

36) This algorithm uses O(mtn) time for BFS algorithm. 3c) In order for there to be acycle, you would have to go from U to V, every other node in U then V. When you complete the cycle, you're back to the node in W. Compare each node from U as in V. Bipartite group has two groups, so you can only have ever groups C = (v1, 1/2, ..., vn, v1) n is odd v1 EX, 1/2 EY, 1/3 EX, ... ViEX It i is odd Bipartite graphs have no odd cycle Vn E X VnV, EE(G)