SOA 07c-01

# Map-Reduce

Service-Oriented Architecture Jeremy Gibbons

## I Google's Map-Reduce architecture

- 'simplified data processing on large clusters'
- massively parallel data analysis applications
- programming model and library
- distribution, fault tolerance, I/O scheduling, monitoring
- inspired by FP notions ('map' and 'reduce')
- paper by Jeffrey Dean and Sanjay Ghemawat at OSDI2004
- see also http://labs.google.com/papers/mapreduce.html
- more recent:

http://www.slideshare.net/jhammerb/mapreduce-pact06-keynote

#### 2 Problem

- 20+ billion web pages, 20KB each: 400+TB
- disk read speed 30–35MB/s: four months to read data
- 1000 disks even to store it, but then only 3 hours to read
- however, programming is very difficult
- stuff breaks: 10,000 servers, MTBF  $\approx$  3 years each...

# 3 Typical task structure

- 1. read a lot of data
- 2. extract something of interest from each record ('map')
- 3. shuffle and sort
- 4. aggregate, summarize, filter, or transform ('reduce')
- 5. write results

## 4 Programming model

Based on key-value pairs. Programmer writes just two functions.

*Map function* takes input key-value to set of intermediate key-values:

```
map :: (InKey, InValue) -> List (OutKey, OutValue)
```

*Reduce function* combines values for each intermediate key (usually producing just a single value):

```
reduce :: (OutKey, List OutValue) -> List OutValue
```

Sometimes also a *partition function*:

```
partition :: (OutKey, NumPartitions) -> PartitionId
```

(by default, hashing modulo number of partitions).

### 4.1 Example: word frequencies (pseudocode)

#### 4.2 Many more examples

distributed grep: mapper emits matching lines, reducer is identity

term vectors: like wordcount, but discarding infrequent words

access statistics: wordcount on web server logs

inverted index: mapper emits (Word, DocId) pairs

reverse web-links: mapper emits (Target, Source) pairs for hyperlinks

distributed sort: depends on ordering properties of partition, output

Lots of *extract-transform-load* tasks work well.

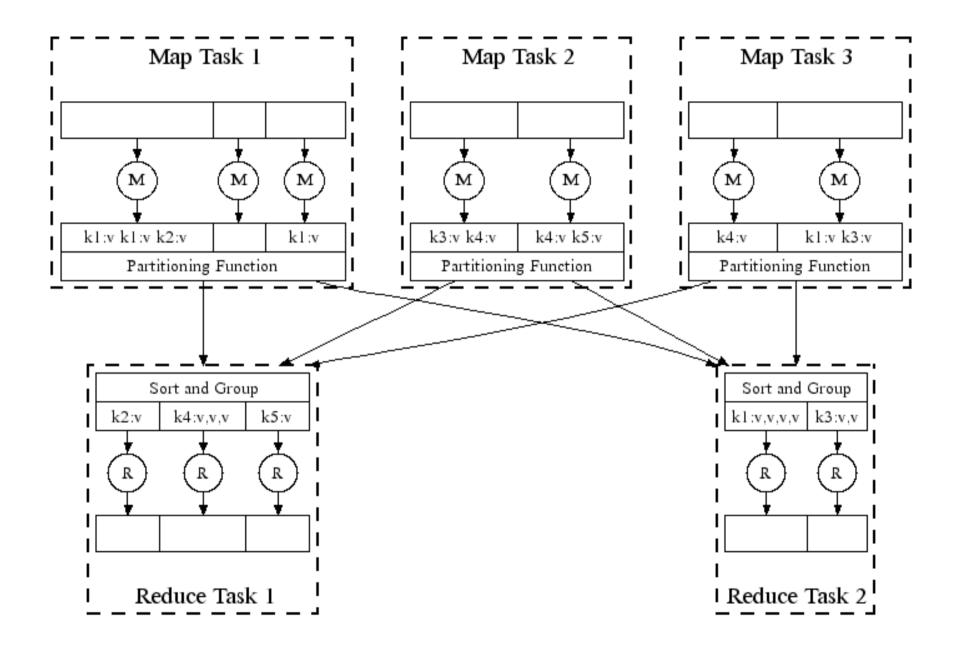
Thousands of applications in Google source tree (6000+ in 2006). Production indexing system is a sequence of 24 MR instances.

## 5 Implementation

- commodity PCs (eg dual-core x86 running linux)
- commodity networking (eg 100Mb/s or 1Gb/s)
- cheap disks attached directly to computers
- custom redundant distributed file system
- cluster of hundreds or thousands of machines (eg 1800 in OSDI paper)

#### 5.1 Execution

- divide input data into M splits, typically about 64MB
- apply map function to each split
- arrange intermediate values into *R partitions*
- apply reduce function to each partition
- often M = 200,000 and R = 5,000 on N = 2,000 workers: fine-grain
- *master* process assigns tasks to workers, taking locality into account



#### 5.2 Fault tolerance

- regular polling of worker machines
- failure handled by restarting map tasks and in-progress reduce tasks
- (complete reduce tasks stored on global server)
- master failure is unlikely, because only a single machine
- soundness depends on lack of side-effects (and on transactional commits)

### 5.3 Backup tasks

- straggler tasks lengthen overall computation time
- eg from bad disk (slowing reads from 30MB/s to 1MB/s)
- eg faulty initialization, disabling processor cache (100-fold slowdown)
- when nearly done, master process spawns backup executions of remaining in-progress tasks

#### **5.4 Refinements**

- guaranteed ordering by key within each partition
- optional combiner function for output of map tasks
- side-effects allowed, if atomic and idempotent
- skipping bad records (eg for buggy user code)
- local execution mode for debugging
- master process runs internal web server for statistics

## 6 Hadoop

- Apache open-source framework for scalable distributed computing
- MapReduce, after Google
- HDFS distributed file system
- other subprojects: Avro, Chukwa, Hive, Pig...
- http://hadoop.apache.org/