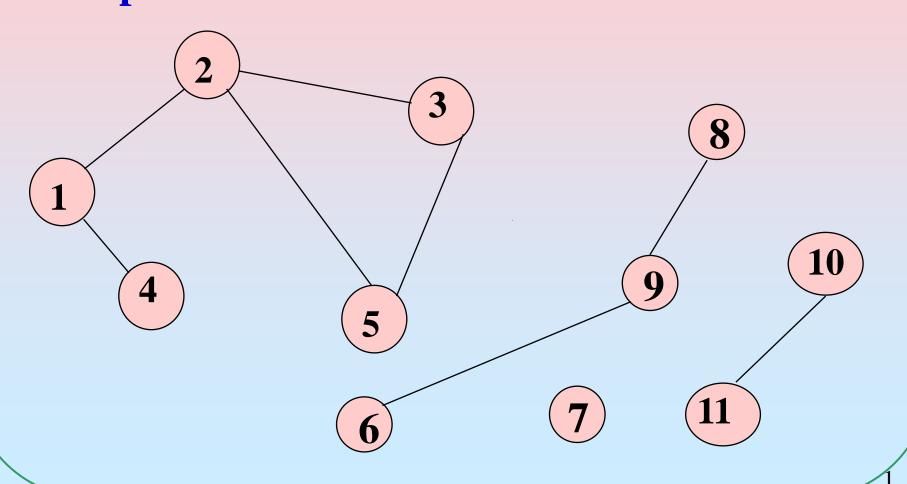
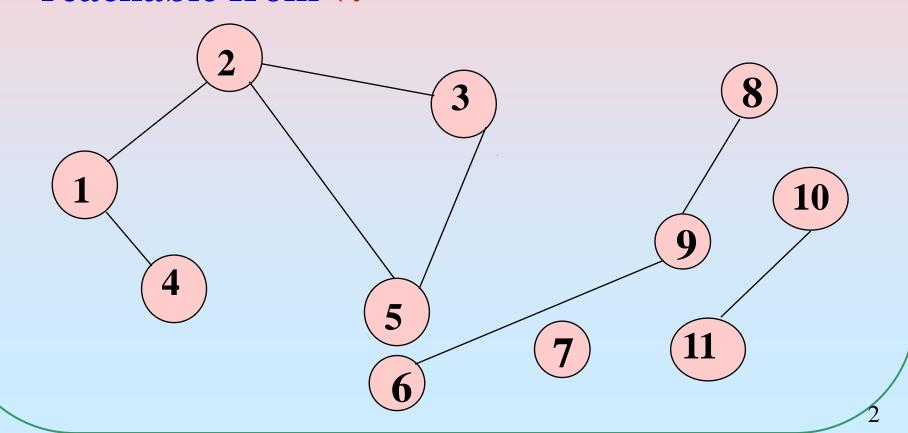
Graph Search Methods

• A vertex u is reachable from vertex v iff there is a path from v to u.



Graph Search Methods

• A search method starts at a given vertex v and visits/labels/marks every vertex that is reachable from v.



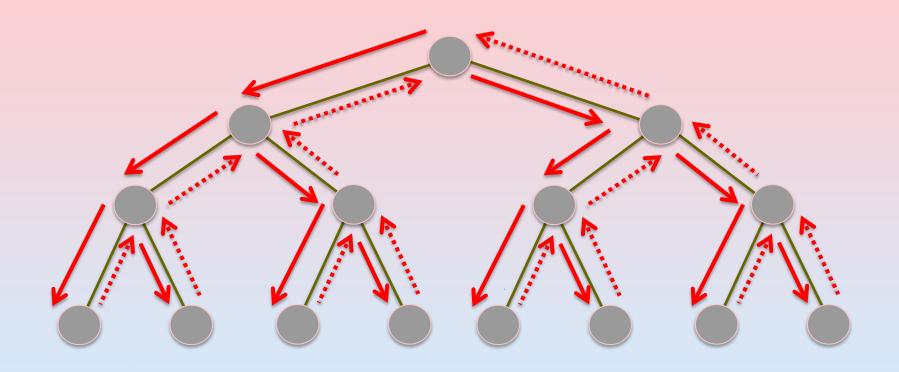
Graph Search Methods

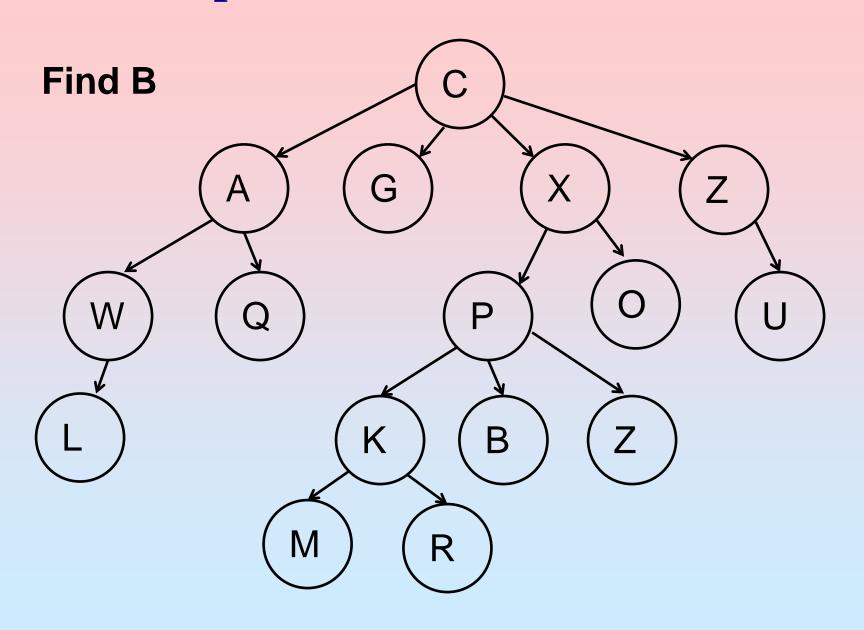
- Many graph problems are solved using a search method.
 - Path from one vertex to another.
 - Is the graph connected?
 - Find a spanning tree, etc.
- Commonly used search methods:
 - Depth-first search(DFS).
 - Breadth-first search(BFS).

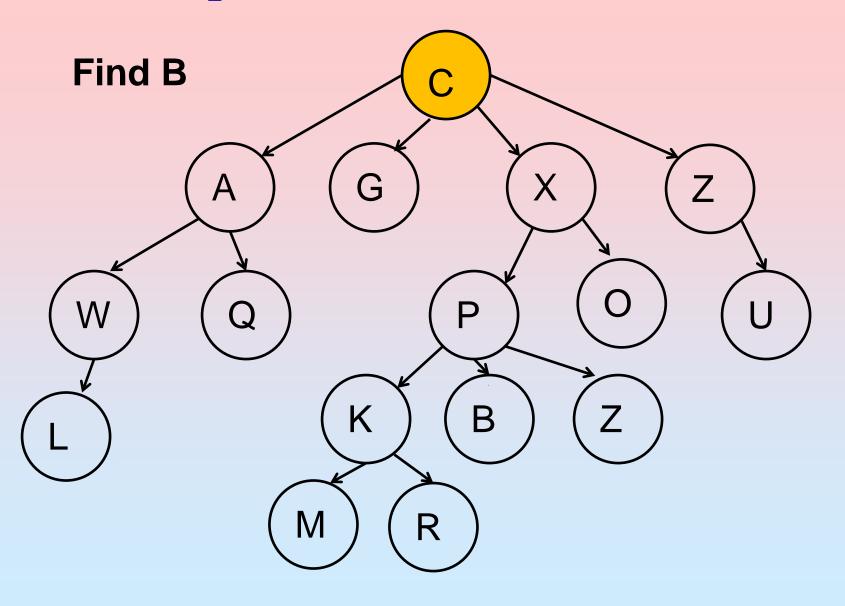
Depth-First Search

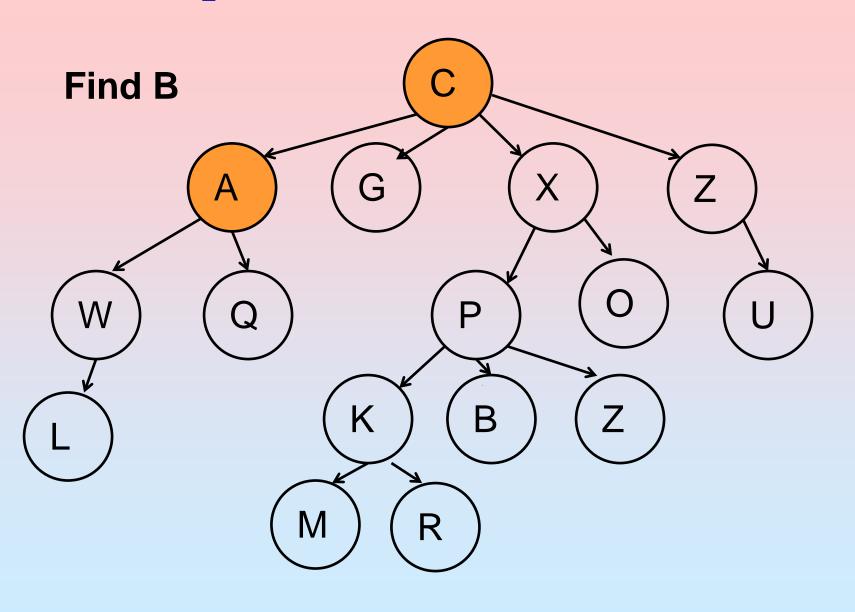
```
dfs (v)
{ Label vertex v as reached.
  for (each unreached vertex u
                 adjacenct from v)
    dfs(u);
```

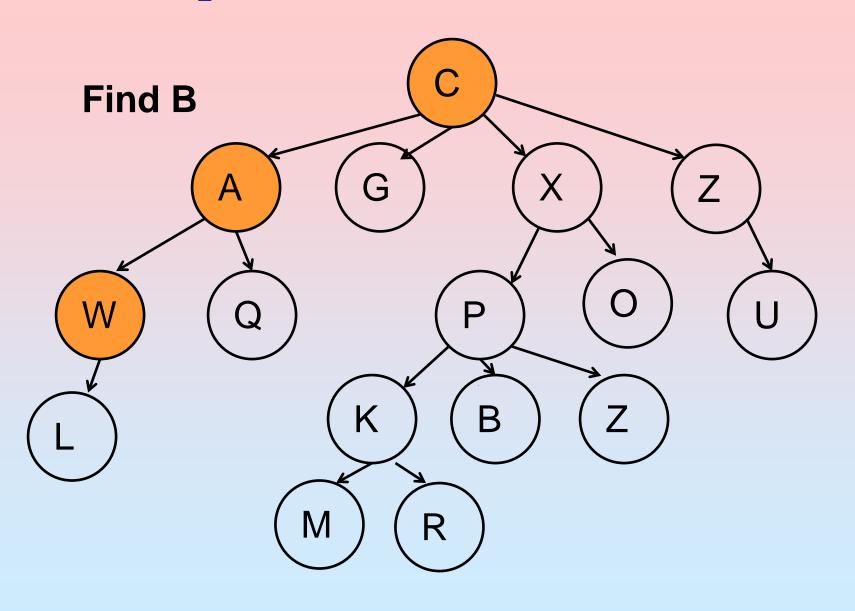
Depth-First Search (DFS) in a Binary Tree

