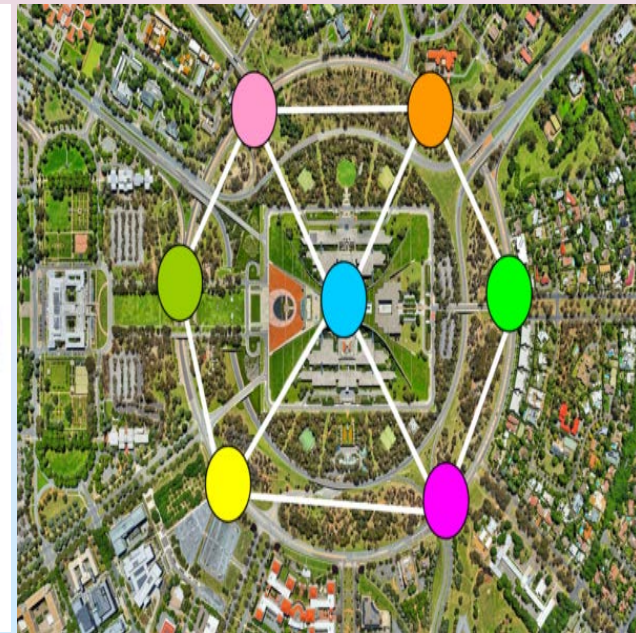
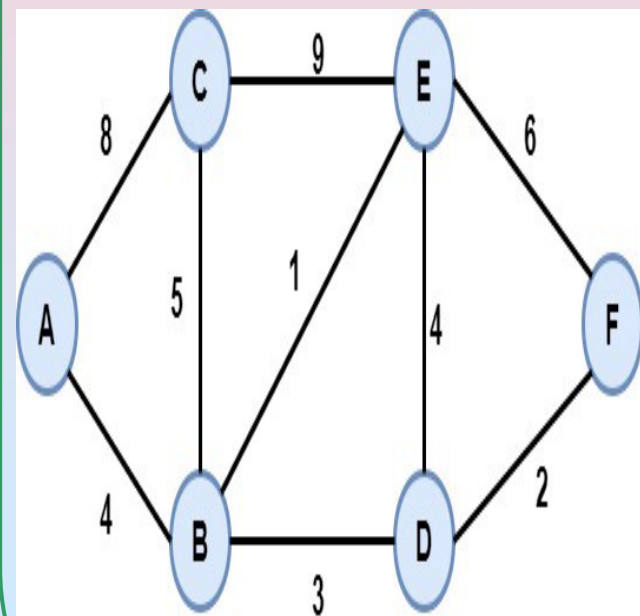
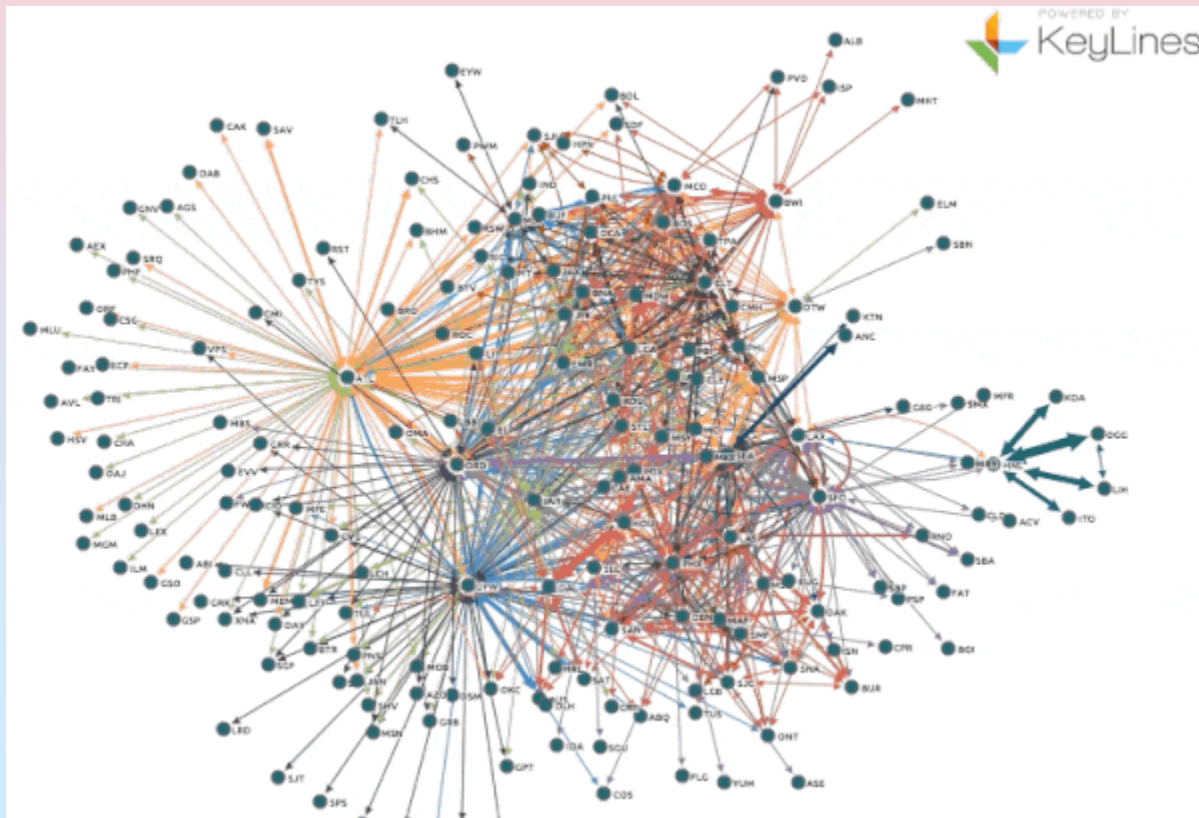


Graphs



Graphs



Graphs

- ❖ $G = (V, E)$
- ❖ V is the vertex set.
- ❖ Vertices are also called nodes and points.
- ❖ E is the edge set.
- ❖ Each edge connects two different vertices.
- ❖ Edges are also called arcs and lines.
- ❖ Directed edge has an orientation (u, v) .



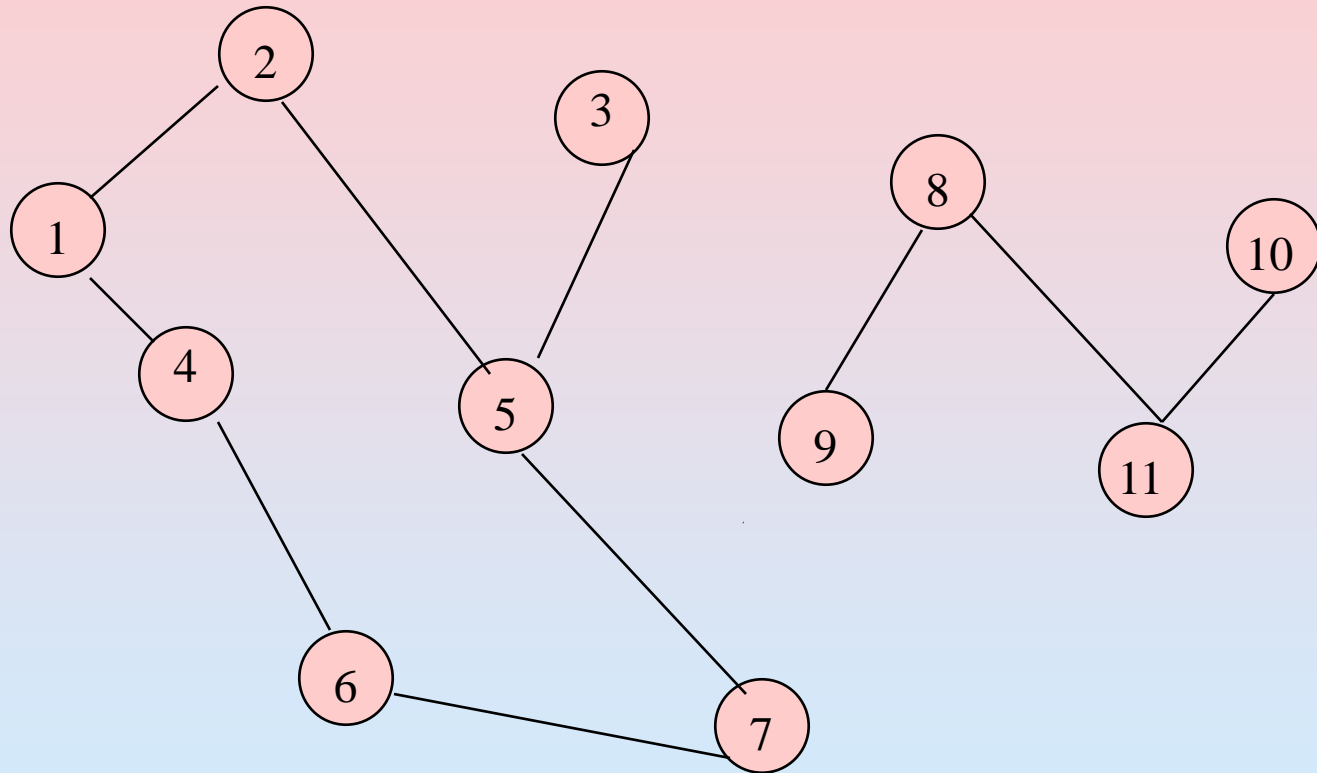
Graphs

❖ Undirected edge has no orientation (u,v).

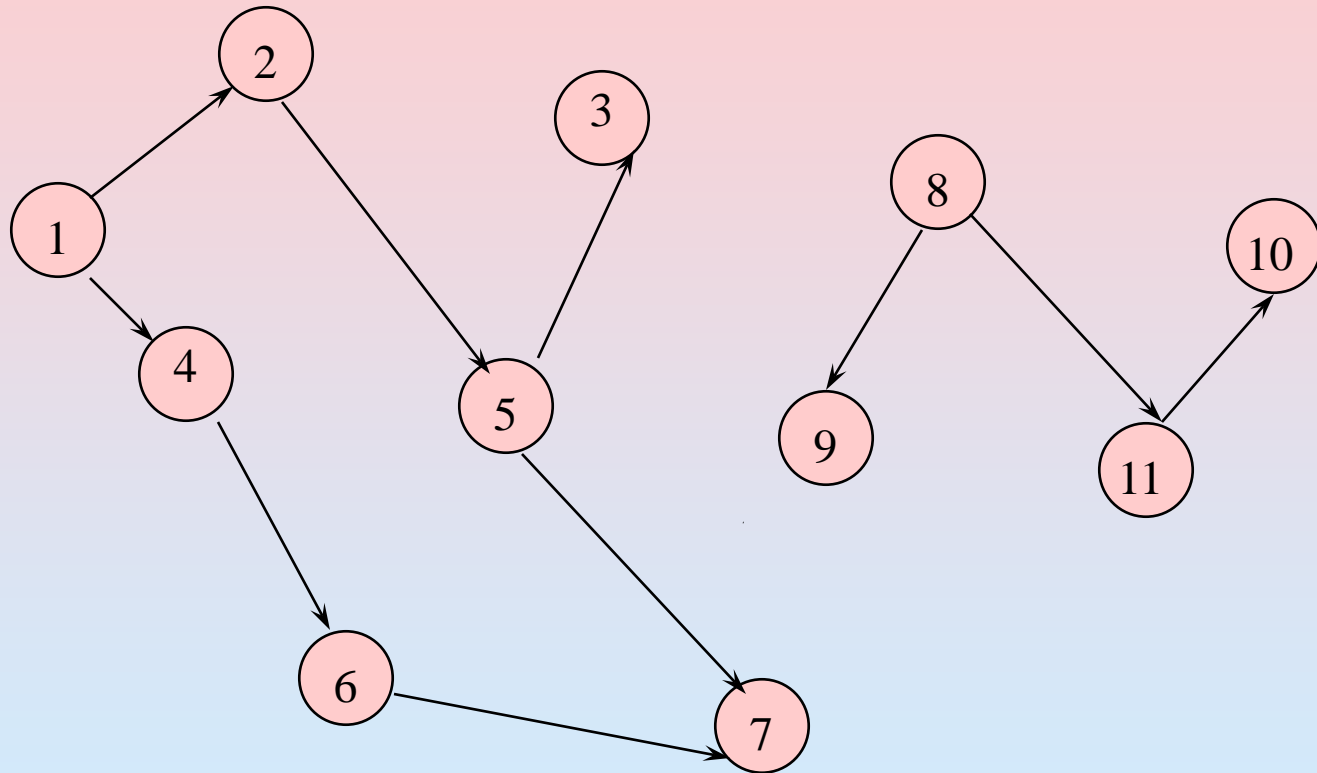
u — v

- Undirected graph \Rightarrow no oriented edge.
- Directed graph \Rightarrow every edge has an orientation.

