

# Universe in the Natural Model of Type Theory

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## 1 Types

Assume an inaccessible cardinal  $\lambda$ . Write **Set** for the category of all sets. Say that a set  $A$  is  $\lambda$ -small if  $|A| < \lambda$ . Write **Set** $_{\lambda}$  for the full subcategory of **Set** spanned by  $\lambda$ -small sets.

Let  $\mathbb{C}$  be a small category, i.e. a category whose class of objects is a set and whose hom-classes are sets.

We write **Psh**( $\mathbb{C}$ ) for the category of presheaves over  $\mathbb{C}$ ,

$$\mathbf{Psh}(\mathbb{C}) =_{\text{def}} [\mathbb{C}^{\text{op}}, \mathbf{Set}]$$

The Natural Model associated to a presentable map  $\text{tp}: \mathsf{Tm} \rightarrow \mathsf{Ty}$  consists of

- contexts as objects  $\Gamma, \Delta, \dots \in \mathbb{C}$ ,
- a type in context  $y(\Gamma)$  as a map  $A: y(\Gamma) \rightarrow \mathsf{Ty}$ ,
- a term of type  $A$  in context  $\Gamma$  as a map  $a: y(\Gamma) \rightarrow \mathsf{Tm}$  such that

$$\begin{array}{ccc} & \mathsf{Tm} & \\ & \uparrow a & \downarrow \text{tp} \\ \Gamma & \xrightarrow{A} & \mathsf{Ty} \end{array}$$

commutes,

- an operation called “context extension” which given a context  $\Gamma$  and a type  $A: y(\Gamma) \rightarrow \mathsf{Ty}$  produces a context  $\Gamma \cdot A$  which fits into a pullback diagram below.

$$\begin{array}{ccc} y(\Gamma \cdot A) & \longrightarrow & \mathsf{Tm} \\ \downarrow & & \downarrow \\ y(\Gamma) & \xrightarrow{A} & \mathsf{Ty} \end{array}$$

**Remark.** Sometimes, we first construct a presheaf  $X$  over  $\Gamma$  and observe that it can be classified by a map into  $\mathsf{Ty}$ . We write

$$\begin{array}{ccc} X & \longrightarrow & \mathsf{Tm} \\ \downarrow & & \downarrow \\ y(\Gamma) & \xrightarrow{\ulcorner X \urcorner} & \mathsf{Ty} \end{array}$$

to express this situation, i.e.  $X \cong y(\Gamma \cdot \ulcorner X \urcorner)$ .

## 2 A type of small types

We now wish to formulate a condition that allows us to have a type of small types, written  $\mathsf{U}$ , not just *judgement* expressing that something is a type. With this notation, the judgements that we would like to derive is

$$\mathsf{U} : \mathsf{Ty} \quad \frac{a : \mathsf{U}}{\mathsf{El}(a) : \mathsf{Ty}}$$

(A sufficient and natural condition for this seems to be that we now have another inaccessible cardinal  $\kappa$ , with  $\kappa < \lambda$ .)

In the Natural Model, a universe  $\mathsf{U}$  is postulated by a map

$$\pi : \mathsf{E} \rightarrow \mathsf{U}$$

In the Natural Model:

- There is a pullback diagram of the form

$$\begin{array}{ccc} \mathsf{U} & \longrightarrow & \mathsf{Tm} \\ \downarrow & & \downarrow \\ 1 & \xrightarrow{\ulcorner \mathsf{U} \urcorner} & \mathsf{Ty} \end{array}$$

- There is an inclusion of  $\mathsf{U}$  into  $\mathsf{Ty}$

$$\mathsf{El} : \mathsf{U} \rightarrowtail \mathsf{Ty}$$

- $\pi : \mathsf{E} \rightarrow \mathsf{U}$  is obtained as pullback of  $\mathsf{tp}$ ; There is a pullback diagram

$$\begin{array}{ccc} \mathsf{E} & \twoheadrightarrow & \mathsf{Tm} \\ \downarrow & & \downarrow \\ \mathsf{U} & \twoheadrightarrow_{\mathsf{El}} & \mathsf{Ty} \end{array}$$

With the notation above, we get

$$\begin{array}{ccccc}
 y(\Gamma, \text{El}(a)) & \longrightarrow & E & \longrightarrow & Tm \\
 \downarrow & & \downarrow & & \downarrow \\
 y(\Gamma) & \xrightarrow{a} & U & \xrightarrow{\text{El}} & Ty \\
 & \searrow \scriptstyle A & \nearrow & & 
 \end{array}$$

Both squares above are pullback squares.