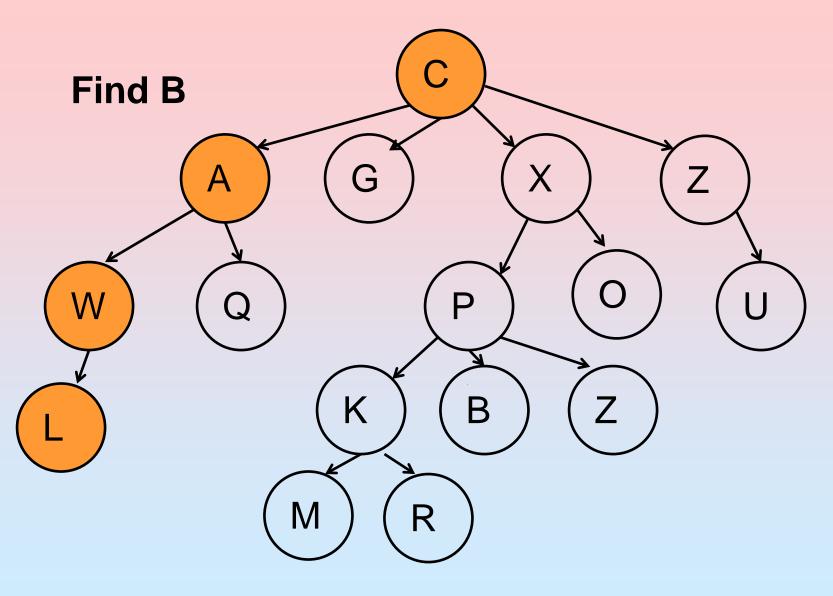


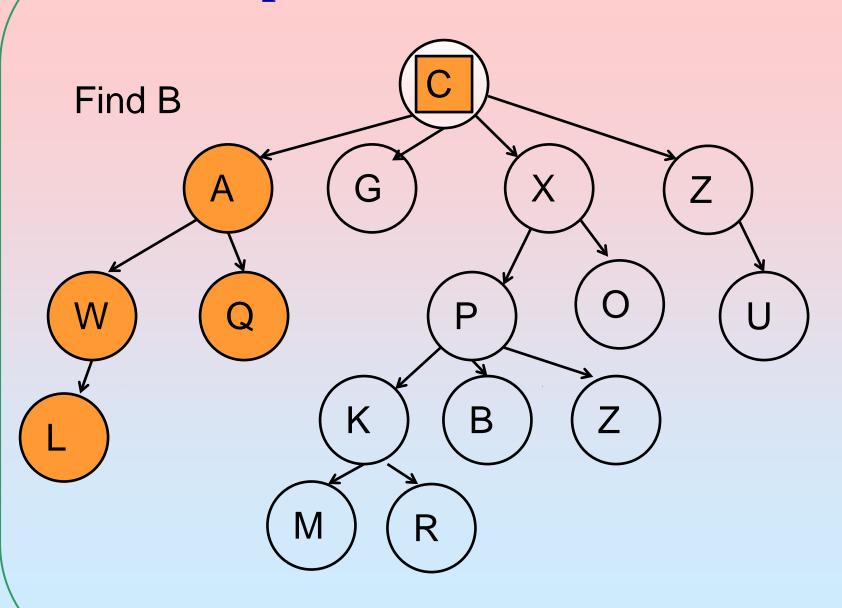


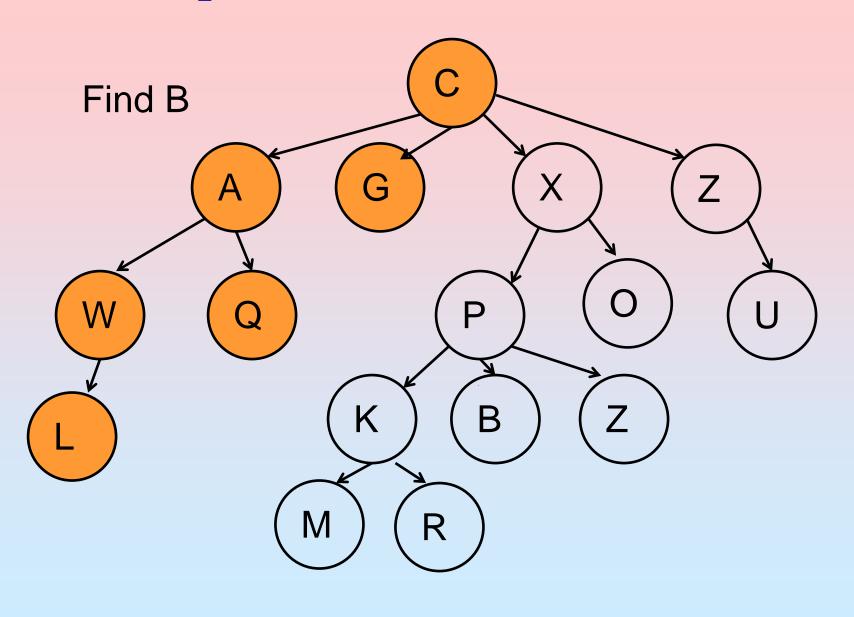


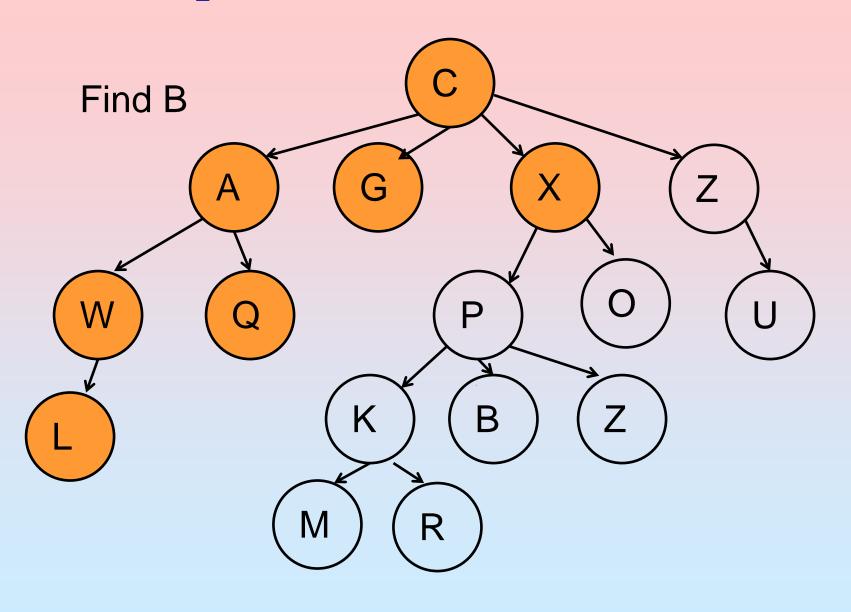


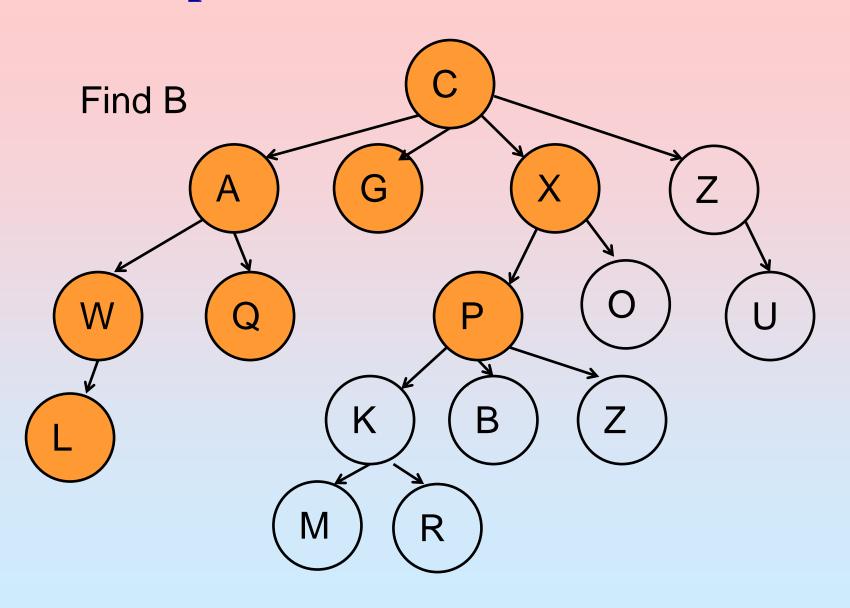


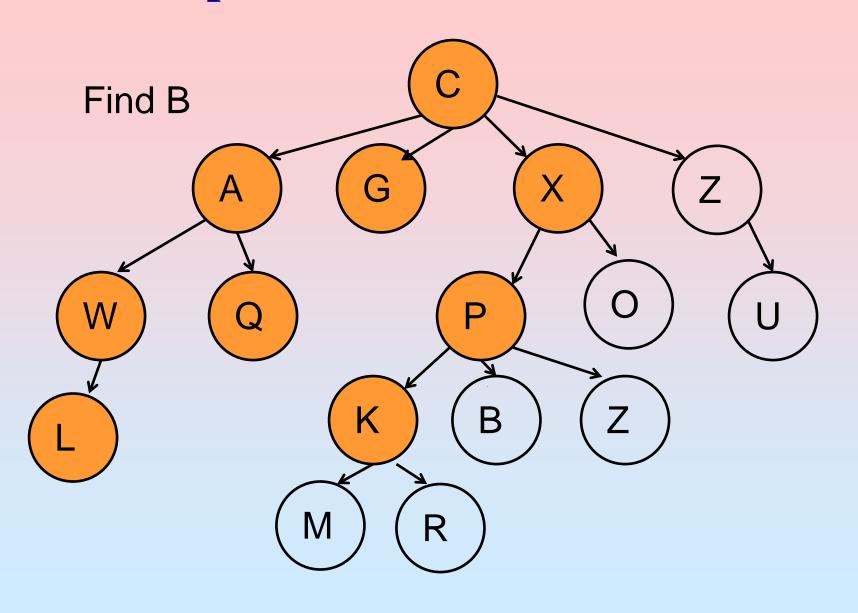


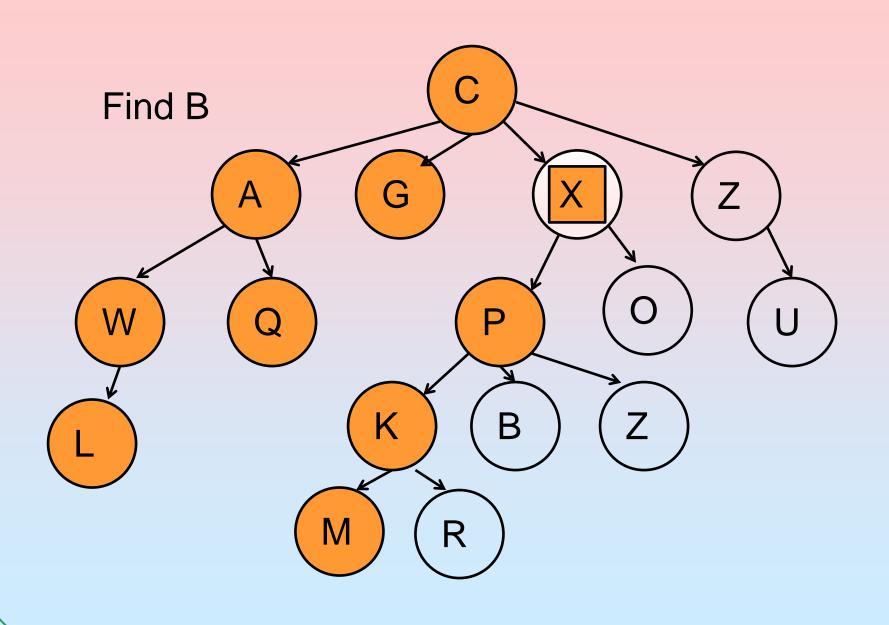


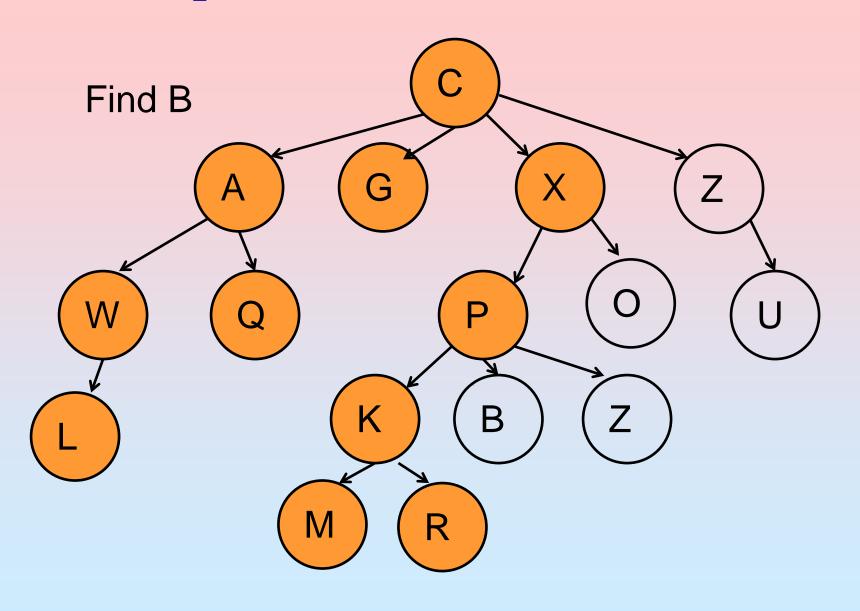


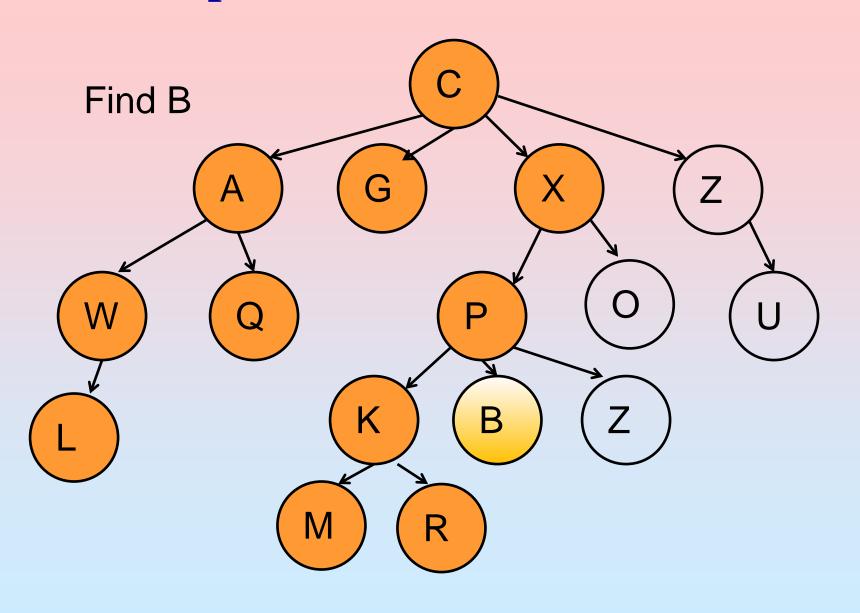


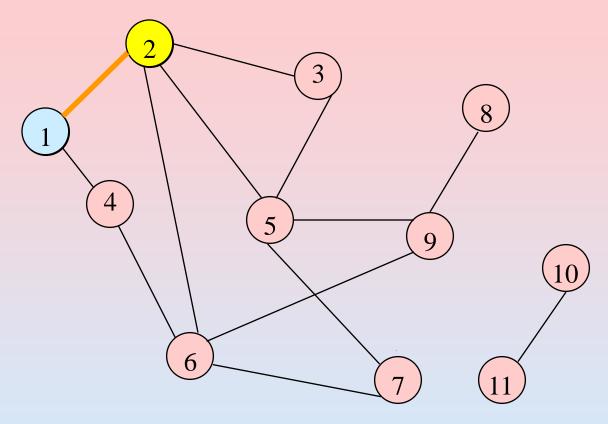








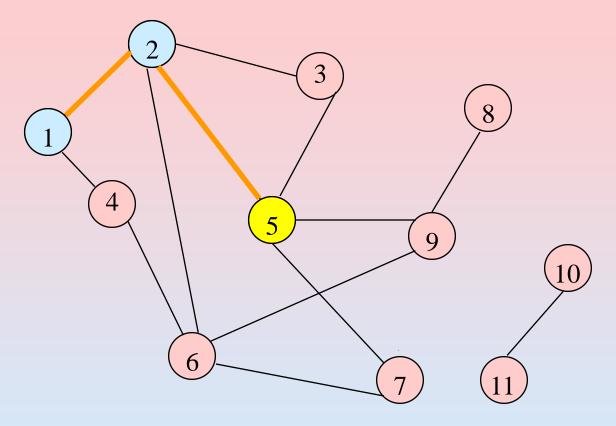




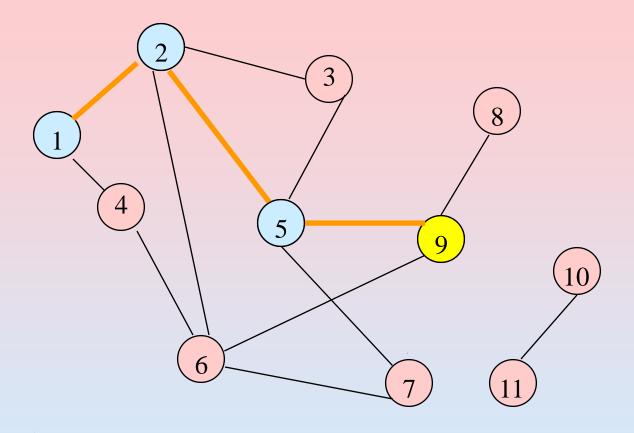
Start search at vertex 1.

