



- Quick Sort is based on the Divide and Conquer approach
- Fastest known sorting algorithm in practice
- Average case: O(N log N)
- Worst case: O(N²)

Quick Sort: Main Idea



- Pick any element v called the pivot in S.
- 2. Partition the elements in S except v into two disjoint groups:

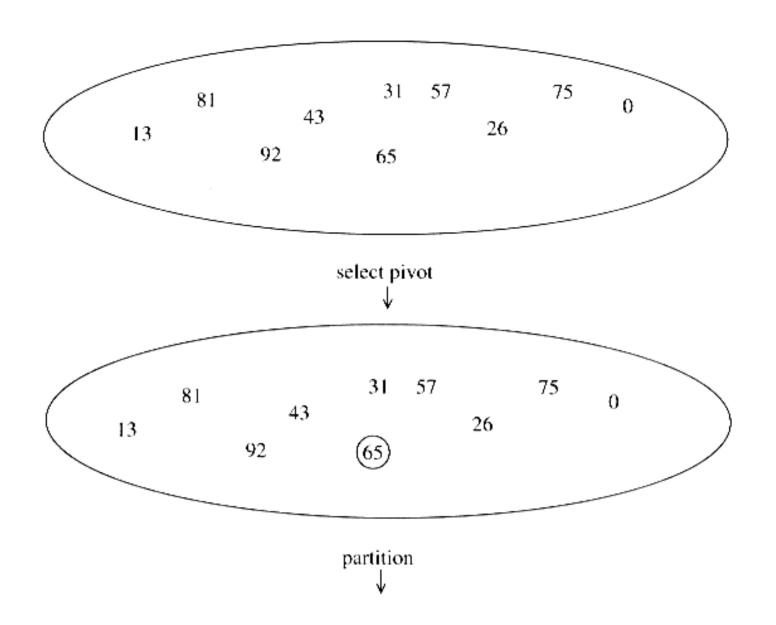
1.
$$S_1 = \{ x \in S - \{ v \} \mid x \le v \}$$

2.
$$S_2 = \{ x \in S - \{ v \} \mid x \ge v \}$$

3. QuickSort $(S_1) + v + QuickSort (S_2)$

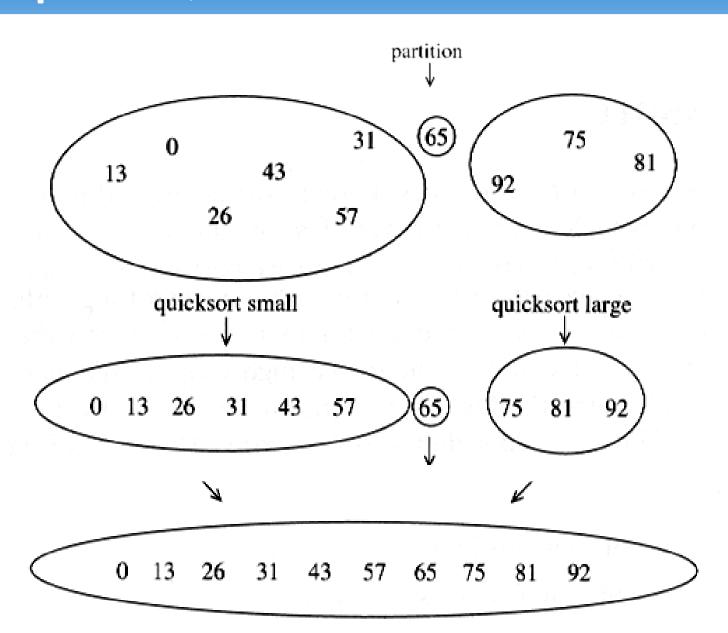
Quick Sort: Example





Example of Quick Sort...