### 3 The Universe in Embedded Type Theory (HoTT0) and the relationship to the Natural Model

### 4 Groupoid Model of HoTT

In this section we construct a natural model in  $\mathbf{Psh}(\mathbf{grpd})$  the presheaf category indexed by the category  $\mathbf{grpd}$  of (small) groupoids. We will build the classifier for display maps in the style of Hofmann and Streicher [HS98] and Awodey [Awo23]. To interpret the type constructors, we will make use of the weak factorization system on  $\mathbf{grpd}$  - which comes from restricting the “classical Quillen model structure” on  $\mathbf{cat}$  [Joy] to  $\mathbf{grpd}$ .

*Notation.* We will have two universe sizes - one small and one large. We denote the category of small sets as  $\mathbf{set}$  and the large sets as  $\mathbf{Set}$  (in the previous sections this would have been  $\mathbf{Set}_\lambda$  and  $\mathbf{Set}$  respectively). We denote the category of small categories as  $\mathbf{cat}$  and the large categories as  $\mathbf{Cat}$ . We denote the category of small groupoids as  $\mathbf{grpd}$ . The category of small pointed groupoids will be  $\mathbf{grpd}_\bullet$ .

We are primarily working in the category of large presheaves indexed by small groupoids, which we will denote by

$$\mathbf{Psh}(\mathbf{grpd}) = [\mathbf{grpd}^{\text{op}}, \mathbf{Set}]$$

In this section,  $\mathbf{Tm}$  and  $\mathbf{Ty}$  and so on will refer to the natural model semantics in this specific model.

**Definition 4.1** (The display map classifier). We would like to define a natural transformation in  $\mathbf{Psh}(\mathbf{grpd})$

$$\mathbf{tp}: \mathbf{Tm} \rightarrow \mathbf{Ty}$$

with representable fibers.

Consider the functor that forgets the point

$$U: \mathbf{grpd}_\bullet \rightarrow \mathbf{grpd} \quad \text{in} \quad \mathbf{Cat}.$$

If we apply the Yoneda embedding  $y: \mathbf{Cat} \rightarrow \mathbf{Psh}(\mathbf{Cat})$  to  $U$  we obtain

$$U \circ [-, \mathbf{grpd}_\bullet] \rightarrow [-, \mathbf{grpd}] \quad \text{in} \quad \mathbf{Psh}(\mathbf{Cat}).$$

Since any small groupoid is also a large category  $\mathbf{grpd} \hookrightarrow \mathbf{Cat}$ , we can restrict  $\mathbf{Cat}$  indexed presheaves to be  $\mathbf{grpd}$  indexed presheaves. We define  $\mathbf{tp}: \mathbf{Tm} \rightarrow \mathbf{Ty}$  as the image of  $U \circ$  under this restriction.

$$\begin{array}{ccc} \mathbf{Cat} & \xrightarrow{y} & \mathbf{Psh}(\mathbf{Cat}) \xrightarrow{\text{res}} \mathbf{Psh}(\mathbf{grpd}) \\ \mathbf{grpd} & \longmapsto & [-, \mathbf{grpd}] \longmapsto \mathbf{Ty} \end{array}$$

Note that  $\mathbf{Tm}$  and  $\mathbf{Ty}$  are not representable in  $\mathbf{Psh}(\mathbf{grpd})$ .

