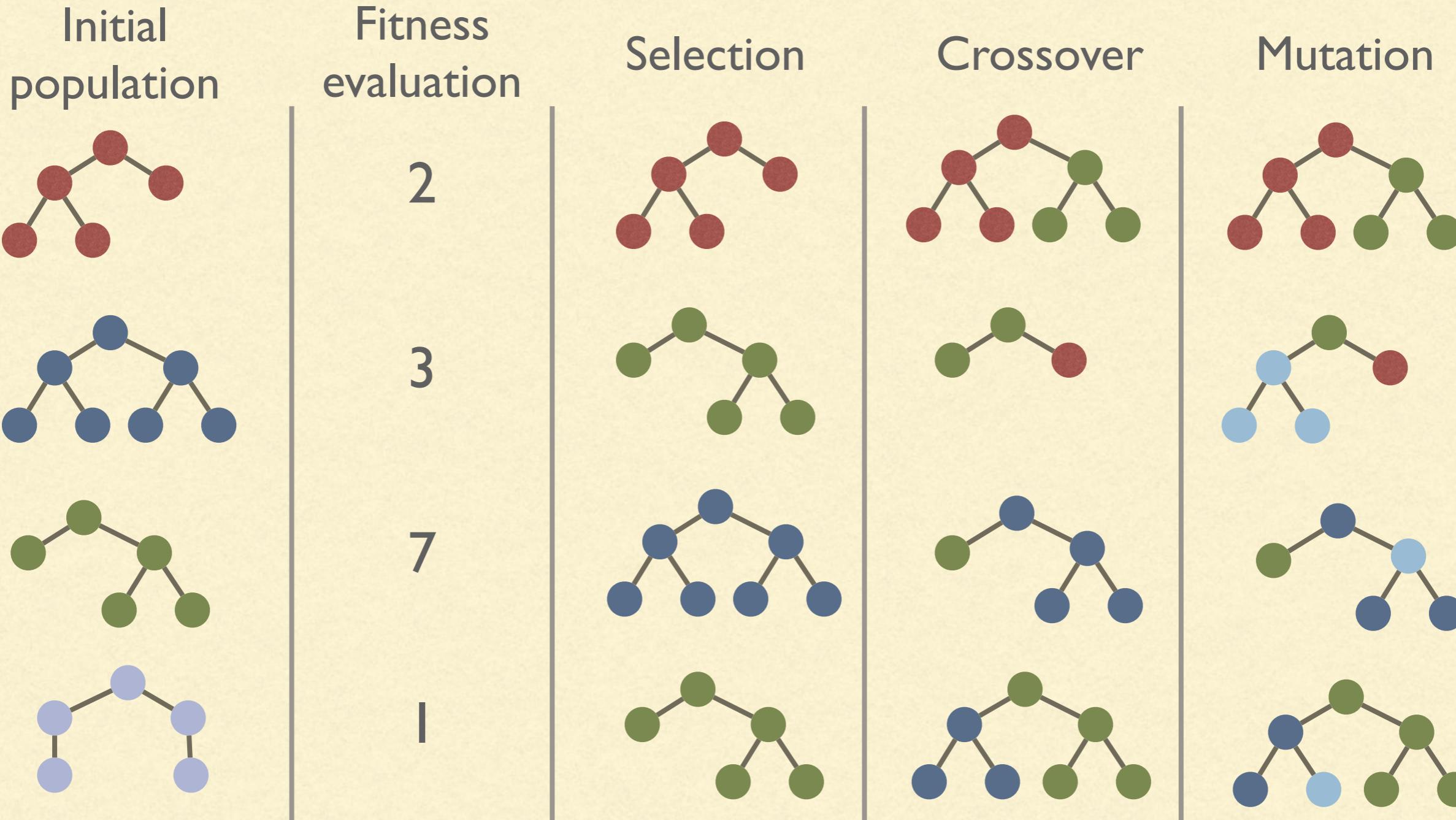
Used to reduce
program size

PERMUTATION/SWAP MUTATION

Apply a permutation to the arguments of a functional.
Swap mutation is a special case when only non-commutative
binary operators have their arguments swapped