- Array is divided into subarrays by selecting a pivot element(element selected from the array).
- While dividing the array, the pivot element should be positioned in such a way that elements less than pivot are kept on the left side and elements greater than pivot are on the right side of the pivot.
- The left and right subarrays are also divided using the same approach.
- This process continues until each subarray contains a single element.
- At this point, elements are already sorted.





1. Select the Pivot Element

- There are different variations of quicksort where the pivot element is selected from different positions.
 - Rightmost element of the array
 - Leftmost element of the array
 - Any random element from the array
 - Median element of the array

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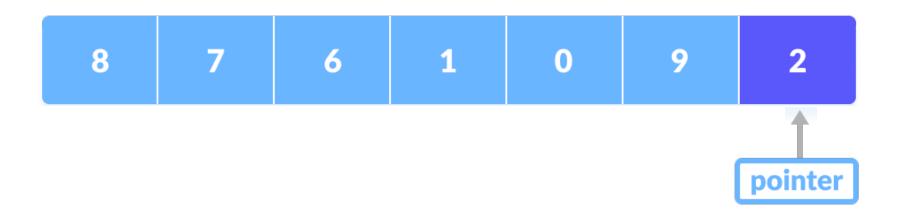
2. Rearrange the Array

a) Elements of the array are rearranged so that elements that are smaller than the pivot are put on the left and the elements greater than the pivot are put on the right.



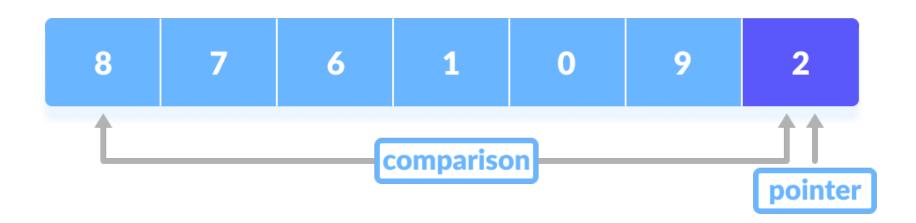


- Select the Pivot Element
- A pointer is fixed at the pivot element.



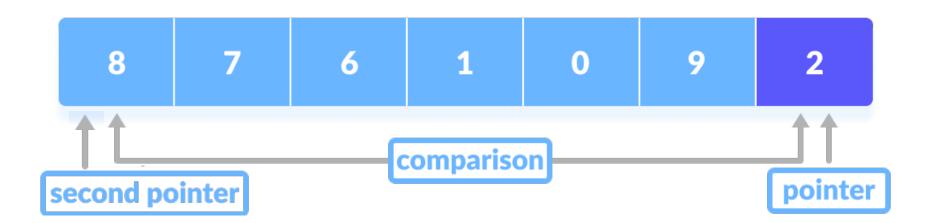


 The pivot element is compared with the elements beginning from the first index.



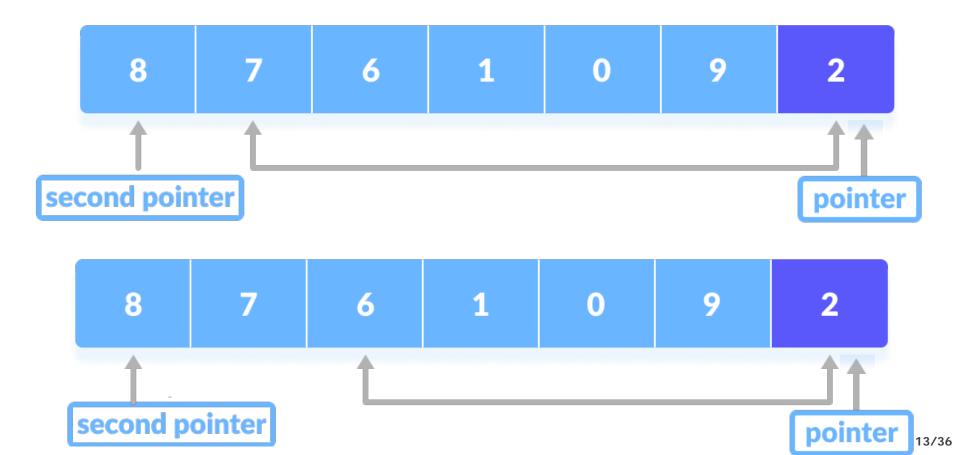


 If the element is greater than the pivot element, a second pointer is set for that element.