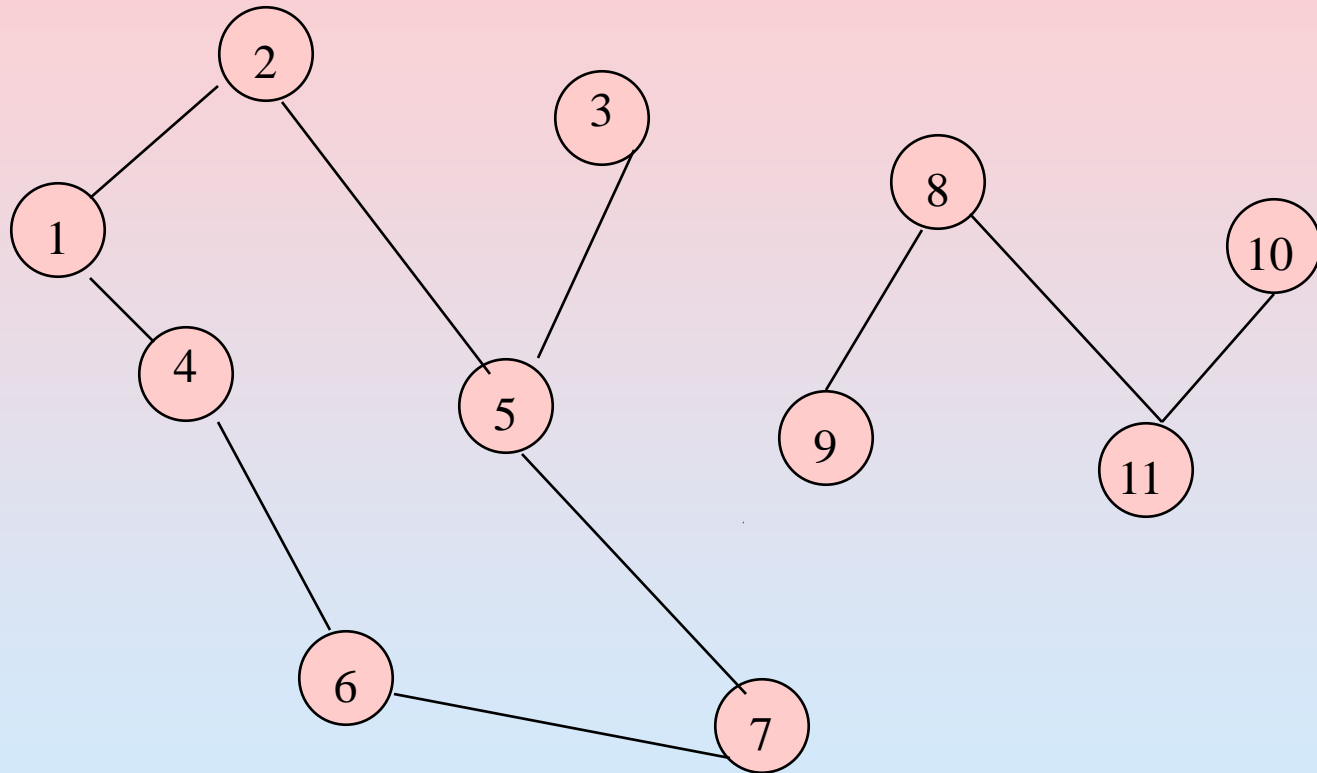
Undirected Graph



Directed Graph (Digraph)

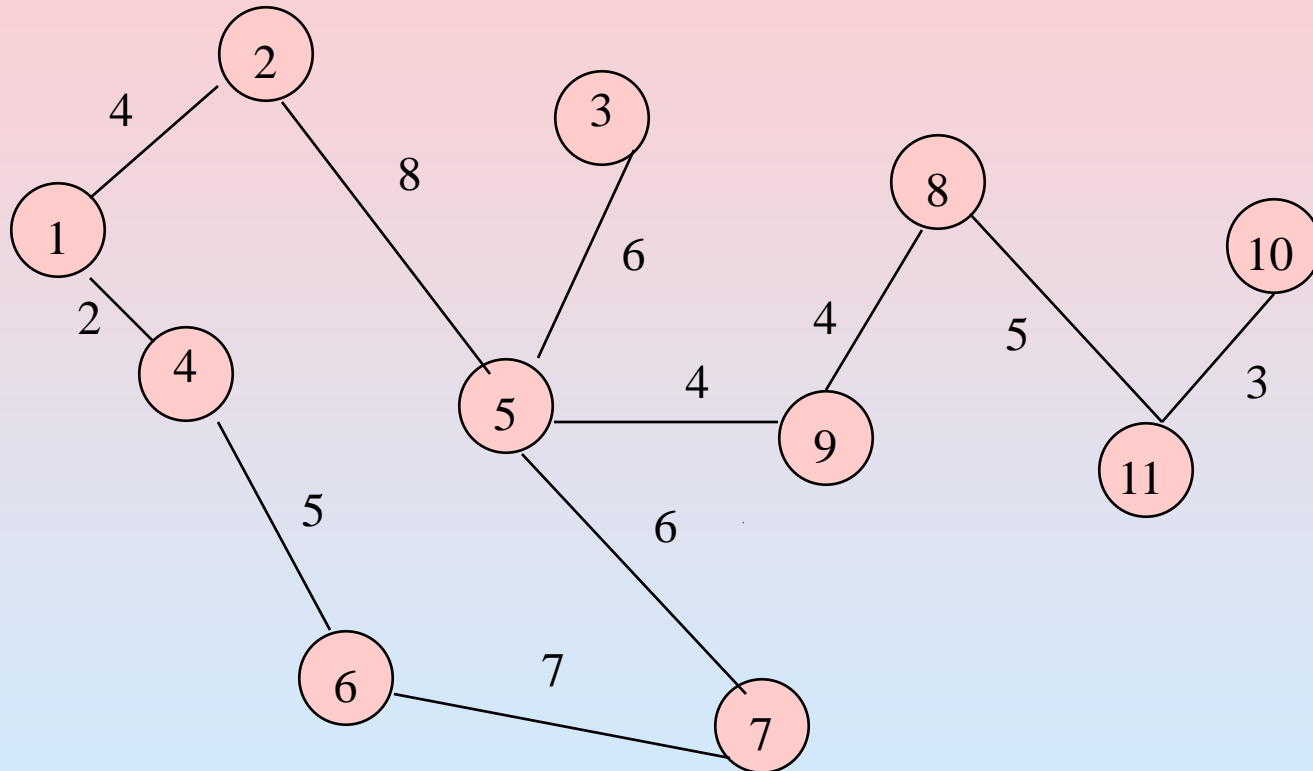


Applications—Communication Network



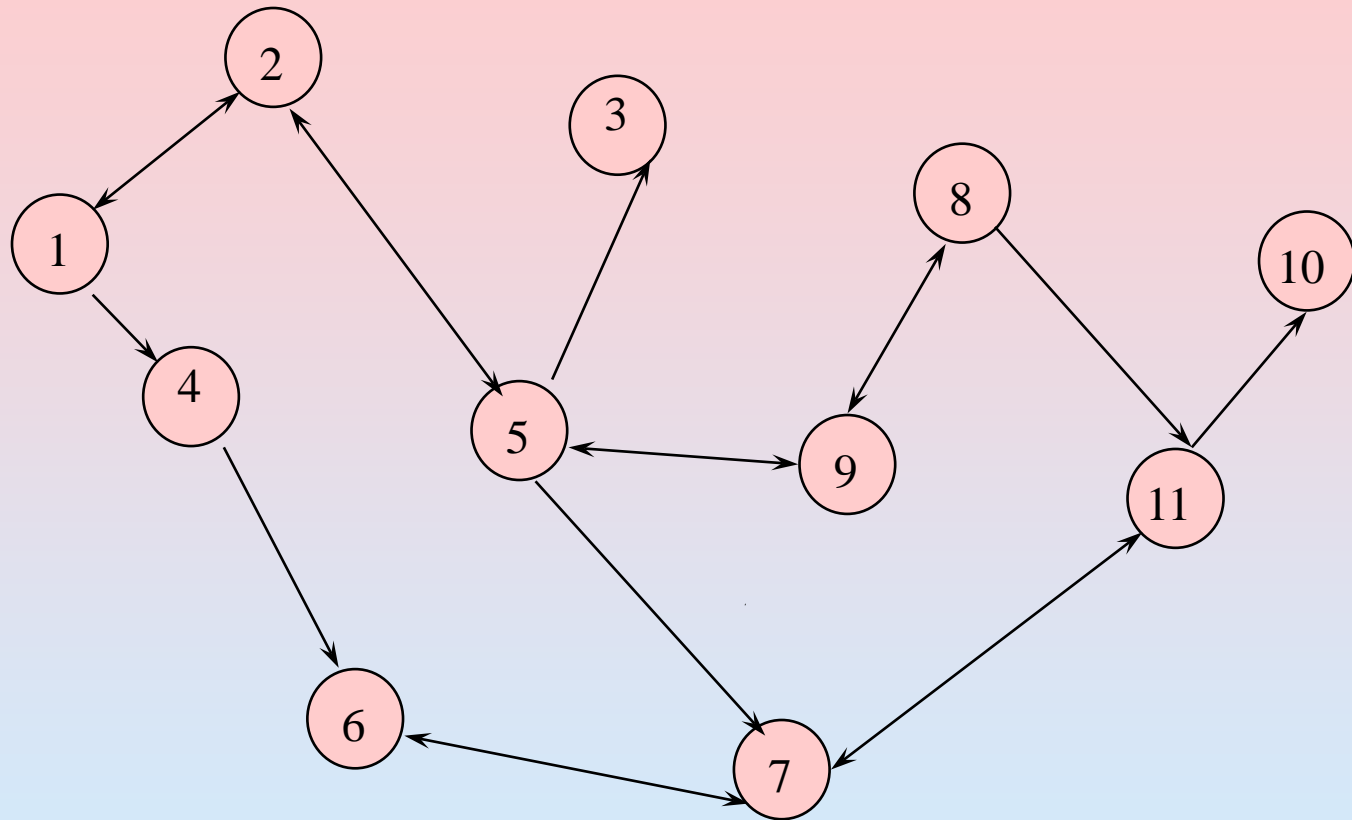
❖ **Vertex = city, edge = communication link.**

Driving Distance/Time Map



❖ Vertex = city, edge weight = driving distance/time.

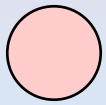
Street Map



❖ **Some streets are one way.**

Complete Undirected Graph

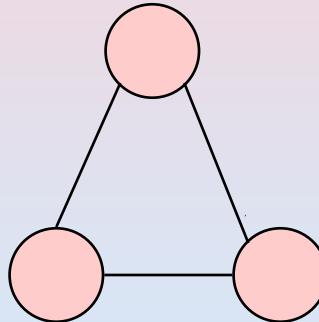
Has all possible edges.



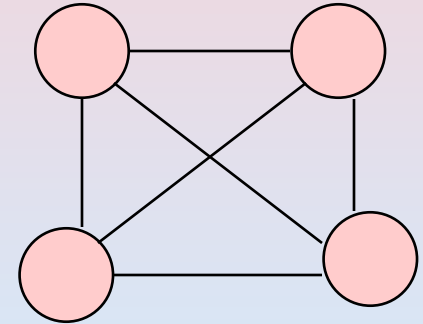
$n = 1$



$n = 2$



$n = 3$



$n = 4$

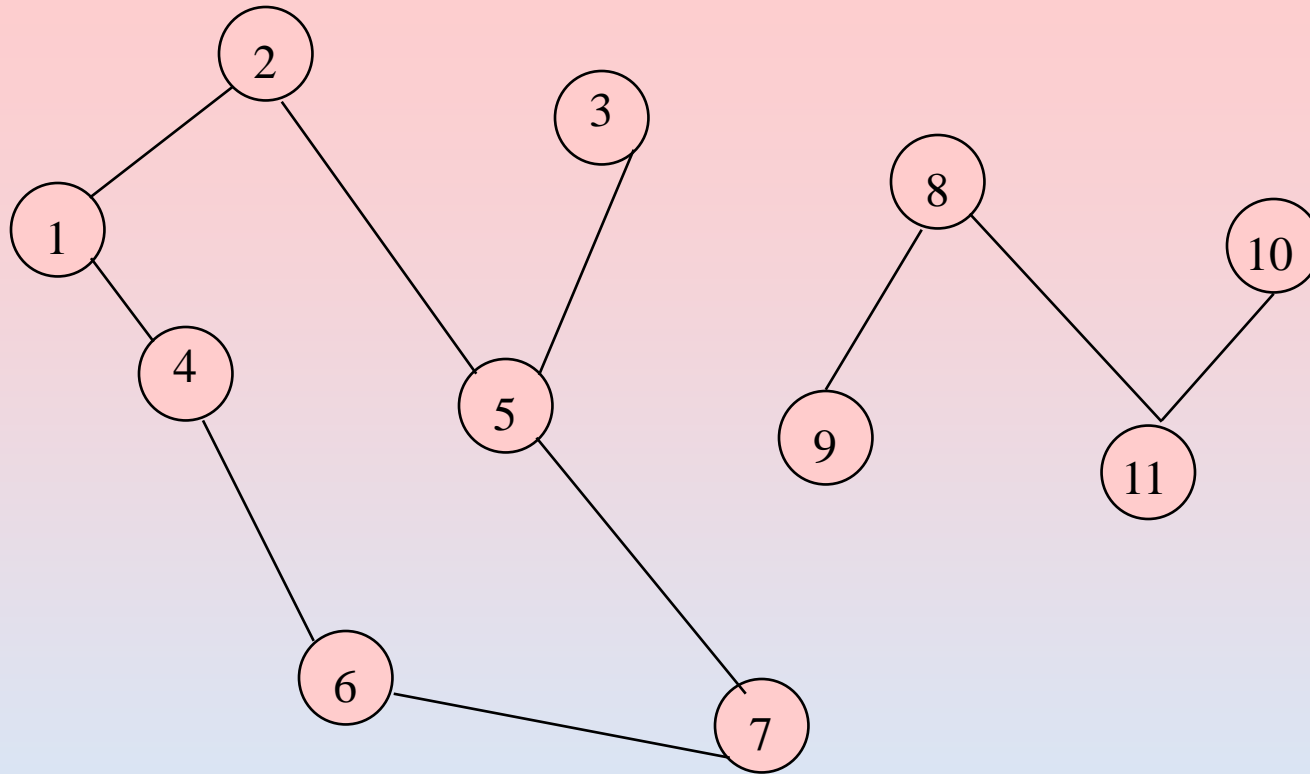
Number Of Edges—Undirected Graph

- ❖ Each edge is of the form (u, v) , $u \neq v$.
- ❖ Number of such pairs in an n vertex graph is $n(n-1)$.
- ❖ Since edge (u, v) is the same as edge (v, u) , the number of edges in a complete undirected graph is $n(n-1)/2$.
- ❖ Number of edges in an undirected graph is $\leq n(n-1)/2$.

