

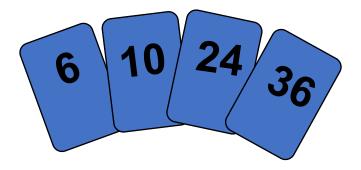
# SORTING (Insertion and Merge)





- Idea: like sorting a hand of playing cards
  - Start with an empty left hand and the cards facing down on the table.
  - Remove one card at a time from the table, and insert it into the correct position in the left hand
    - compare it with each of the cards already in the hand, from right to left
  - The cards held in the left hand are sorted
    - these cards were originally the top cards of the pile on the table

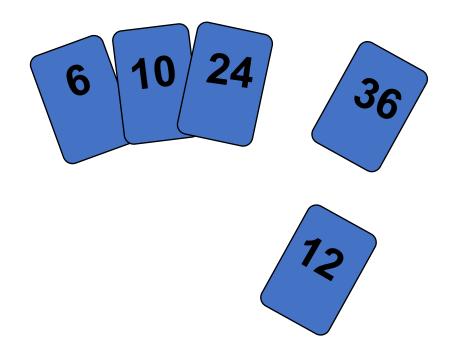




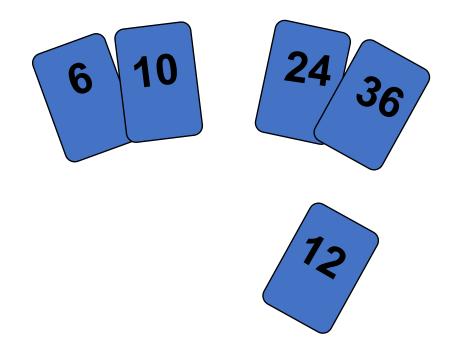


To insert 12, we need to make room for it by moving first 36 and then 24.

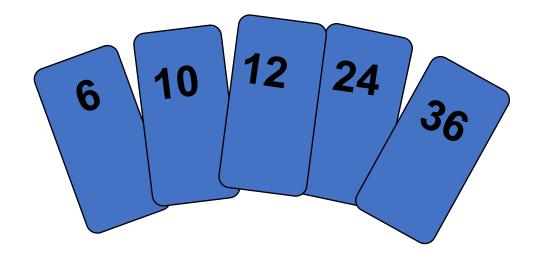














#### Insertion Sort Execution Example









- Iteration i. Repeatedly swap element i with the one to its left if smaller.
- After  $i^{th}$  iteration, a [0] through a [i] contain first i+1 elements in ascending order.

Array index	0	1	2	3	4	5	6	7	8	9
Value	2.78	7.42	0.56	1.12	1.17	0.32	6.21	4.42	3.14	7.71

Iteration 0: step 0.



Iteration i. Repeatedly swap element i with the one to its left if smaller.

❖ After i<sup>th</sup> iteration, a[0] through a[i] contain first i+1 elements in ascending order.

Array index	O	1	2	3	4	5	6	7	8	9
Value	2.78	7.42	0.56	1.12	1.17	0.32	6.21	4.42	3.14	7.71

Iteration 1: step 0.



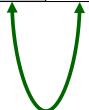
Array index	0	1	2	3	4	5	6	7	8	9
Value	2.78	0.56	7.42	1.12	1.17	0.32	6.21	4.42	3.14	7.71



Iteration 2: step 0.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.56	2.78	7.42	1.12	1.17	0.32	6.21	4.42	3.14	7.71



Iteration 2: step 1.

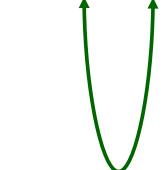


Array index	0	1	2	3	4	5	6	7	8	9
Value	0.56	2.78	7.42	1.12	1.17	0.32	6.21	4.42	3.14	7.71

Iteration 2: step 2.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.56	2.78	1.12	7.42	1.17	0.32	6.21	4.42	3.14	7.71



Iteration 3: step 0.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	2.78	7.42	1.17	0.32	6.21	4.42	3.14	7.71

Iteration 3: step 1.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	2.78	7.42	1.17	0.32	6.21	4.42	3.14	7.71

Iteration 3: step 2.



