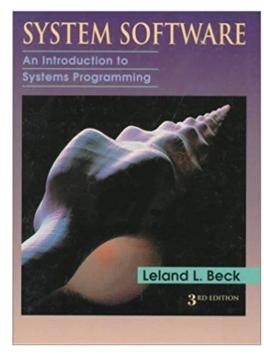
# CSCI 30000 Systems Programming Course Project



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MW 12:30-1:30 SL228

#### Course Project: Introduction

- These slides will serve as a basic guide to help you get started on the course project
- This project will take a considerable amount of time to complete successfully
  - If you have not already, please quickly make a plan with your group detailing how you will complete this project
  - There is no way you will be able to successfully complete this project if you put if off until the end of the semester
- This project will enhance both your understanding of System Software and your skills as a programmer

#### Course Project: Goal

- Design and Implementation Project: implement an assembler using Java, C++, Perl, Tcl/Tk, others (any one is OK)
  - I would suggest Java or Python, but you may choose whatever language reflects your team members' strengths
  - Irrespective of choice, you will be responsible for clearly outlining the details
    of how to successfully compile and run your code
- You will implement an assembler which can process SIC/XE assembly programs and generate the corresponding object codes and object programs

#### Course Project: Overview

- You will implement an assembler which can process SIC/XE assembly programs and generate the corresponding object codes and object programs
  - What does this actually entail?
  - We will see some specific details in a few slides, but the overall challenge is to...
- In homework one we saw how use an object code to calculate the corresponding SIC/XE instruction's target addresses, addressing modes, opcode and instruction format (object code -> instruction details)
- In this project you are tasked with discerning a SIC/XE program's instructions' information and working to create the corresponding object codes (instruction details -> object code)

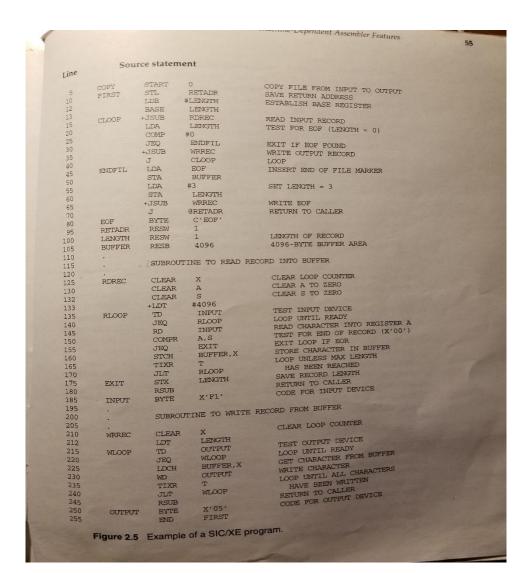
#### Course Project: First Steps

- You need to start by reading the first two chapters of the textbook
- Your ability to contribute to your team's effort will be extremely limited if you have not read and fully understand the first two chapters of the textbook
- These first two chapters outline many of the operations, considerations, data structures and implementation details that this project will require
- I will give a few details here, but you will need to review the text to clearly understand what you must do

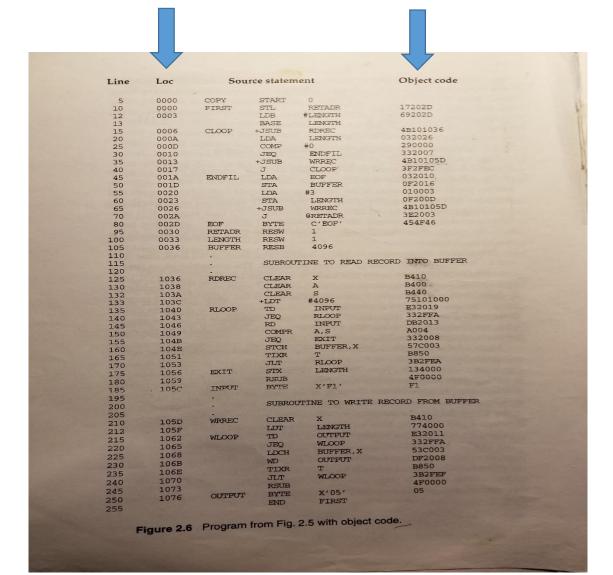
- Let us take a look at some assembler implementation details
- In an email you were given six sample SIC/XE assembly programs
  - 1. basic.txt
  - 2. functions.txt
  - 3. literals.txt
  - 4. program\_blocks.txt
  - 5. control\_sections.txt
  - 6. macros.txt
- You need to implement an assembler which can successfully process all these files
- Let us start by taking a closer look at functions.txt