Number Of Edges—Directed Graph

- ❖ Each edge is of the form (u, v) , $u \neq v$.
- ❖ Number of such pairs in an n vertex graph is $n(n-1)$.
- ❖ Since edge (u, v) is **not** the same as edge (v, u) , the number of edges in a complete directed graph is $n(n-1)$.
- ❖ Number of edges in a directed graph is $\leq n(n-1)$.

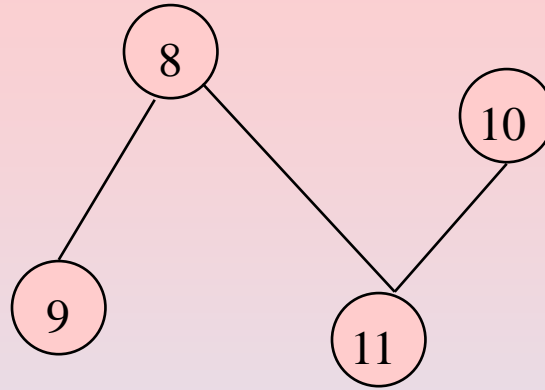
Vertex Degree



Number of edges incident to vertex.

$\text{degree}(2) = 2$, $\text{degree}(5) = 3$, $\text{degree}(3) = 1$

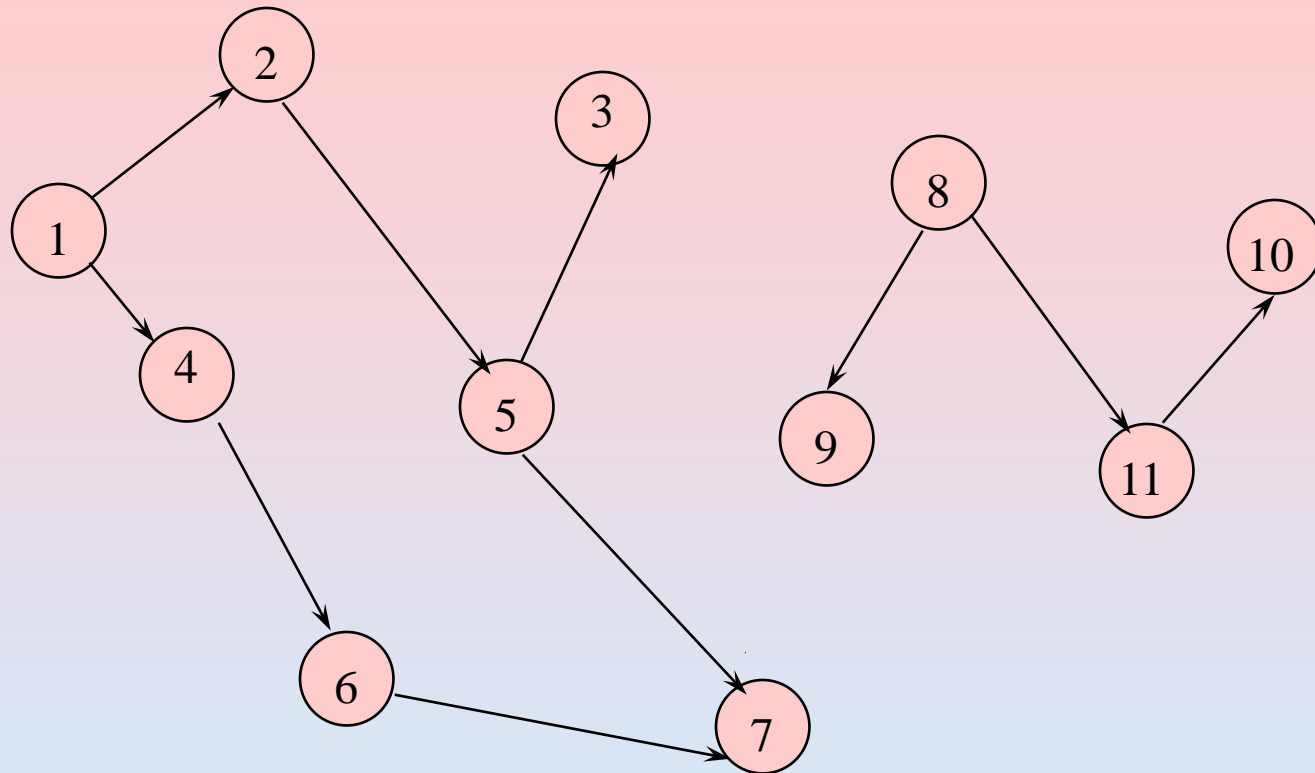
Sum Of Vertex Degrees



Sum of degrees = $2e$ (e is number of edges)

$$= 2 * 3 = 6$$

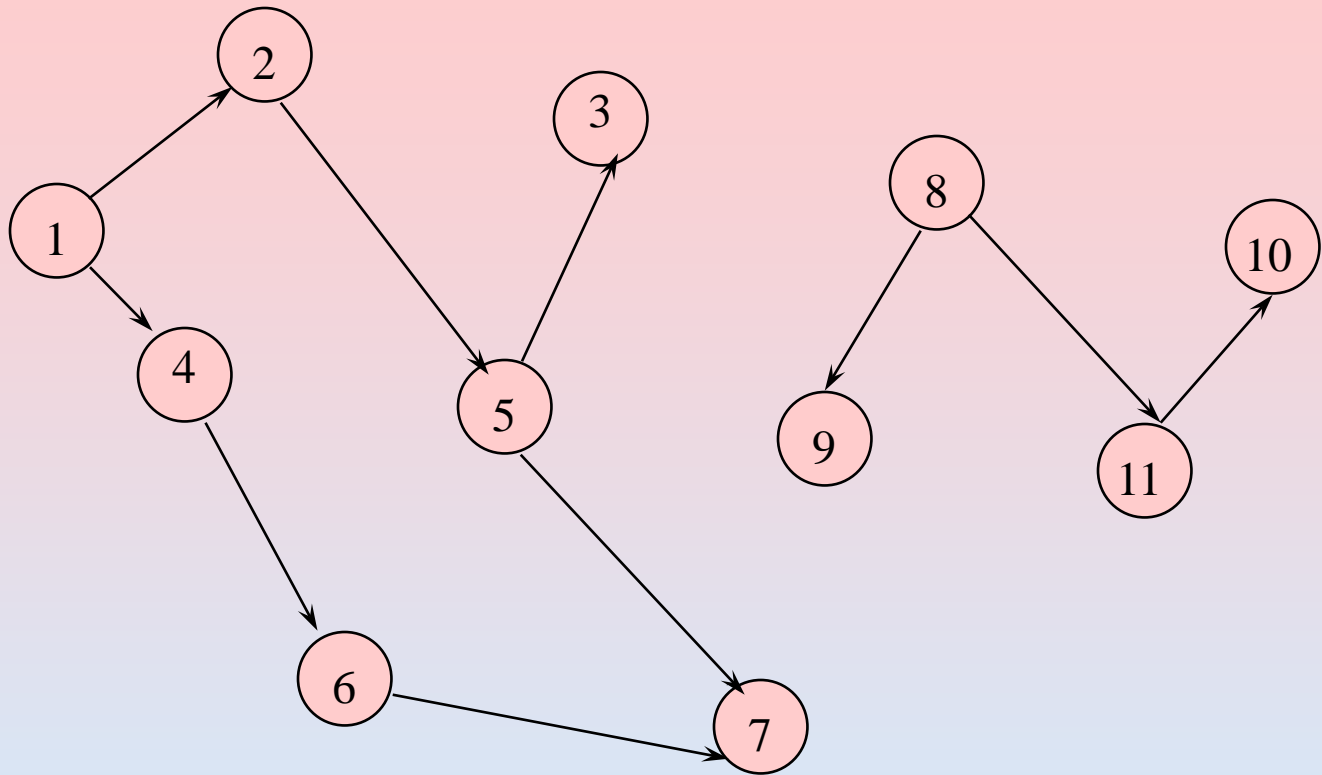
In-Degree Of A Vertex



In-degree is number of incoming edges

$\text{indegree}(2) = 1, \text{indegree}(8) = 0, \text{indegree}(7) = 2$

Out-Degree of a Vertex



out-degree is number of outbound edges

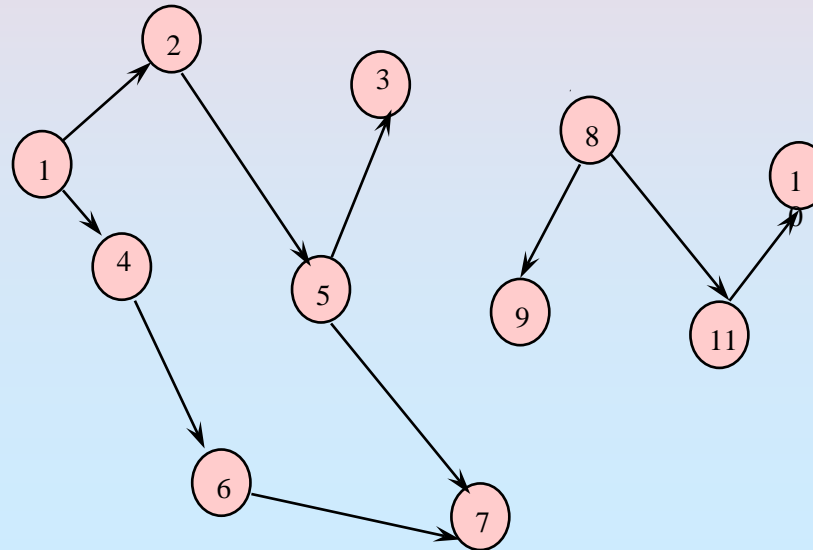
$\text{outdegree}(2) = 1$, $\text{outdegree}(8) = 2$, $\text{outdegree}(7) = 0$

Sum of In- and Out-Degrees

Each edge contributes **1** to the in-degree of some vertex and **1** to the out-degree of some other vertex

sum of in-degrees = sum of out-degrees = **e**,

where **e** is the number of edges in the digraph



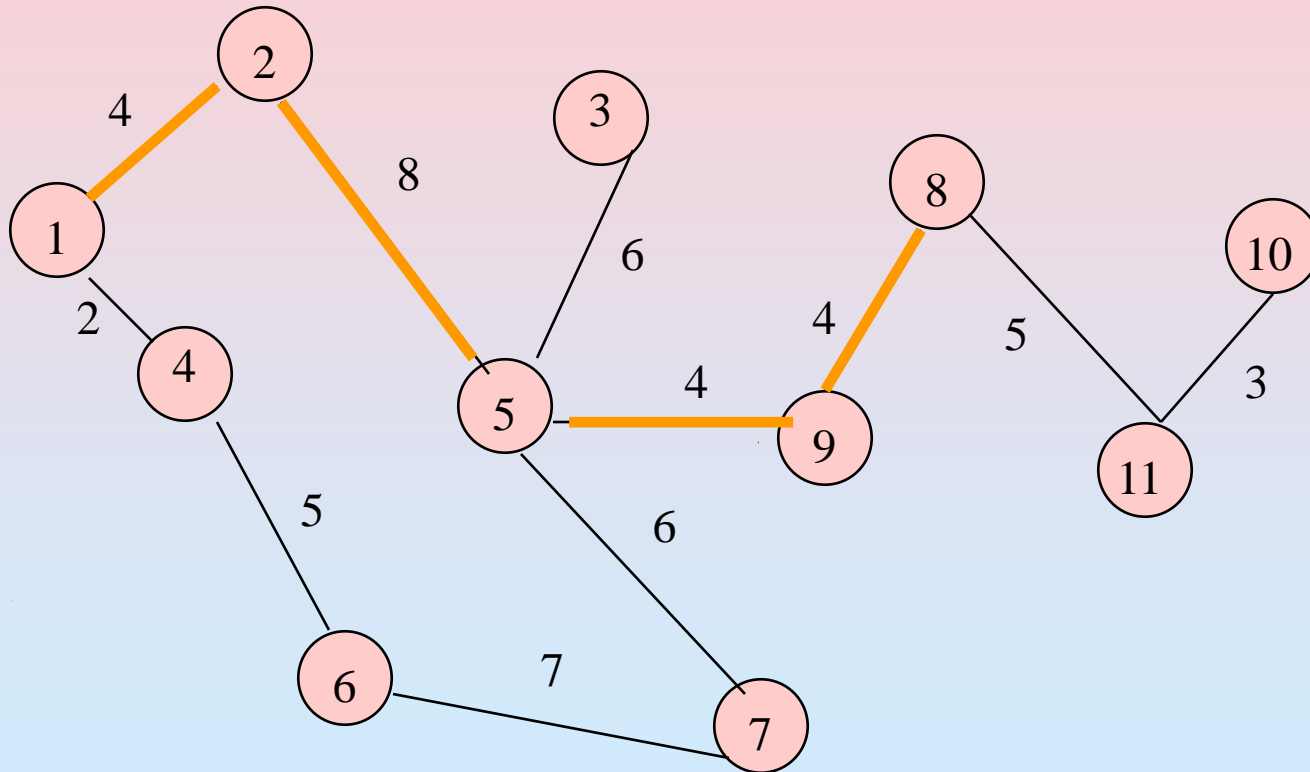
Graph Operations and Representation

Sample Graph Problems

- ❖ **Path problems.**
- ❖ **Connectedness problems.**
- ❖ **Spanning tree problems.**