Label vertex 1 and do a depth first search from either 2 or 4.

Suppose that vertex 2 is selected.

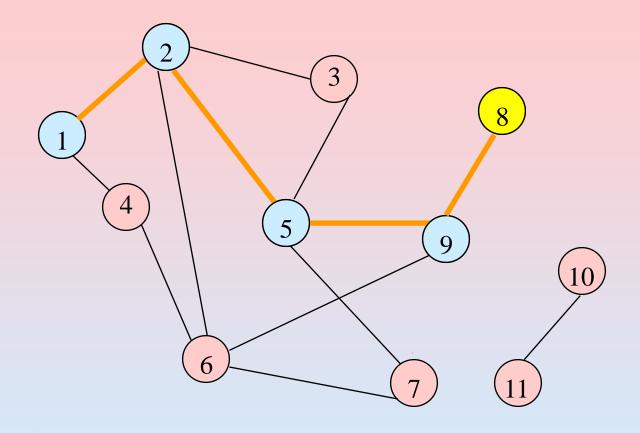


Label vertex 2 and do a depth first search from either 3, 5, or 6. Suppose that vertex 5 is selected.



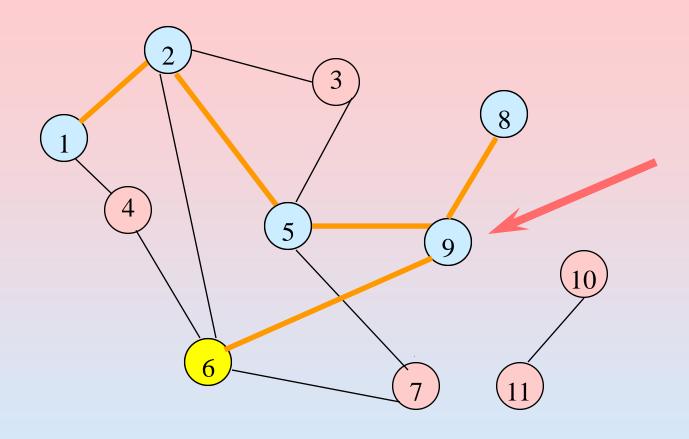
Label vertex 5 and do a depth first search from either 3, 7, or 9.

Suppose that vertex 9 is selected.



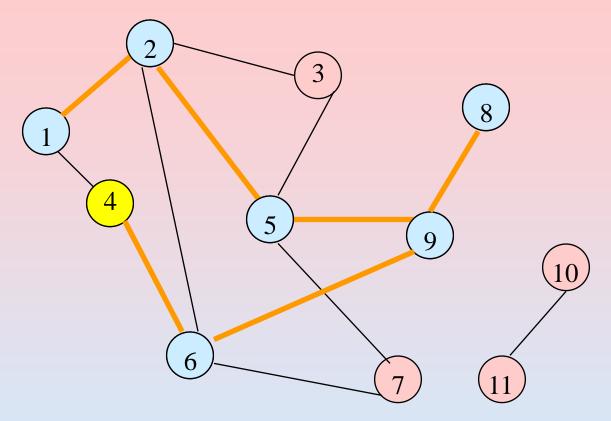
Label vertex 9 and do a depth first search from either 6 or 8.

Suppose that vertex 8 is selected.



Label vertex 8 and return to vertex 9.

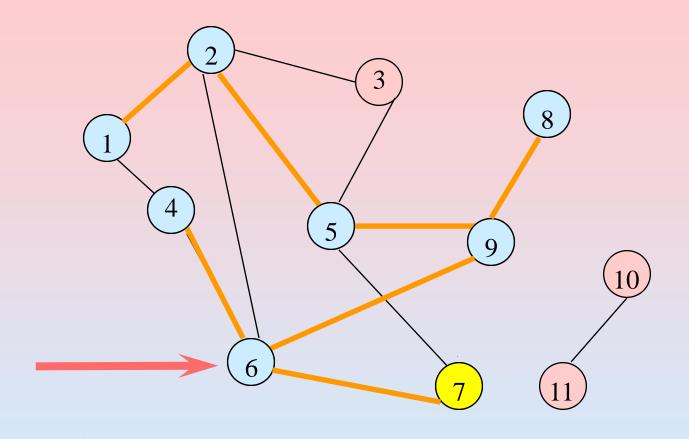
From vertex 9 do a DFS(6).



* Label vertex 6 and do a depth first search from either

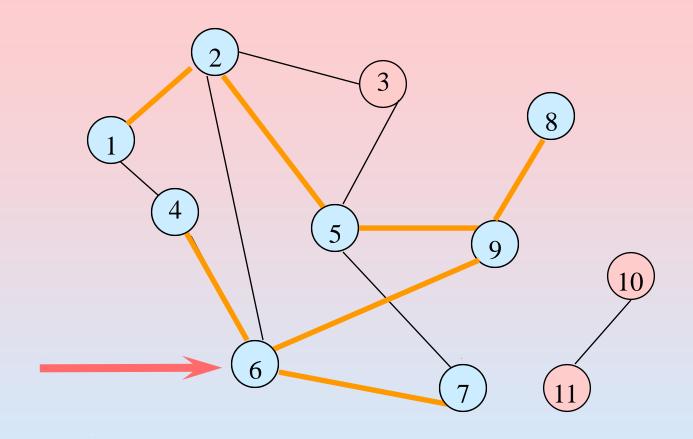
4 or 7.

Suppose that vertex 4 is selected.



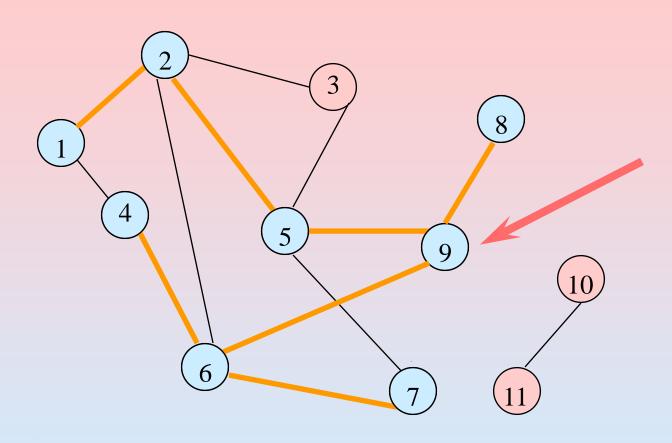
Label vertex 4 and return to 6.

From vertex 6 do a dfs(7).

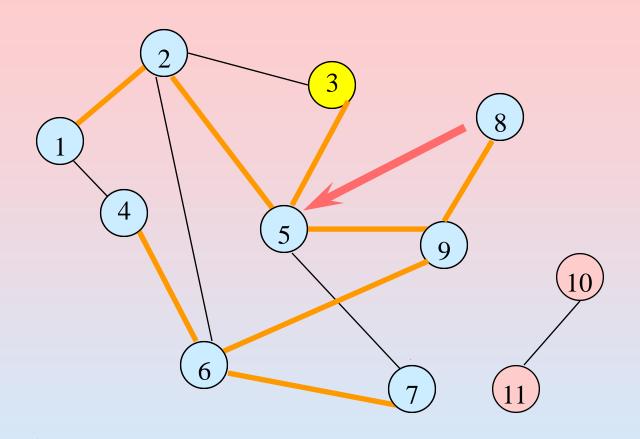


Label vertex 7 and return to 6.

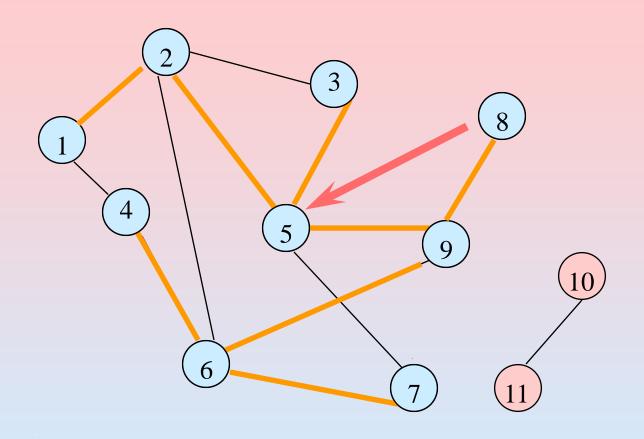
Return to 9.



Return to 5.

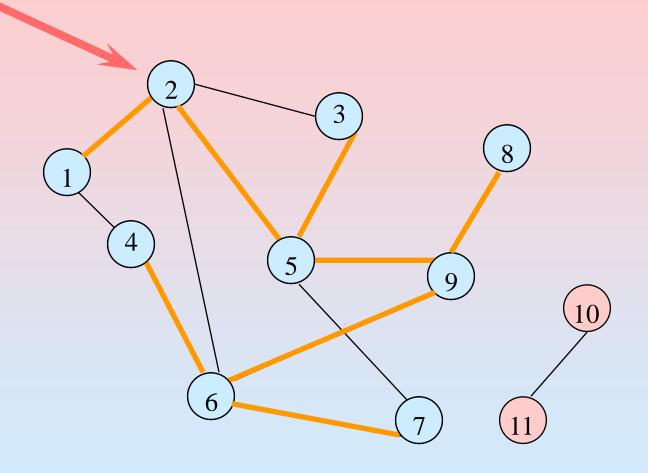


Do a dfs(3).

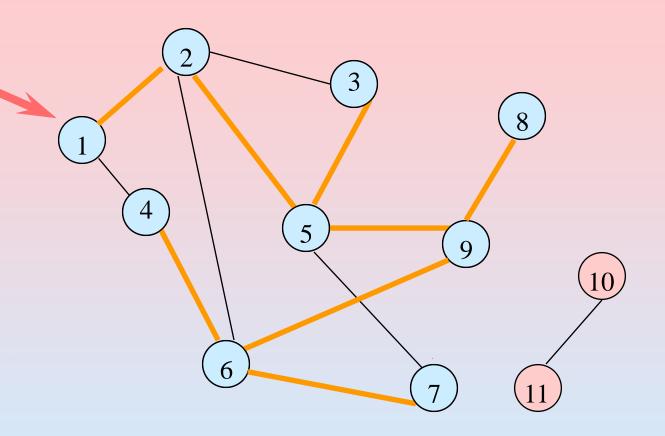


Label 3 and return to 5.

Return to 2.



Return to 1.



Return to invoking method.

Path from Vertex v to Vertex u

- Start a depth-first search at vertex v.
- ***** Terminate when vertex **u** is visited or when **dfs** ends (whichever occurs first).

***** Time Complexity:

• $O(n^2)$ when adjacency matrix used: If the graph is implemented as an adjacency matrix (a n x n array), then, for each node, need to traverse an entire row of length n in the matrix to discover all its outgoing edges. Note that each row in an adjacency matrix corresponds to a node in the graph, and the said row stores information about edges stemming from the node. So, the complexity of DFS is $O(n * n) = O(n^2)$.

Path from Vertex v to Vertex u

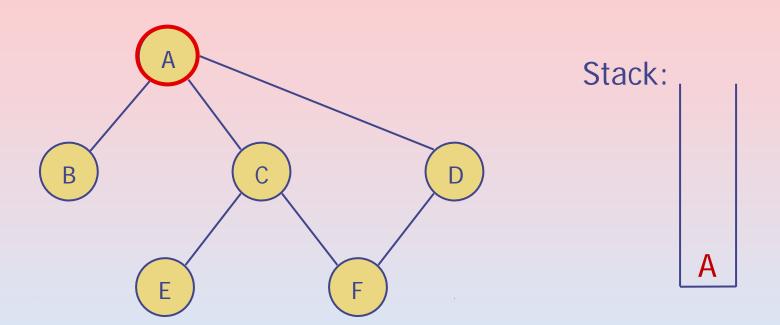
***** Time Complexity:

- O(n+e) when <u>adjacency lists used</u> (e is number of edges): If the graph is implemented using adjacency lists, wherein each node maintains a list of all its adjacent edges, then, for each node, need to discover all its neighbors by traversing its adjacency list just once in linear time. For a directed graph, the sum of the sizes of the adjacency lists of all the nodes is e (total number of edges). So, the complexity of DFS is O(n) + O(e) = O(n + e).
- For an undirected graph, each edge will appear twice in the adjacency list: for an edge ab, a would appear in adjacency list of b, and b would appear in adjacency list of a. So, the overall complexity will be $O(n) + O(2e) \sim O(n + e)$.

DFS implemetation

- Depth first search typically implemented with stack, implicit with recursion or iteratively with an explicit stack
- Start with a node.
- Push that node onto the stack.
- Each time a node is popped off the stack, push all of the new neighbors of that node onto the stack.

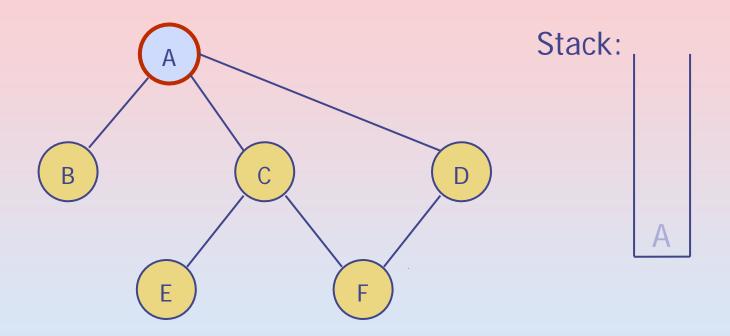
DFS implemetation



Start with a node. Let's start at A!