COPI	C		TARREST OF THE STATE OF THE STA
COPY	START STL	0	.copy the file form INPUT to OUTPUT
FIRST	LDB	RETADR #LENGTH	.Save Return address .establish base register
	BASE	#LENGIH LENGTH	.establish base register
CLOOP	+JSUB	RDREC	.Read INPUT Record
CLOOP	LDA	LENGTH	.Test for EOF (Length = 0)
	COMP	#0	.lest for Eor (Length - 0)
	JEO	ENDFIL	.Exit if EOF found
	+JSUB	WRREC	.Write OUTPUT Record
	J	CLOOP	.Loop
ENDFIL	LDA	EOF	.Insert End Of File Marker
LINDIIL	STA	BUFFER	.Inscis End of Tile narker
	LDA	#3	.Set LENGTH = 3
	STA	LENGTH	.Set LENGIN - S
	+JSUB	WRREC	.write EOF
	J	@RETADR	.Return to Caller
EOF	BYTE	C'EOF'	incouri do durrer
RETADR	RESW	1	
LENGTH	RESW	1	.LENGTH of Record
BUFFER	RESB	4096	.4096-BYTE Buffer area
	SUBROUTINE	TO READ RECORD	INTO BUFFER
1			
RDREC	CLEAR	x	.Clear Loop counter
	CLEAR	A	.Clear A to Zero
	CLEAR	s	.Clear S to Zero
	+LDT	#4096	
RLOOP	TD	INPUT	.Test INPUT device
	JEQ	RLOOP	.LOOP until ready
	RD	INPUT	.READ character into Register A
	COMPR	A,S	.TEST for End Of Record (X'00')
	JEQ	EXIT	.EXIT loop if EOR
	STCH	BUFFER, X	.STORE character in BUFFER
	TIXR	T	.LOOP unless max Length has been reached
	JLT	RLOOP	
EXIT	STX	LENGTH	.SAVE record length
	RSUB		.Return to caller
INPUT	BYTE	X'Fl'	.CODE for Input device
-	SUBROUTINE	TO WRITE RECORD	FROM BUFFER
-			
WRREC	CLEAR	x	.CLEAR loop counter
	LDT	LENGTH	
WLOOP	TD	OUTPUT	.Test OUTPUT Device
	JEQ	WLOOP	.Loop Until Ready
	LDCH	BUFFER, X	.GET character form BUFFER
	WD	OUTPUT	.WRITE character
	TIXR	T	.LOOP until all characters have been written
	JLT	WLOOP	
	RSUB		.Return to caller
OUTPUT	BYTE	X'05'	.CODE for OUPUT device
	END	FIRST	

This SIX/XE program should look familiar, it is from you book!



• functions.txt (pg. 55)



- functions.txt solution (pg. 58)
- Loc and Object Code columns now present!

