csci 210: Data Structures

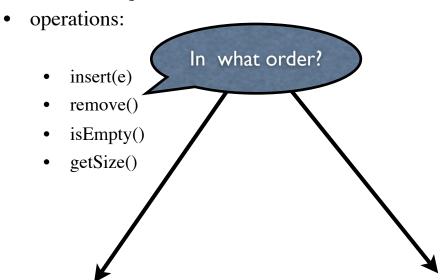
Stacks and Queues

Summary

- Topics
 - stacks and queues as abstract data types
 - implementations
 - arrays
 - linked lists
 - analysis and comparison
 - application: searching with stacks and queues
 - Problem: missionary and cannibals
 - Problem: finding way out of a maze
 - depth-first and breadth-first search
- READING:
 - GT textbook chapter 5

Stacks and Queues

- Fundamental "abstract" data types
 - abstract, i.e. we think of their interface and functionality; the implementation may vary
- Interface:
 - stacks and queues handle a collection of elements

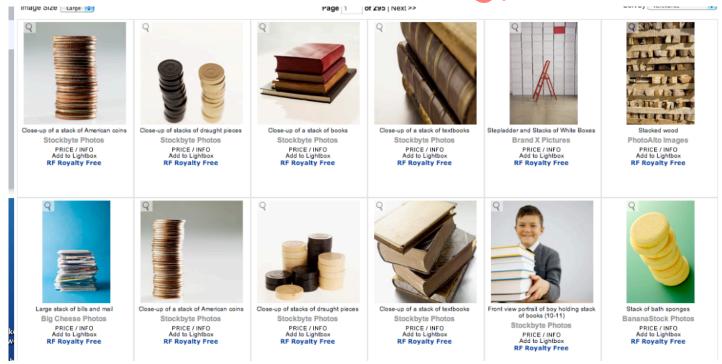


- Stacks
 - only last element can be deleted
 - ==>insert and delete at one end
 - last-in-first-out (LIFO)

Queues

- only first element can be deleted
- ==> insert at one end, delete from the other end
- First-in-first-out (FIFO)

Stack analogy



Stack interface

- push(e)
 - insert element e
- pop()
 - delete and return the last inserted element
- size()
 - return the number of elements in the queue
- isEmpty()
 - return true if queue is empty

Queue Analogy

Queue interface

- enqueue(e)
 - insert element e
- dequeue()
 - delete and return the first inserted element
- size()
 - return the number of elements in the queue
- isEmpty()
 - return true if queue is empty



Applications

- Are stacks and queues useful?
 - YES. They come up in many problems.

Stacks

- Internet Web browsers store the addresses of recently visited sites on a stack. Each time the visits a new site ==> pushed on the stack. Browsers allow to "pop" back to previously visited site.
- The undo-mechanism in an editor. The changes are kept in a stack. When the user presses "undo" the stack of changes is popped.
- The function-call mechanism
 - the active (called but not completed) functions are kept on a stack
 - each time a function is called, a new frame describing its context is pushed onto the stack
 - when the function returns, its frame is popped, and the context is reset to the previous method (now on top of the stack)

Queues

- queue of processes waiting to be processed; for e.g. the queue of processes to be scheduled on the CPU
- round-robin scheduling: iterate through a set of processes in a circular manner and service each element:
 - e.g. the process at front is dequeued, allowed to run for some CPU cycles, and then enqueued at the end of the queue

Using Stacks

• java.util.Stack

Constructor Summary

Stack()

Creates an empty Stack.

Method Summary

ı			
	boolean	empty() Tests if this stack is empty.	
	<u>Object</u>	Looks at the object at the top of this stack without removing it from the stack.	
	<u>Object</u>	Removes the object at the top of this stack and returns that object as the value of this function.	
	<u>Object</u>	Pushes an item onto the top of this stack.	
	int	Returns the 1-based position where an object is on this stack.	