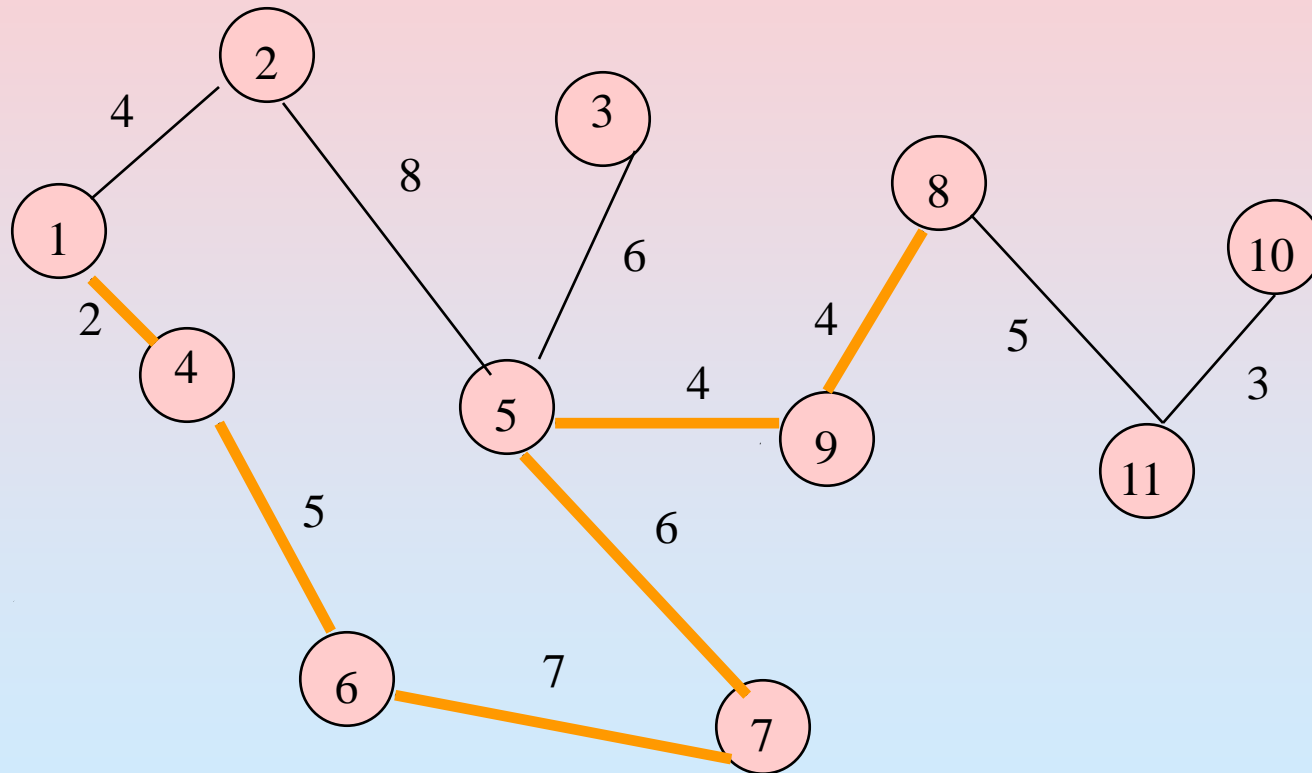
Path Finding

Path between 1 and 8.



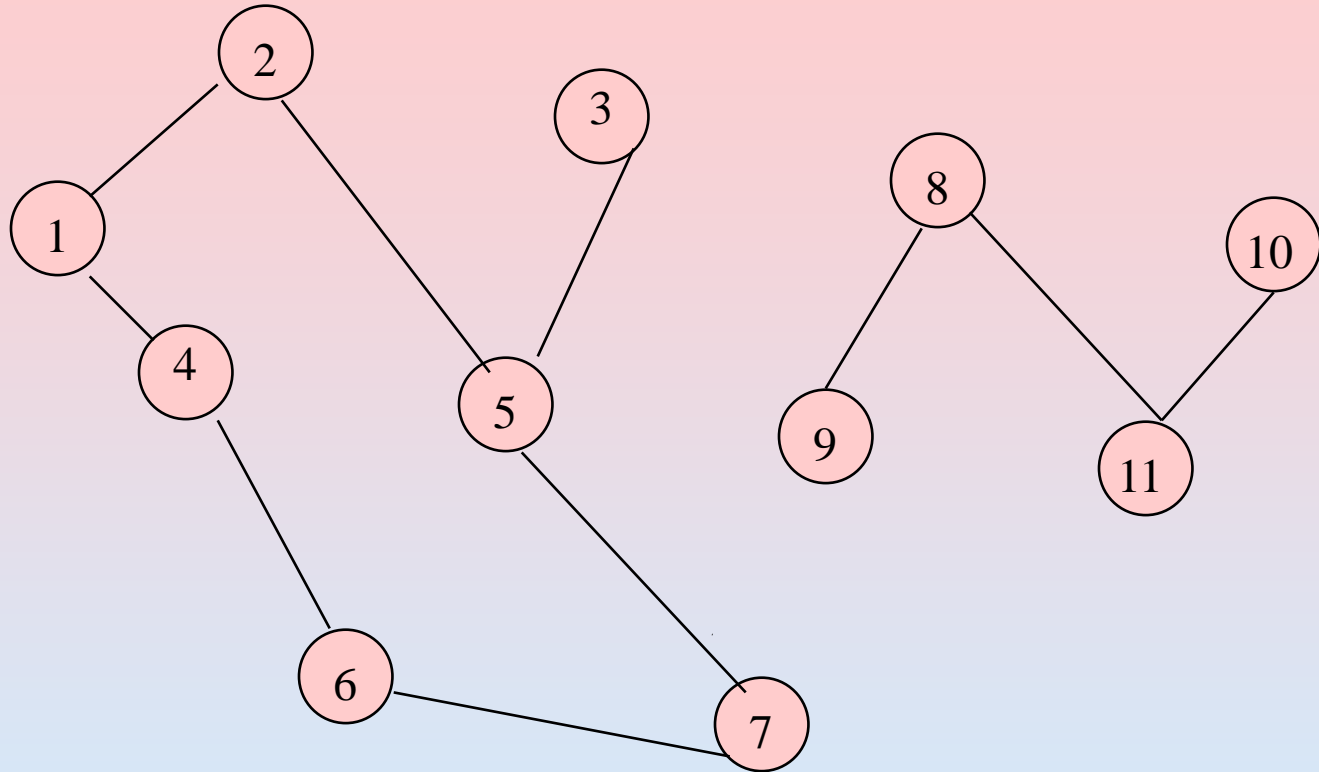
Path length is 20.

Another Path Between 1 and 8



Path length is 28.

Example of No Path



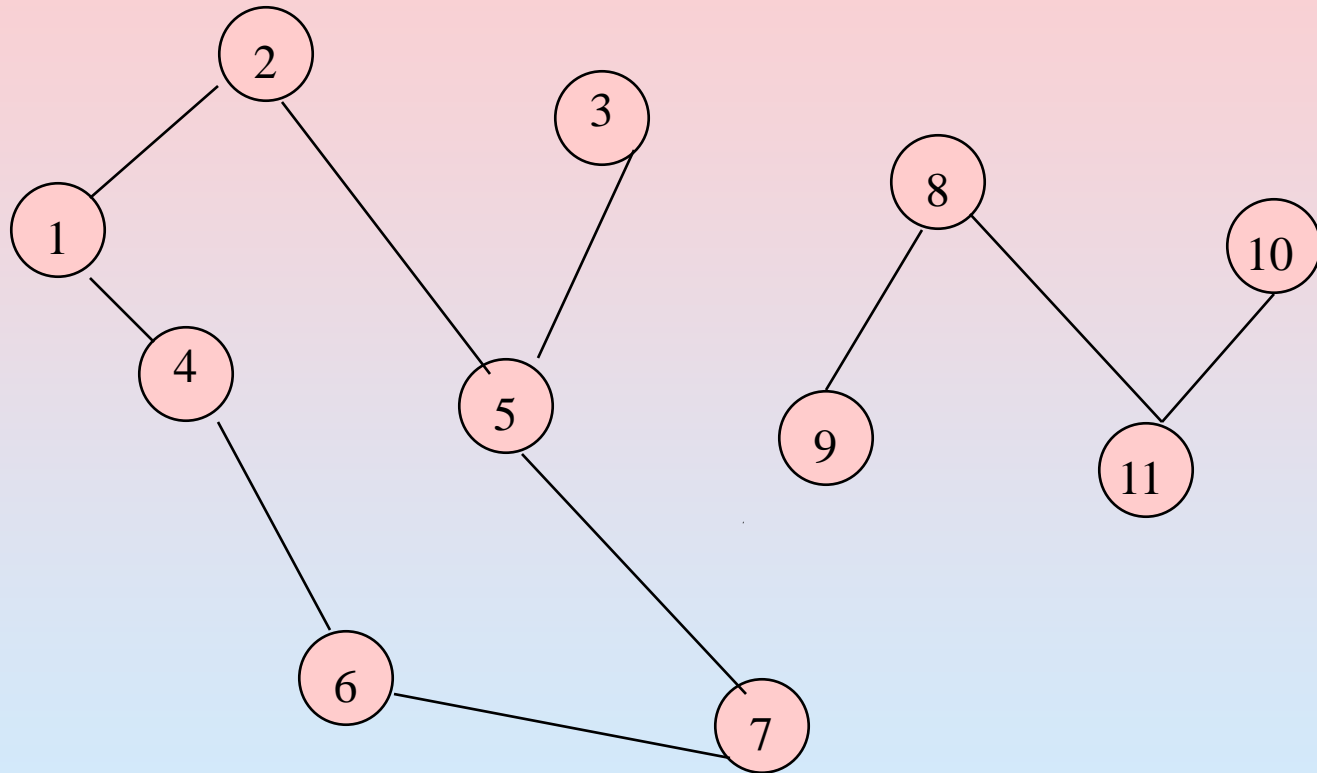
No path between 2 and 9.