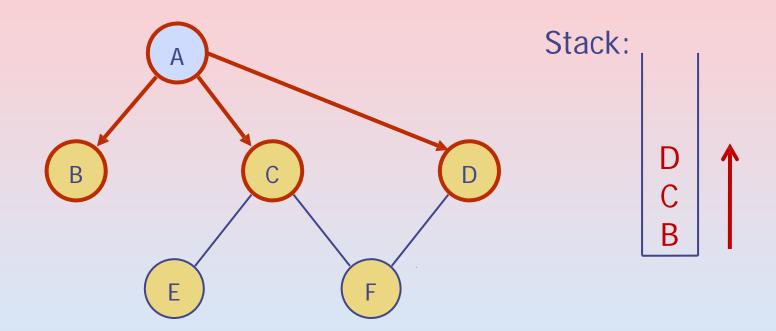
Push the A onto the stack.

DFS implemetation



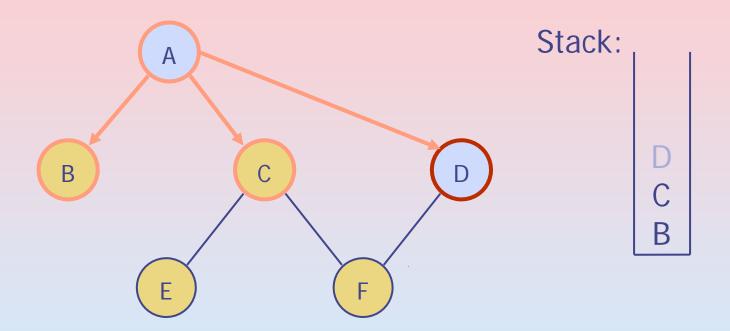
Pop a node off the stack.

A



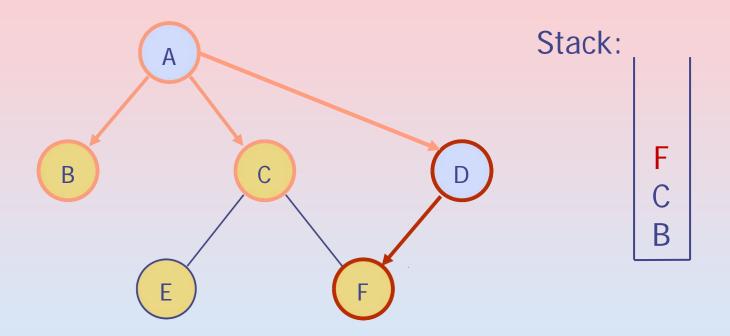
Push the new neighbors of root A onto the stack.

A



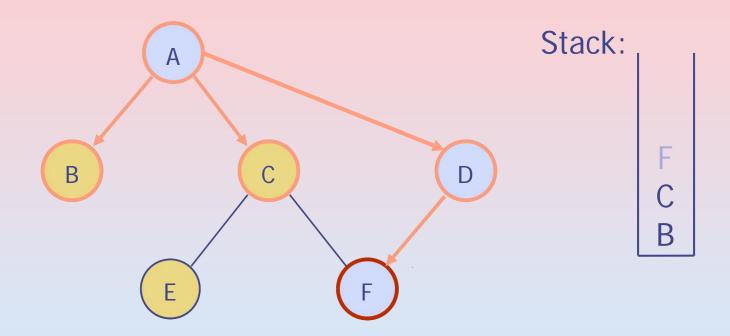
Pop a node off the stack.

A D



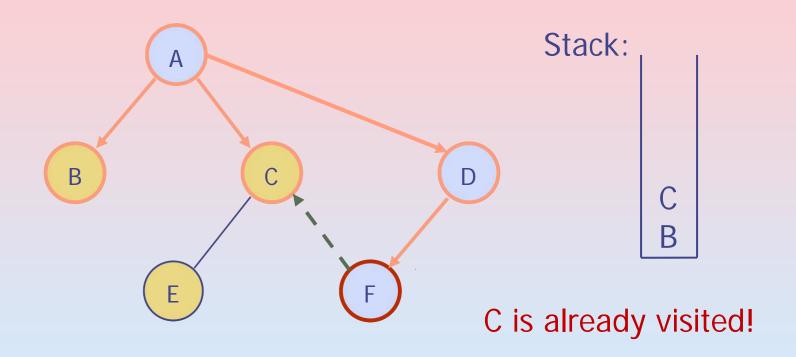
Push the new neighbors of D onto the stack.

A D



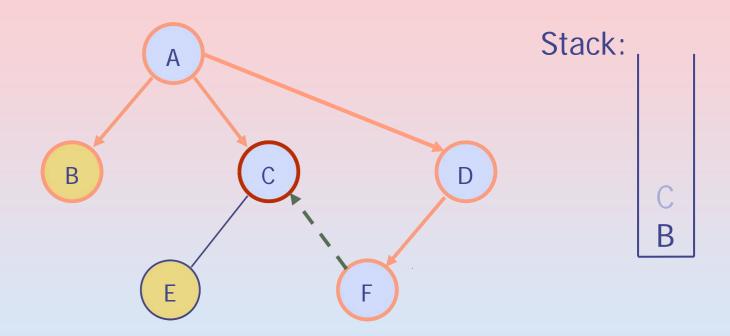
Pop a node off the stack.

ADF



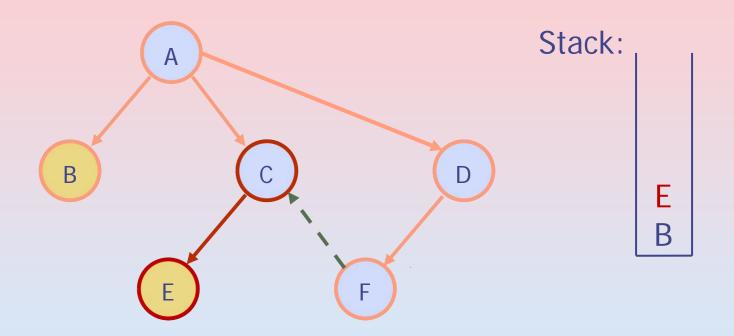
Push the new neighbors of F onto the stack.

A D F



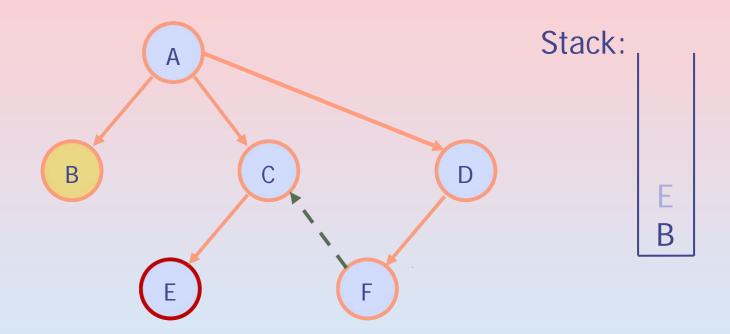
Pop a node off the stack.

ADFC



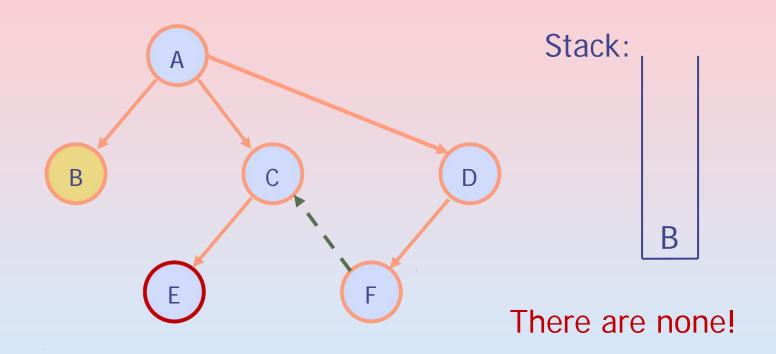
Push the new neighbors of C onto the stack.

ADFC



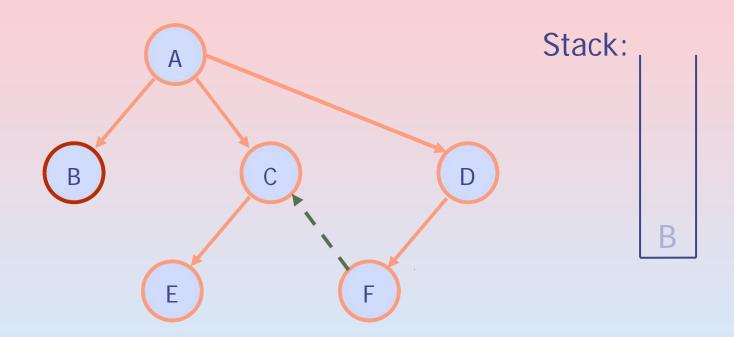
Pop a node off the stack.

ADFCE



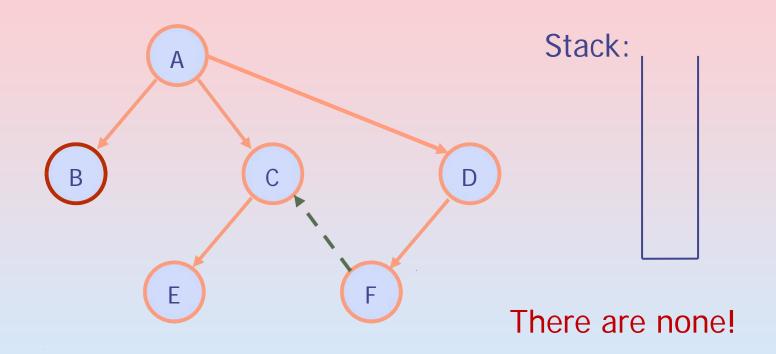
Push the new neighbors of E onto the stack.

ADFCE



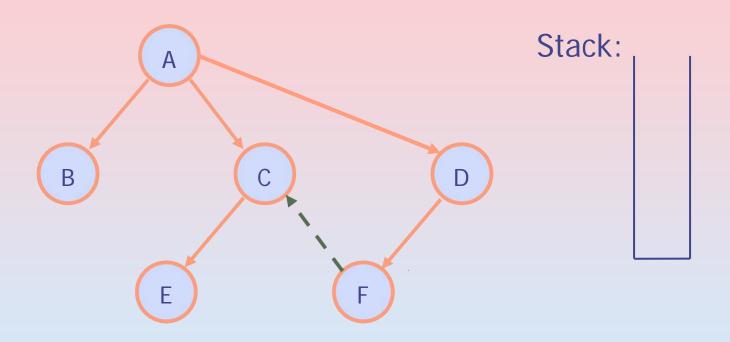
Pop a node off the stack.

ADFCEB



Push the new neighbors of B onto the stack.

ADFCEB



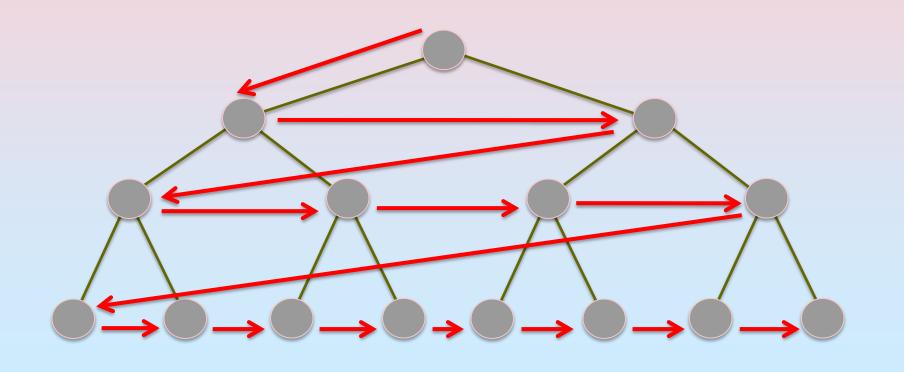
The next step would be to pop a node off the stack.

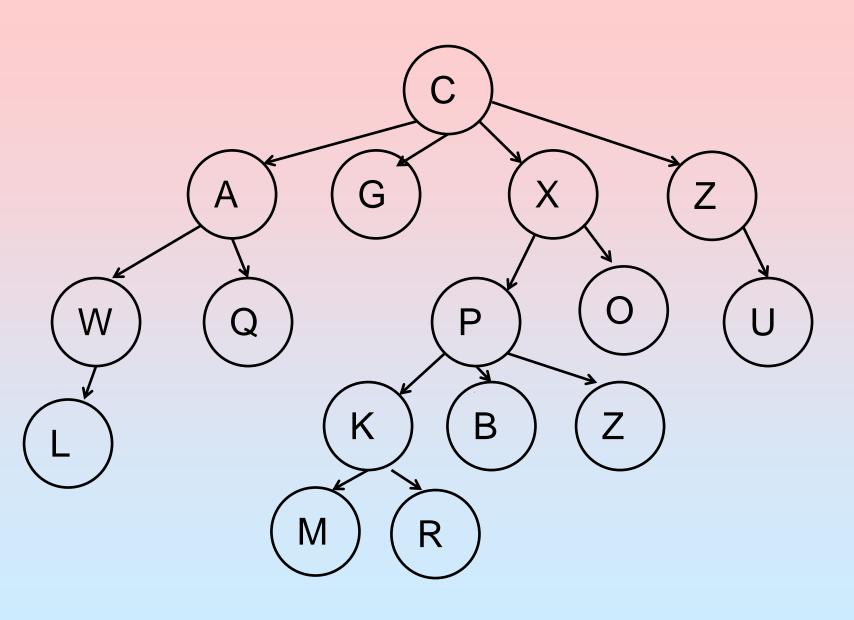
A D F C E B But since the stack is empty, we're done!