Figure 2.8 Object program corresponding to Fig. 2.6.

 functions.txt object program solution (pg. 65)

- Cool! We can already see the answers ©
- For many of the sample programs the solutions are given in the book
- As you implement your assembler and add functionalities to successfully process each of the six sample programs, you can reference these solutions to see if you are on the right track

- This still doesn't tell us much about how to actually generate the displayed object code or location columns though ☺
- Let us try to figure out

- Our given functions.txt file has four tab delimited columns
  - Symbol Column, Operation Column, Operand Column and Comments Column
- Comments are indicated with a "."
- In the solution a Loc Coulmn and Object Code Column are added

	Sou	rce stater	nent	
Line				
5 10 12	COPY FIRST	START STL LDB BASE	0 RETADR #LENGTH LENGTH	COPY FILE FROM INPUT TO OUTPUT SAVE RETURN ADDRESS ESTABLISH BASE REGISTER
13 15 20	CLOOP	+JSUB LDA COMP	RDREC LENGTH #0	READ INPUT RECORD TEST FOR EOF (LENGTH = 0)
25 30 35 40		JEQ +JSUB J LDA	ENDFIL WRREC CLOOP EOF	EXIT IF EOF FOUND WRITE OUTPUT RECORD LOOP
45 50 55	ENDFIL	STA LDA STA	BUFFER #3 LENGTH	INSERT END OF FILE MARKER SET LENGTH = 3
60 65 70 80	EOF RETADR	+JSUB J BYTE RESW	WRREC @RETADR C'EOF' 1	WRITE EOF RETURN TO CALLER
95 100 105	LENGTH BUFFER	RESW RESB	1 4096	LENGTH OF RECORD 4096-BYTE BUFFER AREA

Line Loc	Source state	ement	Object code
5 0000	COPY START	0	
10 0000	FIRST STL	RETADR	17202D
12 0003	LDB	#LENGTH	69202D
13	BASE	LENGTH	
15 0006	CLOOP +JSUB	RDREC	4B101036
20 000A	LDA	LENGTH	032026
25 000D	COMP	#0	290000
30 0010	JEQ	ENDFIL	332007
35 0013	+JSUB	WRREC	4B10105D
40 0017	J	CLOOP .	3F2FEC
45 001A	ENDFIL LDA	EOF	032010
50 001D	STA	BUFFER	0F2016
55 0020	LDA	#3	010003
60 0023	STA	LENGTH	0F200D
65 0026	+JSUB	WRREC	4B10105D
70 002A	J	@RETADR	3E2003
80 002D	EOF BYTE	C'EOF'	454F46
95 0030	RETADR RESW	1	
100 0033	LENGTH RESW	1	
105 0036	BUFFER RESB	4096	

- First we need to perform our assembler's first pass which will serve mainly to create the Loc Column and Symbol Table (SYMTAB)
- The Loc (Location) Column can be created by simply finding the program's starting address and tracking the amount of bytes that have been used by instructions
  - How many bytes per instruction?
  - We will need to use our OPTAB
     (we will see what this is in the next slide)
- The SYMTAB will keep track of the values in the Loc Column corresponding to each symbol in the Symbol Column

Line	Loc	Sou	rce staten	nent	Object code
Line	Loc	Sou	ree states		object code
5	0000	COPY	START	0	
10	0000	FIRST	STL	RETADR	17202D
12	0003		LDB	#LENGTH	69202D
13			BASE	LENGTH	
15	0006	CLOOP	+JSUB	RDREC	4B101036
2.0	000A		LDA	LENGTH	032026
25	000D		COMP	#0	290000
30	0010		JEQ	ENDFIL	332007
35	0013		+JSUB	WRREC	4B10105D
40	0017		J	CLOOP'	. 3F2FEC
45	001A	ENDFIL	LDA	EOF	032010,
50	001D		STA	BUFFER	0F2016
55	0020		LDA	#3	010003
60	0023		STA	LENGTH	0F200D
65	0026		+JSUB	WRREC	4B10105D
70	002A		J	GRETADR	3E2003
80	002D	EOF	BYTE	C'EOF'	454F46
95	0030	RETADR	RESW	1	
100	0033	LENGTH	RESW	1	
105	0036	BUFFER	RESB	4096	

- Next, we need to perform our second pass in which we generate the Object Code column
- Take a second to re-read (or read ⊗) 2.2.1 Instruction Formats and Addressing Modes (pg. 57-61) to get an idea of what we will need to do
- First we should create an Operation Table (OPTAB)
  - The OPTAB will contain an entry for each instruction belonging to the SIC/XE instruction set
  - Each entry should contain the instruction's corresponding mnemonic, opcode, and format(s)
  - This information is all given in the back of the book

