Lab Page Tables

Learn how xv6 manages process' virtual memory space.

Environment Preparation

Run the following commands:

```
$ git fetch
$ git checkout pgtbl
$ make clean
```

Lab Tasks

speed up system calls

Some operating systems (e.g., Linux) speed up certain system calls by sharing data in a read-only region between userspace and the kernel. This eliminates the need for kernel crossings when performing these system calls. To help you learn how to insert mappings into a page table, your first task is to implement this optimization for the <code>getpid()</code> system call in xv6.

When each process is created, map one read-only page at USYSCALL (a virtual address defined in memlayout.h). At the start of this page, store a struct usyscall (also defined in memlayout.h), and initialize it to store the PID of the current process. For this lab, ugetpid() has been provided on the userspace side and will automatically use the USYSCALL mapping. You will receive full credit for this part of the lab if the ugetpid test case passes when running pgtbltest.

Some hints:

- You can perform the mapping in proc_pagetable() in kernel/proc.c.
- Choose permission bits that allow userspace to only read the page.
- You may find that mappages() is a useful utility.
- Don't forget to allocate and initialize the page in allocproc().
- Make sure to free the page in freeproc().

print a page table

To help you visualize RISC-V page tables, and perhaps to aid future debugging, your second task is to write a function that prints the contents of a page table.

Define a function called vmprint(). It should take a pagetable_t argument, and print that
pagetable in the format described below. Insert if(p->pid==1) vmprint(p->pagetable) in exec.c
just before the return argc, to print the first process's page table. You receive full credit for this part
of the lab if you pass the pte printout test of make grade.

Now when you start xv6 it should print output like this, describing the page table of the first process at the point when it has just finished exec() ing init:

```
page table 0x000000001fd9c01 pa 0x0000000087f67000
...0: pte 0x0000000021fd9801 pa 0x0000000087f66000
...0: pte 0x0000000021fda01b pa 0x0000000087f68000
....0: pte 0x0000000021fda01b pa 0x0000000087f68000
....1: pte 0x0000000021fd9417 pa 0x0000000087f65000
....2: pte 0x0000000021fd9007 pa 0x0000000087f64000
.....3: pte 0x0000000021fd8c17 pa 0x0000000087f63000
...255: pte 0x0000000021fda801 pa 0x0000000087f69000
....511: pte 0x0000000021fda401 pa 0x0000000087f73000
....509: pte 0x0000000021fdcc13 pa 0x0000000087f74000
....510: pte 0x0000000021fdd007 pa 0x0000000087f74000
....511: pte 0x0000000021fdd007 pa 0x0000000087f74000
....511: pte 0x0000000021fdd007 pa 0x0000000080007000
init: starting sh
```

The first line displays the argument to <code>vmprint</code>. After that there is a line for each PTE, including PTEs that refer to page-table pages deeper in the tree. Each PTE line is indented by a number of ".." that indicates its depth in the tree. Each PTE line shows the PTE index in its page-table page, the pte bits, and the physical address extracted from the PTE. Don't print PTEs that are not valid. In the above example, the top-level page-table page has mappings for entries 0 and 255. The next level down for entry 0 has only index 0 mapped, and the bottom-level for that index 0 has entries 0, 1, and 2 mapped.

Your code might emit different physical addresses than those shown above. The number of entries and the virtual addresses should be the same.

Some hints:

- You can put vmprint() in kernel/vm.c.
- Use the macros at the end of the file kernel/riscv.h.
- The function freewalk may be inspirational.
- Define the prototype for vmprint in kernel/defs.h so that you can call it from exec.c.
- Use %p in your printf calls to print out full 64-bit hex PTEs and addresses as shown in the example.

detect which pages have been accessed

Some garbage collectors (a form of automatic memory management) can benefit from information about which pages have been accessed (read or write). In this part of the lab, you will add a new feature to xv6 that detects and reports this information to userspace by inspecting the access bits in the RISC-V page table. The RISC-V hardware page walker marks these bits in the PTE whenever it resolves a TLB miss.

Your job is to implement pgaccess(), a system call that reports which pages have been accessed. The system call takes three arguments. First, it takes the starting virtual address of the first user page to check. Second, it takes the number of pages to check. Finally, it takes a user address to a buffer to store the results into a bitmask (a datastructure that uses one bit per page and where the first page corresponds to the least significant bit). You will receive full credit for this part of the lab if the pgaccess test case passes when running pgtbltest.

Some hints:

- Read pgaccess_test() in user/pgtlbtest.c to see how pgaccess is used.
- Start by implementing sys_pgaccess() in kernel/sysproc.c.

- You'll need to parse arguments using argaddr() and argint().
- For the output bitmask, it's easier to store a temporary buffer in the kernel and copy it to the user (via copyout()) after filling it with the right bits.
- It's okay to set an upper limit on the number of pages that can be scanned.
- walk() in kernel/vm.c is very useful for finding the right PTEs.
- You'll need to define PTE_A, the access bit, in kernel/riscv.h. Consult the RISC-V privileged architecture manual to determine its value.
- Be sure to clear PTE_A after checking if it is set. Otherwise, it won't be possible to determine if the page was accessed since the last time pgaccess() was called (i.e., the bit will be set forever).
- vmprint() may come in handy to debug page tables.

Submit the Lab

Report

Put your report about this lab in a new file report.txt. You can write down the errors you met when doing the lab, or some techniques you found that helps you to finish the lab better. You can use both **English** and **Chinese**.

Don't forget to git add and git commit the file.

Time spent

Create a new file, time.txt, and put in it a single integer, the number of hours you spent on the lab. Don't forget to git add and git commit the file.

Check

- Please run make grade to ensure that your code passes all of the tests
- Commit any modified source code before running make tarball

Submit

Run make tarball, you should see a new file in the root directory called [lab-pgtbl-handin.tar.gz]. Upload this file to the corresponding section in course <u>canvas page</u>.