Using Stacks

```
Stack<Integer> st = new Stack<Integer>();
s.push (3);
s.push (5);
s.push (2);
//print the top
System.out.print(s.peek());
s.pop();
s.pop();
s.pop();
```

A Stack Interface

- stack can contain elements of arbitrary type E
- use generics: define Stack in terms of a generic element type E

```
Stack<E> {
}...
```

• when instantiating Stack, specify E

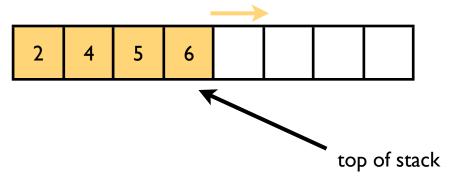
```
Stack<String> st;
```

• Note: could use Object, but then need to cast every pop()

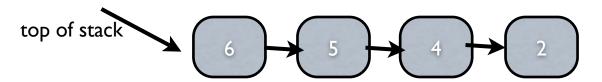
```
/**
  * Interface for a stack: a collection of objects that are inserted
  * and removed according to the last-in first-out principle. This
  * interface includes the main methods of java.util.Stack.
  */
public interface Stack<E> {
 /**
  * Return the number of elements in the stack.
  * @return number of elements in the stack.
  * /
 public int size();
 /**
  * Return whether the stack is empty.
  * @return true if the stack is empty, false otherwise.
  */
 public boolean isEmpty();
 /**
  * Inspect the element at the top of the stack.
  * @return top element in the stack.
  * @exception EmptyStackException if the stack is empty.
  */
 public E top()
    throws EmptyStackException;
 /**
  * Insert an element at the top of the stack.
  * @param element to be inserted.
  */
 public void push (E element);
 /**
  * Remove the top element from the stack.
  * @return element removed.
  * @exception EmptyStackException if the stack is empty.
  */
 public E pop()
    throws EmptyStackException;
}
```

Stack Implementations

- Stacks can be implemented efficiently with both
 - arrays
 - linked lists
- Array implementation of a Stack



- Linked-list implementation of a stack
 - a linked list provides fast inserts and deletes at head
 - ==> keep top of stack at front



Next..

- Provide sketch for each implementation
- Analyze efficiency
- Compare

Arrays vs Linked-List Implementations

Array

- simple and efficient
- assume a fixed capacity for array
 - if CAP is too small, can reallocate, but expensive
 - if CAP is too large, space waste

• Lists

- no size limitation
- extra space per element

• Summary:

• when know the max. number of element, use arrays

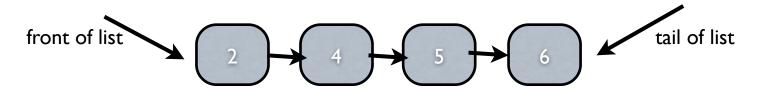
Method	Time
size()	O(I)
isEmpty()	O(I)
top	O(I)
push	O(I)
рор	O(I)

A Queue Interface

```
public interface Queue<E> {
 /**
  * Returns the number of elements in the queue.
  * @return number of elements in the queue.
  * /
 public int size();
 /**
  * Returns whether the queue is empty.
  * @return true if the queue is empty, false otherwise.
  */
 public boolean isEmpty();
 /**
  * Inspects the element at the front of the queue.
  * @return element at the front of the queue.
  * @exception EmptyQueueException if the queue is empty.
  */
 public E front() throws EmptyQueueException;
 /**
  * Inserts an element at the rear of the queue.
  * @param element new element to be inserted.
  */
 public void enqueue (E element);
 /**
  * Removes the element at the front of the queue.
  * @return element removed.
  * @exception EmptyQueueException if the queue is empty.
  */
 public E dequeue() throws EmptyQueueException;
}
```

Queue Implementations

- Queue with arrays
 - say we insert at front and delete at end
 - need to shift elements on inserts ==> insert not O(1)
- Queue with linked-list
 - in a singly linked-list can delete at front and insert at end in O(1)



• Exercise: sketch implementation

Method	Time
size()	O(I)
isEmpty()	O(I)
front	O(I)
enqueue	O(I)
dequeue	O(I)

Queue with a Circular Array

• A queue can be implemented efficiently with a circular array if we know the maximum number of elements in the queue at any time