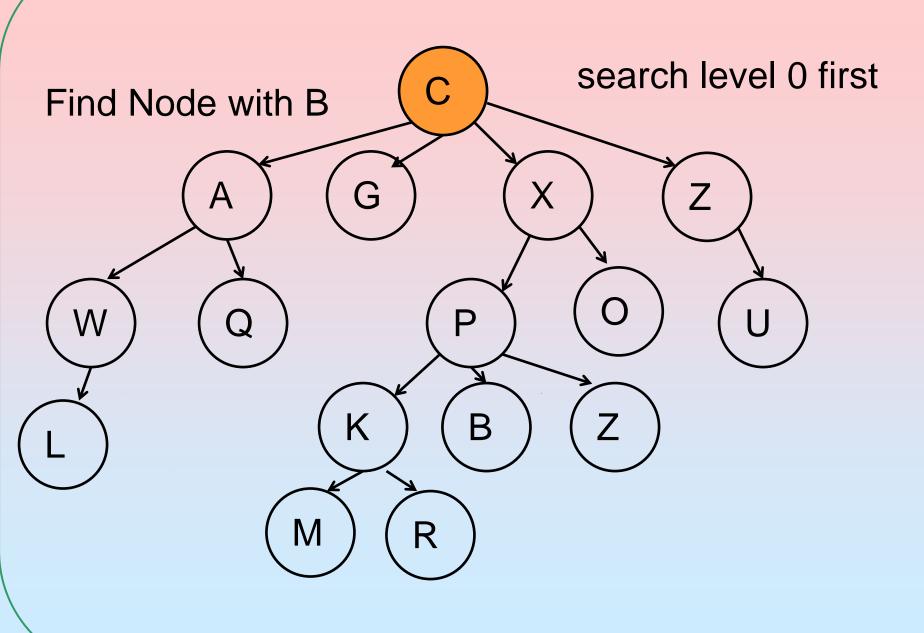
- **❖** We got DFS A D F C E B
- *Is it the only one DFS for the given graph?
 - >NO
- *Another DFS for the given graph?
 - >ABCEFD
- **❖** If we want the search to end up **A B C E F D**, push the neighbors into the stack from right to left (e.g. **D**, **C**, **B**).

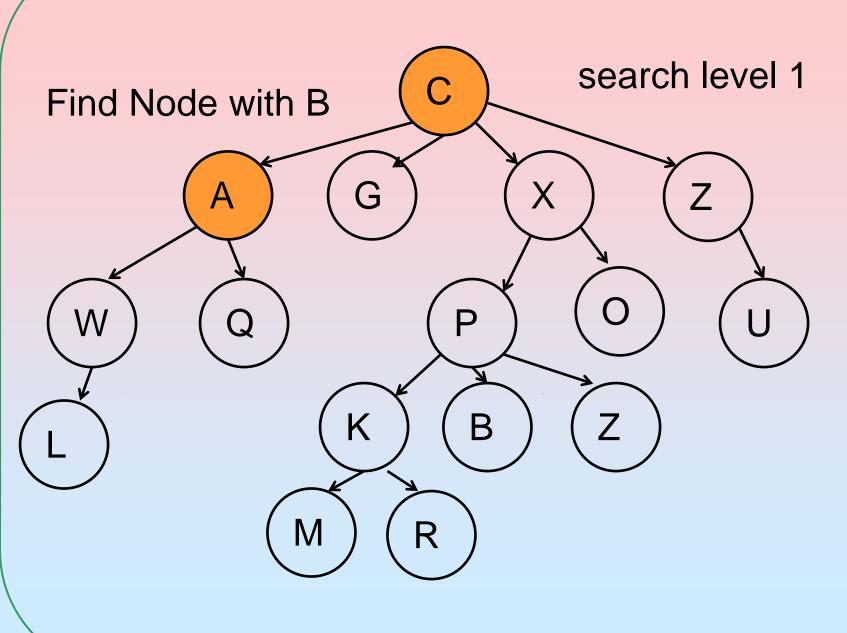
Breadth First

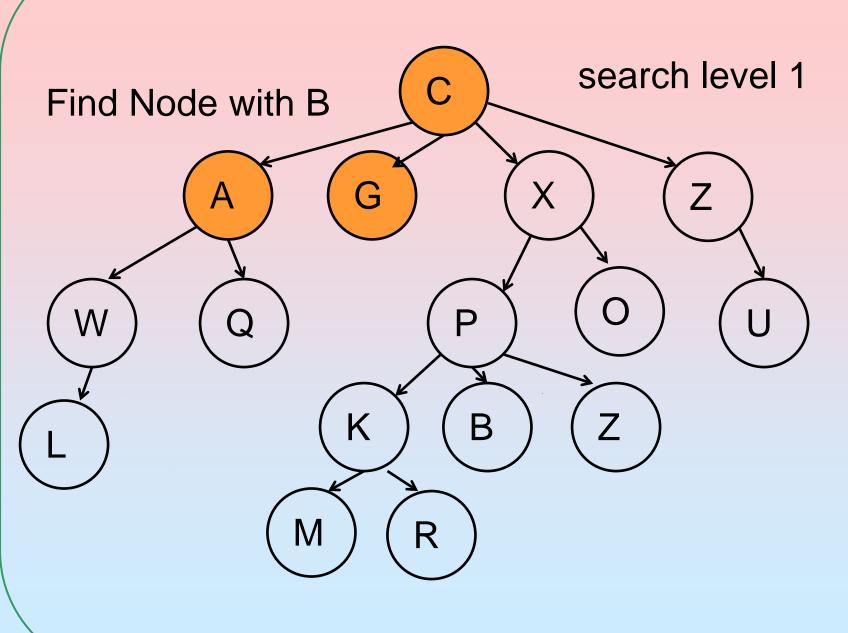
- **❖** A level order traversal of a tree could be used as a breadth first search
- **❖** Search all nodes in a level before going down to the next level