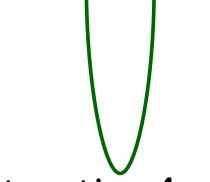
Array index	O	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	2.78	1.17	7.42	0.32	6.21	4.42	3.14	7.71



Iteration 4: step 0.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	1.17	2.78	7.42	0.32	6.21	4.42	3.14	7.71



Iteration 4: step 1.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	1.17	2.78	7.42	0.32	6.21	4.42	3.14	7.71

Iteration 4: step 2.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	1.17	2.78	0.32	7.42	6.21	4.42	3.14	7.71

Iteration 5: step 0.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	1.17	0.32	2.78	7.42	6.21	4.42	3.14	7.71

Iteration 5: step 1.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.56	1.12	0.32	1.17	2.78	7.42	6.21	4.42	3.14	7.71



Iteration 5: step 2.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.56	0.32	1.12	1.17	2.78	7.42	6.21	4.42	3.14	7.71

Iteration 5: step 3.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	7.42	6.21	4.42	3.14	7.71



Iteration 5: step 4.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	7.42	6.21	4.42	3.14	7.71

Iteration 5: step 5.



Array index	0	1	2	3	4	15	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	6.21	7.42	4.42	3.14	7.71

Iteration 6: step 0.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	6.21	7.42	4.42	3.14	7.71

Iteration 6: step 1.

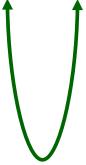


Array index	0	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	6.21	4.42	7.42	3.14	7.71

Iteration 7: step 0.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	4.42	6.21	7.42	3.14	7.71



Iteration 7: step 1.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	4.42	6.21	7.42	3.14	7.71

Iteration 7: step 2.

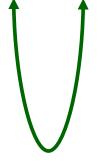


Array index	O	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	4.42	6.21	3.14	7.42	7.71

Iteration 8: step 0.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	4.42	3.14	6.21	7.42	7.71



Iteration 8: step 1.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	3.14	4.42	6.21	7.42	7.71



Iteration 8: step 2.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	3.14	4.42	6.21	7.42	7.71

Iteration 8: step 3.



Array index	O	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	3.14	4.42	6.21	7.42	7.71

Iteration 9: step 0.



Array index	0	1	2	3	4	5	6	7	8	9
Value	0.32	0.56	1.12	1.17	2.78	3.14	4.42	6.21	7.42	7.71

Iteration 10: DONE.

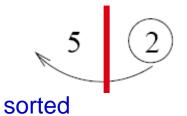
### **Insertion Sort**



input array

at each iteration, the array is divided in two sub-arrays:

