

DMTH237 Discrete Mathematics II — Assignment 3

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1. (a) We have $26 \times 2 + 10 + 6 = 68$ available characters. Hence, there are $68^8 + 68^9 + 68^{10} + 68^{11} + 68^{12} = 9,920,671,339,261,325,541,376$ different passwords.
- (b) With only letters and numbers we have $26 \times 2 + 10 = 62$ characters so that there are $62^8 + 62^9 + 62^{10} + 62^{11} + 62^{12}$ passwords which *do not* contain at least one special character. Subtracting this from the number of passwords found in part 1a we have

$$\begin{aligned} & (68^8 + 68^9 + 68^{10} + 68^{11} + 68^{12}) - (62^8 + 62^9 + 62^{10} + 62^{11} + 62^{12}) \\ &= 6,641,514,961,387,068,437,760 \end{aligned}$$

which is the number of passwords not containing at least one special character.

- (c) Checking each password would take 314.6 billion years (short scale)—about 23 times the age of the universe!
2. To solve this problem we note the following.
- Without restricting the value of each digit we have 10^9 possible strings.
 - There are 9^9 strings which do not contain the digit 1 but may contain 3 or 7.
 - There are $9^9 - 8^9$ strings which do not contain the digit 3, contain at least a single 1 and may contain a 7. This is a problem of the same kind as in part 1b.
 - There are $9^9 - 2(8^9 - 7^9) - 7^9$ strings which do not contain the digit 7 but contain at least a single 1 and at least a single 3. This is because there are $8^9 - 7^9$ strings which do not contain the digits 1 and 7 but contain at least a single 3; there are $8^9 - 7^9$ strings which do not contain the digits 3 and 7 but contain at least a single 1 and; there are 7^9 strings which do not contain the digits 1, 3 or 7.

From this we see that the number of strings which have 1, 3 and 7 appearing at least once is given by

$$\begin{aligned} & (10^9) - (9^9) - (9^9 - 8^9) - (9^9 - 2(8^9 - 7^9) - 7^9) \\ &= 200,038,110. \end{aligned}$$

3. (a) A teacher has n students and asks them all m questions. Any student who gets a question correct gets a piece of candy. (It is possible that no student is correct, so that no one gets candy.)
- Now, if r questions are answered correctly there are

$$\binom{n+r-1}{r}$$

ways the candy might be distributed among the students. By the rule of sum,

$$\sum_{r=0}^m \binom{n+r-1}{r}$$

is the number of different distributions of candy for a varying number of correct answers.

The RHS,

$$\binom{n+m}{m},$$

is then the number of ways to distribute the students among the candy, including the case where no students answer correctly.

(b) First we show that the base case (where $m = 0$) holds.

$$\begin{aligned} \sum_{r=0}^m \binom{n+r-1}{r} &= \binom{n+m}{m} \\ \sum_{r=0}^0 \binom{n+r-1}{r} &= \binom{n+0}{0} \\ \binom{n+0-1}{0} &= \binom{n+0}{0} \\ 1 &= 1 \end{aligned} \tag{1}$$

Now, in order to undertake the inductive step we will use

$$\sum_{j=0}^n \binom{j}{k} = \binom{n+1}{k+1}, \quad \text{for every } n, k \in \mathbb{N} \tag{2}$$

which, for completeness, we prove as follows:

$$\begin{aligned} \sum_{j=0}^n \binom{j}{k} &= \sum_{j=0}^n \left[\binom{j+1}{k+1} - \binom{j}{k+1} \right] \\ &= \left[\binom{1}{k+1} - \binom{0}{k+1} \right] + \left[\binom{2}{k+1} - \binom{1}{k+1} \right] + \\ &\quad \left[\binom{3}{k+1} - \binom{2}{k+1} \right] + \cdots + \left[\binom{n+1}{k+1} - \binom{n}{k+1} \right] \\ &= -\binom{0}{k+1} + \left[\binom{1}{k+1} - \binom{1}{k+1} \right] + \left[\binom{2}{k+1} - \binom{2}{k+1} \right] \\ &\quad + \cdots + \left[\binom{n}{k+1} - \binom{n}{k+1} \right] + \binom{n+1}{k+1} \\ &= \binom{n+1}{k+1}. \end{aligned}$$