A GENERATION OF GP



AUTOMATICALLY DEFINED FUNCTIONS - ADF

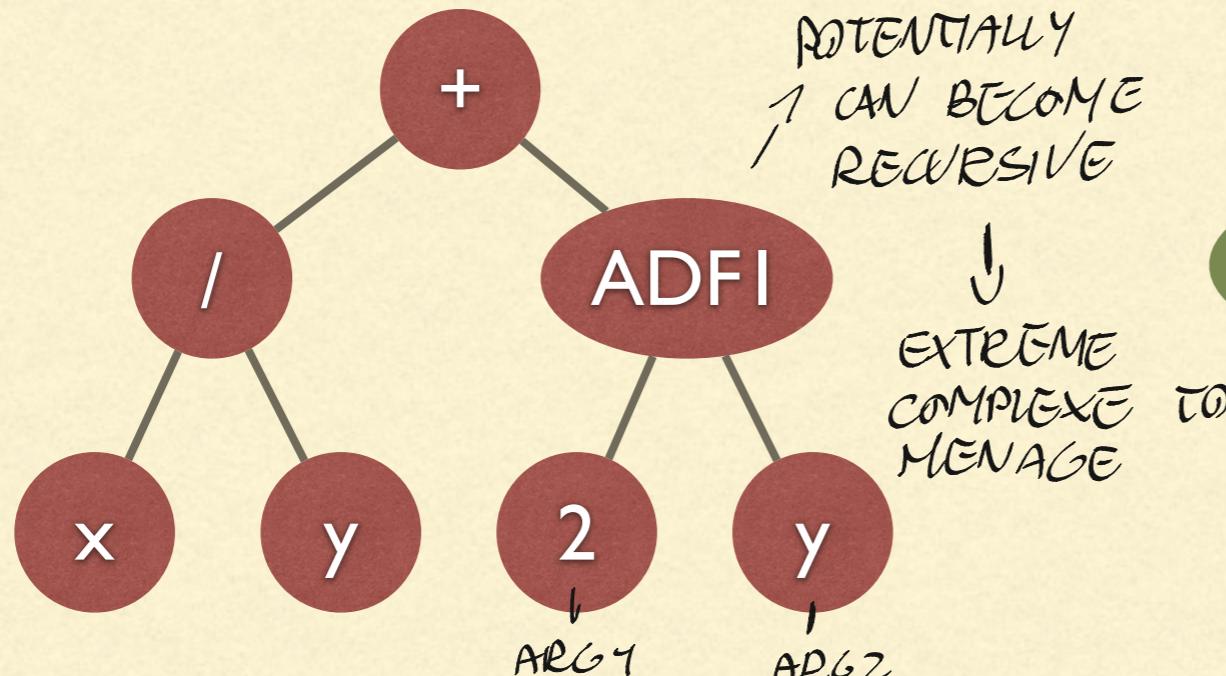
MOTIVATIONS

- One of the most useful (and older) ideas in programming is the use of *subroutines*
→ COPY PASTE SOME PREMADE CODE
- However, instead of calling a function k times, a GP individual must evolve the same code k times
- Since this is unlikely we can add “subroutines” and calling subroutines to GP
WHY? BECAUSE THEY ARE GOING TO BE GENERATED AUTOMATICALLY
- In GP those are called Automatically Defined Functions or ADF

AUTOMATICALLY DEFINED FUNCTIONS

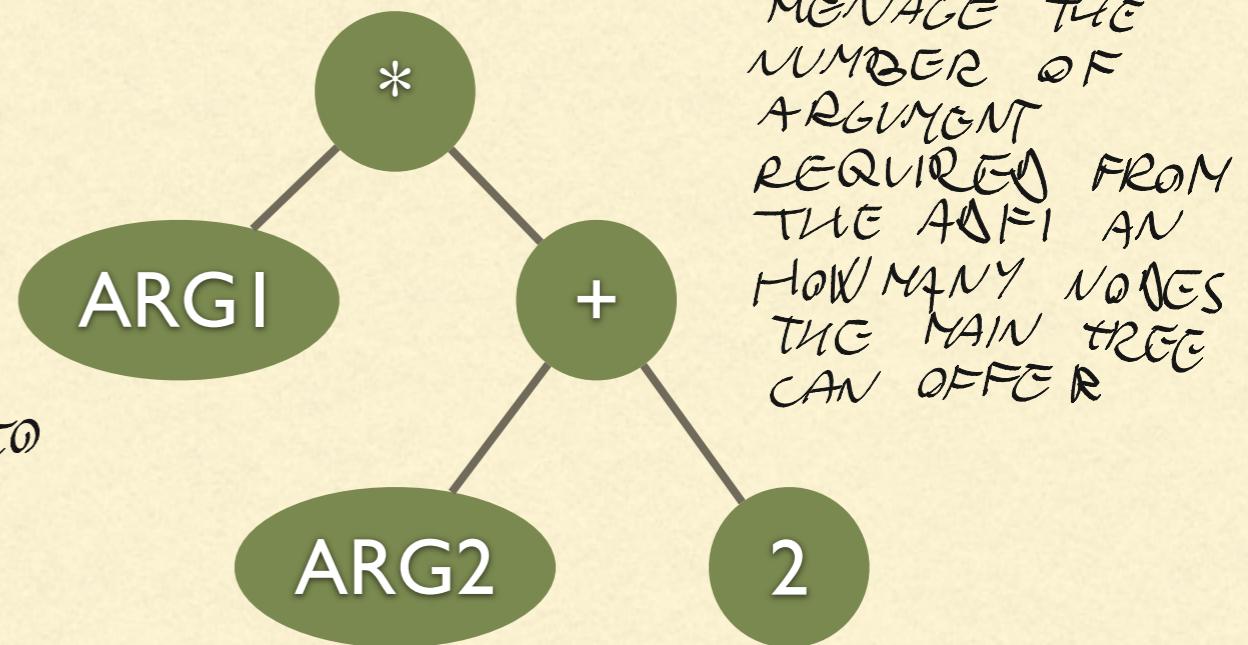
We select how many ADF we need and for each one the number of arguments that it accepts

