Universe in the Natural Model of Type Theory

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1 Types

Assume an inaccessible cardinal λ . Write **Set** for the category of all sets. Say that a set A is λ -small if $|A| < \lambda$. Write **Set** $_{\lambda}$ for the full subcategory of **Set** spanned by λ -small sets.

Let $\mathbb C$ be a small category, i.e. a category whose class of objects is a set and whose hom-classes are sets.

We write $\mathbf{Psh}(\mathbb{C})$ for the category of presheaves over \mathbb{C} ,

$$\mathbf{Psh}(\mathbb{C}) =_{\mathrm{def}} [\mathbb{C}^{\mathrm{op}}, \mathbf{Set}]$$

The Natural Model associated to a presentable map $tp: Tm \to Ty$ consists of

- contexts as objects $\Gamma, \Delta, \ldots \in \mathbb{C}$,
- a type in context $y(\Gamma)$ as a map $A: y(\Gamma) \to \mathsf{Ty}$,
- a term of type A in context Γ as a map $a: y(\Gamma) \to Tm$ such that



commutes,

• an operation called "context extension" which given a context Γ and a type $A \colon \mathsf{y}(\Gamma) \to \mathsf{T}\mathsf{y}$ produces a context $\Gamma \cdot A$ which fits into a pullback diagram below.

$$\begin{array}{ccc} \mathsf{y}(\Gamma.A) & \longrightarrow \mathsf{Tm} \\ \downarrow & & \downarrow \\ \mathsf{y}(\Gamma) & \longrightarrow_A & \mathsf{Ty} \end{array}$$

Remark. Sometimes, we first construct a presheaf X over Γ and observe that it can be classified by a map into Ty. We write

$$\begin{array}{c} X \longrightarrow \mathsf{Tm} \\ \downarrow \\ \mathsf{y}(\Gamma) \xrightarrow{\Gamma X \to \mathsf{Ty}} \mathsf{Ty} \end{array}$$

to express this situation, i.e. $X \cong y(\Gamma \cdot \lceil X \rceil)$.

2 A type of small types

We now wish to formulate a condition that allows us to have a type of small types, written U, not just *judgement* expressing that something is a type. With this notation, the judgements that we would like to derive is

$$\mathsf{U}\colon \mathsf{Ty} \qquad \frac{a\colon \mathsf{U}}{\mathsf{El}(a)\colon \mathsf{Ty}}$$

(A sufficient and natural condition for this seems to be that we now have another inaccessible cardinal κ , with $\kappa < \lambda$.)

In the Natural Model, a universe U is postulated by a map

$$\pi \colon \mathsf{E} \to \mathsf{U}$$

In the Natural Model:

• There is a pullback diagram of the form

$$\begin{array}{c} \mathsf{U} \longrightarrow \mathsf{Tm} \\ \downarrow \\ \downarrow \\ 1 \xrightarrow{} \mathsf{Ty} \end{array}$$

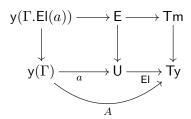
• There is an inclusion of U into Ty

$$\mathsf{EI}\colon\mathsf{U}\rightarrowtail\mathsf{Ty}$$

• $\pi: E \to U$ is obtained as pullback of tp; There is a pullback diagram

$$\begin{array}{c} E {\longmapsto} \operatorname{Tm} \\ \downarrow \qquad \qquad \downarrow \\ \operatorname{U} {\longmapsto} \operatorname{Ty} \end{array}$$

With the notation above, we get



Both squares above are pullback squares.

- 3 The Universe in Embedded Type Theory (HoTT0) and the relationship to the Natural Model
- 4 Groupoid Model of HoTT

In this section we construct a natural model in **Psh**(**grpd**) the presheaf category indexed by the category **grpd** of (small) groupoids. We will build the classifier for display maps in the style of Hofmann and Streicher [HS98] and Awodey [Awo23]. To interpret the type constructors, we will make use of the weak factorization system on **grpd** - which comes from restricting the "classical Quillen model structure" on **cat** [Joy] to **grpd**.

Notation. We will have two universe sizes - one small and one large. We denote the category of small sets as **set** and the large sets as **Set** (in the previous sections this would have been \mathbf{Set}_{λ} and \mathbf{Set} respectively). We denote the category of small categories as **cat** and the large categories as **Cat**. We denote the category of small groupoids as **grpd**. The category of small pointed groupoids will be \mathbf{grpd}_{\bullet} .

We are primarily working in the category of large presheaves indexed by small groupoids, which we will denote by

$$Psh(grpd) = [grpd^{op}, Set]$$

In this section, Tm and Ty and so on will refer to the natural model semantics in this specific model.