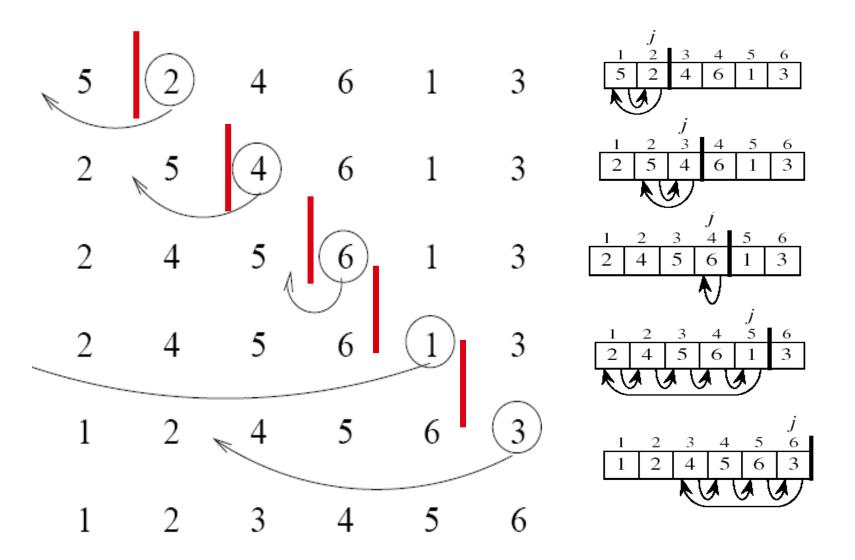
left sub-array right sub-array



unsorted

# **Insertion Sort**





#### **INSERTION SORT ALGORITHM**



 $\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \end{bmatrix}$   $\begin{bmatrix} a_1 & a_2 & a_3 & a_4 & a_5 & a_6 & a_7 & a_8 \end{bmatrix}$ 

Sorts an array 'A' with 'n' elements using the INSERTION sort.

Step1: Repeat steps 2 to 4 for i=2 to n

Step2: Set **Temp = A[i]** and k = i-1

Step3: Repeat While **Temp<A[k] AND k>0** 

Set A[k+1] = A[k]

Set **k** = **k-1** 

[End loop]

Step4: Set A[k+1] = Temp

[End loop]

Step5: Exit

# Is it right code for Insertion Sort?



```
for (i = 1; i < n; i++)
    key = A [ i ];
    j = i - 1;
   while(j \ge 0 \&\& A[j] > key)
    \{ A[j+1] = A[j];
       j--;
    A[j+1] = key;
```



```
#include <stdio.h>
main()
{ int i,n=12,j,key, A[12] = \{ 97,556,43,-5,6,404,55,3,22,-122,4,35 \};
for( i=0;i<n;i++) printf(" %d ",A[i]);
for (i = 1; i < n; i++)
{ key = A [ i ];
 j = i - 1;
 while(j \ge 0 \&\& A [j] > key)
 \{ A[j+1] = A[j];
    j--;
  A[j+1] = key;
```



```
printf("\n");
for( i=0;i<n;i++)
  printf(" %d ",A[i]);
}</pre>
```



Which is the WORST CASE for Insertion Sort ?
 WORST CASE for Insertion Sort is when elements of the list are in the reverse order of the desired one.

Desired order is Ascending

Input List is: 896, 467, 235, 198, 89,75, 64, 15, 6



- Which is the BEST CASE for Insertion Sort ?
  - BEST CASE for Insertion Sort is when elements of the list are already sorted in desired order.
- Desired order is Descending
- Input List is: 896, 467, 235, 198, 89,75, 64, 15, 6



- The worst case for Insertion sorting is, when elements of the list are in the reverse order of the desired one.
- In the worst case, each element of the unsorted part will be compared with all the elements of the sorted part of the list. Thus, in worst case,
- To position the 2<sup>nd</sup> element of the list at proper position, number of comparisons = 1
- To position the 3<sup>rd</sup> element of the list at proper position, number of comparisons = 2
- To position the 4<sup>th</sup> element of the list at proper position, number of comparisons = 3



•:

•

- To position the n<sup>th</sup> element of the list at proper position, number of comparisons = n-1
- So, the total number. of comparisons in the worst case.

$$f(n) = 1+2+3+4+5+6+....+(n-2)+(n-1)$$
$$= n*(n-1)/2$$
$$= O(n2)$$

- Thus, for the worst case the complexity of the insertion sort is O (n²).
- The best case for this algorithm is when elements are already sorted. Then the complexity will be O(n).
- However, for the average case complexity will be O(n²).





- Merge sort is one of the best sorting techniques which is based on the divide and conquer strategy.
- It is an in-place sorting algorithm; It uses no auxiliary data structures (extra space) while sorting.
- Partition the n > 1 elements into two smaller instances.
- •First (n/2) elements define one of the smaller instances; remaining (n/2) elements define the second smaller instance.
- Merge Sort reiterates on both the newly formed partitions by dividing them in half and continue the process.



