Exercise 1: punti 18

Simulate the Counting Sort algorithm (stable version) on sorting the array 5 5 6 1 2 1 4 3 5.

Exercise 2: punti 7

Consider a binary tree T with n nodes. Each node u has a color $u.\mathsf{color}$ which could be either white or black.

Design and analyze an efficient algorithm to compute the number of nodes u such that the subtree rooted at u contains more white nodes than black nodes.

Exercise 3: punti 5

Given an undirected graph $G = (V, E)$, the measure $d(G)$ equals the largest smallest distance between to pairs of nodes in the graph. More preciselly, let $m(u, v)$ be the length of the shortest path from u to v, $d(G) = max_{(u,v)\in V^2} m(u, v)$.

If the graph is not connected, $d(G) = +\infty$.

Design an algorithm to compute $d(G)$.

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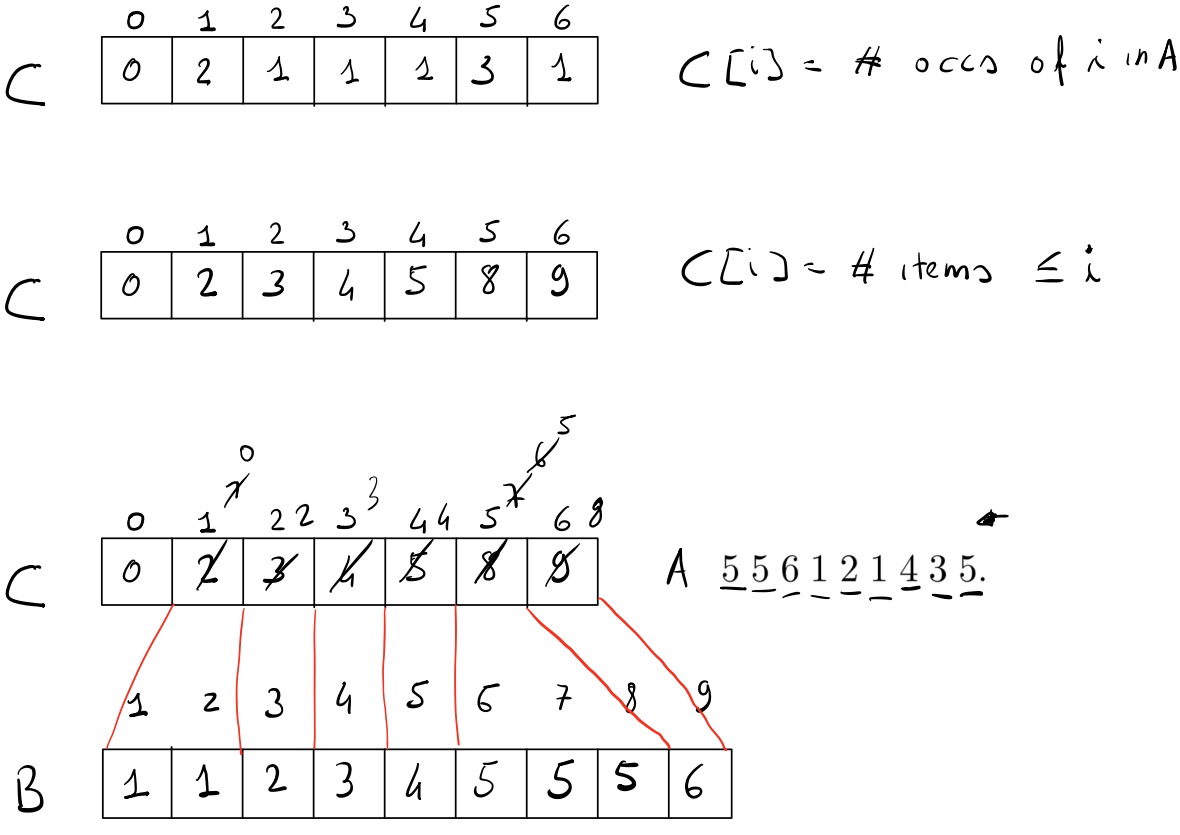
Simulate the Counting Sort algorithm (stable version) on sorting the array 5 5 6 1 2 1 4 3 5.



Exercise 2: punti 7

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Design and analyze an efficient algorithm to compute the number of nodes u such that the subtree rooted at u contains more white nodes than black nodes.

```
                                    // returns   # of node satisfying the property
                                                 # of white nodes
                                                 # of black nodes
count(x)
    if x == NIL : return 0, 0, 0
    rl, wl, bl = count(x.left)
    rr, wr, br = count(x.right)
    rx = rl + rr
    wx = wl + wr
    bx = bl + br
    if x.color == "white" : wx += 1
    else                  : bx += 1
    if wx > bx : rx += 1
    return rx, wx, bx
```

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If the graph is not connected, $d(G) = +\infty$.

Design an algorithm to compute $d(G)$.

$\max_{u \in V} \max_{v \in V} m(u, v)$

```
diameter(G)
    d = -∞
    for u ∈ V
        BFS(u)
        for v ∈ V
            if v.color == white
                return +∞
            d = max(d, v.d)
    return d
```

complexity: $\Theta(|V| \cdot (|V| + |E|))$ time