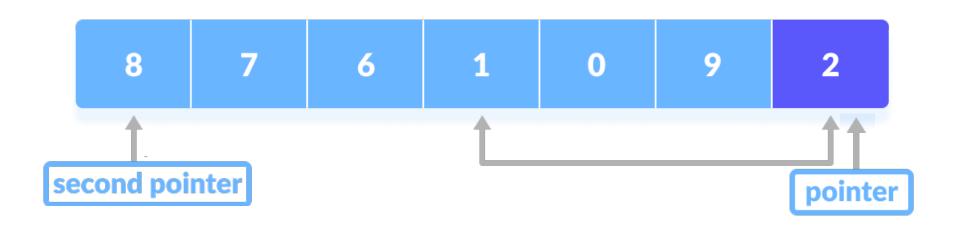


- Now, pivot is compared with other elements.
 - If the element greater than the pivot element, and second pointer is already set then go to the next element.



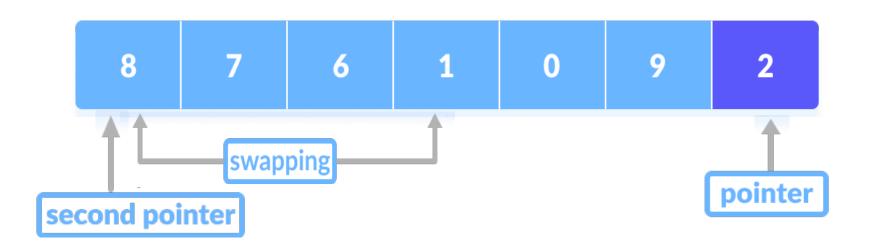


 If an element smaller than the pivot element is found, then the smaller element is swapped with the greater element found earlier.



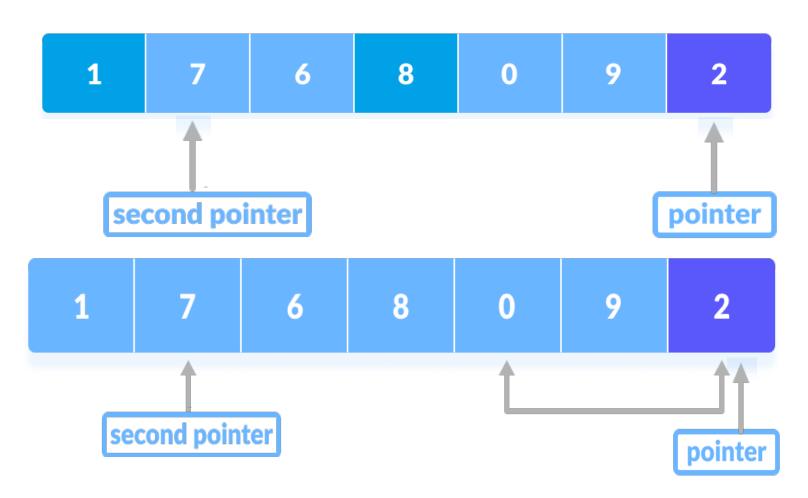


Smaller element is swapped with the greater element found earlier.