- This process of dividing stop when the partition size reaches to one item.
- •At this point it has created many one-item lists. Any one-item list is naturally in sorted order.
- •The next step is to merge these one-item lists together for creating the sorted list.
- •To combine two sorted lists, the merge sort compare successive pairs of elements, one from each list.
- •If the list A has any element < than all the elements of list B, then it will be chosen to be append to the aggregated list or vice versa.



- And when all the elements of one list are added to the aggregated list then all the remaining elements of other list will be append directly to the aggregated list.
- Merge Sort repeats the process of combining sorted sub-lists into larger sorted sub-lists until all have been successfully integrated into a single sorted list.
- This merging process takes at least (n/2) comparisons but not more than (n-1) comparisons.
- Complexity is O(n log n).

# **Merge Two Sorted Lists**



$$A = (2, 5, 6)$$

$$B = (1, 3, 8, 9, 10)$$

$$C = ()$$

Compare smallest elements of A and B and merge smaller into C.

$$A = (2, 5, 6)$$

$$B = (3, 8, 9, 10)$$

$$C = (1)$$

# **Merge Two Sorted Lists**



$$A = (5, 6)$$

$$B = (3, 8, 9, 10)$$

$$C = (1, 2)$$

$$A = (5, 6)$$

$$B = (8, 9, 10)$$

$$C = (1, 2, 3)$$

$$A = (6)$$

$$B = (8, 9, 10)$$

$$C = (1, 2, 3, 5)$$

# **Merge Two Sorted Lists**



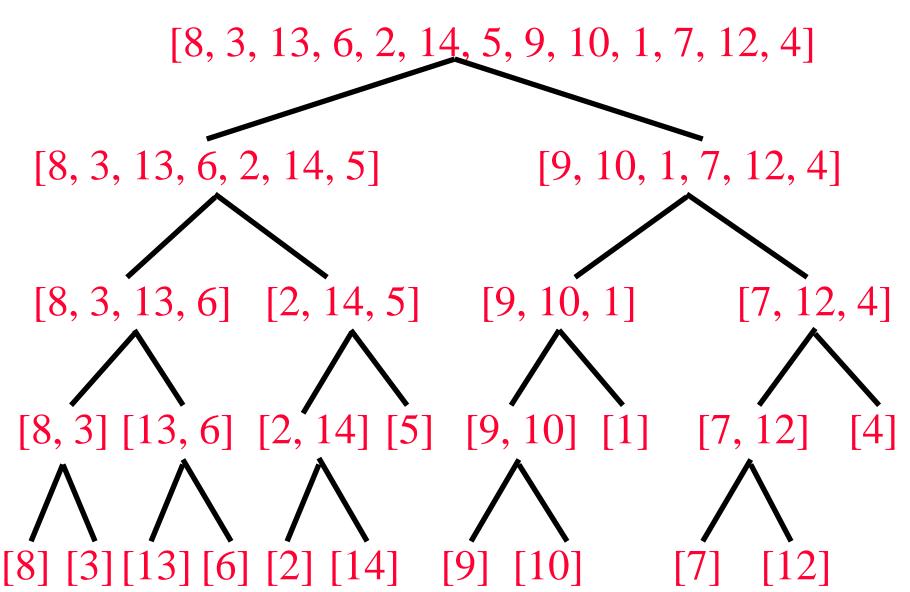
When one of A and B becomes empty, append the other list to C.

O(1) time needed to move an element into C.

Total time is O(n + m), where n and m are, respectively, the number of elements initially in A and B.