*Remark 4.2.* By Yoneda we can identify maps with representable domain into the type classifier

$$A: y\Gamma \rightarrow \mathbf{Ty} \quad \text{in} \quad \mathbf{Psh}(\mathbf{grpd})$$

with functors

$$A: \Gamma \rightarrow \mathbf{grpd} \quad \text{in} \quad \mathbf{Cat}$$

**Definition 4.3** (Grothendieck construction). From  $\mathbb{C}$  a small category and  $F : \mathbb{C} \rightarrow \mathbf{Cat}$  a functor, we construct a small category  $\int F$ . For any  $c$  in  $\mathbb{C}$  we refer to  $Fc$  as the fiber over  $c$ . The objects of  $\int F$  consist of pairs  $(c \in \mathbb{C}, x \in Fc)$ , and morphisms are those in  $\mathbb{C}$  preserving the points in the fiber. This makes the following pullback

$$(c, x) \longmapsto (Fc, x)$$

$$\begin{array}{ccccc} (c, x) & & \int F & \longrightarrow & \mathbf{Cat}_\bullet & & (C, c) \\ \downarrow & & \downarrow & \lrcorner & \downarrow & & \downarrow \\ c & & \mathbb{C} & \xrightarrow{F} & \mathbf{Cat} & & C \end{array}$$

where  $\mathbf{Cat}_\bullet$  is the category of large pointed categories.

**Definition 4.4** (Grothendieck construction for groupoids). Let  $\Gamma$  be a groupoid and  $A : \Gamma \rightarrow \mathbf{grpd}$  a functor, we can compose  $F$  with the inclusion  $i : \mathbf{grpd} \hookrightarrow \mathbf{Cat}$  and form the Grothendieck construction which we denote as

$$\Gamma \cdot A := \int i \circ A \quad \text{disp}_A : \Gamma \cdot A \rightarrow \Gamma$$

This is also a small groupoid since the underlying morphisms are from the groupoid  $\Gamma$ . Furthermore the pullback factors through (pointed) groupoids.

$$\begin{array}{ccccc} \Gamma \cdot A & \longrightarrow & \mathbf{grpd}_\bullet & \longrightarrow & \mathbf{Cat}_\bullet \\ \downarrow & \lrcorner & \downarrow & \lrcorner & \downarrow \\ \Gamma & \xrightarrow{A} & \mathbf{grpd} & \longrightarrow & \mathbf{Cat} \end{array}$$

**Corollary 4.5** (The display map classifier is presentable). *For any small groupoid  $\Gamma$  and  $A : y\Gamma \rightarrow \mathbf{Ty}$ , the pullback of  $\mathbf{tp}$  along  $A$  can be given by the representable map  $y\text{disp}_A$ .*

$$\begin{array}{ccc} y\Gamma \cdot A & \longrightarrow & \mathbf{Tm} \\ \text{disp}_A \downarrow & \lrcorner & \downarrow \mathbf{tp} \\ y\Gamma & \xrightarrow{A} & \mathbf{Ty} \end{array}$$

*Proof.* Consider the pullback in **Cat**

$$\begin{array}{ccc} \Gamma \cdot A & \longrightarrow & \mathbf{grpd}_\bullet \\ \downarrow & \lrcorner & \downarrow \\ \Gamma & \xrightarrow{A} & \mathbf{grpd} \end{array}$$

We send this square along  $\mathbf{res} \circ \mathbf{y}$  in the following

$$\begin{array}{ccc} \mathbf{Cat} & \xrightarrow{\mathbf{y}} & \mathbf{Psh}(\mathbf{Cat}) \\ \uparrow & \searrow & \downarrow \mathbf{res} \\ \mathbf{grpd} & \xrightarrow{\mathbf{y}} & \mathbf{Psh}(\mathbf{grpd}) \end{array}$$

The Yoneda embedding  $\mathbf{y} : \mathbf{Cat} \rightarrow \mathbf{Psh}(\mathbf{Cat})$  preserves pullbacks, as does  $\mathbf{res}$ , since it is a right adjoint (left Kan extension  $\iota_! \dashv \mathbf{res}_\iota$ ).  $\square$

## References

- [Awo23] Steve Awodey. On hofmann-streicher universes, 2023.
- [HS98] Martin Hofmann and Thomas Streicher. The groupoid interpretation of type theory. In *Twenty-five years of constructive type theory (Venice, 1995)*, volume 36 of *Oxford Logic Guides*, pages 83–111. Oxford Univ. Press, New York, 1998.
- [Joy] André Joyal. Model structures on cat. <https://ncatlab.org/joyalscatlab/published/Model+structures+on+Cat>.