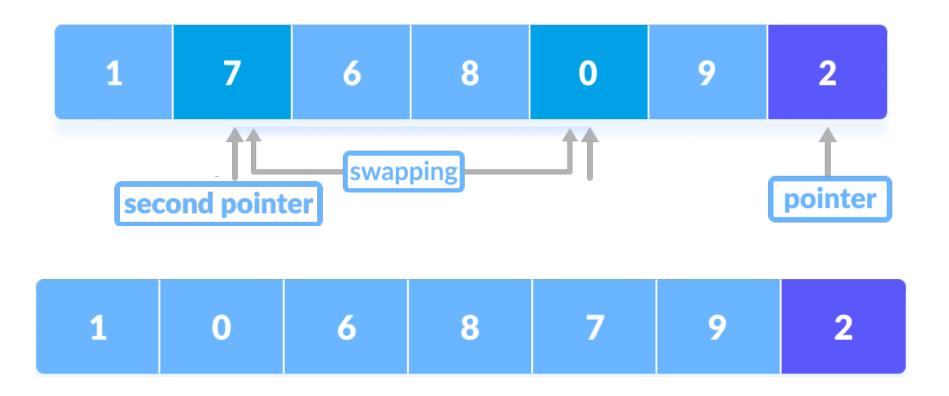


 The process is repeated to set the next greater element as the second pointer. And, swap it with another smaller element.



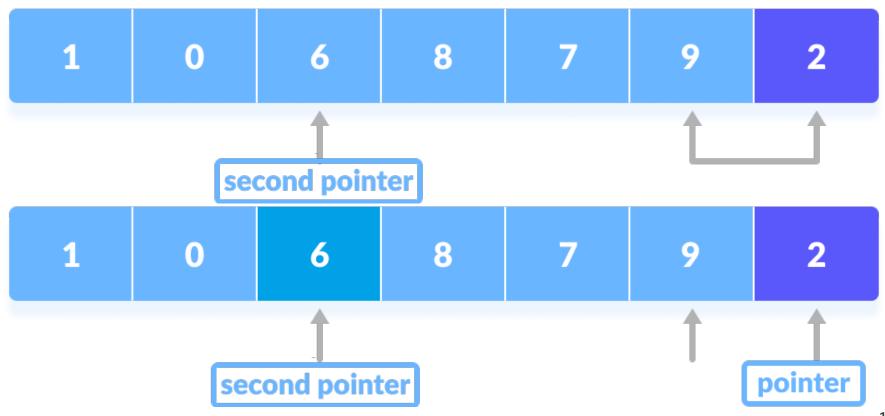


 The process is repeated to set the next greater element as the second pointer. And, swap it with another smaller element.



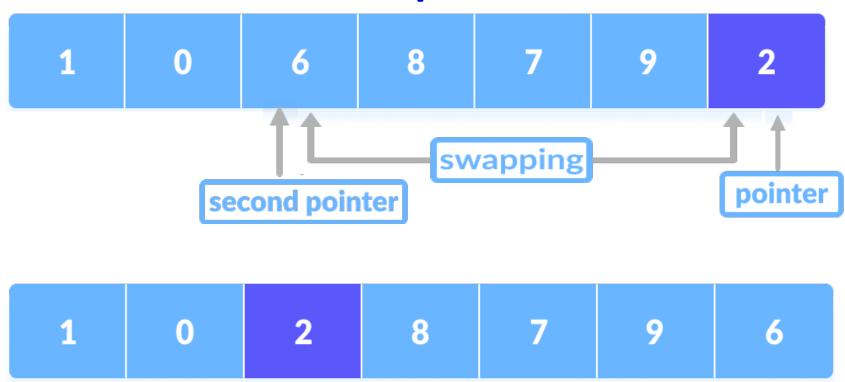


 The process goes on until the second last element is reached





 Finally, the pivot element is swapped with the second pointer.





3. Divide Subarrays

Pivot elements are again chosen for the left and the right sub-parts separately and step 2 is repeated.