Connected Graph

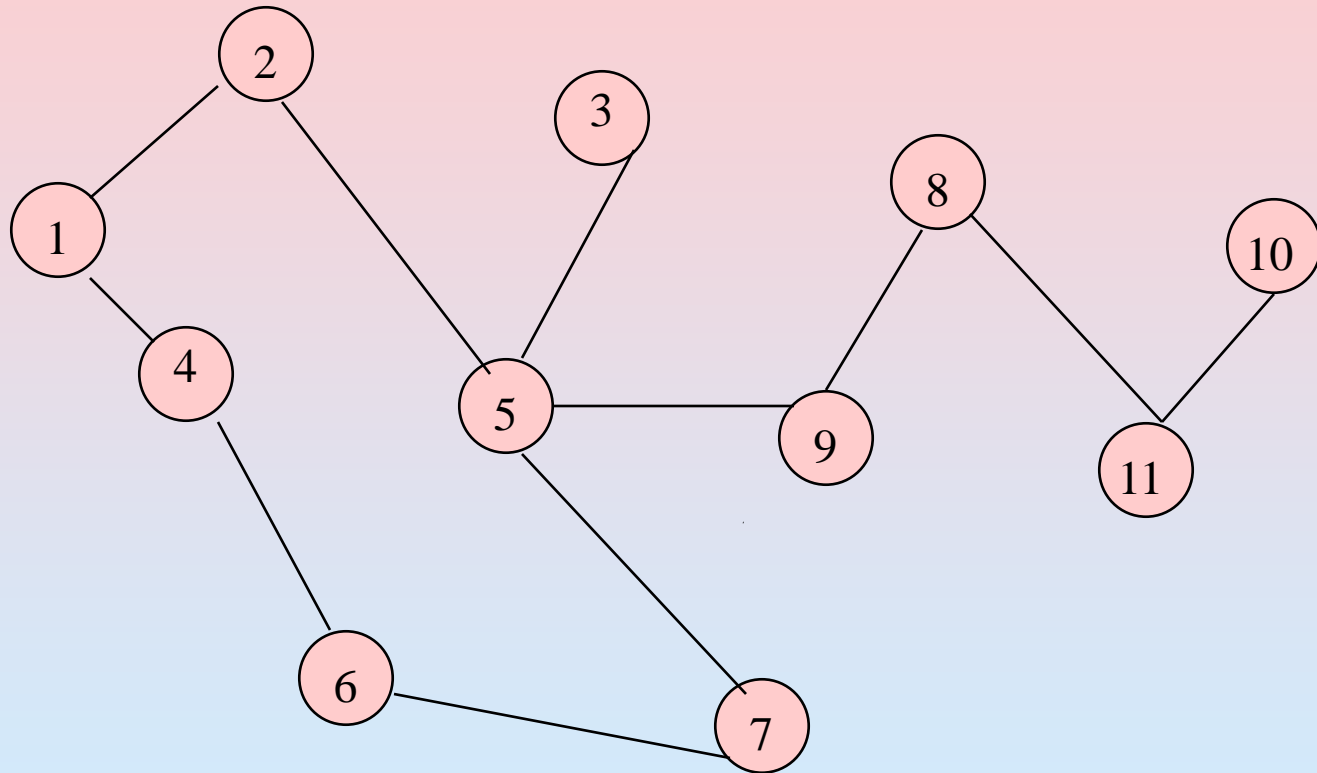
- ❖ **Undirected graph.**

- ❖ **There is a path between every pair of vertices.**

Example of Not Connected



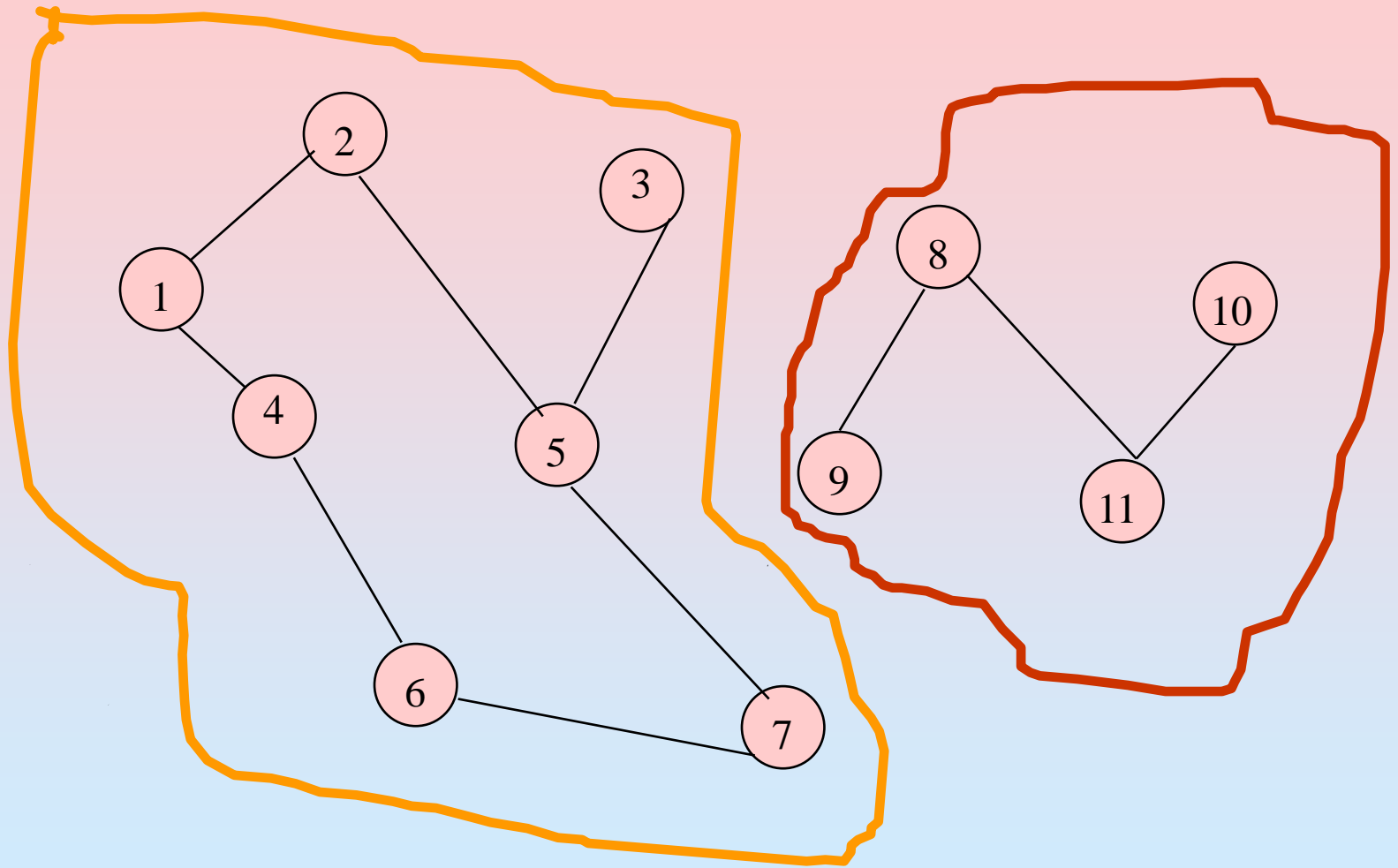
Connected Graph Example



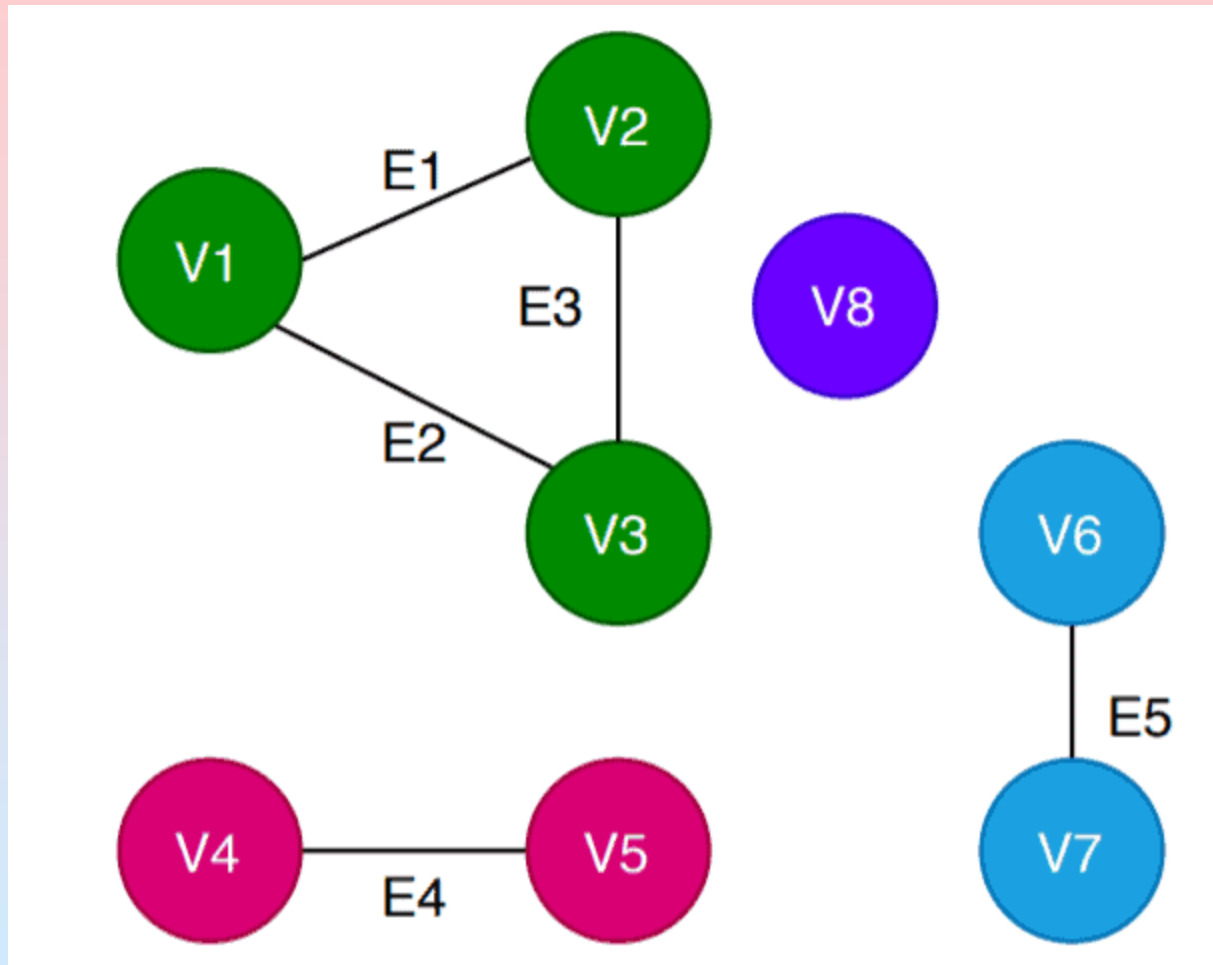
Connected Component

- ❖ A connected component is a set of vertices in a graph that are connected to each other.
- ❖ Inside a component, each vertex is reachable from every other vertex in that component.
- ❖ A maximal subgraph that is connected.
- ❖ A graph can have multiple connected components.
- ❖ A connected graph has exactly 1 component.

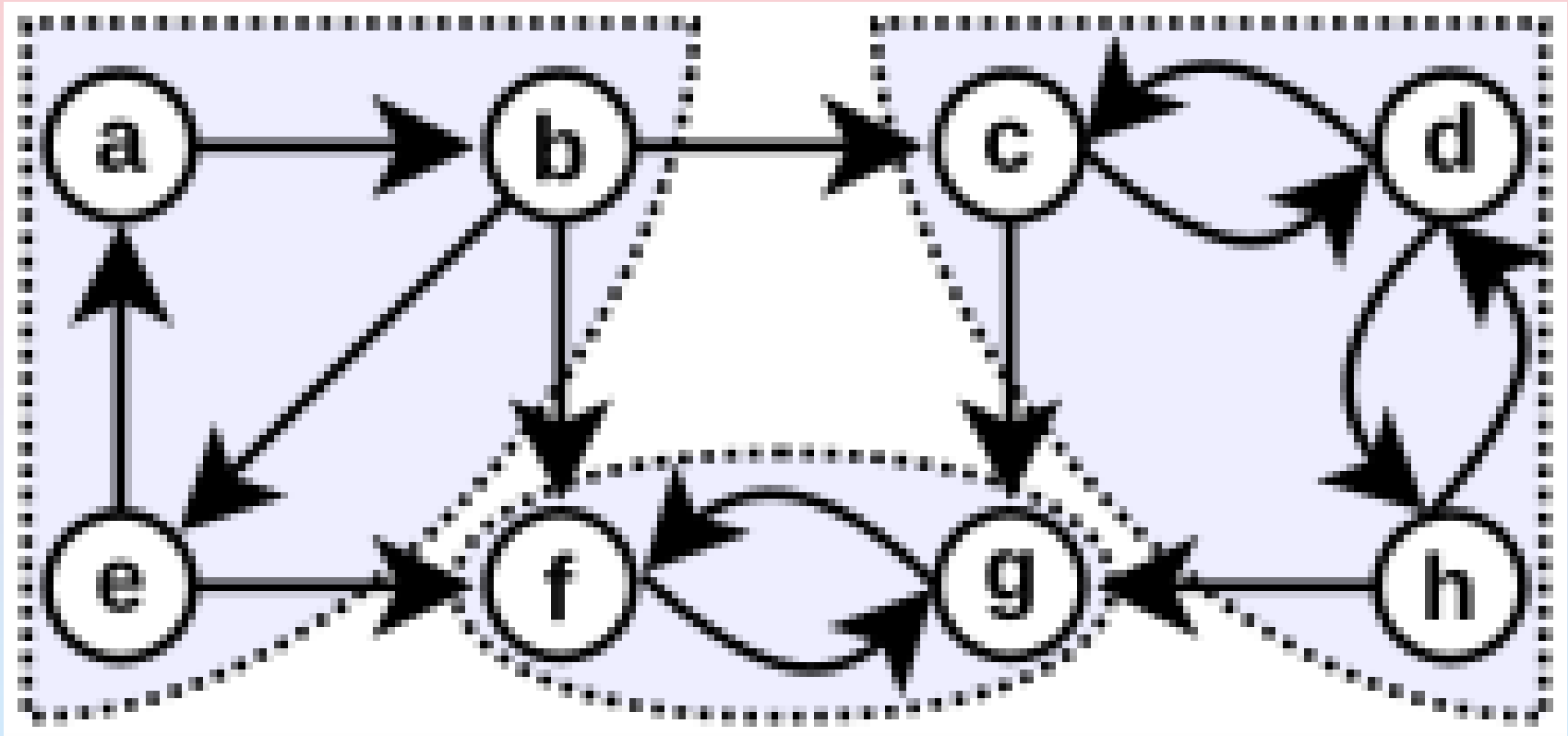
Connected components in the Graph - Example



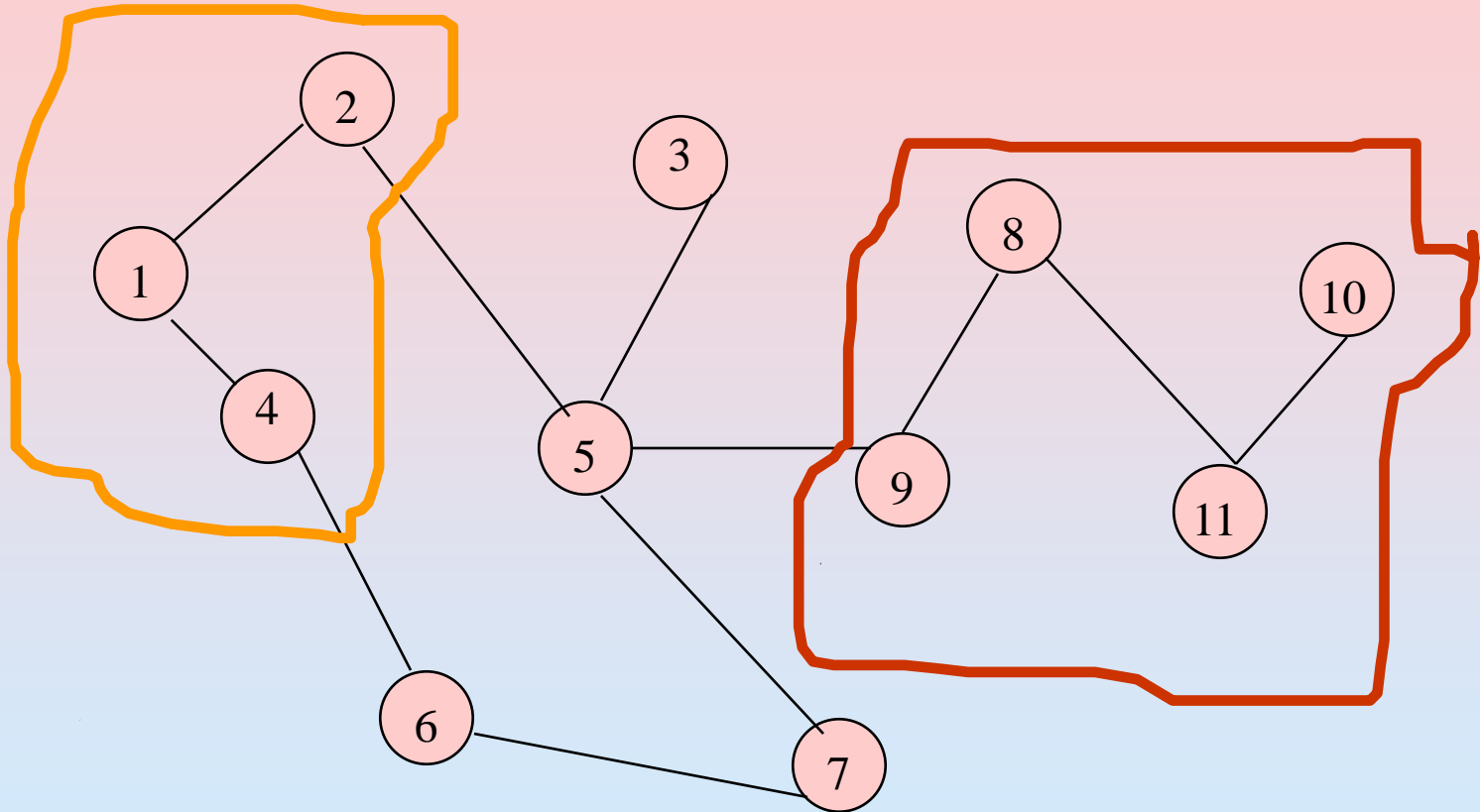
Connected components in the Graph - Example



