

# Universe in the Natural Model of Type Theory

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June 28, 2024

## 1 Types

Assume an inaccessible cardinal  $\lambda$ . Write **Set** for the category of all sets. Say that a set  $A$  is  $\lambda$ -small if  $|A| < \lambda$ . Write **Set** $_{\lambda}$  for the full subcategory of **Set** spanned by  $\lambda$ -small sets.

Let  $\mathbb{C}$  be a small category, i.e. a category whose class of objects is a set and whose hom-classes are sets.

We write  $\mathbf{Psh}(\mathbb{C})$  for the category of presheaves over  $\mathbb{C}$ ,

$$\mathbf{Psh}(\mathbb{C}) =_{\text{def}} [\mathbb{C}^{\text{op}}, \mathbf{Set}]$$

The Natural Model associated to a presentable map  $\text{tp}: \mathbf{Tm} \rightarrow \mathbf{Ty}$  consists of

- contexts as objects  $\Gamma, \Delta, \dots \in \mathbb{C}$ ,
- a type in context  $y(\Gamma)$  as a map  $A: y(\Gamma) \rightarrow \mathbf{Ty}$ ,
- a term of type  $A$  in context  $\Gamma$  as a map  $a: y(\Gamma) \rightarrow \mathbf{Tm}$  such that

$$\begin{array}{ccc} & \mathbf{Tm} & \\ & \uparrow a & \downarrow \text{tp} \\ \Gamma & \xrightarrow{A} & \mathbf{Ty} \end{array}$$

commutes,

- an operation called “context extension” which given a context  $\Gamma$  and a type  $A: y(\Gamma) \rightarrow \mathbf{Ty}$  produces a context  $\Gamma \cdot A$  which fits into a pullback diagram below.

$$\begin{array}{ccc} y(\Gamma \cdot A) & \longrightarrow & \mathbf{Tm} \\ \downarrow & & \downarrow \\ y(\Gamma) & \xrightarrow{A} & \mathbf{Ty} \end{array}$$

**Remark.** Sometimes, we first construct a presheaf  $X$  over  $\Gamma$  and observe that it can be classified by a map into  $\mathsf{Ty}$ . We write

$$\begin{array}{ccc} X & \longrightarrow & \mathsf{Tm} \\ \downarrow & & \downarrow \\ y(\Gamma) & \xrightarrow{\ulcorner X \urcorner} & \mathsf{Ty} \end{array}$$

to express this situation, i.e.  $X \cong y(\Gamma \cdot \ulcorner X \urcorner)$ .

## 2 A type of small types

We now wish to formulate a condition that allows us to have a type of small types, written  $\mathsf{U}$ , not just *judgement* expressing that something is a type. With this notation, the judgements that we would like to derive is

$$\mathsf{U} : \mathsf{Ty} \quad \frac{a : \mathsf{U}}{\mathsf{El}(a) : \mathsf{Ty}}$$

(A sufficient and natural condition for this seems to be that we now have another inaccessible cardinal  $\kappa$ , with  $\kappa < \lambda$ .)

In the Natural Model, a universe  $\mathsf{U}$  is postulated by a map

$$\pi : \mathsf{E} \rightarrow \mathsf{U}$$

In the Natural Model,

- There is a pullback diagram of the form

$$\begin{array}{ccc} \mathsf{U} & \longrightarrow & \mathsf{Tm} \\ \downarrow & & \downarrow \\ 1 & \xrightarrow{\ulcorner \mathsf{U} \urcorner} & \mathsf{Ty} \end{array}$$

- There is a map

$$\mathsf{El} : \mathsf{U} \rightarrow \mathsf{Ty}$$

- $\pi : \mathsf{E} \rightarrow \mathsf{U}$  is obtained as pullback of  $\mathsf{tp}$  along  $\mathsf{El}$ ; There is a pullback diagram

$$\begin{array}{ccc} E & \longrightarrow & \mathsf{Tm} \\ \pi \downarrow & & \downarrow \mathsf{tp} \\ \mathsf{U} & \xrightarrow{\mathsf{El}} & \mathsf{Ty} \end{array} \tag{2.1}$$

With the notation above, we get

$$\begin{array}{ccccc}
y(\Gamma.\text{El}(a)) & \longrightarrow & E & \longrightarrow & \mathsf{Tm} \\
\downarrow & & \downarrow & & \downarrow \\
y(\Gamma) & \xrightarrow{a} & U & \xrightarrow{\text{El}} & \mathsf{Ty} \\
& \searrow & \nearrow & & \\
& & A & & 
\end{array}$$

Both squares above are pullback squares.

### 3 Defining $\Pi$ and $\Sigma$ type formers for the universe $\pi: E \rightarrow U$

Take the pullback diagram eq. (2.1). That is a morphism in the category of polynomials. We have a cartesian natural transformation  $P_\pi \rightarrow P_{\text{tp}}$  induced by the pullback eq. (2.1). This cartesian natural transformation induces the diagrams of the left cube in below; all of the squares in the left cube are pullback squares.

Now, consider the right cube only. The right-side face of the cube is a pullback square due to the definition of  $\pi$ . The left-side is also because polynomials are left exact functors and therefore they preserve pullbacks. Observe also that the front face is a pullback square by the definition of  $\Pi$  for  $\mathsf{Ty}$ .

$$\begin{array}{ccccccc}
& & P_\pi E & \longrightarrow & P_{\text{tp}} E & \longrightarrow & E \\
& \swarrow & \downarrow & & \downarrow & & \downarrow \\
P_\pi \mathsf{Tm} & \longrightarrow & P_{\text{tp}} \mathsf{Tm} & \xrightarrow{\lambda} & \mathsf{Tm} & & \\
\downarrow & & \downarrow & & \downarrow & & \downarrow \\
& & P_\pi U & \longrightarrow & P_{\text{tp}} U & \xrightarrow{\quad \quad} & U \\
& \swarrow & \downarrow & & \downarrow & & \downarrow \\
P_\pi \mathsf{Ty} & \longrightarrow & P_{\text{tp}} \mathsf{Ty} & \xrightarrow{\Pi} & \mathsf{Ty} & & 
\end{array}$$

We define  $\Pi_U: P_\pi U \rightarrow U$  as a dashed arrow which makes the bottom square in the right cube commute, that is the following square commutes.

$$\begin{array}{ccc}
P_{\text{tp}} U & \longrightarrow & P_{\text{tp}} \mathsf{Ty} \\
\downarrow \Pi_U & & \downarrow \Pi_{\mathsf{Ty}} \\
U & \xrightarrow{\text{El}} & \mathsf{Ty}
\end{array} \tag{3.1}$$

By the universal property of the pullback which defines  $\pi$  we get a unique arrow  $P_{\text{tp}}E \rightarrow E$  which we shall name  $\lambda_U$ . By the pullback pasting lemma it follows that the square involving  $\Pi_U$  and  $\lambda_U$  is a pullback square.

$$\begin{array}{ccc}
 P_{\text{tp}}E & \xrightarrow{\lambda_U} & E \\
 P_{\text{tp}}\pi \downarrow & & \downarrow P_{\text{tp}}\pi \\
 P_{\text{tp}}U & \xrightarrow{\Pi_U} & U
 \end{array} \tag{3.2}$$

This concludes the construction of  $\Pi$  type former for the universe  $U$ . The only data we needed to supply for the definition of  $\Pi_U$  was a lift of  $\Pi: P_{\text{tp}}\text{Ty} \rightarrow \text{Ty}$  to  $U$ .

## 4 The Universe in Embedded Type Theory (HoTT0) and the relationship to the Natural Model