Select pivot element in each half and put at correct place using recursion







```
// Quick sort in C
#include <stdio.h>
// function to swap elements
void swap(int *a, int *b)
   int t = *a;
   *a = *b;
   *b = t;
```



```
// function to find the partition position
int partition(int array[], int low, int high)
{ // select the rightmost element as pivot
 int pivot = array[high];
 // pointer for greater element
 int i = (low - 1);
```



```
// traverse each element of the array
// compare them with the pivot
for (int j = low; j < high; j++)
{ if (array[j] <= pivot)
 { // if element smaller than pivot is found
    // swap it with the greater element
    // pointed by i
   i++;
```



```
// swap element at i with element at i
   swap(&array[i], &array[j]);
//swap pivot element with greater element at i
 swap(&array[i + 1], &array[high]);
  // return the partition point
 return (i + 1);
```



```
void quickSort(int array[], int low, int high) {
 if (low < high) {
  // find the pivot element such that
  // elements smaller than pivot are on
  // left of pivot
  // elements greater than pivot are on
  // right of pivot
  int pi = partition(array, low, high);
```



```
// recursive call on the left of pivot
quickSort(array, low, pi - 1);
// recursive call on the right of pivot
quickSort(array, pi + 1, high);
```



```
// function to print array elements
void printArray(int array[], int size)
{ for (int i = 0; i < size; ++i) {
  printf("%d ", array[i]);
 printf("\n");
```



```
int main()
 int data[] = \{8, 7, 2, 1, 0, 9, 6\};
  int n = sizeof(data) / sizeof(data[0]);
  printf("Unsorted Array\n");
  printArray(data, n);
  // perform quicksort on data
  quickSort(data, 0, n - 1);
  printf("Sorted array in ascending order: \n");
  printArray(data, n);
```



Quick Sort Algorithm in Python



```
# Quick sort in Python
# function to find the partition position
def partition(array, low, high):
 # choose the rightmost element as pivot
 pivot = array[high]
 # pointer for greater element
 i = low - 1
 # traverse through all elements
 # compare each element with pivot
```



```
for j in range(low, high):
  if array[j] <= pivot:</pre>
   # if element smaller than pivot is found
   # swap it with the greater element pointed
   # by i
   i = i + 1
   # swapping element at i with element at j
   (array[i], array[j]) = (array[j], array[i])
```



```
# swap the pivot element with the greater
# element specified by i
(array[i + 1], array[high]) = (array[high],
                                array[i + 1])
# return the position from where partition
# is done
return i + 1
```



```
# function to perform quicksort
def quickSort(array, low, high):
 if low < high:
  # find pivot element such that element smaller than
  # pivot are on the left element greater than pivot
  # are on the right
  pi = partition(array, low, high)
  # recursive call on the left of pivot
  quickSort(array, low, pi - 1)
  # recursive call on the right of pivot
  quickSort(array, pi + 1, high)
```



```
data = [38,47,942,61,10,90,46,66,53,83,16]
print("Unsorted Array")
print(data)
size = len(data)
quickSort(data, 0, size - 1)
print('Sorted Array in Ascending Order:')
print(data)
```



Unsorted Array

[38, 47, 942, 61, 10, 90, 46, 66, 53, 83, 16]

Sorted Array in Ascending Order:

[10, 16, 38, 46, 47, 53, 61, 66, 83, 90, 942]

Analysis of Quick Sort



Worst Case Complexity [Big-O]: O(n²)

- It occurs when the pivot element picked is either the greatest or the smallest element.
- This condition leads to the case in which the pivot element lies in an extreme end of the sorted array.
- One sub-array is always empty and another sub-array contains n 1 elements.
- However, the quicksort algorithm has better performance for scattered pivots.

Analysis of Quick Sort



Best Case Complexity [Big-omega]: O(n*log n)

•It occurs when the pivot element is always the middle element or near to the middle element.

Average Case Complexity [Big-theta]: O(n*log n)

•It occurs when the above conditions do not occur.

Analysis of Quick Sort



Space Complexity

 The space complexity for quicksort is O(log n).

Visual Demo of Quick Sort



https://visualgo.net/en/sorting