Connected components in a Directed Graph - Example



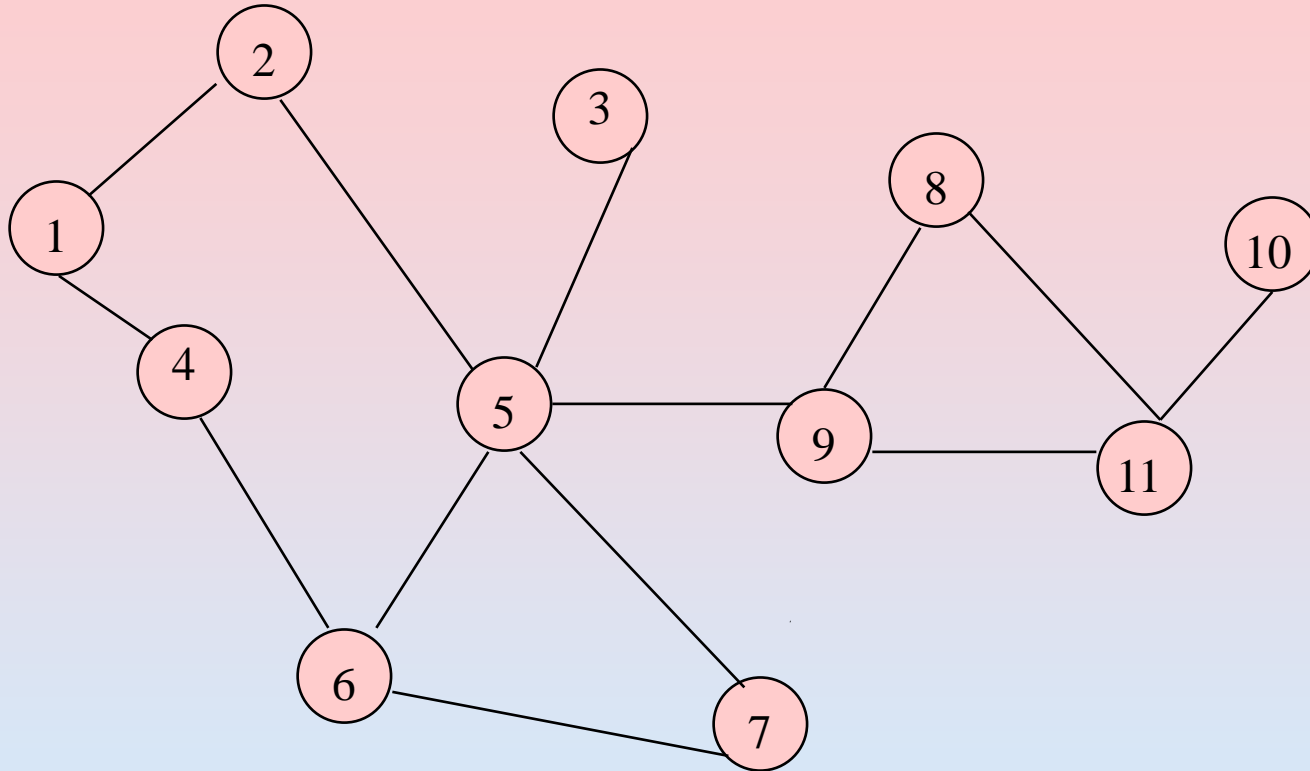
Not a Component



Applications of Connected Component

- **Graph Theory:** It is used to find subgraphs or clusters of nodes that are connected to each other.
- **Computer Networks:** It is used to discover clusters of nodes or devices that are linked and have similar qualities, such as bandwidth.
- **Image Processing:** Connected components also have usage in automated image analysis applications.

Communication Network



Each edge is a link that can be constructed (i.e., a feasible link).

Communication Network Problems

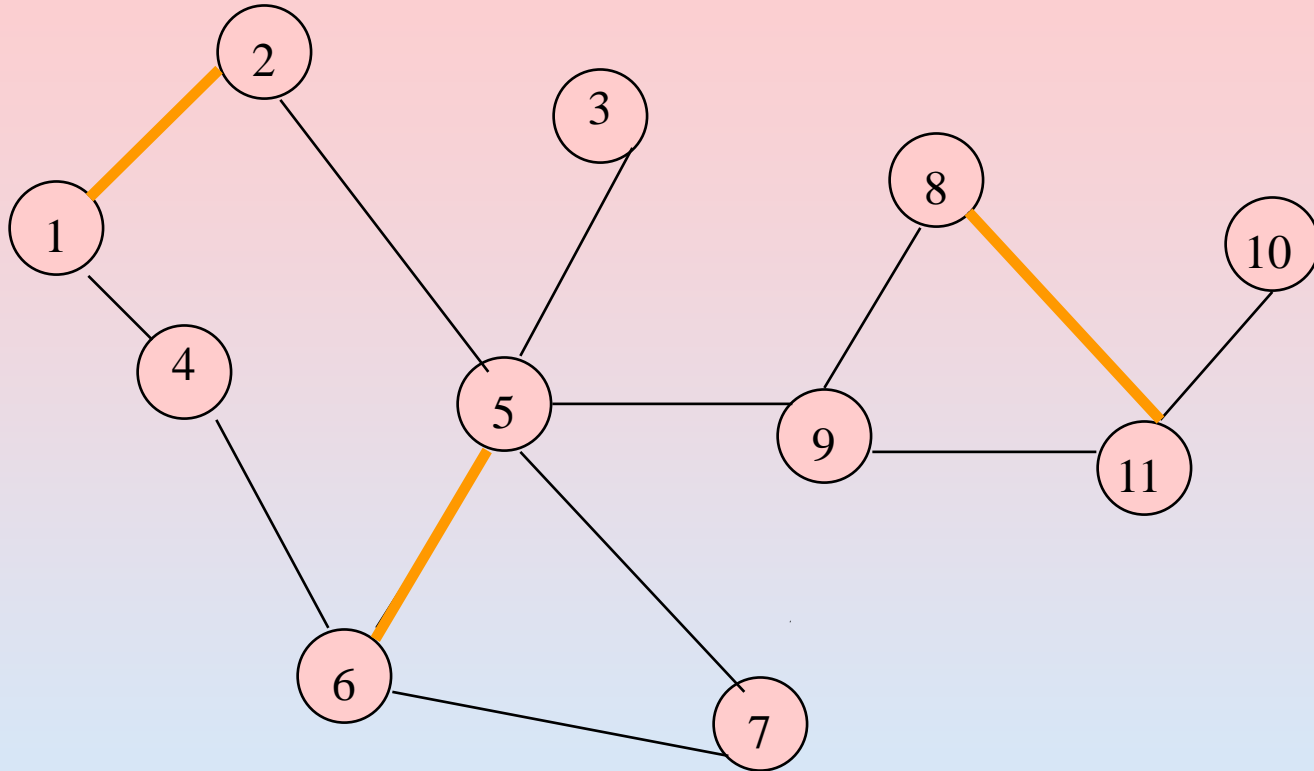
❖ **Is the network connected?**

- **Can we communicate between every pair of cities?**

❖ **Find the components.**

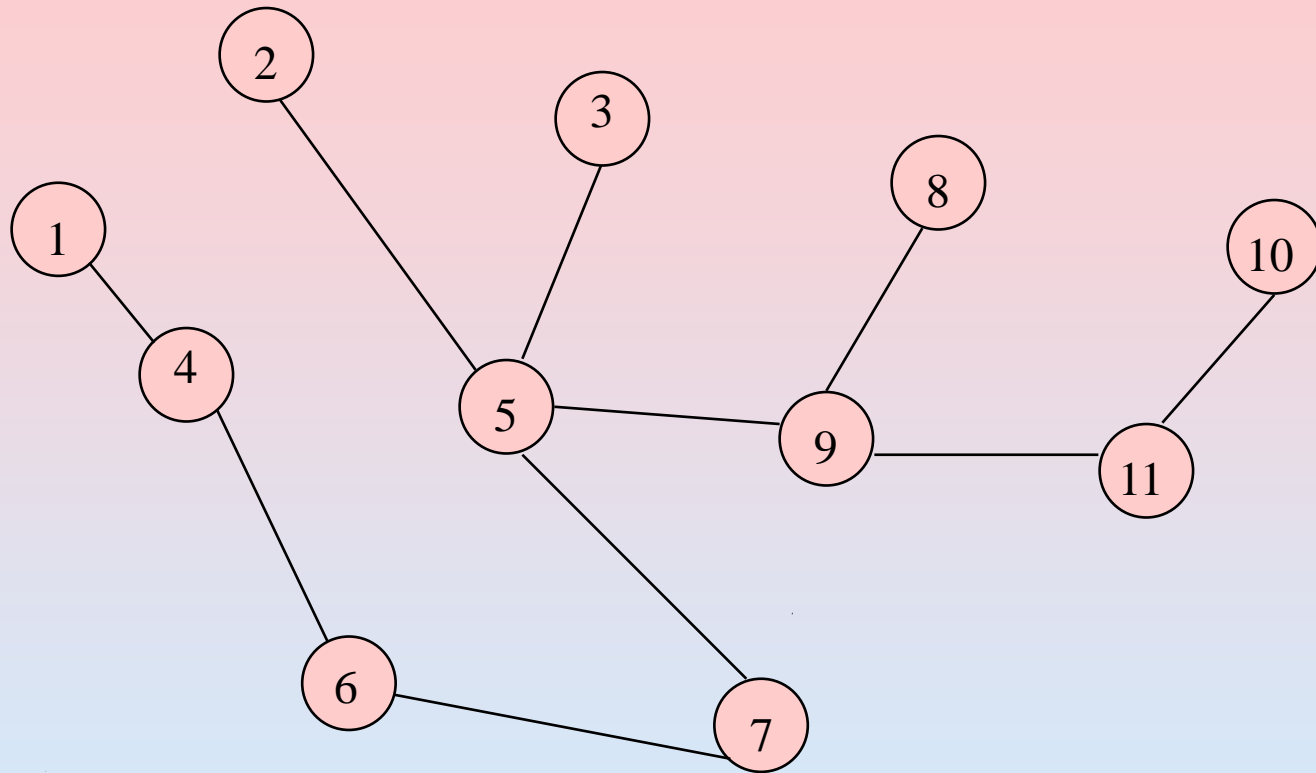
❖ **Want to construct smallest number of feasible links so that resulting network is connected.**

Cycles and Connectedness



Removal of an edge that is on a cycle does not affect connectedness.

Cycles and Connectedness



Connected subgraph with all vertices and minimum number of edges has no cycles.