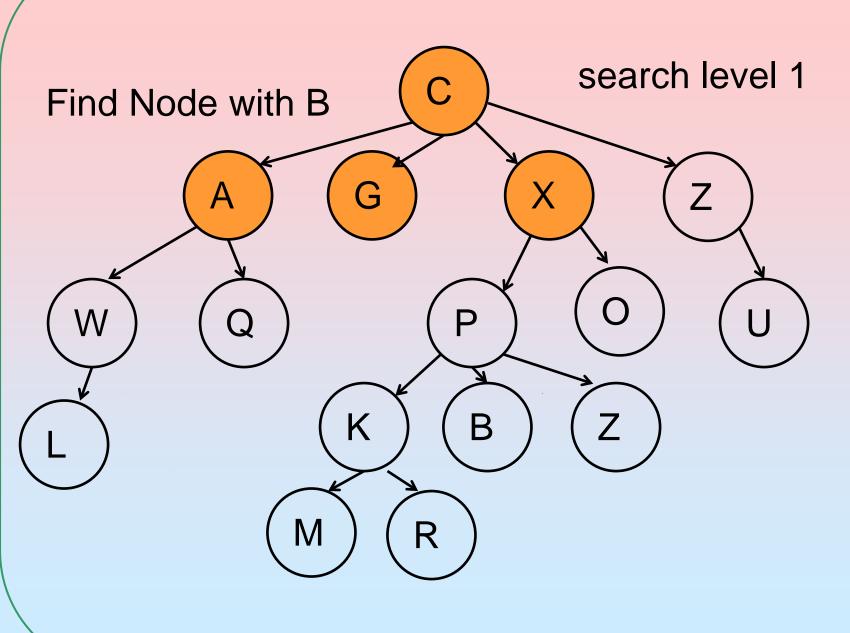


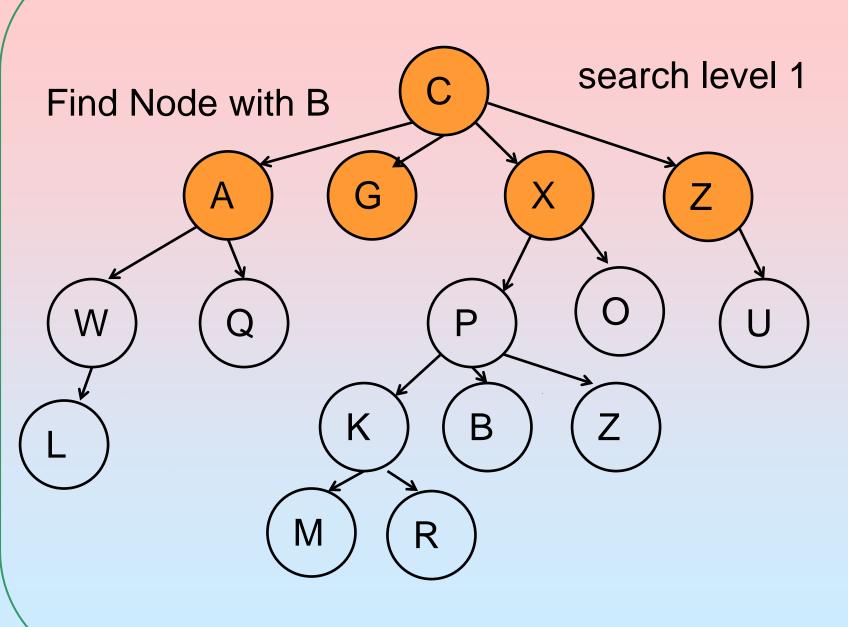


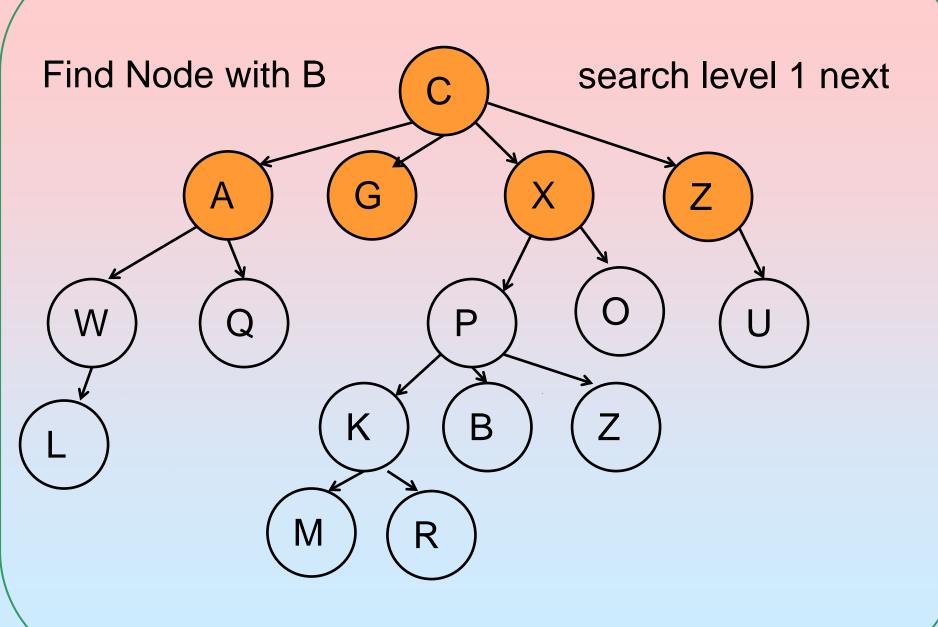


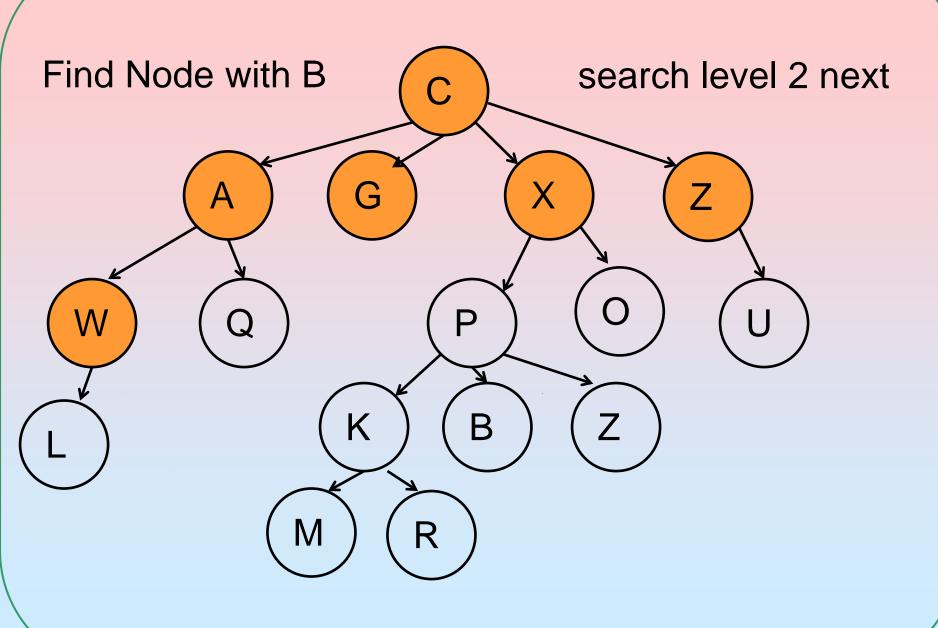


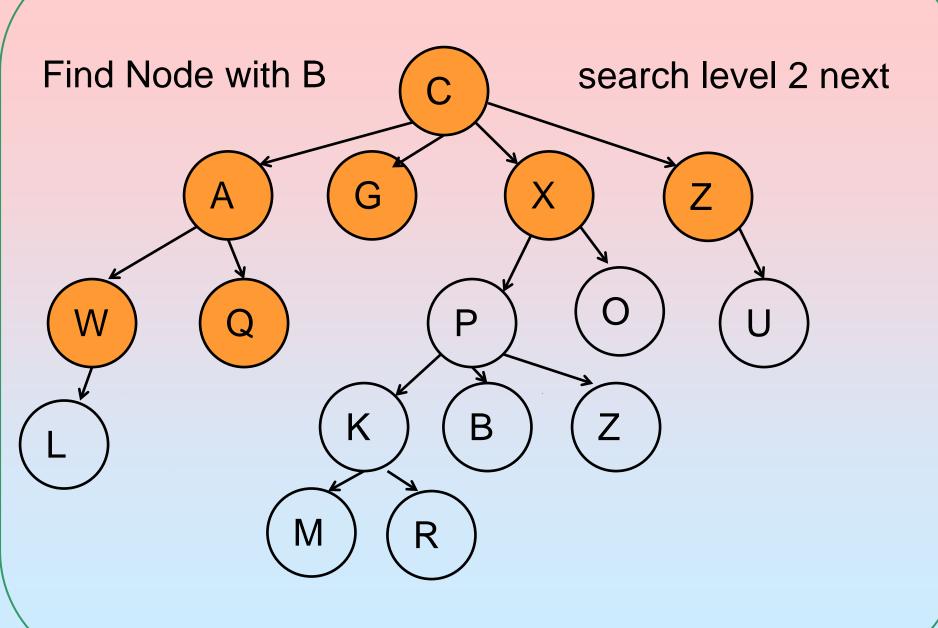


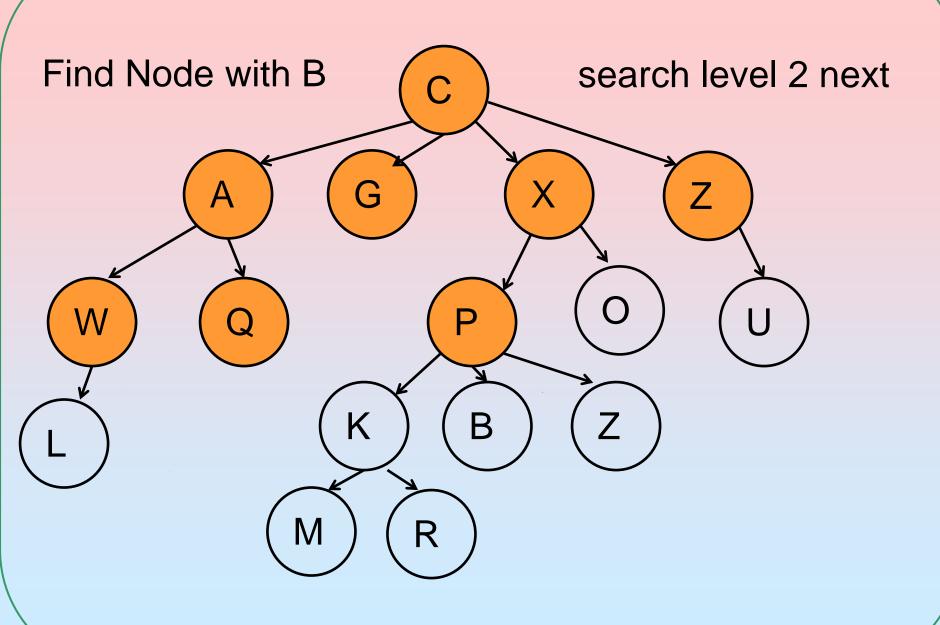


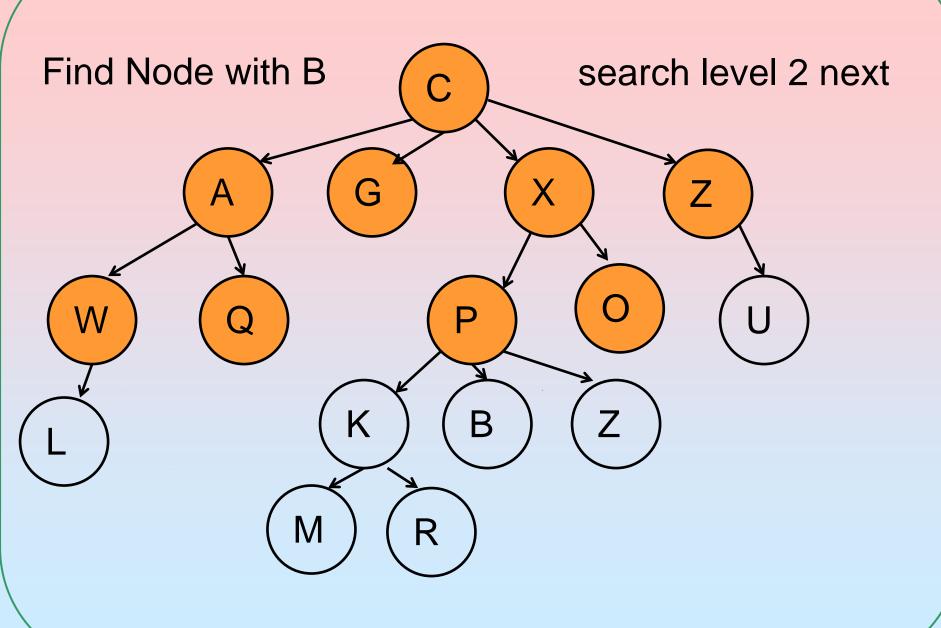


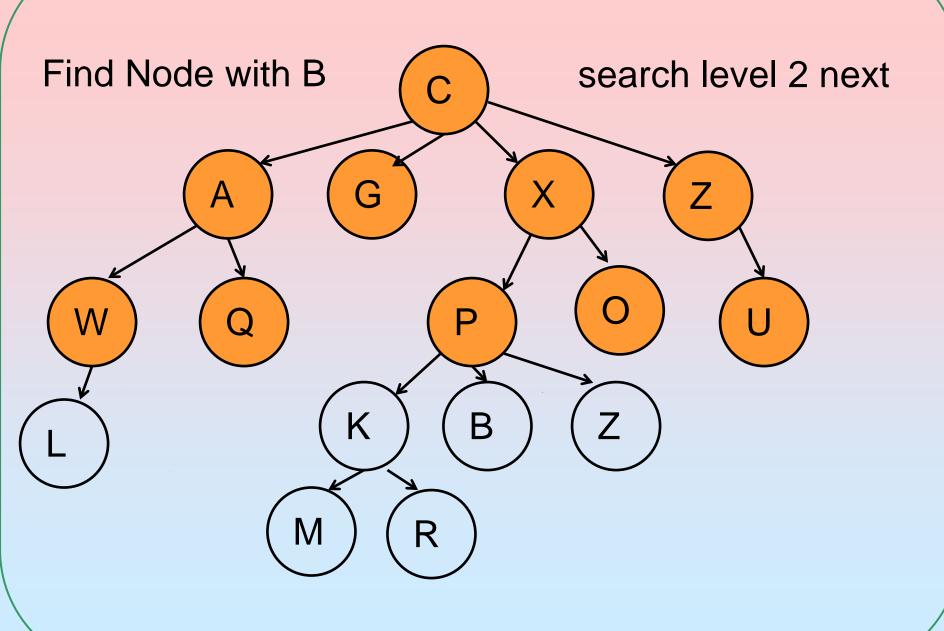


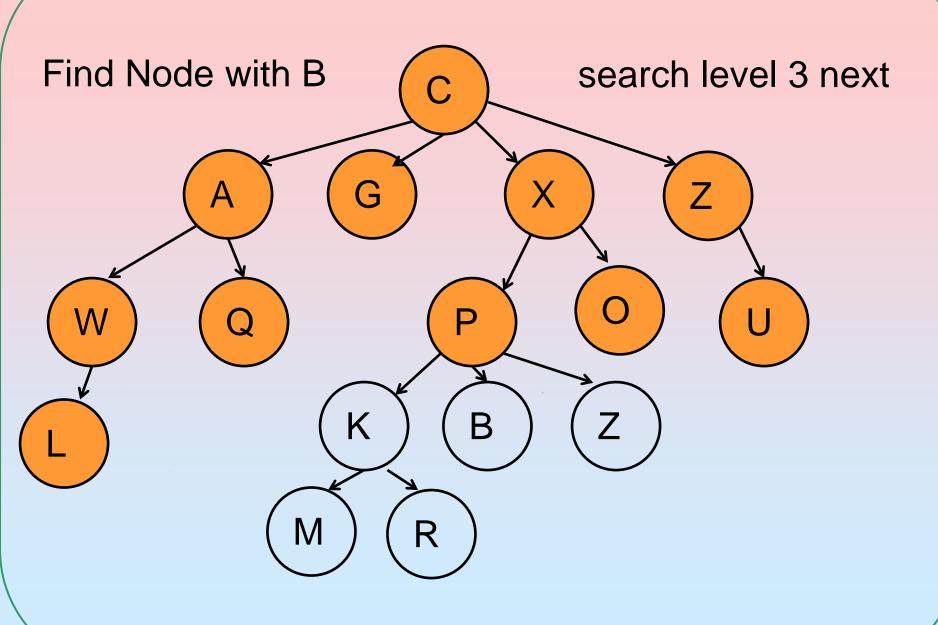


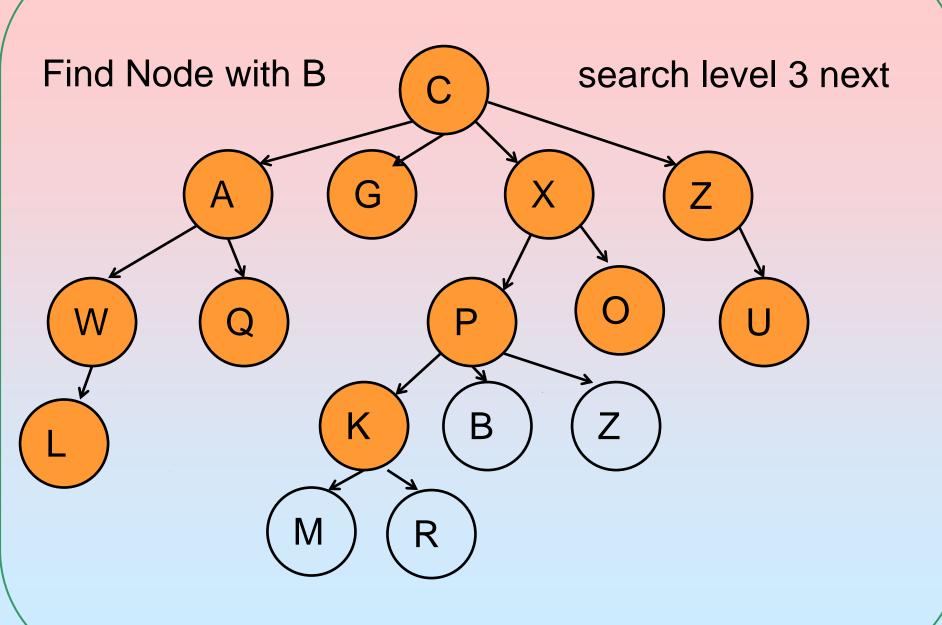


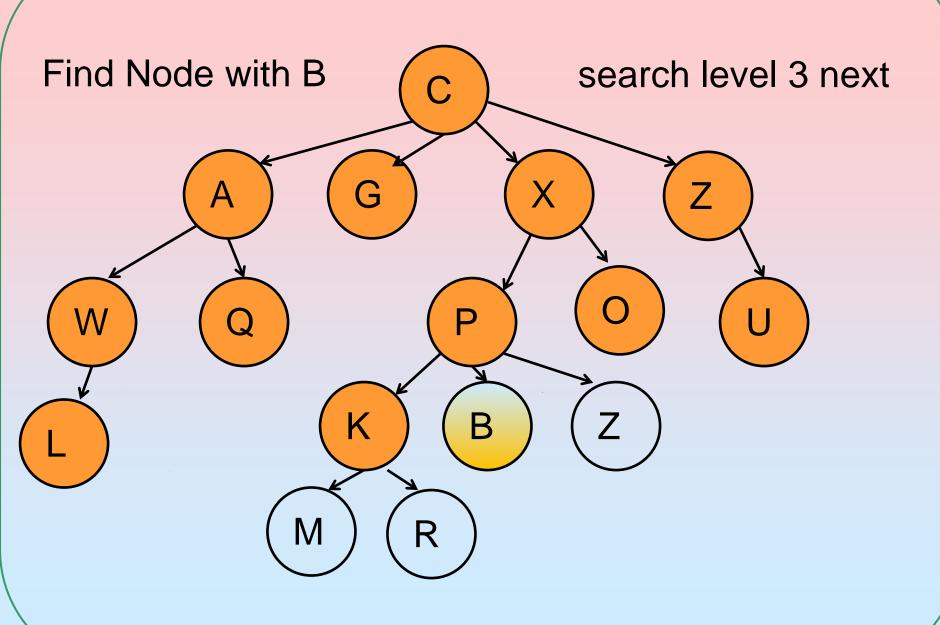






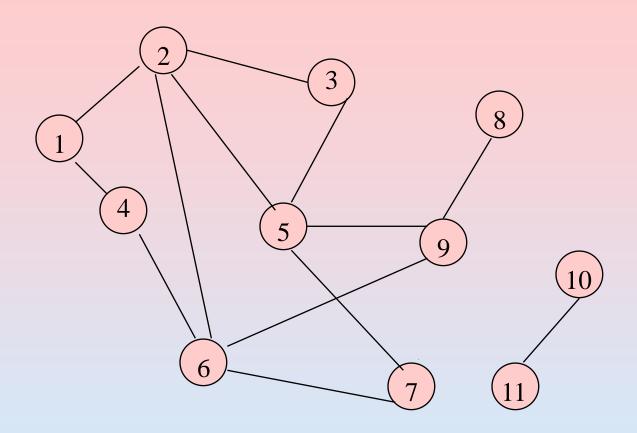




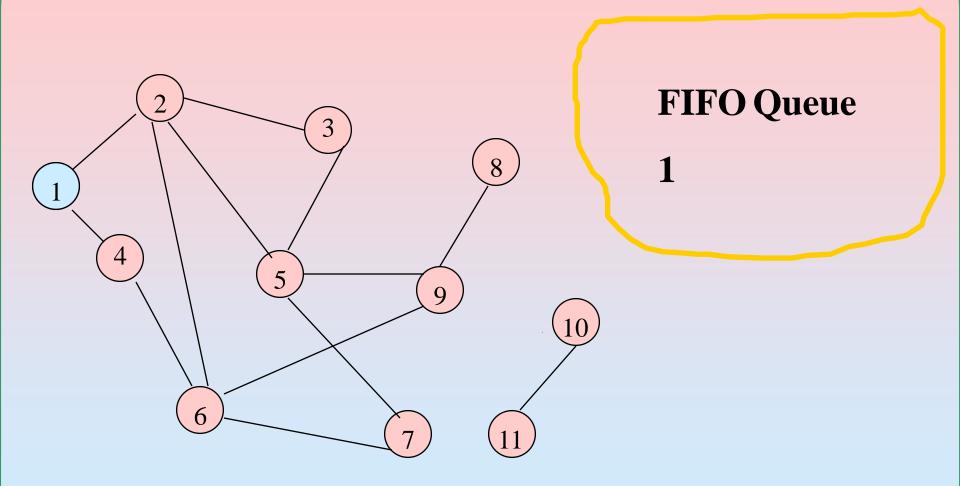


Breadth-First Search Implementaion

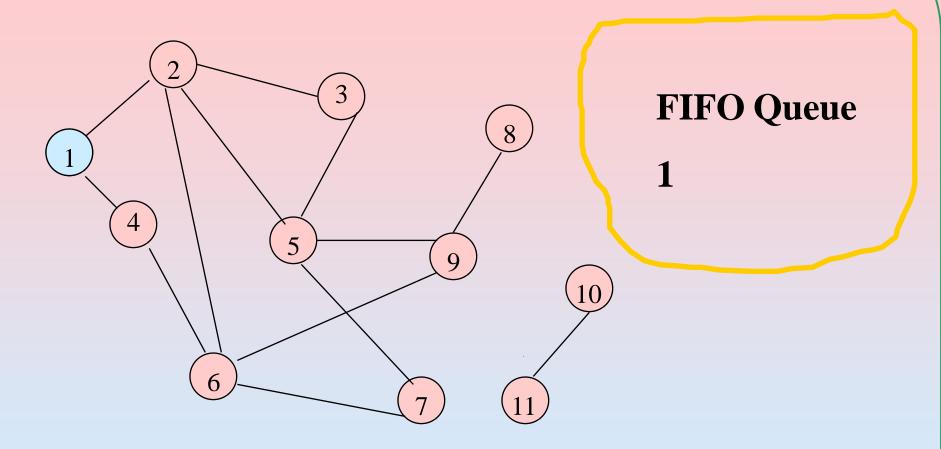
- Visit start vertex and put into a FIFO queue.
- *Repeatedly remove a vertex from the queue, visit its unvisited adjacent vertices, put newly visited vertices into the queue.