Completing the inductive step, we must have

$$\begin{aligned}
\sum_{r=0}^{m+1} \binom{n+r-1}{r} &= \binom{n+m+1}{m+1} && \text{(by the induction hypothesis)} \\
\sum_{r=0}^{m+1} \binom{n+r-1}{n-1} &= \binom{n+m+1}{m+1} && (n \text{ choose } k \text{ equals } n \text{ choose } n-k) \\
\sum_{s=n-1}^{n+m} \binom{s}{n-1} &= \binom{n+m+1}{m+1} && (\text{letting } s = n+r-1) \\
\binom{n+m+1}{n-1+1} &= \binom{n+m+1}{m+1} && (\text{by equation (2)}) \\
\binom{n+m+1}{m+1} &= \binom{n+m+1}{m+1}, && (n \text{ choose } k \text{ equals } n \text{ choose } n-k)
\end{aligned}$$

which was to be proved.

4. Two integers a_i, a_{i+1} are consecutive if $a_{i+1} > a_i + 1$. Hence we have

$$1 \leq a_1 < a_2 - 1 < a_3 - 2 < a_4 - 3 < a_5 - 4 < a_6 - 5 < a_7 - 6 \leq 50 - 6.$$

Writing $b_i = a_i - (i - 1)$, we have

$$1 \leq b_1 < b_2 < \dots < b_7 \leq 44.$$

Now there clearly exists a bijection between every a_i, b_i . Thus, the number of ways to choose $a_i, 1 \leq i \leq 7$ is equal to the number of ways to choose $b_i, 1 \leq i \leq 7$.

The number of ways to choose the seven different numbers is therefore

$$\binom{44+1}{7} = 45,379,620.$$

5. The left and right machines accept the languages \mathcal{L} and \mathcal{M} . Their respective state tables are show below.

		0	1	
→ A		B	C	
B		D	E	
C		F	Z	
D		Z	Z	*
E		G	G	*
F		Z	Z	*
G		Z	A	
Z		Z	Z	

		0	1	
→ R		S	R	*
S		T	S	
T		R	T	

The state tables for the machines which accept the languages $\overline{\mathcal{L}}$ and $\overline{\mathcal{M}}$ are shown below.

		0	1	
→ A		B	C	*
B		D	E	*
C		F	Z	*
D		Z	Z	
E		G	G	
F		Z	Z	
G		Z	A	*
Z		Z	Z	*

		0	1	
→ R		S	R	*
S		T	S	*
T		R	T	*

Combining the above tables to find a finite state acceptor for the language $\overline{\mathcal{L}} + \overline{\mathcal{M}}$ yields the mutually equivalent machines shown below.

	0	1	λ			0	1			0	1	
$\rightarrow A_0$			AR		$\rightarrow AR$	BS	CR	*	$\rightarrow 0$	1	2	*
A	B	C		*	BS	CR	ES	*	1	2	3	*
B	D	E		*	CR	CR	CR	*	2	2	2	*
C	F	Z		*	ES	GT	GS	*	3	4	5	*
D	Z	Z			GT	CR	AT	*	4	2	6	*
E	G	G			GS	CR	AS	*	5	2	7	*
F	Z	Z			AT	BR	CT	*	6	8	9	*
G	Z	A		*	AS	BT	CR	*	7	10	2	*
Z	Z	Z		*	BR	CR	ER	*	8	2	11	*
R	S	R			CT	FR	CR	*	9	12	2	*
S	T	S		*	BT	FR	ET	*	10	12	13	*
T	R	T		*	ER	GS	GR		11	5	14	
					FR	CR	CR		12	2	2	
					ET	GR	GT	*	13	14	4	*
					GR	CR	AR	*	14	2	0	*

Hence, we have the machine accepting strings in the language $\mathcal{L} \cap \mathcal{M}$, shown below.

	0	1	
$\rightarrow 0$	1	2	
1	2	3	
2	2	2	
3	4	5	
4	2	6	
5	2	7	
6	8	9	
7	10	2	
8	2	11	
9	12	2	
10	12	13	
11	5	14	*
12	2	2	*
13	14	4	
14	2	0	