Tree

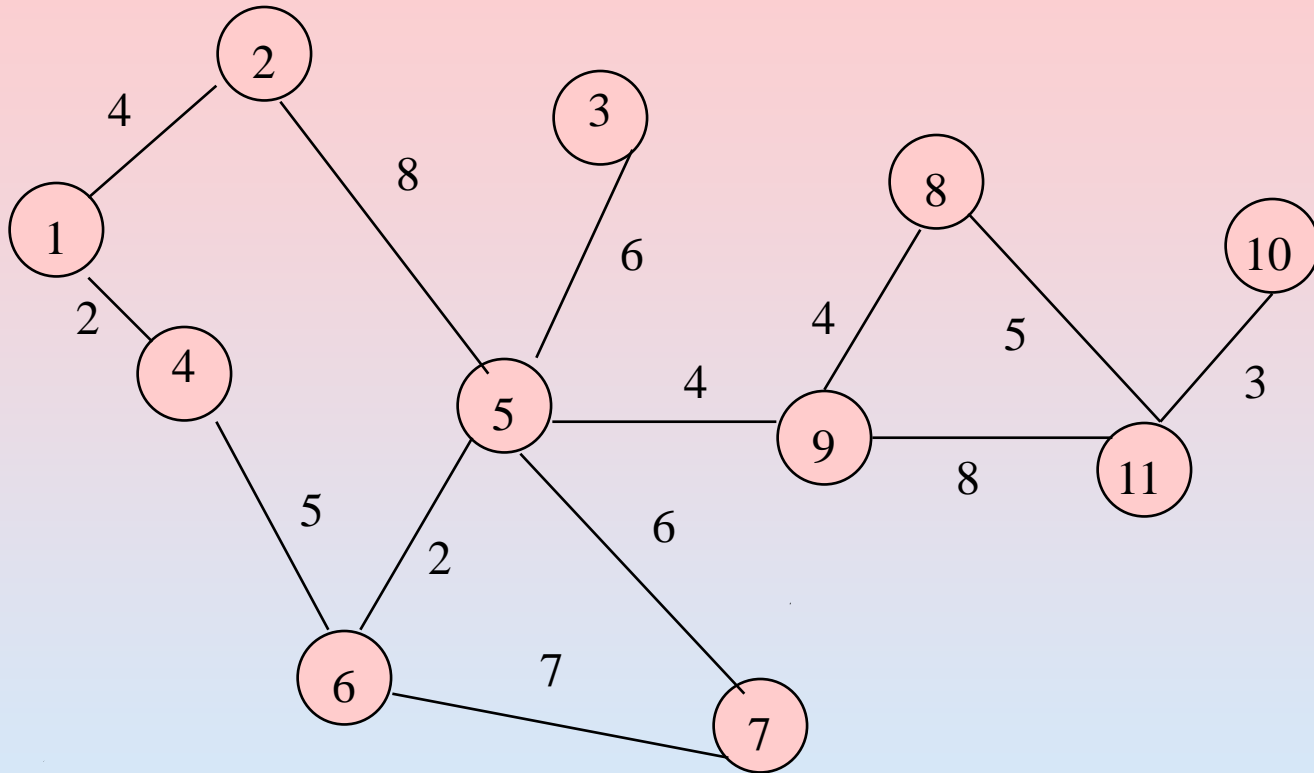


- ❖ **Connected graph that has no cycles.**
- ❖ **n vertex connected graph with $n-1$ edges.**

Spanning Tree

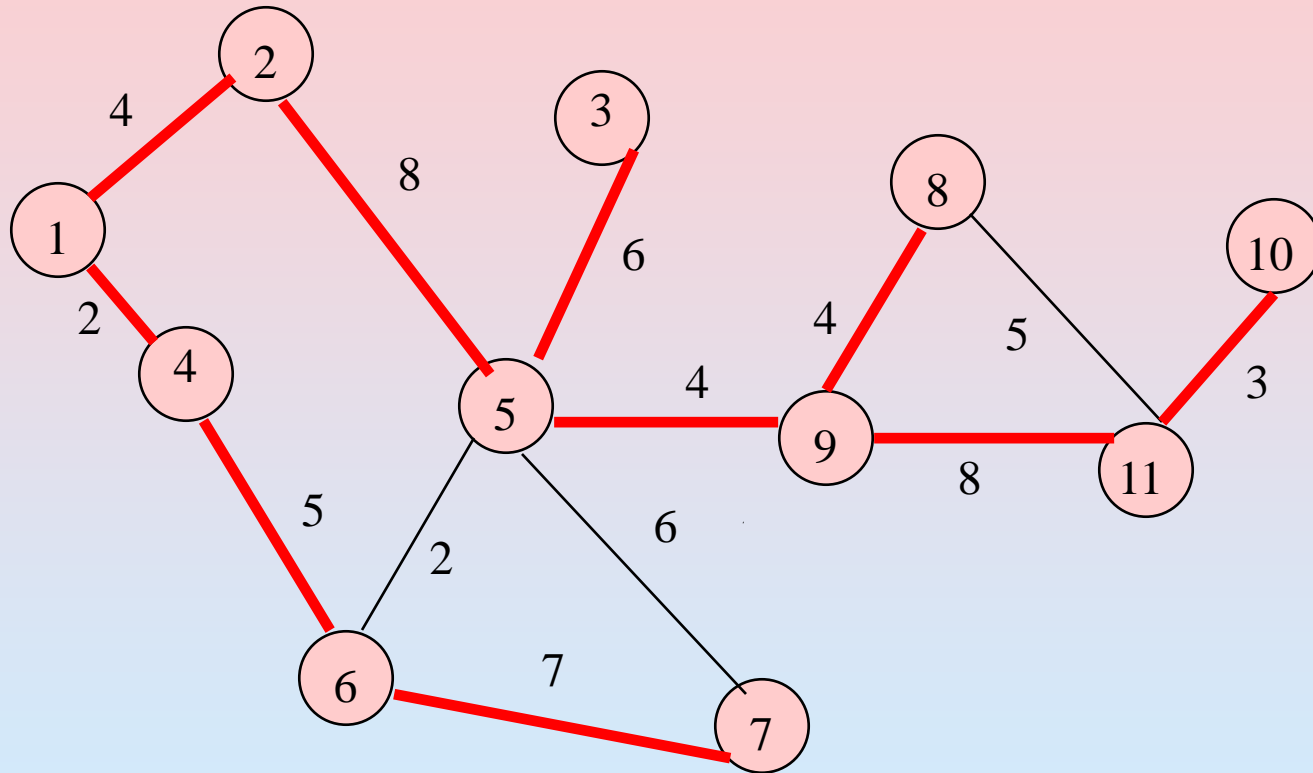
- ❖ Subgraph that includes all vertices of the original graph.
- ❖ Subgraph is a tree.
 - If original graph has n vertices, the spanning tree has n vertices and $n-1$ edges.

Minimum Cost Spanning Tree



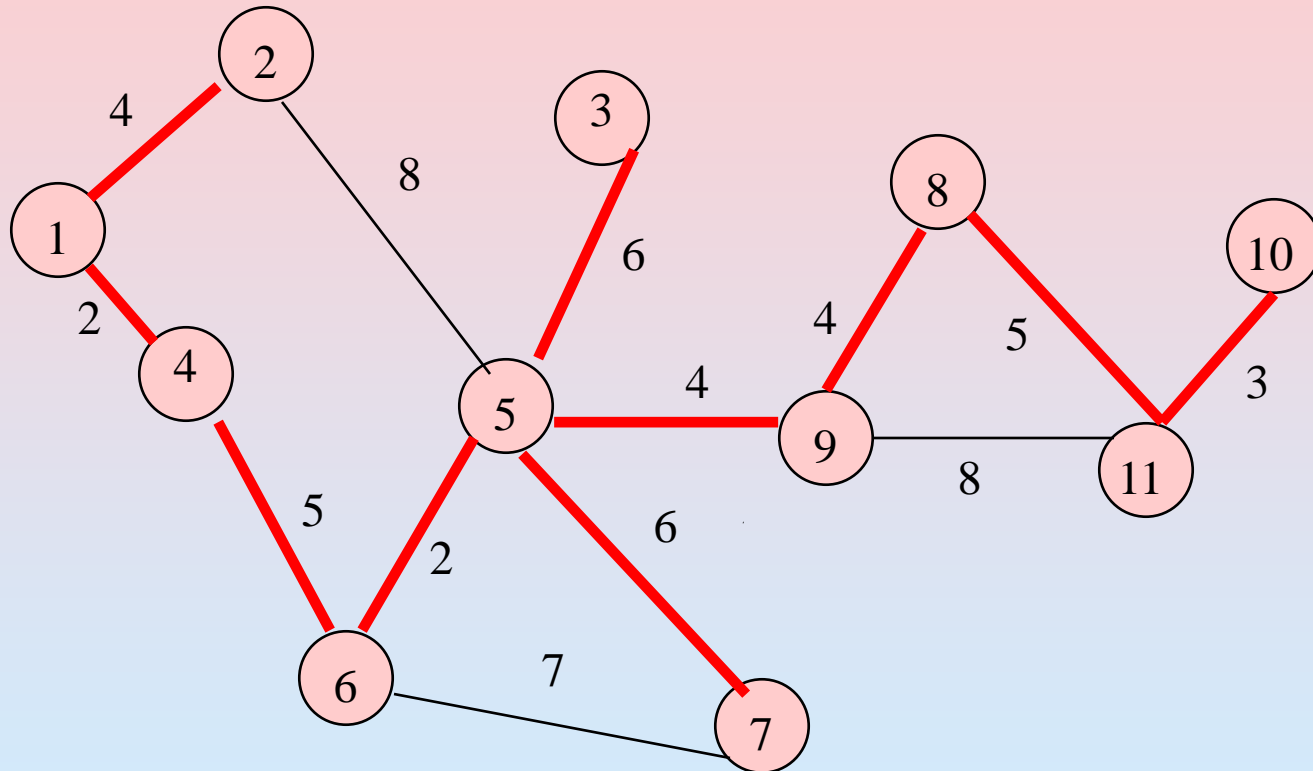
❖ **Tree cost is sum of edge weights/costs.**

A Spanning Tree



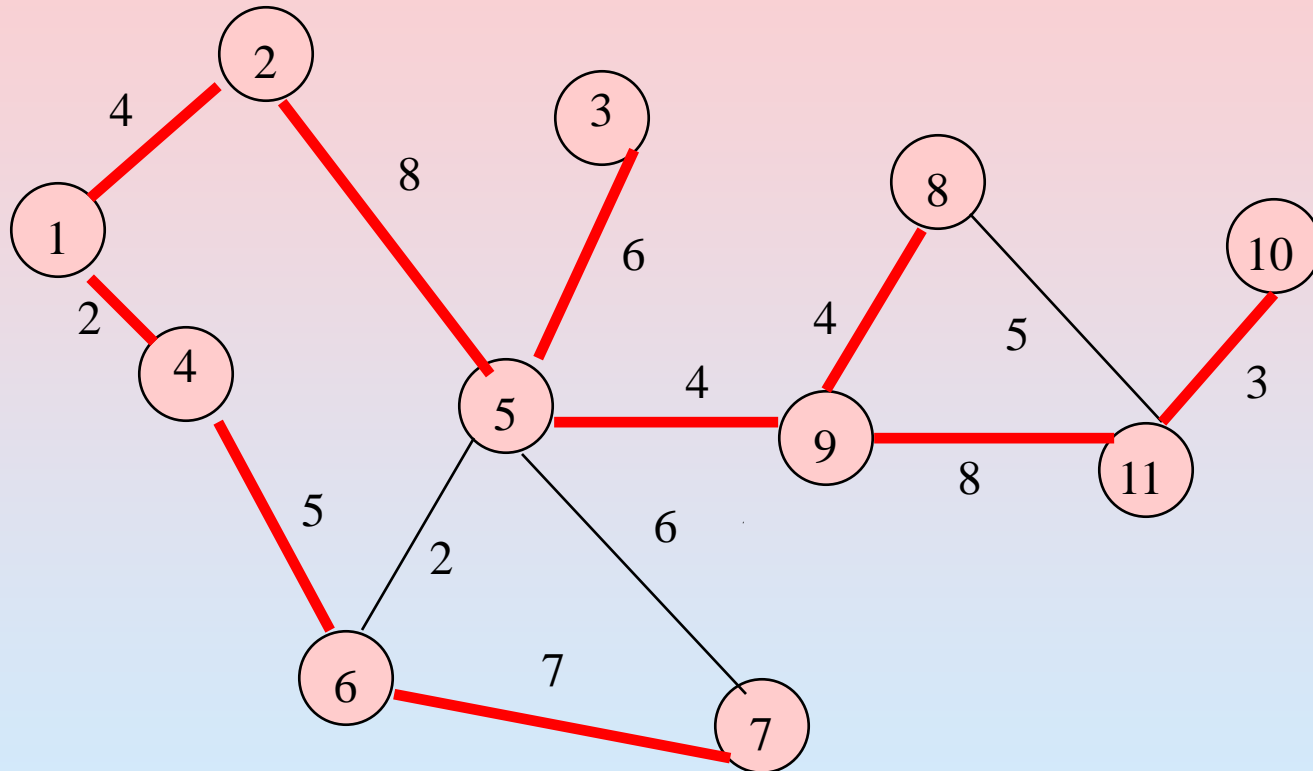
Spanning tree cost = 51.

Minimum Cost Spanning Tree



Spanning tree cost = 41.

A Wireless Broadcast Tree



Source = 1, weights = needed power.

Cost = 4 + 8 + 5 + 6 + 7 + 8 + 3 = 41.