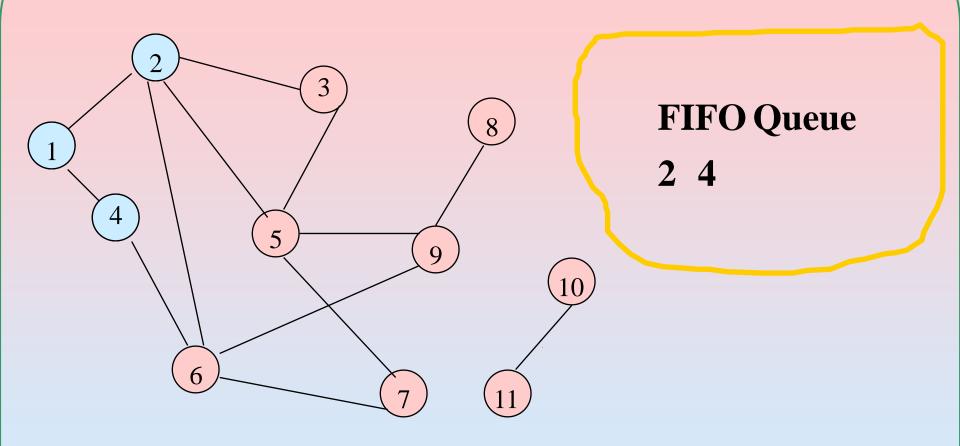
Start search at vertex 1.



Visit/mark/label start vertex and put in a FIFO queue.

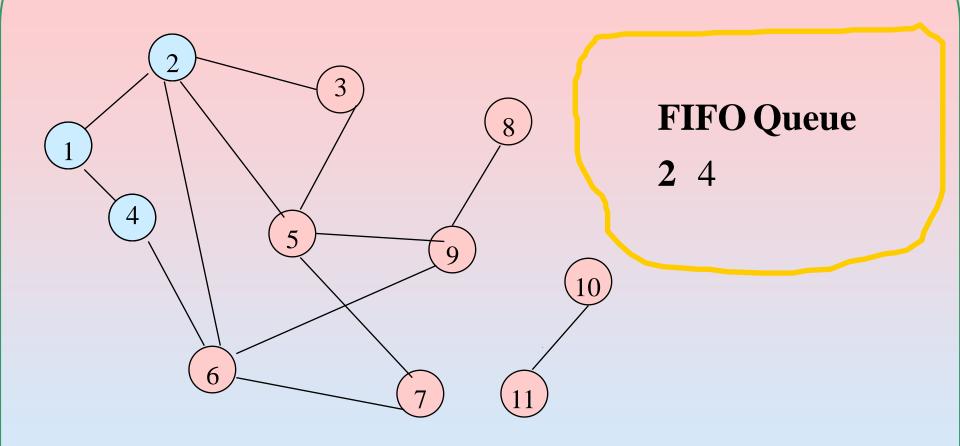


Remove 1 from Queue; visit adjacent unvisited vertices; put in Queue.



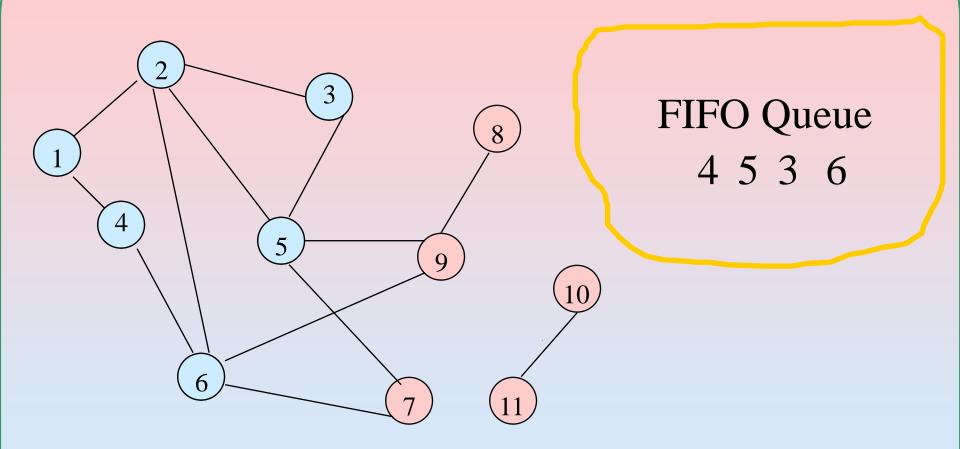
Remove 1 from Queue; visit adjacent unvisited vertices;

put in Queue.

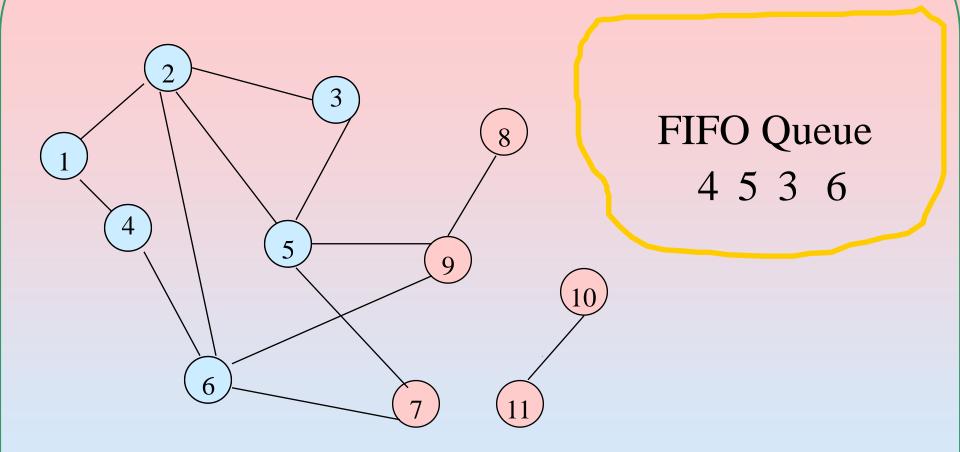


Remove 2 from Queue; visit adjacent unvisited vertices;

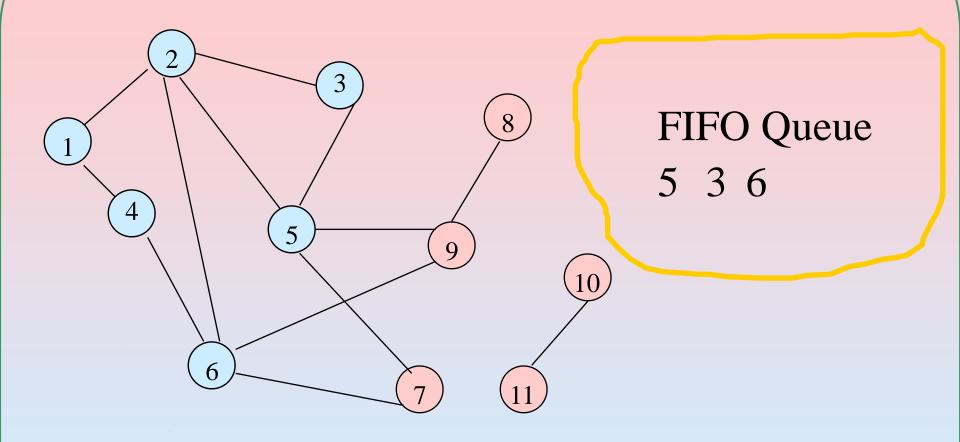
put in Queue.



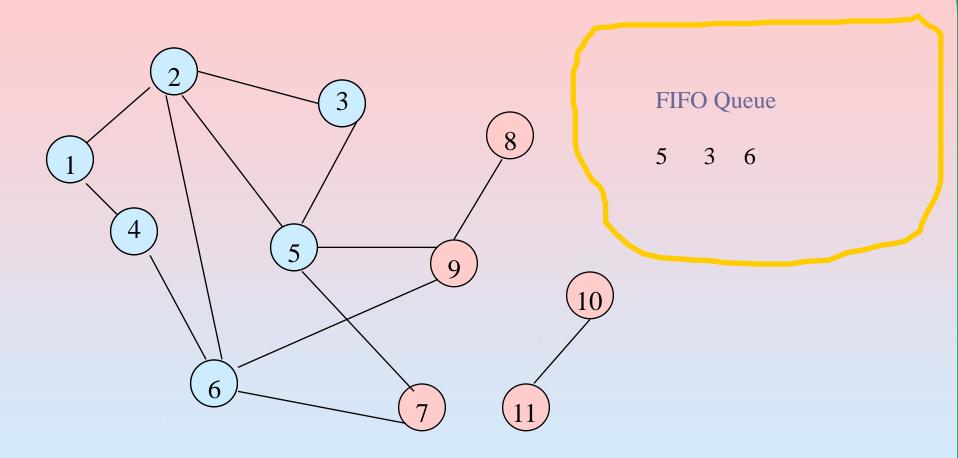
Remove 2 from Queue; visit adjacent unvisited vertices; put in Queue.



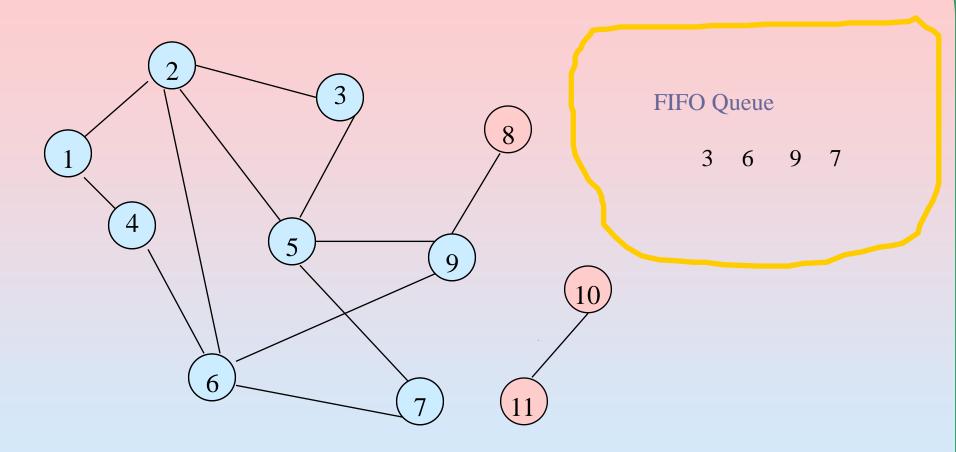
Remove 4 from Queue; visit adjacent unvisited vertices; put in Queue.



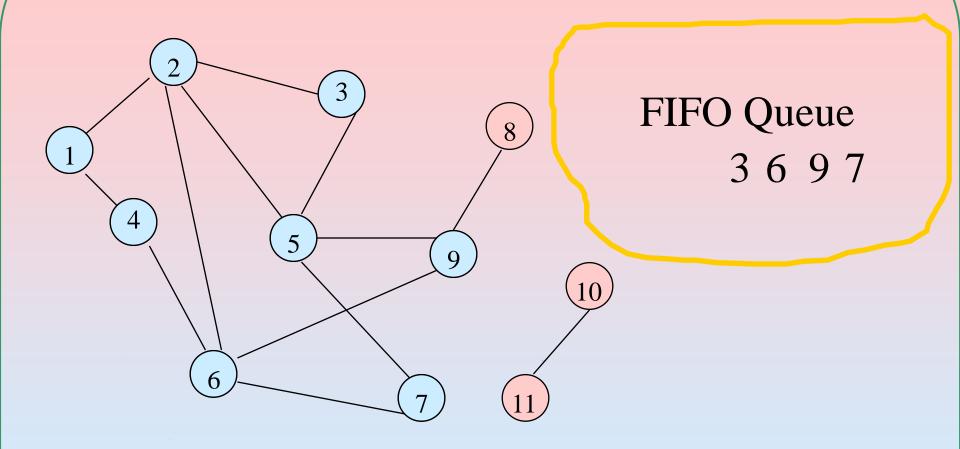
Remove 4 from Queue; visit adjacent unvisited vertices; put in Queue.