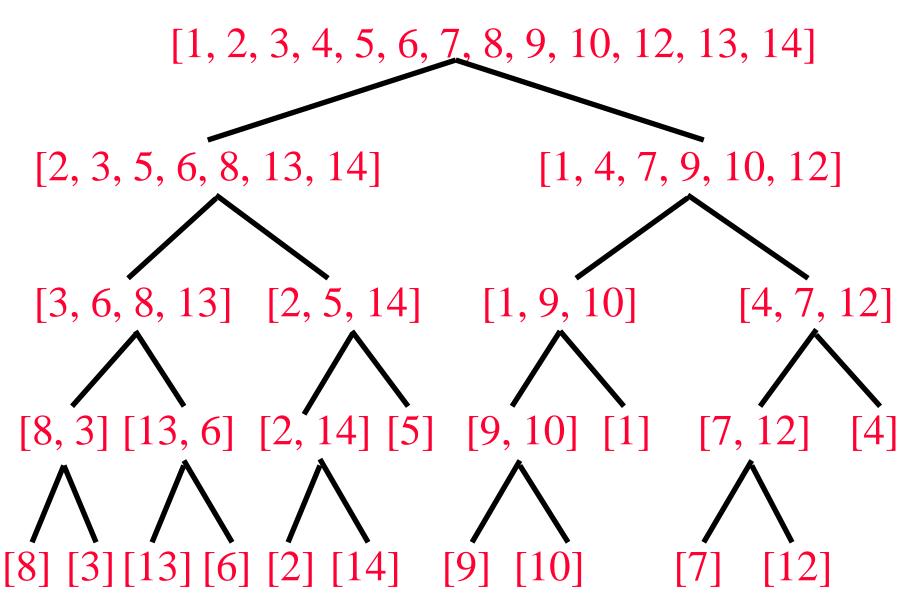
# **Merge Sort**





# **Merge Sort**





# **Time Complexity**



- Let t(n) be the time required to sort n elements.
- •t(0) = t(1) = c, where c is a constant.
- •When n > 1,
  - $\bullet t(n) = t(ceil(n/2)) + t(floor(n/2)) + dn,$
  - where d is a constant.
- To solve the recurrence, assume n is a power of 2 and use repeated substitution.
- •t(n) = O(n log n).

# **Merge Sort**



- Downward pass over the recursion tree.
  - Divide large instances into small ones.
- Upward pass over the recursion tree.
  - Merge pairs of sorted lists.
- Number of leaf nodes is n.
- Number of nonleaf nodes is n-1.

# **Time Complexity**



- Downward pass.
  - •O(1) time at each node.
  - O(n) total time at all nodes.
- Upward pass.
  - O(n) time merging at each level that has a nonleaf node.
  - Number of levels is O(log n).
  - Total time is O(n log n).

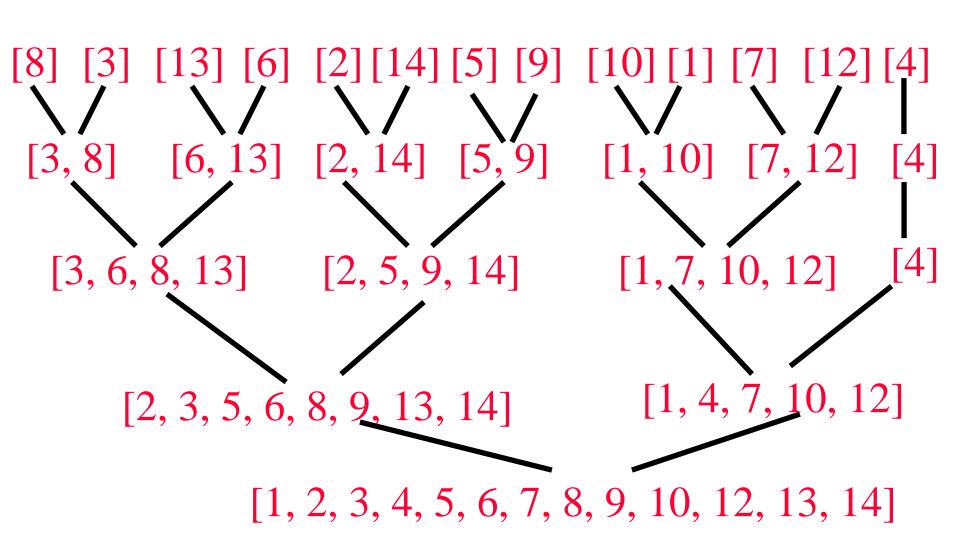
#### **Nonrecursive Version**



- Eliminate downward pass.
- Start with sorted lists of size 1 and do pairwise merging of these sorted lists as in the upward pass.

