

Full Name: \_\_\_\_\_

*“On my honor as a University of Colorado at Boulder student I have neither given nor received unauthorized assistance on this work.”*

## CSCI 2400, Fall 2017

### Final Exam

#### Instructions:

- Make sure that your exam is not missing any sheets, then write your full name on the front.
- Write your answers in the space provided below the problem. Show your work. If you make a mess, clearly indicate your final answer. Feel free to use the back of pages, but indicate that you have done so.
- This exam is **CLOSED BOOK** and you can use a *single page* of notes along with our reference sheets. You can not use a calculator. Your computer can only be used to upload answers to the moodle.

Problem	Possible	Score
1	12	
2	12	
3	28	
4	20	
5	25	
6	10	
<b>Total</b>	<b>107</b>	

#### 1. [ 12 Points ]

```
#include <stdlib.h>
#include <stdio.h>

void main() {

    if(fork() && fork()){
        fork();
    }

    if(fork() || fork()){
        fork();
    }
    printf("Hello World\n");

}
```

How many "Hello World" output lines does this program print?

(Hint: Notice there is "short-circuiting" in C language. For example, for &&, "0 && WhoCares" will always be false so C just skips "WhoCares". Similar to "1 || WhoCares".)

2. [ 12 Points ] Solve the following questions about signal handlers:

```
#include <stdlib.h>
#include <signal.h>

int count1 = 0;
int count2 = 0;

void handler(int sig) {
    count1++;
    printf("count1 = %d\n", count1);
    printf("count2 = %d\n", count2);
}

void handler2(int sig) {
    count2++;
    printf("count1 = %d\n", count1);
    printf("count2 = %d\n", count2);
}

int main() {

    signal(SIGCHLD, handler);
    int pid = fork();

    if(pid == 0){
        signal(SIGCHLD, handler2);
        exit(0);
    }

    if(!fork()) {
        count2++;
        exit(0);
    }

    signal(SIGINT, handler);
    signal(SIGTSTP, handler2);

    while (1)
    {
        sleep(1);
    }

    exit(0);
}
```

(a) How many SIGCHLD signals get generated as this program executes?

**for parts (b) and (c) assume that child process exits before the parent sleeps in the while loop.**

(b) For the parent process that started the main() function running in while loop, If the user input ctrl+c , what are the printed values for the count1 and count2 variables?

(c) For the parent process that started the main() function running in while loop, If the user input ctrl+z , what are the printed values for the count1 and count2 variables? Assuming **independent** run of the program from part (b).

3. [ 28 Points ] The following problem concerns the way virtual addresses are translated into physical addresses.

- The memory is byte addressable.
- Memory accesses are to **1-byte words** (not 4-byte words).
- The TLB is 4-way set associative with 8 total entries.
- The L1 Cache is 2-way set associative, with a 4-byte block size and 64 total bytes.
- Virtual addresses are 13 bits wide.
- Physical addresses are 11 bits wide.
- The page size is 32 bytes.

In the following tables, **all numbers are given in hexadecimal**. The contents of the TLB and a portion of the page tables are as follows:

TLB				Page Table			Cache			
Index	Tag	PPN	Valid	VPN	PPN	Present	Index	Valid	Tag	Data
0	1a	06	1	031	000	1	0	1	1C	6021130E
	–	–	0	0a2	00c	1		1	00	DCAEB820
	15	02	1	032	009	1	1	0	12	1DFE0C46
	–	–	0	051	004	1		0	0B	29E5DBF8
1	14	0f	1	03f	00d	1	2	1	1F	DFFBCC85
	–	–	0	02f	00a	1		1	02	CB570940
	0a	0c	1	00e	008	1	3	1	08	57A84A44
	07	04	1	02c	003	1		1	3C	8E85761F
				021	00f	1	4	1	0D	DF2C1CE2
				012	00b	1		1	07	BE10CEA4
				01a	00e	1	5	1	04	579C4AB6
				03d	002	1		1	0C	A11D81A1
				006	001	1	6	1	13	B250AE92
				034	006	1		1	15	7751E21A
				017	005	1	7	0	0C	6AA3E19A
				003	007	1		1	09	6AC09E41

(a) [ 6 Points ] Calculate the number of bits for the following elements:

*VPO* The virtual page offset \_\_\_\_\_ *TLBI* The TLB index \_\_\_\_\_  
*VPN* The virtual page number \_\_\_\_\_ *TLBT* The TLB tag \_\_\_\_\_  
*PPO* The physical page offset \_\_\_\_\_ *PPN* The physical page number \_\_\_\_\_

(b) [ 22 Points ] (1 points each) For the given virtual addresses, indicate the TLB entry accessed and the physical address. **Indicate whether the TLB misses and whether the entry is or is not in the page table.** If the physical page number and address can not be determined, write “N/A”. Then if a physical address exists indicate the cache translation parts, if its a cache hit, and a value if applicable. If any part can’t be determined just write “N/A”.