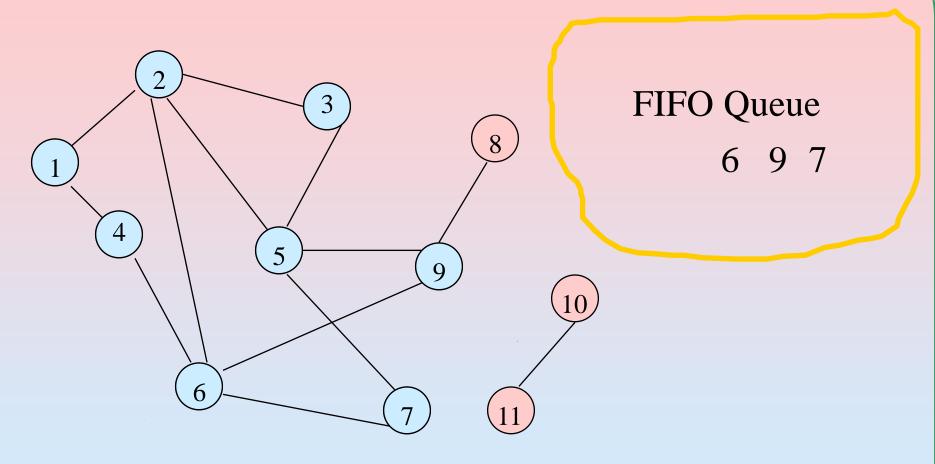
Remove 5 from Queue; visit adjacent unvisited vertices; put in Queue.



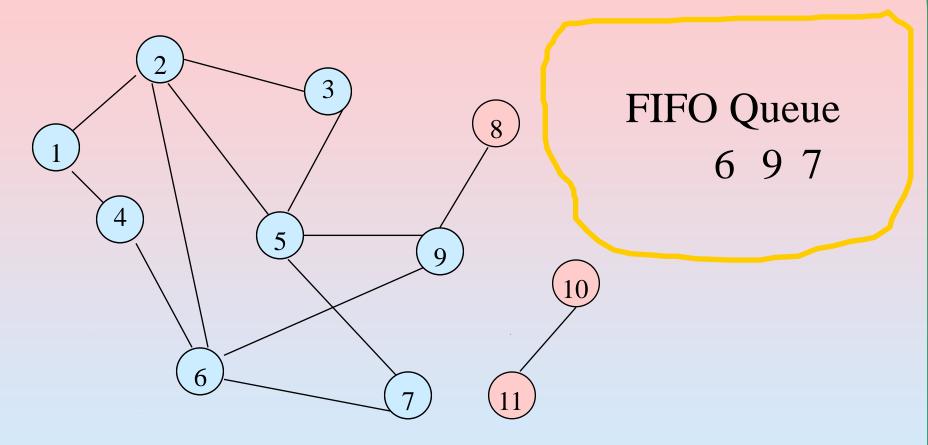
Remove 5 from Queue; visit adjacent unvisited vertices; put in Queue.



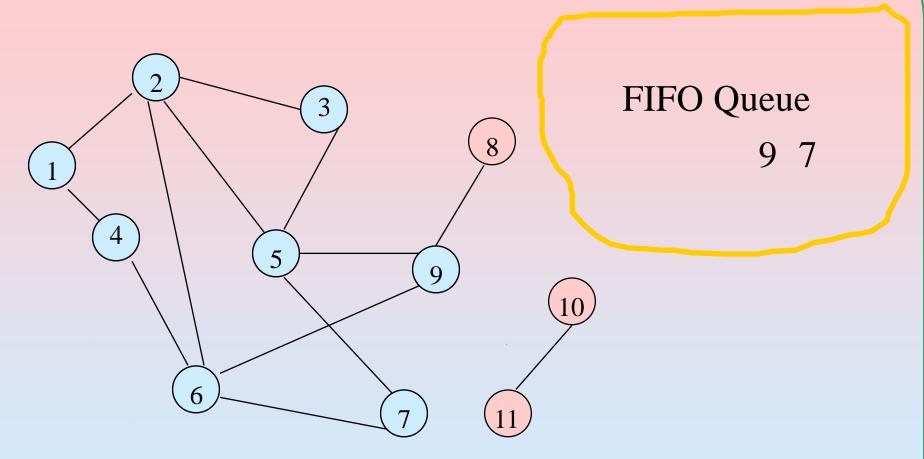
Remove 3 from Queue; visit adjacent unvisited vertices; put in Queue.



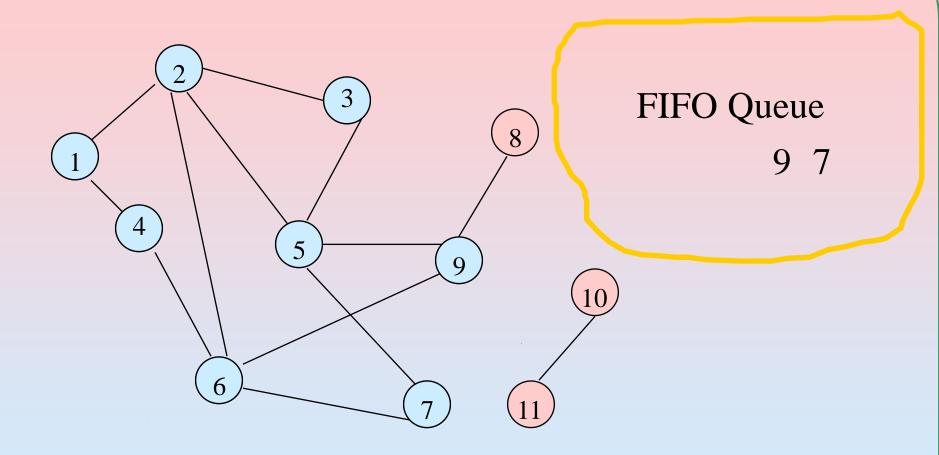
Remove 3 from Queue; visit adjacent unvisited vertices; put in Queue.



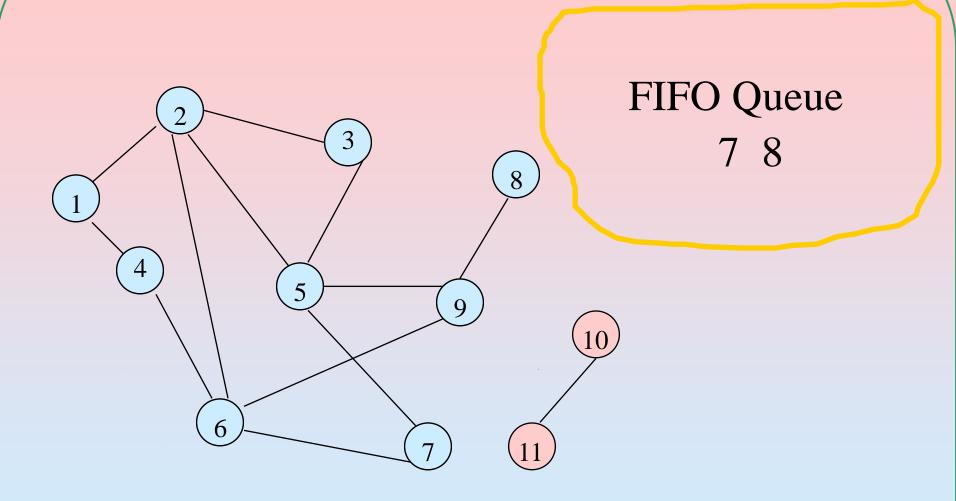
Remove 6 from Queue; visit adjacent unvisited vertices; put in Queue.



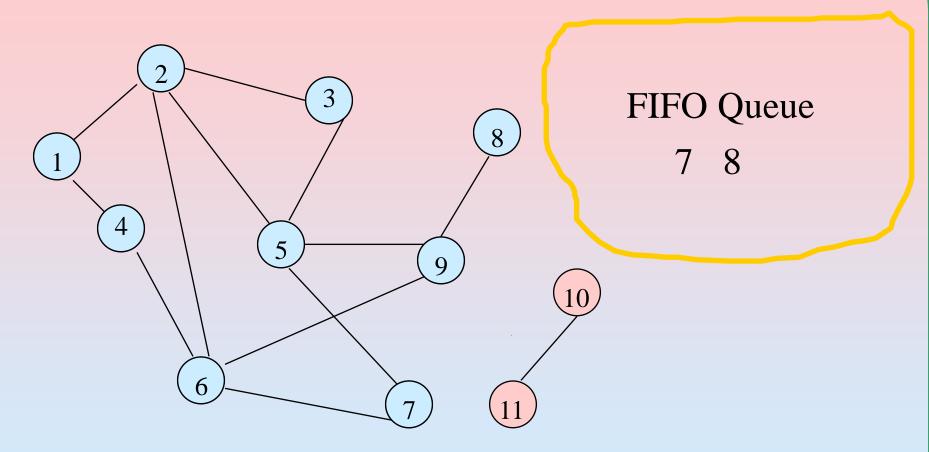
Remove 6 from Queue; visit adjacent unvisited vertices; put in Queue.



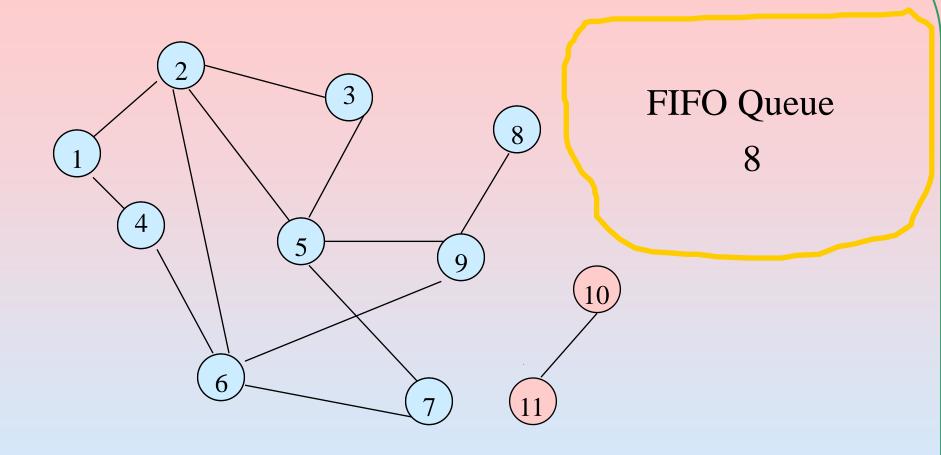
Remove 9 from Queue; visit adjacent unvisited vertices; put in Queue.



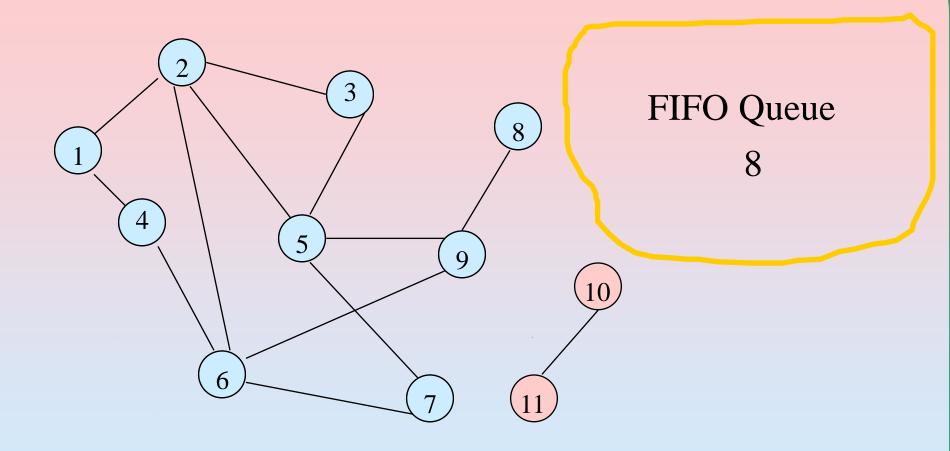
Remove 9 from Queue; visit adjacent unvisited vertices; put in Queue.



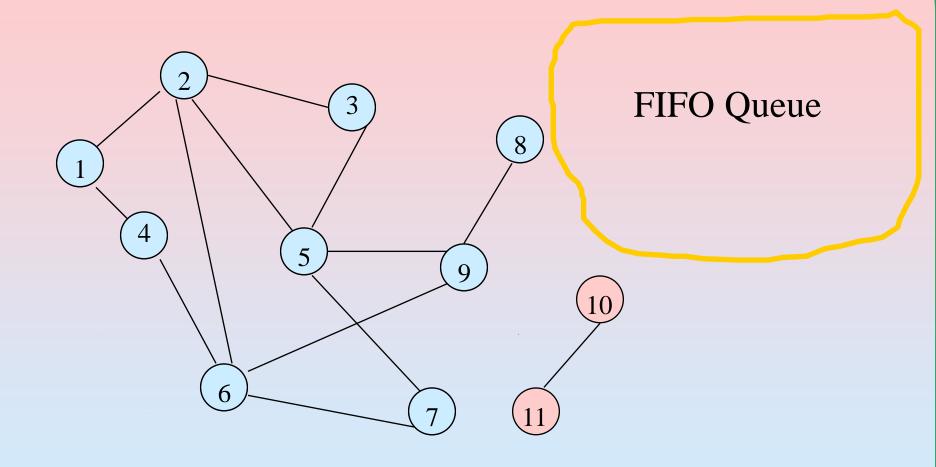
Remove 7 from Queue; visit adjacent unvisited vertices; put in Queue.



Remove 7 from Queue; visit adjacent unvisited vertices; put in Queue.



Remove 8 from Queue; visit adjacent unvisited vertices; put in Queue.



Queue is empty. Search terminates.

BFS - DFS

❖Breadth first search typically implemented with a Queue

❖ Depth first search typically implemented witha stack, implicit with recursion or iteratively
with an explicit stack