# **Nonrecursive Merge Sort**





# Complexity



- Sorted segment size is 1, 2, 4, 8, ...
- Number of merge passes is ceil(log<sub>2</sub>n).
- Each merge pass takes O(n) time.
- Total time is O(n log n).
- Need O(n) additional space for the merge.
- Merge sort is slower than insertion sort when n <= 15 (approximately). So define a small instance to be an instance with n <= 15.</li>
- Sort small instances using insertion sort.
- Start with segment size = 15.

# **Natural Merge Sort**



- Initial sorted segments are the naturally occurring sorted segments in the input.
- Input = [8, 9, 10, 2, 5, 7, 9, 11, 13, 15, 6, 12, 14].
- Initial segments are:
  - [8, 9, 10] [2, 5, 7, 9, 11, 13, 15] [6, 12, 14]
- 2 (instead of 4) merge passes suffice.
- Segment boundaries have a[i] > a[i+1].



```
// Merge Sort C Program
#include<stdlib.h>
#include<stdio.h>
int main()
{ //int arr[] = {85, 24, 63, 45, 17, 31, 96, 50};
int arr[] = \{-5,406,3,425,317,31,6,50,55,54,66,78,99,43,72\};
int arr_size = sizeof(arr)/sizeof(arr[0]);
 printf("\n Given array is : ");
 printArray(arr, arr_size);
 mergeSort(arr, 0, arr size - 1);
 printf("\n Sorted array is : ");
 printArray(arr, arr_size);
```



```
void printArray(int A[], int size)
{ int i;
 for (i=0; i < size; i++)
    printf("%d ", A[i]);
 printf("\n");
void mergeSort(int arr[], int I, int r)
\{ if (l < r) \}
 \{ int m = l+(r-l)/2; \}
  mergeSort(arr, I, m);
  mergeSort(arr, m+1, r);
  merge(arr, I, m, r);
```



```
// Merge Function
void merge(int arr[], int I, int m, int r)
{ int i, j, k;
 int n1 = m - l + 1;
 int n2 = r - m;
 int L[n1], R[n2]; // Create temp arrays
 for (i = 0; i < n1; i++)
   L[i] = arr[l + i]; // Copy data to temp array
 for (j = 0; j < n2; j++)
   R[j] = arr[m + 1 + j];
```



```
// Merge temp arrays
i = 0;
j = 0;
k = I;
while (i < n1 && j < n2)
{ if (L[i] <= R[j])
 { arr[k] = L[i];
  i++;
 else
  { arr[k] = R[j];
   j++;
 k++;
```



```
while (i < n1) // Copy the remaining elements of L[]
{ arr[k] = L[i];
 i++;
 k++;
while (j < n2)// Copy the remaining elements of R[]
{ arr[k] = R[j];
 j++;
 k++;
```



Given array is: 85 24 63 45 17 31 96 50

Sorted array is: 17 24 31 45 50 63 85 96

Given array is: -5 406 3 425 317 31 6 50 55 54 66 78 99 43 72

Sorted array is: -5 3 6 31 43 50 54 55 66 72 78 99 317 406 425



Given array is: -55 -44 -33 0 3 6 10 15 24 26 37 49 53 66 70 72

Sorted array is: -55 -44 -33 0 3 6 10 15 24 26 37 49 53 66 70 72

Given array is: 555 333 244 100 93 86 60 55 44 36 27 25 23 16 7 2

Sorted array is: 2 7 16 23 25 27 36 44 55 60 86 93 100 244 333 555