**Virtual address:** 0x0549

## (i) Address translation

Parameter	Value
VPN	0x
TLB Index	0x
TLB Tag	0x
TLB Hit? (Y/N)	
In Page Table? (Y/N)	
PPN	0x

## (ii) Cache Translation

Parameter	Value
Cache Offset	0x
Cache Index	0x
Cache Tag	0x
Cache Hit? (Y/N)	
Byte Value	0x

**Virtual address: 0x0244**

## (i) Address translation

Parameter	Value
VPN	0x
TLB Index	0x
TLB Tag	0x
TLB Hit? (Y/N)	
In Page Table? (Y/N)	
PPN	0x

## (ii) Cache Translation

Parameter	Value
Cache Offset	0x
Cache Index	0x
Cache Tag	0x
Cache Hit? (Y/N)	
Byte Value	0x

## 4. [ 20 Points ]

Suppose our memory allocator uses an implicit free list with both header and footer. Assume a word size of eight bytes, and that all blocks are aligned to addresses divisible by eight. You should assume that the addresses you see span the entire heap, and that a block is marked as allocated by setting the least significant bit of the header and footer to 1. Similarly, a block is marked as free by setting the least significant bit of the header and footer to 0. Note that each row in the heap pictured below represents one eight-byte word. **Assume a best-fit placement policy.**

Address	Value
FF00	00 ... 00 18
FF08	?? . . . ??
FF10	00 ... 00 18
FF18	00 ... 00 29
FF20	?? . . . ??
FF28	?? . . . ??
FF30	?? . . . ??
FF38	00 ... 00 29
FF40	00 ... 00 19
FF48	?? . . . ??
FF50	00 ... 00 19
FF58	00 ... 00 21
FF60	?? . . . ??
FF68	?? . . . ??
FF70	00 ... 00 21
FF78	00 ... 00 30
FF80	?? . . . ??
FF88	?? . . . ??
FF90	?? . . . ??
FF98	?? . . . ??
FFA0	00 ... 00 30
FFA8	00 ... 00 21
FFB0	?? . . . ??
FFB8	?? . . . ??
FFC0	00 ... 00 21
FFC8	00 ... 00 39
FFD0	?? . . . ??
FFD8	?? . . . ??
FFE0	?? . . . ??
FFE8	?? . . . ??
FFF0	?? . . . ??
FFF8	00 ... 00 39

**Unless clearly marked otherwise, assume all numbers are in hexadecimal!**

Suppose that, after some sequence of malloc's and free's, the state of the heap is as you see it on the left. Then, assume that the following calls to malloc and free are made:

```

10: void* p1 = malloc(0x08);
11: free(0xff48);
12: free(0xff60);
13: free(0xffd0);
14: void* p2 = malloc(0x30);

```

And answer the following:

- [ 5 Points ] How much space on the heap does the smallest valid block size take up, in bytes?
- [ 5 Points ] What is p1?
- [ 5 Points ] Which of the five lines above will cause the allocator to perform a 'coalesce' operation? (ie, 10, 11, 12, 13, or 14?)
- [ 5 Points ] What is p2?

**Unless clearly marked otherwise, assume all numbers are in hexadecimal!**

## 5. [ 25 Points ]

Answer these questions on linking

- (a) [ 15 Points ] For the following code, identify the symbols listed in the symbol table of the ELF relocatable object files (.o), whether that symbol is defined or undefined, and if defined, then in which section of the corresponding ELF file that the symbol would be defined.

**main.c**

```
extern void func();
int p=7;
int q;
int main()
{
    int m=500000;
    q=sqrt(m);
    func(q);
    return 0;
}
```

**func.c**

```
int n=10;
int temp;
int func(int x)
{
    if(x>100)
        temp=x;
    else
        temp=-x;
    return temp;
}
```

**main.o**

Symbol Name	Defined/undefined	Section
	undefined	.data .bss

**func.o**

Symbol Name	Defined/undefined	Section
n		.bss

- (b) [ 5 Points ] For the code in Question (a), the sizes of the .text and .data sections of the .o relocatable object files are listed below. The two object files above are then linked together with the command line `ld -o p main.o func.o`. Assume the object files are combined similar to the order shown in the lecture slides and the starting address of the .text section of the unified executable object file starts at 0x8048501. What is the relocated address of p?

File+Section	Size (Byte)
main.o's .text	32
main.o's .data	8
func.o's .text	58
func.o's .data	4

- (c) [ 5 Points ] When the two .o files above are linked together with the command line `ld -o p main.o func.o`, the virtual addresses of the merged and relocated various subsections follow what kind of orderings, from lowest to highest addresses (circle)? (can circle more than one correct answer):

- (a) `.data(func.o) < .text(main.o) < .text(func.o)`      (b) `.text(func.o) < .data(func.o) < .data(main.o)`  
(c) `.text(main.o) < .data(main.o) < .data(func.o)`      (d) `.data(func.o) < .data(main.o) < .bss(func.o)`  
(e) `.data(func.o) < .bss(main.o) < .bss(func.o)`      (f) None of the above

6. [ 10 Points ] For each of the following, answer True or False:

- (a) \_\_\_\_\_ For a binary number, left shift by 1 corresponds to division by 2.  
(b) \_\_\_\_\_ The stack pointer in 64-bit x86 systems is stored in the `%rsp` register.  
(c) \_\_\_\_\_ There is no difference between binary encoding of integers and floating point.  
(d) \_\_\_\_\_ Each Y86 instruction can be divided into 6 stages of execution: Fetch, Decode, Execute, Memory, Write, Update PC  
(e) \_\_\_\_\_ `for(i=0;i≤100;i++)sum+=a[0];` For the given code, cache helps on spatial locality.