Main Tree



ADFs are used as functionals

ADFI tree



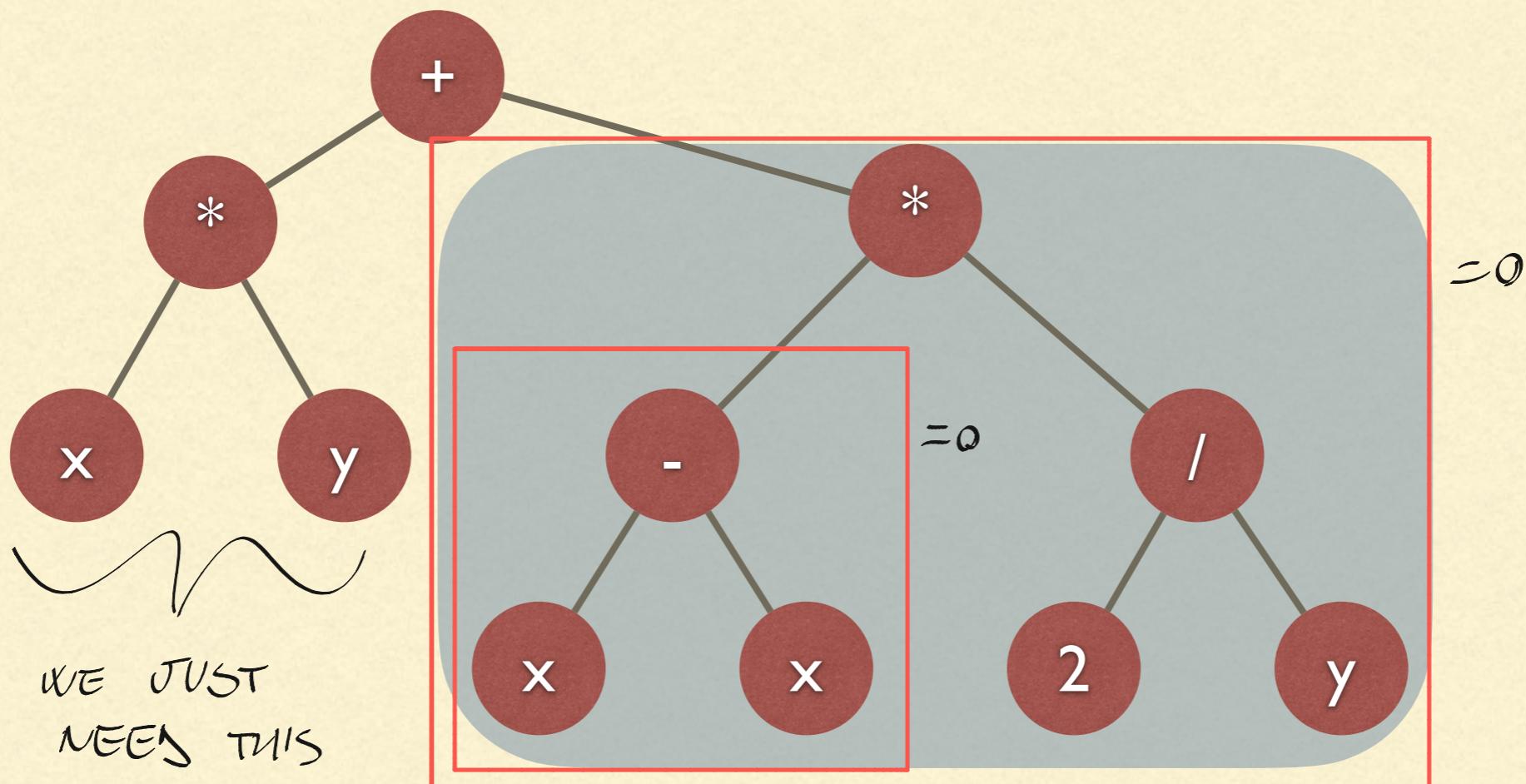
But each of them is an entire tree

AUTOMATICALLY DEFINED FUNCTIONS

- The individuals are now vectors of trees (forests)
- We can still apply the usual mutation and crossover to the elements of the vector
 - ↗ UP TO A REASONABLE LIMITS
- We can have as many ADF as we want...
- ...and ADF can call each other (nested subroutines/functions calls)
- recursion might be problematic (evolving the base case might not be easy)

BLOAT

SPOT THE PROBLEM



This entire region
is non-coding

WHAT IS BLOAT?