Graph Representation

- ❖ **Adjacency Matrix**

- ❖ **Adjacency Lists**

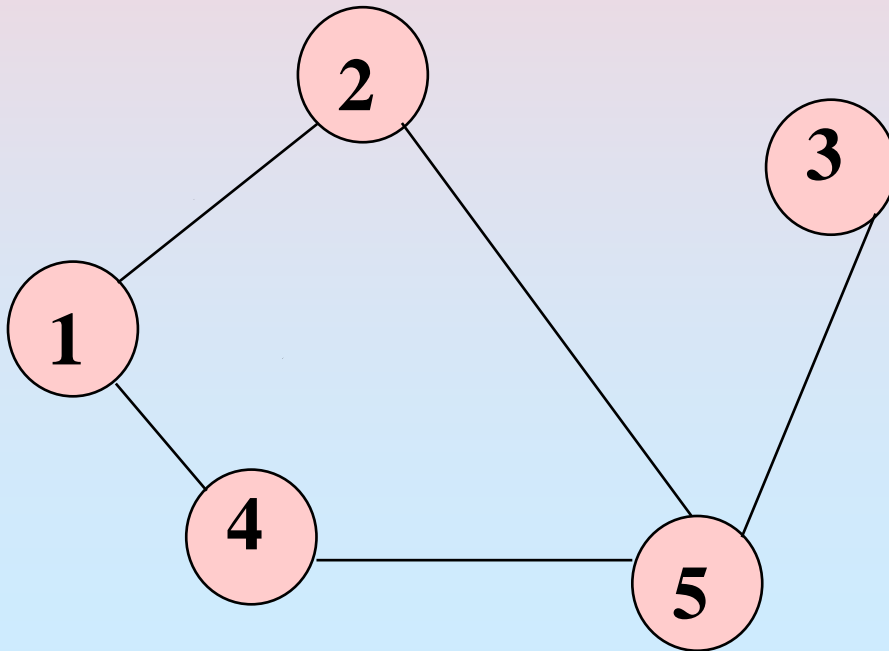
- **Linked Adjacency Lists**

- **Array Adjacency Lists**

Adjacency Matrix

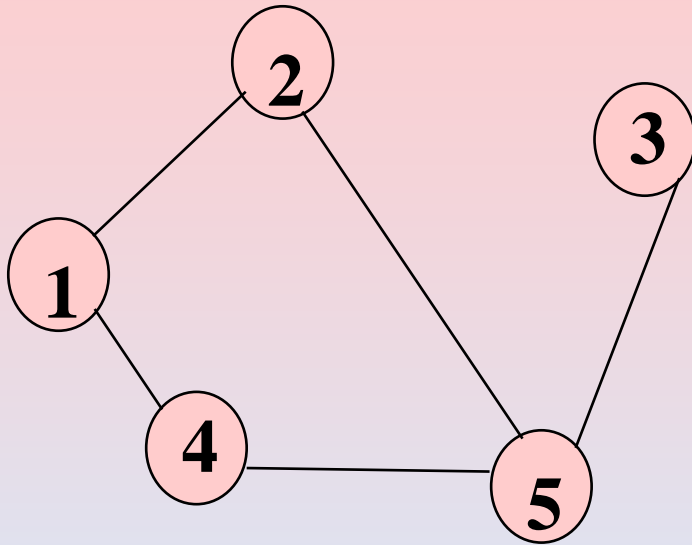
❖ **0/1 $n \times n$ matrix**, where **n is number of vertices**

❖ **$A(i,j) = 1$ iff (i,j) is an edge**



	1	2	3	4	5
1	0	1	0	1	0
2	1	0	0	0	1
3	0	0	0	0	1
4	1	0	0	0	1
5	0	1	1	1	0

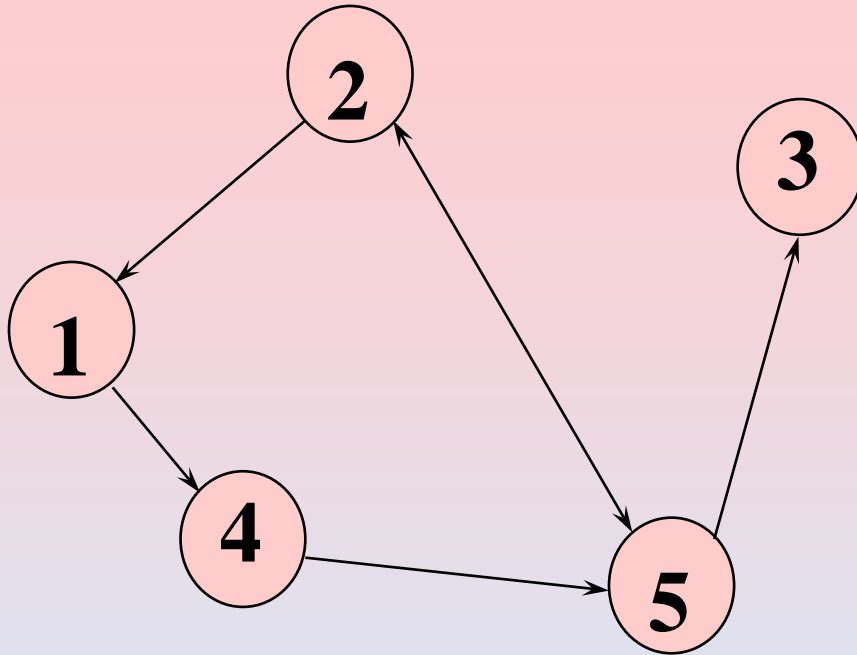
Adjacency Matrix Properties



	1	2	3	4	5
1	0	1	0	1	0
2	1	0	0	0	1
3	0	0	0	0	1
4	1	0	0	0	1
5	0	1	1	1	0

- Diagonal entries are zero.
- Adjacency matrix of an undirected graph is symmetric.
 - $A(i,j) = A(j,i)$ for all i and j .

Adjacency Matrix (Digraph)



	1	2	3	4	5
1	0	0	0	1	0
2	1	0	0	0	1
3	0	0	0	0	0
4	0	0	0	0	1
5	0	1	1	0	0

- Diagonal entries are zero.
- Adjacency matrix of a digraph need not be symmetric.

Adjacency Matrix

- ❖ n^2 bits of space

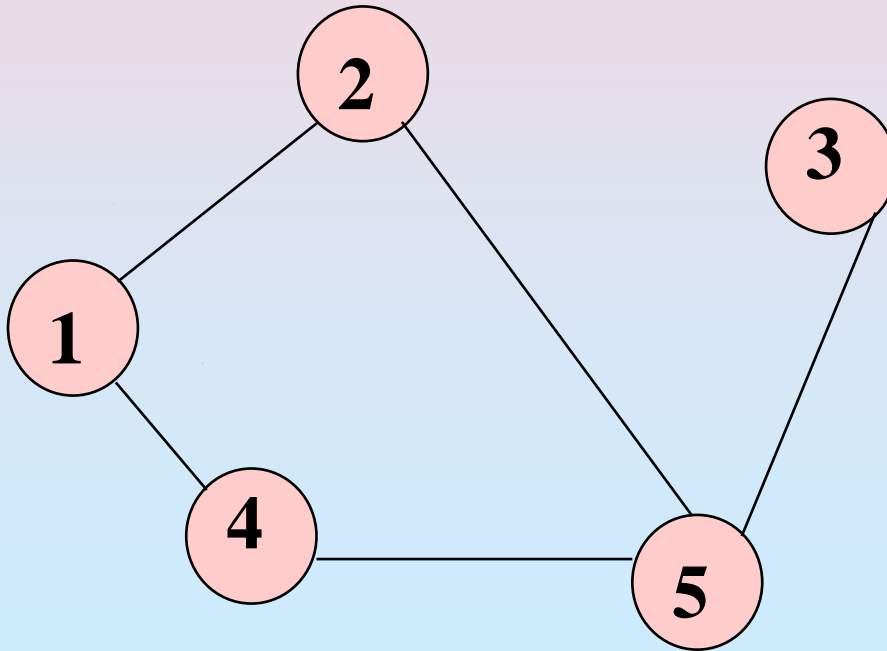
- ❖ For an undirected graph, may store only lower or upper triangle (exclude diagonal).

 - $(n-1)n/2$ bits

- ❖ $O(n)$ time to find vertex degree and/or vertices adjacent to a given vertex.

Adjacency Lists

- ❖ Adjacency list for vertex **i** is a linear list of vertices adjacent from vertex **i**.
- ❖ An array of **n** adjacency lists.



$\text{aList}[1] = (2,4)$

$\text{aList}[2] = (1,5)$

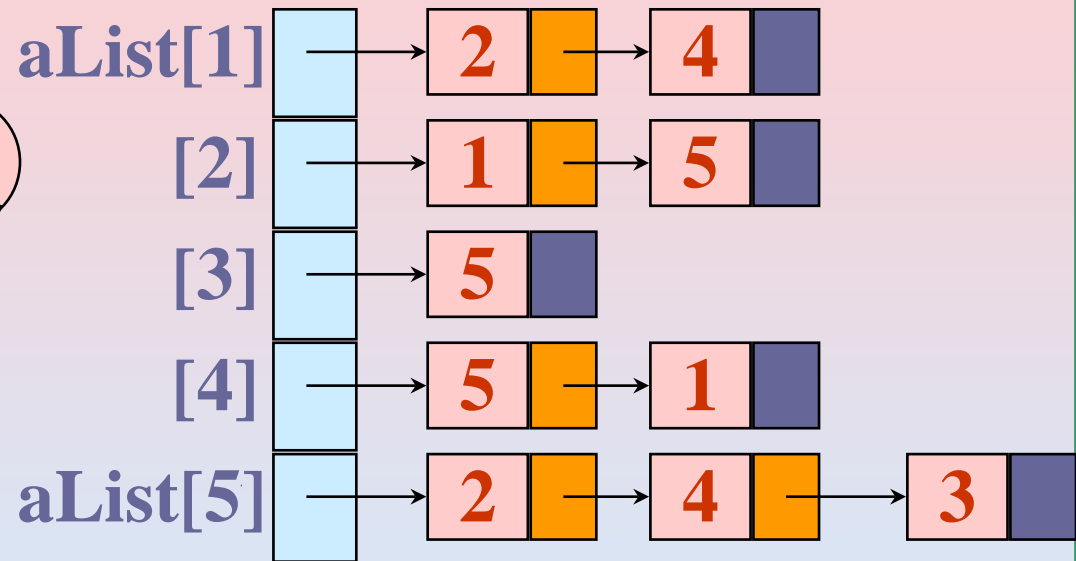
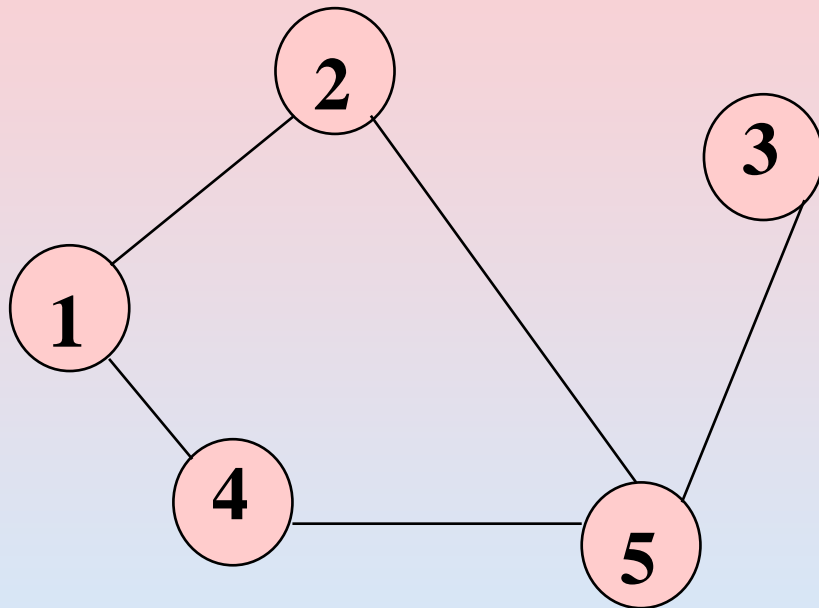
$\text{aList}[3] = (5)$

$\text{aList}[4] = (5,1)$

$\text{aList}[5] = (2,4,3)$

Linked Adjacency Lists

❖ Each adjacency list is a chain.



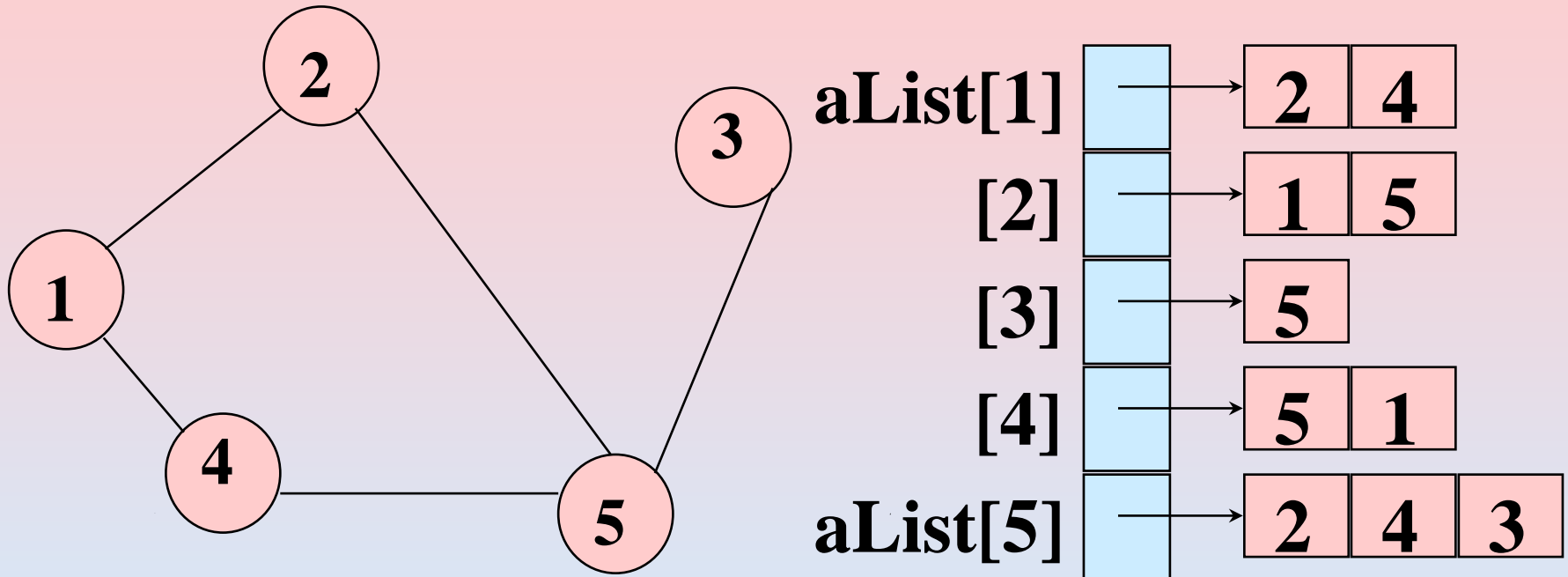
Array Length = n

of chain nodes = $2e$ (undirected graph)

of chain nodes = e (digraph)

Array Adjacency Lists

❖ Each adjacency list is an array list.



Array Length = n

of list elements = $2e$ (undirected graph)

of list elements = e (digraph)

Weighted Graphs

❖ Cost adjacency matrix.

▪ $C(i, j)$ = cost of edge (i, j)

❖ Adjacency lists \Rightarrow each list element is a pair (adjacent vertex, edge weight)