- Once we have the OPTAB, we will use it to determine an instruction's opcode and format
- Then (if necessary) we need to calculate the addressing mode information (nixbpe bits)
- Finally (if necessary) we need to calculate the disp/target address information based on the addressing mode information
- All this information will be used to calculate the instructions object code
- Let us see a few examples

Take a look in functions.txt at this line:

FIRST STL RETADR .SAVE RETURN ADDRESS

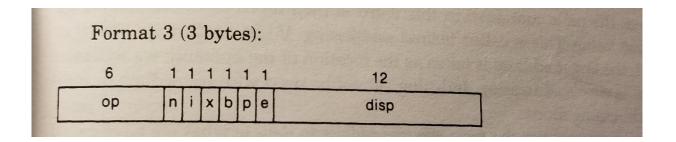
- First, from our OPTAB we can see STL has an opcode of 14 (hexadecimal) and has a format of 3/4 (3 or 4)
- This means the first byte of our opcode will contain the value 14 (0001 0100) plus the n and i bit values we find
- This also means we will either have a format 3 or 4 instruction

• Remember our SIC/XE addressing modes? (pg. 8 and 9)

Format 1	(1 byte):
8	
ор	The second second
Format 2 (2 bytes):	Participant Charles and Charles
8 4	4
op r1	r2
Format 3 (3 bytes):	
6 11111	1 12
op nixbpe	e disp
Format 4 (4 bytes):	THE CHARLES OF STREET
6 111111	1 20
op nixbpe	e address

• How can we tell if it is format 3 or 4?

- To indicate an instruction is format 4, the instruction will be preceded by a '+'
  - For example, in functions.txt take a look at the +JSUB WRECC line
- Since STL has no preceding '+' it must be format 3
- So we have the first six bits opcode of our 3 byte format 3 instruction, what about the rest?

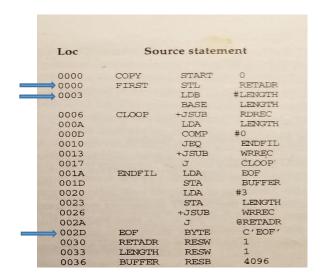


- Next we need to calculate the nixbpe bits
- Remember:
  - n=1, i=1 indicates simple addressing
    - We have this by default if we don't have indirect or immediate addressing
  - n=1, i=0 indicates indirect addressing
    - This will be signaled by an operand preceded by a '@'
    - For example, in functions.txt find the J @RETADR line
  - n=0, i=1 indicates immediate addressing
    - This will be signaled by an operand preceded by a '#'
    - For example, in functions.txt find the LDB #LENGTH line
  - x=1 indicates indexed addressing
    - This will be signaled by an operand followed by ', X'
    - For example, in functions.txt find the STCH BUFFER, X line
  - b=1, p=0 indicates for base relative addressing
  - b=0, p=1 indicates for program-counter relative addressing
  - e=1 indicates a format 4 instruction (e=0 indicates format 3 instruction)

- The base relative vs. program-counter relative distinction is a bit harder to discern
- In the reading we learned program-counter relative is used by default, but, if the displacement (disp) we calculate is out of range, we resort to base relative

Mode	Indication	Target address ca	Iculation
Base relative	b = 1, p = 0	TA = (B) + disp	$(0 \le \text{disp} \le 4095)$
Program-counter relative	b = 0, p = 1	TA = (PC) + disp	$(-2048 \le \text{disp} \le 2047)$

- So, in order to see whether we should use program-counter relative or base relative we need to calculate the displacement (disp)
- Its been a few slides, remember we are looking at the line: FIRST STL RETADR .SAVE RETURN ADDRESS
- From out first pass, we know this line's corresponding Loc Column value is 0000
- From the reading, we know that the PC register will contain the next Loc Column value (0003)
- We also know from our first pass-generated SYMTAB that the value corresponding to RETADR is 0030 (this is the target address (TA))