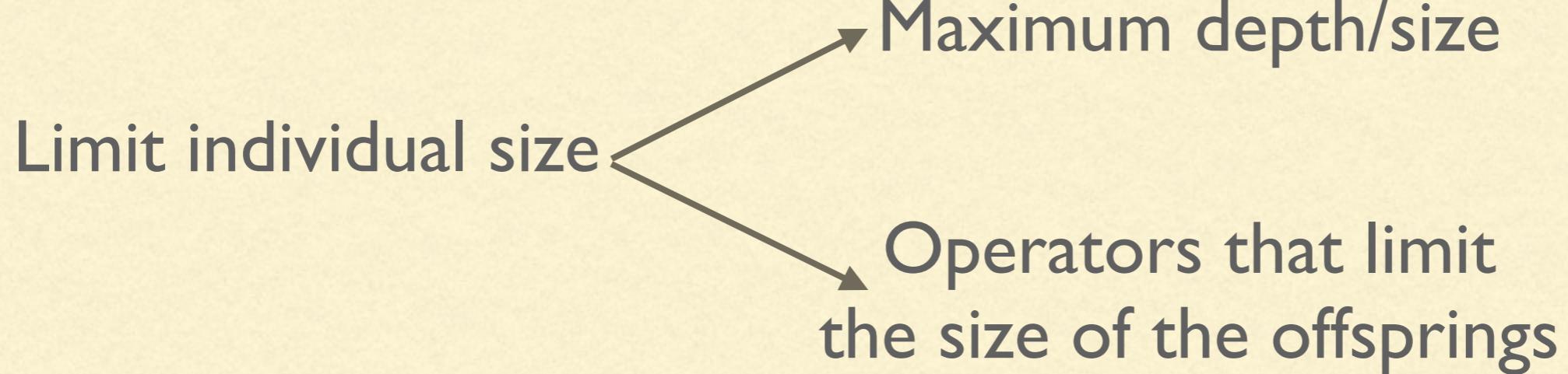
IS A PARTICULAR FORM OF OVER FITTING

- Not a simple definition, lets focus on what can be observed:
- Large non-coding regions: many operations on the tree do not change the function that it represents
- Increase in the size of the trees without a noticeable increase in the fitness
- As a consequence, fitness evaluation is slower, slowing down the entire GP process

WE WANT THE SMALLEST TREE WITH
THE BEST RESULT

HOW CAN WE AVOID OR IDENTIFY THE PROBLEM?

ATTACKING BLOAT

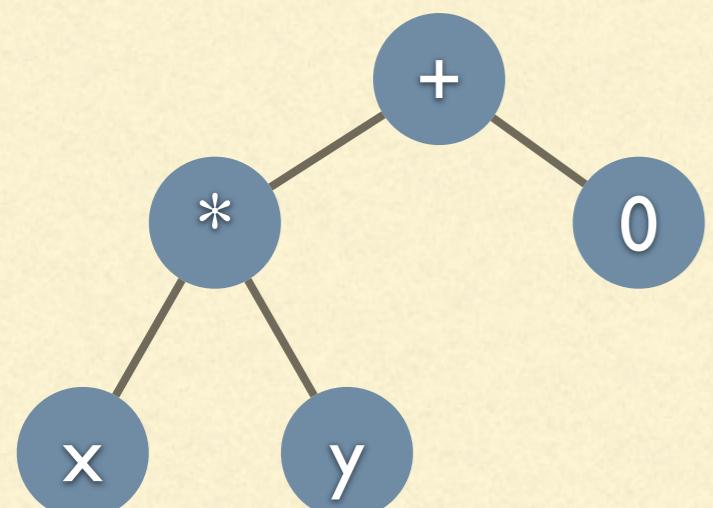


Remove non-coding regions

Punish the individuals that are too big

LINEAR PARSIMONY PRESSURE

Real fitness



5

Adjusted fitness

$$0.9 \times 5 - 0.1 \times 5 = 4$$

5

$$0.9 \times 5 - 0.1 \times 11 = 3.4$$

Real fitness

$$\alpha f - (1 - \alpha)s$$

Size

$$\alpha \in [0,1]$$