Definition 4.1 (The disply map classifier). We would like to define a natural transformation in **Psh**(**grpd**)

$$\mathsf{tp}\colon \mathsf{Tm}\to \mathsf{Ty}$$

with representable fibers.

Consider the functor that forgets the point

$$U : \mathbf{grpd}_{\bullet} \to \mathbf{grpd}$$
 in Cat.

If we apply the Yoneda embedding y: $Cat \rightarrow Psh(Cat)$ to U we obtain

$$U \circ : [-, \mathbf{grpd}_{\bullet}] \to [-, \mathbf{grpd}]$$
 in $\mathbf{Psh}(\mathbf{Cat})$.

Since any small groupoid is also a large category $\mathbf{grpd} \hookrightarrow \mathbf{Cat}$, we can restrict \mathbf{Cat} indexed presheaves to be \mathbf{grpd} indexed presheaves. We define $\mathsf{tp} \colon \mathsf{Tm} \to \mathsf{Ty}$ as the image of $U \circ \mathsf{under}$ this restriction.

$$\begin{array}{ccc} \mathbf{Cat} & \xrightarrow{y} & \mathbf{Psh}(\mathbf{Cat}) & \xrightarrow{\mathsf{res}} & \mathbf{Psh}(\mathbf{grpd}) \\ \\ \mathbf{grpd} & \longmapsto & [-,\mathbf{grpd}] & \longmapsto & \mathsf{Ty} \end{array}$$

Note that Tm and Ty are not representable in $\mathsf{Psh}(\mathsf{grpd})$.

Remark 4.2. By Yoneda we can identify maps with representable domain into the type classifier

$$A: y\Gamma \to \mathsf{Ty}$$
 in $\mathbf{Psh}(\mathbf{grpd})$

with functors

$$A:\Gamma\to\mathbf{grpd}$$
 in Cat

Definition 4.3 (Grothendieck construction). From \mathbb{C} a small category and $F: \mathbb{C} \to \mathbf{Cat}$ a functor, we construct a small category $\int F$. For any c in \mathbb{C} we refer to Fc as the fiber over c. The objects of $\int F$ consist of pairs $(c \in \mathbb{C}, x \in Fc)$, and morphisms are those in \mathbb{C} preserving the points in the fiber. This makes the following pullback

$$(c,x) \longmapsto (Fc,x)$$

where Cat_• is the category of large pointed categories.

Definition 4.4 (Grothendieck construction for groupoids). Let Γ be a groupoid and $A \colon \Gamma \to \mathbf{grpd}$ a functor, we can compose F with the inclusion $i \colon \mathbf{grpd} \hookrightarrow \mathbf{Cat}$ and form the Grothendieck construction which we denote as

$$\Gamma \cdot A := \int i \circ A \qquad \operatorname{disp}_A \colon \Gamma \cdot A \to \Gamma$$

This is also a small groupoid since the underlying morphisms are from the groupoid Γ . Furthermore the pullback factors through (pointed) groupoids.

$$\begin{array}{cccc} \Gamma \cdot A & \longrightarrow \mathbf{grpd}_{\bullet} & \longrightarrow \mathbf{Cat}_{\bullet} \\ \downarrow & & \downarrow & & \downarrow \\ \Gamma & \longrightarrow \mathbf{grpd} & \longrightarrow \mathbf{Cat} \end{array}$$

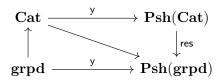
Corollary 4.5 (The display map classifier is presentable). For any small groupoid Γ and $A: y\Gamma \to Ty$, the pullback of tp along A can be given by the representable map $ydisp_A$.

$$\begin{array}{ccc} \mathsf{y}\Gamma \cdot A & \longrightarrow \mathsf{Tm} \\ \mathsf{disp}_A & & \mathsf{tp} \\ & \mathsf{y}\Gamma & \longrightarrow & \mathsf{Ty} \end{array}$$

Proof. Consider the pullback in Cat

$$\begin{array}{ccc} \Gamma \cdot A & \longrightarrow \mathbf{grpd}_{\bullet} \\ \downarrow & & \downarrow \\ \Gamma & \longrightarrow \mathbf{grpd} \end{array}$$

We send this square along $res \circ y$ in the following



The Yoneda embedding $y : \mathbf{Cat} \to \mathbf{Psh}(\mathbf{Cat})$ preserves pullbacks, as does res. \square

References

- [Awo23] Steve Awodey. On hofmann-streicher universes, 2023.
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- [Joy] André Joyal. Model structures on cat. https://ncatlab.org/joyalscatlab/published/Model+structures+on+Cat.