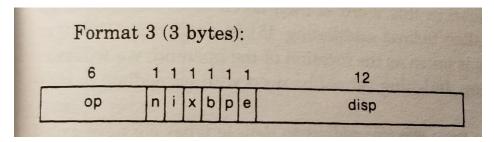
- From the book we see, for pc-relative addressing:
   TA=(PC)+disp -> disp=TA-(PC)
- This means our disp=0030-0003, disp=002D (hexadecimal values)
- This is within our program-counter relative disp bounds (-2048 <= disp <= 2047) (decimal values)

Mode	Indication	Target address ca	Iculation
Base relative	b = 1, p = 0	TA = (B) + disp	$(0 \le \text{disp} \le 4095)$
Program-counter relative	b = 0, p = 1	TA = (PC) + disp	$(-2048 \le \text{disp} \le 2047)$

 We finally have all the information we need to form the object code for line:

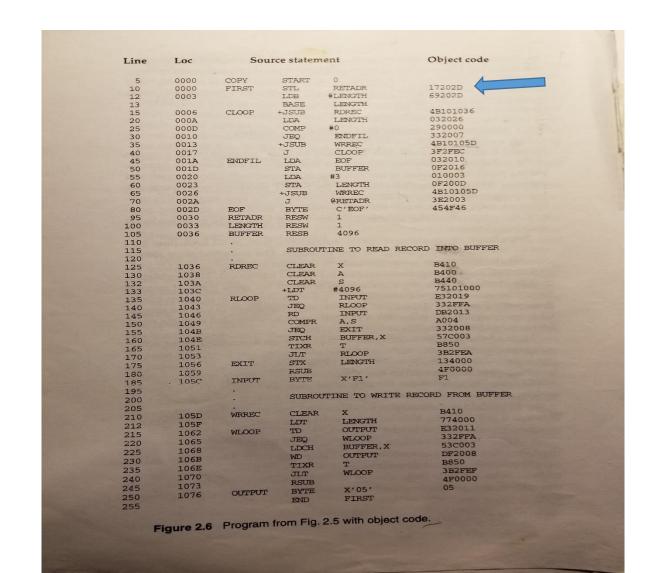
FIRST STL RETADR .SAVE RETURN ADDRESS

- op (opcode)=14 or 0001 0100
- n=1,i=1,x=0,b=0,p=1,e=0 or 11 0010
  - We have simple addressing because we didn't have immediate or indirect addressing!
- disp=02D or 0000 0011 1101



#### FIRST STL RETADR .SAVE RETURN ADDRESS

- Object Code Generation:
  - First byte = opcode + ni -> 0001 0100 + 11 (or 14 + 3) -> 0001 0111 (or 17)
     So, our first byte = 0001 0111 (or 17)
  - Second byte first half = xbpe -> 0010 (or 2)
     So, our second byte first half = 0010 (or 2)
  - Second byte second half and third byte = disp -> 0000 0011 1101 (or 02D)
     So, our second byte second half and third byte = 0000 0011 1101 (or 02D)
  - Altogether, our 3-byte format 3 instruction is:
     0001 0111 0010 0000 0011 1101 (or 17202D)



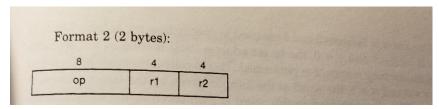
- The overall process will be quite similar for format 4 instructions
- This process will differ slightly with the different addressing modes
- Base relative addressing will require you to keep track of what is loaded into the B register for use in your disp calculation
- Immediate and indirect addressing effects also need to be considered!

- Let us take a look at a format 2 instruction
- Look at functions.txt and find the line:

COMPR A,S .TEST for End Of Record (X'00')

 From our OPTAB we see COMPR has opcode of A0 (hexadecimal value) and that it is a format 2 instruction

- So, we already have the op section ready (A0 or 1010 0000) for the first byte
- We need to get the r1 and r2 values for the second byte

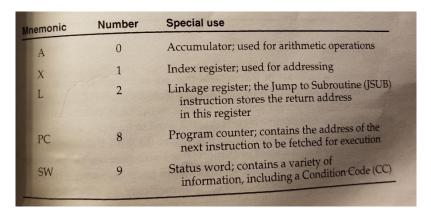


The book lists all the register values for us in the first chapter

Inemonic	Number	Special use
A	0	Accumulator; used for arithmetic operations
X	1	Index register; used for addressing
i /	2	Linkage register; the Jump to Subroutine (JSUB) instruction stores the return address in this register
PC	8	Program counter; contains the address of the next instruction to be fetched for execution
SW	9	Status word; contains a variety of information, including a Condition Code (CC)

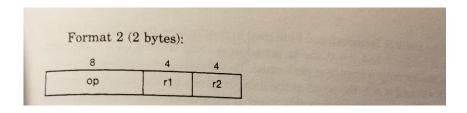
Mnemonic	Number	Special use
В	3	Base register; used for addressing
S	4	General working register—no special use
T	5	General working register—no special use
F	6	Floating-point accumulator (48 bits)

#### COMPR A,S .TEST for End Of Record (X'00')



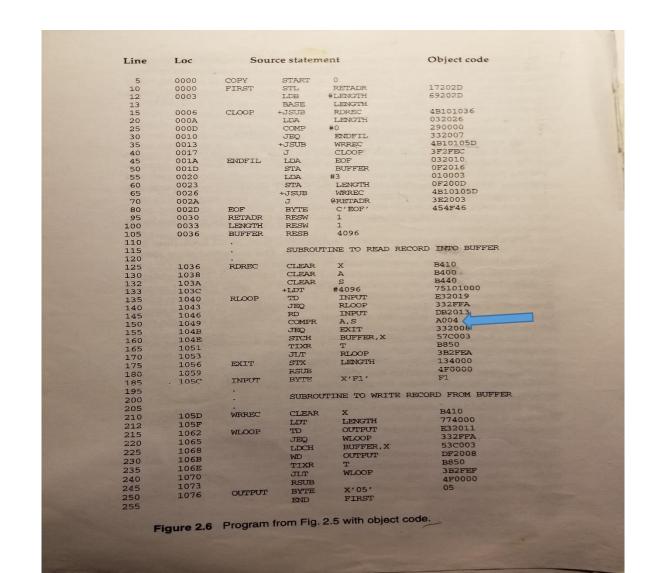
Mnemonic	Number	Special use
В	3	Base register; used for addressing
S	4	General working register—no special use
T	5	General working register—no special use
F	6	Floating-point accumulator (48 bits)

• We need the A (0 hexadecimal) and S (4 hexadecimal) register values for r1 and r2 respectively



COMPR A,S .TEST for End Of Record (X'00')

- Object Code Generation:
  - First byte = opcode -> 1010 0000 (or A0)
     So, our first byte = 1010 0000 (or A0)
  - Second byte first half = r1 -> A -> 0000 (or 0)
     So, the second byte first half = 0000 (or 0)
  - Second byte second half = r2 -> S -> 0100 (or 4)
     So, the second byte second half = 0100 (or 4)
  - Altogether, our 2-byte format 2 instruction is:
     1010 0000 0000 0100 (or A004)



- With each of the sample programs there will need to be added functionality your assembler will need to support
  - literals.txt adds literals
  - program-blocks.txt adds program blocks
  - etc.
- You also should create object programs using the object code you generate
  - I would suggest leaving this part until you can support all six test files
- I will leave figuring out these parts to you
- Your book is very straightforward and will be your friend as you are figuring out these details
- You can come to my office hours and ask me questions

#### Course Project: Testing

- Your assembler should support the six test programs
- To test for robustness, your assembler will be tested using files you do not have access to

- Good luck and have fun!
- Email: phillity@iu.edu
- Office Hour: MW 12:30-1:30 SL228

